MONITORING NODE SELECTION ALGORITHM FOR INTRUSION DETECTION IN CONGESTED SENSOR NETWORK

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Abstract: Since wireless resources are limited, an efficient way of utilizing wireless resources is needed in selecting IDSs. We propose a monitoring node selection scheme for intrusion detection in congested wireless sensor network. Network congestion is an important issue in mobile network. The network congestion does not guarantee a proper detection rate and congested networks should cause an unreliable network. We consider congested intrusion detection tasks by queuing theory. We confirm that proposed algorithm guarantee QoS of monitoring tasks and reliable sensor networks.

1   INTRODUCTION

In recent years, as wireless network techniques develop rapidly, security becomes one of the major problems in wireless network. Wireless networks are notably vulnerable to intrusions, as they operate in open medium and don't have any centralized security systems. In contrast to wired networks, wireless networks need a special method to detect intrusions.

An Intrusion Detection System (IDS) for wireless networks is widely employed for security purpose to detect illegal intrusions. An IDS is a system that tries to detect intrusions in the network using statistical behavior models. To detect intrusions, all nodes should be monitored by IDS nodes. Only considering security, every node had better perform intrusion detection. However, this approach is inefficient in terms of wireless resources such as network bandwidth and node battery. If all nodes implement intrusion detection processing, the wireless resources are inefficiently used and some nodes are suffered from battery depletion, because a node acting as an IDS which overhears and analyses all packets within monitoring range consumes additional resources. Since wireless network resources such as battery and bandwidth are limited, an efficient way of utilizing these wireless resources is needed in construction IDSs [1], [2].

In this paper, we apply an IDS node selection scheme to a wireless sensor network (WSN). A wireless sensor network is a wireless network consisting of distributed sensors to cooperatively monitor physical or environmental conditions, such as temperature, sound, vibration, pressure, motion or pollutants, at different locations [3], [4]. Wireless sensor networks are now used in many industrial and civilian application areas [3], [5]. Wireless sensor networks have constraints on resources such as energy, memory, computational speed and bandwidth [3]. Therefore, an efficient IDS selection algorithm is required in wireless sensor network.

There are two related works to select monitoring nodes for intrusion detection before, distributed IDS (DIDS) and lifetime-enhancing monitoring node selection (LES) [1], [2]. These schemes efficiently use wireless resources such as network lifetime or battery consumption of whole network. However, previous researches did not consider congestion of monitored packets in a buffer of IDS. Our work differs from others by considering congestion of monitoring tasks. In recent years, data transmission speed becomes faster and faster. If the data transmission rate is faster than monitoring rate, the unmonitored packets increase in a buffer of IDS. When congestion occurs in a node due to high packet arrival rates, the IDS cannot afford to monitor the all arrival packets and the buffer overflow happens. So the IDS becomes ineffective and deteriorated. Hence, it is desirable to assign IDS

node to the network properly so that prevent the denial of IDS service due to the overflow.

In this paper, we propose a monitoring node selection algorithm for intrusion detection in congested sensor network. By using queuing theory, we prevent congested monitoring situations. In congested system, our proposed scheme has the longest network lifetime and the smallest total battery consumption than other existing schemes. By preventing congested monitoring tasks, our algorithm guarantees QoS of monitoring tasks and reliable sensor networks.

2 ALGORITHM

We propose an IDS node selection scheme based on two requirements. First, to monitor all nodes in the network, every node should be in the monitoring coverage of IDS. Second, to prevent the congestion of monitored packets, we design the IDS nodes placement scheme considering congested systems. For satisfying the first requirement, we apply a set covering problem (SCP) to the IDS nodes selection scheme.

The SCP is a classical question in computer science and complexity theory. The SCP selects a minimum number of sets that contain all elements and additionally minimizes the cost of the sets. Therefore, the SCP guarantees that every element is covered by at least one server at minimal total cost. To cover all nodes by minimal IDS, we propose a formulation using the SCP. The formulation of the IDS node placement scheme using the SCP is as follows;

$$\min \quad \sum_{j=1}^{n} c_j x_j$$

$$s.t \quad \sum_{j=1}^{n} a_{ij} x_j \geq 1 \quad \forall i \in N$$

$$x_j \in \{0,1\} \quad \forall j \in N \quad . \ (1)$$

Formulation (1) is a typical SCP formulation. Binary variable x_j is one if node j is IDS node, and zero otherwise. Like figure 1, binary variable $\underline{a_{ij}}$ is one if node j is in the transmission range of node i, and zero otherwise. In the typical SCP, c_j is the cost which is needed to select server j. In this problem, every node has same weight. Therefore, we define c_j of every node as 1. We define the set N as the set of all nodes in the network.
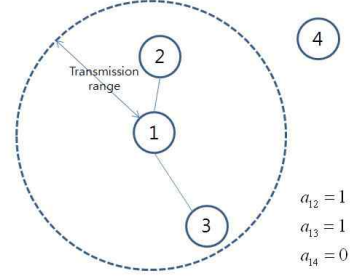

Figure 1: Transmission range of a wireless node

Formulation (1) satisfies the first requirement. Then, we discuss a constraint considering congested systems. An implicit assumption in traditional SCP is that each node in the coverage of a server always receive satisfied service. However, when a server suffers from congestion by excessive demand, some users are not able to receive satisfied service in the real situation. Especially, if an IDS node suffers from congestion of monitored packets, intrusion detection efficiency is reduced and battery consumption of the IDS is high. In order to guarantee high detection rate and use efficiently limited wireless resources, considering congested systems is important. To prevent congestions, any packet should not stand in waiting line in the buffer of IDS nodes for a time longer than a given time-limit [6]. The constraint which considers congestion is as follows;

$$P[\ waiting\ time\ at\ IDS\ node\ j \leq \tau\ ] \geq \alpha \quad \forall j \ . \ (2)$$

Constraint (2) makes the total time spent by a packet at the IDS node shorter than equal to τ with probability of at least α. The variables τ and α are predefined time and probability. In order to express constraint (2) as a numerical formula, we use the queuing theory [6]. In this paper, we make an assumption that an packet arrival rate from node i to j appears according to a poisson process with intensity f_{ij}. Also, we assume an exponentially distributed monitoring service time, with an average rate of μ_j. This is a reasonable assumption, since IDS systems behave as M/M/1 queuing systems. As we assume a M/M/1 queuing system, we are able to use the well known results for a M/M/1 queuing system for each IDS and its allocated nodes [6]. Rewriting constraint (2) as a numerical formula, we get

$$\sum_{i=1}^{n} f_{ij} a_{ij} x_j \leq \mu_j + \frac{1}{\tau} \ln(1-\alpha) \quad \forall j \in N \ . \ (3)$$

Adding constraint (3) to formulation (1), we finally get our proposed formulation as follows;

$$\min \quad \sum_{j=1}^{n} c_j x_j$$

$$s.t \quad \sum_{j=1}^{n} a_{ij} x_j \geq 1 \quad \forall i \in N$$

$$\sum_{i=1}^{n} f_{ij} a_{ij} x_j \leq D_j$$

$$where, D_j = \mu_j + \frac{1}{\tau} \ln(1-\alpha) \quad \forall j \in N$$

$$x_j \in \{0,1\} \quad \forall j \in N \tag{4}$$

3   PERFORMANCE MEASURE

The normal node uses their energy for transmission and receiving. Additionally, the monitoring node overhears packets within monitoring range, and analyzes them to detect intrusions. The battery consumed by a monitoring node is calculated as:

$$E = (m^t s^t + b^t) + (m^r s^r + b^r) + (m^o s^o + b^o) + (m^m s^m + b^m) \tag{5}$$

where, s^t, s^r, s^o and s^m, respectively, represent the packet sizes for operations : transmission, receiving, overhearing and monitoring. Factors m and b are variable and fixed energy costs for each operation [7]. We will show numerical simulations using equation (5).

For comparison with other schemes, we analysis two performance measures; network lifetime and sum of remain battery. The network lifetime defined as the duration of time until the first node runs out of battery. The network time is a meaningful performance measure in a sense that a node which has no battery cannot communicate with any other nodes and it causes unreliable network [1]. The sum of remaining battery indicates total battery consumption of whole network. By analysing this measure, we should know efficiency of utilizing wireless resources

4   SIMULATION RESULTS

In the simulation, we consider a congested wireless sensor network consisting of 20 nodes by MATLAB. For the purpose of comparison, we consider DIDS and LES mentioned above. We generate certain amount of packets destined for random nodes with a packet size of 512 bytes. Also, we assume that all nodes have the same initial battery. The simulation results are the average of 100 trials.
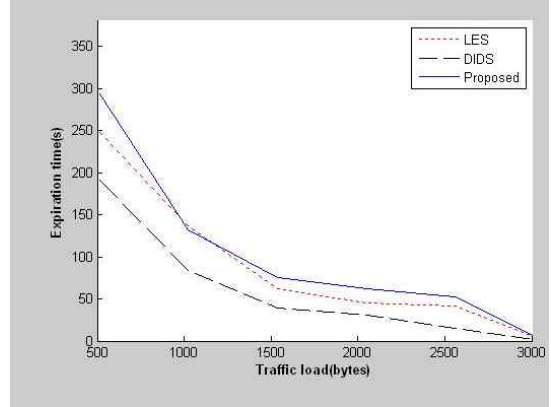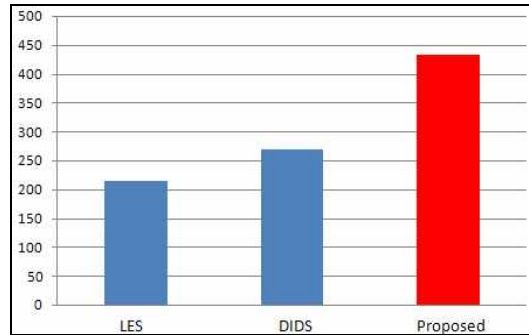

Figure 2: Network lifetime

Figure 3: Sum of remain battery

Figure 2 shows the average network lifetime in congested systems and figure 3 shows the sum of remain battery in congested network. In figure 2, the network lifetime of proposed scheme is longer than or similar to those of LES and DIDS. The total battery consumption in proposed scheme is the lowest among three schemes. By experiments, proposed scheme has superior performance than other existing schemes in congested network.

5 CONCLUSIONS

In this paper, we propose an IDS selection algorithm for intrusion detection in congested sensor network. In order to detect intrusions, the packets be buffered in the queue of a monitoring node. If the buffer is full, intrusion detection is not possible and it causes the decline of detection rate. It means that congested network cannot guarantee QoS of intrusion detection. Moreover, the IDS node which is suffered from congestion consumes more battery than other normal IDS nodes.

To prevent congested IDS node, we use queuing theory in set covering problem. Proposed algorithm does not select IDS node which is suffered from congestion of monitoring tasks. By experiment, in congested situation, our proposed scheme has superior performance than other existing schemes. Moreover, we confirm that proposed algorithm guarantee QoS of monitoring tasks and reliable sensor network.

In the future works, the verification of our algorithm is still required to use in the real situation. Moreover, to reduce calculation time, heuristic algorithms are required.

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