

Bounds on some van der Waerden numbers

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Abstract

For positive integers s and k_1, k_2, \dots, k_s , the van der Waerden number $w(k_1, k_2, \dots, k_s; s)$ is the minimum integer n such that for every s -coloring of set $\{1, 2, \dots, n\}$, with colors $1, 2, \dots, s$, there is a k_i -term arithmetic progression of color i for some i . We give an asymptotic lower bound for $w(k, m; 2)$ for fixed m . We include a table of values of $w(k, 3; 2)$ that are very close to this lower bound for $m = 3$. We also give a lower bound for $w(k, k, \dots, k; s)$ that slightly improves previously-known bounds. Upper bounds for $w(k, 4; 2)$ and $w(4, 4, \dots, 4; s)$ are also provided.

1 Introduction

Two fundamental theorems in combinatorics are van der Waerden's theorem [18] and Ramsey's theorem [16]. The theorem of van der Waerden says that for all positive integers s and k_1, k_2, \dots, k_s , there exists a least positive integer $n = w(k_1, k_2, \dots, k_s; s)$ such that whenever $[1, n] = \{1, 2, \dots, n\}$ is s -colored (i.e., partitioned into s sets), there is a k_i -term arithmetic progression with color i for some i , $1 \leq i \leq s$.

Similarly, Ramsey's theorem has an associated "threshold" function $R(k_1, k_2, \dots, k_s; s)$ (which we will not define here). The order of magnitude of $R(k, 3; 2)$ is known to be $\frac{k^2}{\log k}$ [11], while the best known upper bound on $R(k, m; 2)$ is fairly close to the best known lower bound. In contrast, the order of magnitude of $w(k, 3; 2)$ is not known, and the best known lower and upper bounds on $w(k, k; 2)$ are

$$(k-1)2^{(k-1)} \leq w(k, k; 2) < 2^{2^{2^{2^{(k+9)}}}},$$

the lower bound known only when $k-1$ is prime. The lower bound is due to Berlekamp [1] and the upper bound is a celebrated result of Gowers [6], which answered a long-standing question of Ron Graham. Graham currently offers 1000 USD for a proof or disproof of $w(k, k; 2) < 2^{k^2}$ [2]. Several other open problems are stated in [14].

Recently, there have been some further breakthroughs in the study of the van der Waerden function $w(k, m; 2)$. One was the amazing calculation that $w(6, 6; 2) = 1132$ by Kouril [12], extending the list

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of previously known values $w(3, 3; 2) = 9$, $w(4, 4; 2) = 35$, and $w(5, 5; 2) = 178$. A list of other known exact values of $w(k, m; 2)$ appears in [15]. Improved lower bounds on several specific values of $w(k, k; s)$ are given in [3] and [10].

In another direction, Graham [7] gives an elegant proof that if one defines $w_1(k, 3)$ to be the least n such that every 2-coloring of $[1, n]$ gives either k consecutive integers in the first color or a 3-term arithmetic progression in the second color, then

$$k^{c \log k} < w_1(k, 3) < k^{dk^2},$$

for suitable constants $c, d > 0$. This immediately gives $w(k, 3; 2) < k^{dk^2}$ since we trivially have $w(k, 3; 2) \leq w_1(k, 3)$. In view of Graham's bounds on $w_1(k, 3)$, it would be desirable to obtain improved bounds on $w(k, 3; 2)$. Of particular interest is the question of whether or not there is a non-polynomial lower bound for $w(k, 3; 2)$.

In this note we give a lower bound of $w(k, 3; 2) > k^{(2-o(1))}$. Although this may seem weak, we do know that $w(k, 3; 2) < k^2$ for $5 \leq k \leq 16$ (i.e., for all known values of $w(k, 3; 2)$ with $k \geq 5$; see Table 2). More generally, we give a lower bound on $w(k, m; 2)$ for arbitrary fixed m . We also present a lower bound for the classical van der Waerden numbers $w(k, k, \dots, k; s)$ that is a slight improvement over previously published bounds. In addition, we present an upper bound for $w(k, 4; 2)$ and an upper bound for $w(4, 4, \dots, 4; s)$.

2 Upper and lower bounds for certain van der Waerden functions

We shall need several definitions, which we collect here.

For positive integers k and n ,

$$r_k(n) = \max_{S \subseteq [1, n]} \{|S| : S \text{ contains no } k\text{-term arithmetic progression}\}.$$

For positive integers k and m , denote by $\chi_k(m)$ the minimum number of colors required to color $[1, m]$ so that there is no monochromatic k -term arithmetic progression.

The function $w_1(k, 3)$ has been defined in Section 1. Similarly, we define $w_1(k, 4)$ to be the least n such that every 2-coloring of $[1, n]$ has either k consecutive integers in the first color or a 4-term arithmetic progression in the second color.

We begin with an upper bound for $w_1(k, 4)$. The proof is essentially the same as the proof given by Graham [7] of an upper bound for $w_1(k, 3)$. For completeness, we include the proof here. We will make use of a recent result of Green and Tao [9], who showed that for some constant $c > 0$,

$$r_4(n) < ne^{-c\sqrt{\log \log n}} \tag{1}$$

for all $n \geq 3$.

Proposition 1. *There exists a constant $c > 0$ such that $w_1(k, 4) < e^{k^{c \log k}}$ for all $k \geq 2$.*

Proof. Suppose we have a 2-coloring of $[1, n]$ (assume $n \geq 4$) with no 4-term arithmetic progression of

the second color and no k consecutive integers of the first color. Let $t_1 < t_2 < \dots < t_m$ be the integers of the second color. Hence, $m < r_4(n)$. Let us define $t_0 = 0$ and $t_{m+1} = n$. Then there must be some i , $1 \leq i \leq m$, such that

$$t_{i+1} - t_i > \frac{n}{2r_4(n)}.$$

(Otherwise, using $r_4(n) \geq 3$, we would have $n = \sum_{i=0}^m (t_{i+1} - t_i) \leq \frac{n(m+1)}{2r_4(n)} \leq \frac{n(r_4(n)+1)}{2r_4(n)} \leq \frac{n}{2} + \frac{n}{6}$.)

Using (1), we now have an i with

$$t_{i+1} - t_i > \frac{n}{2r_4(n)} > \frac{1}{2}e^{c\sqrt{\log \log n}}.$$

If $n \geq e^{k^{d \log k}}$, $d = c^{-2}$, then $\frac{1}{2}e^{c\sqrt{\log \log n}} \geq k$ and we have k consecutive integers of the first color, a contradiction. Hence, $n < e^{k^{d \log k}}$ and we are done. \square

Clearly $w(k, 4; 2) \leq w_1(k, 4)$. Consequently, we have the following result.

Corollary 2. *There exists a constant $d > 0$ such that $w(k, 4; 2) < e^{k^{d \log k}}$ for all $k \geq 2$.*

Using Green and Tao's result, it is not difficult to obtain an upper bound for $w(4, 4, \dots, 4; s)$.

Proposition 3. *There exists a constant $d > 0$ such that $w(4, 4, \dots, 4; s) < e^{s^{d \log s}}$ for all $s \geq 2$.*

Proof. Consider a $\chi_4(m)$ -coloring of $[1, m]$ for which there is no monochromatic 4-term arithmetic progression. Some color must be used at least $\frac{m}{\chi_4(m)}$ times, and hence $\frac{m}{\chi_4(m)} \leq r_4(m)$ so that $\frac{m}{r_4(m)} \leq \chi_4(m)$. Let $c > 0$ be such that (1) holds for all $n \geq 3$, and let $m = e^{s^{d \log s}}$, where $d = c^{-2}$. Then $\chi_4(m) \geq \frac{m}{r_4(m)} > e^{c\sqrt{\log \log m}} = s$. This means that every s -coloring of $[1, m]$ has a monochromatic 4-term arithmetic progression. Since $m = e^{s^{d \log s}}$, the proof is complete. \square

It is interesting that the bounds in Corollary 2 and Proposition 3 have the same form.

The following theorem gives a lower bound on $w(k, k, \dots, k; s)$. It is deduced without too much difficulty from the Symmetric Hypergraph Theorem as it appears in [8], combined with an old result of Rankin [17]. To the best of our knowledge it has not appeared in print before, even though it is better, for large s , than the standard lower bound $\frac{cs^k}{k}(1 + o(1))$ (see [8]), as well as the lower bounds $s^{k+1} - \sqrt{c(k+1)\log(k+1)}$ and $\frac{ks^k}{e^{(k+1)^2}}$ due to Erdős and Rado [4], and Everts [5], respectively. We give the proof in some detail. The proof makes use of the following facts:

$$\chi_k(n) < \frac{2n \log n}{r_k(n)}(1 + o(1)), \quad (2)$$

which appears in [8] as a consequence of the Symmetric Hypergraph Theorem; and

$$r_k(n) > ne^{-c(\log n)^{\frac{1}{\lfloor \log_2 k \rfloor + 1}}}, \quad (3)$$

which, for some constant $c > 0$, holds for all $n \geq 3$ (this appears in [17]).

Theorem 4. *Let $k \geq 3$ be fixed, and let $z = \lfloor \log_2 k \rfloor$. There exists a constant $d > 0$ such that $w(k, k, \dots, k; s) > s^{d(\log s)^z}$ for all sufficiently large s .*

Proof. We make use of the observation that for positive integers s and m , if $s \geq \chi_k(m)$, then $w(k, k, \dots, k; s) > m$, which is clear from the definitions. For large enough m , (2) gives

$$\chi_k(m) < \frac{2m \log m}{r_k(m)} \left(1 + \frac{1}{2}\right) = \frac{3m \log m}{r_k(m)}. \quad (4)$$

Now let $d = \left(\frac{1}{2c}\right)^{z+1}$, where c is from (3), and let $m = s^{d(\log s)^z}$, where s is large enough so that (4) holds. By (3), noting that $\log m = d(\log s)^{z+1} = \left(\frac{\log s}{2c}\right)^{z+1}$, we have

$$\frac{m}{r_k(m)} < e^{c(\log m)^{\frac{1}{z+1}}} = e^{c \cdot \frac{\log s}{2c}} = \sqrt{s}.$$

Therefore,

$$\frac{3m \log m}{r_k(m)} < 3d\sqrt{s}(\log s)^{z+1} < s$$

for sufficiently large s . Thus, for sufficiently large s ,

$$\chi_k(m) < \frac{3m \log m}{r_k(m)} < s.$$

According to the observation at the beginning of the proof, this implies that $w(k, k, \dots, k; s) > m = s^{d(\log s)^z}$, as required. \square

We now give a lower bound on $w(k, m; 2)$. We make use of the Lovász Local Lemma (see [8] for a proof), which will be implicitly stated in the proof.

Theorem 5. *Let $m \geq 3$ be fixed. Then for all sufficiently large k ,*

$$w(k, m; 2) > k^{m-1 - \frac{1}{\log \log k}}.$$

Proof. Given m , choose $k > m$ large enough so that

$$k^{\frac{1}{2m \log \log k}} > \left(m - \frac{1}{2 \log \log k}\right) \log k \quad (5)$$

and

$$6 < \frac{\log k}{\log \log k}. \quad (6)$$

Next, let $n = \left\lfloor k^{m-1 - \frac{1}{\log \log k}} \right\rfloor$. To prove the theorem, we will show that there exists a (red, blue)-coloring of $[1, n]$ for which there is no red k -term arithmetic progression and no blue m -term arithmetic progression.

For the purpose of using the Lovász Local Lemma, randomly color $[1, n]$ in the following way. For each $i \in [1, n]$, color i red with probability $p = 1 - k^{\alpha-1}$ where

$$\alpha := \frac{1}{2m \log \log k},$$

and color it blue with probability $1 - p$.

Let \mathcal{P} be any k -term arithmetic progression. Then, since $1 + x \leq e^x$, the probability that \mathcal{P} is red is

$$p^k = (1 - k^{\alpha-1})^k \leq \left(e^{-k^{\alpha-1}}\right)^k = e^{-k^\alpha}.$$

Hence, applying (5), we have

$$p^k < \left(\frac{1}{e}\right)^{\left(m - \frac{1}{2\log\log k}\right)\log k} = \frac{1}{k^{m - \frac{1}{2\log\log k}}}.$$

Also, for any m -term arithmetic progression \mathcal{Q} , the probability that \mathcal{Q} is blue is

$$(1 - p)^m = (k^{\alpha-1})^m = \frac{1}{k^{m - \frac{1}{2\log\log k}}}.$$

Now let $\mathcal{P}_1, \mathcal{P}_2, \dots, \mathcal{P}_t$ be all of the arithmetic progressions in $[1, n]$ with length k or m . So that we may apply the Lovász Local Lemma, we form the “dependency graph” G by setting $V(G) = \{\mathcal{P}_1, \mathcal{P}_2, \dots, \mathcal{P}_t\}$ and $E(G) = \{\{\mathcal{P}_i, \mathcal{P}_j\} : i = j, \mathcal{P}_i \cap \mathcal{P}_j \neq \emptyset\}$. For each $\mathcal{P}_i \in V(G)$, let $d(\mathcal{P}_i)$ denote the degree of the vertex \mathcal{P}_i in G , i.e., $|\{e \in E(G) : \mathcal{P}_i \in e\}|$. We now estimate $d(\mathcal{P}_i)$ from above. Let $x \in [1, n]$. The number of k -term arithmetic progressions \mathcal{P} in $[1, n]$ that contain x is bounded above by $k \cdot \frac{n}{k-1}$, since there are k positions that x may occupy in \mathcal{P} and since the gap size of \mathcal{P} cannot exceed $\frac{n}{k-1}$. Similarly, the number of m -term arithmetic progressions \mathcal{Q} in $[1, n]$ that contain x is bounded above by $m \cdot \frac{n}{m-1}$.

Let \mathcal{P}_i be any k -term arithmetic progression contained in $[1, n]$. The total number of k -term arithmetic progressions \mathcal{P} and m -term arithmetic progressions \mathcal{Q} in $[1, n]$ that may have nonempty intersection with \mathcal{P}_i is bounded above by

$$k \left(k \cdot \frac{n}{k-1} + m \cdot \frac{n}{m-1} \right) < kn \left(2 + \frac{2}{m-1} \right), \quad (7)$$

since $k > m$. Thus, $d(\mathcal{P}_i) < kn(2 + \frac{2}{m-1})$ when $|\mathcal{P}_i| = k$. Likewise, $d(\mathcal{P}_i) < mn(2 + \frac{2}{m-1})$ when $|\mathcal{P}_i| = m$. Thus, for all vertices \mathcal{P}_i of G , we have $d(\mathcal{P}_i) < kn(2 + \frac{2}{m-1})$.

To finish setting up the hypotheses for the Lovász Local Lemma, we let X_i denote the event that the arithmetic progression \mathcal{P}_i is red if $|\mathcal{P}_i| = k$, or blue if $|\mathcal{P}_i| = m$, and we let $d = \max_{1 \leq i \leq t} d(\mathcal{P}_i)$. We have seen above that for all i , $1 \leq i \leq t$, the probability that X_i occurs is at most

$$q := \frac{1}{k^{m - \frac{1}{2\log\log k}}},$$

while from (7) we have $d < 2kn(1 + \frac{1}{m-1})$.

We are now ready to apply the Lovász Local Lemma, which says that in these circumstances, if the condition $eq(d+1) < 1$ is satisfied, then there is a (red, blue)-coloring of $[1, n]$ such that no event X_i occurs, i.e., such that there is no red k -term arithmetic progression and no blue m -term arithmetic progression. This will imply

$$w(k, m; 2) > n = k^{m-1 - \frac{1}{\log\log k}},$$

as desired. Thus, the proof will be complete when we verify that $eq(d+1) < 1$. Using $m \geq 3$, we have

Table 1: Small values of $w(k, 3)$ and $w_1(k, 3)$

k	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
$w(k, 3; 2)$	6	9	18	22	32	46	58	77	97	114	135	160	186	218	238
$w_1(k, 3)$	9	23	34	73	113	193	?	?	?	?	?	?	?	?	?

$d < 3kn$, so that $d + 1 < 3kn + 1 < e^2kn$. Hence, it is sufficient to verify that

$$e^3qkn < 1. \quad (8)$$

Since $q = \frac{1}{k^{m-2\log\log k}}$ and $n \leq k^{m-1-\frac{1}{\log\log k}}$, inequality (8) may be reduced to (6), and the proof is now complete. \square

Remark. As long as $k > e^{m^6}$, the inequality of Theorem 5 holds. To show this, we need only to show that conditions (5) and (6) hold if $k > e^{m^6}$. That (6) holds is obvious. For (5), it suffices to have $k^{\frac{1}{2m\log\log k}} > m\log k$; that is $\log k > 2m\log\log k(\log m + \log\log k)$. When $k \geq e^{m^6}$, we have $2m\log\log k(\log m + \log\log k) \leq 2(\log k)^{1/6}\log\log k(\frac{1}{6}\log\log k + \log\log k) = \frac{7}{3}(\log k)^{1/6}(\log\log k)^2$. Since $(\log\log k) < (\log k)^{7/20}$ for $k \geq e^{m^6}$ we have $2m\log\log k(\log m + \log\log k) \leq \frac{7}{3}(\log k)^{13/15}$. Finally, since $(\log k)^{2/15} \geq \frac{7}{3}$ for $k \geq e^{m^6}$, condition (5) is satisfied.

We end with a table of computed values. These were all computed with a standard backtrack algorithm except for $w(14, 3; 2)$, $w(15, 3; 2)$, and $w(16, 3; 2)$, which are due to Michal Kouril [13]. The values $w(k, 3; 2)$, $k \leq 12$, appeared previously in [15].

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