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THE FAST HARTLEY TRANSFORM ON A PARALLEL PROCESSOR

BY

LIM, BOON POCK

**A thesis submitted in partial fulfillment
of the requirements for the degree
Master of Science
Major in Engineering
South Dakota State University**

1988

THE FAST HARTLEY TRANSFORM ON A PARALLEL PROCESSOR

This thesis is approved as a creditable and independent investigation by a candidate for the degree, Master of Science, and is acceptable for meeting the thesis requirements for this degree. Acceptance of this thesis does not imply that the conclusions reached by the candidate are necessarily the conclusions of the major department.

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ACKNOWLEDGEMENTS

I would like to thank Dr. Douglas B. Miron, Associate Professor of Electrical Engineering, for his guidance and suggestions. Special thanks go to Cheng, Tone Vei for allowing me to use his PANASONIC KX-P10801 printer.

Many have assisted me in my study. I would like to express my deep appreciation to Dr. Ernest L. Buckley, Dean of Engineering, for his encouragement and assistance in obtaining a Lee Foundation scholarship. My sincere appreciation is also extended to the Naito family (Japan), Mr. Lim Chin Lin, Ronnie Chua, the Ang family, and the Koh family (Singapore), for their tangible and/or intangible support.

Finally, I am grateful to my family, especially my two sisters, for their assistance in many ways.

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CHAPTER ONE

INTRODUCTION

The transform kernel of the Fourier transform is $\exp(-j\omega t)$, i.e. $\cos(\omega t) - j\sin(\omega t)$; where $j = \sqrt{-1}$. Obviously the complex arithmetic is needed to obtain the required Fourier transform. In 1942, Hartley[1] proposed a new formulation of the Fourier integral identity by using $\text{cas}(\omega t)$ as the transform kernel, where $\text{cas}(\omega t) = \cos(\omega t) + \sin(\omega t)$, is an abbreviation for cosine and sine. It is called the Hartley transform and has many properties similar to those of the Fourier transform. The important distinctions are that the Hartley transform of a real-valued function is also real valued, and the inverse transformation is the same integral operation as the direct transformation. Furthermore, its evaluation does not involve complex functions. This is a potential advantage if the transform is to be explicitly computed.

The discrete Fourier transform (DFT) has the same transform kernel as the Fourier transform. In 1983, Bracewell[2] introduced the discrete Hartley transform (DHT) by using Hartley's transform kernel, i.e. $\text{cas}(\omega t)$. The DHT can apply to numerical spectral analysis and convolution. Unlike the DFT, no additional program is required for the inverse DHT as it is the same as the direct transformation. If the real and imaginary parts of the DFT are expressly

required then they are directly obtainable as the even and odd parts of the DHT. The power spectrum can also be obtained directly from the DHT without first calculating the real and imaginary parts of the DFT as in the usual way of calculating power spectra.

In 1984, Bracewell[3] worked out a fast algorithm for performing the DHT of a data sequence of N elements in a time proportional to $N \log_2 N$. He proved that the fast Hartley transform (FHT) is as fast as or faster than the fast Fourier transform (FFT) and serves for all uses such as spectral analysis, digital signal processing, and convolution to which the FFT is at present applied. Since then, many discussions on the FHT [5 - 10] and its applications [16], [18] were aroused.

The vector processor (VP) was initiated by Miron in 1985. It is an adjunct to an IBM personal computer in which a sequence of data is passed from the host's memory to a row of mathematical coprocessors, operated on simultaneously by each of a sequence of the coprocessor's instructions, and then the results are passed back to the main memory [21]. It is a single-instruction, multiple-data stream (SIMD) computer system [23]. The purpose of this device is to economically achieve maximum speed in the computer execution of a group of scientific calculations.

Since the FHT is applied on a data sequence of N elements, where N is usually very large, it is a good

candidate to be performed on the vector processor and theoretically, a great improvement on the computer execution time should be obtained.

The Hartley transform, the DHT, and their properties will be presented in the subsequent chapters. The derivation of the DFT from the DHT will also be illustrated. Finally, a software system for computing the DHT on a VP will be discussed and its performance will also be evaluated.

CHAPTER TWO
THE HARTLEY TRANSFORM

2-1 Definition

The old version of the Fourier transform and the inverse Fourier transform were written as

$$S(\omega) = (2\pi)^{-1/2} \int_{-\infty}^{\infty} V(t) \exp(-i\omega t) dt ,$$

$$V(t) = (2\pi)^{-1/2} \int_{-\infty}^{\infty} S(\omega) \exp(i\omega t) d\omega .$$

In order to achieve a symmetrical appearance, Hartley [1] introduced the different pair of formulas

$$\tilde{Q}(\omega) = (2\pi)^{-1/2} \int_{-\infty}^{\infty} V(t) \text{cas}(\omega t) dt ,$$

$$V(t) = (2\pi)^{-1/2} \int_{-\infty}^{\infty} \tilde{Q}(\omega) \text{cas}(\omega t) d\omega ,$$

where cas is the sum of the cosine and sine, i. e. $\text{cas}(x) \equiv \cos(x) + \sin(x)$. However, the new version of the Fourier transform uses $(2\pi)^{1/2}S(w)$ instead of $S(w)$. The result is that the factor $(2\pi)^{-1/2}$ disappears from the transform definition but a factor $(2\pi)^{-1}$ appears in the inversion formula. In terms of frequency, the Fourier transform can be written as

$$F(f) = \int_{-\infty}^{\infty} V(t) \exp(-i2\pi ft) dt, \quad (2-1)$$

$$V(t) = \int_{-\infty}^{\infty} F(f) \exp(i2\pi ft) df. \quad (2-2)$$

Adopting this practice leads to

$$H(f) = \int_{-\infty}^{\infty} V(t) \text{cas}(2\pi ft) dt, \quad (2-3)$$

$$V(t) = \int_{-\infty}^{\infty} H(f) \text{cas}(2\pi ft) df. \quad (2-4)$$

The proof of equation (2-4) can be found in [10]. Therefore, the inverse Hartley transform is indistinguishable from the direct transform.

2-2 Relationship Between $H(f)$ And $F(f)$

Rewriting equation (2-1) and (2-3), we have

$$\begin{aligned}
 F(f) &= \int_{-\infty}^{\infty} V(t) [\cos(2\pi ft) - i\sin(2\pi ft)] dt \\
 &= \int_{-\infty}^{\infty} V(t) \cos(2\pi ft) dt - \\
 &\quad i \int_{-\infty}^{\infty} V(t) \sin(2\pi ft) dt, \quad (2-5)
 \end{aligned}$$

and

$$\begin{aligned}
 H(f) &= \int_{-\infty}^{\infty} V(t) [\cos(2\pi ft) + \sin(2\pi ft)] dt \\
 &= \int_{-\infty}^{\infty} V(t) \cos(2\pi ft) dt + \\
 &\quad \int_{-\infty}^{\infty} V(t) \sin(2\pi ft) dt. \quad (2-6)
 \end{aligned}$$

Equation (2-5) and (2-6) imply that

$$H(f) = \text{Re } F(f) - \text{Im } F(f). \quad (2-7)$$

Let $H(f) = E(f) + O(f)$, where $E(f)$ and $O(f)$ are the even and odd parts of $H(f)$ respectively. Then

$$\text{Re } F(f) = E(f) = \int_{-\infty}^{\infty} V(t) \cos(2\pi ft) dt \quad (2-8)$$

$$\text{Im } F(f) = -O(f) = - \int_{-\infty}^{\infty} V(t) \sin(2\pi ft) dt. \quad (2-9)$$

By definition, $E(f) = E(-f)$

and $O(f) = -O(-f)$.

Therefore

$$\begin{aligned} E(f) &= \frac{H(f) + H(-f)}{2} \\ &= \int_{-\infty}^{\infty} V(t) \cos(2\pi ft) dt \end{aligned}$$

and

$$\begin{aligned} O(f) &= \frac{H(f) - H(-f)}{2} \\ &= \int_{-\infty}^{\infty} V(t) \sin(2\pi ft) dt. \end{aligned}$$

These two integrals are known as the Fourier cosine transform and the Fourier sine transform respectively and have been extensively tabulated[11]. It is true that $E(f)$, $O(f)$, $\text{Re } F(f)$, and $\text{Im } F(f)$ are real only when $V(t)$ is real.

The discussion above is summarized as follows: The Fourier transform is the even part of the Hartley transform minus i times the odd part; conversely, the Hartley transform is the real part of the Fourier transform minus the imaginary part.

Some good examples can be found in [4].

2-3 Power Spectrum

The power spectrum is well presented in [4].

By definition, the power spectrum $P(f)$ is

$$\begin{aligned} P(f) &= |F(f)|^2 \\ &= [\text{Re } F(f)]^2 + [\text{Im } F(f)]^2. \end{aligned}$$

It is also possible to obtain the power spectrum directly from the Hartley transform. Thus

$$\begin{aligned} P(f) &= [\text{Re } F(f)]^2 + [\text{Im } F(f)]^2 \\ &= [E(f)]^2 + [O(f)]^2 \\ &= \frac{[H(f) + H(-f)]^2 + [H(f) - H(-f)]^2}{4} \\ &= \frac{[H(f)]^2 + [H(-f)]^2}{2} \end{aligned}$$

Thus in lieu of squaring the real and imaginary parts and summing the two values at a given value f , we square and

sum the two values of the Hartley transform at $+f$ and $-f$, and then divide the result by a factor of 2.

2-4 Theorems

Two sorts of theorem are discussed in [4]. The first pertains to operations such as modulation, convolution, and other common operations that may be carried out on a function. This sort of theorem tells what corresponding operation goes on simultaneously in the transform domain. The second kind of theorem deals with relations between functions and their transforms that can typically be expressed in the form of an equation. They are reproduced from [4] and listed in Table 2.1 and Table 2.2 respectively. A worthy point to note is the convolution theorem. If one or both of the functions entering into the convolution are even, then the Hartley theorem is the same as the Fourier theorem, i. e. $H_1(f) H_2(f)$.

Table 2.1 Theorems for the Fourier and Hartley transforms

| Theorem | $V(t)$ | $F(f)$ | $H(f)$ |
|-----------------|-----------------------|------------------------------------|--|
| Similarity | $V(t/T)$ | $ T F(Tf)$ | $ T H(Tf)$ |
| Addition | $V_1(t) + V_2(t)$ | $F_1(f) + F_2(f)$ | $H_1(f) + H_2(f)$ |
| Reversal | $V(-t)$ | $F(-f)$ | $H(-f)$ |
| Shift | $V(t-T)$ | $e^{-j2\pi Tf}F(f)$ | $\cos(2\pi Tf)H(f) + \sin(2\pi Tf)H(-f)$ |
| Modulation | $V(t)\cos 2\pi f_0 t$ | $\frac{1}{2}[F(f-f_0) + F(f+f_0)]$ | $\frac{1}{2}[H(f-f_0) + H(f+f_0)]$ |
| Convolution | $V_1(t) * V_2(t)$ | $F_1(f)F_2(f)$ | $\frac{1}{2}[H_1(f)H_2(f) - H_1(-f)H_2(-f) + H_1(f)H_2(-f) + H_1(-f)H_2(f)]$ |
| Product | $V_1(t)V_2(t)$ | $F_1(f) * F_2(f)$ | $\frac{1}{2}[H_1(f) * H_2(f) - H_1(-f) * H_2(-f) + H_1(f) * H_2(-f) + H_1(-f) * H_2(f)]$ |
| Autocorrelation | $V(t) \otimes V(t)$ | $ F(f) ^2$ | $\frac{1}{2}([H(f)]^2 + [H(-f)]^2)$ |
| Derivative | $V'(t)$ | $j2\pi fF(f)$ | $-2\pi fH(-f)$ |
| 2nd derivative | $V''(t)$ | $-4\pi^2 f^2 F(f)$ | $-4\pi^2 f^2 H(f)$ |

Table 2.2 Theorems for relations between domains

| Theorem | Property | Fourier relation | Hartley relation |
|-------------------|--|------------------------|------------------------|
| Infinite integral | $\int_{-\infty}^{\infty} V(t) dt$ | $= F(0)$ | $= H(0)$ |
| First moment | $\int_{-\infty}^{\infty} tV(t) dt$ | $= F'(0)/(-12\pi)$ | $= -H'(0)/2\pi$ |
| Second moment | $\int_{-\infty}^{\infty} t^2V(t) dt$ | $= -F''(0)/4\pi^2$ | $= -H''(0)/4\pi^2$ |
| Centroid | $\frac{\int_{-\infty}^{\infty} tV(t) dt}{\int_{-\infty}^{\infty} V(t) dt}$ | $= 1F'(0)/(2\pi F(0))$ | $= -H'(0)/(2\pi H(0))$ |

CHAPTER THREE

THE DISCRETE HARTLEY TRANSFORM

3-1 Definition

The discrete Fourier transform (DFT) and its inverse have the standard form

$$F(k) = N^{-1} \sum_{n=0}^{N-1} f(n) \exp(-j2\pi nk/N) \quad (3-1)$$

$$f(n) = \sum_{k=0}^{N-1} F(k) \exp(j2\pi nk/N) \quad (3-2)$$

The function $f(n)$ may be the discrete representation of an underlying continuous waveform or may be a function of a variable that is basically discrete.

Bracewell [2] defined the discrete Hartley transform (DHT) of a finite length sequence and its inverse as

$$H(k) = N^{-1} \sum_{n=0}^{N-1} f(n) \text{cas}(2\pi nk/N) \quad , \quad (3-3)$$

$$f(n) = \sum_{k=0}^{N-1} H(k) \text{cas}(2\pi nk/N) \quad , \quad (3-4)$$

$$0 \leq n \leq (N-1) \quad ,$$

where $\text{cas}(x) \equiv \cos(x) + \sin(x)$. The proof of equation (3-4) can be found in [4], pp. 28. The DHT is real if $f(n)$ is real. Comparing these equations, one can see that the DHT and the DFT are closely related.

3-2 Relationship Between The DHT And The DFT

Rewriting equation (3-1) and (3-3), we have

$$\begin{aligned}
 F(k) &= N^{-1} \sum_{n=0}^{N-1} f(n) [\cos(2\pi nk/N) - i \sin(2\pi nk/N)] \\
 &= N^{-1} \sum_{n=0}^{N-1} f(n) \cos(2\pi nk/N) - \\
 &\quad i N^{-1} \sum_{n=0}^{N-1} f(n) \sin(2\pi nk/N)
 \end{aligned}$$

and

$$\begin{aligned}
 H(k) &= N^{-1} \sum_{n=0}^{N-1} f(n) [\cos(2\pi nk/N) + i \sin(2\pi nk/N)] \\
 &= N^{-1} \sum_{n=0}^{N-1} f(n) \cos(2\pi nk/N) + \\
 &\quad i N^{-1} \sum_{n=0}^{N-1} f(n) \sin(2\pi nk/N).
 \end{aligned}$$

As in the continuous case, the DHT possesses even and odd parts

$$H(k) = E(k) + O(k)$$

$$\text{where } E(k) = N^{-1} \sum_{n=0}^{N-1} f(n) \cos(2\pi nk/N) ,$$

$$O(k) = N^{-1} \sum_{n=0}^{N-1} f(n) \sin(2\pi nk/N) .$$

If k is replaced by $(N - k)$,

$$E(N-k) = N^{-1} \sum_{n=0}^{N-1} f(n) \cos(2\pi n - 2\pi nk/N)$$

$$= N^{-1} \sum_{n=0}^{N-1} f(n) \cos(2\pi nk/N)$$

$$= E(k)$$

and

$$O(N-k) = N^{-1} \sum_{n=0}^{N-1} f(n) \sin(2\pi n - 2\pi nk/N)$$

$$= -N^{-1} \sum_{n=0}^{N-1} f(n) \sin(2\pi nk/N)$$

$$= -O(k) .$$

Thus

$$E(k) = \frac{[H(k) + H(N-k)]}{2}$$

and

$$O(k) = \frac{[H(k) - H(N-k)]}{2}$$

In the case of $k = 0$ and $k = N/2$, one can prove that

$$E(0) = H(0) \quad , \quad O(0) = 0$$

$$E(N/2) = H(N/2) \quad , \quad O(N/2) = 0.$$

From the definition of the DFT, it is apparent that $F(k)$ can be formed from the DHT's even and odd parts by

$$\begin{aligned} F(k) &= E(k) - iO(k) \\ &= \frac{[H(k) + H(N-k)]}{2} - i \frac{[H(k) - H(N-k)]}{2} \\ &= \frac{[H(N-k) + H(k)]}{2} + i \frac{[H(N-k) - H(k)]}{2} . \end{aligned}$$

Conversely, to form $H(k)$ when $F(k)$ is available

$$H(k) = \text{Re } F(k) - \text{Im } F(k) .$$

These relations are strictly analogous to those obtained previously for the continuous variable.

A number of examples that illustrates the characteristics of the DHT can also be found in [4].

3-3 Theorems

The theorems of the DHT discussed here are based upon [4]. Just as there is a Hartley transform theorem for every theorem that applies to the Fourier transform, there are also corresponding theorems for the discrete transforms. The theorems such as convolution, autocorrelation, and first value, etc. are listed in Table 3.1 and Table 3.2.

The proofs and discussions on some theorems such as reversal, addition, shift, product, and convolution are well presented in [4].

One may note that the mean value of the sequence $f(n)$ is given by $H(0)$ and that the mean square value of $f(n)$ is given by ΣH^2 .

Some of the theorems for the two different transforms correspond exactly, as is the case with $\Sigma H(k) = f(0)$ and $\Sigma f(n) = N \times H(0)$, but some exhibit differences.

The only theorem not listed in the table is the

stretch or similarity theorem. It states that if a sequence $f(n)$ is stretched to double its length by inserting a zero element after each given element, then the elements of the original DHT are repeated. For example,

{1 2 3 4} has the DHT {2.5 -1 -0.5 0} ,

then {1 0 2 0 3 0 4 0} has the DHT

{2.5 -1 -0.5 0 2.5 -1 -0.5 0}.

The properties of the DHT commend themselves for application to numerical analysis. The fact that the transform values are real is a convenience in managing calculations. In addition, the reversibility of the transform is helpful as one does not need to keep track of which domain one is in. Furthermore, several of the theorems for the Fourier transform have different forms according to the domain, a concern that is obviated with the DHT. The factor N is domain-dependent and could be a normalization factor or calibration factor to be applied at the end of a numerical calculation. Experience shows that the last step is the place to consolidate proportionality factors and so the departure from strict reversibility represented by the factor N is not computationally important [4].

Table 3.1 Theorems for operations on discrete transforms

| Theorem | Function $f(n)$ | DFT $F(k)$ | DHT $H(k)$ |
|-----------------|---------------------|------------------------|---|
| Reversal | $f(-n)$ | $F(-k)$ | $H(-k)$ |
| Addition | $f_1(n) + f_2(n)$ | $F_1(k) + F_2(k)$ | $H_1(k) + H_2(k)$ |
| Shift | $f(n-T)$ | $e^{-j2\pi Tk/N} F(k)$ | $\cos(2\pi Tk/N) H(k) + \sin(2\pi Tk/N) H(N-k)$ |
| Convolution | $f_1(n) * f_2(n)$ | $N F_1(k) F_2(k)$ | $\frac{1}{N} [H_1(k) H_2(k) - H_1(-k) H_2(-k) + H_1(k) H_2(-k) + H_1(-k) H_2(k)]$ |
| Product | $f_1(n) f_2(n)$ | $F_1(k) * F_2(k)$ | $\frac{1}{N} [H_1(k) * H_2(k) - H_1(-k) * H_2(-k) + H_1(k) * H_2(-k) + H_1(-k) * H_2(k)]$ |
| Autocorrelation | $f(n) \otimes f(n)$ | $\frac{1}{N} F(k) ^2$ | $N \{ [H(k)]^2 + [H(-k)]^2 \}$ |
| Derivative | $f'(n)$ | $j2\pi k F(k)$ | $2\pi k H(-k)$ |
| 2nd derivative | $f''(n)$ | $-4\pi^2 k^2 F(k)$ | $-4\pi^2 k^2 H(k)$ |

Table 3.2 Theorems on relations for discrete transforms

Sum of sequence $\sum_{n=0}^{N-1} f(n) = NF(0) = NH(0)$

First value $f(0) = \sum_{k=0}^{N-1} F(k) = \sum_{k=0}^{N-1} H(k)$

Quadratic content $\sum_{n=0}^{N-1} [f(n)]^2 = N \sum_{k=0}^{N-1} |F(k)|^2 = N \sum_{k=0}^{N-1} |H(k)|^2$

CHAPTER FOUR

THE FAST HARTLEY TRANSFORM

Since the discrete Hartley transform (DHT) is closely related to the discrete Fourier transform (DFT), one can predict that the fast Fourier transform (FFT) algorithm can be converted to a fast Hartley transform (FHT) with some modifications. The FHT can be expected to transform one real array of length N in half the time that it takes the FFT to process a complex array of length N . Buneman [10] presented the similarities between DHT and DFT and showed that the role played by the imaginary part in the complex Fourier transform is taken on by the real Hartley transform recorded backwards so that one can convert an FFT program into an FHT program with only a few indexing changes.

About one year after his proposal of the DHT, Bracewell [3] published the first fast algorithm for performing the DHT which was basically the same as Cooley and Tukey's method [12] for the FFT, or so-called decimation-in-time (DIT) algorithm (It is also called a decimation-in-time radix-2 algorithm or just a radix-2 algorithm.) Since then, the various methods for computation of the FFT were examined and attempted for the FHT. The methods such as decimation-in-frequency (DIF) FHT [5], split-radix (SR) FHT [9], in-place FHT [6], and radix-4 FHT [8] had been worked out for computing the DHT. The total numbers of the operations of

the first two methods and the DIT are summarized as follows [9]:

| | Additions | Multiplications |
|-----|---------------------------|-----------------------|
| DIT | $2N(\log_2 N - 1) + 2$ | $N(\log_2 N - 2) + 2$ |
| DIF | $3/2 N(\log_2 N - 1) + 2$ | $N(\log_2 N - 3) + 4$ |
| SR | $(2N - 4)(\log_4 N)$ | $(N - 4)(\log_4 N)$ |

For ease of comparison, various values of N are calculated and shown in Table 4.1 and Table 4.2.

Table 4.1 Number of real multiplications to compute an N -point DHT

| N | DIT | DIF | SR |
|------|------|------|------|
| 16 | 34 | 20 | 24 |
| 32 | 98 | 68 | 84 |
| 64 | 258 | 196 | 180 |
| 128 | 642 | 516 | 496 |
| 256 | 1538 | 1284 | 1008 |
| 512 | 3586 | 3076 | 2540 |
| 1024 | 8194 | 7172 | 5100 |

Table 4.2 Number of real additions to compute an N -point DHT

| N | DIT | DIF | SR |
|------|-------|-------|-------|
| 16 | 98 | 74 | 56 |
| 32 | 258 | 194 | 180 |
| 64 | 642 | 482 | 372 |
| 128 | 1538 | 1152 | 1008 |
| 256 | 3586 | 2690 | 2032 |
| 512 | 8194 | 6146 | 5100 |
| 1024 | 18434 | 13826 | 10220 |

A comparison of operation counts on various methods of the FHT can be found in [7], [8], and [9]. The application of the same method but with different approach will cause a minor difference in operation counts. Although [7] and [8] concluded that the FHT algorithms do not give any increase in performance over the existing fastest real-valued FFT algorithms such as Winograd's FFT algorithm [14] and Bergland and Dolan's algorithm [15], one may predict that the FHT has a great potential in some situations as its real-valued function nature and the equivalence of the forward and inverse DHT may justify the cost. Hou[13] presented his algorithm and showed that in lower order transforms the number of nontrivial, real arithmetic operations in his FHT is about the same as in the Winograd's FFT.

One may note that N^{-1} in equation (3-3) usually is regarded as the normalization factor in the FHT algorithm and is neglected until the very end.

Equation (3-3) can be written in matrix form as

$$[H] = N^{-1} [CAS] [X]. \quad (4-1)$$

Where $[H]$ and $[X]$ are $N \times 1$ matrices, N^{-1} is a scalar, and

$$\begin{array}{c}
 \text{CAS} \\
 \text{[CAS]} = \begin{array}{c}
 \parallel \\
 0 \\
 1 \\
 2 \\
 \vdots \\
 \vdots \\
 N-1
 \end{array} \left[\begin{array}{c}
 \cos(2\pi nk/N) + \sin(2\pi nk/N) \\
 \vdots \\
 \vdots \\
 \cos(2\pi nk/N) + \sin(2\pi nk/N)
 \end{array} \right]
 \end{array}$$

k n = 0, 1, 2, , N-1

For a computer language which has powerful matrix manipulations, e.g. APL, equation (4-1) is very useful and should be considered.

If the cas function is carefully analyzed, one may notice that

when n is replaced by $(N - n)$

$$\cos [2\pi (N-n)k/N] = \cos [2\pi nk/N] ;$$

$$\sin [2\pi (N-n)k/N] = -\sin [2\pi nk/N] ,$$

and when k is replaced $(N - k)$

$$\cos [2\pi n(N-k)/N] = \cos [2\pi nk/N] ;$$

$$\sin [2\pi n(N-k)/N] = -\sin [2\pi nk/N] .$$

Also, when $n = 0$ or $k = 0$, the values of the cosine terms are all '1's and the sine terms are all '0's.

When $n = N/2$ or $k = N/2$, the values of the sine terms are also all '0's but the values of the cosine terms are varied, that is

$$\cos(\pi k) = (-1)^k \quad \text{and}$$

$$\cos(\pi n) = (-1)^n.$$

These imply that

$$\begin{matrix}
 & 0 & 1 & \dots & N/2 & \dots & N-1 \\
 \begin{matrix} 0 \\ \vdots \\ N/2 \\ \vdots \\ N-1 \end{matrix} & \left[\begin{array}{cc}
 1 & 1 & 1 & \dots & 1 & \dots & 1 & 1 \\
 1 & \boxed{\begin{matrix} Z_1 & Z_2 & \dots \\ Z_1 & Z_2 & \dots \end{matrix}} & -1 & \dots & \dots & \dots & \boxed{\begin{matrix} Z_2 & Z_1 \\ Z_2 & Z_1 \end{matrix}} \\
 2 & 1 & \boxed{\begin{matrix} Z_3 & Z_4 & \dots \\ :3 & :4 & \dots \end{matrix}} & 1 & \dots & \dots & \boxed{\begin{matrix} Z_4 & Z_3 \\ :4 & :3 \end{matrix}} \\
 \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
 N/2 & 1 & -1 & 1 & -1 & \dots & 1 & \dots & 1 & -1 \\
 \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
 \vdots & 1 & \boxed{\begin{matrix} : & : & \dots \\ Z_3 & Z_4 & \dots \end{matrix}} & -1 & \dots & \dots & \boxed{\begin{matrix} \dots & : & : \\ \dots & Z_4 & Z_3 \end{matrix}} \\
 \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
 N-1 & 1 & \boxed{\begin{matrix} Z_1 & Z_2 & \dots \\ Z_1 & Z_2 & \dots \end{matrix}} & 1 & \dots & \dots & \boxed{\begin{matrix} \dots & Z_2 & Z_1 \\ \dots & Z_2 & Z_1 \end{matrix}}
 \end{array} \right]
 \end{matrix}$$

and

$$\begin{matrix}
 & 0 & \dots & N/2 & \dots & N-1 \\
 \begin{matrix} 0 \\ \vdots \\ N/2 \\ \vdots \\ N-1 \end{matrix} & \left[\begin{array}{cc}
 0 & 0 & 0 & \dots & 0 & \dots & 0 & 0 \\
 0 & \boxed{\begin{matrix} Y_1 & Y_2 & \dots \\ Y_1 & Y_2 & \dots \end{matrix}} & 0 & \dots & \dots & \dots & \boxed{\begin{matrix} \dots & -Y_2 & -Y_1 \\ \dots & -Y_2 & -Y_1 \end{matrix}} \\
 2 & 0 & \boxed{\begin{matrix} Y_3 & Y_4 & \dots \\ :3 & :4 & \dots \end{matrix}} & 0 & \dots & \dots & \boxed{\begin{matrix} \dots & -Y_4 & -Y_3 \\ \dots & -Y_4 & -Y_3 \end{matrix}} \\
 \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
 N/2 & 0 & 0 & 0 & 0 & \dots & 0 & \dots & 0 & 0 \\
 \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
 \vdots & 0 & \boxed{\begin{matrix} : & : & \dots \\ -Y_3 & -Y_4 & \dots \end{matrix}} & 0 & \dots & \dots & \boxed{\begin{matrix} \dots & : & : \\ \dots & Y_4 & Y_3 \end{matrix}} \\
 \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
 N-1 & 0 & \boxed{\begin{matrix} -Y_1 & -Y_2 & \dots \\ -Y_1 & -Y_2 & \dots \end{matrix}} & 0 & \dots & \dots & \boxed{\begin{matrix} \dots & Y_2 & Y_1 \\ \dots & Y_2 & Y_1 \end{matrix}}
 \end{array} \right]
 \end{matrix}$$

Recall that

$$[CAS] = [COS] + [SIN],$$

which means a quarter of each sine and cosine matrix shall be calculated so that the complete matrix can be formed via move and negate instructions. Since the time

taken for calculating the trigonometric function is much longer than moving the data to the various locations, a time reduction should be achieved by using this method although it increases the complexity of the program.

Since $\text{cas}(x) = \cos(x) + \sin(x)$,

$$\cos(\pi/4) = \sin(\pi/4) = 2^{-1/2},$$

it may be faster to replace $\text{cas}(x)$ by $2^{1/2}\sin[x+(\pi/4)]$ or $2^{1/2}\cos[x-(\pi/4)]$ on some computers.

CHAPTER FIVE

THE VECTOR PROCESSOR

The vector processor (VP) was initiated by Miron in 1985. It is an adjunct to a personal computer in which a sequence of data is passed from the host's memory to a row of mathematical coprocessors (MCs), operated on simultaneously by each of a sequence of coprocessor instructions, and then the results are passed back to main memory. It is a single-instruction, multiple-data (SIMD) auxiliary computer. Its purpose is to speed up all those operations which can be expressed in vector-matrix forms by parallel execution of the data operations [21].

The general arrangement of the major elements of the VP is shown in figure 5.1. The vector control unit (VCU) in the vector processor is as important as the general purpose processor (GPP) in the personal computer. It consists of a vector instructions decoder, a sequential load/store control unit, and a parallel execution control unit. Its function is to co-ordinate the operation between the GPP and the row of MCs. If the vector operation (parallel mode) is not required, the VCU will connect one of the MCs to the GPP as in the conventional architecture of figure 5.2 so that faster speed in executing ordinary arithmetic operations can be achieved. When a vector instruction appears, it is decoded by the VCU and appropriate action taken. The effect of these

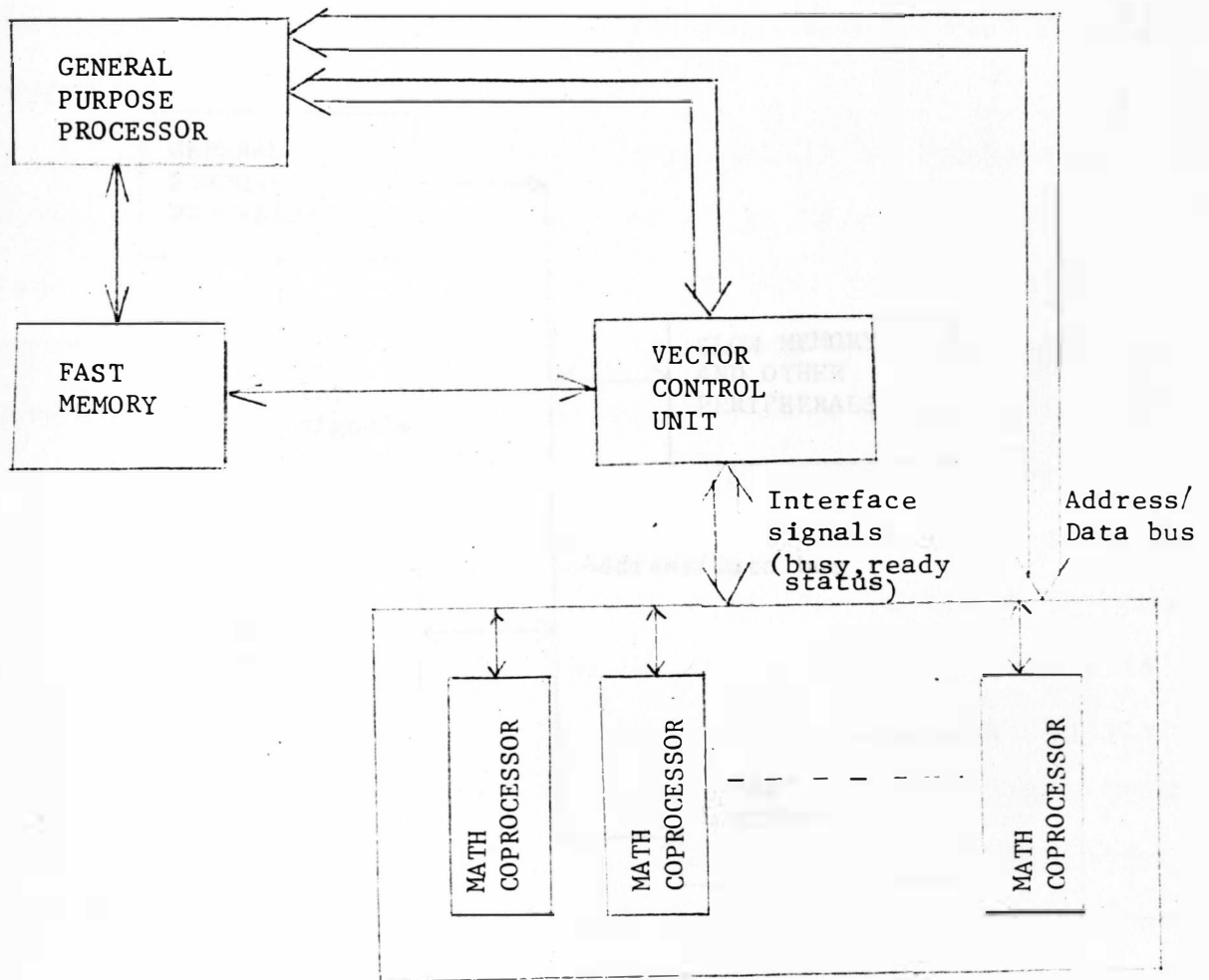


Figure 5.1 Vector Processor Architecture

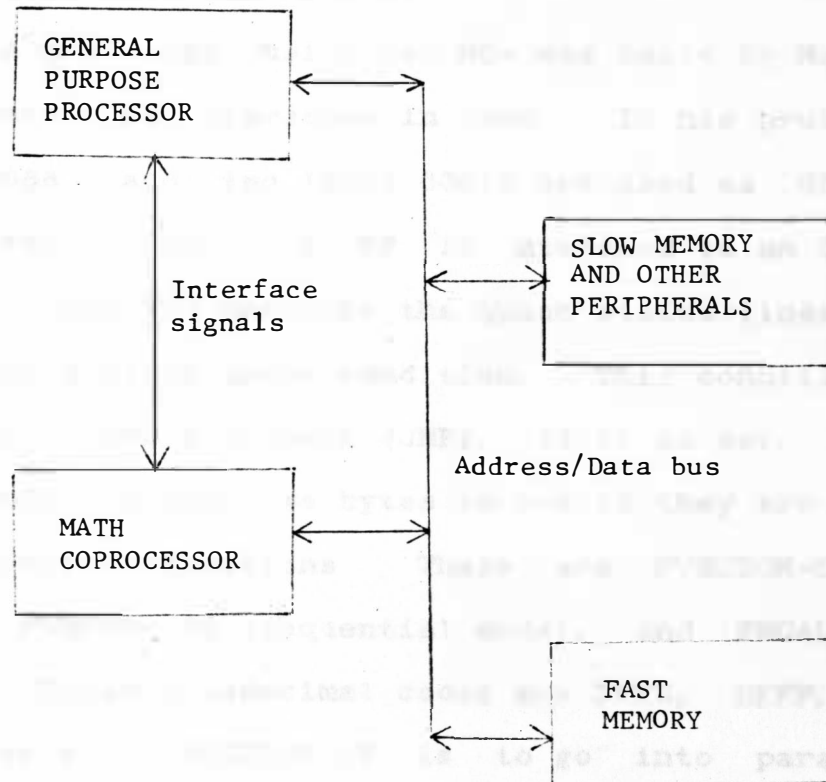


Figure 5.2 A Conventional Microcomputer Architecture

actions is to load the elements of the required vectors sequentially from fast memory into the MCs, do the required operations simultaneously, and then store the results sequentially into fast memory.

A prototype using two MCs was built by Manja [22] to demonstrate the operation in 1986. In his prototype, an Intel 8088 and two Intel 8087s are used as GPP and MCs respectively, and the VP is attached to an IBM personal computer. The VCU monitors the queue status lines of the GPP to detect a clear-queue condition. This condition will be set after a jump statement (JMP). If it is set, the decoder will decode the next two bytes to see if they are one of the three vector instructions. These are FVECTOR-OP (parallel mode), FVECTOR-SQ (sequential mode), and FSCALAR (scalar mode). Their hexadecimal codes are DFFD, DFFF, and DFFE respectively. FVECTOR-OP is to go into parallel mode. FVECTOR-SQ is used for loading/storing the data to/from the MCs, whereas FSCALAR is for returning to the conventional mode and also acting as a reset between the other two modes. In each mode, the VCU selects the MC by its READY and queue status lines. The functional blocks and signals managed in the VCU are shown in figure 5.3.

As mentioned earlier the purpose of the VP is to speed up the processing time of the data. A measure of the improvement to be expected is defined as the ratio of the total time to process the data with a single MC to the time

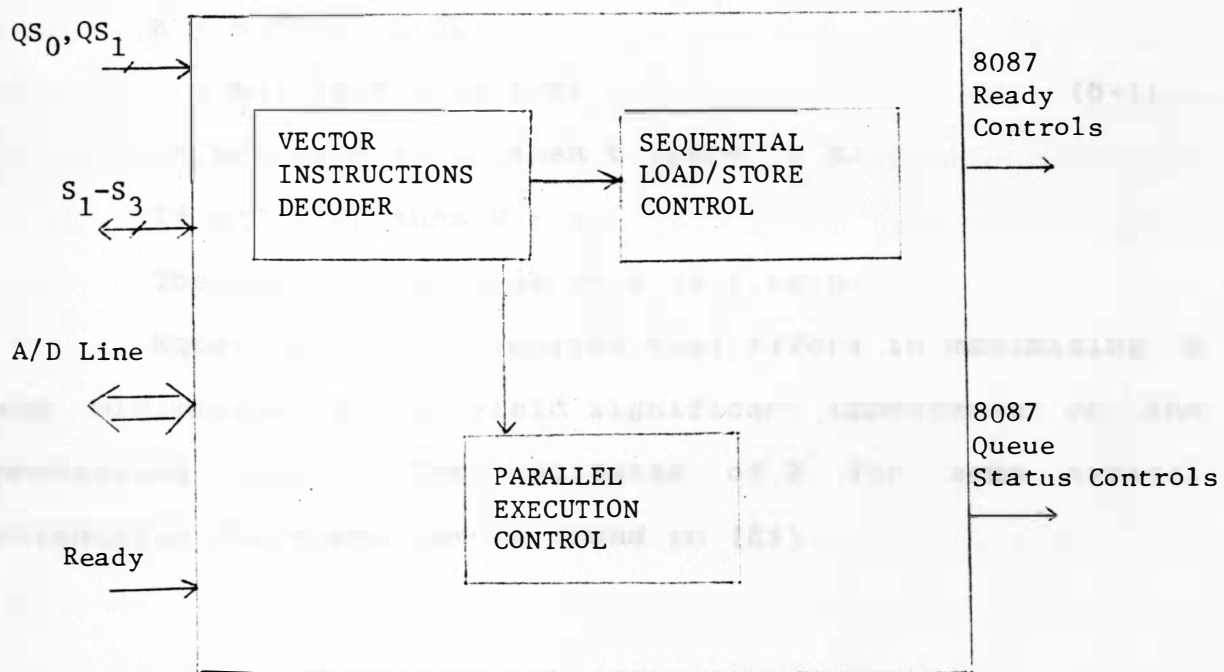


Figure 5.3 Vector Controller

with M MCs in a VP [21]. For a given function, let

E = execution time for one result element,

T = total data transfer time for that result element,

M = the number of MCs,

and R = improvement ratio.

By definition,

$$\begin{aligned} R &= M(E+T)/(E+MT) \\ &= M(1+(E/T))/(M+E/T) \end{aligned} \quad (5-1)$$

If $E/T \gg M \gg 1$, then R approx. = M .

If $E/T = 0$, then $R = 1$.

Therefore, the range of R is 1 to M .

Equation (5-1) implies that effort in maximizing E and minimizing T will yield significant improvement on the processing time. The estimates of R for some typical scientific functions can be found in [21].

CHAPTER SIX

THE FHT ON A VECTOR PROCESSOR

6-1 Selection of an Algorithm

A program is written for computing the discrete Hartley transform (DHT) on a vector processor (VP). Some fast Hartley transform (FHT) algorithms that have been published are for a single processor [3 - 10]. These algorithms can be implemented on the VP with some modifications. An algorithm with less processing time for a single processor may not necessarily be the best algorithm for the VP. This is due to the VP's configuration. An Intel 8088 and some Intel 8087s are assumed to be the general purpose processor (GPP) and the mathematical coprocessors (MCs) respectively for a VP. As proved in chapter five, the greatest improvement on the processing time will be achieved if the execution time for one result element is much greater than the total data transfer time for that result element. This means the data set involved in parallel processing should be arranged in such a way that a result can be derived from the original data set and not the results previously computed from the data set. This is called data independence. If the elements of the data set are not independent of each other, parallel processing cannot be applied. Furthermore, comparison of the intermediate result must be avoided. Comparison can only be done in scalar or sequential

mode. It will defeat the purpose of parallel processing. The Intel 8087 has eight internal registers [24], [25]. Fully using these registers to keep the intermediate results will definitely reduce the data transfer time. Although sometimes it is done at the expense of the memory storage, data independence, avoiding comparison, and reducing the data transfer time are the key considerations for programming the parallel processing routine on a VP.

The result from the above discussion is the use of the algorithm of Bold [7] for computing the DHT on a VP. A detailed discussion on the algorithm will be presented in the next section.

6-2 The FHT Algorithm

Consider a data sequence

$$f(n) = \{a_1 \ a_2 \ b_1 \ b_2 \ c_1 \ c_2 \ d_1 \ d_2\},$$

we can rewrite it as

$$f(n) = g_1(n) + g_2(n-1),$$

$$\text{where } g_1(n) = \{a_1 \ 0 \ b_1 \ 0 \ c_1 \ 0 \ d_1 \ 0\},$$

$$g_2(n) = \{a_2 \ 0 \ b_2 \ 0 \ c_2 \ 0 \ d_2 \ 0\},$$

$$\text{and } g_2(n-1) = \{0 \ a_2 \ 0 \ b_2 \ 0 \ c_2 \ 0 \ d_2\}.$$

Suppose $\{a_1 \ b_1 \ c_1 \ d_1\}$ has the DHT

$$\{\alpha_1 \ \beta_1 \ \tau_1 \ \theta_1\}$$

and $\{a_2 \ b_2 \ c_2 \ d_2\}$ has the DHT

$$\{\alpha_2 \ \beta_2 \ \tau_2 \ \theta_2\}.$$

By applying the stretch theorem and the shift theorem for the DHT, we can show that

$$G_1(k) = \{\alpha_1 \beta_1 \tau_1 \delta_1 \alpha_1 \beta_1 \tau_1 \delta_1\},$$

$$G_2(k) = \{\alpha_2 \beta_2 \tau_2 \delta_2 \alpha_2 \beta_2 \tau_2 \delta_2\},$$

$$\text{and } G_2(k-1) = \cos(2\pi k/N)G_2(k) + \sin(2\pi k/N)G_2(N-k).$$

Therefore, the DHT of a sequence $f(n)$ can be written as

$$G(k) = G_1(k) + \cos(2\pi k/N)G_2(k) + \sin(2\pi k/N)G_2(N-k).$$

Adopting this approach, Bracewell [4] shows that the general decomposition formula that produces the discrete Hartley transform (DHT) for a sequence $f(n)$ of N elements, where $N = 2^P$, is

$$H(k) = H_{\text{odd}}(k) + H_{\text{even}}(k)\cos(2\pi k/N) + H_{\text{even}}(N-k)\sin(2\pi k/N), \quad (6-1)$$

where $k = 0, 1, 2, 3, \dots, (N-1)$, $H_{\text{odd}}(k)$ and $H_{\text{even}}(k)$ are the DHT of the odd and the even numbered terms in the sequence $f(n)$ respectively.

$$\text{Since } \cos(2\pi((N/2)+k)/N) = -\cos(2\pi k/N),$$

$$\sin(2\pi((N/2)+k)/N) = -\sin(2\pi k/N),$$

and the terms in H_{odd} and H_{even} repeat modulo $N/2$ [4], equation (6-1) can be rewritten as

$$H(k) = H_{\text{odd}}(k) + H_{\text{even}}(k)\cos(2\pi k/N) + H_{\text{even}}(N-k)\sin(2\pi k/N) \quad (6-2)$$

$$\text{and } H((N/2)+k) = H_{\text{odd}}(k) - H_{\text{even}}(k)\cos(2\pi k/N) -$$

$$H_{\text{even}}(N-k)\sin(2\pi k/N), \quad (6-3)$$

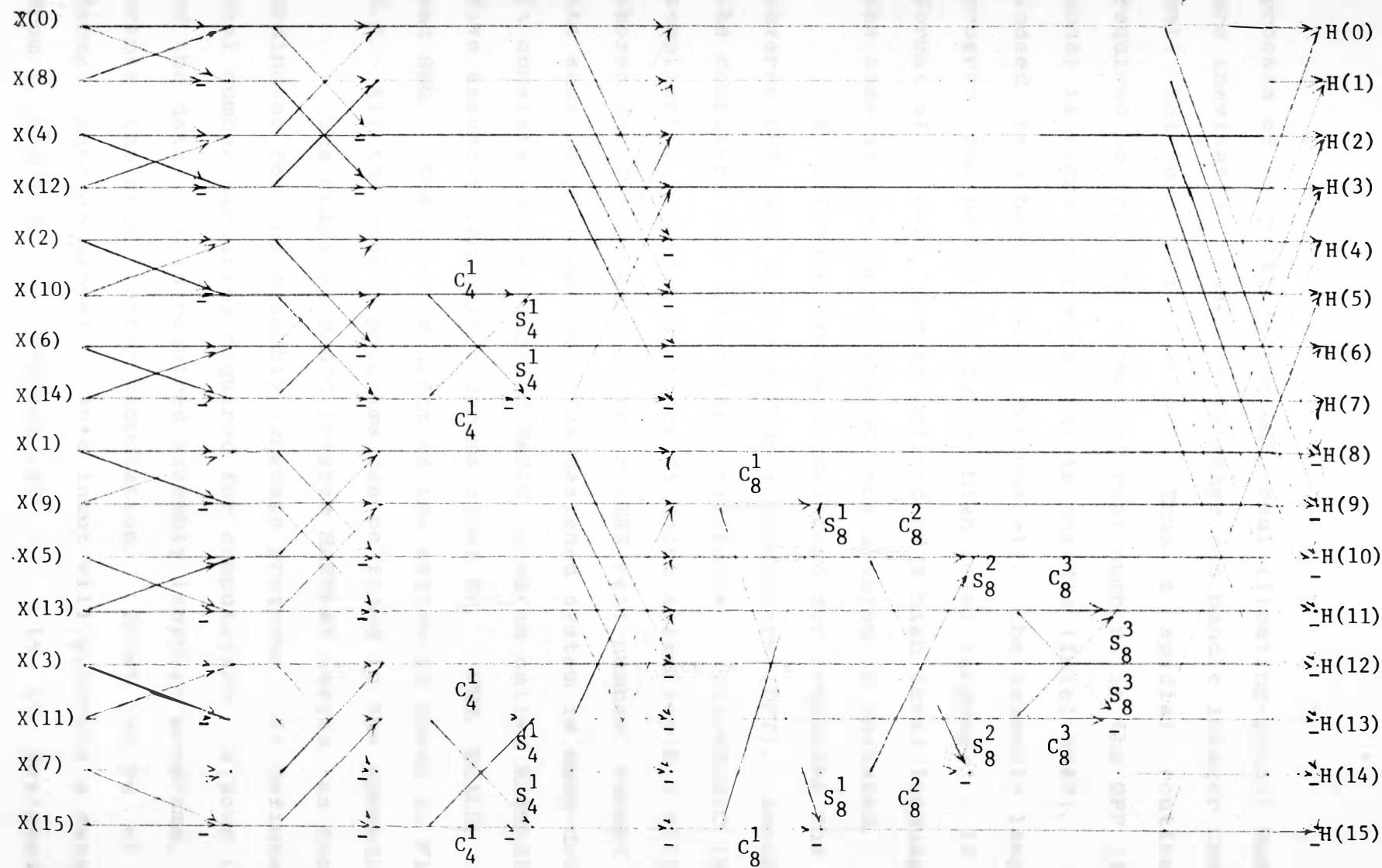
where $k = 0, 1, 2, 3, \dots, ((N/2)-1)$.

Since $H_{\text{odd}}(k)$ and $H_{\text{even}}(k)$ can be derived by continued decomposition until two-element sequences are reached, the complete breakdown is expressible in terms of the original data. This is similar to the form used by Cooley and Tukey [12] for computing the DFT. Consequently, a bit-reversing operation is required before the computation begins. The basic algorithm for a bit-reversing operation can be found in [26]. The signal flow graph as developed from equation (6-2) and equation (6-3) for a 16-element sequence is shown in figure 6.1.

6-3 System Design

The algorithm used for computing the FHT on a VP is based on the work of Bold [7]. However, some modifications have been made so that optimization of processing time can be achieved. The algorithm for the bit-reversing operation in Bold's work involves comparison of the data. Hence it is not suitable for programming on the VP and is replaced by the work done by Brigham [26]. Because the VP has its own three unique instructions, i.e. parallel mode, sequential mode, and scalar mode, which cannot be recognized by the high level computer languages (such as FORTRAN, BASIC, PASCAL, and C), assembly language must be used. By using assembly language, the VP's three instructions can be coded as data constants, using their corresponding hexadecimal numbers. In the

Figure 6.1 FHT signal flow graph for N = 16



$$S_N^n = \sin((2\pi/N)n) , C_N^n = \cos((2\pi/N)n)$$

process of calculating the DHT, real (floating-point) numbers are inevitable. Assembly language can handle integer numbers well but not real numbers. Thus a special routine is required so that the format of real numbers in the GPP (Intel 8088) is compatible with that in the MCs (Intel 8087). This indeed is a hard task! Fortunately, the assembly language program can be called by a high level language. If the format of a real number defined by a high level language is the same as the one in the MC, the problem is obviated.

A software system is designed for computing the DHT, inverse DHT, and discrete Fourier transform (DFT). Among all the compilers and assemblers available, QuickBASIC (BASIC compiler) 3.0 [28] and Microsoft Macro Assembler 5.0 [27] are chosen for use because of their IEEE real number format and the ease of programming. The designed system is menu-driven. It consists of one compiled BASIC program called HARTLEY and five assembly language programs named P2, FHT, SCALE, DFT, and SEQ. The block diagram of the system is shown in Figure 6.2. All the source programs can be found in the appendix.

The Compiled BASIC program HARTLEY serves as the coordinator for the assembly language programs. It defines the real number variables required for computation, allows input of the data, calls required assembly language programs, and prints the result after computation. Input can be of two forms: auto or manual. Auto input will generate a data set from 1 to N with an interval of 1. It is designed for

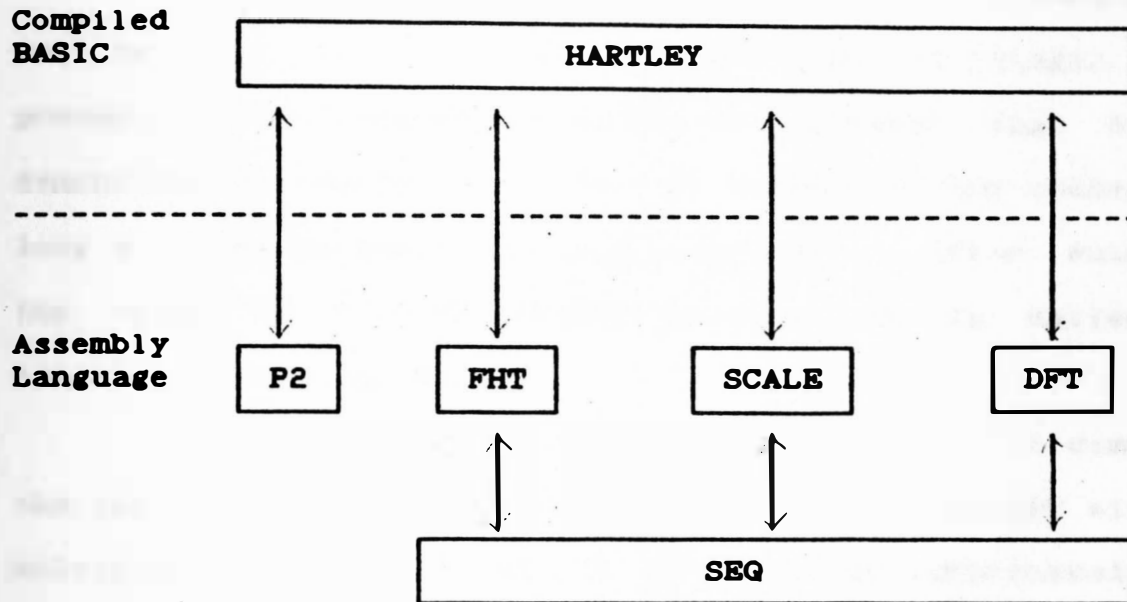


Figure 6.2 The DHT computation system

testing and demonstration. After computing the DHT, the result can be shown on the terminal or sent to the printer.

As mentioned before, for a sequence of N elements, N must be equal to 2^P , where P is a positive integer. At present P is limited to being not greater than 8 for demonstration purpose. It can be changed to any number as long as there is sufficient memory storage. After entering the value of P in the HARTLEY program, P2 is called to calculate the value of N .

Program FHT is the heart of the system. It computes the inverse DHT which is the same as the direct DHT without multiplying by a factor of $1/N$. For ease of comprehension of the program, it is translated into a BASIC program as shown in figure 6.3. One point to note is the calculation of the trigonometric functions. The partial tangent instruction (FPTAN) is the only trigonometric function available in the Intel 8087. It computes $\tan A$, where A , in radians, is the top stack element and is between 0 and $\pi/4$. The result is a ratio Y/X , with Y replacing A and X being pushed onto the stack. When $0 < A < \pi/4$, with the aid of this function, $\sin A$ and $\cos A$ can be calculated by $Y/\text{SQRT}(X^2 + Y^2)$ and $X/\text{SQRT}(X^2 + Y^2)$ respectively [25]. When A is outside its acceptable range, some procedures are used to obtain a suitable range that is inside the 0 to $\pi/4$ range. The program first computes all the angles required for the FHT and reduces them to the range between 0 and $\pi/4$, i.e. $\text{MOD}(\pi/4)$. It also keeps

```

10 GOSUB 60 'bit-reversal rtn
20 GOSUB 220 'cal. angles required
30 GOSUB 350 'cal. FHT
40 END
50 'BIT-REVERSAL RTN
60 FOR I=0 TO L-1
70 K=0
80 J=I
90 FOR Q=1 TO P
100 H=J
110 K=K*2+(J-2*H)
120 J=H
130 NEXT Q
140 IND(I)=K
150 NEXT I
160 FOR I=0 TO L-1
170 Y(I)=X(IND(I))
180 NEXT I
190 FOR I=0 TO L-1 : X(I)=Y(I) : NEXT I
200 RETURN
210 ' CAL. ANGLES REQUIRED
220 K=1 : BX=0 : P2=6.283185 : J=P
230 N2=K : K=K+K : FK=P2/K
240 R=N2 : I=0
250 F=FK*I
260 ANG(BX)=F
270 BX=BX+1
280 I=I+1 : R=R-1 : IF R<>0 THEN 250
290 J=J-1 : IF J<>0 THEN 230
300 Q=BX : BX=0
310 C(BX)=COS(ANG(BX)) : S(BX)=SIN(ANG(BX))
320 BX=BX+1 : Q=Q-1 : IF Q<>0 THEN 310
330 RETURN
340 REM FHT TRANSFORM SECTION
350 K=1 : BX=0 : J=P
360 GOSUB 500
370 N2=K : K=K+K : NM = L-K : Q=NM\K+1
380 R=N2 : I=0
390 T=Q : BP=I : SI=I+N2 : DI=K -I
400 U=X(SI)*C(BX) + X(DI)*S(BX)
410 W=X(BP)
420 Y(BP)=W+U
430 Y(SI)=W-U
440 BP=BP+K : T=T-1 : SI=SI+K : DI=DI+K
450 IF T<>0 THEN 400
460 BX=BX+1 : I=I+1 : R=R-1 : IF R<>0 THEN 390
470 J=J-1 : IF J<>0 THEN 360
480 FOR M=0 TO L-1 : X(M)=Y(M)/L : NEXT M
490 RETURN
500 FOR M=0 TO L-1 : X(M) = Y(M) : NEXT M
510 RETURN

```

Figure 6.3 Program FHT in BASIC

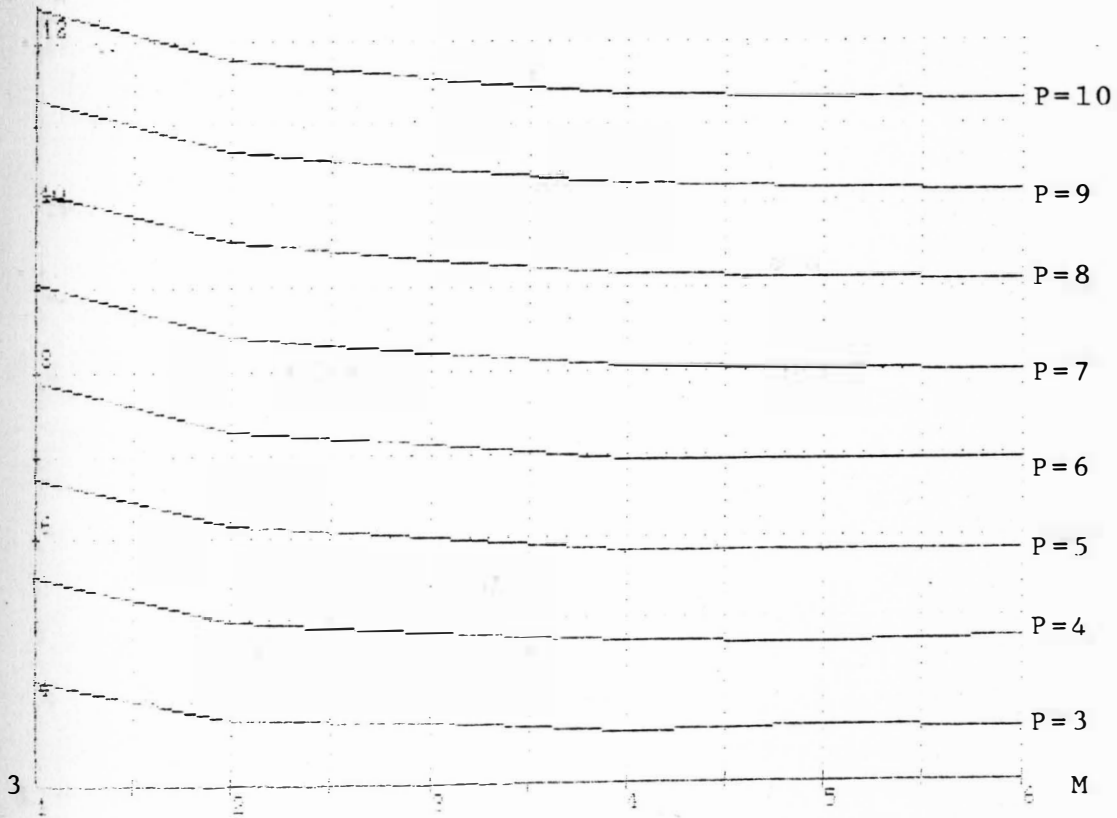
$\log_2 T$ (ms)

Figure 6.5 Run times for the inverse DHT

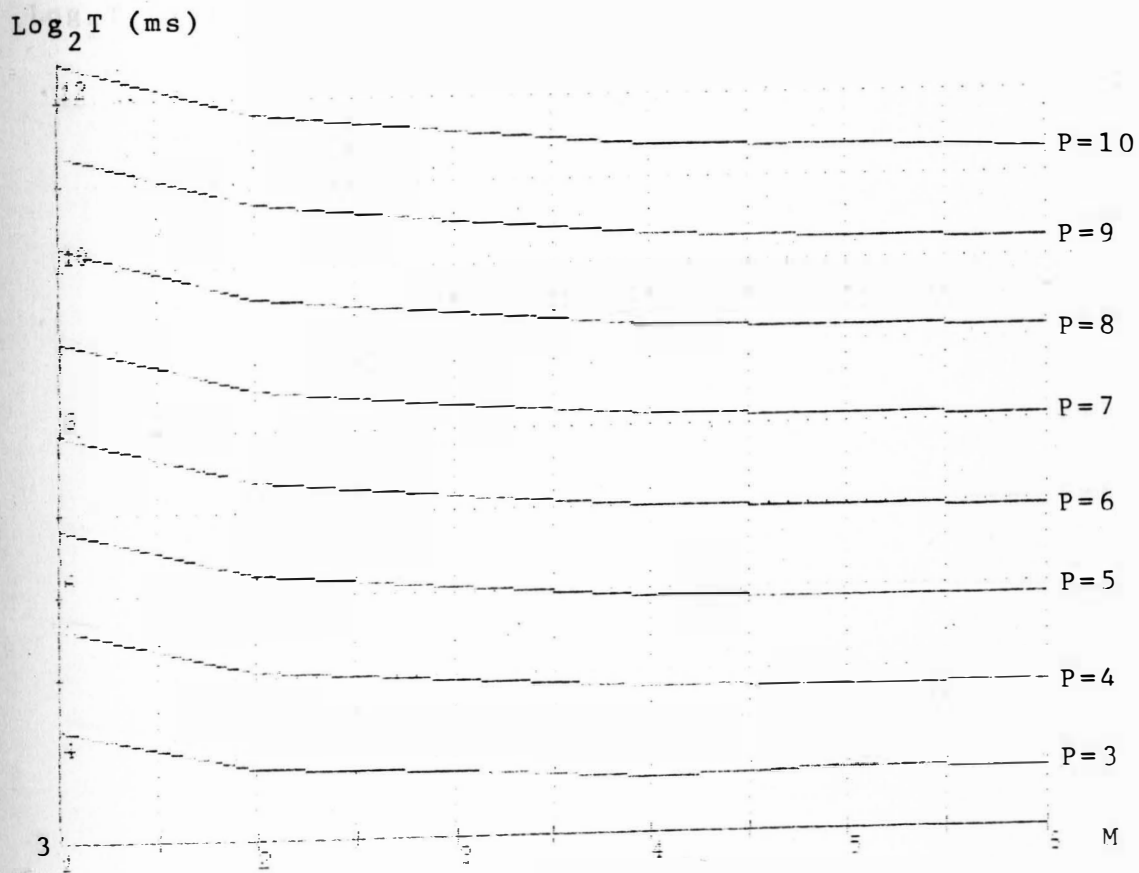


Figure 6.6 Run times for the DHT

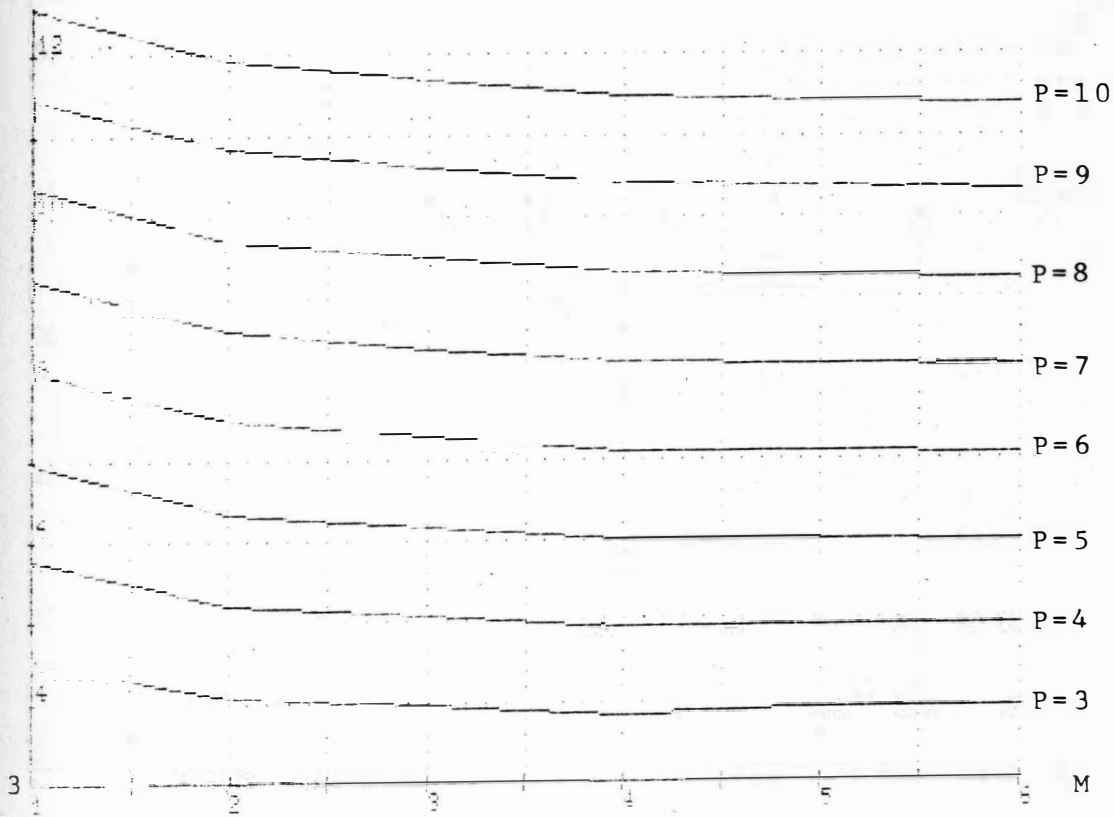
$\text{Log}_2 T$ (ms)

Figure 6.7 Run times for the DFT

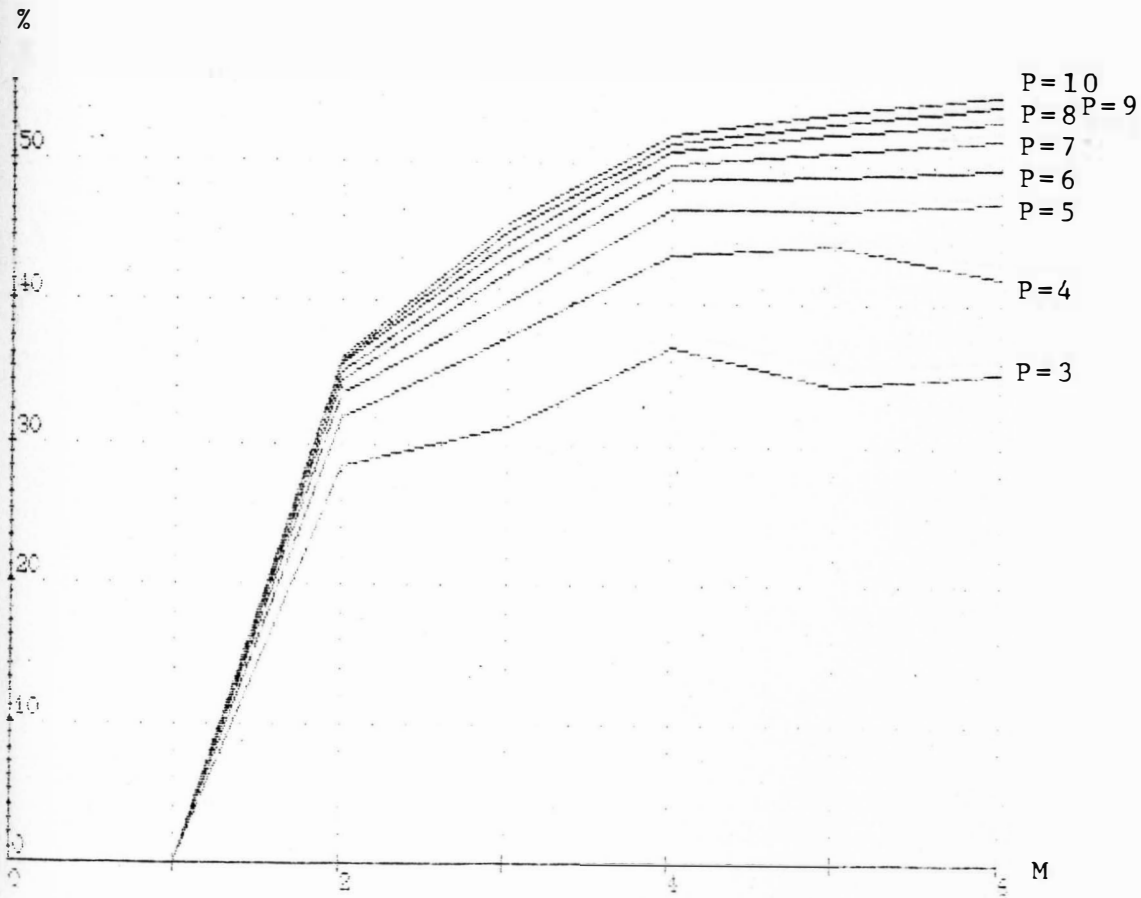


Figure 6.8 Improvement factor for the inverse DHT

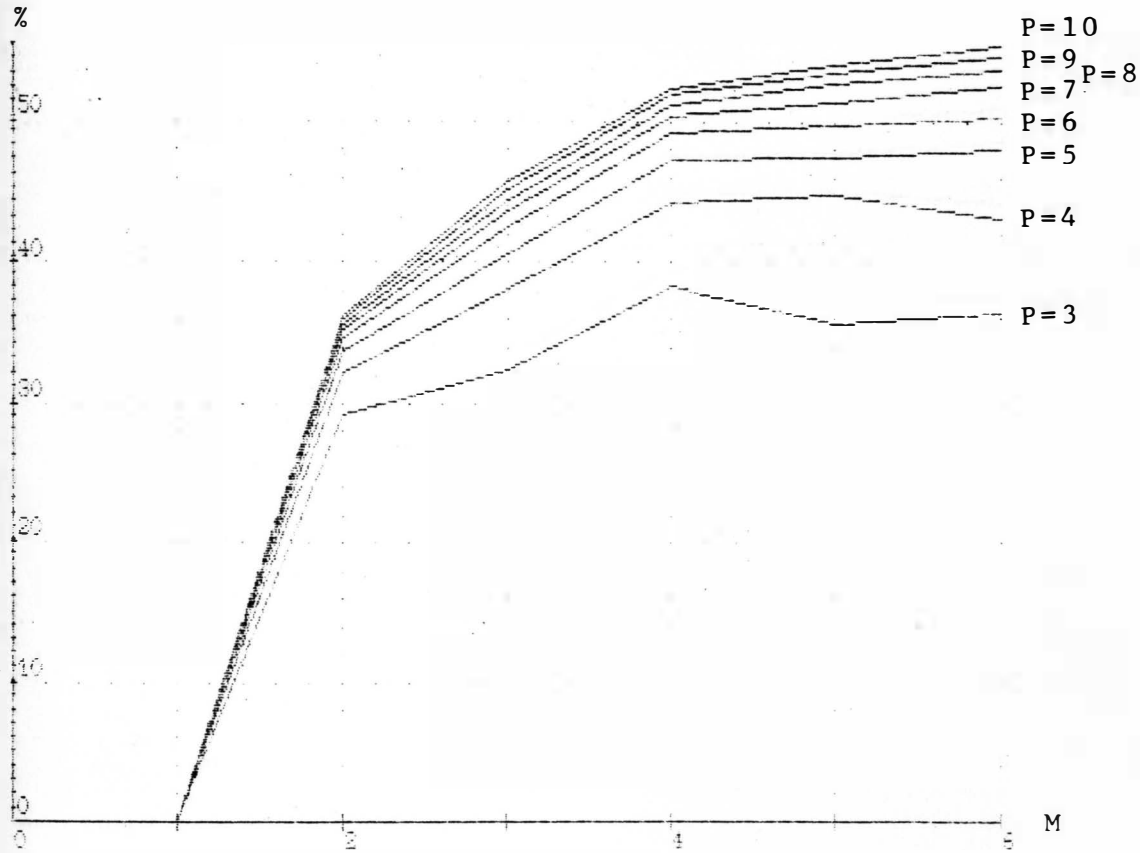


Figure 6.9 Improvement factor for the DHT

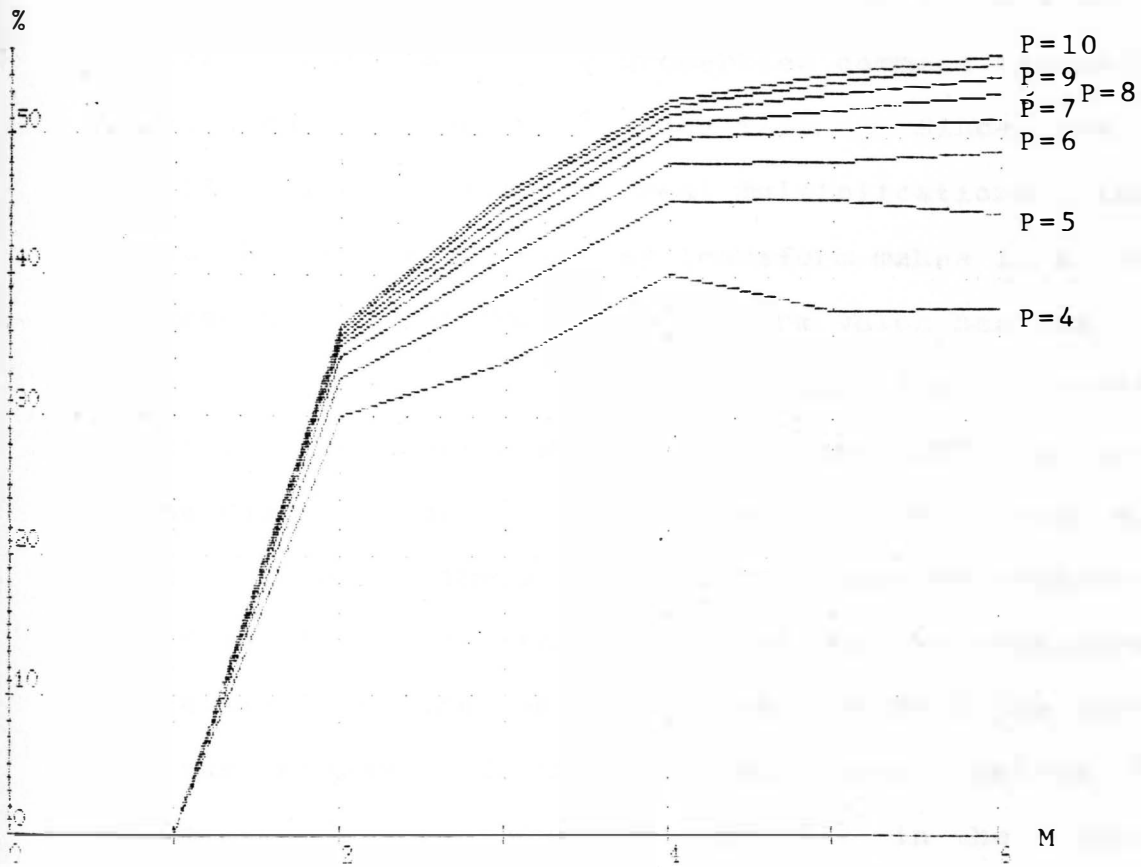


Figure 6.10 Improvement factor for the DFT

CHAPTER SEVEN

CONCLUSION

The Hartley transform is closely related to the Fourier transform. Its properties commend themselves for application to numerical analysis. Since one complex multiplication equals four real multiplications, the real-valued kernel of the Hartley transform makes it a potential replacement to the Fourier transform which has the complex-valued kernel. Most of the existing fast algorithms for computing the discrete Fourier transform (DFT) can be applied to the discrete Hartley transform (DHT) with some modifications. However, the speed of a fast Hartley transform (FHT) should be about twice that of a fast Fourier transform (FFT). In addition, one program is required for both the forward FHT and the inverse. In the FFT, additional control must be implemented for reversing the sign of i in the exponential $\exp(-i2\pi kn/N)$ during the operation of changing from a forward transformation to an inverse transformation or vice-versa. The FHT can also serve as the intermediate stage for the computation of other transforms such as the DFT and the discrete cosine transform (DCT) [13] so that time reduction can be achieved. Although the Hartley transform is a new term in signal processing, many researches on this subject have been carried out [16]-[20]. The growing interest will generate more researches and publications in the near future.

The vector processor (VP) discussed here is a single-instruction, multiple-data stream computer system. The purpose of this device is to speed up the processing of a group of scientific calculations. The greatest improvement can be obtained by minimizing the total data transfer time for each result element. Data independence, avoiding branching on comparison, and reducing data transfer time are the key considerations for programming the parallel processing routine on a VP. Sometimes this is done at the expense of the memory storage. However, the memory storage is relatively less expensive nowadays. A program running on a VP usually consists of scalar mode and parallel mode operations. It should be designed in such a way that scalar mode procedure is at minimum level. If scalar mode procedures dominate the whole program, overall performance of the system will be less significantly improved although parallel mode procedures are at the desired level.

A software system has been designed for computing the DHT, the inverse DHT, and the DFT on a VP. In this particular system, the improvement factor increases with both the number of the data points and the number of the mathematical coprocessors.

Finally, loading a dummy one to the unused MCs is an unproductive process in the VP. Hence it should be eliminated if the improvement of the VP is considered in the future.

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LIST OF SYMBOLS

| | |
|------------|-----------------------------------|
| DFT | discrete Fourier transform |
| DHT | discrete Hartley transform |
| DIF | decimation-in-frequency |
| DIT | decimation-in-time |
| FFT | fast Fourier transform |
| FHT | fast Hartley transform |
| Hz | Hertz |
| GPP | general purpose processor |
| MC | mathematical coprocessor |
| SR | split-radix |
| VCU | vector control unit |
| VP | vector processor |

APPENDIX
COMPUTER PROGRAMS

A software system which consists of one compiled BASIC program and five assembly language programs is designed for computing the fast Hartley transform (FHT) on a vector processor (VP). The block diagram of the system is shown in figure 6.2. Their source programs are listed in this appendix and can be found as follows:

| | | | |
|----------------|---|------------|---|
| HARTLEY | - | Figure A-1 | , |
| FHT | - | Figure A-2 | , |
| P2 | - | Figure A-3 | , |
| SCALE | - | Figure A-4 | , |
| DFT | - | Figure A-5 | , |
| SEQ | - | Figure A-6 | . |

At present the number of the data points N is limited to 256 for demonstration purpose. However, if the memory storage is sufficient, it can be expanded to any number by modifying the arrays' size in line 250 of **HARTLEY** and the sizes of **Y**, **SIN**, **COS**, **ANGLE**, and **LINDEX** as defined in the data segment of the program **FHT**. Note that if the VP is available, **NUM_PROC** which is defined in **FHT** should be changed to the actual number of the mathematical coprocessors (MCs) and the semi-colon preceded the VP's three instructions should be removed.

```

100 '*****'
110 ' This program is to calculate direct and inverse '
120 ' Hartley transform by calling other assembly '
125 ' language subroutines. '
130 ' Assembly language subroutines required: '
140 '           P2 '
150 '           FHT '
160 '           SEQ '
170 '           SCALE '
180 '           DFT '
190 ' Author : Boon Pock Lim '
200 ' Program Name : HARTLEY '
210 ' Date : March 14, 1988 '
220 '*****'
230 DEFINT N,P
240 DEFSNG I,R,X,Y
250 DIM X(256), Y(256), RE(256), IM(256)
255 '
260 ' MAIN RTN
265 '
280 GOSUB 2100 'DISPLAY HEADING
290 INPUT "N = 2^P ,      Input P = ",P
300 IF P < 1 OR P > 8 THEN 280
310 CALL P2(P,N) 'CAL. N
315 LOCATE 5,1 : PRINT SPACE$(10);"P =";P,"N =";N
320 GOSUB 2200 'INPUT DATA
330 GOSUB 2500 'SAVE INPUT DATA
340 GOSUB 2600 'TRANSFORMATION
350 INPUT "New data set ? (Y/N) ",A$
360 IF A$ = "Y" OR A$ = "y" THEN 280
370 END
2095 '
2100 'DISPLAY HEADING
2105 '
2110 CLS
2120 LOCATE 1,1 : PRINT "Date : ";DATE$
2130 LOCATE 1,35 : PRINT "S D S U"
2140 LOCATE 2,26 : PRINT "Electrical Engineering Dept."
2150 LOCATE 4,30 : PRINT "The Hartley Transform"
2160 RETURN
2199 '
2200 ' INPUT DATA
2205 '
2210 LOCATE 7,20 : PRINT "Input data set"
2220 LOCATE 9,20 : PRINT "(1) Auto"
2230 LOCATE 11,20 : PRINT "(2) Manual"
2240 LOCATE 13,20 : INPUT "Option : ",OPT
2250 IF OPT = 1 OR OPT = 2 THEN 2270

```

Figure A-1 HARTLEY program listing

```
2260 LOCATE 13,1 : PRINT SPACE$(80) : GOTO 2240
2270 ON OPT GOSUB 2300, 2400
2280 RETURN
2295 '
2300 ' AUTO INPUT
2305 '
2310 FOR J = 0 TO N-1
2320 X(J) = J+1
2330 NEXT J
2340 FOR J = 7 TO 13 STEP 2
2350 LOCATE J,1 : PRINT SPACE$(80)
2360 NEXT J
2370 RETURN
2395 '
2400 ' MANUAL INPUT
2405 '
2410 FOR J = 9 TO 13 STEP 2
2420 LOCATE J,1 : PRINT SPACE$(80)
2430 NEXT J
2440 LN = 9
2445 FOR J = 0 TO N-1
2447 IF LN <> 23 THEN 2460
2450 FOR K = 9 TO 23
2452 LOCATE K,1 : PRINT SPACE$(80)
2454 NEXT K
2456 LN = 9
2460 LOCATE LN,1 : PRINT "X (;J;) = "; : INPUT "",X(J)
2470 LN = LN + 1
2480 NEXT J
2482 FOR J = 7 TO LN
2484 LOCATE J,1 : PRINT SPACE$(80)
2486 NEXT J
2490 RETURN
2495 '
2500 ' SAVE INPUT DATA
2505 '
2510 FOR J = 0 TO N-1
2520 Y(J) = X(J)
2530 NEXT J
2540 RETURN
2595 '
2600 ' TRANSFORMATION
2605 '
2610 LOCATE 7,20 : PRINT "(1) Direct FHT"
2620 LOCATE 9,20 : PRINT "(2) Inverse FHT"
2630 LOCATE 11,20 : INPUT "Option : ",OPT
2640 IF OPT = 1 OR OPT = 2 THEN 2660
2650 LOCATE 11,1 : PRINT SPACE$(80) : GOTO 2630
```

Figure A-1 (continue)

```

2660 ON OPT GOSUB 2700, 2800
2665 GOSUB 2900                                'OUTPUT
2670 RETURN
2695 '
2700 ' DIRECT FHT
2705 '
2710 CALL FHT(X(0),P)
2720 CALL SCALE(X(0),N)
2730 CALL DFT(X(0),N,RE(0),IM(0))
2740 RETURN
2795 '
2800 ' INVERSE FHT
2805 '
2810 CALL FHT(X(0),P)
2820 RETURN
2895 '
2900 ' OUTPUT
2905 '
2910 LOCATE 7,20 : PRINT "Output           "
2920 LOCATE 9,20 : PRINT "(1) Display      "
2930 LOCATE 11,20 : PRINT "(2) Print        "
2940 LOCATE 13,20 : INPUT "Option : ",T
2950 IF T = 1 OR T = 2 THEN 2970
2960 LOCATE 13,1 : PRINT SPACE$(80) : GOTO 2940
2970 FOR J = 9 TO 13 STEP 2
2972 LOCATE J,1 : PRINT SPACE$(80)
2974 NEXT J
2980 ON T GOSUB 3000, 3300
2990 RETURN
2995 '
3000 ' DISPLAY OUTPUT
3005 '
3010 LOCATE 9,1
3020 ON OPT GOSUB 3200,3250                    'SUB HEADING
3040 LN = 9
3050 FOR J = 0 TO N-1
3060 IF LN <> 22 THEN 3110
3070 LOCATE 23,1 : PRINT "PRESS ANY KEY TO CONTINUE"
3080 K$ = INKEY$ : IF K$ = "" THEN 3080
3090 FOR K = 10 TO 23
3092 LOCATE K,1 : PRINT SPACE$(80)
3094 NEXT K
3100 LN = 9
3105 LOCATE 10,1
3110 LN = LN + 1
3120 PRINT USING "###";J;TAB(10);
3130 ON OPT GOSUB 3220,3260                    'DISPLAY OUTPUT
3150 NEXT J

```

Figure A-1 (continue)

```

3160 RETURN
3195 '
3200 ' DISPLAY DIRECT FHT SUB HEADING
3205 '
3207 PRINT " n,k";TAB(18);"f(n);TAB(30);"H(k)";
3210 PRINT TAB(41);"RE[F(k)]";TAB(54);"IM[F(k)]"
3212 RETURN
3215 '
3220 ' DISPLAY DIRECT FHT
3222 '
3230 PRINT USING " ####.#####";Y(J),X(J),RE(J),IM(J)
3232 RETURN
3245 '
3250 ' DISPLAY INVERSE FHT SUB HEADING
3252 '
3254 PRINT " n,k";TAB(18);"H(k);TAB(30);"f(n)"
3256 RETURN
3258 '
3260 ' DISPLAY INVERSE FHT
3262 '
3264 PRINT USING " ####.#####";Y(J),X(J)
3266 RETURN
3295 '
3300 ' PRINT OUTPUT
3305 '
3310 PG = 0 : GOSUB 4000 'PRINT HEADING
3320 FOR J = 0 TO N-1
3330 IF LN > 56 THEN GOSUB 4000
3340 LN = LN + 1
3350 LPRINT USING "###";J;TAB(10);
3360 ON OPT GOSUB 4150,4250 'PRINT OUTPUT
3380 NEXT J
3390 RETURN
3995 '
4000 ' PRINT HEADING
4005 '
4010 LPRINT CHR$(12) : PG = PG + 1 : LPRINT
4020 LPRINT "Date : ";DATE$;TAB(35);"S D S U";TAB(71);"Page";PG
4030 LPRINT TAB(26);"Electrical Engineering Dept."
4040 LPRINT
4045 LPRINT "The Hartley Transform : P =";P;" , N = 2^P =";N
4047 LPRINT
4050 ON OPT GOSUB 4100,4200 'SUB HEADING
4070 LN = 7
4080 RETURN
4095 '
4100 ' PRINT DIRECT FHT SUB HEADING
4105 '

```

Figure A-1 (continue)

```
4110 LPRINT " n,k";TAB(18);"f(n)";TAB(30);"H(k)";
4115 LPRINT TAB(41);"RE[F(k)]";TAB(54);"IM[F(k)]"
4120 RETURN
4145 '
4150 ' PRINT DIRECT FHT
4155 '
4160 LPRINT USING " ####.#####";Y(J),X(J),RE(J),IM(J)
4170 RETURN
4195 '
4200 ' PRINT INVERSE FHT SUB HEADING
4205 '
4210 LPRINT " n,k";TAB(18);"H(k)";TAB(30);"f(n)"
4220 RETURN
4245 '
4250 ' PRINT INVERSE FHT
4255 '
4260 LPRINT USING " ####.#####";Y(J),X(J)
4270 RETURN
```



```

.8087
.MODEL MEDIUM
.CODE
;-----;
; Program Name : FHT ;
; This subroutine(Procedure) calculates the Fast Hartley ;
; Transform by using radix-2 method (decimation in time). ;
; From compiled BASIC program : CALL FHT(X(0),P) ;
; Subroutines(Procedures) required : BIT_REV ;
; ; REORDER_Y ;
; ; Y_TO_X ;
; ; ANGLE_TB ;
; ; H_TRAN ;
; Author : Boon Pock Lim ;
; Date : Feb 22, 1988 ;
;-----;

EXTRN SEQ_LD_2:FAR
EXTRN SEQ_ST_2:FAR
EXTRN SEQ_LD_4:FAR
EXTRN SEQ_ST_4:FAR

PUBLIC FHT
FHT PROC FAR
PUSH BP
MOV BP,SP
PUSH AX ;Save the registers
PUSH BX
PUSH CX
PUSH DX ;-----
MOV BX,[BP]+8 ;1st Arg - Store the
MOV X,BX ;address in X
MOV BX,[BP]+6 ;2nd Arg
MOV CX,[BX]
MOV POWER,CX
MOV AX,1 ;Calculate # of
SHL AX,CL ;elements N
MOV N,AX ;N = 2^POWER
MOV CX,AX ;Set CX = N
XOR BX,BX ;Set BX = 0
XOR DX,DX ;Set DX = 0

MAIN_1:
MOV LINDEX[BX],DX ;Set up index table
ADD BX,2 ;0,1,2,.....,N-1
INC DX
LOOP MAIN_1 ;-----
;
; JMP MAIN_D_1 ;Dummy jump
;MAIN_D_1:

```

Figure A-2 FHT program listing

```

;           DB           0DFh,0FDh           ;Parallel mode
;
MOV        DX,POWER           ;Set DX = POWER
CALL       BIT_REV           ;Bit reversing
CALL       REORDER_Y        ;Y = new order of X
CALL       Y_TO_X           ;Set X = Y
CALL       ANGLE_TB         ;set up angle table
CALL       H_TRAN           ;FHT
;
;           JMP          MAIN_D_2           ;Dummy jump
;MAIN_D_2:
;           DB           0DFh,0FEh           ;Scalar mode
;
POP        DX                 ;Restore the
POP        CX                 ;registers
POP        BX
POP        AX
POP        BP                 ;-----
RET        4
FHT          ENDP

```

```

PUBLIC BIT_REV
BIT_REV PROC NEAR
;-----;
; This procedure performs bit reversing operation ;
; Input parameters : AX = # of elements ;
;                   DX = # of bits ;
; Reads : LINDEX, BASE, TEMP ;
; Writes : LINDEX, TEMP ;
; Algorithm : Let J be the original index ;
;             K = 0 ;
;             FOR Q = 1 TO (# of bits) ;
;                 L = J MOD 2 ;
;                 J = INT(J/2) ;
;                 K = K * 2 + L ;
;             NEXT Q ;
;             new index = K ;
; Contents of the 8087's registers : ;
; ST(7) = -1 ;
; ST(6) = J = original index ;
; ST(5) = 2 = BASE ;
; ST(4) = L = J MOD 2 ;
; ST(3) = 1 ;
; ST(2) = K ;
;-----;
PUSH     AX                   ;Save the registers
PUSH     BX

```

Figure A-2 (continue)

```

                PUSH    CX
                PUSH    DX
                DEC     AX
                DEC     AX
                OR      AX,AX
                JNZ    BIT
                JMP    BIT_4
BIT:
                MOV     BX,2
BIT_1:
                POP     DX
                PUSH    DX
                FINIT
                FSTCW   TEMP
                FWAIT
                OR      TEMP,0C00h
                FLDCW   TEMP
                FLD1
                FCHS
                PUSH    BX
                LEA    BX,LINDEX[BX]
                CALL   SEQ_LD_2
                POP     BX
                FILD   BASE
                FDECSTP
                FLD1
                FLDZ
BIT_2:
                FINCSTP
                FINCSTP
                FSTP   ST(0)
                FLD   ST(1)
                FPREM
                FINCSTP
                FINCSTP
                FSCALE
                FRNDINT
                FDECSTP
                FDECSTP
                FDECSTP
                FDECSTP
                FSCALE
                FADD   ST,ST(2)
                FWAIT
                DEC    DX
                OR     DX,DX
                JNZ   BIT_2
                PUSH   BX
                ;-----
                ;AX = N - 2 , 1st
                ;and last elements
                ;are the same.
                ;JZ
                ;Begin with 2nd
                ;element
                ;Get # of bits
                ;Store
                ;Initialize 8087
                ;Set 8087 rounding
                ;control -
                ;round down
                ;ST(7)=-1 , Scale
                ;factor
                ;Load index
                ;sequentially
                ;-----
                ;ST(5)=2 , Base
                ;SET ST=4
                ;ST(3)=1
                ;ST(2)=K ,
                ;init. value = 0
                ;SET ST=4
                ;-----
                ;ST(4)=L=J
                ;L=J MOD 2
                ;SET ST=6
                ;-----
                ;J=INT(J/2)
                ;-----
                ;SET ST=2
                ;-----
                ;K=K*2
                ;K=K*2+L
                ;remainder
                ;Do POWER times
                ;Store new index

```

Figure A-2 (continue)

```

LEA     BX,LINDEX[BX]    ;sequentially
CALL    SEQ_ST_2
POP     BX                ;-----
MOV     CX,NUM_PROC      ;Set next address
CMP     CX,AX
JBE     BIT_3
MOV     CX,AX

BIT_3:
ADD     BX,2
LOOP    BIT_3            ;-----
SUB     AX,NUM_PROC      ;Do (N-2) times
CMP     AX,0
JLE     BIT_4            ;JG will cause out
JMP     BIT_1            ;of range error

BIT_4:
POP     DX                ;Restore the
POP     CX                ;registers
POP     BX
POP     AX                ;-----
RET

BIT_REV      ENDP

PUBLIC REORDER_Y
REORDER_Y    PROC        NEAR
;-----;
; This procedure makes Y equal to bit-reversal order of X. ;
; Input parameter : AX = N ;
; Reads : LINDEX, X ;
; Writes : Y ;
;-----;

PUSH     AX                ;Save the registers
PUSH     BX
PUSH     CX
PUSH     SI
PUSH     DI                ;-----
MOV     CX,AX              ;Set CX = N
XOR     BX,BX              ;Set BX = 0
XOR     DI,DI              ;Set DI = 0

;
;
;REORDER_D_1:
DB      0DFh,0FEh          ;Scalar mode
;
FINIT                          ;Initialize 8087

```

Figure A-2 (continue)

```

REORDER:
        MOV     SI,LINDEX[BX]    ;Find the actual
        ADD     SI,SI            ;address of X
        ADD     SI,SI            ;Offset = index * 4
        ADD     SI,X             ;-----
        FLD     DWORD PTR [SI]   ;Move X to Y
        FSTP    Y[DI]           ;-----
        ADD     DI,4             ;Next addr of Y
        ADD     BX,2             ;Next addr of index
        LOOP    REORDER         ;Do N times
;
;
;REORDER_D_2:
        DB     0DFh,0FDh        ;Parallel mode
;
;
        POP     DI               ;Restore the
        POP     SI               ;registers
        POP     CX
        POP     BX
        POP     AX               ;-----
        RET
REORDER_Y ENDP

                PUBLIC  Y_TO_X
Y_TO_X         PROC    NEAR
;-----;
; This procedure moves Y to X.
; Input parameter : AX = N
; Reads : Y
; Writes : X
;-----;
        PUSH    AX               ;Save the registers
        PUSH    CX
        PUSH    SI
        PUSH    DI               ;-----
        MOV     CX,AX            ;Set CX = N
        XOR     SI,SI           ;Set SI = 0
        MOV     DI,X            ;DI = 1st addr of X
;
;
        JMP     Y_X_D_1         ;Dummy jump
;Y_X_D_1:
        DB     0DFh,0FEh        ;Scalar mode
;
;
        FINIT                    ;Initialize 8087
Y_X_MOVE:      FLD     Y[SI]     ;Transfer Y to X

```

Figure A-2 (continue)

```

FSTP    DWORD PTR [DI] ;-----
ADD     SI,4           ;Next addr - 2 for
ADD     DI,4           ;word, 4 for S. real
LOOP    Y_X_MOVE       ;and 8 for D. real
;
;
;Y_X_D_2:
;
;
;
POP     DI             ;Restore the
POP     SI             ;registers
POP     CX
POP     AX             ;-----
RET
Y_TO_X  ENDP

```

```

PUBLIC  ANGLE_TB
ANGLE_TB PROC NEAR
;-----;
; This procedure calculates sine and cosine of the angles ;
; required for the FHT. ;
; Reads : POWER, ANGLE, LINDEX, SIN, COS ;
; Writes : ANGLE, LINDEX, SIN, COS ;
; Intermediate variables : K, N2, I ;
; Refer to the BASIC program for more information ;
;-----;
PUSH    AX             ;Store the register
PUSH    BX
PUSH    CX
PUSH    DX
PUSH    SI             ;-----;
;-----;
; This routine calculates all the angles required for the ;
; FHT and reduce them to the range between 0 and PI/4 ;
; i.e. MOD PI/4. It also calculates the corresponding ;
; octant (0 - 7) of the angles. ;
;-----;
MOV     K,1           ;Set K = 1
XOR     BX,BX         ;Set BX = 0
XOR     SI,SI         ;Set SI = 0
FINIT                          ;Initialize 8087
FLD1                          ;ST(7)=1
FLDPI                          ;ST(6)=PI
FSCALE                          ;ST(6)=2*PI
FSTP   ST(1)           ;ST(7)=2*PI
FILD   BASE            ;ST(6)=2
FCHS                          ;ST(6)=-2

```

Figure A-2 (continue)

```

        FLDPI                ;ST(5)=PI
        FSCALE               ;ST(5)=PI/4
        FSTP                 ST(1) ;ST(6)=PI/4
        FWAIT
        XOR                  DI,DI  ;Set DI=0
        MOV                  CX,POWER ;1st loop, CX=POWER

ANG_POWER:
        MOV                  AX,K    ;N2 = K
        SHL                  K,1    ;K = K + K
        FILD                 K      ;ST(5)=K
        FDIVR                ST,ST(2) ;ST(5)=2*PI/K
        FWAIT
        XOR                  DX,DX   ;Set DX = 0

ANG_I_LP:
        CALL                 ANG_LD_I ;Load I sequentially
        FMUL                 ST,ST(1) ;ST(4)=2*PI*I/K
        FLD                  ST(0)   ;ST(3)=ST(4)
        FINCSTP              ;Set ST=4
        FXCH                 ST(2)   ;ST(4)=PI/4
        FDECSTP              ;Set ST(3)
        FPREM                ;ST(3)=ST(3) MOD PI/4
        PUSH                 BX      ;Store ANGLE
                                   ;sequentially
        LEA                  BX,ANGLE[BX] ;0 <= ANGLE <= PI/4
        CALL                 SEQ_ST_4
        POP                  BX      ;-----
        FXCH                 ST(2)   ;Exchange
        FINCSTP              ;Set ST=5
        FXCH                 ST(2)   ;ST(5)=2*PI
        FDECSTP              ;Set ST=4
        FPREM                ;ST(4)=ST(4) MOD 2*PI
        FDIV                 ST,ST(2) ;ST(4)=ST(4)/(PI/4)
        FSTCW                TEMP    ;Set 8087 rounding control
        FWAIT                ;control
        OR                   TEMP,0C00h ;Round down
        FLDCW                TEMP
        PUSH                 BX      ;Store the octant
        LEA                  BX,LINDEX[SI] ;(0-7)
        CALL                 SEQ_ST_2
        POP                  BX
        FXCH                 ST(2)   ;Exchange
        AND                  TEMP,0F3FFh ;Set 8087 rounding
        FLDCW                TEMP    ;control - round to
        FWAIT                ;nearest or even
        PUSH                 CX      ;Set next address
        MOV                  CX,NUM_PROC
        CMP                  CX,AX
        JBE                  ANG_INC_1

```

Figure A-2 (continue)

```

ANG_INC_1:      MOV      CX,AX
                ADD      BX,4
                ADD      SI,2          ;-----
                INC      DI          ;DI=DI+1
                INC      DX          ;DX=DX+1
                LOOP     ANG_INC_1
                POP      CX
                MOV      I,DX        ;I=DX
                SUB      AX,NUM_PROC
                CMP      AX,0
                JLE     ANG_A
                JMP      ANG_I_LP

ANG_A:
                FSTP     ST(0)
                FWAIT
                DEC      CX          ;LOOP will cause out
                OR       CX,CX       ;of range error
                JZ       ANG_NEXT   ;-----
                JMP      ANG_POWER   ;1st loop end

ANG_NEXT:
                MOV      DX,DI        ;DX=# of elements in
                                     ;ANGLE
; ~~~~~~;
; This routine changes the angle A in octant 1,3,5,7 to ;
; (PI/4 - A). ;
; ~~~~~~;
;
;
;                JMP      ANG_D_1    ;Dummy jump
;ANG_D_1:
;                DB      0DFh,0FEh   ;Scalar mode
;
                FINIT     ;Reset 8087
                FILD     BASE       ;ST(7)=2
                FCHS     ;ST(7)=-2
                FLDPI    ;ST(6)=PI
                FSCALE   ;ST(6)=PI/4
                FSTP     ST(1)      ;ST(7)=PI/4
                XOR      BX,BX      ;BX=0
                XOR      SI,SI      ;SI=0
                MOV      CX,DX      ;CX=DX
ANG_ADJ:
                MOV      AX,LINDEX[SI] ;Loop begin
                SHR      AX,1       ;Check the octant
                JNC     ANG_OUT     ;(PI -angle) if it
                FLD     ANGLE[BX]  ;is odd
                FSUBR    ST(0),ST(1)
                FSTP     ANGLE[BX]  ;-----

```

Figure A-2 (continue)


```

ANG_OUT:      FWAIT

              ADD     BX,4          ;Set next addresses
              ADD     SI,2          ;-----
              LOOP    ANG_ADJ       ;Loop end
;
;
;ANG_D_2:     JMP     ANG_D_2       ;Dummy jump
;
;              DB     0DFh,0FDh     ;Parallel mode
;
;~~~~~;
; This routine calculates sine and cosine of the angle ;
; with a proper sign.                                  ;
;~~~~~;
              XOR     BX,BX         ;BX=0
              XOR     SI,SI         ;SI=0
              MOV     AX,DX         ;AX=DX

ANG_CAL:      FINIT                ;Reset 8087
              FLD1                ;ST(7)=Sign of sine
              PUSH    BX           ;=1
              MOV     BX,SI
              CALL    SIGN_S
              POP     BX
              FLD1                ;ST(6)=Sign of
              PUSH    BX           ;cosine=1
              MOV     BX,SI
              CALL    SIGN_C
              POP     BX

ANG_CONT:     PUSH    BX           ;Load ANGLE
              LEA    BX,ANGLE[BX] ;sequentially
              CALL    SEQ_LD_4
              POP     BX
              FPTAN                ;ST(5)=Y, ST(4)=X
              FMUL    ST,ST(2)     ;Set sign for cosine
              FXCH    ST(1)        ;ST(5)=X, ST(4)=Y
              FMUL    ST,ST(3)     ;Set sign for sine
              FLD     ST(0)        ;ST(3)=ST(4)=Y
              FMUL    ST,ST(1)     ;ST(3)=Y*Y
              FLD     ST(2)        ;ST(2)=X
              FMUL    ST,ST(3)     ;ST(2)=X*X
              FADD    ST,ST(1)     ;ST(2)=(X*X) + (Y*Y)
              FSQRT                ;ST(2)=SQRT(ST(2))
              FLD     ST(0)        ;ST(1)=ST(2)
              FDIVR   ST,ST(3)     ;ST(1)=ST(4)/ST(1)
              PUSH    BX           ;=sine
              LEA    BX,SIN[BX]    ;Store SIN

```

Figure A-2 (continue)

```

CALL      SEQ_ST_4      ;sequentially
POP       BX            ;-----
FDIVR    ST,ST(3)      ;ST(2)=ST(5)/ST(2)
PUSH     BX            ;=cosine
LEA     BX,COS[BX]     ;Store COS
CALL     SEQ_ST_4      ;sequentially
POP      BX            ;-----
PUSH     CX
MOV      CX,NUM_PROC
CMP      CX,AX
JBE     ANG_INC_2
MOV      CX,AX

ANG_INC_2:
ADD      BX,4          ;Set next addresses
ADD      SI,2          ;-----
LOOP    ANG_INC_2
POP      CX
SUB      AX,NUM_PROC
CMP      AX,0
JG      ANG_CAL        ;Loop end
; ~~~~~~;
; This routine swaps the values of sine and cosine for ;
; the angle in octant 1,2,5,6. ;
; ~~~~~~;
;
;
;          JMP      ANG_D_3      ;Dummy jump
;ANG_D_3:
;          DB      0DFh,0FEh     ;Scalar mode
;
;          FINIT     ;Reset 8087
;          XOR      BX,BX        ;BX=0
;          XOR      SI,SI        ;SI=0
;          MOV      CX,DX        ;CX=DX
ANG_SWAP:
;          ;Loop begin
;          CMP      LINDEX[SI],0 ;Swap sin and cos
;          JE      ANG_DONE     ;if octant is
;          CMP      LINDEX[SI],3 ;1,2,5,6
;          JE      ANG_DONE
;          CMP      LINDEX[SI],4
;          JE      ANG_DONE
;          CMP      LINDEX[SI],7
;          JE      ANG_DONE
;          FLD     SIN[BX]
;          FLD     COS[BX]
;          FSTP    SIN[BX]
;          FSTP    COS[BX]     ;-----

ANG_DONE:
ADD      BX,4          ;Set next addresses

```

Figure A-2 (continue)

```

                ADD     SI,2           ;-----
                LOOP    ANG_SWAP      ;Loop end
;
;
;ANG_D_4:      JMP     ANG_D_4        ;Dummy jump
;
;                DB     0DFh,0FDh     ;Parallel mode
;
                POP     SI           ;Restore the
                POP     DX           ;registers
                POP     CX
                POP     BX
                POP     AX           ;-----
                RET
ANGLE_TB      ENDP

ANG_LD_I      PROC     NEAR
;-----;
; This procedure loads integer I sequentially. ;
; Input parameters : AX = # of data to be processed ;
;                   DX = 1st # of I ;
; Reads : NUM_PROC ;
;-----;
                PUSH    AX           ;Save the registers
                PUSH    BX
                PUSH    CX
                PUSH    DX
                PUSH    SI           ;-----
;
;
;ANG_LD_D1:    JMP     ANG_LD_D1      ;Dummy jump
;
;                DB     0DFh,0FFh     ;Sequential mode
;
                MOV     BX,DX         ;BX=DX
                XOR     DX,DX         ;If AX < # of
                MOV     CX,NUM_PROC   ;processors then
                MOV     SI,CX         ;CX=AX else
                CMP     CX,AX         ;CX=# or processors
                JBE     ANG_LD_1      ;DX : Flag
                MOV     CX,AX
                MOV     DX,1

ANG_LD_1:     MOV     I,BX           ;Set I
                FILD   I             ;Load I
                INC    BX           ;BX=BX+1
                LOOP   ANG_LD_1
                OR     DX,DX
                JZ     ANG_LD_3
                MOV     CX,SI

```

Figure A-2 (continue)

```

ANG_LD_2:      SUB      CX,AX
               FLD1          ;Load dummy 1 to
               LOOP      ANG_LD_2      ;other processors.
ANG_LD_3:
;
;
;ANG_LD_D2:    JMP      ANG_LD_D2      ;Dummy jump
;
;               DB      0DFh,0FEh      ;Scalar mode
;               JMP      ANG_LD_D3      ;Dummy jump
;ANG_LD_D3:    DB      0DFh,0FDh      ;Parallel mode
;
;               POP      SI          ;Restore the
               POP      DX          ;registers
               POP      CX
               POP      BX
               POP      AX          ;-----
               RET
ANG_LD_I      ENDP

               PUBLIC  SIGN_S
SIGN_S        PROC    NEAR
;-----;
; This procedure checks the octant [LINDEX] sequentially ;
; and changes the sign of ST(7) (sine) if required.      ;
; Input parameters : AX = # of data to be processed      ;
;                   BX = index of LINDEX                 ;
; Reads : NUM_PROC, LINDEX                               ;
; Writes : ST(7)                                         ;
;-----;
               PUSH     AX          ;Save the registers
               PUSH     BX
               PUSH     CX          ;-----
;
;
;SIGN_S_D_1:   JMP      SIGN_S_D_1      ;Dummy jump
;
;               DB      0DFh,0FFh      ;Sequential mode
;
;               MOV     CX,NUM_PROC      ;Set the count
               CMP     CX,AX
               JBE     SIGN_S_1
               MOV     CX,AX
SIGN_S_1:     CMP     LINDEX[BX],2      ;If octant = 0,1,3,6
               JB     SIGN_S_2          ;then sine = +ve
               CMP     LINDEX[BX],3      ;else sine = -ve

```

Figure A-2 (continue)

```

                JE      SIGN_S_2
                CMP    LINDEX[BX],6
                JE      SIGN_S_2
SIGN_S_2:      FCHS                                ;-----
                ADD    BX,2                        ;Set next address
                LOOP   SIGN_S_1                    ;-----
;
;
;SIGN_S_D_2:   JMP     SIGN_S_D_2                  ;Dummy jump
;
;              DB     0DFh,0FEh                  ;Scalar mode
;              JMP     SIGN_S_D_3                  ;Dummy jump
;SIGN_S_D_3:   DB     0DFh,0FDh                  ;Parallel mode
;
;              POP    CX                          ;Restore the
                POP    BX                          ;registers
                POP    AX
SIGN_S         RET
                ENDP

                PUBLIC SIGN_C
SIGN_C         PROC    NEAR
;-----;
; This procedure checks the octant [LINDEX] sequentially ;
; and changes the sign of ST(6) (cosine) if required. ;
; Input parameters : AX = # of data to be processed ;
;                   BX = index of LINDEX ;
; Reads : NUM_PROC, LINDEX ;
; Writes : ST(6) ;
;-----;
                PUSH   AX                          ;Save the registers
                PUSH   BX
                PUSH   CX                          ;-----
;
;
;SIGN_C_D_1:   JMP     SIGN_C_D_1                  ;Dummy jump
;
;              DB     0DFh,0FFh                  ;Sequential mode
;
;              MOV    CX,NUM_PROC                  ;Set the count
                CMP    CX,AX
                JBE    SIGN_C_1
                MOV    CX,AX
SIGN_C_1:     CMP    LINDEX[BX],3                ;If octant = 0,1,2,7
                JB     SIGN_C_2                    ;then cosine = +ve
                CMP    LINDEX[BX],7                ;else cosine = -ve

```

Figure A-2 (continue)

```

SIGN_C_2:      JE      SIGN_C_2
               FCHS
               ;-----
               ADD     BX,2      ;Set next address
               LOOP   SIGN_C_1  ;-----
;
;
;SIGN_C_D_2:   JMP     SIGN_C_D_2   ;Dummy jump
;
;               DB     0DFh,0FEh ;Scalar mode
;               JMP     SIGN_C_D_3 ;Dummy jump
;SIGN_C_D_3:   DB     0DFh,0FDh ;Parallel mode
;
;               POP     CX      ;Restore the
               POP     BX      ;registers
               POP     AX      ;-----
SIGN_C        RET
               ENDP

```

```

PUBLIC H_TRAN
H_TRAN PROC NEAR
;-----
; This procedure calculates the FHT.
; Reads : X, Y, COS, SIN, POWER, N
; Writes : X, Y
; Intermediate variables : TEMP, K, N2, NM, T
; Refer to the BASIC program for more information
;-----
               PUSH    AX      ;Save the registers
               PUSH    BX
               PUSH    CX
               PUSH    DX
               PUSH    BP
               PUSH    SI
               PUSH    DI
               MOV     AX,N     ;T=N
               MOV     T,AX
               MOV     K,1     ;Set K = 1
               XOR     BX,BX    ;Set BX = 0
               MOV     CX,POWER ;Set 1st loop,
TRAN_POWER:   ;CX = POWER
               PUSH    CX      ;Save CX
               MOV     AX,N     ;Set input para. for
               CALL    Y_TO_X   ;subroutine and call

```

Figure A-2 (continue)

```

MOV      AX,K                ;N2 = K
MOV      N2,AX              ;-----
MOV      CL,2               ;2 for Single P. real
SHL      AX,CL              ;3 for Double P. real
MOV      N2_ADDRESS,AX     ;Convert N2 to addr
SHL      K,1                ;K = K + K
MOV      AX,K              ;-----
SHL      AX,CL              ;Convert K to addr
MOV      K_ADDRESS,AX     ;-----
SHR      T,1                ;T=T/2
XOR      DX,DX              ;Set DX = 0
MOV      CX,N2              ;Set 2nd loop,
TRAN_I_LP:                   ;CX = N2
MOV      BP,DX              ;BP = DX
MOV      SI,N2_ADDRESS     ;SI = N2 + DX (Addr)
ADD      SI,DX              ;-----
MOV      DI,K_ADDRESS     ;DI = K - DX (Addr)
SUB      DI,DX              ;-----
MOV      AX,T              ;Set 3rd loop,
TRAN_T_LP:                   ;CX = T
FINIT      ;Initialize 8087
PUSH      BX                ;Load W=X[BP]
MOV      BX,BP              ;sequentially
ADD      BX,X
CALL      SEQ_LD_4
POP      BX                  ;-----
FLD      COS[BX]           ;Load COS
PUSH      BX                ;Load X[SI]
MOV      BX,SI              ;sequentially
ADD      BX,X
CALL      SEQ_LD_4
POP      BX                  ;-----
FLD      SIN[BX]           ;Load SIN
PUSH      BX                ;Load X[DI]
MOV      BX,DI              ;sequentially
ADD      BX,X
CALL      SEQ_LD_4
POP      BX                  ;-----
FMULP    ST(1),ST          ;ST(4)=X[DI]*SIN[BX]
FXCH     ST(2)             ;ST(4)=COS[BX]
FMULP    ST(1),ST          ;ST(5)=X[SI]*COS[BX]
FADDP    ST(1),ST          ;ST(6)=U
FLD      ST(1)             ;ST(5)=W
FADD     ST,ST(1)          ;ST(5)=W+U
PUSH     BX                ;Store Y[BP]
LEA     BX,Y[BP]           ;sequentially
CALL    SEQ_ST_4
POP     BX                  ;-----

```

Figure A-2 (continue)

```

FSUBR    ST,ST(1)          ;ST(6)=W-U
PUSH     BX                ;Store Y[SI]
LEA      BX,Y[SI]         ;sequentially
CALL     SEQ_ST_4
POP      BX                ;-----
PUSH     AX
MOV      AX,K_ADDRESS     ;Set next addr for
PUSH     CX                ;BP,SI,DI
MOV      CX,NUM_PROC
CMP      CX,AX
JBE      TRAN_INC_1
MOV      CX,AX

TRAN_INC_1:
ADD      BP,AX
ADD      SI,AX
ADD      DI,AX
LOOP     TRAN_INC_1       ;-----
POP      CX
POP      AX
SUB      AX,NUM_PROC
CMP      AX,0
JLE      TRAN_B
JMP      TRAN_T_LP       ;3rd loop end

TRAN_B:
ADD      BX,4              ;Set next addr for
ADD      DX,4              ;BX, DX
DEC      CX
OR       CX,CX
JZ       TRAN_C
JMP      TRAN_I_LP

TRAN_C:
POP      CX
DEC      CX                ;LOOP will cause
OR       CX,CX             ;out of range error
JZ       TRAN_DONE
JMP      TRAN_POWER      ;3rd loop end

TRAN_DONE:
MOV      AX,N              ;Set input para. for
CALL     Y_TO_X           ;subroutine and call
POP      DI                ;Restore the
POP      SI                ;registers
POP      BP
POP      DX
POP      CX
POP      BX
POP      AX                ;-----
RET
H_TRAN   ENDP

```

Figure A-2 (continue)


```
.DATA

POWER          DW          0h
BASE           DW          2h
TEMP           DW          0h
X              DW          0h
Y              DD          256 DUP(0)
SIN            DD          256 DUP(0)
COS            DD          256 DUP(0)
K              DW          0h
K_ADDRESS      DW          0h
N2             DW          0h
N2_ADDRESS     DW          0h
T              DW          0h
I              DW          0h
ANGLE          DD          256 DUP(0)
LINDEX         DW          256 DUP(0)

NUM_PROC       PUBLIC NUM_PROC
                DW          1h
                PUBLIC N
N               DW          0h

.STACK

END
```

Figure A-2 (continue)

```

.MODEL MEDIUM
.CODE
;-----;
; Program Name : P2 ;
; This subroutine (procedure) is to calculate 2 to the ;
; power of P, where P is the input parameter. ;
; From compiled BASIC program : CALL TWO_TO_P(P,N) ;
; where N = 2 ^ P ;
; Author : Boon Pock Lim ;
; Date : March 10, 1988 ;
;-----;

P2 PUBLIC P2
PROC FAR
PUSH BP ;Save BP
MOV BP,SP ;BP=SP
PUSH AX ;Save the registers
PUSH BX
PUSH CX ;-----
MOV BX,[BP]+8 ;1st Arg.= P
MOV CX,WORD PTR [BX] ;CX=P
MOV AX,1
SHL AX,CL ;AX = 2^P
MOV BX,[BP]+6 ;2nd Arg. = N
MOV WORD PTR [BX],AX
POP CX ;Restore the
POP BX ;registers
POP AX
POP BP ;-----
RET 4
P2 ENDP
END

```

Figure A-3 P2 program listing

```

.8087
.MODEL MEDIUM
.CODE
;-----;
; Program Name : SCALE ;
; This subroutine (procedure) is to scale the result from ;
; the FHT by N. ;
; From compiled BASIC program : CALL SCALE(X(0),N) ;
; Author : Boon Pock Lim ;
; Date : March 10, 1988 ;
;-----;
                EXTRN    SEQ_LD_4:FAR
                EXTRN    SEQ_ST_4:FAR

SCALE          PUBLIC   SCALE
              PROC     FAR
              PUSH     BP                ;Save BP
              MOV      BP,SP            ;BP=SP
              PUSH     AX                ;Save the registers
              PUSH     BX
              PUSH     CX                ;-----
              MOV      BX,[BP]+6        ;2nd Arg.
              MOV      AX,WORD PTR [BX]
              MOV      N,AX
              MOV      BX,[BP]+8        ;1st Arg.
;
;
;SC_D_1:       JMP      SC_D_1            ;Dummy jump
;
;              DB      0DFh,0FDh        ;Parallel mode
;
              FINIT
              FILD     N

SC_RTN_1:     CALL     SEQ_LD_4          ;Load X sequentially
              FDIV    ST,ST(1)         ;Scaling
              CALL    SEQ_ST_4         ;Store X sequentially
              MOV     CX,NUM_PROC       ;Set next address
              CMP     CX,AX
              JBE    SC_RTN_2
              MOV     CX,AX

SC_RTN_2:     ADD     BX,4
              LOOP   SC_RTN_2          ;-----
              SUB     AX,NUM_PROC
              CMP     AX,0
              JG     SC_RTN_1

;
;              JMP     SC_D_2            ;Dummy jump

```

Figure A-4 SCALE program listing

```
;SC_D_2:
;
;
    DB      0DFh,0FEh      ;Scalar mode
    POP     CX              ;Restore the
    POP     BX              ;registers
    POP     AX
    POP     BP              ;-----
    RET     4
SCALE    ENDP

.DATA

    EXTRN  NUM_PROC:WORD
    EXTRN  N:WORD

    END
```

Figure A-4 (continue)

```

.8087
.MODEL MEDIUM
.CODE
;-----;
; Program Name : DFT ;
; This subroutine(procedure) is to calculate the DFT ;
; from the DHT ;
; From compiled BASIC program : ;
; CALL DFT(X(0),N,RE(0),IM(0)) ;
; where RE : real part of DFT ;
; IM : imaginary part of DFT ;
; Author : Boon Pock Lim ;
; Date : March 10, 1988 ;
;-----;

                EXTRN    SEQ_LD_4:FAR
                EXTRN    SEQ_ST_4:FAR

DFT             PUBLIC  DFT
                PROC    FAR
                PUSH    BP                ;Save BP
                MOV     BP,SP            ;BP=SP
                PUSH    AX                ;Save the registers
                PUSH    BX
                PUSH    CX
                PUSH    DX
                PUSH    SI
                PUSH    DI                ;-----
                MOV     SI,[BP]+12        ;1st Arg.
                MOV     BX,[BP]+10        ;2nd Arg.
                MOV     AX,[BX]           ;N
                DEC     AX                ;N-1
                PUSH    AX
                MOV     CX,2              ;(N-1)*4 - offset
                SHL     AX,CL             ;-----
                MOV     DI,AX            ;Last element of X
                ADD     DI,SI            ;-----
                POP     AX
                MOV     BX,[BP]+8         ;3rd Arg. - RE
                MOV     BP,[BP]+6         ;4th Arg. - IM
                FINIT                    ;Reset 8087
                FLDZ                       ;ST(7)=0
                FLD     DWORD PTR [SI]    ;ST(6)=X(0)
                FSTP   DWORD PTR [BX]    ;RE = X(0)
                FSTP   DWORD PTR [BP]    ;IM = 0
;
;
;DFT_D_1:      JMP     DFT_D_1            ;Dummy jump

```

Figure A-5 DFT program listing

```

;           DB           0DFh,0FDh           ;Parallel mode
;           FINIT
;
;           FLD1         ;ST(7)=1
;           FCBS        ;ST(7)=-1 (/2)
;           MOV          DX,4                ;Set next addresses
;           ADD          SI,DX
;           ADD          BX,DX
;           ADD          BP,DX              ;-----
DFT_RTN_1:
;           PUSH        BX                  ;Load X(k)
;           MOV         BX,SI              ;sequentially
;           CALL        SEQ_LD_4
;           MOV         BX,DI              ;Load X(n-k)
;           CALL        SEQ_LD_4          ;sequentially
;           POP         BX                  ;-----
;           FLD         ST(0)
;           FINCSTP
;           FXCH        ST(2)
;           FDECSTP
;           FADD        ST,ST(2)
;           FSCALE
;           CALL        SEQ_ST_4          ;Real #
;           FXCH        ST(2)
;           FSUBRP     ST(1),ST
;           FSCALE
;           PUSH        BX
;           MOV         BX,BP              ;Imaginary #
;           CALL        SEQ_ST_4
;           POP         BX
;           MOV         CX,NUM_PROC        ;Set next address
;           CMP         CX,AX
;           JBE        DFT_RTN_2
;           MOV         CX,AX
DFT_RTN_2:
;           SUB         DI,DX              ;Set next addresses
;           ADD         SI,DX
;           ADD         BX,DX
;           ADD         BP,DX
;           LOOP       DFT_RTN_2          ;-----
;           SUB         AX,NUM_PROC
;           CMP         AX,0
;           JG         DFT_RTN_1
;
;           JMP         DFT_D_2           ;Dummy jump
;DFT_D_2:
;           DB           0DFh,0FEh        ;Scalar mode
;

```

Figure A-5 (continue)

```

                POP     DI           ;Restore the
                POP     SI           ;registers
                POP     DX
                POP     CX
                POP     BX
                POP     AX
                POP     BP           ;-----
                RET     8
DFT            ENDP

.DATA

                EXTRN   NUM_PROC:WORD

.STACK

                END
```

Figure A-5 (continue)

```

.8087
.MODEL MEDIUM
.CODE

                PUBLIC SEQ_LD_2
SEQ_LD_2        PROC FAR
;-----;
; This procedure loads 2-byte data sequentially. ;
; Input parameters : AX = # of data to be processed ;
;                   BX = index of data ;
; Reads : NUM_PROC, data ;
;-----;
                PUSH    AX                ;Save the registers
                PUSH    BX
                PUSH    CX
                PUSH    DX
                PUSH    SI                ;-----
;
;
;S2_LD_D1:      JMP     S2_LD_D1                ;Dummy jump
;
;                DB     0DFh,0FFh          ;Sequential mode
;
;                XOR     DX,DX                ;If AX < # of
;                MOV     CX,NUM_PROC          ;processors then
;                MOV     SI,CX                ;CX=AX else
;                CMP     CX,AX                ;CX=# or processors
;                JBE     S_2_LD_1            ;DX : Flag
;                MOV     CX,AX
;                MOV     DX,1
S_2_LD_1:
                FILD    WORD PTR [BX]        ;Load data
                ADD     BX,2                ;sequentially
                LOOP    S_2_LD_1
                OR      DX,DX
                JZ      S_2_LD_3
                MOV     CX,SI
                SUB     CX,AX
S_2_LD_2:
                FLD1
                LOOP    S_2_LD_2            ;Load dummy 1 to
;                                           ;other processors.
S_2_LD_3:
;
;                JMP     S2_LD_D2            ;Dummy jump
;S2_LD_D2:
;                DB     0DFh,0FEh          ;Scalar mode
;                JMP     S2_LD_D3            ;Dummy jump
;S2_LD_D3:
;                DB     0DFh,0FDh          ;Parallel mode

```

Figure A-6 SEQ program listing


```

                POP     SI             ;Restore the
                POP     DX             ;registers
                POP     CX
                POP     BX
                POP     AX             ;-----
                RET
SEQ_LD_2      ENDP

                PUBLIC  SEQ_ST_2
SEQ_ST_2     PROC   FAR
;-----;
; This procedure stores 2-byte data sequentially. ;
; Input parameters : AX = # of data to be processed ;
;                  BX = index of data ;
; Reads : NUM_PROC ;
; Writes : data ;
;-----;
                PUSH   AX             ;Save the registers
                PUSH   BX
                PUSH   CX             ;-----
;
;
;S2_ST_D1:    JMP     S2_ST_D1           ;Dummy jump
;
;                DB     0DFh,0FFh     ;Sequential mode
;
;                MOV    CX,NUM_PROC   ;If AX<# of
                CMP    CX,AX         ;processors then
                JBE    S_2_ST_1       ;CX = AX else
                MOV    CX,AX         ;CX=# of processors
S_2_ST_1:    FISTP   WORD PTR [BX]   ;Store data
                ADD    BX,2           ;sequentially
                LOOP   S_2_ST_1       ;-----
;
;
;S2_ST_D2:    JMP     S2_ST_D2           ;Dummy jump
;
;                DB     0DFh,0FEh     ;Scalar mode
;                JMP    S2_ST_D3       ;Dummy jump
;S2_ST_D3:    DB     0DFh,0FDh     ;Parallel mode
;
;                POP    CX             ;Restore the
                POP    BX             ;registers
                POP    AX             ;-----
                RET
SEQ_ST_2     ENDP

                PUBLIC  SEQ_LD_4

```

Figure A-6 (continue)

```

SEQ_LD_4          PROC    FAR
;-----;
; This procedure loads 4-byte data sequentially.
; Input parameters : AX = # of data to be processed
;                   BX = index of data
; Reads : NUM_PROC, data
;-----;
                PUSH    AX                ;Save the registers
                PUSH    BX
                PUSH    CX
                PUSH    DX
                PUSH    SI                ;-----
;
;
;S4_LD_D1:      JMP     S4_LD_D1                ;Dummy jump
;
;                DB     0DFh,0FFh            ;Sequential mode
;
;                XOR     DX,DX                ;DX : flag
;                MOV     CX,NUM_PROC          ;If AX < # of
;                MOV     SI,CX                ;processors then
;                CMP     CX,AX                ;CX = AX else
;                JBE     S_4_LD_1            ;CX=# of processors
;                MOV     CX,AX
;                MOV     DX,1                ;-----
S_4_LD_1:
                FLD     DWORD PTR [BX]      ;Load data
                ADD     BX,4                ;sequentially
                LOOP    S_4_LD_1            ;-----
                OR     DX,DX
                JZ     S_4_LD_3
                MOV     CX,SI
                SUB     CX,AX
S_4_LD_2:
                FLD1    S_4_LD_2            ;Load dummy 1 to
                LOOP    S_4_LD_2            ;other processors
S_4_LD_3:
;
;
;                JMP     S4_LD_D2                ;Dummy jump
;S4_LD_D2:
;                DB     0DFh,0FEh            ;Scalar mode
;                JMP     S4_LD_D3                ;Dummy jump
;S4_LD_D3:
;                DB     0DFh,0FDh            ;Parallel mode
;
;                POP     SI                ;Restore the
;                POP     DX                ;registers
;                POP     CX
;                POP     BX

```

Figure A-6 (continue)

```

                POP     AX           ;-----
                RET
SEQ_LD_4        ENDP

                PUBLIC  SEQ_ST_4
SEQ_ST_4        PROC    FAR
;-----;
; This procedure stores 4-byte data sequentially. ;
; Input parameters : AX = # of data to be processed ;
;                   BX = index of data             ;
; Reads : NUM_PROC ;
; Writes : data ;
;-----;
                PUSH   AX           ;Save the registers
                PUSH   BX
                PUSH   CX           ;-----
;
;
;S4_ST_D1:      JMP     S4_ST_D1      ;Dummy jump
;
;               DB     0DFh,0FFh    ;Sequential mode
;
;               MOV    CX,NUM_PROC  ;If AX < # of
;               CMP    CX,AX        ;processors then
;               JBE    S_4_ST_1     ;CX=AX else # of
;               MOV    CX,AX        ;CX=# of processors
S_4_ST_1:      FSTP   DWORD PTR [BX] ;Store data
                ADD    BX,4         ;sequentially
                LOOP   S_4_ST_1     ;-----
;
;               JMP    S4_ST_D2     ;Dummy jump
;S4_ST_D2:      DB     0DFh,0FEh    ;Scalar mode
;               JMP    S4_ST_D3     ;Dummy jump
;S4_ST_D3:      DB     0DFh,0FDh    ;Parallel mode
;
;               POP    CX           ;Restore the
;               POP    BX           ;registers
;               POP    AX           ;-----
                RET
SEQ_ST_4        ENDP

.DATA
                EXTRN  NUM_PROC:WORD

                END

```

Figure A-6 (continue)



