# Software for Regional Studies: Adjustment Procedures for Integrated Balances 

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SOFTWARE FOR REGIONAL STUDIES:
ADJUSTMENT PROCEDURES FOR
INTEGRATED BALANCES
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## PREFACE

The Tuscany case study is the last in the sequence of four regional case studies carried out in the Regional and Urban Development Group (RUD) at IIASA. The Tuscany study has developed as a joint effort between the RUD group and researchers from the Regional Institute for Economic Planning of Tuscany (IRPET).

One of the several efforts in the Tuscany study has been the construction of a regional accounting balance table connecting so called material and financial balances. The character of such integrated economic balance tables is discussed by Issaev and Umnov in WP-82-118.

The current paper presents an algorithm which can be used as a means to construct and adjust integrated balances. It also includes a detailed software description. In summary, the paper provides a documentation of work related to the Tuscany study.

> BÖrje Johansson Acting Leader Regional and Urban Development Group

May, 1983

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# Software for Regional Studies: Adjustment. Procedures for Integrated Balances 

Miloslav Lenko, Klim Kim<br>IIASA

## Introduction

This report is devoted to the problem of analysis of sensitivity of balance tables, which arises by mathematical modeling, for example, economic-financial flow within a region, a trade market, a transport system and so on. The report contains both the statement of the problem and the algorithms of solution, including a well detailed software description.

## Statement of the problem

The object of a formal analysis in this problem is a square matrix ( NxN ), elements of which are real positive numbers. An example of this table is shown in Figure 1. Everywhere further we will denote the element of the matrix placed in the i-th row and j-th column as $A_{i j}$.

The matrix A is called balanced, if the following relations

$$
\sum_{j} A_{i j}=\sum_{k} A_{k i}
$$

take place for all $\mathrm{i}=1, \ldots, \mathrm{~N}$.
Generally speaking, considered matrixes need not be balanced, but a procedure to make them balanced is a main aim of the analysis.

Let us introduce vectors $c=\left(c_{1}, c_{2} \cdots c_{N}\right)$ and $r=\left(r_{1}, r_{2}, \cdots, r_{N}\right)$ defined as row and column sums

$$
\begin{align*}
& c_{i}=\sum_{j} A_{i j}  \tag{1}\\
& r_{i}=\sum_{k} A_{k i} \tag{1}
\end{align*}
$$

for all $i=1, \ldots, N$ and a vector $B=\left(b_{1}, b_{2}, \cdots b_{N}\right)$, where

$$
\begin{equation*}
b_{i}=r_{i}-c_{i} \tag{2}
\end{equation*}
$$

It is obvious, that for a balanced matrix $c_{i}=r_{i}$ or $b_{i}=0$ are valid.
Now we are able to state the following problem:
Problem 1: if the given matrix A is not balanced, find such values $X_{i j}$, that will balance the new matrix with elements $A_{i j}=A_{i j}+X_{i j}$.

It is necessary to note that the problem does not have a unique solution. Moreover, this problem is not of practical value, because of the absence of any relations between new and old matrixes. For example, the simplest trivial solution of problem 1 is the following. Let elements of the new matrix have sufficiently great numbers $H$, then $X_{i j}=H-A_{i j}$ are a solution of the problem.

The statement of the problem should be defined more exactly.
The most natural additional condition is a result of a natural trend to introduce the smallest possible changes in the initial matrix.

There are two versions of the concept of the smallest change. Either the number of elements to be changed is kept to a minimum, or any changes in the relative values of the elements are kept as small as possible. Currently, only the second version is used in practice, although the first one may also be interesting from a theoretical point of view.

A more general description for the nonlinear case is given in (B.Issaev, A.Umnov, 1982).

Further we will consider only the following concrete formalization of the second concept, called below uniform balancing. Its mathematical statement is

Minimized $Y$ with respect all $X_{i j}$.
Subject to

$$
\begin{gathered}
a b s\left(X_{i j}\right) \leq Y * X_{i j} \\
\text { for all } \mathrm{i}, \mathrm{j}=1,2, \ldots, \mathrm{~N} \text { and } \\
\sum_{j}\left(A_{i j}+X_{i j}\right)-\sum_{k}\left(A_{k i}+X_{k i}\right)=0 \\
\text { for all } \mathrm{i}=1,2, \ldots, \mathrm{~N} .
\end{gathered}
$$

Using notatations (1) and (2), finally the main problem can be formulated as follows

$$
\begin{equation*}
\text { Minimized } Y \text { with respect } X_{i j} \tag{3}
\end{equation*}
$$

Subject to

$$
\begin{gather*}
\sum_{j} X_{i j}-\sum_{k} X_{k i}=b_{i}  \tag{4}\\
X_{i j} \geq-A_{i j} * Y  \tag{5}\\
X_{i j} \leq A_{i j} * Y \tag{6}
\end{gather*}
$$

It is easy to see, that the problem (3) - (6) is a linear programming problem (LP), which can be solved by a standard software tool. On the other hand, this LP-problem has a special structure. Therefore, it can also be solved by a more effective special algorithm.

In practice the choice of the software tool should be made by compromising. We have to provide possibilities of including some extra relations in the statement of the main problem. A general tool is required. But use of the model with many repetitions (for example, an interactive regime) is also possible. In these cases it is desirable to have a high speed software.

There is no doubt about what kind of general software tool may be used for the problem of the uniform balance. It is the standard simplex method.

A special structure of the problem (namely, a modified problem about the minimal flow in a network), may be a basis for developing both new, and more effective algorithm and software units.

Before begining discussion of the network interpretation of the problem to be solved, let us formulate a modified problem of the uniform balance.

We will suppose that the user is interested in solutions of the problem where some elements in the balanced matrix are fixed with a priori given values. Hence, we have

Problem 2: the problem of the uniform balance with fired elements.
It is not necessary to give the mathematical statement of this new problem, because the statement will repeat the statement (3) - (6) of problem 1 with additional constraints $X_{i j}=0$, subject to the initial matrix having $A_{i j}=\mathrm{f}$, where f is the given value.

## Network interpretation of the problem

The suggested approach to problem 2 is a generalization of an algorithm for solving flow problems, which are formulated for a network.

Now we give a network interpretation for problem 2.
We associate each pair of rows and columns possessing the same index with a node of a network with this same index. Each nonzero element of the matrix $A_{i j}$ we associate with two oriented arcs ( $\mathrm{i}, \mathrm{j}$ ) and ( $\mathrm{j}, \mathrm{i}$ ). The nonnegative value $X_{i j}$ of the flow along the arc ( $\mathrm{i}, \mathrm{j}$ ) we will consider as increasing the value of $A_{i j}$. and the nonnegative value $X_{i j}$ of the flow along the arc ( $\mathrm{j}, \mathrm{i}$ ) - as decreasing the value of $A_{i j}$.

It will be convenient to enumerate all nonzero elements of the matrix $A$. The numbers will be $k=1,2, \ldots, M$. In the same manner we will enumerate all arcs. The increasing arcs will have indexes $k=1,2, \ldots, M$ and the decreasing arcs numbers $k+M$. The index $k$ is the same for the pair of arcs belonging to the same element of the matrix $A$.

Therefore, for $M$ elements of $A$ we consider $2 M$ oriented arcs, the first $M$ of which are increasing and the last M are decreasing. Notice, that we will not take into account the diagonal elements, because they do not change balance constraints at all.

Now let us consider a small example. Let us have the following matrix

Table 1

| $\dot{j}$ | 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: | :---: |
| 1 |  | 50. | 0. | 0. |
| 2 | 10. |  | 40. | 0. |
| 3 | 0. | 0. |  | 30. |
| 4 | 20. | 0. | 0. | - |

Denote as K, IB, IE, A the arrays of indexes of nonzero elements of the matrix A , the beginning and ending of the arcs, and the values of the elements. Then the total information about the network corresponding to the matrix can be collected in table 2 .

Table 2

| K | IB | IE | A |
| :---: | :--- | :--- | :--- |
| 1 | 1 | 2 | 50. |
| 2 | 2 | 1 | 10. |
| 3 | 2 | 3 | 40. |
| 4 | 3 | 4 | 30. |
| 5 | 4 | 1 | 20. |
| 6 | 2 | 1 | 50. |
| 7 | 1 | 2 | 10. |
| $B$ | 3 | 2 | 40. |
| 9 | 4 | 3 | 30. |
| 10 | 1 | 4 | 20. |

In figure 1 the picture of the network is shown.
Values $b_{i}$ can be interpreted as values of a uniform flow at i-th node. If $b_{i}>0$, then the node is a source. If $b_{i}<0$, then the node is a sink.

In the given example we have $b_{1}=-20, b_{2}=0 ., b_{3}=10$. and $b_{9}=-10$., Hence, nodes 3 and 4 are sources and the node 1 is a sink.

The given rules of data preparing are not of great importance here, of course. But they are sufficiently convenient to explain algorithms and to interprete results. From the formal viewpoint it is sufficient to consider the problem of the uniform balance as the following network problem

Find the minimal value of $Y$, for which the feasible flow saturates all sinks of the network with channel capacities equal to $Y{ }^{*} A_{k}$.
A particular result follows from the interpretation given above. If the network has only a single pair source-sink, then the problem of the uniform balance can be reduced to the classical problem of the maximal flow.

Let us have the solution of the problem where the maximal flow equals $f$, subject to $Y=1$. Then it is possible to equate this value of the flow to the given b. by mean proportional changes of all flows of the arcs $X_{i j}$ and channel capacities $A_{i j}$.

It gives us the solution of the special case of the uniform balance problem with value $Y=b_{i} / \mathbf{f}$.

This result is very useful as we have a balanced matrix with a single fixed element.

One more imortant note.
In the above network interpetation of the uniform balance problem, some additional constraints can be introduced on the values of changes of the matrix elements. Namely, it is possible to forbid decreasing or increasing for some of the elements. To do this, we may simply adjust the corresponding channel capacities to the zero level.

## The network algorithm

When the problem of the balance generates a network problem with several sources and sinks, the direct usage of the above simple algorithm is not possible. Here a transportation problem arises. The main specific feature of the problem is the fact that the channel capacities are linear functions of $Y$. It is necessary to find the minimal value of $Y$, when the problem is still feasible.

A very short description of the approach applied here is given below.
The whole process consists of two steps. At the first step a balancing flow is built without constraints on the channel capacities of the arcs. We may say that the problem is solved by a very large $Y$.

At the second step an optimal ( with respect to $Y$ ) flow has been found. During this step the value of $Y$ is iteratively changed. These changes are connected with a change of the structure of the optimal basis of the transportation problem.

The whole process will be out, if there is no possibility to decrease $Y$ by changing flows in ares.

Without going into details of proofs, we will formulate some geometrical properties of these basis solutions of the transportation problem.

If a value of $Y$ and the corresponding distribution of the flows within the network are a basis solution of the LP-problem, then the set of flows in arcs can be described as follows.

The set of the arcs is divided into three groups
S - nonbasis arcs,
D - basis nonsaturated arcs,
G - basis saturated arcs.
The flows $X_{k}$ will have the following values

$$
\begin{array}{ll}
X_{k}=0 . & \text { if } \mathrm{k} \in \mathrm{~S} \\
0 \leq X_{k} \leq Y^{*} A_{k}, & \text { if } \mathrm{k} \in \mathrm{D} \\
X_{k}=Y^{*} A_{k}, & \text { if } \mathrm{k} \in \mathrm{G}
\end{array}
$$

The most important property is that the subset $D$ contains exactly $N-2$ arcs. The subset consists also of two tree-like components of connectedness $D_{1}$ and $D_{2}$. There exists always an arc from $G$. which connects $D_{1}$ and $D_{2}$. The set, which connects $D_{1}$ and $D_{2}$, we will call pseudo-cut and denote as R.

In figure 2 a possible structure of a current basis is shown. S-arcs are shown by dot lines, D-arcs are shown by thin pointers and G-arcs are shown by thick pointers. R-arcs are shaded.

The above properties permit the formulation of optimal conditions for a basis solution.

A basis solution is optimal, if all nonsaturated arcs from $R$ are oriented in the same direction (between $D_{1}$ and $D_{2}$ ) and all saturated arcs from $R$ are oriented in the opposite direction.

In other words, a basis is optimal if the total channel capacity permitting the increase of the flow from $D_{1}$ to $D_{2}$ is zero. It means that the set $R$ is a cut.

Therefore, to decrease $Y$ (the norm of the channel capacity) it is necessary to decrease flows in all arcs from $G$, including from $R$. But to keep the balance conditions valid, all the changes must form a nonzero circulation. If such a circulation cannot be built, then the current basis solution is optimal. Figure 2 shows just this situation.

If $R$ is not a cut, then the algorithm defines necessary circulations, in which arcs from $G$ and at least one nonbasis arc are involved. The decreasing $Y$ leads to increasing this circulation. It happens as long as the change of $Y$ reaches its limit value. Then the rebuilding of basis becomes necessary.

The work of the algorithm may be illustrated for the network given in Figure 1. The sequence of iterations is shown in Figure 3.

The first balanced flow (with $\mathrm{Y}=+\infty$ ) goes from node 3 to node 1 through node 4. The minimal value of $Y$, granted the feasibility of the flow, is $\max \left(X_{k} / A_{k}\right)=X_{5} / A_{5}=20 . / 20=1$.
There is a nonsaturated 5 -th arc here, which divides the network into two subsets $D_{1}=(3,4)$ and $D_{2}=(2,1)$.

Nonbasis arc 8 permits the formulation of a circulation to decrease the flow through arc 5 . Increasing this circulation leads to the saturation of arc 2, which occured at $Y=0.66$. And so on.

Notice, that for all nonsaturated arcs the inequality $X_{k} \leq Y^{*} A_{k}$ takes place. In the optimal basis solution all arcs between $D_{1}$ and $D_{2}$ consist of a cut, and this does not permit the decrease of the value of $Y$ further.

The computer codes of the algorithm were elaborated by K.Kim in cooperation with a laboratory of mathematical programming of The Central Economic Mathematical Institute of the Academy of Science of the USSR and Regional Development Group of IIASA.

Reference
B. Issaev, A. Umnov INTEGRATED ECONOMIC BALANCE: ADJUSTMENT PROCEDURES, IIASA, WP-82-118, 1982

## User's manual for balance programs

This manual describes how to use BALK, BA山M programs for The Integrated Economic Balance model of Tuscany.

The manual describes how to use these programs on a VAX 11/780 computer under the UNDX operating system. There are only several changes for using these programs on an IBM computer.

The manual contains several chapters:

1. Definition of task and some items
2. Usage, compilation and linkage
3. I/O files, their formats
4. Error messages and other recommendations

## 1. Definition of task and some items

Let's have a quadratic matrix, which contains $n$ rows and $n$ columns and which can be sparse. Let's call this matrix balance table.

For reasons of flexibility, divide this table equally rowwise and columnwise into several parts, let's call them groups, where each group can contain 1 to m sequential rows (and columns too).

The whole table consists now of several groups. Each group has its own size (nro of rows/columns). For the marking of each group we can use 1 (nonnumeric) character, so the table will contain e.g. groups A,B,C etc.

Each row/column of our table - let's call it account should be marked by absolute number of this row/column or by the name of group and the relative number in it, e.g. BO2, CO1 etc. Let's call row income account and column outcome account. Similarly, each element of this table, let's call it flow, can be marked with the help of two numbers - number of the row and number of the column; or alternatively with the help of relative numbers in groups, e.g. B02C04, A01E04 etc. The places where groups are crossing these submatrices (which should not be quadratic) we call blocks.

Our first task is as follows:
Let's have one balance table, which is created from blocks. All these data create data base. For computing we use only some blocks. Let's call them included blocks according to our scheme. All flows from these included blocks must have the following property: The sum of data for each input account must be equal to the sum of data for the corresponding output account. Let's call this property balancing.

Generally, the input balance table need not be balanced. Our task is to get this balance in such a way, that the maximum relative changes for all flows must be minimized.

All this structure (division of the table into groups and so into blocks) was done from practical definition of the task. Each group and block has its own economic interpretation for this model. In order to make this construction more flexible, it is possible to change the size of groups with minor changes in input data etc.

The second task is to balance this table if we do some exogenous changes. These changes may be of the following type:
a/ let the resulting flow have exactly the initial defined value
b/ let the resulting flow increase only
$\mathrm{c} /$ let the resulting flow decrease only (but still be $>=0$.)
At the solution of this task (and also the first task) we get new values of flows: and so a new table - let's call it result table. The differences between new and old values create a so called correction table. Of course, should the situation occur that the solution is infeasible in both cases, then the table cannot be balanced.

The third task is like the first two, but we have another additional property of data in the table. Some of this data must correspond to some other so called equations of proportionality.

For the solution of all these three types of problem a software package has been developed which contains, as a main part, 2 different programs for optimal balancing (because this task can be formulated as a linear programming problem). The first program is a network algorithm for optimization of flows in the network. The program was developed at the laboratory of applied software at CEMI, Moscow. The second program is the MINOS linear programming package from Stanford University.

What must the user define and prepare, and what results will he get?

1. The user must prepare all input data files in appropriate form. Here the user must divide the table into groups, giving names and sizes of each group. The user can also give the name for each group ( 40 chars max) which will be printed on output files. Then the user must prepare the whole data base - all blocks which contain nonzero flows. Then the user must prepare a list of included blocks and file with names for each account (also 40 chars long).
2. A so-called specification file will contain a description of our task which means names of all input and output files, some important parameters and type of task (balance only, balance with restriction, proportionality, compressed form of table etc.). If the user wants to fix some variables, he must prepare a so-called restriction file too.
3. The user gets results stored on output files, according to the options he has chosen. All output files are optional. There are e.g. input balance table file; table of correction file; results table file. All results can be stored on data file in the same format as data base file and used for further computing.

The user can also display data in so called graphical form, where each result flow is displayed as 1 character; for instance, . if it did not change; + or if a maximum change was made, and so on (see chapter 3). Looking at this file the user can get an overview of how corrections data are distributed in the resulting table.

The next possibility is to get the file with some chosen accounts which are interesting for the user (not a complete table); e.g. input, correction and result income and outcome accounts. Finally, the user can define and get a file where input blocks will be stored in a form of matrix with subtotals for rows and columns.
4. The user has the possibility of choosing some of the following functions:

- balancing input table
- balancing input table with restrictions
- as before but with proportionality
- work with compressed version of table (first two tasks only), where each
block is interpreted as 1 flow which is the sum of all flows from this block.


## 2. Usage of programs

All programs are written in FORTRAN-77 language for VAX 11/780 computer under the UNIX operating system. The programs were written in such a way that their portability is very good, as generally, they are written in FORTRAN IV, when only some statements must be changed. It is necessary to use virtual machine for their computing because of large volumes of arrays of variables (which depends on the type of work and amount of data) and the MINOS programming package.

Generally the user has 2 different program packages for solving his problems in IIASA's computer. There is a package using network algorithm (BALK) and a package using MINOS for solution (BALM).

User must translate all the programs with the f 77 compiler and link them to executable form. There are several FORTRAN subroutines in one source file.

For package BALK, user must translate and link following files: balnew.f bdata.f comput.f flow.f maxfld.f subr1.f vstup.f vystup.f There is file link.balk where commands can be found on how to link programs.

For package BALM, user must translate the following files: balprp.f bdata.f subri.f vstup.f vystup.f gmpsp.f subp.f and link them together with files where the MINOS package is. File on how to link these is link.balm.

Submitting the job for computing:
Before submitting, user must prepare all input files (some of them are optional), according to type of job. All files besides spec file are opened internally in program, so user must submit the job simply by typing

## BALK 1=spec or BALM 1=spec

Choice of appropriate algorithm:
network algorithm enables solution of tasks of following types only:

- balancing non-balanced table
- balancing with restriction of type equality (fixing result flows)

This algorithm is very fast (approximately 30 times faster than MINOS). It serves for quick preliminary balancing of job.

MINOS enables solving of all kinds of tasks (fixing of all types of restrictions, proportionality), by generating an mps-file but solving is relatively slow.

General description of program:
Program consists of 3 parts:

- inputing of data
- optimization
- producing results.

First, program reads specification file. According to defined task program does all necessary functions, reads groups and included-blocks file. Then program reads database file. If the user chooses option to compress table, data are compressed during reading input block and each block is considered as consisting of one flow. At the moment of reading data, the output file with output blocks is also prepared (if user wishes to have such a file). Reading the names of
accounts follows (if user did not choose compress options - in that case names of accounts are equal to names of groups). If the user wants to have restrictions, restriction file is read, the same as for input account file and proportional file.

Then two programs can be used:
a/ network program - Description of this algorithm can be found in the first part of this paper. Program does the following: at first it checks whether the table is balanced. Then, according to user's wishes, it balances input matrix and/or balances it with restrictions.
b/ MINOS program - For using MINOS, a so called mps file must be ready before calling. This file is created automatically by a program genmps. If the user wants to use proportionality, the mps file is created slightly differently, First, a file mpsfile is created, where all data for nonproportional variables are stored. After that a file mpsprop is created where data from proportional equations only are stored. Finally, the program merge merges these two files producing a file mpsall which will be standard input mps file for MINOS. After calling MINOS results are stored in working array $z$ and a program solut transfers these results from this array into the appropriate positions in array res. Output part prepares all other output files (according to user demands). First the balance file, then the correction file, result file, output account file, output database file and finally the output graph file are produced.

## 3. 1/0 files

There is a difference in managing the files in IBM and VAX computers. For VAX all files are opened internally within the programs. The user gives their numbers in specification file and their numbers are allocated at generation time. For IBM computers, files are opened externally and the numbers of files are given as they were made at the installation. The names of files in the specification file have here only information meaning. They are not used to open files.

It is possible to connect several files into one file (input as well as output files), which must be done during the installation, when the same numbers for several files are assigned. User must connect these files together and use blank line as EOF indicator of each file. Concerning input files, the user must know the sequence in which input files are read. That is: spec, groups, include, database, names, restriction(optional), accounts (optional) and proport(optional).

All necessary files could be divided into three parts:
a/ Input files
b/ Output files
c/ Working files

### 3.1. Input files:

1. specification file
2. database file
3. groups file
4. included blocks file
5. names file
6. restriction file [optional]
7. accounts file [optional]
8. proportional file [optional]
9. specfile for MINOS [optional]

### 3.2. Output files

All output files except fort. 9 are optional

1. standard output file fort. 9 and tty output
2. balance table file
3. corrections table file
4. results table file
5. graph file
6. output accounts file
7. output blocks file
8. output database file
9. outfile for MINOS

### 3.3. Working files

These files are used when MINOS algorithm is chosen only

## 1. mpsfile

2. mpsprop [only if proportionality]
3. mpsall [only if proportionality]
4. scratch file for MINOS fort. 8
5. optional basis files for MINOS
a1/Specification file:
The following is an example of such a file:

| name of task | Tuscany II. |
| :---: | :---: |
| type of algorithm | : MINOS |
| database file | : database |
| groups file | : group |
| names file | : names |
| included blocks file: | : incl |
| restriction file | : rest751 |
| out bal. table file | : outbal |
| out correction file | : outcor |
| out result file | : outres |
| out graph file | : outgrp |
| inp \& out account f. | .: acc outacc |
| out database file | : no |
| proportional file | : propor |
| out blocks file | no |
| compress | no |
| balance of input | : no |
| restr \& bal. of inp. | : yes |
| tolerance eps | $1.0000 \mathrm{e}-5$ |
| tolerance seps | $1.0000 \mathrm{e}-3$ |
| tolerance epstol | $1.0000 \mathrm{e}-3$ |

The structure and position of elements are fixed. Each line starts with an explanatory text which is 22 chars long after which follow the parameter(s). All names of files are 12 characters long. Name of task (1.line) has a max of 40 chars, this occurs on output files. Type of algorithm should be keyword Kim or minos (2.line). If some file is optional and is not chosen, just type no. Two names of account files (input and output) must be given on line nro 12, name of output account file starts in column 35. Parameters compress (16.line), balance of input (17.line) and restr. \& balance of input (17.line) must be keywords no or yes. There is a check to some extent e.g. if user does not choose balancing and wants to have output result or output correction file then error message - wrong spec file-occurs. Tolerances eps and seps serve as criterion for accuracy when balancing with network algorithm, epstol serves as criterion if the variable is on upper or lower bound at preparing graph file.
formats are: for lines 1-18 (22x,10a4)
for lines 19-21 (22x, e12.5)
a2/ database file
All data are stored as a sequence of data belonging to blocks. Database can also contain blocks which will not be included in scheme for computing (will not be in the list of incl. blocks). Blocks contain all nonzero flows. It is not possible to have 2 blocks with the same name in the database. This error is not checked by the program and if left it may cause strange results.

Block contains the head of block where name of block is 2 characters. If this name is not in list of included blocks, the data of this block are skipped. After head all nonzero flows follow for this block in a mode:
row index, column index, value of flow
Indexes are compared with values got from the table - dimensions of groups. All
values must be greater than 0 . Values need not be sorted.
The tail contains the sequence: $0 \quad 0 \quad 0$. Block can be empty.
Formats are for the head: (6x,2a1)
for flows:(2i5,e12.5)
An example follows:
block CD
1,2.14.
2,3,3.
4,1,11.
1.1.7.
$0,0,0$.
block CS
4,4,1.
$0,0,0$.

## a3/group file

File contains identifier of the group (1 character), dimension of the group and title of the group ( 40 chars).
Format is ( $\mathrm{a} 1, \mathrm{i} 3,1 \mathrm{x}, 10 \mathrm{a} 4$ )
Sequence of groups defines division of table.

## a4/ included blocks file

This file contains identifiers of all blocks ( 2 characters), which must be taken from database and included in the table. It is possible to put character $\mathbf{Y}$ at the end of the row if you want to print this block on the output file of blocks (in matrix mode with subtotals). Blocks could be in the file in any order.
Format is (2a1,1x,a1)
a5/ names files
This file contains names (titles) for each account ( 40 chars max). There must be the account name at first, e.g. name of group ( 1 char), than index in this group, and then the title for account ( 40 chars.) on each line. Accounts may not be sorted. If user chooses compress option this file is not read and titles for accounts are the same as for groups.
Format is ( $\mathrm{a} 1, \mathrm{i} 2,1 \mathrm{x}, 10 \mathrm{a} 4$ )
a6/ restriction file
The file has similar structure to the database file. The data are grouped together according to blocks, which they belong to. For each restricted datum it is necessary to give type of restriction, which can be one of three possibilities:

- fixing datum for exact value, character is $=$
- character > means, that flow can increase only
- character < means, that flow can decrease only

It is not necessary to give any value at the end of the line if modes < or > were chosen. Then at each line user must define:
index of row, index of column, restr, type( 1 ch ), restr, value (if type is ' $=$ ')
Format is ( $\mathrm{i} 3, \mathrm{i} 3,1 \mathrm{x}, \mathrm{a} 1,1 \mathrm{x}, \mathrm{e} 12.5$ )
Example:

```
block AB
1,4,=15.6
2,4,>
0,0,0.
block UA
    31 2 = 1234.56
    212<
        0 0 0.
```


## a7/ accounts file

Here user defines which accounts (income and outcome) he wants to write in output accounts file. For each account it is possible to have its value for input flows, correction flows and result fows, what user defines by giving characters $\mathbf{I}$, $\mathbf{C}, \mathbf{R}$ to the fields for type of acc. The user defines on each line number of account (by means of identifier of the group and index in this group) and the type of account. User can, of course, choose $\mathbf{C}$ and $R$ type only if he chooses an option for balancing input matrix (or bal with restr). If the user wants to write all accounts from one group, he uses char 0 (zero) instead of index within group. If the user wants to print sum of all accounts for the whole table, he uses character * as name of group.
Format is (a1,i2,1x,3a1)
Example:

* O ICR (prints sums of all accounts for input, corrections and output)

A 2 I
A 3 I
A 4 I (from group A user wants to print input accounts for rows 2,3.4 only) C O CR (print all accounts for group C, and so correction and result)
a8/ proportional file
This file has fixed structure.
Example:

$$
\begin{array}{llllll}
\text { A } 0131 \mathrm{~K} & 01 & 09 \\
\text { A } 01 & 31 & \mathrm{~L} & 01 & 08 \\
\text { A } 01 & 31 & \text { A } 01 & 31 \\
\text { C } 01 & 01 & \text { A } & 01 & 31 \\
\text { C } & 02 & 02 & \text { A } & 01 & 31 \\
\text { P } & 01 & 01 & \text { A } & 01 & 31
\end{array}
$$

The file has 6 lines, for each proportionality one line. These lines show which variables are included in the relations of proportionality. If user does not choose a proportionality, he must put name $\mathbf{X}$ for the name of group. There are 6 various types of proportionality, first three for columnwise proportionality are for tables $A K$, AL and AA respectively. Columnwise proportionality means, that if in input balance table value of one flow was $1 / 16$. of the sum of this column of this block, the resulting flow must also be $1 / 16$.-th of new sum of this column. The next three types of proportionality are concerning table CA (first and second row) and table PA (first row). The values of these flows are defined by the set of equations. Numbers in prop files give lower and upper limits for proportional variables, these values are checked with dimension of groups.

If user has chosen the MINOS package, he must prepare a so called specification file (see MINOS reference manual). Its name is always specfile. Here user can change number of rows, columns and nro of elements according to his task. User can also define usage of basis files. Number of rows, cols and elements is approximately:

```
if our table has N accounts and M nonzero flows, then
nro of rows =2 * M + N+1
nro of cols = N +1
nro of elements=6*M
```


## b/ Output files

Output files are mostly created after solution. All files (except fort.9) are optional and their names are given in spec file. The fort. 9 and the standard tty output contain information about solution, e.g. number of accounts, flows, sum of flows, debalance, obj. function.
The source balance table file contains the whole table with all data from included blocks. Rows and columns are signed absolutely and relatively within groups. Only nonzero flows are displayed. Blocks are separated by lines. It is possible to glue together the table into one large table.
The correction and result table files (their format is equal to format of balance table) contain corrections and result flows respectively.
The graph table file contains the picture of the table, where each flow is presented by one of the following characters:

```
. - there was no change or flow was equal to zero
e - element was fixed exogenously
+ - change of this flow is on the upper bound
-- change of this flow is on the lower bound
o-change of this element is somewhere between lower and upper bounds.
```

It is possible to get the graph only after balancing.
If nro of accounts is $>128$, graph is too big and it cannot be written.
The output accounts file contains accounts according to specifications from the input account file.
The output blocks file has all blocks written separately, with sums for all rows/columns.
The output database file has the same format as input database file. The values stored there are resulting flows. In this way this file can be used as input database file for next iteration.
The outfile is standard output file when MINOS is used. It is used for information purposes only.

## c/ Working files

These files are created at the time of computing. All files have given names and all are used if MINOS algorithm only was chosen.
The mpsfile is mps file for MINOS
The mpsprop is also mps file, which is slightly different. It has only variables (rows, cols, ...), which are used in proportional equations only. This file must be merged with file mpsfile. It produces the file mpsall which will be the input mps file for MINOS if proportionality was chosen.
The mpsall is produced as merging of files mpsfile and mpsprop
The fort.B is a scratch file for MINOS.

## 4. Configuration, error messages, remares

a/ Program has defined some maximal sizes for some arrays:
maximal number of groups - mxing - 30
maximal number of accounts - mxacc - 200
maximal number of incl. blocks - mxinc - 150
maximal number of flows - mxvar - 2000
maximal number of restrictions - mxres - 1000
maximal number of graph size --128
In the case of a larger job it is necessary to change the dimensions of these arrays and also the value of the corresponding mx... variables (see programmers' guide), to translate programs and to link them together
b/ Names of files are 12 characters long. Names of groups and accounts have a max of 40 chars. All character variables are declared as integer* 4 (it means 4 characters per 1 variable). Hence for storing names of files it is array of 3 variables. For storing names of groups it is array of 10 variables. Some variables are initiated by data statements.
c/ For usage of MINOS single large array (of type real*B) with name $z$ is necessary. The default length of this array is 50000 (variable nwcore). If this is too small, error diagnostic 28, 40 occurs.
d/ Real arrays and variables are declared as type real* 8 for accuracy at inputing and computing.

## e / Tolerances are of three types:

eps, seps ( $1, e-3,1, e-4$ ) are used by network algorithm. They are used to check whether the difference income-outcome is small enough. User can change them in the spec file,
toleps (1.e-3) is used at computing variables for graph file to check if the flow is close to the upper or to the lower bound.
f/ objective function gives maximal relative difference for any flow in the table. This value can be read from the file fort. 9 or the graph file.
g/ error messages are displayed on tty output as 2 numbers. First number is code of error. Second number depends on the first one. Generally when reading input files it is the number of the line where the error occurs. At code 28 it is a program error (see Programmer's guide or MINOS manual). A list of error codes is at the end of this chapter. After you get an error message, you must either correct some data file or make changes in source programs. Look in Programmers guide to see what action may be taken.
h / The following restrictions must be taken into account:

- when preparing database file all data belonging to one block must be stored in this one block
- for proportionality purposes:
- group A must be the first group in list of groups
- max size of group A is 31
- max size of group $K$ is 10
- max size of group L is 9
- max size of group $U$ is 1
- all names for groups $A, K, L, C, P, S, R$ must be the upper case letters
i/ Restrictions for compressed form are like those for the compressed flows (not source flows).
j/ How to fix element which is not in database file (which is zero there)? You may add this element into database with a very small value.


## BUGS:

1. All output files are generated on finishing the optimization.
2. The user should beware of not having data of one block in different parts of database file.
3. There are some restrictions for the proportionality option as well.

## ERROR CODES

1 [in
2 [i
3 i
[inpgrp] - wrong data on group file (size of group)
[inpgrp] - group name is already used
[inpinc] - such a group does not exist
[inpinc] - block name is already used
3 [inpdat] - parameters in base file are not correct(out of range or $v=0$.)
[inpinc] - number of included blocks $>$ mxinc
[inpgrp] - number of groups $=0$
[inpdat] - number of variables $>$ mxvar
[inpnam] - wrong data on name file
[inpnam] - group not found
[inpinc] - number of included blocks $=0$
[vstup ] - number of accounts $>$ mxacc
[inpdat] - last block was not closed
[inpdat] - number of variables $=0$
[inpres] - number of restrictions $>$ mxres
[inpres] - bad data on restriction file
[inpres] - number of restrictions $=0$
[absind\} - element does not match
[inpres] - restriction has no variable
pripr ]- error at computing graph elements
[maluj] - more than 128 accounts for preparing graph
[comput] - number of variables $>$ mmax
comput]-fatal error bad program
comput] - $n$ or nel are too big, change dimensions
comput] - error in flow
comput] - number of fixed elements $>$ nmax
[inpspc] - error in spec file
[inpinf] - error in data in input for acc file
[otvor] - open error
outgrp] - obj < 0.00000001
isrest] - something strange with data in restriction file \& source data
inpprp] - proport tables too small
inpprp] - description outside limits
[inpprp] - such a group was not found
inpprp]-such a block is not included
inpinf] - groups outside limits for input acc file
inpinf] - not such a group for input acc file
inpprp] - error on proport. file
[inpspc] - algorithm is not corresponding to program used
inpres] - restr. of such a type should not be solved with this algorithm.
[inpdat] - size of working arrays too small for storing block at preparing block output
[whatva] - restr, of proport. not found
[pprop] - bad proportional data or program
[inpres] - type of restriction does not match any type of $<=>$

## Programmer's Guide for Balance Programs

## Programming conventions

All programs are written in FORTRAN-77 programming language, mostly in FORTRAN IV subset. Only some subroutines use new features of FORTRAN-77 language. The portability of programs is possible. Because of large arrays and large MINOS package computer for implementation must be virtual. Character variables are stored in variables of type integer (resp. integer*4 which is the same) as 4 characters per variable. Some of them are initialized in data statement. Comparison for equality is often used. Names of files are 12 characters long ( 3 variables), texts are stored as 40 characters ( 10 vars) generally.

Real data are stored in real* 8 variables for the reason of maximum possible accuracy at computing and reading input data.

For opening and closing files standard FORTRAN commands open, close are used. In some subroutines declaration character* x is used, also commands of type encode, decode (which is implemented as read or write from character variables) are used. It is used in functions inpinf (decode and character*Z declaration), presn (encode and ch*8) and otvor (ch*12). There is also an equivalence statement in routine presn and some small inconsistencies with types of parameters.

All these inconsistencies were corrected for the IBM version. Here commands marked as JBM or VAX must be chosen at implementation of programs for a particular computer.

There are 2 versions of programs now:
BALK - version where network algorithm for balancing is used. This is a very quick version but user can fix elements only.
link: file link.balk
files: balnew.o, bdata.o, comput.o, flow.o, maxfld.o, subri.o, vstup.o, vystup.o
BALM - version where MINOS is used for balancing. The user can use all types of restrictions and proportionalities.
link: file link.balm
files: balprp.o, bdata.o, subri.o, vstup.o, vystup.o. gmpsp.o, subp.o and MINOS subroutines

## PROGRAMS:

Main programs, subroutines and functions are grouped into several files. The following files contain the following subroutines:
balnew.f - balnew, err
balprp.f - balprp, err
bdata.f - block data
comput.f - comput
flow.f - flow
maxfld.f - maxfld
subr1.f - absind, indabs, sums, cpf
vstup.f - vstup, inpspe, inpgrp, inpnam, inpinc, inpdat, inpres, inpinf, inpipr, otvor, zatvor
vystup.f - vystup, outbal, outcor, outres, zaptab, outgrp, pripr, maluj, outsum, outtab, head, indexy, vyber, presn, tlacri, outace
gmpsp.f - ries, genmps, isrest, solut, calcfg
subp.f - pprop, isprop, genprp, merge, isresp, whatva

## COMMON

common /prop/lna,izc1,izcR,ikcr,izk,ikk,izl,ikl,izp1,iu,

* izs,iks,izr,ikr,iprflg (6)
in pgms: genprp,pprop,isresp,whatma,inpipr, solut
Ina - length of group A
izc1 - absolute index starting of group $C$
izc2 - absolute index for account C2
ikcr - absolute index end of group C
izk - absolute index for start of group K
ikk - absolute index for end of group K
izl - absolute index for start of group $L$
ikl - absolute index for end of group $L$
izp1 - absolute index for start of group P
iu - absolute index for start of group $U$
izs - absolute index for start of group $S$
iks - absolute index for end of group S
izr - absolute index for start of group $R$
ikr - absolute index for end of group $R$
iprflg(6)- lags which proportionality is chosen $1=y e s / 0=$ no
common /dsi/ ndspc,ndmpsa,ndoldb, ndnewb,ndspef,ndoutf,ndtty, ndf9,ndgrp,ndincl,ndnam,ndprop,ndfB,nddatb,ndrest,
* ndacc,ndobal,ndocor,ndores,ndogrp,ndoacc,ndoblk.
* ndodat,ndmps,ndmpsp
in pgms: balnew,balprp,block data,err, comput,fow,maxfld,cpf,inpspc,inpgrp, inpnam,inpinc,inpdat,inpres,inpinf,inpipr,outbal,outcor,outres, outgrp,outacc,ries, genmps,genprp,merge

All variables contain identifiers of files used in read/write statements. These values are assigned to them with the help of data statement and must be prepared at the installation of both programs. Identifiers are for following files:

```
ndspc - spec fle for balans;
ndmpsa - mpsall file;
ndoldb - old basis file for MINOS;
ndnewb - new basis file for MINOS;
ndspcf - specfile for MINOS, it must be equal to 5;
ndoutf - outfile for MINOS, it must be equal to 6;
ndtty - standard tty output, it is usually 6;
ndf9 - list output fle fort.9;
ndgrp - groups file;
ndincl-included blocks file;
ndnam - names file;
ndprop - proportional file;
ndf8 - scratch file for MINOS, it must be 8;
nddatb - data base file;
ndrest - restrictions file
ndacc - input accounts file;
ndobal - output balance table;
ndocor - output corrections table;
ndores - output results table;
ndogrp - output graph table;
ndoacc - output accounts file;
ndoblk - output blocks file;
ndodat - output database file;
ndmps - mpsfile
ndmpsp - mpsprop file
```

next 4 commons are from MINOS subroutines and are used in program ries only
common /lpcom / krhs,ns1,maxr,ierr,jr1
common/intcom/itn,itnitm,nphs,kmodlu,kmodpi
common /fxcom / fx,sinf,wtobj,minimz, nfx, ninf,obj, nprob
common/core / kz1,kz2,kz3

## ARRAYS AND VARIABLES

ind1 [mxvar ]-absolute row index of flow
ind2 [mxvar ]-absolute column index of flow
val [mxvar]-input flow
res [mxvar]-result flow
group [mxlng ] - identifiers of groups (1 char)
$\operatorname{lgr}$ [mxlng ] - length of groups
grpnam [10, mxlng] - names of group parts 40 chars
blinc 1 [mxinc ] - identifiers of included blocks 1.part (1 char)
blinc2 [mxinc ]-identifiers of included blocks 2.part (1 char)
inclpr [mxinc ] - output this block to file or not (1 yes/0 no)
accnam [10,mxacc] - names of accounts ( 40 characters each)
infacc [mxacc] - information which account to print
( $1=$ input balance, 2 =correction, $4=$ results, other their sum)
sincom [mxacc]-income of account (sum of the row)
soucom [mxacc] - outcome of account (sum of the column)
subory [3,15] - names of all files (12 character each)
solnam [10] - name of our task (40 characters)

lnginc - actual number of included blocks
lngacc - actual number of accounts
lngres - actual number of restrictions
lngip - actual length of propor
minkim - what algorithm is used $1=$ MINOS $2=$ network program

## ERRORS:

code of error [subroutine from which is called] - /subcode/ explanation A: Action to correct error

1 [inpspc] - /line/short file or bad data on input spec file A: correct spec file
2 [inpgrp] - /line/ number of groups $>$ mxgrp
A: change dimension of arrays group(), lgr(), grpnam() and variable mxlng in main program
3 [inpgrp] - /line / size in group is $<0$ or $>99$
A: correct groups file
4 [inpgrp] - /line/ wrong data on groups file (size of group) A: correct groups file
5 [inpgrp] - /line / group name is already used A: correct groups file
6 [inpinc] - /line/ such a group does not exist
A: correct included-blocks file or add such group into groups file
7 [inpinc] - /line/ block name is already used
A: correct included-blocks file
B [inpdat] - /line / parameters in database file are not correct (indexes out of range or $v=0$.)
A: correct database file
9 [inpinc] - /line/ number of included blocks $>$ mxinc
A: change dimension of arrays blinc 1(), blinc2(), inclpr() and mxinc in main program
10 [inpgrp] - /line / number of groups $=0$
A: correct groups file (it is empty)
11 [inpdat] - /line / number of variables $>\mathrm{mxvar}$
A: change dimension of arrays ind1(), ind2(), val(), res() value mxvar and other arrays depending on it for network algorithm in main program
12 [inpnam] - /line/ wrong data on names file
A: correct names file
13 [inpnam] - /line / group not found A: correct names file
14 [inpinc] - /line / number of included blocks $=0$
A: correct included-blocks file (it is empty)
15 [vstup ] - /0/ number of accounts $>$ mxacc
A: change dimension of arrays accnam(), infacc(), sincom(), soucom() and mxacc and other arrays for network algorithm in main program
16 [inpdat] - /line / last block at reading database file was not closed
A: correct database file
17 [inpdat] - /line / number of variables $=0$
A: correct database file (no input datum is included)
18 [inpres] - /line / number of restrictions $>$ mures
A: change dimension of arrays irest(), jrest(), nrest(), vrest() and mxres and also some arrays for network algorithm in main program
19 [inpres] - /line / bad data on restriction file (e.g. block not closed or fixing at negative number)
A: correct restriction file

20 [inpres] - /line/ number of restriction $=0$
A: correct restriction file (it was empty)
21 [absind\} - /line/ element does not match (wrong group name or outside limits)
A: correct database or restriction file
$2 R$ [inpres] - /line / you want to make restriction on non-existent flow
A: correct restriction file
23 [pripr ] - /0/ error at computing graph elements division by zero
A: consult with author
24 [maluj]-10/more than 128 accounts for preparing graph
A: not possible to have graph file, remove from spec file
25 [comput] - /0/ number of variables $>\operatorname{mmax}$
A: consult with author
26 [comput] - /0/ fatal error bad program
A: something strange with program, consult with author
27 [comput] - /0/n or nel are too big, change dimensions
A: consult with author
28 [comput] - error in optimization algorithm, second number gives code of err for network program:

1 -problem is unbounded
2 - problem cannot be solved
3-y is too big
4-wrong data
A: change task (change restrictions etc)

## [ries ]

## for MINOS

1 - infeasible (change task)
2-unbounded (change task)
3-too many iterations (increase number of iterations on specfile)
$>=4$ - another error condition (look into MINOS manuals)
$30-$ not enough core at reading mps file (correct dimension of $z()$ and nwcore in ries program)
40 - fatal error at reading mps file (perhaps increase number of rows, columns, elements ...)
30 [comput] - $10 /$ number of fixed elements $>$ nmax
A: consult with author
31 [inpspc] - /0/ error in spec file definition of task is incorrect
A: correct spec file
32 [inpinf] - /line / you want to have output correction or result accounts but no balancing function was chosen
A: correct input accounts file
33 [otvor] - $/ 0 /$ open error
A: file cannot be opened, look for name of file, consult with system programmer
34 [outgrp] - / $0 / \mathrm{obj}<0.00000001$, it is not possible get graph file
A: release out graph filename from spec file
35 [isrest] - $/ 0 /$ something strange with data in restriction file $\&$ source data
A: consult with author
36 [inpipr] - /line/ proportional tables too small
A: consult with author
37 [inpipr] - /line/ index in group outside limits at reading proportional file
A: correct proportional file
38 [inpipr] - /line/ such a group was not found at reading proportional file
A: correct proportional file
39 [inpipr] - /line/ such a block is not included at reading proportional file

A: correct proportional file
40 [inpinf] - /line / index of group outside limits at reading input account file
A: correct input accounts file
41 [inpinf] - /line/ such a group does not exist at reading input accounts file
A: correct input accounts file
42 [inpipr] - /line / error on proportional file - wrong number
A: correct proportional file
43 [inpspc] - /line/ algorithm is not corresponding to program used
A: change algor. on spec file of used BALM instead of BALK or vice versa
44 [inpres] - /line / restr. of such a type should not be solved with this algorithm
A: correct restriction file
45 [inpdat] - /line / size of working arrays too small for storing block at preparing block output at reading database file
A: change dimension of arrays tb()$, \mathrm{s} 1(), \mathrm{s} 2()$ and imt in program inpdat
46 [whatva] - /0/ restr of proport not found; strange error
A: consult with author
47 [pprop] - /code/ bad proportional data or program
0 - group A is not first in list
1 - size of group A is > 31
2 - size of group C is $<2$
3 - size of group K is $>9$
4 - size of group L is $>10$
$5-$ size of group $U$ is not $=1$
6-11-limits on proportional variables does not match
A: change task (size of groups) or change program (consult with author) or do not use proportionality
48 [inpres] - /line / type of restriction does not match to any type of < = >
A: correct restriction file

## FILES:

O- open (identifier of file), C - close, R-read, W - write
All files are opened in subroutine otvor, but 0 means the subroutine from which otvor is called; same for close.

Identifiers are assigned by user at installation time. There is standard assigning for VAX version in this paragraph. Files are opened and closed internally for VAX version, but externally for IBM version. In this case if more files have the same identifier for VAX, they are still separate files. For IBM it means that several files are concatenated to one file, where blank line serves as EOF separator.

For VAX version there are some files where names are given in program (user cannot define their names in spec file). They are files: specfile, fort.8, fort.9, mpsfile, mpsprop, mpsall.
spec file:
O: user (1), C,R: inpspe
database file O,C,R: inpdat (1)
groups file O,C,R: inpgrp (1)
incl.blocks file O,C,R: inpine (1) names file O,C,R: inpnam (1)
restriction file O,C,R: inpres (1)
accounts file
O,C,R: inpinf (1)
proportional file
O,C,R: inpipr (1)
specfile for MINOS O,C: ries (5), R: MINOS
standard tty output
O: system (6): W: flow,maxfld,ries,err, comput
standard output file fort. 9
O,C: system (9), W: flow,err,maxfld, ries, comput
balance table file
O,C: outbal (2), W: outtab,head,tlacri
correction table file
O,C: outcor (2), W: outtab,head,tlacri
results table file
O,C: outres (2), W: outtab,head,tlacri
graph file
O,C: outgrp (2), W: maluj,head
output accounts file
O,C: outacc (2). W: outacc,head,outsum
output blocks file
O,C,W; inpdat (13)
output database file O,C,W: zaptab (2)

```
outfile for MINOS
        O,C: ries (6),W: MINOS
mpsfile for MINOS
        O,C,W: genmps (1), R:MINOS, O,C,R: merge (2), O,C,R: cpif (1)
mpsprop file for MINOS
    O,C,W: genprp (3), O,C,R: merge (3)
mpsall file for MINOS
    O,C,W: merge (4), O,C; ries (2), R: MINOS, O,C,W; cpf (\Omega)
scratch file for MINOS fort.8
    O,C: system (B), W,R: MINOS
```

Some remarks concerning how output files are prepared:
a/ preparing output balance, correction and result files: calling sequence of subroutines for preparing these files is the same, only input data are different. For balance file input data are got from array val, for result file it is array res and for correction it is array res, which is difference res=res- val and which is corrected back to previous values after preparing correction file. Calling sequence is: sums prepares sums for whole columns. Then outtab for each page (which contains 13 output accounts) prints heading (name of task), page number - subroutine head) and names for each output accounts (it prepares subroutine indexy, name of account grupy and relative numbers within these groups are irel and also separators after end of group in array oddel which is either blank or $\mid$. Then it prints for each input account values for corresponding 13 output accounts. These output accounts for each row are formed with a help of subroutine vyber. Then these numbers are stored in maximal possible accuracy into output array prem with a help of routine presn and are printed within a subroutine tlacri. The whole cycle is repeated for each page after the whole table is printed.
b/ preparing graphical output: sequence of routines is pripr, which prepares array res in such a way, that it stores values $0,1,2,4,8$ for each variable according to whether this variable is no change(0), is on upper band(1), on lower bound(2), between(4) or fixed(8). Then routine maluj is called. Within this routine at first heading is printed and then for each row values are taken for these variables for the whole row from array res with the help of routine vyber. Then each variable according to its code is replaced with its corresponding graphical symbol (.+-oe) and is printed on line. Because of length of line of prints, where a maximum of 132 characters is possible, the maximal size of 128 accounts was chosen, so as not to split one row between several lines on the printer

## Description of programs:

For Parameter section: most parameters are described in ARRAYS \& VARIABLES section, here we give only parameters which are not described there, or need more explanation.

For Method section: there is no description if method is trivial.

## BALPRP

```
        program balprp
Call: vstup ries vystup
Called from:
Purpose: main program for BALM where MINOS is used
```


## ABSIND

subroutine absind(bl1,bl2,in1,in2,iab1,iab2,group,lgr,lnggr,line)
Call: err
Called from: inpdat inpres
Purpose: from relative indexes and group names gives absolute indexes for flow
Parameters: bl1,bl2 - group names (1 char) for row and column in1,in2 - relative indexes within group iab1,iab2 - absolute indexes
Method: At first computes absolute index of starting row for group, then index for particular row; the same for column

## BLOCK DATA

block data
Purpose: initiates variables in common /dsi/

## CALCFG

subroutine calcfg(mode, $n, x, f, g$, nstate, $n p r o b$ )
Call:
Called from: MINOS subroutine
Purpose: dummy subroutine for MINOS, it will never be called
Parameters: all are dummy variables

## CPF

> subroutine cpf

Call:
Called from: ries
Purpose: copies file mpsfile to mpsall if no proportionality was chosen

## ERR

subroutine err(i,j)
Call:
Called from: absind inpdat inpgrp inpinc inpinf inpipr inpnam
inpres inpspe isrest maluj otvor pprop pripr
ries vstup comput
Purpose: prints error codes on fort. 9 and standard tty output and STOP
Parameters: i - code of error; j-subcode of error

## GENMPS

subroutine genmps(ind1,ind2, val,lngvar,irest,jrest,nrest, vrest,lngres, sincom, soucom, lngacc,ipflag,ip1z,ip1k,ip2z,ip2k,lngip)
Call: otvor isprop isrest zatvor
Called from: ries
Purpose: prepares mps file mpsfile for MINOS
Method: row, column, rhs and bound parts are prepared sequentially. If flow is proportional (call of routine isprop is $=1$ ) skip commands which generate part of mpsfile for these variables. Similarly skip commands if flow is fixed.
For flow with indexes $i, j$ the column 'aiii. jjj ' is generated (here iii $=100+\mathrm{ijjj}=100+\mathrm{j}$ ). Two rows 'loiiijjj' \& 'upiiijjj' are also generated for each flow. Sequence of flows is the same as if they are stored in arrays ind1,ind2, val. For each account i a row is generated with name 'bal.iii'. All input flows are going with coef 1. into this row and -1 . for output accounts.
Objective function has name 'aim' objective variable has name 'u'. Rhs has name 'side' and bounds 'bnd'.
For bounds part for restricted variables we generate following rows:
if type < then 2 rows: lo bnd aiii.jjj 0 .
up bnd aiii.jjj fixed_yalue
if type $=$ then 1 row: fx bnd aiii.jjj fixed_yalue
if type > then 1 row: lo bnd aiii.jjj fixed_yalue

## GENPRP

subroutine genprp(ind1,ind2,val,lngvar,ip1z,ip1k,ip2z,ip2k,lngip,alfa, beta,gama,sa,sk,sl,c1a,c2a,p1a,zzz,x,y,irest,jrest,
nrest,vrest,lngres)
Call: otvor isresp isrest isprop whatva zatvor
Called from: ries
Purpose: prepares proportional mps file mpsprop
Method: File mpsprop will contain only data for proportional variables Also here row, column, rhs \& bounds are prepared. This subroutine was written especially for our 6 types of proportionality cases. For each type of proportionality corresponding rows, columns and bounds are generating only if particular type of proportionality was chosen. For propor. in blocks AA AK AL names of variables will be a100.jjj which is sum for the whole column of this block.

GO

> subroutine go(z.nwcore,lpi)
> Call: MINOS subroutines
> Called from: ries
> Purpose: main program for MINOS package

HEAD
subroutine head(odsi,sub, solnam, page)
Call:
Called from: maluj outace outtab
Purpose: Prints heading to output balance, correction, result and graph file Parameters: odsi-identifier of file
page - page number of output file

## INDABS

subroutine indabs(iab1,iab2,bl1,bl2,in1,in2,group,lgr,lnggr)
Call:
Called from: maluj zaptab
Purpose: Converts absolute indexes of flow to indexes relative in block
Parameters: see routine absind
Method: at first it finds group to which indexes belong and then indexes

## INDEXY

subroutine indexy(iab2z,iab2k,group,lgr,lnggr,grupy,oddel,irel)
Call:
Called from; outtab
Purpose:for each page it computes names of groups, accounts and their relative separators
Parameters: grupy - names of corresponding 13 accounts for this page oddel - separators - for end of group it is | otherwise blank irel - array of relative indexes for corresponding 13 accounts iab2z - starting absolute index of column iab2k - ending absolute index of column

## INPDAT

subroutine inpdat(sub,ind1,ind2,val,lngvar, mxvar, blinc1,blinc2,
lnginc, group,lgr, lnggr,iflag,iflgb, sub1,inclpr)
Call: otvor err absind zatvor
Called from: vstup
Purpose: Reads data from database file and prints output block file if necessary
Parameters: sub1-name of output block file

## INPGRP

subroutine inpgrp(sub,group,lgr,lnggr,mxlng,lngacc,grpnam)
Call: otvor err zatvor
Called from: vstup
Purpose: Reads data from groups file

## INPINC

subroutine inpinc(sub,group,lnggr,blinc 1,blinc2,lnginc,mxinc,inclpr)
Call: otvor err zatvor
Called from: vstup
Purpose: Reads data from include file

## INPINF

subroutine inpinf(sub,infacc,lngacc,infall, group,lgr,lnggr,ifg)
Call: otvor err zatvor
Called from: vstup
Purpose: Reads data from input accounts file

## INPIPR

subroutine inpipr(sub,ip1z,ip1k,ip2z,ip2k,mxip,lngip,group,lgr,
lnggr, blinc 1, blinc2,lnginc)
Call: otvor err zatvor
Called from: vstup
Purpose: Reads data from input proportionality file

## INPNAM

subroutine inpnam(sub,group,lgr,lnggr,accnam,lngacc)
Call: otvor err zatvor
Called from: vstup
Purpose: Reads data from names file

## INPRES

subroutine inpres(sub,group,lgr,lnggr,ind1,ind2,lngvar,irest,jrest, nrest, vrest,lngres,mxres,natype)
Call: otvor absind err zatvor
Called from: vstup
Purpose: Reads restriction file
Parameters: natype - type of algorithm user wants to use for solving his problem ( 1 for MINOS, 2 for network program)

## INPSPC

subroutine inpspc(subory,solnam,iflg,eps,seps,epstol,minkim,natype)
Call: err zatvor
Called from: vstup
Purpose: Reads spec file

## ISPROP

function isprop(ii,jj,ip1z,ip1k,ip2z,ip2k,lngip)
Call:
Called from: genmps genprp solut
Purpose: If flow is of proportional type returns 1, otherwise 0
Parameters: ii,jj-absolute indexes for flow
subroutine isresp(ii,jj,irest,jrest, nrest,vrest,lngres,if1,ifl2,value)
Call:
Called from: genprp
Purpose: If flow is in list of restricted \& it is proport., returns ifl1=1
and ifl2 \& value according to parameters in list of restricted variables
for proportionality for types 1,2 or 3 it returns also index
Parameters: $\mathrm{ii}, \mathrm{jj}$ - absolute indexes for flow ifl1 - returns 1 if restricted, 0 otherwise ifl2 - returns type of restriction value - returns restricted value

## ISREST

subroutine isrest(ii,jj,irest,jrest,nrest,vrest,lngres,if1,ifl2,
value,ind1,ind2,val,lngvar)
Call: err
Called from: genmps genprp
Purpose: tests if flow is in list of restricted variables, output is same as for routine isresp

## MALUJ

subroutine maluj(odsi, ind1,ind2,res,lngvar,lngacc, solnam,sub,obj,
epstol, group,lgr,lnggr, blinc1,blinc2,lnginc)
Call: err head vyber indabs
Called from: outgrp
Purpose: routine prints output graph file
Method: at first prepares head, gets absolute indexes.
For each accounts it takes all flows(subr. vyber). According to their values assign one character for each flow (, $+,-, e, o,$. ) and prints one line. Because of the limitation of paper width the whole graph can be on one page (not divided into several pages) the maximum number of accounts is limited to 128 . If there are more accounts, it is not possible to have a graph output.

## MERGE

subroutine merge
Call: otvor zatvor
Called from: ries
Purpose: Merges mpsfile with mpsprop producing thus mpsall as input file for MINOS

## OTVOR

subroutine otvor(idsi,menof)
Call: err
Called from: genmps genprp inpdat inpgrp inpinc inpinf inpipr inpnam inpres merge outacc outbal outcor outgrp outres ries zaptab
Purpose: Opens file with identifier idsi and name menof. This routine is empty for 3IBM version.

## OUTACC

subroutine outacc(sub,ind1,ind2,val,res,lngvar,sincom,soucom,Ingace,
infacc, infall,acenam,group,lgr,Inggr,epstol,solnam,grpnam)
Call: otvor head outsum zatvor
Called from: vystup
Purpose: writes ouput account file

## OUTBAL

subroutine outbal(sub,ind1,ind2,val,lngvar, sincom, soucom,lngace, solnam, acenam, group.lgr.lnggr)
Call: otvor sums outtab zatvor
Called from: vystup
Purpose: writes output balance file

## OUTCOR

subroutine outcor(sub,ind1,ind2,val,res,lngvar,sincom,soucom,Ingace, solnam,accnam,group,lgr,lnggr)
Call: otvor sums outtab zatvor
Called from: vystup
Purpose: writes output correction table

## OUTGRP

subroutine outgrp(sub,ind1,ind2,val,res,lngvar,irest,jrest,nrest,
vrest,lngres,solnam,lngacc,obj, epstol, group,lgr,
lnggr,blinc 1, blinc2,lnginc)
Call: otvor pripr maluj zatvor
Called from: vystup
Purpose: writes output graph file

## OUTRES

subroutine outres(sub,ind1,ind2,res,Ingvar,sincom, soucom,Ingace, solnam,accnam,group,lgr,lnggr)
Call: otvor sums outtab zatvor
Called from: vystup
Purpose: writes output results file

## OUTSUM

subroutine outsum(ins,odsi,in1,in2,var, Ingvar, sincom, soucom,lngacc, acenam,group,lgr,Inggr,epstol,grpnam)
Call: sums
Called from: outace
Purpose: routine prints sums or income/outcome accounts; it is called separately for input data, for results and corrections
Parameters: ins - if it is 0 then produces sums of all accounts, otherwise there are accounts for row/column ins
in1,in2,var - indexes and values for input, correction or results data

## OUTTAB

subroutine outtab(odsi,in1,in2,var,lngvar,group,lgr,lnggr,sub, accnam, Ingace,solnam,sincom,soucom)
Call: head indexy vyber presn tlacri
Called from: outbal outcor outres
Purpose: prints output table
Method: Array var contains values to print, which may be input data, corrections or results (it depends from what subroutine it is called). A first routine prints heading. It must get absolute and relative addresses for each output account and separators after flows (call routine indexy). Then sequentially 13 output accounts are printed on one page. For each row it takes flows (call vyber), prepares values in maximum possible accuracy (call presun) \& prints one row (call tlacri).

## PPROP

subroutine pprop(ind1,ind2,val, Ingvar,soucom, lngace, alfa, beta, gama, sa, sk,sl,c1a,c2a,p1a,zzz,x,y,group,lgr,lnggr,ip1z,ip1k, ip2z,ip2k,lngip)
Call: err
Called from: ries
Purpose: prepares necessary arrays \& variables if proportionality was chosen for later usage

## PRESN

subroutine presn(prvky,prem,kon)
Call:
Called from: outtab
Purpose: prepares data from real*B mode to formated mode in maximum possible accuracy
Parameters: prvky - max 13 values of flows (real*B)
prem - array $(2,13)$ of 4 characters (implemented as integer*4)
where values in formated mode will be
kon - number of arrays prvky ( $<=13$ )
Method: This routine is strongly dependent on FORTRAN dialect. The best way to transfer data from real*B form to formated form is with a decode statement, which is not available for IBM FORTRAN. For IBM it is done with the help of file 8 (Data are written to file and read as character

## PRIPR

subroutine pripr(ind1,ind2,val,res,Ingvar,irest,jrest, nrest,vrest, Ingres,obj,epstol)
Call: err
Called from: outgrp
Purpose: prepares array res for graph output in such a way that it stores numbers depending on whether the flow exists, if the change was to the upper or lower bound, if it was fixed etc.

## RIES

subroutine ries(ind1,ind2,val, res, lngvar,irest,jrest, nrest,vrest, Ingres, istate, obj, sincom, soucom, Ingace, ipflag, ip $1 \mathrm{z}, \mathrm{ip} 1 \mathrm{k}, \mathrm{ip} 2 \mathrm{z}$, ip2k,lngip,group,lgr,lnggr)
Call: sums genmps pprop genprp merge otvor go zatvor solut err cpf
Called from: balprp
Purpose: controls solution with MINOS algorithm
Method: at first it prepares file mpsfile. If proportionality was chosen it prepares file mpsprop. Then these two files merge to produce thus mpsall, which is input mps file for MINOS. It copies mpsfile to mpsall if no proportionality. Then MINOS is called and after solution gets results (call solut).

SOLUT
subroutine solut(zz,ind1,ind2,res,val,lngvar,ipflag,alfa,beta,gama, sa,
sk,sl,ip1z,ip1k,ip2z,ip2k,lngip,c1a,c2a,p1a)
Call: isprop
Called from: ries
Purpose: gets results from working array $z()$

## SUMS

subroutine sums(ins,in1,in2,var, Ingvar, sincom,soucom, Ingace,sinp,sout) Call:
Called from: outbal outcor outres outsum ries
Purpose: returns accounts or sum of accounts
Parameters: ins - 0 for sum of accounts, integer number for corresp. account sinp - sum of the whole input account sout - sum of the whole output account
subroutine tlacri(odsi,i,grp,irl,prem,oddel,kon)
Call:
Called from: outtab
Purpose: prints one line of output data, correction or results tables
Parameters: i - row number = absolute account number
grp-relative account name irl - relative account number prem - values of flows in character mode oddel - separators after each flow (blank of | after group) kon-number of columns

## VSTUP

subroutine vstup(subory,solnam,iflg,group,lgr, blinc1,blinc2,accnam, ind 1,ind2, val,irest,jrest, nrest, vrest,infacc,infall, mxvar, mxlng, mxinc, mxace, mxres, lngvar, lnggr, lngine, lngacc,lngres,eps,seps,epstol,ip1z,ip1k,ip2z,ip2k, mxip,lngip,grpnam,minkim,inclpr)
Call: inpspe inpgrp err inpinc inpdat inpnam inpres inpinf inpipr
Called from: balprp balnew
Purpose: calls sequentially all input routines (if necessary) to read input data

## VYBER

subroutine vyber(in1,in2,var,lngvar,i1,ia2z,ia2k, prvky)
Call:
Called from: maluj outtab
Purpose: searches in arrays in1, in2, var for flows between interval ia2z ia2k for input account i1
Parameters: i1 - absolute input account number ia2z-absolute output account starting interval ia2k-absolute output account ending interval prvky - values moved from array var

## VYSTUP

subroutine vystup(ind1,ind2,val,res,lngvar, group,lgr,lnggr, blinc1, blinc2, Inginc, accnam, sincom, soucom, lngacc, infacc, infall,irest,jrest,nrest, vrest,lngres,subory,solnam, iflg,epstol,obj,grpnam)
Call: outbal outcor outres outacc zaptab outgrp
Called from: balprp balnew
Purpose: Calls main output subroutines for preparing output files

## FHATVA

subroutine whatva(ii,jj,ind1, ind2,val,lngvar,alfa,beta,gama,oldval,value)
Call:
Called from: genprp
Purpose: returns value of restricted variables in tables AA, AK, AL

## ZAPTAB

subroutine zaptab(sub,ind1,ind2,res,lngvar,group,lgr, lngg)
Call: otvor indabs zatvor
Called from: vystup
Purpose: writes data (results flows) to output database file

## ZATVOR

subroutine zatvor(idsi)
Call:
Called from: genmps genprp inpdat inpgrp inpinc inpinf inpipr inpnam inpres inpspe merge outacc outbal outcor outgrp outres ries zaptab
Purpose: closes file with identifier idsi - empty for IBM version

## BALNEW

program balnew
Call: vstup comput vystup
Called from:
Purpose: main program for network version, it calls input module, then subroutines for solving problem with network algorithm and then output module

## COMPUT

subroutine comput(cut,head,headr,stk, dfr,df0, df, a, arc,list,fi, ind1, ind2,val,res,lngvar,irest,jrest,nrest,vrest,lngres, num,cc,y,eps,seps,nmax,mmax,mmax2,mmax4)
Call: err flow
Called from: balnew
Purpose: controls solution with network algorithm - checks balance, balances input data without or with restrictions

## FLOW

subroutine flow( $n, e$, edge, $f, a, d f, y, c c, d f 0, c u t, h e a d, l i s t, d f r$, stk, headr,
nmax,mmax,mmax2,mmax4,eps,seps)
Call: maxfld
Called from: comput
Purpose: finds minimax flow
Parameters \& method: see comments in program and another part of this manual

## MAXFLD

subroutine maxfld(n,e,edge,fldf,cut,flow,r,code,head,list, dfr, stk,
headr,nmax,mmax,mmax2,mmax4,eps)
Call:
Called from: flow
Purpose: constructs maximal flow in directed network
Parameters \& method: see comments in program and another part of this manual

## Remarks concerning MINOS

Changes into standard version of MINOS 3.4

1. instead of \#implicit real*B(c-g,o-z) \# use \#implicit real*B(a-g,o-z)\# in all MINOS subroutines.
2. in subroutine initlz set \#nwordr=1\#
3. change \#abs\# to \#dabs\# in routines: crash, factor, forms, insert, invert, loadn, mps, rgitn, setx, solprt
4. in subroutine unpack set \#real*B a\#
5. in subroutine mpsin change parameters ne, m,krhs to ne2,m2,krhs2 before \#call mps ( $3, \ldots$ \# instead of last 3 lines insert lines:
```
ne2=2* ne
mm2 = 2* m
krhs2 = 2 * krhs
call move (...,0,ne,ne2)
call move (...,krhs2,n,mm2)
call move (...,krhs2,n,mm2)
```

6. in subroutine go description is \#subroutine go(z,nwcore,lpi)\#
7. in subroutine go save ierr parameter:
before label 100 insert line: icnt=0
after label 100 add line : icnt=icnt+1
after call minos add line : if(ient.eq.1) iersol=ierr
after label 900 add line : ierr=iersol
B. remove tabulator from all programs
8. remove commands concerning measuring of time consumed by MINOS


Figure 1.


Figure 2.

Initial basis

$$
X(5)=20=Y \cdot A(5)
$$



$$
\mathrm{Y}=0.66
$$

$$
\begin{aligned}
& X(5)=13.3=Y \cdot A(5) \\
& X(2)=6.6=Y \cdot A(2)
\end{aligned}
$$



Figure 3.

