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### Difference Covering Arrays and Pseudo-Orthogonal Latin Squares

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Abstract A pair of Latin squares, A and B, of order n, is said to be pseudoorthogonal if each symbol in A is paired with every symbol in B precisely once, except for one symbol with which it is paired twice and one symbol with which it is not paired at all. A set of t Latin squares, of order n, are said to be mutually pseudo-orthogonal if they are pairwise pseudo-orthogonal. A special class of pseudo-orthogonal Latin squares are the mutually nearly orthogonal Latin squares (MNOLS) first discussed in 2002, with general constructions

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A. Rao School of Mathematical and Geospatial Sciences RMIT University Melbourne, VIC 3000, Australia E-mail: asha@rmit.edu.au given in 2007. In this paper we develop row complete MNOLS from difference covering arrays. We will use this connection to settle the spectrum question for sets of 3 mutually pseudo-orthogonal Latin squares of even order, for all but the order 146.

**Keywords** Difference covering array  $\cdot$  Latin squares  $\cdot$  pseudo-orthogonal Latin squares  $\cdot$  mutually nearly orthogonal Latin squares

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#### **1** Introduction

Difference matrices are a fundamental tool used in the construction of combinatorial objects, generating a significant body of research that has identified a number of existence constraints. These difference matrices have been used for diverse applications, for instance, in the construction of authentication codes without secrecy [13], software testing [3,4] and data compression [9]. This diversity of applications, coupled with existence constraints, has motivated authors to generalise the definition to holey difference matrices, difference covering arrays and difference packing arrays, to mention just a few.

Some orthogonal Latin squares are combinatorially equivalent to difference matrices [8] with specific properties, hence many of the applications for difference matrices apply to orthogonal Latin squares. However in the design of certain experiments, designs that are close to being mutually orthogonal are preferred [1,2]. Thus constructing sets of pseudo-orthogonal Latin squares, (and their combinatorially equivalent difference matrices) are of interest to experimental scientists, in addition to satisfying mathematical curiosity.

In the current paper we are interested in constructing subclasses of cyclic difference covering arrays and exploiting these structures to emphasize new connections with other combinatorial objects, such as pseudo-orthogonal Latin squares. We use this connection to settle the existence spectrum for sets of 3 mutually pseudo-orthogonal Latin squares of even order, in all but one case. We begin with the formal definitions.

A difference matrix (DM) over an abelian group (G, +) of order n is defined to be an  $\lambda n \times k$  matrix Q = [q(i, j)] with entries from G such that, for all pairs of columns  $0 \le j, j' \le k - 1, j \ne j'$ , the difference multi-set

$$\Delta_{j,j'} = \{q(i,j) - q(i,j') \mid 0 \le i \le \lambda n - 1\}$$

contains every element of G equally often, say  $\lambda$  times. E.g. [5,7,8]. Note that we label the rows from 0 to  $\lambda n - 1$  and the columns 0 to k - 1. Also to be consistent with later sections involving Latin squares and covering arrays our definition uses the transpose of the matrix given in [5] and [7]. Since the addition of a constant, over G, to any row and a constant to columns does not alter the set  $\Delta_{j,j'}$ , we may assume that one row and one column contain only 0, the identity element of G. More precisely, to simplify later calculations, we will assume that all entries in the last row and last column of Q are 0. A difference matrix will be denoted  $DM(n, k; \lambda)$ . If (G, +) is a cyclic group we refer to a *cyclic difference matrix*.

#### **Theorem 1** [5, Thm 17.5, $p \downarrow 11$ ] A DM $(n, k; \lambda)$ does not exist if $k > \lambda n$ .

In the main, we will use difference matrices with k = 4,  $\lambda = 1$  and where possible we will work with cyclic difference matrices. In Section 4 we list a number of existence results.

A holey difference matrix (HDM) [16] over an abelian group (G, +) of order n with a subgroup H of order h is defined to be an  $\lambda(n-h) \times k$  matrix Q = [q(i, j)] with entries from G such that, for all pairs of columns  $0 \leq j, j' \leq k-1$ ,  $j \neq j'$ , the difference multi-set

$$\Delta_{j,j'} = \{q(i,j) - q(i,j') \mid 0 \le i \le \lambda(n-h) - 1\}$$

contains every element of  $G \setminus H$  equally often, say  $\lambda$  times. A holey difference matrix will be denoted HDM(k, n; h), where |G| = n and |H| = h. If G is a cyclic group then we refer to a cyclic holey difference matrix.

As before, a constant may be added to any column without affecting  $\Delta_{j,j'}$ so we may assume that all entries in the last column of Q are equal to 0. However since H is a subgroup, 0 belongs to the hole. Consequently, 0 does not occur in  $\Delta_{j,j'}$ , and thus there will be no row containing two or more 0's. Further since  $\Delta_{j,k-1} = \lambda(G \setminus H), 0 \leq j \leq k-2$ , the entries of H do not occur in the first k-1 columns of Q.

A difference covering (packing) array over an abelian group (G, +) of order n is defined to be an  $\eta \times k$  matrix Q = [q(i, j)] with entries from G such that, for all pairs of distinct columns  $0 \leq j, j' \leq k - 1$ , the difference multi-set

$$\Delta_{j,j'} = \{q(i,j) - q(i,j') \mid 0 \le i \le \eta - 1\}$$

contains every element of G at least (at most) once, see for example [16, 17]. A difference covering array will be denoted DCA $(k, \eta; n)$  and a difference packing array will be denoted DPA $(k, \eta; n)$ . If (G, +) is the cyclic group, then the difference covering (packing) array is said to be *cyclic*. As before we may assume that the last row and last column of a DCA $(k, \eta; n)$  contain only 0.

In the papers [16] and [17], Yin constructs cyclic DCA(4, n + 1; n) for all even integers n, with similar results for cyclic difference packing arrays. Yin documents a number of product constructions for difference covering arrays, some of which will be reviewed in Section 4 and then adapted to construct difference covering arrays with specific properties; properties that build connections with pseudo-orthogonal Latin squares.

The additional properties that we seek are that 0 (the identity element of G) occurs at least twice in each column of the DCA(k, n + 1; n) and for pairs of columns, not including the last column, the repeated difference is not the element 0. Formally, we are interested in DCA(k, n + 1; n), Q = [q(i, j)],  $(0 \le i \le n, 0 \le j \le k - 1)$  satisfying the properties:

**P1.** the entry  $0 \in G$  occurs at least twice in each column of Q, and

- **P2.** for all pairs of distinct columns j and j',  $j \neq k 1 \neq j'$ ,  $\Delta_{j,j'} = \{q(i, j) q(i, j') \mid 0 \le i \le n 1\} = G \setminus \{0\},\$ 
  - Note that **P2** implies that  $\Delta_{j,j'}$  contains a repeated difference that is not 0. The following example is a cyclic DCA(4,7;6) satisfying **P1** and **P2** [12].

$$B^{T} = \begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 0 \\ 1 & 3 & 5 & 0 & 2 & 4 & 0 \\ 3 & 0 & 4 & 1 & 5 & 2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

In the next lemma we show that if G is the cyclic group  $\mathbb{Z}_n$ , then these conditions imply that for all distinct columns j and j',  $j \neq k - 1 \neq j'$ ,

$$\Delta_{j,j'} = \{0, 1, 2, \dots, n/2, n/2, \dots, n-1\}$$

with repetition retained.

**Lemma 1** If there exists a cyclic DCA(k, n + 1; n), Q = [q(i, j)],  $(0 \le i \le n, 0 \le j \le k - 1)$  satisfying properties **P1** and **P2**, then n is even. Further, given  $d_0$  such that  $d_0 = q(i, j) - q(i, j') = q(i', j) - q(i', j')$ , for  $i \ne i'$  and  $k - 1 \ne j \ne j' \ne k - 1$ , then  $d_0 = n/2$ .

Proof Let Q = [q(i, j)]  $(0 \le i \le n, 0 \le j \le k - 1)$  represent the difference covering array. The definition requires that  $\mathbb{Z}_n \subseteq \Delta_{j,j'}$  and since column k-1of Q contains all zeros, property **P1** implies that the remaining columns are permutations of the multi-set  $\{0, 0, 1, 2, \ldots, n-1\}$ .

Let  $d_0 \in \mathbb{Z}_n \setminus \{0\}$  represent the repeated difference in  $\Delta_{j,j'}$ . Suppose *n* is odd and, without loss of generality, that column 0 is in standard form. Then, for all  $0 < j \leq k - 2$ ,  $\sum_{i=0}^{n-1} q(i,j) = \frac{(n-1)n}{2}$  and

$$\sum_{i=0}^{n-1} (i - q(i,j)) \equiv \frac{(n-1)n}{2} + d_0 \mod n.$$

Consequently,  $2d_0 = n(2\ell - n + 1)$ , for some integer  $\ell$ , or equivalently  $n|2d_0$ . But since n is odd, this leads to the contradiction,  $d_0 \in \mathbb{Z}_n$  and  $n|d_0$ . Thus n is 2p for some integer p, where p divides  $d_0$ , implying  $d_0 = p$ .

The remainder of this paper is organised as follows. In Section 2 we will draw the connection between DCA(k, n + 1; n) and sets of mutually pseudo-orthogonal Latin squares. In Section 3 we give three new constructions for DCA(4, n+1; n)'s and consequently new families of mutually pseudo-orthogonal Latin squares. In Section 4 we review some of the general constructions for difference covering arrays and show that these constructions can be used to construct DCA(k, n + 1; n) that satisfy Properties **P1** and **P2**, and hence mutually pseudo-orthogonal Latin squares. This leads to showing that 3 mutually pseudo-orthogonal Latin squares exist for every order except possibly 146.

The notation  $[a, b] = \{a, a + 1, \dots, b - 1, b\}$  refers to the closed interval of integers from a to b.

#### 2 Pseudo-orthogonal Latin squares and difference covering arrays

In this section we verify that cyclic difference covering arrays can be used to construct pseudo-orthogonal Latin squares.

A Latin square of order n is an  $n \times n$  array in which each of the symbols of  $\mathbb{Z}_n$  occurs once in every row and once in every column. Two Latin squares A = [a(i, j)] and B = [b(i, j)], of order n, are said to be orthogonal if

$$O = \{(a(i,j), b(i,j)) \mid 0 \le i, j \le n-1\} = \mathbb{Z}_n \times \mathbb{Z}_n.$$

A set of t Latin squares is said to be mutually orthogonal, t-MOLS(n), if they are pairwise orthogonal. A set of t idempotent MOLS(n), denoted t-IMOLS(n), is a set of t-MOLS(n) each of which is idempotent; that is, the cell (i, i) contains the entry i, for all  $0 \le i \le n - 1$ . It is well known that difference matrices can be used to construct sets of mutually orthogonal Latin squares, see for instance [8, Lemma 6.12].

Raghavarao, Shrikhande and Shrikhande [12] and Bate and Boxall [1] slightly vary the orthogonality condition to that of pseudo-orthogonal. A pair of Latin squares, A = [a(i,j)] and B = [b(i,j)], of order n, is said to be *pseudo-orthogonal* if given  $O = \{(a(i,j),b(i,j)) \mid 0 \leq i,j \leq n-1\}$ , for all  $c \in \mathbb{Z}_n$ 

$$|\{(c, b(i, j)) \mid (c, b(i, j)) \in O\}| = n - 1.$$

That is, each symbol in A is paired with every symbol in B precisely once, except for one symbol with which it is paired twice and one symbol with which it is not paired at all. A set of t Latin squares, of order n, are said to be mutually pseudo-orthogonal if they are pairwise pseudo-orthogonal.

Mutually nearly orthogonal Latin squares (MNOLS) are a special class of pseudo-orthogonal Latin squares, in that the set O does not contain the pair (c, c), for any  $c \in \mathbb{Z}_n$  and the pair  $(i, i + n/2 \mod n)$  appears twice in O, where  $0 \le i \le n - 1$ . Mutually nearly orthogonal Latin squares (MNOLS) were first discussed in a paper by Raghavarao, Shrikhande and Shirkhande in 2002 [12].

A natural question to ask is: Can we use difference techniques to construct mutually pseudo-orthogonal Latin squares? Raghavarao, Shrikhande and Shirkhande did precisely this and constructed mutually pseudo-orthogonal Latin squares from cyclic DCA(k, n+1; n) termed (k, n)-difference sets in [12].

**Theorem 2** [12] If there exists a cyclic DCA(t+1, 2p+1; 2p), Q' = [q'(i, j)], that satisfies **P1** and **P2**, then there exists a set of t mutually nearly orthogonal Latin squares of order 2p.

Proof Recall that without loss of generality we may assume that the last row and column of Q' contain all zeros. Construct a new matrix Q = [q(i, j)] by removing the last row and last column from Q' and define a set of t arrays,  $L_s = [l_s(i, j)], 0 \le s \le t - 1$ , of order 2p, by

$$l_s(i,j) = q(i,s) + j \pmod{2p}, \ 0 \le i,j \le 2p - 1.$$
(1)

It is easy to see that each column of  $L_s$  is a permutation of  $\mathbb{Z}_{2p}$  and so  $L_s$  is a Latin square. By Lemma 1

$$\Delta_{j,j'} = \{q'(i,j) - q'(i,j') \mid 1 \le i \le 2p\} = (\mathbb{Z}_{2p} \setminus \{0\}) \cup \{p\}$$

implying that when any two Latin squares are superimposed we obtain the set of ordered pairs  $(\{\mathbb{Z}_{2p} \times \mathbb{Z}_{2p}\} \setminus \{(x,x) \mid 0 \le x \le 2p-1\}) \cup \{(x,x+p) \mid 0 \le x \le 2p-1\}$  with repetition retained.

If there exists a pair of pseudo-orthogonal Latin squares generated from cyclic difference covering arrays satisfying **P1** and **P2**, then there exists a pair of nearly orthogonal Latin squares. Conversely, a pair of nearly orthogonal Latin squares are necessarily pseudo-orthogonal Latin squares. Given this and the connection with [10] and [12] we will state all results in terms of MNOLS.

Raghavarao, Shrikhande and Shirkhande established bounds on the maximum number of Latin squares in a set of MNOLS [12]. This result provides bounds on t for DCA(t, n + 1; n) that satisfy **P1** and **P2**.

**Lemma 2** Let  $p \ge 2$  be a positive integer. If there exists a DCA(t + 1, 2p + 1; 2p) that satisfies **P1** and **P2**, then  $t \le p + 1$ . Further if p is even and there exists a DCA(t, 2p + 1; 2p), then t .

*Proof* Theorem 5.2 of [12] gives bounds on NMOLS(2p). Theorem 2 shows that this can be applied to DCA(t + 1, 2p + 1; 2p) and DCA(t, 2p + 1; 2p).

An  $n \times n$  Latin square  $L = [\ell(i, j)]$  with entries from  $\mathbb{Z}_n$  is row complete if the ordered pairs  $(\ell_{i,j}, \ell_{i,j+1})$  are all distinct for all  $1 \le i \le n$  and  $1 \le j \le n-1$ . Williams showed that the Latin square corresponding to the cyclic group of even order is row complete, [15]. Since the MNOLS(2p) constructed from cyclic difference covering arrays are essentially copies of the cyclic group, they are row complete. In addition, they are *bachelor squares*, in that they have no orthogonal mate.

# 3 Construction of difference covering arrays DCA(4, 2m + 1; 2m)

This section is devoted to giving new constructions for families of cyclic DCA(4, 2m + 1; 2m) when m = 2k + 1, m = 8k + 4 and m = 3k + 2, respectively. In each of these cases the difference covering arrays satisfy **P1** and **P2** and so they can be used to construct MNOLS(2m).

When using a cyclic DCA(4, 2m + 1; 2m) to construct nearly orthogonal Latin squares we strip off the last row and last column of zeros. Thus to reduce the complexity of the notation and to avoid confusion, we will assume that we are constructing a  $2m \times 3$  array Q = [q(i, j)] that satisfies:

– each column is a permutation of  $\mathbb{Z}_{2m}$  and

 $-\Delta_{j,j'} = \{q(i,j) - q(i,j') \mid 0 \le i \le 2m - 1\} = \{1, 2, \dots, m, m, \dots, 2m - 1\},\$ with repetition retained.

Also we will use the following notation: q(a,0) = a (or  $q(\alpha,0) = a(\alpha)$ ), q(a,1) = b(a) (or  $q(\alpha,1) = b(\alpha)$ ), and q(a,2) = c(a) (or  $q(\alpha,2) = c(\alpha)$ ).

The following lemmas document some well known results, stated without proof, which will be used extensively in the proof of subsequent results.

**Lemma 3** For all integers x, y, z, gcd(x + yz, z) = gcd(x, z).

**Lemma 4** Let g and p be positive integers and h a non-negative integer. Working modulo 2p, if gcd(g, 2p) = 1 then

$$\{gx + h \mid 0 \le x \le 2p - 1\} = \mathbb{Z}_{2p},$$

or if gcd(g, 2p) = r and  $h \equiv s \mod r$  then

$$\{gx+h \mid 0 \le x \le 2p/r-1\} = \{rx+s \mid 0 \le x \le 2p/r-1\}.$$

3.1 Construction for general families DCA(4, 2m + 1; 2m) for some odd m

In this subsection we give a general construction for a difference covering array DCA(4, 2m + 1; 2m), for m odd. The proof that such a difference covering array exists uses the results presented in the following lemma. Note that in this section unless otherwise stated all arithmetic is modulo 2m. In particular, for  $i \neq 2 \mod 3$  a non-negative integer, and  $k = 2i^2 + 7i + 6$ , we present an infinite family of DCA(4, 2m + 1; 2m) for m = 2k + 1.

**Lemma 5** Let f and m be integers such that gcd(f, 2m) = 2, gcd(f+2, 2m) = 2, and  $f^2 + f + 1 \equiv m \mod 2m$ . Then

$$gcd(f,m) = 1, (2)$$

$$gcd(f+1,m) = 1, (3)$$

$$gcd(f-1,m) = 1, (4)$$

$$\gcd(2f+1,m) = 1\tag{5}$$

$$mf \equiv 0 \bmod 2m. \tag{6}$$

*Proof* Eq 2: Since f is even,  $f^2 + f + 1$  is odd, implying m is odd and hence gcd(f, m) = 1.

Eq 3: Since  $f + 1 \equiv -f^2 \mod m$  and gcd(f,m) = 1, we have  $1 = gcd(f,m) = gcd(f^2,m) = gcd(f+1,m)$ .

Eq 4: Since  $f-1 \equiv f(f+2) \mod m$ ,  $\gcd(f,m) = 1$  and  $\gcd(f+2,m) = 1$ , we have  $1 = \gcd(f(f+2),m) = \gcd(f-1,m)$ .

Eq 5: Since  $2f+1 \equiv -f(f-1) \mod m$ , gcd(f,m) = 1 and gcd(f-1,m) = 1, we have 1 = gcd(f(f-1),m) = gcd(2f+1,m).

Eq 6: This follows from the fact that f is even.

For a suitable choice of f, we partition the domain of a into the subintervals [0, m + f], [m + f + 1, m - 1], [m, m - f - 1] and [m - f, 2m - 1] where all endpoints are included. (Additions are all modulo 2m.)

Intervals	$I_1$	$I_2$	$I_3$	$I_4$
q(a,0) = a	[0, m + f]	[m+f+1, m-1]	[m, m - f - 1]	[m-f, 2m-1]
q(a,1) = b(a)	af + m	af + m	(a+1)f	(a+1)f
			+m - 1	+m - 1
q(a,2) = c(a)	-(a-1)(f+1)	-(a-1)(f+1)	-a(f+1)	-a(f+1)
	-2	+m-2	+m	
b(a) - a	a(f-1)	a(f-1)	(a+1)(f-1)	(a+1)(f-1)
b(a) - a	a(f-1) + m	a(f-1) +m	$\begin{array}{c} (a+1)(f-1) \\ +m \end{array}$	$\begin{array}{c} (a+1)(f-1) \\ +m \end{array}$
b(a) - a $c(a) - a$				
	+m	+m	+m	+m
	+m -(a - 1)(f + 2)	+m -(a-1)(f+2)	$\frac{+m}{-a(f+2)}$	+m

**Fig. 1** Entries are elements of  $\mathbb{Z}_{2m}$ , where q(a,0) = a, q(a,1) = b(a) and q(a,2) = c(a) in the array Q = [q(i,j)]. Rows 5 to 7 give the differences.

Example 1 To aid understanding we begin with an example constructing a DCA according to Figure 1 with m = 13 and f = 16. We give the transpose of the difference covering array DCA(4, 27; 26). The key to understanding the proof is to recognise that within the subintervals  $I_1 = [0,3], I_2 = [4,12], I_3 = [13,22]$  and  $I_4 = [23,25]$ , the value of a, b(a) and c(a) increases by a constant "jump", respectively 1, f = 16 and -(f + 1) = 9. This implies that the differences will also increase by a constant. By carefully choosing the start value on each subinterval it is possible to obtain the required values in  $\mathbb{Z}_{2m}$ . The value m = 13 is boldfaced in the differences.

		Ι	1						$I_2$									1	3						$I_4$	
				3																						
b(a)	13	3	19	9	25	15	5	21	11	1	17	$\overline{7}$	23	2	18	8	24	14	4	20	10	0	16	6	22	12
c(a)	15	24	$\overline{7}$	16	12	21	4	13	22	5	14	23	6	0	9	18	1	10	19	<b>2</b>	11	20	3	25	8	17
b(a) - a																										
c(a) - a	15	23	5	<b>13</b>	8	16	24	6	14	22	4	12	20	13	21	3	11	19	1	9	17	25	$\overline{7}$	2	10	18
c(a) - b(a)	2	21	14	$\overline{7}$	<b>13</b>	6	25	18	11	4	23	16	9	24	17	10	3	22	15	8	1	20	<b>13</b>	19	12	<b>5</b>

**Theorem 3** Let f and m be natural numbers such that gcd(f, 2m) = 2, gcd(f+2, 2m) = 2,  $f^2 + f + 1 \equiv m \mod 2m$ , and  $m+3 \leq f \leq 2m-4$ . Then a cyclic DCA(4, 2m+1; 2m) satisfying **P1** and **P2** exists.

*Proof* The proof is by construction with the values of DCA(4, 2m + 1; 2m) as given in Figure 1, with the 3 columns of Q = [q(i, j)] given by q(a, 0) = a, q(a, 1) = b(a), q(a, 2) = c(a). We will show that Q has the required properties.

If  $m + 3 \le f \le 2m - 4$  and f is even then, working in the least residue system modulo 2m, we have  $3 \le m + f \le m - 4$  and so the intervals [0, m + f]and [m + f + 1, m - 1] are non-empty. Further,  $m + 4 \le m - f \le 2m - 3$  and so the intervals [m, m - f - 1] and [m - f, 2m - 1] are non-empty.

Given that f is even and m is odd, then by Lemma 4

$$af + m \equiv 1 \mod 2 \Longrightarrow \{b(a) \mid a \in I_1 \cup I_2\} = \{2g + 1 \mid 0 \le g \le m - 1\},\$$

 $(a+1)f + m - 1 \equiv 0 \mod 2 \Longrightarrow \{b(a) \mid a \in I_3 \cup I_4\} = \{2g \mid 0 \le g \le m - 1\},\$ and so  $\{b(a) \mid 0 \le a \le 2m - 1\} = [2m].$  On each of the subintervals c(a) takes the form ag+h, where g = -(f+1), so the "jump" size is -(f+1) and

$$\begin{split} c(m) &= 0, \\ c(m+f+1) - c(m-f-1) &= -m - f(f+1) + m - 2 + m \\ &- f(f+1) - (f+1) - m \\ &= -2f^2 - 2f - 2 - (f+1) = -(f+1), \\ c(0) - c(m-1) &= -(f+1), \\ c(m-f) - c(m+f) &= -(f+1), \\ c(m) - c(2m-1) &= -(f+1). \end{split}$$

Thus by reordering the subintervals as  $I_3, I_2, I_1, I_4$ , and noting that c(m - f - 1) - (f + 1) = c(m + f + 1), we get  $\{c(a) \mid 0 \le a \le 2m - 1\} = [2m]$ .

For b(a)-a, c(a)-a and c(a)-b(a) we are required to show that for  $a \in [2m]$  the differences cover the multi-set  $\{1, 2, \ldots, m-1, m, m, m+1, \ldots, 2m-1\} = ([2m] \setminus \{0\}) \cup \{m\}.$ 

For b(a) - a, the subintervals are taken in natural order  $I_1, I_2, I_3, I_4$ . Starting at a = 0 and finishing at a = 2m - 1, we have b(0) - 0 = m = (2m - 1 + 1)(f - 1) + m = b(2m - 1) - (2m - 1), so the difference m occurs twice. Further, the gcd(f - 1, 2m) = 1 implies that  $a(f - 1) + m, 0 \le a \le m - 1$ , are all distinct, as are  $(a + 1)(f - 1) + m, m \le a \le 2m - 1$  and

$$b(m) - m = (m+1)(f-1) + m = f - 1,$$
  
$$b(m-1) - (m-1) = (m-1)(f-1) + m = -(f-1).$$

Thus there is a "jump" of -2(f-1) between a = m - 1 and a = m and the difference 0 is omitted, implying  $\{b(a) - a \mid 0 \le a \le 2m - 1\} = ([2m] \setminus \{0\}) \cup \{m\}$ .

For c(a) - a, since f + 2 is even and gcd(f + 2, m) = 1, these values are all distinct on each of the subintervals,  $|I_3 \cup I_1| = m + 1$ ,  $|I_4 \cup I_2| = m - 1$ , and

$$\begin{split} c(m) - m &= -m(f+2) + m = m, \\ c(m+f) - (m+f) &= -(m+f-1)(f+2) - 3 = m, \\ (c(0) - 0) - (c(m-f-1) - (m-f-1)) &= m - f(f+2) - 3 = -(f+2), \\ c(2m-1) - (2m-1) &= -(2m-1)(f+2) = f+2, \\ c(m+f+1) - (m+f+1) &= -f^2 - 2f + m - 3 = -(f+2). \end{split}$$

Thus -(a-1)(f+2) - 3,  $-a(f+2) + m \equiv 1 \mod 2$ , hence,

$$\{c(a) - a \mid a \in I_3 \cup I_1\} = \{2g + 1 \mid 0 \le g \le m - 1\} \cup \{m\}.$$

In addition,  $-a(f+2) + m - 1, -a(f+2) \equiv 0 \mod 2$  implies that

$$\{c(a) - a \mid a \in I_4 \cup I_2\} = \{2g \mid 1 \le g \le m - 1\},\$$

giving  $\{c(a) - a \mid 0 \le a \le 2m - 1\} = ([2m] \setminus \{0\}) \cup \{m\}.$ 

For c(a) - b(a), since gcd(2f + 1, 2m) = 1, these values are all distinct on the subintervals,  $c(m + f + 1) - b(m + f + 1) = -2f^2 - 2f - 2 + m = m = m + 2f^2 + 2f + 2 = c(m - f - 1) - b(m - f - 1)$ , and

$$c(0) - b(0) - (c(m-1) - b(m-1)) = -(2f+1),$$
  

$$c(m+f) - b(m+f) = 2f+1,$$
  

$$c(m-f) - b(m-f) = -(2f+1).$$

Thus when the subintervals are reordered to  $I_2, I_1, I_4, I_3$  we may verify that  $\{c(a) - b(a) \mid 0 \le a \le 2m - 1\} = ([2m] \setminus \{0\}) \cup \{m\}$ . Note that the values of c(a) - b(a) start and finish on m and the value 0 is omitted between a = m + f and a = m - f.

**Corollary 1** For every non-negative integer  $i \not\equiv 2 \mod 3$  there exists a cyclic DCA(4, 2m + 1; 2m) satisfying **P1** and **P2**, where  $m = 2(2i^2 + 7i + 6) + 1$ .

*Proof* Taking f = m + 3 + 2i, then f is even. In addition  $m + 3 \le f \le 2m - 4$ , since  $i \ge 0$  and  $2m - 4 = (m + 3) + (m - 7) = (m + 3) + (4i^2 + 14i + 6) \ge m + 3 + 2i = f$ . Now

$$f^{2} + f + 1 \equiv 4i^{2} + 14i + 13 \mod 2m$$
$$= 2(2i^{2} + 7i + 6) + 1 = m \mod 2m.$$

Further, applying Lemma 3 repeatedly,

$$gcd(f, 2m) = 2(gcd(2i^{2} + 8i + 8, (2i^{2} + 8i + 8) + 2i^{2} + 6i + 5)),$$
  
= 2(gcd(2i + 3, 2i^{2} + 4i + 2 + (2i + 3))),  
= 2(gcd(2i + 3, 2(i + 1)^{2})) = 2(gcd(2i + 3, i + 1)) = 2.

Also,

$$gcd(f+2,2m) = 2(gcd(2i^2+8i+9,(2i^2+8i+9)+2i^2+6i+4)),$$
  
= 2(gcd(2i+5,2(i+2)(i+1))),  
= 2(gcd(2i+5,(i+2)(i+1))).

Now gcd(2i+5, i+2) = gcd(2(i+2)+1, i+2) = 1. Whereas

$$gcd(2i+5, i+1) = gcd(2(i+1)+3, i+1)) = gcd(3, i+1)$$
  

$$\neq 1 \text{ when } i+1 \equiv 0 \mod 3 \text{ or equivalently } i \equiv 2 \mod 3.$$

Thus taking  $i \not\equiv 2 \mod 3$  we can construct a DCA(4, 2m + 1; 2m) as per Theorem 3.

Intervals	$I_1$	$I_2$	$I_3$	$I_4$
q(a,0) = a	[0, m-1]	[m, 2m - 1]	[2m, 3m-1]	[3m, 4m - 1]
q(a,1) = b(a)	(a+1)f - 1	(a+1)f - 1	af	af
q(a,2) = c(a)	(a+1)(2m-f+2)	a(2m-f+2)	(a+1)(2m-f+2)	a(2m-f+2)
	-1	-m	+m-1	
b(a) - a	(a+1)(f-1)	(a+1)(f-1)	a(f-1)	a(f-1)
c(a) - a	(a+1)(2m-f+1)	a(2m-f+1)	(a+1)(2m-f+1)	a(2m-f+1)
		-m	+m	
c(a) - b(a)	(a+1)(2m-2f+2)	a(2m-2f+2)	a(2m-2f+2)	a(2m-2f+2)
$\int c(u) = o(u)$	(a + 1)(2m 2j + 2)	(2 <i>inc</i> 2 <i>j</i> 12)		

Fig. 2 Entries are elements of  $\mathbb{Z}_{4m}$ , where q(a,0) = a, q(a,1) = b(a) and q(a,2) = c(a) in the array Q = [q(i,j)].

3.2 Construction of difference covering arrays DCA(4, 4m + 1; 4m)

In this subsection we give a general construction for a difference covering array DCA(4, 4m + 1; 4m). It will be shown that for all non-negative integers k, such that  $3 \nmid (2k + 1)$ , this construction gives an infinite family of DCA(4, 16k + 9; 16k + 8). The proof that such a difference covering array exists uses the results presented in the following lemma. Note that in this section unless otherwise stated all arithmetic is modulo 4m.

**Lemma 6** Let f and m be natural numbers such that  $m \equiv 2 \mod 4$ , gcd(f, 4m) = 2, gcd(f-1, 4m) = 1, and  $f^2 + f - 2 \equiv 2m \mod 4m$ . Then

$$\gcd(2m+2-f,4m) = 4,$$
(7)

$$gcd(2m - f + 1, 4m) = 1, (8)$$

$$\gcd(2m - 2f + 2, 4m) = 2, (9)$$

$$mf \equiv 2m \bmod 4m. \tag{10}$$

Proof Eq 7: Rewriting  $2m + 2 - f = f^2 + f - 2 + 2 - f = f^2 \mod 4m$  and assuming gcd(f, 4m) = 2 gives  $gcd(f^2, 4m) = 4$ .

Eq 8: Since 2m - f + 1 is odd, the gcd(2m - f + 1, 4m) is odd. Assume there exists an odd x such that x|4m and x|(2m - f + 1), then x|m and so x|(f - 1). But the gcd(f - 1, 4m) = 1, so x = 1.

Eq 9: Assume that there exists x such that x|(m - f + 1) and x|2m. Since m - f + 1 is odd, x is odd and so x|m. Consequently, x|(f-1) and x|4m, implying x = 1.

Eq 10:  $mf \equiv m(2m - f^2 + 2) = 2m^2 - mf^2 + 2m \equiv 2m \mod 4m.$ 

**Theorem 4** Let f be a natural number and m = 4k + 2, where k is a nonnegative integer, such that gcd(f, 4m) = 2, gcd(f-1, 4m) = 1, and  $f^2+f-2 \equiv 2m \mod 4m$ . Then a cyclic DCA(4, 4m + 1; 4m) satisfying **P1** and **P2** exists.

*Proof* The proof is by construction with the values of DCA(4, 4m + 1; 4m) as given in Figure 2, with the 3 columns of Q = [q(i, j)] given by q(a, 0) = a, q(a, 1) = b(a), q(a, 2) = c(a). We will show that Q has the required properties.

For b(a), since f is even, Lemma 4 implies that

$$(a+1)f - 1 \equiv 1 \mod 2 \Longrightarrow \{b(a) \mid a \in I_1 \cup I_2\} = \{2g+1 \mid 0 \le g \le 2m-1\}, \\ af \equiv 0 \mod 2 \Longrightarrow \{b(a) \mid a \in I_3 \cup I_4\} = \{2g \mid 0 \le g \le 2m-1\},$$

and  $\{b(a) \mid 0 \le a \le 4m - 1\} = \mathbb{Z}_{4m}$ . For c(a), since gcd(2m - f + 2, 4m) = 4, Lemma 4 implies that

$$\begin{aligned} (a+1)(2m-f+2)+m-1 &\equiv 1 \mod 4 \Longrightarrow \{c(a) \mid a \in I_3\} = \{4g+1 \mid 0 \le g \le m-1\}, \\ a(2m-f+2) &\equiv 0 \mod 4 \Longrightarrow \{c(a) \mid a \in I_4\} = \{4g \mid 0 \le g \le m-1\}, \\ (a+1)(2m-f+2)-1 &\equiv 3 \mod 4 \Longrightarrow \{c(a) \mid a \in I_1\} = \{4g+3 \mid 0 \le g \le m-1\}, \\ a(2m-f+2)-m &\equiv 2 \mod 4 \Longrightarrow \{c(a) \mid a \in I_2\} = \{4g+2 \mid 0 \le g \le m-1\}. \end{aligned}$$

Thus the set of values  $\{c(a) \mid a \in \mathbb{Z}_{4m}\} = \mathbb{Z}_{4m}$ .

For b(a) - a, c(a) - a and c(a) - b(a) we are required to show that for  $a \in \mathbb{Z}_{4m}$  the differences cover the multi-set  $\{1, 2, \ldots, 2m - 1, 2m, 2m, 2m + 1, \ldots, 4m - 1\} = (\mathbb{Z}_{4m} \setminus \{0\}) \cup \{2m\}.$ 

For b(a) - a, the gcd(f - 1, 4m) = 1, and

$$b(2m) - 2m = 2m(f - 1) = 2m,$$
  

$$b(2m - 1) - (2m - 1) = (2m - 1 + 1)f - 1 - (2m - 1) = (2m)(f - 1) = 2m,$$
  

$$b(4m - 1) - (4m - 1) = -(f - 1),$$
  

$$b(0) - 0 = f - 1.$$

So using a "jump" of f - 1 and ordering the subintervals as  $I_3, I_4, I_1, I_2$  we obtain the difference 2m twice and the difference 0 is omitted between a = 4m - 1 and a = 0 implying that  $\{b(a) - a \mid 0 \le a \le 4m - 1\} = (\mathbb{Z}_{4m} \setminus \{0\}) \cup \{2m\}.$ 

For c(a) - a, the gcd(2m - f - 1, 4m) = 1, and

$$\begin{split} c(m) - m &= m(2m - f + 1) - m = 2m, \\ c(3m - 1) - (3m - 1) &= (3m - 1 + 1)(2m - f + 1) + m \\ &= 2m, \\ c(3m) - 3m - (c(2m - 1) - (2m - 1)) &= (m + 1)(2m - f + 1) + m \\ &= 2m - f + 1, \\ c(0) - 0 &= 2m - f + 1, \\ c(4m - 1) - (4m - 1) &= (4m - 1)(2m + f - 1) \\ &= -(2m - f + 1), \\ c(2m) - 2m - (c(m - 1) - (m - 1)) &= (m + 1)(2m - f + 1) + m \\ &= 2m - f + 1. \end{split}$$

So using a "jump" of 2m - f + 1 and ordering the subintervals as  $I_2, I_4, I_1, I_3$  we obtain the difference 2m twice and the difference 0 is omitted between a = 4m - 1 and a = 0, implying  $\{c(a) - a \mid 0 \le a \le 4m - 1\} = (\mathbb{Z}_{4m} \setminus \{0\}) \cup \{2m\}$ .

For c(a) - b(a), and

$$\begin{split} c(3m)-b(3m) &= m(2m-2f+2) = 2m,\\ c(m-1)-b(m-1) &= (m-1+1)(2m-2+2) = 2m\\ c(4m-1)-b(4m-1) &= -(2m-2f+2),\\ c(0)-b(0) &= 2m-2f+2,\\ c(2m)-b(2m)-(c(2m-1)-b(2m-1)) &= 2m-2f+2. \end{split}$$

Then since

$$2m - 2f - 2 \equiv 0 \mod 2 \Longrightarrow \{c(a) - b(a) \mid a \in I_4 \cup I_1\} = \{2g \mid 1 \le g \le 2m - 1\} \cup \{2m\}, \\ -f + 1 \equiv 1 \mod 2 \Longrightarrow \{c(a) - b(a) \mid a \in I_2 \cup I_3\} = \{2g + 1 \mid 0 \le g \le 2m - 1\},$$

implying  $\{c(a) - b(a) \mid 0 \le a \le 4m - 1\} = (\mathbb{Z}_{4m} \setminus \{0\}) \cup \{2m\}.$ 

**Corollary 2** For  $k \ge 0$  such that  $k \not\equiv 1 \mod 3$  a cyclic DCA(4, 4m + 1; 4m), where m = 4k + 2, satisfying **P1** and **P2** can be constructed as described in Theorem 4.

Proof Given m = 4k + 2, take f = 2m - 2. Then f = 8k + 2, and

$$gcd(f, 4m) = 2(gcd(4k+1, 8k+4)) = 2(gcd(4k+1, 2(4k+1)+2)) = 2$$

In addition

$$gcd(f - 1, 4m) = gcd(2m - 3, 4m) = gcd(8k + 1, 16k + 8)$$
  
= gcd(8k + 1, 2k + 1) since 2\frac{1}{8k}(8k + 1)  
= gcd(6k, 2k + 1) = 1, if 3\frac{1}{2k}(2k + 1).

Also

$$f^{2} + f - 2 = 64k^{2} + 32k + 4 + 8k + 2 - 2 = 4k(16k + 8) + 8k + 4$$
  

$$\equiv 2m \mod 4m.$$

Hence, f = 2m-2 satisfies the assumptions of Theorem 4 and we can construct a DCA(4, 16k + 9; 16k + 8) for k such that  $3 \nmid (2k + 1)$ .

3.3 Construction of difference covering arrays  $\mathrm{DCA}(4,2m+1;2m),$  where  $m=3\mu+2$ 

In this subsection we give a general construction for a difference covering array DCA(4, 2m + 1; 2m), where  $m = 3\mu + 2$ . Note that in this section unless otherwise stated all arithmetic is modulo 12k + 10,  $k \ge 0$ .

**Theorem 5** Let  $\mu$  be an odd positive integer. Then there exists a cyclic  $DCA(4, 6\mu + 5; 6\mu + 4)$  satisfying **P1** and **P2**.

Intervals for	$I_1$	$I_2$	$I_3$
α	$[0, \mu - 1]$	$[\mu, 2\mu]$	$[2\mu + 1, 3\mu + 1]$
$q(\alpha, 0) = a(\alpha)$	$3\alpha + 3\mu + 4$	$3\alpha + 2$	$3\alpha + 3\mu + 4$
$q(\alpha, 1) = b(\alpha)$	$3\alpha(\mu+1) + 2\mu + 2$	$3\alpha(\mu+1) + 2\mu + 2$	$3\alpha(\mu+1) + 2\mu + 2$
$q(\alpha, 2) = c(\alpha)$	$\alpha(3\mu+4) + 5\mu + 4$	$\alpha(3\mu+4) + 5\mu + 4$	$\alpha(3\mu+4) + 5\mu + 4$
$b(\alpha) - a(\alpha)$	$3\alpha\mu - \mu - 2$	$3\alpha\mu + 2\mu$	$3\alpha\mu - \mu - 2$
$c(\alpha) - a(\alpha)$	$\alpha(3\mu+1)+2\mu$	$\alpha(3\mu+1) + 5\mu + 2$	$\alpha(3\mu+1)+2\mu$
$c(\alpha) - b(\alpha)$	$\alpha + 3\mu + 2$	$\alpha + 3\mu + 2$	$\alpha + 3\mu + 2$
Intervals for	Т	Т	T
Intervals for	$I_4$	$I_5$	$I_6$
$\alpha$	$[3\mu + 2, 4\mu + 2]$	$[4\mu + 3, 5\mu + 2]$	$[5\mu + 3, 6\mu + 3]$
			-
α	$[3\mu + 2, 4\mu + 2]$	$[4\mu + 3, 5\mu + 2]$	$[5\mu + 3, 6\mu + 3]$
$\frac{\alpha}{q(\alpha,0) = a(\alpha)}$	$\frac{[3\mu + 2, 4\mu + 2]}{3\alpha + 3\mu + 3}$	$\frac{[4\mu+3,5\mu+2]}{3\alpha+1}$	$\frac{[5\mu + 3, 6\mu + 3]}{3\alpha + 3\mu + 3}$
$\begin{aligned} \alpha \\ q(\alpha,0) &= a(\alpha) \\ q(\alpha,1) &= b(\alpha) \end{aligned}$	$\begin{array}{c} [3\mu + 2, 4\mu + 2] \\ \hline 3\alpha + 3\mu + 3 \\ \hline 3\alpha(\mu + 1) + 2\mu + 1 \end{array}$	$\frac{[4\mu + 3, 5\mu + 2]}{3\alpha + 1}$ $\frac{3\alpha(\mu + 1) + 2\mu + 1}{3\alpha(\mu + 1) + 2\mu + 1}$	$ \begin{array}{c} [5\mu + 3, 6\mu + 3] \\ \hline 3\alpha + 3\mu + 3 \\ \hline 3\alpha(\mu + 1) + 2\mu + 1 \end{array} $
$\begin{array}{c} \alpha \\ \hline q(\alpha,0) = a(\alpha) \\ \hline q(\alpha,1) = b(\alpha) \\ \hline q(\alpha,2) = c(\alpha) \end{array}$	$\begin{array}{c} [3\mu + 2, 4\mu + 2] \\ \hline 3\alpha + 3\mu + 3 \\ \hline 3\alpha(\mu + 1) + 2\mu + 1 \\ \hline \alpha(3\mu + 4) + 5\mu + 4 \end{array}$	$\begin{array}{c} [4\mu + 3, 5\mu + 2] \\ \hline 3\alpha + 1 \\ \hline 3\alpha(\mu + 1) + 2\mu + 1 \\ \hline \alpha(3\mu + 4) + 5\mu + 4 \end{array}$	$ \begin{array}{c} [5\mu + 3, 6\mu + 3] \\ \hline 3\alpha + 3\mu + 3 \\ \hline 3\alpha (\mu + 1) + 2\mu + 1 \\ \hline \alpha (3\mu + 4) + 5\mu + 4 \end{array} $

**Fig. 3** Entries are elements of  $\mathbb{Z}_{6\mu+4}$ , where  $q(\alpha, 0) = a(\alpha)$ ,  $q(\alpha, 1) = b(\alpha)$  and  $q(\alpha, 2) = c(\alpha)$  in the array Q = [q(i, j)].

Proof First note that  $n = 6\mu + 4 \ge 10$  and  $gcd(3\mu + 4, 6\mu + 4) = gcd(3\mu, 3\mu + 4) = gcd(3\mu, 4) = 1$  since  $\mu$  is odd.

Let  $\mu = 2k + 1, k \ge 0$ , then

$$\mu(3\mu + 2) = (2k+1)(6k+5) = 3\mu + 2$$
, and  
 $(3\mu + 2)^2 = 3\mu + 2$ .

That is,  $(n/2)^2 \equiv n/2 \mod n$ .

The proof is by construction with the values for DCA(4,  $6\mu + 5$ ;  $6\mu + 4$ ) as given in Figure 3, with the three columns of Q = [q(i, j)] given by  $q(\alpha, 0) = a(\alpha)$ ,  $q(\alpha, 1) = b(\alpha)$  and  $q(\alpha, 2) = c(\alpha)$ .

Since gcd(3, n) = 1, by Lemma 4,  $\{3\alpha \mid 0 \le \alpha \le n - 1\} = \mathbb{Z}_n$ . Further  $a(3\mu+2) = 3(3\mu+2) + 3\mu+3 = 1$  and  $a(2\mu) = 6\mu+2$  and there is a "jump" of 3 between  $a(4\mu+2)$  and a(0);  $a(\mu-1)$  and  $a(5\mu+3)$ ;  $a(6\mu+3)$  and  $a(2\mu+1)$ ;  $a(3\mu+1)$  and  $a(4\mu+3)$ ;  $a(5\mu+2)$  and  $a(\mu)$ .

Thus reordering the subintervals as  $I_4, I_1, I_6, I_3, I_5, I_2$  gives  $\{a(\alpha) \mid 0 \le \alpha \le n-1\} = \mathbb{Z}_n$ .

For  $b(\alpha)$ ,

$$\begin{aligned} 3\alpha(\mu+1) + 2\mu + 2 &\equiv 0 \mod 2 \Longrightarrow \{b(\alpha) \mid \alpha \in I_1 \cup I_2 \cup I_3\} \\ &= \{2g \mid 0 \le g \le 3\mu + 2\}, \\ 3\alpha(\mu+1) + 2\mu + 1 &\equiv 1 \mod 2 \Longrightarrow \{b(\alpha) \mid \alpha \in I_4 \cup I_5 \cup I_6\} \\ &= \{2g + 1 \mid 0 \le g \le 3\mu + 1\}. \end{aligned}$$

For  $c(\alpha)$ , since  $gcd(3\mu + 4, 6\mu + 4) = 1$  Lemma 4 implies that  $\{c(\alpha) \mid 0 \le \alpha \le 6\mu + 3\} = \mathbb{Z}_{6\mu+4}$ .

For  $b(\alpha) - a(\alpha)$ , since  $gcd(3\mu, 6\mu + 4) = 1$ ,

$$\begin{split} b(\mu)-a(\mu)&=3\mu+2\\ b(4\mu+2)-a(4\mu+2)&=3\mu+2\\ b(5\mu+3)-a(5\mu+3)-(b(2\mu)-a(2\mu))&=3\mu,\\ b(4\mu+3)-a(4\mu+3)-(b(\mu-1)-a(\mu-1))&=6\mu,\\ b(2\mu+1)-a(2\mu+1)-(b(5\mu+2)-a(5\mu+2))&=3\mu,\\ b(0)-a(0)-(b(6\mu+3)-a(6\mu+3)&=3\mu. \end{split}$$

So using a "jump" of  $3\mu$  and ordering the subintervals as  $I_2, I_6, I_1, I_5, I_3, I_4$ , we obtain the difference  $3\mu + 2$  twice and since there is  $6\mu$  between  $b(\alpha) - a(\alpha)$ for  $\alpha = \mu - 1$  and  $\alpha = 4\mu + 3$  the difference 0 is omitted implying that  $\{b(\alpha) - a(\alpha) \mid 0 \le \alpha \le 6\mu + 3\} = (\mathbb{Z}_n \setminus \{0\}) \cup \{n/2\}.$ 

For  $c(\alpha) - b(\alpha)$ , since  $gcd(3\mu + 1, 6\mu + 4) = 2$ , we have

$$\begin{split} c(5\mu+3)-a(5\mu+3)&=3\mu+2,\\ c(2\mu)-a(2\mu)&=3\mu+2,\\ c(3\mu+2)-a(3\mu+2)-(c(6\mu+3)-a(6\mu+3))&=-(3\mu+3),\\ c(\mu)-a(\mu)-(c(4\mu+2)-a(4\mu+2))&=-(3\mu+3),\\ c(2\mu+1)-a(2\mu+1)&=3\mu+1,\\ c(5\mu+2)-c(5\mu+2)&=3\mu+3,\\ c(0)-a(0)-(c(3\mu+1)-a(3\mu+1))&=-(3\mu+3),\\ c(\mu-1)-a(\mu-1)-(c(4\mu+3)-a(4\mu+3))&=-(3\mu+3). \end{split}$$

Reordering the intervals as  $I_6, I_4, I_2$  and  $I_3, I_1, I_5$  and using a regular "jump" of  $-(3\mu + 3)$ , with the jump of  $6\mu + 2$  between  $\alpha = 2\mu + 1$  and  $\alpha = 5\mu + 2$ , being the exception, we have the difference  $3\mu + 2$  twice and the difference 0 omitted, thus  $\{c(\alpha) - a(\alpha) \mid 0 \le \alpha \le 6\mu + 3\} = (\mathbb{Z}_{6\mu+4} \setminus \{0\}) \cup \{3\mu + 2\}$  with repetition retained.

For  $c(\alpha) - b(\alpha)$ , we note that

$$c(0) - b(0) = 3\mu + 2,$$
  
$$c(3\mu + 1) - b(3\mu + 1) = -1,$$

and so the values of  $c(\alpha) - a(\alpha)$  on the subinterval  $I_1 \cup I_2 \cup I_3$  cover the set  $\{3\mu + 2, \ldots, -1\}$ . Also

$$\begin{split} c(3\mu+2) &- b(3\mu+2) = 1, \\ c(6\mu+3) &- b(6\mu+3) = 3\mu+2, \end{split}$$

and so the values of  $b(\alpha) - c(\alpha)$  on the subinterval  $I_4 \cup I_5 \cup I_6$  cover the set  $\{3\mu + 2, \ldots, 6\mu - 1\}$ . Consequently  $\{b(\alpha) - c(\alpha) \mid 0 \le \alpha \le 6\mu + 3\} = ([6\mu + 4] \setminus \{0\}) \cup \{3\mu + 2\}$  with repetition retained.

3.4 Infinite families

The construction of Theorem 5 constructs sets of three MNOLS of orders 10, 22, 34, 46 mod 48. The construction of Corollary 2 constructs sets of three MNOLS of orders 8, 40 mod 48. Combined with the constructions of [6], there is a construction of three MNOLS for 8, 10, 14, 22, 34, 38, 40, 46 mod 48. There are infinite families constructed from Corollary 1 and from results of Li and van Rees [10], but these cannot be described mod 48.

## 4 The spectrum for sets of 3 mutually nearly orthogonal Latin squares

In 2007, Li and van Rees [10] continued the study of 3-MNOLS(n) conjecturing that they exist for all even  $n \ge 6$ . Here we make a slightly stronger conjecture.

Conjecture 1 There exists DCA(4, 2p + 1; 2p) satisfying **P1** and **P2** for all positive integers  $p \ge 3$ .

In a partial solution, Li and van Rees proved the existence for orders less than 22 and orders greater than 356, (see also [11]).

**Theorem 6** [10, Thm 4.8] If  $2p \ge 358$ , then there exists a 3-MNOLS(2p).

This work was extended in 2014, when Demirkale, Donovan and Khodkar [6] developed further constructions for cyclic DCA(4, 2p + 1; 2p) proving:

**Theorem 7** [6] There exist 3-MNOLS(2p), where  $2p \equiv 14, 22, 38, 46 \mod 48$ .

The next result lists known values for cyclic DCA(4, 2p + 1; 2p) satisfying **P1** and **P2**, with  $2p \leq 356$ .

**Lemma 7** There exists cyclic DCA(4, 2p+1; 2p) for  $2p \in \{6, 8, \dots, 20, 22, 38, 46, 62, 70, 86, 94, 110, 118, 134, 142, 158, 166, 182, 190, 206, 214, 230, 238, 254, 262, 278, 286, 302, 310, 326, 334, 350\}.$ 

*Proof* The existence of orders 6 and 8 was given in [12] and orders 10, 12, 14, 16, 18, 20 in [10]. All the remaining cases were shown to exist in [6].

van Rees recently summarised these results and indicated that 3-MNOLS(2p) exist for all orders except possibly those given below.

**Lemma 8** [14] A set of 3-MNOLS(2p) exists except possibly when  $2p \in \{24, 26, 28, 30, 34, 36, 42, 50, 52, 54, 58, 66, 74, 82, 92, 102, 106, 114, 116, 122, 124, 130, 138, 146, 148, 170, 172, 174, 178\}.$ 

In this section we show that 3-MNOLS(2*p*) exist for all even orders except possibly 2p = 146. For completeness we list all values less than 2p = 358 and, given the connection with row complete Latin squares, where possible we will use cyclic difference covering arrays to construct the Latin squares. We begin this section by reviewing relevant results from [7], [16] and [17], and adapting these to construct cyclic DCA(4, 2m + 1; 2m) that satisfy **P1** and **P2**.

We begin with the following straightforward result that is analogous to [7, Lem 2.3].

**Lemma 9** Suppose that there exists a HDM(k, n; h) over the group (G, +) with a hole over the subgroup H. Further suppose there exists a DCA(k, h+1; h) over H satisfying **P1** and **P2**. Then there exists a DCA(k, n + 1; n) over G satisfying **P1** and **P2**. Further suppose that the HDM(k, n; h) and DCA(k, h + 1; h) are cyclic. Then there exists a cyclic DCA(k, n + 1; n) satisfying **P1** and **P2**.

Proof In the cyclic case, let A = [a(i, j)]  $(0 \le i \le n - 1 - h, 0 \le j \le k - 1)$ represent the cyclic HDM(k, n; h) and B = [b(i, j)]  $(0 \le i \le h, 0 \le j \le k - 1)$ represent the cyclic DCA(k, h + 1; h). The definition of cyclic implies that  $H = \{0, u, 2u, \dots (h - 1)u\}$ , where n = uh, and Lemma 1 implies that h is even and the repeated difference in  $\Delta_{j,j'}, j \ne k - 1 \ne j'$ , of B is hu/2 = n/2.

Set Q = [q(i, j)]  $(0 \le i \le n, 0 \le j \le k)$  to be the concatenation of A with an isomorphic copy of B and we obtain a cyclic DCA(k, n+1; n) that satisfies **P1** and **P2**.

The non-cyclic case follows similarly.

Next we give a general product type construction taken from [7] and adapt it to construct cyclic difference covering arrays that satisfy **P1** and **P2**.

**Lemma 10** [7, Lem 2.6] If both a cyclic HDM(k, n; h) and a cyclic DM(n', k; 1) exist, then so does a cyclic HDM(k, nn'; hn'). In particular, if there exists a cyclic DM(n, k; 1) and a cyclic DM(n', k; 1) then there exists cyclic DM(nn', k; 1).

The first statement of Lemma 10 coupled with Lemma 9 leads to the following straightforward result.

**Corollary 3** Suppose that there exists a cyclic HDM(k, n; h), a cyclic DM(n', k; 1) and a cyclic DCA(k, hn' + 1; hn') that satisfies **P1** and **P2**. Then there exists a cyclic DCA(k, nn' + 1; nn') that satisfies **P1** and **P2**.

The second statement of Lemma 10 can also be adapted.

**Lemma 11** Suppose a cyclic DM(n, k; 1), a cyclic DM(n', k; 1) and a cyclic DCA(k, n' + 1; n') satisfying **P1** and **P2** exist. Then there exists a cyclic DCA(k, nn' + 1; nn') that satisfies **P1** and **P2**.

*Proof* This result can be obtained by taking a hole of size 1 in Corollary 3.

This result can be generalised to construct non-cyclic difference covering arrays as in [16].

We now combine these results with various results of [5], [7] and [17] to obtain results for cyclic DCA(4, 2p + 1; 2p), satisfying **P1** and **P2**, implying new existence results for row complete 3-MNOLS(2p), where 2p < 358. In the lists below \* indicates that the existence of 3-MNOLS(2p) was previously unknown.

In doing this we will settle the remaining cases in Lemma 8, with the exception of 2p = 146. Here we believe that there exists a DCA(4, 147; 146) satisfying **P1** and **P2** but have been unable to verify it.

**Theorem 8** [5, Thm 17.6,  $p \downarrow 11$ ] If n is a prime greater than or equal to k, then there exists a cyclic DM(n, k; 1).

**Lemma 12** [17, Lem 2.5] Let  $n \ge 5$  be prime. Then there exists a cyclic HDM(4, 2n; 2).

**Lemma 13** There exists a cyclic DCA(4, 2p + 1; 2p) for  $2p = 50^*$ , 98, 170<sup>\*</sup>, 242, 290, 338.

*Proof* Corollary 3 together with Theorem 8 and Lemma 12 can be used first to construct a cyclic HDM(4, 2p; h) with 2p(h) = 50(10), 98(14), 170(10), 242(22), 290(10), 338(26) and then the required DCA(4, 2p + 1; 2p).

Lemmas 15, 16, 17 do not document any new existence results, however they do verify that for the given orders cyclic DCA(4, 2p+1; 2p) satisfying **P1** and **P2** exist.

**Lemma 14** [17, Thm 2.1] Let n be an odd positive integer satisfying  $gcd(n,9) \neq 3$ . Then there exists a cyclic HDM(4, 2n; 2).

**Lemma 15** There exists a cyclic DCA(4, 2p + 1; 2p) for 2p = 90, 126, 198, 234, 306, 342.

*Proof* Corollary 3 together with Theorem 8 and Lemma 14 can be used first to construct a cyclic HDM(4, 2p; h) with 2p(h) = 90(10), 126(14), 198(22), 234(26), 306(34), 342(38) and then the required DCA(4, 2p + 1; 2p).

**Theorem 9** [17, Thm 2.3] Let  $n \ge 4$  and  $n = 2^{\alpha}3^{\beta}p$ , where p is coprime to 6 and  $(\alpha, \beta) \ne (1, 0)$ . Then there exists a cyclic HDM(4, 2n; 2).

**Lemma 16** There exists cyclic DCA(4, 2p + 1; 2p) for 2p = 60, 80, 84, 100, 112, 120, 132, 156, 160, 168, 176, 180, 204, 208, 224, 228, 240, 252, 264, 272, 276, 300, 304, 312, 320, 336, 352.

*Proof* Corollary 3 together with Theorem 8 and Theorem 9 can be used first to construct a cyclic HDM(4, 2p; h) with 2p(h) = 60(10), 80(10), 84(14), 100(10), 112(14), 120(10), 132(22), 156(26), 160(10), 168(14), 176(22), 180(10), 204(34), 208(26), 224(14), 228(38), 240(10), 252(14), 264(22), 272(34), 276(46), 300(10), 304(38), 312(26), 320(10), 336(14), 352(22) and then the required DCA(4, 2p + 1; 2p).

**Theorem 10** [17, Thm 2.4] Let n be coprime to 6. Then there exists a cyclic HDM(4, 4n; 4).

**Lemma 17** There exists a cyclic DCA(4, 2p + 1; 2p) for 2p = 140, 196, 220, 260, 308, 340.

*Proof* Corollary 3 together with Theorem 8 and Theorem 10 can be used first to construct a cyclic HDM(4, 2p; h) with 2p(h) = 140(28), 196(28), 220(20), 260(20), 308(28), 340(20) and then the required DCA(4, 2p + 1; 2p).

**Theorem 11** [7, Thm 3.10] A cyclic  $DM(3^i, 5; 1)$  exists for all  $i \ge 3$ .

**Lemma 18** There exists a cyclic DCA(4, 2p + 1; 2p) for 2p = 216, 270, 324.

*Proof* Corollary 3 together with Theorem 11 and Theorem 9 can be used to first construct a cyclic HDM(4, 2p; h) with 2p(h) = 216(54), 270(54), 324(54) and then the required DCA(4, 2p + 1; 2p).

The next result from Yin's paper [17] is interesting in that it allows us to construct difference covering arrays and so mutually nearly orthogonal Latin squares of order 6n, and gives many values that were previously unresolved (the obstruction was the non-existence of MOLS of order 6).

**Theorem 12** [17, Thm 2.2] Let n be coprime to 6. Then there exists a cyclic HDM(4, 6n; 6).

**Corollary 4** Let n be an integer of the form  $p_1^{\alpha_1} p_2^{\alpha_2} \dots p_t^{\alpha_t}$ , where  $\alpha_i \ge 0$  and the prime factors  $p_i \ge 5$  for  $1 \le i \le t$ . Then there exists a cyclic DCA(4, 6n + 1; 6n) satisfying **P1** and **P2**. Consequently, there exists a cyclic DCA(4, 2p + 1; 2p) for  $2p = 30^*, 42^*, 66^*, 78, 102^*, 114^*, 138^*, 150, 174^*, 186, 210, 222, 246, 258, 282, 294, 318, 330, 354.$ 

The following result can be verified using direct constructions given in Section 3.

**Lemma 19** There exists a cyclic DCA(4, 2p + 1; 2p) for  $2p = 26^*$ , 266,  $2p = 40, 56, 88, 104, 136, 152, 184, 200, 232, 248, 280, 296, 328, 344 and <math>2p = 34^*$ ,  $58^*$ ,  $82^*$ ,  $106^*$ ,  $130^*$ , 154,  $178^*$ , 202, 226, 250, 274, 298, 322, 346.

*Proof* Corollary 1 verifies the existence of cyclic DCA(4, 2p + 1; 2p) with  $2p = 26^*, 266$ .

Corollary 2 verifies the existence of cyclic DCA(4, 2p+1; 2p) with 2p = 40, 56, 88, 104, 136, 152, 184, 200, 232, 248, 280, 296, 328, 344.

Theorem 5 verifies the existence of cyclic DCA(4, 2p + 1; 2p) with  $2p = 34^*, 58^*, 82^*, 106^*, 130^*, 154, 178^*, 202, 226, 250, 274, 298, 322, 346.$ 

**Lemma 20** There exists a cyclic DCA(4, 2p + 1; 2p) for  $2p = 24^*, 28^*, 32, 36^*, 44, 48, 52^*, 54^*$ .

*Proof* These results have been verified by computer searches. The first column of the DCA(4, 2p + 1, 2p) is given by  $[0, 1, 2, \ldots, 2p - 1]$ , the second column by  $[1, 3, \ldots, 2p - 1, 2, 4, \ldots, 2p - 2]$  and the third column by

 $2p=24 . \ [2, \ 0, \ 3, \ 1, \ 14, \ 21, \ 20, \ 19, \ 23, \ 15, \ 6, \ 18, \ 16, \ 10, \ 17, \ 8, \ 11, \ 22, \ 5, \ 13, \ 4, \ 9, \ 7, \ 12]$ 

2p = 28: [2, 0, 3, 1, 11, 16, 22, 25, 20, 23, 4, 8, 21, 5, 18, 10, 19, 13, 24, 27, 7, 26, 15, 9, 6, 14, 17, 12]

 $2p=32; \, [2,\,0,\,3,\,6,\,1,\,13,\,22,\,30,\,21,\,25,\,28,\,26,\,7,\,5,\,23,\,20,\,12,\,10,\,24,\,17,\,31,\,15,\,29,\,27,\,11,\,14,\,4,\,9,\,8,\,19,\,18,\,16]$ 

2p = 36: [5, 35, 13, 20, 11, 9, 1, 31, 10, 2, 30, 33, 4, 34, 32, 25, 28, 16, 27, 22, 3, 29, 19, 24, 18, 15, 6, 23, 17, 7, 0, 8, 14, 12, 21, 26]

2p = 44: [39, 13, 26, 21, 35, 3, 17, 16, 40, 28, 38, 25, 6, 10, 34, 5, 18, 30, 43, 15, 19, 36, 7, 24, 32, 14, 4, 0, 31, 12, 2, 9, 23, 37, 11, 42, 41, 29, 20, 1, 33, 27, 8, 22]

2p = 48: [5, 41, 23, 40, 1, 39, 34, 25, 28, 8, 4, 9, 21, 30, 43, 18, 12, 2, 42, 45, 32, 37, 33, 0, 26, 15, 13, 22, 10, 35, 44, 7, 36, 16, 27, 19, 46, 38, 3, 47, 31, 29, 17, 14, 11, 24, 20, 6]

 $\begin{array}{l} 2p=52; \ [18,\ 12,\ 50,\ 37,\ 16,\ 6,\ 45,\ 4,\ 31,\ 34,\ 47,\ 21,\ 29,\ 2,\ 5,\ 22,\ 38,\ 3,\ 39,\ 27,\\ 0,\ 15,\ 51,\ 7,\ 28,\ 24,\ 42,\ 40,\ 48,\ 32,\ 9,\ 26,\ 20,\ 11,\ 1,\ 41,\ 19,\ 35,\ 43,\ 13,\ 49,\ 33,\\ 14,\ 17,\ 46,\ 8,\ 36,\ 23,\ 10,\ 30,\ 25,\ 44] \end{array}$ 

 $\begin{array}{l} 2p=54; \ [6,\ 5,\ 31,\ 27,\ 20,\ 38,\ 19,\ 4,\ 30,\ 51,\ 3,\ 52,\ 49,\ 14,\ 48,\ 23,\ 41,\ 12,\ 25,\ 0,\\ 32,\ 40,\ 21,\ 50,\ 9,\ 45,\ 16,\ 1,\ 46,\ 11,\ 28,\ 42,\ 47,\ 35,\ 39,\ 2,\ 22,\ 13,\ 34,\ 33,\ 24,\ 44,\\ 15,\ 53,\ 7,\ 17,\ 37,\ 36,\ 26,\ 18,\ 10,\ 43,\ 29,\ 8]. \end{array}$ 

For the remaining values 64, 68, 72, 74, 76, 92, 96, 108, 116, 122, 124, 128, 144, 146, 148, 162, 164, 172, 188, 192, 194, 212, 218, 236, 244, 256, 268, 284, 288, 292, 314, 316, 332, 348, 356 we were unable to construct cyclic difference covering arrays however for completeness and to answer questions about the spectrum we give full details verifying existence. It should be noted that it is possible to construct DCA(4, 2p + 1; 2p) satisfying **P1** and **P2** for some of these orders however our construction does not give cyclic difference covering arrays and so the details have been omitted here.

To verify existence of certain sizes of DCA we require two more results. The second result uses group divisible designs: A  $\mathcal{K}$ -group divisible design of type  $g_1^{a_1}g_2^{a_2}\ldots g_s^{a_s}$  is a partition  $\mathcal{G}$  of a finite set  $\mathcal{V}$ , of cardinality  $v = \sum_{i=1}^s a_i g_i$ , into  $a_i$  groups of size  $g_i$ ,  $1 \leq i \leq s$ , together with a family of subsets (blocks)  $\mathcal{B}$  of  $\mathcal{V}$  such that: 1) if  $B \in \mathcal{B}$ , then  $|B| \in \mathcal{K}$ , 2) every pair of distinct elements of  $\mathcal{V}$  occurs in 1 block of  $\mathcal{B}$  or 1 group of  $\mathcal{G}$  but not both, and 3)  $|\mathcal{G}| > 1$ .

**Theorem 13** [10, Thm 4.1] Suppose there exists a k-MNOLS(2p), a k-MOLS(2p), and a k-MOLS(n). Then there exists a k-MNOLS(2pn).

**Theorem 14** [10, Thm 4.5] Suppose there exists a  $\mathcal{K}$ -GDD of type  $g_1^{a_1} \dots g_s^{a_s}$ . Further suppose that for any group size  $g_i$  there exists a s-MNOLS $(g_i)$  and for any block size  $k \in \mathcal{K}$  there exists a s-IMOLS(k). Then there are s-MNOLS(t), where  $t = \sum_{i=1}^{s} a_i g_i$ . **Lemma 21** There exists a 3-NMOLS(2p) for  $2p = 76, 92^*, 96, 108, 116^*, 124^*, 128, 144, 148^*, 164, 172^*, 188, 192, 212, 236, 244, 256, 268, 284, 288, 292, 316, 332, 348 and 356.$ 

 $\begin{array}{l} Proof \mbox{ The 3-MNOLS}(2p), 2p = 76(12^{5}16^{1}), 92(20^{4}12), 96(16^{5}16^{1}), 108(20^{5}8^{1}), \\ 116(20^{5}, 16), 124(20^{5}24^{1}), 128(24^{5}8^{1}), 144(24^{5}24^{1}), 148(24^{5}28), 164(28^{5}24^{1}), \\ 172(32^{5}12), 188(32^{5}28^{1}), 192(32^{5}36^{1}), 212(36^{5}32^{1}), 236(40^{5}36^{1}), 244(40^{5}44^{1}), \\ 256(44^{5}36^{1}), 268(44^{5}48^{1}), 284(48^{5}44^{1}), 288(48^{5}48^{1}), 292(48^{5}52^{1}), 316(52^{5}56^{1}), \\ 332(56^{5}52^{1}), 348(56^{5}68^{1}) \mbox{ and } 356(60^{5}56^{1}) \mbox{ constructed applying 5-GDD}, \\ that exist by [5, \mbox{ Thm 4.17, p 258], in Theorem 14. Here the bracketed information gives the type of the GDD. \end{array}$ 

**Lemma 22** There exists a 3-NMOLS(2p) for  $2p = 64, 68, 72, 74^*, 122, 162, 194, 218 and 314.$ 

*Proof* 3-NMOLS(2*p*) for 2*p* = 64, 68, 72, 74<sup>\*</sup>, 122, 162, 194 and 218 can be constructed by applying, respectively, 8-GDD(8<sup>8</sup>), {7,8,9}-GDD(8<sup>7</sup>6<sup>2</sup>), 9-GDD(8<sup>9</sup>), {7,8,9}-GDD(8<sup>7</sup>6<sup>3</sup>), {7,8}-GDD(16<sup>7</sup>10<sup>1</sup>), {11, 12, 13}-GDD(12<sup>12</sup> 10<sup>1</sup>8<sup>1</sup>), {11, 12, 13}-GDD(16<sup>11</sup>10<sup>18</sup>), {13, 14}-GDD(16<sup>13</sup>10<sup>1</sup>) and {8, 9, 10, 11}-GDD(32<sup>8</sup>14<sup>2</sup>10<sup>1</sup>) in Theorem 14. All of the required group divisible designs exist by the standard finite field constructions for MOLS.

All the results of this paper combine to the following theorem.

**Theorem 15** There exists a 3-MNOLS(2p) for each positive integers  $p \ge 3$ , except possibly p = 73.

This leaves us tantalisingly close to Conjecture 1.

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