





University of Nottingham Ningbo China, 199 Taikang East Road, Ningbo, 315100, China

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Dynamic Random Testing of Web Services: A Methodology and Evaluation

Chang-ai Sun[®], *Senior Member, IEEE*, Hepeng Dai, Guan Wang, Dave Towey[®], *Member, IEEE*, Tsong Yueh Chen[®], *Member, IEEE*, and Kai-Yuan Cai

Abstract—In recent years, service oriented architecture (SOA) has been increasingly adopted to develop distributed applications in the context of the Internet. To develop reliable SOA-based applications, an important issue is how to ensure the quality of web services. In this article, we propose a dynamic random testing (DRT) technique for web services, which is an improvement over the widely-practiced random testing (RT) and partition testing (PT) approaches. We examine key issues when adapting DRT to the context of SOA, including a framework, guidelines for parameter settings, and a prototype for such an adaptation. Empirical studies are reported where DRT is used to test three real-life web services, and mutation analysis is employed to measure the effectiveness. Our experimental results show that, compared with the three baseline techniques, RT, Adaptive Testing (AT) and Random Partition Testing (RPT), DRT demonstrates higher fault-detection effectiveness with a lower test case selection overhead. Furthermore, the theoretical guidelines of parameter setting for DRT are confirmed to be effective. The proposed DRT and the prototype provide an effective and efficient approach for testing web services.

Index Terms—Software testing, random testing, dynamic random testing, web service, service oriented architecture, software cybernetics

16 **1** INTRODUCTION

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CERVICE oriented architecture (SOA) [1] defines a loosely cou-17 **D**pled, standards-based, service-oriented application devel-18 opment paradigm in the context of the Internet. Within SOA, 19 three key roles are defined: service providers (who develop 20 21 and own services); service requestors (who consume or invoke services); and a service registry (that registers services from 22 23 providers and returns services to requestors). Applications are built upon services that present functionalities through pub-24 lishing their interfaces in appropriate repositories, abstracting 25 away from the underlying implementation. Published interfa-26 27 ces may be searched by other services or users, and then invoked. Web services are the realization of SOA based on 28 open standards and infrastructures [2]. Ensuring the reliability 29 of SOA-based applications can become critical when such 30 applications implement important business processes. 31

Software testing is a practical method for ensuring the quality and reliability of software. However, some SOA features can pose challenges for the testing of web services [3],

- E-mail: casun@ustb.edu.cn, daihepeng@sina.cn, jayjaywg@qq.com. D. Towey is with the School of Computer Science, University of Nottingham
- Ningbo China, Ningbo 315100, China. E-mail: dave.towey@nottingham.edu.cn.
- T.Y. Chen is with the Department of Computer Science and Software Engineering, Swinburne University of Technology, Hawthorn, VIC 3122, Australia. E-mail: tychen@swin.edu.au.
- K.-Y. Cai is with the School of Automation Science and Electrical Engineering, Beihang University, Beijing 100191, China. E-mail: kycai@buaa.edu.cn.

Manuscript received 22 Jan. 2019; revised 26 Oct. 2019; accepted 10 Dec. 2019. Date of publication 0 . 0000; date of current version 0 . 0000. (Corresponding author: Chang-ai Sun.) Digital Object Identifier no. 10.1109/TSC.2019.2960496 [4]. For instance, service requestors often do not have access 35 to the source code of web services which are published and 36 owned by another organization, and, consequently, it is not 37 possible to use white-box testing techniques. Testers may, 38 therefore, naturally turn to black-box testing techniques. 39

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Random Testing (RT) [5] is one of the most widely- 40 practiced black-box testing techniques. Because test cases in 41 RT are randomly selected from the input domain (which 42 refers to the set of all possible inputs of the software under 43 test), it can be easy to implement. Nevertheless, because RT 44 does not make use of any information about the software 45 under test (SUT), or the test history, it may be inefficient in 46 some situations. In recent years, many efforts have been made 47 to improve RT in different ways [6], [7], [8]. Adaptive random 48 testing (ART) [7], [9], for example, has been proposed to 49 improve RT by attempting to have a more diverse distribution 50 of test cases in the input domain.

In contrast to RT, partition testing (PT) attempts to generate 52 test cases in a more "systematic" way, aiming, to use fewer test 53 cases to reveal more faults. When conducting PT, the input 54 domain of the SUT is divided into disjoint partitions, with test 55 cases then selected from each and every one. Each partition is 56 expected to have a certain degree of homogeneity—test cases 57 in the same partition should have similar software execution 58 behavior. Ideally, a partition should also be homogeneous in 59 fault detection: If one input can reveal a fault, then all other 60 inputs in the same partition should also be able to reveal 61 a fault.

RT and PT are based on different intuitions, and each 63 have their own advantages and disadvantages. Because it is 64 likely that they can be complementary to each other, detect- 65 ing different faults, it is intuitively appealing to investigate 66 their integration in random partition testing (RPT). 67

[•] C.-A. Sun, H. Dai, and G. Wang are with the School of Computer and Communication Engineering, University of Science and Technology Beijing, Beijing 100083, China.

In traditional RPT [6], the partitions and corresponding 68 test profiles remain constant throughout testing, which may 69 not be the best strategy. Independent researchers have 70 observed that fault-revealing inputs tend to cluster into 71 "continuous regions" [10], [11]-there is similarity in the 72 execution behavior of neighboring software inputs. Based on 73 74 software cybernetics that explores the interplay between software engineering and control theory, Cai et al. proposed adap-75 tive testing (AT) to control the testing process [12], however, 76 AT's decision-making incurs an extra computational over-77 head. To alleviate this, dynamic random testing (DRT) [6] was 78 proposed by Cai et al., aiming to improve on both RT and RPT. 79 In practice, web services have usually been tested by the 80 service providers, and simple or easy-to-test faults have been 81

removed, meaning that the remaining faults are normally
hard to detect. For ensuring a higher reliability of the web services, a simple RT strategy may not be appropriate [13], especially when the scale is large, or there are some stubborn faults.

In this paper, we present a DRT approach for web serv-86 87 ices, as an enhanced version of RT adapted to the context of SOA. We examine key issues of such an adaptation, and, 88 accordingly, propose a framework for testing web services 89 that combines the principles of DRT [6] and the features of 90 web services. To validate the fault detection effectiveness 91 and efficiency of the proposed DRT method in the context 92 of SOA, we conduct a comprehensive empirical study. We 93 also explore the impact factors of the proposed DRT, and 94 provide guidelines for setting DRT parameters based on a 95 theoretical analysis. Finally, we compare the performance 96 of the proposed DRT with other baseline techniques. 97

98 This paper extends our previous work [14] in the following aspects. First, this paper extensively examines the challenges 99 100 and practical situations related to testing web services (Section 2.2). It also extensively discusses the limitations of RT, 101 102 PT, RPT, and AT, when they are used for testing web services (Section 1). Second, although previous work [14] provided a 103 coarse-grained framework for DRT of web services, PT was 104 not studied. In contrast, this paper provides a comprehensive 105 solution based on partitioning (Section 4.4.1). Third, based on a 106 theoretical analysis (Section 3.2), this paper provides guidelines 107 for setting DRT parameters. Such guidelines are crucial to 108 enhance the practical application of DRT, which was not cov-109 ered in previous work [14]. Fourth, previous work [14] only 110 111 evaluated the fault detection effectiveness and efficiency of the proposed approach (DRT) in terms of the F-measure and T-112 113 measure, and only two small web services (ATM Service and Warehouse Service) were used in the evaluation of its perfor-114 mance. This paper, in contrast, provides a comprehensive eval-115 uation that not only evaluates the fault detection effectiveness 116 of the proposed approach in terms of the F-measure, F2-mea-117 118 sure, and T-measure (Section 5.1), but also evaluates its efficiency in terms of F-time, F2-time, and T-time (Section 5.3). 119 Furthermore, we also examine three real-life web services, 120 comparing the fault-detection effectiveness and efficiency of 121 122 the proposed approach with those of RT, RPT, and AT. Statistical analysis is used to validate the significance of the empirical 123 evaluations and comparisons (Sections 5.1 and 5.3), which was 124 not covered in previous work [14]. Extending again the previ-125 ous work [14], we also examine the relationship between the 126 number of partitions and the optimal control parameter set-127 tings for DRT, evaluating the usefulness of guidelines provided 128

by the theoretical analysis (Section 5.2). The contributions of 129 this work, combined with previous work [14], include: 130

- We develop an effective and efficient testing method 131 for web services. This includes a DRT framework that 132 addresses key issues for testing web services, and a 133 prototype that partly automates the framework. 134
- We evaluate the performance of DRT through a series 135 of empirical studies on three real web services. These 136 studies show that DRT has significantly higher fault-137 detection efficiency than RT and RPT. That is, to detect 138 a given number of faults, DRT uses less time and fewer 139 test cases than RT and RPT.
- We provide guidelines for the DRT parameter set- 141 tings, supported by theoretical analysis, and vali- 142 dated by the empirical studies. 143

The rest of this paper is organized as follows. Section 2 144 introduces the underlying concepts for DRT, and web serv-145 ices. Section 3 presents the DRT framework for web services, 146 guidelines for its parameter settings, and a prototype that par-147 tially automates DRT. Section 4 describes an empirical study 148 where the proposed DRT is used to test three real-life web 149 services, the results of which are summarized in Section 5. 150 Section 6 discusses related work and Section 7 concludes the 151 paper. 152

2 BACKGROUND

In this section, we present some of the underlying concepts 154 for DRT, and web services. 155

2.1 Dynamic Random Testing (DRT)

DRT combines RT and PT, with the goal of benefitting from 157 the advantages of both. Given a test suite *TS* classified into 158 *m* partitions (denoted s_1, s_2, \ldots, s_m), suppose that a test case 159 from s_i $(i = 1, 2, \ldots, m)$ is selected and executed. If this test 160 case reveals a fault, $\forall j = 1, 2, \ldots, m$ and $j \neq i$, we then set 161

$$p'_{j} = \begin{cases} p_{j} - \frac{\varepsilon}{m-1} & \text{if } p_{j} \ge \frac{\varepsilon}{m-1} \\ 0 & \text{if } p_{j} < \frac{\varepsilon}{m-1} \end{cases}, \tag{1}$$

where ε is a probability adjusting factor, and then

$$p'_{i} = 1 - \sum_{j=1, j \neq i}^{m} p'_{j}.$$
 (2)

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Alternatively, if the test case does not reveal a fault, we set 167

$$p'_{i} = \begin{cases} p_{i} - \varepsilon & \text{if } p_{i} \ge \varepsilon \\ 0 & \text{if } p_{i} < \varepsilon \end{cases}$$
(3)

and then for $\forall j = 1, 2, \dots, m$ and $j \neq i$, we set

$$p'_{j} = \begin{cases} p_{j} + \frac{\varepsilon}{m-1} & \text{if } p_{i} \ge \varepsilon\\ p_{j} + \frac{p'_{i}}{m-1} & \text{if } p_{i} < \varepsilon \end{cases}$$
(4)

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Algorithm 1 describes DRT. In DRT, the first test case is 174 taken from a partition that has been randomly selected 175 according to the initial probability profile $\{p_1, p_2, \ldots, p_m\}$ 176

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177 (Lines 2 and 3 in Algorithm 1). After each test case execution, the test profile $\{\langle s_1, p_1 \rangle, \langle s_2, p_2 \rangle, \dots, \langle s_m, p_m \rangle\}$ is updated by 178 changing the values of p_i : If a fault is revealed, Formulas (1) 179 and (2) are used; otherwise, Formulas (3) and (4) are used. The 180 updated test profile is then used to guide the random selection 181 of the next test case (Line 8). This process is repeated until a 182 termination condition is satisfied (Line 1). Examples of possi-183 ble termination conditions include: "testing resources have 184 been exhausted"; "a certain number of test cases have been 185 executed"; and "a certain number of faults have been 186 detected". 187

188 Algorithm 1. DRT

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]	input: $\varepsilon, p_1, p_2, \dots, p_m$
1:	while termination condition is not satisfied
2:	Select a partition s_i according to the testing profile $\{\langle s_1, p_1 \rangle$
	$\langle s_2, p_2 \rangle, \dots, \langle s_m, p_m \rangle \}.$
3:	Select a test case t from s_i .
4:	Test the software using t .
5:	if a fault is revealed by <i>t</i> then
6:	Update p_j ($j = 1, 2,, m$ and $j \neq i$) and p_i according to
	Formulas 1 and 2.
7:	else
8:	Update p_j ($j = 1, 2,, m$ and $j \neq i$) and p_i according t
	Formulas 3 and 4.
9:	end if
10.	end while

As can be seen from Formulas (1) to (4), updating the test profile involves m simple calculations, thus requiring a constant time. Furthermore, the selection of partition s_i , and subsequent selection and execution of the test case, all involve a constant time. The execution time for one iteration of DRT is thus a constant, and therefore the overall time complexity for DRT to select n test cases is $O(m \cdot n)$.

210 2.2 Web Services

A web service is a platform-independent, loosely coupled, 211 self-contained, programmable, web-enabled application that 212 can be described, published, discovered, coordinated and 213 configured using XML artifacts for the purpose of developing 214 distributed interoperable applications [1]. A web service con-215 sists of a description (usually specified in WSDL) and imple-216 mentation (written in any programming language). Web 217 services present their functionalities through published inter-218 faces, and are usually deployed in a service container. Invoca-219 tion of a web service requires analysis of the input message in 220 its WSDL, test data generation based on its input parameters, 221 and wrapping of test data in a SOAP message. 222

A web service is a basic component of SOA software, and, accordingly, the reliability of such SOA software depends heavily on the quality of the component web services. While testing is an obvious potential activity to help assuring the quality of web services, due to the unique features of SOA, web service testing can be more challenging than traditional software testing [4]. Some of these features include:

Lack of access to service implementation: Normally, web
 service owners will not make the source code of their
 web services accessible. Typically, service users only
 have access to the service interface defined in a WSDL

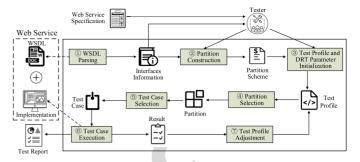


Fig. 1. DRT for web services framework.

file, which means that white-box testing approaches 234 are not possible. 235

- Incomplete documentation or specification: A service pro- 236 vider may only offer an incomplete or inaccurate 237 description of a service's functional and non-functional 238 behavior. This makes it difficult to decide whether or 239 not a test passes, especially when details about behav- 240 ior or restrictions on implementations are missing [15]. 241
- *Lack of control:* Unlike traditional software testing 242 where testers can control the execution of the software 243 under test, there is usually no opportunity to intervene 244 in the execution of the web service under test, which is 245 often deployed in a remote service container. 246
- *Side effects caused by testing:* A large number of tests 247 may introduce an additional communication load, 248 and hence impact on the performance of the web ser- 249 vice under test. This suggests that the number of 250 tests should be kept as low as possible [16]. 251

3 DRT FOR WEB SERVICES

In this section, we describe a framework for applying DRT 253 to web services, discuss guidelines for setting DRT's param-254 eters, and present a prototype that partially automates DRT 255 for web services. 256

3.1 Framework

Considering the principles of DRT and the features of web 258 services, we propose a DRT for web services framework, as 259 illustrated in Fig. 1. In the figure, the DRT components are 260 inside the box, and the web services under test and testers 261 are located outside. Interactions between DRT components, 262 the web services, and testers are depicted in the framework. 263 We next discuss the individual framework components. 264

- WSDL Parsing. Web services are composed of serv- 265 ices and the relevant WSDL documents. By parsing 266 the WSDL document, we can get the input informa- 267 tion for each operation in the services. This includes 268 the number of parameters, their names and types, 269 and any additional requirements that they may have. 270
- 2) Partition Construction. Partition testing (PT) refers to a 271 class of testing techniques that classify the input 272 domain into a number of partitions [17]. Because DRT 273 is a black-box testing technique, combining RT and PT, 274 the PT approaches used are at the specification level. 275 Various approaches and principles for achieving con- 276 venient and effective partitions have been discussed in 277

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the literature [17], [18]. The input domain of the web 278 service under test (WSUT) can be partitioned based on 279 the WSUT specifications and the parsed parameters. 280 Once partitioned, testers can assign probability distri-281 butions to the partitions as an initial testing profile. 282 This initial testing profile can be assigned in different 283 284 ways, including using a uniform probability distribution, or one that sets probabilities according to the 285 importance of the partition: For example, a partition 286 within which faults were previously detected should 287 be given higher priority. 288

Test Profile and DRT Parameter Initialization. Testers 3) 289 need to initialize the test profile, a simple way of doing 290 which would be the use of a uniform probability distri-291 bution $(p_1 = p_2 = \ldots = p_k = 1/k)$, where k denotes the 292 293 number of partitions, and p_i (i = 1, 2, ..., k) denotes the probability of selecting the *i*th partition). They also 294 295 need to set the DRT parameters (guidelines for which are introduced in Section 3.2). 296

- 297 4) *Partition Selection*. DRT randomly selects a partition 298 s_i according to the test profile.
- 5) *Test Case Selection*. DRT selects a test case from the selected partition s_i according to a uniform distribution.
- 3026)Test Case Execution. The relevant DRT component303receives the generated test case, converts it into an304input message, invokes the web service(s) through305the SOAP protocol, and intercepts the test results306(from the output message).
- Test Profile Adjustment. Upon completion of each test, 7) 307 308 its pass or fail status is determined by comparing the actual and expected results (a pass status if both are 309 310 the same). The pass or fail status is then used to adjust the (partition) probability distribution accordingly. Sit-311 uations where determination of the test outcome status 312 is not possible (i.e., in the presence of the oracle prob-313 lem [19], [20]) may potentially be addressed using 314 metamorphic testing [21]. 315

Generally speaking, DRT test case generation is influenced 316 by both the probability distribution (for selection of the rele-317 vant partition), and the principles of RT-combining the effec-318 tiveness of PT with the ease of RT. Because our technique is 319 based on PT, it is necessary that the partition details be pro-320 321 vided (by the tester), which can easily be done through analysis of the input parameters and their constraints, as described 322 323 in the specification of the web service under test. Once the partition details are available, then it is not difficult to set an initial 324 test profile. The tester can, for example, simply use a uniform 325 probability distribution ($p_1 = p_2 = \ldots = p_m = 1/m$, where m 326 denotes the number of partitions, and p_i (i = 1, 2, ..., m)327 328 denotes the probability of selecting the *i*th partition). In Section 3.2, we provide some guidelines for how to set the DRT 329 parameters. Furthermore, many of the components in the DRT 330 for web services framework can be automated. To make it eas-331 332 ier for potential adopters of DRT for web services, we have also developed a prototype application (described in Section 3.3). 333

334 3.2 Guidelines for Parameter Setting

Our previous work [14] found that the DRT performance depends on the number of partitions and the parameter ε . We next explore these impacts through a theoretical analysis, which, to be mathematically tractable, has the following assumptions: 339

- 1) The failure rate θ_i of each partition s_i (i = 1, 2, ..., m, 340and m > 1) is unknown, but can be estimated. 341
- 2) Each failure rate θ_i (i = 1, 2, ..., m, and m > 1) 342 remains unchanged throughout the testing process 343 (faults are not removed after their detection). 344
- Test cases are selected with replacement, which means 345 that some test cases may be selected more than once. 346

A principle of the DRT strategy is to increase the selection 347 probabilities (by amount ε) of partitions with larger failure 348 rates. In addition to the impact of the parameter ε , the number 349 of partitions also influences the speed of updating the test profile (Formulas (1) to (4)). Therefore, for a given number of partitions, we are interested in investigating what values of ε 352 yield the best DRT performance. 353

Letting θ_M denote the maximum failure rate, and s_M 354 denote partitions with that failure rate, then p_i^n denotes the 355 probability of executing the *n*th test case from partition s_i . 356 As testing proceeds, the probability p_M of partition s_M being 357 selected is expected to increase: 358

$$p_M^{n+1} > p_M^n.$$
 (5) 360

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In order to achieve the best performance, the probability 362 of selecting the partition s_M (which has the maximum failure rate) is expected to increase. To achieve this, the increase 364 in probability of s_M being selected for the next round should 365 be larger than that for other partitions. We further analyze 366 sufficient conditions for this goal, and can accordingly 367 derive an interval for ε . 368

Initially, the test profile is $\{\langle s_1, p_1^0 \rangle, \langle s_2, p_2^0 \rangle, \dots, \langle s_m, p_m^0 \rangle\}$, 369 which, after *n* test cases have been executed, is then updated 370 to $\{\langle s_1, p_1^n \rangle, \langle s_2, p_2^n \rangle, \dots, \langle s_m, p_m^n \rangle\}$. During the testing process, 371 p_i^n is increased or decreased by the value ε , which is relatively 372 small (set to 0.05 in previous studies [22], [23]). Because the 373 initial p_i^0 is larger than ε , and the adjustment of p_i is relatively 374 small (Formulas (1) to (4)), the following two situations are 375 rare, and thus not considered here: $p_i < \varepsilon/(m-1)$ or $p_i < \varepsilon$ 376 $(i = 1, 2, \dots, m)$.

To explore the relationship between p_i^{n+1} and p_i^n , we calculate the conditional probability, $p(i|\delta)$, of the following four situations (denoted as $\delta_1, \delta_2, \delta_3$, and δ_4): 380

1) If $t_n \notin s_i$ and a fault is detected by t_n , then $p(i|\delta_1)$ is cal- 381 culated according to Formula (1): 382

$$p(i|\delta_1) = \sum_{i \neq j} \theta_j \left(p_i^n - \frac{\varepsilon}{m-1} \right).$$
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2) If $t_n \in s_i$ and a fault is detected by t_n , then $p(i|\delta_2)$ is calculated according to Formula (2): 387

$$p(i|\delta_2) = \theta_i(p_i^n + \varepsilon).$$
³⁸⁹
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3) If $t_n \in s_i$ and no fault is detected by t_n , then $p(i|\delta_3)$ is 391 calculated according to Formula (3): 392

$$p(i|\delta_3) = (1 - \theta_i)(p_i^n - \varepsilon).$$
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4) If
$$t_n \notin s_i$$
 and no fault is detected by t_n , then $p(i|o_4)$ is concluded according to Formula (4):

$$p(i|\delta_4) = \sum_{i \neq j} (1 - \theta_j) \left(p_i^n + \frac{\varepsilon}{m - 1} \right)$$

401 Therefore, p_i^{n+1} for all cases together is:

$$p_i^{n+1} = p_i^n + Y_i^n, (6)$$

403 404 where

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 $Y_i^n = \frac{\varepsilon}{m-1} (2p_i^n \theta_i m - p_i^n m - 2p_i^n \theta_i + 1)$ $- \frac{2\varepsilon}{m-1} \sum_{j \neq i} p_j^n \theta_j.$ (7)

From Formula (7), we have:

$$Y_M^n - Y_i^n = \frac{2\varepsilon}{m-1} \left(m(p_M^n \theta_M - p_i^n \theta_i) - \frac{m(p_M^n - p_i^n)}{2} \right).$$
(8)

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Before presenting the final guidelines, we need the following lemma.

414 **Lemma 1.** If $p_i^n - p_M^n > 2(p_i^n \theta_i - p_M^n \theta_M)$, then $p_M^{n+1} > p_M^n$.

415 Proof. See Appendix A, which can be found on the Computer
416 Society Digital Library at http://doi.ieeecomputersociety.
417 org/TSC.2019.2960496.

Accordingly, we can now present the following theorem that states a sufficient condition for achieving $p_M^{n+1} > p_M^n$.

420 **Theorem 1.** For failure rate $\theta_{min} = min\{\theta_1, \dots, \theta_m\}, \theta_M >$ 421 θ_{min} , if $0 < \theta_{min} < \frac{1}{2}$, the following condition is sufficient to 422 guarantee that $p_M^{n+1} > p_M^n$:

$$\frac{2m\theta_{\min}^2}{1-2\theta_{\min}} < \varepsilon < \frac{(m-1)m\theta_{\min}}{2(m+1)}.$$
(9)

426 **Proof.** See Appendix A, available in the online supplemental material.

In summary, when $\frac{1}{2} < \theta_M < 1$, there is always an interval \mathcal{I} :

$$\varepsilon \in \left(\frac{2m\theta_{\min}^2}{1-2\theta_{\min}}, \frac{(m-1)m\theta_{\min}}{2(m+1)}\right) \tag{10}$$

432 where $\theta_{min} \leq \theta_i, i \in \{1, 2, ..., m\}$, and $\theta_i \neq 0$, which can 433 guarantee $p_M^{n+1} > p_M^n$.

From the proof in Appendix A, available in the online supplemental material, it is clear that the value of θ_M affects the upper bound (\mathcal{I}_{upper}) of \mathcal{I} . When $\theta_{min} < \theta_M < \frac{1}{2}$, the value of \mathcal{I}_{upper} should be close to the lower bound of \mathcal{I} . In practice, we should set

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$$\varepsilon \approx \frac{2m\theta_{min}^2}{1-2\theta_{min}}.$$
 (11)

For a given partition scheme, with a total of m partitions, identification of the partition with the minimum failure rate (θ_{min}) first requires calculation of the failure rates of each partition, then identification of the minimum. Each parti- 445 tion's failure rate can be obtained in two ways: 446

- 1) It can be calculated directly as F/E (F is the number 447 of failures and E is the number of executed tests), if 448 the test history of the web service under test is 449 available. 450
- 2) It can be approximated by $1/k_i$, where k_i is the total 451 number of test cases executed before revealing a fault. 452

3.3 Prototype

This section describes a tool that partially automates DRT for 454 web services, called DRTester¹. DRTester supports the follow-455 ing tasks in testing web services: a) WSDL parsing; b) partition 456 construction; c) setting DRT parameters (probability adjusting 457 factor ε and test profiles); and d) test case preparation and execution. The details of DRTester are as follows: 459

- 1) *Guidance*. This feature describes the steps the tester 460 should follow when testing a web service. 461
- 2) *Configuration.* This feature, as shown in Fig. 2, inter- 462 acts with the testers to obtain and set the information 463 related to testing the web service, including: the 464 address of the web service under test; the DRT 465 parameters and partition details; and the test case 466 preparation. The detailed steps are as follows: 467
 - Inputting and parsing the URL (Fig. 2a): We integrate 468 the WSDL parsing functionality provided by 469 MT4WS [24]. This enables all the (WSDL) parame- 470 ters and their types to be automatically obtained. 471
 - Parameter setting (Fig. 2a): The tester is responsi- 472
 ble for selecting which operations of the current 473
 web service under test are to be tested, and for 474
 partitioning each parameter into disjoint choices. 475
 - *Partition setting (Fig. 2b)*: The tester is responsible 476 for specifying the partitions by combining the 477 choices associated with each parameter. 478
 - *Test case generation (Fig. 2b)*: The tester is respon- 479 sible for specifying the mode of test case genera- 480 tion (either randomly generating test cases based 481 on the parameters, or uploading test cases gener- 482 ated using other techniques). 483
- Execution. This feature presents a summary of the 484 testing results, including details of the test case exe- 485 cution (input, expected output, partition, and result 486 (pass or fail)). For randomly generated tests, the tes- 487 ter has to check each individual result. Otherwise, 488 when all tests have been completed, a report is gen- 489 erated in a downloadable file. 490

The back-end logic is composed of several Restful APIs and 491 Java classes: The APIs are responsible for communicating 492 HTTP messages to and from the front-end interface. The controller class is responsible for updating the test profile according to the test results, and for selecting test cases from the 495 partitions. The selected test cases are wrapped in SOAP messages and sent to the web service under test through the proxy class, which also intercepts the test results. 498

1. The prototype tool, together with a number of accompanying resources, has been made available at: https://githup.com/phantomDai/DRTester.git

Please enter th	e address of the we	b service under	test
A http://localhos	st:8081/services/acms	?wsdl	뗮 Parse
Parameters Setting	1		
T didificiolo Octaing	9		
Please select a	n operator:		
feeCalculation			
Index	Parameter	Туре	Options
1	airClass	int	1-1:{0};1-2:{1};1-3:{2}
2	area	int	2-1:{0};2-2:{1}
3	isStudent	boolean	3-1:{true};3-2:{false}
4	luggage	double	4-1:{0};4-2:(0,3000]
5	economicfee	double	5-1:(0,10000]
5	economictee	aouple	
			⊘ Save

Partition	Option Combination	Test Profile	Adjusting Factor				
partition1	1-1;2-1	0.167					
partition2	1-1;2-2	0.167					
partition3	1-2;2-1	0.167	0.05				
partition4	1-2;2-2	0.167	0.05				
partition5	1-3;2-1	0.167					
partition6	1-3;2-2	0.165					
est Cases Preparatio		+ Add - Da	iete ⊘ Save				
lease select a me	ethod to generate a test suite						
 Randomly Generate Test 		Upload Test Suite File]				
	r Suite		ML file that				

(a) WSDL Parsing and Parameters Setting

(b) Partition and DRT Parameter Setting

Fig. 2. DRTester configuration snapshots.

499 4 EMPIRICAL STUDY

We conducted a series of empirical studies to evaluate theperformance of DRT.

502 4.1 Research Questions

505

In our experiments, we focused on addressing the followingthree research questions:

506RQ1How effective is DRT at detecting web service faults?507Fault-detection effectiveness is a key criterion for508evaluating the performance of a testing technique.509This study used three popular real-life web services510as subject programs, and applied mutation analysis511to evaluate the effectiveness.

TABLE 1
Subject Web Services

Web service	LOC	Number of mutants
ACMS	116	3
CUBS	131	11
PBS	129	4

RQ2 How do the number of partitions and the DRT 513 parameter ε impact on the failure detection effective-514 ness and efficiency of DRT? 515 In our earlier work [14], we found that the DRT 516 parameter ε had a significant effect on DRT effi-517 ciency, and that the optimal value of the parameter 518 could be related to the number of partitions. The 519 relationship between ε and the number of partitions 520 is examined through theoretical analysis, and veri-521 fied through the empirical studies. 522

RQ3 Compared with the baseline techniques, how efficient 524 is DRT at detecting web service faults in terms of time? 525 Compared with RT and PT, DRT incorporates the selection of partitions and test cases within a partition. Compared with AT, which also introduces feedback and 528 adaptive control principles to software testing, DRT has 529 a simple but efficient control strategy. Thus, we are 530 interested in comparing the fault detection efficiency of 531 DRT, RT, PT, and AT in terms of their time costs. 532

4.2 Subject Web Services

Although a number of web services are publicly available, for 534 various reasons, their implementations are not. This renders 535 them unsuitable for our experiments, which involve the crea- 536 tion of faulty mutants (requiring access to the implementa- 537 tions). We therefore selected three web services as the subject 538 programs for our study, and implemented them ourselves, 539 based on real-life specification:² Aviation Consignment 540 Management Service (ACMS); China Unicom billing 541 service (CUBS); and Parking billing service (PBS). 542 We used the tool MuJava [25] to conduct mutation analysis 543 [26], [27], generating a total of 1563 mutants. Each mutant was 544 created by applying a syntactic change (using one of all appli-545 cable mutation operators provided by MuJava) to the original 546 program. Equivalent mutants, and those that were too easily 547 detected (requiring less than 20 randomly generated test 548 cases), were removed. To ensure the statistical reliability, we 549 obtained 50 different test suites using different random seeds, 550 then tested all mutants with all test suites, calculating the 551 average number of test cases needed to kill (detect) a mutant. 552 Table 1 summarizes the basic information of the used web 553 services and their mutants. A detailed description of each 554 web service is given in the following. 555

4.2.1 Aviation Consignment Management Service (ACMS)

512

533

ACMS helps airline companies check the allowance (weight) 558 of free baggage, and the cost of additional baggage. Based 559

2. The implementations have been made available at: https://github.com/phantomDai/subjects4tsc.git

TABLE 2 ACMS Baggage Allowance and Pricing Rules

		Domestic fligh	its	International flights						
	First class	Business class	Economy class	First class	Business class	Economy class				
Carry on (kg)	5	5	5	7	7	7				
Free checked-in (kg) Additional baggage pricing (kg)	40	30	20 weight $ imes$ pr	$\begin{array}{c} 40\\ ice \times 1.5\%\end{array}$	30	20/30				

on the destination, flights are categorised as either domestic 560 or international. For international flights, the baggage allow-561 ance is greater if the passenger is a student (30 kg), otherwise 562 it is 20 kg. Each aircraft offers three cabin classes from which 563 to choose (economy, business, and first), with passengers in 564 565 different classes having different allowances. The detailed price rules are summarized in Table 2, where price means 566 567 economy class fare and *weight* is the weight that exceeds the weight of the free carry. 568

569 4.2.2 China Unicom Billing Service (CUBS)

CUBS provides an interface through which customers can
know how much they need to pay according to cell-phone
plans, calls, and data usage. The details of several cellphone plans are summarized in Tables 3, 4, and 5.

574 4.2.3 Parking Billing Service (PBS)

Consider a parking billing service that accepts the park-575 ing details for a vehicle, including the vehicle type, day 576 of the week, discount coupon, and hours of parking. 577 This service rounds up the parking duration to the next 578 full hour, and then calculates the parking fee according 579 to the hourly rates in Table 6. If a discount voucher is 580 presented, a 50 percent discount off the parking fee is 581 applied. 582

To facilitate better parking management, at the time of 583 parking, customers may provide an estimation of parking 584 duration, in terms of three different time ranges ((0.0,2.0]), 585 (2.0,4.0], and (4.0,24.0]). If the estimation and actual parked 586 hours fall into the same time range, then the customer will 587 588 receive a 40 percent discount; but if they are different 589 ranges, then a 20 percent markup is applied. A customer may choose to either use a discount coupon, or provide an 590 estimation of parking duration, but may not do both. No 591 vehicles are allowed to remain parked for two consecutive 592 days on a continuous basis. 593

TABLE 3	
Plan A	

Plan d	etails	M	onth charge (C	NY)
		op_A^1	op_A^2	op_A^3
Basic	Free calls (min)	260	380	550
	Free data (MB)	40	60	80
	Free incoming calls	Domesti	c (including vic	leo calls)
Extra	Incoming calls (CNY/min)	0.25	0.20	0.15
	Data (CNY/KB)	3E-4	3E-4	3E-4
	Video calls (CNY/min)	0.60	0.60	0.60

4.3 Variables

4.3.1 Independent Variables

The independent variable is the testing technique. RT, RPT, 596 DRT, and AT [12] were used for comparison. 597

4.3.2 Dependent Variables

The dependent variable for RQ1 is the metric for evaluat- 599 ing the fault-detection effectiveness. Several effectiveness 600 metrics exist, including: the P-measure [28] (the probabil- 601 ity of at least one fault being detected by a test suite); the 602 E-measure [29] (the expected number of faults detected 603 by a test suite); the F-measure [30] (the expected number 604 of test case executions required to detect the first fault); 605 and the T-measure [31] (the expected number of test cases 606 required to detect all faults). Since the F- and T-measures 607 have been widely used for evaluating the fault-detection 608 efficiency and effectiveness of DRT-related testing techni- 609 ques [6], [8], [22], [23], [31], [32], they are also adopted in 610 this study. We use F and T to represent the F-measure 611 and the T-measure of a testing method. As shown in 612 Algorithm 1, the testing process may not terminate after 613 the detection of the first fault. Furthermore, because the 614 fault detection information can lead to different probabil- 615 ity profile adjustment mechanisms, it is also important to 616 see what would happen after revealing the first fault. 617 Therefore, we introduce the F2-measure [30] as the num- 618 ber of additional test cases required to reveal the second 619 fault after detection of the first fault. We use F2 to repre-620 sent the F2-measure of a testing method, and $SD_{measure}$ to 621 represent the standard deviation of metrics (where 622 measure can be F, F2, or T). 623

An obvious metric for RQ3 is the time required to detect 624 faults. Corresponding to the T-measure, in this study we 625 used *T-time*, the time required to detect all faults. *F-time* 626 and *F2-time* denote the time required to detect the first 627 fault, and the additional time needed to detect the second 628 fault (after detecting the first), respectively. For each of these 629 metrics, smaller values indicate a better performance. 630

TABLE 4 Plan B

Plan c	letails	Month charge (CNY)										
		op_B^1	op_B^2	op_B^3	op_B^4	op_B^5	op_B^6					
Basic	Free calls (min) Free data (MB) Free incoming calls	120 40 Don	200 60 nestic	450 80 (inclue	680 100 ding v	920 120 rideo c	1180 150 calls)					
Extra	Incoming calls (CNY/min) Data (CNY/KB) Video calls (CNY/min)	3E-4	3E-4	0.15 3E-4 0.60	3E-4	3E-4	0.15 3E-4 0.60					

595

598

TABLE 5 Plan C

Plan de	etails					Montl	nly charg	ge (CNY))			
		op_C^1	op_C^2	op_C^3	op_C^4	op_C^5	op_C^6	op_C^7	op_C^8	op_C^9	op_C^{10}	op_C^{11}
Basic	Free calls (min) Free data (MB) Free incoming calls	50 150	50 300	240 300	320 400 Do:	420 500 mestic (i	510 650 ncludinş	700 750 g video c	900 950 alls)	1250 1300	1950 2000	3000 3000
Extra	Incoming calls (CNY/min) Data (CNY/KB) Video calls (CNY/min)	0.25 3E-4 0.60	0.20 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60	0.15 3E-4 0.60

TABLE 6 Hourly Parking Rates

Actual parking hours		Weekday		Saturday and sunday							
1 0	Motorcycle	Car: 2-door coupe	Car: others	Motorcycle	Car: 2-door coupe	Car: others					
(0.0,2.0]	\$4.00	\$4.50	\$5.00	\$5.00	\$6.00	\$7.00					
(2.0,4.0]	\$5.00	\$5.50	\$6.00	\$6.50	\$7.50	\$8.50					
(4.0,24.0]	\$6.00 \$6.50		\$7.00	\$8.00	\$9.00	\$10.00					

TABLE 7 Decision Table for ACMS

													_											
	R_1	R_2	R_3	R_4	R_5	R_6	R_7	R_8	R_9	R_{10}	R_{11}	R_{12}	R_{13}	R_{14}	R_{15}	R_{16}	R_{17}	R_{18}	R_{19}	R_{20}	R_{21}	R_{22}	R_{23}	R_{24}
class	0	1	2	0	1	2	0	1	2	0	1	2	0	1	2	0	1	2	0	1	2	0	1	2
destination	0	0	0	1	1	1	0	0	0	1	1	1	0	0	0	1	1	1	0	0	0	1	1	1
isStudent	Ν	Ν	Ν	Ν	Ν	Ν	Υ	Υ	Υ	Y	Y	Y	Ν	Ν	Ν	Ν	Ν	Ν	Y	Y	Y	Y	Y	Υ
isOverload	Ν	Ν	Ν	Ν	Ν	Ν	Ν	Ν	Ν	Ν	Ν	Ν	Y	Y	Υ	Y	Y	Y	Y	Y	Y	Y	Y	Y
$f_{1,1}$	\checkmark																							
$f_{1,2}$													\checkmark						\checkmark					
$f_{1,3}$														\checkmark						\checkmark				
$f_{1,4}$															\checkmark						\checkmark			
$f_{1,5}$																\checkmark						\checkmark		
$f_{1,6}$																	\checkmark						\checkmark	\checkmark
$f_{1,7}$																		\checkmark						

631 4.4 Experimental Settings

632 4.4.1 Partitioning

In our study, we set the partitions by making use of a decision
table (DT) [33]. A DT presents a large amount of complex decision
sions in a simple, straightforward manner, representing a set
of decision rules under all exclusive conditional scenarios in a
pre-defined problem. Typically, a DT consists of four parts:

- 6381)The upper-left part lists the conditions denoted C_i 639 $(i = 1, \dots, n,$ where n is the number of conditions in640the pre-defined problem, and $n \ge 1$). Each condition641 C_i contains a set of possible options $O_{i,q} \in CO_i =$ 642 $\{O_{i,1}, \dots, O_{i,t_i}\}$, where t_i is the number of possible643options for C_i , and $q = \{1, \dots, t_i\}$.
- 6442)The upper-right part shows the condition space, which645is a Cartesian product of all the CO_i (SP(C) =646 $CO_1 \times CO_2 \times \ldots \times CO_n$). Each element in the SP(C)647is a condition entry (CE) with the ordered *n*-tuple.
- 648 3) The lower-left part shows all possible actions, repre-649 sented A_j (j = 1, ..., m, where m is the number of pos-650 sible actions and $m \ge 1$). Similar to CO_i , an action A_j 651 contains a set of possible options $O'_{j,p} \in AO_j =$

 $\{O'_{j,1}, \dots, O'_{j,k_j}\}$, where k_j is the number of alternatives 652 for A_j , and $p = \{1, \dots, k_j\}$.

4) The lower-right part shows the action space SP(A), 654 which is also a Cartesian product of all the AO_j 655 $(SP(A) = AO_1 \times AO_2 \times ... \times AO_m)$. Similar to *CE*, 656 each element in the SP(A) is an action entry (*AE*) 657 with the ordered *m*-tuple. 658

A DT *rule* is composed of a *CE* and its corresponding *AE*. 659 With DT, it is possible to obtain partition schemes with dif-660 ferent granularities. For fine-grain partition schemes, each 661 *CE* of a DT *rule* corresponds to a partition; while for coarse-662 grained schemes, a partition corresponds to the union of a 663 group of partitions for which all *CE* of DT *rules* have the 664 same *AE*. The decision tables for ACMS, CUBS, and PBS are 665 shown in Tables 7, 8 and 9, respectively. In the tables, R_i 666 (i = 1, 2, ..., n) denotes the identified *i*th *rule*; *n* is the total 667 number of *rules*; and the checkmark (\checkmark) under each *rule* indi-668 cates that the corresponding action should be taken. The 669 details of actions are provided in Table 10, where *w* is the 670 weight of baggage; *price* means economy class fare; *op* 671 means the monthly charge; *call* and *data* mean the call dura-672 tion and data usage, respectively; *freeCall* and *freeData* 673

TABLE 8	
Decision Table for CUBS	

	R_1	R_2	R_3	R_4	R_5	R_6	R_7	R_8	R_9	R_{10}	R_{11}	R_{12}	R_{13}	R_{14}	R_{15}	R_{16}	R_{17}	R_{18}	R_{19}	R_{20}
plan option	$\mathop{A}\limits_{op^1_A}$	$\mathop{A}\limits_{op^2_A}$	$A \\ op_A^3$	$B \\ op_B^1$	$B \\ op_B^2$	$B o p_B^3$	$B o p_B^4$	$B o p_B^5$	${B \atop op_B^6}$	$\begin{array}{c} C \\ op_C^1 \end{array}$	$\begin{array}{c} C \\ op_C^2 \end{array}$	$\begin{array}{c} C \\ op_C^3 \end{array}$	$\begin{array}{c} C \\ op_C^4 \end{array}$	$\begin{array}{c} C \\ op_C^5 \end{array}$	$C \\ op_C^6$	$\begin{array}{c} C \\ op_C^7 \end{array}$	$\begin{array}{c} C \\ op_C^8 \end{array}$	$\begin{array}{c} C \\ op_C^9 \end{array}$	$\begin{array}{c} C \\ op_C^{10} \end{array}$	$\begin{array}{c} C \\ op_C^{11} \end{array}$
$f_{2,1} \\ f_{2,2} \\ f_{2,3}$	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark
								De		ABLE 1 Table	9 e for PI	35								
	F	\mathbf{l}_1	R_2	R_3	R_4	R_5	R_6	R_7	R_8	R_9	R_{10}	R_{11}	R_{12}	$_2$ F	13	R_{14}	R_{15}	R_{16}	R_{17}	R_{18}
vehicle time discount	((t ()	1 0 0	2 0 0	0 1 0	1 1 0	2 1 0	0 0 1	1 0 1	2 0 1	0 1 1	1 1 1	2 1 1		0 0 2	1 0 2	2 0 2	0 1 2	1 1 2	2 1 2
$f_{3,1} \\ f_{3,2}$	v	/	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark							

mean the free calls and free data, respectively; and *baseFee* 674 means the cost before the discount. In Table 7, the condition 675 for calculating the cost of the baggage includes class (0: first 676 class; 1: business class; and 2: economy class), isStudent (Y: 677 the passenger is a student; and N: the passenger is not a stu-678 dent), isOverload (Y: the baggage exceeds the free carry-on 679 weight limit; and N: the baggage does not exceed the free 680 carry-on weight limit.), and Destination (0: domestic flight; 681 and 1: international flight). In Table 8, conditions that influ-682 683 ence cell-phone bills include *plan* (A: plan A; B: plan B; and C: plan C) and option y under plan x, represented as op_x^y , 684 685 where $x \in \{A, B, C\}$, and $y \in \{w | 1 \le w \le 11 \land w \in \mathcal{Z}\}$. In Table 9, conditions that affect the parking fee include the 686 type of vehicle (0: motorcycle; 1: 2-door coupe; and 2: others), 687 day of week (0: weekday; and 1: saturday or sunday), and dis-688 count information (0: customers provide a discount coupon; 689 1: the estimated hours of parking and the actual hours of 690

TABLE 10 Formulas of the Actions in Table 7 \sim 9

Web Service	Formulas
ACMS	$ \begin{array}{l} f_{1,1} = 0 \\ f_{1,2} = (w-25) \times price \times 1.5\% \\ f_{1,3} = (w-35) \times price \times 1.5\% \\ f_{1,4} = (w-25) \times price \times 1.5\% \\ f_{1,5} = (w-47) \times price \times 1.5\% \\ f_{1,6} = (w-37) \times price \times 1.5\% \\ f_{1,7} = (w-27) \times price \times 1.5\% \end{array} $
CUBS	$ \begin{array}{l} f_{2,1} = op + (call - freeCall) \times 0.25 \\ + (data - freeData) \times 0.0003 \\ f_{2,2} = op + (call - freeCall) \times 0.20 \\ + (data - freeData) \times 0.0003 \\ f_{2,3} = op + (call - freeCall) \times 0.15 \\ + (data - freeData) \times 0.0003 \end{array} $
CUBS	$\begin{array}{l} f_{3,1} = baseFee \times 50\% \\ f_{3,2} = baseFee \times (1-40\%) \\ f_{3,3} = baseFee \times (1+20\%) \end{array}$

parking fall into the same time range; and 2: estimated 691 hours and the actual hours are in different time ranges). 692

As can be seen from the description above, because the DT 693 considers all parameters, and identifies their invalid combinations, it can provide a systematic and efficient way to partition 695 an input domain into disjoint subdomains, and then generate 696 test cases. In practice, each DT *rule* condition entry corresponds to a partition in which test cases cover some paths— 698 thus, the faults in those paths have a chance of being detected. 699

4.4.2 Initial Test Profile

Because test cases may be generated randomly during the test 701 process, a feasible method is to use a uniform probability dis-702 tribution as the initial testing profile. On the other hand, test-703 ers may also use past experience to help guide selection of a 704 different probability distribution as the initial profile. In our 705 experiment, we used a uniform probability distribution for 706 the initial test profile. The initial test profiles of each web ser-707 vice are summarized in Table 11, where $< s_i, p_i >$ means 708 that the probability of selecting partition s_i is p_i .

4.4.3 Constants

710

700

In the experiments, we were interested in exploring the rela- 711 tionship between the number of partitions and the DRT 712

TABLE 11 Initial Test Profile for Subject Web Services

Actual parkin	g Hourly parking	Initial test
hours	rates	profile
ACMS	24 7	$\begin{cases} \langle s_1, \frac{1}{24} \rangle, \langle s_2, \frac{1}{24} \rangle, \dots, \langle s_{24}, \frac{1}{24} \rangle \\ \langle s_1, \frac{1}{7} \rangle, \langle s_2, \frac{1}{7} \rangle, \dots, \langle s_7, \frac{1}{7} \rangle \end{cases}$
CUBS	20 3	$\begin{cases} \langle s_1, \frac{1}{20} \rangle, \langle s_2, \frac{1}{20} \rangle, \dots, \langle s_{20}, \frac{1}{20} \rangle \\ \langle s_1, \frac{1}{3} \rangle, \langle s_2, \frac{1}{3} \rangle, \dots, \langle s_3, \frac{1}{3} \rangle \end{cases}$
PBS	18 3	$\begin{cases} \left\langle s_1, \frac{1}{18} \right\rangle, \left\langle s_2, \frac{1}{18} \right\rangle, \dots, \left\langle s_{18}, \frac{1}{18} \right\rangle \\ \left\langle s_1, \frac{3}{2} \right\rangle, \left\langle s_2, \frac{1}{3} \right\rangle, \dots, \left\langle s_3, \frac{1}{3} \right\rangle \end{cases} \end{cases}$

strategy parameter ε , and therefore selected a set of parame-713 ter values: $\varepsilon \in \{1.0E-05, 5.0E-05, 1.0E-04, 5.0E-04, 1.0E-03, \ldots, 0.0E-04, 1.0E-03, \ldots, 0.0E-04, \ldots, 0.0E-03, \ldots, 0.0E-04, \ldots, 0.0E-03, \ldots, 0.0E-04, \ldots, 0.$ 714 5.0E-03, 1.0E-02, 5.0E-02, 1.0E-01, 2E-01, 3E-01, 4E-01, 5E-01. 715 It should be noted that $\varepsilon = 5E$ -01 is already a large value. 716 Consider the following scenario. For PBS, when the test is 717 carried out under partition scheme 2, if $\varepsilon = 7.5E$ -01 and a 718 719 uniform probability distribution is used as the testing profile (that is, $p_i = 1/3$), then suppose that the first test case 720 belonging to c_1 is executed and does not reveal any faults, 721 then, according to Formula (3), the value of p_1 would 722 become 0. It is important, therefore, that the initial value of 723 ε should not be set too large. 724

725 4.5 Experimental Environment

Our experiments were conducted on a virtual machine 726 running the Ubuntu 11.06 64-bit operating system, with 727 two CPUs, and a memory of 2GB. The test scripts were 728 written in Java. To ensure statistically reliable values [34] 729 of the metrics (F-measure, F2-measure, T-measure, F-time, 730 F2-time, and T-time), each testing session was repeated 731 30 times with 30 different seeds, and the average value 732 calculated. 733

734 4.6 Threats To Validity

735 4.6.1 Internal Validity

A threat to internal validity is related to the implementations
of the testing techniques, which involved a moderate amount
of programming work. However, our code was cross-checked
by different individuals, and we are confident that all techniques were correctly implemented.

741 4.6.2 External Validity

The possible threat to external validity is related to the 742 subject programs and seeded faults under evaluation. 743 744 Although the three subject web services are not very complex, they do implement real-life business scenarios of 745 diverse application domains. Furthermore, 18 distinct 746 faults were used to evaluate the performance. These faults 747 cover different types of mutation operators and require an 748 average of more than 20 randomly generated test cases to 749 be detected. Although we have tried to improve the gener-750 751 alisability of the findings by applying different partitioning granularities, and 13 kinds of parameters, we anticipate 752 that the evaluation results may vary slightly with different 753 754 subject web services.

755 4.6.3 Construct Validity

The metrics used in our study are simple in concept and
straightforward to apply, and hence there should be little
threat to the construct validity.

759 4.6.4 Conclusion Validity

As reported for empirical studies in the field of software engineering [34], at least 30 observations are necessary to ensure the statistical significance of results. Accordingly, we have run a sufficient number of trials to ensure the reliability of our experimental results. Furthermore, as will be discussed in Section 5, we also conducted statistical tests to confirm the significance of the results.

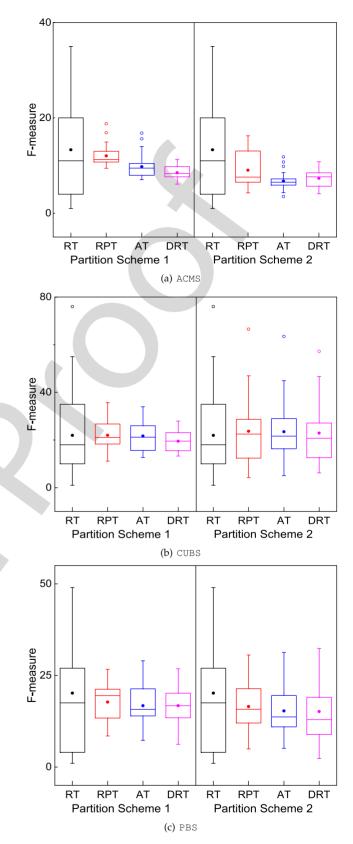


Fig. 3. F-measure boxplots for each web service.

5 EXPERIMENTAL RESULTS

5.1 RQ1: Fault Detection Effectiveness

F-, F2-, and T-measure results for ACMS, CUBS, and PBS are 769 shown using boxplots in Figs. 3, 4, and 5, where the DRT 770

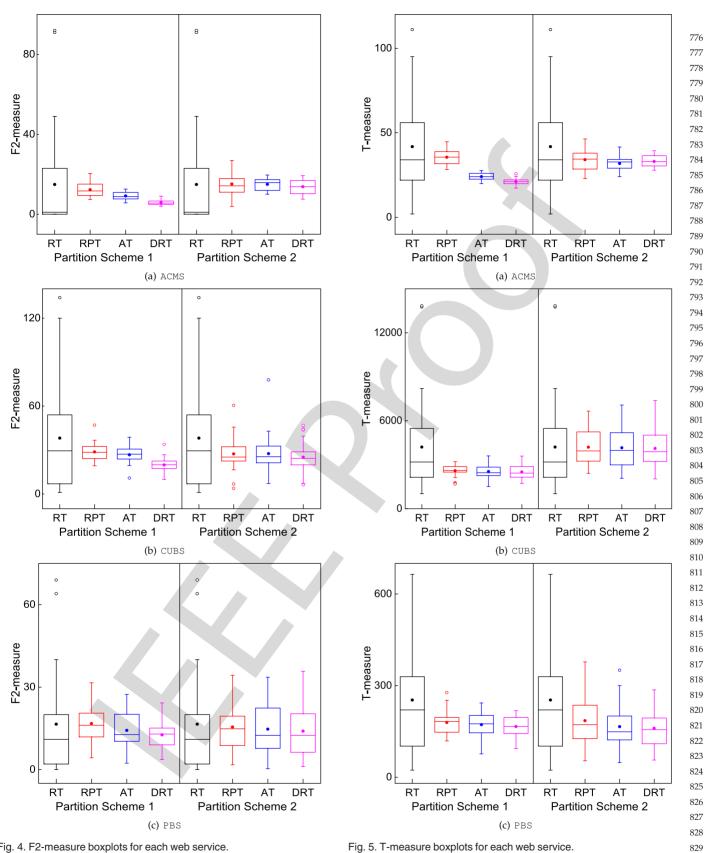


Fig. 4. F2-measure boxplots for each web service.

parameter ε was set to the optimal values, as described in 771 Section 5.2. The experimental results of DRT with other val-772 ues of ε are shown in Appendix B, available in the online 773 supplemental material. In each boxplot, the upper and 774 lower bounds of the box represent the third and first 775

quartiles of the metric, respectively; the middle line repre- 832 sents the median value; the upper and lower whiskers 833 mark, respectively, the largest and smallest data within the 834 range of $\pm 1.5 \times IQR$ (where IQR is the interquartile range); 835 outliers beyond the IQR are denoted with hollow circles; 836

TABLE 12Number of Scenarios Where the Technique on the
Top Row has a Lower Metric (F-/F2-/T-Measure)
Score Than the Technique on the Left Column

		F-me	easui	re		F2-m	easu	re	T-measure				
	RT	RPT	AT	DRT	RT	RPT	AT	DRT	RT	RPT	AT	DRT	
RT	_	4	5	5	_	4	5	6	_	6	6	6	
RPT	2	—	6	6	2	—	5	6	0	—	6	6	
AT	1	0		4	1	1	—	6	0	0		5	
DRT	1	0	2	—	0	0	0	—	0	0	1	—	

837 and each solid circle represents the mean value of the 838 metric.

839 It can be observed from the figures that, in an overwhelming majority of cases, DRT was the best performer in terms of 840 F-, F2-, and T-measure, followed by AT, RPT, and RT. On the 841 other hand, RT may be the best performer occasionally or the 842 worst performer in terms of F-, F2-, and T-measure, which 843 means that the fault detection effectiveness of RT is not stable. 844 In contrast, DRT and AT show a relatively stable fault detec-845 tion effectiveness. We also conducted statistical testing to 846 verify the significance of this observation, using the Holm-847 848 Bonferroni method [30] (with p-value equal to 0.05) to determine which pairs of testing techniques had significant 849 differences. The statistical data are shown in Table 12, where 850 each cell gives the number of scenarios where the technique 851 above (in the table) performed better than one to the left. For 852 example, the "6" in the top right cell of Table 12 indicates that, 853 of 6 scenarios (two partition schemes \times three web services), 854 DRT had lower T-measure scores than RT for 6, with the fault-855 detection capabilities of these two techniques being signifi-856 cantly different. 857

Table 12 clearly shows that the differences between pairsof testing techniques are all significant.

860 5.2 RQ2: Relationship between Partition Number 861 and ε

In Section 3.2, we analyzed the relationship between the number of partitions and the DRT parameter ε . In this section, we show that our theoretical analysis provides useful guidance to testers to set the value of ε .

We used three web services to validate our theoretical anal-866 ysis. Before starting the test, it is necessary to know the failure 867 868 rate θ_i of partition s_i . From Tables 2, 3, 4 and 5, it can be observed that the values of some parameters (such as the bag-869 gage weight, the call duration, and parking duration) are such 870 that the total number of test case values in a partition could be 871 infinite. For such a situation, we approximate the failure rate 872 θ_i of s_i by $1/k_i$ (where k_i is the total number of test cases exe-873 cuted before revealing a fault). According to Formula (19), the 874 theoretically optimal values of ε in each scenario for each web 875 service are shown in Table 13, where ε^* denotes the theoretical 876 877 value of ε . We ran a series of experiments with the parameters set according to those in Table 13: The F-, F2-, and T-measure 878 results for each program are shown in Fig. 6, where ε_1^* and ε_2^* 879 denote the theoretical values of parameter ε in the two differ-880 ent partition schemes, respectively. For ease of presentation 881 and understanding, we used $log_{100}(1.0E05 \times \varepsilon)$ for the 882

TABLE 13 Theoretical Optimal Values of DRT Parameter

Web	Partition	$ heta_{min}$	$arepsilon^*$
service	scheme		
ACMS	1	5.452E-2	1.601E-1
	2	2.797E-3	1.102E-4
CUBS	1	1.193E-3	5.702E-5
	2	1.397E-3	1.734E-5
PBS	1	1.760E-3	1.118E-4
	2	1.492E-3	1.340E-5

horizontal axis in Fig. 6. Apart from the DRT strategy parame- 883 ter ε , all other experimental settings remained the same as in 884 Section 5.1. 885

From Fig. 6, we have the following observations:

• In most scenarios, the DRT strategy with theoreti- 887 cally optimum parameter value performs best. Fur- 888 thermore, the DRT strategy performs better when 889 the parameter values are near the theoretically optimum value than when not. 891

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• From Fig. 6a, it can be observed that the DRT strategy 892 with larger parameter values performs better than 893 with the theoretically optimum value, in terms of the 894 F-measure. The main reason for this is that, for this 895 scenario, the maximum failure rate ($\theta_M = 4.781E - 3$) 896 is large and the number of partitions is small: When 897 the parameter value is large, the probability of select-898 ing partitions with lower failure rates is quickly 899 reduced, and the probability of selecting partitions 900 with larger failure rates is quickly increased, according 901 to Formulas (3) and (4). 902

5.3 RQ3: Fault Detection Efficiency

The F-, F2-, and T-time results for ACMS, CUBS, and PBS are 904 summarized in Table 14, where the values of DRT parame-905 ters for the subject web services are the same as those in 906 Section 5.1. The F-, F2-, and T-time results of DRT with dif-907 ferent parameter values are summarized in Appendix B, 908 available in the online supplemental material. It can be 909 observed from the table that, in general, DRT had the best 910 performance; RPT marginally outperforms RT; and AT had 911 the worst performance. 912

As was done for the F-, F2-, and T-measure data, we used 913 the Holm-Bonferroni method to check the difference between 914 each pair of testing strategies in terms of F-time, F2-time, and 915 T-time, as shown in Table 15. Table 15 shows that: a) DRT was 916 significantly better than AT in terms of F-/F2-/T-time; b) DRT 917 was significantly better than RT and RPT in terms of 918 F2-/T-time; and c) DRT marginally outperformed RT and 919 RPT in terms of F-time. In other words, the additional computation incurred in DRT by updating the test profile is compensated for in terms of test execution savings. 922

In summary, the DRT strategy is considered the best testing 923 technique across this three metrics, RPT marginally outper-924 formed RT, and DRT, RPT, and RT significantly outperformed 925 AT. 926

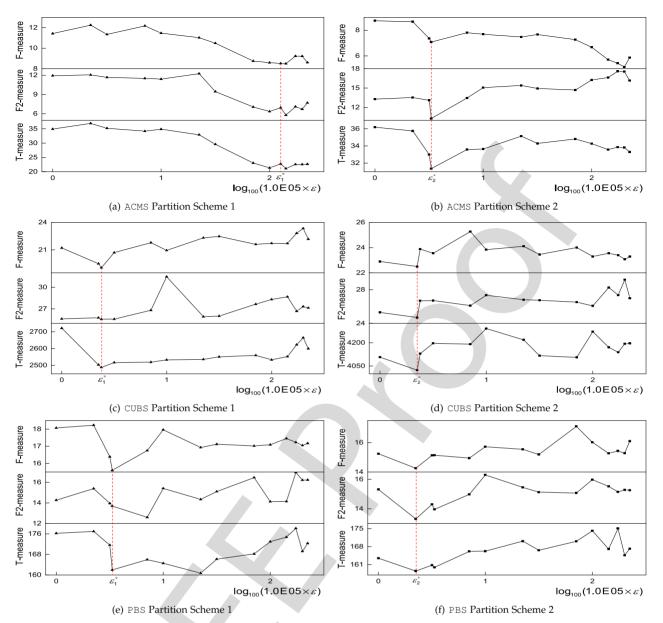


Fig. 6. Line charts of F-measure, F2-measure, and T-measure values for each web service (for both the theoretically optimum parameter value, and other values.

927 **5.4 Summary**

928 Based on the evaluation, we have the following observations:

- DRT outperformed RT, RPT, and AT, according to all 929 the applied metrics for all three studied web services. 930 DRT marginally outperformed AT in terms of the 931 F-, F2-, and T-measure, for all the studied web services. 932 Moreover, AT incurs heavier computational overhead, 933 and takes a significantly longer time. For instance, AT 934 required 32429.07ms to select and execute sufficient 935 test cases to detect all faults in CUBS, while DRT only 936 needed 30.21ms (Table 14). This indicates that among 937 RT, RPT, and AT, DRT should be chosen. 938
- DRT is more effective in terms of the F-, F2-, and T-measure when the parameter settings are optimal (according to the theoretical analysis): In most cases, DRT has the best performance for all three web services, according to these three metrics (F-measure,

F2-measure, and T-measure) when following the 944 guidelines for the parameter settings. This highlights 945 the usefulness of the parameter-setting guidelines. 946 We also note the following limitations: 947

- While DRT outperformed RT and RPT in terms of 948 fault detection effectiveness and efficiency, this was 949 achieved at the cost of the additional effort required 950 to set the partitions and test profiles. 951
- Applying DRT involves setting parameters, which 952 may not be trivial. Even when following the theoretical 953 guidelines. 954

6 RELATED WORK

In this section, we describe related work from two perspectives: related to testing techniques for web services; and related to improving RT and PT. 958

TABLE 14
F-time, F2-time, and T-time in ms for Subject Web Services

Partition	Metric	letric ACMS					CUBS					PBS			
Scheme		RT	RPT	AT	DRT	RT	RPT	AT	DRT	RT	RPT	AT	DRT		
	F-time	0.43	0.57	0.49	0.23	0.82	0.91	140.13	0.95	0.81	0.85	22.25	0.68		
1	F2-time	0.29	0.31	2.47	0.12	1.14	0.87	172.27	0.86	0.42	0.52	25.40	0.34		
	T-time	0.85	1.08	2.47	0.43	34.69	30.54	32429.07	30.21	4.12	3.83	289.04	3.20		
	F-time	0.43	0.33	15.53	0.24	0.82	0.75	16.82	0.87	0.81	0.66	12.99	0.49		
2	F2-time	0.29	0.45	363.47	0.28	1.14	0.79	15.76	0.83	0.42	0.35	17.44	0.34		
	T-time	0.85	0.78	459.67	0.65	34.69	34.59	2666.17	36.49	4.12	2.98	200.54	2.26		

959 6.1 Testing Techniques for Web Services

In recent years, a lot of effort has been made to test web 960 services [4], [13], [35], [36]. Test case generation or selection 961 is core to testing web services, and model-based [37] and 962 specification-based [38] techniques are two common app-963 roaches. Before making services available on the Internet, test-964 ers can use model-based techniques to verify whether or not 965 the behavior of the WSUT meets their requirements. In these 966 techniques, test data can be generated from a data model that 967 specifies the inputs to the software—this data model can be 968 built before, or in parallel to, the software development 969 process. Verification methods using technologies such as 970 theorem-proving [39], models [40] and Petri-Nets [41] also exist. 971

All of the above approaches aim to generate test cases with-972 out considering the impact of test case execution order on test 973 efficiency. In contrast, Bertolino et al. [42] proposed using the 974 category-partition method [43] with XML schemas to perform 975 XML-based partition testing. Because PT aims to find subsets of 976 all possible test cases to adequately test a system, it can help 977 reduce the required number of test cases. Our approach 978 involves software cybernetics and PT: In DRT, selection of a 979 partition is done according to the testing profile, which is 980 updated throughout the testing process. An advantage of DRT 981 is that partitions with larger failure rates have higher probabili-982 ties of selection. Zhu and Zhang [44] proposed a collaborative 983 testing framework, where test tasks are completed using collab-984 orating test services-a test service is a service assigned to per-985 986 form a specific testing task. Our framework (Section 3.1) aims to find more faults in the WSUT, with the result of the current test 987 case execution providing feedback to the control system so that 988 the next test case selected has a greater chance to reveal faults. 989

Most web service testing techniques assume that the computed output for any test case is verifiable, which is, however, not always true in practice (a situation known as the oracle problem [19]). Thus, many testing techniques may not be applicable in some cases. To address the oracle problem for

TABLE 15 Number of Scenarios Where the Technique on the Top Row has a Lower Metric (F-/F2-/T-Time) Score Than the Technique on the Left Column

	F-time					F2-	time		T-time				
	RT	RPT	AT	DRT	RT	RPT	AT	DRT	RT	RPT	AT	DRT	
RT	_	3	0	4	_	3	0	6	_	5	0	6	
RPT	3	—	1	4	3	_	0	5	1	—	0	6	
AT	6	5	—	6	6	6	—	6	6	6	—	6	
DRT	2	2	0	—	0	1	0	—	0	0	0		

testing web services, Sun *et al.* [21] proposed a metamorphic 995 testing [45], [46] approach that not only alleviates the oracle 996 problem, but is also a practical and efficient option for testing 997 web services. They conducted a case study that showed that 998 up to 94.1 percent of seeded faults could be detected without 999 the need for oracles. 1000

6.2 Improving RT and PT

Based on the observation that failure-causing inputs tend to 1002 cluster into contiguous regions in the input domain [10], [11], 1003 much work has been done to improve RT [6], [7], [9]. Adaptive 1004 random testing [7], [9] is a family of techniques based on ran- 1005 dom testing that aim to improve the failure-detection effective- 1006 ness by evenly spreading test cases throughout the input 1007 domain. One well-known ART approach, FSCS-ART, selects a next test input from the fixed-size candidate set of tests that is 1009 farthest from all previously executed tests [47]. Many other 1010 ART algorithms have also been proposed, including RRT [48], 1011 DF-FSCS [49], and ARTsum [50], with their effectiveness exam-1012 ined and validated through simulations and experiments. 1013

Adaptive testing (AT) [8], [51], [52] takes advantage of 1014 feedback information to control the execution process, and 1015 has been shown to outperform RT and RPT in terms of the 1016 T-measure and the number of detected faults, which means that AT has higher efficiency and effectiveness than RT and 1018 RPT. However, AT may require a rather long execution time in practice. To alleviate this, Cai et al. [6] proposed 1020 DRT, which uses testing information to dynamically adjust 1021 the testing profile. There are several things that can impact 1022 on DRT's test efficiency. Yang et al. [32] proposed A-DRT, 1023 which adjusts parameters during the testing process. 1024

7 CONCLUSION

In this paper, to address the challenges of testing SOA-based 1026 applications, we have presented a dynamic random testing 1027 (DRT) method for web services. Our method uses random 1028 testing to generate test cases, and selects test cases from differ- 1029 ent partitions in accordance with a testing profile that is 1030 dynamically updated in response to the test data collected. In 1031 this way, the proposed method enjoys benefits from both ran-030 dom testing and partition testing. 1033

We proposed a framework that examines key issues when 1034 applying DRT to test web services, and developed a prototype 1035 to make the method practical and effective. To guide testers to 1036 correctly set the DRT parameters, we used a theoretical analysis to study the relationships between the number of partitions 1038 (*m*) and the probability adjusting factor (ε). Three real web 1039 services were used as experimental subjects to validate the 1040

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feasibility and effectiveness of our approach. Our experimental results show that, in general, DRT has better performance than both RT and RPT, in terms of the F-, F2-, and T-measures, and always outperforms when the ε settings follow our guidelines. In other words, our theoretical analysis can provide genuinely useful guidance to use DRT.

In our future work, we plan to conduct experiments on
more web services to further validate the effectiveness, and
identify the limitations of our method.

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Chang-ai Sun received the bachelor degree in computer science from the University of Science and Technology Beijing, China, and the PhD degree in computer science from Beihang University (Beijing University of Aeronautics and Astronautics), China. He is currently a professor with the School of Computer and Communication Engineering, University of Science and Technology Beijing. Before that, he was an assistant professor at Beijing Jiaotong University, China, a postdoctoral fellow at the Swinburne University of

Technology, Australia, and a postdoctoral fellow at the University of Groningen, The Netherlands. His research interests include software testing, program analysis, and service-oriented computing. He is a senior member of the IEEE.



Hepeng Dai received the bachelor degree in information and computing sciences from the China University of Mining and Technology, Beijing, and the master degree in software engineering from the University of Science and Technology Beijing, China. He is currently working toward the PhD degree in the School of Computer and Communication Engineering, University of Science and Technology Beijing, China. His current research interests include software testing and debugging.



Guan Wang received the bachelor degree in computer science from the University of Science and1238puter science from the University of Science and1249Technology Beijing. He is currently working toward1240the master degree at the School of Computer and1241Communication Engineering, University of Science1242and Technology Beijing. His current research interests include software testing and service-oriented1244computing.1245



Dave Towey received the BA and MA degrees1246from the University of Dublin, Trinity College,1247Ireland; PgCertTESOL from the Open University of1248Hong Kong; the MEd degree from the University of1249Bristol, United Kingdom; and the PhD degree from1250associate professor with the School of Computer1252Science, University of Nottingham Ningbo China.1253His current research interests include technology-1254enhanced teaching and learning, and software test-1255ing, especially metamorphic testing and adaptive1256random testing. He is a member of both the IEEE12571258125712591254125012541255125712541256125512571256125712571258125812571259125712501257125112571251125712511257125212571254125712551257125612571257125812581257125912581259125912501259125112571251125712511257125212571254125712551257125612571257125812581257125912571259



Tsong Yueh Chen received the BSc and MPhil1259degrees from the University of Hong Kong,1260Hong Kong, the MSc and DIC degrees from the1261Imperial College of Science and Technology,1262London, United Kingdom, and the PhD degree in1263computer science from the University of Melbourne,1264Australia. He is currently a professor of software1266engineering at the Department of Computer Science1266and Software Engineering, Swinburne University of1267Technology. He is the inventor of metamorphic test-1268ing and adaptive random testing. He is a member of1269the IEEE.1270



Kai-Yuan Caireceived the BS, MS, and PhD1271degrees from Beihang University (Beijing University1272of Aeronautics and Astronautics), Beijing, China, in12731984, 1987, and 1991, respectively. He has been a1274full professor at Beihang University since 1995. He1275is currently a Cheung Kong scholar (chair professor)1276appointed by the Ministry of Education of China in12771999. His main research interests include software1278testing, software reliability, reliable flight control,127912811281

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