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Memetic chicken swarm algorithm for job shop scheduling problem

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ABSTRACT

This paper presents a Memetic Chicken swarm optimization (MeCSO) to solve job shop scheduling problem (JSSP). The aim is to find a better solution which minimizes the maximum of the completion time also called Makespan. In this paper, we adapt the chicken swarm algorithm which take into consideration the hierarchical order of chicken swarm while seeking for food. Moreover, we integrate 2-opt method to improve the movement of the rooster. The new algorithm is applied on some instances of OR-Library. The empirical results show the forcefulness of MeCSO comparing to other metaheuristics from literature in term of run time and quality of solution.

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1. INTRODUCTION

The job-shop scheduling problem (JSSP) was formulated for the first time by Muth and Thompson in 1963. The JSSP is one of the NP-Hard problems [1] and the most known of the classical scheduling problems in the context of manufacturing [2], which help to improve competitiveness of many companies and organizations. The aim purpose of the job-shop scheduling problem is to find a schedule which minimizes the time required to complete a group of jobs (the makespan).

Historically, several algorithms are proposed in literature to solve the job shop scheduling problem by optimizing the makespan such as: branch and bound (BB)[3], simulated annealing (SA) [4], Tabu search (TS)[5][6], genetic algorithms (GA)[7][8][9], neural networks (NN)[10], ant colony optimization (ACO)[11], Particle swarm optimization (PSO)[12], Bee colony optimization (BCO)[13] and firefly algorithm(FA) [14]. Additionally, some researchers have developed an hybrid optimization strategy for JSSP such as parallel GRASP with path-relinking[15] and new hybrid genetic algorithm [16].

2. JOB-SHOP SCHEDULING PROBLEM

The JSSP can be briefly introduced [17] as a sequential allocation of a production schedule for a given set of jobs and resources that optimizes the completion time of all jobs which helps to minimize the makespan. As result, the makespan (the maximum job completion time) C_{max} is the duration between the time of completion of last job and the starting time of the first job (1).

$$C_{max} = max_{t_{ij}}(t_{ij} + p_{ij}) \tag{1}$$

Where t_{ij} is denoted as the starting time and p_{ij} as the uninterrupted processing time.

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The JSSP can be formulated by assigning a set of n jobs $J = \{J_1, \ldots, J_n\}$ to a set of m machines $M = \{M_1, \ldots, M_m\}$, each machine can process at most one operation at time. As well, each job consists of a set of O_{ik} , which contains m operations where i denotes the job of a specific operation and k represents the current machine M_k . Each operation must be processed during an uninterrupted period of time on a given machine. In the jssp, the order and the uninterrupted processing time must be take into consideration.

The schedule as a solution for the JSSP can be modeled as a vector of a sequence of operation $(C_{11}, C_{ji}, ..., C_{nm+1})$ then the main goal is to find the minimum time of all processes, the problem is formulated as follows:

$$minC_{nm+1}$$
 (2)

Where

$$C_{kl} \le C_{ji} - d_{kl}; j = 1, \dots, n; i = 1, \dots, m; kl \in P_{ji}$$
 (3)

$$\sum_{ji \in O(t)} r_{ji} \le 1; i \in M; t \ge 0 \tag{4}$$

$$C_{ii} \ge 0; j = 1, \dots, n; i = 1, \dots, m$$
 (5)

The constraint (2) minimizes the finish time of operation o_{nm+1} (the makespan).

The constraint (3) represents the fact that between operations the precedence relations should be respected.

The constraint (4) describes that each machine can process one operation at a each time.

The constraint (5) guarantees that the finish times to be positive.

The remainder of this paper is organised as follows: The section 2 represents the literature review of the problem. The Section 3 describes the proposed memetic-CSO algorithm. The Section 4 presents the results of the experimental study . The Section 5 gives a discussion of the empirical results. Finally, the Section 6 gives the conclusion and the prospects for further works.

3. FORMULATION OF THE PROBLEM

In the job-shop scheduling problem (JSSP),the solution can be depicted as a sequence of $n \times m$ operations,which optimizes the completion time of all jobs and then helps tp find a schedule with minimum makespan.

let's consider the following example with m=3 machines and n=3 jobs, where: $J=\{Job0, Job1, Job2\}$ and $M=\{0,1,2\}$

$$Job0 = \{(0;3), (1;2), (2;2)\}$$
$$Job1 = \{(0;2), (2;1), (1;4)\}$$

 $Job2 = \{(1;4), (2;3)\}$

The representation of the matrix will be as bellow:

The first line contains the operation number, the second line contains the job number, the third line contains the sequence number, the forth line contains the machine number and the last line contains the processing time of each operation.

As indicated in the Gantt chart representation Figure 1, the solution $S = \{0, 6, 3, 4, 1, 5, 2, 7\}$ is given by a permutation of a set of operations on each machine, in this example the minimum makespan Cmax=11. In this paper, the chicken can search food in a set of solutions S defined as the search space.

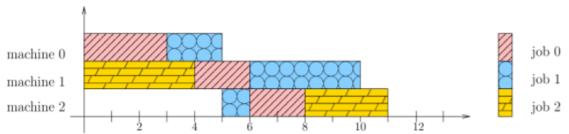


Figure 1. Gantt chart representation

4. CHICKEN SWARM OPTIMIZATION

The Chicken swarm optimization (CSO) was introduced by Meng, X.B. And al. [18] and inspired by the behavior of a chicken swarm while searching for food. Each swarm is divided into several groups, which comprises one rooster,hens and chicks. The hierarchical order in the swarm is established by the fitness value. We refer the number of roosters, hens, chicks and mother hens by RN, HN, CN and MN.

The position update equation of the rooster can be formulated as:

$$x_{i,j}^{t+1} = x_{i,j}^t * (1 + Randn(0, \sigma^2))$$
(6)

$$\sigma^2 = \begin{cases} 1, & if \ f_i \le f_k, \\ \exp\left(\frac{f_k - f_i}{|f_i| - \varepsilon}\right) & otherwise \ k \in [1, N], k \ne i \end{cases}$$
 (7)

where

 $Randn(0, \sigma^2)$ is a Gaussian distribution

 σ^2 is a standard deviation

The rooster index k is randomly selected from the rooster's group.

f is the fitness value of the corresponding x.

The position update equation of the hen can be formulated as bellow:

$$x_{i,j}^{t+1} = x_{i,j}^{t} + S1 * Rand * (x_{r1,j}^{t} - x_{i,j}^{t}) + S2 * Rand * (x_{r2,j}^{t} - x_{i,j}^{t})$$
(8)

and

S1=
$$\exp(\left(\frac{f_i - f_{r_1}}{|f_i| + \varepsilon}\right))$$
 and S2= $\exp((f_{r_2} - f_i))$

where $Rand \in [0, 1]$, r_1 is the index of the rooster and r_2 is the index of a random chicken from the swarm (where $r1 \neq r2$).

Finally, the position update equation of the chick is formulated in [19] as follows:

$$x_{i,j}^{t+1} = W * x_{i,j}^t + FL * (x_{m,j}^t - x_{i,j}^t)) + C * (x_{r,j}^t - x_{i,j}^t))$$

$$(9)$$

Where W is a self-learning factor for chicks, $FL \in [0,2]$ is a randomly selected parameter to refer to the relationship between the chicks and its mother with the index m where $m \in [1,N]$. Otherwise,C is a learning- factor from the rooster with the index r.

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5. ADAPTATION OF CHICKEN SWARM ALGORITHM TO JOB SHOP SCHEDULING PROB-LEM

During the discretization of the original version of the chicken swarm algorithm in order to solve the jop shop scheduling problem, the redefinition of operators is represented by the subtraction \ominus , the multiplication \otimes and the addition \oplus used in the original version [19].

Furthermore, we used the uniform crossover (UX) [20] in the position update equation of hens and chicks for the movement towards the leaders of groups and the sequential constructive crossover (SCX) [21] to simulate the movement towards the neighbors. \ominus operator represents the crossover operator and \otimes operator as applying the chosen crossover to the equation.the addition operator indicates that the randomly chosen crossover is applied to the movement. The application of UX and SCX ensure the competition between groups in the swarm.

As well, we integrate the 2-opt neighborhood operator to realize the auto-improvement mechanism in the position equation of the roosters and the chicks. In this new adaptation each schedule of a group is chosen randomly. The MeCSO in pseudo-code is represented by algorithm 1.

Algorithm 1: MeCSO for jssp

- 1. Initialize P the size of swarm
- 2. Generate P chickens
- 3. initialize parameters: P, G, FL, C and w.
- 4. Evaluate the fitness values at t=0 for each chicken
- 5. Rank and establish a hierarchal order
- 6. Create groups and assign chicks to mother-hens
- 7. Update the position by equations 6,8 and 9
- 8. Update the new solution if the fitness value is better.
- 9. Rank if G is reatched until stop criterion.
- 10. Return results of MeCSO

6. EXPERIMENTAL RESULTS AND DISCUSSION

6.1. Experimental environment

The proposed algorithm MeCSO was coded in python and run on a DELL in visual studio 2017 and simulated with Intel(R) Core(TM) i7-6500 U CPU 2.5GHZ (4 CPUs) 2.6 GHz and 16.00 GB of RAM and Microsoft Windows 10 Professional (64-bit) operating system. The performance of MeCSO was tested on different instances of OR-Library [22] 20 times in 100 iterations.

6.2. Default parameters

The table 1 shows the parameter values used in the new adaptation MeCSO. We execute different tests on instances Abz5 and Orb1 in order to choose the values which guarantee to obtain good results and converge towards the global optimum \cdot

Table 1. The Parameters for the Memetic-CSO Algorithm

Parameters of MeCSO	Values
P : Population size	500
RN: Number of roosters (%)	12
HN: Number of hens (%)	25
CN: Number of chicks (%)	63
G: Number of iterations to update the algorithm	10
W : Self-learning factor	0.5
FL: Learning factor from the mother hens	0.4
C: Learning factor from the rooster	0.65

where

$$CN = P - (NR + NH) \tag{10}$$

6.3. RESULTS AND DISCUSSION:

We applied MeCSO on some instances of OR-library, the table 2 summarizes the obtained results of 20 runs. The first column represents different instances instance in OR-Library, the second column indicates the best Known solution (BKS) ,the third column describes the average of the best found solution δ_{avg} , the remaining columns represent the measures use to perform the quality of the solution. The proposed algorithm MeCSO allows to find the best-known solution about 51.08 % from all tested instances.

Table 2. Numerical Results by MeCSO Applied to Some Instances of OR-library

$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	Numerical Results by MeCSO Applied to Some Instances of OR-						
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	$n \times m$			δ_{avg}	$\overline{T_{avg}(s)}$	$\overline{\operatorname{Err}(\%)}$	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Abz5	1234		139	0.068	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Abz6	943	949	521	0.821	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb1	1059	1093	807	1.045	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb2	888	907	157	0.981	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb3	1005	1011	408	0.056	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb4	1005	1024	101	0.328	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb5	887	891	633	0.766	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb6	1010	1016	412	0.831	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb7	397	402	595	0.907	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb8	899	907	276	0.837	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Orb9	934	944	166	0.741	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	6×6	Ft06	55	55	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×10	Ft10	930	939	102	0.801	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×5	LA01	666	666	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×5	LA02	655	655	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×5	LA03	597	599	21	0.086	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×5	LA04	590	590	3	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	10×5	LA05	593	593	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	15×5	LA06	926	926	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	15×5	LA07	890	890	2	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	15×5	LA08	863	863	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	15×5	LA09	951	951	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	15×5	LA10	958	958	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	20×5	LA11	1222	1222	2	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	20×5	LA12	1039	1039	1	0	
20×5 LA15 1207 1207 3 0 10×10 LA16 945 950 12 0.551 10×10 LA17 784 784 84 0 10×10 LA18 848 851 534 0.021 10×10 LA19 842 850 126 0.352 10×10 LA20 902 911 782 0.045	20×5	LA13	1150	1150	1	0	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	20×5	LA14	1292	1292	1	0	
10×10 LA17 784 784 84 0 10×10 LA18 848 851 534 0.021 10×10 LA19 842 850 126 0.352 10×10 LA20 902 911 782 0.045	20×5	LA15	1207	1207	3	0	
10×10 LA17 784 784 84 0 10×10 LA18 848 851 534 0.021 10×10 LA19 842 850 126 0.352 10×10 LA20 902 911 782 0.045	10×10	LA16	945	950	12	0.551	
10×10 LA19 842 850 126 0.352 10×10 LA20 902 911 782 0.045	10×10		784	784	84	0	
10×10 LA20 902 911 782 0.045	10×10	LA18	848	851	534	0.021	
10×10 LA20 902 911 782 0.045	10×10	LA19	842	850	126	0.352	
15 × 10 LA21 1046 1085 477 0.755	10×10		902	911	782	0.045	
	15 × 10	LA21	1046	1085	477	0.755	

The mathematical formulation of the percentage of error ERR (11) is represented as bellow:

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[H]

$$Err = \frac{(\delta_{avg} - BKS)}{BKS} \times 100 \tag{11}$$

where BKS is the best known value , δ_{avg} the average of the best found solution.

The proposed algorithm MeCSO seems to be promising to solve jssp in a reasonable time compared to GB algorithm [23] as represented in Figure 2. Furthermore, the algorithm allows to obtain good results in term of the global optimum compared to other algorithms from literature, such as [24] and [25] as represented in table 3 and GB algorithm [23] as represented in Figure 3.

Table 3. Average of the BFS of Some Algorithms in the Literature Compared to MeCSO

Instance	MeCSO	Bondal(GA) [25]
Udomsakdigool and Kachitvichyanukul [24]		
Abz5	1236	1339
-		
Abz6	949	1043
-		
Ft06	55	55
55		
Ft10	939	944
1099		
LA01	666	666
666		
LA02	655	658
716		
LA03	599	603
638		
LA04	590	590
619		
LA05	593	593
593		
LA16	950	977
1033		

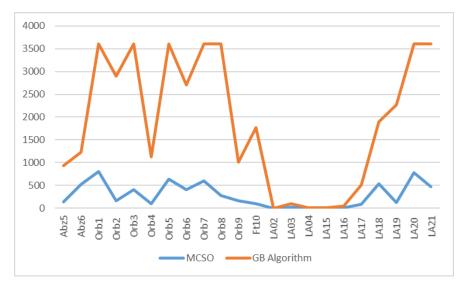


Figure 2. Average time (s) of MeCSO and GB algorithm

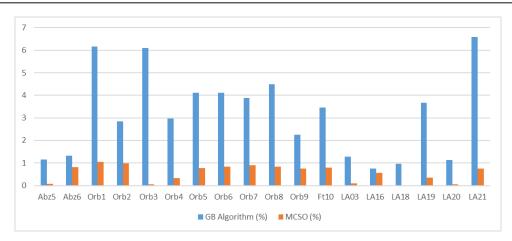


Figure 3. Err (%) of MeCSO and GB algorithm

7. CONCLUDING REMARKS

In this paper,we proposed a Memetic Chicken swarm optimization algorithm based on the original version of chicken swarm optimization (CSO) and 2-opt mechanism in order to solve the job shop scheduling problem. The empirical results show that MeCSO algorithm is efficient to solve this type of problem than the other algorithms from literature such as GB algorithm and GA in term of the quality of solutions and the computing time. In further research, we suggest to integrate the simulating annealing with the chicken swarm algorithm to ensure the redistribution of the swarm.

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