# Minimal cellular resolutions of the edge ideals of forests

Margherita Barile

Antonio Macchia\*

Dipartimento di Matematica Università degli Studi di Bari Aldo Moro Via Orabona 4, 70125 Bari, Italy Fachbereich Mathematik und Informatik Freie Universität Berlin Arnimallee 2, 14195 Berlin, Germany

margherita.barile@uniba.it

macchia.antonello@gmail.com

Submitted: Jun 20, 2019; Accepted: May 21, 2020; Published: May 29, 2020 © The authors. Released under the CC BY-ND license (International 4.0).

#### Abstract

We present an explicit construction of minimal cellular resolutions for the edge ideals of forests, based on discrete Morse theory. In particular, the generators of the free modules are subsets of the generators of the modules in the Lyubeznik resolution. This procedure allows us to ease the computation of the graded Betti numbers and the projective dimension.

Mathematics Subject Classifications: 13A15, 13C10, 13D02, 05C05

# 1 Introduction

Cellular resolutions (free resolutions supported on a CW-complex) were introduced for monomial modules by Bayer and Sturmfels in [4]. Since then, they have been studied by several authors (among others, [2], [5], [7], [8], [10], [12], [16], [19], [22]). In some cases they turn out to be minimal (see, e.g., [5] for generic and shellable monomial modules, [8] for the powers of the edge ideals of paths, [16] for the well-known Eliahou-Kervaire resolution for stable ideals, [19] for the matroid ideal of a finite projective space). We also know, however, that a minimal cellular free resolution need not exist: a class of counterexamples was provided by Velasco in [22]. A special type of (regular) cellular resolutions are the simplicial resolutions, first considered by Bayer, Peeva, and Sturmfels [3]. Two famous (in general, non minimal) examples for monomial ideals are the *Taylor resolution* [20] and its refinement called the *Lyubeznik resolution* [15]: an essential overview of the topic is contained in [17]. In [5], Batzies and Welker use Chari's reformulation [6] of Forman's discrete Morse Theory [11] as a tool for obtaining new cellular resolutions from the Taylor

<sup>\*</sup>Supported by the Einstein Foundation Berlin under grant EVF-2015-230.

complex, where they also show that the Lyubeznik resolution can be derived in this way. Later on, Àlvarez Montaner, Fernández-Ramos, and Gimenez [2] applied similar techniques for developing an algorithmic transformation (called *pruning*) of the Taylor resolution, which, however, does not always produce a minimal resolution.

In our paper we present an explicit discrete Morse theoretical construction of a minimal cellular free resolution for any edge ideal of a forest, i.e., for any ideal in a polynomial ring over a field whose generators are the squarefree quadratic monomials corresponding to the edges of a simple, undirected, and acyclic graph. The sets of generators of the free modules are explicitly described as special subsets of the sets of generators of the modules in the Lyubeznik resolution. They can be determined in two ways: we present both a selection procedure (see the steps (I)-(V) in Section 3) and a combinatorial characterization (see Proposition 13). These two methods relevantly facilitate the computation of the (multi)graded Betti numbers and the projective dimension, and, in particular, they allow a transparent and constructive approach to some of the formulas given by Jacques [13] and Kimura [14]. Also note that, unlike the algorithmic approach applied in [21], our method is direct, not recursive, and also totally different and independent with respect to the one developed for the quadratic monomial ideals considered by Horwitz [12]. The resolution we determine is the same over any field.

Since, for edge ideals, the minimal free resolutions are additive with respect to the connected components of the graph, we can restrict our attention to the edge ideals of trees.

# 2 Preliminaries

Let T be a tree on the vertex set  $\{x_1, \ldots, x_n\}$ , which is a set of indeterminates over the field K. Let  $R = K[x_1, \ldots, x_n]$  and let  $I = I(T) \subset R$  be the *edge ideal* of T, i.e., the ideal generated by the (edge) monomials  $x_i x_j$  such that  $\{x_i, x_j\}$  (also denoted by  $x_i x_j$ ) is an edge of T.

Fix a vertex, say  $x_1$ . Let d be the maximum distance that a vertex of T can have from  $x_1$ . For all  $i = 0, \ldots, d$ , call  $x_1^{(i)}, \ldots, x_{s_i}^{(i)}$  the vertices lying at distance i from  $x_1$  (so that, in particular,  $s_0 = 1$  and  $x_1^{(0)} = x_1$ ). We will say that these vertices have rank i: each of these vertices is connected to  $x_1$  by a (unique) path of length i (formed by i + 1 vertices).

Neighbors always have consecutive ranks. In the line graph  $\overline{T}$ , let  $((i)_p, (i+1)_q)$  denote the vertex that, in  $\overline{T}$ , corresponds to the edge  $x_p^{(i)} x_q^{(i+1)}$  of T. The vertices of  $\overline{T}$  of the form  $((i)_p, -)$  or  $(-, (i)_p)$  form a non-empty complete subgraph of  $\overline{T}$ , whose vertices represent the edges of T that contain the vertex  $x_p^{(i)}$ . It will be called a K-subgraph of  $\overline{T}$ , denoted by  $K_p^{(i)}$ , and  $(i)_p$  will be called its *index*. Similar considerations apply to  $(i+1)_q$ . The K-subgraphs are the maximal non-empty complete subgraphs of  $\overline{T}$ . Note that, if  $x_p^{(i)} x_q^{(i+1)}$ is a non-terminal edge monomial (i.e., if none of its endpoints,  $x_p^{(i)}$  and  $x_q^{(i+1)}$ , has degree 1 in T), then  $((i)_p, (i+1)_q)$  is the only common vertex of  $K_p^{(i)}$  and  $K_q^{(i+1)}$ .

Remark 1. Let  $i \ge 0$ . For every index  $(i+1)_q$  there is exactly one index of the form  $(i)_p$  such that  $((i)_p, (i+1)_q)$  is a vertex of  $\overline{T}$  (i.e., such that  $x_p^{(i)} x_q^{(i+1)}$  is an edge of T). In fact,

if  $x_p^{(i)}$  and  $x_{p'}^{(i)}$  are different vertices of T of rank i, there are paths of length i connecting  $x_1$  to each of them. If these vertices were both adjacent to the vertex  $x_q^{(i+1)}$  (of rank i+1), then T would contain a cycle. We will say that  $x_p^{(i)}$  is the only predecessor of  $x_q^{(i+1)}$ .

On the monomials of R fix the lexicographic order induced by the following arrangement of the indeterminates:

$$x_1^{(0)} > x_1^{(1)} > \dots > x_{s_1}^{(1)} > x_1^{(2)} > \dots > x_{s_2}^{(2)} > \dots > x_{s_d}^{(d)}.$$

The arrangement thus obtained will be called the *sequence of generators*. The same order will be considered on the set of indices of the K-subgraphs.

Given a minimal graded free resolution of R/I

$$0 \longrightarrow \oplus_j R(-j)^{\beta_{s,j}} \longrightarrow \cdots \longrightarrow \oplus_j R(-j)^{\beta_{2,j}} \longrightarrow \oplus_j R(-j)^{\beta_{1,j}} \longrightarrow R \longrightarrow 0$$

the coefficients  $\beta_{i,j}(R/I) = \beta_{i,j}$  are called the graded Betti numbers of R/I. Its length s is called the projective dimension, pd(R/I), of R/I.

#### 3 The resolution

#### 3.1 The generators of the modules

Let S be the sequence of generators of I. Any subsequence of S will be called a *symbol*, and will be written in parentheses as a list of pairwise and distinct elements, which are separated by commas.

Given a symbol  $u = (\mu_1, \ldots, \mu_r)$ , r will be called the *length* of u, and denoted by |u|. Every subsequence of u will be called a *subsymbol* of u (or a symbol *contained* in u). We will also say that  $\mu_1, \ldots, \mu_r$  are the *elements* of (belong to) u. We will thus treat u as an ordered set. We will set  $lcm(u) = lcm(\mu_1, \ldots, \mu_r)$ .

We recall the following definition, which Lyubeznik [15] gave for an arbitrary monomial ideal.

**Definition 2.** A symbol  $u = (\mu_1, \ldots, \mu_r)$  is called *L*-admissible if  $\mu_q$  does not divide  $\operatorname{lcm}(\mu_{i_h}, \mu_{i_{h+1}}, \ldots, \mu_{i_t})$  for any h < t such that  $q < i_h$ . It is called *reduced* if  $\mu_q$  does not divide  $\operatorname{lcm}(\mu_1, \ldots, \hat{\mu_q}, \ldots, \mu_r)$  for all q such that  $1 \leq q \leq r$ .

Remark 3. In the special case that we are considering here, i.e., when I is generated by squarefree monomials of degree two, the condition of L-admissibility can be restated in the following simpler form:  $\mu_q$  does not divide any product  $\mu_{i_h}\mu_{i_k}$  for any h, k such that  $h < k \leq t$  and  $q < i_h$ . Similarly, the condition of being reduced translates into:  $\mu_q$  does not divide any product  $\mu_{i_h}\mu_{i_k}$  for any h, k such that  $q \neq i_h, q \neq i_k$ . Thus being reduced implies L-admissibility.

Set  $L_0 = R$  and, for all r = 1, ..., |S|, let  $L_r$  be the free *R*-module generated by all *L*-admissible symbols of length *r*. Define the map  $\delta_r : L_r \to L_{r-1}$  by setting

$$\delta_r((\mu_{i_1},\ldots,\mu_{i_r})) = \sum_{j=1}^r (-1)^{j+1} \frac{\operatorname{lcm}(\mu_{i_1},\ldots,\mu_{i_r})}{\operatorname{lcm}(\mu_{i_1},\ldots,\widehat{\mu_{i_j}},\ldots,\mu_{i_r})} (\mu_{i_1},\ldots,\widehat{\mu_{i_j}},\ldots,\mu_{i_r}).$$

Then one has the following

**Theorem 4** ([15], p. 193). The complex

$$0 \longrightarrow L_s \xrightarrow{\delta_s} L_{s-1} \xrightarrow{\delta_{s-1}} \cdots \xrightarrow{\delta_1} L_0 \longrightarrow 0 \tag{(\star)}$$

is a free resolution of R/I.

The resolution  $(\star)$  is called a *Lyubeznik resolution* of *I*.

The Lyubeznik resolution is a refinement of the Taylor resolution, whose rth module is generated by all symbols of length r, and in general it is not minimal. For all indices r we will determine a submodule  $F_r$  of  $L_r$  such that the resulting resolution is always minimal for the edge ideals of trees. The differentials, however, will have to be redefined by means of discrete Morse theory. The submodules  $F_r$  are generated by special Ladmissible symbols, which we will call F-admissible. We now start the description of a procedure for selecting these symbols. The first steps are presented here below.

- (I) Select a descending sequence of indices  $(i_1)_{p_1}, \ldots, (i_t)_{p_t}$  that does not contain any pair of indices corresponding to adjacent vertices in T.
- (II) Pick all edge monomials corresponding to the vertices of the K-subgraphs  $K_{p_1}^{(i_1)}, \ldots, K_{p_t}^{(i_t)}$ .
- (III) In the symbol thus obtained, cancel all monomials  $\mu$  fulfilling the following condition:  $\mu$  corresponds to a vertex of  $K_{p_h}^{(i_h)}$  and is not coprime with respect to an element  $\nu$ of the symbol corresponding to a vertex of  $K_{p_k}^{(i_k)}$  for some k > h.
- (IV) Consider all subsymbols of the symbols thus obtained.

**Definition 5.** A symbol arising from the procedure (I)–(IV) will be called *almost F*admissible. The set of the monomials of an almost *F*-admissible symbol that are divisible by  $x_{p_j}^{(i_j)}$  - i.e., those corresponding to the vertices of  $K_{p_j}^{(i_j)}$  - will be called the  $(i_j)_{p_j}$ -block of u. The number  $(i_j)_{p_j}$  will be called its *index*.

Remark 6. Note that the sequence of the indices of the blocks of an almost F-admissible symbol u is not always uniquely determined. This is particularly evident in the case where u is formed by one single monomial  $x_p^{(i)} x_q^{(i+1)}$  not corresponding to a terminal edge of T; then u can be, of course, indifferently assigned to the  $(i)_p$ -block or to the  $(i+1)_q$ -block of u.

Remark 7. As a consequence of the selection performed at step (III), any two monomials of an almost F-admissible symbol u belonging to different blocks are coprime. In particular, if two distinct monomials of u are divisible by  $x_p^{(i)}$ , then they belong to the  $(i)_p$ -block of u.

**Proposition 8.** Every almost *F*-admissible symbol is reduced. In particular it is *L*-admissible.

Proof. Let u be an almost F-admissible symbol. In view of Remark 7, if the edge monomial  $\mu = x_p^{(i)} x_q^{(i+1)}$  belongs to the  $(i)_p$ -block of u, no other monomial of u is divisible by  $x_q^{(i+1)}$ . If it belongs to the  $(i+1)_q$ -block of u, no other monomial of u is divisible by  $x_p^{(i)}$ . Thus  $\mu$  cannot divide the product of any other two monomials of u.

**Definition 9.** Let x, y, z, w be four distinct vertices of T such that xy and zw are elements of a symbol u, where xy > zw. If  $xz \in I$ , we say that xz is the *bridge* between xy and zw. In this case, we will say that xy and zw form a gap in u if  $xz \notin u$ , no other monomial of u other than zw is divisible by w, and no monomial smaller than zw is divisible by y. The monomial xz will be called the *bridge* of the gap.

We will say that a symbol contains a bridge if it contains a triple of monomials xy, zw, xz, which occurs if and only if the symbol is not reduced.

Remark 10. (a) In the assumption of Definition 9, the condition  $xz \notin u$  is always fulfilled if u is reduced.

(b) In Definition 9 we are not assuming that x > y. In this case the inequality xy > zw implies that x > z. The same is true, however, if y > x, since y is then the only predecessor of x. Hence x is always the predecessor of z, which, in turn, implies that z > w. Finally, we deduce that xz > zw.

We now present the last step of the procedure.

(V) Discard all symbols that contain a gap.

**Definition 11.** The symbols obtained after completion of the procedure (I)–(V) will be called *F*-admissible (for T).

Remark 12. One may ask whether the steps (III) and (IV) could be interchanged. This would enlarge the set of almost F-admissible symbols: if the indices  $(i_h)_{p_h}$  and  $(i_k)_{p_k}$  with k > h have been selected at step (I), and  $\mu = xy$  (where  $x = x_{p_h}^{(i_h)}$ ) and  $\nu = xz$  (where  $z = x_{p_k}^{(i_k)}$ ) are monomials corresponding to vertices of  $K_{p_h}^{(i_h)}$  and  $K_{p_k}^{(i_k)}$ , then, according to step (III),  $\mu$  should be deleted. This would not occur, however, if, after selecting a subsymbol (as in step (IV)),  $\nu$  no longer appeared. Suppose that the  $(i_k)_{p_k}$ -block is still represented in the subsymbol by some monomial zw (otherwise we could regard the index  $(i_k)_{p_k}$  as not selected in step (I)). We would thus have a symbol that contains  $\mu$  but not  $\nu$ , a possibility that is forbidden by the original procedure. But the point is that xy and zw form a gap, so that the subsymbol would be discarded in step (V). In other words, performing steps (III) and (IV) in this order makes the selection procedure more efficient. The F-admissible symbols can be characterized combinatorially in the following way.

**Proposition 13.** A symbol is *F*-admissible if and only if it is reduced and has no gaps.

*Proof.* The "only if" part is a consequence of Definition 11 and Proposition 8. For the "if" part, let u be a symbol that is reduced and has no gaps. We show that its blocks fulfil the requirements set by steps (I) and (III) of the above procedure. First suppose that the elements of u are pairwise coprime. Then the requirement in (III) is fulfilled. Let x, y, z, w be vertices of T such that xy and zw are two distinct elements of u, where xy > zw. Then neither x nor y is a neighbor of z or w, which immediately yields the requirement in (I): if this were not true, and, say, xz were an edge monomial, then xy and zw would form a gap (against our assumption) because no other monomial of u is divisible by y or w.

Now assume that, for some indeterminate  $z = x_p^{(i)}$ , u contains more than one monomial divisible by z, say  $zw_1$  and  $zw_2$  are two distinct elements of u. Then no other monomial of u is divisible by  $w_1$  or  $w_2$ , because otherwise u would not be reduced. We show that, up to changes in the way in which the monomials of u are assigned to its blocks, no monomial of u belonging to a block preceding the  $(i)_p$ -block can be divisible by z. Suppose by contradiction that some element of the  $(i')_{p'}$ -block of u, where  $(i')_{p'} > (i)_p$ , is divisible by z. Let  $x = x_{p'}^{(i')}$ . Then  $xz \in u$ , and no other monomial of u can be divisible by x, because otherwise u would not be reduced. Hence the monomial xz can be assigned to the  $(i)_p$ -block of u, and the  $(i')_{p'}$ -block disappears. If this transformation is applied repeatedly, scanning the blocks of u from right to left, then after a finite number of steps the condition required at step (III) is achieved and, in particular, the monomials belonging to different blocks are coprime.

Now assume that the sequence of indices of the blocks of u contains two indices  $(i')_{p'} > (i)_p$  corresponding to consecutive vertices x and z. Let xy and zw be elements of the  $(i')_{p'}$ -block and the  $(i)_p$ -block, respectively. Then no other monomial of u can be divisible by y or w, so that xy and zw form a gap. But this contradicts our assumption. Hence u also fulfils the requirement in (I).

#### 3.2 The complex

For all indices r, let  $F_r$  be the free R-module generated by the F-admissible symbols of length r. We show that  $F_r$  is the rth module of a minimal graded free resolution of R/I: this is the cellular resolution derived from the Taylor resolution by means of the construction described in Section 1 of [5]. Note that the Taylor resolution can be viewed as a simplicial complex, and therefore as a  $\mathbb{Z}^n$ -graded regular CW-complex  $(X, \operatorname{gr})$ , where, for all indices r, the r-cells are the symbols of length (dimension) r, and the  $\mathbb{Z}^n$ -grading is defined as follows:  $\operatorname{gr}(\mu_1, \ldots, \mu_r) = \sum_{k=1}^s e_{j_k}$  if  $\operatorname{lcm}(\mu_1, \ldots, \mu_r) = \prod_{k=1}^s x_{j_k}$  (here  $e_j$ denotes the *j*th element of the canonical basis of  $\mathbb{Z}^n$ ). We also endow the Cartesian product  $\mathbb{Z}^n$  with the usual termwise defined partial order, with respect to which, for any two symbols u and v, we have  $\operatorname{gr}(u) \leq \operatorname{gr}(v)$  if and only if  $\operatorname{lcm}(u)$  divides  $\operatorname{lcm}(v)$ .

We then consider the directed graph  $G_X$  on X whose set of edges  $E_X$  is formed by the directed edges  $u \to u'$  such that  $u' \subset u$  and the lengths of u and u' differ by one. We consider the set of symbols that are not F-admissible, i.e., according to Proposition 13, that either contain a gap or are not reduced.

With respect to the notation of Definition 9, given a symbol u containing the gap formed by xy > zw (whose bridge is xz), we say that a bridge  $\lambda$  of u follows this gap if  $xz > \lambda$ .

We will say that a symbol is of type 1 if it contains a gap that is not followed by any bridge. We say that a symbol is of type 2 if it is non-*F*-admissible and not of type 1. Note that any symbol of type 2 has no gaps (and then it is not reduced) or it contains a bridge that follows all its gaps. In any case it contains a bridge.

Let A be the set of directed edges  $u \to u'$  such that u' is of type 1 and u is obtained from u' by inserting the monomial (say xz) which is the bridge of one of its gaps (say xy > zw) and is the smallest among the bridges of its gaps. We will call this operation bridge insertion.

#### Lemma 14. With the notation from above, the following properties hold:

- The symbol u is of type 2 (thus u can never appear as the first vertex of an edge of A).
- 2) The symbol u' is obtained from u by omitting its smallest bridge (thus u appears in exactly one edge of A).
- 3) If a symbol is of type 2, it can be obtained by bridge insertion from a symbol of type 1 (thus any symbol of type 2 appears in some edge of A).

*Proof.* Recall that, u is obtained from u' by inserting the monomial xz which is the bridge of one of its gaps, say xy > zw, and is the smallest among the bridges of its gaps.

1) Suppose by contradiction that u contains some gap ab > cd (with bridge bc) which is not followed by any bridge. Thus xz > bc. Moreover, by definition of gap, no monomial of the form de (where d > e) can belong to u. If both ab and cd belonged to u', then they would form a gap also in u', which would contradict the assumption on u'. Hence one of ab and cd is xz. But  $cd \neq xz$ , because, by Remark 10 (b), bc > cd. Thus  $cd \in u'$ and we must have that ab = xz. Then b = x or b = z. Recall that by Remark 10 (b) we have x > z > w and b > c > d. If bc = xc, then from xz > xc we deduce that z > c, so that zw > cd and, by definition of gap (the one formed by xy > zw in u'), no monomial containing y can follow cd. This implies that xy > cd is a gap in u'. But its bridge xc is smaller than xz, a contradiction. Now suppose that bc = zc, so that xz > zc. Note that x > z implies z > c. Thus zw > cd is a gap in u' (recall that no monomial of u' other than zw is divisible by w), and its bridge zc is smaller than xz, a contradiction.

2) By assumption, xz is smaller than all bridges of u'. We have to show that it is smaller than all monomials of u' that are bridges in u, but are not such in u'. The point is that the insertion of xz can produce new bridges, since a monomial  $\mu$  of u' of the form axor az could become a bridge between xz and some other monomial ab of u'. We show that in this case  $\mu > xz$ . First suppose that  $\mu = az$ . Note that  $a \neq w$ , because, by definition of gap, zw is the only monomial of u' containing w, and  $ab \in u'$ . But then az is already a bridge in u' (between zw and ab). Now suppose that  $\mu = ax$ . If  $a \neq y$ , then ax is a bridge in u' between xy and ab. So assume that a = y, i.e.,  $\mu = xy$ . If x > y, then y > b. On the other hand, by definition of gap (applied to xy > zw), we have yb > zw, which implies that y > z, so that  $\mu > xz$ . If y > x, then x > z implies, once again, that  $\mu > xz$ .

3) We show that, if in some symbol u of type 2 we cancel the smallest bridge, say xz, which is the bridge between xy and zw (where xy > zw), then, in the resulting monomial u', these two monomials form a gap (which, of course, is not followed by any bridge). We have to verify the following two conditions:

- i) u does not contain any other monomial of the form wb. Suppose by contradiction that u contains such a monomial. Then zw is the bridge between xz and wb. But from Remark 10 (b) we know that xz > zw, against the minimality of xz.
- ii) u does not contain any monomial of the form ya that is smaller than zw. Suppose that u contains some monomial ya. We show that ya > zw. This is clear if y > x, because x > z. So assume that x > y, whence y > a. Since xy is the bridge between xz and ya, we have that xy > xz. Thus y > z, which immediately implies that ya > zw, as desired.

**Proposition 15.** The graph  $G_X^A$  with edge set

$$E_X^A = (E_X \setminus A) \cup \{u' \to u \mid u \to u' \in A\}$$

does not contain any directed cycle.

*Proof.* The (directed) edges of  $E_X^A$  are of the following two types:

- (a) the edges  $u_1 \rightarrow u_2$  not belonging to A, where  $u_2$  is obtained from  $u_1$  by deleting a monomial;
- (b) the edges  $u' \to u$  where u' is of type 1, u is of type 2 and u is obtained from u' by bridge insertion.

Note that, with respect to the above notation,  $|u_2| = |u_1| - 1$ , whereas |u| = |u'| + 1. It follows that any directed cycle of  $G_X^A$  must contain at least one edge of each type. More precisely, since every edge of type (b) (along which the length grows by 1) is followed by an edge of type (a) (along which the length drops by 1), every directed cycle consists of an alternating sequence of edges of types (a) and (b), and an alternating sequence of vertices of types 1 and 2.

Moreover,  $\operatorname{gr}(u') = \operatorname{gr}(u)$ , whereas  $\operatorname{gr}(u_1) \ge \operatorname{gr}(u_2)$  and equality holds if and only if  $u_2$ is obtained from  $u_1$  by deleting a bridge. Thus, in any directed cycle of  $G_X^A$  all vertices have the same degree and two consecutive vertices always differ by a bridge. Let C be a directed cycle of  $G_X^A$ . Let  $u' \to u$  be a directed edge of C of type (b), where u is obtained from u' by inserting the bridge xz of the gap xy > zw. Moreover, assume that, among all edges of type (b) of C, this edge is one for which xz is maximum. The cycle C also contains an edge  $v \to u'$  of type (a), where v is obtained from u' by inserting a bridge other than xz (because  $v \neq u$ ). In particular, xz does not belong to v. This implies that at some point of the directed path of C from u to v the bridge xz is deleted. The first edge of this path, say  $u \to u_1$ , is of type (a), and  $u_1 \neq u'$ , so that  $u_1$  is of type 1 and  $xz \in u_1$ . Hence in  $u_1$  there are two monomials ab > cd forming a gap with bridge bc (the smallest), where, by choice of u', bc < xz. Since by Remark 10 (b) we have bc > cd, we also have that  $cd \neq xz$ , whence  $cd \in u'$ , because  $cd \in u_1 \subset u$ .

Moreover, as we have seen in the proof of Lemma 14 1), the condition ab = xz is incompatible with the fact that xz is the smallest of the bridges of the gaps contained in u'. Hence ab and cd both belong to u', where, of course, they cannot form a gap, because xz > bc. Since  $u' \subset u$  (so that  $ab, cd \in u$ ), for the same reason we have that  $bc \notin u$  (otherwise bc would be a bridge of u smaller than xz), whence  $bc \notin u'$ . Thus the obstruction preventing ab > cd from forming a gap in u' must be due to the presence of some other monomials forbidden by the definition of gap, i.e., in u' there is some other monomial  $\mu = de$  (with d > e) or some monomial  $\mu = af$  that is smaller than cd.

Since  $\mu$  does not belong to  $u_1$ , it must have been deleted along the path from u to  $u_1$ , hence  $\mu$  is a bridge in u, so that  $\mu > xz > bc$ . These inequalities exclude the case  $\mu = de$ . If  $\mu = af$ , since cd > af, we have c > a, whence b > a, and a > f. Thus  $bc > \mu$ , a contradiction. This shows that  $G_X^A$  cannot contain any directed cycle.

As a consequence of Lemma 14 and Proposition 15, A is a so-called *acyclic matching* for  $G_X$ . According to [5, Proposition 1.2], this implies that there is a  $\mathbb{Z}^n$ -graded CW-complex (called the *Morse complex* of the Taylor resolution), which has the following two properties: its *r*-cells are in one-to-one correspondence with the *r*-cells of X not belonging to any edge of A (which are called *critical cells* in [5]), i.e., with the *F*-admissible symbols of length r, and it is homotopy equivalent to the Taylor resolution. This Morse complex supports a cellular resolution of R/I. An explicit description of its differentials will be provided later on. We first show its minimality. In view of [5, Corollary 7.6], our claim will follow from the next result.

#### **Lemma 16.** Let u and v be distinct F-admissible symbols. Then $gr(u) \neq gr(v)$ .

*Proof.* First note that whenever some reduced symbol is contained in some other reduced symbol, the grading of the former is strictly smaller than the grading of the latter. So assume that u and v are incomparable by inclusion. Note that it suffices to prove the claim under the assumption that  $gr(u) \ge gr(v)$ , i.e., that lcm(v) divides lcm(u). For every monomial ab of u (or v) such that a > b, we call b a successor of u (of v). Since every vertex of the tree T has at most one predecessor, whenever b is a successor of u, ab belongs to u. Therefore, since u is not contained in v, there is some successor of u that is not a successor of v. Let b (whose predecessor is a) be the smallest such successor of u. Then  $ab \in u$ , but  $ab \notin v$ . We show that ab does not divide lcm(v). This will immediately yield the claim.

Suppose by contradiction that ab divides lcm(v). Since b divides lcm(v), it follows that  $bc \in v$  for some c < b. Thus c divides lcm(u). If  $bc \notin u$ , then  $cd \in u$  for some d < c. Since ab > cd and these two monomials do not form a gap in u, we must have  $de \in u$  for some

e < d. Note that d and e are successors of u smaller than b, whence  $cd, de \in v$ . But this, together with  $bc \in v$ , would imply that v is not reduced. We thus conclude that  $bc \in u$ . On the other hand, a divides lcm(v), which implies that  $xa \in v$  for some  $x \neq b$ . The vertex x may be greater or smaller than a, but in any case xa > bc. Since these two monomials do not form a gap in v, we have one of the following cases:

- i)  $cd \in v$ , for some d < c. Then, as above, we conclude that  $cd \in u$ : it suffices to apply the argumentation developed for the monomials  $ab \in u$  and  $bc \in v$  to the monomials  $bc \in u$ , and  $cd \in v$ . But then all three monomials ab, bc, and cd belong to u, which is impossible.
- ii) b > x (in which case a > x) and  $xy \in v$  for some y < x. Then u contains some monomial divisible by x, either ax or xz, with z < x. If  $ax \in u$ , we conclude as above that  $xy \in u$ , which is incompatible with the fact that  $ab \in u$ . So suppose that  $xz \in u$ . Since ab > xz and these monomials do not form a gap in u, we have that  $zw \in u$  for some w < z. Since both z and w are successors of u smaller than b, it follows that xz and zw belong to v, which, together with  $ax \in v$ , would once again cause v to be not reduced.

In any case we have a contradiction.

*Remark* 17. The preceding lemma yields the so-called *dual version of Hochster's formula* (see, e.g., [18, Corollary 1.40]) on squarefree monomial ideals in the special case of edge ideals of trees: it proves that the nonzero multigraded Betti numbers all lie in squarefree degrees. It also gives the result in [9, Theorem 3.5], which Erey and Faridi proved for the more general class of simplicial forests: for every multidegree there is at most one nonzero Betti number, and this is equal to 1.

The following notation is taken from [5]. For every pair (u, u') of F-admissible symbols with r = |u| = |u'| + 1, call [u : u'] the coefficient of u' in  $\delta_r(u)$ . If the directed edge  $u \to u'$ belongs to A, we then set  $m(\{u, u'\}) = -[u : u']$ , otherwise we set  $m(\{u, u'\}) = [u : u']$ . Given a directed path  $P : u_0 \to u_1 \to \cdots \to u_t$  in  $G_X^A$  (a so-called gradient path from  $u_0$ to  $u_t$ ), we also set  $m(P) = \prod_{i=0}^{t-1} m(\{u_i, u_{i+1}\})$ . Note that  $gr(u_0) \ge gr(u_t)$ . We can now define the *r*th differential  $\partial_r : F_r \to F_{r-1}$  of our resolution. According to [5, Lemma 7.7], for every F-admissible symbol u of length r,

$$\partial_r(u) = \sum_{\substack{u' \subset u \\ |u'| = r-1}} [u:u'] \sum_{\substack{u'' F - \text{admissible,} \\ |u''| = r-1}} \sum_{\substack{P \text{ gradient path} \\ \text{from } u' \text{ to } u''}} m(P) \underline{x}^{\operatorname{gr}(u) - \operatorname{gr}(u'')} u''.$$
(1)

We have just established the following:

**Theorem 18.** The cellular resolution  $(F_r, \partial_r)$  is a minimal graded free resolution of R/I.

*Proof.* We just have to observe that minimality is ensured by Lemma 16. By [5, Corollary 7.6] and [5, Proposition 7.3], it is sufficient to show that  $\operatorname{gr}(u) \neq \operatorname{gr}(v)$  for all *F*-admissible symbols u and v such that |u| = |v| + 1 and either  $v \subset u$  or there exists a gradient path from a symbol  $u' \subset u$  of length |v| to v.

Remark 19. Note that, according to (1), the morphism  $\partial_r$  sends every *F*-admissible symbol u of length r to a linear combination (with monomial coefficients) of *F*-admissible symbols u'' of length r-1 such that  $\operatorname{lcm}(u'')$  divides  $\operatorname{lcm}(u)$ . Among these we find all u'' that are contained in u (those appearing in the definition of the morphism  $\delta_r$  of the Lyubeznik resolution, and for which the gradient path P consists of one single directed edge  $u \to u''$ ), but possibly some others. This will become evident in Example 21.

### 4 Examples and further remarks

We will first give a concrete example of how to determine the modules of the cellular minimal free resolution of the edge ideal of a tree. We will apply the selection procedure described in Section 3. In order to simplify our notation, we will replace the symbol  $x_p^{(i)}$  by the number *i* with *p* apices. In this way we will also write the index of the *K*-subgraph  $K_p^{(i)}$  and of the  $(i)_p$ -block.

**Example 20.** Let T be the tree on the vertices 0 > 1 > 1' > 2 > 2' > 2'' > 3 whose edges are 01, 01', 12, 12', 1'2'', 23, see Figure 1.



Figure 1: The tree T of Example 20

Its K-subgraphs are:

$$[0]: 01, 01'$$
  $[1]: 01, 12, 12'$   $[1']: 01', 1'2''$   $[2]: 12, 23.$ 

In the next table we list, for any choice of the sequences of indices in step (I) (written in square brackets), the edges of the corresponding K-subgraphs, as prescribed in step (II), and perform on them the cancellations indicated in (III). Then we consider, as in step (IV), all their subsymbols. Each horizontal section refers to the sequences of a given length (one in each column). In each column, the subsymbols will be arranged in descending order of length (denoted by r). We will avoid repetitions by omitting the subsymbols already obtained in the preceding sections, and replacing the others by a reference to the number of the column of the same section (expressed by (\*) or (\*\*)) where (in the same row) they appear for the first time. The monomials forming a gap are boxed, and the vertices of the bridge are in bold. According to step (V), the corresponding subsymbols have to be discarded.

	[0]		[1]	[1']	[2]				
	r = 3		(01, 12, 12')						
	r = 2	(01, 01')	(01, 12)	(01', 1'2'')	'') (12, 23)				
			(01, 12')						
			(12, 12')						
	r = 1	(01)	(*)						
		(01')	(12)	(*)	(**)				
			(12')	(1'2'')	(23)				
	[0,2]		[1, 1']		[1', 2]				
r = 4			01,01',12,1	12', 1'2'')	(01', 12, 1'2'', 23)				
r = 3	3 (Ø1,0	(11, 12, 23)	(01', 12, 1)	.2')	(*)				
			$(\overline{01', 12, 1'2''})$	)	(**)				
			(01', 12', 1'2')	′)	(01', 1'2'', 23)				
			(12, 12', 1'2'')	)	(12, 1'2'', 23)				
r = 2	2 (01'	, 12)	(*)		(*)				
	(01',	23)	(01', 12')		(*)				
			(12, 1'2'')		(**)				
			(12', 1'2'')		(1'2'', 23)				

If we count lengths r and degrees d in order to compute the graded Betti numbers  $\beta_{r,d}(R/I)$ , we exactly obtain the numerical resolution provided by CoCoA [1]:

$$0 \longrightarrow R(-6)^2 \longrightarrow R(-4) \oplus R(-5)^6 \longrightarrow R(-3)^6 \oplus R(-4)^4 \longrightarrow R(-2)^6 \longrightarrow 0.$$

Note that the vertex 0 has been chosen at random. The result is independent of this choice.

We next present two gradient paths from a given F-admissible symbol u of length r to two F-admissible symbols  $u''_1$  and  $u''_2$  of length r-1.

**Example 21.** Let T be the tree on the vertices 0 > 1 > 2 > 3 > 4 > 4' > 5 > 6, whose edges are 01, 12, 23, 34, 34', 45, 56, see Figure 2.



Figure 2: The tree T of Example 21

Then u = (01, 23, 34', 45, 56) is an *F*-admissible symbol of length 5. In the following table, the cancellations correspond to directed edges of type (a), the insertions to directed edges of type (b), according to the classification introduced at the beginning of the proof of Proposition 15. The first step is a cancellation that takes u to a symbol of type 1 (which is reduced and contains a gap). The following symbols are alternatively of types 2 and

1, of lengths r and r - 1, and the procedure stops eventually, when an edge of type (a) leads to an F-admissible symbol of length r - 1. The vertices in bold are those forming the bridge of the smallest gap in a symbol of type 1 (and thus forming the monomial that will be inserted in the subsequent step).

	u		01	L		2	23	3	84'			4	45	E.	66	
														$\downarrow$		
	type 1		01	1		2	23		<b>3</b> 4′		4		<b>1</b> 5	[		
											$\downarrow$					
	type 2		01	L	( 2		23	34'		(	(34)   4		45			
											Ŭ		$\downarrow$			
	$u_1''$		01	L		2	23	3	84'	ę	34	[				
ı	l	(	)1				23	3	34	/			45	5	56	3
									↓							
t	ype 1	(	)1				2	3					45	5	56	3
					$\downarrow$											
t	ype 2	(	)1	(	$\widehat{12}$	)	23	3					45	5	56	3
					$\smile$										$\downarrow$	
t	ype 1	(	)1	1	12		23	3					45	5		]
											$\downarrow$					
t	ype 2	(	)1	]	12		23	3			(34	)	45	<u>;</u>		
								Ļ								
ı	$\iota_2''$	(	)1	1	12			]			34		45	5		

From the description of the minimal cellular resolution of R/I(T) we can derive a formula to compute its projective dimension. Recall that a *leaf* or *free vertex* of a graph is a vertex of degree one.

Lemma 22. Let T be a tree, and let u be a maximal F-admissible symbol for T. Then

 $|u| = \#\{\text{leaves of } T \text{ in the blocks of } u\} + \#\{\text{blocks not in } u\}.$ 

*Proof.* Let u be a maximal F-admissible symbol. We show that u has length equal to the right-hand side of the formula in the statement.

Let (a) be a block of u that has been selected at step (I) and let ab be a free vertex of the block (a) in the line graph of T, i.e., b is a leaf of T. Then the monomial ab is not canceled at step (III) since no monomial divisible by a or b appears in any of the following blocks (c), with a > c.

Now, let (a) be a block not in u, i.e., that has not been selected at step (I). Let  $\mu_1, \ldots, \mu_k$  be the vertices of the block (a) that belong also to some other block that has been selected at step (I), say,  $\mu_i$  belongs to the block  $(b_i)$ , i.e.,  $\mu_i = ab_i$  for  $i = 1, \ldots, k$ , with  $b_1 > b_2 > \cdots > b_k$ . (Notice that if a vertex is not free and belongs to some block of u, then it also belongs to some non-selected block.) Then  $ab_1, \ldots, ab_{k-1}$  are removed at

step (III), while  $ab_k$  is not canceled (notice that  $ab_k$  is not a free vertex of the line graph, since it belongs to two blocks).

The case in which all vertices of the block (a) do not belong to any block of u cannot happen, otherwise the symbol u would not be maximal since adding the block (a) to u would produce a larger symbol.

Remark 23. Notice that the characterization in Lemma 22 identifies the maximal F-admissible symbols for T with the strongly disjoint set of bouquets of T introduced by Kimura in [14, Definition 2.3]. According to [14, Theorem 4.1], the maximum number of their flowers gives the projective dimension of R/I(T). In fact, the indices of the blocks of the maximal F-admissible symbols correspond to the roots of the bouquets and the other vertices (different from the roots) of the monomials remaining after the steps (I)-(V) correspond to the flowers of the bouquets. In particular, step (III) in the procedure above, together with the absence of gaps, ensures that the bouquets are pairwise strongly disjoint.

From the previous remark and [14, Theorem 4.1] it follows that:

**Corollary 24.** Let T be a tree. Then for every i, j,  $\beta_{i,i+j}(R/I(T))$  is the number of subsets  $W \subseteq V$  such that the induced subgraph  $T_W$  contains an F-admissible symbol of length i consisting of j blocks.

We can also recover Jacques' recursive formulas for the graded Betti numbers of a tree [13, Theorem 9.3.15], and for the projective dimension [13, Theorem 9.4.17], for which we can present a new, short, elementary, non-homological proof. We first retrieve the necessary notation, which is the one introduced in [13] on page 96, under (9.1.2). In [13, Proposition 9.1.1] it is proven that any tree T contains a vertex v such that among its neighbors  $v_1, \ldots, v_n$  at most one (say  $v_n$ ) has degree greater than 1. Call  $w_1, \ldots, w_m$  the neighbors of  $v_n$  other than v. Then call T' the subgraph of T obtained by eliminating the vertex  $v_1$ , and T'' the one obtained by eliminating all vertices  $v, v_1, \ldots, v_n$ . We can now prove:

**Theorem 25.** [13, Chapter 9] For all indices r, d we have

$$\beta_{r,d}(R/I(T)) = \beta_{r,d}(R/I(T')) + \sum_{j=0}^{n-1} \binom{n-1}{j} \beta_{r-(j+1),d-(j+2)}(R/I(T'')).$$

Moreover,

$$pd(R/I(T)) = max\{pd(R/I(T')), pd(R/I(T'')) + n\}.$$

*Proof.* We fix, on the vertices of T, an ordering with respect to which the smallest vertices are  $v_n > v > v_1 > \cdots > v_{n-1}$ . Consider the set of F-admissible symbols for T having length r and degree d. It includes the set of all F-admissible symbols for T' having length r and degree d. The complementary set is formed by the F-admissible symbols containing the monomial  $vv_1$ . Let u be such a symbol. Then  $u' = u \setminus \{v_n v, vv_1, \ldots, vv_{n-1}\}$  (if nonempty) is an F-admissible symbol for T''. Note that it is certainly reduced, and the elimination of  $v_n v$  could produce a gap only for n > 1 and if u contained some monomial of the form  $w_i v_n$ . But u cannot contain both  $w_i v_n$  and  $v_n v$  because by assumption it contains  $vv_1$ . Conversely, if u' is an F-admissible symbol for T'' of length r - (j + 1) with  $0 \leq j \leq n - 1$ , then, for every j-subset V of  $\{vv_2, \ldots, vv_n\}$ ,  $u = u' \cup \{vv_1\} \cup V$  is an Fadmissible symbol for T of length r. Note that u is reduced (because none of the vertices of the monomials added is a vertex of T'') and, moreover, if n > 1 and some monomial of the form  $w_i a$  appears in u', then the presence of  $vv_1$  prevents  $w_i a$  and  $v_n v$  from forming a gap. The symbol u has degree d if and only if u' has degree d - (j + 2). Hence, given an F-admissible symbol u' for T'' of length r - (j + 1) and degree d - (j + 2), we can construct the F-admissible symbol u in  $\binom{n-1}{j}$  different ways (as many as the possible choices of V). This shows the first identity. For the second identity it suffices to observe that the F-admissible symbol u has maximum length if it is of maximum length among those not containing  $vv_1$  or it is obtained by adding the whole set  $\{v_nv, vv_1, \ldots, vv_{n-1}\}$ to an F-admissible symbol u' for T'' having maximum length.

# Acknowledgements

The first author is gratefully indebted to Giandomenico Boffi for raising the question and for useful discussions. Both authors thank Volkmar Welker for his valuable hints on discrete Morse theory.

# References

- J. ABBOTT, A.M. BIGATTI, G. LAGORIO, CoCoA-5: a system for doing Computations in Commutative Algebra. Available at http://cocoa.dima.unige.it/ cocoalib.
- [2] J. ALVAREZ MONTANER, O. FERNÁNDEZ-RAMOS, P. GIMENEZ, Pruned cellular free resolutions of monomial ideals. J. Algebra 541 (2020), 126–145.
- [3] D. BAYER, I. PEEVA, B. STURMFELS, Monomial resolutions. Math. Res. Lett. 5 (1998), 31–46.
- [4] D. BAYER, B. STURMFELS, Cellular resolutions of monomial modules. J. Reine Angew. Math. 503 (1998), 123–140.
- [5] E. BATZIES, V. WELKER, Discrete Morse theory for cellular resolutions. J. Reine Angew. Math. 543 (2002), 147–168.
- [6] M. K. CHARI, On discrete Morse functions and combinatorial decompositions. Discrete Math. 217 (2000), 101–113.
- [7] A. DOCHTERMANN, A. ENGSTRÖM, Cellular resolutions of cointerval ideals. Math. Z. 270 (2012), 145–163.
- [8] A. ENGSTRÖM, P. NORÉN, Cellular resolutions of powers of monomial ideals. Preprint (2012), arXiv:1212.2146.

- [9] N. EREY, S. FARIDI, Multigraded Betti numbers of simplicial forests. J. Pure Appl. Algebra 218 (2014), 1800–1805.
- [10] G. FLØYSTAD, Cellular resolutions of Cohen-Macaulay monomial modules. J. Commut. Algebra 1 (2009), 57–89.
- [11] R. FORMAN, Morse theory for cell complexes. Adv. Math. **134** (1998), 1, 90–145.
- [12] N. HORWITZ, Linear resolutions of quadratic monomial ideals. J. Algebra 318 (2007), 981–1001.
- [13] S. JACQUES, *Betti numbers of graph ideals*. Ph.D. thesis, University of Sheffield, 2004.
- [14] K. KIMURA, Non-vanishingness of Betti numbers of edge ideals. In: Harmony of Gröbner bases and the modern industrial society, 153–168, World Sci. Publ., Hackensack, NJ, 2012.
- [15] G. LYUBEZNIK, A new explicit finite free resolution of ideals generated by monomials in an R-sequence. J. Pure Appl. Algebra 51 (1988), 1-2, 193–195.
- [16] J. MERMIN, The Eliahou-Kervaire resolution is cellular. J. Commut. Algebra (2010), 55–78.
- [17] J. MERMIN, Three simplicial resolutions. Progress in commutative algebra 1, 127– 141, de Gruyter, Berlin, 2012.
- [18] E. MILLER, B. STURMFELS, *Combinatorial commutative algebra*. Springer, New York, 2005.
- [19] I. NOVIK, Lyubeznik's resolution and rooted complexes. J. Alg. Comb. 16 (2002), 97–101.
- [20] D. TAYLOR, *Ideals generated by an R-sequence*. Ph.D. thesis, University of Chicago, 1966.
- [21] M-L. TORRENTE, M. VARBARO, Computing the Betti table of a monomial ideal: a reduction algorithm, J. Symb. Comp. 87 (2018), 87–98.
- [22] M. VELASCO, Minimal free resolutions that are not supported by a CW-complex. J. Algebra **319** (2008), 102–114.