Uniquely and 2-Uniquely Hamiltonian Graphs

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A) Preface

- Quick introduction to graph theory
- What is a Hamilton cycle?
 - Maximum uniquely Hamiltonian Graphs

B) A look at Maximum uniquely Hamiltonian graphs

• Verifying Hamiltonicity in Polynomial time

C) Extending the maximum uniquely Hamiltonian graph

- Adding cycles to a max unique Ham graph
- Proof of maximal edge case for 2-unique Hamiltonian graphs

Preface:

1) Graph theory developed by Leonhard Euler (1707 – 1783)

- Konigsberg bridge problem (represented Konigsberg as a graph) (1736)
- Simple representation of groups of objects sharing some relationship
- Found that counting node degrees (number of edges incident on a node) determines the existence of an Euler tour.





Travelling Salesman Problem (1930)

- Euler circuits are not efficient for a travelling salesman

1) Introducing Hamilton cycles

- Touch every node just once and return to the node from which you started.
- Analogous to Euler circuits, but concerned with nodes instead of edges.

2) Unfortunately for our salesman, Hamilton cycles are very hard to find

- NP-complete problem
- Arbitrary graphs hard to solve
- Easier for graphs with fewer cycles?



3) Maximum uniquely Hamiltonian graphs

- As many edges as we can fit into a graph and maintaining a single Hamilton cycle.

A look at maximum uniquely Hamiltonian graphs (MUHG's)

- 1) Sheehan, Entringer, and Barefoot's work
 - Sheehan identified maximal edge case for graphs with a single cycle (1977)
 - Entringer and Barefoot discovered family of MUHG's of size $2^{([n/2]-4)}$ (1981)
 - The configuration of these graphs permits a polynomial computational solving-time

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2) AAAlgorithm (AAAlg):
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while (any_deg > 2) {
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find(maxDegreeNode);
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if (maxDegreeNode isAdjacentTo 2 Deg2Nodes) {
 delete(all_other_edges);
 update(degree_values);
} end if;

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} end loop;
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Proof of correctness by induction:

Take the 7-node base case graph from the previous slide, we've seen that our algorithm works for this graph. It works similarly for the 8-node base case.

Now consider the n-2 case for graph G of order n > 8:

- Step 1: Locate global max degree node
- Step 2: remove all edges not connecting degree 2 nodes
- Step 3: Retract edges connecting previous global max node and nodes of degree 2.



We have reduced our graph to the (n-2) MUHG, and the proof is complete.

Extending the MUHG:

What about other graphs with few Hamilton cycles?

- Can we modify the MUHG to obtain a graph with 2 cycles?
- Adding edge A (carefully) adds an additional cycle to our graph:

Actually: the new number of edges, N + 1, is extremal for 2 cycles where N is the number of edges in a MUHG: N = $([n^2/4] + 1)$.

Proof for lower bound:

Take a look at the MUHG of size n = 8.

Now notice we can add one particular edge (edge A) to an extremal G_1 graph to produce a graph with exactly 2 Ham cycles. Therefore: Max(E(G₂)) \geq N + 1



Proof for upper bound:

Take the graph G_2 with exactly 2 Ham cycles $H_1 \& H_2$.Remove an edge e from G_2 that's in H_1 but not H_2 We get: $E(G_2) \le N + 1$.

We've proved both the upper bound: $E(G_2) \le N + 1$, as well as the lower bound: Max($E(G_2) \ge N + 1$. Therefore $E(G_2) = N + 1$ and our proof is complete.

An interesting discovery for graphs containing 2 cycles:

- Configuration for the n case carries up to n + 2 case!
- This fact lends these graphs to a polynomial time algorithm. But we won't cover that today.

G₂ : Blue and Red edges introduce 2 cycles.



Works cited and thanks!

- Sheehan, J. (1977), Graphs with exactly one hamiltonian circuit. J. Graph Theory, 1: 37-43. doi:<u>10.1002/jgt.3190010110</u>

-Barefoot, C.A. and Entringer, R.C. (1981), A census of maximum uniquely hamiltonian graphs. J. Graph Theory, 5:315-321. doi:10.1002/jgt.3190050313

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