Centrum voor Wiskunde en Informatica
Centre for Mathematics and Computer Science

J.W. de Bakker, J.-J.Ch. Meyer

Order and metric in the stream semantics of elemental concurrency

Computer Science/Department of Software Technology Report CS-R8638 December

The Centre for Mathematics and Computer Science is a research institute of the Stichting Mathematisch Centrum, which was founded on February 11, 1946, as a nonprofit institution aiming at the promotion of mathematics, computer science, and their applications. It is sponsored by the Dutch Government through the Netherlands Organization for the Advancement of Pure Research (Z.W.O.).

69 D41, 69 F32, 69 F33

ORDER AND METRIC IN THE STREAM SEMANTICS OF ELEMENTAL CONCURRENCY

J.W. de Bakker
Centre for Mathematics and Computer Science,
P.O. Box 4079, 1009 AB Amsterdam, The Netherlands

J.-J.Ch. Meyer
Department of Mathematics and Computer Science,
Free University of Amsterdam,
P.O. Box 7161, 1007 MC Amsterdam, The Netherlands

Two denotational semantics for a language with simple concurrency are presented. The language has parallel composition in the form of the shuffle operation, in addition to the usual sequential concepts including full recursion. Two linear time models, both involving sets of finite and infinite streams, are given. The first model is order-theoretic and is based on the Smyth order. The second model employs complete metric spaces. Various technical results are obtained relating the order-theoretic and metric notions. The paper culminates in the proof that the two semantics for the language considered coincide. The paper completes previous investigations of the same language, establishing the equivalence of altogether four semantic models for it.

1980 Mathematics Subject Classification: 68B10, 68C01
1982 CR Categories: D.3.1, F.3.2, F.3.3.
Keywords & Phrases: concurrency, denotational semantics, Smyth order, metric spaces, merge, infinitary languages
Notes: (i) The research of J.W. de Bakker is partially supported by ESPRIT project 415: Parallel Architectures and Languages for AIP: a VLSI Directed Approach
(ii) This report will be submitted for publication elsewhere.

INTRODUCTION

We present two denotational semantics for a language with simple concurrency, and prove their equivalence. The first semantics has an order-theoretic, the second a metric structure as underlying model. In the course of proving the equivalence theorem, a number of results are obtained relating the two structures which may be of some independent interest.

The first model will be based on the so-called Smyth order between sets of *streams* (in the sense of, e.g., [Br1,Br2]). This model was first developed in [M1,M2]. The second model introduces a *distance* between streams. In this way, the set of all streams is turned into a complete metric space, and familiar tools such as Banach's fixed point theorem become available. The metric model was first presented in [BBKM]; essential inspiration for it was provided by [Ni].

Both models are of what has been called the 'linear time' variety. They are built on (sets of) sequences rather than on tree (-like) objects. For an overview of situations where the latter - also called 'branching time' - approach is preferable or even necessary, we refer to [BKMOZ]. Briefly, once notions such as deadlock or global nondeterminacy are covered, branching time models or variations along the lines of ready or failure sets (see [OH] for a systematic treatment) are required.

Report CS-R8638
Centre for Mathematics and Computer Science
P.O. Box 4079, 1009 AB Amsterdam, The Netherlands

In the present paper we restrict ourselves to a very simple setting. The language $\mathcal{L}$ which we investigate has the familiar sequential notions (elementary or atomic actions, sequential composition), and in addition recursion, nondeterministic choice and parallel composition specifying the interleaving or *merge* of (sequences of) elementary actions. No forms of synchronization or communication are included: $\mathcal{L}$ is, indeed quite elementary. The motivation for its study is primarily that we are able to obtain an exhaustive analysis of its various semantic models - more about this in a moment -, rather than its intrinsic semantic interest. Still, we believe that the notions of recursion and merge are both fundamental in (the nature of) parallel computation, justifying our terminology of *elemental* concurrency.

Our paper may in fact be seen as the third in a series, completing the comparison of altogether four semantic models, viz. one operational, one metric denotational and two order-theoretic denotational semantics. The precise picture is the following:

1. In [BMOZ 1,2] we have developed an operational ($\mathcal{O}$) and a metric denotational ($\mathcal{M}$) model (the same one as the one described below), and proved their equivalence. The operational semantics uses the transition systems of Hennessy and Plotkin ([HP], [P]); as we saw already, the metric model goes back to [BBKM].
2. In [M1, M2] the Smyth order-theoretic semantics $\mathcal{S}$ for $\mathcal{L}$ was first proposed. A second order-theoretic semantics, $\mathcal{F}$, building upon ideas in [OH], was designed by Olderog, see [BMO 1,2] for details. This model uses sets of *finite* so-called observations rather than sets of possibly infinite streams; as order between the sets simple (reverse) set inclusion is used. In [BMO 1,2] it was proved that the two order-theoretic structures - subject to certain conditions specification of which we omit here - are isomorphic. As an easy consequence, we obtain that $\mathcal{S} = \mathcal{F}$. (Roughly; the precise statement involves the isomorphism between the two structures.)
3. Altogether, we have four semantics for $\mathcal{L}$, viz. $\mathcal{O}$, $\mathcal{M}$, $\mathcal{S}$ and $\mathcal{F}$, and we know that $\mathcal{O} = \mathcal{M}$ and $\mathcal{S} = \mathcal{F}$. There remains the natural question whether $\mathcal{M} = \mathcal{S}$, and our paper answers this question affirmatively, thus completing (this branch of) the comparative semantics for elemental concurrency.
4. As a side remark pertaining to the relationship with branching time models, we recall that in [BBKM] we also designed a branching time model for $\mathcal{L}$ (in terms of the *processes* as in [BZ]). Calling this semantics $\mathcal{B}$, we showed that, by applying the *trace* operation to $\mathcal{B}$ - collecting all *paths* in the tree-like object resulting from application of $\mathcal{B}$ to a statement -, we obtain $\mathcal{M}$. Thus, we proved that $\mathcal{M} = \mathit{trace} \circ \mathcal{B}$.

Section 2 contains a few mathematical preliminaries, covering elementary definitions for metric spaces and complete partially ordered sets (cpo's). This section is almost as in [BKMOZ]. Section 3 develops various basic semantic definitions: We define the set of streams as a cpo and as a metric space and similarly for the power set of the set of streams. Moreover, we define, for sets of streams (satisfying certain restrictions) the semantic operators of sequential composition, union and merge. The section culminates in the definitions of $\mathcal{S}$ and $\mathcal{M}$. In section 4 we prove a number of technical results concerning the order-theoretic and metric structures, and their mutual relationship. Maybe the most important fact is the following: Let $(X_i)_i$ be a Smyth-ordered chain of sets of streams (satisfying certain conditions). Then $(X_i)_i$ is also a Cauchy sequence in an appropriate metric space, and the order-theoretic and topological limits coincide. For the proof of this the *compactness* of the spaces concerned - a direct consequence of the finiteness of the alphabet of elementary actions - is necessary. In section 5 we establish the main result of the paper, viz. that $\mathcal{M} = \mathcal{S}$. The proof uses the properties relating metric and order obtained in section 4. In addition, a proof technique closely resembling a method used in [BMOZ 2] (in theorem 2.4.1 of that paper) is applied.

ACKNOWLEDGEMENT. Jeff Zucker has scrutinized a first version of this paper and suggested various corrections and improvements. We are grateful to the members of the CWI concurrency project, Frank de Boer, Joost Kok, Jan Rutten, Anton Eliens, and to Erik de Vink for their comments on a first presentation of the material in this paper.

2. Mathematical Preliminaries

In this section we collect some basic definitions and properties concerning (i) metric spaces and (ii) complete partially ordered sets. Both structures will play a role in the denotational models to be presented in section 3 and analyzed in sections 4 and 5.

2.1. Elementary definitions.

Let X be any set. $\mathcal{P}(X)$ denotes the powerset of X, i.e., the set of all subsets of X. $\mathcal{P}_{...}(X)$ denotes the set of all subsets of X which have property .... A sequence $x_0, x_1, \ldots$ of elements of X is usually denoted by $(x_i)_{i=0}^{\infty}$ or, briefly $(x_i)_i$. Often, we shall have occasion to use the *limit, supremum* (sup), *least upper bound* (lub), etc, of a sequence $(x_i)_i$. We then use the notations $\lim_{i \to \infty} x_i$, or, briefly, $lim_i\, x_i$, $sup_i\, x_i$, $lub_i\, x_i$, etc. The notation $f : X \to Y$ expresses that f is a function with domain X and range Y. If $X = Y$ and, for $x \in X$, $f(x) = x$, we call x a *fixed point* of f. We use $\mathbb{N}$ to denote the set of nonnegative integers.

2.2. Metric spaces.

Definition 2.1. A metric space is a pair (M, d) with M a set and d (for *distance*) a mapping $d : M \times M \to [0,1]$ which satisfies the following properties:

a. $d(x, y) = 0$ iff $x = y$
b. $d(x, y) = d(y, x)$
c. $d(x, y) \leq d(x, z) + d(z, y)$

If clause a. is replaced by the weaker a' : $d(x, y) = 0$ if $x = y$, we call (M, d) a pseudo-metric space.

Definition 2.2. Let (M, d) be a metric space.

a. Let $(x_i)_i$ be a sequence in M. We say that $(x_i)_i$ *converges to* an element x in M called its limit, whenever we have:

$$\forall \epsilon > 0\ \exists N \in \mathbb{N}\ \ \forall n > N\ [\ d(x, x_n) < \epsilon\ ]$$

A sequence $(x_i)_i$ in M is a convergent sequence if it converges to x for some $x \in X$

b. A sequence $(x_i)_i$ is called a *Cauchy sequence* whenever we have

$$\forall \epsilon > 0\ \exists N \in \mathbb{N}\ \ \forall n, m > N\ [\ d(x_n, x_m) < \epsilon\ ]$$

c. The space (M, d) is called *complete* whenever each Cauchy sequence converges to an element in M.

d. A subset X of a complete space (M, d) is called *closed* whenever each Cauchy sequence in X converges to an element of X.

Definition 2.3.

a. Let (M_1, d_1) and (M_2, d_2) be two metric spaces. We call the spaces isometric if there exists a bijection $f : M_1 \to M_2$ such that, for all $x, y \in M_1$, $d_2(f(x), f(y)) = d_1(x, y)$.
b. Let (M_1, d_1) and (M_2, d_2) be two metric spaces. We call the function $f : M_1 \to M_2$ *continuous*, whenever, for each sequence $(x_i)_i$ with limit x in M_1, we have that $lim_i\, f(x_i) = f(x)$.
c. Let (M,d) be a metric space and $f : M \to M$. We call f contracting if there exists a real constant c, $0 \leq c < 1$, such that, for all $x, y \in M$, $d(f(x), f(y)) \leq c \cdot d(x, y)$.

Proposition 2.4.

a. *Each contracting function is continuous.*
b. *(Banach's fixed point theorem). Let (M, d) be complete and $f : M \to M$ contracting. Then f has a unique fixed point, which can be obtained as the limit of the (Cauchy) sequence*

x_0, $f(x_0)$, $f(f(x_0))$, ... *for arbitrary* x_0.

For each metric space (M, d) it is possible to define a complete metric space $(\tilde{M}, \tilde{d})$ such that (M, d) is isometric to a (dense) subspace of $(\tilde{M}, \tilde{d})$. In fact, we may take for $(\tilde{M}, \tilde{d})$ the pseudo-metric space of all Cauchy sequences $(x_i)_i$ im M with distance $\tilde{d}((x_i)_i, (y_i)_i) = lim_i d(x_i, y_i)$ which is turned into a metric space by taking equivalence classes with respect to the equivalence relation $(x_i)_i \equiv (y_i)_i$ iff $\tilde{d}((x_i)_i, (y_i)_i) = 0$. M is embedded into $\tilde{M}$ by identifying each $x \in M$ with the constant Cauchy sequence $(x_i)_i$ with $x_i = x$, $i = 0, 1, ...$ in $\tilde{M}$. For each metric space (M, d) we can define a metric $\hat{d}$ on the collection of its nonempty closed subsets, denoted by $\mathcal{P}_{nc}(M)$, as follows:

DEFINITION 2.5 (Hausdorff distance).
Let (M, d) be a metric space, and let X, Y be nonempty subsets of M. We put

a. $d'(x, Y) = inf_{y \in Y} d(x, y)$.

b. $\hat{d}(X, Y) = max\,(sup_{x \in X} d'(x, Y)\,,\ sup_{y \in Y} d'(y, X))$.

We have the following theorem which is quite useful in our metric denotational models:

PROPOSITION 2.6. *Let (M, d) be a metric space and $\hat{d}$ as in definition 2.5.*

a. *$(\mathcal{P}_{nc}(M), \hat{d})$ is a metric space.*

b. *If (M, d) is complete then $(\mathcal{P}_{nc}(M), \hat{d})$ is complete. Moreover, for $(X_i)_i$ a Cauchy sequence in $(\mathcal{P}_{nc}(M), \hat{d})$ we have*

$$lim_i\, X_i \;=\; \{\, lim_i\, x_i : x_i \in X_i,\ (x_i)_i \text{ a Cauchy sequence in } M \,\}$$

Proofs of proposition 2.6 can be found e.g. in [Du] or [En]. The proposition is due to Hahn [Ha]; the proof is also repeated in [BZ]. We close this subsection with a few definitions and properties relating to *compact* spaces and sets. First some terminology. A subset X of a space (M, d) is *open* if its complement $M \setminus X$ is closed. An (open) *cover* of a set X is a family of (open) sets Y_i, $i \in I$, such that $X \subseteq \bigcup_{i \in I} Y_i$.

PROPOSITION 2.7. Let (M, d) be a metric space.

a. (M, d) is called compact whenever each open cover of M has a finite subcover.

b. A subset X of M is called compact whenever each open cover of X has a finite subcover.

PROPOSITION 2.8.

a. *Each closed subset of a compact space is compact.*

b. *If X is compact and f is continuous then $f(X)$ is compact.*

c. *X is compact iff there is a Cauchy sequence $(X_i)_i$ (with respect to the metric of definition 2.5) of finite sets such that $X = lim_i\, X_i$.*

d. *(M, d) is compact whenever each infinite sequence $(x_i)_i$ has a convergent subsequence.*

e. *A subset X of a metric space (M, d) is compact whenever each infinite sequence $(x_i)_i$, $x_i \in X$, has a subsequence converging to an element of X.*

In the final definition and proposition of this subsection we suppress explicit mentioning of the metrics involved. For f a function: $M_1 \to M_2$ we define $\hat{f} : \mathcal{P}_{nc}(M_1) \to \mathcal{P}_{nc}(M_2)$ by $\hat{f}(X) = \{f(x) : x \in X\}$. We have the following result from Rounds ([Ro]):

PROPOSITION 2.9. *Let f be a function from a compact metric space M_1 to a compact metric space M_2. The following three statements are equivalent:*

a. *f is continuous.*
b. *$f : \mathcal{P}_{nc}(M_1) \to \mathcal{P}_{nc}(M_2)$ is continuous with respect to the Hausdorff metric(s).*
c. *For $X \in \mathcal{P}_{nc}(M_1)$, $f(X) \in \mathcal{P}_{nc}(M_2)$ and, for $(X_i)_i$ a decreasing ($X_i \supseteq X_{i+1}$, $i = 0,1,2,\ldots$) chain of elements in $\mathcal{P}_{nc}(M_1)$ we have*

$$\hat{f}(\cap_i X_i) = \cap_i \hat{f}(X_i).$$

2.3. Complete partially ordered sets.

DEFINITION 2.10.
a. A partial order (po) is a pair $(C, \sqsubseteq)$ where C is a set and $\sqsubseteq$ a relation on C (subset of $C \times C$) satisfying
   1 $x \sqsubseteq x$
   2 if $x \sqsubseteq y$ and $y \sqsubseteq x$ then $x = y$
   3 if $x \sqsubseteq y$ and $y \sqsubseteq z$ then $x \sqsubseteq z$

If $\sqsubseteq$ satisfies only 1 and 3 it is called a *preorder*.
b. An (ascending) chain in $(C, \sqsubseteq)$ is a sequence $(x_i)_i$ such that $x_i \sqsubseteq x_{i+1}$, $i = 0,1,\ldots$. The chain is called infinitely often increasing if $x_i \neq x_{i+1}$ for infinitely many i.
c. For $X \subseteq C$ we call $y \in C$ the *least upperbound* (lub) of X if
   1 $\forall x \in X[x \sqsubseteq y]$
   2 $\forall z \in C[\forall x \in X[x \sqsubseteq z] \Rightarrow y \sqsubseteq z]$

DEFINITION 2.11. A *complete partially ordered* set (cpo) is a triple $(C, \sqsubseteq, \bot)$ with $(C, \sqsubseteq)$ a po and $\bot \in C$ such that
a. $\forall x \in C[\bot \sqsubseteq x]$
b. Each chain $(x_i)_i$ in C has a lub in C.

For "the cpo $(C, \sqsubseteq, \bot)$" we often simply write "the cpo C".

DEFINITION 2.12 (continuity). Let C_1 and C_2 be cpo's.
a. A function $f : C_1 \to C_2$ is called monotonic whenever for all $x_1, x_2 \in C_2$, if $x_1 \sqsubseteq x_2$ then $f(x_1) \sqsubseteq f(x_2)$.
b. A function $f : C_1 \to C_2$ is called *continuous* whenever it is monotonic and, for each chain $(x_i)_i$ in C_1 we have $f(\text{lub}_i\, x_i) = \text{lub}_i\, f(x_i)$.

PROPOSITION 2.13 *Let f be a continuous mapping from a cpo C into itself. f has a least fixed point μf satisfying*
1 $f(\mu f) = \mu f$
2 *if $f(y) \sqsubseteq y$ then $\mu f \sqsubseteq y$.*
3 $\mu f = \text{lub}_i\, f^i(\bot)$.

DEFINITION 2.14
a. A subset X is called *flat* whenever, for all $x,y \in X$, $x \sqsubseteq y$ implies $x = y$.
b. A subset X of a cpo C is called *closed* whenever, for each infinitely often increasing chain $(x_i)_i$ of elements in C such that, for all $i = 0,1,\ldots$ we have that $x_i \sqsubseteq y_i$ for some $y_i \in X$, it follows that

$\text{lub}_i\, x_i \in X$.

This definition of closed appears in [Ba] and [Ku]. We now introduce a number of preorders on $\mathcal{P}(C)$, for $(C, \sqsubseteq, \perp)$ a cpo.

DEFINITION 2.15.

a. The Smyth preorder $\sqsubseteq_S$: $X \sqsubseteq_S Y$ iff $\forall y \in Y \exists x \in X [x \sqsubseteq y]$

b. The Hoare preorder $\sqsubseteq_H$: $X \sqsubseteq_H Y$ iff $\forall x \in X \exists y \in Y [x \sqsubseteq y]$

c. The Egli-Milner preorder $\sqsubseteq_{EM}$: $X \sqsubseteq_{EM} Y$ iff $X \sqsubseteq_S Y$ and $X \sqsubseteq_H Y$.

None of the three preorders is, in general, a partial order. In fact, we may take the two sets $X = \{x,y,z\}$ and $Y = \{x,z\}$ with $x \sqsubseteq y$ and $y \sqsubseteq z$ as a counterexample. In subsequent sections, only $\sqsubseteq_S$ will be used. The other preorders are included for completeness' sake.

3. STREAM SEMANTICS FOR ELEMENTAL CONCURRENCY

We introduce a simple language $\mathcal{L}$ with concurrency and design two denotational semantics for it. The first semantic function is called $\mathcal{S}$ (for Smyth - like order - theoretic) and the second $\mathcal{M}$ (for metric). In subsequent sections we shall develop the tools for proving the equivalence $\mathcal{S} = \mathcal{M}$.

We recall from the introduction that we already showed in previous papers:

(i) For $\mathcal{F}$ the denotational semantics based on the cpo of (sets of) finite observations, $\mathcal{F} = \mathcal{S}$ (modulo the isomorphism linking the two cpo's).

(ii) For $\mathcal{O}$ the operational semantics based on transition systems, $\mathcal{O} = \mathcal{M}$.
(In addition, we know that

(iii) For $\mathcal{B}$ the (metric) branching time semantics, $trace \circ \mathcal{B} = \mathcal{M}$.)

We start the section with a description of the syntax of $\mathcal{L}$. Elements of $\mathcal{L}$ will be called statements or, occasionally, processes, and we use s, t to range over $\mathcal{L}$. The language $\mathcal{L}$ is what we like to call a *uniform* language: its elementary actions are left uninterpreted. No constructs such as (individual) variables, assignments or tests are present in the syntax, and neither do we employ notions such as states in the semantics. In fact, statements in $\mathcal{L}$ may well be seen as (pieces of) grammar which prescribe the generation of finite or infinite sequences of symbols (or actions), and our semantic studies may shed light on questions in formal language theory as well.

For the syntax of $\mathcal{L}$ we need two classes of terminal elements:

1. The class A, with typical elements $a,b,\ldots$, of elementary actions. For A we take an arbitrary (but finite!) alphabet.
2. The class $\mathcal{P}var$, with typical elements $x,y,\ldots$, of process variables. For $\mathcal{P}var$ we take some infinite set of symbols: it is convenient to have an infinite supply of fresh process variables. Process variables play a role in the syntactic construct for recursion as we shall see in a moment.

We now give, in a self-explanatory notation,

DEFINITION 3.1 (syntax for $\mathcal{L}$).

$$s ::= a \mid s_1 \,;\, s_2 \mid s_1 \cup s_2 \mid s_1 \parallel s_2 \mid x \mid \mu x[s]$$

A statement s is of one of the following six forms:

- an *elementary action* a
- the *sequential composition* $s_1 \,;\, s_2$ of statements s_1 and s_2
- the *nondeterministic choice* $s_1 \cup s_2$: it is executed by executing s_1 or s_2 chosen nondeterministically
- the *concurrent execution* $s_1 \parallel s_2$, modelled by arbitrarily interleaving the elementary actions of s_1 and s_2

- a *process variable* x which is (normally) used in
- the *recursive construct* $\mu x[s]$: its execution amounts to execution of s where occurrences of x in s are executed by (recursively) executing $\mu x[s]$. For example, with the definitions to be proposed presently, the intended meaning of $\mu x[(a\ ;\ x)\cup b]$ is the set $a^*.\ b\ \cup\ \{a^\omega\}$. (Here a^ω denotes the infinite sequence of a's.)

The prefix $\mu x\ .\ \cdots$ binds occurrences of x in $\cdots$ in the usual way, inducing the familiar notions of free and bound (occurrences of) process variables. We shall call a statement *closed* if it has no free occurrences of process variables.

We continue with the development of the two semantic models. For both of them we need various basic definitions which we may use to build the structures in which our semantics are defined. Apart from an occasional point of presentation, no new material is presented here: the definitions stem originally from [M 1,2] and [BBKM], and are included also in papers such as [BMO 1,2], [BMOZ 1,2], [BKMOZ].

We begin with the definition of the set of *streams* over A, denoted by A^{st} (cf. e.g. [Br1,Br2]). Let $\perp$ be a symbol not in A.

DEFINITION 3.2 (streams). $A^{st} = A^* \cup A^*.\ \{\perp\} \cup A^\omega$.

Here $A^*(A^\omega)$ denotes the set of all finite (infinite) words over A. We use ϵ to denote the empty sequence. $A^*\ .\ \{\perp\}$ is the collection of all finite words over A, followed by the $\perp$-symbol. We use u, v, w to range over A^{st}. We recall (from section 2.1) the notation $\mathcal{P}_{...}(A^{st})$ for the collection of all subsets of A^{st} with property $\cdots$. Usually, we abbreviate $\mathcal{P}_{...}(A^{st})$ to $\mathfrak{S}_{...}$. We shall use X, Y, Z to range over $\mathfrak{S}$.

The first group of basic definitions is assembled in

DEFINITION 3.3.

a. The function *strip*: $A^{st} \rightarrow A^* \cup A^\omega$. We put *strip* $(u) = u$ for $u \in A^* \cup A^\omega$, and *strip* $(u) = u'$ for $u = u'\perp$, with $u' \in A^*$.
b. The prefix order $\leqslant$. We put $u \leqslant v$ whenever one of the following three conditions is satisfied
   (i) $u = v$
   (ii) $u, v \in A^* \cup A^\omega$ and $\exists w[u\ .\ w = v]$
   (iii) $v \in A^*\ .\ \{\perp\}$ and $u \leqslant$ *strip* (v)
c. The function *length*: $A^{st} \rightarrow \mathbb{N} \cup \{\infty\}$. We put *length* (u) as usual for $u \in A^*$, *length* $(u) = \infty$ for $u \in A^\omega$, and *length* $(u) =$ *length* $(u')+1$ for $u = u'\perp$, $u' \in A^*$.
d. A $\leqslant$-chain $(u_i)_i$ is a sequence $u_0, u_1, \ldots,$ such that $u_i \leqslant u_{i+1}$, $i = 0,1,....$The least upper bound of the $\leqslant$-chain $(u_i)_i$ is denoted by $\sup_i u_i$.
e. The $\leqslant$-truncation $u(n)$: if *length* $(u) \geqslant n$, $u(n)$ denotes the prefix of u of length n. If *length* $(u) < n$, $u(n) = u$.
f. The stream order $\sqsubseteq$: We put $u \sqsubseteq v$ whenever one of the following two conditions is satisfied:
   (i) $u = v$
   (ii) $u \in A^*\ .\ \{\perp\}$ and *strip* $(u) \leqslant v$.
g. A $\sqsubseteq$-chain $(u_i)_i$ is a sequence $u_0, u_1, \ldots,$ such that $u_i \sqsubseteq u_{i+1}$, $i = 0,1,....$ The least upper bound of the $\sqsubseteq$-chain $(u_i)_i$ is denoted by $\mathrm{lub}_i\, u_i$.
h. The $\sqsubseteq$-truncation $u[n]$. If *length* $(u) \geqslant n$, we put $u[n] = u(n)$, if $u(n) \in A^*\ .\ \{\perp\}$, and $u[n] = u(n)\ .\ \perp$, otherwise. If *length* $(u) < n$, we put $u[n] = u$.

REMARKS

1. Properly speaking, the concatenation of two streams as used in b (ii) has not yet been defined. It is in fact implicit in definition 3.10 below.

2. A chain $(u_i)_i$ (either $\leqslant$ - or $\sqsubseteq$ -) such that $u_i \neq u_{i+1}$, for infinitely many i, is called infinitely often increasing (i.o.i.). A chain which is not i.o.i. is called *stabilizing*. In that case, there is an index i_0 such that $u_i = u_{i_0}$, all $i \geqslant i_0$, and we say that $(u_i)_i$ stabilizes in u_{i_0}.

The following first results are easily shown:

LEMMA 3.4.

a. *$(A^{st}, \leqslant, \epsilon)$ is a cpo. For a $\leqslant$-chain $(u_i)_i$, we have $u = \sup_i u_i$ iff either $(u_i)_i$ is i.o.i. and $u \in A^{\omega}$ is such that $u_i \leqslant u$, $i \geqslant 0$, or $(u_i)_i$ stabilizes in u.*

b. $\forall u, v, w\ [((u \leqslant w) \wedge (v \leqslant w)) \Rightarrow ((u \leqslant v) \vee (v \leqslant u))]$

c. *$(A^{st}, \sqsubseteq, \perp)$ is a cpo. For a $\sqsubseteq$-chain $(u_i)_i$, we have $u = \mathrm{lub}_i\, u_i$ iff either $(u_i)_i$ is i.o.i., (then) $u_i = u'_i \,.\, \perp$ for all i, $(u'_i)_i$ is a $\leqslant$-chain and $u = \sup_i u'_i$, or $(u_i)_i$ stabilizes in u.*

d. $u = \sup_n u(n) = \mathrm{lub}_n\, u[n]$.

We proceed with the definition of the *distance* d between streams:

DEFINITION 3.5. The mapping $d : A^{st} \times A^{st} \rightarrow [0, 1]$ is defined by

$$d(u, v) = 2^{-\sup\{n\, :\, u(n) = v(n)\}}$$

with the convention that $2^{-\infty} = 0$.

The following theorem is fundamental for the metric framework:

THEOREM 3.6 ([Ni]). *(A^{st}, d) is a complete and compact metric space.*

We next turn to the development of an order - theoretic and metric structure for *sets* of streams

DEFINITION 3.7. Let $X, Y \in \mathfrak{S}$.

a. $X(n) = \{u(n) : u \in X\}$, $X[n] = \{u[n] : u \in X\}$.

b. $X \sqsubseteq_S Y$ is the Smyth preorder (definition 2.15) induced by the stream order $\sqsubseteq$ on A^{st}.

c. $\min(X) = \{u : u \in X$ and for all $v \in X[v \sqsubseteq u \Rightarrow v = u]\}$.

d. Let $\mathfrak{S}_{nc}$ denote the collection of all nonempty closed sets of streams. $\hat{d}(X, Y)$ denotes the Hausdorff distance (definition 2.5) on $\mathfrak{S}_{nc}$.

e. We use $\mathfrak{S}_f$ and $\mathfrak{S}_{ncf}$ to denote the collection of all flat (definition 2.14a) and of nonempty closed (definition 2.14b) and flat sets of streams, respectively.

f. For a $\sqsubseteq$-chain $(X_i)_i$ we denote its least upper bound by $\bigsqcup_i X_i$.

The following theorem states, essentially, that $\mathfrak{S}_f$ and $\mathfrak{S}_{nc}$ are the structures we want. (Note, however, that we shall later specialize $\mathfrak{S}_f$ to $\mathfrak{S}_{ncf}$ to ensure continuity of the semantic operators.)

THEOREM 3.8.

a. *X is $\sqsubseteq$-closed in $(A^{st}, \sqsubseteq, \perp)$ iff X is d-closed in (A^{st}, d).*

b. *For any X, X_1, X_2 in $\mathfrak{S}$ we have*
   (i) *$X \sqsubseteq_S min(X)$ and $min(X) \sqsubseteq_S X$*
   (ii) *$X_1 \sqsubseteq_S X_2 \Rightarrow min(X_1) \sqsubseteq_S min(X_2)$*
   (iii) *$(min(X))[n] = min(X[n])$*

c. *$(\mathfrak{S}_f, \sqsubseteq_S, \{\perp\})$ is a cpo. For $(X_n)_n$ a $\sqsubseteq_S$-chain we have*

$\bigsqcup_n X_n = \{u : u = \text{lub}_n u_n, (u_n) \text{ a } \sqsubseteq\text{-chain with } u_n \in X_n\}$

d. *$(\mathfrak{S}_{nc}, \hat{d})$ is a complete (and compact) metric space.*
e. For $X, Y \in \mathfrak{S}_{nc}$, $\hat{d}(X, Y) = 2^{-\sup\{n : X(n) = Y(n)\}}$ with the convention that $2^{-\infty} = 0$.
f. *For $X \in \mathfrak{S}_{nc}$, $(X(n))_n$ is a Cauchy sequence in $(\mathfrak{S}_{nc}, \hat{d})$, and $X = \lim_n X(n)$.*
g. *For $X \in \mathfrak{S}_f$, $(X[n])_n$ is a $\sqsubseteq_S$-chain in $(\mathfrak{S}_f, \sqsubseteq_S, \{\perp\})$*

PROOF. These result are, essentially, from [M1, M2] and [BBKM]; see also [BKMOZ], [MV] for related references and results. □

Having defined our fundamental structures, we next arrive at the definition of the various semantic operators which we will have as counterparts of the syntactic operators ; , ∪, ||. Once these have been defined satisfactorily, we shall have completed the preparations for the semantic definitions. Recursion will be dealt with by the familiar (least) fixed point technique, for which the relevant apparatus will then be available.

We define the semantic operators directly for $X, Y \in \mathfrak{S}$, rather than going through a two stage process in which the operators are first defined on A^{st}. This is for convenience rather than out of necessity.

We first deal with the case that X, Y consist of finite words only. Let $\mathfrak{S}_{fin}$ be short for $\mathcal{P}(A^* \cup A^* . \{\perp\})$.

DEFINITION 3.9. We define $\mathbf{op}^{fin} : \mathfrak{S}_{fin} \times \mathfrak{S}_{fin} \to \mathfrak{S}_{fin}$, where $\mathbf{op}^{fin} \in \{\cdot, \cup, \|\}$. We let X, Y range over $\mathfrak{S}_{fin}$.

a. We assume as known the operator of *prefixing* which for $a \in A$, $u \in A^* \cup A^* . \{\perp\}$, delivers $a . u$.
b. $a . X = \{a . u : u \in X\}$
c. $X . Y = \bigcup \{u . Y : u \in X\}$, where $u . Y$ is defined (inductively) by $\epsilon . Y = Y$, $\perp . Y = \{\perp\}$, $(au) . Y = a . (u . Y)$.
d. $X \cup Y$ is the set-theoretic union of X and Y
e. $X \| Y = (X \mathbin{\|\!\!\_} Y) \cup (Y \mathbin{\|\!\!\_} X)$; moreover, $X \mathbin{\|\!\!\_} Y = \bigcup \{u \mathbin{\|\!\!\_} Y : u \in X\}$, where $u \mathbin{\|\!\!\_} Y$ is defined (inductively) by $\epsilon \mathbin{\|\!\!\_} Y = Y$, $\perp \mathbin{\|\!\!\_} Y = \{\perp\}$, $(au) \mathbin{\|\!\!\_} Y = a . (\{u\} \| Y)$.

REMARK. $\mathbin{\|\!\!\_}$ stems from ACP, cf. [BK].

Next, we define the metric and (Smyth-) order-theoretic operators $\mathbf{op}^{\mathfrak{M}}$ and $\mathbf{op}^{\mathfrak{S}}$, where $\mathbf{op}^{\mathfrak{M}}, \mathbf{op}^{\mathfrak{S}} \in \{\cdot, \cup, \|\}$, for the general case, i.e., for X, Y which do not necessarily consist of finite words only. Note that $\mathbf{op}^{\mathfrak{S}}$ is defined on $\mathfrak{S}_{ncf}$ rather than on all of $\mathfrak{S}_f$. This is necessary to ensure continuity of $\mathbf{op}^{\mathfrak{S}} \in \{\cdot, \|\}$ (see below).

DEFINITION 3.10.

a. $\mathbf{op}^{\mathfrak{M}} : \mathfrak{S}_{nc} \times \mathfrak{S}_{nc} \to \mathfrak{S}_{nc}$ is defined by

$$X \,\mathbf{op}^{\mathfrak{M}}\, Y = \lim_n \left[X(n) \,\mathbf{op}^{fin}\, Y(n) \right]$$

b. $\mathbf{op}^{\mathfrak{S}} : \mathfrak{S}_{ncf} \times \mathfrak{S}_{ncf} \to \mathfrak{S}_{ncf}$ is defined by

$$X \,\mathbf{op}^{\mathfrak{S}}\, Y = \mathit{min}\,(X \,\mathbf{op}^{fin}\, Y)\,, \text{ for } X, Y \in \mathfrak{S}_{fin} \cap \mathfrak{S}_{ncf}$$

$$X \,\mathbf{op}^{\mathfrak{S}}\, Y = \bigsqcup_n (X[n] \,\mathbf{op}^{\mathfrak{S}}\, Y[n])\,, \text{ for } X, Y \in \mathfrak{S}_{ncf}$$

The following theorem expresses well-definedness, (monotonicity and) $\sqsubseteq_S$- and d-continuity of the respective operators.

THEOREM 3.11.

a. *The operators* $\mathbf{op}^{\mathfrak{M}}$ *and* $\mathbf{op}^{\mathbb{S}}$ *are well-defined. In particular, they take (pairs of) nonempty closed (and flat) sets to nonempty closed (and flat) sets.*

b. *The operators* $\mathbf{op}^{\mathbb{S}}$ *are* $\sqsubseteq_S$*-monotonic.*

c. *The operators* $\mathbf{op}^{\mathbb{S}}$ *are* $\sqsubseteq_S$*-continuous mappings:*

$$\mathfrak{S}_{ncf} \times \mathfrak{S}_{ncf} \rightarrow \mathfrak{S}_{ncf}$$

d. *The operators* $\mathbf{op}^{\mathfrak{M}}$ *are d-continuous mappings*

$$\mathfrak{S}_{nc} \times \mathfrak{S}_{nc} \rightarrow \mathfrak{S}_{nc}$$

PROOF. The results for $\mathbf{op}^{\mathbb{S}}$ are from [M 1,2]. For $\mathbf{op}^{\mathfrak{M}}$ the result follows from [BBKM] and proposition 2.9 (equivalence of b and c).

REMARK. The sets $(X_n)_n$, $(Y_n)_n$ defined by $X_n = \{u \in a^*: \text{length}(u) \geq n\}$, $n = 0,1,\ldots$, and $Y_n = \{a^\omega\}$, $n = 0,1,\ldots$, show that the operators $\mathbf{op}^{\mathbb{S}} \in \{\cdot, \|\}$ are, in general, discontinuous in the case that they are not restricted to $\mathfrak{S}_{ncf} \times \mathfrak{S}_{ncf}$.

We are almost ready to present the definitions of the semantic functions $\mathbb{S}$ and $\mathfrak{M}$. As final preparation, we need one further syntactic notion, viz. that of *guarded* statements. The reason for this is that the semantics based on the metric approach is valid only for statements satisfying the guardedness requirement. (Specifically, the metric treatment of the recursive construct requires this condition to be satisfied.) Intuitively, a statement s is guarded when all its recursive substatements $\mu x[t]$ satisfy the condition that (recursive) occurrences of x in t are 'semantically preceded' by some statement. More precisely, we have

DEFINITION 3.12 (guarded statements).

a. We first define the notion of an occurrence of a variable being *exposed* in s. The definition is by structural induction on s
   1. x is exposed in x
   2. If an occurrence of x is exposed in s_1, then it is exposed in $s_1; s_2$, $s_1 \| s_2$, $s_2 \| s_1$, $s_1 \cup s_2$, $s_2 \cup s_1$ and $\mu y[s_1]$ for $y \not\equiv x$.

b. A statement s is defined to be guarded if for all its recursive substatements $\mu x[t]$, t contains no exposed occurrences of x.

EXAMPLES.

1. In the statement $x;a \cup b;x$ the first occurrence of x is exposed and the second is not.
2. $\mu x[a\ ;\ (x \| b)]$ is guarded, but $\mu x[x]$, $\mu y[y \| b]$ and $\mu y[\mu x[y]]$, as well as any statement containing these, are not.

We have now arrived at the definition of the two semantics for $\mathfrak{L}$. Let $\Gamma_{\ldots} = \mathfrak{P}_{var} \rightarrow \mathfrak{S}_{\ldots}$, and let $\gamma \in \Gamma_{\ldots}$. (Here ... ranges over $\{nc, ncf\}$.) We use the notation $(\gamma' =)\gamma\langle X/x\rangle$ for a *variant* of γ, which is like γ but for its value in x which equals X (*i.e.*, $\gamma'(y) = \gamma(y)$ for $y \not\equiv x$ and $\gamma'(x) = X$). We use $\mathbf{op}$ without superscript to range over the syntactic operators $\{;, \cup, \|\}$ and $\mathbf{op}^{\ldots}$ with superscript ... to range over the corresponding semantic operators.

DEFINITION 3.13 (two denotational semantics).

a. The mapping $\mathbb{S} : \mathfrak{L} \rightarrow (\Gamma_{ncf} \rightarrow \mathfrak{S}_{ncf})$ is defined by
   (i) $\mathbb{S}[\![a]\!](\gamma) = \{a\}$
   (ii) $\mathbb{S}[\![s_1\, \mathbf{op}\, s_2]\!](\gamma) = \mathbb{S}[\![s_1]\!](\gamma)\, \mathbf{op}^{\mathbb{S}}\, \mathbb{S}[\![s_2]\!](\gamma)$
   (iii) $\mathbb{S}[\![x]\!](\gamma) = \gamma(x)$

(iv) $\mathcal{S}[\![\, \mu x[s]]\!](\gamma) = \bigsqcup_n X_n$, where $X_0 = \{\perp\}$ and $X_{n+1} = \mathcal{S}[\![\, s \,]\!](\gamma\langle X_n/x\rangle)$

b. The mapping $\mathcal{M} : \mathcal{L} \rightarrow (\Gamma_{nc} \rightarrow \mathfrak{S}_{nc})$ is defined by

(i) $\mathcal{M}[\![\, a \,]\!](\gamma) = \{a\}$

(ii) $\mathcal{M}[\![\, s_1 \, \mathbf{op} \, s_2]\!](\gamma) = \mathcal{M}[\![\, s_1]\!](\gamma) \, \mathbf{op}^{\mathcal{M}} \, \mathcal{M}[\![\, s_2]\!](\gamma)$

(iii) $\mathcal{M}[\![\, x \,]\!](\gamma) = \gamma(x)$

(iv) $\mathcal{M}[\![\, \mu x[s]]\!](\gamma) = \lim_n X_n$, where $X_0 = \{\perp\}$ and $X_{n+1} = \mathcal{M}[\![\, s \,]\!](\gamma\langle X_n/x\rangle)$

The following facts support this definition

THEOREM 3.14

a. *The function* $\Phi = \lambda X \,.\, \mathcal{S}[\![\, s \,]\!](\gamma\langle X/x\rangle)$ *is a* $\sqsubseteq_S$*-continuous mapping:* $\mathfrak{S}_{ncf} \rightarrow \mathfrak{S}_{ncf}$ *and, for* $(X_n)_n$ *as in clause a (iv),* $\bigsqcup_n X_n = \mu\Phi$.

b. *Assume s guarded. The function* $\Psi = \lambda X \,.\, \mathcal{M}[\![\, s \,]\!](\gamma\langle X/x\rangle)$ *is a contracting mapping:* $\mathfrak{S}_{nc} \rightarrow \mathfrak{S}_{nc}$, *and, for* $(X_n)_n$ *as in clause b (iv),* $\lim_n X_n$ *yields the unique fixed point of* Ψ.

REMARK. For the contractivity property in part b of this theorem, the guardedness of s is necessary. For the semantic function $\mathcal{S}$, the situation is the following:

(i) For (closed and) guarded s we have that $\mathcal{S}[\![\, s \,]\!](\gamma) \subseteq A^* \cup A^\omega$. This is a consequence of an analogous fact for $\mathcal{M}$ (see end of section 5) and the equality $\mathcal{S} = \mathcal{M}$ (theorem 3.15 below).

(ii) For unguarded s, $\mathcal{S}[\![\, s]\!](\gamma)$ will involve streams ending in $\perp$. For example $\mathcal{S}[\![\, (a;\, \mu x[b \| x]) \cup c]\!](\gamma) = \{a\perp, c\}$ and $\mathcal{S}[\![\, (a;\, \mu x[b \| x]) \cup (a;\, c)]\!](\gamma) = \{a\perp\}$. This follows from (the treatment of recursion and) the flattening operator *min* in the definition of $\mathbf{op}^{\mathcal{S}}$ (in the clause $X \, \mathbf{op}^{\mathcal{S}} \, Y = min(X \, \mathbf{op}^{fin} \, Y)$, X, Y with finite words only).

Our aim in the next section will be to prove the

THEOREM 3.15. *For each closed and guarded* $s \in \mathcal{L}$

$$\mathcal{S}[\![\, s \,]\!] = \mathcal{M}[\![\, s \,]\!] \,.$$

In order to establish this result, we have to study the relationship between the two structures $\mathfrak{S}_{ncf}$ as a cpo are $\mathfrak{S}_{nc}$ as a metric space in more detail, as we shall do in section 4.

4. RELATING THE SEMANTIC DOMAINS

The first main result of this section states that, for $(X_i)_i$ a $\sqsubseteq_S$-chain in $\mathfrak{S}_{ncf}$, $(X_i)_i$ is also a Cauchy sequence (in $\mathfrak{S}_{nc}$), and $\lim_i X_i = \bigsqcup_i X_i$. This result is, clearly, fundamental for the proof of

$$\mathcal{M}[\![s]\!] = \mathcal{S}[\![s]\!], \qquad (*)$$

for s a recursive construct. The second part of the section is devoted to a number of properties of the *min*-operator. We first prove that *min* is *d*-continuous. Next, we use this - and various other properties of *min*- to prove that, if $min(X_i) = Y_i$, $X_i \in \mathfrak{S}_{nc}$, $Y_i \in \mathfrak{S}_{ncf}$, $i = 1,2$, then $min(X_1 \, \mathbf{op}^{\mathcal{M}} \, X_2) = Y_1 \, \mathbf{op}^{\mathcal{S}} \, Y_2$. The latter result is crucial for the derivation of (*) for s of the form $s_1 \, \mathbf{op} \, s_2$.

We begin with an auxiliary lemma.

LEMMA 4.1 (interpolation)

a. *Let* $(X_i)_i$ *be a* $\sqsubseteq_S$*-chain in* $\mathfrak{S}_{ncf}$*. For each* $\sqsubseteq$*-chain* $(u_{i_j})_j$, *with* $u_{i_j} \in X_{i_j}, j = 0,1,\ldots$, *there exists a* $\sqsubseteq$*-chain* $(u_i)_i$, *with* $u_i \in X_i, i = 0,1,\ldots$, *which has* $(u_{i_j})_j$ *as a subsequence.*

b. *Let* $(X_i)_i$ *be a* $\sqsubseteq_S$*-chain in* $\mathfrak{S}_{ncf}$*. For each convergent sequence* $(u_{i_j})_j$, *with* $u_{i_j} \in X_{i_j}, j = 0,1,\ldots$, *there exists a convergent sequence* $(u_i)_i, u_i \in X_i$, *containing* $(u_{i_j})_j$ *as a subsequence (and, consequently,*

$\lim_j u_{i_j} = \lim_i u_i$).

Proof

a. It is, clearly, sufficient to prove that, if $X \sqsubseteq_S Y \sqsubseteq_S Z$, $X, Y, Z \in \mathfrak{S}_{ncf}$, and $u \in X, w \in Z$ with $u \sqsubseteq w$, then there exists $v \in Y$ with $u \sqsubseteq v \sqsubseteq w$. By the definition of $\sqsubseteq_S$ we find $v_1 \in Y$ such that $v_1 \sqsubseteq w$ and $u_1 \in X$ such that $u_1 \sqsubseteq v_1 \sqsubseteq w$. Since both $u \sqsubseteq w$ and $u_1 \sqsubseteq w$ we have $u_1 \sqsubseteq u$ or $u \sqsubseteq u_1$. Since X is flat we have $u_1 = u$, and we see that v_1 is the desired element in Y.

b. Let $u_{i_j} = v_j w_j$, where (v_j) is a $\leqslant$-chain and $sup_j v_j = \lim_j u_{i_j}$. Consider, for some fixed j, u_{i_j} and $u_{i_{j+1}}$, and suppose $i_{j+1} - i_j > 1$. So, for some i, $i_j < i < i_{j+1}$. We can find an element u_i such that $u_i = v_j w'_j$ for some w'_j. This can be seen as follows: Since $X_i \sqsubseteq_S X_{i+1}$ there must be an element u_i such that $u_i \sqsubseteq u_{i_{j+1}} = v_{j+1} w_{j+1} = v_j w'_{j+1}$, for some w'_{j+1}. If $u_i \in A^* \cup A^\omega$, the result is immediate. Now let $u_i = u\bot$. If u is such that $v_j \leqslant u$, we have finished. If $u < v_j$ we argue as follows: Since $X_{i_j} \sqsubseteq_S X_i$, there must be some u'_{i_j} such that $u'_{i_j} \sqsubseteq u\bot \sqsubseteq v_j w_j$. By flatness of X_{i_j}, $u'_{i_j} = v_j w_j$. So $u_i = u\bot = u'_{i_j} = v_j w_j$ as well. Hence in this case we have also found an element u_i as desired. Consequently, we are always able to interpolate the converging sequence $(u_{i_j})_j$ to one of the form $(u_i)_i$, where $u_i \in X_i, i = 0, 1, \ldots$. □

The next lemma is also auxiliary, and relies essentially on the compactness of (A^{st}, d).

Lemma 4.2. *Let $X_1, X_2, \in \mathfrak{S}_{nc}$. (At least) one of the following two conditions holds:*

1. *There exists $\bar{u}_1 \in X_1$ such that*
   $\hat{d}(X_1, X_2) = sup_{u_1 \in X_1} d'(u_1, X_2) = d'(\bar{u}_1, X_2)$
   (See definition 2.5 for d').
2. *Symmetric.*

Proof Direct from the fact that a (real-valued) continuous function on a compact set attains its maximum. □

We next state two important properties which relate $\bigsqcup_i X_i$ and $\lim_i X_i$.

Lemma 4.3. *Each $\sqsubseteq_S$-chain $(X_i)_i$, with X_i in $\mathfrak{S}_{ncf}$, is a Cauchy sequence (in $\mathfrak{S}_{nc}$).*

Proof (cf. [R]). Let $(X_j)_j$ be a $\sqsubseteq_S$-chain in $\mathfrak{S}_{ncf}$. We define the set $\lim_j X_j$ by

$$\lim_j X_j = \{u \mid u = \lim_j u_j, u_j \in X_j \text{ and } (u_j)_j \text{ a Cauchy sequence}\}$$

(Note that this definition does not require that $(X_j)_j$ is a Cauchy sequence.) We first prove that the set $\lim_j X_j$ is nonempty and closed. By [M1], p. 91-93, the set

$$\bigsqcup_j X_j = \{u \mid u = \text{lub}_j\ u_j,\ u_j \in X_j, (u_j)_j \text{ a } \sqsubseteq\text{-chain}\}$$

is nonempty if all X_j are nonempty. Clearly, $\bigsqcup_j X_j \subseteq \lim_j X_j$; hence, $\lim_j X_j$ is nonempty. In order to prove that $\lim_j X_j$ is closed. assume that $(u_i)_i$ is a Cauchy sequence in $\lim_j X_j$. Then, for each $i, u_i = \lim_j u_{i,j}$ for $(u_{i,j})_j$ a Cauchy sequence with $u_{i,j} \in X_j$, $j = 0, 1, \ldots$. Following an argument as in [J, proposition 4.3, p. 303] we can find a sequence $(n_j)_j$ of indices such that $(u_{j,n_j})_j$ is also a Cauchy sequence, and $\lim_i u_i = \lim_j u_{j,n_j} \in \lim_j X_{n_j} \subseteq \lim_j X_j$ (the inclusion holds by interpolation). We shall now show that

$$\hat{d}(X_i, \lim_j X_j) \to 0 \text{ as } i \to \infty,$$

thus proving that $(X_i)_i$ is a Cauchy sequence. We shall only exhibit the proof that

$$\sup_{u \in X_i} d'(u, \lim_j X_j) \to 0 \text{ as } i \to \infty.$$

By Lemma 4.2 there exist u_i such that $sup_{u \in X_i} d'(u, \lim_j X_j) = d'(u_i, \lim_j X_j)$. By compactness, $(u_i)_i$ has a converging subsequence $(u_{i_j})_j$. Suppose $\lim_j u_{i_j} = \bar{u}$. By Lemma 4.1b there are interpolating $u'_i \in X_i$ such that $(u'_i)_i$ contains $(u_{i_j})_j$ as a subsequence. So $\bar{u} \in \lim_j X_j$. Now let $\epsilon > 0$ and choose i_k such that $d(u_{i_k}, \bar{u}) < \epsilon$. Then, for each $j \geqslant i_k$,

$$\sup_{u \in X_j} d'(u, \lim_j X_j) \leqslant \text{ (since } X_{i_k} \sqsubseteq_S X_j)$$

$$\sup_{u \in X_{i_k}} d'(u, \lim_j X_j) =$$

$$d'(u_{i_k}, \lim_j X_j) \leqslant \text{ (since } \bar{u} \in \lim_j X_j)$$

$$d(u_{i_k}, \bar{u}) \leqslant \epsilon. \quad \square$$

For $(X_j)_j$ a $\sqsubseteq_S$-chain in $\mathfrak{S}_{ncf}$, we now know that $(X_j)_j$ is also a Cauchy sequence. The next theorem answers the natural question 'is it the case that $\bigsqcup_j X_j = \lim_j X_j$?' affirmatively.

THEOREM 4.4. *Let $(X_j)_j$ be a $\sqsubseteq_S$-chain in $\mathfrak{S}_{ncf}$. Then*

$$\bigsqcup_j X_j = \lim_j X_j$$

PROOF. Recall that

$$\lim_j X_j = \{u \mid u = \lim_j u_j,\ u_j \in X_j,\ (u_j)_j \text{ a Cauchy sequence}\}$$

$$\bigsqcup_j X_j = \{u \mid u = \text{lub}_j u_j,\ u_j \in X_j,\ (u_j)_j \text{ a } \sqsubseteq\text{-chain}\}$$

Since every $\sqsubseteq$-chain $(u_j)_j$ in A^{st} is also a Cauchy sequence such that $\lim_j u_j = \text{lub}_j u_j$, we clearly have that $\bigsqcup_j X_j \subseteq \lim_j X_j$. There remains the proof that $\lim_j X_j \subseteq \bigsqcup_j X_j$. Take some $u = \lim_j u_j \in \lim_j X_j$. First we assume that the sequence $(u_j)_j$ stabilizes at some u_{i_0}. By the definition of $\sqsubseteq_S$, there is a $\sqsubseteq$-chain $u'_0 \sqsubseteq u'_1 \sqsubseteq \ldots \sqsubseteq u'_i = u'_{i+1} = \ldots$ with $u'_i = u_{i_0}$, $i \geqslant i_0$. Thus, $u = u'_{i_0} = \text{lub}_i u'_i$. Now take the case that $(u_j)_j$ does not stabilize. Thus, $u \in A^\omega$. We consider, for some fixed j, the set X_j. Since $X_j \sqsubseteq_S X_{j+i}$, $i \geqslant 0$, there must be elements $u_j^{(i)} \in X_j$ such that $u_j^{(i)} \sqsubseteq u_{j+i}$. Let $V_j =^{df} \{u_j^{(i)} \mid i \geqslant 0\}$. If V_j is infinite, it must contain an infinite (i.o.i.) convergent subsequence $(u_j^{(i_k)})_k$. Since $V_j \subseteq X_j$ and X_j is closed, X_j must contain $\lim u_j^{(i_k)}$. Since, for each k, $u_j^{(i_k)} \sqsubseteq u_{j+i_k}$, and since the sequence $(u_j^{(i_k)})_k$ is i.o.i., we have that $\lim_k u_j^{(i_k)} = u$. Thus, for V_j infinite we infer that $u \in X_j$. We now distinguish two cases:

CASE 1.
V_j is infinite for almost all j, say for all $j \geqslant j_0$. We can then construct the chain

$$u'_0 \sqsubseteq u'_1 \sqsubseteq \cdots \sqsubseteq u'_{j_0} \sqsubseteq u'_{j_0+1} \sqsubseteq \cdots$$

with $u'_{j_0+l} = u$, $l \geqslant 0$. Thus, $u = \text{lub}_n u'_n$, and we are done.

CASE 2.
There are infinitely many finite V_j, say V_j is finite for all j in the index set J. Consider such a finite V_j. Since V_j contains a finite number of elements approximating an infinite number of streams $(u_{j+i},$

all $i \geqslant 0$), V_j must contain a stream of the form $u_j\bot$ which approximates an infinite number of the $u_{j+i}(i \geqslant 0)$. This must be the case for all $j \in J$. Clearly, $u_j\bot \sqsubseteq u$ for all $j \in J$. Thus, for $j<j'$, either $u_j\bot \sqsubseteq u_{j'}\bot$ or $u_{j'}\bot \sqsubseteq u_j\bot$. However, since the $(X_i)_i$ form a $\sqsubseteq_S$-chain and all X_i are flat, $u_{j'}\bot \sqsubseteq u_j\bot$ implies $u_{j'}\bot = u_j\bot$. Consequently, $(u_j\bot)_{j\in J}$ is a $\sqsubseteq$-chain. We again distinguish two cases.

SUBCASE 2.1.
The chain $(u_j\bot)_{j\in J}$ is i.o.i. Then, after applying the interpolation lemma, we obtain the chain

$$u'_0 \sqsubseteq u'_1 \sqsubseteq \cdots \sqsubseteq u_{j_1}\bot = u'_{j_1} \sqsubseteq \cdots \sqsubseteq u_{j_2}\bot = u'_{j_2} \sqsubseteq \cdots$$

with $u'_n \in X_n$ and $u = \text{lub}_n u'_n$.

SUBCASE 2.2.
The chain $(u_j\bot)_{j\in J}$ stabilizes at some $\bar{j}$:

$$u_{j_1}\bot \sqsubseteq u_{j_2}\bot \sqsubseteq \ldots \sqsubseteq u_{\bar{j}}\bot = \cdots$$

This implies that there must be some $k \geqslant \bar{j}$, where X_k contains both $u_{\bar{j}}\bot$ and u_k, and $u_{\bar{j}}\bot \sqsubsetneq u_k$, contradicting the flatness of X_k.
Altogether, if $(u_j)_j$ is i.o.i. and u is infinite there must be a chain $(u'_j)_j$ with $u'_j \in X_j$ and $\text{lub}_j u'_j = u$, i.e., we have found $u \in \bigsqcup_j X_j$. □

The second part of section 4 is devoted to an analysis of various properties of the *min*-operator. We begin with an easy result :

LEMMA 4.5. *For $X, Y \in \mathfrak{S}$ and* **op** *any $\sqsubseteq_S$-monotonic operator: $\mathfrak{S} \times \mathfrak{S} \to \mathfrak{S}$, we have*

$$min\,(X \textbf{ op } Y) = min\,(\,min\,(X) \textbf{ op } min\,(Y))$$

PROOF. Since $X \sqsubseteq_S min\,(X) \sqsubseteq_S X$, and similarly for Y, we have, by the monotonicity of **op**, that

$$X \textbf{ op } Y \sqsubseteq_S min\,(X) \textbf{ op } min\,(Y) \sqsubseteq_S X \textbf{ op } Y$$

Thus, by the monotonicity of *min* , $min\,(X \textbf{ op } Y) \sqsubseteq_S min\,(min(X) \textbf{ op } min\,(Y)) \sqsubseteq min(X\ op\ Y)$.

Since $\sqsubseteq_S$ is an order on flat sets, we have the desired result. □

Next, we prove the *d*-continuity of *min*:

THEOREM 4.6. *Let $(X_i)_i$ be a Cauchy sequence in $(\mathfrak{S}_{nc}, \hat{d})$. Then $min\,(\lim_i X_i) = \lim_i min(X_i)$.*

PROOF. We prove two inclusions.

PART 1.
$\lim_i min\,(X_i) \subseteq min\,(\lim_i X_i)$. Take some $u \in \lim_i min(X_i)$, i.e., $u = \lim_i u_i$, $u_i \in min\,(X_i) \subseteq X_i$. Thus, $u \in \lim_i X_i$. We show that u is a minimal element in $\lim_i X_i$. Assume that there exists some u', $u' \sqsubsetneq u$, and $u' \in \lim_i X_i$. Then $u' = \lim_i u'_i$, $u'_i \in X_i$. We distinguish two cases:
(i) $u' = \lim_i u'_i$ is infinite. This is impossible since $u' \sqsubsetneq u$.

(ii) $u' = \lim_i u'_i$ is finite. Then $u' = u'_{i_0}$ for some u'_{i_0}. If $u' \in A^*$, $u' \sqsubsetneq u$ is impossible.

There remains the case that $u' = \bar{u}\bot$ for some $\bar{u} \in A^*$. If $u \in A^\omega$ then $\exists i > i_0[(u_i \in X_i) \wedge (u'_{i_0} =) u'_i \sqsubsetneq u_i]$.

This contradicts the minimality of u_i. If $u \in A^* \cup A^* \cdot \{\perp\}$, then $\lim_i u_i = u_{j_0}$ for some j_0. Now take $k_0 = \max(i_0, j_0)$. Then $u'_{k_0} \sqsubsetneq u_{k_0}$, which again yields a contradiction.

PART 2.
We prove $min\,(\lim_i X_i) \subseteq \lim_i min(X_i)$. Take $u \in min\,(\lim_i X_i)$. Thus, $u = \lim_i u_i$, $u_i \in X_i$, and u is minimal. We now take $u'_i \in min(X_i)$ such that $u'_i \sqsubseteq u_i$, and consider $\lim_i u'_i$.

SUBCASE 1.
$\lim_i u'_i$ is infinite. We can find a prefix chain $(v_i)_i$ such that $u'_i = v_i w'_i$, $u_i = v_i w_i$ and $w'_i \sqsubseteq w_i$. Moreover, $u = \lim_i u_i = \sup_i v_i = \lim_i u'_i = u'$. Thus, in this case $u \in \lim_i\, min\,(X_i)$.

SUBCASE 2.
$\lim_i u'_i$ is finite, say $\lim_i u'_i = u'_{i_0}$. If $\exists i \forall j \geq i [u_j = u'_j]$ then $u = \lim_i u_i = \lim_i u'_i \in \lim_i\, min(X_i)$. Otherwise, $\forall i \exists j \geq i [u'_{j'} \sqsubsetneq u_j]$. Since $u'_j = u'_{i_0}$ for $j \geq i_0$, we now have that $u_i \neq u'_{i_0}$ for infinitely many i, so $u = \lim_i u_i \neq u'_{i_0} = \lim_i u'_i = u'$. We show that this leads to a contradiction. Once more, we distinguish two subcases:

SUBCASE 2.1.
$u = \lim_i u_i$ is finite, say $\lim_i u_i = u_{j_0}$. Take $k_0 = \max(i_0, j_0)$, Then $u' = \lim_i u'_i = u'_{k_0} \sqsubseteq u_{k_0} = \lim_i u_i = u$. The two facts $u' \sqsubseteq u$ and $u' \neq u$ contradict the minimality of u.

SUBCASE 2.2.
$u = \lim_i u_i$ is infinite. Then there exist v_i, w_i such that $u_i = v_i w_i$, $(v_i)_i$ is a prefix chain, and $\lim_i u_i = \sup_i v_i$. Since $\sup_i v_i$ is infinite we have $\exists j_0 \forall j \geq j_0 [u'_{i_0} \sqsubseteq v_j]$. So $u' = \lim_i u'_i = u'_{i_0} \sqsubseteq \sup_i v_i = u$.

Again, we have $u' \sqsubseteq u$ and $u' \neq u$, a contradiction as in subcase 2.1.

We are now in the position to establish the main technical result relating the operators $\mathbf{op}^{\mathfrak{M}}$ and $\mathbf{op}^{\mathfrak{S}}$.

THEOREM 4.7. *Let* $\mathbf{op}^{\mathfrak{M}}$, $\mathbf{op}^{\mathfrak{S}}$ *be as in definition 3.10, let* $X_1, X_2 \in \mathfrak{S}_{nc}$ *and* $Y_1, Y_2 \in \mathfrak{S}_{ncf}$, *and assume*

$$min\,(X_i) = Y_i, \quad i = 1,2.$$

Then $min\,(X_1\ \mathbf{op}^{\mathfrak{M}} X_2) = Y_1\ \mathbf{op}^{\mathfrak{S}}\, Y_2$.

PROOF. We have, successively,

$$\begin{aligned}
& min\,(X_1\ \mathbf{op}^{\mathfrak{M}} X_2) = && (X_1, X_2 \text{ closed and theorem 3.8f})\\
& min\,(\lim_n X_1(n)\ \mathbf{op}^{\mathfrak{M}} \lim_n X_2(n)) = && (\text{clear})\\
& min\,(\lim_n X_1[n]\ \mathbf{op}^{\mathfrak{M}} \lim_n X_2[n]) = && (\text{d-cont. of } \mathbf{op}^{\mathfrak{M}})\\
& min\, \lim_n (X_1[n]\ \mathbf{op}^{fin} X_2[n]) = && (\text{d-cont. of } min)\\
& \lim_n\, min\,(X_1[n]\ \mathbf{op}^{fin} X_2[n]) = && (\text{lemma 4.5})\\
& \lim_n\, min\,(\, min\,(X_1[n])\ \mathbf{op}^{fin}\ min\,(X_2[n])) = && (\text{theorem 3.8b})\\
& \lim_n\, min\,(\, min\,(X_1)[n]\ \mathbf{op}^{fin}\ min\,(X_2)[n]) = && (\text{assumption})\\
& \lim_n\, min\,(Y_1[n]\ \mathbf{op}^{fin} Y_2[n]) = && (\text{def. } \mathbf{op}^{\mathfrak{S}})\\
& \lim_n\, (Y_1[n]\ \mathbf{op}^{\mathfrak{S}} Y_2[n]) = && (\text{theorem 4.4})
\end{aligned}$$

$$\bigsqcup_n (Y_1[n] \ \mathbf{op}^{\mathbb{S}} Y_2[n]) = \text{(def. } \mathbf{op}^{\mathbb{S}})$$
$$Y_1 \ \mathbf{op}^{\mathbb{S}} \ Y_2 .$$

5. Proof of the equivalence theorem

In section 4 we have collected all results necessary to prove the main result of this paper which we repeat here for convenience:

Theorem 3.15. *For closed and guarded* $s \in \mathcal{L}$

$$\mathcal{M}[\![s]\!] = \mathbb{S}[\![s]\!] .$$

Proof. We first prove a more general result - following a similar pattern as in [BMOZ 2], proof of theorem 2.4.1 - in which s is not necessarily syntactically closed (but still guarded), viz.

$$(*) \quad min\,(\mathcal{M}[\![s]\!]\,(\gamma<X_i\,/x_i>_{i=1}^{n})) = \mathbb{S}[\![s]\!](\gamma<Y_i\,/x_i>_{i=1}^{n})$$

where
(i) $\{x_1, \ldots, x_n\}$ is the set of free process variables in s
(ii) $min\,(X_i) = Y_i$, $i=1,2,\ldots, n$.
We prove (*) by induction on the complexity of s. If $s \equiv a$ the result is obvious and if $s \equiv x$ then $x \equiv x_i$ for some $i \in \{1,\ldots, n\}$ and the desired result follows from (ii). Next, we consider the case that $s \equiv s_1 \ \mathbf{op}\ s_2$, for $\mathbf{op} \in \{;, \cup, \|\}$. Then

$$min\,(\mathcal{M}[\![s]\!]\,(\gamma<X_i/x_i>_i)) =$$
$$min\,(\mathcal{M}[\![s_1 \ \mathbf{op}\ s_2]\!]\,(\gamma<X_i/x_i>_i)) =$$
$$min\,\left[\mathcal{M}[\![s_1]\!]\,(\gamma<X_i/x_i>_i)\ \mathbf{op}^{\mathcal{M}}\mathcal{M}[\![s_2]\!]\,(\gamma<X_i/x_i>_i)\right] =$$

(by the induction hypothesis and theorem 4.7)

$$\mathbb{S}[\![s_1]\!]\,(\gamma<Y_i/x_i>_i)\ \mathbf{op}^{\mathbb{S}}\mathbb{S}[\![s_2]\!]\,(\gamma<Y_i/x_i>_i) =$$
$$\mathbb{S}[\![s_1 \ \mathbf{op}\ s_2]\!]\,(\gamma<Y_i/x_i>_i) =$$

Finally, consider the case that $s \equiv \mu y[s_0]$, for some y and s_0. Without lack of generality, we assume $y \notin \{x_1, \ldots, x_n\}$. Let $Z_0 = U_0 = \{\perp\}$ and

$$Z_{k+1} = \mathcal{M}[\![s_0]\!]\,(\gamma<X_i/x_i\,,\ Z_k/y>_{i=1}^{n})$$
$$U_{k+1} = \mathbb{S}[\![s_0]\!]\,(\gamma<Y_i/x_i\,,\ U_k/y>_{i=1}^{n})$$

Then $\mathcal{M}[\![\mu y[s_0]]\!]\,(\gamma<X_i/x_i>_{i=1}^{n}) = \lim_k Z_k$, and $\mathbb{S}[\![\mu y[s_0]]\!]\,(\gamma<Y_i/x_i>_{i=1}^{n}) = \bigsqcup_k U_k$. We shall prove that (**) $min\,(\lim_k Z_k) = \bigsqcup_k U_k$. By d-continuity of min, the fact that $(U_k)_k$ is a Cauchy sequence and theorem 4.4, we replace (**) by $\lim_k min\,(Z_k) = \lim_k U_k$. Thus, it is sufficient to prove (***) $min(Z_k) = U_k$, $k=0,1,\ldots$. We use induction on k. The case $k=0$ is clear. Next assume (***), to prove $min(Z_{k+1}) = U_{k+1}$, i.e.,

$$min\,(\mathcal{M}[\![s_0]\!]\,(\gamma<X_i/x_i\,,\ Z_k/y>_{i=1}^{n})) =$$
$$\mathbb{S}[\![s_0]\!]\,(\gamma<X_i/x_i\,,\ U_k/y>_{i=1}^{n})$$

Now this follows from the main induction hypothesis (for (*)), with s_0 replacing s and $n+1$

replacing n, and using (***) to establish the $(n+1)$-st part of condition (ii).

We are almost finished with the proof: for closed s, the set of its free variables is empty, and (*) specializes to

$$min\,(\mathcal{M}[\![s]\!](\gamma)) = \mathcal{S}[\![s]\!](\gamma)$$

By the definition of $\mathcal{M}[\![s]\!]$ it is easily seen that, for s closed and guarded, $\mathcal{M}[\![s]\!](\gamma) \subseteq A^* \cup A^\omega$. This follows from definition 3.14b, after varying its clause 3.14b (iv) by taking for X_0 an arbitrary subset of $A^* \cup A^\omega$. (The choice for X_0 is immaterial anyway (see proposition 2.4b); the choice $X_0 = \{\perp\}$ is convenient in the proof just given where we showed $Z_0 = U_0$.) It is then straightforward to show that $\mathcal{M}[\![s]\!](\gamma) \subseteq A^* \cup A^\omega$ by structural induction on s. Thus, $min\,(\mathcal{M}[\![s]\!](\gamma)) = \mathcal{M}[\![s]\!](\gamma)$. Altogether, we have established that, for s closed and guarded, $\mathcal{M}[\![s]\!] = \mathcal{S}[\![s]\!]$, as was to be shown. □

REFERENCES

[Ba] R.J. BACK, *A continuous semantics for unbounded nondeterminism,* TCS 23 (1983) 187-210.

[BBKM] J.W. DE BAKKER, J.A. BERGSTRA, J.W. KLOP, J.-J.CH. MEYER, *Linear time and branching time semantics for recursion with merge,* TCS 34 (1984) 135-156.

[BKMOZ] J.W. DE BAKKER, J.N. KOK, J.-J.CH. MEYER, E.-R. OLDEROG, J.I. ZUCKER, *Contrasting themes in the semantics of imperative concurrency,* in Current Trends in Concurrency: Overviews and Tutorials (J.W. DE BAKKER, W.P. DE ROEVER, G. ROZENBERG, eds.), LNCS 224, Springer (1986) 51-121.

[BMO1] J.W. DE BAKKER, J.-J.CH. MEYER, E.-R. OLDEROG, *Infinite streams and finite observations in the semantics of uniform concurrency,* in: Proceedings 12th ICALP (W. BRAUER, ed.), LNCS 194, Springer (1985) 149-157.

[BMO2] J.W. DE BAKKER, J.-J.CH. MEYER, E.-R. OLDEROG, *Infinite streams and finite observations in the semantics of uniform concurrency,* Report CS-R8512, Centre for Mathematics and Computer Science, 1985 (full version of [BMO1], to appear in TCS).

[BMOZ1] J.W. DE BAKKER, J.-J.CH. MEYER, E.-R. OLDEROG, J.I. ZUCKER, *Transition systems, infinitary languages and the semantics of uniform concurrency,* in: Proceedings 17th ACM STOC, Providence, R.I. (1985) 252-262.

[BMOZ2] J.W. DE BAKKER, J.-J.CH. MEYER, E.-R. OLDEROG, J.I. ZUCKER, *Transition systems, metric spaces and ready sets in the semantics of uniform concurrency,* Report CS-R8601, Centre for Mathematics and Computer Science, 1986 (full version of [BMOZ1]).

[BZ] J.W. DE BAKKER, J.I. ZUCKER, *Processes and the denotational semantics of concurrency,* Inform. and Control 54 (1982) 70-120.

[BK] J.A. BERGSTRA, J.W. KLOP, *Process algebra for synchronous communication,* Inform. and Control 60 (1984) 109-137.

[Br1] M. BROY, *Fixed point theory for communication and concurrency,* IFIP TC2 Working Conference '82 (D. BJØRNER, ed.), Garmisch-Partenkirchen, North-Holland (1983).

[Br2] M. BROY, *A theory for nondeterminism, parallelism, communication and concurrency,* TCS 45 (1986), 1-62.

[Du] J. DUGUNDJI, *Topology,* Allen and Bacon, Rockleigh, N.J. 1966.

[En] R. ENGELKING, *General topology,* Polish Scientific Publishers 1977.

[Ha] H. HAHN, *Reelle Funktionen,* Chelsea, New York, 1948.

[HP] M. HENNESSY, G.D. PLOTKIN, *Full abstraction for a simple parallel programming language,* in: Proceedings 8th MFCS (J. BECVAR ed.), LNCS 74 Springer (1979) 108-120.

[J] K.D. JOSHI, *Introduction to General Topology,* Wiley Eastern ltd, New Delhi, 1983.

[Ku] R. KUIPER, *An operational semantics for bounded nondeteminism equivalent to a denotational one,* IFIP TC2-MC Symp. on Algorithmic Languages (J.W. DE BAKKER & J.C. VAN VLIET, eds.) North-Holland, Amsterdam (1981) 373-398.

[M1] J.-J.CH. MEYER, *Programming calculi based on fixed point transformations: semantics and applications,* dissertation, Free University of Amsterdam, 1985.

[M2] J.-J.CH. MEYER, *Merging regular processes by means of fixed point theory,* TCS 45 (1986).

[MV] J.-J.CH. MEYER, E.P. DE VINK, *Applications of compactness in the Smyth powerdomain of streams,* Report IR-110, Free University, Amsterdam, 1986, to be presented at CAAP 87.

[Ni] M. NIVAT, *Infinite words, infinite trees, infinite computations,* Foundations of Computer Science III.2, Mathematical Centre Tracts 109 (1979) 3-52.

[OH] E.-R. OLDEROG, C.A.R. HOARE, *Specification-oriented semantics for communicating processes,* Acta Informatica 23 (1986) 9-66.

[P] G.D. PLOTKIN, *A structural approach to operational semantics,* Report DAIMI FN-19, Comp.Sci.Dept., Aarhus Univ. 1981.

[R] W.C. ROUNDS, *On the relationship between Scott domains, synchronization trees and metric spaces,* Report Univ. of Michigan CRL-TR-25-83, 1983.