

A new lower bound for the de Bruijn-Newman constant

H.J.J. te Riele

Centre for Mathematics and Computer Science, P.O. Box 4079, 1009 AB Amsterdam,
The Netherlands

Received August 20, 1989/July 30, 1990

Summary. Strong numerical evidence is presented for a new lower bound for the so-called de Bruijn-Newman constant. This constant is related to the Riemann hypothesis. The new bound, -5 , is suggested by high-precision floating-point computations, with a mantissa of 250 decimal digits, of i) the coefficients of a so-called Jensen polynomial of degree 406, ii) the so-called Sturm sequence corresponding to this polynomial which implies that it has two complex zeros, and iii) the two complex zeros of this polynomial. A *proof* of the new bound could be given if one would repeat the computations i) and iii) with a floating-point accuracy of at least 2600 decimal digits.

Subject classification: AMS(MOS): 65E05, 30DXX, 11M26; CR: G1.5.

1 Introduction

Recently, Csordas et al. [2] have introduced the so-called *de Bruijn-Newman constant* A as follows. Let the function $H_\lambda(x)$, $\lambda \in \mathcal{R}$, be defined by

$$(1) \quad H_\lambda(x) := \int_0^\infty e^{\lambda t^2} \Phi(t) \cos(xt) dt,$$

where

$$(2) \quad \Phi(t) = \sum_{n=1}^\infty (2n^4 \pi^2 e^{9t} - 3n^2 \pi e^{5t}) \exp(-n^2 \pi e^{4t}).$$

The function Φ satisfies the following properties:

- i) $\Phi(z)$ is analytic in the strip $-\pi/8 < \Im z < \pi/8$;
- ii) $\Phi(t) = \Phi(-t)$, and $\Phi(t) > 0 (t \in \mathcal{R})$;
- iii) for any $\varepsilon > 0$, $\lim_{t \rightarrow \infty} \Phi^{(n)}(t) \exp[(\pi - \varepsilon) e^{4t}] = 0$, for each $n = 0, 1, \dots$

The function H_λ is an entire function of order one, and $H_\lambda(x)$ is real for real x . From results of de Bruijn [1] it follows that if the Riemann hypothesis is true, then $H_\lambda(x)$ must possess only real zeros for any $\lambda \geq 0$. C.M. Newman has shown [6] that there exists a real number A , $-\infty < A \leq \frac{1}{2}$, such that

- i) $H_\lambda(x)$ has only real zeros when $\lambda \geq A$, and
- ii) $H_\lambda(x)$ has some non-real zeros when $\lambda < A$.

This number A was baptized the *de Bruijn-Newman constant* in [2]. The truth of the Riemann hypothesis would imply that $A \leq 0$, whereas Newman [6] conjectures that $A \geq 0$. In [2] it was proved that $A > -50$.

In this note we will describe high-precision computations which provide strong numerical evidence for the new bound $A > -5$. Moreover, our computations show that trying to prove this result, or improve upon it, would be a formidable task, unless the algorithm used could be improved substantially.

The computations were carried out on the CDC Cyber 995 (about 2 h CPU time for testing), and on the CDC Cyber 205 (about 30 h CPU time for 'production') of SARA (The Academic Computer Centre Amsterdam). Brent's MP package was an indispensable tool for the high-precision floating-point computations. Since this package has not been vectorized, we used the Cyber 205 just as an extremely fast scalar machine.

This note will rely heavily on [2]. We assume the reader to have a copy of [2] at hand (slight change: in the present paper we write b_m and β_m instead of \hat{b}_m and $\hat{\beta}_m$).

2 Algorithm and results

If we expand the cosine in (1) in its Taylor series, we obtain

$$(3) \quad H_\lambda(x) = \sum_{m=0}^{\infty} \frac{(-1)^m b_m(\lambda) x^{2m}}{(2m)!},$$

where

$$b_m(\lambda) = \int_0^{\infty} t^{2m} e^{\lambda t^2} \Phi(t) dt,$$

$m=0, 1, \dots; \lambda \in \mathbb{R}$. The n -th degree Jensen polynomial $G_n(t; \lambda)$ associated with H_λ is defined by

$$(4) \quad G_n(t; \lambda) := \sum_{k=0}^n \binom{n}{k} \frac{k! b_k(\lambda)}{(2k)!} t^k,$$

and it is shown in [2] that if there exists a positive integer m and a real number $\hat{\lambda}$ such that $G_m(t; \hat{\lambda})$ possesses a non-real zero, then $\hat{\lambda} < A$. The problem is to find m , given $\hat{\lambda}$.

In [2] a lower bound for A was constructively obtained as follows. For $\lambda = -50$ and $n = 16$ the moments $b_k(\lambda)$, $k = 0, \dots, n$ were computed with a known precision, by means of Romberg quadrature. The approximate Jensen poly-

mial which we obtain by using this numerical approximation of $b_k(\lambda)$ in (4) is denoted by $g_n(t; \lambda)$. The so-called Jenkins algorithm was used to compute *all* the zeros of g_n (including one complex zero and its complex conjugate). A theorem of Ostrowski was then invoked to find an upper bound for the distance of each of these zeros to the corresponding zeros of G_n . This error bound was small enough to guarantee that the complex zero of $g_{16}(t; -50)$ found by high-precision computation indeed is an approximation of a complex zero of the Jensen polynomial $G_{16}(t; -50)$.

The sensitivity of the zeros of polynomials to errors in their coefficients required that the computations were performed in very high precision. Csordas et al. [2] used 110 digits of precision for their proof that $-50 < \lambda$. As a partial check, we repeated their computations in double precision on a CDC Cyber 995 (which means an accuracy of about 28 decimal digits) and could reproduce the complex zero of $g_{16}(t; -50)$ with an accuracy of about 20 decimal digits. This illustrates the large amount of extra work needed to provide a *proof* of the existence of complex zeros of the Jensen polynomial $G_n(t; \lambda)$.

In order to improve the result of Csordas et al., we realized that the degree of the Jensen polynomial $G_n(t; \lambda)$ which possesses complex zeros, might grow very fast with λ . Consequently, finding *all* the zeros of G_n , $n=1, 2, \dots$ (in order to prove the existence of complex ones) might become very expensive. Therefore, we decided to use so-called *Sturm sequences* [4] to get an indication of the existence of any complex zeros of the given Jensen polynomial. The computation of a Sturm sequence is much simpler than computing all the zeros of a given polynomial.

A Sturm sequence associated with a given polynomial $p_0(x)$ of degree m is a sequence of polynomials $p_0(x), p_1(x), \dots$ of strictly decreasing degree which can be defined as follows:

$$\begin{aligned} p_1(x) &:= p_0'(x), \\ p_{i-1}(x) &:= q_i(x)p_i(x) - p_{i+1}(x), \quad i=1, 2, \dots, \end{aligned}$$

where $q_i(x)$ is found by the Euclidean algorithm, such that the degree of $p_{i+1}(x)$ is less than the degree of $p_i(x)$. If $p_0(x)$ has only simple zeros, $p_i(x)$ has degree $m-i$, and the Sturm sequence consists of $m+1$ polynomials $p_0(x), \dots, p_m(x)$. Let $v(a)$ be the number of sign changes in the sequence $\{p_i(a)\}_{i=0}^m$ (where zero values are skipped). Then $v(a) - v(b)$ is the number of real zeros of the polynomial $p_0(x)$ on the interval $[a, b]$. Note that $v(\pm\infty)$ can be determined by inspection of the signs of the highest degree coefficients of the polynomials in the Sturm sequence.

Our algorithm now works as follows. Suppose we know λ_0 and $m=m(\lambda_0)$, which is the smallest value for which $g_m(t; \lambda_0)$ has complex zeros (to start with, we take $\lambda_0 = -50$ and $m=16$ from [2]). Then for a new value of λ which is somewhat larger than λ_0 we compute $\beta_i(\lambda)$, $i=0, 1, \dots$, and for each new β_i we compute the coefficients of the associated Jensen polynomial, and by means of the associated Sturm sequence, its number of real zeros on the interval $(-\infty, 0]$. This is continued until we have found n for which $g_n(t; \lambda)$ should have complex zeros. Then we try to compute a complex zero of this polynomial by means of the Newton process, where the starting value is chosen as follows. Let $z = z(\lambda_0)$ be the known complex zero of $g_m(t; \lambda_0)$. We tabulate the values

Table 1. Minimal degrees $m(\lambda)$ of Jensen polynomials with complex zeros

λ	$m(\lambda)$	Complex zeros of $g_m(t; \lambda)$		Accuracy used
		\Re	$\pm \Im$	
-50	16	-220.9191117	7.092565255	28D
-49	16	-217.9076244	5.773253615	28D
-48	16	-214.9084360	4.111013736	28D
-47	16	-211.9217860	1.006843660	28D
-46	17	-202.2196553	5.677704348	28D
-45	17	-199.3211883	3.991036911	28D
-44	17	-196.4360833	0.462709708	28D
-43	18	-187.4386728	4.830351149	28D
-42	18	-184.6425759	2.749091911	28D
-41	19	-176.2289375	4.969975476	28D
-40	19	-173.5216696	3.024436421	28D and 40D
-30	27	-116.8258164	2.400595686	28D and 50D
-20	41	-111.0654985	1.322239430	50D
-10	97	-45.53019819	0.156978360	75D
-5	406	-24.34071458	0.031926616	250D

of $g_n(t; \lambda)$ and its derivative, for some real values of t around $\Re(z)$, and we look for a local *positive minimum*, or a local *negative maximum*. In our experience, such a minimum, or maximum, is easy to find if λ is not too far away from λ_0 . Then we take $c + di$ as starting value for the Newton process where c is the value of t for which $g_n(t; \lambda)$ assumes its local minimum or maximum, and where $d = \Im(z)$.

In this way we found complex zeros of $g_n(t; \lambda)$ for $\lambda = -50(1) - 40, -30, -20, -10, -5$. Table 1 presents the values for which we have determined the polynomial $g_n(t; \lambda)$ of smallest degree with complex zeros by means of the associated Sturm sequence. This degree is denoted by $m = m(\lambda)$. In all cases this Jensen polynomial has $m(\lambda) - 2$ real zeros. Table 1 also lists the complex zeros found, truncated to 10 decimal digits, and the accuracy used. For λ close to -50 , the degree of the Jensen polynomial with complex zeros does not increase too quickly with λ . However, from $\lambda \approx -20$ this pattern changes drastically, as Table 1 shows. As λ increases, the imaginary parts of the complex zeros found seem to tend to zero.

Our computations do not provide a *mathematical* proof of the existence of complex zeros of $G_n(t; \lambda)$, although there is strong numerical evidence. A proof of the new bound $-5 < \mathcal{A}$ along the lines of Csordas et al. would require an extension of the accuracy we used (250 decimal digits) to at least 2600 decimal digits.

The following simple error analysis, carried out for the case $\lambda = -5$, may help to convince the reader. If we define $g_{406}(t; -5) =: \sum_{k=0}^{406} a_k t^{406-k}$, then, since the a_k are numerical approximations of the coefficients of G_{406} , computed with an accuracy of 250 decimal digits, we may write

$$G_{406}(t) = g_{406}(t) + \varepsilon h(t),$$

where $\varepsilon = 10^{-250}$ and $h(t) = \sum_{k=0}^{406} \delta_k a_k t^{406-k}$ with $|\delta_k| < 1$. If θ is a simple zero of g_{406} , and θ_ε the corresponding zero of G_{406} , then we have, as an approximation of the first order in ε (cf. [7, formula (5.8.1)]):

$$\theta_\varepsilon = \theta - \varepsilon \frac{h(\theta)}{g'_{406}(\theta)}.$$

Hence,

$$|\theta_\varepsilon - \theta| < \varepsilon \frac{|h(\theta)|}{|g'_{406}(\theta)|} < \varepsilon \frac{g_{406}(|\theta|)}{|g'_{406}(\theta)|},$$

since the coefficients of g_{406} are positive. For $\theta = -24.3407 \dots + 0.03192 \dots i$ we found $g_{406}(|\theta|) = 4.837 \dots \times 10^{12}$ and $|g'_{406}(\theta)| = 7.824 \dots \times 10^{-17}$, so that $|\theta_\varepsilon - \theta| < 10^{-221}$ and θ_ε is indeed a complex zero of $G_{406}(t; -5)$.

3 Computational details

In this section we shall explain the main details of how we computed $\beta_m(\lambda)$ and the Sturm sequences of $g_n(t; \lambda)$.

We write $b_m(\lambda)$ as the sum

$$(5) \quad b_m(\lambda) = \int_0^a t^{2m} e^{\lambda t^2} \Phi(t) dt + \int_a^\infty t^{2m} e^{\lambda t^2} \Phi(t) dt$$

(Csordas et al. used $a = 1$). An upper bound for the second integral of (5) is found as follows. The function $t^{2m} e^{\lambda t^2}$ has maximum value

$$\exp \left[m \left(-1 + \log \frac{m}{-\lambda} \right) \right]$$

(for $t = (-m/\lambda)^{\frac{1}{2}}$), so that

$$\begin{aligned} \int_a^\infty t^{2m} e^{\lambda t^2} \Phi(t) dt &< \exp \left[m \left(-1 - \log \frac{m}{-\lambda} \right) \right] \int_a^\infty \Phi(t) dt \\ &< \frac{\pi}{2} \exp \left[m \left(-1 + \log \frac{m}{-\lambda} \right) + 5a - \pi e^{4a} \right] \end{aligned}$$

(cf. [3], in. (3.7)). This bound is used, for given λ and m , to choose a such that the contribution of the second integral in (5) to the value of $b_m(\lambda)$ is negligible, in view of the precision used. E.g. for $\lambda = -5$, we chose $a = 1.65$. For $m = 406$, this yields an upper bound of 10^{-400} on the value of the second integral in (5). The smallest $b_m(-5)$ we computed is $b_{344}(-5) = 1.46822 \dots \times 10^{-73}$. Since we worked with a precision of 250 decimal digits, it follows that the contribution of the second integral in (5) to $\{b_m(-5)\}_{m=0}^{406}$ is indeed negligible.

Let $\Phi_N(t)$ denote the sum of the first N terms of (2); then we have (cf. [3], eq. (4.6))

$$0 < \Phi(t) - \Phi_N(t) < \pi N^3 \exp(5t - \pi N^2 e^{4t}) \quad (0 \leq t < \infty).$$

Given t , the number N is chosen such that the right hand side is less than 10^{-4} where A is the number of decimal digits of precision employed in the computations plus $\{-\log_{10}(\text{first term of } \Phi(t))\}$ (since this first term determines the size of $\Phi(t)$). Since the N in the exp is dominating, it is sufficient (most of the time) to choose N to be the smallest integer larger than $\sqrt{(e^{-4t}(A \log 10 + 5t)/\pi)}$.

Using the same notation as in [2], we now have to compute the integral

$$(6) \quad b_m^{(2)}(\lambda) = \int_0^a t^{2m} e^{\lambda t^2} \Phi_N(t) dt$$

to sufficient accuracy. In [2] this was done by Romberg quadrature. However, by inspecting the Romberg table for $b_m^{(2)}(\lambda)$, we noticed that when going from left to right, i.e., when comparing T_{ij} with $T_{i,j+1}$, the accuracy did *decrease* (rather than increase, as one would expect: cf., e.g., [7, p. 141]). Moreover, the most accurate results were found in the first column of the Romberg table (just the trapezoidal rule results for step $a, a/2, a/4, \dots$), and the convergence in this column was much faster than quadratic. An explanation is given by the fact that the integrand in $b_m(\lambda)$ is an *even* function, and under certain conditions given in Theorem 2.2 of [5] the convergence of the trapezoidal rule for such functions is *exponential*. The integrand $b_m(\lambda)$ happens to satisfy these conditions. Therefore, it is unnecessary to apply Romberg quadrature. We just applied the composite trapezoidal rule with step $a, a/2, a/4, \dots$, until a sufficiently small correction was obtained. For the computation of $\beta_m(-5)$ we never needed to work with a step less than $a/1024$. Before applying the trapezoidal rule, a table of values of $e^{\lambda t^2} \Phi(t)$ was precomputed for $t = ja/1024, j=0, \dots, 1024$, since (a selection of) these values are needed for each $\beta_m(\lambda)$. In the final steps, we always observed an approximate doubling of the number of correct digits upon halving the step.

The Sturm sequence associated with the polynomial $g_m(t; \lambda)$ was computed as follows. Let $p_0(x) := g_m(x; \lambda)$ and let

$$p_i(x) := \sum_{j=0}^{m-i} c_{ij} x^{m-i-j}, \quad i=0, \dots, m.$$

The coefficients c_{0j} of $p_0(x)$ are computed by means of the relation (which follows from (4)):

$$c_{0j} = \frac{m!}{j!(2(m-j))!} \beta_{m-j}(\lambda), \quad j=0, \dots, m.$$

Since $p_1(x) = p_0'(x)$ we have $c_{1j} = (m-j) c_{0j}, j=0, \dots, m-1$. Let $q_i(x) := q_{i0}x + q_{i1}, i=1, \dots, m-1$. Then, by applying the definition of a Sturm sequence given in § 2, we find q_{i0} and $q_{i1}, i=1, \dots, m-1$, from,

$$\begin{aligned} q_{i0} c_{i0} - c_{i-1,0} &= 0, \\ q_{i0} c_{i1} + q_{i1} c_{i0} - c_{i-1,1} &= 0, \end{aligned}$$

and $c_{i+1,j}, j=0, \dots, m-i-2$ from

$$c_{i+1,j} = q_{i0} c_{i,j+2} + q_{i1} c_{i,j+1} - c_{i-1,j+2}$$

and $c_{i+1,m-i-1}$ from

$$c_{i+1,m-i-1} = q_{i1} c_{i,m-i} - c_{i-1,m-i+1}.$$

The number of negative real zeros of $g_m(t; \lambda)$ is just the number of sign changes in the sequence $\{(-1)^{m-i} c_{i0}\}_{i=0}^m$ minus the number of sign changes in the sequence $\{c_{i,m-i}\}_{i=0}^m$. In some instances, when going from g_m to g_{m+1} , this difference dropped down from m sharply. It turned out that this was caused always by insufficient precision used in the computation of the Sturm sequence associated with g_{m+1} . The normal pattern (i.e., finding $v(-\infty) - v(0) = m+1$, or $m-1$) could be restored easily, by increasing the accuracy.

We have used the Sturm sequence technique to get an easy *numerical indication* for the existence of complex zeros of $g_m(t; \lambda)$. The actual computation of these complex zeros by means of the Newton process provides the numerical evidence for their existence. The error analysis given at the end of § 2 is meant to convince the reader who would rightly notice that these numerical computations do not provide a rigorous mathematical proof.

Acknowledgements. We like to thank Dr. N.M. Temme for his help and, in particular, for pointing out [5] to us.

We also acknowledge the provision of computer time on the Cyber 205 by the Dutch Working Group on the Use of Supercomputers.

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