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MEDIAN GRAPHS AND HELLY HYPERGRAPHS

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Median graphs and Helly hypergraphs^{*)}

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ABSTRACT

One-to-one correspondences are established between the following combinatorial structures: (i) median interval structures (or median segments, introduced by SHOLANDER); (ii) maximal Helly hypergraphs such that with each edge also its complement is in the hypergraph; and (iii) median graphs (connected graphs such that for any three vertices u, v, w there is exactly one vertex x such that d(u,v) = d(u,x) + d(x,v), d(v,w) = d(v,x) + d(x,w)and d(w,u) = d(w,x) + d(x,u), where d is the distance function of the graph).

KEY WORDS & PHRASES: median, interval structure, Helly hypergraph, copair, median graph.

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O. INTRODUCTION

In this paper one-to-one correspondences will be established between three at first sight fairly distinct concepts. These concepts are:

- (i) median interval structures introduced by M. SHOLANDER [7], [8] under the name of median segments (cf. 1.1);
- (ii) maximal Helly copair hypergraphs (i.e. simple Helly hypergraphs, the edge-set of which contains with each edge its complement, and which are maximal with respect to this property; see 1.2); and

(iii) median graphs, introduced in section 1.3.

The one-to-one correspondences are established in section 2.

In section 3 is elaborated how to construct a maximal Helly copair hypergraph from a median graph, using results of SHOLANDER [9].

With minor adaptations we adopt the terminology of BERGE [1] on hypergraphs, of WILSON [10] on graphs and of BIRKHOFF [2] on lattice theory.

1. DEFINITIONS AND PRELIMINARIES

Throughout this paper V denotes a fixed finite set.

1.1. INTERVAL STRUCTURES. A function $I: V \times V \rightarrow P(V)$ is called an *interval* structure on V if

(I1) $x, y \in I(u, v)$ iff $I(x, y) \subset I(u, v)$ $(x, y, u, v \in V)$, (I2) $I(u, v) \cap I(v, w) \cap I(w, u) \neq \emptyset$ $(u, v, w \in V)$.

Each set I(u,v) is called an *interval*. A subset U of V is I-convex if for all u, $v \in V$ the interval I(u,v) is contained in U. The notion of interval structure was introduced in [3]. Examples of interval structures on V can be obtained from trees with vertex-set V (then take I(u,v) = $\{w \in V \mid w \text{ lies on the shortest } u,v-path\}$), and from lattices (V,\leq) (in this case $I(u,v) = \{w \in V \mid u \land v \leq w \leq u \lor v\}$).

If I satifies condition (I1) and the following condition

(I2')
$$|I(u,v) \cap I(v,w) \cap I(w,u)| = 1$$
 (u,v,w $\in V$),

then I is called a *median interval structure* on V. Interval stuctures obtained from trees as indicated above are median interval structures. An interval structure obtained from a lattice is a median interval structure iff the lattice is distributive (cf. [2]). SHOLANDER [8] has given the following characterization of median interval structures (he used the term median segments):

<u>THEOREM 1</u>. (SHOLANDER [8]) A function $I: V \times V \rightarrow P(V)$ is a median interval structure on V iff

1.2.<u>HYPERGRAPHS</u>. In this paper a hypergraph H = (V, E) consists of a vertex-set V and a family $E \subset P(V)$ of nonvoid subsets of V, the members of which are called *edges*. Occasionally we will write E instead of (V,E).

A hypergraph is a *Helly hypergraph* if it satisfies the *Helly property*, i.e. every subfamily of E, any two members of which meet, has a non-empty intersection. For vertices u and v of the hypergraph (V,E) define

 $I_{\mathcal{E}}(u,v) = \cap \{B \in \mathcal{E} \mid u,v \in B\}.$

A theorem of P.C. GILMORE (see [5], or [1] p. 396) can be formulated as follows:

<u>THEOREM 2</u>. (GILMORE) A hypergraph (V,E) satisfies the Helly property iff I_F is an interval structure on V.

As a consequence of GILMORE's theorem we have: Let I be an interval structure on V. Any family E of nonvoid I-convex subsets of V satisfies the Helly property.

A hypergraph (V,E) with the property that $V \setminus B \in E$ for all $B \in E$ will be called a *copair hypergraph*. We call the set {B,V\B} a copair of V and { ϕ ,V} the trivial copair. A Helly copair hypergraph of course is a copair hypergraph, which satisfies the Helly property. Finally a *maximal Helly copair hypergraph* (V,E) is a Helly copair hypergraph such that: if {A,V\A} is a non-trivial copair and $E \cup \{A, V \setminus A\}$ satifies the Helly property then $A \in E$.

A hypergraph (V,E) is said to separate vertices if for any two distinct vertices $u, v \in V$ there exists an edge $A \in E$ such that $u \in A$ and $v \notin A$.

<u>LEMMA 3.</u> Let (V,E) be a Helly copair hypergraph. Then (V,E) is maximal iff (V,E) separates vertices.

PROOF. Note that (V,E) separates vertices iff $I_F(v,v) = \{v\}$ for all $v \in V$.

Assume that E does not separate vertices. That is there exists a vertex $v \in V$ such that $I_E(v,v)$ contains besides v another vertex. Using GILMORE's theorem it can be verified that in this case $E \cup \{\{v\}, V \setminus \{v\}\}\$ satisfies the Helly property. Therefore E is not maximal.

To prove sufficiency of vertex separation let {A,V\A} be a nontrivial copair of V not in E. Take a vertex $u \in A$ and a vertex $v \in V\setminus A$ such that $|I_E(u,v)|$ is as small as possible. We assert that $I_E(u,v) \cap A = \{u\}$ and $I_F(u,v)\setminus A = \{v\}$.

For suppose $I_E(u,v) \cap A \neq \{u\}$ and let $w \in I_E(u,v) \cap A$ with $w \neq u$. Since E separates vertices, there exists an edge $C \in E$ such that $w \in C$ and $u \notin C$. Then we have that $v \in C$. So $u \notin I_E(w,v) \subset I_E(u,v)$, contradicting the minimality of $I_E(u,v)$. In the same way we prove $I_E(u,v) \setminus A = \{v\}$. Hence $I_E(u,v) = \{u,v\}$.

Let B ϵ E be an edge such that v ϵ B and u \notin B. Then A \cap B $\neq \emptyset$ or (V\A) \cap (V\B) $\neq \emptyset$, since A \notin {B,V\B} \subset E; say A \cap B $\neq \emptyset$. Now the set of edges, which contain both u and v, together with A and B forms a family of subsets of V, any two members of which meet. The intersection of this family equals

$$I_{E}(u,v) \cap A \cap B = \{u,v\} \cap A \cap B,$$

which clearly is empty. Thus $E \cup \{A, V \setminus A\}$ does not satisfy the Helly property.

COROLLARY 4. Let (V,E) be a maximal Helly copair hypergraph. Then

$$|E| \ge 2 \quad \boxed{2} \log |V|$$

1.3 <u>MEDIAN GRAPHS</u>. Let G be a simple loopless graph with vertex-set V and distance function d. G will be called a *median graph* if it is connected and satisfies the graph median property, i.e. for any $u,v,w \in V$ there exists precisely one vertex $x \in V$, called the graph median of u,v and w, such that

 $\begin{cases} d(u,x) + d(x,v) = d(u,v) \\ d(v,x) + d(x,w) = d(v,w) \\ d(w,x) + d(x,u) = d(w,u). \end{cases}$

Note that all trees and the n-cubes are median graphs. It is easy to see that each median graph is bipartite.

2. THE THEOREM

<u>THEOREM 5.</u> There exists a one-to-one correspondence between the median interval structures on V, the maximal Helly copair hypergraphs with vertexset V and the median graphs with vertex-set V. The one-to-one correspondences are indicated in the following diagram, which commutes in all directions.



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The proof of the theorem amounts to the following propositions (the direct correspondence between median graphs and maximal Helly copair hypergraphs will be explained in section 3).

For vertices u and v of the graph G = (V, E) define

 $I_{G}(u,v) = \{w \in V \mid w \text{ lies on a shortest } u,v-path in G\}.$

<u>PROPOSITION 6</u>. Let G = (V, E) be a median graph. Then I_G is a median interval structure on V.

<u>PROOF.</u> I_C satisfies the conditions mentioned in theorem 1. \Box

<u>PROPOSITION 7.</u> Let I be a median interval structure on V. Define the graph $G_{\rm T}$ with vertex-set V by

 $uv \in E(G_{\tau})$ iff $u \neq v$ and $I(u,v) = \{u,v\}$ $(u,v \in V)$.

Then $\boldsymbol{G}_{_{\boldsymbol{T}}}$ is a median graph.

<u>PROOF</u>. We will prove that G_I is connected and that $I_{G_I} = I$. Then clearly G_T is a median graph.

First observe that for $u, v, w \in V$ we have

$$w \in I(u,v)$$
 iff $I(u,w) \cap I(w,v) = \{w\}$.

Thus for $w \in I(u,v) \setminus \{u,v\}$ holds $u \notin I(w,v) \subset I(u,v)$ and $v \notin I(u,w) \subset I(u,v)$. Using this it is easily verified by induction on |I(u,v)| that I(u,v)induces a connected subgraph of G_T for all $u,v \in V$. Hence G_T is connected.

To prove that $I(u,v) = I_{G_{I}}(u,v)$ for all $u,v \in V$ we use induction on d(u,v). Clearly $I(u,v) = I_{G_{I}}(u,v)$ for all $u,v \in V$ with $d(u,v) \leq 1$. So take vertices $u,v \in V$ with d(u,v) > 1.

Let w $\in I_{G_{I}}(u,v) \setminus \{u,v\}$. Then d(u,w) < d(u,v) and d(w,v) < d(u,v), so $I_{G_{I}}(u,w) = I(u,w)$ and $I_{G_{I}}(w,v) = I(w,v)$. Since clearly $I_{G_{I}}(u,w) \cap I_{G_{I}}(w,v) = \{w\}$, we have w $\in I(u,v)$ and thus $I_{G_{I}}(u,v) \subset I(u,v)$. Assume $I(u,v) \setminus I_{G_{I}}(u,v) \neq \emptyset$.

For any vertex w \in I(u,v)\I_{GT}(u,v) we must have I(u,w) \cap I_{GT}(u,v)={u},

and similarly $I(w,v) \cap I_{G_{I}}(u,v) = \{v\}$. For if $w' \in I(u,w) \cap I_{G_{I}}(u,v)$, with $w' \neq u$, then $w \in I(w',v)$ and by the induction hypothesis $I(w',v) = I_{G_{I}}(w',v) \subset I_{G_{I}}(u,v)$. Hence $w \in I_{G_{I}}(u,v)$, contradicting the choice of w.

Since I(u,v) induces a connected subgraph of G_I , there exists a path P from u to v, all the internal vertices of which lie in $I(u,v) \setminus I_{GI}(u,v)$. Clearly the length of P exceeds d(u,v) so P has at least two distinct internal vertices, say x and y.

Since $d(u,v) \ge 2$, there exists a vertex $z \in I_{G_I}(u,v) \setminus \{u,v\}$. By the induction hypothesis we have $I(u,z) = I_{G_I}(u,z)$ and $I(z,v) = I_{G_I}(z,v)$. Now

$$u \in I(u,z) \cap I(u,x) = I_{G_I}(u,z) \cap I(u,x) \subset I_{G_I}(u,v) \cap I(u,x) = \{u\}.$$

So $u \in I(z,x)$. Similarly $v \in I(z,x)$ and thus $I(u,v) \subset I(z,x) \subset I(u,v)$. In the same way it follows that I(u,v) = I(z,y). But then

$$x, y \in I(x, y) = I(z, x) \cap I(x, y) \cap I(y, z),$$

contradicting the fact that I is a median interval structure. Conclusion: $I(u,v) = I_{G_T}(u,v)$.

In the proof of the preceding proposition we have seen that for a median interval structure I holds: $I_{GI} = I$. Furthermore from propositions 6 and 7 follows immediately that, when G is a median graph, we have $G_{I_G} = G$.

<u>PROPOSITION 8.</u> Let (V, E) be a maximal Helly copair hypergraph. Then I_E is a median interval structure on V.

<u>PROOF</u>. Assume that there exist vertices $u, v, w \in V$ such that $x, y \in I_E(u, v) \cap I_E(v, w) \cap I_E(w, u)$ for vertices $x, y \in V$, with $x \neq y$. According to lemma 3 there is an edge $B \in E$ such that $x \in B$ and $y \notin B$. Then one of the edges B and $V \setminus B$, say B, must contain at least two of the three vertices u, v and w, say u and v. But then $y \notin I_F(u, v)$. Contradiction. \Box

PROPOSITION 9. Let I be a median interval structure on V and let

 $E_{I} = \{B \subset V | \phi \neq B \neq V, B \text{ and } V \setminus B \text{ are } I-convex\}.$

Then (V, E_{τ}) is a maximal Helly copair hypergraph.

<u>PROOF</u>. Clearly (V, E_I) is a Helly copair hypergraph. By lemma 3 it suffices to show that E_I separates vertices. So suppose that for vertices $u, v \in V$, with $u \neq v$, there is no edge B such that $u \in B$ and $v \notin B$. Assume futhermore that u and v are such that |I(u,v)| is as small as possible.

We first prove that $I(u,v) = \{u,v\}$. Suppose $w \in I(u,v) \setminus \{u,v\}$. Since |I(u,w)| < |I(u,v)|, there exists an edge A such that $u \in A$ and $w \notin A$. It follows that $v \in A$ (u and v cannot be separated). So $w \in I(u,v) \subset A$, for A is I-convex, contradicting $w \notin A$. Therefore $I(u,v) = \{u,v\}$.

Now let $B = \{z \in V | v \notin I(u,z)\}.$

Then $V \setminus B = \{z \in V | u \notin I(z,v)\}$, since $I(u,z) \cap I(z,v) \cap \{u,v\}$ is a singleton. We assert that B and $V \setminus B$ are I-convex, that is $B \in E_I$. Since $u \in B$ and $v \notin B$ this contradicts our assumption that E_I does not separate vertices.

We only prove that B is I-convex (the I-convexity of $V \setminus B$ can be treated similarly).

Note that for each $z \in B$ we have $I(u,z) \subseteq B$, since $v \notin I(u,z)$. Let $x,y \in B$ and suppose $I(x,y) \notin B$. Take $w \in I(x,y) \setminus B$. Since $I(u,x) \subseteq I(v,x)$ and $I(u,y) \subseteq I(v,y)$ we have that

 $\{z\} = I(u,x) \cap I(x,y) \cap I(y,u) = I(v,x) \cap I(x,y) \cap I(y,v)$

for some z \in B. Now also

 $\{z\} \subset I(z,w) \cap I(z,v) \subset I(x,y) \cap I(x,v) \cap I(y,v) = \{z\},\$

since $z, w \in I(x, y)$ and $z \in I(u, x) \cap I(u, y) \subset I(v, x) \cap I(v, y)$. This implies $z \in I(w, v)$ according to the observation made at the beginning of the proof of proposition 7. So $I(z, v) \ll I(w, v)$. But, since $w \notin B$, $u \notin I(w, v)$ and thus $u \notin I(z, v)$, that is $z \in V \setminus B$, contradicting the fact that $z \in B$.

From propositions 8 and 9 we deduce: let I be a median interval structure on V, then $I_{E_I} = I$; and let (V,E) be a maximal Helly copair hypergraph, then $E_{IF} = E$.

3. MEDIAN GRAPHS AND HELLY HYPERGRAPHS

In this section the direct correspondence between median graphs and maximal Helly copair hypergraphs with vertex-set V, mentioned in the theorem, is further elaborated.

3.1. MEDIAN SEMILATTICES. Let (V, \leq) be a partially ordered set (poset). v is said to *cover* u (u, v \in V), if u \leq v and there is no w \in V such that u < w < v. A *semilattice* (V, \leq) is a poset, in which any two elements u,v have a greatest lower bound u \wedge v. For u, v \in V set $[u,v] = \{w \in V | u \leq w \leq v\}$. The semilattice (V, \leq) is called *distributive* if ([u,v], \leq) is a distributive lattice for all u, v \in V. The semilattice is said to satisfy the *coronation property* if for any three elements u,v, w \in V, such that the three least upper bounds u \vee v, v \vee w, w \vee u exist, there exists a least upper bound u \vee v \vee w.

A median semilattice is a distributive semilattice, which satisfies the coronation property. This concept was introduced by SHOLANDER [9].

On a median semilattice (V, \leq) the ternary operation $(u, v, w) = (u \land v) \lor \lor (u \land w) \lor (w \land u) \in V$ can be defined, called the *median* of u,v and w (SHOLANDER [9] also characterized medians).

We review some results of SHOLANDER [9] reformulating them in our terminology:

(A) Each median semilattice (V, \leq) yields a median interval structure $I_{<}$ on V, where

$$I_{\leq}(u,v) = \{w | w \text{ is the median of } u,v,w\} \qquad (u,v \in V)$$

(B) Let I be a median interval structure on V and $u \in V$. Define an ordering $\leq_{I,U}$ on V by

$$v \leq I, u \quad \text{w iff } v \in I(u, w) \quad (v, w \in V).$$

Then $(V, \leq_{I,u})$ is a median semilattice. Furthermore the correspondences given in (A) and (B) commute.

(C) Let (V, \leq) be a median semilattice. Then (V, \leq) can be embedded in a Boolean algebra by an order preserving mapping, which also perserves the covering relation in (V, \leq) .

3.2 <u>CUTSET COLOURINGS</u>. A *cutset colouring* of a connected graph is a colouring of the edges in such a way that the edges of any colour form a matching as well as a cutset (i.e. a minimal disconnecting edge-set). If we want to establish a cutset colouring of a graph we are forced to colour non-adjacent edges in each circuit of length four with the same colour. So the n-cube admits a cutset colouring with n colours, which is uniquely determined up to the labelling of the colours. Deleting the edges with a given colour from the n-cube breaks the graph up into two components, which both are (n-1)-cubes.

Note that not all connected graphs admit a cutset colouring. Necessary conditions for the existence of a cutset colouring of the edges of a connected graph are for instance that the graph is simple, loopless and bipartite and that it does not contain $K_{2,3}$ as a subgraph.

3.3 <u>MEDIAN GRAPHS AND MAXIMAL HELLY COPAIR HYPERGRAPHS</u>. The *diagraph* of a poset (V, \leq) is the graph with vertex-set V, in which two vertices are joined by an edge iff one of the two covers the other in the poset. Clearly, the diagraph of the Boolean algebra on 2^n elements is the n-cube. As a consequence of (A) and (B) and propositions 6 and 7 we have

<u>**PROPOSITION 10.**</u> Let G be a graph. Then G is a median graph iff G is the diagraph of a median semilattice.

<u>PROPOSITION 11</u>. Let G be a graph. Then G is a median graph iff G is a connected induced subgraph of an n-cube such that with any three vertices of G their graph median in the n-cube also is a vertex of G.

PROOF. The only if part follows from proposition 10 and (C).

The if part follows as soon as we have proved that the distance in G between two vertices equals their distance in the n-cube. Let d be the distance function of G and e that of the n-cube. Assume that there are vertices u,v of G with $d(u,v) \neq e(u,v)$ and let k:= d(u,v) be as small as possible.

Note that k > 2.

Let w be a vertex of G with d(u,w) = 2 and d(w,v) = k - 2. Then e(u,w) = 2 and e(w,v) = k - 2. Let z be the graph median of u,v and w in the n-cube. Thus z is a vertex of G.

If z = w then e(u,v) = e(u,w) + e(w,v) = 2 + k - 2. So $z \neq w$. But then, since e(u,w) = 2 = d(u,w), z is a common neighbour of u and w. Now e(z,v) = e(w,v) - e(w,z) = k - 2 - 1 = k - 3. Thus $d(u,v) \le d(u,z)+d(z,v)=$ = 1 + e(z,v) = k - 2 < k, which is a contradiction.

Let G be a median graph with vertex-set V. Embed G in an n-cube K with n as small as possible. Since G is connected G has at least one edge of each colour from the cutset colouring of K.

The cutset colouring of K induces an edge colouring of G. According to proposition 10 with any two vertices u and v of G a shortest u,v-path of u lies entirely in G. So the induced edge colouring of G in fact is a cutset colouring. Any cutset from this colouring induces a copair of V: after deleting the cutset from G the graph breaks up into two components, the vertex-sets of which form the complementary subsets of the copair. In this way the cutset colouring of G induces a copair hypergraph (V, E_G). Since G is an induced subgraph of K it follows that E_G consists of I_G -convex subsets of V. Besides it follows that E_G separates vertices. And thus according to lemma 3 (V, E_G) is a maximal Helly copair hypergraph. Furthermore uv is an edge in G iff u \neq v and \cap {B $\in E_G | u, v \in B$ } = {u,v}. That is $G_{I_{E_G}} = G$.

Starting with a maximal Helly copair hypergraph (V,E) then $G_E = G_{IE}$ is a median graph with vertex-set V. Moreover E consists of I_{GE} -convex subsets of V. But also E_{GE} is a Helly copair hypergraph consisting of I_{GE} convex subsets of V. Since both E and E_{GF} are maximal, we have that $E = E_{GF}$.

The preceding observations imply that a median graph G, with vertexset V, admits only one cutset colouring which induces a maximal Helly copair hypergraph. Let us call the copairs of V induced by this cutset colouring of G the *canonical copairs* of G. (In fact it can be proved that up to the labelling of the colours a median graph admits exactly one cutset colouring of its edges, cf. [6].)

Recapitulating we have proved:

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<u>PROPOSITION 12</u>. The hypergraph (V,E) is a maximal Helly copair hypergraph iff E consists of the canonical copairs of a median graph with vertex-set V.

3.4 <u>CONCLUDING REMARKS</u>. Let G be a connected graph with n vertices, which admits a cutset colouring. Since each cutset contains edges of a spanning tree, the number of colours in the cutset colouring is at most n - 1.

<u>LEMMA 13</u>. Let G be a connected graph with n vertices admitting a cutset colouring. Then the number of colours in the cutset colouring is n - 1 iff G is a tree.

<u>PROOF.</u> The if part of this lemma is trivial. To prove the only if part let T be a spanning tree of G. Then T has n - 1 edges, so the edges of T all have different colours. Thereby every edge of T determines exactly one cutset of the colouring. Assume that there is an edge joining u and v in G, which is not in T. The u,v-path in T must contain at least two edges, say f_1, f_2, \ldots But then the edge uv is in the cutset determined by f_1 and in the cutset determined by f_2 , which is a contradiction.

The term *maximum* will be used in the sense of: with a maximal number of edges.

<u>PROPOSITION 14</u>. The hypergraph (V, E) is a maximum Helly copair hypergraph iff E consists of the canonical copairs of a tree with vertex-set V.

COROLLARY 15. Let (V, E) be a Helly copair hypergraph. Then

$$|E| \leq 2(|V| - 1).$$

COROLLARY 16. (E.C. MILNER, cf. [4]). Let (V,E) be a Helly hypergraph. Then

$$\begin{vmatrix} \mathbf{V} &- \mathbf{1} \\ |\mathbf{E}| \leq 2 + |\mathbf{V}| - \mathbf{1}. \end{vmatrix}$$

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