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# RAPID APPROXIMATION OF BILINEAR FORMS INVOLVING MATRIX FUNCTIONS THROUGH ASYMPTOTIC ANALYSIS OF

## GAUSSIAN NODE PLACEMENT

by

Elisabeth Marie Palchak

A Thesis Submitted to the Graduate School of The University of Southern Mississippi in Partial Fulfillment of the Requirements for the Degree of Master of Science



Approved:

Dean of the Graduate School

### ABSTRACT

## RAPID APPROXIMATION OF BILINEAR FORMS INVOLVING MATRIX FUNCTIONS THROUGH ASYMPTOTIC ANALYSIS OF GAUSSIAN NODE PLACEMENT

#### by Elisabeth Marie Palchak

#### May 2012

Technological advancements have allowed computing power to generate high resolution models. As a result, greater stiffness has been introduced into systems of ordinary differential equations (ODEs) that arise from spatial discretization of partial differential equations (PDEs). The components of the solutions to these systems are coupled and changing at widely varying rates, which present problems for time-stepping methods. Krylov Subspace Spectral methods, developed by Dr. James Lambers, bridge the gap between explicit and implicit methods for stiff problems by computing each Fouier coefficient from an individualized approximation of the solution operator. KSS methods demonstrate a high order of accuracy, but their efficiency needs to be improved. We will carry out an asymptotic study to determine how these approximations behave at high frequencies to develop a formula to reduce the computation of each node while still achieving a high level of accuracy. Our numerical results will reveal that our method does prove to increase the efficiency as well as the accuracy of KSS methods.

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## TABLE OF CONTENTS

ABSTRACT	ii
ACKNOWLEDGMENTS	iii
LIST OF ILLUSTRATIONS	v
LIST OF TABLES	vi
NOTATION AND GLOSSARY	vii
INTRODUCTION	1
2 BACKGROUND	3
2.1 Polynomial Approximation of the Solution Operator	3
2.2 Matrices	4
2.3 Block Gaussian Quadrature	5
2.4 Block KSS Methods	6
3 ASYMPTOTIC BLOCK LANCZOS ITERATION	8
3.1 Constant Leading Coefficient	8
3.2 Variable Leading Coefficient	10
4 NUMERICAL RESULTS	16
4.1 One-Dimensional Parabolic PDE	16
4.2Two-Dimensional Parabolic PDE	18
5 CONCLUSION	22
BIBLIOGRAPHY	23

## LIST OF ILLUSTRATIONS

## Figure

3.1	Quadrature Nodes and Block Gaussian Nodes		•	•	•		•		•		•	•		9
3.2	Variable Leading Coefficient Fails		•	•	•	•			•	•			•	10
3.3	Well-Approximated Nodes for Variable Coefficient Case		•	•	•	•	•	•	٠	•	•	•		15
4.1	4-Node Comparison on 128 Point Grid				•		•					8		17
4.2	4-Node Comparison on 256 Point Grid				•	•		•			•		•	18
4.3	6-Node Comparison on 128 Point Grid				•	•	•						•	19
4.4	6-Node Comparison on 256 Point Grid						•		•					20
4.5	2D, 4-Node Comparison on 16 Points per Dimension													21
4.6	2D, 4-Node Comparison on 32 Points per Dimension	×			•									21

## LIST OF TABLES

## Table

4.1	Relative Error Estimates of 4-node Comparisons	17
4.2	Relative Error Estimates of 6-node Comparisons	18
4.3	Relative Error Estimates of 2D Comparisons	20

## NOTATION AND GLOSSARY

#### **General Usage and Terminology**

The notation used in this text represents fairly standard mathematical and computational usage. In many cases these fields tend to use different preferred notation to indicate the same concept, and these have been reconciled to the extent possible, given the interdisciplinary nature of the material. In particular, the notation for partial derivatives varies extensively, and the notation used is chosen for stylistic convenience based on the application. While it would be convenient to utilize a standard nomenclature for this important symbol, the many alternatives currently in the published literature will continue to be utilized.

The blackboard fonts are used to denote standard sets of numbers:  $\mathbb{R}$  for the field of real numbers,  $\mathbb{C}$  for the complex field. The capital letters,  $A, B, \cdots$  are used to denote matrices. Functions which are denoted in boldface type typically represent vector valued functions, and real valued functions usually are set in lower case roman or greek letters. Lower case letters such as *i*, *j*, *k*, *l*, *m*, *n* and sometimes *p* and *d* are used to denote indices.

Vectors are typset in square brackets, e.g.,  $[\cdot]$ , and matrices are typeset in parenthesese, e.g.,  $(\cdot)$ . In general the norms are typeset using double pairs of lines, e.g.,  $||\cdot||$ , and the abolute value of numbers is denoted using a single pairs of lines, e.g.,  $|\cdot|$ . Single pairs of lines around matrices indicates the determinant of the matrix. Denote  $\langle \cdot, \cdot \rangle$  as the standard inner product on  $[0, 2\pi]$ . Define  $[f(x) \ g(x)]$  as the limit of a block on *n* rows and 2 columns of vectors of values of f(x) and g(x) as  $n \to \infty$ .

### **Chapter 1**

### INTRODUCTION

In recent years, technology has advanced dramatically. These advancements in technology have allowed computing power to generate high resolution models, but with new advancements comes new challenges. The components of the solution of the ordinary differential equations (ODEs) that arise from spatial discretization of partial differential equations (PDEs) are coupled and changing at widely varying rates. Therein lies the challenge. This coupling has increased the stiffness [2] of the ODEs which presents problems for time-stepping methods.

We first discuss the Hockbruck-Lubich method [16, 17] which is a representative approach to using a polynomial approximation to compute  $e^{-At}\mathbf{v}$ . Since solutions to stiff systems have components that are changing at widely varying rates, this method is not efficient because the time step needs to be chosen very small to compute each component to achieve an acceptable order of accuracy. To handle these components differently, we introduce the next approach which applies Krylov Subspace Spectral (KSS) methods, developed by Dr. James Lambers [20]. KSS methods compute each Fourier coefficient from an individualized approximation of the solution operator. As a result, these KSS methods demonstrate a high order of accuracy, but are lacking in efficiency.

To improve these KSS methods, we will carry out an asymptotic study to determine how these approximations behave at high frequencies. In our anaylsis, we want to recognize any patterns to develop a formula to approximate the largest node to assign as a sharp lower bound to alleviate the expense of computing each node at a high level of accuracy.

In Chapter 2 we discuss the background material. In Section 1 of Chapter 2 we discuss the Hockbruck-Lubich method for solving the problem. In Section 2, we describe the background material in detail as it relates to matrices, quadrature, and partial differential equations. In Section 3, we discuss the block approach, and in Section 4 we introduce block Krylov Subspace Spectral methods. In Chapter 3, we discuss the solution to the problem through the asymptotic block Lanczos iteration. Section 1 of Chapter 3 consider the case when we have a constant leading coefficient, and Section 2 considers the case when we have a variable leading coefficient. Chapter 4 presents our numerical results. In Section 1 of Chapter 4, we examine a one-dimensional parabolic partial differential equation. In Section

2, we examine a two-dimensional example of a parabolic partial differential equation. In Chapter 5, we state our conclusions based on these results.

## Chapter 2 BACKGROUND

### 2.1 Polynomial Approximation of the Solution Operator

For simplicity, we consider the one-dimensional case which can be generalized to higher dimensions [13, 15]. Consider the parabolic PDE

$$u_t + Lu = 0, \quad t > 0,$$
 (2.1)

with  $\exp[-Lt]$  as the solution operator on  $(0, 2\pi)$ , where L is a second order, self-adjoint, positive definite differential operator. The first approach to approximate the solution, developed by Hochbruck and Lubich [16, 17], follows. Consider

$$\mathbf{u}'(t) + A\mathbf{u} = 0, \quad \mathbf{u}(t_0) = \mathbf{u}_0, \tag{2.2}$$

where A is an  $N \times N$  matrix,  $\mathbf{u}(t)$  and  $\mathbf{u}_0$  are N-vectors. This system is a result of the spatial discretization of the previous PDE (2.1). We find the exact solution of the ODE to be  $\mathbf{u}(t) = e^{-At}\mathbf{u}_0$ , which results in an approximate solution of the PDE.

We use a polynomial of A produced by the Lanczos iteration as an approximation of exponential. For example, consider  $\mathbf{w} = e^{-At}\mathbf{v}$  for a given symmetric matrix A and vector  $\mathbf{v}$ . If we apply the Lanczos iteration to the matrix A with initial vector  $\mathbf{v}$ , j times, where  $j \ll N$ , at the end of the jth iteration, the Lanczos iteration produces an  $N \times j$  orthogonal matrix  $X_j$  and a  $j \times j$  tridiagonal matrix  $T_j$  such that  $X_j^T A X_j = T_j$ . Then we can compute the approximation

$$\mathbf{w}_j = \|\mathbf{v}\|_2 X_j e^{-T_j t} \mathbf{e}_1, \tag{2.3}$$

where each column  $\mathbf{x}_k$ , where k = 1, ..., j, of  $X_j$  is  $\mathbf{x}_k = p_{k-1}(A)\mathbf{v}$ . The polynomial  $p_n(A)$  is of degree *n* in *A*, the polynomials  $p_j$  are orthonormal, and  $\mathbf{w}_j$  is a product of a polynomial in *A* of degree j - 1 and  $\mathbf{v}$  [21, 17].

It is not possible to approximate exponential functions on a large interval with a polynomial without using a large number of terms. Similarly, it is not possible to approximate a matrix exponential with widely varying eigenvalues without many Lanczos iterations, which is the case if *A* arises from a stiff system. So if the eigenvalues are not clustered, this approach is not practical. A new approach is needed that approximates the high frequency components and low frequency components separately. Specifically, we need to apply a different polynomial approximation of the exponential for each component.

#### 2.2 Matrices

A componentwise technique is used in the next approach. We execute this componentwise approach by using an approximation of the solution operator that is ideal for each Fourier coefficient of the solution. Given the solution  $\tilde{u}(x,t_n)$  at time  $t_n$ , the solution at time  $t_{n+1}$  is

$$\hat{u}(\omega, t_{n+1}) = \left\langle \frac{1}{\sqrt{2\pi}} e^{i\omega x}, \exp[-L\Delta t] \tilde{u}(x, t_n) \right\rangle.$$
(2.4)

The result from discretizing the (2.4) is

$$\mathbf{u}^T f(A) \mathbf{v},\tag{2.5}$$

where  $\mathbf{u} = \frac{1}{\sqrt{2\pi}} e^{i\omega x}$  and  $\mathbf{v} = \tilde{u}(x,t_n)$  are *N*-vectors,  $A = L_N$  is an  $N \times N$  symmetric positive definite matrix that comes from discretizing the operator *L*, and  $f(\lambda) = \exp(-\lambda \Delta t)$ . Methods for computing the bilinear form in (2.4) are discussed by Golub and Meurant in [10]. As given by properties of symmetric positive definite matrices, the matrix *A* has real eigenvalues

$$b = \lambda_1 \ge \lambda_2 \ge \dots \ge \lambda_N = a > 0, \tag{2.6}$$

and corresponding orthonormal eigenvectors  $\mathbf{q}_j$ , where j = 1, ..., N. Thus we can rewrite  $\mathbf{u}^T f(A)\mathbf{v}$  in terms of its spectral decomposition,

$$\mathbf{u}^T f(A) \mathbf{v} = \sum_{j=1}^N f(\lambda_j) \mathbf{u}^T \mathbf{q}_j \mathbf{q}_j^T \mathbf{v}.$$
 (2.7)

The matrix function (2.7) can also be regarded as the Riemann-Stieltjes integral

$$\mathbf{u}^{T} f(A) \mathbf{v} = I[f] = \int_{a}^{b} f(\lambda) d\alpha(\lambda)$$
(2.8)

where

$$\boldsymbol{\alpha}(\boldsymbol{\lambda}) = \begin{cases} 0, & \text{if } \boldsymbol{\lambda} < a \\ \sum_{j=i}^{N} \alpha_{j} \beta_{j}, & \text{if } \lambda_{i} \leq \boldsymbol{\lambda} < \lambda_{i-1} \\ \sum_{j=1}^{N} \alpha_{j} \beta_{j}, & \text{if } b \leq \boldsymbol{\lambda} \end{cases}, \quad \boldsymbol{\alpha}_{j} = \mathbf{u}^{T} \mathbf{q}_{j}, \quad \boldsymbol{\beta}_{j} = \mathbf{q}_{j}^{T} \mathbf{v}.$$
(2.9)

We can approximate the integral I[f] by Gaussian, Gauss-Radau, or Gauss-Lobatto Quadrature rules [3, 8, 9, 10], but the scope of my research only considers Gaussian. Approximating I[f] yields

$$I[f] = \sum_{j=1}^{K} w_j f(\lambda_j) + R[f]$$
(2.10)

where we use the Lanczos algorithm to acquire the nodes  $\lambda_j$ , j = 1, ..., K and weights  $w_j$ , j = 1, ..., K [4, 6, 7, 12]. In the case where  $\mathbf{u} = \mathbf{v}$ , we apply the Lanczos algorithm [32]:

$$\beta_0 = 0, \mathbf{p} = 0, \mathbf{u} = given, \mathbf{q} = \mathbf{u}/\|\mathbf{u}\|_2$$

for 
$$n = 1, 2, ..., K$$
  
 $\mathbf{v} = A\mathbf{q}$   
 $\alpha_n = \mathbf{q}^T \mathbf{v}$   
if  $n < K$   
 $\mathbf{v} = \mathbf{v} - \beta_{n-1}\mathbf{p} - \alpha_n \mathbf{q}$   
 $\beta_n = \|\mathbf{v}\|_2$   
 $\mathbf{p} = \mathbf{q}$   
 $\mathbf{q} = \mathbf{v}/\beta_n$   
end

#### end

The  $\alpha_j$ s and  $\beta_j$ s produced by the Lanczos algorithm form a tridiagonal matrix  $\mathcal{T}_K$ , where the diagonal elements are the  $\alpha$ s and the super and sub diagonal elements are the  $\beta$ s. The resulting tridiagonal matrix produces the nodes and the weights [12] in our approximation. The nodes  $\lambda_j$  are the eigenvalues of

$$\mathfrak{T}_{K} = \begin{bmatrix}
\alpha_{1} & \beta_{1} & & & \\
\beta_{1} & \alpha_{2} & \beta_{2} & & \\
& \ddots & \ddots & \ddots & \\
& & \beta_{K-2} & \alpha_{K-1} & \beta_{K-1} \\
& & & & \beta_{K-1} & \alpha_{K}
\end{bmatrix},$$
(2.11)

and the weights are obtained by squaring the first components of the eigenvectors.

### 2.3 Block Gaussian Quadrature

Now we consider the case when  $\mathbf{u} \neq \mathbf{v}$ . As a consequence, the weights could be negative and destabilize the quadrature rule [1, 21]. Thus, we consider the block approach

$$\begin{bmatrix} \mathbf{u} & \mathbf{v} \end{bmatrix}^T f(A) \begin{bmatrix} \mathbf{u} & \mathbf{v} \end{bmatrix}.$$
(2.12)

The block matrix function (2.12) can be regarded as the Riemann-Stieltjes integral

$$\int_{a}^{b} f(\lambda) d\mu(\lambda) = \begin{bmatrix} \mathbf{u}^{T} f(A) \mathbf{u} & \mathbf{u}^{T} f(A) \mathbf{v} \\ \mathbf{v}^{T} f(A) \mathbf{u} & \mathbf{v}^{T} f(A) \mathbf{v} \end{bmatrix}$$
(2.13)

where  $\mu(\lambda)$  is a 2 × 2 matrix, each entry of which is a measure of the form  $\alpha(\lambda)$ . Using the most general quadrature formula with *K* matrix nodes [10, 21] we get

$$\int_{a}^{b} f(\lambda) d\mu(\lambda) = \sum_{j=1}^{K} W_{j} f(T_{j}) W_{j} + error, \qquad (2.14)$$

where  $T_j$  and  $W_j$  are 2 × 2 matrices. If we diagonalize each  $T_j$ , the resulting quadrature formula is given by

$$\int_{a}^{b} f(\lambda) d\mu(\lambda) = \sum_{j=1}^{2K} f(\lambda_j) \mathbf{v}_j \mathbf{v}_j^T + error, \qquad (2.15)$$

where  $\lambda_j$  is a scalar and  $\mathbf{v}_j$  is a 2-vector. The nodes and weights for our quadrature formula are obtained by applying the block Lanczos algorithm:

$$X_0 = 0, R_0 = \begin{bmatrix} \mathbf{u} & \mathbf{v} \end{bmatrix}, R_0 = X_1 B_0$$

for 
$$n = 1, 2, ..., K$$
  
 $V = AX_n$   
 $M_n = X_n^T V$   
if  $n < K$   
 $R_n = V - X_{n-1}B_{n-1}^T - X_n M_n$   
 $R_n = X_{n+1}B_n$   
end

end

The block Lanczos algorithm produces  $2 \times 2$  matrices  $M_j$  and  $B_j$ , as stated in [11], that form a block tridiagonal matrix  $T_K$ , where each  $B_j$  is upper triangular. The resulting block tridiagonal matrix produces the nodes and the weights. The nodes  $\lambda_j$  are the eigenvalues of

$$\mathfrak{T}_{K} = \begin{bmatrix}
M_{1} & B_{1}^{T} & & & \\
B_{1} & M_{2} & B_{2}^{T} & & & \\
& \ddots & \ddots & \ddots & & \\
& & B_{K-2} & M_{K-1} & B_{K-1}^{T} \\
& & & B_{K-1} & M_{K}
\end{bmatrix},$$
(2.16)

and the weights are the  $2 \times 2$  matrices obtained by taking the outer product of the first two components of each eigenvector with itself.

#### 2.4 Block KSS Methods

The general method for Block KSS [19, 24] begins by defining

$$R_0 = \begin{bmatrix} \hat{\mathbf{e}}_{\boldsymbol{\omega}} & \mathbf{u}(t_n) \end{bmatrix}, \qquad (2.17)$$

where  $\hat{\mathbf{e}}_{\omega}$  is a discretization of  $\frac{1}{\sqrt{2\pi}}e^{i\omega x}$  and

$$\mathbf{u}_{\omega}^{n} = \mathbf{u}^{n} - \hat{\mathbf{e}}_{\omega} \hat{\mathbf{e}}_{\omega}^{H} \mathbf{u}^{n}.$$
(2.18)

Next we compute the *QR* factorization of  $R_{j-1}$  for j = 1, such that

$$R_0 = X_1(\omega)B_0(\omega) \tag{2.19}$$

which outputs

$$X_1(\boldsymbol{\omega}) = \begin{bmatrix} \hat{\mathbf{e}}_{\boldsymbol{\omega}} & \frac{\mathbf{u}_{\boldsymbol{\omega}}^n}{\|\mathbf{u}_{\boldsymbol{\omega}}^n\|_2} \end{bmatrix}$$
(2.20)

and

$$B_0(\boldsymbol{\omega}) = \begin{bmatrix} 1 & \hat{\mathbf{e}}_{\boldsymbol{\omega}}^H \mathbf{u}^n \\ 0 & \|\mathbf{u}_{\boldsymbol{\omega}}^n\|_2 \end{bmatrix}.$$
 (2.21)

Then we apply the block Lanczos algorithm to the operator  $L_N$ , which comes from the discretization of L, with initial block  $X_1(\omega)$ . The result produces a block tridiagonal matrix  $T_K$  like (2.16), where every entry of  $T_K$  is a function of  $\omega$ . Then, at time  $t_{n+1}$ , each Fourier coefficient of the solution is

$$[\hat{\mathbf{u}}^{n+1}]_{\omega} = [B_0^H E_{12}^H exp[\mathcal{T}_K(\omega)\Delta t] E_{12} B_0]_{12}, \quad E_{12} = \begin{bmatrix} \mathbf{e}_1 & \mathbf{e}_2 \end{bmatrix}.$$
(2.22)

For the parabolic problem we will examine in Chapter 4, we will show that block KSS methods have a high-order of accuracy  $O(\Delta t^{2K-1})$  [20] for *K* block Gaussian nodes, whereas the Hochbruck and Lubich method only has  $O(\Delta t^{K-1})$  [17].

Compared to the Hochbruck and Lubich method, it would seem that KSS methods have a high amount of computation expense. However, this is not actually the case. We compute an approximation of the solution operator for each Fourier coefficient, and all of the subspaces generated are closely related by the wave number  $\omega$  [21]. Although KSS methods take less time computationally in comparison to the Hochbruck and Lubich method, we still do not achieve a desired level of efficiency. In Chapter 3, we will discuss how we address this issue.

### **Chapter 3**

## ASYMPTOTIC BLOCK LANCZOS ITERATION

#### 3.1 Constant Leading Coefficient

Previously in our discussion of block KSS methods, we stated a Jacobi matrix is computed for each Fourier coefficient by using an approximation that is ideal for that component [16, 17, 18, 31, 33]. Then each component is approximated using a different polynomial of *A*. Thus, the computed solution including all the components is defined by

$$\mathbf{u}^{n+1} = \tilde{f}(L_N; \Delta t) \mathbf{u}^n = \sum_{j=0}^{2K} D_j(\Delta t) A^j \mathbf{u}^n,$$
(3.1)

where  $\mathbf{u}^n$  is the solution from the previous time step and  $D_j(\Delta t)$  is a diagonal matrix in the Fourier basis.

Block KSS methods are more flexible than time-stepping methods based on polynomial or rational approximations of  $\exp[-At]$  and can be used to improve accuracy and stability for stiff problems. However, compared to other time-stepping methods with the same number of grid points, more computation time is required per time step to compute the nodes and weights. Given the high order of accuracy and stability of KSS methods [14, 19, 20, 22, 23, 24, 25, 26, 27, 28, 29, 30], we wish to improve their efficiency without sacrificing any accuracy. We now examine how we can take advantage of this flexibility.

First we will examine the case where we have a constant leading coefficient. We use (2.20) to form the initial block recursion coefficient

$$M_{1}(\boldsymbol{\omega}) = \begin{bmatrix} \hat{\mathbf{e}}_{\boldsymbol{\omega}}^{H} L_{N} \hat{\mathbf{e}}_{\boldsymbol{\omega}} & \hat{\mathbf{e}}_{\boldsymbol{\omega}}^{H} L_{N} \mathbf{u}_{\boldsymbol{\omega}}^{n} \\ [\mathbf{u}_{\boldsymbol{\omega}}^{n}]^{H} L_{N} \hat{\mathbf{e}}_{\boldsymbol{\omega}} & [\mathbf{u}_{\boldsymbol{\omega}}^{n}]^{H} L_{N} \mathbf{u}_{\boldsymbol{\omega}}^{n} \end{bmatrix}.$$
(3.2)

Suppose L is a second-order differential operator where

$$Lu = -pu_{xx} + q(x)u \tag{3.3}$$

and p is constant. Then

$$M_1(\boldsymbol{\omega}) = \begin{bmatrix} p\boldsymbol{\omega}^2 + \bar{q} & \hat{\mathbf{e}}^H_{\boldsymbol{\omega}} \tilde{\mathbf{q}} \mathbf{u}^n_{\boldsymbol{\omega}} \\ [\mathbf{u}^n_{\boldsymbol{\omega}}]^H \tilde{\mathbf{q}} \hat{\mathbf{e}}_{\boldsymbol{\omega}} & [\mathbf{u}^n_{\boldsymbol{\omega}}]^H L_N \mathbf{u}^n_{\boldsymbol{\omega}} \end{bmatrix}, \qquad (3.4)$$

where  $\overline{q}$  is the mean of a function q(x) defined on  $[0, 2\pi]$  and  $\tilde{q}(x) = q(x) - \overline{q}$ . Denote  $\tilde{\mathbf{q}}$  by the vector with components  $[\tilde{\mathbf{q}}]_j = \tilde{q}(x_j)$ .

As  $|\omega|$  increases, the Fourier coefficients of a function go to zero and  $\mathbf{u}_{\omega}^{n} \rightarrow \mathbf{u}^{n}$ . So the off-diagonal entries are virtually the Fourier coefficient of  $\tilde{q}u(\cdot, t_{n})$ , corresponding to  $\omega$ , which decay to 0 as  $|\omega| \rightarrow \infty$ . So,

$$M_1(\boldsymbol{\omega}) \sim \begin{bmatrix} p \boldsymbol{\omega}^2 + \bar{q} & 0\\ 0 & \frac{[\mathbf{u}^n]^H L_N \mathbf{u}^n}{[\mathbf{u}^n]^H \mathbf{u}^n} \end{bmatrix}, \qquad (3.5)$$

and the non-diagonal entries become negligible. Thus as  $\omega$  increases,  $B_1(\omega)$  becomes small so the overall matrix  $\mathcal{T}_K(\omega)$  approximately decouples. Thus, for higher frequencies, the diagonal entries of  $M_1(\omega)$  are approximately two eigenvalues of  $\mathcal{T}_K(\omega)$ .



Figure 3.1: Quadrature Nodes and Block Gaussian Nodes

Figure 3.1 above shows the prescribed quadrature nodes by a 4-node KSS method with the blue circles. The red crosses show a 2-node block KSS method with 4 block Gaussian nodes. The previous approximation is a good approximation of the Gaussian nodes, but the work that goes into computing the actual Gaussian nodes is excessive. It has been observed in numerical experiments for the 2-node Block KSS method, that half of the four scalar nodes  $\lambda_i$  are clustered around the values of the diagonal entries of  $M_1(\omega)$ . So we prescribe the smallest and largest nodes

$$\lambda_1(\omega) = \frac{[\mathbf{u}^n]^H L_N \mathbf{u}^n}{[\mathbf{u}^n]^H \mathbf{u}^n}, \quad \lambda_{2K}(\omega) = \frac{\hat{\mathbf{e}}_{\omega}^H L_N \hat{\mathbf{e}}_{\omega}}{\hat{\mathbf{e}}_{\omega}^H \hat{\mathbf{e}}_{\omega}}$$
(3.6)

as seen in Figure 3.1. The rest of the nodes are equally spaced between  $\lambda_1(\omega)$  and  $\lambda_{2K}(\omega)$  [21].

### 3.2 Variable Leading Coefficient

The previous approach does not work when the operator contains a variable leading coefficient as demonstrated Figure 3.2. Different from the case in (3.3) where the operator contains a constant leading coefficient, we define the operator

$$Lu = -(p(x)u_x)_x + q(x)u.$$
 (3.7)



Figure 3.2: Variable Leading Coefficient Fails

We now propose to assign the largest node for each component to serve as a sharper lower bound, as opposed to the lower bound we could obtain from the (1,1) entry of (3.5). To achieve our approximation of the largest node, we examine the columns of the  $R_j$  blocks generated by block Lanczos, only considering the highest powers of  $\omega$ , to develop a formula that maximizes the Rayleigh quotient. From block Lanczos, we compute  $M_2(\omega)$  in the following manner. We calculate

$$R_1 = LX_1 - X_1 M_1 \tag{3.8}$$

Thus,

$$R_{1} = \left[ L(\frac{1}{\sqrt{2\pi}}e^{i\omega x}) - [\frac{1}{\sqrt{2\pi}}e^{i\omega x}(M_{1})_{11} + u_{\omega}(x)(M_{1})_{21}], \\ L(u_{\omega}(x)) - [\frac{1}{\sqrt{2\pi}}e^{i\omega x}(M_{1})_{12} + u_{\omega}(x)(M_{1})_{22}] \right] \\ = \left[ L(\frac{1}{\sqrt{2\pi}}e^{i\omega x}) - [\frac{1}{\sqrt{2\pi}}e^{i\omega x}(\overline{p}\omega^{2} + \overline{q}) + u_{\omega}(x)[\frac{\omega^{2}}{\sqrt{2\pi}}\hat{p}u - \frac{i\omega}{\sqrt{2\pi}}\hat{p}^{2}u + \frac{1}{\sqrt{2\pi}}\hat{q}u]], \\ L(u_{\omega}(x)) - [\frac{1}{\sqrt{2\pi}}e^{i\omega x}[\frac{\omega^{2}}{\sqrt{2\pi}}\hat{p}u - \frac{i\omega}{\sqrt{2\pi}}\hat{p}^{2}u + \frac{1}{\sqrt{2\pi}}\hat{q}u] + \\ u_{\omega}(x)(-puu'' - p'uu' + qu^{2}] \right] \\ = \left[ \frac{\omega^{2}}{\sqrt{2\pi}}pe^{i\omega x} - \frac{i\omega}{\sqrt{2\pi}}p'e^{i\omega x} + \frac{1}{\sqrt{2\pi}}qe^{i\omega x} - \frac{\omega^{2}}{\sqrt{2\pi}}\overline{p}e^{i\omega x} - \frac{1}{\sqrt{2\pi}}\overline{q}e^{i\omega x} - \frac{\omega^{2}}{\sqrt{2\pi}}\hat{p}uu + \\ \frac{\omega}{\sqrt{2\pi}}\hat{p}^{2}uu - \frac{1}{\sqrt{2\pi}}\hat{q}uu, -pu'' - p'u' + qu - \frac{\omega^{2}}{\sqrt{2\pi}}\hat{p}ue^{i\omega x} + \frac{i\omega}{\sqrt{2\pi}}\hat{p}^{2}ue^{i\omega x} - \\ \frac{1}{\sqrt{2\pi}}\hat{q}ue^{i\omega x} + pu^{2}u'' + p'u^{2}u' + qu^{3} \right].$$

$$(3.9)$$

After only considering the highest powers of  $\omega$  in each column, we determine

$$R_1 = \left[ \begin{array}{cc} \frac{\omega^2}{\sqrt{2\pi}} \tilde{p} e^{i\omega x} & \frac{-\omega^2}{2\pi} \hat{p} u(x) e^{i\omega x} \end{array} \right] + \text{lower order terms.}$$
(3.10)

Neglecting the constants we conclude

$$R_1 = \begin{bmatrix} e^{i\omega x} & \tilde{p}(x)e^{i\omega x} \end{bmatrix}.$$
(3.11)

From  $R_1$  we determine

$$X_2 = \left[\begin{array}{cc} \frac{1}{\sqrt{2\pi}} e^{i\omega x} & \tilde{p}(x) e^{i\omega x} \end{array}\right], \tag{3.12}$$

which leads to

$$M_2(\omega) = X_2^H L X_2. (3.13)$$

We calculate

$$M_{2}(\omega) = \begin{bmatrix} \frac{1}{\sqrt{2\pi}} e^{i\omega x} & \tilde{p}(x)e^{i\omega x} \end{bmatrix}^{H} L \begin{bmatrix} \frac{1}{\sqrt{2\pi}} e^{i\omega x} & \tilde{p}(x)e^{i\omega x} \end{bmatrix}$$
$$= \begin{bmatrix} \left\langle \frac{1}{\sqrt{2\pi}} e^{i\omega x}, L\frac{1}{\sqrt{2\pi}} e^{i\omega x} \right\rangle & \left\langle \frac{1}{\sqrt{2\pi}} e^{i\omega x}, L\tilde{p}(x)e^{i\omega x} \right\rangle \\ \left\langle \tilde{p}(x)e^{i\omega x}, L\frac{1}{\sqrt{2\pi}} e^{i\omega x} \right\rangle & \left\langle \tilde{p}(x)e^{i\omega x}, L\tilde{p}(x)e^{i\omega x} \right\rangle \end{bmatrix}.$$
(3.14)

After computing the inner products of the entries we get

$$M_{2}(\omega) = \begin{bmatrix} \frac{\overline{p}\omega^{2} + \overline{q}}{\frac{\omega^{2}}{\sqrt{2\pi}} \|\tilde{p}\|_{2} + \frac{1}{\sqrt{2\pi}} \frac{\langle \tilde{q}, \tilde{p} \rangle}{\|\tilde{p}\|_{2}}} & \frac{\omega^{2}}{\sqrt{2\pi}} \|\tilde{p}\|_{2}^{2} + \frac{\|\tilde{p}'\|_{p}^{2}}{\|\tilde{p}\|^{2}} + \frac{\|\tilde{p}'\|_{q}^{2}}{\|\tilde{p}\|^{2}} \end{bmatrix}.$$
(3.15)

Using  $M_2(\omega)$ , we compute

$$R_2 = LX_2 - X_2M_2 - X_1B_1^T. (3.16)$$

Since we determined  $B_1^T \sim 0$  is negligible, we only need to compute

$$R_2 = LX_2 - X_2M_2 \tag{3.17}$$

where

$$X_2 = \left[\begin{array}{cc} \frac{1}{\sqrt{2\pi}} e^{i\omega x} & \tilde{p}(x) e^{i\omega x} \end{array}\right].$$
(3.18)

Thus,

$$R_{2} = \left[ L(\frac{1}{\sqrt{2\pi}}e^{i\omega x}) - \left[\frac{1}{\sqrt{2\pi}}e^{i\omega x}(M_{2})_{11} + \tilde{p}(x)e^{i\omega x}(M_{2})_{21}\right], \\ L(\tilde{p}(x)e^{i\omega x}) - \left[\frac{1}{\sqrt{2\pi}}e^{i\omega x}(M_{2})_{12} + \tilde{p}(x)e^{i\omega x}(M_{2})_{22}\right] \right] \\ = \left[ L(\frac{1}{\sqrt{2\pi}}e^{i\omega x}) - \left[\frac{1}{\sqrt{2\pi}}e^{i\omega x}(\bar{p}\omega^{2} + \bar{q}) + \tilde{p}(x)e^{i\omega x}(\frac{\omega^{2}}{\sqrt{2\pi}}\|\tilde{p}\|_{2} + \frac{1}{\sqrt{2\pi}}\frac{\langle \bar{q}, \bar{p} \rangle}{\|\bar{p}\|_{2}})\right], \\ L(\tilde{p}(x)e^{i\omega x}) - \left[\frac{1}{\sqrt{2\pi}}e^{i\omega x}(\frac{\omega^{2}}{\sqrt{2\pi}}\|\tilde{p}\|_{2} + \frac{1}{\sqrt{2\pi}}\frac{\langle \bar{q}, \bar{p} \rangle}{\|\bar{p}\|_{2}})\right] \right] \\ = \left[ \frac{\omega^{2}}{\sqrt{2\pi}}pe^{i\omega x} - \frac{i\omega}{\sqrt{2\pi}}p'e^{i\omega x} + \frac{1}{\sqrt{2\pi}}qe^{i\omega x} - \left[\frac{1}{\sqrt{2\pi}}e^{i\omega x}(\bar{p}\omega^{2} + \bar{q}) + \tilde{p}(x)e^{i\omega x}(\frac{\omega^{2}}{\sqrt{2\pi}}\|\tilde{p}\|_{2} + \frac{1}{\sqrt{2\pi}}\frac{\langle \bar{q}, \bar{p} \rangle}{\|\bar{p}\|_{2}})\right], p_{0}\omega^{2}\tilde{p}e^{i\omega x} - p_{0}i\omega\tilde{p}'e^{i\omega x} + \omega^{2}\tilde{p}^{2}e^{i\omega x} - 2i\omega\tilde{p}\tilde{p}'e^{i\omega x} - p_{0}i\omega\tilde{p}'e^{i\omega x} - p_{0}\tilde{p}''e^{i\omega x} - \tilde{p}\tilde{p}''e^{i\omega x} - i\omega\tilde{p}\tilde{p}''e^{i\omega x} + \frac{1}{\sqrt{2\pi}}e^{i\omega x}(\frac{\omega^{2}}{\sqrt{2\pi}}\|\tilde{p}\|_{2}) + \tilde{p}(x)e^{i\omega x}(\omega^{2}\frac{\|\tilde{p}\|_{2}^{2}}{\|\tilde{p}\|_{2}^{2}} + \frac{\|\tilde{p}'\|_{2}^{2}}{\|\tilde{p}\|_{2}^{2}})\right] \right].$$
(3.19)

After only considering the highest powers of  $\omega$  in each column, we determine

$$R_2 = \left[ \tilde{p}\omega^2 \frac{1}{\sqrt{2\pi}} e^{i\omega x} + \tilde{p}(x)e^{i\omega x} \frac{\omega^2}{\sqrt{2\pi}} \|\tilde{p}\|_2 \quad \tilde{p}^2 \omega^2 e^{i\omega x} - \frac{\omega^2}{2\pi} e^{i\omega x} \|\tilde{p}\|_2 \right].$$
(3.20)

Neglecting the constants, we conclude

$$R_2 = \left[ \begin{array}{cc} \tilde{p}(x)e^{i\omega x} & \tilde{p}(x)^2e^{i\omega x} - e^{i\omega x} \end{array} \right].$$
(3.21)

Using  $R_2$ , we compute  $X_3$ 

$$X_3 = \begin{bmatrix} \tilde{p}(x)e^{i\omega x} & \tilde{p}(x)^2e^{i\omega x} - e^{i\omega x} \end{bmatrix}$$
(3.22)

which yields

$$M_3(\omega) = X_3^H L X_3. \tag{3.23}$$

So

$$M_{3}(\omega) = \begin{bmatrix} \tilde{p}(x)e^{i\omega x} & \tilde{p}(x)^{2}e^{i\omega x} - e^{i\omega x} \end{bmatrix}^{H} L \begin{bmatrix} \tilde{p}(x)e^{i\omega x} & \tilde{p}(x)^{2}e^{i\omega x} - e^{i\omega x} \end{bmatrix}.$$
(3.24)

We note that the columns of the  $X_j$  blocks are linear combinations of

$$e^{i\omega x}, \tilde{p}e^{i\omega x}, \tilde{p}^2 e^{i\omega x}, \dots, \tilde{p}^{j-1}e^{i\omega x}$$
(3.25)

+ lower order terms of  $\omega$ . We recall our goal is to estimate the largest eigenvalue of  $T_k$  such that

$$\lambda_1(T_{k,\omega}) = \lambda_1(Q_{k,\omega}^H L Q_{k,\omega}) \tag{3.26}$$

where

$$Q_{k,\omega} = \left[ X_1 X_2 \dots X_k \right]. \tag{3.27}$$

The largest eigenvalue of a symmetric matrix maximizes the Raleigh quotient, where

$$\max_{\mathbf{v}\in\mathrm{range}(\mathcal{Q}_{k,\omega})}\frac{\mathbf{v}^{H}L\mathbf{v}}{\mathbf{v}^{H}\mathbf{v}}.$$
(3.28)

By computing  $R_1$  and  $R_2$ , we observed a nice pattern that will allow us to consider a much smaller space. Since we only consider the high frequency components, the useful part of the range $(Q_{k,\omega})$  is well approximated by the span $\{e^{i\omega x}, \tilde{p}e^{i\omega x}, \tilde{p}^2e^{i\omega x}, \dots, \tilde{p}^{k-1}e^{i\omega x}\}$ . As  $|\omega| \to \infty$ , the largest eigenvalue is well approximated by

$$\max_{\mathbf{v}\in\mathrm{range}(P_{k,\omega})}\frac{\mathbf{v}^{H}L\mathbf{v}}{\mathbf{v}^{H}\mathbf{v}}$$
(3.29)

where

$$P_{k,\omega} = \begin{bmatrix} e^{i\omega x} & \tilde{p}e^{i\omega x} & \cdots & \tilde{p}^{k-1}e^{i\omega x} \end{bmatrix}$$
(3.30)

All the vectors  $\mathbf{v}$  are a linear combination of

$$\mathbf{v} = c_1 \mathbf{p}_1 + c_2 \mathbf{p}_2 + \dots + c_k \mathbf{p}_k \tag{3.31}$$

SO

$$\mathbf{v} = P_{k,\omega} \mathbf{c}.\tag{3.32}$$

Since **v** is a linear combination of the columns of  $P_{k,\omega}$ , we determine

$$\max_{\mathbf{v}\in \mathrm{range}(Q_{k,\omega})} \frac{\mathbf{v}^{H} L \mathbf{v}}{\mathbf{v}^{H} \mathbf{v}} \approx \max_{\mathbf{c}\in\mathbb{C}} \frac{\mathbf{c}^{H}(P_{k,\omega}^{H} L P_{k,\omega}) \mathbf{c}}{\mathbf{c}^{H}(P_{k,\omega}^{H} P_{k,\omega}) \mathbf{c}} = \max_{\mathbf{c}\in\mathbb{C}} \frac{\mathbf{c}^{H} B \mathbf{c}}{\mathbf{c}^{H} F \mathbf{c}} = \lambda_{1}(B,F)$$
(3.33)

where  $\lambda_1$  is the largest generalized eigenvalue [5] of the matrix pair (B, F). For

$$\mathbf{u}_k = \tilde{p}^{k-1} e^{i\omega x} \tag{3.34}$$

and

$$\mathbf{u}_j = \tilde{p}^{j-1} e^{i\omega x} \tag{3.35}$$

we determine formulas for B and F. We compute F in the following manner.

$$F_{kj} = \langle \mathbf{u}_{k}, \mathbf{u}_{j} \rangle$$

$$= \langle \tilde{p}^{k-1} e^{i\omega x}, \tilde{p}^{j-1} e^{i\omega x} \rangle$$

$$= \int_{0}^{2\pi} (\overline{\tilde{p}^{k-1} e^{i\omega x}}) (\tilde{p}^{j-1} e^{i\omega x}) dx$$

$$= \int_{0}^{2\pi} (\tilde{p}^{k-1}) (\tilde{p}^{j-1}) dx$$

$$= \int_{0}^{2\pi} (\tilde{p}^{k+j-2}) dx \qquad (3.36)$$

We compute our formula for *B* similarly

$$\begin{split} B_{kj} &= \langle \mathbf{u}_{k}, L\mathbf{u}_{j} \rangle \\ &= \langle \tilde{p}^{k-1} e^{i\omega x}, L(\tilde{p}^{j-1} e^{i\omega x}) \rangle \\ &= \int_{0}^{2\pi} (\overline{\tilde{p}^{k-1} e^{i\omega x}}) (L(\tilde{p}^{j-1} e^{i\omega x})) dx \\ &= \int_{0}^{2\pi} (\overline{\tilde{p}^{k-1} e^{i\omega x}}) (p_{0} \tilde{p}^{j-1} \omega^{2} e^{i\omega x} + \tilde{p}^{j} \omega^{2} e^{i\omega x} - 2(j-1) p_{0} \tilde{p}^{j-2}(\tilde{p})' i\omega e^{i\omega x} \\ &- 2(j-1) \tilde{p}^{j-1}(\tilde{p})' i\omega e^{i\omega x} - (j-1)(j-2) p_{0} \tilde{p}^{j-3}((\tilde{p})')^{2} e^{i\omega x} - (j-1)(j-2) \tilde{p}^{j-2}(\tilde{p})'' e^{i\omega x} - (j-1)(j-2) \tilde{p}^{j-1}(\tilde{p})' e^{i\omega x} - (j-1) p_{0} \tilde{p}^{j-2}(\tilde{p})'' e^{i\omega x} - (j-1) \tilde{p}^{j-1}(\tilde{p})' e^{i\omega x} - (j-1) \tilde{p}^{j-1}(\tilde{p})' e^{i\omega x} - p' \tilde{p}^{j-1} i\omega e^{i\omega x} - (j-1) p' \tilde{p}^{j-2}(\tilde{p})' e^{i\omega x} + q(\tilde{p}^{j-1} e^{i\omega x})) dx \\ &= \int_{0}^{2\pi} (p_{0} \omega^{2} \tilde{p}^{k+j-2} + \omega^{2} \tilde{p}^{k+j-1} - (j-1)(j-2) p_{0} \tilde{p}^{k+j-4} \tilde{p}'^{2} - (j-1)(j-2) \tilde{p}^{k+j-3} \tilde{p}'^{2} - (j-1) p_{0} \tilde{p}^{k+j-3} \tilde{p}'' - (j-1) \tilde{p}^{k+j-2} \tilde{p}'' - (j-1) \tilde{p}^{k+j-3} \tilde{p}'^{2} + \tilde{p}^{k+j-2} q) dx \end{split}$$

$$(3.37)$$

As shown in Figure 3.3, this method is very accurate for the 4-node case in approximating the external nodes when p is variable. To summarize, in the case where p is variable, we prescribe the largest node accoring to the aforementioned method. When p is constant, we follow the given procedure in (3.5) and (3.6) to use the quadrature rule. Our numerical results are shown in Chapter 4.



Figure 3.3: Well-Approximated Nodes for Variable Coefficient Case

### **Chapter 4**

## NUMERICAL RESULTS

### 4.1 One-Dimensional Parabolic PDE

In this chapter we solve examples of the one-dimensional and two-dimensional heat equation using the Hochbruck-Lubich method of approximating  $e^{-LT}$ , an accelerated KSS method with nodes prescribed in the same way as when p is constant, Dr. Lambers' original KSS method, and our new method. All experiments are conducted on 128 and 256 point grids with 4- or 6- dimensional Krylov subspaces. The exact solutions are unknown, so we compute the error by comparing each solution to the solution at the smallest time step and taking the 2-norm of the relative difference.

To compare our results to previous methods, we solve (2.1), (3.7) using the coefficients

$$p(x) = 1 + \frac{1}{2}\cos x - \frac{1}{4}\sin 2x + \frac{1}{8}\cos 3x$$
  

$$q(x) = 1 + \frac{1}{4}\sin x - \frac{1}{4}\cos 2x + \frac{1}{8}\sin 3x - \frac{1}{8}\cos 4x$$
(4.1)

with initial conditions

$$u(x,0) = 1 + \frac{3}{10}\cos x - \frac{3}{20}\sin 2x + \frac{3}{40}\cos 3x, \quad 0 < x < 2\pi.$$
(4.2)

First we compare our methods using 4-node KSS methods. The results are given in Figure 4.1, Figure 4.2, and Table 4.1. Figure 4.1 gives the relative error for the solution of (2.1) using four different methods. The black curve demonstrates our improved KSS method. The blue curve demonstrates Dr. Lambers' original KSS method. The red curve is an accelerated KSS method with nodes prescribed in the manner as when p is constant. The green curve demonstrates the Hockbruck-Lubich method of approximating  $e^{-Lt}$ . All methods should be 3rd-order accurate, but only the our new KSS method achieves this level of accuracy. Similarly, Figure 4.2 shows our 4-node comparison on a 256 point grid. Table 4.1 gives estimates of the relative error of the four methods we discussed.

Then we compare our methods using 6-node KSS methods. The results are given in Figure 4.3, Figure 4.4, and Table 4.2. Figure 4.3 gives the relative error for the solution of (2.1) using the four different methods we discussed. The black curve demonstrates our improved KSS method. The blue curve demonstrates Dr. Lambers' original KSS method.



Figure 4.1: 4-Node Comparison on 128 Point Grid

N	$\Delta t$	Improved KSS	KSS-p constant	Original KSS	Hockbruck-Lubich
	1	4.560e-002	2.440e-001	2.605e-002	4.272e-002
	1/2	6.375e-003	1.964e-001	7.848e-003	4.077e-002
128	1/4	1.073e-003	1.891e-001	2.877e-003	3.694e-002
	1/8	2.321e-004	3.967e-001	1.069e-003	1.947e-002
	1/16	4.083e-005	1.983e-001	2.935e-004	3.157e-003
256	1	4.560e-002	2.440e-001	2.604e-002	4.272e-002
	1/2	6.375e-003	1.964e-001	7.839e-003	4.077e-002
	1/4	1.073e-003	1.891e-001	2.868e-003	3.694e-002
	1/8	2.321e-004	3.967e-001	1.134e-003	1.947e-002
	1/16	4.083e-005	1.983e-001	3.089e-004	3.157e-003

Table 4.1: Relative Error Estimates of 4-node Comparisons

The red curve is an accelerated KSS method with nodes prescribed in the manner as when p is constant. The green curve demonstrates the Hockbruck-Lubich method of approximating  $e^{-Lt}$ . All methods should be 5th-order accurate, but only the our new KSS method achieves this level of accuracy. Similarly, Figure 4.4 shows our 6-node comparison on a 256 point grid. Table 4.2 gives estimates of the relative error of the four methods we discussed. Again we find that all of these methods should be 5th-order accurate, but only our new method is.



Figure 4.2: 4-Node Comparison on 256 Point Grid

N	$\Delta t$	Improved KSS	KSS-p constant	Original KSS	Hockbruck-Lubich
128	1	1.090e-001	6.085e+002	1.534e-002	4.665e-002
	1/2	2.323e-002	1.002e+008	2.070e-003	1.755e-002
	1/4	9.684e-004	1.351e+005	1.150e-003	6.993e-003
	1/8	4.323e-005	1.280e+003	4.171e-004	1.441e-002
1000	1/16	4.416e-006	3.848e-002	8.049e-005	2.407e-002
256	1	1.090e-001	1.407e+002	1.533e-002	4.665e-002
	1/2	2.323e-002	4.898e+007	2.065e-003	1.755e-002
	1/4	9.684e-004	6.936e+002	1.506e-003	6.993e-003
	1/8	4.323e-005	7.051e+004	4.763e-004	1.441e-002
	1/16	4.416e-006	1.244e-001	8.375e-005	2.407e-002

Table 4.2: Relative Error Estimates of 6-node Comparisons

### 4.2 Two-Dimensional Parabolic PDE

We now consider high dimensions of the heat equation. To test our methods, we use the two-dimensional coefficients

$$p(x,y) = 1 + \frac{1}{2}\cos(x+y) - \frac{1}{4}\sin(2(x-y)) + \frac{1}{8}\cos(3(x+y))$$
  

$$q(x,y) = 1 + \frac{1}{4}\sin x - \frac{1}{4}\cos 2y + \frac{1}{8}\sin 3x - \frac{1}{8}\cos 4y$$
(4.3)



Figure 4.3: 6-Node Comparison on 128 Point Grid

with initial conditions

$$u(x, y, 0) = 1 + \frac{3}{10}\cos x - \frac{3}{20}\sin(2(x+y)) + \frac{3}{40}\cos 3x, \quad 0 < x, y < 2\pi.$$
(4.4)

The results are given in Figure 4.5, Figure 4.6, and Table 4.3. Figure 4.5 gives the relative error for the solution of (2.1) using the four different methods we discussed. The black curve demonstrates our improved KSS method. The blue curve demonstrates Dr. Lambers' original KSS method. The red curve is an accelerated KSS method with nodes prescribed in the manner as when p is constant. The green curve demonstrates the Hockbruck-Lubich method of approximating  $e^{-Lt}$ . All methods should be 3rd-order accurate, but only the our new KSS method achieves this level of accuracy. Similarly, Figure 4.6 shows our 4-node comparison on 32 points per dimension. Table 4.3 gives estimates of the relative error of the four methods we discussed, with 4 nodes, on a grid with 16 points per dimension.



Figure 4.4: 6-Node Comparison on 256 Point Grid

N	$\Delta t$	Improved KSS	KSS-p constant	Original KSS	Hockbruck-Lubich
16	1	6.723e-002	5.510e-001	2.644e-002	1.351e-001
	1/2	1.443e-002	3.301e-001	8.910e-003	9.480e-002
	1/4	3.191e-003	6.699e-001	2.956e-003	4.589e-002
	1/8	4.519e-004	4.469e-001	7.558e-004	6.060e-003
	1/16	1.101e-004	1.133e-002	1.360e-004	1.270e-003
32	1	6.724e-002	6.013e-001	2.646e-002	1.351e-001
	1/2	1.320e-002	4.232e-001	8.593e-003	9.480e-002
	1/4	1.511e-003	7.755e-001	3.273e-003	4.589e-002
	1/8	3.186e-004	5.897e-001	1.090e-003	6.060e-003
	1/16	1.189e-004	2.141e-001	2.662e-004	1.270e-003

Table 4.3: Relative Error Estimates of 2D Comparisons



Figure 4.5: 2D, 4-Node Comparison on 16 Points per Dimension



Figure 4.6: 2D, 4-Node Comparison on 32 Points per Dimension

## **Chapter 5**

## CONCLUSION

We examined the problem of stiffness that has been introduced into systems of ordinary differential equations that arise from the spatial discretization of partial differential equations. We discussed the Hockbruck-Lubich method of approximating  $e^{-Lt}$ , and determined that it takes too many Lanczos iterations to achieve a suitable level of accuracy. We discussed Krylov Subspace Spectral methods to address this stiffness issue. Although KSS methods are more scalable than the Hockbruck-Lubich method, they are too computationally expensive per time step. From our asymptotic study, we were led to develop a formula that uses the largest generalized eigenvalue to approximate the largest quadrature node for each component. This method has allowed us improve the efficiency of the existing KSS methods without sacrificing any accuracy. In this thesis, we solved a one-dimensional and two-dimensional paraboic PDE to analyze our results. We found that in addition to not sacrificing any accuracy, we actually improved the accuracy of the existing KSS methods. Now we have set the foundation to develop an efficient method that is more practical for stiff problems.

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