# Shortest Route at Dynamic Location with Node Combination-Dijkstra Algorithm 

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#### Abstract

Online transportation has become a basic requirement of the general public in support of all activities to go to work, school or vacation to the sights. Public transportation services compete to provide the best service so that consumers feel comfortable using the services offered, so that all activities are noticed, one of them is the search for the shortest route in picking the buyer or delivering to the destination. Node Combination method can minimize memory usage and this methode is more optimal when compared to $A^{*}$ and Ant Colony in the shortest route search like Dijkstra algorithm, but can't store the history node that has been passed. Therefore, using node combination algorithm is very good in searching the shortest distance is not the shortest route. This paper is structured to modify the node combination algorithm to solve the problem of finding the shortest route at the dynamic location obtained from the transport fleet by displaying the nodes that have the shortest distance and will be implemented in the geographic information system in the form of map to facilitate the use of the system.


Keywords- Shortest Path, Algorithm Dijkstra, Node Combination, Dynamic Location (key words)

## I. InTRODUCTION

In the world of technology which is very rapid development, a lot of information contains in an internet both spatial and textual information. The information can be combined into spatial queries or called Geographic In-formation Systems (GIS) [1]. Geographic Information Systems are widely used in information technology such as information on flood-prone areas, land match-ing to plants, tourism, searching for company location or searching for a person's location [2].

Geographic Information System is a very useful tool for spatial analysis that provides functionality for stor-ing, capturing, analyzing and displaying geographic information [3]. However, GIS can also manipulate and display all forms of geo-referenced information efficient-ly so that it can give a decision to the user [4]. One ex-ample is the information used by the user in showing the decision to guide the trip by showing which route or path to go through to reach the destination by display-ing information through the website or mobile device [5]. There are two points to be considered in the process of making decisions about the route of the road on GIS [6], namely: location type and the shortest route search.

A static location is a location that has unchanging coordinates such as gas station locations, hotels, hospi-tals and companies. While the dynamic location is the location where the coordinates may change at any time such as the car or the location of someone who is al-ways on the move [7]. One company that needs this de-cision is a public transport company like a taxi, where the administrator can monitor the fleet and guide the route to the taxi customer. In addition, the shortest route search is now also very popularly used by companies to assist in the delivery of products to consumers [8], con-struction of apartments to regulate the exit points of de-velopment materials [9] as well as a private person who travels.

With the increasing number of users search the shortest route then some researchers use algorithms to overcome the problem including algorithms A*, Dijkstra algorithm and Ant Colony algorithm. Based on existing shortest path search algorithms, the researchers com-pete each other in providing the shortest route to static locations, especially dynamic locations, where the algo-rithms are tested by detailing the shortest route infor-mation with the fastest time and accuracy in the search for the shortest route where the destination location moves constantly.

The shortest route search by using Dijkstra algo-rithm is very popular than other algorithms where Dijkstra algorithm uses greedy principle, that is the steps chosen with minimum weights that will connect between the selected vertices with the other node [10] so that it can provide information the fastest and shortest route. The previous research concludes that Dijkstra's algorithm provides the shortest path information with faster search than the A* algorithm and the Ant Colony algorithm [11],[12],[13],[14].

Dijkstra's algorithm is very fast in searching the shortest route on the undigraph model (not directed) [15]. Dijkstra's algorithm, however, has algorithmic changes for some improvements in optimizing the search for the diagraph model, including: 1. increasing the effectiveness of the diagraph output to avoid infinite looping 2. the algorithm will notice adjacent nodes in the shortest path search, and 3. considering many vertices that can separate labels simultaneously [16]. However, this algorithm does not discuss the memory used in the shortest route search so that researchers have an alternative in an effective and efficient search by com-bining Dijkstra algorithm with Node

Combination algo-rithm, where Node Combination is an algorithm that has easy to understand steps and smaller memory us-age compared to Dijkstra's algorithm, and the Node Combination algorithm can be applied to Dijkstra's algorithm [17].

Based on the background that has been described, the researcher will modify the node combination algo-rithm to solve the problem in the storage of the nodes passed in the search for the shortest route.

## A. Geolocation

Geolocation is a technique of identifying geograph-ical location in the real world that comes from an inter-netconnected object, such as a mobile phone, radar or computer. [18],[19]. In geolocation itself there is the simplest form, where geolocation involves a set of geographical coordinates that are closely related to the use of the system in positioning such as a street address at a location [20].

For geolocation or position, search engines often use radio frequency location (RF) methods, such as Time Difference Of Arrival (TDOA) for precision. TDOA systems often use mapping or other geographic infor-mation systems. If GPS signals are not available, geolo-cation apps may use information from cell towers to triangulate position estimates, methods that are not as accurate as GPS but have improved considerably in re-cent years. This is different from previous location radio technology, such as Direction Finding where the bear-ing line to the transmitter is reached as part of the pro-cess. The word geolocation also refers to the latitude and longitude coordinates of certain locations standard-ized by ISO / IEC 19762-5: 2008 [21].

## B. Google Maps API

Google Map API (Application Programing Interface) is one of the advantages provided by google to access data from google map and google local search. Application Programing Interface is a documentation that consists of interface, function, class, structure and so on to build a software. With this API, it allows the programmer to develop a software to be integrated with other software. API can be said to link an application with other applications that allow programmers to use the system function. This process is managed through the operating system. The advantage of this is to allow an application with other applications to interact and interact. The programming language used by google map consisting of html, javascript, ajax, and xml, allows to display google maps on web pages.

By using the google map API, we can display the digital map on its own website page. In order for this API application to appear, it is necessary to have a key API that is a unique code that is generalized by google map to recognize a number of utilities for manipulating maps and adding content in maps through various services, allowing the creation of powerful map applications on web pages [22],[23].

## C. Node Combination Algorithm

Node combination has a search system by combining all nodes and still maintaining a labeling set such as Dijkstra's algorithm which will be depicted in "Fig. 1" [24],[17] Suppose all nodes are connected to a rope and will be pulled
slowly, over time all nodes will converge into one where the string will see the shortest strap.

Fig. 1. An example of a combination of nodes

(a) conditions before being combined

(b) conditions after being combined

The shortest path search steps using a combination node can be explained by pseudocode to be described in Table 1

TABLE I. PSEUDOCODE ALGORITMA NODE COMBINATION
Algoritma Node Combination (G,s)

```
\(1 \mathrm{~W}[\mathrm{~s}, \mathrm{u}]:=0, \mathrm{vu}:=\mathrm{vs}, \mathrm{V}:=\mathrm{V}-\{\mathrm{s}\} \quad / *\) Initialization
*/
2 while \(\mathrm{W}[\mathrm{s}, \mathrm{u}]<\infty\) and \(|\mathrm{V}|>0\)
\(3 \mathrm{~V}:=\mathrm{V}-\{\mathrm{u}\} \quad / *\) Node Combination */
4 for each j in V
\(5 \mathrm{~W}[\mathrm{~s}, \mathrm{j}]:=\min \{\mathrm{W}[\mathrm{s}, \mathrm{j}], \mathrm{W}[\mathrm{s}, \mathrm{u}]+\mathrm{W}[\mathrm{u}, \mathrm{j}]\} / *\) updating
edge weights */
\(6 \mathrm{vu}:=\) the nearest neighbor of s in \(\mathrm{V} / *\) finding the
nearest neighbor */
/* at the end of the algorithm, the sth row in W contains
the corresponding distances */
```


## D. Node Combination-Dijkstra Algorithm

Node Combination-Dijkstra algorithm is algorithm developed by writer in search of shortest route by adopting Node Combination search method that is combining one node with other node and modification in case of storage node which will be explained by pseudocode in Table 2.

TABLE II. PSEUDOCODE ALGORITMA NODE COMBINATION-DIJSKTRA
Algoritma Node Combination-Dijkstra (G,s)

1. $W[s, u]:=0, v_{u}:=v_{s}, P[u]:=[s], V:=V-\{s\} / *$

Initialization */
2. while $W[s, u]<\infty$ and $|V|>0$
3. $V:=V-\{u\} / *$ Node Combination */
4. for each $j$ in $V$
5. If $W[s, u]+W[u, j]<W[s, j]$
6. Then $P[\mathrm{j}]:=P[u] \cup\{j\}, / *$ updating path */
7. $W[s, j]:=W[s, u]+W[u, j] / *$ updating edge weights */
8. $v u:=$ the nearest neighbor of $s$ in $V / *$ finding the nearest neighbor */
/* at the end of the algorithm, the sth row in W contains the corresponding distances */

## II. DATA

The data used for this study is the location of places in the city surabaya which is interconnected between one point with another point with information weight / distance in meters, latitude longitude coordinates, and address which will be explained in Table 3.

TABLE III. CORDINAT POINTS

| No | From | To | Coordinat | Weight | Addres |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | A | B | "-7.274982997930746, 112.72668129821636" | 1692.91 | J. Pasar Kembang Kec. Sawahan, Kota SBY, Jawa Timur |
|  |  | D |  | 954.0145 |  |
|  |  | E |  | 728.3023 |  |
|  |  | G |  | 2103.285 |  |
| 2 | B | A | "-7.272939639599, 112.741873330076" | 1692.91 | JI. Basuki Rahmat Embong Kaliasin, Genteng, Kota Surabaya |
|  |  | C |  | 826.8011 |  |
|  |  | M |  | 2451.245 |  |
|  |  | J |  | 1692.327 |  |
| 3 | k | 1 | "-7.283092639519452,$112.75612122436382 "$ | 785.2675 | Bank Antardaerah. PT - Pucang Anom Timur J. Pucang Anom Tim. No.19, Surabaya |
|  |  | J |  | 483.4498 |  |
|  |  | N |  | 1374.426 |  |
| 4 | L | M | "-7.269129757840544,$112.77423149963238 "$ | 1624.072 | J. Raya Dharma Husada Indah Mulyorejo, Kota SBY, Jawa Timur |
|  |  | 1 |  | 1759.133 |  |
| 5 | M | B | "-7.262233296342604, | 2451.245 | JI. Kedung Tarukan No.86C Pacar Kembang, Tambaksar |
|  |  | L | 112.76127106567242" | 1624.072 |  |
| 6 | N | K | "-7.290867201464144, 112.74645163518085" | 1374.426 | STE YOUTH SURABAYA JL Bung Tomo, No. 08, Kavling 8 Surabaya |
|  |  | G |  | 1388.481 |  |
|  |  | 0 |  | 727.7807 |  |
| 7 | 0 | N | "-7.297337545443966, | 727.7807 | J. Ngagel Rejo Kidul $54-38$ Ngagelrejo, Wonokromo, Kota Surabaya |
|  |  | P | 112.74739577275409" | 1439.804 |  |
| 8 | P | 0 | "-7.309256355149782, | 1439.804 | Bendul Merisi Wonocolo, Kota Surabaya |
|  |  | Q | 112.742331762133971 | 760.3496 |  |
| 9 | Q | P | "-7.306617217538221, | 760.3496 | J. Raya Malang - Surabaya Wonokromo, Kota Surabaya |
|  |  | H | 112.7359802911867" | 790.3205 |  |
| 10 | c | B | "-7.266213519414126, | 826.8011 | JI. Tegalsari No. 25 Tegalsari, Kota Surabaya |
|  |  | D | 112.73869759460308" | 1670.288 |  |
| 11 | D | C | "-7.26697979135388, | 1670.288 | II. Petemon No. 5 Petemon, Kec. Sawahan, Kota Surabaya |
|  |  | A | 112.7235913934312" | 954.0145 |  |
| 12 | E | A | "-7.278729130708208, | 728.3023 | J. Putat Jaya 30 Putat Jaya, Kec. Sawahan, Kota Surabaya |
|  |  | F | 112.72127396484234" | 1539.035 |  |
| 13 | F | E | "-7.292266031440466, | 1539.035 | Jl. Mayjen Sungkono 75-143 Gn. Sari, Dukuh Pakis, Kota Surabaya |
|  |  | G | 112.72410637756207" | 1090.059 |  |
| 14 | G | A | "-7.292436304316622, 112.73397690673687" | 2103.285 | Bubur Ayam Rasa Darmo, Wonokromo, Surabaya |
|  |  | F |  | 1090.059 |  |
|  |  | N |  | 1388.481 |  |
|  |  | H |  | 807.5721 |  |
| 15 | H | Q | "-7.299672841676093, | 790.3205 | JI. Joyoboyo Sawunggaling, Wonokromo, Kota Surabaya |
|  |  | G | 112.73499189086773" | 807.5721 |  |
| 16 | 1 | L | "-7.2796019598365245, 112.76230103393414" | 1759.133 | II. Raya Menur Manyar Sabrangan, Mulyorejo, Kota Surabaya |
|  |  | 1 |  | 698.3279 |  |
|  |  | K |  | 785.2675 |  |
| 17 | J | B | "-7.278750570426501, 112.75603539367535" | 1692.327 | II. Dharmawangsa No. 97 Kertajaya, Gubeng, Kota Surabaya |
|  |  | 1 |  | 698.3279 |  |
|  |  | K |  | 483.4498 |  |

## III. Result and discussion

## A. Dynamic Location Search

Determination of coordinates on the built system uses a geolocation where the identification or estimation of the real-world geographical location of an object such as a mobile phone. Determination of coordinates on this system is the basic system built to sync the object coordinates with the coordinates that already exist in the system.

Locations generated from geolocation are not all coordinates listed on the location data in this shortest route search information system so that the determination of coordinates generated from the geolocation with the coordinates present in the system will be calculated for the closest search of the coordinates to the coordinates present in system.

The algorithm for calculating the coordinate distance of one with the other coordinates will be written in the Table 4 about pseudocode distance calculation.

TABLE IV. PSEUDOCODE DISTANCE CALCULATION

```
function jarak (\$lat1, \$lon1, \$lat2, \$lon2, \$hasil)
    \$kordinat \(=\$ l o n 1-\$ l o n 2\);
2 \$jarak \(=\sin (\operatorname{deg} 2 \operatorname{rad}(\$ l a t 1)) \quad *\)
    \(\sin (\operatorname{deg} 2 \operatorname{rad}(\$ l a t 2))+\cos (\operatorname{deg} 2 \operatorname{rad}(\$ \operatorname{lat} 1)) \quad *\)
    \(\cos (\) deg \(2 \operatorname{rad}(\$ \operatorname{lat} 2)) * \cos (\) deg \(2 \operatorname{rad}(\$ k o d i n a t))\);
    \$ jarak = acos(\$jarak);
    \$ jarak \(=\) rad2deg (\$jarak);
    \$mili \(=\$\) jarak \(* 60 * 1.1515\);
    \$hasil = strtoupper (\$hasil);
    if (\$hasil \(==\) " \(K\) ") \{
    return (\$mili * 1.609344 )
    \} else if (\$hasil \(==" N ")\) \{
    return (\$mili * 0.8684);
    \} else \{
    return \$mili; \}
    end
```

After the calculation of the algorithm it will get the distance between the coordinates of one with the coordinates in the system then, the location is given the coordinates that have the closest distance so that the location will be initialized as a starting point in the search for the shortest route. "Fig. 2" is an example of the position of the geolocation represented by the letter "P" with the graph contained in the system.


Fig. 2. Position of geolocation with graph.
From "Fig. 2" we can conclude that the position of the geolocation is initialized starting point is at point "A" because that point has the most shortest distance.

## B. Implementation of Node Combination-Dijsktra Algorithm

The Node Combination-Dijkstra algorithm has a base that the initial node will find the closest vertices to the destination node, when it finds the nearest node, the node will be joined with the next node until the end result of this system will be a node that informs the shortest distance between the initial node with the destination node along with the points that have been passed.

For the application of Node Combination-Dijkstra algorithm will be described in Table 5 where the authors will calculate into the assessment matrix starting from point I with the goal to point $C$. with this matrix the writer can show that the Node Combination-Dijkstra algorithm has the same calculation pattern with Node Combination but the algorithm can also save the history passed.

[^0]TABLE V. The Route Search Matrix From Point I to point C

| Step | Start | From I to C |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | A | B | C | D | E | F | G | H | I | J | K | L | M | N | O | P | Q |
| 1 | I $=698$ |  |  |  |  |  |  |  |  |  | $\begin{aligned} & 698 \\ & \{I, J\} \end{aligned}$ | $\begin{aligned} & 785 \\ & \{\mathrm{I}, \mathrm{~K}\} \end{aligned}$ | $\begin{aligned} & 1759 \\ & \{\mathrm{I}, \mathrm{~L}\} \end{aligned}$ |  |  |  |  |  |
| 2 | J $=785$ |  | $\begin{gathered} 698+1692 \\ =2390 \\ \{\mathrm{I}, \mathrm{~J}, \mathrm{~B}\} \end{gathered}$ |  |  |  |  |  |  |  |  | $\begin{gathered} \hline 698+483 \\ =1181> \\ 785 \\ \text { Then } \\ 785 \\ \{\mathrm{I}, \mathrm{~K}\} \end{gathered}$ | 1759 |  |  |  |  |  |
| 3 | $\begin{gathered} K \\ =1759 \end{gathered}$ |  | 2390 |  |  |  |  |  |  |  |  |  | 1759 |  | $\begin{gathered} 785+1374 \\ =2159 \\ \{I, K, N\} \end{gathered}$ |  |  |  |
| 4 | $=2159$ |  | 2390 |  |  |  |  |  |  |  |  |  |  | $\begin{gathered} 1759+1624 \\ =3383 \\ \\ \{I, L, M\} \end{gathered}$ | 2159 |  |  |  |
| 5 | $\begin{aligned} & N \\ = & 2390 \end{aligned}$ |  | $\begin{gathered} 2390 \\ \{\mathrm{I}, \mathrm{~J}, \mathrm{~B}\} \end{gathered}$ |  |  |  |  |  |  |  |  |  |  | 3383 |  | $\begin{gathered} \begin{array}{c} 2159+727 \\ =2886 \end{array} \\ \{\mathrm{I}, \mathrm{~K}, \mathrm{~N}, \mathrm{O}\} \end{gathered}$ |  |  |
| 6 | $\begin{gathered} B \\ =2886 \end{gathered}$ | $\begin{gathered} 2390+1692 \\ =4083 \\ \{\mathrm{I}, \mathrm{~J}, \mathrm{~B}, \mathrm{~A}\} \end{gathered}$ |  | $\begin{gathered} \begin{array}{c} 2390+826 \\ =3216 \end{array} \\ \{I, J, B, C\} \end{gathered}$ |  |  |  |  |  |  |  |  |  | 3383 |  | 2886 |  |  |
| 7 | $\begin{gathered} O \\ =3216 \end{gathered}$ | 4083 |  | $3216$ $\{\mathrm{I}, \mathrm{~J}, \mathrm{~B}, \mathrm{C}\}$ |  |  |  |  |  |  |  |  |  | 3383 |  | 4325 | $\begin{gathered} 2886+1439 \\ =4325 \\ \{I, K, N, O, P\} \end{gathered}$ |  |

From Table 5 it can be concluded that the system of finding the shortest route from point I to point C can be reached through the point I-J-B-C with a distance of 3216 meters.

## C. System Analysis

To prove that the system has chosen the shortest route, a few alternate paths from the starting point are located at point I with the end point at point C. Alternative routes that can be passed will be explained in Table 6.

TABLE VI. Route Search Alternatives

| Nomor Alternatif | Start | End | Rute | Jarak yang ditempuh |
| :---: | :---: | :---: | :---: | :---: |
| Alt. 1 | I | C | I-J-B-C | 3216 (m) |
| Alt. 2 |  |  | I-K-J-B-C | 3786 (m) |
| Alt. 3 |  |  | I-L-M-B-C | 6660 (m) |
| Alt. 4 |  |  | $\begin{aligned} & \text { I-L-M-B-A-D- } \\ & \text { C } \end{aligned}$ | 10150 (m) |
| Alt. 5 |  |  | $\begin{aligned} & \text { I-K-N-O-P-Q- } \\ & \text { H-G-A-D-C } \end{aligned}$ | 11409 (m) |

In Table 6 we can see that there are several alternatives in route determination and we can conclude that the shortest route search system using the Node Combination-Dijkstra Algorithm selects the first alternative with a distance of 3216 meters which is the shortest distance among other alternatives, and also, the system can record the history of the nodes passed the

## Conclusion

This research has presented the shortest route search system by modifying the combination node algorithm to the dynamic location used in general such as mobile phone online transportation. From the test result of the system, the node combination-djikstra algorithm can give optimal result in the shortest route search on the dynamic location applied in surabaya city that shows the route search from point I to point C is 3216 meters for the distance with I-J-B-C route.

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