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Statistics on parallelogram polyominoes and a q, t -analogue of the Narayana numbers

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Abstract

We study the statistics **area**, **bounce** and **dinv** on the set of parallelogram polyominoes having a rectangular m times n bounding box. We show that the bi-statistics (**area**, **bounce**) and (**area**, **dinv**) give rise to the same q, t -analogue of Narayana numbers which was introduced by two of the authors in [4]. We prove the main conjectures of that paper: the q, t -Narayana polynomials are symmetric in both q and t , and m and n . This is accomplished by providing a symmetric functions interpretation of the q, t -Narayana polynomials which relates them to the famous diagonal harmonics.

Keywords:

q, t -Narayana, parallelogram polyominoes, parking functions, bounce, **dinv**
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1. Introduction

A *parallelogram polyomino* having an $m \times n$ bounding box is a polyomino in a rectangle consisting of $m \times n$ cells that is formed by cutting out two (possibly empty) non-touching standard Young tableaux which have corners at $(0, n)$ and $(m, 0)$. An example of a parallelogram polyomino having a 12×7 bounding box is illustrated in Figure 1. Let $\text{Poly}_{m,n}$ be the set of all parallelogram polyominoes having a rectangular $m \times n$ bounding box. The cardinality of $\text{Poly}_{m,n}$ is known to be $N(m+n-1, m)$ where for positive integers a and b ,

$$N(a, b) := \frac{1}{a} \binom{a}{b} \binom{a}{b-1}$$

are the famous *Narayana numbers*. Two authors of this work introduced [4] two statistics on these combinatorial objects, **area** and **bounce**, which led to a q, t -analogue of the Narayana numbers $N(m+n-1, m)$, namely

$$\text{Nara}_{m,n}(q, t) := \sum_{P \in \text{Poly}_{m,n}} q^{\text{area}(P)} t^{\text{bounce}(P)},$$

that they called the q, t -Narayana polynomial. In that same work it was conjectured that the q, t -Narayana polynomials are symmetric in q and t , and as expressions were also symmetric in m and n . We introduce a new statistic **dinv** which gives a new q, t -analogue of the same numbers

$$\widetilde{\text{Nara}}_{m,n}(q, t) := \sum_P q^{\text{dinv}(P)} t^{\text{area}(P)}.$$

The following theorem establishes a relation between these two polynomials.

Theorem 4.1. *For all $m \geq 1$ and $n \geq 1$, we have*

$$\text{Nara}_{m,n}(q, t) = \widetilde{\text{Nara}}_{n,m}(q, t).$$

We give two proofs of this result, one using an explicit bijection, and another one using a recursion. The main result of this paper is the proof of the symmetries conjectured in [4].

Theorem 6.2. *For all $m \geq 1$ and $n \geq 1$, we have*

$$\text{Nara}_{m,n}(q, t) = \text{Nara}_{m,n}(t, q)$$

and

$$\text{Nara}_{m,n}(q, t) = \text{Nara}_{n,m}(q, t).$$

In particular

$$\text{Nara}_{m,n}(q, t) = \widetilde{\text{Nara}}_{m,n}(q, t).$$

In order to prove this result, we will give a symmetric functions interpretation of our q, t -Narayana numbers:

Theorem 6.1. *For all $m \geq 1$ and $n \geq 1$ we have*

$$\text{Nara}_{m,n}(q, t) = (qt)^{m+n-1} \cdot \langle \nabla e_{m+n-2}, h_{m-1} h_{n-1} \rangle,$$

where e_k and h_k are the elementary and the homogeneous symmetric functions of degree k respectively, ∇ is the well known nabla operator introduced by Bergeron and Garsia (see [2, Section 9.6]), and the scalar product is the usual Hall inner product on symmetric functions.

This result establishes a remarkable link between the q, t -Narayana polynomials and the well-known *diagonal harmonics* DH_n , since ∇e_n is the Frobenius characteristic of this important module of the symmetric group \mathfrak{S}_n , as shown by Haiman in [7].

Haglund [5] gave a combinatorial interpretation of the particular polynomial $\langle \nabla e_{m+n-2}, h_{m-1} h_{n-1} \rangle$ in terms of parking functions. In fact Haglund's result would be an easy consequence of the famous *shuffle conjecture*, which predicts a combinatorial interpretation of ∇e_n in terms of parking functions (see [6, Chapter 6]), if a proof of it could be found.

In order to prove Theorem 6.1, we use the results of Section 5, proving that the combinatorial polynomials in Haglund's result and our q, t -Narayana polynomials both satisfy the same recursion. This paper is organized in the following way:

- In Section 2 we define three statistics on parallelogram polyominoes and two q, t -analogues of Narayana numbers.
- In Section 3 we establish a bijection between our parallelogram polyominoes and a set of Dyck paths. We classify those words that are area words of members of $\text{Polyo}_{m,n}$. Area words are important in the definition of several statistics mentioned in Section 2.

- In Section 4 we present a bijection from $\text{Polyo}_{m,n}$ to $\text{Polyo}_{n,m}$ which sends the bi-statistic $(\text{area}, \text{bounce})$ to the bi-statistic $(\text{div}, \text{area})$, thereby establishing Theorem 4.1.
- In Section 5 we prove a recursion satisfied by both of our q, t -Narayana polynomials, which gives another proof of Theorem 4.1.
- In Section 6 we provide the necessary background to state Theorem 6.1, and we show how Theorem 6.2 follows from it. Theorem 6.1 is then proven.

2. Three statistics on parallelogram polyominoes

We may give an alternative characterisation of parallelogram polyominoes in terms of non-intersecting paths in the plane. This alternative characterisation will prove useful in defining the statistics and mappings used in the paper.

Consider a rectangular grid in \mathbb{Z}^2 of width m and height n . On this grid consider two paths, both starting from the Southwest corner and arriving at the Northeast corner, travelling on the grid, performing only North or East steps, with the further restriction that they touch each other only at the starting point and at the ending point. Such a pair of paths uniquely defines a parallelogram polyomino. The region between the two paths is called the *interior* of the (parallelogram) polyomino. The two paths defining the parallelogram polyomino of Figure 1 are coloured in red and green, and the interior has been shadowed.

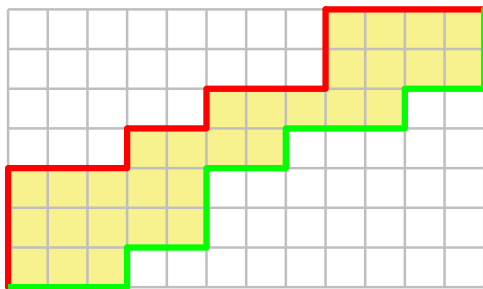


Figure 1: A parallelogram polyomino having a 12 times 7 bounding box.

In what follows we will encode a parallelogram polyomino as an *area word* consisting of natural numbers (*unbarred numbers*) and natural numbers with

a bar on top (*barred numbers*), in the following way. We will label every North step of the upper (red) path with a barred number, and every East step of the lower (green) path with an unbarred number. This is done in two stages.

First, for each East step of the lower path we draw a line starting with the East endpoint and going Northwest until reaching the upper path: we label this step with the number of squares crossed by this line. Second, we label each North step of the upper path with the number of squares in the interior of the polyomino to the East of it which were not crossed by any of the lines that we drew during the previous stage. An example of this labelling is shown in Figure 2, where we put a black dot in the non-crossed squares.

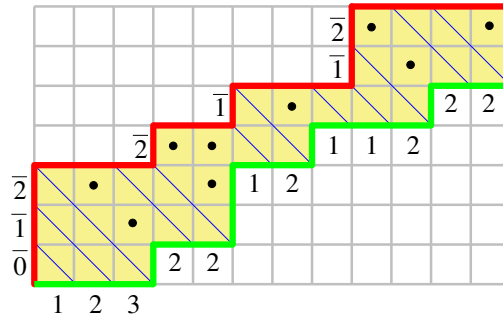


Figure 2: The parallelogram polyomino of Figure 1 with its perimeter labelled.

Once we have done this labelling, we read the labels in the following order: starting from Southwest and going to Northeast imagine moving a straight line of slope -1 over the polyomino. When we encounter vertical steps of the upper path or horizontal steps of the lower path we write the corresponding labels. If we encounter both types of steps at the same time then we write the label of the upper path first. The area word of the example in Figure 2 is $\bar{0}1\bar{1}2\bar{2}322\bar{2}1\bar{1}211\bar{1}2\bar{2}22$.

Notice that the sum of these numbers (disregarding the bars) gives the **area** of the polyomino, which is the number of squares between the two paths. This is the first of the statistics that are relevant to us. In the example the **area** is 30.

Next we will define the **dinv** statistic. Consider the total order on the labels

$$\bar{0} < 1 < \bar{1} < 2 < \bar{2} < 3 < \bar{3} < 4 < \bar{4} < \dots .$$

Given a polyomino with area word $a_1 a_2 \dots a_k$, we define its **dinv** as the number

and

$$\widetilde{\mathbf{Nara}}_{m,n}(q, t) := \sum_{P \in \text{Polyo}_{m,n}} t^{\text{bounce}(P)} q^{\text{area}(P)}.$$

are studied in this paper. The polynomials $\mathbf{Nara}_{m,n}(q, t)$ were first introduced in [4] by two of the authors of the present work. In the same paper, it was conjectured that these were polynomials symmetric in q and t , and as expressions symmetric in m and n .

3. A bijection with Dyck paths

In this section we present a bijection \mathbf{ptd} between the set $\text{Polyo}_{m,n}$ and a set of Dyck paths having length $2(m+n)$. We then prove a result which shows how to read the area word of a parallelogram polyomino from its corresponding Dyck path under \mathbf{ptd} . From this we will get a characterization of the area words of polyominoes from $\text{Polyo}_{m,n}$ which will be used in the proof of Theorem 4.1. We finally observe that this description provides a way to computationally work with the set of area words of $\text{Polyo}_{m,n}$ by working with the easier to construct set of Dyck paths.

Recall that a *Dyck path* can be thought of as a path consisting of North-east or Southeast steps lying between parallel horizontal lines, such that the path starts with a Northeast step, it never crosses the starting horizontal line, and returns to it at the end. Its *length* is simply the number of its steps it contains. Figure 4 shows an example of a Dyck path having length 38.

Notice that a Dyck path is uniquely determined by the sequence of rises and falls we encounter as we move along the path from left to right.

We will next describe a bijection between the polyominoes in $\text{Polyo}_{m,n}$ and the set of Dyck paths of length $2(m+n)$ with m rises in even positions and n rises in odd positions, which do not return to the starting horizontal line until the end. This bijection appears in [3] in a somewhat different language.

The idea is to read the steps of the upper and lower paths of a parallelogram polyomino P alternately and form the Dyck path $D = \mathbf{ptd}(P)$ by using two rules. We perform a rise of the Dyck path for either a North step of the upper path or an East step of the lower path, and perform a descent of the Dyck path for either an East step of the upper path or a North step of the lower path. Using this construction, the polyomino in Figure 1 is sent to the Dyck path shown in Figure 4.

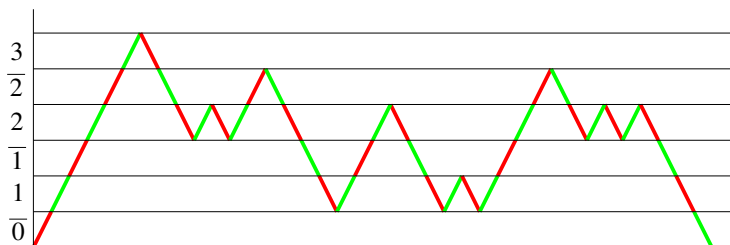


Figure 4: The Dyck path corresponding to the polyomino of Figure 1.

It should be clear that this mapping sends the parallelogram polyominoes to the stated subset of Dyck paths. The fact that the Dyck path does not return to the starting horizontal line before the end corresponds to the fact that the upper and the lower paths do not intersect each other between the starting and ending points. The inverse operation is straightforward to describe and verify.

We can easily read the area word of a parallelogram polyomino P from the corresponding Dyck path $\text{ptd}(P)$ as we will now describe. Consider the Dyck path in Figure 4 when reading the next proposition. It consists of Northeast and Southeast steps lying between parallel lines which determine certain rows.

Proposition 3.1. *Let $P \in \text{Polyo}_{m,n}$ with $D = \text{ptd}(P)$. If we label the rows of D with $\bar{0}, 1, \bar{1}, 2, \bar{2}, 3, \bar{3}, \dots$ from bottom to top, then reading the labels of the rows of the rises from left to right we get the area word of the polyomino P .*

To prove this proposition, we will use induction on the number of pairs of steps of the upper and lower paths starting from the Southwest corner. At each step of this induction we will consider the *partial box* that includes the partial paths, i.e. the smallest rectangle that includes them (see Figure 5). Then we imagine to complete the paths inside the partial box by moving along the edges to reach the Northeast corner, and we read the labels of the resulting polyomino on the partial paths.

We claim that this gives exactly the corresponding part of the area word of the original polyomino.

The key observation is the following claim.

Claim. *In the last pair of steps, the label of a North step of the upper path is always the distance from the right edge of the previous partial box, with a bar*

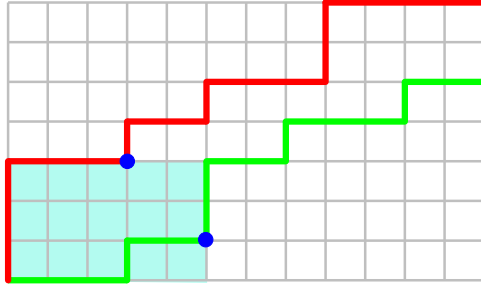


Figure 5: The highlighted area is the partial box after the first 6 steps of both the green path and the red path.

on top; while the label of a East step of the lower path is always the distance from the upper edge of the partial box.

After proving this claim, it remains only to observe that the distance from the right edge of the previous partial box of the North steps of the upper path corresponds to the number of odd rows from the bottom line in the corresponding Dyck path; while the distance from the upper edge of the partial box of the East steps of the lower path corresponds to the number of even rows from the bottom line. This completes the proof of the proposition.

Proof of the Claim. At the beginning the upper path is forced to go North and the lower path is forced to go East. The partial box at this point consists of a single square, and we clearly have the partial area word $\bar{0}1$: this is always the beginning of an area word for a polyonimo, and it corresponds to the first two rises in the corresponding Dyck path, as it should be.

Now suppose that everything works up to a certain pair of steps, and let us make the next pair of steps. We have four cases (see Figure 6):

Case 1: The upper path moves East, and the lower path moves North. Then there are no labels to add, and the previous labels remain unchanged, since the partial box remains unchanged.

Case 2: Both the upper and the lower paths move North. Then the label of the North step of the upper path is clearly the distance from the right edge of the partial box, which is the same distance from the one of the previous partial box. The previous labels clearly remain unchanged.

Case 3: Both the upper and the lower paths move East. Then the label of the East step of the lower path is the distance to the upper edge of the partial box. The previous labels of the upper path remain unchanged, since

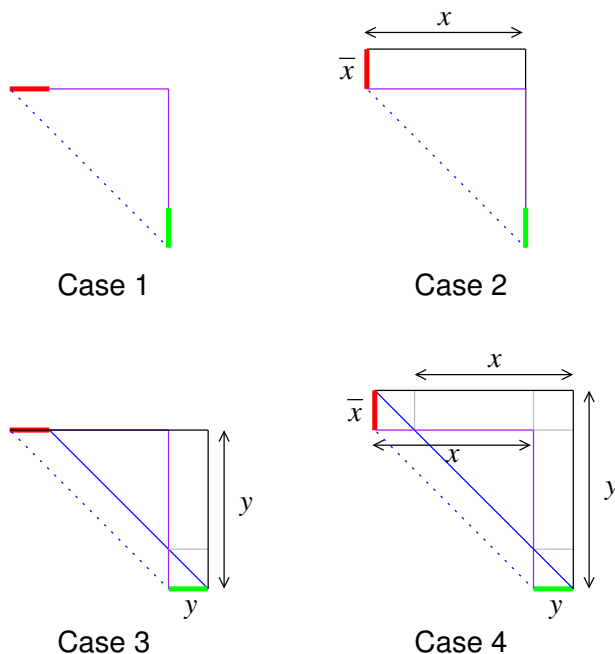


Figure 6: The four cases. The previous partial box is marked by violet lines, while the new one is marked by black lines.

in each row we added just a box crossed by the diagonal corresponding to the new East step of the lower path. The previous labels of the lower path also remain unchanged, since we did not move the upper edge of the partial box.

Case 4: The upper path moves North, and the lower path moves East. Then the label of the North step of the upper path is the distance from the right edge of the partial box minus 1, since the first box becomes crossed by the diagonal of the East step of the lower path. But this is equal to the distance from the right edge of the previous partial box. The label of the East step of the lower path is clearly the distance from it to the upper edge of the partial box. The previous labels of the upper path remain unchanged, since in each row we added just a box crossed by the diagonal corresponding to the new East step of the lower path. The previous labels of the lower path also remain unchanged, since the diagonals of the previous horizontal steps all hit the upper path in the same spots as before. \square

As an immediate consequence, we get a characterization of the words in

the ordered alphabet $\bar{0} < 1 < \bar{1} < 2 < \bar{2} < 3 < \bar{3} < \dots$ which are area words of elements of $\text{Polyo}_{m,n}$.

We state this characterization here as a corollary.

Corollary 3.2. *Consider the alphabet $\bar{0} < 1 < \bar{1} < 2 < \bar{2} < 3 < \bar{3} < \dots$, with the letters in the given order. A word $a_1 a_2 \dots a_r$ in this alphabet is the area word of an element of $\text{Polyo}_{m,n}$ if and only if the following conditions hold:*

1. $a_1 = \bar{0}$, and this is the only $\bar{0}$ that appears in the word;
2. there are exactly m of the a_i 's which are from the set of numbers without a bar $\{1, 2, 3, \dots\}$, and exactly n of the a_i 's which are from the set of numbers with a bar $\{\bar{0}, \bar{1}, \bar{2}, \dots\}$ (in particular $r = m + n$);
3. for all $i = 1, 2, \dots, m + n - 1$, the letter a_{i+1} is less than or equal to the immediate successor of the letter a_i , in the given order on the alphabet.

We mention here that this bijection also gives an easy way to construct the polyomino from its area word: draw the corresponding Dyck path (this is immediate), and then look at the odd and even steps to construct the polyomino.

4. The bi-statistics (area, bounce) and (dinv, area)

This section is dedicated to proving the following theorem.

Theorem 4.1. *For all $m \geq 1$ and $n \geq 1$,*

$$\text{Nara}_{m,n}(q, t) = \widetilde{\text{Nara}}_{n,m}(q, t).$$

In order to prove this theorem it suffices to give a bijection from $\text{Polyo}_{m,n}$ to $\text{Polyo}_{n,m}$ which sends the bi-statistic (area, bounce) to the bi-statistic (dinv, area).

The bijection that we will now describe is similar in spirit to the one used in the proof of the analogous [6, Theorem 3.15].

Let $P \in \text{Polyo}_{m,n}$. Starting from P , we read the labels of its bounce path, getting a word consisting of barred and unbarred numbers. Then, starting from the bottom-left corner, for each turn of the bounce path, we look at the part of the path (upper or lower) that includes it. For example in the polyomino of Figure 3, the first turn of the bounce path is between $\bar{0}$ and the next 1 in the labelling of the bounce path. The containing path consists

of the first 4 steps (counted from the Southwest corner) of the upper path. We label the vertical steps of the containing path with the labels used for the vertical steps in that part of the bounce path, and the horizontal steps of the containing path with the labels used for the horizontal steps in that part of the bounce path. See Figure 7 for an example.

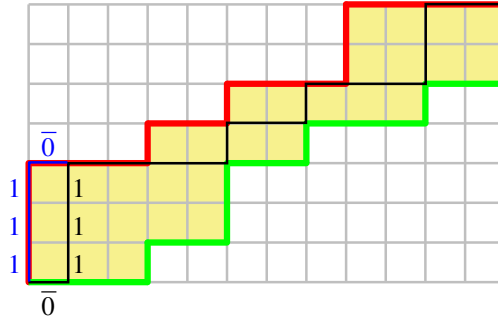


Figure 7: The containing path is the blue line from $(0,0)$ to $(1,3)$ and the new labels for each of its steps are also blue.

We then read the new labels by following the containing path from North-east down to Southwest. In the example we read $\bar{0}111$.

During the remainder of the construction we will preserve the relative positions of these labels.

We then repeat the algorithm with the second turn of the bounce path of P . In the example this occurs between the last 1 and the first $\bar{1}$ in the bounce path. This time the containing path consists of the steps of the lower path between the second and the eighth. We repeat the procedure that we used before, and the word that we get reading the new labels will prescribe the relative positions of the 1's and the $\bar{1}$'s. In the example (see Figure 8) we get the prescriptions $1\bar{1}\bar{1}\bar{1}\bar{1}\bar{1}$. This together with the other prescription gives a partial word $\bar{0}11\bar{1}\bar{1}\bar{1}\bar{1}$.

In general we will construct this partial word in a way that it can be the word of a parallelogram polyomino while respecting all the prescriptions. This will always be possible since the first step of the containing path that we read will always be labelled by the smallest of the two types of labels that we are considering: this is due to the definition of the bounce path.

We keep doing this until all the labels of the bounce path of P have been included. At the end we will get a word of another parallelogram poly-

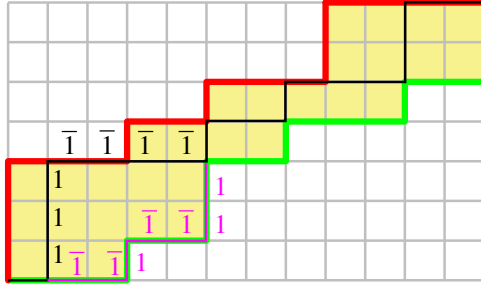


Figure 8: The containing path is the new violet line along the Southeast border with new violet labels added in the same corner.

omino. In the example, at the next step we get the prescriptions $\overline{112\overline{11}}$, which gives the partial word $\overline{011\overline{11}2\overline{11}}$; then we get the prescriptions $\overline{222}$, which gives the partial word $\overline{011\overline{11}2\overline{22}1\overline{11}}$; then we get the prescriptions $\overline{223}$, which gives the partial word $\overline{011\overline{11}2\overline{22}3\overline{11}}$; then we get the prescriptions $\overline{3333}$, which gives the partial word $\overline{011\overline{11}2\overline{22}3\overline{33}3\overline{11}}$; then we get the prescriptions $\overline{3344\overline{3}}$, which gives the partial word $\overline{011\overline{11}2\overline{22}3\overline{33}4\overline{43}1\overline{11}}$; and finally we get the prescriptions $\overline{4444}$, which gives the final word $\overline{011\overline{11}2\overline{22}3\overline{33}4\overline{44}4\overline{3}1\overline{11}}$.

It is clear from the construction and the characterization of Corollary 3.2, that in this way we get the area word of a polyomino $F(P)$ in $\text{Polyo}_{n,m}$. Moreover $F(P)$ clearly has **area** equal to the **bounce** of the original polyomino P , again by construction. Figure 9 illustrates $F(P)$ for when P is the polyomino of Figure 1.

We need to show that the **dinv** of $F(P)$ is equal to the **area** of P .

To see this, recall how we constructed the word of the new polyomino: for consecutive types of labels, we prescribed the relative positions by reading the corresponding containing path. But in the containing path, those pairs of vertical and horizontal steps which contribute to the **dinv** of the polyomino correspond each to a square in its area.

It remains to show that F is a bijection. To see this, we can consider the inverse function: given a parallelogram polyomino, write in weakly increasing order its area word, and draw it as a bounce path with labels. Then reading the relative positions of consecutive types of labels you can reconstruct piecewise both the upper and lower paths. This completes the proof.

Let us observe some remarkable consequences of this result. First of all, notice that iterating this bijection a second time, we get a bijection $F \circ F$

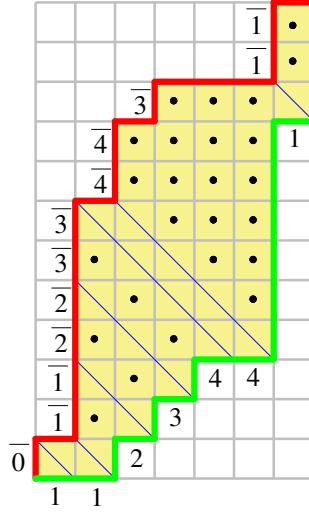


Figure 9: The outcome of applying F to the polyomino of Figure 1.

from $\text{Polyo}_{m,n}$ to itself which sends bounce to divv . Moreover, applying the inverse and composing it with the flip along the Southwest to Northeast line that pass through the Southwest corner (which obviously preserves the area) we get a bijection from $\text{Polyo}_{m,n}$ to itself which sends divv to area .

In conclusion, we see that all our three statistics are equidistributed both inside the same m times n rectangle and with the polyominoes in the flipped n times m rectangle.

5. Recursions for $\text{Nara}_{m,n}(q, t)$ and $\widetilde{\text{Nara}}_{n,m}(q, t)$

In this section we prove that both $\text{Nara}_{m,n}(q, t)$ and $\widetilde{\text{Nara}}_{n,m}(q, t)$ satisfy a certain recursion. As an immediate byproduct we get another proof of the identity $\text{Nara}_{m,n}(q, t) = \widetilde{\text{Nara}}_{n,m}(q, t)$ stated in Theorem 4.1.

Let $\widetilde{\text{Polyo}}_{m,n}^{(r,s)}$ be the set of polyominoes in $\text{Polyo}_{m,n}$ whose labelled bounce path has r many 1's and s many $\bar{1}$'s. In other words, r is the number of steps between the first and the second bounce of the bounce path, while s is the number of steps between the second and the third bounce. Define

$$\widetilde{\text{Nara}}_{m,n}^{(r,s)}(q, t) := \sum_{P \in \widetilde{\text{Polyo}}_{m,n}^{(r,s)}} t^{\text{bounce}(P)} q^{\text{area}(P)},$$

so that $\widetilde{\text{Nara}}_{m,n}(q, t)$ is the sum over all r and s of $\widetilde{\text{Nara}}_{m,n}^{(r,s)}(q, t)$. Also, we define the q -analogue of the non-negative integers by setting $[0]_q := 1$, and for all positive integers n ,

$$[n]_q := \frac{1 - q^n}{1 - q} = 1 + q + q^2 + \cdots + q^{n-1}.$$

We define the q -analogue of the factorial of a non-negative integer by setting $[0]_q! := 1$, and for all positive integers n ,

$$[n]_q! := \prod_{i=1}^n [i]_q.$$

Finally, for $0 \leq k \leq n$,

$$\begin{bmatrix} n \\ k \end{bmatrix}_q := \frac{[n]_q!}{[n-k]_q! [k]_q!}$$

denotes the q -analogue of the binomial $\binom{n}{k}$.

Theorem 5.1. *For all m, n, r and s such that $1 \leq r \leq n$ and $0 \leq s \leq m-1$, we have the recursion*

$$\widetilde{\text{Nara}}_{m,n}^{(r,s)}(q, t) = t^{m+n-1} q^{r+s} \sum_{h=1}^{n-r} \sum_{k=0}^{m-s-1} \begin{bmatrix} s+r-1 \\ s \end{bmatrix}_q \begin{bmatrix} s+h-1 \\ h \end{bmatrix}_q \widetilde{\text{Nara}}_{m-s,n-r}^{(h,k)}(q, t),$$

with initial conditions

$$\widetilde{\text{Nara}}_{m,n}^{(n,s)}(q, t) = \begin{cases} (qt)^{m+n-1} \begin{bmatrix} m+n-2 \\ m-1 \end{bmatrix}_q & \text{if } s = m-1 \\ 0 & \text{if } s < m-1, \end{cases}$$

and $\widetilde{\text{Nara}}_{1,n}^{(r,0)}(q, t) = 0$ for $r < n$.

Proof. The argument in this proof is best understood by referring to Figure 10.

The figure shows a typical element of $\widetilde{\text{Poly}}_{m,n}^{(r,s)}$. The orange grid cuts out an element of $\widetilde{\text{Poly}}_{m-s,n-r}^{(h,k)}$: its lower-left corner is placed at the beginning of the rightmost step of the bounce path labelled by $\bar{1}$.

Observe that the labels of the bounce path in the orange grid are the same as the labels of the corresponding small path all increased by 1. Hence,

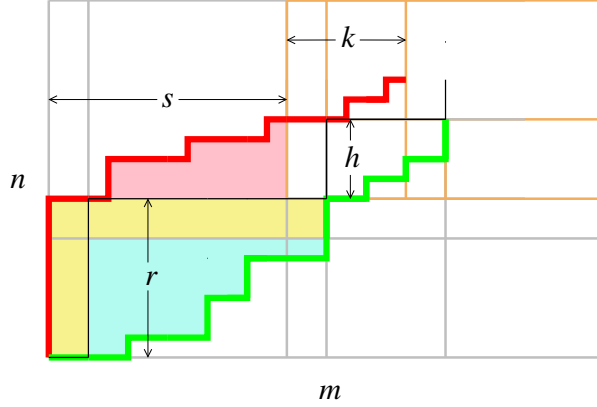


Figure 10: The black path is the bounce path.

together with the 1's and the $\bar{1}$'s of the bounce path outside of the orange grid, we see that the bounce of the larger polyomino is $m + n - 1$ more than the bounce of the small polyomino in the orange grid. This shift is taken care of by the factor t^{m+n-1} .

The area of the larger polyomino is equal to the area of the small polyomino in the orange grid plus the yellow area, which is taken care of by the factor q^{r+s} , the light blue area, which is counted by the factor $\begin{bmatrix} s+r-1 \\ s \end{bmatrix}_q$, and the pink area, which is counted by the factor $\begin{bmatrix} s+h-1 \\ h \end{bmatrix}_q$. This explains the recursion formula. \square

Let us denote by $\text{Polyo}_{n,m}^{(r,s)}$ the set of parallelogram polyominoes in an $n \times m$ rectangle whose area word has r many 1's and s many $\bar{1}$'s. Define

$$\text{Nara}_{n,m}^{(r,s)}(q, t) := \sum_{P \in \text{Polyo}_{n,m}^{(r,s)}} t^{\text{area}(P)} q^{\text{dinv}(P)},$$

so that $\text{Nara}_{n,m}(q, t)$ is the sum over all r and s of $\text{Nara}_{n,m}^{(r,s)}(q, t)$.

These polynomials satisfy the same recursion satisfied by the $\widetilde{\text{Nara}}_{m,n}^{(r,s)}(q, t)$'s:

Theorem 5.2. *For all m, n, r and s with $1 \leq r \leq n$ and $0 \leq s \leq m - 1$, we have the recursion*

$$\text{Nara}_{n,m}^{(r,s)}(q, t) = t^{m+n-1} q^r \sum_{h=1}^{n-r} \sum_{k=0}^{m-s-1} q^s \begin{bmatrix} s+r-1 \\ s \end{bmatrix}_q \begin{bmatrix} s+h-1 \\ h \end{bmatrix}_q \text{Nara}_{n-r, m-s}^{(h,k)}(q, t),$$

with initial conditions

$$\text{Nara}_{n,m}^{(n,s)}(q, t) = \begin{cases} (qt)^{m+n-1} \begin{bmatrix} m+n-2 \\ m-1 \end{bmatrix}_q & \text{if } s = m - 1 \\ 0 & \text{if } s < m - 1, \end{cases}$$

and $\text{Nara}_{n,1}^{(r,0)}(q, t) = 0$ for $r < n$.

Proof. Given an element of $\text{Polyo}_{n,m}^{(r,s)}$ with h many 2's and k many $\bar{2}$'s, we construct an element of $\text{Polyo}_{n-r,m-s}^{(h,k)}$ by subtracting 1 from all the letters in the area word, then removing all the resulting 0's and $\bar{0}$'s and replacing the only $\bar{-1}$ (which comes from the only $\bar{0}$) by $\bar{0}$.

For example, if we start with the word $\bar{0}1\bar{1}2\bar{2}\bar{1}2\bar{2}\bar{2}1\bar{1}\bar{1}$, which is an element of $\text{Polyo}_{6,7}^{(3,4)}$ with 3 many 2's and 2 many $\bar{2}$'s, then we first get $\bar{-1}0\bar{0}011\bar{0}1\bar{1}\bar{1}0\bar{0}\bar{0}$, and hence we finally get $\bar{0}111\bar{1}\bar{1}$, which is an element of $\text{Polyo}_{6-3,7-4}^{(3,2)} = \text{Polyo}_{3,3}^{(3,2)}$.

Now the area of this new element is clearly $m + n - 1$ less than the area of the original polyomino, since we subtracted 1 from all the letters of the area word different from $\bar{0}$. This is taken care of by the factor t^{m+n-1} .

The dinv of the original polyomino is equal to the dinv of this smaller polyomino, plus the dinv coming from the original $\bar{0}$ and the 1's, which is taken care of by the factor q^r , the dinv coming from the 1's and the $\bar{1}$'s, which is counted by the factor $q^s \begin{bmatrix} s+r-1 \\ s \end{bmatrix}_q$ (the 1's and the $\bar{1}$'s form a word which always starts with 1), and the dinv coming from the $\bar{1}$'s and the 2's, which is counted by the factor $\begin{bmatrix} s+h-1 \\ h \end{bmatrix}_q$ (as before, the $\bar{1}$'s and the 2's form a word which always starts with $\bar{1}$, but the dinv coming from this first letter is already counted by the $\bar{0}$ that we insert in the new area word!).

This explains the recursion. □

As already mentioned, these recursions give immediately $\text{Nara}_{n,m}^{(r,s)}(q, t) = \widetilde{\text{Nara}}_{m,n}^{(r,s)}(q, t)$, and hence another proof of the identity $\text{Nara}_{m,n}(q, t) = \widetilde{\text{Nara}}_{n,m}(q, t)$.

6. Symmetric functions interpretation

In this section we will use some tools from the theory of Macdonald polynomials. For a quick survey of what we need (and more), we refer to the book [2], in particular Chapters 3 and 9. In what follows we will recall only some basic facts, mostly to fix the notation.

Let $\Lambda = \bigoplus_{n \geq 0} \Lambda^n$ be the space of symmetric functions with coefficients in $\mathbb{C}(q, t)$, where q and t are variables, with its natural decomposition in components of homogeneous degree. Recall the fundamental bases of symmetric functions: *elementary* $\{e_\mu\}_{\mu \in \mathcal{P}}$, *homogeneous* $\{h_\mu\}_{\mu \in \mathcal{P}}$, *power* $\{p_\mu\}_{\mu \in \mathcal{P}}$, *monomial* $\{m_\mu\}_{\mu \in \mathcal{P}}$ and *Schur* $\{s_\mu\}_{\mu \in \mathcal{P}}$, where \mathcal{P} is the set of all partitions.

A scalar product is defined on Λ by declaring the Schur basis to be orthonormal:

$$\langle s_\lambda, s_\mu \rangle = \chi(\lambda = \mu),$$

where χ is the indicator function, which is 1 when its argument is true, and 0 otherwise. Another fundamental basis of Λ is $\{\tilde{H}_\mu\}_\mu$, the *modified Macdonald polynomial* basis.

The fundamental ingredient of the theory is the nabla operator ∇ acting on Λ . This is an homogeneous invertible operator introduced by Bergeron and Garsia in the study of the diagonal harmonics DH_n of \mathfrak{S}_n . In fact, it turns out that ∇e_n gives precisely the bigraded Frobenius characteristic of DH_n .

The so-called *shuffle conjecture* predicts a combinatorial interpretation of ∇e_n in terms of parking functions. Special cases of this conjecture have been proven by several authors. In particular, Haglund [5] proved the combinatorial interpretation of $\langle \nabla e_n, h_j h_{n-j} \rangle$ for $1 \leq j \leq n$ predicted by the shuffle conjecture.

Surprisingly, this same polynomial provides the symmetric functions interpretation of our q, t -Narayana numbers. More precisely, we have the following theorem, which is the main result of this paper.

Theorem 6.1. *For $m, n \geq 1$ we have*

$$\text{Nara}_{m,n}(q, t) = (qt)^{m+n-1} \cdot \langle \nabla e_{m+n-2}, h_{m-1} h_{n-1} \rangle.$$

Before proving this theorem, we give here an immediate corollary.

Theorem 6.2. *For all $m \geq 1$ and $n \geq 1$, we have*

$$\text{Nara}_{m,n}(q, t) = \text{Nara}_{m,n}(t, q)$$

and

$$\text{Nara}_{m,n}(q, t) = \text{Nara}_{n,m}(q, t).$$

Moreover, we have

$$\text{Nara}_{m,n}(q, t) = \widetilde{\text{Nara}}_{m,n}(q, t).$$

Proof of the Theorem 6.2. The symmetry in q and t comes from a general property of the nabla operator, which is an immediate consequence of the well-known identity [2, Equation (9.8)]: nabla applied to any Schur function is symmetric in q and t .

The second equation, symmetry in m and n , is obvious from the formula in Theorem 6.1. Finally, the fact that $\mathbf{Nara}_{m,n}(q, t) = \widehat{\mathbf{Nara}}_{m,n}(q, t)$ is a direct consequence of the symmetries and of Theorem 4.1. \square

6.1. Proof of Theorem 6.1

In order to prove Theorem 6.1, we need to make use of Haglund's combinatorial interpretation of $\langle \nabla e_{m+n-2}, h_{m-1} h_{n-1} \rangle$. To do this we first require some definitions.

For us a *Dyck path of order k* will be given by an *area word* which is a sequence of non-negative integers $b_1 b_2 \cdots b_k$ such that $b_1 = 0$, and $b_{i+1} \leq b_i + 1$ for all $1 \leq i < k$. A *domino* is a pair of values (a, b) written as the first above the second $\begin{array}{|c|} \hline a \\ \hline b \\ \hline \end{array}$.

A *parking function* PF of size k is a sequence of k dominoes $\begin{array}{|c|c|c|c|} \hline a_1 & a_2 & \cdots & a_k \\ \hline b_1 & b_2 & \cdots & b_k \\ \hline \end{array}$ such that $b_1 b_2 \cdots b_k$ is the area word of a Dyck path of order k , and the a_i 's are a permutation of the integers $\{1, \dots, k\}$ with the property $a_i < a_{i+1}$ if $b_i < b_{i+1}$ (and hence $b_i = b_{i+1} - 1$).

Example 6.1. PF = $\begin{array}{|c|c|c|c|c|c|c|c|c|c|} \hline 5 & 11 & 1 & 9 & 6 & 8 & 3 & 4 & 7 & 10 & 2 \\ \hline 0 & 1 & 1 & 2 & 0 & 1 & 0 & 1 & 2 & 3 & 3 \\ \hline \end{array}$ is a parking function of size 11.

Remark 6.1. Parking functions are often represented by a diagram like the one in Figure 11. In this diagram the red path represents the underlying Dyck path, where the number of the squares between the vertical steps of the Dyck path and the (green) diagonal are given by the lower numbers in the dominoes. The numbers that label the vertical steps of the Dyck path in the diagram are simply the upper numbers in the dominoes.

Given a parking function, we can reorder its dominoes by comparing first the bottom numbers, from the biggest to the smallest, and then, we place the dominoes with the same bottom number in order as we read them from right to left in the parking function.

The *reading word* $\sigma(\text{PF})$ associated to a parking function PF is the permutation that we obtain by reading the upper entries of this reordered sequence of dominoes.

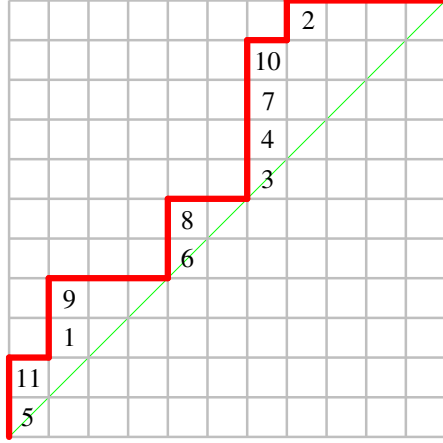


Figure 11: The parking function $\text{PF} = \begin{bmatrix} 5 & 11 & 1 & 9 & 6 & 8 & 3 & 4 & 7 & 10 & 2 \\ 0 & 1 & 1 & 2 & 0 & 1 & 0 & 1 & 2 & 3 & 3 \end{bmatrix}$.

Example 6.2. If $\text{PF} = \begin{bmatrix} 5 & 11 & 1 & 9 & 6 & 8 & 3 & 4 & 7 & 10 & 2 \\ 0 & 1 & 1 & 2 & 0 & 1 & 0 & 1 & 2 & 3 & 3 \end{bmatrix}$ then we reorder the dominoes as

$$\begin{bmatrix} 2 & 10 & 7 & 9 & 4 & 8 & 1 & 11 & 3 & 6 & 5 \\ 3 & 3 & 2 & 2 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \end{bmatrix},$$

and the corresponding reading word is $\sigma(\text{PF}) = 2\ 10\ 7\ 9\ 4\ 8\ 1\ 11\ 3\ 6\ 5$.

Given a parking function $\text{PF} = \begin{bmatrix} a_1 & a_2 & \cdots & a_k \\ b_1 & b_2 & \cdots & b_k \end{bmatrix}$, we define its *area* to be $\text{area}(\text{PF}) = b_1 + \dots + b_k$, and its *divv* $\text{divv}(\text{PF})$ as the number of pairs (i, j) with $1 \leq i < j \leq k$ such that either $b_i = b_j$ and $a_i < a_j$, or $b_i = b_j + 1$ and $a_i > a_j$. For example the area of the parking function of Example 6.2 is 14, while its *divv* is 8.

Given two disjoint sequences of numbers A and B , we denote by $A \sqcup B$ the set of *shuffles* of A and B , i.e. the sequences consisting of the numbers from $A \cup B$ in which all the elements of A and B appear in their original order, so that $|A \sqcup B| = \binom{|A|+|B|}{|A|}$.

For any a and b in \mathbb{N} , we call $\text{Park}_{a,b}$ the set of parking functions PF of size $a + b$ such that $\sigma(\text{PF}) \in (1, 2, \dots, a) \sqcup (a + 1, a + 2, \dots, a + b)$. Finally, we set

$$\text{Para}_{a,b}(q, t) := \sum_{\text{PF} \in \text{Park}_{a,b}} t^{\text{area}(\text{PF})} q^{\text{divv}(\text{PF})}.$$

We may state now the result of Haglund (see [5] for a proof, and [6] for the necessary background).

Theorem 6.3 (Haglund). *For all $m \geq 1$ and $n \geq 1$, we have*

$$\langle \nabla e_{m+n-2}, h_{m-1} h_{n-1} \rangle = \text{Para}_{n-1, m-1}(q, t).$$

This theorem reduces the problem of proving Theorem 6.1 to proving the following:

$$\text{Nara}_{m, n}(q, t) = (qt)^{m+n-1} \text{Para}_{n-1, m-1}(q, t). \quad (6.1)$$

In order to show the validity of this equation we do as follows. For $0 \leq r < n$, $0 \leq s < m$ with $r + s \geq 1$, let $\text{Park}_{n-1, m-1}^{(r, s)}$ be the set of parking functions PF of size $m + n - 2$ such that

$$\sigma(\text{PF}) \in A \sqcup B, \quad |D_0(\text{PF}) \cap A| = r, \quad \text{and} \quad |D_0(\text{PF}) \cap B| = s,$$

where $A = (1, 2, \dots, n-1)$, $B = (n, n+1, \dots, m+n-2)$, and $D_0(\text{PF})$ is the set of upper numbers of dominoes of PF whose bottom numbers equal 0. Define the polynomial

$$\text{Para}_{n-1, m-1}^{(r, s)}(q, t) := \sum_{\text{PF} \in \text{Park}_{n-1, m-1}^{(r, s)}} t^{\text{area}(\text{PF})} q^{\text{dinv}(\text{PF})},$$

and set $\text{Para}_{n-1, m-1}^{(0, 0)}(q, t) = 0$. Clearly the sum of all these polynomials as r and s range over their possible values is equal to $\text{Para}_{n-1, m-1}(q, t)$. With this new generalization in mind, it is clear that equation 6.1 holds true if

$$(qt)^{m+n-1} \text{Para}_{n-1, m-1}^{(r, s)}(q, t) = \text{Nara}_{m, n}^{(s+1, r)}(q, t).$$

Our proof of this identity is similar to what Haglund did in [5].

We will show that $(qt)^{m+n-1} \text{Para}_{n-1, m-1}^{(r, s)}(q, t)$ also satisfies the recursion in Theorem 5.2, with the 4-tuple (m, n, r, s) replaced with $(n, m, s+1, r)$, i.e.

$$\begin{aligned} & \text{Para}_{n-1, m-1}^{(r, s)}(q, t) = \\ & t^{m+n-r-s-2} \sum_{h=0}^{n-r-1} \sum_{k=1}^{m-s-1} \begin{bmatrix} r+s \\ r \end{bmatrix}_q \begin{bmatrix} r+k-1 \\ k \end{bmatrix}_q \text{Para}_{n-r-1, m-s-2}^{(h, k-1)}(q, t). \end{aligned} \quad (6.2)$$

The initial conditions are

$$\text{Para}_{n-1,m-1}^{(r,m-1)}(q,t) = \frac{\text{Nara}_{m,n}^{(m,r)}(q,t)}{(qt)^{m+n-1}} = \begin{cases} \begin{bmatrix} m+n-2 \\ n-1 \end{bmatrix}_q & \text{if } r = n-1, \\ 0 & \text{if } r < n-1, \end{cases}$$

and

$$\text{Para}_{0,m-1}^{(r,s)}(q,t) = \frac{\text{Nara}_{m,1}^{(s+1,0)}(q,t)}{(qt)^m} = 0 \text{ for } s < m-1.$$

In order to see how the recursion (6.2) works for parking functions, we first make a simplification. It follows immediately from the definitions that, since the reading word of the parking functions we are interested in is a shuffle of the sequences $A = (1, 2, \dots, n-1)$ and $B = (n, n+1, \dots, n+m-2)$, the pairs of dominoes with both upper numbers in A or both in B do not contribute to the dinv . The only pairs that contribute are the ones where one of the upper numbers is in A and the other is in B . Since all the elements of A are smaller than all the elements of B , we can simply consider dominoes in which the upper number is 1 (if the corresponding element was in A) or 2 (if the corresponding element was in B), with the dinv defined in the same way.

For example the parking function $\text{PF} \in \text{Park}_{9,9}^{(3,1)}$

3	13	6	15	8	7	16	12	5	14	9	2	11	1	10	4
0	1	1	2	2	2	3	1	1	2	0	0	1	0	1	1

whose reading word is

$$\sigma(\text{PF}) = 16 \ 14 \ 7 \ 8 \ 15 \ 4 \ 10 \ 11 \ 5 \ 12 \ 6 \ 13 \ 1 \ 2 \ 9 \ 3,$$

would correspond to

1	2	1	2	1	1	2	2	1	2	2	1	2	1	2	1
0	1	1	2	2	2	3	1	1	2	0	0	1	0	1	1

whose reading word is

$$\sigma(\text{PF}) = 2 \ 2 \ 1 \ 1 \ 2 \ 1 \ 2 \ 2 \ 1 \ 2 \ 1 \ 2 \ 1 \ 1 \ 2 \ 1.$$

In both cases the dinv is 32, and the area is 18.

Using this identification, we do as follows: given an element PF of $\text{Park}_{n-1,m-1}^{(r,s)}$, we remove the dominoes whose lower number is 0, and we decrease the lower number of the remaining dominoes by 1, keeping them in the given order.

In our example, applying this procedure to PF we get

2	1	2	1	1	2	2	1	2	2	2	1
0	0	1	1	1	2	0	0	1	0	0	0

In doing this, observe that we will always get a parking function which starts with a domino $\begin{bmatrix} 2 \\ 0 \end{bmatrix}$, which is always followed by a domino with lower entry 0, since PF cannot contain a sequence of three consecutive dominoes with strictly increasing lower numbers. This first domino contributes 0 to both area and dinv and we can therefore remove it. In doing this we get an element of $\text{Park}_{n-r-1, m-s-2}^{(h, k-1)}$, where h is the number of $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ dominoes in PF, and k is the number of $\begin{bmatrix} 2 \\ 1 \end{bmatrix}$ dominoes in PF.

Remark 6.2. Conversely, given an element of $\text{Park}_{n-r-1, m-s-2}^{(h, k-1)}$, we can prepend it with a $\begin{bmatrix} 2 \\ 0 \end{bmatrix}$ domino, increase all the lower numbers by 1, and then insert r $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ dominoes and s $\begin{bmatrix} 2 \\ 0 \end{bmatrix}$ dominoes. This gives us an element of $\text{Park}_{n-1, m-1}^{(r, s)}$. In doing so, we are forced to put a $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ in front of the first $\begin{bmatrix} 2 \\ 1 \end{bmatrix}$ which we just prepended. Other inserted dominoes must satisfy the following constraints: a $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ domino is followed by a domino in $\left\{ \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 2 \\ 0 \end{bmatrix}, \begin{bmatrix} 2 \\ 1 \end{bmatrix} \right\}$, if any, and a $\begin{bmatrix} 2 \\ 0 \end{bmatrix}$ domino is followed by a domino in $\left\{ \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 2 \\ 0 \end{bmatrix} \right\}$, if any.

Let us now look at how the area and the dinv change with respect to the operation that we have just described. The area of the new parking function is equal to the area of PF minus $(n-1-r) + (m-1-s)$, which is taken care of by the factor $t^{n-r+m-s-2}$ on the right hand side of (6.2).

The dinv is going to be the dinv of PF minus the dinv created by the dominoes that we have removed. First of all, there are the pairs of dominoes $\begin{bmatrix} 2 \\ 0 \end{bmatrix}$ and $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ in PF, whose relative position creates dinv : this dinv is taken care of by the factor $\begin{bmatrix} r+s \\ r \end{bmatrix}_q$ on the right hand side of (6.2). Then there is the dinv created by the dominoes $\begin{bmatrix} 2 \\ 1 \end{bmatrix}$ in PF with the dominoes $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ in PF: this is taken care by the factor $\begin{bmatrix} r+k-1 \\ k \end{bmatrix}_q$, since the first domino $\begin{bmatrix} 2 \\ 1 \end{bmatrix}$ is necessarily preceded by a $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$.

The initial conditions are obvious. This completes the proof of Theorem 6.1.

Remark. The link between parallelogram polyominoes and Macdonald polynomials seems to be deep as suggested by recent investigations on *labelled* parallelogram polyominoes [1], in which it appears that the Frobenius characteristic, valued by area, may be conjecturally expressed through a ∇ -like operator.

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