Decode-and-Forward Buffer-Aided Relay Selection

in Cognitive Relay Networks

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Abstract

This paper investigates decode-and-forward (DF) buffer-aided relay selection for underlay cognitive relay networks in the presence of both primary transmitter and receiver. We propose a novel buffer aided relay selection scheme for the cognitive relay network, where the best relay is selected with the highest signal-to-interference-ratio (SIR) among all available source-to-relay and relay-to-destination links while keeping the interference to the primary destination within a certain level. A closed-form expression for the outage probability of the proposed relay selection scheme is obtained. Both simulation and theoretical results are shown to confirm performance advantage over the conventional max-min relay selection scheme, making the proposed scheme attractive for cognitive relay networks.

Index Terms

Cognitive relay networks, relay selection, buffer-aided decode-and-forward relay

I. INTRODUCTION

Cognitive relay networks (CRNs) provide a promising way to exploit the advantages of both cognitive radio and cooperative relay networks [1]. While spectrum sharing in a CRN can be realized through various approaches including spectrum underlay, overlay and interweave [2], the underlay approach has most practical interest as the interference from the secondary users to the primary users is strictly limited. In a typical underlay CRN, beside the primary users, there are secondary users including secondary source, destination and a number of relay nodes. Relay selection provides an efficient way to achieve diversity gain in the CRN, because when only the best relay (rather than all relays) is selected for transmission, the

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interference to the primary users is also limited. The system with relay selection generally works in two phases: in the first phase, the source transmits data to the selected relay; in the second phase, the selected relay forwards the data to the secondary destination. In a CRN, the best relay is selected to maximize the transmission rate between the secondary and destination nodes while keeping the interference to the primary users within a pre-required level. In the first phase, in particular, because the transmission power from source to every potential relay is limited according to the same interference constraint at the primary users, the received signal-to-noise (SNR) at the relays becomes correlated [3]. This may imply that the relay selection process among all candidates is mutually dependent, and so full diversity cannot always be achieved even when all relevant channel coefficients are identically and independently distributed (i.i.d.). This is very different from the conventional relay selection where the best relay is usually selected among independent candidates.

Nonetheless, relay selection in CRNs has attracted much attention recently. In [4], a max-min like selection decode-and-forward (SDF) relay scheme was proposed for a CRN, and outage analysis was based on the assumption that there are multiple independent links between the secondary source and primary user. This, however, contradicts the assumption that there is only one secondary source and primary user in the system. Moreover, although this simplified the outage analysis since the relay selection process can be assumed (wrongly) to be independent, the analytical result is not accurate. Some early works (e.g. [5]) on CRNs also failed to consider the dependence in the relay selection. The dependency in cognitive relay selection was identified in [3], and a "half" selection decode-and-forward (SDF) relay scheme was proposed to break the dependency in the relay selection. In the first phase of this approach the source broadcasts data to all relays and only in the second phase applies the relay selection. A similar relay selection technique was also considered in [6] so that the outage performance could be analyzed. The "half" relay selection (e.g. [3], [6]) is not the most efficient in making use of the relays, and generates security risk, because all relays (rather than only the selected relay) are involved in transmission in the first phase. Alternatively, the relay selection dependency problem can be avoided by assuming the link between the secondary source and primary user is constant, but this only applies to some specialized systems such as when the secondary source and primary user have little mobility (e.g. [7]).

Most current relay selection approaches (including the aforementioned) are for CRNs with no primary transmitters. In practice, both a primary transmitter and receiver may be present [8], [9], for which the interference from the primary transmitter to the secondary users cannot be ignored. This motivates us

to investigate relay selection in a more general CRN with both a primary transmitter and receiver being present. It has been recently recognized that the performance of conventional relay selection can be further improved by relaxing the constraint that the best source-to-relay and relay-to-destination links for a packet transmission must be determined altogether. This is achieved by introducing a data buffer at the relay nodes [10]–[15]. Such buffer-aided relay selection is even more useful in the CRN: because now the best source-to-relay and relay-to-destination links are selected separately, the dependency in the conventional max-min based relay selection can thereby be de-correlated.

Of particular interest is the max-link relay selection where the best link is always selected with the highest signal-to-noise (SNR) among all available source-to-relay and relay-to-destination links [10]. In this paper, considering the interference from/to the primary users, we propose a so-called max-SIR-link relay selection scheme for the CRN, where the best relay is selected with the highest signal to the primary interference ratio at the corresponding receiving nodes while satisfying the interference constraint at the primary receivers. The main contributions of this paper are listed as follows:

- Proposing DF buffer-aided max-SIR-link relay selection in the underlay CRN. As the proposed relay selection only lets the selected relay join the transmission at any one time, it is more efficient at de-correlating the relay selection process than the aforementioned "half" relay selection (e.g. [3], [6]), an important issue in cognitive relay selection. To the best of our knowledge, this is the first relay selection scheme for a CRN with both primary transmitter and receiver available.
- Deriving the closed-form expression of the outage probability for the proposed relay selection scheme. With the presence of both the primary transmitter and receiver, the analysis is much more involved than those for both the conventional and the existing cognitive relay selection schemes. The analysis not only provides deep insight in understanding the proposed scheme but also shows a potential approach to analyze similar systems in the future.

We next introduce the proposed relay selection scheme in the context of a CRN.

II. MAX-SIR-LINK RELAY SELECTION

A. System model

The cognitive relay network with buffers at the relays is shown in Fig. 1, where there is one secondary source node (SS), one secondary destination node (SD), one primary source (PS), one primary destination (PD) and the number of DF relays $SR_k, k \in (1, 2, ..., K)$. All nodes are half-duplex and do not transmit

and receive simultaneously. Each relay is equipped with a data buffer Q_k $(1 \le k \le K)$ of finite size L (in the number of data packets), and the data packets in the buffer follow the "first-in first-out" rule. For simplicity of exposition, we assume no direct link between the secondary source (SS) and the secondary destination (SD) as path loss or shadowing is assumed to render it unusable [16].



Fig. 1. The system model of the CRN within buffered relay selection.

All channels in Fig. 1 related to secondary transmission can be divided into three groups: secondary transmission channels for $SS \to SR_k$ and $SR_k \to SD$ with channel coefficients as $h_{sr_k}(t)$ and $h_{r_kd}(t)$, secondary interfering channels for $PS \to SR_k$ and $PS \to SD$ with coefficients as $h_{pr_k}(t)$, $h_{pd}(t)$, and primary interfering channels for $SS \to PD$ and $SR_k \to PD$ with coefficients as $h_{sp}(t)$ and $h_{r_kp}(t)$ respectively. The instantaneous and average channel gains are defined as $\gamma_{ab}(t) = |h_{ab}(t)|^2$ and $\lambda_{ab} =$ $E|h_{ab}(t)|^2$ respectively, where $ab \in \{sr_k, sp, r_kd, r_kp, pr_k, pd, pp\}$. For convenience in development, the time index t is ignored in the rest of the paper unless necessary.

We assume all channels are quasi-static Rayleigh fading so that the channel coefficients remain unchanged during one packet duration but independently vary from one packet time to another. We also assume that channels within every group are i.i.d. fading, but channels for different groups may have different average gains, or we have $\lambda_{sr_k} = \lambda_{r_kd}$, $\lambda_{pr_k} = \lambda_{pd}$ and $\lambda_{sp} = \lambda_{r_kp}$ for all k. This is a more practical assumption than those in many existing approaches where all channels are assumed to be i.i.d. fading (e.g. [?], [6], [10]). Exact knowledge of all instantaneous channels is assumed to be available at the secondary relay and destination nodes¹. All channel noises are assumed to be zero mean

¹The CSI is usually estimated through pilots and feedback (e.g. [17]), and the CSI estimation without feedback may also be applied (e.g [18]). Further detail of the CSI estimation is beyond the scope of this paper.

additive-white-Gaussian-noise (AWGN).

In the underlay cognitive system, the secondary transmission nodes including SS and SR_k are only allowed to share the spectrum with the primary user PD if the corresponding interfering power to PDis below a pre-defined level I_{th} , so that we have

$$P_{ss}\gamma_{sp} \le I_{th}$$
 and $P_{sr_k}\gamma_{r_kp} \le I_{th}, \quad k = 1, \cdots, K,$ (1)

where P_{ss} and P_{sr_k} are the transmission powers for SS and SR_k respectively.

If the relay SR_k is selected to receive data from the secondary source SS, due to the interference from the primary source PS, the received signal at SR_k is given by

$$\mathbf{y}_{sr_k} = \sqrt{P_{ss}} h_{sr_k} \mathbf{s} + h_{pr_k} \sqrt{P_{ps}} \mathbf{s}' + \mathbf{n}_{r_k}, \tag{2}$$

where s and s' are transmission vectors from SS and PS respectively, P_{ps} is the transmission power of the primary source which is assumed to be unity without losing generality and \mathbf{n}_{r_k} is the noise vector at SR_k . From (2), and with the power constraint as in (1), the received SIR at SR_k is obtained as

$$\operatorname{SIR}_{sr_k} = \frac{P_{ss}|h_{sr_k}|^2}{P_{ps}|h_{pr_k}|^2} = \frac{I_{th}\gamma_{sr_k}}{\gamma_{sp}\gamma_{pr_k}}.$$
(3)

As in [9], we focus on the interference-limited scenario wherein the interference power from the primary source is dominant relative to the noise so that the noise effects can be ignored. Therefore the instantaneous capacity for $SS \rightarrow SR_k$ is approximated as $C_{sr_k} \approx (1/2)\log_2(1 + SIR_{sr_k})$.

On the other hand, if the relay SR_k is selected to forward data to the secondary destination SD, the received signal at SD is given by

$$\mathbf{y}_{r_k d} = \sqrt{P_{sr_k}} h_{r_k d} \mathbf{s} + h_{pd} \sqrt{P_{ps}} \mathbf{s}' + \mathbf{n}_d, \tag{4}$$

where \mathbf{n}_d is the noise vector at SD. From (4) and (1), the SIR at the secondary destination is obtained as

$$SIR_{r_kd} = \frac{P_{sr_k} |h_{r_kd}|^2}{P_{ps} |h_{pd}|^2} = \frac{I_{th} \gamma_{r_kd}}{\gamma_{r_kp} \gamma_{pd}}.$$
(5)

Similarly, with the interference dominating the noise, the instantaneous capacity for $SR_k \rightarrow SD$ is approximated as $C_{r_kd} \approx (1/2)\log_2(1 + SIR_{r_kd})$.

B. Selection rule

In the max-SIR-link relay selection for the CRN, at any time, the best transmission link with the highest SIR is selected among all available source-to-relay and relay-to-destination links. A source-to-relay or a relay-to-destination link is considered available when the buffer of the corresponding relay node is not full nor empty respectively. To be specific, if a source-to-relay link is selected, the source node transmits one data packet to the corresponding relay node. If the selected relay can successfully decode the data, the decoded packet is stored in the buffer and the number of data packets in the buffer is increased by one. On the other hand, if a relay-to-source link is selected, the corresponding relay transmits the earliest stored packet in the buffer to the destination. If the destination can successfully decode the packet, the number of packets in the buffer is decreased by one.

The best selected relay (either for transmission or reception) in the max-SIR-link scheme can be obtained as $R_{best} = \arg \max_{SR_k} \{SIR_{sr_k}, SIR_{r_kd}\}$. While the SIRs for the source-to-relay and relay-to-destination link are given by (3) and (5) respectively, we have

$$R_{best} = \arg \max_{SR_k} \left\{ \frac{\max_{SR_k:\Psi(Q_k)\neq L} \{\frac{I_{th}\gamma_{sr_k}}{\gamma_{pr_k}}\}}{\gamma_{sp}} , \frac{\max_{SR_k:\Psi(Q_k)\neq L} \{\frac{I_{th}\gamma_{r_kd}}{\gamma_{r_kp}}\}}{\gamma_{pd}} \right\},$$
(6)

where $\Psi(Q_k)$ gives the number of data packets in the buffer Q_k .

The outage probability can be defined as the probability that the selected link is in outage as

$$P_{out} \triangleq \begin{cases} \mathbb{P}\{(1/2)\log_2(1 + \operatorname{SIR}_{sr_k}) < C_{th}\} & \text{for relay reception,} \\ \mathbb{P}\{(1/2)\log_2(1 + \operatorname{SIR}_{r_kd}) < C_{th}\} & \text{for destination reception,} \end{cases}$$
(7)

where C_{th} is the target rate, and the factor 1/2 captures the fact that it takes two time slots to transmit any packet from the source to the destination. Next, we perform the outage probability analysis.

III. OUTAGE PROBABILITY ANALYSIS

This section analyzes the outage probability of the max-SIR-link relay selection in the CRN. At any time, the numbers of data packets in every buffer form a "state". Because there are K available relays and every relay is equipped with a buffer of size L, there are $(L+1)^K$ states in total. The *l*-th state vector is defined as

$$\mathbf{s}_{l} = [\Psi_{l}(Q_{1}), \cdots, \Psi_{l}(Q_{K})]^{\mathrm{T}}, \qquad l = 1, \cdots, (L+1)^{K}$$
(8)

where $\Psi_l(Q_k)$ gives the number of data packets in buffer Q_k at state s_l . It is clear that $0 \leq \Psi_l(Q_k) \leq L$.

We assume that state s_l corresponds to the pair of (K_1, K_2) , where K_1 and K_2 are the numbers of available links for source-to-relay and relay-to-destination transmission at state s_l respectively. By considering all possible available links for K_1 and K_2 , the outage probability of the overall system can be obtained as

$$P_{out} = \sum_{l=1}^{(L+1)^K} \pi_l \overline{p}_{s_l}^{(K_1, K_2)},\tag{9}$$

where $\overline{p}_{s_l}^{(K_1,K_2)}$ is the outage probability when the state is at s_l , and π_l is the stationary probability for the state s_l . The following two sub-sections show the calculation of $\overline{p}_{s_l}^{(K_1,K_2)}$ and π_l respectively.

A. $\overline{p}_{s_l}^{(K_1,K_2)}$: outage probability for state s_l

According to (6) and the theory of order statistics [19], if there are K_1 source-to-relay links available, the cumulative distribution function (CDF) of $X_1 = \max_{SR_k:\Psi(Q_k)\neq L} \{\frac{\gamma_{sr_k}}{\gamma_{pr_k}}\}$ is given by

$$F_{X_1}(x) = \left(\frac{x}{L_1 + x}\right)^{K_1},\tag{10}$$

where $L_1 = \frac{I_{th}\lambda_{sr_k}}{\lambda_{pr_k}}$. Then the CDF of $X = X_1/\gamma_{sp}$ is given by

$$F_{X}(x) = \begin{cases} 1, & \text{if } K_{1} = 0, \\ 1 - \frac{L_{1}}{\lambda_{sp}x} e^{\frac{L_{1}}{\lambda_{sp}x}} \text{Ei}(1, \frac{L_{1}}{\lambda_{sp}x}), & \text{if } K_{1} = 1, \\ \left(\frac{\lambda_{sp}x}{L_{1}}\right)^{K_{1}-1} \frac{\mathcal{MG}\left([[0], [\]], [[K_{1}-1, K_{1}], [\]], \frac{L_{1}}{\lambda_{sp}x}\right)}{\Gamma(K_{1})}, & \text{elsewhere,} \end{cases}$$
(11)

where $\operatorname{Ei}(1, a) = \int_{1}^{\infty} \frac{\exp(-ta)}{a} dt, a > 0, \ \Gamma(\bullet)$ is the Gamma function, and $\mathcal{MG}([[], []], [[\bullet, \bullet], []], \bullet)$ is the Meijer G function [20].

Proof: See Appendix I.

Similarly, the CDF of $Y = \frac{\max_{SR_k:\Psi(Q_k)\neq L} \{\frac{I_{th}\gamma_{r_kd}}{\gamma_{r_kp}}\}}{\gamma_{pd}}$ is given by

$$F_{Y}(y) = \begin{cases} 1, & \text{if } K_{2} = 0, \\ 1 - \frac{L_{2}}{\lambda_{pd}y} e^{\frac{L_{2}}{\lambda_{pd}y}} \text{Ei}(1, \frac{L_{2}}{\lambda_{pd}y}), & \text{if } K_{2} = 1, \\ \left(\frac{\lambda_{pd}y}{L_{2}}\right)^{K_{2}-1} \frac{\mathcal{MG}\left([[0], [\]], [[K_{2}-1, K_{2}], [\]], \frac{L_{2}}{\lambda_{pd}y}\right)}{\Gamma(K_{2})}, & \text{elsewhere.} \end{cases}$$
(12)

Because X and Y are independent, the CDF of $Z = \max(X, Y)$ is obtained as $F_Z(z) = F_X(z)F_Y(z)$. It is then from (6) and (7) that

$$\overline{p}_{s_l}^{(K_1, K_2)} = F_Z(\gamma_{th}) = F_X(\gamma_{th}) F_Y(\gamma_{th}), \tag{13}$$

where $\gamma_{th} = 2^{2C_{th}-1}$, and C_{th} is defined in (7) which is the target data rate.

B. π_l : stationary distribution probability for state s_l

The Markov chain can be used to model the transitions between the buffer states. Suppose at time t, the state is at s_l . At time t + 1, if the received data can be successfully decoded, there must be one relay either receiving or transmitting a data packet, so that the number of packets in the corresponding buffer is increased or decreased by one respectively. Depending on which relay receives or transmits data, at time t + 1, the buffers may move from state s_l to several possible states. We assume the set U_l contains all states which can be reached from s_l in one step.

Because the channels within secondary transmission, secondary interfering and primary interfering groups are i.i.d. fading, it is clear from (3) and (5) that the SIRs for all channels are i.i.d. so that the probability to select any link is $1/(K_1 + K_2)$. Further noting that the state remains unchanged if outage occurs (or the decoding is not successful), the probabilities that the state s_l moves to a state in U_l is given by

$$p_{s_l} = \frac{1 - \overline{p}_{s_l}^{(K_1, K_2)}}{K_1 + K_2}.$$
(14)

We denote **A** as the $(L+1)^K \times (L+1)^K$ state transition matrix, where the entry $\mathbf{A}_{n,l} = P(X_{t+1} = s_n | X_t = s_l)$ which is the transition probability to move from state s_l at time t to state s_n at time (t+1). With the above analysis, we have

$$\mathbf{A}_{n,l} = \begin{cases} \overline{p}_{s_l}^{(K_1, K_2)}, & \text{if } s_n \notin U_l, \\ p_{s_l}, & \text{if } s_n \in U_l, \\ 0, & \text{elsewhere}, \end{cases}$$
(15)

Because the transition matrix A is column stochastic, irreducible and aperiodic², the stationary state

²Column stochastic means all entries in any column sum up to one, *irreducible* means that it is possible to move from any state to any state, and *aperiodic* means that it is possible to return to the same state at any steps [21], [22].

probability vector is obtained as (see [22], [23])

$$\boldsymbol{\pi} = (\mathbf{A} - \mathbf{I} + \mathbf{B})^{-1} \mathbf{b},\tag{16}$$

where $\pi = [\pi_1, \cdots, \pi_{(L+1)^K}]^T$, $\mathbf{b} = (1, 1, ..., 1)^T$, **I** is identity matrix and $\mathbf{B}_{n,l} = 1, \forall n, l$.

Finally, from (9), the outage probability for the max-SIR-link scheme is given by

$$P_{out} = \operatorname{diag}(\mathbf{A})\boldsymbol{\pi}.\tag{17}$$

Particularly, if the relay buffer size $L \to \infty$, similar to that in [10], it can be shown that probabilities for $K_1 = K$ and $K_2 = K$ are one. Thus we have

$$\lim_{L \to \infty} P_{out} = \left(\frac{\lambda_{sp}\lambda_{pd}z^2}{L_1L_2}\right)^{K-1} \frac{\mathcal{MG}\left([[0], [\]], [[K-1, K], [\]], \frac{L_1}{\lambda_{sp}z}\right) \mathcal{MG}\left([[0], [\]], [[K-1, K], [\]], \frac{L_2}{\lambda_{pd}z}\right)}{\Gamma^2(K)}.$$
(18)

These results are next verified by Monte Carlo simulations.

IV. NUMERICAL SIMULATIONS

In the simulations below, the pre-defined level $I_{th} = 1$, and the average channel gains are set as $\lambda_{sr_k} = \lambda_{r_kd} = 30$ dB, $\lambda_{sp} = \lambda_{r_kp} = 10$ dB and $\lambda_{pr_k} = \lambda_{pd} = 10$ dB. The transmission powers of the primary transmitter and channel noise are normalized to unity.



Fig. 2. Theoretical and simulation outage probability vs target rate for the proposed max-SIR-link relay selection.

Fig. 2 verifies the theoretical analysis for the proposed max-SIR-link scheme with simulations. We have

performed extensive simulations with different number of relays and buffer sizes. While all simulation results match the theoretical analysis, only a few are shown in Fig. 2 for better illustration. It is clearly shown that the outage probability decreases as the number of relays and buffer size increases. For example, for target rate $\gamma_{th} = 0.5$ bits per channel use (BPCU), when the number of relays and buffers (K, L)increase from (2, 2) to (5, 5), the outage probability drops about 40 dB. It is not surprising that higher diversity is obtained with more relays and higher coding gain is obtained with larger buffer size. For better illustration, only theoretical results for the proposed scheme are shown in the following simulation.

Fig. 3 compares the outage probabilities of the proposed max-SIR-link, conventional max-min and no relay selection schemes, where the number of relays is set as K = 3, different relay buffer sizes for the proposed approach are applied which are set as $L = 1, 5, 50, \infty$, respectively. It is clearly shown that the proposed relay selection (even with L = 1) has significantly better outage performance than the conventional max-min scheme, while both relay selection schemes are superior to the no-relay scheme in outage performance. Fig. 3 also shows that, for the proposed approach, the outage performance improves with larger buffer size, but the improvement becomes less significant when the buffer size is large enough. Particularly with L = 50, the outage performance is almost the same as that for $L \to \infty$.



Fig. 3. Outage probability comparison for the proposed max-SIR-link, conventional max-min and no-relay selection schemes.

V. CONCLUSIONS

This paper proposed DF buffer-aided max-SIR-link relay selection for an underlay CRN, in the presence of both primary source and destination. In the proposed scheme, the best relay corresponds to the highest SIR among all available source-to-relay and relay-to-destination links while keeping the interference at the primary user within a pre-defined level. The closed-form expression of the outage probability of the proposed scheme was obtained, which matches exactly the simulation results. Both theoretical and simulation results showed that the proposed scheme has significantly better outage performance than the conventional max-min scheme, making it an attractive scheme in a CRN.

APPENDIX I - PROOF OF (11)

Proof: From (6), we define $X = \frac{\max_{SR_k:\Psi(Q_k)\neq L} \{\frac{I_{th}\gamma_{sr_k}}{\gamma_{pr_k}}\}}{\gamma_{sp}}$ and $Y = \frac{\max_{SR_k:\Psi(Q_k)\neq L} \{\frac{I_{th}\gamma_{r_kd}}{\gamma_{r_kp}}\}}{\gamma_{pd}}$. The CDF of $X_1 = \max_{SR_k:\Psi(Q_k)\neq L} \{\frac{\gamma_{sr_k}}{\gamma_{pr_k}}\}$ is obtained as in (10). The PDF of exponentially distributed γ_{sp} is given by $f_{\gamma_{sp}}(\gamma) = (1/\gamma_{sp})e^{-\gamma/\gamma_{sp}}$.

Because X_1 and γ_{sp} are independent, the CDF of $X = X_1/\gamma_{sp}$ is obtained as

$$F_X(x) = \int_0^\infty \left(\frac{x\gamma}{L_1 + x\gamma}\right)^{K_1} \frac{1}{\gamma_{sp}} e^{-\frac{\gamma}{\gamma_{sp}}} d\gamma.$$
(19)

For (19), if $K_1 = 0$, we have

$$F_X(x) = \int_0^\infty \frac{1}{\gamma_{sp}} e^{-\frac{\gamma}{\gamma_{sp}}} d\gamma = 1,$$
(20)

if $K_1 = 1$, we have

$$F_X(x) = \int_0^\infty \left(\frac{x\gamma}{L_1 + x\gamma}\right) \frac{1}{\gamma_{sp}} e^{-\frac{\gamma}{\gamma_{sp}}} d\gamma = 1 - \frac{L_1}{\lambda_{sp}x} e^{\frac{L_1}{\lambda_{sp}x}} \operatorname{Ei}(1, \frac{L_1}{\lambda_{sp}x}),$$
(21)

and if $K_1 > 1$, we have

$$F_X(x) = \int_0^\infty \left(\frac{x\gamma}{L_1 + x\gamma}\right)^{K_1} \frac{1}{\gamma_{sp}} e^{-\frac{\gamma}{\gamma_{sp}}} d\gamma = \left(\frac{\lambda_{sp}x}{L_1}\right)^{K_1 - 1} \frac{\mathcal{MG}\left([[0], [\]], [[K_1 - 1, K_1], [\]], \frac{L_1}{\lambda_{sp}x}\right)}{\Gamma(K_1)}.$$
 (22)

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