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Programmable Logic Device with an 8-stage cascade of 64K-bit Asynchronous SRAMs

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The first implementation of a new programmable logic device using LUT(Look-Up Table) cascade architecture is developed in 0.35um CMOS logic process. Eight 64Kb asynchronous SRAMs are simply connected to form an LUT cascade with a few additional circuits. Benchmark results show that it has a competitive performance to FPGAs.

1. Introduction

RAMs and PLAs (Programmable Logic Array) are used for programmable logic devices that realize multiple-output combinational logic functions. However, when the number of inputs and/or outputs for the target function is large, these devices require excessive amounts of hardware. Thus, FPGAs (Field Programmable Gate Arrays) are often used. However, the area and delay for interconnections among logic cells are much larger than those for logic elements. Therefore the prediction of the performance of the FPGA is difficult without complete physical design. To solve these problems, an LUT cascade architecture that is composed of a serial connection of large-scale memories has been developed [1]. The LUT cascade uses relatively larger LUTs (1kb - 1Mb), and the interconnections between LUTs are limited to the adjacent cells in the cascade. This is quite different from the two-dimensional structure of FPGAs with smaller (16b - 64b) LUTs. The large area for the interconnections in an FPGA is absorbed in the larger LUTs in the cascade. So, the cascade is re-configured by only changing the contents of the LUTs.

2. Design of LUT Cascade LSI

An LUT block is mainly composed of an asynchronous 64kbit SRAM with 13bit address inputs and 8bit data I/Os. Since the memory should be operated as a data-path in LUT cascade architecture, an asynchronous SRAM is employed. In our design, each LUT block has 17 inputs: 8 bits are connected to the outputs of the preceding LUT, and 9 bits are from the external inputs (X). Then 13 of the 17 inputs are selected by crossover switches to form the actual address inputs of an LUT, and the unselected 4 signals can be used as intermediate outputs through Y terminals. This configuration enables to extend the number of allowable input/output signals from 48/8 (without crossover switches) to 76/36 (with crossover switches) for an 8-LUT cascade. This achieves an increase in the number of functions realizable with the LUT cascade.

The LUT cascade LSI is simply realized by a cascade connection of LUT blocks. Each LUT block also has connections to the common address lines used for programming and testing. When a TEST (Program) signal is High, an LUT is selected by the block select signals (BS) and all address inputs of the LUT can directly be controlled by the external address inputs (X and ADDL). In this mode, this chip is compatible with a conventional memory.

In order to realize a variety of functions and high-frequency operation, our design supports both parallel and pipeline

operations. The cascade can be split into two to operate 4 upper and 4 lower LUTs in parallel. In this case, we can use this device as a dual 4-LUT cascade. Since asynchronous SRAM is used as memory, by only providing address registers, an eight-stage pipeline operation can be realized, where each LUT corresponds to one stage of pipeline.

3. Measurement Result

The LUT cascade LSI was developed by 0.35um CMOS logic process. Although it looks like a conventional large-scale memory, the layout is point-symmetry as opposed to conventional line-symmetry. We verified that conversion of a simple memory into the LUT cascade requires small additional circuits, where the area overhead is +0.4% and the transistor count overhead is +0.2%. The latency of an internal LUT is 3.8ns. A total latency of 11.6ns for a 2-LUT cascade, 34.4ns for an 8-LUT cascade in asynchronous operation, and the operating frequency of 200MHz in an 8-stage pipeline operation were experimentally confirmed.

4. Performance Comparison with FPGA

To compare the performance (area, delay, power) of LUT cascades with FPGAs, we mapped simple benchmark functions to the LUT cascade and a commercial FPGA (Xilinx XCV50: 0.22um, 5-Layer metal, 2.5V, 384 CLBs) [2][3]. We used commercial logic synthesis and layout tools for the design of FPGA. On the other hand, for the design of LUT cascade, we used our newly developed logic synthesis tool that converts BDDs (Binary Decision Diagrams) into LUT cascades [1]. By taking the difference of process technology into account, we can conclude that, by using the same process technology as the FPGA for the LUT cascade, we can achieve a comparable layout area with less delay time and less power dissipation.

5. Conclusion

We first implemented the LUT cascade in CMOS process, and experimentally confirmed its competitive performance to FPGAs. The appearance of the chip is quite similar to the conventional memory. The design cost is much lower than FPGAs, and we can apply the test and redundancy methodologies for memories. The LUT cascade LSI is a new and promising reconfigurable logic device for future sub-100nm LSIs.

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