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8-8-2019

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Matthew C. Fickus

*Air Force Institute of Technology*

John Jasper

*University of Cincinnati*

Dustin G. Mixon

*Air Force Institute of Technology*

Jesse D. Peterson

*Air Force Institute of Technology*

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#### Recommended Citation

Fickus, M., Jasper, J., Mixon, D. G., & Peterson, J. D. (2019). Hadamard equiangular tight frames. *Applied and Computational Harmonic Analysis*. <https://doi.org/10.1016/j.acha.2019.08.003>

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# Hadamard Equiangular Tight Frames

Matthew Fickus<sup>a</sup>, John Jasper<sup>b</sup>, Dustin G. Mixon<sup>a</sup>, Jesse D. Peterson<sup>a</sup>

<sup>a</sup>*Department of Mathematics and Statistics, Air Force Institute of Technology, Wright-Patterson AFB, OH 45433*

<sup>b</sup>*Department of Mathematical Sciences, University of Cincinnati, Cincinnati, OH 45221*

## Abstract

An equiangular tight frame (ETF) is a type of optimal packing of lines in Euclidean space. They are often represented as the columns of a short, fat matrix. In certain applications we want this matrix to be flat, that is, have the property that all of its entries have modulus one. In particular, real flat ETFs are equivalent to self-complementary binary codes that achieve the Grey-Rankin bound. Some flat ETFs are (complex) Hadamard ETFs, meaning they arise by extracting rows from a (complex) Hadamard matrix. These include harmonic ETFs, which are obtained by extracting the rows of a character table that correspond to a difference set in the underlying finite abelian group. In this paper, we give some new results about flat ETFs. One of these results gives an explicit Naimark complement for all Steiner ETFs, which in turn implies that all Kirkman ETFs are possibly-complex Hadamard ETFs. This in particular produces a new infinite family of real flat ETFs. Another result establishes an equivalence between real flat ETFs and certain types of quasi-symmetric designs, resulting in a new infinite family of such designs.

*Keywords:* Hadamard, flat, equiangular, tight, frame

*2010 MSC:* 42C15

## 1. Introduction

An equiangular tight frame is a type of optimal packing of lines in Euclidean space. To be precise, if  $\{\varphi_j\}_{j=1}^n$  is any sequence of nonzero, equal-norm vectors in  $\mathbb{F}^d$  where  $n \geq d$  and  $\mathbb{F}$  is either  $\mathbb{R}$  or  $\mathbb{C}$ , the *coherence* of  $\{\varphi_j\}_{j=1}^n$  is bounded below by the *Welch bound* [35]:

$$\max_{j \neq j'} \frac{|\langle \varphi_j, \varphi_{j'} \rangle|}{\|\varphi_j\| \|\varphi_{j'}\|} \geq \left[ \frac{n-d}{d(n-1)} \right]^{\frac{1}{2}}. \quad (1)$$

In the case where  $\mathbb{F} = \mathbb{R}$ , each vector  $\varphi_j$  spans a real line, and the coherence is the cosine of the smallest interior angle between any pair of these lines. In this case, if equality in (1) is achieved then this smallest pairwise angle is as large as possible, meaning the lines are optimally packed.

For any vectors  $\{\varphi_j\}_{j=1}^n$  in  $\mathbb{F}^d$ , the corresponding *synthesis operator*  $\Phi$  is the  $d \times n$  matrix whose  $j$ th column is  $\varphi_j$ . It is well known [31] that a sequence of nonzero equal-norm vectors  $\{\varphi_j\}_{j=1}^n$  in  $\mathbb{F}^d$  achieves equality in (1) if and only if they form an *equiangular tight frame* (ETF) for  $\mathbb{F}^d$ , that is, if and only if the rows of its synthesis operator  $\Phi$  are equal-norm and orthogonal (tightness) while

*Email address:* Matthew.Fickus@gmail.com (Matthew Fickus)

$|\langle \varphi_j, \varphi_{j'} \rangle|$  is constant over all  $j \neq j'$  (equiangularity). These conditions are restrictive, and ETFs are not easy to find. That said, a growing number of explicit constructions of them are known [19]. Real ETFs in particular are known to be equivalent to a special class of *strongly regular graphs* (SRGs) [31, 23, 34]. Beyond these, the most popular ETFs are *harmonic ETFs* [31, 36, 12], which are obtained by restricting the characters of a finite abelian group  $\mathcal{G}$  to a *difference set*, namely a subset  $\mathcal{D}$  of  $\mathcal{G}$  with the property that the cardinality of  $\{(i, i') \in \mathcal{D} \times \mathcal{D} : g = i - i'\}$  is constant over all nonzero  $g \in \mathcal{G}$ . For example,  $\{0001, 0101, 0010, 1010, 0011, 1111\}$  is a difference set in  $\mathbb{Z}_2^4$ , and the corresponding six rows of the canonical Hadamard matrix of size 16 yields the following matrix whose columns give an optimal packing of 16 lines in  $\mathbb{R}^6$ :

$$\Phi = \begin{bmatrix} 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 \\ 1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 \\ 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 \\ 1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & 1 & -1 & -1 & 1 \end{bmatrix}. \quad (2)$$

Every harmonic ETF is *unital*, meaning each entry of its synthesis operator  $\Phi$  is a root of unity; number-theoretic necessary conditions on the existence of unital ETFs are given in [32]. Every harmonic ETF also has the following two special properties:

**Definition 1.** An ETF  $\{\varphi_j\}_{j=1}^n$  for  $\mathbb{F}^d$  is (*complex*) *Hadamard* if its  $d \times n$  synthesis operator  $\Phi$  is a submatrix of an  $n \times n$  (complex) Hadamard matrix, and is *flat* if each entry of  $\Phi$  is unit modulus.

In this paper, we study Hadamard ETFs and flat ETFs in general. Flat ETFs in particular arise in several applications, as detailed below. Because of this, we would like to have many ways to construct them. However, to date we only have two ways: besides harmonic ETFs, the only known flat ETFs are *Kirkman ETFs* [24], which are a special class of *Steiner ETFs* [20]. In this paper, we give some new results about flat ETFs which we hope will better inform future searches for them. In particular, we show that every Kirkman ETF is a possibly-complex Hadamard ETF (Theorem 2) and give a new characterization of all real flat ETFs (Theorem 3).

Flat ETFs are especially attractive for certain applications involving coding theory and waveform design. In particular, real flat ETFs are equivalent to self-complementary binary codes whose minimum-pairwise-Hamming distance is as large as possible, achieving the *Grey-Rankin bound* [24]. More generally, ETFs have been proposed as waveforms for wireless communication [31]. In that setting, flat ETFs allow the transmitted signals to have maximal energy (2-norm) subject to real-world bounds on transmitter power ( $\infty$ -norm), while still interfering with each other as little as possible. This same rationale is part of the reason why Hadamard matrices are used in traditional CDMA, and why *constant amplitude zero-autocorrelation* waveforms have been proposed as radar waveforms [2]. It also helped motivate the investigation into unital ETFs given in [32]; a real ETF is flat if and only if it is unital. Recently, certain complex flat ETFs also have been used to construct tight frames for  $\mathbb{C}^d$  that consist of  $d^2 + 1$  vectors and have minimal coherence, meeting the *orthoplex bound* [5]. These frames present a reasonable alternative to ETFs for  $\mathbb{C}^d$  that consist of  $d^2$  vectors; though such ETFs have been much sought after in quantum information theory, their existence remains unsettled in all but a finite number of cases [28, 37]. Real flat sensing matrices with low coherence also arise in certain compressed sensing applications like the Single Pixel Camera [13], though so far, all known real flat ETFs with  $n - 1 > d > 1$  are known to be inferior to random  $\{\pm 1\}$ -valued matrices with respect to the *restricted isometry property* [1, 24].

In the next section, we establish notation and discuss some well-known concepts that we need, such as *Naimark complements* of ETFs. In Section 3, we construct an explicit Naimark complement of any Steiner ETF, and use it to show that every Kirkman ETF is a possibly-complex Hadamard ETF. This in particular yields a new infinite family of real flat ETFs. In Section 4, we characterize all real flat ETFs in terms of combinatorial designs known as *quasi-symmetric designs* (QSDs). When combined with the results of Section 3, this characterization provides a new infinite family of QSDs. When combined with [7], this characterization shows that there exists a real flat ETF when  $(d, n)$  is either  $(66, 144)$  or  $(78, 144)$ . This characterization also leads to new necessary conditions on the existence of real flat ETFs, showing for example that real ETFs with  $(d, n) = (15, 36)$  are not flat. In the fifth and final section, we give some other miscellaneous results about Hadamard ETFs.

One may also consider flat representations of higher dimension, such as when the columns of an  $m \times n$  flat matrix  $\Phi$  form an ETF for their  $d$ -dimensional span where  $d < m$ . Such ETFs have arisen only recently [18, 17], and we leave a more thorough investigation of them for future work. Preliminary versions of parts of the material presented in Sections 3 and 4 appeared in the conference proceedings [14] and [15], respectively.

## 2. Preliminaries

As above, let  $\Phi$  denote the  $d \times n$  synthesis operator of  $\{\varphi_j\}_{j=1}^n$ , and let  $\Phi^*$  be its  $n \times d$  conjugate transpose. The corresponding *Gram matrix* is the  $n \times n$  matrix  $\Phi^* \Phi$  whose  $(j, j')$ th entry is  $(\Phi^* \Phi)(j, j') = \langle \varphi_j, \varphi_{j'} \rangle$ . We say  $\{\varphi_j\}_{j=1}^n$  is *equal-norm* if there exists  $\beta > 0$  such that  $\|\varphi_j\|^2 = \beta$  for all  $j$ , and say it is *equiangular* if we further have  $\gamma \geq 0$  such that  $|\langle \varphi_j, \varphi_{j'} \rangle| = \gamma$  for all  $j \neq j'$ . That is,  $\{\varphi_j\}_{j=1}^n$  is equiangular when the diagonal entries of  $\Phi^* \Phi$  are constant and its off-diagonal entries have constant modulus. We say  $\{\varphi_j\}_{j=1}^n$  is a *tight frame* for  $\mathbb{F}^d$  if there exists  $\alpha > 0$  such that  $\Phi \Phi^* = \alpha \mathbf{I}$ , namely when the rows of  $\Phi$  are nonzero, equal-norm and orthogonal. As mentioned above, it is well known that a sequence of equal-norm vectors  $\{\varphi_j\}_{j=1}^n$  achieves equality in (1) if and only if it is both equiangular and a tight frame for  $\mathbb{F}^d$ , namely when it is an ETF for  $\mathbb{F}^d$ ; see [24] for a short, modern proof of this fact. In this case, the tightness constant  $\alpha$  is necessarily  $\frac{n\beta}{d}$ , and the coherence  $\frac{\gamma}{\beta}$  is necessarily the Welch bound  $[\frac{n-d}{d(n-1)}]^{\frac{1}{2}}$ .

For any ETF  $\{\varphi_j\}_{j=1}^n$  for  $\mathbb{F}^d$  with  $n > d$ , there exists an ETF  $\{\tilde{\varphi}_j\}_{j=1}^n$  for  $\mathbb{F}^{n-d}$ . Here, the  $n - d$  rows of the corresponding synthesis operator  $\tilde{\Phi}$  are formed by completing the  $d$  rows of  $\Phi$  to an equal-norm orthogonal basis for  $\mathbb{F}^n$ . This ensures that

$$\alpha^{-\frac{1}{2}} \begin{bmatrix} \Phi \\ \tilde{\Phi} \end{bmatrix}$$

is unitary, implying that  $\tilde{\Phi} \tilde{\Phi}^* = \alpha \mathbf{I}$  and that  $\tilde{\Phi}^* \tilde{\Phi} = \alpha \mathbf{I} - \Phi^* \Phi$ , meaning that the columns  $\{\tilde{\varphi}_j\}_{j=1}^n$  of  $\tilde{\Phi}$  form a tight frame for  $\mathbb{F}^{n-d}$  with the property that  $\|\tilde{\varphi}_j\|^2 = \alpha - \beta$  for all  $j$  and that  $|\langle \tilde{\varphi}_j, \tilde{\varphi}_{j'} \rangle| = |-\langle \varphi_j, \varphi_{j'} \rangle| = \gamma$  for all  $j \neq j'$ . Any such sequence  $\{\tilde{\varphi}_j\}_{j=1}^n$  is called a *Naimark complement* for  $\mathbb{F}^{n-d}$ . Since the rows of  $\tilde{\Phi}$  can be any appropriately-scaled orthogonal basis for the orthogonal complement of the row space of  $\Phi$ , Naimark complements are not unique. This is a key point in distinguishing Hadamard ETFs from those that are simply flat, as expressed in the following restatement of Definition 1, whose proof is immediate:

**Proposition 1.** *A real flat ETF is Hadamard if and only if it has a real flat Naimark complement. A complex flat ETF is complex Hadamard if and only if it has a flat Naimark complement.*

For a simple example of these ideas, consider *regular simplices*, namely ETFs for  $\mathbb{F}^d$  that consist of  $d+1$  vectors. The Naimark complement of a regular simplex is an ETF for  $\mathbb{F}^1$ , namely a sequence of  $d+1$  nonzero scalars of equal modulus. Conversely, any such sequence is an ETF for  $\mathbb{F}^1$ , meaning each of its Naimark complements is a regular simplex. In particular, for any positive integer  $d$ , there exists a regular simplex for  $\mathbb{F}^d$ . When  $\mathbb{F} = \mathbb{C}$ , this ETF can always be chosen to be flat: we can form its synthesis operator by removing any row from a complex Hadamard matrix of size  $d+1$ , such as a discrete Fourier transform (DFT). This same method produces a flat regular simplex for  $\mathbb{R}^d$  if there exists a Hadamard matrix of size  $d+1$ . For example, removing the all-ones row from the canonical  $4 \times 4$  Hadamard matrix gives a flat regular simplex for  $\mathbb{R}^3$  (a tetrahedron):

$$\begin{bmatrix} 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix}. \quad (3)$$

In fact, since the Naimark complement of a regular simplex is automatically flat, this is the only way to produce a real flat regular simplex, meaning each such ETF is Hadamard. In particular, by the known half of the Hadamard conjecture, a real flat regular simplex for  $\mathbb{R}^d$  can only exist if  $d+1$  is either 2 or divisible by 4.

For another example of these ideas, note that every harmonic ETF is possibly-complex Hadamard, being a  $d \times n$  submatrix of an  $n \times n$  character table of some finite abelian group. In this case, the remaining  $n-d$  rows of the character table give a flat Naimark complement. Some of these ETFs are real, such as those that arise from a *McFarland* difference set in  $\mathbb{Z}_2^{2(e+1)}$  where  $e$  is a positive integer [12, 24]. These ETFs have parameters  $d = \frac{1}{2}(n - \sqrt{n})$  where  $n = 2^{2(e+1)}$ ; taking  $e = 1$  gives the real flat ETF given in (2). By Corollary 2 of [24], every real harmonic ETF is either a regular simplex or has these parameters.

Thus, the two most popular constructions of flat ETFs—regular simplices and harmonic ETFs—also happen to be possibly-complex Hadamard. This leads one to ask whether this is true in general:

**Conjecture 1.** Every real flat ETF is Hadamard.

**Conjecture 2.** Every complex flat ETF is complex Hadamard.

To be clear, it is reasonable to believe these conjectures are false, since apart from the above examples, there is no compelling reason to believe that every flat ETF necessarily has a flat Naimark complement. That said, we have not been able to prove that either one of these conjectures is false. In fact, in the next section we show to the contrary that the other known class of flat ETFs—the Kirkman ETFs of [24]—are possibly-complex Hadamard. Both the proof of that result as well as the characterization of real flat ETFs given in Section 4 depend on certain well-known combinatorial designs.

To elaborate, given integers  $v, k, \lambda, r$  and  $b$  with  $v > k > 0$  and  $b > 0$ , a corresponding *balanced incomplete block design* BIBD( $v, k, \lambda, r, b$ ) is a set  $\mathcal{V}$  of cardinality  $v$ —whose elements are called *vertices*—along with  $b$  subsets of  $\mathcal{V}$ —called *blocks*—with the property that each block contains exactly  $k$  vertices, each vertex is contained in exactly  $r$  blocks, and every pair of distinct vertices is contained in exactly  $\lambda$  blocks. Letting  $\mathbf{X}$  be a corresponding  $b \times v$  incidence matrix, this is equivalent to having

$$\mathbf{X}\mathbf{1} = k\mathbf{1}, \quad \mathbf{X}^T\mathbf{1} = r\mathbf{1}, \quad \mathbf{X}^T\mathbf{X} = (r - \lambda)\mathbf{I} + \lambda\mathbf{J}, \quad (4)$$

where  $\mathbf{1}$  and  $\mathbf{J}$  denote all-ones vectors and matrices, respectively. For example, taking all two-element subsets of  $\mathcal{V} = \{1, 2, 3, 4\}$  gives a BIBD(4, 2, 1, 3, 6) with an incidence matrix of

$$\mathbf{X} = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}. \quad (5)$$

The parameters of a BIBD are dependent, satisfying

$$bk = vr, \quad (v - 1)\lambda = r(k - 1). \quad (6)$$

Since  $v > k > 0$  and  $b > 0$ , these relationships imply  $0 \leq \lambda < r < b$ . This in turn implies that  $\mathbf{X}^T \mathbf{X} = (r - \lambda)\mathbf{I} + \lambda\mathbf{J}$  is positive-definite, implying *Fisher's inequality*, namely that  $b \geq v$ .

### 3. Naimark complements of Steiner and Kirkman equiangular tight frames.

We begin this section by constructing an explicit Naimark complement for any Steiner ETF. Every Steiner ETF arises from a BIBD( $v, k, \lambda, r, b$ ) with  $\lambda = 1$ ; such a combinatorial design is also known as a *Steiner system*  $S(t, k, v)$  with  $t = 2$ . To be precise, let  $\mathbf{X}$  be the  $b \times v$  incidence matrix of a BIBD( $v, k, 1, r, b$ ), let  $\Psi$  be the  $r \times (r + 1)$  synthesis operator of a flat regular simplex. The synthesis operator  $\Phi$  of the corresponding Steiner ETF is the  $b \times v(r + 1)$  matrix obtained by replacing each 1-valued entry in any given column of  $\mathbf{X}$  with a distinct row of  $\Psi$ , and replacing each 0-valued entry of  $\mathbf{X}$  with a  $1 \times (r + 1)$  submatrix of zeros. For example, combining (5) and (3) in this way gives the following Steiner ETF with parameters  $(d, n) = (6, 16)$ :

$$\Phi = \begin{bmatrix} 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 & 0 & 0 & 0 & 0 & 1 & 1 & -1 & -1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & -1 & -1 & 0 & 0 & 0 & 0 & 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & -1 & -1 & 1 \\ 0 & 0 & 0 & 0 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 0 & 0 & 0 & 0 \end{bmatrix}. \quad (7)$$

This construction was introduced in [21] as a method for obtaining SRGs. In [20] it was rediscovered and recognized as a method for constructing ETFs. These ETFs are real if the flat regular simplex is real, namely when it is obtained from a Hadamard matrix of size  $r + 1$ . This method of combining two matrices is unlike other common methods for doing so, as it explicitly relies on the fact that  $\mathbf{X}$  is  $\{0, 1\}$ -valued and has constant column sums. Because of the unusualness of this construction, the standard proof of the fact that it yields ETFs is “wordy,” see [21, 20]. This is frustrating, since  $\Phi$  seems tantalizingly similar to the tensor product of  $\mathbf{X}$  and  $\Psi$ . In fact, as we now explain, an explicit tensor-product-based expression for  $\Phi$  can be found, provided we first “lift” the incidence matrix of a BIBD to a permutation matrix:

**Definition 2.** Let  $\mathbf{X}$  be the  $b \times v$  incidence matrix of a BIBD( $v, k, 1, r, b$ ). A corresponding *BIBD permutation matrix* is a permutation matrix  $\Pi$  of size  $bk = vr$  such that

$$\mathbf{X} = (\mathbf{I}_b \otimes \mathbf{1}_k^T) \Pi (\mathbf{I}_v \otimes \mathbf{1}_r). \quad (8)$$

Here,  $\mathbf{I}_b$  and  $\mathbf{I}_v$  are  $b \times b$  and  $v \times v$  identity matrices, respectively, while  $\mathbf{1}_k$  and  $\mathbf{1}_r$  are all-ones vectors of size  $k \times 1$  and  $r \times 1$ , respectively. Regarding  $\mathbf{\Pi}$  as a  $b \times v$  array of  $\{0, 1\}$ -valued submatrices  $\{\mathbf{\Pi}_{i,j}\}_{i=1, j=1}^{b, v}$  of size  $k \times r$ , this means that for any  $i = 1, \dots, b$  and  $j = 1, \dots, v$ ,

$$\mathbf{X}(i, j) = \mathbf{1}_k^\top \mathbf{\Pi}_{i,j} \mathbf{1}_r = \sum_{p=1}^k \sum_{q=1}^r \mathbf{\Pi}_{i,j}(p, q). \quad (9)$$

Thus,  $\mathbf{\Pi}_{i,j} = \mathbf{0}$  if  $\mathbf{X}(i, j) = 0$  and contains exactly one 1-valued entry if  $\mathbf{X}(i, j) = 1$ . The permutation matrix of a BIBD is not unique, but one does always exist. For example, we can define  $\mathbf{\Pi}_{i,j}(p, q)$  to be 1 if and only if  $\mathbf{X}(i, j)$  is both the  $p$ th “1” in the  $i$ th row of  $\mathbf{X}$  and the  $q$ th “1” in the  $j$ th column of  $\mathbf{X}$ ; for the incidence matrix (5), this gives

$$\mathbf{\Pi} = \left[ \begin{array}{ccc|ccc|ccc|ccc} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ \hline 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ \hline 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ \hline 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \end{array} \right]. \quad (10)$$

This method produces a permutation matrix since the corresponding mappings from  $[b] \times [k]$  and  $[v] \times [r]$  into  $\{(i, j) : \mathbf{X}(i, j) = 1\}$  are both invertible, establishing a permutation  $(i, p) \mapsto (j, q)$ .

A Steiner ETF is obtained by taking a BIBD’s permutation matrix  $\mathbf{\Pi}$ , and replacing the vector  $\mathbf{1}_r$  in (8) with an  $r \times (r + 1)$  unimodular regular simplex  $\mathbf{\Psi}$ , that is, by letting

$$\mathbf{\Phi} = (\mathbf{I}_b \otimes \mathbf{1}_k^\top) \mathbf{\Pi} (\mathbf{I}_v \otimes \mathbf{\Psi}). \quad (11)$$

Indeed, regarding  $\mathbf{I}_b \otimes \mathbf{1}_k^\top$ ,  $\mathbf{\Pi}$  and  $\mathbf{I}_v \otimes \mathbf{\Psi}$  as  $b \times b$ ,  $b \times v$  and  $v \times v$  arrays of submatrices of size  $1 \times k$ ,  $k \times r$  and  $r \times (r + 1)$ , respectively, gives that  $\mathbf{\Phi}$  is a  $b \times v$  array of submatrices of size  $1 \times (r + 1)$ . In particular, the  $(i, j)$ th submatrix of  $\mathbf{\Phi}$  is  $\mathbf{\Phi}_{i,j} = \mathbf{1}_k^\top \mathbf{\Pi}_{i,j} \mathbf{\Psi}$ ; when  $\mathbf{X}(i, j) = 0$ , this is a row of zeros; when  $\mathbf{X}(i, j) = 1$ , this is the row of  $\mathbf{\Psi}$  corresponding to the column index of the nonzero entry of  $\mathbf{\Pi}_{i,j}$ . As part of the next result, we directly verify that (11) indeed defines the synthesis operator of an ETF with parameters  $(d, n) = (b, v(r + 1))$ . In fact, we show this holds even if  $\mathbf{1}_k$  is replaced with any  $k \times 1$  vector with unimodular entries; this turns out to be the key to constructing an explicit Naimark complement of  $\mathbf{\Phi}$ .

Motivating by example, note that when  $\mathbf{\Pi}$  and  $\mathbf{\Psi}$  are (10) and (3), respectively, the matrix  $(\mathbf{I}_6 \otimes [1 \ 1]) \mathbf{\Pi} (\mathbf{I}_4 \otimes \mathbf{\Psi})$  gives the ETF in (7). Meanwhile, replacing  $[1 \ 1]$  with the orthogonal flat

vector  $[1 \ -1]$  yields the  $6 \times 16$  synthesis operator of a second ETF:

$$(\mathbf{I}_6 \otimes [1 \ -1])\mathbf{\Pi}(\mathbf{I}_4 \otimes \mathbf{\Psi}) = \begin{bmatrix} 1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 \\ 1 & 1 & -1 & -1 & 0 & 0 & 0 & 0 & -1 & -1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & -1 & -1 & 0 & 0 & 0 & 0 & -1 & -1 & 1 & 1 \\ 1 & -1 & -1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1 & 1 & 1 & -1 \\ 0 & 0 & 0 & 0 & 1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 & 0 & 0 & 0 & 0 \end{bmatrix}. \quad (12)$$

In essence, this second ETF is obtained by negating the second nonzero  $1 \times (r+1)$  submatrix in each row of (7). As explicitly verified in the next result, this ensures that the row spaces of the synthesis operators of these two ETFs are mutually orthogonal. In particular, (12) gives six rows of a  $10 \times 16$  Naimark complement of (7). The remaining four rows can be obtained by tensoring  $\mathbf{I}_v = \mathbf{I}_4$  with the  $1 \times (r+1) = 1 \times 4$  Naimark complement of  $\mathbf{\Psi}$ , and scaling it appropriately:

$$\sqrt{k}(\mathbf{I}_v \otimes \mathbf{1}_{r+1}^*) = \sqrt{2}(\mathbf{I}_4 \otimes \mathbf{1}_4^*) = \sqrt{2} \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix}.$$

In the following result, we generalize this example to construct an explicit Naimark complement for any Steiner ETF.

**Theorem 1.** *Let  $\mathbf{\Pi}$  be the permutation matrix of a BIBD( $v, k, 1, r, b$ )—see Definition 2—and let  $\mathbf{F}$  and  $\mathbf{G}$  be possibly-complex Hadamard matrices of size  $k$  and  $r+1$ , respectively. Write  $\mathbf{G} = [\mathbf{G}_1 \ \mathbf{g}_2]$  where  $\mathbf{G}_1$  is  $(r+1) \times r$ . For any  $l = 1, \dots, k$ , let*

$$\mathbf{\Phi}_l := (\mathbf{I}_b \otimes \mathbf{f}_l^*)\mathbf{\Pi}(\mathbf{I}_v \otimes \mathbf{G}_1^*), \quad (13)$$

where  $\mathbf{f}_l$  is the  $l$ th column of  $\mathbf{F}$ . Then, the  $v(r+1)$  columns of each  $\mathbf{\Phi}_l$  form an ETF for  $\mathbb{F}^b$ . Moreover, the row spaces of the matrices  $\{\mathbf{\Phi}_l\}_{l=1}^k$  are mutually orthogonal and the columns of

$$\tilde{\mathbf{\Phi}}_1 = \begin{bmatrix} \mathbf{\Phi}_2 \\ \vdots \\ \mathbf{\Phi}_k \\ \sqrt{k}(\mathbf{I}_v \otimes \mathbf{g}_2^*) \end{bmatrix} \quad (14)$$

form a Naimark complement for the ETF formed by the columns of  $\mathbf{\Phi}_1$ . Here, both  $\mathbf{F}$  and  $\mathbf{G}$  can be chosen to be real if and only if  $k = 2$  and there exists a Hadamard matrix of size  $v$ . In this case, the columns of  $\mathbf{\Phi}_1$  and  $\tilde{\mathbf{\Phi}}_1$  form real ETFs of  $v^2$  vectors for  $\mathbb{R}^{\frac{1}{2}v(v-1)}$  and  $\mathbb{R}^{\frac{1}{2}v(v+1)}$ , respectively.

*Proof.* For any  $l, l' = 1, \dots, k$ , the fact that  $\mathbf{G}_1^* \mathbf{G}_1 = (r+1)\mathbf{I}_r$  gives

$$\begin{aligned} \mathbf{\Phi}_l \mathbf{\Phi}_{l'}^* &= [(\mathbf{I}_b \otimes \mathbf{f}_l^*)\mathbf{\Pi}(\mathbf{I}_v \otimes \mathbf{G}_1^*)][(\mathbf{I}_v \otimes \mathbf{G}_1)\mathbf{\Pi}^T(\mathbf{I}_b \otimes \mathbf{f}_{l'})] \\ &= (\mathbf{I}_b \otimes \mathbf{f}_l^*)\mathbf{\Pi}(\mathbf{I}_v \otimes \mathbf{G}_1^* \mathbf{G}_1)\mathbf{\Pi}^T(\mathbf{I}_b \otimes \mathbf{f}_{l'}) \\ &= (\mathbf{I}_b \otimes \mathbf{f}_l^*)\mathbf{\Pi}[\mathbf{I}_v \otimes (r+1)\mathbf{I}_r]\mathbf{\Pi}^T(\mathbf{I}_b \otimes \mathbf{f}_{l'}) \\ &= (r+1)(\mathbf{I}_b \otimes \mathbf{f}_l^*)(\mathbf{I}_b \otimes \mathbf{f}_{l'}) \\ &= (r+1)(\mathbf{I}_b \otimes \langle \mathbf{f}_l, \mathbf{f}_{l'} \rangle) \\ &= \begin{cases} k(r+1)\mathbf{I}_b, & l = l', \\ \mathbf{0}_b, & l \neq l'. \end{cases} \end{aligned} \quad (15)$$



Thus, the columns of each  $\Phi_l$  form a tight frame for  $\mathbb{F}^b$ , and the row spaces of  $\Phi_l$  and  $\Phi_{l'}$  are orthogonal for any  $l \neq l'$ , as claimed.

To see that  $\tilde{\Phi}_1$  is a Naimark complement of  $\Phi_1$ , note that  $\mathbf{G}_1^* \mathbf{g}_2 = \mathbf{0}$  and so for any  $l = 1, \dots, k$ ,

$$\Phi_l [\sqrt{k}(\mathbf{I}_v \otimes \mathbf{g}_2^*)]^* = \sqrt{k}[(\mathbf{I}_b \otimes \mathbf{f}_l^*) \mathbf{\Pi}(\mathbf{I}_v \otimes \mathbf{G}_1^*)](\mathbf{I}_v \otimes \mathbf{g}_2) = \sqrt{k}(\mathbf{I}_b \otimes \mathbf{f}_l^*) \mathbf{\Pi}(\mathbf{I}_v \otimes \mathbf{G}_1^* \mathbf{g}_2) = \mathbf{0}.$$

Also,  $[\sqrt{k}(\mathbf{I}_v \otimes \mathbf{g}_2^*)][\sqrt{k}(\mathbf{I}_v \otimes \mathbf{g}_2^*)]^* = k(\mathbf{I}_v \otimes \langle \mathbf{g}_2, \mathbf{g}_2 \rangle) = k(r+1)\mathbf{I}_v$ . When combined with (15), these facts imply that the rows of  $\Phi_1$  are orthogonal to each other as well as to all rows of  $\tilde{\Phi}_1$ , and that all rows of  $\Phi_1$  and  $\tilde{\Phi}_1$  have the same norm. Since  $\Phi_1$  is  $b \times v(r+1)$  while  $\tilde{\Phi}_1$  is  $[b(k-1)+v] \times v(r+1)$  where (6) gives  $b + [b(k-1) + v] = bk + v = vr + v = v(r+1)$ , this implies that the columns of  $\tilde{\Phi}_1$  indeed form a Naimark complement of the tight frame formed by the columns of  $\Phi_1$ .

Next, for any  $l = 1, \dots, k$  we show that the columns of  $\Phi_l$  form an ETF for  $\mathbb{F}^b$ . To be precise, we show that the all off-diagonal entries of  $\Phi_l^* \Phi_l$  have modulus 1, while all diagonal entries have value  $r$ . By (13),  $\Phi_l$  can be regarded as a  $b \times v$  array of submatrices of size  $1 \times (r+1)$ , namely  $\{(\Phi_l)_{i,j}\}_{i=1}^b, j=1}^v = \{\mathbf{f}_l^* \mathbf{\Pi}_{i,j} \mathbf{G}_1^*\}_{i=1}^b, j=1}^v$ . Thus, the Gram matrix  $\Phi_l^* \Phi_l$  is a  $v \times v$  array of submatrices of size  $(r+1) \times (r+1)$ . Specifically, for any  $j, j' = 1, \dots, v$ , the  $(j, j')$ th submatrix of  $\Phi_l^* \Phi_l$  is

$$(\Phi_l^* \Phi_l)_{j,j'} = \sum_{i=1}^b (\Phi_l^*)_{j,i} (\Phi_l)_{i,j'} = \sum_{i=1}^b (\mathbf{f}_l^* \mathbf{\Pi}_{i,j} \mathbf{G}_1^*)^* (\mathbf{f}_l^* \mathbf{\Pi}_{i,j'} \mathbf{G}_1^*). \quad (16)$$

Here, for any  $j = 1, \dots, v$  and  $s = 1, \dots, r+1$ , the  $(1, s)$ th entry of  $\mathbf{f}_l^* \mathbf{\Pi}_{i,j} \mathbf{G}_1^*$  is

$$(\mathbf{f}_l^* \mathbf{\Pi}_{i,j} \mathbf{G}_1^*)(1, s) = \sum_{p=1}^k \sum_{q=1}^r \overline{\mathbf{f}_l(p)} \mathbf{\Pi}_{i,j}(p, q) \mathbf{G}_1^*(q, s) = \sum_{p=1}^k \sum_{q=1}^r \mathbf{\Pi}_{i,j}(p, q) \overline{\mathbf{F}(p, l) \mathbf{G}(s, q)}. \quad (17)$$

Combining (16) and (17) gives that for any  $j, j' = 1, \dots, v$  and  $s, s' = 1, \dots, r+1$ , the  $(s, s')$ th entry of the  $(j, j')$ th submatrix of  $\Phi_l^* \Phi_l$  is

$$\begin{aligned} (\Phi_l^* \Phi_l)_{j,j'}(s, s') &= \sum_{i=1}^b \overline{(\mathbf{f}_l^* \mathbf{\Pi}_{i,j} \mathbf{G}_1^*)(1, s)} (\mathbf{f}_l^* \mathbf{\Pi}_{i,j'} \mathbf{G}_1^*)(1, s') \\ &= \sum_{i=1}^b \sum_{p,p'=1}^k \sum_{q,q'=1}^r \mathbf{\Pi}_{i,j}(p, q) \mathbf{\Pi}_{i,j'}(p', q') \mathbf{F}(p, l) \overline{\mathbf{F}(p', l)} \mathbf{G}(s, q) \overline{\mathbf{G}(s', q')}. \end{aligned} \quad (18)$$

To proceed, recall that the definition of the BIBD permutation matrix implies (9) where  $\mathbf{X}$  is the  $b \times v$  incidence matrix of the underlying BIBD( $v, k, 1, r, b$ ). In particular, if  $j \neq j'$  then

$$1 = (\mathbf{X}^T \mathbf{X})(j, j') = \sum_{i=1}^b \mathbf{X}(i, j) \mathbf{X}(i, j') = \sum_{i=1}^b \sum_{p,p'=1}^k \sum_{q,q'=1}^r \mathbf{\Pi}_{i,j}(p, q) \mathbf{\Pi}_{i,j'}(p', q').$$

Since  $\mathbf{\Pi}$  is  $\{0, 1\}$ -valued, this means that when  $j \neq j'$ , there exists exactly one choice of  $(i, p, p', q, q')$  such that  $\mathbf{\Pi}_{i,j}(p, q) \mathbf{\Pi}_{i,j'}(p', q') = 1$ . (Here,  $i$  corresponds to the unique block in the design that contains both the  $j$ th and  $j'$ th vertices, whereupon  $(i, j)$  uniquely determines  $(p, q)$  and  $(i, j')$  uniquely determines  $(p', q')$ .) As such, when  $j \neq j'$ , there is exactly one nonzero summand in (18), that is,  $(\Phi_l^* \Phi_l)_{j,j'}(s, s') = \mathbf{F}(p, l) \overline{\mathbf{F}(p', l)} \mathbf{G}(s, q) \overline{\mathbf{G}(s', q')}$  for this unique choice of  $(i, p, p', q, q')$ . In particular, if  $j \neq j'$  then  $|(\Phi_l^* \Phi_l)_{j,j'}(s, s')| = 1$  for all  $s, s'$ .

In the remaining case where  $j = j'$ , note that since each submatrix  $\Phi_{i,j}$  contains at most one entry that has value 1,  $\Pi_{i,j}(p, q)\Pi_{i,j}(p', q') = 1$  only when  $p' = p$  and  $q' = q$ . As such, when  $j = j'$ , (18) simplifies to

$$\begin{aligned} (\Phi_l^* \Phi_l)_{j,j}(s, s') &= \sum_{i=1}^b \sum_{p=1}^k \sum_{q=1}^r \Pi_{i,j}(p, q) \mathbf{F}(p, l) \overline{\mathbf{F}(p, l)} \mathbf{G}(s, q) \overline{\mathbf{G}(s', q)} \\ &= \sum_{q=1}^r \mathbf{G}(s, q) \overline{\mathbf{G}(s', q)} \sum_{i=1}^b \sum_{p=1}^k \Pi_{i,j}(p, q). \end{aligned}$$

Here,  $\sum_{i=1}^b \sum_{p=1}^k \Pi_{i,j}(p, q)$  is the sum of all of the entries in a column of the permutation matrix  $\Pi$ , namely 1. When combined with the fact that  $\mathbf{G}$  is a Hadamard matrix, this implies

$$(\Phi_l^* \Phi_l)_{j,j}(s, s') = \sum_{q=1}^r \mathbf{G}(s, q) \overline{\mathbf{G}(s', q)} = (r+1)\mathbf{I}_{r+1}(s, s') - \mathbf{G}(s, r+1) \overline{\mathbf{G}(s', r+1)},$$

a quantity which equals  $r$  when  $s = s'$  and has modulus 1 when  $s \neq s'$ .

For the final conclusions, note that if both  $\mathbf{F}$  and  $\mathbf{G}$  are (real) Hadamard matrices, then the known half of the Hadamard conjecture implies that  $k$  and  $r+1$  are either 2 or are divisible by 4. In particular, if  $k \neq 2$ , then  $k$  and  $r+1 \geq k+1$  are divisible by 4. Since  $b = \frac{vr}{k}$  is an integer and  $r$  is odd, this implies that 4 divides  $v$ , contradicting the fact that  $v = r(k-1)+1 \equiv (-1)^2+1 \equiv 2 \pmod{4}$ . Thus, if  $\mathbf{F}$  and  $\mathbf{G}$  are both Hadamard matrices, then  $k = 2$ . Conversely, if  $k = 2$  and there exists a Hadamard matrix of size  $v$ , we note that there is a unique BIBD on  $v$  vertices with this  $k$  and with  $\lambda = 1$ , namely the BIBD( $v, 2, 1, v-1, \frac{1}{2}v(v-1)$ ) that consists of all 2-element subsets of  $[v]$ . Taking this BIBD, letting  $\mathbf{F}$  be the canonical Hadamard matrix of size 2, and letting  $\mathbf{G}$  be the given Hadamard matrix of size  $v$ , the resulting matrices  $\Phi_1$  and  $\tilde{\Phi}_1$  are real flat matrices of size  $\frac{1}{2}v(v-1) \times v^2$  and  $\frac{1}{2}v(v+1) \times v^2$ , respectively.  $\square$

For any BIBD( $v, k, 1, r, b$ ) we note there always exists complex Hadamard matrices of size  $k$  and  $r+1$ , such as DFTs. In fact, if  $k > 2$  and there exists a Hadamard matrix of size  $r+1$ , then taking  $\mathbf{G}$  to be that matrix and letting  $\mathbf{F}$  be a DFT (or any other complex Hadamard matrix whose first column is all ones), the Steiner ETF  $\Phi_1$  is real and its Naimark complement  $\tilde{\Phi}_1$  is complex. Of course, any real ETF has a real Naimark complement. In fact, if  $\mathbf{G}$  is Hadamard and  $\mathbf{F}$  is any real multiple of an orthogonal matrix whose first column is all ones, then the first part of the proof of Theorem 1 is still valid and (14) is still a Naimark complement of  $\Phi_1$ . However, in this case,  $\Phi_l$  is not necessarily an ETF for  $l = 2, \dots, k$ . For example, the *Fano plane* is a well-known BIBD(7, 3, 1, 3, 7), and letting  $\mathbf{G}$  be the canonical  $4 \times 4$  Hadamard matrix and letting either

$$\mathbf{F} = \begin{bmatrix} 1 & \sqrt{2} & 0 \\ 1 & -\frac{1}{\sqrt{2}} & \frac{\sqrt{3}}{\sqrt{2}} \\ 1 & -\frac{1}{\sqrt{2}} & -\frac{\sqrt{3}}{\sqrt{2}} \end{bmatrix} \quad \text{or} \quad \mathbf{F} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & \omega & \omega^2 \\ 1 & \omega^2 & \omega \end{bmatrix}, \quad \omega = \exp\left(\frac{2\pi i}{3}\right)$$

gives either a real or complex Naimark complement  $\tilde{\Phi}_1$  for the same  $7 \times 28$  real ETF  $\Phi_1$ . In the second case however,  $\Phi_2$  and  $\Phi_3$  are themselves complex Steiner ETFs; as we now discuss, this is important when trying to modify these ETFs in a way that makes them complex Hadamard.

### 3.1. Naimark complements of Kirkman equiangular tight frames

A BIBD is *resolvable*—denoted an RBIBD—if its  $b$  blocks can be arranged as  $r$  collections of  $\frac{b}{r} = \frac{v}{k}$  blocks apiece—called *parallel classes*—with each parallel class forming a partition of the vertex set. For example, (5) is the incidence matrix of an RBIBD(4, 2, 1, 3, 6) since its six blocks can be arranged as three parallel classes—blocks one and two, blocks three and four, and blocks five and six—with each class giving a partition of the vertex set  $\{1, 2, 3, 4\}$ .

In [24], it is shown that every Steiner ETF arising from a RBIBD( $v, k, 1, r, b$ ) can be made flat by applying a scaled unitary operator to it. Such ETFs are dubbed Kirkman ETFs, in honor of *Kirkman's schoolgirl problem*, a foundational problem in combinatorial design regarding the existence of an RBIBD(15, 3, 1, 7, 35). To be precise, let  $\mathbf{X}$  be the incidence matrix of an RBIBD( $v, k, 1, r, b$ ), arranged without loss of generality as an  $r \times 1$  array of submatrices of size  $\frac{v}{k} \times v$ , each corresponding to a parallel class. For any  $l = 1, \dots, k$ , and any possibly-complex Hadamard matrices  $\mathbf{F}$  and  $\mathbf{G}$  of size  $k$  and  $r + 1$ , respectively, let  $\Phi_l$  be the corresponding Steiner ETF (13). Due to the structure of  $\mathbf{X}$ , every column of  $\Phi_l$  is a direct sum of  $r$  vectors in  $\mathbb{F}^{\frac{v}{k}}$ , each having a single entry of modulus one with all other entries being zero. Multiplying any such vector by a possibly-complex Hadamard matrix  $\mathbf{E}$  of size  $\frac{v}{k}$  produces a vector with all unimodular entries. In particular,  $(\mathbf{I}_r \otimes \mathbf{E})\Phi_l$  is flat. Moreover, since  $(\mathbf{I}_r \otimes \mathbf{E})$  is a scalar multiple of a unitary matrix, the columns of  $(\mathbf{I}_r \otimes \mathbf{E})\Phi_l$  still form an ETF for  $\mathbb{F}^b$ . When combined with Theorem 1, these ideas lead to the following result:

**Theorem 2.** *Let  $\Pi$  be the permutation matrix of an RBIBD( $v, k, 1, r, b$ ), and let  $\mathbf{E}$ ,  $\mathbf{F}$  and  $\mathbf{G}$  be possibly-complex Hadamard matrices of size  $\frac{v}{k}$ ,  $k$  and  $r + 1$ , respectively. Write  $\mathbf{G} = [\mathbf{G}_1 \ \mathbf{g}_2]$  where  $\mathbf{G}_1$  is  $(r + 1) \times r$ . For any  $l = 1, \dots, k$ , let*

$$\Psi_l := (\mathbf{I}_r \otimes \mathbf{E})(\mathbf{I}_b \otimes \mathbf{f}_l^*)\Pi(\mathbf{I}_v \otimes \mathbf{G}_1^*),$$

where  $\mathbf{f}_l$  is the  $l$ th column of  $\mathbf{F}$ . Then, the  $v(r + 1)$  columns of each  $\Psi_l$  form a flat ETF for  $\mathbb{F}^b$ . Moreover, the row spaces of the matrices  $\{\Psi_l\}_{l=1}^k$  are mutually orthogonal and the matrix

$$\tilde{\Psi}_1 = \begin{bmatrix} \Psi_2 \\ \vdots \\ \Psi_k \\ (\mathbf{E} \otimes \mathbf{F})(\mathbf{I}_v \otimes \mathbf{g}_2^*) \end{bmatrix}$$

is a flat Naimark complement of  $\Psi_1$ , meaning  $\Psi_1$  is a possibly-complex Hadamard ETF.

In particular, if there exists a Hadamard matrix of size  $u$  then there exists a Hadamard ETF with parameters  $(d, n) = (u(2u - 1), 4u^2)$ .

*Proof.* As noted above, for each  $l = 1, \dots, j$ , the columns of  $\Psi_l = (\mathbf{I}_r \otimes \mathbf{E})\Phi_l$  form a flat ETF for  $\mathbb{F}^b$ . Moreover, for each  $l = 1, \dots, j$ , the row space of  $\Psi_l$  equals that of  $\Phi_l$ . By Theorem 1, these row spaces are mutually orthogonal. They are also orthogonal to the row space of  $\sqrt{k}(\mathbf{I}_v \otimes \mathbf{g}_2^*)$ , which equals that of  $(\mathbf{E} \otimes \mathbf{F})(\mathbf{I}_v \otimes \mathbf{g}_2^*)$ . Here,  $\mathbf{E} \otimes \mathbf{F}$  is a possibly-complex Hadamard matrix of size  $v$ , implying that  $(\mathbf{E} \otimes \mathbf{F})(\mathbf{I}_v \otimes \mathbf{g}_2^*)$  has unimodular entries and orthogonal rows. Altogether, these facts imply that  $\tilde{\Psi}_1$  is indeed a flat Naimark complement for  $\Psi_1$ . By Proposition 1, this means  $\Phi_1$  is a possibly-complex Hadamard ETF.

For the final conclusion, assume there exists a Hadamard matrix  $\mathbf{E}$  of size  $u$ , and let  $v = 2u$ . Since  $v$  is even, the unique BIBD( $v, 2, 1, v - 1, \frac{1}{2}v(v - 1)$ ) is resolvable using the well-known *round-robin tournament schedule*. As such, letting  $\mathbf{F}$  be the canonical Hadamard matrix of size 2, and

noting that  $\mathbf{G} = \mathbf{E} \otimes \mathbf{F}$  is a Hadamard matrix of size  $2u = v = r + 1$ , the columns of  $\Psi_1$  and  $\tilde{\Phi}_1$  form Naimark complementary real flat ETFs with parameters  $(d, n) = (\frac{1}{2}v(v-1), v^2) = (u(2u-1), 4u^2)$  and  $(n-d, n) = (\frac{1}{2}v(v+1), v^2) = (u(2u+1), 4u^2)$ , respectively. Thus, both ETFs are Hadamard.  $\square$

Generally speaking, the significance of Theorem 2 is that it shows that the fundamental idea of harmonic ETFs—that ETFs can be constructed by carefully extracting rows from a possibly-complex Hadamard matrix—can be successful even when the matrix is not the character table of an abelian group. For example, taking  $u = 12$  gives a Hadamard matrix of size 576 from which 276 special rows can be extracted to form a real ETF; in contrast, the character table of any abelian group of order  $576 = 2^6 3^2$  contains cube roots of unity. This realization hopefully better informs searches for new ETFs in the future.

That said, the most immediate contribution of Theorem 2 is that it gives a new infinite family of real flat ETFs. Indeed, prior to this result, the only known real flat ETFs with parameters of the form  $(u(2u+1), 4u^2)$  had  $u = 2^e$  for some  $e \geq 0$ ; such ETFs arise, for example, from the complements of McFarland difference sets in  $\mathbb{Z}_2^{2(e+1)}$  [12]. We now know they exist whenever there exists a Hadamard matrix of size  $u$ . An infinite number of these ETFs are new: for example, there are an infinite number of prime powers  $q$  that are congruent to 1 modulo 4, and for each of these, Paley’s method yields a Hadamard matrix of size  $u = 2(q+1)$ , which is not a power of 2.

#### 4. Real flat equiangular tight frames and quasi-symmetric designs

In this section, we combine ideas from [27] and [24] to give a new characterization of real flat ETFs in terms of particular types of BIBDs known as *quasi-symmetric designs* (QSDs). As we shall see, this characterization further generalizes to a particular class of real ETFs that are not necessarily flat, namely those that arise from the *block graphs* of QSDs [21] via a well-known equivalence between real ETFs and certain *strongly regular graphs* (SRGs) [31, 23, 34].

To be precise, a  $\text{BIBD}(v, r, \lambda, r, b)$  with  $b > v$  is a QSD with *block intersection numbers*  $y > x \geq 0$  if the cardinality of the intersection of any two distinct blocks is either  $x$  or  $y$ . That is, a BIBD with  $b > v$  is a  $\text{QSD}(v, k, \lambda, r, b, x, y)$  if its  $b \times v$  incidence matrix  $\mathbf{X}$  satisfies

$$\mathbf{X}\mathbf{X}^T = (k-x)\mathbf{I} + (y-x)\mathbf{A} + x\mathbf{J} \quad (19)$$

for some  $\{0, 1\}$ -valued matrix  $\mathbf{A}$ . Here,  $\mathbf{A}$  is the adjacency matrix of a graph known as the *block graph* of the QSD. One simple example of a QSD is a BIBD with  $b > v$  and  $\lambda = 1$ : in such a design, any two distinct blocks have at most one vertex in common, and so it is a QSD with  $x = 0$  and  $y = 1$ . Another common way to construct a QSD is to take the *complementary design* of another QSD, namely to consider the incidence matrix  $\mathbf{J} - \mathbf{X}$  instead of  $\mathbf{X}$ .

In our arguments below, we use the well-known fact that any QSD realizes both of its intersection numbers  $\{x, y\}$ , namely that its corresponding block graph is neither complete nor empty. Indeed, if we instead have a BIBD in which any two distinct blocks intersect in exactly  $x$  vertices where  $x < k$ , then  $\mathbf{X}\mathbf{X}^T = (k-x)\mathbf{I} + x\mathbf{J}$  is positive-definite, contradicting the underlying assumption that  $b > v$ . (BIBDs with  $b = v$  are instead called *symmetric designs*.) We shall also make use of the fact that the parameters of a QSD are dependent [29], satisfying

$$k(r-1)(x+y-1) - xy(b-1) = k(k-1)(\lambda-1). \quad (20)$$

It is known that the block graph of any QSD is an SRG [21, 29]. To elaborate, a graph on  $b$  vertices is an  $\text{SRG}(b, a, c, \mu)$  if any vertex has  $a$  neighbors, any two adjacent vertices have  $c$

neighbors in common, and any two nonadjacent vertices have  $\mu$  neighbors in common, namely when its adjacency matrix  $\mathbf{A}$  satisfies  $\mathbf{A}\mathbf{1} = a\mathbf{1}$  and  $\mathbf{A}^2 = (c - \mu)\mathbf{A} + (a - \mu)\mathbf{I} + \mu\mathbf{J}$ . Using (4), it is straightforward to show that the matrix  $\mathbf{A}$  defined by (19) satisfies such a relationship where

$$a = \frac{k(r-1)-x(b-1)}{y-x}, \quad \theta_1 = \frac{(r-\lambda)-(k-x)}{y-x}, \quad \theta_2 = -\frac{k-x}{y-x}, \quad c = a + \theta_1 + \theta_2 + \theta_1\theta_2, \quad \mu = a + \theta_1\theta_2. \quad (21)$$

As with any SRG, these parameters are dependent [8], satisfying

$$a(a - c - 1) = \mu(b - a - 1). \quad (22)$$

It is also known that real ETFs are equivalent to a certain class of SRGs [31, 23, 34]. In particular, as detailed in [16], real ETFs correspond to SRGs whose parameters  $(b, a, c, \mu)$  satisfy  $a = 2\mu$ ; in this case, the corresponding real ETF has parameters

$$n = b + 1, \quad d = \frac{b+1}{2} \left[ 1 + \frac{b-2a-1}{\sqrt{(b-2a-1)^2+4b}} \right]. \quad (23)$$

Putting all of these facts together, it is natural to ask which class of QSDs leads to SRGs which equate to real ETFs. By searching tables of known SRGs [8, 9], we find some instances where this occurs. For example, there exists a BIBD(15, 3, 1, 7, 35); having  $\lambda = 1$ , this design is also a QSD with  $(x, y) = (0, 1)$ . By (21), the corresponding block graph is an SRG(35, 18, 9, 9). Since the parameters of this SRG satisfy  $a = 18 = 2(9) = 2\mu$ , it indeed corresponds to a real ETF; by (23), this ETF has parameters  $(d, n) = (15, 36)$ . Note here that  $d = 15 = v$ . This is not a coincidence: using (6), (20), (21), (22) and (23), the interested reader can show that  $d = v$  for any real ETF that arises from an SRG( $b, a, c, \mu$ ) with  $a = 2\mu$  that itself arises from a QSD( $v, k, \lambda, r, b, x, y$ ). This inspires us to seek a relationship between the  $b \times v$  incidence matrix  $\mathbf{X}$  of such a QSD, and the  $v \times (b+1)$  synthesis operator of the ETF synthesis operator it generates. In particular, in Theorem 3 we characterize when there exists scalars  $\delta$  and  $\varepsilon$  such that  $\Phi = [\mathbf{1} \ \delta\mathbf{J} + \varepsilon\mathbf{X}^T]$  is the synthesis operator of an ETF.

Combining results from the existing literature reveals a second connection between QSDs and ETFs. In particular, [27] establishes an equivalence between a certain class of QSDs and self-complementary binary codes that achieve equality in the Grey-Rankin bound. In [24], it was shown that these same codes are equivalent to real flat ETFs. Together, these two results imply an equivalence between real flat ETFs and a particular class of QSDs. As we shall see, it turns out that this equivalence is a special case of that described in the previous paragraph, namely when the scalars  $\delta$  and  $\varepsilon$  can be chosen to be 1 and  $-2$ , respectively. The analysis here is delicate: a QSD(15, 3, 1, 7, 35, 0, 1) exists and yields a real ETF with  $(d, n) = (15, 36)$ . However, as we shall see, such an ETF cannot be flat. Meanwhile, a QSD(6, 2, 1, 5, 15, 0, 1) exists and yields a real flat ETF with  $(d, n) = (6, 16)$ .

**Theorem 3.** *For any  $\delta, \varepsilon \in \mathbb{R}$  and any  $\{0, 1\}$ -valued  $b \times v$  matrix  $\mathbf{X}$  with  $b > v > 1$ , let*

$$\Phi = [\mathbf{1} \ \delta\mathbf{J} + \varepsilon\mathbf{X}^T].$$

*There exists a choice of  $\delta, \varepsilon \in \mathbb{R}$  such that the columns  $\{\varphi_j\}_{j=1}^{b+1}$  of  $\Phi$  form an ETF for  $\mathbb{R}^v$  with  $\langle \varphi_1, \varphi_j \rangle > 0$  for all  $j$  if and only if  $\mathbf{X}$  is the incidence matrix of a QSD( $v, k, \lambda, r, b, x, y$ ) with  $0 < k < v$  whose parameters satisfy the following relationships:*

$$w = \left[ \frac{v(b+1-v)}{b} \right]^{\frac{1}{2}}, \quad r = \frac{bk}{v}, \quad \lambda = \frac{r(k-1)}{v-1}, \quad x = k - \frac{(v+w)(r-\lambda)}{b+1}, \quad y = k - \frac{(v-w)(r-\lambda)}{b+1}. \quad (24)$$

Specifically, there are two choices for  $(\delta, \varepsilon)$  here:

$$\delta = \frac{1}{v}[w \pm k(\frac{b+1}{r-\lambda})^{\frac{1}{2}}], \quad \varepsilon = \frac{1}{k}(w - \delta v) = \mp(\frac{b+1}{r-\lambda})^{\frac{1}{2}}. \quad (25)$$

In particular, if  $n-1 > d > 1$  and  $\Phi$  is any  $\{\pm 1\}$ -valued  $d \times n$  matrix whose rows and columns have been signed to assume without loss of generality that  $\Phi = [\mathbf{1} \mathbf{J} - 2\mathbf{X}^T]$  where  $\mathbf{X}$  is  $\{0, 1\}$ -valued and  $\langle \varphi_1, \varphi_j \rangle \geq 0$  for all  $j$ , then the columns of  $\Phi$  form an ETF for  $\mathbb{R}^d$  if and only if  $\mathbf{X}$  is the incidence matrix of a QSD( $v, k, \lambda, r, b, x, y$ ) with  $v = d$ ,  $b = n - 1$ , and  $k = \frac{v-w}{2}$  where  $w, r$  and  $\lambda$  are given by (24) and  $x = \frac{v-3w}{4}$ ,  $y = \frac{v-w}{4}$ . In this case,  $d$  and  $w$  are necessarily even, and  $n$  is necessarily divisible by 4.

*Proof.* To simplify notation, let  $\mathbf{Z} = \mathbf{X}^T$ . For any  $\delta, \varepsilon \in \mathbb{R}$ , the fact that  $\|\varphi_1\|^2 = \|\mathbf{1}\|^2 = v$  implies that  $\{\varphi_j\}_{j=1}^{b+1}$  is an ETF for  $\mathbb{R}^v$  if and only if

$$|\langle \varphi_j, \varphi_{j'} \rangle| = \begin{cases} v, & j = j', \\ w, & j \neq j', \end{cases} \quad \forall j, j' = 1, \dots, b+1, \quad (26)$$

where  $w$ , as defined in (24), is obtained by scaling the Welch bound for  $n = b + 1$  and  $d = v$  by a factor of  $\|\varphi_1\|^2 = v$ . Letting  $\{\mathbf{z}_j\}_{j=1}^b$  denote the columns of  $\mathbf{Z}$ , we have  $\varphi_{j+1} = \delta \mathbf{1} + \varepsilon \mathbf{z}_j$  for all  $j = 1, \dots, b$ . Having the additional property that  $\langle \varphi_1, \varphi_j \rangle > 0$  for all  $j$  equates to having

$$w = \langle \varphi_1, \varphi_{j+1} \rangle = \langle \mathbf{1}, \delta \mathbf{1} + \varepsilon \mathbf{z}_j \rangle = \delta v + \varepsilon \langle \mathbf{1}, \mathbf{z}_j \rangle, \quad \forall j = 1, \dots, b.$$

If  $\varepsilon = 0$ , the columns of  $\Phi$  are collinear, meaning they are not a tight frame for  $\mathbb{R}^v$  since  $v > 1$ . When  $\varepsilon \neq 0$ , the above equation gives that  $\langle \mathbf{1}, \mathbf{z}_j \rangle = \frac{1}{\varepsilon}(w - \delta v)$  for all  $j = 1, \dots, b$ , meaning in particular that each column of  $\mathbf{Z}$  contains exactly  $k = \frac{1}{\varepsilon}(w - \delta v)$  ones. Here,  $k$  is an integer satisfying  $0 < k < v$  since having either  $k = 0$  or  $k = v$  again implies that the columns of  $\Phi$  are collinear. Moreover, in this case, the vectors  $\{\varphi_j\}_{j=1}^{b+1} = \{\mathbf{1}\} \cup \{\delta \mathbf{1} + \varepsilon \mathbf{z}_j\}_{j=1}^b$  satisfy (26) if and only if for every  $j, j' = 1, \dots, b, j \neq j'$  we have

$$v = \|\delta \mathbf{1} + \varepsilon \mathbf{z}_j\|^2 = \delta^2 \|\mathbf{1}\|^2 + 2\delta\varepsilon \langle \mathbf{1}, \mathbf{z}_j \rangle + \varepsilon^2 \|\mathbf{z}_j\|^2 = \delta^2 v + (2\delta\varepsilon + \varepsilon^2)k, \quad (27)$$

$$w = |\langle \delta \mathbf{1} + \varepsilon \mathbf{z}_j, \delta \mathbf{1} + \varepsilon \mathbf{z}_{j'} \rangle| = |\delta^2 v + 2\delta\varepsilon k + \varepsilon^2 \langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle|. \quad (28)$$

To summarize, there exists  $\delta, \varepsilon \in \mathbb{R}$  such that  $\{\varphi_j\}_{j=1}^{b+1} = \{\mathbf{1}\} \cup \{\delta \mathbf{1} + \varepsilon \mathbf{z}_j\}_{j=1}^b$  is an ETF for  $\mathbb{R}^v$  with  $\langle \varphi_1, \varphi_j \rangle > 0$  for all  $j$  if and only if there exists  $\delta \in \mathbb{R}$  and an integer  $k, 0 < k < v$  such that letting  $\varepsilon = \frac{1}{k}(w - \delta v)$  we have that  $\delta$  and  $\varepsilon$  satisfy (27), that  $\{\mathbf{z}_j\}_{j=1}^b$  satisfies  $\langle \mathbf{1}, \mathbf{z}_j \rangle = k$  for all  $j$ , and that  $\{\mathbf{z}_j\}_{j=1}^b$  satisfies (28) for all  $j \neq j'$ . Continuing, note that since  $\varepsilon = \frac{w-\delta v}{k}$ , (27) becomes

$$v = \delta^2 v + (2\delta\varepsilon + \varepsilon^2)k = \delta^2 v + 2\delta(w - \delta v) + \frac{1}{k}(w - \delta v)^2,$$

which is equivalent to having

$$0 = \frac{1}{2}v\delta^2 - w\delta - \frac{vk-w^2}{2(v-k)}. \quad (29)$$

To find the roots of this equation, it helps to define  $r$  and  $\lambda$  as in (24); since  $0 < k < v$ , these definitions imply  $0 \leq \lambda < r < b$  with

$$0 < r - \lambda = (1 - \frac{k-1}{v-1})r = \frac{v-k}{v-1}r = \frac{bk(v-k)}{v(v-1)}.$$

This fact and the definition of  $w$  imply that the discriminant of (29) is the positive quantity

$$w^2 + v \frac{vk-w^2}{v-k} = \frac{k(v^2-w^2)}{v-k} = \frac{k}{v-k} [v^2 - \frac{v}{b}(b+1-v)] = \frac{kv(v-1)(b+1)}{b(v-k)} = k^2 \frac{b+1}{r-\lambda}, \quad (30)$$

meaning (29) has two real roots, each of which leads to a corresponding value of  $\varepsilon$ , namely those paired values of  $\delta$  and  $\varepsilon$  given in (25).

Having characterized when (27) is satisfied, we turn to (28): in light of (25) and (29),

$$\begin{aligned} \delta^2 v + 2\delta\varepsilon k + \varepsilon^2 \langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle &= \delta^2 v + 2\delta(w - \delta v) + \frac{b+1}{r-\lambda} \langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle \\ &= -2(\frac{1}{2}v\delta^2 - w\delta) + \frac{b+1}{r-\lambda} \langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle \\ &= -\frac{vk-w^2}{v-k} + \frac{b+1}{r-\lambda} \langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle. \end{aligned} \quad (31)$$

Repurposing (30) further gives

$$\frac{vk-w^2}{v-k} = \frac{v[v-(v-k)]-w^2}{v-k} = \frac{v^2-w^2}{v-k} - v = k \frac{b+1}{r-\lambda} - v.$$

Substituting this into (31), we see that (28) is satisfied when

$$v - k \frac{b+1}{r-\lambda} + \frac{b+1}{r-\lambda} \langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle = \delta^2 v + 2\delta\varepsilon k + \varepsilon^2 \langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle \in \{-w, w\},$$

namely when  $\langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle \in \{x, y\}$  where  $x$  and  $y$  are defined in (25).

To summarize, there exists  $\delta, \varepsilon \in \mathbb{R}$  such that  $\{\varphi_j\}_{j=1}^{b+1} = \{\mathbf{1}\} \cup \{\delta\mathbf{1} + \varepsilon\mathbf{z}_j\}_{j=1}^b$  is an ETF for  $\mathbb{R}^v$  with  $\langle \varphi_1, \varphi_j \rangle > 0$  for all  $j$  if and only if there exists an integer  $k$ ,  $0 < k < v$  such that  $\langle \mathbf{1}, \mathbf{z}_j \rangle = k$  for all  $j$  and such that  $\langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle \in \{x, y\}$  for all  $j \neq j'$ , under the definitions given in (25). In this case, there are two choices for  $(\delta, \varepsilon)$ , namely the values given in (25). This characterization is complete. However, it differs from the characterization given in the statement of the result since we have not yet used the fact that any vectors that meet the Welch bound are necessarily tight, a fact which pertains to the rows of  $\mathbf{Z}$ . In particular, under the above hypotheses, we have  $\Phi\Phi^* = \alpha\mathbf{I}$  where  $\alpha v = \text{Tr}(\alpha\mathbf{I}) = \text{Tr}(\Phi\Phi^*) = \text{Tr}(\Phi^*\Phi) = \sum_{j=1}^{b+1} \|\varphi_j\|^2 = (b+1)v$ , implying  $\alpha = b+1$ . As such,

$$(b+1)\mathbf{1} = \Phi\Phi^*\mathbf{1} = \sum_{j=1}^{b+1} \langle \varphi_j, \mathbf{1} \rangle \varphi_j = v\mathbf{1} + w \sum_{j=1}^b (\delta\mathbf{1} + \varepsilon\mathbf{z}_j) = (v + bw\delta)\mathbf{1} + w\varepsilon \sum_{j=1}^b \mathbf{z}_j.$$

This implies that each row of  $\mathbf{Z}$  sums to  $r' := \frac{1}{w\varepsilon}(b+1-v-bw\delta)$ ; since  $\mathbf{Z}$  is a  $\{0,1\}$ -valued  $v \times b$  matrix whose columns sum to  $k$ , we have  $vr' = bk = vr$  and so  $r' = r$ . This in turn implies that for any  $i, i' = 1, \dots, v$ ,  $i \neq i'$ ,

$$0 = (\Phi\Phi^*)(i, i') = 1 + \sum_{j=1}^b [\delta + \varepsilon\mathbf{Z}(i, j)][\delta + \varepsilon\mathbf{Z}(i', j)] = 1 + \delta^2 b + 2\delta\varepsilon r + \varepsilon^2 (\mathbf{Z}\mathbf{Z}^T)(i, i'),$$

meaning that the dot product of any two distinct rows of  $\mathbf{Z}$  is  $\lambda' = -\frac{1}{\varepsilon^2}(1 + \delta^2 b + 2\delta\varepsilon r)$ . As such,  $\mathbf{Z}$  is the  $v \times b$  incidence matrix of a BIBD( $v, k, \lambda'$ ) implying  $\lambda'(v-1) = k(r-1) = \lambda(v-1)$  and so  $\lambda' = \lambda$ . (The interested reader can also use (25) to directly show that  $r' = r$  and  $\lambda' = \lambda$ .)

As such, if there exists  $\delta, \varepsilon \in \mathbb{R}$  such that  $\{\varphi_j\}_{j=1}^{b+1} = \{\mathbf{1}\} \cup \{\delta\mathbf{1} + \varepsilon\mathbf{z}_j\}_{j=1}^b$  is an ETF for  $\mathbb{R}^v$  with  $\langle \varphi_1, \varphi_j \rangle > 0$ , then  $\mathbf{X} = \mathbf{Z}^T$  is the incidence matrix of a QSD( $v, k, \lambda, r, b, x, y$ ) with  $0 < k < v$ . Conversely, if  $\mathbf{X} = \mathbf{Z}^T$  is the incidence matrix of such a design, then  $\langle \mathbf{1}, \mathbf{z}_j \rangle = k$  for all  $j$  and

$\langle \mathbf{z}_j, \mathbf{z}_{j'} \rangle \in \{x, y\}$  for all  $j \neq j'$ , implying that  $\{\varphi_j\}_{j=1}^{b+1} = \{\mathbf{1}\} \cup \{\delta \mathbf{1} + \varepsilon \mathbf{z}_j\}_{j=1}^b$  is such an ETF provided  $\delta$  and  $\varepsilon$  are chosen according to (25).

For the next set of conclusions, let  $\Phi$  be any  $\{\pm 1\}$ -valued  $n \times d$  matrix where  $n - 1 > d > 1$ . To determine when the columns  $\{\varphi_j\}_{j=1}^n$  of  $\Phi$  form an ETF for  $\mathbb{R}^d$ , note that by signing the rows of  $\Phi$  we can assume without loss of generality that  $\varphi_1 = \mathbf{1}$ . Moreover, by signing the columns of  $\Phi$  we can further assume without loss of generality that  $\langle \varphi_1, \varphi_j \rangle \geq 0$  for all  $j$ . Any such matrix is of the form  $\Phi = [\mathbf{1} \mathbf{J} - 2\mathbf{Z}]$  where  $\mathbf{Z}$  is a  $\{0, 1\}$ -valued  $v \times b$  matrix where  $v = d$ ,  $b = n - 1$ . By what we have already seen,  $\{\varphi_j\}_{j=1}^n = \{\mathbf{1}\} \cup \{\mathbf{1} - 2\mathbf{z}_j\}_{j=1}^b$  is an ETF for  $\mathbb{R}^d$  if and only if there exists an integer  $k$ ,  $0 < k < v$  such that  $\mathbf{Z}^T$  is the incidence matrix of a QSD whose remaining parameters are given by (24), provided (25) allows  $(\delta, \varepsilon) = (1, -2)$ , that is, if and only if

$$1 = \delta = \frac{1}{v}[w + k(\frac{b+1}{r-\lambda})^{\frac{1}{2}}], \quad -2 = \varepsilon = -(\frac{b+1}{r-\lambda})^{\frac{1}{2}}.$$

This occurs precisely when  $\frac{b+1}{r-\lambda} = 4$  and  $v = w + 2k$ . These two conditions are redundant: if  $v = w + 2k$  then again repurposing (30) gives

$$\frac{b+1}{r-\lambda} = \frac{1}{k^2} \frac{k(v^2 - w^2)}{v - k} = \frac{v^2 - w^2}{k(v - k)} = \frac{(w + 2k)^2 - w^2}{k(w + k)} = 4.$$

As such, the columns of  $\Phi = [\mathbf{1} \mathbf{J} - 2\mathbf{Z}]$  form an ETF for  $\mathbb{R}^d$  if and only if letting  $k = \frac{v-w}{2}$  we have that  $\mathbf{Z}^T$  is the incidence matrix of a QSD whose remaining parameters are given by (24). Here, since  $k = \frac{v-w}{2}$  and  $\frac{b+1}{r-\lambda} = 4$ , our expressions for  $x$  and  $y$  simplify to

$$x = k - \frac{(v+w)(r-\lambda)}{b+1} = \frac{v-w}{2} - \frac{v+w}{4} = \frac{v-3w}{4}, \quad y = k - \frac{(v-w)(r-\lambda)}{b+1} = \frac{v-w}{2} - \frac{v-w}{4} = \frac{v-w}{4}.$$

In this case, we have  $n = b + 1 = 4(r - \lambda)$  is divisible by 4. Moreover, since  $x$  and  $y$  are both integers, then so are  $y - x = \frac{w}{2}$  and  $y + x = \frac{v}{2} - w$ , implying  $w$  and  $v$  are even integers.  $\square$

We note that Theorem 3 specifically excludes ETFs in which  $n = d + 1$ , that is, regular simplices. We did this because regular simplices are already well understood, and because the corresponding incidence matrices are square. In particular, as discussed in Section 2, the existence of a flat regular simplex for  $\mathbb{R}^d$  is equivalent to that of a Hadamard matrix of size  $d + 1$ . In this case, the arguments of the proof of Theorem 3 do not produce a QSD, but rather a symmetric design with  $k = \frac{v-1}{2}$  and  $x = \lambda = \frac{v-3}{4}$ . For example, to put the real flat tetrahedron of (3) in the form where Theorem 3 applies, we negate all but its first column, giving the following real flat ETF and corresponding symmetric design:

$$\Phi = \begin{bmatrix} 1 & 1 & -1 & 1 \\ 1 & -1 & 1 & 1 \\ 1 & 1 & 1 & -1 \end{bmatrix}, \quad \mathbf{X}^T = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}.$$

In this situation, no two distinct blocks intersect in  $y = \frac{v-1}{4}$  points, giving no guarantee that this is an integer and by extension no guarantee that  $v$  is even. In fact, since in this case there necessarily exists a Hadamard matrix of size  $v + 1$ , we have that  $v$  is necessarily odd. This explains how Theorem 3 is consistent with the dichotomy of the statement of Theorem A of [27]: for any real flat ETF for  $\mathbb{R}^d$ , we either have  $d$  is odd—meaning the ETF is a regular simplex arising from a Hadamard matrix of size  $d + 1$ —or  $d$  is even, meaning the ETF equates to a QSD whose parameters are given by Theorem 3. This in particular means that not all of the ETFs produced by Theorem 3



are flat: applying it to a QSD(15, 3, 1, 7, 35, 0, 1) gives a real ETF with  $(d, n) = (15, 36)$ , but since 15 is odd, this ETF is not flat.

We also mention that if a real ETF arises from the block graph of a QSD, then Theorem 3 can be applied to that QSD to produce an explicit ETF with those parameters. To be precise, for any QSD whose block graph's SRG parameters satisfy  $a = 2\mu$ , one can use (6), (20), (21) and (22) to show that  $y + x = 2[k - \frac{v(r-\lambda)}{b+1}]$  and  $(y - x)^2 = [\frac{2w(r-\lambda)}{b+1}]^2$ . (These calculations are nontrivial, and we performed them with the aid of a computer algebra system.) This implies  $y$  and  $x$  are of the form given in (24), meaning Theorem 3 can be applied to that QSD to produce an ETF with parameters  $(d, n) = (v, b + 1)$ . As noted above, these parameters match those of the ETF that arises from the QSD's strongly regular block graph via (23). In particular, one can show that a QSD yields a real ETF via its block graph if and only if its parameters satisfy (24).

#### 4.1. Quasi-symmetric designs from Kirkman ETFs and their complements.

We now combine Theorem 3 with previously-known results to produce other new results.

**Corollary 1.** *Let  $u$  be an integer with  $u \geq 2$ . There exists a real flat ETF with parameters  $(d, n) = (u(2u - 1), 4u^2)$  if and only if there exists a QSD with parameters*

$$(v, k, \lambda, r, b, x, y) = (2u^2 - u, u^2 - u, u^2 - u - 1, 2u^2 - u - 1, 4u^2 - 1, \frac{u(u-2)}{2}, \frac{u(u-1)}{2}). \quad (32)$$

*Such ETFs exist whenever there exists a Hadamard matrix of size  $u$ , or alternatively, whenever there exists a Hadamard matrix of size  $2u$  and  $u - 2$  mutually orthogonal Latin squares (MOLS) of size  $2u$ , such as when  $u = 6$ .*

*Similarly, there exists a real flat ETF with parameters  $(d, n) = (u(2u + 1), 4u^2)$  if and only if there exists a QSD with parameters*

$$(v, k, \lambda, r, b, x, y) = (2u^2 + u, u^2, u^2 - u, 2u^2 - u, 4u^2 - 1, \frac{u(u-1)}{2}, \frac{u^2}{2}). \quad (33)$$

*Such ETFs exist whenever there exists a Hadamard matrix of size  $u$ , or alternatively, whenever there exists a Hadamard matrix of size  $2u$  and  $u - 1$  MOLS of size  $2u$ , such as when  $u = 6$ .*

*In either case,  $u$  is necessarily even. Also, when there exists a Hadamard matrix of size  $u$ , these two ETFs can be chosen to be Naimark complements.*

*Proof.* If there exists a real flat ETF with parameters  $(d, n) = (u(2u - 1), 4u^2)$ , then applying Theorem 3 to it produces a QSD with parameters (32) and  $w = u$ . Conversely, applying Theorem 3 to any QSD with parameters (32) produces a real flat ETF with parameters  $(d, n) = (u(2u - 1), 4u^2)$ . Similarly, there exists a real flat ETF with parameters  $(u(2u + 1), 4u^2)$  if and only if there exists a QSD with parameters (33); here we again have  $w = u$ . In either case, note Theorem 3 requires  $w = u$  to be even.

If there exists a Hadamard matrix of size  $u$ , [4] gives a QSD with parameters (32). For any such  $u$ , [24] gives an independent method for constructing a real flat Kirkman ETF with parameters  $(u(2u - 1), 4u^2)$ . The real flat Naimark complement of it constructed in Theorem 2 has parameters  $(u(2u + 1), 4u^2)$ . As noted above, that ETF implies the existence of a QSD with parameters (33). If we instead have a Hadamard matrix of size  $2u$ , then Theorems 1 and 2 of [7] give QSDs with parameters (32) and (33) provided we also have  $u - 2$  or  $u - 1$  MOLS of size  $2u$ , respectively. These conditions hold, for example, when  $u = 6$ .  $\square$

As seen from this proof, the true novelty of Corollary 1 is the existence of QSDs with parameters (33) whenever there exists a Hadamard matrix of size  $u$ . To our knowledge, the only previously known QSDs with parameters (33) had either  $u = 2^e$  for some  $e \geq 0$  or  $u = 6$  [7]. As such, this result provides a new infinite family of such designs.

To be clear, Theorem 3 can be applied to any real flat ETF, and these do not necessarily have parameters of the form  $(u(2u - 1), 4u^2)$  or  $(u(2u + 1), 4u^2)$ . In particular, if there exists an RBIBD $(\hat{v}, \hat{k}, 1, \hat{r}, \hat{b})$  and Hadamard matrices of size  $\hat{r} + 1$  and  $\frac{\hat{v}}{\hat{k}}$ , then [24] gives a real flat Kirkman ETF with parameters  $(d, n) = (\hat{b}, \hat{v}(\hat{r} + 1))$ . When  $\hat{k} > 2$ , such an ETF is neither of the types characterized in Corollary 1. Applying Theorem 3 to it and simplifying the expressions for the parameters of the resulting QSD gives the following result:

**Corollary 2.** *If there exists an RBIBD $(\hat{v}, \hat{k}, 1, \hat{r}, \hat{b})$  and Hadamard matrices of size  $\hat{r} + 1$  and  $\frac{\hat{v}}{\hat{k}}$ , then there exists a QSD with parameters*

$$(v, k, \lambda, r, b, x, y) = \left( \hat{b}, \frac{\hat{v}(\hat{r}-1)}{2\hat{k}}, \frac{\hat{v}(\hat{r}-1)-2\hat{k}}{4}, \frac{(\hat{r}-1)(\hat{v}+\hat{k}-1)}{2}, \hat{r}(\hat{v} + \hat{k} - 1), \frac{\hat{v}(\hat{r}-3)}{4\hat{k}}, \frac{\hat{v}(\hat{r}-1)}{4\hat{k}} \right), \quad w = \frac{\hat{v}}{\hat{k}}.$$

Moreover, using a class of RBIBDs that were overlooked in [24], one can show that instances of such ETFs with  $\hat{k} > 2$  exist. To be precise, for any positive integer  $i$  and any prime power  $q$ , the finite projective space  $\text{PG}(2^{i+1} - 1, q)$  is resolvable [3] and so is an RBIBD with parameters

$$(\hat{v}, \hat{k}, \hat{\lambda}, \hat{r}, \hat{b}) = \left( \frac{q^{2^{i+1}}-1}{q-1}, q+1, 1, \frac{q^{2^{i+1}-1}-1}{q-1}, \frac{(q^2)^{2^i}-1}{q^2-1} \frac{q^{2^{i+1}-1}-1}{q-1} \right).$$

In order for Hadamard matrices of size  $\frac{\hat{v}}{\hat{k}}$  and  $\hat{r} + 1$  to exist, these quantities are necessarily divisible by 4, which happens when  $i \geq 2$  and  $q \equiv 1 \pmod{4}$ . In particular, taking  $i = 2$  and  $q = 5$  gives an RBIBD(97656, 6, 1, 19531, 317886556); since Hadamard matrices of size  $\frac{\hat{v}}{\hat{k}} = 16276$  and  $\hat{r} + 1 = 19532$  exist [11], this implies the existence of a real flat Kirkman ETF with parameters (317886556, 1907416992). Though these parameters are admittedly large, this to our knowledge is the first example of a real flat ETF whose parameters are not of the form  $(u(2u - 1), 4u^2)$  or  $(u(2u + 1), 4u^2)$ . Applying Theorem 3 to this ETF à la Corollary 2 gives a new QSD with parameters (317886556, 158943278, 476829831, 953659665, 1907416991, 79459432, 79467570), which is not of the form (32) or (33). Applying Theorem A of [27] to this QSD or alternatively Theorem 3 of [24] to this ETF gives a new code that achieves equality in the Grey-Rankin bound.

It is also noteworthy that Theorem 3 sometimes yields non-flat real ETFs with the same  $(d, n)$  parameters as flat ones. For example, one can show that QSD parameters  $(v, k, \lambda, r, b, x, y)$  with  $\lambda = 1$  satisfy (24) if and only if  $v = 2k^2 - k$ . An infinite number of such QSDs exist, being instances of known BIBD $(2k^2 - k, k, 1, 2k + 1, 4k^2 - 1)$ . This includes whenever  $k = 2^e$  for some  $e \geq 1$ —a type of *Denniston design*—and also whenever  $k = 3, 5, 6, 7$  [26]. For any such design with  $k > 2$ , applying Theorem 3 to it yields a non-flat ETF with parameters  $(d, n) = (2k^2 - k, 4k^2)$  and entries valued either 1,  $\delta = \frac{1}{2k-1}(1 \pm \sqrt{2k})$  or  $\delta + \varepsilon = \frac{1}{2k-1}[1 \mp (2k - 2)\sqrt{2k}]$ . At the same time, taking  $u = k$  in Corollary 1 gives a real flat ETF that also has  $(d, n) = (2k^2 - k, 4k^2)$  whenever there exists a Hadamard matrix of size  $k$  or when  $k = 6$ . In particular, there are an infinite number of instances in which two different QSDs give two different ETFs—one flat and the other not—with the same  $(d, n)$  parameters. We also see here that when  $k = 3, 5, 7$ , Theorem 3 produces ETFs with  $(d, n)$  being (15, 36), (45, 100) and (91, 196), respectively; since these values of  $d$  are odd, these ETFs cannot be flat.

#### 4.2. Necessary integrality conditions for Hadamard ETFs

We conclude this section by strengthening the necessary conditions on the existence of real flat ETFs given in Theorem 3 by combining them with other known necessary conditions on real and unital ETFs in the literature, specifically those given in [32]:

**Corollary 3.** *If  $n - 1 > d > 1$  and there exists a real flat ETF with parameters  $(d, n)$ , then*

$$\left[\frac{d(n-1)}{n-d}\right]^{\frac{1}{2}}, \quad \left[\frac{(n-d)(n-1)}{n-d}\right]^{\frac{1}{2}}, \quad \left[\frac{d(n-d)}{n-1}\right]^{\frac{1}{2}}, \quad (34)$$

are integers, and are necessarily odd, odd and even, respectively. Moreover,  $n$  is divisible by 16.

*Proof.* Under these assumptions we have  $n \neq 2d$ : otherwise Theorem 3 gives that

$$4w^2 = \frac{4d(n-d)}{n-1} = \frac{4d^2}{2d-1} = 2d + 1 + \frac{1}{2d-1}$$

is an integer divisible by 16, implying  $d = 1$ . Since  $n - 1 > d > 1$  where  $n \neq 2d$ , Theorem A of [32] gives that the first two quantities in (34) are indeed odd integers. Moreover, Theorem B of [32] gives that the third quantity in (34) is an integer. In fact, in the notation of Theorem 3, this third quantity is the parameter  $w$ , and so is necessarily even. (Thus Theorem 3 shows, for example, that real ETFs with parameters  $(d, n) = (15, 36)$  cannot be flat, despite satisfying all necessary conditions on such ETFs given in [32].) For the final conclusion, note that being an odd square,  $\frac{d(n-1)}{n-d}$  is congruent to either 1 or 9 modulo 16. In the first case, we thus have  $d(n-1) \equiv n-d \pmod{16}$  and so  $(d-1)n \equiv 0 \pmod{16}$ ; since  $d$  is even,  $d-1$  is a unit in  $\mathbb{Z}_{16}$ , implying  $n \equiv 0 \pmod{16}$ . Similarly, in the second case we have  $d(n-1) \equiv 9(n-d) \pmod{16}$  and so  $(d-9)n \equiv 8d \pmod{16}$ ; since  $d$  is even, this again implies  $n \equiv 0 \pmod{16}$ .  $\square$

## 5. Miscellanea

In this section, we present two other results regarding Hadamard ETFs. The first of these results show how, in a special case, we can take tensor products of two Hadamard ETFs to produce another. The second result generalizes the well-known Gerzon bound to provide new necessary conditions on such ETFs.

In Definition 1, we define a Hadamard ETF to be an ETF whose synthesis operator is a submatrix of a Hadamard matrix. One may also consider ETFs whose Gram matrices are related to Hadamard matrices. In particular, there is a well-known equivalence between symmetric Hadamard ETFs with constant diagonal and real ETFs whose parameters  $(d, n)$  satisfy  $d = \frac{1}{2}(n - \sqrt{n})$  [23]; see [10] for a review of the literature of such Hadamard matrices. This idea has also been generalized to the complex setting [6, 30].

Here, the idea is that if  $\{\varphi_j\}_{j=1}^n$  is an ETF for  $\mathbb{F}^d$  with  $d = \frac{1}{2}(n - \sqrt{n})$  and if we assume without loss of generality that  $\|\varphi_j\|^2 = \frac{2d}{\sqrt{n}}$  for all  $j$ , then  $\mathbf{H} = \sqrt{n}\mathbf{I} - \Phi^*\Phi$  is a self-adjoint possibly-complex Hadamard matrix whose diagonal entries are one. Indeed, the diagonal entries of  $\mathbf{H}$  are  $\sqrt{n} - \frac{2d}{\sqrt{n}} = 1$ , and the off-diagonal entries have modulus  $\frac{2d}{\sqrt{n}} \left[\frac{n-d}{d(n-1)}\right]^{\frac{1}{2}} = 1$ . At the same time, the fact that  $\Phi\Phi^* = \frac{n}{d} \frac{2d}{\sqrt{n}} \mathbf{I} = 2\sqrt{n}\mathbf{I}$  implies that the eigenvalues of  $\mathbf{H}$  are  $\pm\sqrt{n}$ , implying  $\mathbf{H}\mathbf{H}^* = n\mathbf{I}$ . Conversely, if  $\mathbf{H}$  is any such Hadamard matrix, then  $\mathbf{G} = \sqrt{n}\mathbf{I} - \mathbf{H}$  is positive-semidefinite having eigenvalues 0 and  $2\sqrt{n}$  and trace  $n(\sqrt{n} - 1)$ . Diagonalizing  $\mathbf{G}$  thus reveals it to be the Gram matrix of an ETF for  $\mathbb{F}^d$  where  $d$  is the multiplicity of  $2\sqrt{n}$ , namely  $d = \frac{1}{2}(n - \sqrt{n})$ .

Given any two self-adjoint possibly-complex Hadamard matrices whose diagonal entries are one, we can take their tensor product to construct another such matrix. As noted in [6, 30, 22], this fact along with the above equivalence implies that if there exists ETFs  $\{\varphi_j\}_{j=1}^{n_1}$  and  $\{\psi_j\}_{j=1}^{n_2}$  for  $\mathbb{F}^{d_1}$  and  $\mathbb{F}^{d_2}$ , respectively, where  $d_1 = \frac{1}{2}(n_1 - \sqrt{n_1})$  and  $d_2 = \frac{1}{2}(n_2 - \sqrt{n_2})$ , then there exists an ETF consisting of  $n_1 n_2$  vectors for  $\mathbb{F}^{d_3}$  where  $d_3 = \frac{1}{2}(n_1 n_2 - \sqrt{n_1 n_2})$ . We now present an alternative proof of this fact that constructs the synthesis operator of the resulting ETF explicitly. In the special case where  $\{\varphi_j\}_{j=1}^{n_1}$  and  $\{\psi_j\}_{j=1}^{n_2}$  are Hadamard ETFs, this construction implies that the  $n_1 n_2$ -vector ETF is as well.

**Theorem 4.** *Let  $\{\varphi_j\}_{j=1}^{n_1}$  and  $\{\psi_j\}_{j=1}^{n_2}$  be ETFs for  $\mathbb{F}^{d_1}$  and  $\mathbb{F}^{d_2}$ , respectively, where the ETF parameters satisfy  $d_1 = \frac{1}{2}(n_1 - \sqrt{n_1})$  and  $d_2 = \frac{1}{2}(n_2 - \sqrt{n_2})$ . Let  $\{\tilde{\varphi}_j\}_{j=1}^{n_1}$  and  $\{\tilde{\psi}_j\}_{j=1}^{n_2}$  be any Naimark complements of  $\{\varphi_j\}_{j=1}^{n_1}$  and  $\{\psi_j\}_{j=1}^{n_2}$  in  $\mathbb{F}^{n_1-d_1}$  and  $\mathbb{F}^{n_2-d_2}$ , respectively. Then*

$$\{(\varphi_j \otimes \tilde{\psi}_{j'}) \oplus (\tilde{\varphi}_j \otimes \psi_{j'})\}_{j=1, j'=1}^{n_1, n_2}, \quad \{(\varphi_j \otimes \psi_{j'}) \oplus (\tilde{\varphi}_j \otimes \tilde{\psi}_{j'})\}_{j=1, j'=1}^{n_1, n_2}, \quad (35)$$

are Naimark complementary ETFs for  $\mathbb{F}^{d_3}$  and  $\mathbb{F}^{n_1 n_2 - d_3}$ , respectively, where  $d_3 = \frac{1}{2}(n_1 n_2 - \sqrt{n_1 n_2})$ . In particular, if  $\{\varphi_j\}_{j=1}^{n_1}$  and  $\{\psi_j\}_{j=1}^{n_2}$  are Hadamard, then so are the ETFs in (35).

*Proof.* Without loss of generality,  $\|\varphi_j\|^2 = d_1$ ,  $\|\psi_j\|^2 = d_2$ ,  $\|\tilde{\varphi}_j\|^2 = n_1 - d_1$ , and  $\|\tilde{\psi}_j\|^2 = n_2 - d_2$  for all  $j$ . Letting  $\Phi$ ,  $\Psi$ ,  $\tilde{\Phi}$  and  $\tilde{\Psi}$  denote the synthesis operators for  $\{\varphi_j\}_{j=1}^{n_1}$ ,  $\{\psi_j\}_{j=1}^{n_2}$  and their Naimark complements, respectively, we have that

$$\Phi\Phi^* = n_1\mathbf{I}, \quad \tilde{\Phi}\tilde{\Phi}^* = n_1\mathbf{I}, \quad \Phi\tilde{\Phi}^* = \mathbf{0}, \quad \Psi\Psi^* = n_2\mathbf{I}, \quad \tilde{\Psi}\tilde{\Psi}^* = n_2\mathbf{I}, \quad \Psi\tilde{\Psi}^* = \mathbf{0}. \quad (36)$$

The synthesis operators of the vector sequences in (35) are the  $d_3 \times n_1 n_2$  and  $d_4 \times n_1 n_2$  matrices

$$\begin{bmatrix} \Phi \otimes \tilde{\Psi} \\ \tilde{\Phi} \otimes \Psi \end{bmatrix}, \quad \begin{bmatrix} \Phi \otimes \Psi \\ \tilde{\Phi} \otimes \tilde{\Psi} \end{bmatrix},$$

respectively, where the fact that  $d_1 = \frac{1}{2}(n_1 - \sqrt{n_1})$  and  $d_2 = \frac{1}{2}(n_2 - \sqrt{n_2})$  implies

$$\begin{aligned} d_3 &= d_1(n_2 - d_2) + (n_1 - d_1)d_2 = \frac{1}{2}(n_1 n_2 - \sqrt{n_1 n_2}), \\ d_4 &= d_1 d_2 + (n_1 - d_1)(n_2 - d_2) = \frac{1}{2}(n_1 n_2 + \sqrt{n_1 n_2}). \end{aligned}$$

Here, (36) implies

$$\begin{bmatrix} \Phi \otimes \tilde{\Psi} \\ \tilde{\Phi} \otimes \Psi \\ \Phi \otimes \Psi \\ \tilde{\Phi} \otimes \tilde{\Psi} \end{bmatrix} \begin{bmatrix} \Phi \otimes \tilde{\Psi} \\ \tilde{\Phi} \otimes \Psi \\ \Phi \otimes \Psi \\ \tilde{\Phi} \otimes \tilde{\Psi} \end{bmatrix}^* = \begin{bmatrix} \Phi\Phi^* \otimes \tilde{\Psi}\tilde{\Psi}^* & \Phi\tilde{\Phi}^* \otimes \tilde{\Psi}\Psi^* & \Phi\Phi^* \otimes \tilde{\Psi}\Psi^* & \Phi\tilde{\Phi}^* \otimes \tilde{\Psi}\tilde{\Psi}^* \\ \tilde{\Phi}\tilde{\Phi}^* \otimes \Psi\Psi^* & \tilde{\Phi}\tilde{\Phi}^* \otimes \Psi\Psi^* & \tilde{\Phi}\tilde{\Phi}^* \otimes \Psi\Psi^* & \tilde{\Phi}\tilde{\Phi}^* \otimes \Psi\Psi^* \\ \Phi\Phi^* \otimes \tilde{\Psi}\tilde{\Psi}^* & \Phi\tilde{\Phi}^* \otimes \tilde{\Psi}\Psi^* & \Phi\Phi^* \otimes \tilde{\Psi}\tilde{\Psi}^* & \Phi\tilde{\Phi}^* \otimes \tilde{\Psi}\tilde{\Psi}^* \\ \tilde{\Phi}\tilde{\Phi}^* \otimes \Psi\Psi^* & \tilde{\Phi}\tilde{\Phi}^* \otimes \Psi\Psi^* & \tilde{\Phi}\tilde{\Phi}^* \otimes \tilde{\Psi}\tilde{\Psi}^* & \tilde{\Phi}\tilde{\Phi}^* \otimes \tilde{\Psi}\tilde{\Psi}^* \end{bmatrix} = n_1 n_2 \mathbf{I}.$$

Since  $d_3 + d_4 = n_1 n_2$ , this shows that (35) indeed defines Naimark complementary tight frames. What remains is to show that one of these two sequences of vectors is equiangular. For the second

sequence in particular, Naimark complementarity implies

$$\begin{aligned}
& \langle (\varphi_j \otimes \psi_{j'}) \oplus (\tilde{\varphi}_j \otimes \tilde{\psi}_{j'}), (\varphi_{j''} \otimes \psi_{j''}) \oplus (\tilde{\varphi}_{j''} \otimes \tilde{\psi}_{j''}) \rangle \\
&= \langle \varphi_j, \varphi_{j''} \rangle \langle \psi_{j'}, \psi_{j''} \rangle + \langle \tilde{\varphi}_j, \tilde{\varphi}_{j''} \rangle \langle \tilde{\psi}_{j'}, \tilde{\psi}_{j''} \rangle \\
&= \langle \varphi_j, \varphi_{j''} \rangle \langle \psi_{j'}, \psi_{j''} \rangle + \begin{cases} n_1 - \langle \varphi_j, \varphi_{j''} \rangle, & j = j'' \\ -\langle \varphi_j, \varphi_{j''} \rangle, & j \neq j'' \end{cases} \begin{cases} n_2 - \langle \psi_{j'}, \psi_{j''} \rangle, & j' = j'' \\ -\langle \psi_{j'}, \psi_{j''} \rangle, & j' \neq j'' \end{cases} \\
&= \begin{cases} d_1 d_2 + (n_1 - d_1)(n_2 - d_2), & j = j'', j' = j'' \\ -(n_1 - 2d_1) \langle \psi_{j'}, \psi_{j''} \rangle, & j = j'', j' \neq j'' \\ -\langle \varphi_j, \varphi_{j''} \rangle (n_2 - 2d_2), & j \neq j'', j' = j'' \\ 2 \langle \varphi_j, \varphi_{j''} \rangle \langle \psi_{j'}, \psi_{j''} \rangle, & j \neq j'', j' \neq j'' \end{cases}
\end{aligned}$$

As such, this sequence is equiangular if  $|\langle \varphi_j, \varphi_{j''} \rangle| = \frac{1}{2}(n_1 - 2d_1) = \frac{1}{2}\sqrt{n_1}$  for all  $j \neq j''$  and  $|\langle \psi_{j'}, \psi_{j''} \rangle| = \frac{1}{2}(n_2 - 2d_2) = \frac{1}{2}\sqrt{n_2}$  for all  $j' \neq j''$ ; these hold since  $\{\varphi_j\}_{j=1}^{n_1}$ ,  $\{\psi_j\}_{j=1}^{n_2}$  are equiangular and achieve the Welch bound with  $d_1 = \frac{1}{2}(n_1 - \sqrt{n_1})$  and  $d_2 = \frac{1}{2}(n_2 - \sqrt{n_2})$ .  $\square$

As an example of the previous result, note that for  $n_1 = n_2 = 4$  and  $d_1 = d_2 = \frac{1}{2}(4 - \sqrt{4}) = 1$ , we can take  $\Phi = \Psi = [1 \ 1 \ 1 \ 1]$  and take  $\tilde{\Phi} = \tilde{\Psi}$  to be the  $3 \times 4$  matrix (3) formed by the remaining three rows of the canonical  $4 \times 4$  Hadamard matrix. Applying Theorem 4 then produces a Hadamard ETF with  $(d, n) = (6, 16)$ :

$$\begin{bmatrix} \Phi \otimes \tilde{\Phi} \\ \tilde{\Phi} \otimes \Phi \end{bmatrix} = \begin{bmatrix} 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 & 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 & -1 & -1 & -1 & -1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 & -1 & -1 & -1 & -1 & 1 & 1 & 1 & 1 \end{bmatrix}.$$

Repeatedly applying Theorem 4 to this and other resulting ETFs yields real flat ETFs with  $d = \frac{1}{2}(n - \sqrt{n})$  where  $n = 2^{2(e+1)}$  for any  $e \geq 0$ . Though Hadamard ETFs of this size are already well-known, this method of construction is trivial, making no use of the theory of difference sets or BIBDs. In light of Corollary 1, one may also be tempted to apply Theorem 4 to real flat ETFs arising from MOLS. This is challenging, since it seems the ETFs arising from the QSDs of [7] with parameters (32) and (33) are not Naimark complements by default. Moreover, this has little payoff: at best, such a combination of Theorem 4 and Corollary 1 gives a way to combine two QSDs with parameters (32) for some even integers  $u_1, u_2$  so as to produce another such QSD with  $u_3 = 2u_1u_2$ ; since 8 divides  $u_3$ , such QSDs are probably more easily obtained via a  $u_3 \times u_3$  Hadamard matrix.

### 5.1. The Gerzon bound for Hadamard ETFs

It has long been known [25] that if there exists  $n$  equiangular noncollinear lines in  $\mathbb{R}^d$  then we necessarily have  $n \leq \binom{d+1}{2}$ ; in the complex case we necessarily have  $n \leq d^2$ . An alternative proof of these facts is discussed in [33]. In [23] it is noted that in order for an ETF to exist its Naimark complement must also satisfy these restrictions, that is, we also need  $n \leq \binom{n-d+1}{2}$  and  $n \leq (n-d)^2$  in the real and complex cases, respectively. Moreover, in the case of complex unital ETFs, this upper bound can be strengthened so as to require  $n \leq d^2 - d + 1$  [32]. We now refine these ideas to obtain necessary conditions on the existence of possibly-complex Hadamard ETFs.

Given any  $\{\varphi_j\}_{j=1}^n$  in  $\mathbb{F}^d$ , consider their outer products  $\{\varphi_j\varphi_j^*\}_{j=1}^n$  which lie in the real Hilbert space of all self-adjoint matrices in  $\mathbb{F}^{d \times d}$ . The Frobenius inner product of any two such outer products is  $\langle \varphi_j\varphi_j^*, \varphi_{j'}\varphi_{j'}^* \rangle_{\text{Fro}} = \text{Tr}(\varphi_j\varphi_j^*\varphi_{j'}\varphi_{j'}^*) = \text{Tr}(\varphi_j^*\varphi_{j'}\varphi_{j'}^*\varphi_j) = |\langle \varphi_j, \varphi_{j'} \rangle|^2$ . In particular, if  $\{\varphi_j\}_{j=1}^n$  is any sequence of noncollinear equiangular vectors, then the  $n \times n$  Gram matrix of  $\{\varphi_j\varphi_j^*\}_{j=1}^n$  is  $(\beta^2 - \gamma^2)\mathbf{I} + \gamma^2\mathbf{J}$  for some scalars  $0 \leq \gamma^2 < \beta^2$ . Such a Gram matrix has rank  $n$ , implying  $\{\varphi_j\varphi_j^*\}_{j=1}^n$  is linearly independent. This implies that  $n$  is at most the dimension of  $\{\mathbf{B} \in \mathbb{F}^{d \times d} : \mathbf{B}^* = \mathbf{B}\}$ , which is  $\binom{d+1}{2}$  or  $d^2$  depending on whether  $\mathbb{F} = \mathbb{R}$  or  $\mathbb{F} = \mathbb{C}$ , respectively.

Now assume that  $\{\varphi_j\}_{j=1}^n$  is a flat ETF for  $\mathbb{F}^d$  where  $1 < d < n - 1$ . Here,  $\{\varphi_j\varphi_j^*\}_{j=1}^n$  is a sequence of linearly independent vectors in the subspace of  $\{\mathbf{B} \in \mathbb{F}^{d \times d} : \mathbf{B}^* = \mathbf{B}\}$  that consists of all such matrices with constant diagonals. Computing the dimensions of this subspace, we thus have  $n \leq \frac{1}{2}d^2 - \frac{1}{2}d + 1$  or  $n \leq d^2 - d + 1$  depending on whether  $\mathbb{F} = \mathbb{R}$  or  $\mathbb{F} = \mathbb{C}$ , respectively. If we further assume that  $\{\varphi_j\}_{j=1}^n$  is possibly-complex Hadamard, then letting  $\{\tilde{\varphi}_j\}_{j=1}^n$  in  $\mathbb{F}^{n-d}$  be a flat Naimark complement for it, we necessarily have  $n \leq \frac{1}{2}(n-d)^2 - \frac{1}{2}(n-d) + 1$  or  $n \leq (n-d)^2 - (n-d) + 1$  when  $\mathbb{F} = \mathbb{R}$  or  $\mathbb{F} = \mathbb{C}$ , respectively. Solving for  $n$  in these inequalities gives the following result:

**Theorem 5.** *Let  $\{\varphi_j\}_{j=1}^n$  be an ETF for  $\mathbb{F}^d$  where  $1 < d < n - 1$ .*

- (a) *If  $\mathbb{F} = \mathbb{C}$  and  $\{\varphi_j\}_{j=1}^n$  is flat then  $n \leq d^2 - d + 1$ .*
- (b) *If  $\mathbb{F} = \mathbb{C}$  and  $\{\varphi_j\}_{j=1}^n$  is complex Hadamard then  $d + d^{\frac{1}{2}} + 1 \leq n \leq d^2 - d + 1$ .*
- (c) *If  $\mathbb{F} = \mathbb{R}$  and  $\{\varphi_j\}_{j=1}^n$  is flat then  $n \leq \frac{1}{2}d^2 - \frac{1}{2}d + 1$ ;*
- (d) *If  $\mathbb{F} = \mathbb{R}$  and  $\{\varphi_j\}_{j=1}^n$  is Hadamard then  $d + (2d + \frac{1}{4})^{\frac{1}{2}} + \frac{3}{2} \leq n \leq \frac{1}{2}d^2 - \frac{1}{2}d + 1$ .*

We remark that the above bounds are achieved infinitely often in the complex setting: for any prime power  $q$ , there exists a  $d \times n$  harmonic ETF arising from a Singer difference set that has  $d = q + 1$  and  $n = q^2 + q + 1$  [31, 36], meaning  $n = d^2 - d + 1$ ; its  $q^2 \times (q^2 + q + 1)$  flat Naimark complement achieves equality in the lower bound.

## Acknowledgments

This work was partially supported by NSF DMS 1321779, AFOSR F4FGA06060J007 and AFOSR Young Investigator Research Program award F4FGA06088J001. The views expressed in this article are those of the authors and do not reflect the official policy or position of the United States Air Force, Department of Defense, or the U.S. Government.

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