

# A Fast Iterative Method for Solving the Eikonal Equation on Triangulated Surfaces

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## Motivation

In this project, we consider the numerical solution of the Eikonal equations, a special case of nonlinear Hamilton-Jacobi partial differential equations (PDEs), defined on a three dimensional surface with a scalar speed function:

$$H(\mathbf{x}, \Delta\phi) = |\Delta\phi(\mathbf{x})|^2 - \frac{1}{f^2(\mathbf{x})} = 0 \quad \forall \mathbf{x} \in S \subset \Omega$$

$S$  is a surface domain. The solution of this equation simulates travel time of the wave propagation with speed  $f$  at  $\mathbf{x}$  from some source points whose values are zero. The Eikonal equation appears in various Applications, such as computer vision, image processing, computer graphics, geoscience, and medical image analysis.

## Background

### 1.Mesh Fast Iterative Method(meshFIM) [1]

◆An iterative computational technique to solve the Eikonal equation efficiently on parallel architectures.

◆This method relies on a modification of a label-correcting method.

◆The core elements for our FIM based method are:

- (1) Upwind scheme: calculate the value at a vertex with the values of the solved vertices.
- (2) Active list management: Active list contains the patches which has wave front vertices. If a active patch is convergent, it is removed from the Active list and its neighbor patches are added to this list.
- (3) Patch-based iteration: divide the whole mesh into patches to fit into GPU cores.
- (4) Triangle-based Jacobi update: update all the triangles inside a patch concurrently with parallel threads and each thread updates values of the three triangle vertices.

### 2.Method description

- (1) Firstly, partition the mesh into patches.
- (2) Add the patches which contain the source vertices to active list.
- (3) Assign each patch to a GPU stream processor and iterate multiple times for each patch.
- (4) Then check if a patch is convergent which means all the vertices of this patch are convergent. Remove a convergent patch from the active list and add its neighbor patches.
- (5) Check if the patches in active list are already convergent, if so remove.
- (6) Iterate again.

### 3.Suitability for GPU

- ◆Each vertex updates independently
- ◆According to the algorithm, update operation can be completed concurrently
- ◆Computing only depends on the neighbors of same facet at every time step

## Implementation

### 1.Partition

◆In the process of partitioning, we will use edges instead of coordinates, thus our partition can be viewed as the graph-based partition

◆We use METIS [2] as partition tool (See the figure below for a partition result of a dragon)

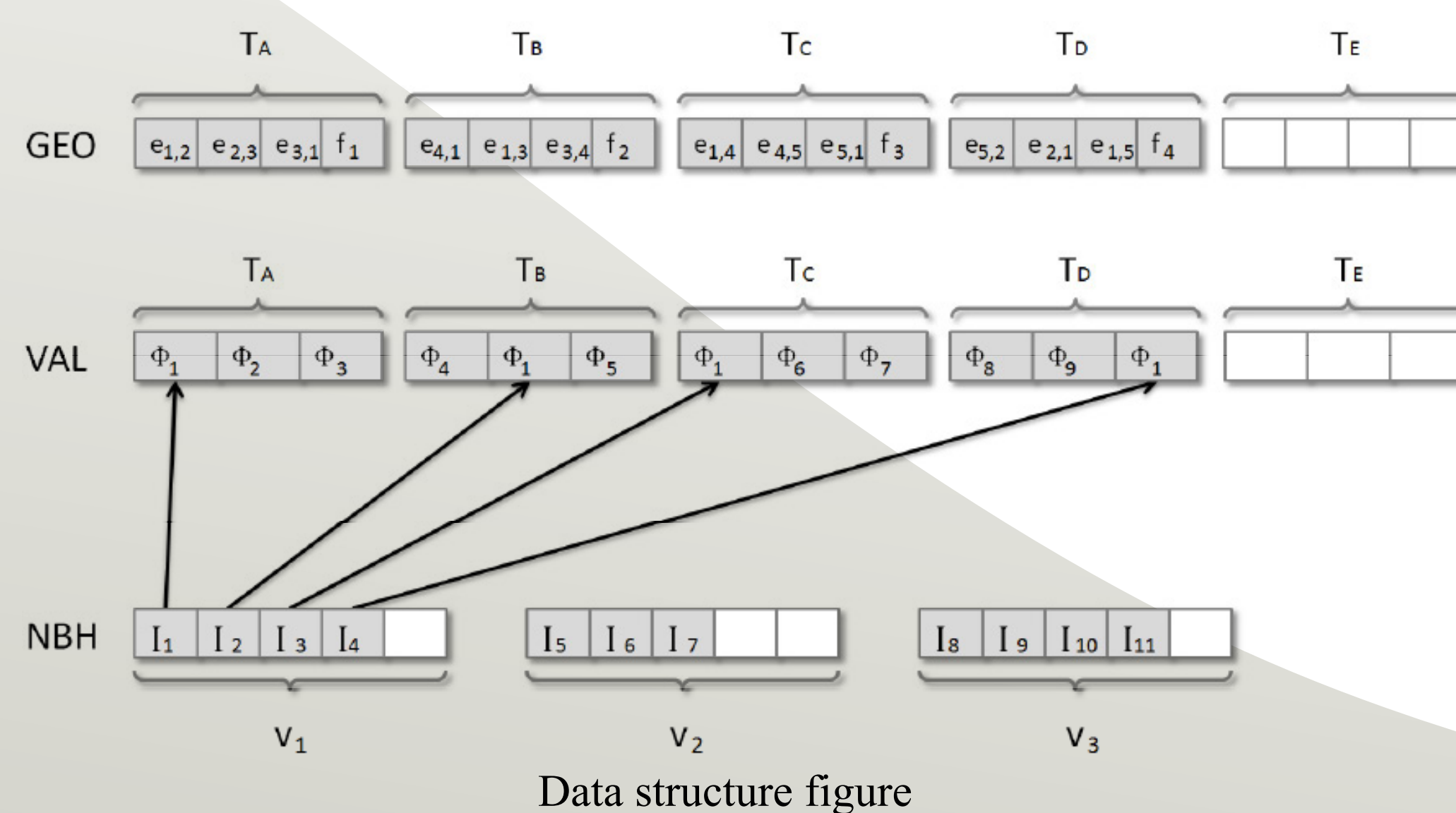


### 2.Triangle-based data structure

◆GEO: divided into sub segment for each patch and each patch subsegment contains geometric data and speed information for each triangle: three floats for edge lengths of the triangle and one float for speed.

◆VAL: hold all the vertex values(float) of all triangles patch by patch.

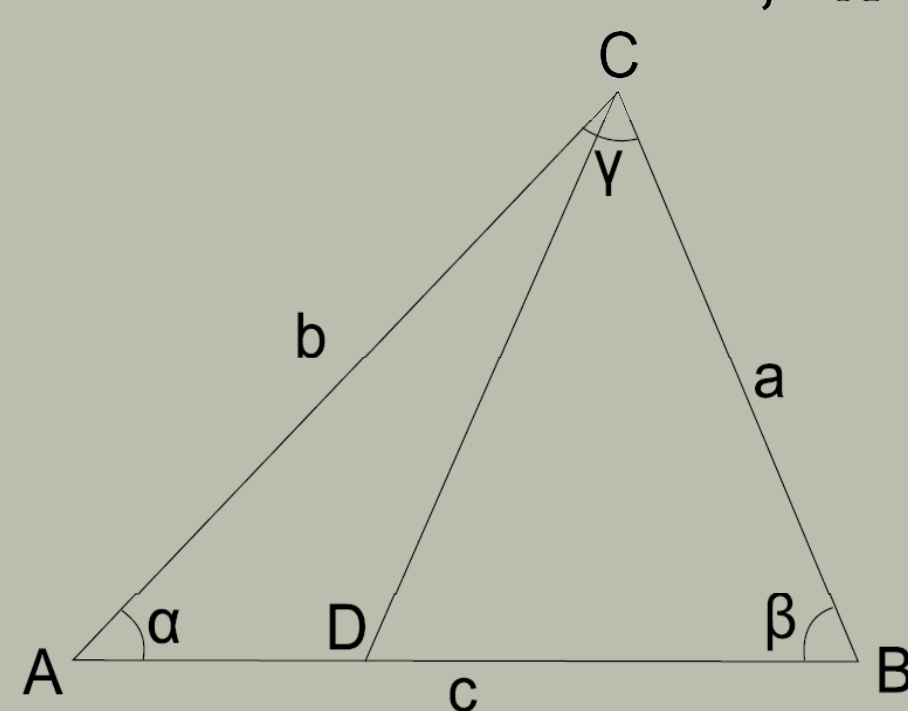
◆NBH: an integer array with each integer element representing an index of a vertex value in the value array.



◆GEO is in global memory and VAL and NBH are copied into shared memory for multiple updates.

### 3.Localsolver

As in the figure below, local solver calculate the value of a vertex of the triangle  $\Delta ABC$  from the other two vertices. Without loss of generality we only talk about calculating value of  $C$ ,  $T_C$ , from values of  $A$  and  $B$ ,  $T_A, T_B$ .  $f$  is the speed.



Assume wave propagation direction is from  $D$  to  $C$ , and  $T_D = \lambda(T_{AB}) + T_A$  we get:

$$\begin{aligned} T_C &= T_D + T_{DC} \\ &= \lambda(T_{AB}) + T_A + f\|DC\| \\ &= \lambda(T_{AB}) + T_A + f\|\vec{AC} - \lambda\vec{AB}\| \end{aligned}$$

And the location of  $D$  must minimize  $T_C$ , so let:  $\frac{dT_C(\lambda)}{d\lambda} = 0$ , We can solve for  $\lambda$  and then substitute into above equation to get  $T_C$ .

## Algorithm

Algorithm 3.2: PATCHFIM( $VAL_{in}, VAL_{out}, L, P$ )

comment:  $L$ : active list of patches,  $P$ : set of all patches  
while  $L$  is not empty  
do {MainUpdate( $L, C_v, VAL_{in}, VAL_{out}$ )  
  do {CheckNeighbor( $L, C_v, C, VAL_{in}, VAL_{out}$ )  
    do {UpdateActiveList( $L, P, C$ )

Algorithm 3.3: MAINUPDATE( $L, C_v, VAL_{in}, VAL_{out}$ )

comment: 1. Main iteration  
for each  $p \in L$  in parallel  
  for  $i = 1$  to  $n$   
    do {for each  $t \in p$  in parallel  
      do { $VAL_{out}(t) \leftarrow$  LocalSolver( $VAL_{in}(t)$ )  
        do {reconcile solutions in  $t$   
          do {update  $C_v(p)$   
            do {swap  $VAL_{in}(t)$  and  $VAL_{out}(t)$   
              do {reconcile solutions in  $p$

Algorithm 3.4: CHECKNEIGHBOR( $L, C_v, C, VAL_{in}, VAL_{out}$ )

comment: 2. Check neighbors  
for each  $p \in L$  in parallel  
  do { $C(p) \leftarrow$  reduction( $C_v(p)$ )  
  for each  $p \in L$  in parallel  
    do {if  $C(p) = \text{TRUE}$   
      then {for each adjacent neighbor of  $p_{nb}$  of  $p$   
        do {add  $p_{nb}$  to  $L$

for each  $p \in L$  in parallel  
  for each  $t \in p$  in parallel  
    do { $VAL_{out}(t) \leftarrow$  LocalSolver( $VAL_{in}(t)$ )  
      do {reconcile solutions in  $t$   
      do {update  $C_v(p)$   
      do {swap  $VAL_{in}(t)$  and  $VAL_{out}(t)$   
      do {reconcile solutions in  $p$

for each  $p \in L$  in parallel  
  do { $C(p) \leftarrow$  reduction( $C_v(p)$ )

Algorithm 3.5: UPDATEACTIVELIST( $L, P, C$ )

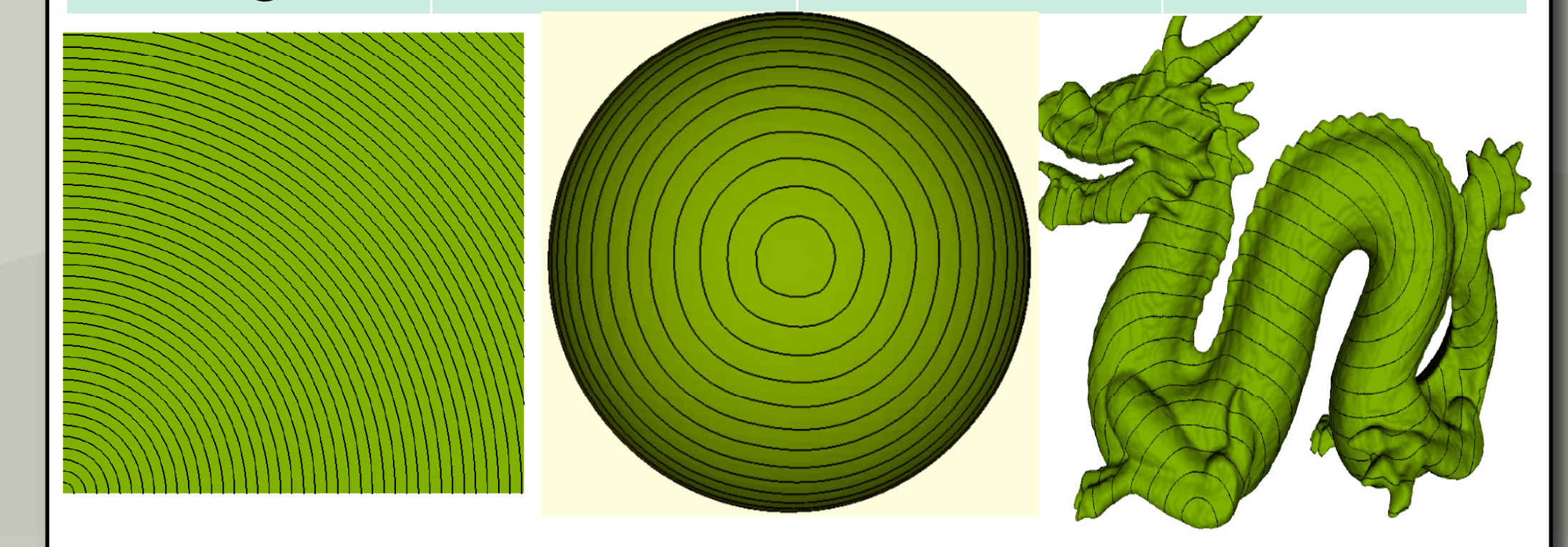
comment: 3. Update active list

clear( $L$ )  
for each  $p \in P$   
  do {if  $C(p) = \text{FALSE}$   
    then insert  $p$  to  $L$

## Result

- ◆CPU: Intel i7 920, 2.66GHz, 8M cache
  - ◆GPU: Nvidia GTX 275, 1.404GHz, 240 core
- We test running time(ms) for a CPU version of meshFIM to compare with GPU version on three different meshes:

Mesh	CPU	GPU	Speedup
Square	6562	201	33x
Sphere	8591	415	21x
Dragon	4331	287	15x



## References

1. A Fast Iterative Method For Eikonal Equations. Won-Ki Jeong, Ross Whitaker.
2. METIS: A Family of Multilevel Partitioning Algorithms. <http://glaros.dtc.umn.edu/gkhome/views/metis>