# Data Security Using Multi-bit LSB and Modified Vernam Cipher 

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#### Abstract

Data security is one of the most important aspects of today's information era. Some methods are used to secure important data from hackers. The LSB is a steganographic algorithm that is often used to store data in the last bit. In order to improve the security, we combine steganography with cryptography enables. In this research LSB is modified using the multi-bit LSB model. Modifications are made to the bits of each character, the rotation by a certain amount can randomize the plaintext content before cryptographic algorithm, Vernam is performed. The bit on LSB can be inserted data as much as $1,2,3$ or 4 - bit information. The calculation results of MSE and PSNR values indicate that the use of 1-bit LSB is superior to that of 2-, 3-, or 4-bit LSB.


Keyword: Cryptography, LSB, Steganographic, Vernam Encryption


#### Abstract

Abstrak. Keamanan data adalah salah satu aspek terpenting dari era informasi saat ini. Beberapa metode digunakan untuk mengamankan data penting dari peretas. LSB adalah algoritma steganografi yang sering digunakan untuk menyimpan data dalam bit terakhir. Untuk meningkatkan keamanan, kami menggabungkan steganografi dengan memungkinkan kriptografi. Dalam penelitian ini LSB dimodifikasi menggunakan model LSB multi-bit. Modifikasi dilakukan pada bit masing-masing karakter, rotasi dengan jumlah tertentu dapat mengacak konten plaintext sebelum algoritma kriptografi, Vernam dilakukan. Bit pada LSB dapat dimasukkan data sebanyak 1, 2, 3 atau 4 - bit informasi. Hasil perhitungan nilai MSE dan PSNR menunjukkan bahwa penggunaan LSB 1-bit lebih unggul daripada LSB 2-, 3-, atau 4-bit.


Kata Kunci: Cryptography, LSB, Steganographic, Vernam Encryption
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## 1. Introduction

According to the Cisco Visual Networking Index (VNI), the total IP traffic that occurred in 2016 is expected to reach 91.3 exabytes per month. Given this amount of traffic, data security is becoming increasingly important, particularly for sensitive data, such as company data and state security information [1]. Digital information is indirectly transmitted through a data network via

[^0]a small electric current that is used as a link to analog signals [2]. Data often comprise important information that must be properly secured to prevent theft by certain parties [3]-[5].

Steganography and cryptography are two methods of securing data. Steganography conceals data, whereas cryptography secures data by encoding the plaintext. Acts of crime are becoming more advanced, and consequently, these methods often fail. Various types of data theft are performed to solve well-encrypted ciphertexts; hence, the power of steganography and cryptography must be increased through various methods.

The least significant bit (LSB) steganography technique functions by replacing the rightmost bit of each pixel according to the message bit. The vulnerability of LSB is the storage of plaintext bits in a row, which allows intruders to easily extract bits of confidential information. Herein, three color elements are alternately used from the red, green, and blue layers. The aim of this study is to improve data security by combining steganography and cryptography techniques.

The Vernam algorithm encrypts data by performing exclusive-OR (XOR) operations on each plaintext character [6], [7] in a given set of data. This algorithm is modified by rotating bits. After obtaining the cryptographic results, the ciphertext is hidden via the multi-bit LSB steganography technique on 24-bit imagery using several storage bit models. Steganography encrypts each bit of information into RGB color elements [9]-[12], and multi-bits can be used to increase security by creating variations of the LSB technique [12]-[15]. Each sample is tested for the mean square error (MSE) and peak signal-to-noise ratio (PSNR) values from stego-images, which are used to determine the most superior image with the smallest error rate; this superior image is used as the storage image. In this manner, the above described hybrid technique can be used to improve data security.

## 2. Cryptographic and Steganography Algorithms

### 2.1 Cryptographic Algorithm: Vernam Algorithm

The Vernam cipher is an algorithm based on the principle that each character in the plaintext is encrypted using the XOR process against a certain generated key [16], [17]. The provided key is generated repeatedly or extended to many plaintext lengths. In the Vernam cipher algorithm, a symmetric key type signifies that the same key is used for both encryption and decryption. The Vernam cipher's vulnerability is the use of XOR for plaintext encryption and decryption because it is very simple. The following is an example of message encryption performed using the Vernam algorithm.

## Encryption process:

```
Plaintext: "algorithm"
Row of bits:
```

```
0 1 1 0 0 0 0 1 0 1 1 0 1 1 0 0 ~ 0 1 1 0 0 1 1 1 ~ 0 1 1 0 1 1 1 1 ~ 0 1 1 1 0 0 1 0 0 1 1 0 1 0 0 1 ~ 0 1 1 1 0 1 0 0 ~ 0 1 1 0 1 0 0 0 0 1 1 0 1 1 ,
01
Key:
1 0 1 1 1 0 0 1 0 1 1 1 0 1 0 0 ~ 0 0 1 1 1 0 1 0 ~ 0 0 0 1 0 1 0 1 ~ 0 1 1 0 1 0 1 0 ~ 1 1 0 1 0 1 0 0 ~ 1 0 1 0 1 0 1 0 ~ 1 0 0 0 1 1 1 0 ~ 1 0 1 0 1 0 ~
10
Ciphertext:
1 1 0 1 1 0 0 0 ~ 0 0 0 1 1 0 0 0 ~ 0 1 0 1 1 1 0 1 ~ 0 1 1 1 1 0 1 0 ~ 0 0 0 1 1 0 0 0 ~ 1 0 1 1 1 1 0 1 ~ 1 1 0 1 1 1 1 0 ~ 1 1 1 0 0 1 1 0 ~ 1 1 0 0 0 1 ~
11
```


## Decryption process:

```
Ciphertext:
1 1 0 1 1 0 0 0 0 0 0 1 1 0 0 0 ~ 0 1 0 1 1 1 0 1 ~ 0 1 1 1 1 0 1 0 ~ 0 0 0 1 1 0 0 0 ~ 1 0 1 1 1 1 0 1 ~ 1 1 0 1 1 1 1 0 ~ 1 1 1 0 0 1 1 0 ~ 1 1 0 0 0 1 )
11
Key:
1 0 1 1 1 0 0 1 0 1 1 1 0 1 0 0 ~ 0 0 1 1 1 0 1 0 ~ 0 0 0 1 0 1 0 1 ~ 0 1 1 0 1 0 1 0 1 1 0 1 0 1 0 0 ~ 1 0 1 0 1 0 1 0 1 0 0 0 1 1 1 0 1 0 1 0 1 0 ~
10
-----------------------------------------------------------------------------XOR
Row of bits:
0 1 1 0 0 0 0 1 0 1 1 0 1 1 0 0 0 1 1 0 0 1 1 1 0 1 1 0 1 1 1 1 0 1 1 1 0 0 1 0 0 1 1 0 1 0 0 1 0 1 1 1 0 1 0 0 0 1 1 0 1 0 0 0 0 1 1 0 1 1 )
01
Plaintext: "algorithm"
```


### 2.2 Steganography Least Significant Bit (LSB)

Steganography, which is a Greek term that means "closed writing," is the technique of communicating by hiding the existence of messages. Steganography plays an important role in information security, and a steganography system comprises three elements: the cover image, secret message, and stego-image. Digital images are represented using X and Y coordinates that contain three color elements in each pixel. Typically, a gray image uses 8 bits, whereas a color image uses 24 bits to describe the RGB color model.

Several techniques exist for hiding information in the cover image. Spatial domain techniques manipulate pixel bit values to embed confidential information, and message bits are directly encoded into the pixel bits of the cover image. Thus, spatial domain techniques can be easily implemented [18].

LSB is the position of a bit in a binary integer that provides a unit value to determine whether a number is even or odd. Bit numbering is performed from 0 to 7 to represent the intensity value of a color. The LSB is the smallest or rightmost bit. The LSBs can quickly change if the amount changes slightly. For example, if a 1 -bit LSB changes $(00000000)$ to $(00000001)$, the result is a 1-bit jump, and if a 2-bit LSB changes $(0000000)$ to ( 00000011 ), the result is a 3-bit jump. If a 3-
bit LSB changes $(00000000)$ to $(00000111)$, the result is a 7 -bit jump, whereas if a 4 -bit LSB changes $(00000000)$ to $(00001111)$, the result is a 15 -bit jump.

### 2.3 MSE and PSNR Measurement

Fidelity is a security element in the practice of hiding secret messages. The measurement of fidelity steganography can be performed by calculating the values of the MSE and PSNR. The PSNR is usually measured in decibels (dB). To determine the PSNR, first, the average value of the square of the MSE error must be determined. If the MSE value is small, a large PSNR value is produced, and vice versa. Images with a PSNR value $\geq 40 \mathrm{~dB}$ are considered high-quality images [18].

The MSE value can be calculated using the following formula:

$$
\begin{equation*}
M S E=\sum_{i=1}^{M} \sum_{j=1}^{N}(f(i, j)-g(i, j))^{2} / M * N \tag{1}
\end{equation*}
$$

Where:

MSE $=$ MSE value from image
$M=$ Length of image (in pixels)
$\mathrm{N}=$ width of image (in pixels)
$f(i, j)=$ Value of the pixel coordinates of the image before the message is inserted .After calculating the MSE value, the PSNR value can be calculated as follows

$$
\begin{equation*}
P S N R=10 * \log _{10}\left(\frac{C_{m a x}^{2}}{M S E}\right) \tag{2}
\end{equation*}
$$

Where:
PSNR $=\mathrm{PSNR}$ value of the image, and
$C_{\text {max }}=$ the maximum pixel value.
For RGB color images with values in each pixel, the MSE value is calculated for each color, and then all such values are summed and divided by 3 [19].

## 3. Methodology

In this section, after analyzing the two abovementioned data security techniques- the modified Vernam cipher cryptography and LSB multi-bit steganography techniques-the stages of system design are discussed as follows.

### 3.1 Data Input



Novell NetWare adalah sebuah sistem operasi jaringan Novell Open Server (NOS) yang menyediakan akses file remote transparan dan berbagai layanan jaringan terdistribusi, termasuk sharing printer dan dukungan untuk berbagai aplikasi seperti surat elektronik transfer dan akses database. NetWare menentukan atas lima lapisan model referensi OSI dan berjalan pada media apapun-protokol akses (Layer 2). Selain itu, NetWare berjalan pada hampir semua jenis sistem komputer, dari PC ke mainframe.

Figure 1. Retrieval of .docx Content
Input in the form of a.docx file is limited to characters that can be changed in ASCII code, however the system cannot encrypt tables, images, or symbols. As illustrated in Fig. 1,.docx content refers to the contents of a.docx file in the form of words. The key input is also limited to $\mathrm{a}-\mathrm{Z}$ and $\mathrm{A}-\mathrm{Z}$ characters.

### 3.2 Vernam Chiper Modification

This section discusses the modification of the Vernam cipher algorithm through bit rotation. Before the encryption process is completed, bit rotates the character and key bit positions to the left and right, respectively, and then performs the XOR operation. One character contains 8 bits, and bit rotation is performed from the $1^{\text {st }}$ bit to the $8^{\text {th }}$ bit (this can be adjusted as needed). Because one character contains 8 bits, if the rotation is performed more than eight times, the bit position repeats; i.e., 9 -bit rotation is the same as 1 -bit rotation, whereas 0 -bit rotation is the same as 8 -bit rotation.

During decryption, the first step is the XOR operation between the ciphertext and the key, which is performed via right-bit rotation. The XOR results are then rotated to the right for the message to return to its original sequence (plaintext).

1) Process of modifying the Vernam algorithm for encryption

## Plaintext:

| P 0 | P 1 | P 2 | P 3 | P 4 |
| :--- | :--- | :--- | :--- | :--- |

Bit plaintext:

| B07 | B06 | B05 | B04 | B03 | B02 | B01 | B00 |  | P0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B17 | B16 | B15 | B14 | B13 | B12 | B11 | B10 |  | P1 |
| B27 | B26 | B25 | B24 | B23 | B22 | B21 | B20 |  | P2 |
| B37 | B36 | B35 | B34 | B33 | B32 | B31 | B30 |  | P3 |
| B47 | B46 | B45 | B44 | B43 | B42 | B41 | B40 | $\longrightarrow$ | P4 |

1-bit rotation to left:

| B06 | B05 | B04 | B03 | B02 | B01 | B00 | B07 | P0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B16 | B15 | B14 | B13 | B12 | B11 | B10 | B17 | P1 |
| B26 | B25 | B24 | B23 | B22 | B21 | B20 | B27 | P2 |
| B36 | B35 | B34 | B33 | B32 | B31 | B30 | B37 | P3 |
| B46 | B45 | B44 | B43 | B42 | B41 | B40 | B47 | P4 |

Key:


Bit key:

| K07 | K06 | K05 | K04 | K03 | K02 | K01 | K00 | $\longrightarrow$ | K0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| K17 | K16 | K15 | K14 | K13 | K12 | K11 | K10 |  |  | K1 |
| K27 | K26 | K25 | K24 | K23 | K22 | K21 | K20 | $\longrightarrow$ |  | K2 |
| K37 | K36 | K35 | K34 | K33 | K32 | K31 | K30 | $\longrightarrow$ | K3 |  |
| K47 | K46 | K45 | K44 | K43 | K42 | K41 | K40 | $\longrightarrow$ | K4 |  |

1-bit rotation to the right:

| K00 | K07 | K06 | K05 | K04 | K03 | K02 | K01 | $\longrightarrow$ | K0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| K10 | K17 | K16 | K15 | K14 | K13 | K12 | K11 | $\longrightarrow$ |  |
| K20 | K27 | K26 | K25 | K24 | K23 | K22 | K21 | $\longrightarrow$ |  |
| K30 | K37 | K36 | K35 | K34 | K33 | K32 | K31 | $\longrightarrow$ |  |
| K2 |  |  |  |  |  |  |  |  |  |
| K40 | K47 | K46 | K45 | K44 | K43 | K42 | K41 | $\longrightarrow$ | K3 |

Ciphertext / XOR Operation:

| $\mathrm{B} 06 \oplus \mathrm{~K} 00$ | $\mathrm{~B} 05 \oplus \mathrm{~K} 07$ | $\mathrm{~B} 04 \oplus \mathrm{~K} 06$ | $\mathrm{~B} 03 \oplus \mathrm{~K} 05$ | $\mathrm{~B} 02 \oplus \mathrm{~K} 04$ | $\mathrm{~B} 01 \oplus \mathrm{~K} 03$ | $\mathrm{~B} 00 \oplus \mathrm{~K} 02$ | $\mathrm{~B} 07 \oplus \mathrm{~K} 01$ | $\longrightarrow$ | C 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{~B} 16 \oplus \mathrm{~K} 10$ | $\mathrm{~B} 15 \oplus \mathrm{~K} 17$ | $\mathrm{~B} 14 \oplus \mathrm{~K} 16$ | $\mathrm{~B} 13 \oplus \mathrm{~K} 15$ | $\mathrm{~B} 12 \oplus \mathrm{~K} 14$ | $\mathrm{~B} 11 \oplus \mathrm{~K} 13$ | $\mathrm{~B} 10 \oplus \mathrm{~K} 12$ | $\mathrm{~B} 17 \oplus \mathrm{~K} 11$ | $\longrightarrow$ | $\longrightarrow \mathrm{C} 1$ |
| $\mathrm{~B} 26 \oplus \mathrm{~K} 20$ | $\mathrm{~B} 25 \oplus \mathrm{~K} 27$ | $\mathrm{~B} 24 \oplus \mathrm{~K} 26$ | $\mathrm{~B} 23 \oplus \mathrm{~K} 25$ | $\mathrm{~B} 22 \oplus \mathrm{~K} 24$ | $\mathrm{~B} 22 \oplus \mathrm{~K} 23$ | $\mathrm{~B} 20 \oplus \mathrm{~K} 22$ | $\mathrm{~B} 27 \oplus \mathrm{~K} 21$ | $\longrightarrow$ | $\longrightarrow \mathrm{C} 2$ |


| $\mathrm{B} 36 \oplus \mathrm{~K} 30$ | $\mathrm{~B} 35 \oplus \mathrm{~K} 37$ | $\mathrm{~B} 34 \oplus \mathrm{~K} 36$ | $\mathrm{~B} 33 \oplus \mathrm{~K} 35$ | $\mathrm{~B} 32 \oplus \mathrm{~K} 34$ | $\mathrm{~B} 31 \oplus \mathrm{~K} 33$ | $\mathrm{~B} 30 \oplus \mathrm{~K} 32$ | $\mathrm{~B} 37 \oplus \mathrm{~K} 31$ | $\longrightarrow$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

2) Process of modifying the Vernam algorithm for decryption

Ciphertext:

| C0 | C1 | C2 | C3 | C4 |
| :--- | :--- | :--- | :--- | :--- |

Bit ciphertext:

| C 07 | C 06 | C 05 | C 04 | C 03 | C 02 | C 01 | C 00 | $\longrightarrow$ | C 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| C 17 | C 16 | C 15 | C 14 | C 13 | C 12 | C 11 | C 10 | $\longrightarrow$ | C 1 |
| C 27 | C 26 | C 25 | C 24 | C 23 | C 22 | C 21 | C 20 | $\longrightarrow$ | C 2 |
| C 37 | C 36 | C 35 | C 34 | C 33 | C 32 | C 31 | C 30 | $\longrightarrow$ | C 3 |
| C 47 | C 46 | C 45 | C 44 | C 43 | C 42 | C 41 | C 40 | $\longrightarrow$ | C 4 |

Key:

| K0 | K1 | K2 | K3 | K4 |
| :--- | :--- | :--- | :--- | :--- |

Bit Key:

| K07 | K06 | K05 | K04 | K03 | K02 | K01 | K00 | $\longrightarrow$ | K0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| K17 | K16 | K15 | K14 | K13 | K12 | K11 | K10 | $\longrightarrow$ | K1 |
| K27 | K26 | K25 | K24 | K23 | K22 | K21 | K20 | $\longrightarrow$ | $\longrightarrow$ |
| K37 | K36 | K35 | K34 | K33 | K32 | K31 | K30 | $\longrightarrow$ | $\longrightarrow$ |
| K47 | K46 | K45 | K44 | K43 | K42 | K41 | K40 | $\longrightarrow$ | $\longrightarrow$ |

1-bit rotation to Right:

| K00 | K07 | K06 | K05 | K04 | K03 | K02 | K01 | $\longrightarrow$ | K0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| K10 | K17 | K16 | K15 | K14 | K13 | K12 | K11 | $\longrightarrow$ | K1 |
| K20 | K27 | K26 | K25 | K24 | K23 | K22 | K21 | $\longrightarrow$ | $\longrightarrow$ |
| K30 | K37 | K36 | K35 | K34 | K33 | K32 | K31 | $\longrightarrow$ | K3 |
| K40 | K47 | K46 | K45 | K44 | K43 | K42 | K41 | $\longrightarrow$ | $\longrightarrow$ |

Plaintext / XOR Operation:

| $\mathrm{C} 07 \oplus \mathrm{~K} 00$ | $\mathrm{C} 06 \oplus \mathrm{~K} 07$ | $\mathrm{C} 05 \oplus \mathrm{~K} 06$ | $\mathrm{C} 04 \oplus \mathrm{~K} 05$ | $\mathrm{C} 03 \oplus \mathrm{~K} 04$ | $\mathrm{C} 02 \oplus \mathrm{~K} 03$ | $\mathrm{C} 01 \oplus \mathrm{~K} 02$ | $\mathrm{C} 00 \oplus \mathrm{~K} 01$ | $\longrightarrow \mathrm{P} 0$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{C} 17 \oplus \mathrm{~K} 10$ | $\mathrm{C} 16 \oplus \mathrm{~K} 17$ | $\mathrm{C} 15 \oplus \mathrm{~K} 16$ | $\mathrm{C} 14 \oplus \mathrm{~K} 15$ | $\mathrm{C} 13 \oplus \mathrm{~K} 14$ | $\mathrm{C} 12 \oplus \mathrm{~K} 13$ | $\mathrm{C} 11 \oplus \mathrm{~K} 12$ | $\mathrm{C} 10 \oplus \mathrm{~K} 11$ | $\longrightarrow \mathrm{P} 1$ |
| $\mathrm{C} 27 \oplus \mathrm{~K} 20$ | $\mathrm{C} 26 \oplus \mathrm{~K} 27$ | $\mathrm{C} 25 \oplus \mathrm{~K} 26$ | $\mathrm{C} 24 \oplus \mathrm{~K} 25$ | $\mathrm{C} 23 \oplus \mathrm{~K} 24$ | $\mathrm{C} 22 \oplus \mathrm{~K} 23$ | $\mathrm{C} 21 \oplus \mathrm{~K} 22$ | $\mathrm{C} 20 \oplus \mathrm{~K} 21$ | $\longrightarrow \mathrm{P} 2$ |
| $\mathrm{C} 37 \oplus \mathrm{~K} 30$ | $\mathrm{C} 36 \oplus \mathrm{~K} 37$ | $\mathrm{C} 35 \oplus \mathrm{~K} 36$ | $\mathrm{C} 34 \oplus \mathrm{~K} 35$ | $\mathrm{C} 33 \oplus \mathrm{~K} 34$ | $\mathrm{C} 32 \oplus \mathrm{~K} 33$ | $\mathrm{C} 31 \oplus \mathrm{~K} 32$ | $\mathrm{C} 30 \oplus \mathrm{~K} 31$ | $\longrightarrow \mathrm{P} 3$ |
| $\mathrm{C} 47 \oplus \mathrm{~K} 40$ | $\mathrm{C} 46 \oplus \mathrm{~K} 47$ | $\mathrm{C} 45 \oplus \mathrm{~K} 46$ | $\mathrm{C} 44 \oplus \mathrm{~K} 45$ | $\mathrm{C} 43 \oplus \mathrm{~K} 44$ | $\mathrm{C} 42 \oplus \mathrm{~K} 43$ | $\mathrm{C} 41 \oplus \mathrm{~K} 42$ | $\mathrm{C} 40 \oplus \mathrm{~K} 41$ | $\longrightarrow \mathrm{P} 4$ |

XOR result bits:

| B07 | B 06 | B 05 | B 04 | B 03 | B 02 | B 01 | B 00 | $\longrightarrow$ | P 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| B 17 | B 16 | B 15 | B 14 | B 13 | B 12 | B 11 | B 10 | $\longrightarrow$ | P 1 |
| B 27 | B 26 | B 25 | B 24 | B 23 | B 22 | B 21 | B 20 | $\longrightarrow$ | P 2 |
| B 37 | B 36 | B 35 | B 34 | B 33 | B 32 | B 31 | B 30 | $\longrightarrow$ | P 3 |
| B 47 | B 46 | B 45 | B 44 | B 43 | B 42 | B 41 | B 40 | $\longrightarrow$ | P 4 |

Plaintext 1-bit Rotation to Right:

| B 00 | B 07 | B 06 | B 05 | B 04 | B 03 | B 02 | B 01 |  |  |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| B 10 | B 17 | B 16 | B 15 | B 14 | B 13 | B 12 | B 11 |  |  |
| B 0 |  |  |  |  |  |  |  |  |  |
| B 20 | B 27 | B 26 | B 25 | B 24 | B 23 | B 22 | B 21 |  | P 1 |
| B 30 | B 37 | B 36 | B 35 | B 34 | B 33 | B 32 | B 31 |  |  |
| B 40 | B 47 | B 46 | B 45 | B 44 | B 43 | B 42 | B 41 |  |  |

### 3.3 Multi-bit LSB

Multi-bit LSB is performed using several bit insertion models. The steps performed in the multibit LSB process for inserting messages, in the form of three elements of color (RGB), in specified bits within each pixel are as follows.

Inserting the character as 1-bit LSB storage:

| $R$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $G$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| B | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| $R$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| G | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| B | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| $R$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| G | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |

Inserting the character as 2-bits LSB storage:

| $R$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $G$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| $B$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| $R$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |

Inserting the character as 3-bits LSB storage:

| $R$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| G | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| $B$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |

Inserting the character as 4-bits LSB storage:

| $R$ | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $G$ | 7 | 6 | 5 | 4 | 4 | 2 | 1 | 0 |

The four patterns presented above store up to one character (i.e., 8 bits). The more are the number of bits used, the fewer are the pixels used for information storage: 1-bit insertion uses less than 3 pixels; 2-bit insertion uses less than 2 pixels; 3-bit insertion uses exactly 1 pixel; and 4-bit insertion uses less than 1 pixel. In addition, unused color elements can be used to save the next character.

### 3.4 Calculation of MSE and PSNR

Herein, a 24-bit color image was used for the experiment. The MSE and PSNR values were calculated to determine the extent to which an image changes post message insertion. Storage from 1-bit up to 4-bit LSBs was performed, and the MSE and PSNR values were calculated for each stego-image to determine which image is superior and how many bits should be used to store information or messages. The following formula was used to calculate the MSE and PSNR values.
Using (2.2) and (2.3)

Example:
Table 1. RGB Pixels

| Red |  |  |
| :--- | :--- | :--- |
| R1 | $\ldots$ | Rn |


| Green |  |  |
| :---: | :---: | :---: |
| G1 | $\ldots$ | Gn |


| Blue |  |  |
| :--- | :--- | :--- |
| B1 | $\cdots$ | Bn |

$$
\begin{aligned}
& M S E=\frac{1}{M * N}\left(\frac{\left(R_{1}^{\prime}-R_{1}\right)^{2}+\cdots+\left(R_{n}^{\prime}-R_{n}\right)^{2}+\left(G_{1}^{\prime}-G_{1}\right)^{2}+\cdots+\left(G_{n}^{\prime}-G_{n}\right)^{2}+\left(B_{1}^{\prime}-B_{1}\right)^{2}+\cdots+\left(B_{n}^{\prime}-B_{n}\right)^{2}}{3}\right) \\
& \quad P S N R=10 \log _{10}\left(\frac{C_{\max }^{2}}{M S E}\right)
\end{aligned}
$$

## 4. Result and Discussion

An example of an image used herein is one containing $4 \times 4$ pixels, which can accommodate up to 48 bits. If converted to characters, the image can accommodate up to 6 characters for the use of 1 bit. The following are the steps for 1 -bit plaintext rotation and the 1 -bit key.

Plaintext : Secret

| 83 | 101 | 99 | 114 | 101 | 116 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 01010011 | 01100101 | 01100011 | 01110010 | 01100101 | 01110100 |

Rotation : 1-bit to the left

| 10100110 | 11001010 | 11000110 | 11100100 | 11001010 | 11101000 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 166 | 202 | 198 | 228 | 202 | 232 |

Key : World

| 87 | 111 | 114 | 108 | 100 |
| :--- | :--- | :--- | :--- | :--- |
| 01010111 | 01101111 | 01110010 | 01101100 | 01100100 |

Rotation : 1-bit to the right

| 10101011 | 10110111 | 00111001 | 00110110 | 00110010 |
| :--- | :--- | :--- | :--- | :--- |
| 171 | 183 | 57 | 54 | 50 |

Ciphertext : $-\jmath y \hat{\mathrm{E}}^{\circ} \mathrm{C}$

| 13 | 125 | 255 | 210 | 248 | 67 |
| :--- | :--- | :--- | :--- | :--- | :--- |

The results of the ciphertext are stored in an image using several bit storage models, such as that illustrated below. In the test conducted herein, a $4 \times 4$-pixel image was used as a cover image; this could accommodate $4 \times 4 \times 3$ bits $=48$ bits (i.e., six characters) using 1-bit LSB.

RGB bit for the $4 \times 4$ pixel cover image:

| Red | Green | Blue |
| :---: | :---: | :---: |


| 65 | 58 | 187 | 190 | 87 | 176 | 149 | 241 | 116 | 249 | 167 | 229 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 222 | 171 | 203 | 63 | 223 | 74 | 221 | 68 | 150 | 134 | 191 | 182 |
| 246 | 101 | 32 | 89 | 205 | 204 | 143 | 53 | 128 | 84 | 243 | 163 |
| 192 | 151 | 36 | 11 | 125 | 130 | 167 | 60 | 226 | 101 | 153 | 23 |

RGB binary for a $4 \times 4$ pixel cover image:

| Red | Green | Blue |
| :---: | :---: | :---: |
| 01000001001110101011101110111110 | 01010111101100001001010101111001 | 01110100111110011010011111100101 |
| 11011110101010111100101100111111 | 11011111010010101101110101000100 | 10010110100001101011111110110110 |
| 11110110011001010010000001011001 | 11001101110011001000111100110101 | 10000000010101001111001110100011 |
| 11000000100101110010010000001011 | 01111101100000101010011100111100 | 11100010011001011001100100010111 |

Ciphertext bit to be inserted: $\quad-\} \ddot{\mathrm{y}} \hat{\mathrm{E}}^{\circ} \mathrm{C}$

| 00001101 | 01111101 | 11111111 | 11010010 | 11111000 | 01000011 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Stego-image results obtained using 1-bit LSB:

| Red |  |  |  | Green |  |  |  | Blue |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\underline{\mathbf{6 4}}$ | 58 | $\underline{\mathbf{1 8 6}}$ | $\underline{\mathbf{1 9 1}}$ | $\underline{\mathbf{8 6}}$ | $\underline{\mathbf{1 7 7}}$ | 149 | 241 | 116 | 249 | $\underline{\mathbf{1 6 6}}$ | 229 |
| $\underline{\mathbf{2 2 3}}$ | 171 | 203 | 63 | 223 | $\underline{\mathbf{7 5}}$ | 221 | $\underline{\mathbf{6 9}}$ | 150 | $\underline{\mathbf{1 3 5}}$ | 191 | $\underline{\mathbf{8 3 3}}$ |
| $\underline{\mathbf{2 4 7}}$ | 101 | $\underline{\mathbf{3 3}}$ | 89 | 205 | 204 | $\underline{\mathbf{1 4 2}}$ | 53 | 128 | 84 | 243 | 163 |
| $\underline{\mathbf{1 9 3}}$ | $\underline{\mathbf{1 5 0}}$ | 36 | $\underline{\mathbf{1 0}}$ | $\underline{\mathbf{1 2 4}}$ | 130 | $\underline{\mathbf{1 6 6}}$ | $\underline{\mathbf{6 1}}$ | 226 | 101 | $\underline{\mathbf{1 5 2}}$ | 23 |

$$
\begin{aligned}
\text { MSE } & =\frac{1}{16}(7) \\
& =0,4375 \\
\text { PSNR } & =10 * \log 10\left(\frac{249^{2}}{0,4375}\right) \\
& =51,5142
\end{aligned}
$$

Stego-image results obtained using 2-bit LSB:

| Red |  |  |  | Green |  |  |  | Blue |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\underline{\mathbf{6 4}}$ | $\underline{\mathbf{5 7}}$ | 187 | $\underline{\mathbf{1 9 1}}$ | $\underline{\mathbf{8 4}}$ | $\underline{\mathbf{1 7 7}}$ | 149 | $\underline{\mathbf{2 4 3}}$ | $\underline{\mathbf{1 1 9}}$ | $\underline{\mathbf{2 5 1}}$ | 167 | $\underline{\mathbf{2 3 1}}$ |
| $\underline{\mathbf{2 2 3}}$ | $\underline{\mathbf{1 7 0}}$ | $\underline{\mathbf{2 0 2}}$ | $\underline{\mathbf{6 0}}$ | $\underline{\mathbf{2 2 1}}$ | $\underline{\mathbf{7 5}}$ | $\underline{\mathbf{2 2 0}}$ | 68 | $\underline{\mathbf{1 4 8}}$ | $\underline{\mathbf{1 3 5}}$ | $\underline{\mathbf{1 8 9}}$ | $\underline{\mathbf{1 8 3}}$ |
| 246 | 101 | 32 | 89 | 205 | 204 | 143 | 53 | 128 | 84 | 243 | 163 |
| 192 | 151 | 36 | 11 | 125 | 130 | 167 | 60 | 226 | 101 | 153 | 23 |

$$
\begin{aligned}
\text { MSE } & =\frac{1}{16}(10,67) \\
& =0,6667
\end{aligned}
$$

$$
\begin{aligned}
\operatorname{PSNR} & =10 * \log 10\left(\frac{249^{2}}{0,6667}\right) \\
& =49,6846
\end{aligned}
$$

Stego-image results obtained using 3-bit LSB:

| Red |  |  |  | Green |  |  |  | Blue |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\underline{\mathbf{6 4}}$ | $\underline{\mathbf{6 3}}$ | $\underline{\mathbf{1 9 1}}$ | $\underline{\mathbf{1 8 8}}$ | $\underline{\mathbf{8 3}}$ | $\underline{\mathbf{1 8 2}}$ | $\underline{\mathbf{1 5 1}}$ | $\underline{\mathbf{2 4 6}}$ | $\underline{\mathbf{1 1 3}}$ | $\underline{\mathbf{2 5 5}}$ | $\underline{\mathbf{1 6 6}}$ | $\underline{\mathbf{2 3 1}}$ |
| $\underline{\mathbf{2 2 0}}$ | 171 | 203 | 63 | $\underline{\mathbf{2 1 8}}$ | 74 | 221 | 68 | $\underline{\mathbf{1 4 4}}$ | 134 | 191 | 182 |
| $\mathbf{2 4 6}$ | 101 | 32 | 89 | 205 | 204 | 143 | 53 | 128 | 84 | 243 | 163 |
| $\mathbf{1 9 2}$ | 151 | 36 | 11 | 125 | 130 | 167 | 60 | 226 | 101 | 153 | 23 |

$$
\begin{aligned}
\text { MSE } & =\frac{1}{16}(18) \\
& =1,125 \\
\text { PSNR } & =10 * \log 10\left(\frac{249^{2}}{1,125}\right) \\
& =47,4142
\end{aligned}
$$

Stego-image results obtained using 4-bit LSB:

| Red |  |  |  | Green |  |  |  | Blue |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\underline{\mathbf{6 4}}$ | $\underline{\mathbf{6 1}}$ | $\underline{\mathbf{1 8 9}}$ | $\underline{\mathbf{1 8 4}}$ | $\underline{\mathbf{9 3}}$ | $\underline{\mathbf{1 9 1}}$ | $\underline{\mathbf{1 4 6}}$ | $\underline{\mathbf{2 4 4}}$ | $\underline{\mathbf{1 1 9}}$ | $\underline{\mathbf{2 5 5}}$ | $\underline{\mathbf{1 7 5}}$ | $\underline{\mathbf{2 2 7}}$ |
| 222 | 171 | 203 | 63 | 223 | 74 | 221 | 68 | 150 | 134 | 191 | 182 |
| 246 | 101 | 32 | 89 | 205 | 204 | 143 | 53 | 128 | 84 | 243 | 163 |
| 192 | 151 | 36 | 11 | 125 | 130 | 167 | 60 | 226 | 101 | 153 | 23 |

MSE $=\frac{1}{16}(19,333)$
$=1,2083$
PSNR $=10 * \log 10\left(\frac{249^{2}}{1,2083}\right)$
$=47,1022$
Bold and underlined pixel bits are pixel bits that have been inserted into the message. In 1-bit LSB storage more image pixels are used because each RGB color element can only store 1-bit.

For storing 4-bit LSB fewer pixels are used because each RGB color element can store 4-bit ciphertex.

## 5. Conclusion

The results of this study indicate that the level of data security can be increased by combining steganography methods with cryptography algorithms. The resolution of the cover image and character length of the encrypted message considerably affect the MSE and PSNR values. The calculation results of MSE and PSNR values indicate that the use of 1-bit LSB is superior to that of 2-, 3-, or 4-bit LSB. However, according to the results, the use of 4-bit LSB is feasible because the PSNR value for 4-bit LSB is above 40 db . The smaller percentage of pixel usage in 4-bit LSB storage provides satisfactory MSE and PSNR values. The application of the multi-bit LSB method in steganographic activities is thus advantageous, in the sense that each pixel can hold more message bits compared with the conventional LSB method.

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