

# READOUT DRIVER FOR THE ATLAS LIQUID ARGON CALORIMETERS

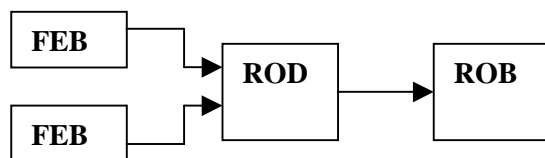
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## Abstract

The Readout Drive (ROD) for the Liquid Argon calorimeter front-end electronics of the ATLAS detector is described. The ROD receives triggered data from 256 calorimeter cells. It must derive the precise energy and timing of calorimeter signals from discrete samplings of the pulse. In addition, it performs monitoring and formats the digital stream for the succeeding element in the readout chain. Data arrive over two 1.28 Gbit/s fiber optics links at a 100 kHz event rate (25Kbit/event). Principals of the design are discussed, along with simulations of data processing.

## 1. INTRODUCTION

In ATLAS [1] there are three basic types of liquid argon calorimeters: the em calorimeters (barrel and end-cap), the hadronic end-cap calorimeter and the forward calorimeters. The front-end electronics for all of these calorimeters is essentially identical, the differences being confined to the amplification stage upstream of the shaping amplifier. After shaping, the signals are stored in a switched capacitor array (SCA), and upon receipt of a Level 1 trigger, the samples relevant to the event are digitized on the front-end board (FEB). These digitizations are transmitted to the Readout Driver (ROD) module, whose function is to extract the parameters of interest for each calorimeter cell and pass these data to the Readout Buffer (ROB) module, the first element in the data acquisition chain. A simplified diagram of this part of the readout chain is shown in Figure 1.



**Figure 1: Simplified diagram of the portion of the readout chain involving the ROD.**

In the ATLAS FEB, which treats 128 calorimeter channels, the signals are amplified, shaped and then stored as analog levels in the SCA. Upon receipt of a Level 1 trigger signal, a 12-bit ADC digitises the appropriate samples. Three gain scales are employed, requiring three shaper channels and three SCA channels for each calorimeter cell. All samples are

digitized on a common gain scale, which is chosen event-by-event by examining the amplitude of the sample closest to the peak of the signal. Each ADC digitises the signals from 8 calorimeter channels, and the results are sent over an optical fiber to the ROD module. The fiber contains data from all 16 ADCs in the FEB; 32 bits (2 bits/ADC) are sent every 25 ns, giving a transmission rate of 1.28 Gbit/s.

## 2. THE ROD MODULE

### 2.1 Overview

A single ROD module receives data from two front-end boards, consisting of (typically) 5 samples from 256 channels. The processors in the module calculate energy and timing information from these data, and in most cases, discard the raw data. The ROD also performs monitoring tasks, and during calibration runs, it executes a signal averaging task and sends averaged data to a local processor, which then calculates calibration constants for the channels belonging to that module. The ROD modules will be housed in a Readout Crate, which will in all likelihood will be a 9U VME crate with a dedicated host processor. The ROD system will be a highly specialized distributed computing resource for the ATLAS detector. It will consist of about 800 modules, each of which services up to 256 calorimeter channels. The total computing power of this resource will be approximately  $4 \times 10^{15}$  arithmetic operations per second.

### 2.2 The Basic Algorithm

The most important function of the ROD is to determine the energy  $E$  and time  $T$  (relative to the nominal bunch crossing timing) from the digitized samples, along with a parameter  $Q$  (such as chi-square), which indicates how closely the samples follow the known waveform. Secondary functions include updating histograms of these quantities and performing certain monitoring functions for some small fraction of the events.

A typical waveform of the shaped liquid argon waveform is shown in Figure 2, along with samples spaced by 25 ns. A general technique to estimate  $E$  and  $T$  in an accurate and computationally efficient manner is that of optimal filtering [2], in which the desired quantity is expressed as a sum of the samples multiplied by predetermined weights. The weights are found by requiring that the standard deviation of the

quantity be minimized while satisfying certain constraints. In our case, where there are two quantities to be determined, the procedure involves simultaneously minimising the uncertainty in both quantities. The expressions are:

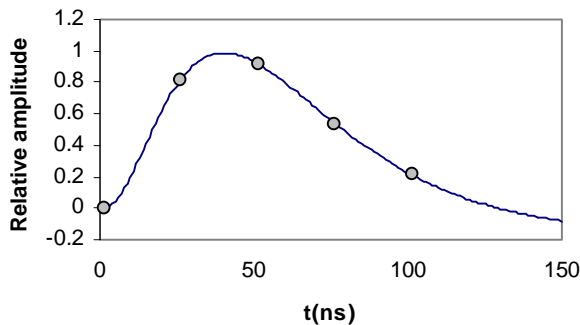
$$E = \sum a_i S_i$$

$$E \cdot T = \sum b_i S_i,$$

where the sum extends over all of the samples  $S_i$ , and where  $a_i$  and  $b_i$  are the weights. From the structure of these formulae, one sees that the error in the amplitude is amplitude independent, whereas the error in the time varies inversely with the amplitude. For this reason it only makes sense to calculate  $T$  only for those channels with  $E$  above some threshold value  $E_{th}$ . The quality-of-fit parameter will most likely be a simplified expression for chi-square (*i.e.*, one that ignores correlations between the different terms):

$$Q = \sum (S_i - E(g_i + \dot{g}_i T))^2,$$

in which  $g_i$  is the expected waveform normalized to unity. Since this calculation involves knowing both  $E$  and  $T$ , it will also be performed only for the case where  $E$  is above the threshold value. Once  $E$ ,  $T$ , and  $Q$  are found, the corresponding bins for general histograms of these quantities are calculated and incremented. In addition, if special histograms are required for specific calorimeter cells, which are being monitored, these histograms are also incremented.



**Figure 2: Typical shaped calorimeter signal with samples (dots) spaced by 25 ns.**

Thus the basic algorithm is as follows:

- calculate  $E$  for all channels
- if  $E > E_{th}$ , calculate  $T$  and  $Q$

- update general histograms
- calculate quantities for monitoring specific channels

In simulation, these steps have been implemented for the C6202 DSP, and execution times for each stage of the algorithm are given in Section 5 below.

### 2.3 Design Considerations

The maximum Level 1 trigger rate for ATLAS is currently specified as 75 kHz, but an upgrade to 100 KHz is considered a strong possibility, and hence we use the latter figure as a design parameter for the ROD. Since the processing time per event depends on the fraction of cells with energy above the threshold, there can be considerable fluctuations from module to module. For this reason, a derandomizing buffer will be required to reduce the system dead time. One model of the process indicates that in order to keep the dead time below 0.5%, it is necessary to buffer 7 events [3]. We plan for each ROD to serve two FEBS, or 256 calorimeter channels. With a Level 1 trigger rate of 100 KHz, the average processing time per event cannot exceed 10 microseconds. Fortunately only a small fraction of the cells contain significant energy deposits in each event, which reduces considerably the computing power required. As mentioned above, the energy is calculated for each channel, but the time value is computed only for events with energy significantly above the noise, since it is only for these channels that the measurement is meaningful.

Given the design criteria listed above, there are several possible approaches to the ROD design. One could imagine using a multiply-accumulate chip, which is optimized to perform the calculation we require. Or one can even imagine designing and building a special-purpose ASIC to carry out the task. Our preference, however, is to evaluate solutions using programmable, commercially available processors which can perform our algorithm efficiently but have limited general computing capability. A natural choice is the Digital Signal Processor (DSP), a device whose technology is advancing at a rapid pace. We plan to study such devices in detail to ascertain if they are able to perform the task in the required time before examining more ambitious solutions. Likewise, we plan to use off-the-shelf components for all ancillary elements in the ROD, in order to minimise design effort.

### 2.4 Integer vs. Floating Point DSP

Since we are dealing with quantities which cover many orders of magnitude in size ( $E$ , for example can range from tens of MeV to several TeV), it is both natural and convenient to use floating point DSPs. However, a

detailed investigation of the system of digitization we plan to use (12 bit ADC operating on 3 gain scales) indicate that integer DSPs, which are in general both faster and cheaper, are also adequate in this case. As long as 16-bit constants are used in the calculation, the effects of rounding in the determination of both the  $E$  and  $T$  are completely dominated by ADC quantization. The division of  $E \times T$  by  $E$  to obtain  $T$  is an inconvenience of course, but this can be handled to adequate accuracy by table lookup. Thus we plan to investigate both integer and floating point DSPs for possible use in the ROD.

### 3. ROD DEMONSTRATOR PROJECT

#### 3.1 Purpose and Scope

In order to demonstrate the capability of candidate DSPs and to understand more clearly the design problems of the ROD, ATLAS has decided to pursue a ROD Demonstrator Project. The project involves the construction of a motherboard in the 9U VME64x format, into which can be plugged up to four daughterboard processing units (PUs). These processing units will contain one or more DSPs and will be used to process calorimeter data, either fed in from an artificial source or from a calorimeter module running in the test beam. The plan calls for implementing the FEB-ROD-ROB chain, so that all of

the required functionality can be tested

#### 3.2 Hardware

The ROD Demonstrator Board is a VME64x 9U board with a custom P3 backplane. It accepts up to four Programming Units, which may be of different types. Input data may be supplied to the board from up to two FEBs or through the VME backplane. An input is also provided for the timing and trigger information (TTC) in the format that will be used in ATLAS. Output can be to a ROB module or through the VME backplane. In Figure 3 a schematic diagram of the board is shown.

The Processing Unit (PU) is a small (85x185 mm) daughterboard containing one or more DSPs plus any external memory required, input and output buffers, and an interface to the motherboard. The tasks performed by the PUs are expected to be very similar to the tasks required of the ROD in ATLAS. Currently two PUs are being designed, one based on a floating point DSP (the SHARC of Analog Devices) and the other based on an integer DSP (the C6202 of Texas Instruments). As newer DSPs, which look promising, become available, we expect to add them to the project.

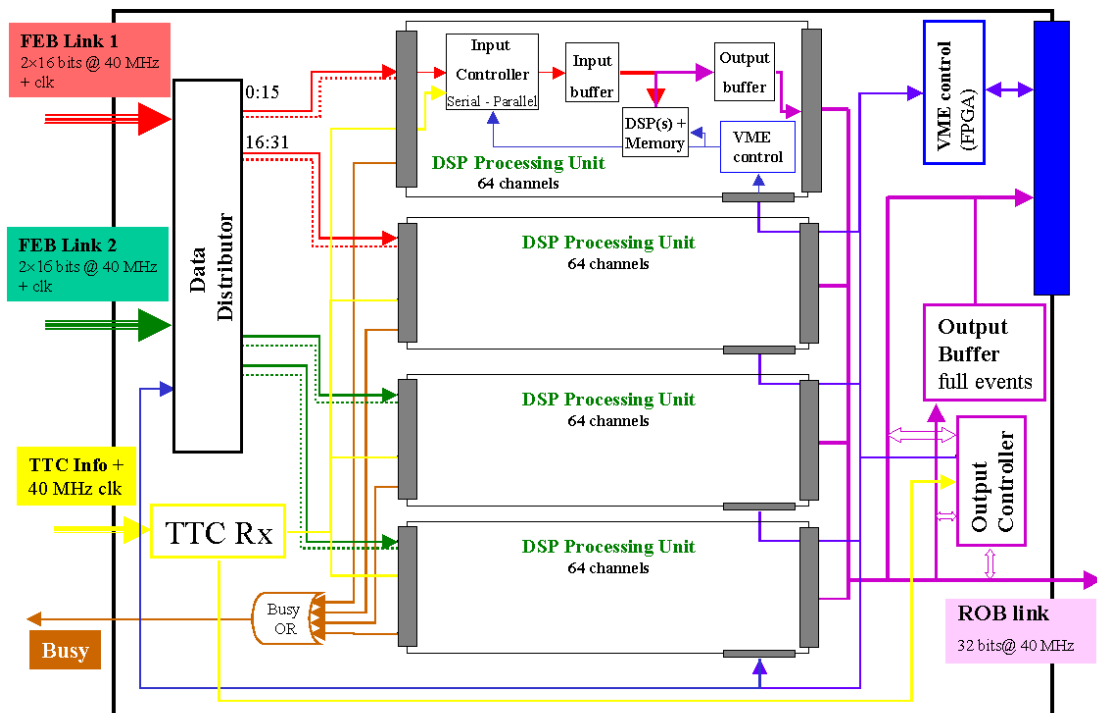


Figure 3: Schematic diagram of the ROD Demonstrator board.

## 4. A DESIGN EXAMPLE

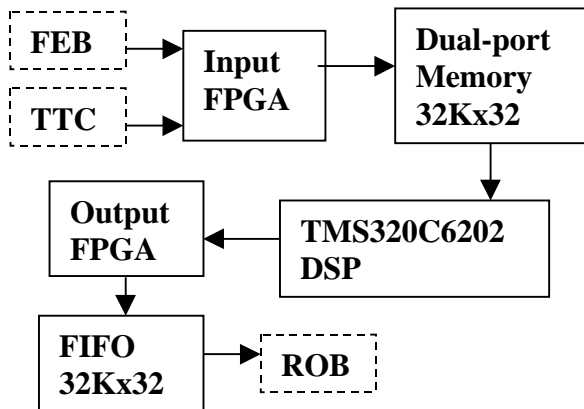
### 4.1 Motivation

In order to illustrate the type of studies that are needed to qualify a DSP for adoption for the ROD, we discuss a concrete example in some detail. For this we choose the design of the PU based on the Texas Instruments C6202 DSP. It is not unlikely that another DSP will eventually be used, given the time scale for the ATLAS experiment, which begins in 2005.

### 4.2 The TI C6202 DSP

The C6202 is an integer processor with many of the features that are required by the ATLAS ROD. It can operate at clock speeds up to 250 MHz and has eight independent functional units (6 ALUs and 2 16-bit multipliers), permitting it to execute eight 32-bit instructions per cycle. The internal 128 Kilobyte data memory is somewhat limited for our purposes, since we need to (a) buffer the input data, (b) carry out calculations, and (c) store histograms, so we plan to augment it with an external dual-port memory. The unit has both an external memory interface and an expansion bus, which offers a convenient interface to an FPGA.

### 4.3 Design Sketch of the Processing Unit



**Figure 4: Simplified schematic diagram of the Processing Unit based on the Texas Instruments C6202 DSP. The items shown in dashed boxes are external to the Processing Unit. The interface to the VME system is not shown.**

In Fig. 3 we show a block diagram of the processing unit. The data from the FEB and the TTC are brought in on the left, where they enter an FPGA. This device performs certain routine checks on the input data and transfers them into a fast dual-port memory, which is connected to the external memory bus of the DSP. This configuration permits the data to be processed in the

DSP without performing a transfer to its internal memory.

### 4.4 Ancillary Circuitry

There are four logical units in the Processing Unit in addition to the DSP: (1) Input FPGA, (2) Dual-port Memory, (3) Output FPGA, and (4) FIFO. The Input FPGA receives input data from both the FEB and the TTC (Trigger, Timing and Control) modules. The former contains the digitized data from the calorimeter whereas the latter contains information about the trigger (bunch crossing, trigger type, etc.). The two types of data are combined into one record by the FPGA, which also performs parity checks, and stored in the dual-port memory (the input buffer) until it is processed by the DSP. Once the DSP finishes the processing of the event, the results are written to the output FPGA, which formats the output data stream and puts it into the FIFO, which is the output buffer memory for the ROD. Because this unit will be used in a test situations, where there may be large fluctuations in event size, the input and output buffers are larger (corresponding to about 100 events of average size) than will probably be required for ATLAS.

## 5. QUANTITATIVE STUDIES

### 5.1 Benchmark Code

We have developed a set of tasks which each of the DSPs should perform in order to be able to make comparisons between them and also establish their viability as candidates for the final ROD system. This code does not include all of the tasks that will be performed by the DSPs in ATLAS, only the most time-critical ones. First, the event header is checked for errors in data format or parity. Then the basic operations performed are:

- check event header
- read five input samples and gain scale
- fetch weights
- calculate  $E$
- if  $E > E_{th}$ , calculate  $T$  and  $Q$

The above steps are performed for each channel assigned to the DSP. In our current model of the ROD, four DSPs are used in each module, so 64 calorimeter channels are assigned to each DSP.

### 5.2 Timing Studies for TI C6202

The benchmark code described above has been implemented for the TI C6202 processor, using the simulator provided by Texas Instruments. To improve the calculation time, it was found necessary to break the code into small loops, minimising branching and conditional statements. Hand coding was used to optimise the parallelism available in the processor. Let

$N_E$  represent the number of cells which have  $E > E_{th}$  (for this study  $E_{th}$  was chosen to be twice the noise of the calorimeter cell, or about 100 MeV). In Table 1 are given the number of DSP cycles taken for each section of the algorithm, as a function of  $N_E$ , and in Table 2 we give an estimate of how this translates into execution times for events with different numbers of cells above threshold. Here we see that even if one-third of the cells are above threshold (a rare occurrence), the execution time is equal to the average spacing between events at the design trigger rate. This indicates that the TI C6202 would in fact meet the requirements for the ROD. More recent results, arising from further optimisation of the algorithm, show that the execution times given here are overestimates, which strengthen this conclusion.

**Table 1: Number of cycles required by TI C6202 for each stage of the algorithm.**

Calculation	Cycles
Preliminary (determine $N_E$ )	311
Energy	$26+22*N_E$
Time	$27+22*N_E$
Chi-square	$89+15*N_E$
General histograms	$39+16*N_E$
Monitoring selected cells	$48+13*N_E$

**Table 2: Execution times of the TI C6202 DSP (in microseconds) as a function of number of cells above threshold, out of a total of 64 cells, and whether or not monitoring functions are being performed.**

$N_E$	Monitoring?	Execution time
0	no	1.6
1	no	2.8
10	no	6.2
20	no	10.0
64	no	26.4
1	yes	9.7
10	yes	13.7
64	yes	37.0

### 5.3 Monte Carlo Studies

The ATLAS Monte Carlo chain is used to produce artificial output data for the FEB, and this output has been fed to the benchmark code for the ROD described above. The first case to be studied is for the shower of

a 50 GeV electron, for the case of only thermal noise (zero luminosity) and for pileup noise at low and high luminosity. It was found that the energy and width of the reconstructed distribution are consistent with the expected values for each case, and that the introduction of integer arithmetic has no measurable effect on the results. We expect to use this program for a variety of purposes, such as checking the algorithm in the reconstruction of different types of showers, estimating with more precision the parameters, which enter into the execution time of the algorithm, and to understand how these parameters depend upon event type.

## 6. SUMMARY

We have described the technical requirements for the Readout Driver for the liquid argon calorimeters in ATLAS. From our studies to date, it appears that commercial DSPs can meet the needs for this device. We are carrying out a demonstration project in which several DSPs, both integer and floating point will be evaluated for their suitability in ATLAS. We give an example of the conceptual design and the simulation results for one of the processing units being built for this project, which is based on the Texas Instruments C6202 DSP. Simulation of benchmark code indicates that the speed of this processor is close to meeting our requirements, and initial Monte Carlo results indicate that the algorithm used produces acceptable results for the case studied.

## 7. REFERENCES

1. ATLAS Liquid Argon Calorimeter TDR, N/LHCC/96-41 ATLAS TDR 2, 15 Dec. 1996
2. W. E. Cleland and E.G. Stern, Nuclear Instruments and Methods A338 (1994) 467
3. ATLAS LAr TDR, *ibid*, p. 425