

Symmetric Cryptosystem Based on Petri Net

Hussein A. Lafta Rand Abdul-Wahid M. Ali

*College of Science for Women, University of Babylon College of Information
Technology, University of Babylon*

wsci.husein.attia@uobabylon.edu.iq

albeerrand@yahoo.com

Abstract

In this work, a novel approach based on ordinary Petri net is used to generate private key . The reachability marking of petri net is used as encryption/decryption key to provide more complex key . The same ordinary Petri Nets models are used for the sender(encryption) and the receiver(decryption).The plaintext has been permuted using look-up table ,and XOR-ed with key to generate cipher text

Keyword: pseudo random number , symmetric cryptosystem ,petri net

الخلاصة

يتضمن هذا البحث طريقة جديدة تعتمد على شبكة بترى لتوليد مفتاح سري خاص . يستخدم مؤشر الوصول الذي يشير الى البيانات الموجودة في الشبكة كمفتاح للتشفير وفك التشفير للحصول على مفتاح معقد بشكل جيد . يستخدم كلا الطرفين (المرسل والمستقبل) للتشفير وفك التشفير شبكة ذات تصميم مطابق لتوليد نفس المفتاح . يتم اعادة ترتيب النص الصريح باستخدام جدول معين قبل عملية التجميع مع المفتاح لتوليد النص المشفر .

الكلمات المفتاحية: بترى نت ، تشفير تناظري ، توليد رقم شبه عشوائي .

1.Introduction

One of the most used in information security is Cryptography ,one of its services is protected sensitive information to discover by unauthorized people or made modification on it during sending or storage. Cryptography depends upon two basic elements: cryptography algorithm and key generator algorithm . key stream generator has a short input(seed) and a long output, the two party(sender and receiver) have the same input and generate the same key, our project use petri net as generator. Petri net is mathematical graphical model ,consists three component ,places , transitions and arc that connect one place to one transition or one transition to one place, each arc labels with its weight(positive integer), each place may hold either zero or a positive number of tokens or mark. A marking of net is denoted by M, it consists tokens of all places, short key(seed) is used as initial marking to the first places of petri net ,firing of enabled transition change the token in petri net according to the weights of arc and number of firing which can be change according to the agreement between sender and receiver ,this generated random long number . The plaintext has been permuted using look-up table, each plaintext is permuted in differ form according to the key . And finally the plaintext and key have been mixed to generate ciphertext.

2. Related works

There are few paper that used petri net as a tool to generate public key sequence based on the average of token or complicated of more than one petri net .

1-“A Public-Key Cryptosystem Based On Stochastic Petri Net”(Zuohua Ding, Hui Zhou, Hui Shena, Qi-wei Geb)2014. They have developed a public-key cryptosystem based on scope of initial marking and average number of token in places of stochastic petri net, in decryption process used CPN. The markings in coverability tree are used in key generation, and the plaintext can be encrypted in many steps. This system has higher security than other like RSA.

2-“Construction of Petri nets and Calculation of Elementary T invariants for Multi-stage-Encryptions Public-Key Cryptography: MEPKC”(Ryo Yamaguchi , Qi-Wei Ge and Mitsuru Nakata,Graduate School of Education, Yamaguchi University)

2008. They have proposed a method, by collecting two Petri nets, in order to increase complex of petri nets to be used as a key generator of a public-key cryptography.

3. Symmetric Cryptography

Cryptography is the study of information security and provide possibility of connecting over an insecure channel to protect information during transmission. Information security is become most important in general modern business and technology for privacy of communications and transmission. Good cryptography gets its security by using complex keys . [Shafi ,Mihir 08]

Cryptography can be divided into asymmetric cipher and symmetric cipher ,A symmetric cipher which use public key to encrypt message and secret key to decrypt it , while symmetric cipher use the same secret key to encrypted and decrypted the message [Paar & Pelzl 10]

In symmetric cipher , the sender and receiver shared secret key by secure channel which will be used in the encryption/decryption algorithm .Symmetric cipher consists three algorithm the randomized key generation algorithm, encryption algorithm and decryption algorithm .The randomized key generation algorithm are used to generated random number, it considers as basic role in the use of cryptography in various security application, the security of algorithms which use it are based on the assumption that it is infeasible to distinguish when a random sequence is being used . A generator began with short random bit string (as a seed) and extend it to become long complex bit string, it should be as long as plaintext to maintain high security . Random number generator are deterministic functions, so it based on the seed, sender and receiver have been agreed on the same seed which must be randomly and shared it over secure channel to get same secret key, which will be xor-ed with the plaintext. [Andrea 2005]

4.Petri Nets

Petri nets are graphical and mathematical tool, it has been used in many applications including the modeling and analysis of discrete-event systems ,concurrent systems , fault tolerant systems , communication protocols, distributed-software systems , control systems and parallel systems . [Jiacun]

A Petri net is a weighted directed graphs, it consists two types of nodes, the place(represented by circle) and transition(represented by bar), place and transition connected by directed arcs either from place to transition or transition to place ,each arc are labeled with weight. [Bab 01]

A marking of net is labeled by M , (m vector where m is the total number of places), the p th component of M denoted by $M(p)$, is the number of tokens in place p . Transition t is enabled when each input place is marked with at least w tokens, where w is the weight of the input arc of that transition).

A firing on an enabled transition deletes token from each input place p , and adds token to each output place according to the weight of associative arc, see fig (1). Token-flow occurs via the firing of transitions. The system achieves a new marking via the firing of a transition. [Mutata 89]

A formal definition of Petri net, is represented by five-element $PN = (P, T, F, W, M_0)$ where:

- $P = \{p_1, p_2, p_3, \dots, p_n\}$ is a finite set of places,
- $T = \{t_1, t_2, t_3, \dots, t_n\}$ is a finite set of transitions,
- $F \subseteq (P * T) \cup (T * P)$ is a set of arcs (flow relation),
- $W: F \rightarrow \{1, 2, 3, \dots\}$ is a weight function,
- $M_0: \{0, 1, 2, 3, \dots\}$ is the initial marking

A Petri net structure $N = (P, T, F, W)$ without initial marking is denoted by N . A Petri net with the given initial marking is denoted by (N, M_0) [Mutata 89]

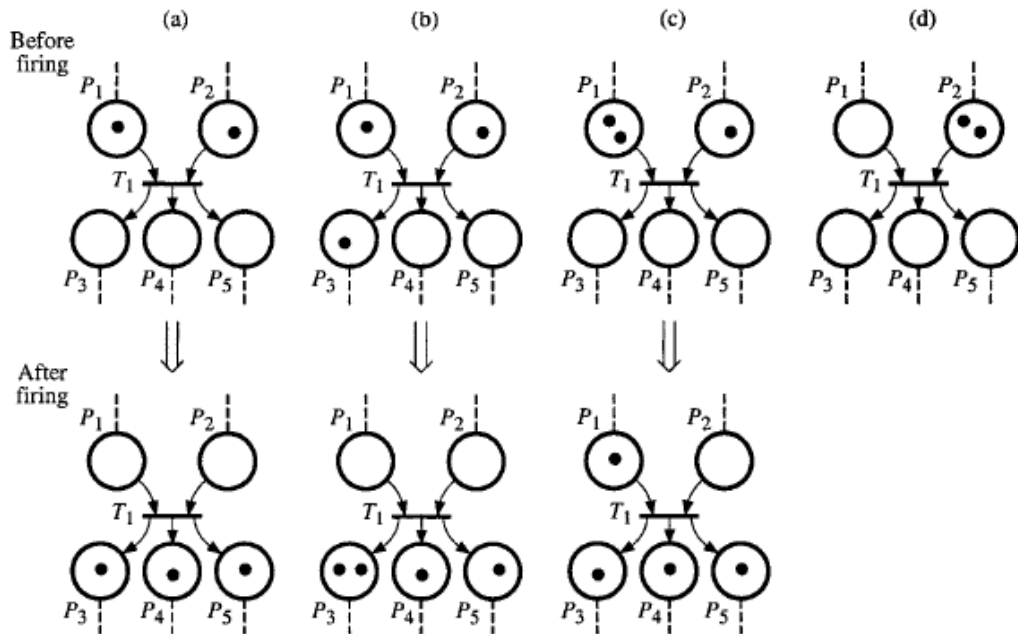


Figure (1) :Firing of transition

5. Design and Implementation of the Proposed System

The Petri net was used as good tools in cryptography. We have developed cryptosystem based on petri net . We can used it to generate nonlinear random number based on the marking of petri net and flow of token according to firing of transitions. this system can supply us with complex key to be used as private key for cryptosystem . the figure (2) depict the suggested petri net model that used in our research. The algorithm

(5.1) is used to generated pseudo random key sequences. The key sequences are used in encryption algorithm (5.2) to generate cipher text and decryption algorithm (5.3) to return the original message.

5.1 algorithm to generate random number .

Input : short secret key(seed) .

Output : long complex key .

Processing includes the following steps:

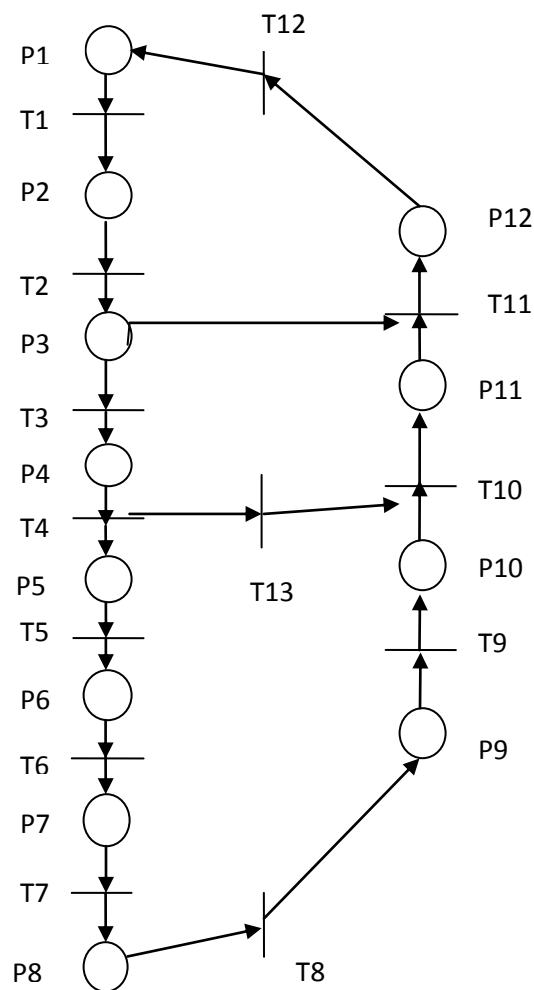
1- Obtain ASCII code of the seed then convert it to binary code and divided into sequence of five bit ,we get sequences $S=\{S_0,S_1,S_2,S_3,S_4,S_5,\dots\}$

2-Initialize marking of petri net with S, set P1 with S1,P2 with S2,.....,p7 with S7, while p8-p11 set to zero.

3-Set the weight of arc according to table(1),

4- Firing the enabled transitions of petri net n time.

5-Take the marking of petri net as private key



Figure(2): Suggested PN_model for generating a random number

Table(1):Suggested weight of Petri net

1 st node	2 nd node	Weight	1 st node	2 nd node	weight	1 st node	2 nd node	Weight
P1	T1	2	T5	P6	3	P10	T10	4
T1	P2	3	P6	T6	1	T10	P11	1
P2	T2	2	T6	P7	1	P11	T11	1
T2	P3	1	P7	T7	2	T11	P12	1
P3	T3	3	T7	P8	4	P12	T12	3
T3	P4	4	P8	T8	2	T12	P1	1
P4	T4	3	T8	P9	4	T3	P11	2
T4	P5	2	P9	T9	2	P5	T13	1
P5	T5	4	T9	P10	1	T13	P10	1

5.2 Algorithm for encryption

Input : plaintext message

Output: cipher text message

Processing including the following:

- 1-Obtain ASCII code of the plaintext then convert it to binary code
- 2-Divided the binary code into 16-bit sequence and permuted the binary code according to table(2),which indexed by S0 .
- 3- XORed the result with the private key .

5.3 Algorithm for decryption

Input : ciphertext message

Output: plain text message

Processing including the following:

- 1-Obtain ASCII code of the cipher then convert it to binary code
- 2- XORed the result with the private key .
- 3-Applied inverse permutation to binary code according to table(3) which index by S0 to return to original order

Table (2) permutation of plain

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	3	13	7	4	9	8	0	12	1	2	14	15	5	10	11	6
1	10	5	13	1	7	3	0	6	12	14	2	8	9	4	15	11
2	11	6	14	2	8	4	1	7	13	15	3	9	10	5	0	12
3	12	7	15	3	9	5	2	8	14	0	4	10	11	6	1	13
31	8	3	11	15	5	1	14	4	10	12	0	6	7	2	13	9

Table (3) revers permutation of cipher

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	6	8	9	0	3	12	15	2	5	4	13	14	7	1	10	11
1	6	3	10	5	13	1	7	4	11	12	0	15	8	2	9	14
2	14	6	3	10	5	13	1	7	4	11	12	0	15	8	2	9
3	9	14	6	3	10	5	13	1	7	4	11	12	0	15	8	2
31	10	5	13	1	7	4	11	12	0	15	8	2	9	14	6	3

Case study

The word (absyter) used as secret key ,and (my name is ahmed)used as plaintext.
 First ,applied random number algorithm according to secret key (absyter) to generate the complex key, use n=32 we get the following key
 (00001110110011000100111010111011100101111111101010000101111010100001011
 011001100111110110101001001001010101110011011011110111101001101101001011
 010100100111110101011000101011010111010110110010001011010001000001111100
 110011101101100010111101010110001111000101001101010001101111011001000100
 11001101011010010011101111111000101011101010111011100101100011101110011
 011100110011101011111000011110100100101011101101111010100010110111001010
 00001110010011111010011101110000111101110100100010101000111011111111110
 011101010000010110100100110010101000001011110111000100011011000100110010
 010011000011111110101111001010011010100100110001110100010111111100110100
 10111100100110010010101111110100000111010011111110100100110010011000000
 100001110110011110000010011001000111011101110001010010100101010001010010
 011001010110101010001110010011100100100100110010010011101101001001110011
 101001101001001100011101100011000101101001011101000100110010010101010110
 100000101010011010101001001100100110110011110101100111110110010011001000
 111010011111111010100110101110100100110010101000010111010010100101100001
 001100100100110001011011111001110011001010010011000111011101110011101001
 111101001001100100101010111010111101010011110110100100110010011100010101
 101101001111000100110010001110100111011001010011101110101001001100101011
 010101000010100100111100100110010010011110100101110100101111101001001100
 011101001100001101001111100000010011001001010110111011010011111000000100
 1100100101011011101)

Obtain binary code of plaintext and applied permutation we get
 (10111111000111101110000110011000100011010110110100000010011110110111101
 00100010001001100110001111011010100111000)
 Xor_ed the private key with the plaintext we get (
 101100011101001010101111001000110001101010010111100001111001000101101100
 1000100010110111100101011111111110000001)

And the cipher text is (XtUr j/ H[<W□). in other side ,the same petri net will be used to generate the same key , XOR-ed with cipher text and applied inverse permuted to get plaintext.

6. Conclusions and future works

In this work, a novel approach based on Petri net is used to encrypted messages. This a new approach is used to generated the private key. The proposed system consists two part, the first one is used to generate a random number based on the token of places, which generate complex long nonlinear random number used to encrypted message, the second part is used to permuted the plaintext and XORed the result with private key. There are some future works like private key cryptosystem based on coloured petri net, public key cryptosystem based on petri net, using coloured petri net for secure

transmission of message between sender and receiver and using petri net in parallel processing.

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