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Implementation of Latency by Using Distributed Load Balancing Algorithm for Logistics

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Abstract

Cloud computing provides a data center at any location in the world. There are so many resources used in a distributed data center that a logistics suppliers and users can buy, sell and rent like execution time, bandwidth, cost, storage and memory. The logistics suppliers and users do not need to know where the data center is located and how to operate or maintain their resources by using cloud computing. They only need to know how to connect to these resources and use of applications needed to perform their jobs. Many companies want to services by using its own local data centers. In logistics information system, information sharing, transfer mode of information using cloud service provider to different logistics partners is a big challenge. Every logistics user wants that information sharing should be in real time within minimum cost. The number of logistics partners and logistics users is growing and they need services for balancing the network traffic and load. The paper has different sections as follows: Section 2 discusses related works. Section 3 discusses scheduling and load balancing algorithms. Section 4 provides our proposed algorithm's design. Section 5 provides experimental result and performance analysis. Section 6 concludes our research work and future work.

Keywords: Cloud Computing, Logistics information system, Load balancing algorithms, DSBP

I. INTRODUCTION

A data center is centralized repository for information distribution and storage of information. Cloud computing provides a data center at any location in the world. There are so many resources used in centralized data center that a logistics suppliers and users can buy or rent like execution time, response time, cost, storage and memory. The logistics suppliers and users do not need to know where the data center is located and how to operate or maintain their resources by using cloud computing environment. They only need to know how to connect to their resources and use of applications needed to performance their jobs, so many companies want to services by using its own local data centers. In logistics information system, information sharing, transfer mode of information using cloud service provider to different logistics partners is a big challenge. Every logistics user wants that information sharing should be in real time. Since centralized data center provides highly efficient data transfer and its services for managing logistics over cloud.

Today, distribution of information between the suppliers and users is very important issue by logistics information system because number of users is increasing to use internet so there is need to balance this overloading and traffic for providing better availability and high performance. There are so many load balancing algorithms but each and every have some problem to reduce response time, overloading and cost. Cloud based distributed algorithms for logistics information system can be used as one of the best solution for avoiding delay in response time and cost due to multiple requests at a time on distributed data center. This paper presents evaluation of existing load balancing algorithm and proposed load balancing algorithm using cloud simulator (Cloud Analyst). The paper also defines improving and comparing the proposed load balancing algorithm against existing load balancing algorithms.

II. RELATED WORKS

Load balancing is a method or technique, which is used for distributing workload on the multiple computers cluster across the network links to achieve recourse utilization and overall time. Load balancing avoids lot of overloading on the resources and divides the traffic between servers and data, so the data can be sent and received without minimum delay.

For new transmission to survive in the competitive market logistic information system is using cloud based service for developing distributed data center. Load balancing helps to balance the increased user's requests on the distributed data center so it is important to use load balancing on web servers in logistics information system. Meenakshi Sharma (2012) proposed a central load balancing policy for virtual machine (CLBVM) to balance the load in distributed environment but it does not consider fault tolerant system.

Bhupendra Verma (2012) analyzed virtual machine load balancing and proposed new virtual machine load balancing algorithm for IaaS framework, implemented of weighted Active Monitoring load balancing algorithm to achieve better performance in response time and data processing time by using Cloud Sim. Dhaval Limbani and Bhavesh Oza (2012) proposed the extended service proximity based routing policy algorithm for routing of user requests such that cost effective data center selection by using Cloud Analyst. Neeraj Bhargava, Rita Bhargava et.al (2013) analyzed and discussed round robin algorithm used by different data center and calculated the overall response time in better performance. S. Mohana Priya and B. Subramani (2013) proposed algorithm by using active monitoring load balancing algorithm and resource aware scheduling algorithm for high performance in cloud systems. The experiment result of proposed algorithm is the efficient virtual machine is selected for process and minimum execution time of task. It increases the performance and reduces the response time and cost. S.Mohana Priya and B.Subramani (2013) proposed algorithm by using active monitoring load balancing algorithm and resource aware scheduling algorithm for improving resource utilization and scheduled load balancing for high performance in cloud systems. The experiment result of proposed algorithm is the efficient virtual machine is selected for process and minimum execution time of task, it increases the performance and reduces the response time and cost. Harvinder Singh and Rakesh Chandra Gangwar (2014) presented comparative study of load balancing algorithms based on the parameters like; response time and cost which are reduced by round robin and active monitoring load balancing algorithm. Rakesh kr. Mishra and Sreenu Naik Bhukya (2014) proposed priority and extended priority based round robin service broker algorithms which distribute the requests based on the rating of data centers and gives better performance. Deepak Kapgate (2014) proposed and compared the new service broker (DC selection) algorithm with existing service broker algorithm. The proposed algorithm reduced service response time and improved performance of data center. Slesha Nayak and Prangesh Patel (2015) presented comparative study for existing throttled algorithms and proposed throttled algorithm of load balancing in cloud computing. Both are tested and compared in terms of response time, data center service request time and cost by using Cloud Analyst. Reena Panwar and Bhawna Mallick (2015) discussed the various load balancing algorithms and compared them based on parameters like; data processing time and response time etc. by using round robin and throttled scheduling algorithms in Cloud Analyst. Sarika Vasantrao Bodke (2016) presented a comparison of assorted policies which is used for load reconciliation by using throttled, round robin and first come first serve (FCFS) and load aqualization algorithms. This comparison shows that response time was effectively reduced but they did not consider the cost. Mamta Khanchi and Sanja tyagi (2016) proposed and implemented a hybrid approach for virtual machine level load balancing. This algorithm distribute workload among the virtual machines that are available in data center at the same time to minimize the overall response time and data center processing time.

In this paper we propose a distributed service broker policy algorithm (DSBP) with the best possible response time, delay and minimum cost in selecting the most suitable data center. DSBP is the mainly implementation of throttled algorithm for taking bandwidth, latency and size of job which achieves minimum response time and minimum cost.

III. PROBLEMS IN EXISTING ALGORITHMS

In distributed data center, the main propose of service brokers is to direct the user request to the best data center with better performance because service broker policy has to efficiently select the best data center for the job considering many issues like response time, service time and cost. For directing the user request to the best data center there are so many load balancing policy algorithms like network latency based, service proximity based routing, etc. (Sharma, 2012) modified the throttled load balancing algorithm and proposed virtual machine load balancing algorithm which reduce response time effectively but not reduce cost. (Bhargava, 2013) analyzed round robin algorithm by using different data center and user base (UB). This result shows that response time is to be minimize but cost should need to be reduce also. (Priya, 2013) proposed a new load balancing algorithm for virtual machine, in which virtual machine is selected for process and minimum execution time of task. It reduces the response time and total cost but there is some issues to improve the response time efficiency parallel to the cost performance. (Harvinder, 2014) analyzed the comparison of three existing load balancing algorithm; round robin, active monitoring, throttled algorithm. The experimental result of theses algorithms shows the reduction in response but increment in cost. (Rakesh, 2014) . proposed priority and extended priority based round robin service broker algorithms which distribute the requests based on the rating of data centers and gives better performance in response time but cost is not reduces similar to response time. (Deepak Kapgate, 2014) proposed and compared the new service broker (DC selection) algorithm with existing service broker algorithm. The proposed algorithm reduced service response time but cost is not considered. (Slesha, 2015) proposed throttled algorithm and compared it with existing throttled algorithm. The proposed algorithm reduces response time and cost in better performance but if data center will be increased then response time cost will be increased. (Reena ,2015) discussed the various load balancing algorithms and compared them based on parameters like; data processing time and response time etc. by using round robin and throttled scheduling algorithms in Cloud

Analyst. The result of both algorithms shows reduction in response time and cost but if data centers will be increased then response time and cost will increase by using separately round robin and throttled algorithm. (Bodke, 2016) also presented round robin algorithm. In this algorithm the time slits into multiple slices and every node is given a specific time slice. There are multiple requests in multiple processes. Each process is given time slot. If user request completes among time then user must not wait otherwise user need to watch for its next time slot. This may create slow process and overloading. (Mamta, 2016) proposed virtual machine load balancing algorithm which is a combination of round robin, throttled, (ESCE) equally spread current execution and hybrid algorithm. In the below table (Tabel.1) we can see so many highlights of different algorithms and drawbacks. we analyze the results of various algorithms in terms of response time and cost to identify the improvement of round robin, active monitoring and throttled algorithms.

Authors (Year)	Algorithms								
	Round Robin	Active Monitoring	Throttled Load	Average Response	Total				
	Algorithm	Algorithm	balancing Algorithm	Time (ms)	Cost (\$)				
Limabani & Oza (2012)	*	*	*	*	967.27				
Jasmin James (2012)	\checkmark	\checkmark	\checkmark	\checkmark	694.82				
Harvinder Singh (2013)	*	\checkmark	\checkmark	112.18	*				
S. Mohana Priya (2013)	*	*	*	171.43	242.05				
Rakesh kr. (2014)	*	*	*	98.18	144.52				
Kunal Kishore (2014)	*	*	*	481.54	*				
Ali Naser (2014)	*	*	\checkmark	139.98	59.50				
Sunny Nandwani	✓	✓	√	92.10	*				
(2015)									
Slesha (2015)	*	*	\checkmark	50.1004	1.324				
Reena Panwar (2015)	\checkmark	*	\checkmark	101.61	50.27				
Sarika Vasantrao	\checkmark	*	\checkmark	*	10.10				
Bodke (2016)									
Mamta Khanchi	\checkmark	*	\checkmark	92.45	*				
(2016)									
Simar Preet Singh	\checkmark	*	\checkmark	57.68	*				
(2016)									
Er. Pooja (2016)	*	*	*	57.08	12.85				
Pradeep Singh (2016)	\checkmark	*	*	484.25	*				
Saurabh Shukla (2016)	\checkmark	*	✓ 1:0: 1 1 1	150.87	*				

Table.1: Existing Algorithm simulation results

A new DSBP algorithm has been proposed from modifying the throttled load balancing algorithm in Virtual Machine environment of cloud computing in order to achieve better response time, processing time and cost.

IV. IMPLEMENTATION OF EXISTING AND PROPOSED ALGORITHMS

The proposed service broker policy selects the data center based on the job size, the expected processing time, the network latency and the available bandwidth to minimize the overall response and processing time. The routing policy considers different factors like; request size, user size, number of virtual machines, number of data centers, available bandwidth, response time, data transfer cost etc. The processing time is based on the data center specification such as Ram, CPU and VM configuration. In the below table we have implementation the method of round robin, active monitoring and throttled load balancing algorithm. DSBP is the modification of throttled load balancing algorithm. As a result, the proposed policy accommodates the current needs by taking real-time values to calculate the processing time to minimize the time needed to make the forwarding decision by the broker. Note that the job processing time can vary depending on the computational task to be performed. For instance, a smaller job requires less processing time if there was no I/O operation involved. However, since it is not the service broker functionality to analyze the jobs and determine their complexity, we considered the size of the job as an indication to the needed processing time with a positive relation between them. Here we have all the implemented methods of some existing algorithms. In the below table we can see all the methods of existing and proposed DSBP algorithms. The proposed algorithm is used to improve throttled algorithm. This improved throttled algorithm works well even though underlying capacity of each virtual machines are different because the hardware configuration of virtual machines are different.

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<pre>vmId=thisVmId; } allocatedVm(vmId);</pre>	ount){	if (currCount < minCount){	return vmId;	vmId=thisVmId;		
} allocatedVm(vmId);	it;	minCount = currCount;	}	}		
		vmId=thisVmId;		}		
		}		}		
return vmId;		}		allocatedVm(vmId);		
		}		return vmId;		
allocatedVm(vmId):		allocatedVm(vmId);				
return vmId;						
3		}				
		,				

Tabel.2: Methods of Implemented Load Balancing Algorithms

V. EXPERIMENTAL RESULTS AND PERFORMANCE ANALYSIS

The performance of the proposed DSBP algorithm has been analyzed based on the result of simulation done in Cloud Analyst. The classes of Cloud Analyst simulator have been extended to utilize newly proposed DSBP algorithm. In the below illustrations, the response time and load are analyzed by using methods of round robin, active monitoring, throttled and DSBP algorithm for distributed data center.

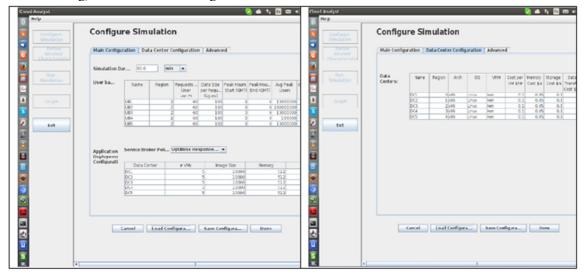


Figure (a) Figure (b) Figure.1(a): Main Configuration Screen of UserBase and Application Deployment. (b) Main Configuration Screen of Data Center

Cloud Analyst

c)

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centers: (Equivalent to number of simultaneous requests a single applicaiton	10				

	Run Simulation	Bases: (Equivalent to number of simultaneous	10
a S Ø	Graph Exit	Centers: (Equivalent to number of simultaneous requests a single application server	10
2		Flot in all the instant in the inger per Focues: Flucture:	100
		Luau valaining puing across VM's in a single Data Center	DS8P
		Cancel Load Configur	ra Save Configura Done

Figure 2: Configure Simulation Screen for proposed Distributed Service Broker Policy Algorithm

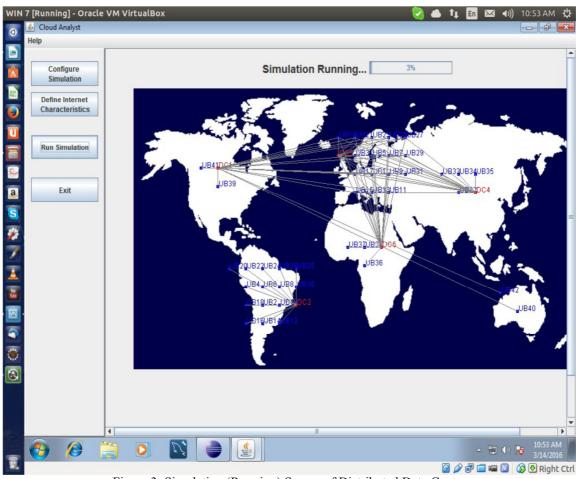


Figure.3: Simulation (Running) Screen of Distributed Data Center

Simulation Setup and analysis of result s are carried out for a period of 60 hrs by taking different number of users and 6 data centers from DC1, DC2,DC6. Each data center has 5 virtual machines separately. The other parameters are fixed according to table no.

Table.3. Parameters and values

Parameters	Values				
VM image size	10000				
VM Memory	1024 MB				
VM Bandwidth	1000				
Data Center (DC)	6				
Virtual Machine (VM) in each data center	5				
DC Memory per machine	1000 MBPS				
DC storage per machine	1GB				
DC available bandwidth per machine	1000000				

After performance different experiments by cloud analyst successfully, we get overall response time of different load balancing algorithms as given in the table. Analytical result of overall response time and load based on various algorithms in cloud computing environment is shown below:

Table.4: Comparison of Overall Response Time and Data Center (DC) processing Time Summary

Distributed Data Center Over All Response Time (RT)												
Algorithm	hm Round Robin Algorithm		Active Monitoring Algorithm		Throttled Algorithm		DSBP Algorithm					
Response Time	Avg(ms)	Min(ms)	Max(ms)	Avg(ms)	Min(ms)	Max(ms)	Avg(ms)	Min(ms)	Max(ms)	Avg(ms)	Min(ms)	Max(ms)
Over all RT	1188.11	14.58	5780.21	117.17	10.50	4933.89	185.50	14.62	5780.05	58.52	0.00	65.00
DC Processing Time	960.43	0.00	36061.25	520.79	0.00	30765.56	946.23	0.01	36059.00	6.68	0.00	100.56

Table.5: Comparison of Total Cost Summary

Distributed Data Center data transfer Cost										
Cost	Round Robin	Active Monitoring	Throttled	DSBP						
	Algorithm	Algorithm	Algorithm	Algorithm						
Total VM cost (\$)	0.04	0.42	0.42	0.40						
Total Data Transfer cost (\$)	0.85	0.85	0.85	0.08						
Grand Total (\$)	1.25	1.25	1.27	0.48						

The above table represents the values of all existing algorithms and proposed DSBP algorithm for overall response time from different data centers which shows DSBP has better performance showing in below graph (figure. no. 2).

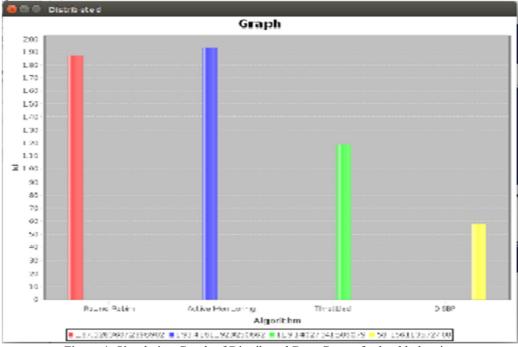


Figure.4: Simulation Graph of Distributed Data Center for load balancing

VI. CONCLUSION

This paper presents the improvement of load balancing algorithm as servers used to get overloaded. Hence we

conclude that DSBP algorithm is an effective and efficient for distributed data center in logistics operations. After analyzing the performance results, its seen that DSBP performs better than the round robin, active monitoring and throttled load balancing algorithms in the same time. The proposed algorithm DSBP not only processes more transactions, it also reduces overall response time and load on the distributed data center. Performance of proposed algorithm confirms that it will be beneficial for logistics information system to track and manage the operation for sharing information from different location at same time.

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