

離散最適化による最適保護区連結ネットワークの探求

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連結ネットワーク集約制約:

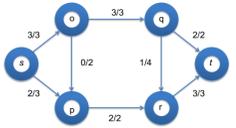
ある特定の林分に対し、その連結の集約に対する制約
ここでは、保護区連結ネットワーク形成に応用

本研究の目的:

制約式により保護区連結ネットワークを形成し、最適保護区連結ネットワークを探求できる離散最適化モデルを構築

空間的制約:

ユニットの連結を伴う制約。特に隣接に関わる制約。集約化あるいは団地化。最大伐採面積の規制。



1. Spatial component for unit aggregation (MFC)

$y_{ij} = \begin{cases} 1 & \text{if an arc is connected between } i\text{-th and } j\text{-th nodes} \\ 0 & \text{otherwise} \end{cases}$

$w_{i,j}$: flow from i - to j -th node

$y_{ij} + y_{ji} \leq 1, \forall i, j \in NB_i$

$y_{i0} + \sum_{j \in NB_i} y_{ij} = 1, \forall i$

$w_{ij} \leq L \cdot y_{ij}, \forall i, j \in NB_i$

$w_{i0} + \sum_{j \in NB_i} w_{ij} = \sum_{j \in NB_i} w_{ji} + L_i, \forall i$

$\sum_{i=1}^m w_{i0} = \sum_{i=1}^m L_i$

$w_{ij} \leq L_{ij} \cdot y_{ij}, \forall i, j$

$\{y_{ij}\} \in \{0,1\}, \{w_{ij}\} \geq 0, \forall i, j \in (NB_i \cup \{0\})$

50ha Aggregation

$w_{3,0} = 32.5$
 $w_{4,0} = 48.9$
 $w_{11,0} = 46.9$
 $w_{14,0} = 46.5$

2. Sequential Triangle Connection (STC)

$y_{ij}^{B(n)} = \begin{cases} 1 & \text{if } n\text{-th process connects a boundary between } i\text{- and } j\text{-th nodes} \\ 0 & \text{otherwise} \end{cases}$

$u_{ijk}^{(n)} = \begin{cases} 1 & \text{if the triangle by nodes } (i, j, k) \text{ is created at the } n\text{-th process} \\ 0 & \text{otherwise} \end{cases}$

$y_{ij}^{B(1)} = y_{ij} + y_{ji}, \forall i, j \in NB_i$

$2 \cdot u_{ijk}^{(n)} + 3 \cdot \sum_{l=1}^{n-1} u_{ijk}^{(l)} \leq \sum_{l=1}^n \{y_{ij}^{B(l)} + y_{ik}^{B(l)} + y_{jk}^{B(l)}\} \leq 2 \cdot u_{ijk}^{(n)} + 3 \cdot \sum_{l=1}^{n-1} u_{ijk}^{(l)} + 1, \forall i, j \in MB_i, k \in (MB_i \cap MB_j), n = 1, \dots, N$

$3 \cdot u_{ijk}^{(n)} \leq \sum_{l=1}^{n+1} \{y_{ij}^{B(l)} + y_{ik}^{B(l)} + y_{jk}^{B(l)}\}, \forall i, j \in MB_i, k \in (MB_i \cap MB_j), n = 1, \dots, N$

$y_{ij}^{B(n+1)} \leq \sum_{k \in (MB_i \cap MB_j)} u_{ijk}^{(n)}, \forall i, j \in MB_i, n = 1, \dots, N$

$\sum_{n=1}^N y_{ij}^{B(n)} \leq 1, \forall i, j \in MB_i$

$\{y_{ij}^{B(n)}\}, \{u_{ijk}^{(n)}\} \in \{0,1\}$

3. Temporal & Integrated component

Linking with Harvest Scheduling
Aggregation => Spatial
Harvesting => Temporal

Table 1: Example of treatments

Treatment No.	Decision Variable	Coefficient	Period																	
			1	2	3	4	5	6	7	8	9	10								
1	$x_{1,1}$	$c_{1,1}$	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
2	$x_{1,2}$	$c_{1,2}$	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
3	$x_{1,3}$	$c_{1,3}$	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
4	$x_{1,4}$	$c_{1,4}$	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
5	$x_{1,5}$	$c_{1,5}$	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
6	$x_{1,6}$	$c_{1,6}$	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
7	$x_{1,7}$	$c_{1,7}$	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
8	$x_{1,8}$	$c_{1,8}$	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
9	$x_{1,9}$	$c_{1,9}$	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
10	$x_{1,10}$	$c_{1,10}$	0	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
11	$x_{1,11}$	$c_{1,11}$	0	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
12	$x_{1,12}$	$c_{1,12}$	0	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
13	$x_{1,13}$	$c_{1,13}$	0	0	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0
14	$x_{1,14}$	$c_{1,14}$	0	0	0	0	X	0	0	0	0	0	0	0	0	0	0	0	0	0
15	$x_{1,15}$	$c_{1,15}$	0	0	0	0	0	X	0	0	0	0	0	0	0	0	0	0	0	0
16	$x_{1,16}$	$c_{1,16}$	0	0	0	0	0	0	X	0	0	0	0	0	0	0	0	0	0	0
17	$x_{1,17}$	$c_{1,17}$	0	0	0	0	0	0	0	X	0	0	0	0	0	0	0	0	0	0
18	$x_{1,18}$	$c_{1,18}$	0	0	0	0	0	0	0	0	X	0	0	0	0	0	0	0	0	0
19	$x_{1,19}$	$c_{1,19}$	0	0	0	0	0	0	0	0	0	X	0	0	0	0	0	0	0	0
20	$x_{1,20}$	$c_{1,20}$	0	0	0	0	0	0	0	0	0	0	X	0	0	0	0	0	0	0

Note: X denotes harvesting while 0 denotes no harvesting

$\sum_{k=0}^K x_{ik} = 1, \forall i$

$(1-\alpha) \cdot v_0 \leq \sum_{k=1}^K v_{ik} \cdot x_{ik} \leq (1+\alpha) \cdot v_0, \forall t$

$a_{kl}^v \cdot x_{ik} + a_{kl}^v \cdot x_{jk} \leq 1, \forall i, j \in NB_i, k, l (k \neq l)$

$\{x_{ij}\} \in \{0,1\}$

$2 \cdot z_{ij}^k \leq x_{ik} + x_{jk} \leq 2 \cdot z_{ij}^k + 1, \forall i, j \in NB_i, k$

$\sum_{n=1}^N y_{ij}^{B(n)} = \sum_{k=1}^m z_{ij}^k, \forall i, j \in NB_i$

$\{x_{ij}\}, \{y_{ij}^{B(n)}\}, \{z_{ij}^k\} \in \{0,1\}$

$a_{kl}^v = \begin{cases} 1 & \text{if } k\text{-th and } l\text{-th treatments have any concurrent harvest} \\ 0 & \text{otherwise} \end{cases}$

4. Habitat connection

Arc connection

$y_{ij} + y_{ji} \leq 1, \forall i, j \in NB_i$

$y_{i0} + \sum_{j \in NB_i} y_{ij} = 1, \forall i$

Maximum flow network

Corridor flow constraints

No connection, no flow

$w_{ij}^c \leq N^c \cdot y_{ij}, \forall i, j \in NB_i$

$w_{i0}^c \leq N^c \cdot y_{i0}, \forall i \in S^c$

N^c : the number of habitats

Flow balance, no leakage

$w_{i0}^c + \sum_{j \in NB_i} w_{ij}^c = \sum_{j \in NB_i} w_{ji}^c + 1, \forall i \in S^c$

$\sum_{j \in NB_i} w_{ij}^c = \sum_{j \in NB_i} w_{ji}^c, \forall i \notin S^c$

One connection to super node

$\sum_{i \in S^c} w_{i0}^c = N^c$

$\sum_{i \in S^c} y_{i0} = 1$

結果

Model I

Max NPV

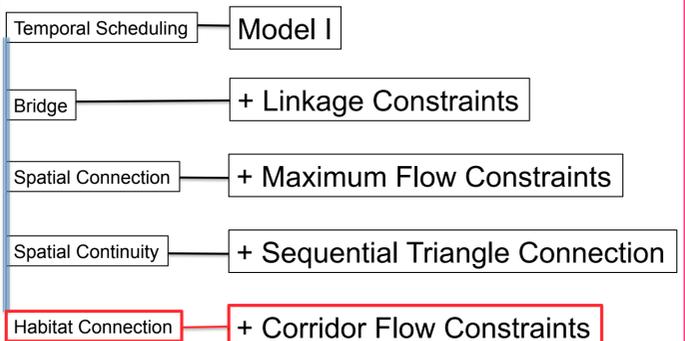
$Z = \max_{\{x_{ij}\}} \sum_{i=1}^m \sum_{j=1}^K c_{ij} \cdot x_{ij}$

st.

$\sum_{k=0}^K x_{ik} = 1, \forall i$ Land Account Constraints

$(1-\alpha)v_0 \leq \sum_{i=1}^m \sum_{j=1}^K v_{ij}^p \cdot x_{ij} \leq (1+\alpha)v_0, p = 1, \dots, T$ Harvest Flow Constraints

MF-Model I for Corridor



Optimal habitat connection

