

# Algorithmic analogies to Kamae-Weiss theorem on normal numbers

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## 1 Selection function

Example

$$\begin{array}{rcccccc} x = & 0 & 1 & 0 & 0 & 1 & \cdots \\ y = & 1 & 0 & 1 & 0 & 1 & \cdots \\ x/y = & 0 & & 0 & & 1 & \cdots \end{array}$$

$$x, y \in \{0, 1\}^\infty.$$

## 2 Kamae-Weiss theorem

*Definition 1.*  $x = x_1x_2\cdots$  is called normal number if

$$\forall s \in \{0, 1\}^* \lim_{n \rightarrow \infty} \#\{1 \leq i \leq n \mid x_i \cdots x_{i+|s|-1} = s\} / n = 2^{-|s|}.$$

Let  $\mathcal{N}$  be the set of normal numbers.

*Definition 2.*  $p$  is called cluster point if there is a sequence  $\{n_i\}$

$$\forall s \ p(s) = \lim_{i \rightarrow \infty} \#\{1 \leq j \leq n_i \mid x_j \cdots x_{j+|s|-1} = s\} / n_i.$$

Let  $V(x)$  be the set of cluster point of  $x$ .

$V(x) \neq \emptyset$  for all  $x$ .

Kamae entropy is defined by

$$h(x) = \sup\{h(p) \mid p \in V(x)\}.$$

**Theorem 1** (Kamae[1]). *Suppose that  $\liminf \frac{1}{n} \sum_{i=1}^n y_i > 0$  then (i) and (ii) are equivalent:*

- (i)  $h(y) = 0$ .
- (ii)  $\forall x \in \mathcal{N} \ x/y \in \mathcal{N}$ ,

Note: The part (i)  $\Rightarrow$  (ii) is appeared in Weiss [3].

## 3 van Lambalgen's conjecture

$K$  : prefix Kolmogorov complexity.

$\mathcal{R}$  : the set of Martin-Löf random sequences with respect to  $(1/2, 1/2)$ -i.i.d. process.

In van Lambalgen [2] the following equivalence is conjectured:

- (i)  $\lim_{n \rightarrow \infty} K(y_1^n) / n = 0$ .
- (ii)  $\forall x \in \mathcal{R} \ x/y \in \mathcal{R}$ .

## 4 Proposition 1

**Proposition 1.** *Suppose that  $y$  is Martin-Löf random with respect to some computable probability  $P$  and  $\sum_{i=1}^\infty y_i = \infty$ . Then the following*

*two statements are equivalent:*

- (i)  $y$  is computable.
- (ii)  $\forall x \in \mathcal{R} \ x/y \in \mathcal{R}^y$ .

$\mathcal{R}^y$ : the set of Martin-Löf random sequences with respect to  $(1/2, 1/2)$ -i.i.d. process relative to  $y$ .

## 5 Weak randomness

*Definition 3.*  $y$  is called weakly random with respect to a computable  $P$  if

$$\lim_{n \rightarrow \infty} K(y_1^n) / n = \lim_{n \rightarrow \infty} -\frac{1}{n} \log P(y_1^n).$$

$y$  is weakly random with respect to  $(1/2, 1/2)$ -i.i.d. process if

$$\lim_{n \rightarrow \infty} K(y_1^n) / n = 1.$$

Note

$$y \in \mathcal{R} \rightarrow \lim_{n \rightarrow \infty} K(y_1^n) / n = 1 \rightarrow y \in \mathcal{N}.$$

None of the converse is true.

## 6 Proposition 2

**Proposition 2.** *Suppose that  $y$  is weakly random with respect to a computable measure and  $\liminf \frac{1}{n} \sum_{i=1}^n y_i > 0$ .*

*Then the following two statements are equivalent:*

- (i)  $\lim_{n \rightarrow \infty} K(y_1^n) / n = 0$ .
- (ii)  $\forall x \ \lim_{n \rightarrow \infty} K(x_1^n) / n = 1 \rightarrow \lim_{n \rightarrow \infty} \frac{1}{|x_1^n / y_1^n|} K(x_1^n / y_1^n | y_1^n) = 1$ .

Example: computable sequences and sturmian sequences satisfy (i).

## References

- [1] T. Kamae. Subsequences of normal numbers. *Israel J. Math.*, 16:121–149, 1973.
- [2] M. van Lambalgen. *Random sequences*. PhD thesis, Universiteit van Amsterdam, 1987.
- [3] B. Weiss. Normal sequences as collectives. In *Proc. Symp. on Topological Dynamics and Ergodic Theory*. Univ. of Kentucky, 1971.