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A Fractal which violates the Axiom of Determinacy

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Abstract

By use of the axiom of choice I construct a symmetrical and selfsimilar subset $A \subseteq [0, 1] \subseteq \mathbb{R}$. Then by an elementary strategy stealing argument it is shown that A is not determined. The (possible) existence of fractals like A clarifies the status of the controversial Axiom of Determinacy.¹

In this note, I present an argument against the unrestricted axiom of determinacy (=AD).

Fix $A \subseteq [0,1] \subseteq \mathbb{R}$. We define an infinite game \mathcal{G}_A as follows. The initial position is I := [0,1]. Each position will be an interval $[a,b] \subseteq [0,1]$, $a, b \in \{\frac{p}{2^k} : p, k \in \mathbb{N}\}$. In each position [a,b] there are always two legal moves. The player who has the turn can move "left" or can move "right". If the player moves "left" the new position is $[a, \frac{a+b}{2}]$. If the player moves "right" the new position is $[a, \frac{a+b}{2}]$. If the player moves "right" the new position is $[a, \frac{a+b}{2}]$. If the first move. In each actual game, successively the players construct a sequence $I = [0, 1] \supseteq I_1 \supseteq I_2 \supseteq \dots$ of closed intervals. The interval I_{j+1} is either the left or the right closed interval of I_j . Each actual game produces a point $p := \bigcap_{j \in \mathbb{N}} I_j \in [0, 1]$.

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According to the rules of \mathcal{G}_A player \mathcal{A} wins if $p \in A$. Otherwise (i.e. when $p \notin A$) player \mathcal{B} wins.

A player has a winning strategy if there is a protocol (i.e. a map from the set of positions to { 'left', 'right'}) which guarantees victory. The set of positions can be divided into 3 classes. The positions which are won for player \mathcal{A} , the positions which are won for player \mathcal{B} , and the controversial class of the positions which are undetermined. According to the axiom of determinacy the last class is always empty. Each actual game has a winner. So if the players are clever enough it must be determined who will win the game. The intuition behind **AD** is that if \mathcal{A} and \mathcal{B} have infinite powers the same player will win each game. If for instance \mathcal{A} wins the first game, \mathcal{A} ought also to win the second game. The argument is that if player \mathcal{B} wins this new game, sometimes in the first game \mathcal{B} could not have possibly played optimally ².

It is well-known that \mathbf{AC} (the axiom of choice) and \mathbf{AD} are contradictory. The usual proof uses a diagonal argument combined with the fact that the number of strategies is 2^{\aleph_0} [1], [2]. The status of \mathbf{AD} has been examined in great depth [4],[5]. There seems to be two approaches. One can accept \mathbf{AC} and ask which sets are determined. This leads to questions which are independent of the usual axiomatization of set theory [2],[4], [5]. The other and more radical approach is to discard the axiom of choice [3]. The main argument is that \mathbf{AD} is deductively strong and has many nice consequences, [2],[4]. Still there is no doubt that most of us prefer \mathbf{AC} .

Let \mathbb{Q} denote the rational numbers.

Theorem (AC) There exists a set $A \subseteq [0, 1]$ such that A and $A \cap [0, \frac{1}{2}]$ are isomorphic under the map $x \to \frac{x}{2}$, and $A \setminus \mathbb{Q}$ and $A^c \setminus \mathbb{Q}$ are isomorphic under the map $x \to 1 - x$.

Proof: Consider the collection **J** of pairs (A, B) where $A, B \subseteq [0, 1] \setminus \mathbb{Q}$, where $A \cap B = \emptyset$ and where $A = A \cdot \mathbb{Q} \cap [0, 1] \setminus \mathbb{Q}$ and $B = B \cdot \mathbb{Q} \cap [0, 1] \setminus \mathbb{Q}$. The set **J** of such pairs are ordered inductively under inclusion. According to Zorn's lemma (which is equivalent to **AC**) there must be a pair $(A, B) \in \mathbf{J}$ which is maximal with respect to inclusion. We claim that $A \cup B = [0, 1] \setminus \mathbb{Q}$. Otherwise there would be $x \in ([0, 1] \setminus \mathbb{Q}) \setminus (A \cup B)$. Notice that, $(1-x) \cdot \mathbb{Q} \cap A = \emptyset$, $x \cdot \mathbb{Q} \cap (1-x) \cdot \mathbb{Q} = \emptyset$ and $x \cdot \mathbb{Q} \cap B = \emptyset$. Thus $(A \cup x \cdot \mathbb{Q}, B \cup (1-x) \cdot \mathbb{Q})$

²Unless of course \mathcal{A} first deviate from the line of play in the first play. But to deviate and lose does not seem wise.

is well-defined and belongs to **J**. This violates the maximality of (A, B). \Box **Theorem** All the positions $(I', A \cap I')$ are isomorphic, when the points in \mathbb{Q} are ignored. No move can make any difference to the outcome because all positions are isomorphic. Each game produces a winner.

Proof: We ignore the points in \mathbb{Q} . Notice that the two positions which can be reached from I both are isomorphic to I (when the role of A is changed with that of A^c). This is because $f_1: x \to 1-2x$ maps $A^c \cap [0, \frac{1}{2}]$ isomorphic onto A, and because $f_2: x \to 2-2x$ maps $A^c \cap [\frac{1}{2}, 1]$ isomorphic onto A. \Box **Corollary (AC)** \mathcal{G}_A is not determined.

Proof: If player \mathcal{A} has a winning strategy, player \mathcal{B} can steal it. It is not difficult to show that the points in \mathbb{Q} do not affect this argument. \Box

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