Generalized Turán problems for disjoint copies of graphs

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Abstract

Given two graphs H and F, the maximum possible number of copies of H in an F-free graph on n vertices is denoted by ex(n,H,F). We investigate the function ex(n,H,kF), where kF denotes k vertex disjoint copies of a fixed graph F. Our results include cases when F is a complete graph, cycle or a complete bipartite graph.

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1 Introduction

The vertex set of a graph G is denoted by V(G) and its edge set is denoted by E(G). The disjoint union $G \cup H$ of graphs G and H with disjoint vertex sets V(G) and V(H) is the graph with the vertex set $V(G) \cup V(H)$ and edge set $E(G) \cup E(H)$. The join G + H, of graphs G and H with disjoint vertex sets is the graph obtained by taking a copy of G and a copy of G on disjoint vertex sets and adding all the edges between them.

Given a positive integer k and a graph F, the vertex disjoint union of k copies of the graph F is denoted by kF. Let C_l denote a cycle of length l, $K_{s,t}$ denote the complete bipartite graph with parts of sizes s and t and let K_r denote the complete graph on r vertices.

For a set of graphs \mathcal{F} the $Tur\'{a}n$ number of \mathcal{F} , $ex(n,\mathcal{F})$, denotes the maximum number of edges of an n-vertex graph having no member of \mathcal{F} as a subgraph. If \mathcal{F} contains only a single graph F, we simply denote it by ex(n,F). This function has been intensively studied, starting with the theorems of Mantel [15] and Tur\'{a}n [19] that determine $ex(n,K_{r+1})$ for $r \geq 3$. Tur\'{a}n also showed in [19] that a complete r-partite graph on n vertices with as equal parts as possible is the unique extremal graph. This extremal graph is called $Tur\'{a}n$ graph and is denoted by $T_r(n)$. See, for example, [8, 18] for surveys on this topic.

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Simonovits [17] and Moon [16] showed that if n is sufficiently large, then $K_{k-1} + T_r(n - k + 1)$ is the unique extremal graph for $\mathcal{F} = \{kK_{r+1}\}$. In [12] Gorgol initiated the systematic investigation of Turán numbers of disjoint copies of connected graphs and proved the following.

Theorem 1 (Gorgol). For every non-empty graph F and $k \geq 1$, we have

$$ex(n, kF) = ex(n, F) + O(n).$$

In fact, Gorgol proved the following sharper upper bound: If F is an arbitrary connected graph and k is an arbitrary positive integer, then $ex(n, kF) \le ex(n - (k-1)|V(F)|, F) + {(k-1)|V(F)| \choose 2} + (k-1)|V(F)|(n - (k-1)|V(F)|)$ for $n \ge k|V(F)|$. For recent results about Turán numbers of disjoint copies of graphs see [4, 14, 20].

Given a graph H and a set of graphs \mathcal{F} , the maximum possible number of copies of H in an n-vertex graph that does not contain any copy of $F \in \mathcal{F}$ is denoted by $ex(n, H, \mathcal{F})$ and is called Generalized Turán number. If $\mathcal{F} = \{F\}$, we simply denote it by ex(n, H, F). Note that $ex(n, K_2, F) = ex(n, F)$. Erdős [5] determined $ex(n, K_s, K_t)$ exactly. We will later use the following consequence of his result.

Proposition 2 (Erdős). For s < t we have:

$$ex(n, K_s, K_t) = (1 + o(1)) {t-1 \choose s} \left(\frac{n}{t-1}\right)^s.$$

Another notable result is that of Bollobás and Győri [2], who showed that $ex(n, K_3, C_5) = \Theta(n^{3/2})$. The systematic study of the function ex(n, H, F) was initiated by Alon and Shikhelman in [1].

The function ex(n, H, F) is closely related to the area of Berge hypergraphs. A Berge cycle of length k is an alternating sequence of distinct vertices and hyperedges of the form $v_1, h_1, v_2, h_2, \ldots, v_k, h_k, v_1$ where $v_i, v_{i+1} \in h_i$ for each $i \in \{1, 2, \ldots, k-1\}$ and $v_k, v_1 \in h_k$ and is denoted by Berge- C_k . Győri and Lemons [13] showed that any r-uniform hypergraph avoiding a Berge- C_{2l+1} contains $O(n^{1+1/l})$ hyperedges. They also showed that any r-uniform hypergraph avoiding a Berge- C_{2l} contains $O(n^{1+1/l})$ hyperedges. These results easily imply the following.

Theorem 3. We have

(a) For any $r \geq 3$, $l \geq 2$, we have

$$ex(n, K_r, C_{2l+1}) = O(n^{1+1/l}).$$

(b) For any $r, l \geq 2$, we have

$$ex(n, K_r, C_{2l}) = O(n^{1+1/l}).$$

Proof. We will prove (a) and (b) simultaneously. Let G be a C_k -free graph. Replace each clique of size r in it with a hyperedge on the same vertex set as the clique. It is easy to see that the resulting r-uniform hypergraph H does not contain a Berge- C_k , and in both cases k = 2l and k = 2l + 1, H has at most $O(n^{1+1/l})$ hyperedges by the theorem of Győri and Lemons [13] mentioned before. This completes the proof as the number of cliques in G is equal to the number of hyperedges in H.

Alon and Shikhelman [1] noted that while $ex(n, K_3, C_5) = \Theta(n^{3/2})$, we have $ex(n, K_3, 2C_5) = \Theta(n^2)$, showing that ex(n, H, F) and ex(n, H, kF) can have different order of magnitudes, unlike the graph case, where $ex(n, kF) = \Theta(ex(n, F))$ (see Theorem 1).

Our goal in this paper is to explore this phenomenon. Most of our theorems will relate ex(n, H, kF) to ex(n, H, F) for several graphs H and F.

General approach

The most typical example of a kF-free graph is obtained by taking an F-free graph G on n-k+1 vertices, and considering $K_{k-1}+G$. (We will sometimes refer to these k-1 vertices of degree n-1 in K_{k-1} as universal vertices of $K_{k-1}+G$.) Indeed, since any copy of F in $K_{k-1}+G$ must contain a vertex of K_{k-1} and as there are only k-1 vertices in K_{k-1} , it is impossible to find k vertex disjoint copies of F.

For example, let us take a C_5 -free graph G on n-1 vertices, add a new vertex v and connect it to all the vertices of G. This graph shows $ex(n, K_3, 2C_5) = \Omega(n^2)$. In addition to the triangles in G (which are at most $O(n^{3/2})$ by the result of Bollobás and Győri [2] mentioned before), there are triangles which contain v and an edge of G. If G is the Turán graph $T_2(n-1)$, then there are $\Omega(n^2)$ many such triangles. What happens here is that instead of counting the copies of K_3 in a C_5 -free graph, we count the copies of K_2 (which is a subgraph of K_3). As this happens to be of larger order of magnitude, we get more copies of K_3 in a $2C_5$ -free graph than in a C_5 -free graph.

To prove the upper bounds we will need the following operation: for an integer k, a graph F and a kF-free graph G, we consider the maximum number of disjoint copies of F in G. In the rest of the paper we denote the subgraph of G consisting of these copies by G_L , and the set of vertices spanned by G_L is denoted by L(G). We denote by R(G) the set $V(G) \setminus L(G)$ of the remaining vertices, and by G_R the subgraph of G induced by them. We call this partition of the vertices a canonical (k, F)-partition of G. (If it is clear from the context we simply write canonical partition.) Note that G_R is F-free.

Structure of the paper

The rest of this paper is divided into sections based on which graph is forbidden. In Section 2 we prove bounds on ex(n, H, kF) for general F, while in Section 3 one of our main results is to determine the order of magnitude of $ex(n, K_s, kK_t)$ for all $s \ge t \ge 2$ and $k \ge 1$. In Section 4, we obtain bounds on $ex(n, K_r, kC_l)$. In Section 5 we study the case when F is

a complete bipartite graph. We finish our article with some concluding remarks and open problems in Section 6.

2 Forbidding a general F

2.1 Counting arbitrary graphs

For a family of graphs \mathcal{H} , let us define $N(\mathcal{H}, G)$ as the number of copies of members of \mathcal{H} in G. If $\mathcal{H} = \{H\}$, then we simply write N(H, G) instead of $N(\mathcal{H}, G)$.

Let \mathcal{H}^{ind} be the family of all induced subgraphs of a graph H. Let

$$\overline{ex}(n, H, F) := \max\{N(\mathcal{H}^{ind}, G) : G \text{ is an } F\text{-free graph on } n \text{ vertices}\}.$$

Remark 4. Note that if F is a non-empty graph (i.e., contains at least one edge), then

$$\overline{ex}(n, H, F) \ge \binom{n}{\alpha(H)}.$$

Indeed, let G be an F-free graph on n vertices and let $I \in \mathcal{H}$ be an induced subgraph spanned by a largest independent set of H. Then any set of $\alpha(H)$ vertices in G forms a copy of I.

Theorem 5. For any $k \geq 2$ we have,

$$ex(n, H, kF) = O(\overline{ex}(n, H, F)).$$

Moreover, if $k \geq |V(H)|$, then

$$ex(n, H, kF) = \Theta(\overline{ex}(n, H, F)).$$

Proof. For the lower bound, we take an F-free graph G on n-k+1 vertices that contains $\overline{ex}(n-k+1,H,F)$ copies of induced subgraphs of H. Then $K_{k-1}+G$ is obviously kF-free. If $k \geq |V(H)|$, then every copy of an induced subgraph (having at least one vertex) of H in G can be extended to a copy of H in $K_{k-1}+G$, using the vertices of K_{k-1} . (A small technical issue is the following: Let Z be the induced subgraph of H with zero vertices. A copy of Z in G cannot be extended to a copy of H in $K_{k-1}+G$ if k=|V(H)|, but there is only one copy of Z in G.) Thus $K_{k-1}+G$ contains at least $\overline{ex}(n-k+1,H,F)-1$ copies of H. Now using the following standard argument, we conclude that $\overline{ex}(n-k+1,H,F)-1=\Omega(\overline{ex}(n,H,F))$: Consider a graph G on n vertices with $\overline{ex}(n,H,F)$ copies of induced subgraphs of H. Then a subgraph of G induced by a random subset of vertices of size n-k+1, contains at least $(1+o(1))\overline{ex}(n,H,F)$ copies of induced subgraphs of H. On the other hand, this subgraph contains at most $\overline{ex}(n-k+1,H,F)$ copies of induced subgraphs of H. This finishes the proof of the lower bound.

Now we continue with the upper bound. Let us consider a kF-free graph G, and its canonical partition. Then every copy of H in G contains a subgraph in G_R , which contains

an induced subgraph of H (note that this subgraph may have zero vertices). Moreover, each copy of an induced subgraph of H in G_R can be extended to a copy of H in G using vertices from L(G) in at most $2^{|L(G)|} = O(1)$ ways. Therefore, the number of copies of H in G is at most $O(\overline{ex}(n, H, F))$, as desired.

2.2 Counting triangles

Let F_1, \ldots, F_k be graphs different from K_2 and let F be their vertex-disjoint union. (Note that the F_i 's are not necessarily different.)

Theorem 6.

$$ex(n, K_3, F) = \Theta\left(\max_{1 \le i \le k} \{ex(n, K_3, F_i)\} + \max_{1 \le i < j \le k} ex(n, \{F_i, F_j\})\right).$$

Proof of Theorem 6. For the lower bound, consider the following two constructions.

- 1. Take an F_i -free graph containing the largest number of triangles. This graph is obviously F-free, showing that $ex(n, K_3, F) \ge \max_{1 \le i \le k} \{ex(n, K_3, F_i)\}$.
- 2. Now let i, j be two integers with $1 \le i < j \le k$. Take an $\{F_i, F_j\}$ -free graph G_0 on n-1 vertices, and add a universal vertex v. Then any copy of F_i and any copy of F_j in the resulting graph G must contain v, thus there are no vertex-disjoint copies of F_i and F_j in this graph. Therefore it does not contain F. Furthermore, the number of triangles in G is at least the number of edges in G_0 , as these edges form a triangle with v. Thus we have $ex(n, K_3, F) \ge \max_{1 \le i < j \le k} ex(n, \{F_i, F_j\})$.

For the upper bound, we use induction on k. The base case k = 1 is trivial. Let F' be the graph obtained by deleting F_i from F and F'' be the graph obtained by deleting F_j from F. Let us consider an F-free graph G on n vertices. If G is F'-free or F''-free, then we are done by induction. Thus we may assume G contains a copy of F' and a copy of F'', and let E be the union of their vertex sets. Note that these copies share at least one vertex, otherwise there would be a copy of F in G.

Let G' be the subgraph of G induced by $V(G) \setminus L$. Then G' is obviously both F_i -free and F_j -free, hence it contains at most $ex(n, \{F_i, F_j\})$ edges and at most $ex(n, K_3, F_i)$ triangles. Let T_s denote the set of triangles in G which contain exactly s vertices from L. Then we have $|T_3| = O(1), |T_2| = O(n), |T_1| = O(ex(n, \{F_i, F_j\}))$ and $|T_0| \le ex(n, K_3, F_i)$. Adding up these bounds, the proof is complete.

Remark 7. • Note that by Theorem 6, we have $ex(n, K_3, kF) = \Theta(ex(n, K_3, F) + ex(n, F))$, for any integer $k \geq 2$. This shows that when k increases from 1 to 2, the order of magnitude of $ex(n, K_3, kF)$ can increase, but from then on (i.e., for $k \geq 2$), there is no further increase.

• The Compactness conjecture of Erdős and Simonovits [6] states that for any finite family \mathcal{G} of graphs there is a $G \in \mathcal{G}$ such that $ex(n,\mathcal{G}) = \Theta(ex(n,G))$. It is known to be true for several classes of graphs, for example if \mathcal{G} contains at most one bipartite graph.

Let $ex^{sec}(n, F)$ be the second largest of the Turán numbers $ex(n, F_i)$, $1 \le i \le k$. Note that if the Compactness conjecture is true (even if it is true only for families of two graphs), then $\max_{1 \le i < j \le k} ex(n, \{F_i, F_j\}) = \Theta(ex^{sec}(n, F))$. (In particular, if F is non-bipartite, then this is the case.) Thus if the Compactness conjecture is true, then Theorem 6 can be stated as follows:

$$ex(n, K_3, F) = \Theta\left(\max_{1 \le i \le k} \{ex(n, K_3, F_i)\} + ex^{sec}(n, F)\right).$$

Definition 8. Let us define

$$ex^*(n, F) := \max_{G} \{(k-1) | E(G)| + N(K_3, G) : G \text{ is an } n\text{-vertex } F\text{-free } graph\}.$$

Note that we have $ex(n, K_3, F) \leq ex^*(n, F) \leq (k-1)ex(n, F) + ex(n, K_3, F)$. Let us consider an arbitrary F, and let F_u be the graph we get by deleting the vertex u from F.

Theorem 9. Let $|V(F)| \ge 4$. Then for every $u \in V(F)$ we have

$$ex^*(n-k+1,F) \le ex(n,K_3,kF) \le ex^*(n,F) + (k-1)|V(F)|ex(n,F_u) + O(n).$$

Proof. For the lower bound of $ex(n, K_3, kF)$, take an F-free graph G on n - k + 1 vertices for which $(k-1)|E(G)| + N(K_3, G)$ is maximum, and consider $K_{k-1} + G$. Then every edge of G, together with the k-1 universal vertices, gives k-1 triangles. This shows $ex(n, K_3, kF) \ge ex^*(n-k+1, F)$.

For the upper bound we use induction on k, the case k = 1 is trivial. If a kF-free graph G does not contain (k-1)F, then we are done by induction. So we may assume G contains a copy of (k-1)F.

Recall that F_u is the graph obtained by removing the vertex u from F. Let X(u) be the set of neighbors of u in F. Let F^* be the graph we get by taking (k-1)|V(F)|+1 vertex disjoint copies of F_u and an additional vertex v (we call it the *center* of F^*), that is connected to all the vertices in X(u) in each copy of F_u . In other words we take (k-1)|V(F)|+1 vertex-disjoint copies of F and identify the vertices u from each copy of F. So F^* is created by (k-1)|V(F)|+1 copies of F that intersect only in the center. Observe that if we are given a copy F_0 of F and a copy of F^* , such that F_0 does not contain the center of the F^* , then at least one of the copies of F which create F^* is disjoint from F_0 .

Now let us assume G contains k-1 vertex-disjoint copies of F^* and let v_1, \ldots, v_{k-1} be their centers. In that case every copy F_0 of F contains at least one of the centers. Indeed, otherwise each of the k-1 copies of F^* contain a copy of F that is disjoint from F_0 , and these copies are all disjoint from each other. Therefore they form a copy of kF, a

contradiction. Thus by deleting v_1, \ldots, v_{k-1} we get an F-free graph G', and $ex(n, K_3, kF)$ is at most the number of triangles in G' plus the number of triangles t_i with exactly i vertices from $\{v_1, \ldots, v_{k-1}\}$ for each $1 \leq i \leq 3$. t_1 is at most k-1 times the number of edges in G', whereas $t_2 + t_3 = O(n)$. Therefore, the total number of triangles in G is at most $ex^*(n-k+1,F) + O(n)$, as desired.

So we may assume G does not contain $(k-1)F^*$. Let us consider the canonical partition, i.e. a copy G_L of (k-1)F and the graph G_R induced by the remaining vertices. Then there are O(n) triangles containing 0 or 1 vertices from R(G), and $N(K_3, G_R)$ triangles containing 3 vertices from R(G). It remains to bound the number of triangles which contain exactly one vertex from the (k-1)|V(F)| vertices not in R(G).

Let us consider the subgraph H which consists of the largest number (which is at most (k-2)) of vertex-disjoint copies of F^* where each copy has its center in L(G) and all its other vertices in R(G). Let x be a vertex in L(G) that does not belong to H and consider the graph H_0 induced by its neighborhood in $R(G) \setminus V(H)$. Observe that H_0 cannot contain more than (k-1)|V(F)|+1 vertex-disjoint copies of F_u . Thus the number of edges in H_0 is at most

$$ex(n, ((k-1)|V(F)|+1)F_u) = ex(n, F_u) + O(n)$$

using Theorem 1. In addition the number of edges incident to $R(G) \cap V(H)$ is at most

$$(k-2)((k-1)|V(F)|+1)|V(F)|n = O(n).$$

Thus x is in $ex(n, F_u) + O(n)$ triangles. There are at most (k-1)|V(F)| vertices in L(G). Moreover, there are at most k-2 vertices in L(G) that belong to H, and each of them is in at most $|E(G_R)|$ triangles. Therefore, the number of triangles which contain exactly one vertex from the (k-1)|V(F)| vertices not in R(G) is at most $(k-2)|E(G_R)|+(k-1)|V(F)|ex(n,F_u)+O(n)$.

Therefore, the total number of triangles in G is at most

$$N(K_3, G_R) + (k-2)|E(G_R)| + (k-1)|V(F)|ex(n, F_u) + O(n)$$

$$\leq ex^*(n, F) + (k-1)|V(F)|ex(n, F_u) + O(n),$$

completing the proof.

Remark 10. Note that the proof shows a stronger upper bound. In one of the two cases we get an upper bound matching the lower bound. In the other case we obtain an upper bound of the form $(k-2)|E(G)|+N(K_3,G)$ rather than $(k-1)|E(G)|+N(K_3,G)$ as in the definition of $ex^*(n,F)$. In case $ex(n,F_u)$ has smaller order of magnitude than ex(n,F) and n is large enough, this implies that $ex(n,K_3,kF) = ex^*(n-k+1,F)$.

Note that if F_u is not a forest, we also know $ex(n, |V(F)|F_u) = (1 + o(1))ex(n, F_u)$ by Theorem 1. Theorem 9 shows that if $ex(n, F_u) = o(ex^*(n, F))$, then we have

$$ex(n, K_3, kF) = (1 + o(1))ex^*(n, F).$$
 (1)

This implies the following.

Corollary 11. (a) If F has chromatic number r > 3, then we have

$$ex(n, K_3, kF) = (1 + o(1)) {r - 1 \choose 3} \left(\frac{n}{r - 1}\right)^3.$$

(b) For $k \geq 1$, we have

$$ex(n, K_3, kK_{2,t}) = (1 + o(1)) \left(\frac{(k-1)(t-1)^{1/2}}{2} + \frac{(t-1)^{3/2}}{6} \right) n^{3/2}.$$

Proof. First let us prove (a). For any vertex $u \in V(F)$, trivially $ex(n, F_u) = O(n^2)$. Alon and Shikhelman [1] showed that $ex(n, K_3, F) = (1 + o(1))\binom{r-1}{3}\left(\frac{n}{r-1}\right)^3$ which is equal to $ex^*(n, F)$ asymptotically. Thus $ex(n, F_u) = o(ex^*(n, F))$, so by (1), we are done.

Now we prove (b). Alon and Shikhelman [1] showed that

$$ex(n, K_3, K_{2,t}) = (1 + o(1)) \frac{(t-1)^{3/2}}{6} n^{3/2}.$$

In fact, they establish the lower bound $ex(n, K_3, K_{2,t}) \ge (1 + o(1)) \frac{(t-1)^{3/2}}{6} n^{3/2}$ by considering the $K_{2,t}$ -free graph constructed by Füredi [7] and counting the number of triangles in it. This graph contains $\frac{(t-1)^{1/2}}{2} n^{3/2} (1 + o(1)) = ex(n, K_{2,t})$ edges, so it follows that

$$ex^*(n, K_{2,t}) \ge (1 + o(1)) \left(\frac{(k-1)(t-1)^{1/2}}{2} + \frac{(t-1)^{3/2}}{6} \right) n^{3/2}$$

and this bound is sharp since by definition, $ex^*(n, K_{2,t}) \le ex(n, K_3, K_{2,t}) + (k-1)ex(n, K_{2,t})$. Now it can be easily seen that $ex(n, F_u) = O(n) = o(ex^*(n, F))$ if $F = K_{2,t}$ for some u. So, by (1) again, the proof is complete.

2.3 Counting complete graphs

Theorem 12. Let F be a graph and k, r be integers. Let $\max_{m \le r} (ex(n, K_m, F)) = ex(n, K_{m_0}, F)$. Then we have

$$ex(n, K_r, kF) = O(ex(n, K_{m_0}, F)).$$

Moreover, if $k > r - m_0$, then

$$ex(n, K_r, kF) = \Theta(ex(n, K_{m_0}, F)).$$

Proof. Let us consider a kF-free graph G, and its canonical partition. Then every copy of K_r consists of m vertices in R(G) and r-m vertices in L(G) for some integer $m \leq k$. These

latter ones can be chosen in at most $\binom{k|V(F)|}{r-m} = O(1)$ ways, thus we have

$$ex(n, K_r, kF) = O(\sum_{m \le r} ex(n, K_m, F)) \le O(ex(n, K_{m_0}, F)).$$

Let us now consider the graph on $n - r + m_0$ vertices that contains the most copies of K_{m_0} , and add $r - m_0$ universal vertices. The resulting graph contains $\Omega(ex(n, K_{m_0}, F))$ copies of K_r . On the other hand, every copy of F contains at least one of the additional vertices, thus there are at most $r - m_0 < k$ pairwise vertex-disjoint copies of F in it.

3 Forbidding complete graphs

As we mentioned in the introduction, Erdős [5] determined the exact value of $ex(n, K_s, K_t)$ for s < t and Simonovits [17] determined the exact value of $ex(n, kK_t)$ for sufficiently large n.

In this section we investigate the function $ex(n, K_s, kK_t)$. First we present our main result of this section, which determines the order of magnitude for every s, t and k. Then we show two asymptotic results (for special values of s and t).

3.1 $ex(n, K_s, kK_t)$

The following theorem determines the order of magnitude of $ex(n, K_s, kK_t)$ for all s, t and k (as n tends to infinity). Note that if s < t, then the Turán graph shows $ex(n, K_s, kK_t) = \Theta(n^s)$.

Theorem 13. Let $s \ge t \ge 2$ and $k \ge 1$ be arbitrary integers and let $x := \lceil \frac{kt-s}{k-1} \rceil - 1$. Then we have

$$ex(n, K_s, kK_t) = \Theta(n^x).$$

Proof. For the lower bound, consider the Turán graph $K_{s-x}+T_x(n-s+x)$. This graph is kK_t -free as k vertex-disjoint copies of K_t together contain at most kx vertices from $T_x(n-s+x)$ and at most s-x vertices from K_{s-x} . Thus they together contain at most s+(k-1)x < kt vertices. On the other hand, the Turán-graph $T_x(n-s+x)$ contains $\Omega(n^x)$ copies of K_x , and they all can be extended to different copies of K_s .

To prove the upper bound we will repeatedly apply the canonical partition operation.

Step 1:

(1.1) Consider a kK_t -free graph G_1 and its canonical (k, K_t) -partition.

(1.2) Let us fix an arbitrary nonempty $X_1 \subset L(G_1)$ (of size x_1) and let

$$\mathcal{A}(X_1) := \{A : A \text{ is a copy of } K_s \text{ in } G_1 \text{ with } |V(A) \cap L(G_1)| = X_1\}.$$

Note that $V(A) \cap R(G_1)$ spans a (copy of) K_{s-x_1} for all $A \in \mathcal{A}(X_1)$ and let G_2 be the subgraph of G_1 spanned by the union of $\{V(A) \cap R(G_1) : A \in \mathcal{A}(X_1)\}$. So G_2 is a graph on the vertex set $R(G_1)$. We consider two cases:

Case 1: G_2 contains k disjoint copies of K_{s-x_1} . Observe that $s-x_1 < t$ and let us denote the corresponding copies of K_{s-x_1} by A_1, \ldots, A_k . We claim that in this case we have

$$x_1 + k(s - x_1) < kt.$$

Otherwise we could complete A_1, \ldots, A_k from X_1 into k disjoint copies of K_t as every vertex of G_2 is connected to every vertex of X_1 and that would be a contradiction. This inequality implies $s - x_1 \leq x$, and as obviously there are $O(n^{s-x_1})$ copies of K_s in $\mathcal{A}(X_1)$ we stop the application of canonical partitions here, and we are done.

Case 2: G_2 does not contain k disjoint copies of K_{s-x_1} . In this case we jump to Step 2.

Now we describe the *i*th step for $i \geq 2$:

Step i: We have from the (i-1)th step:

- 1) a sequence of subsets $X_1, L(G_1), \ldots, X_{i-1}, L(G_{i-1})$ of the vertex set of our initial graph G_1 , where:
 - 1.1) $X_i \subset L(G_i)$ for all $j \leq i-1$, and
 - 1.2) $L(G_j)$ $(j \le i 1)$ are pairwise disjoint.
 - 2) A set of copies of K_s in G_1 parametrized by X_1, \ldots, X_{i-1} :

$$\mathcal{A}(X_1,\ldots,X_{i-1}):=$$

 ${A : A \text{ is a copy of } K_s \text{ in } G_1 \text{ with } |V(A) \cap L(G_1)| = X_1, \dots, |V(A) \cap L(G_{i-1})| = X_{i-1}}.$

3) G_i , that is the subgraph of G_{i-1} spanned by the (union of the) edges of the copies in $\mathcal{A}(X_1,\ldots,X_{i-1})$ on the vertex set $R(G_{i-1})$, and G_i is a $kK_{s-x_1-\ldots-x_{i-1}}$ -free graph.

We do the following in Step i:

- (i.1) We consider the canonical $(k, K_{s-x_1-...-x_{i-1}})$ -partition of G_i .
- (i.2) We fix an arbitrary nonempty $X_i \subset L(G_i)$ (of size x_i) and let

$$\mathcal{A}(X_1,\ldots,X_i):=$$

 ${A : A \text{ is a copy of } K_s \text{ in } G_1 \text{ with } |V(A) \cap L(G_1)| = X_1, \dots, |V(A) \cap L(G_i)| = X_i}.$

Note that $V(A) \cap R(G_i)$ spans a (copy of) $K_{s-x_1-...-x_i}$ for all $A \in \mathcal{A}(X_1, ..., X_i)$ and let G_{i+1} be the subgraph of G_i spanned by the union of the edges of the elements of $\{V(A) \cap R(G_i) : A \in \mathcal{A}(X_1, ..., X_i)\}$. So G_{i+1} is a graph on the vertex set $R(G_i)$. We consider two cases:

Case 1: G_{i+1} contains k disjoint copies of $K_{s-x_1-...-x_i}$. Let us denote the corresponding copies of $K_{s-x_1-...-x_i}$ by A_1, \ldots, A_k . We claim that in this case we have

$$x_1 + \ldots + x_i + k(s - x_1 - \ldots - x_i) < kt.$$

Otherwise we could complete the sets $V(A_1) \cap R(G_1), \ldots, V(A_k) \cap R(G_2)$ from $X_1 \cup \ldots \cup X_i$ into k disjoint copies of K_t as all the vertices of G_{i+1} are connected to all vertices in $X_1 \cup \ldots \cup X_i$ and that would be a contradiction. This inequality implies $s - x_1 - \ldots - x_i \leq x$, and as obviously there are $O(n^{s-x_1-\ldots-x_i})$ copies of K_s in $A(X_1,\ldots,X_i)$ we stop the application of canonical partitions here.

Case 2: G_{i+1} does not contain k disjoint copies of $K_{s-x_1-...-x_i}$. In this case we jump to Step (i+1).

Note that our algorithm finishes in at most kt steps as we always choose nonempty subsets. There are at most O(1) ways to pick X_1, \ldots, X_s and every copy of a K_s will be an element of some $\mathcal{A}(X_1, \ldots, X_j)$ for some $j \leq s$ (and we stop the algorithm in Step j), so we are done with the proof.

Theorem 14. If t > s, then we have

$$ex(n, K_s, kK_t) = (1 + o(1)) {t-1 \choose s} \left(\frac{n}{t-1}\right)^s.$$

Proof. To prove the lower bound one just considers the Turán graph $T_{t-1}(n)$.

For the upper bound consider a kK_t -free graph G and its canonical partition. First let us count the copies of K_s that have a common vertex with L(G). There are O(1) ways to pick the vertices from L(G) and $O(n^{s-1})$ ways to pick the remaining vertices from R(G).

Now let us count those copies that are in G_R . As G_R is a K_t -free graph on at most n vertices, by Proposition 2 there are at most

$$(1+o(1))\binom{t-1}{s}\left(\frac{n}{t-1}\right)^s$$

copies of K_s in it. Adding these bounds up, the proof is complete.

Theorem 15. If $s \ge t \ge s - k + 2$ then we have

$$ex(n, K_s, kK_t) = (1 + o(1)) {k-1 \choose s-t+1} \left(\frac{n}{t-1}\right)^{t-1}.$$

Proof. For the lower bound consider the graph $K_{k-1} + T_{t-1}(n-k+1)$. Counting the number of K_s 's that contain exactly s-t+1 vertices from K_{k-1} gives the desired lower bound.

For the upper bound consider a kK_t -free graph G and its canonical partition. Then any copy of K_s contains at most t-1 vertices from R(G). The number of those copies that contain exactly i vertices from R(G) is at most $\binom{t(k-1)}{s-i}\binom{n}{i} = O(n^i)$, thus it is enough to only take care of those copies that contain exactly t-1 vertices from R(G).

Let

$$\mathcal{A} := \{A : A \text{ is a copy of } K_s \text{ in } G \text{ with } |V(A) \cap R(G)| = t - 1\}.$$

Then we want to upper bound the cardinality of \mathcal{A} . Note that any element of \mathcal{A} intersects R(G) in t-1 vertices spanning a complete graph in G. Let B_1, \ldots, B_r be the copies of K_t defining G_L with for some r < k.

Fix an arbitrary $A_1 \in \mathcal{A}$ that intersects B_1 . One can easily see that the cardinality of

$$\mathcal{A}'_1 := \{ A \in \mathcal{A} : V(A) \cap V(B_1) \cap R(G) \neq \emptyset \}$$

is $O(n^{t-2})$. Consider any $A'_1 \in \mathcal{A} \setminus \mathcal{A}'_1$ such that $A_{1,R} = V(A) \cap R(G)$ and $A'_{1,R} = V(A'_1) \cap R(G)$ are disjoint and $A'_1 \cap B_1 \neq \emptyset$. The existence of such A'_1 implies that $A_1 \cap B_1$ and $A'_1 \cap B_1$ are the same one element since otherwise B_1 could be replaced by two disjoint copies of K_t in G (namely those spanned by $V(A_R) \cup \{x\}$ and $V(A'_R) \cup \{y\}$ for some $x \neq y \in B_1$), contradicting the definition of canonical partition.

In the same way for every $1 < i \le r$ we can fix $A_i \in \mathcal{A}$, that intersects B_i , and (except a set $\mathcal{A}'_i \subset \mathcal{A}$, whose cardinality is $O(n^{t-2})$) we get that for every $A \in \mathcal{A} \setminus \mathcal{A}'_i$ we have

$$A \cap B_i \subset B_i \cap A_i$$

where $A_i \cap B_i$ contains either 0 or 1 vertex. Altogether we have that for every $A \in \mathcal{A} \setminus \bigcup_{i=1}^r \mathcal{A}_i'$

$$V(A) \cap L(G) \subset \bigcup_{i=1}^{r} (V(B_i) \cap V(A_i)).$$

By Proposition 2 there are at most $(1+o(1))\left(\frac{n}{t-1}\right)^{t-1}$ copies of a K_{t-1} in a K_t -free graph. The statement easily follows.

4 Forbidding cycles

The systematic study of counting substructures in 2k-cycle-free graphs was initiated independently in [9, 10] and [11].

4.1 Counting complete graphs

For odd cycles, we have the following interesting phenomenon, depending on whether the size of the clique is bigger than k or not.

Theorem 16. (a) If $r \le k$, then $ex(n, K_r, kC_{2l+1}) = \Theta(n^2)$.

(b) If
$$r > k+1$$
, then $ex(n, K_r, kC_{2l+1}) = O(n^{1+1/l})$.

Proof. First, let us prove (a). The lower bound follows from Theorem 12, as

$$\max_{m \le r} (ex(n, K_m, C_{2l+1})) \ge ex(n, K_2, C_{2l+1}) = \Theta(n^2).$$

The quadratic upper bound similarly follows from Theorem 12, using that for any $r \geq 2$, $ex(n, K_r, C_{2l+1}) = O(n^2)$ (we used Theorem 3 for $r \geq 3$).

Now we prove (b). Consider a kC_{2l+1} -free graph G, and its canonical partition. Then every copy of K_r consists of m vertices in R(G) and r-m vertices in L(G) for some m. If m>2, then by Theorem 3, there are $O(n^{1+1/l})$ copies of a K_m in G_R , as it is C_{2l+1} -free. Moreover, there are O(1) ways to choose the r-m vertices in L(G), so there are at most $O(n^{1+1/l})$ copies of K_r in G such that m>2. Note that the case $m\leq 1$ only gives linearly many copies of K_r . Thus we only have to deal with the case m=2. In other words, it only remains to show that the number of copies of K_r in G which contain exactly two vertices from R(G) is $O(n^{1+1/l})$.

Let G'_R be the subgraph of G_R consisting of only those edges $xy \in E(G_R)$ such that x, y and some r-2 vertices from L(G) form a copy of K_r in G. Clearly, the number of copies of K_r in G which contain exactly two vertices from R(G) is $O(E(G'_R))$ because each edge of G'_R can be extended to such a copy of K_r in at most $\binom{|L|}{r-2} = O(1)$ ways. Since an edge xy of G'_R appears in a copy of K_r that contains r-2 vertices in L(G), x and y have at least $r-2 \ge k$ common neighbors in L(G). If G'_R contains k vertex-disjoint copies of C_{2l} , then we pick an arbitrary edge from each of these k copies. For these k edges e_1, \ldots, e_k we can greedily pick k vertices v_1, \ldots, v_k in L(G) such that v_i is adjacent to both endpoints of $e_i = x_i y_i$ for every i. Then we replace e_i with $v_i x_i$ and $v_i y_i$, and this way we get k vertex disjoint copies of C_{2l+1} in G, a contradiction. Thus G'_R does not contain k vertex-disjoint copies of C_{2l} , hence

$$|E(G'_R)| \le ex(n, kC_{2l}) = ex(n, C_{2l}) + O(n),$$

by Theorem 1. A well-known theorem of Bondy and Simonovits [3] states that $ex(n, C_{2l}) = O(n^{1+1/l})$, so

$$|E(G'_R)| \le ex(n, kC_{2l}) = O(n^{1+1/l}).$$

Therefore, the number copies of K_r in G which contain exactly two vertices from R(G) is also $O(n^{1+1/l})$, as desired.

For even cycles, we have the following.

Proposition 17. For any $r \geq 2$, $l \geq 2$, we have $ex(n, K_r, kC_{2l}) = O(n^{1+1/l})$.

Proof. Using Theorem 5, we get $ex(n, K_r, kC_{2l}) = O(\overline{ex}(n, K_r, C_{2l})) = O(\sum_{i=1}^r ex(n, K_i, C_{2l}))$, as any induced subgraph of K_r is also a complete graph. By Theorem 3, we have $ex(n, K_i, C_{2l}) = O(n^{1+1/l})$ for any $i \geq 1$, so $ex(n, K_r, kC_{2l}) = O(n^{1+1/l})$, as required.

5 Forbidding bipartite graphs

Let $K_{a,b}$ denote a complete bipartite graph with color classes of sizes a and b with $a \leq b$. Alon and Shikhelman proved the following.

Proposition 18 ([1], Proposition 4.10). If $s \le t$ and $a \le b < s$ then $ex(n, K_{a,b}, K_{s,t}) = O(n^{a+b-ab/s})$.

We will prove that the same upper bound holds for $ex(n, K_{a,b}, kK_{s,t})$. Note that they also gave a constant factor in their proof; our proof would give a worse constant factor for the case k > 1. They also showed that in some range of a and b this order of magnitude is sharp, this immediately implies the same for the case k > 1.

Proposition 19. If $s \le t$ and $a \le b < s$ then $ex(n, K_{a,b}, kK_{s,t}) = O(n^{a+b-ab/s})$.

Proof. Let G be a $kK_{s,t}$ -free graph and consider its canonical partition. A copy of $K_{a,b}$ intersects G_R in a copy of $K_{a',b'}$ for some $0 \le a' \le a$ and $0 \le b' \le b$. For fixed a' and b', there are $O(n^{a'+b'-a'b'/s})$ such copies by Proposition 18 since G_R is $K_{s,t}$ -free. Increasing a' by 1, increases a' + b' - a'b'/s by 1 - b'/s which is positive since b' < s. Applying a similar argument for a', we get that a' + b' - a'b'/s is maximized when a' = a and b' = b. Therefore, $O(n^{a'+b'-a'b'/s}) \le O(n^{a+b-ab/s})$.

Moreover, the number of ways to extend a copy of $K_{a',b'}$ to a copy of $K_{a,b}$ by adding vertices from L(G) is at most a constant (where this constant depends on a, b, s, t, but not on n). Now as $0 \le a' \le a$ and $0 \le b' \le b$, there are only at most (a+1)(b+1) ways to pick a' and b', finishing the proof.

Note that by Theorem 5 and Remark 4 we have the following general lower bound: $ex(n, K_{a,b}, kK_{s,t}) \ge \Omega(n^{\alpha(K_{a,b})}) = \Omega(n^b)$ if $k \ge a+b$ and $a \le b$. If b > s, then a+b-ab/s < b, so the upper bound in the above proposition cannot hold in this case. Instead, we have the following.

Proposition 20. If $a \le b$, $b \ge s$, $s \le t$, then $ex(n, K_{a,b}, kK_{s,t}) = O(n^b)$. Moreover, if k > a, then we have $ex(n, K_{a,b}, kK_{s,t}) = \Theta(n^b)$.

Proof. Let G be a $kK_{s,t}$ -free graph and let us consider its canonical partition. A copy of $K_{a,b}$ intersects G_R in a copy of $K_{a',b'}$ for some $0 \le a' \le a$ and $0 \le b' \le b$ with $a' \le b'$. Let us fix a' and b' and consider two cases.

If b' < s, then by Proposition 18 there are $O(n^{a'+b'-a'b'/s})$ copies of $K_{a',b'}$ in G_R as it is $K_{s,t}$ -free. By the same argument as in the proof of Proposition 19, it is easy to see that a'+b'-a'b'/s increases, if we increase a' as long as b' < s. So $a'+b'-a'b'/s < s+b'-sb'/s = <math>s \le b$. Thus there are at most $O(n^b)$ copies of $K_{a',b'}$ in G_R .

If $b' \geq s$, then we first claim that there are at most $O(n^b)$ copies of $K_{a',b'}$ in G_R . Indeed, there are at most $O(n^{b'})$ ways to pick b' vertices, and they have at most t-1 common neighbors (otherwise we can find a $K_{s,t}$ in G_R). Thus there are at most $\binom{t-1}{a'} = O(1)$ ways to pick the other class of $K_{a',b'}$, so there are at most $O(n^{b'}) \leq O(n^b)$ copies of $K_{a',b'}$ in G_R again.

Each copy of $K_{a',b'}$ in G_R can be extended to a copy of $K_{a,b}$ in G by adding vertices from L(G) in O(1) ways. Thus for any fixed a' and b', there are at most $O(n^b)$ copies of $K_{a,b}$ in G and as there are only at most (a+1)(b+1) choices for a' and b', the proof is complete.

For the moreover part, take a $K_{s,t}$ -free graph G on n-k+1 vertices and consider $K_{k-1}+G$. It is clearly $kK_{s,t}$ -free. Any set of a vertices from K_{k-1} together with any set of b vertices from G forms a copy of $K_{a,b}$, which finishes the proof.

Let us focus now on the case k = 1 and $b \ge s$, as this case was not examined in [1].

Proposition 21. (a) If $s \le a \le b \le t$, then $ex(n, K_{a,b}, K_{s,t}) = O(n^s)$.

(b) If
$$a < s \le b \le t$$
, then $ex(n, K_{a,b}, K_{s,t}) = \Theta(n^b)$.

Proof. For (a) let us consider a $K_{s,t}$ -free graph G and an arbitrary set S of s vertices. We claim that S is contained in the larger color class (i.e., with size b) of at most O(1) copies of $K_{a,b}$. Indeed, the vertices of S have at most t-1 common neighbors, thus there are $\binom{t-1}{a}$ ways to pick the other side, and they have at most t-1 common neighbors, thus there are at most $\binom{t-1-s}{b-s} = O(1)$ copies of $K_{a,b}$ such that their larger color class contains S. So there are at most $O(n^s)$ copies of $K_{a,b}$, as desired.

For (b) let us consider the graph $K_{s-1,n-s+1}$. It is easy to see that it contains $\Omega(n^b)$ copies of $K_{a,b}$. The upper bound $O(n^b)$ follows from Proposition 20.

6 Concluding remarks and open problems

• Our conclusion is that the significant difference between ex(n, H, F) and ex(n, H, kF) mostly comes from the ability to count subgraphs of H due to the universal vertices (vertices of degree n-1). A particular example when this is the case is ex(n, F, kF). We did not deal with this especially interesting case, but for complete graphs Theorem 15 implies

$$ex(n, K_t, kK_t) = (k - 1 + o(1)) \left(\frac{n}{t - 1}\right)^{t - 1}.$$

 \bullet Another natural direction is to count lH instead of H. Here we present two results of this type.

Proposition 22.

$$ex(n, lK_2, K_3) = \frac{1}{l!} \left\lfloor \frac{n^2}{4} \right\rfloor \left\lfloor \frac{(n-2)^2}{4} \right\rfloor \dots \left\lfloor \frac{(n-2l+2)^2}{4} \right\rfloor.$$

Proof. The lower bound is given by the complete bipartite graph $K_{\lfloor n/2\rfloor,\lceil n/2\rceil}$.

Let G be an triangle free graph on n vertices. To prove the upper bound, we first select an edge e_1 from G and then we select an edge e_2 disjoint from e_1 , and then an edge e_3 disjoint from both e_1 and e_2 and so on. By Mantel's theorem, the maximum number of edges in a triangle-free graph is at most $\lfloor n^2/4 \rfloor$ so we can pick $e_1 = u_1v_1$ in at most $\lfloor n^2/4 \rfloor$ ways. Since the subgraph of G induced by $V(G) \setminus \{u_1, v_1\}$ is also triangle-free, we can pick e_2 in at most $\lfloor (n-2)^2/4 \rfloor$ ways and e_3 in at most $\lfloor (n-4)^2/4 \rfloor$ ways (by Mantel's theorem again) and so on, giving a total of $\lfloor n^2/4 \rfloor \lfloor (n-2)^2/4 \rfloor \dots \lfloor (n-2l+2)^2/4 \rfloor$ ordered tuples of l independent edges (e_1, e_2, \dots, e_l) . Since each copy of lK_2 is counted l! times, this implies the desired upper bound.

Using Proposition 22, we prove the following asymptotic result.

Theorem 23. Let l < k. Then

$$ex(n, lK_3, kK_3) = (1 + o(1)) {k-1 \choose l} \left(\frac{n^2}{4}\right)^l.$$

Proof. The lower bound is given by the graph $K_{k-1} + K_{\lfloor (n-k+1)/2 \rfloor, \lceil (n-k+1)/2 \rceil}$.

For the upper bound, let G be a kK_3 -free graph and consider its canonical partition. We say that a triangle is *good* if it has exactly one vertex in L(G).

We claim that it suffices to only count those copies of lK_3 in which every triangle is good. Indeed, no triangle of lK_3 has three vertices in R(G), and if any of them has two vertices in L(G), then we can pick at least l+1 vertices from L(G) in O(1) ways, and at most 2l-1 vertices from R(G) in $O(n^{2l-1}) = o(n^{2l})$ ways, which is covered by the error term in the theorem. So from now on we count the number of copies of lK_3 in which every triangle is good; we will refer to such a copy of lK_3 as a good copy.

We know that the subgraph of G induced by L(G) consists of (maximum possible number of) vertex-disjoint triangles A_1, \ldots, A_r for some $r \leq k-1$. Let a_i, b_i, c_i be the vertices of A_i for each $1 \leq i \leq r$.

A good copy of lK_3 can contain only one of the vertices a_i, b_i, c_i for any $1 \leq i \leq r$, because otherwise we can find more than r vertex-disjoint triangles in G, a contradiction. So in order to count the number of good copies of lK_3 in G, we first pick l of the $r \leq k-1$ triangles -say A_1, A_2, \ldots, A_l without loss of generality - from G_L in $\binom{r}{l} \leq \binom{k-1}{l}$ ways, and then count the number of good copies of lK_3 in which every triangle has a vertex in one of the triangles A_1, A_2, \ldots, A_l . Now we bound this latter number.

First let us assume that there are two good copies of lK_3 in G which use two different vertices of the triangle $A_i = a_i b_i c_i$, say a_i and b_i for some $1 \le i \le l$. Let the corresponding good triangles of these lK_3 's be $a_i xy$ and $b_i pq$. Then the edges xy and pq must share a vertex, because otherwise we can replace A_i with the triangles $a_i xy$ and $b_i pq$ to produce more than r vertex-disjoint triangles in G, a contradiction. Thus it is easy to see that number of good triangles containing a_i is at most 2n since there are at most 2n edges that share a vertex with pq. Similarly, the number of good triangles containing b_i or c_i is also at most 2n each. Therefore, the total number of good triangles which have a vertex in A_i is at most 6n = O(n) in this case. This implies that the number of good copies of lK_3 in which every triangle has a vertex in one of the triangles A_1, A_2, \ldots, A_l is at most $O(n) \cdot O(n^{2l-2}) = o(n^{2l})$, which is covered by the error term of the theorem again.

So we can assume that every such good copy of lK_3 contains only one of the vertices a_i, b_i, c_i for each $1 \leq i \leq l$, say u_1, u_2, \ldots, u_l . Thus we can count the number of those copies by picking a copy of lK_2 from G_R in at most $ex(n, lK_2, K_3) = (1 + o(1))\frac{1}{l!}\left(\frac{n^2}{4}\right)$ ways by Proposition 22 (recall that G_R is triangle-free), and then pairing the l edges of lK_2 with u_1, u_2, \ldots, u_l in at most l! ways.

Therefore, the total number of good copies of lK_3 is at most $(1 + o(1))\binom{k-1}{l}\left(\frac{n^2}{4}\right)^l$, as required.

- We mention some more specific open problems.
- \circ A lower bound of $\Omega(n^s)$ is trivial in Proposition 21 for s=1. However it would be appealing to prove it for all s or even in case s=2.
 - It would be also interesting to improve Theorem 13 and prove an asymptotic result.
- In this article our results mostly obtain the order of magnitude or asymptotics of various quantities. It would be interesting to prove exact results corresponding to them.
- Finally, let us mention that the Turán number of the disjoint union of graphs $F_1, F_2, ..., F_k$, has not been investigated when the F_i 's can be different. (See Theorem 1 and the comment after it, for the case when all the F_i 's are the same.) It is not hard to prove the following proposition. However, it would be interesting to prove a sharper result in this case.

Proposition 24. Let us suppose that we have graphs $F_1, ..., F_k$ and let $F = \bigcup_{1 \le i \le k} F_i$. Then we have

$$ex(n, F) = \max\{ex(n, F_i) : i \le l\} + O(n).$$

Proof of Proposition 24. Let j be an integer such that $ex(n, F_j) = \max\{ex(n, F_i) : i \leq l\}$. Then the lower bound follows by taking an F_j -free graph with maximum possible number of edges.

For the upper bound, consider an F-free graph G. Let F' be a subgraph of G consisting of vertex disjoint copies of F_1, \ldots, F_j where j is an integer which is chosen as large as possible. Clearly j < l as G is F-free. Then, of course, the subgraph of G induced by $V(G) \setminus V(F')$ is F_{j+1} -free, so it contains at most $ex(n, F_{j+1}) \leq \max\{ex(n, F_i) : i \leq l\}$ edges. Moreover, there are at most O(n) edges incident to the vertices of F'. Adding these bounds up, the proof is complete. \square

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