# Notes on the <br> SWTPC MP-N Calculator Interface and the <br> Calc-1 Program 

prepared for.

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Calc-1 Program

I assembled an SWTPC MP-N Calculator Interface and implemented it using the Calc-1 Program supplied with the kit. The following are my corrections to the documentation and my observations about the interface board and software.

This interface was bought to perform floating-point arithmetic and for its function capabilities such as SIN, COS, and $e^{X}$. My application required an integer truncation function that is not performed by this calculator, so $I$ wrote a small assembly language subroutine to do it. A potentially irritating problem is that the calculator chip does not automatically convert to scientific notation if the numbers become too big to display in floating point. The control program must keep track of the display mode.

My application requires fast numerical processing, so, after observing the calculator's computational speed using the Calc-1 program, I decided that its trigonometric functions were too slow. I developed another approach to the problem using pure assembly language rather than partially relying on peripheral hardware processing.

Though I found no major flaws with the Calc-1 software, I did find some mistakes in its documentation. In Table V, the ASCII to Calculator Instruction Lookup Table, there are three errors: the hex value in the table where MSB is 0 and LSB is $D$ should be 21 , where MSB is 1 and LSB is 8 should be $2 F$, and where $M S B$ is 2 and $L S B$ is $E$, the hex value should be OA. I added a one page summary of the Calc-1 instruction set to the documentation because the supplied documentation is somewhat confusing.

The Calc-1 program uses a part of RAM also used by the monitor. This provides problems when Calc-1 is halted to record it on cassette tape and resume after it has been record. This can be alleviated by changing the "LDX PARADR" instructions in lines 250, 850, 1660, and 1980 to "LDX \#\$800c," or, in hex, "FE A002" to "CE 800C." However, the interface must now always be plugged into $1 / 0$ port 3 . The data at hex address 028A is output to the terminal to clear have the screen. The ADM terminals we use a different character to clear the screen, so I changed the data at that address to hex 1A. Calc-1 also uses address modification within itself so that, as a result, it cannot be implemented in ROM.

I tested the MP-N interface with the Calc-1 program by making various operand and operator entries and checking the results against the results I obtained
with my own pocket scientific calculator. I noticed no major discrepancies between the two sets of results.

## SW T PC

## MP -N Calculator Interface

## Documentation



## Assembly Instructions $\mathrm{MP}-\mathrm{N}$ Calculator Interface Kit

The Southwest Technical Products MP-N Calculator Interface interfaces the SWTPC 6800 Computer System thru a peripheral interface Adaptor (PIA) to the National Semiconductor MM57109 Number Oriented Processor. This "processor" is a Reverse Polish Dotation (RPN) calculator chip without the internal keypad interfacing circuitry which has made interfacing to calculator chips so difficult in the past. This chip allows data and instruction entry in conventional binary form and speeds entry with the elimination of the debounce circuitry built into conventional calculator chips. It is called a processor because it has instructions and control lines which allow it to operate in conjunction with ROM and RAM as a stand alone numerical processor. It may however be operated as a computer peripheral for numerical calculation and this is the configuration in which the chip has been implemented.

All interfacing from the 6800 Computer System to the calculator chip has been done thru a 6820 PIA. Both the PIA and calculator chip reside on a $3 \frac{1}{2}$ " $\mathrm{X} 5 \frac{1}{4}$ double sided, plated thru hole circuit board plugged onto one of the seven available interface card positions on the mother board of the 6800 Computer. All data and instructions fed to and all results received from the calculator chip are handled by your own assembler or machine language program. The calculator features reverse Polish notation, floating point or scientific notation, up to an eight digit mantissa and two digit exponent, trig functions, base 10 and natural logarithms, and overflow indicator.

## PC Board Assembly

NOTE: Since all of the holes on the PC board have been plated thru, it is only necessary to solder the components from the bottom side of the board. The plating provides the electrical connection from the "BOTTOM" to the "TOP" foil of each hole. Unless otherwise noted it is important that none of the connections be soldered until all of the components of each group have been installed on the board. This makes it much easter to interchange components if a mistake is made during assembly. Be sure to use a low wattage iron (not a gun) with a small tip. Do not use acld core solder or any type of paste flux. We will not guarantee or repair any kit on which either product has been used. Use only the solder supplied with the kit or a $60 / 40$ alloy resin core equivalent. Remember all of the connections are soldered on the bottom side of the board only. The plated-thru holes provide tine electrical connection to the top foil.
( ) Before installing any parts on the circuit board, check both sides of the board over carefully for incomplete etching and foil "bridges" or "breaks". It is unlikely that you will find any, but should there be one, especially on the "TOP" side of the board, it will be very hard to locate and correct after all of the components have been installed on the board.
( ) Starting from one end of the circuit board install each of the three, 10 pin Molex female edge connectors along the lower edge of board. These connectors must be inserted from the "TOP" side of the board and must be pressed down firmly against the circuit board, so that each pin extends completely into the holes on the circuit board. Not being careful here will cause the board to either wobble and/or be crooked when plugging it onto the mother board. It is
either wobble and/or be crooked when plugging it onto the mother board. It is suggested that you solder only the two end pins of each of the three connectors until all have been installed at which time if everything looks straight and rigid you should solder the as yet unsoldered pins.
( ) Insert the small nylon indexing plug into the lower edge connector pin. indicated by the small triangular arrow on the "BOTTOM" side of the circuit board. This prevents the board from being accidently plugged on incorrectly.
( ) Attach all of the resistors to the board. As with all other components unless noted, use the parts list and component layout drawing to locate each part and install from the "TOP" side of the board bending the leads along the "BOTTOM" side of the board and trimming so that $1 / 16$ " to $1 / 8$ " of wire remains. Solder.
( ) Install the capacitors on the circuit board. Be sure to orient electrolytic capacitor $C 4$ so its polarity matches with that shown on the component layout drawing. Solder.
( ) Install the transistor and diode. These components must be oriented to match the component layout drawing. Solder.
( ) Install integrated circuit IC2 on the circuit board. This component must be oriented so its metal face is facing the circuit board and is secured to the circuit board with a $\# 4-40 \times 1 / 4^{\prime \prime}$ screw, lockwasher and nut. A heatsink is not used. The three leads of the integrated circuit must be bent down into each of their respective holds. Solder.

NOTE: MOS integrated circuits are susceptible to damage by static electricity. Although some degree of protection is provided internally within the in:egrated circuits, their cost demands the utmost in care. Before opening and/or installing any MOS integrated circuits you should ground your body and all metallic tools coming into contact with the leads, thru a 1 M ohm $1 / 4$ watt resistor (supplied with the kit). The ground must be an "earth" ground such as a water pipe, and not the circuit board ground. As for the connection to your body, attach a clip lead to your watch or metal ID bracelet. Make absolutely sure you have the 1 Meg ohm resistor connected between you and the "earth" ground, otherwise you will be creating a dangerous shock hazard. Avoid touching the leads of the integrated circuits any more than necessary when installing them, even if you are grounded. On those MOS IC's being soldered in place, the tip of the soldering iron should be grounded as well(separately from your body ground) either with or without a 1 Meg ohm resistor. Most soldering irons having a three prong line cord plug already have a grounded tip. Static electricity should be an important consideration in cold, dry environments. It is less of a problem when it is warm and humid.
( ) Install MOS integrated circuits IC1, IC3, IC4 and IC5 following the precautions given in the preceding section. As they are installed, make sure they are down firmly against the board before soldering all of their leads. Do not bend the leads on the back side of the board. Doing so makes it very difficult to remove the integrated circuit should replacement ever be necessary. The "dot" or "notch" on the end of the package is used for orientation purposes and must match with that shown on the component layout drawing for the IC. Solder.
() Working from the "TOP" side of the circuit board, fill in all of the feedthru's with molten solder. The feed-thru's are those unused holes on the
board whose internal plating connects the "TOP" and "BOTTOM" circuit connections. Filling these feed-thru's with molten solder guarantees the integrity of the connections and increases the current handling capability.
( ) Now that all of the components have been installed on the board, double check to make sure all have been installed correctly in their proper location.
( ) Check very carefully to make sure that all connections have been soldered. It is very easy to miss some connections when soldering which can really cause some hard to find problems later during checkout. Also look for solder "bridges" and "cold" solder joints which are another common problem.

Since the MP-N circuit board now contains MOS devices, it is susceptible to damage from severe static electrical sources. One should avoid handling the board any more than necessary and when you must, avoid touching or allowing anything to come into contact with any of the conductors on the board.

## Using the Calculator Interface

Table I gives a complete list and descriotion of the calculator chip's instruction set. Remember that some of the instructions are for stand alone processing systems and are not used on this interface. All numerical entry is in Reverse Polish Notation (RPN) and anyone familiar with Hewlett Packard calculators should have no problem with the data entry sequence. For those not familiar with RPN, the following should be helpful:

To add $7+8$, enter the following
7 enter $8+$ (4 entries)
The answer is now stored in the $X$ accumulator within the calculator chip
The OUT instruction may be used to output the answer

To find the inverse sine of 0.5 , enter the following:
0.5 INV SIN ( 5 entries)

The answer is now stored in the $X$ accumulator within the calculator chip. The OUT instruction may be used to output the answer.

In order to simplify the interfacing between your program and the calculator interface, you will probably want to incorporate the following subroutines into your program.

## INITAL SUBROUTINE

The INITAL or inftialize subroutine configures the PTA interfacing to the calculator chip. This subroutine need only be used once; and is best placed somewhere at the beginning of your program. It is responsible for initializing the data direction registers and control registers of the PIA. The subroutine requires that the index register be loaded with the "lowest" address of the PIA interfacing to the calculator chip prior to execution.

This "lowest" address depends upon which interface port position the MP-N calculator card is plugged. The table below give the "lowest" address of each interface card position.

| PORT | $8 \emptyset \emptyset \emptyset$ |
| :--- | :--- |
| PORT1 | $8 \emptyset 04$ |
| PORT2 | $8 \emptyset \emptyset 8$ |
| PORT3 | $8 \emptyset \emptyset \mathrm{C}$ |
| PORT4 | $8 \emptyset 1 \emptyset$ |
| PORT5 | $8 \emptyset 14$ |
| PORT6 | $8 \emptyset 18$ |
| PORT7 | $8 \emptyset 1 C$ |


| 8 EF | INITAL LIAA A | \＃+ TF | INIT A SIEIE BF FIA |
| :---: | :---: | :---: | :---: |
| A7 00 | ETA A | O，$X$ |  |
|  | LIAA A |  |  |
| A7 01 | ETA A | 1，X |  |
| E\％OO | LIIA A | \＃$=00$ | INIT E EIEE OIF FIA |
| A7 02 | $\operatorname{sTA} A$ | $2, \dot{x}$ |  |
| \％ 34 | LIIA A |  | NE，F\％ |
| A7 03 | ETA A | 三i |  |
| $A 602$ | LIIA A | Z，$\times$ | MEEAF $=$ W F－$\because$ |
| $\because$ | FTS |  |  |

## OUTINS SUBROUTINE

The OUTINS or out instruction subroutine is used to get program data and instructions into the calculator．To send a digit or instruction to the calculator chip，use Table II to find the OP code of the instruction you wish to send．Load this OP code into the $A$ accumulator and jump or branch to the OUTINS subroutine． If you have a string of data you wish to send，just recycle thru this subroutine as many times as necessary．The subroutine takes care of all of the READY and HOLD signals to the calculator chip so there is no worry of sending data faster than the calculator chip can accept it．The subrputine destroys the contents of the $B$ accumulator during execution while the contents of the $A$ accumulator and index register are not destroyed．

| $E 601$ | GUTINE | LIIA | E | 1，X | W⿵冂T Figif SEALY |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2 AFC |  | E：F＇L |  | GITINE |  |
| A7 00 |  | STA | A | O，$x$ | FGINHARLI INETFUGTIGN TG |
| E6 00 |  | L［A | E： | $0, X$ | ELEAF FLAIG EIT |
| $E 6$ |  | LIA | E |  |  |
| E701 |  | STA | Ei | 1，X | EFING HHILE LINE LIW |
| Et 01 | WAIT10 | LIA | E | 1，X |  |
| 2 AFC |  | EFFL |  | WAITIG | LDGF：FGF FEALY LOW |
| E6 00 |  | LIA | E | G， X | ELEAF FLAGE EIT |
| $0 \% 36$ |  | LIIA | $E$ | \＃ ¢ $_{\text {S }}$ | HIEH HILLI－FGE FEALY |
| E7 01 |  | STA | E | $1, x$ | FETHFN HIDLI LINE HISH |
| 39 |  | FTE |  |  |  |

## SETMEM SUBROUTINE

The SETMEM or set memory subroutine initializes the memory locations to which the calculator's output data will be stored. This subroutine must be executed immediately before the OUTANS subroutine is used. Although it can be changed, memory locations $\emptyset \emptyset 2 \emptyset$ thru $\emptyset \emptyset 2 B$ have been designated the temporary storage locations for the calculator's output data. The subroutine sets memroy location $\varnothing \varnothing 2 \emptyset$ to a $\emptyset \emptyset$ while locations 21 thru $2 B$ are set to 20 (ASCII spaces). This subroutine destroys the contents of the index register and $B$ accumulator. The contents of the $A$ accumulator are not destroyed.


## OUTANS SUBROUTINE

The OUTANS or output answer subroutine outputs the contents of the $X$ register within the calculator chip in $B C D$ to memory locations 0020 thru 002B. Since the mantissa digit count of the calculator is variable, the previous SETMEM subroutine blanks out any digit location not filled by the OUTANS subroutine. It is very important that the SETMEM subroutine be used each time before executing the OUTANS subroutine. The OUTANS subroutine outputs data in two different formats depending upon whether the calculator chip is in the floating point or scientific mode. The calculator initially starts out in the floating point mode where it will remain until changed by the TOGM (2216) instruction. This calculator does not automatically convert to scientific notation if the numbers become too big to handle in floating point as many do. An MCLR ( $2 \mathrm{~F}_{16}$ ) instruction will always reset the calculator chip to the floating point mode regardless of what mode it was in originally. Since the calculator chip does not tell you what mode it is in when it is outputting data, your program must know so you can process the data accordingly. Table IV shows the format in whicn the data is stored. At the end of the OUTANS subroutine, the $N$ bit of the condition code register is set if an error has transpired since the last execution of the OUTANS subroutine. You may use a BMI instruction to catch and branch to an error routine to note the error. You should then send an ECLR ( $2 \mathrm{~B}_{16}$ ) instruction to the calculator chip to reset the calculator chip's error flag. Disregarding the error flag on the calculator chip will cause no problems. The chip will continue to function regardless of the state of the flag. The subroutine requires that the index register be loaded with the "lowest" address of the PIA interfacing to the calculator chip prior to execution. Since the SETMEM subroutine usually run prior to this destroys the contents of the index register, don't forget to reload the index register before branching to the OUTANS subroutine. The OUTANS subroutine destroys the contents of both the $A$ and $B$ accumulators during execution while the contents of the index register is not changed.

| $E 6$ | 01 | gllans | LILA | E | ：${ }^{\text {x }}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2A | $F:$ |  | E：F＇L |  | EITANE |  |
| A6 | 00 |  | LEA | A | O，$X$ | İEAn FLAG EIT |
| $E$ | 16 |  | LITA | A | \＃5 | SENET AN Bllut |
| AT | O） |  | STA | A | O，$X$ |  |
| に | EE |  | LIA ${ }^{\text {a }}$ | E | \＃ \％$_{\text {E }} \mathrm{E}$ | LOW HigllimFIE FiEALI＇ |
| E7 | 01 |  | STA | $E$ | i，$x$ | EFING HDLCI LINE LIIW |
| Et． | 01 | WAITEO | LIA | E | $i$ i，$X$ | WAIT FIF EECGREI FEACM |
| 2A | $F C$ |  | EF＇L |  | WHITS |  |
| EG | 00 |  | LIIA | Ei | $\theta, x$ | GLEAF FLAT EIT |
| 86 | OF |  | LIIA | A | \＃ $\mathrm{O}_{\mathrm{F}}$ |  |
| A7 | 00 |  | ETA | A | ］，$X$ | EENLi A NOF |
| E6 | 03 | WAITE | Llia | $E$ | $B_{1} X$ | LiGit FCF |
| 2E： | 06 |  | EMI |  | Elliticiou | TFANEFEA GALE LIATA TII METINE |
| E6． | 01 |  | LIA | E | 1，${ }^{\text {a }}$ |  |
| zE | 16 |  | EMI |  | GİNFL | FFINT MEMOFY GMTENTS |
| 20 | F6． |  | EFA |  | WAIT |  |
| Ac | 02 | OUTEIG | LIA | $A$ | $2 \cdot X$ | LIAL Gilit mata inti a |
| 16 |  |  | TAE： |  |  |  |
| 84 | OF |  | ANCI | $A$ | \＃ 50 OF | ELIMANATE LFFEF 4 EITS |
| EA | 30 |  | GFA | $A$ | \＃$\$^{3} \mathrm{O}$ | CDOVEAT TG AGIII LGTH |
| 54 |  |  | LSF＇ | E |  |  |
| 54 |  |  | LEF | E： |  |  |
| 54 |  |  | LGF＇ | $E$ |  |  |
| 54 |  |  | L马Fi | E： |  |  |
| $G A$ | 20 |  | GRA | E； | \＃5 ${ }^{\circ} \mathrm{C}$ | INEFEMENT ALIFESEES EY \＄ 20 |
| F7 | 0166 |  | STA | E： | FGNTET |  |
| 97 | 00 | FGINTZ | STA | $A$ | まG |  |
| 20 | E2 |  | EFIA |  | WAITE |  |
| 86 | 36 | CONFLG | LIA | $A$ | 牛家家 |  |
| A7 | 01 |  | ETA | $A$ | 1，$x$ | EFINA Cigl Li LINE HIEH |
| $A E$ | 00 |  | LELA | ${ }_{4}$ | $\because x$ | CLEAA |
| 39 |  |  | FTE |  |  |  |

## Number Entry Rules

When a digit, decimal point, or $\pi$ is entered with an $0-9$, DP , or PI instruction, the stack is first pushed and the X register cleared: $\mathrm{Z} \rightarrow \mathrm{T}$, $Y \rightarrow Z, X \rightarrow Y, 0 \rightarrow X$. This process is referred to as "initiation of number entry." Following this, the digit and future digits are entered into the $X$ mantissa. Subsequent entry of digits or DP, EE, or CS instructions do not cause initiation of number entry. Digits following the eighth mantissa digit are ignored. This number entry mode is terminated by any instiruciton except 0-9, DP, EE, CS, PI, or HALT. Termination of number entry means two things. First, the number is normalized by adjusting the exponent and decimal point position so that the decimal point is to the right of the first mantissa digit. Second, the next digit, decimal point, or $\pi$ entered will cause initiation of number entry, as already described. There is one exception to the number entry initiation rule. The stack is not pushed if the instruction prior to the entered digit was an ENTER. However, the $X$ register is still cleared and the entered digit put in $X$.

The ENTER key itself terminates number entry and pushes the stack. The OUT instruction terminates number entry and prepares the stack for pushing upon the next entry of data. This means that if you use the ENTER and OUT instrucitons consecutively, the stack gets pushed twice which is not what you want. If you wish to ENTER data and immediately out the result, use only the OUT instruction. The OUT performs the entry. If you do not wish to OUT the ENTER'ed data, fust use the ENTER instruction by itself.

The AIN and IN instructions should not be used for number entry. Provisions have not been made for their use on this interface.

## How It Works

Peripheral Interface Adaptor (PIA) ICl interfaces the MM57109 calculator chip, IC3, to the SWTPC 6800 buss. The first six bits of the A side of the PIA are used to feed instructions to the calculator chip while the eighth is used as an input to monitor the ERROR output of the calculator. Control line CAl outputs HOLD signals to, while control line CA2 inputs READY signals from the calculator chip. The first four bits of the $B$ side of the PIA are used to input BCD digit data while the last four bits input digit addresses. The CBl line tnputs READ/WRITE signals while the CB2 control line is not used. Hex inverter/buffer, IC4, is used primarily as the 320 to 400 Khz single phase oscillator required by the calculator chip. One section is used to invert the HOLD signal going to the calculator. Shift register IC5 generates the $P O R$ signal required for proper startup and initialization. +5 VDC power required by the board is supplied by voltage regulator IC2 while -4 VDC voltage is supplied by transistor $Q 1$ and its associated components. Figure I shows a block diagram for the internal construction of the calculator chip.

## Resistors

| $\cdots \mathrm{Rl}$ |  |
| :---: | :---: |
|  | $\checkmark$ R2 |
|  | $\checkmark$ R3 |
|  | $\checkmark$ R4 |
|  | $\checkmark$ R5 |
|  | $\checkmark$ R6 |
|  | $\checkmark$ R7 |
|  | $\checkmark$ R8 |
|  | $\checkmark \mathrm{R} 9$ |
|  | $\checkmark$ R10 |
|  | ) R11 |
|  | $\checkmark$ R12 |
|  | $\checkmark$ R13 |
|  | $\checkmark$ R14 |
|  | $\checkmark \mathrm{R} 15$ |
|  | $\checkmark$ R16 |
|  | $\checkmark$ R17 |


| 47 K |  |  |  |
| :---: | :---: | :---: | :---: |
| 1 K | 11 | 11 | 11 |





## CALC- 1 Program

In order to see how the calculator chip is used and how to incorporate these subroutines into a program, the CALC-1 program listing is given. CALC-l allows the operator to use the calculator chip just as you would a standard RPN desk calculator with the same features. All communication to the chip is done thru the terminal's keyboard with all results displayed on the terminal's display. Since the terminal's keyboard just has standard ASCII characters rather than the labeling found on calculator keys; selected ASCII characters have been substituted for normal calculator function keys. It is the job of the CALC-1 program to accept all data and instruction commands from the terminal's keyboard, send them to the calculator chip and display all results on the terminal's display. The program resides from memory locations 0020 thru $02 C 0$ which is approximately 700 bytes of code. Since most of the lower 256 bytes are used for the ASCII character lookup table and some of the upper is used for terminal interfacing, you should be able to incorporate the package into your program using somewhat less memory than was used here.

The program starts at line 50 by storing the ASCII lookup table from memory locations $\emptyset \emptyset 8 \emptyset$ thru $\emptyset \emptyset F F$. This table covers the entire 128 character ASCII set. Whenever an ASCII character is received from the keyboard it is OR'ed with 80 , and the resulting address contains the selected command or instruction for the calculator chip. Line 210 ORG's the program at memroy location $\emptyset \emptyset \emptyset$ where the terminal's screen is cleared and titled. Line 250 loads the index register extended with the contents of memory locations $A \emptyset \emptyset 2$ and $A \emptyset \emptyset 3$ with $8 \emptyset \emptyset C$, the starting address of Port 3 . If you wish to plug the calculator board onto an $I / 0$ port other than PORT 3 . Use the table below to find the address to be loaded into memory locations $A \emptyset \emptyset 2$ and AøØ3 prior to executing the program.

| PORT0 | $8 \emptyset \emptyset \emptyset$ |
| :--- | :--- |
| PORT1 | $8 \emptyset \emptyset 4$ (Serial control interface only) |
| PORT2 | $8 \emptyset \emptyset 8$ |
| PORT3 | $8 \emptyset \emptyset C$ |
| PORT4 | $8 \emptyset 1 \emptyset$ |
| YORT5 | $8 \emptyset 14$ |
| PORT6 | $8 \emptyset 18$ |
| PORT7 | $8 \emptyset 1 C$ |

Lines 280 thru 370 contain the INITAL subroutine described in detail earlier. lines 380 thru 410 accept entered keyboard commands, lookup the selected calculator instructions and deposit the data or instruction in the A accumulator. Lines 440 thru 550 contain the OUTINS subroutine described in detail earlier. Lines 550 thru 740 check to see what instruction or data has been entered so the result may be output if appropiate. Line 710 looks for the TOGM instruction so the program knows which display mode to use when outputting data. Lines 770 thru 840 contain the SETMEM subroutine described in detail earlier. Since the SETMEM subroutine destroys the contents of the index register, line 850 reloads it before proceeding to the OUTANS subroutine contained in lines 880 thru 1200. Line 1210 checks to see of the ERROR flag was set during the last output sequence. If so, program control is transferred to lines 1220 thru 1350 where an error message is output and the error flag cleared by sending an ECLR instruction to the calculator chip. Line 1380 tests to see if the calculator is in the floating point or scientific mode. If floating point, control is transferred to lines 1400 thru 1670. If scientific, control is transferred to lines 1680 thru 1990. In both modes the data is output to the display in the selected mode and program control is transferred back to line 380 where new commands or data may be entered.

FAGE 001 CALC-1


```
    OOE2 O2
    OOES OS
    OOE4 04
    OOES 05
    OOBS O6
    00B7 07
00120 OOEE O8
    0089 09
    OOBA OF
    OOBB OF
    OOBC OF
    OOED OF
    OOBE 22
    OOBF OF
00130 00CO OF
    00C.1 IB
    0OC2 36
    00C3 25
    00C4 2D
    00C5 OB
    00C6 2C
    000:7 1C
00140 00C8 1[
    00C9 20
    OOCA OF
    OOCE OF
    OOCC OF
    00C[1E
    OODE OS
    OOCF 23
OO150 00DO OD
    00D1 33
    00[12 37
    00113 24
    00[14 26
    00|5 32
    OODE. 34
    00017 31
00160 OOLE 50
    00[19 2E
    OODA OC
    OONE OF
    OOLC OF
    OODD OF
    OOLE S8
    OODF OF
OO170 OOEO OF
    OOE1 OF
    OOE2 36
    OOES 25
    OOE4 2[I
    OOES OB
    00E6 2C
    OOE7 1C
    FCE $OE,$O%,$OF, कOF,$OF,$OF,$2Z,$OF
    FEE $OF, &1E, कSG,$2G, क2D, कOE, क2L, $1C
```




```
    FEE $EO, कZE, कOT, कOF, कOF, कOF, #SE, कOF
```



```
FAGE OOS CALC-1
```

$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOE } & 20 \\ \text { OOEA } & \text { OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
FEE $\quad \$ 1[1, \$ 20, \$ 0 F, \$ 0 F, \$ 0 F, \$ 1 E, \$ 5, \$ 23$
$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOES } & 20 \\ \text { OOEA OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOES } & 20 \\ \text { OOEA OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOES } & 20 \\ \text { OOEA OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOES } & 20 \\ \text { OOEA OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOES } & 20 \\ \text { OOEA OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOES } & 20 \\ \text { OOEA OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
$\begin{array}{rrr}\text { OOLEO OOES } & 10 \\ \text { OOES } & 20 \\ \text { OOEA OF } \\ \text { OOEE OF } \\ \text { OOEC } & \text { OF } \\ \text { OOED } & 18 \\ \text { OOEE } & 35 \\ \text { OOEF } & 23\end{array}$
00190 OOFO OD FCE $\$ 0[1, \$ 33, \$ 37, \$ 24, \$ 26, \$ 32, \$ 34, \$ 31$
OOF 133
OOF2 37
OOFS 24
$00 F 426$
OOFE 32
OOFE 34
OOF7 31
00200 00FE 30
OOFG 2E
OOFA OC
OOFE OF
OOFC OF
OOFD OF
OOFE OF
OOFF OF
002100100 GRTG \$0100
002200100 EE A047 STAFT L[IE \#\#AO47
002300103 CE 028:7 LIX \#LLFELEN
002400106 ED EO7E JEF FLIATA1
FAFARIR
INITAL
comanci
INITAL
ERA
$002800110867 F$
002900112 A7 00
003000114 Eb 36
003100116 A7 01
0032001188600
00330 011A A7 02
$00340 \quad 011 \mathrm{C} 86 \quad 34$
$00350011 E$ A7 03
003600120 Ab 02
00370012239
003800123 BD EIAC COMANL JER
003900126 8A 80
$00400 \quad 0128 \quad 87 \quad 012 C$
00410012 B 9600 FOINT LLA A $\$ 00$
DRA A
GTA A FGINT+1
00412012 D 8121
00414 O12F 2743
0042001318002
0043001332017
004400135 E6 01
004500137 2A FC
004600139 A7 00
00470013 E E6 00
004800130 C6 3C
FQE $\$ \mathrm{SO}, \mathrm{\$ 2E}, \$ 0 \mathrm{O}, \$ \mathrm{OF}, \$ \mathrm{OF}, \$ \mathrm{OF}, \$ \mathrm{OF}, \$ \mathrm{OF}$
FCE $\$ 0[1,535,537, \$ 24, \$ 26, \$ 32,534, \$ 31$
---

| 00490 | 01 EF | E701 |  | ETA | $E$ | i，$x$ | EFING HILLI LINE LOW |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00500 | 0141 | EG 01 | WAIT10 | L［1A | $E$ | i．$X$ |  |
| 00510 | 0143 | 2A FC |  | EFFL |  | WATTO | Lioir Fiof FEACM LEW |
| 00520 | 0145 | E6 00 |  | LIA | E | ）$x$ | 二LEAF FLAG Eit |
| 00500 | 0147 | CS E6 |  | L．LTA | E： | \＃安三6 | HIGH HOLL－FIS REAE＇Y |
| 00540 | 0147 | E7 01 |  | STA | E： | i，$X$ | FETURN HOLE IINE HISH |
| O0550 | O14E | 39 |  | FTS |  |  |  |
| 00.60 | O14C | $812 F$ | CHFICHE | EMF | f | \＃ 2 $_{\text {2F }}$ |  |
| 00570 | 014 E | 26.03 |  | ENE |  | EKIFTS |  |
| 00575 | 0150 | $7 F$ O2AE |  | $E L F$ |  | FDFIMAT |  |
| 00580 | 015 | $7 \square 102 A F$ | SKIF75 | TET |  | SM［ios | 二HESt FBF FAEVILSE SMLIG INSTR |
| 00590 | 0156 | 2610 |  | ENE |  | ZEFTMEM |  |
| 00600 | 015 | S1 OF | GINTSO | CMF | A | \＃まOF |  |
| 00620 | 015 A | $27 \quad 67$ |  | EES |  | CimANE | GET MGFE［IATA IF NDF |
| 00600 | 0150 | E1 18 |  | C：MF＇ | A | \＃ 土 $_{\text {13 }}$ |  |
| 00640 | $015 E$ | 2605 |  | ENE |  | SFIF゙S |  |
| 006.50 | 0160 | 73 OZAF |  | C0M |  | EMLiE： |  |
| 00660 | 016 | 20 EE |  | EFFA |  | EMMANE | GET MBFE LIATH IF EMEIE |
| 00670 | 0165 | 8120 | SrIḞ゙－ | CMF． | $A$ | \＃まです |  |
| 006.80 | 0167 | 27 EA |  | EES |  | CDMANL | GET MIAE LIATA IF INV |
| 00690 | 0169 | 81 OE： |  | CMF | A | \＃$\ddagger$ OE： |  |
| 00700 | 016 E | 23 EO |  | ELS |  | COMANI | GET MIFE［IATA IF NIMEEFE |
| 00710 | 016 CI | 3122 |  | EMF | A | \＃ $\mathbf{2}_{2} \mathbf{2}$ | LGint FGFe Ting |
| 00720 | 016 F | 26． 03 |  | EINE |  | ZEFVMEM |  |
| 00730 | 0171 | 73 OZAE |  | CO |  | Fi＿FIMAT |  |
| 00740 | 0174 | 7F OZAF | ZEFIMEM | ILF |  | EMIT． | ZEFIG EMEI＊ |
| 00750 | 0177 | EL O2 |  | ESEA |  | EETMEM |  |
| 00760 | $017 \%$ | 2011 |  | EFA |  | Lighami |  |
| 00770 | 017E | 7 F 0020 | EETMEM | ELF |  |  | CLEAF 50020 |
| 00780 | 017E | CE 0020 |  | Lilix |  | \＃ 2 $^{\text {2 }}$ | EIITTIM IF EIIFFERI |
| 00790 | 0181 | C．6 20 |  | LIAA | E | \＃ 520 |  |
| 00800 | 0183 | OS | LIMF＇ 1 | INX |  |  |  |
| OOE：10 | 0184 | E7 00 |  | ETA | E | $0, \mathrm{X}$ | ETIFE A EFALE |
| 008：20 | 0186 | EC 002E |  | EFX |  | \＃主云 | GHEGT FGF TGF GF EUIFFER |
| 00850 | 0187 | 26．FE |  | ENE |  | LiluFi |  |
| 00840 | 018 E | 39 |  | FTE |  |  |  |
| 0085 | 01 CL | FE AOOZ | LGIMALIF | LIX |  | FAFAALA |  |
| OOECO | 01 FF | E［10 |  | ESFi |  | CIITANE |  |
| 00370 | 0191 | 2030 |  | ERIA |  | －BITEHF |  |
| OOEEO | 0175 | EC 01 | CIITANE | LIAA | E | i，$\chi$ |  |
| 00890 | 0195 | ZA FC： |  | EF＇L |  | Gilitanz |  |
| 00700 | 0197 | $A B 00$ |  | LIIA | $A$ | $0, \mathrm{X}$ | ILEAFi FiAm Eit |
| 00910 | 0199 | Eb 16 |  | LIIA | $A$ | \＃${ }^{\text {a }} \mathrm{S}$ | SENLI ANM Init |
| 00920 | O19E | A7 00 |  | ETA | A | O，$X$ |  |
| 00950 | O175 | C6 3 E |  | LIIA | $E$ | \＃5：$\%$ |  |
| 00940 | 01\％F | E7 01 |  | STA | $E$ | 1 ，$X$ | EFING HGLLE LIVE LDW |
| 00750 | O1A1 | $E G 01$ | WAITEO | LIMA | E | i，$X$ | WAIT FGR EEIGINLI FEAEY |
| 00960 | O1AS | 2 AFC |  | EFF＇L |  | WAITEG |  |
| 00970 | $01 A S$ | E6 00 |  | LIIA | E | O，X | CLEAF FLAG EIT |
| 00880 | $01 A 7$ | EB OF |  | L［IA | $A$ | $\because$ まF |  |
| 00990 | $014 \%$ | A7 00 |  | ETA | A | $\cdots x$ | EENLI A NGF |
| 01000 | 01 AE | E6 03 | WAITE | L［IA | E： | $\therefore$ S，$X$ |  |
| 01010 | 01 AD | 2 B 06 |  | EWM I |  | Hiticis | TFANEFEF İALI：LIATA TO MEMIFY |
| 01020 | OIAF | E6 O1 |  | LLIA | F | 1，$x$ | LIOH FGR FEALY ETFIGE |



FAGE OOS CALC-1


029236
029 S8
$0294 \quad 30$
$02 \% 50$
02960
029743
027E 41
029940
027A 43
0Z9E 2
029世 31
029 Cl 20
02のE 43
02ワF 41
02 AO 4 C
O2A1 43
$02 A 255$
$0 \angle A 34 C$
$02 A 441$
$02 A 554$
O2AB 4F
$02 A 752$
02020 O2AE OD
OZA9 OA
$0 \geq A A \quad 00$
OZAB 00
OZAI： 00
OZALI 04
$020: 0$ OZAE OO
02040 OZAF 00
02050 O2EO OLI
OZEI OA
OZE2 OO
OZES 00
$02060 \quad 02 \mathrm{B4} 45$
O2ES 52
02 EB 52
$02 \mathrm{B7} 4 \mathrm{~F}$
OZE 52
02070 O2E＇9 O4 O20EO EOTE 02090 AOO2 02100 E1AD 02110 E1LI
02120 A04E
02130 A04E 0100
$02140 \mathrm{AOO2}$
02150 A002 800C
02160
ETAFT 0100
INITAL 0110
COMANDI O123
FIINT 012 E
GUTINE OISS


FGRMAT FGE कOO
SMLK FGE कOQ


FGT EFFIG
FEE ま〇4

FILATA1 EOLS $\operatorname{tEOTE}$
FAFALIF EDU \＄AOGZ
INEEE EDII \＆EIAB：
OLITEEE EGG कE1［I］
IRG \＄AOAS
FLE $\$ 0100$
BRG $\quad \mathrm{AOO}$
FLIE क心णG：

ENEI

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PAGE OOE CALC-1
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WAIT1O O141
CHFCHK O14C
SKIF75 0153
CONT50 0158
SKIF25 0165
ZEFMEM 0174
SETMEM 017B
LOOF1 0183
LODADR 018 C
OILITANS 0193
WAITSO OLA1
WAITS OLAB
OUITIIG O1BS
FOINT2 OICS
CONFLG OIC9
OUTCHR OIDO WAIT7O O1D2 WAIT71 OIEO CONT 1 OIFO FLOFNT O1FB MINF'NT 0208 FRINT 1 OZOA DFIND 020D DIGLOF O21A ENDCH1 0229
SCINOT 02FA
NEGFNT 0244
FFIINT2 0246
NUMLOF 024C
SKIPLP 025C
SKPSGN 0271
CLRSCN 0287
CRLF O2AB
FORMAT O2AE
SMDC O2AF
ERRMSG 02BO
PDATA1 EOTE
PARALR AOO2
INEEE EIAC
OUTEEE E1D1
TOTAL ERFOFS 00000

# RPN-the only language that lets you "speak"with contidence and consistency to a pocket-sized computer calculator. 

In 1967, Hewlett-Packard embarked on a major new development effort: to design a family of advanced computer calculators powerful enough to solve complex engineering/scientific problems yet simple enough to be used by anyone who works with numbers.
As part of this effort, HP carefully evaluated the strengths and weaknesses of the various languages which an operator might use to communicate with an electronic calculating device. Among those studied were:

- computer languages such as BASIC and FORTRAN,
- various forms of algebraic notation, and
- RPN (Reverse Polish Notation), a parenthesis-free but unambiguous language derived from that developed by the Polish mathematician, Jan Lukasiewicz.

As might be expected, each of these languages was found to excel in a particular application. For its biggest programmable desktop calculators, HP selected BASIC. For its other powerful desktop calculators, with less extensive storage capacity, HP chose algebraic notation.
But, given the design constraints of a pocket-sized scientific computer calculator, RPN proved the simplest, most efficient, most consistent way to solve complex mathematical problems.

## Only RPN offers these powerful advantages

Compared to alternative logic systems, Hewlett-Packard believes that only RPN -- in combination with a 4 -register operational memory stack - gives you these powerful advantages.

1. You can always enter your data the same way, i.e., from left to right-the same way you read an equation. Yet, there is no need for a parenthesis key; nor for a complicated "operational hierarchy."
2. You can always proceed through your problem the same way. Once you've entered a number, you ask: "Can I perform an operation?" If yes, you do it. If no, you press ENTERA and key in the next number.
3. You always see all intermediate answers - as they are calculated - so that you can check the progress of your calculation as you go. As important, you can review all numbers stored in the calculator at any time by pressing a few keys. There is no "hidden" data.
4. You don't have to think your problem all the way through beforehand unless the problem is so complex that it may require simultaneous storage of three or more intermediate answers.
5. You can easily recover from errors since all operations are performed scquentially, immediately after pressing the appropriate key.


The RPN method consists of four, easy-to-remember steps. Once learned, it can be applied to almost any mathematical expression.
6. You don't have to write down and re-enter intermediate answers, a real time-saver when working with numbers of cight or nine digits each.
7. You can communicate with your calculator confidently, consistently because you can always proceed the same way.
If all this sounds too good to be true, bear with us you'll soon get the chance to see for yourself. But first, we need to describe how RPN and the 4-register operational stack operate.

## The RPN method-it takes a few minutes to learn but can save years of trustration.

Yes, the RPN method does take some getting used to. But. once you've learned it, you can use the RPN method to solve almost any mathematical expression-confidently, consistently.
There are only four easy-to-follow steps:

1. Starting at the left side of the problem, key in the first or next number.
2. Determine if any operations can be performed. If so, do all operations possible.
3. If not, press ENTER to save the number for future use.
4. Repeat steps 1 through 3 until your calculation is completed.
A diagram of the RPN method is shown above.

## Simple arithmetc, the RPN way.

Just to show how it works, let's try the RPN method on two simple problems (we'll use them again in the comparisons that begin on the next page).
Problem: $3 \times 4=12$
RPN solution:

| Step | Press | See Displayed |
| :---: | :---: | :---: |
| 1. Key in firat number. | 3 | 3 |
| 2. Since onty one number has been keyed in, no operations are posalbie. Press ENTER4. | ENTER | 3 |
| 3. Key in next number. | $4]$ | 4 |
| 4. Since both numbers are now In calculator, multiplication can be performed. | $x$ | 12 |

Problem: $(3 \times 4)+(5 \times 6)=42$
RPN solution:

| Step | Press | See <br> Dispisyed |
| :--- | :---: | :---: |
| 1. Key in first number. <br> 2. No operations possible. Press <br> ENTERt. | 3 | 3 |
| 3. Key in second number. |  |  |
| 4. Since both numbers are in |  |  |
| calculator, first multiplication |  |  |
| is possible. |  |  |

If you've followed us this far, you've noticed two important facts:

1. Both of these problems were solved in the same, consistent manner, using the same simple set of rules.
2. All intermediate answers were displayed as they were calculated, and stored and retrieved as needed to complete the calculation. With RPN and a 4 -register operational memory stack, there is almost never a need to write down intermediate answers.

## How the operational stack works.

The four registers of HP's exclusive operational stack can be represented by the following diagram.


When a number is keyed in, it goes into the $X$ register for display. Pressing the ENTER 4 key duplicates the contents of the X register into the Y register and moves all other numbers in the stack up one position.
When an operation key $(\square],[-], x,\left[\div, x^{y}\right)$ is pressed the operation is performed on the numbers in the $X$ and Y registers, and the answer appears in the X register for display. Numbers in the other registers automatically drop one position.
To demonstrate these points, we'll show what happens to the stack as we solve the problem: $(3 \times 4)+(5 \times 6)$ $=42$.


As you can see, all numbers are automatically positioned in the stack on a last-in-first-out basis, in the proper order for subsequent use.
Now that we've described how RPN logic operates, we canproceed with our problem-hy-problem comparison of this system versus two others used in today's scientific pocket calculators.

We think you will find it interesting.

## Calc-l Instuction Set

Full Name ASCII Character
0 ..... 0
1 ..... 1
2 ..... 2
3 ..... 3
4 ..... 4
5 ..... 5
6 ..... 6
7 ..... 7
8 ..... 8
9 ..... 9
Decimal point
$\dot{E}$
Enter exponent
2
Change sign
P
Constant PI
M
Set mantissa digit count
$X$ exchange $M$ ..... A
Memory store ..... G
Memory recall ..... H
Inverse mode ..... I
Enter ..... sp or cr
Toggle mode ..... $>$
Roll stack ..... 0
Sine $X$ ..... S
Cosine X ..... C
Tangent $X$ ..... T
Error clear ..... Y
Radians to degrees ..... F
Degrees to radians ..... D
Master clear ..... cntrl $X$
$X$ exchange $Y$ ..... X
E to X ..... W
Ten to $X$ ..... U
Square ..... Q
Square root ..... V
Natural $\log$ of $X$ ..... N
Base 10 log of X ..... B
One divided by $X$ ..... R
Y to X ..... ヘ
Plus ..... $+$
Minus ..... *
Times ..... *
Divide ..... /

Table I

MM57109 Instruction Description Table (Continued) (* Indicates 2-word instruction)

| CLASS | SUBCLASS | MNEMONIC* | $\begin{aligned} & \text { OCTAL OP } \\ & \text { CODE } \end{aligned}$ | FULL NAME | DESCRIPTION |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Branch | Count | IBNZ | 31 | Increment memory and branch if $M \neq 0$ <br> Decrement memory and branch if $M \neq 0$ Multidigit input to $X$ | $M+1 \rightarrow M$. If $M=0$, skip second instruction word. Otherwise, branch to address specified by second instruction word. |
|  |  | DBNZ | 32 |  | $\mathrm{M}-1 \rightarrow \mathrm{M}$. If $\mathrm{M}=0$, skip second instruction word. Otherwise, branch to address specified by second instruction word. |
| 1/0 | Multi-digit | $\mathrm{IN}^{*}$ | 27 |  | The processor supplies a 4 -bit digit address (DA4-DA1) accompanied by a digit address strobe ( $\overline{O A S}$ ) for each digit to be input. The high order address for the number to be input would typically come from the second instruction word. The digit is input on D4-D1, using ISEL $=0$ to select digit data instead of in. structions. The number of digits to be input depends on the calculation mode (scientific notation or floating point) and the mantissa digit count (See Data Formats and Instruction Timing). Data to be input is stored in $X$ and the stack is pushed $(X \rightarrow Y \rightarrow Z \rightarrow T)$. At the con. clusion of the input, DA4--DA1 $=0$. |
|  |  | OUT* | 26 | Multidigit output from $X$ | Addressing and number of digits is identical to IN instruction. Each time a new digit address is supplied, the processor places the digit to be output on DO4-DO1 and pulses's the R/W line active low. At the conclusion of output, $\mathrm{DO} 4-$ $D O 1=0$ and $D A 4-D A 1=0$. |
| 1/0 | Single-digit | AIN | 16 | Asynchronous Input | A single digit is read into the processor on $\mathrm{D4}-$ D1. ISEL $=0$ is used by external hardware to select the digit instead of instruction. It will not read the digit until $A D R=0$ (iSEL $=0$ selects $\overline{A D R}$ instead of $(5)$, indicating data valid. $F 2$ is pulsed active low to acknowledge data just read. |
| 1/0 | Fiags | $\begin{aligned} & \text { SF } 1 \\ & \text { PF } 1 \end{aligned}$ | $\begin{aligned} & 47 \\ & 50 \end{aligned}$ | Set Flag 1 <br> Pulse Flag 1 | Set $\mathrm{F}_{1}$ high, i.e. $\mathrm{F}:=1$. <br> Fl is pulsed active high. If F is aiready high. this results in it being set low. |
|  |  | $\begin{aligned} & \text { SF2 } \\ & \text { PF2 } \end{aligned}$ | $\begin{aligned} & 51 \\ & 52 \end{aligned}$ | Set Flag 2 Pulse Flag 2 | Set F2 high, i.e. $F 2=1$. <br> F2 is pulsed active high. If F2 is already high, this results in it being set low. |
|  |  | PRW 1 | 75 | Pulse R $\bar{N} \mathbf{W} 1$ | Generates $R / \bar{W}$ active low pulse which may be used as a strobe or to clock extra instruction bits into a flip.flop or register. |
|  |  | PRW2 | 76 | Pulse R/W 2 | Identical to PRW1 instruction. Advantage may be taken of the fact that the last 2 bits of the PRW1 op code are 10 and the last 2 bits of the PRW2 op code are 01. Either of these bits can be clocked into a flig.flop using the R/W pulse. |
| Mode Control |  | TOGM | 42 | Potio Made <br> trexnew | Change mode from floating point to scientitic notation or vice versa, depending on present mode. The mode affects only the $I N$ and OUT instructions. Internal calculations are always in 8 -digit scientific notation. |
|  |  | SMDC* INV | 30 40 | FerMmasse bigitwint Vvot Mode <br>  | Mantissa digit count is set to the conterits of the second instruction word ( $=1$ to 8 ). <br> Set inverse mode for trig or memory function instruction that will immediately follow. Inverse mode is for next instruction only. |

Table I

MM57100 Inatruction Description Table (* Indicates 2 -word inatruction)

| CLASS | SUBCLASS | MNEMONIC* | $\begin{aligned} & \text { OCTAL OP } \\ & \text { CODE } \end{aligned}$ | FULL NAME | DESCRIPTION |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Digit Entry |  |  | $\begin{aligned} & 00 \\ & 01 \\ & 02 \\ & 03 \\ & 04 \\ & 05 \\ & 06 \\ & 07 \\ & 10 \\ & 11 \\ & 12 \\ & 13 \\ & 14 \end{aligned}$ $15$ $41$ |  | Mantissa or exponent digits. On first digit (d) the following occurs: $Z \rightarrow T$ $Y \rightarrow Z$ <br> $X \rightarrow Y$ <br> $d \rightarrow X$ <br> See description of number entry on page 11. <br> Digits that follow will be mantissa fraction. Digits that follow will be exponent. <br> Change sign of exponent or mantissa. <br> $X_{m}=X$ mantissa <br> $X_{e}=X_{\text {exponent }}$ <br> CS causes $-X_{m} \rightarrow X_{m}$ or $-X_{e} \rightarrow X_{e}$ depending on whether or not an EE instruction was executed after last number entry initiation. <br> $3.1415927 \rightarrow X$, stack not pushed. <br> Terminates digit entry and pushes the stack. The argument entered will be in $X$ and $Y$. $\begin{aligned} & Z \rightarrow T \\ & Y \rightarrow Z \\ & X \rightarrow Y \end{aligned}$ |
|  |  | NOP HALT | 77 17 | No Operation <br> Halt | Do nothing instruction that will terminate digit entry. <br> External hardware detects HALT op code and generates HOLD $=1$. Processor waits for HOLD $=0$ before continuing. HALT acts as a NOP and may be inserted between digit entry instructions since it does not terminate digit entry. |
| Move |  | ROLL | 43 | - 4, 4, | Roll Stack. |
|  |  | POP | 56 | Pope | Pop Stack. $\begin{aligned} & Y \rightarrow X \\ & Z \rightarrow Y \\ & T \rightarrow Z \\ & O \rightarrow T \end{aligned}$ |
|  |  | XEY | 60 | - | $\begin{aligned} \text { Exchange } X \text { and } Y . \\ X \leftrightarrow Y \end{aligned}$ |
|  |  | XEM | 33 | $19$ | $\begin{gathered} \text { Exchange } X \text { with memory. } \\ X \leftrightarrow M \end{gathered}$ |
|  |  | MS | 34 |  | Store $X$ in Memory. $X \rightarrow M$ |
|  |  | MR | 35 |  | Recall Memory into $X$. $M \rightarrow X$ |
|  |  | LSH | 36 | Left Shift Xm | $X$ mantissa is left shifted while leaving decimal point in same position. Former most significant digit is saved in link digit. Least significant digit is zero. |
|  |  | RSH | 37 | Fight Shift Xm | $X$ mantissa is right shifted while leaving decimal point in same position. Link digit, which is normally zero except after a left shift, is shifted into the most significant digit. Least significant digit is lost. |

Table I

MMEF7109 Instruction Oescription Table (Continued) (* Indicates 2-word instruction)

| CLASS | SUBCLASS | MNEMONIC* | $\begin{aligned} & \text { OCTAL OP } \\ & \text { CODE } \end{aligned}$ | FULL NAME | DESCRIPTION |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Math | $f(X, Y)$ | + | 71 | $2$ | Add $X$ to $Y X+Y \rightarrow X$ On +, - $X$. I and $Y X$ instructions, stack is popped as follows: $\begin{aligned} & Z \rightarrow Y \\ & T \rightarrow Z \\ & O \rightarrow T \end{aligned}$ <br> Former $X, Y$ are lost. |
|  |  | - | 72 | 5rna | Subtract $X$ from $Y$. $Y-X \rightarrow X$ |
|  |  | x | 73 |  | Multiply $X$ times $Y$. $Y \times X \rightarrow X$ |
|  |  | 1 | 74 |  | Divide $X$ into $Y$. $Y \div X \rightarrow X$ |
|  |  | YX | 70 |  | Raise $Y$ to $X$ power. $Y^{X} \rightarrow X$ |
|  | $F(X, M)$ | INV +* | 40, 71 | Memory Plus | Add $X$ to memory. $M+X \rightarrow M$ On INV +, -, $X$ and ; instructions, $X, Y, Z$, and $T$ are unchanged. |
|  |  | INV -* | 40, 72 | Memory Minus | Subtract $X$ from memory. $M-X \rightarrow M$ |
|  |  | INV ${ }^{\text {- }}$ | 40, 73 | Memory Times | Multiply $X$ times memory. $M \times X \rightarrow M$ |
|  |  | INV /* | 40, 74 | Memory Divide | Divide $X$ into memory. $M \div X \rightarrow M$ |
|  | $F(X)$ Math | $1 / \mathrm{X}$ | 67 |  | $1 \div X \rightarrow X$. On all $F(X)$ math instructions $Y, Z$. $T$ and $M$ are unchanged and previous $X$ is lost. |
|  |  | SORT | 64 |  | $\sqrt{x} \rightarrow x$ |
|  |  | SO | 63 |  | $\mathrm{x}^{2} \rightarrow x$ |
|  |  | 10x | 62 | 10, ${ }^{4}$ | $10^{x} \rightarrow x$ |
|  |  | EX | 61 |  | $e^{x} \rightarrow x$ |
|  |  | LN | 65 |  | $\ln x \rightarrow x$ |
|  |  | LOG | 66 |  | $\log x \rightarrow x$ |
|  | $F(X)$ Trig | SIN | 44 | H 4 | $\operatorname{Sin}(X) \rightarrow X$. On all $F(X)$ trig functions. $Y, Z, T$. and $M$ are unchanged and the previous $X$ is lost. |
|  |  | $\cos$ | 45 |  | $\cos (x) \rightarrow x$ |
|  |  | TAN | 46 | 14.403 | $\operatorname{TAN}(\mathrm{X}) \rightarrow \mathrm{x}$ |
|  |  | INV SIN* | 40, 44 | Tnverse sine $X$ | $\operatorname{SiN}^{-1}(X) \rightarrow X$ |
|  |  | INV COS* | 40, 45 | Inverse cosine $X$ | $\cos ^{-1}(x) \rightarrow x$ |
|  |  | INV TAN* | 40. 46 | Inyerse tan $X$ | $\operatorname{TAN}^{-1}(x) \rightarrow x$ |
|  |  | DTR | 55 |  | Convert $X$ from degrees to radians. |
|  |  | RTD | 54 |  | Convert X from radians to degrees. |
| Clear |  | MCLA | 57 |  | Clear all internal registers and memory: initialize $1 / 0$ control signals, $M D C=8 . M O D E=$ floating point. (See initialization.) |
|  |  | ECLR | 53 | tix $\times$ ctam | $0 \rightarrow$ Error flag |
| Branch | Test | JMP* | 25 | Jump | Unconditional branch to address specified by second instruction word. On all branch instruc thons, second word contains branch address to be loaded into external PC. |
|  |  | TJC* | 20 | Test jump condition | Branch to address speaified by second instruction word if JC ( $I_{6}$ ) is true $(=1)$. Otherwise, skip over second word. |
|  |  | TERR* | 24 | Test error | Branch to address specified by second instruction word if error flag is true $(=1)$. Otherwise, skip over second word. May be used for detecting specific errors as opposed to using the automatic error recovery scheme dealt with in the section on Error Control. |
|  |  | $T X=0^{*}$ | 21 | Test $\mathrm{X}=0$ | Branch to address specified by second instruc tion word if $X=0$. Otherwise, skip over second word. |
|  |  | TXF* | 23 | Test $\|\mathrm{X}\|<1$ | Branch to address specified by second instruction word if $\|X\|<1$. Otherwise, skip over second word. (i.e. branch if X is a fraction.) |
|  |  | TXLTO* | 22 | Test $\mathrm{X}<0$ | Branch to address specified by second instruction word if $X<0$. Otherwise, skip over second word. |

MM57109 Instruction Summary Table (* $=2 \cdot$ word instruction)

| $14-11$ | 1615 |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\emptyset$ | 1 | 2 | 3 |
| $\square$ | 40 | TJC* | 1x | , |
| 1 | 4 | TX: $0^{\circ}$ | 5 | , ${ }^{4}$ |
| 2 | \% | TXLTO* |  | \% |
| 3 | * | TXF* | W |  |
| 4 |  | TERR* | - |  |
| 5 |  | JMP* |  |  |
| 6 | , | OUT* | 8- , |  |
| 7 | 箱 | in* | SF 1 |  |
| 8 | \% |  | PF1 |  |
| 9 | 8 | 'BNZ* | SF2 | $M+1$ |
| A | 8 | DBNZ* | PF2 | (19-1 |
| B | \% | 5 5 | - 7exala | (1) ${ }^{(M x)}$ |
| C | 0 | KS | \%14 | - M -1 |
| D | 12 |  | Wx. | PRW 1 |
| E | AIN | LSH | POP | PRW2 |
| F | HALT | RSH |  | NOP |

# SET memory locations $A 002$ - A003 to 800 C prior to romviring program. 

Table III - CALC-1 Instruction to ASCII Character Lookup Table


| NAME | HEX OP CODE | MNEMONIC | ASCII CHARACTER |
| :---: | :---: | :---: | :---: |
| Texchnday | 30 | XEY | X |
| Prox | 31 | EX | W |
| Tentog | 32 | 1QX | U |
| Squard | 33 | SQ | Q |
| Squaregroo | 34 | SQRT | V |
| Hatural 10 do ${ }^{\text {a }}$ | 35 | LN | N |
|  | 36 | LOG | B |
| Dne dividedubyous | 37 | 1/X | R |
| [ to 2 | 38 | $Y \mathrm{X}$ | $\wedge$ |
| Plus | 39 | + | + |
| 4nd | 3A | - | - |
| 1mes | 3B | X | * |
| Divar | 3 C | / | / |
| Pulse R/W I | 3D | PRW1 |  |
| Pulse R/W 2 | 3E | PRW2 |  |
| No Operation | 3 F | NOP |  |


| Memory Location | $\underline{\mathrm{DP}} \mathrm{POS}$ | D7 | D6 | D5 | D4 | D3 | D2 | D1 | D0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 |  | $\emptyset$ | $\emptyset$ | 0 | $\emptyset$ | 0 | $\emptyset$ | 0 | 0 |
| 21 |  | $\emptyset$ | 0 | 1 | 0 | $\emptyset$ | $\emptyset$ | $\emptyset$ | 0 |
| 22 |  | 0 | 0 | 1 | 1 | Sm | $\emptyset$ | $\emptyset$ | 0 |
| 23 |  | $\emptyset$ | 0 | 1 | 1 | Dp | POS |  |  |
| 24 | $\emptyset \mathrm{B}$ | 0 | 0 | 1 | 1 | BCD | digit(left most) |  |  |
| 25 | DA | 0 | 0 | 1 | 1 | BCD | digit |  |  |
| 26 | 09 | $\emptyset$ | 0 | 1 | 1 | BCD | digit |  |  |
| 27 | $\emptyset 8$ | 0 | 0 | 1 | 1 | BCD | digit |  |  |
| 28 | 07 | 0 | 0 | 1 | 1 | BCD | digit |  |  |
| 29 | 06 | 0 | 0 | 1 | 1 | BCD | digit |  |  |
| 2A | ¢5 | 0 | 0 | 1 | 1 | BCD | digit |  |  |
| 2B | 04 | 0 | 0 | 1 | 1 | BCD | digit(right most) |  |  |


| Memory Location | D7 | D6 | D5 | D4 | D3 | D2 D1 | D0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 | 0 | 0 | 1 | 1 | Most | signif. | exp. digit |
| 21 | - | 0 | 1 | 1 | Leas | signif. | . exp. digit |
| 22 | 0 | 0 | 1 | 1 | Sm | 0 - | Se |
| 23 | not | USED |  |  |  |  |  |
| 24 | 0 | 0 | 1 | 1 | BCD | digit (1 | left most) |
| 25 | 0 | 0 | 1 | 1 | BCD | digit |  |
| 26 | 0 | 0 | 1 | 1 | BCD | digit |  |
| 27 | 0 | 0 | 1 | 1 | BCD | digit |  |
| 28 | 0 | 0 | 1 | 1 | BCD | digit |  |
| 29 | 0 | 0 | 1 | 1 | BCD | digit |  |
| 2A | 0 | 0 | 1 | 1 | BCD | digit |  |
| 2B | 0 | 0 | 1 | 1 | BCD | digit (1 | left most) |

Notes:

1) If the Mantissa Digit Count (set by SMDC instruction, initially 8) is less than 8, the unused digit memory locations will be filled with ASCII spaces ( $\mathbf{2 0}_{16}$ )
2) $S_{m}$ is the sign of the mantissa. 0 positive $l=$ negative
3) Se is the sign of the exponent = positive $1=$ negative
4) $D P$ POS is the decimal point position. The decimal point should follow the digit whose address is stored in memory location 24 when in the Scientific mode. In the Floating Point mode AND the data in memory location 23 with $\emptyset F$ and subtract the result from $2 F$ and $O R$ this with $2 \emptyset$. The decimal point should follow the digit whose address is given by the result.

Table $V$ - ASCII to CALCULATOR INSTRUCTION LOOKUP TABLE

| LSB | MSB | $\underline{0}$ | 1 | $\underline{2}$ | 3 | 4 | $\underline{5}$ | $\underline{6}$ | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 |  | OF | OF | 21 | 00 | OF | OD | OF | OD |
| 1 |  | OF | OF | OF | 01 | 1 b | 33 | OF | 33 |
| 2 |  | OF | OF | OF | 02 | 36 | 37 | 36 | 37 |
| 3 |  | OF | OF | OF | 03 | 25 | 24 | 25 | 24 |
| 4 |  | OF | OF | OF | 04 | 2D | 26 | 2D | 26 |
| 5 |  | OF | OF | OF | 05 | OB | 32 | OB | 32 |
| 6 |  | OF | OF | OF | 06 | 2C | 34 | 2C | 34 |
| 7 |  | OF | OF | OF | 07 | 1 C | 31 | 1 C | 31 |
| 8 |  | OF | 2 F | OF | 08 | 1 D | 30 | 1D | 30 |
| 9 |  | OF | OF | OF | 09 | 20 | 2B | 20 | 2B |
| A |  | OF | OF | 3B | OF | OF | OC | OF | OC |
| B |  | OF | OF | 39 | OF | OF | OF | OF | OF |
| C |  | OF | OF | OF | OF | OF | OF | OF | OF |
| D |  | 21 | OF | 3A | OF | 18 | OF | 18 | OF |
| E |  | OF | OF | OA | 22 | 35 | 38 | 35 | OF |
| F |  | OF | OF | 3 C | OF | 23 | OF | 23 | OF |

Example: An ASCII P is a hex 50 which points in the table to a $O D$ which is the constant PI instruction for the calculator chip

## TABLE VI- ERROR CONDITIONS

The ERROR flag on the calculator chip is set when:

1) LN $X$ when $X \leq 0 \quad$ LOG $X$ when $X \leq 0$
2) Any result $<10^{-99} \quad$ Any result $\geq 1099$
3) TAN $90^{\circ}, 270^{\circ}, 450^{\circ}$, etc.
4) $\operatorname{SIN} X, \operatorname{Cos} X, \tan X$ when $|X| \geq 9000^{\circ}$
5) $\sin ^{-1} x, \cos ^{-1} x$ when $|x|>$ or $|x| \leq 10^{-50}$
6) $S Q R T X$ when $X<0$
7) dividing by 0
8) Outputting a number in floating point mode if the number of mantissa digits to the left of the decimal point is greater than the mantissa digit count.

Figure I


ASCII to Hexadecimal Conversion Table

| LSB MSB | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | NUL | DLE | SP | 0 | C | P | - | P |
| 1 | SOH | DC1 | ! | 1 | A | Q | a | q |
| 2 | STX | DC2 | " | 2 | B | R | b | r |
| 3 | ETX | DC3 | * | 3 | C | S | c | $s$ |
| 4 | EOT | DC4 | \$ | 4 | D | T | d | t |
| 5 | END | NAK | $\%$ | 5 | E | U | e | u |
| 6 | ACK | SYN | \& | 6 | F | V | $f$ | $v$ |
| 7 | BEL | ETB | , | 7 | G | W | g | w |
| 8 | BS | CAN | ( | 8 | H | X | h | x |
| 9 | HT | EM | ) | 9 | 1 | Y | i | y |
| A | LF | SUB | * | : | J | Z | j | $z$ |
| B | VT | ESC | + | ; | K | [ | k | \{ |
| C | FF | FS | , | < | L | 1 | 1 | 1 |
| D | CR | GS | - | = | M | J | m | \} |
| E | SO | RS | . | 7 | N | $\wedge$ | n | $\sim$ |
| F | SI | US | 1 | ? | 0 | - | 0 | DEL |

