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# Optimal Burnable Absorber Assignment for PWR Core Reload Design

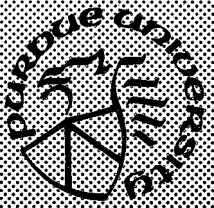
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# Optimal Burnable Absorber Assignment for PWR Core Reload Design

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## ABSTRACT

A new method has been developed to assign burnable poison loadings in the optimization of Pressurized Water Reactor core reload design. The method utilizes successive linear programming to determine the desired burnable poison loading. The optimum loading is selected after the evaluation of all candidate loadings close to the desired loading. The design method was implemented as a sub-program in the nodal core analysis code SIMULATE. The technique was applied to re-design Commonwealth Edison's Zion Unit-1 cycles 9 and 10. Significant improvements were achieved in cycle length, number of BP rods required, and power peaking. The present work completely automates the core reload design problem, significantly decreasing the time and effort required of the designer.

## 1. INTRODUCTION

### 1.1 Existing Method

A method for the optimization of core reload design has been developed and reported by previous researchers.<sup>1</sup> An algorithm based on this method was implemented as a sub-program within the LWR nodal core analysis computer program, SIMULATE-E.<sup>2</sup> The procedure was essentially divided into two separate optimization processes. The first is to determine the fuel loading pattern that yields the longest cycle. Secondly, the burnable poison (BP) loading is determined to control the core power peaking of the optimal pattern. The complete optimization procedure logic is shown in Figure 1. The optimization problem is made separable in this manner through the use of the Haling depletion.<sup>3</sup> By using this constant power depletion, the best loading pattern can be obtained totally independent of the control strategy.

The method used in the fuel loading optimization is a direct search technique which examines all possible two assembly exchanges from a user-input base loading pattern. Assembly exchanges are performed which yield an increase in the cycle length while still meeting peaking constraints. This procedure is repeated until the cycle length can no longer be increased by fuel shuffling. At this point, the direct search has been completed and the optimal loading pattern has been identified. The final loading pattern, and thus the 'optimum' cycle length, are

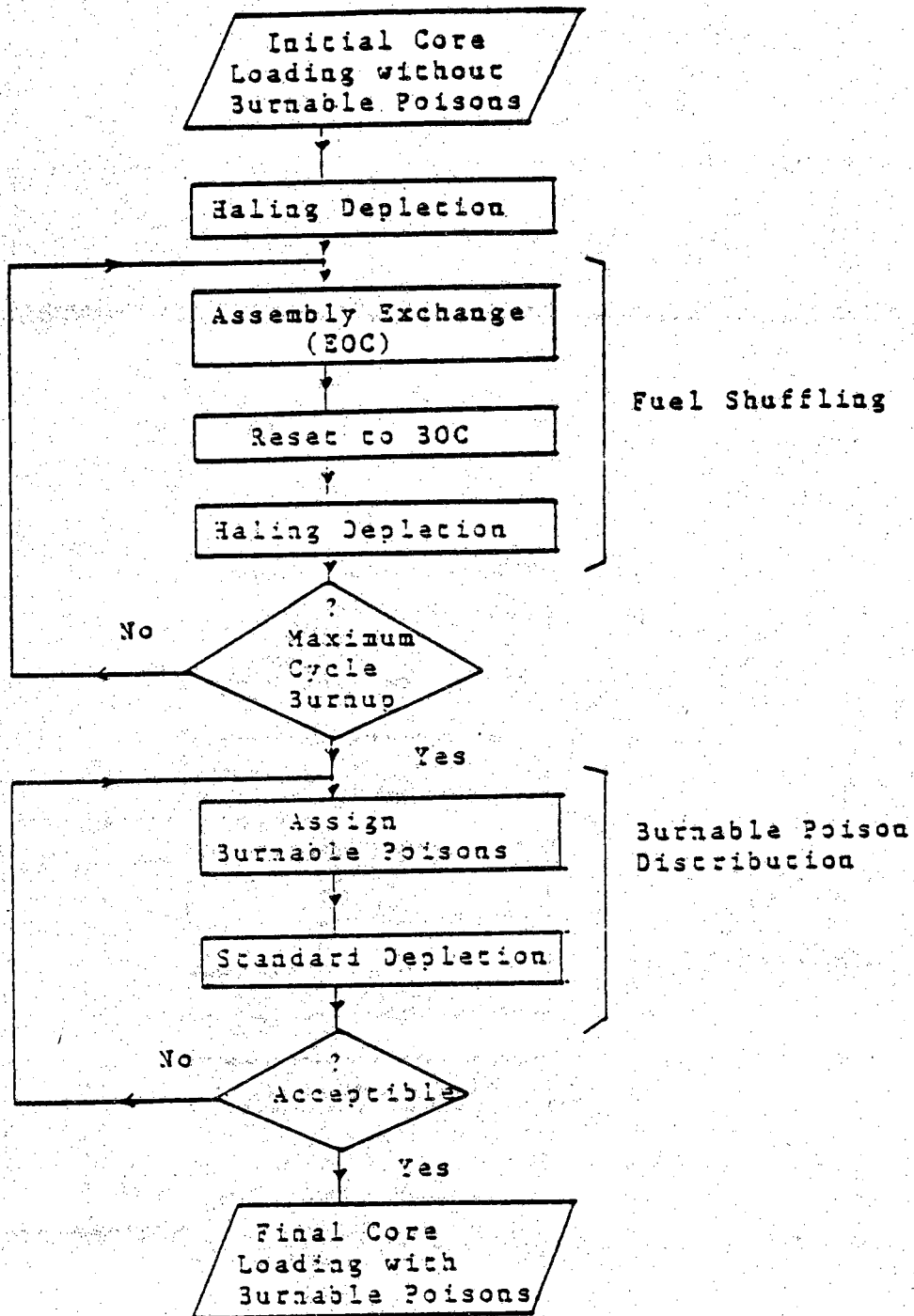


Figure 1. Optimization Procedure for Existing Reload Design Method



strongly influenced by the initial guess loading pattern because the search is not exhaustive, but only a subset of the complete factorial problem.

After the optimal fuel loading has been selected, it remains only to determine the required burnable poison power control. This task is accomplished through a linear programming solution to the second optimization problem. In the direct search algorithm, the objective function was the maximization of the cycle burnup. In this routine, however, the objective was to minimize the difference between the actual power distribution and some 'target' power distribution. For the control to be truly optimum, it would be necessary to have an optimum power shape as a target. This, however, is in itself a fairly significant problem and is not addressed in this work. Instead, the Haling power distribution from the previously determined optimal fuel loading pattern is utilized.

The solution method employs a successive linear programming technique in which the power distribution is represented as a linear function of the burnable absorber loading. This procedure is performed at each burnup step, giving an optimal BP trace throughout the depletion. Typical results of this procedure can be seen in Table 1. In the table, the assembly position is given in terms of the row and column indices of the assembly in the southeast octant of the reactor core (with position 1,1 being the core center). For the current work, as in most core design

TABLE 1. Optimal BP Loading Results from Existing Method

Cycle Burnup (GWD/MT)	Assembly Position								
	(2,2)	(4,1)	(4,3)	(5,4)	(6,5)	(6,6)	(7,2)	(7,3)	(7,4)
0.0	2.2	6.0	5.4	6.3	3.5	1.1	6.2	5.2	2.1
2.0	2.9	6.4	6.1	6.7	5.4	0.7	6.2	6.2	0.6
4.0	2.5	6.5	6.5	7.4	6.1	1.5	5.6	7.2	0.0
6.0	3.3	7.4	7.7	8.5	6.7	2.2	7.0	7.5	0.8
8.0	5.8	7.9	8.7	9.3	7.1	3.3	7.6	7.3	4.2
10.0	4.9	4.8	6.2	6.9	6.7	4.4	6.4	6.6	4.2
12.0	0.0	0.0	0.0	0.1	4.1	5.2	3.0	3.7	3.1

work, octant symmetry of the reactor core is assumed.

The major shortcoming of this present method is that the procedure is not completely automated. In fact, as it is now, the code requires a great deal of user interaction in the design process. This requires that the user possess a significant amount of insight into the core design problem. From data such as that in Table 1, the designer must select a BP distribution using available designs. Current design practice permits only multiples of four BP's per assembly. And, of course, this loading is held constant over the entire duration of the cycle. Thus the engineer is faced with the non-trivial task of selecting a BP loading which (subject to the above criteria) best fits the time varying non-realizable BP distributions returned by the code.

Even after this has been accomplished, the designer's work is far from over. To validate the core design loading, the engineer must perform a series of depletion calculations. If, at any point in the depletion, the core power peaking limits are violated, the BP loading must be adjusted and the procedure repeated. This step in the design process, which is basically a trial and error procedure, is by far the most time consuming and laborious. Finally, even when this manual iteration is complete and a BP loading which controls power peaking has been found, it still remains to be seen whether or not it is the 'optimal' BP distribution.

## 1.2 Objective of Current Work

Basically, the objective of the current work is to completely automate the burnable absorber assignment process. This would make the overall optimization procedure both faster and much easier for the core design engineer. Instead of the code returning a different, non realizable number of BP's for each burnup step, the updated code will return a single BP loading that is physically realizable. In addition, rather than simply accepting the first BP distribution that meets the core peaking constraints, the improved method selects the loading that gives the cycle of greatest length. Finally, the new code is totally automated, requiring no user interaction and a minimum of additional input.

## 2. BP ASSIGNMENT METHODOLOGY

### 2.1 Introduction

The existing BP assignment method requires the designer to perform a large portion of the design work manually and, therefore, achieves a final reload design which is less than optimal. There are basically two separate processes currently being done manually. The first is a selection of an initial BP loading from the linear programming results. The second is the depletion calculation and subsequent alterations to the initial loading due to power peaking violations. A methodology for automating these two problems and achieving a more nearly optimal design will be discussed in the following chapter. The complete logic flow diagram for the improved BP assignment method is given in Figure 2.

Previously, the core designer was required to use personal intuition to select the initial BP distribution from the non-integer linear programming results. One method used was averaging the BP values returned at the various burnup steps and using the closest available number of BP rods. Another method involved selecting the available BP loading closest to the desired BP loading at the point in the cycle requiring the greatest total number of BP rods. Typically, this occurred at the middle of the cycle, at a cycle burnup of approximately 8 GWD/MT. The basis for the current method is that the BP loading is best chosen from the beginning of the cycle (BOC) results.

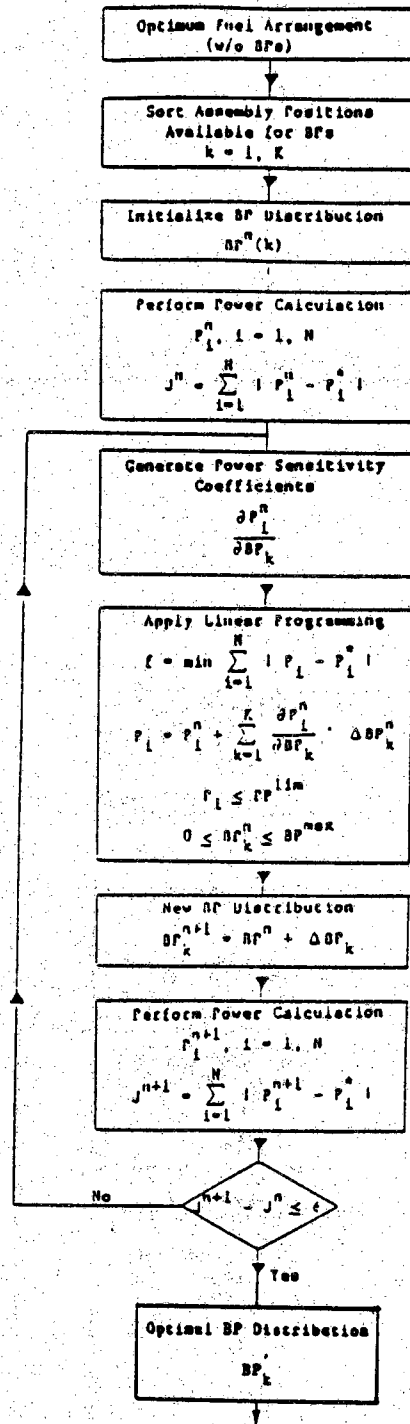


Figure 2. Improved Burnable Absorber Assignment Algorithm

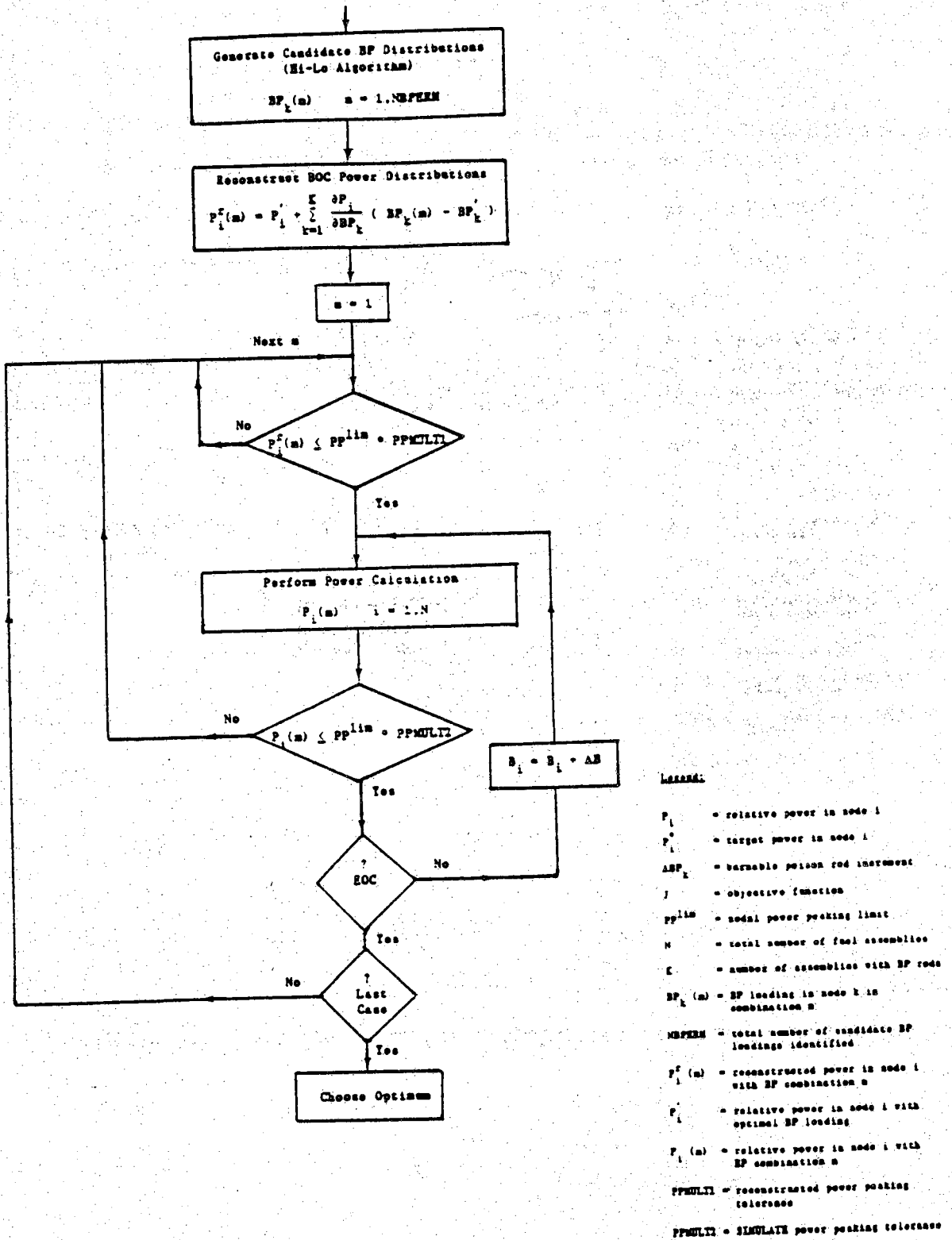


Figure 2. Continued

This can not be proven rigorously because of the complicated physics of the problem, but rather the motivation for this assumption can be seen through the following 'heuristic' rules:

1. If the assembly power is too high throughout the cycle (thus the BOC power,  $P_0$ , is high), then the BP loading should be increased.
2. If the power is too low throughout the cycle ( $P_0$  is too low), the BP loading should be decreased.
3. If the power is too low at BOC ( $P_0$  - low) and thus too high at the middle of the cycle (MOC), the BP loading should be decreased.
4. If the power is too high at BOC ( $P_0$  - high) and too low at MOC, the BP loading should be increased.

From the above implied correlation between the BOC power and the control requirements throughout the cycle, ( $P_0$  high  $\rightarrow$  increase,  $P_0$  low  $\rightarrow$  decrease), it appears that the BP assignment can effectively be performed at the beginning of the cycle. Thus, the linear programming procedure need only be executed at BOC, with the actual BP loading determined from these results.

At this point in the procedure, another major difference arises between the old and new methods. Instead of choosing a single initial guess BP loading, the improved code generates all possible BP distributions close to the optimal BP loading using a Hi-Lo algorithm. This algorithm will



be discussed in more detail in following sections. Each of these loadings is then depleted to end-of-cycle. If, at any point in the depletion, power peaking constraints are violated, the case is terminated and the next case is depleted. Following the depletion of all cases, the BP loading yielding the greatest cycle length is chosen as the optimum.

## 2.2 BOC Linear Programming

The objective for the optimal BP loading search is the minimization of the absolute difference between the actual power distribution and the target power distribution. As mentioned previously, an 'optimum' power shape has yet to be determined. In the absence of an optimum shape, the Haling power shape from the optimal fuel loading is used as the target. This target power does possess some inherent advantages. Haling first proposed<sup>3</sup> that maintaining a constant power shape throughout the cycle would yield the minimum power peaking for a given fuel loading. Maintaining a constant power distribution throughout the cycle is referred to as the Haling depletion.

The key to the solution of the BP loading optimization problem is the accurate prediction of the nodal relative powers,  $P_i$ . If the power distributions were calculated by solving the nodal diffusion equation (which is the standard procedure), the required computation time would be prohibitive. In order to accelerate the solution procedure, a first order perturbation approximation is made of the nodal relative power. This permits the representation of the power as a linear function of the BP loading. This approximation is given below:

$$P_i = P_i^0 + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} \Delta BP_k$$

where,

- $P_i$  = relative power in node  $i$
- $P_i^0$  = initial power in node  $i$  with base BP loading
- $\frac{\partial P_i}{\partial BP_k}$  = first order approximation of the change in the power in node  $i$  due to addition of one BP rod in position  $k$
- $\Delta BP_k$  = difference between the searched BP loading and the base BP loading and
- $K$  = total number of BP positions in the core.

The first derivative,  $\frac{\partial P_i}{\partial BP_k}$ , is determined numerically by perturbing the BP loading and performing standard nodal power calculations. This linear approximation of the nodal relative power is actually a very good one. The core-averaged relative difference between the actual power and the estimated power is less than a percent. The accuracy of the first order approximation will be discussed in further detail in the following sections.

The objective function for the linear programming problem can be written as

$$f = \min \sum_{i=1}^N | P_i - P_i^* |$$

where,

- $P_i$  = actual power in node  $i$
- $P_i^*$  = target power for node  $i$
- $N$  = total number of nodes (assemblies) in the core.

The problem is subject to the following constraints:

$$P_i \leq P_{lim} \quad i = 1, N$$

$$BP_k \geq 0 \quad k = 1, K$$

where,

$P_{lim}$  = core nodal power peaking limit

$BP_k$  = number of rods in BP position k.

The above three relationship are a statement of the optimization problem for the burnable absorber assignment. However, the objective function requires modification for solution by linear programming methods. The transformation of the three equations into the 'standard LP form' is as follows:

By making use of the following substitution of an independent variable,

$d_i$

$$d_i = | P_i - P_i^* |,$$

the objective function can be re-written as

$$f = \min \sum_{i=1}^N d_i.$$

Since it is standard practice to write an optimization problem in terms of the maximization of some quantity, the objective function for this problem can easily be transformed as follows:

$$f = \max \left( - \sum_{i=1}^N d_i \right).$$

Since  $P_i$  was previously defined as

$$P_i = P_i^0 + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} \Delta BP_k,$$

the independent variable  $d_i$  can be re-written by

$$d_i = \left| P_i^0 + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} \Delta BP_k - P_i^* \right|.$$

Re-arranging gives

$$d_i = \left| (P_i^0 - P_i^*) + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} \Delta BP_k \right|.$$

The absolute value in the objective function is not acceptable for LP solution because it is not a 'linear' function. The objective function needs to be re-formulated using additional inequality constraints. This re-formulation can be shown through the following simple example.

Suppose that the objective of a two region optimization problem is the minimization of the variable  $Y$ , subject to the constraint that  $Y = |X|$ . This is shown graphically in Figure 3.

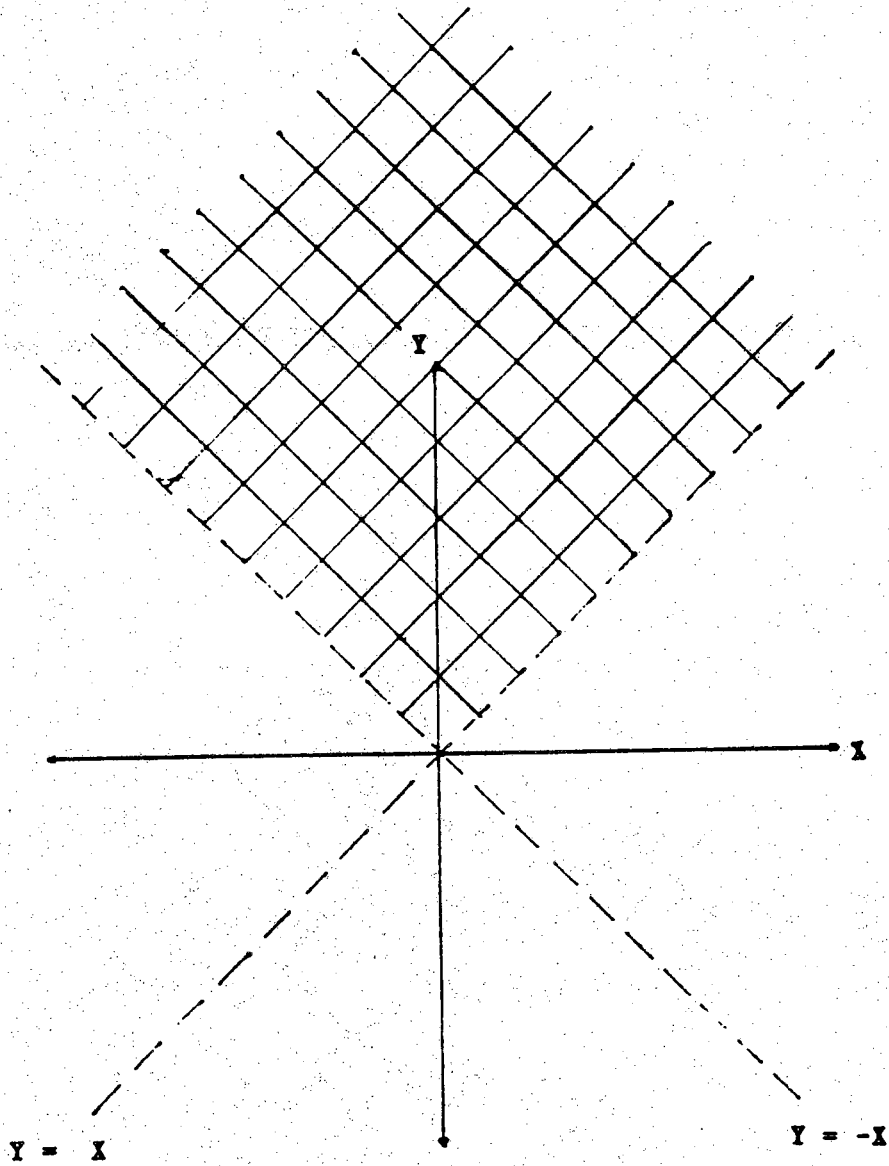


Figure 3. Simple Two Region Optimization Problem

In the figure, two functions are plotted; those being  $Y = X$  and  $Y = -X$ . The function  $Y = |X|$  can be represented as  $Y = X$ , for  $X \geq 0$ , and  $Y = -X$ , for  $X \leq 0$ . Even though these are equality constraints, this is still unsatisfactory for LP formulation. For a linear programming solution procedure, it is necessary to have the objective function (and constraints) in terms of relationships that are valid for all ranges of possible values. The shaded area of the figure is represented by the inequality constraints  $Y \geq X$  and  $Y \geq -X$ . These two inequalities can actually be represented by the single inequality  $Y \geq |X|$ . Now this is not exactly the desired function, since it corresponds to the entire region above the curve  $Y = |X|$ . However, since the objective function is to minimize  $Y$ , the solution procedure would converge to a point on the curve. Thus, the two inequality constraints  $Y \geq X$  and  $Y \geq -X$  can be used to replace  $Y = |X|$  in the LP formulation.

Returning now to the BP optimization problem,  $d_i$  can now be re-written as follows:

$$d_i \geq [ (P_i^0 - P_i^*) + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} \Delta BP_k ]$$

and,

$$d_i \geq - [ (P_i^0 - P_i^*) + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} \Delta BP_k ]$$

The peaking constraint can be re-written making use of the definition of

the nodal power,  $P_i$

$$P_i^0 + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} \Delta BP_k \leq P_{lim}$$

Finally, by replacing the search variable  $\Delta BP_k$  with

$$\Delta BP_k = BP_k^n - BP_k^{n-1},$$

the system of equations can be re-written in their final form (after some re-arranging) as

$$f = \max \left( - \sum_{i=1}^N d_i \right)$$

$$\sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} BP_k^n - d_i^n \leq (P_i^* - P_i^0) + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} BP_k^{n-1} \quad \text{for } i = 1, N$$

$$- \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} BP_k^n - d_i^n \leq - (P_i^* - P_i^0) - \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} BP_k^{n-1} \quad \text{for } i = 1, N$$

$$\sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} BP_k^n \leq (P_{lim} - P_i^0) + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} BP_k^{n-1} \quad \text{for } i = 1, N$$

$$BP_k \geq 0 \quad \text{for } k = 1, K.$$

The linear programming technique used to solve the above problem is the revised simplex method.<sup>4</sup> This method is currently employed by all commercial LP computer codes. The particular code used in this work is



the IMSL library subroutine ZX3LP.

### 2.3 Hi-Lo Algorithm

The solution of the preceding mathematical programming problem prescribes the desired BP loading. However, this desired BP loading is not physically realizable and thus not of great value to the designer. A Hi-Lo algorithm is utilized to determine a practical BP loading closest to the optimal BP loading.

The first task performed by the Hi-Lo algorithm is the determination of the high and low values for the number of BP rods to be placed in each poisonable assembly. Since current design allows BP's to be used in multiples of four only, this corresponds to finding the two multiples of four that bracket the desired number of BP's. For example, if the desired number of BP's is 7.2 for a certain position, the high and low values would be four and eight, respectively.

The next step is the formation of all possible combinations of these high and low values at each BP position. Since this is simply a binary decision (high or low) at each BP position, the total number of these combinations would be  $2^K$ , where K is the total number of BP positions. The identification of all possible combinations is facilitated by the above fact. To generate these combinations, all one has to do is convert the  $2^K$  numbers from 0 to  $2^K-1$  from decimal to binary. This results in a K digit number consisting of nothing but 0's and 1's. In

this representation, a 1 in the i'th position in the binary number corresponds to the placement of the high number of BP's in the i'th core BP position. Similarly, a 0 corresponds to the low number of BP's. Given below is an example of the transformation from decimal to binary to BP loading representation.

Decimal	Binary	BP Loading
100	1100100	Hi Hi Lo Lo Hi Lo Lo

The number of BP combinations formed is expanded when the desired number of BP's for any position is very close to an available loading. For example, if the desired number of BP's is 7.9 for a particular position, then loadings of 4, 8 and 12 BP's in that position are investigated. These additional combinations are formed utilizing a variation of the binary transformation routine described above. Exactly how close the desired number of BP's and the available number of BP's must be is specified by the user. The tolerance used for all of this work and recommended for any subsequent analysis is 0.5 (one half of one BP rod). This tolerance may, however, be adjusted to values ranging from 0 to 2 in order to increase or decrease the number of BP loadings examined. In the instance where no acceptable BP loading is identified, the tolerance should definitely be relaxed. To the other extreme, if an unwieldy number of BP combinations are created, the tolerance may be tightened. The total number of BP combinations is given by

$$N = \sum_{l=0}^n \binom{n}{l} 2^{K-l}$$

where,

$\binom{n}{l}$  = number of combinations of n items taken l at a time

K = total number of BP positions in the core

n = number of positions in the core where 'extra' BP loadings are considered.

Utilizing basic statistical theory for the determination of  $\binom{n}{l}$ , the above equation can be re-written as

$$N = \sum_{l=0}^n \frac{n!}{l!(n-l)!} 2^{K-l}$$

If there are no 'extra' positions (n=0), the previous relationship reduces to  $2^K$ , which is the base number of combinations discussed earlier in the section. On the other hand, if every BP position is an 'extra' position (n=K), then the number of combinations swells to  $3^K$ . A graphical representation of this relationship is given by Figure 4. The upper limit shown in the figure is simply due to the declared array size in the coding itself. Realistically, this upper limit would probably never be approached due to the relatively high cost of depleting such a large number of cases.

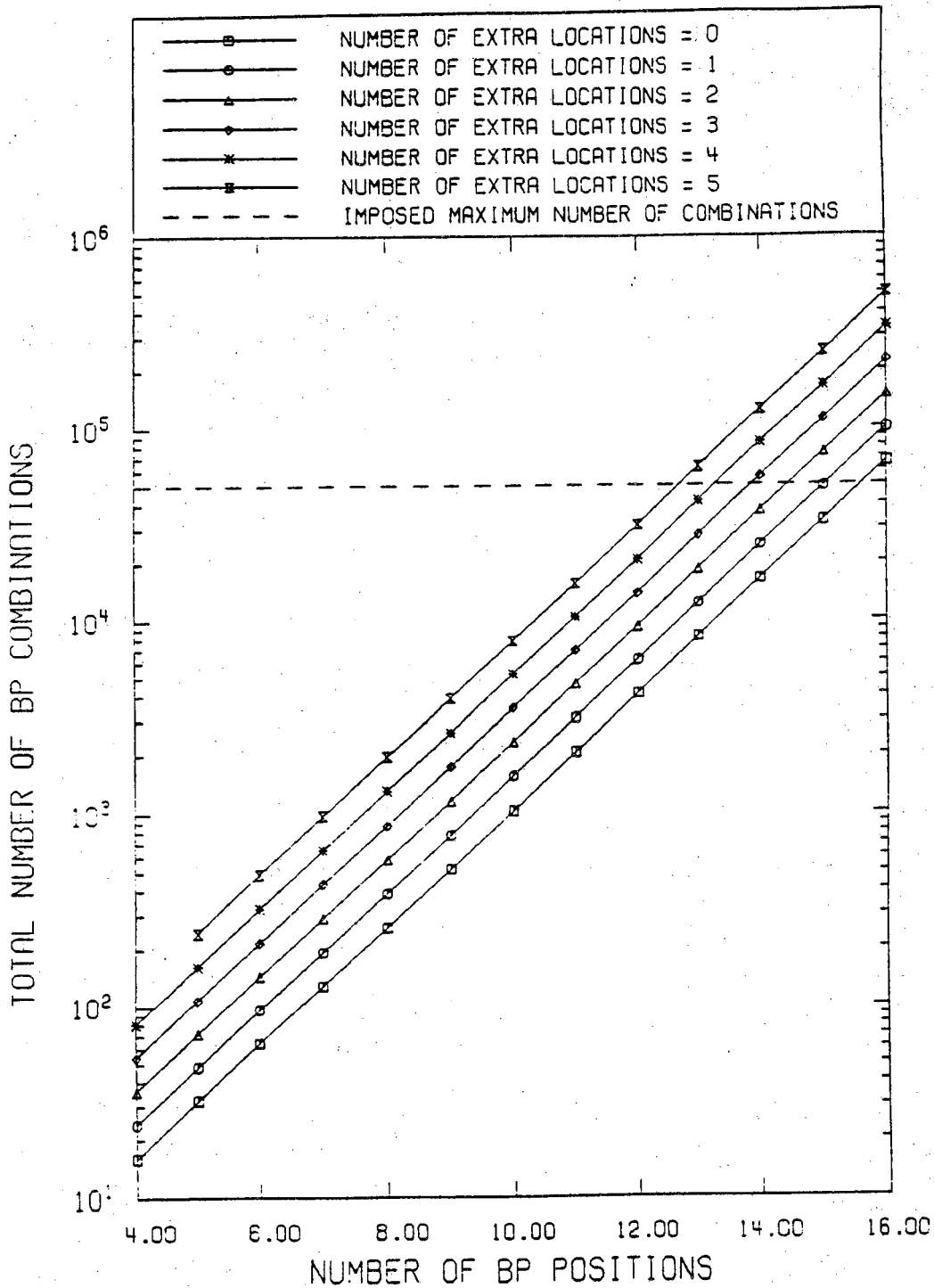


Figure 4. Number of BP Combinations Formed by Hi-Lo Algorithm

## 2.4 Power Reconstruction and Case Pruning

The first order perturbation approximation discussed in earlier sections was fundamental to the solution of the optimization problem using a linear programming technique. In the solution procedure, a matrix of linear sensitivity coefficients,  $\frac{\partial P_i}{\partial BP_k}$ , are calculated numerically.

This sensitivity matrix is used to approximate the nodal power distribution, as the BP loading is iteratively modified. As an additional benefit this matrix may also be used to estimate the power distribution resulting from each of the candidate BP loadings identified by the Hi-Lo algorithm. Similarly to the relationship given in section 2.2, the nodal power is given by

$$P_i^r(m) = P_i' + \sum_{k=1}^K \frac{\partial P_i}{\partial BP_k} (BP_k(m) - BP_k')$$

where,

$P_i^r(m)$  = reconstructed power in node i with BP combination m

$P_i'$  = relative power in node i with optimal BP loading

$BP_k'$  = optimal BP loading for node k

$BP_k(m)$  = BP loading in node k in combination m

$K$  = total number of BP positions in the core.

Thus, without a time consuming normal power calculation, a relatively accurate representation of the nodal power distribution may be obtained. A comparison of the reconstructed power and the SIMULATE calculated

power for a particular BP loading is given in Figure 5. Of particular note is the relatively small average absolute difference of 0.9%.

In order to minimize computation time and thus cost, an excessive number of nodal power calculations should be avoided. This can be accomplished by eliminating infeasible cases before the normal power calculations are performed. The current method incorporates one such pruning procedure in its calculational scheme. The pruning parameter currently used is the nodal power peaking. If the maximum power peaking exceeds a preset limit, the case is discarded and the following case is examined.

The preset peaking limit is based on two multiplicative factors: a base nodal power peaking limit,  $PP_{lim}$ , and a power peaking tolerance multiplier, PPMULT1. The user input base limit,  $PP_{lim}$ , is derived from thermal hydraulic safety considerations and previous core operational data. For this work, a fairly typical value of 1.33 was implemented. The peaking tolerance multiplier is applied to relax the peaking limitations on the reconstructed power distributions. The motivation for this relaxation is the error introduced by the linear approximation made in the power reconstruction. Based on power distribution comparisons for several different burnable poison loadings such as the one given in Figure 5, the overall average percent difference is found to be approximately 0.5% with a standard deviation of about 0.5%. From basic statistical theory, it is 99.7% certain that all data points are

0.953							X.XXX	<- Reference Power
0.982							X.XXX	<- SIMULATE Power
+3.0							+X.X	<- % Difference
							+0.9	<- AVG % Diff
1.110	1.249							
1.139	1.278							
+2.6	+2.3							
1.020	1.227	1.034						
1.040	1.249	1.045						
+2.0	+1.8	+1.1						
1.217	1.338	1.174	1.143					
1.228	1.349	1.173	1.140					
+0.9	+0.8	-0.1	-0.3					
0.868	1.276	1.181	1.285	1.129				
0.872	1.279	1.178	1.278	1.118				
+0.5	+0.2	-0.3	-0.5	-0.9				
0.754	1.077	1.322	1.089	1.255	1.006			
0.752	1.074	1.311	1.081	1.242	0.994			
-0.3	-0.3	-0.8	-0.7	-1.0	-1.2			
0.815	1.194	1.264	1.179	0.858	0.345			
0.810	1.187	1.257	1.179	0.856	0.341			
-0.6	-0.6	-0.6	+0.0	-0.2	-1.2			
0.301	0.387	0.422	0.311					
0.299	0.384	0.419	0.309					
-0.7	-0.8	-0.7	-0.6					

Figure 5. BOC Power Distribution Comparison Cycle 9 - Case 2



within range of  $\bar{x} + 3\sigma_x$ . Thus, all reconstructed nodal powers should be within  $0.5 + 3(0.5)$  or 2% of the actual nodal power calculated by SIMULATE. The peak powers ( $P_i \geq 1.3$ ) that are of most interest here were predicted slightly more accurately, with all reconstructed peak powers expected to be within 1% of the SIMULATE powers.

It follows, then, that a multiplier of 1.02 would effectively account for the uncertainty in the reconstructed power distributions. This multiplier is also user-input, and may be altered if necessary, but this value is recommended for any subsequent work.

This pruning procedure provides a significant reduction in the number of cases which must be evaluated using the normal power calculations. Equally important, however, is the fact that while many cases are discarded, no potentially successful cases are discarded. The relatively loose tolerance insures that all BP loadings that are even close to being acceptable are passed on to the next step in the selection process.

## 2.5 Case Depletion and Optimal Selection

The final step in the core design process is the depletion of all the cases that passed the pruning test described in the previous section. This is by far the most time consuming portion of the code, involving numerous power distribution calculations for each depletion. The optimum BP loading is selected after all cases have been depleted to EOC.

For a particular loading pattern to be acceptable, it must meet power peaking constraints throughout the cycle. Thus, in the core design process, the core power peaking must be checked at each step in the depletion. In addition, the burnup steps must be small enough to insure that power peaking violations do not occur between steps. This checking process is included in the code's depletion procedure. If, at any point in the depletion, the power peaking limits are exceeded, the depletion is terminated and depletion of the succeeding case is begun.

Just as in the power reconstruction pruning, the peaking constraint in the core depletion also includes a power peaking tolerance in the form of a multiplier. In this case, the multiplier, PPMULT2, accounts for the error in the SIMULATE power distribution calculations as compared to the actual power distribution. The quantitative evaluation of this error is beyond the scope of this work. For this reason and basic

conservatism, a multiplier of 1.00 is used here and recommended in all subsequent work. As it applies here, conservatism means that while some potentially successful cases are discarded, the cases that are deemed acceptable by the code are much more likely to satisfy the requirements for loading into the reactor.

After all candidate BP loadings have been depleted, it remains only to select the optimum loading. This selection is based on the maximization of the cycle length. Thus, of those cases that deplete to EOC, the one yielding the longest cycle is selected as the optimum.

### 3. APPLICATION TO CORE RELOAD

#### 3.1 Introduction

The improved core reload design method was employed for a re-design of Commonwealth Edison's Zion Unit-1 cycles 9 and 10. The actual loading patterns for both cycles are given in Figures 6 and 7, respectively. Note: The fuel type descriptor given for each assembly is simply a variable used internally by the program to differentiate between assemblies of different design and/or fuel enrichments. The re-design is done utilizing the original fuel assemblies (i.e. same enrichments) ; the difference between the new and the old designs is the loading pattern and the BP loading. Thus, the result of the optimization is the improvement in the cycle length for a given fuel loading. Conversely, the code could also be used to give a desired cycle length using a lower reload enrichment.

The nodal code SIMULATE has been benchmarked for both Zion-1 cycles 9 and 10, and has been shown to yield sufficiently accurate results for design work of this type. All current work has been performed utilizing the core model formulated through benchmarking procedure from previous work.<sup>1</sup> In addition, all core calculations are performed assuming that equilibrium Xenon and Samarium concentrations are present in the core. This approximation is valid at all burnup steps except BOC, at which point these two fission products have not yet reached their saturation



<table border="1"> <tr><td>3</td></tr> <tr><td>2.90</td></tr> <tr><td>27.530</td></tr> </table>		3	2.90	27.530	<table border="1"> <tr><td>X</td></tr> <tr><td>X.XX</td></tr> <tr><td>XX</td></tr> <tr><td>XX.XXX</td></tr> </table>		X	X.XX	XX	XX.XXX	<p>&lt;- Fuel Type          &lt;- Enrichment          &lt;- # of BP          &lt;- BOC Burnup</p>																		
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2	4	1	4	1																									
2.80	3.20	2.70	3.20	2.70																									
			12																										
10.400	13.060	23.160	0.000	25.100																									
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1	4	5	2																										
2.70	3.20	3.60	2.80																										
20.730	10.200	0.000	15.200																										

Figure 7. Reference Loading Pattern for Cycle 10

level. However, at BOC, the reactor is just beginning its ascent to full power. By a cycle burnup of 150 MWD/MT, the reactor has reached full power and the fission product inventories have attained their equilibrium values. In addition, at this small core burnup, the burnable poison concentration has not changed appreciably from its BOC value. Thus, performing the core design calculations at BOC while assuming equilibrium Xenon and Samarium is equivalent to designing the core reload at a cycle burnup of 150 MWD/MT. Since it is at just this burnup step that almost all core benchmarking is performed, this seems to be an acceptable approximation. The only possible drawback with this method is the lack of a representation of core behavior at the actual beginning-of-cycle conditions. Power peaking at the actual BOC should not pose a problem, however, since the reactor is operating at such a low power level.

The optimization process consists of two separate stages: the fuel loading search and the burnable poison search. Since only the latter has been altered in this work, only the BP assignment results will be discussed in detail. The results of the fuel loading search will be given for completeness. A more detailed analysis of the fuel loading optimization can be found in reference 1.

### 3.2 Optimization Results

In order to achieve the longest cycle for a given fuel loading, it is necessary to move as much of the fresh fuel as possible to the core interior. This both decreases neutron leakage by placing less reactive fuel on the core periphery and increases the worth of the fresh fuel by placing those assemblies in areas of 'high neutron importance'. The limiting factor in how much fresh fuel can be moved inboard is the imposed power peaking limit. This peaking can be partially controlled through the use of a proper burnable poison loading.

Figures 6 and 7 illustrate the loading patterns that were actually used in cycles 9 and 10, respectively. In both cycles, a fairly large percentage of the fresh fuel has been placed in the 56 core peripheral positions. In the cycle 9 loading, 28 of the 68 fresh fuel assemblies are on the periphery, while cycle 10 has 20 out of 60. In addition, the cycle 9 loading has 8 once-burned assemblies on the periphery and cycle 10 has 24. Clearly, there is a large amount of highly reactive fuel on the core periphery. Thus, there seems to be a great deal of room for improvement in the achievable cycle length for both cycles.

This presumption is supported by the results from the direct search procedure. By moving much of the reactive fuel away from the periphery, the shuffling procedure obtained an improvement in cycle length of over



1000 MWD/MT. The improved fuel loading patterns are given in Figures 8 and 9. In striking contrast to the actual loadings, the improved cycle 9 loading has 36 twice-burned fuel assemblies on the periphery and cycle 10 now has 44. This yields improvements in both cycle length and pressure vessel neutron fluence. The latter consideration has recently received a considerable amount of attention.

The next step in the optimal design process is the determination of a BP loading capable of controlling the power peaking in the new fuel loadings. Since so much highly reactive fuel was moved inboard, power peaking control becomes much more difficult. The direct search method did, however, apply a power peaking limit to the fuel loading optimization, so a feasible control strategy should exist. In fact, in the the cycle 9 BP search, 31 different loadings were found to be acceptable. Cycle 10, on the other hand, yielded only 2 acceptable BP loadings. The reason for this apparent discrepancy is that the cycle 9 optimal fuel loading had 10 poisonable positions in the octant while cycle 10 had only 9. Thus, many more candidate BP loadings were identified for cycle 9 than for cycle 10 (2304 compared to 1152). The optimal BP loading was selected from these acceptable loadings for both cycles. The optimal BP loadings determined by the code for cycles 9 and 10 are shown in Figures 10 and 11, respectively.



3 2.90 27.530		X X.XX XX.XXX		← Fuel Type ← Enrichment ← BOC Burnup	
2 2.80 13.980		5 3.60 0.000			
1 2.70 20.730	2 2.80 15.200	2 2.80 14.670			
4 3.20 0.000	4 3.20 17.960	4 3.20 0.000	1 2.70 24.600		
2 2.80 10.400	4 3.20 13.060	4 3.20 14.320	4 3.20 0.000	1 2.70 24.600	
4 3.20 14.990	4 3.20 10.110	2 2.80 14.950	4 3.20 10.200	5 3.60 0.000	4 3.20 0.000
4 3.20 14.990	5 3.60 0.000	5 3.60 0.000	4 3.20 0.000	4 3.20 21.570	1 2.70 25.060
6 2.80 27.200	4 3.20 22.080	4 3.20 25.290	1 2.70 23.160		

Figure 9. Optimal Fuel Loading Pattern for Cycle 10



3 2.90 27.530		<table border="1"> <tr> <td>X</td> <td>&lt;- Fuel Type</td> </tr> <tr> <td>X.XX</td> <td>&lt;- Enrichment</td> </tr> <tr> <td>XX</td> <td>&lt;- # of BP</td> </tr> <tr> <td>XX.XXX</td> <td>&lt;- BOC Burnup</td> </tr> </table>		X	<- Fuel Type	X.XX	<- Enrichment	XX	<- # of BP	XX.XXX	<- BOC Burnup
X	<- Fuel Type										
X.XX	<- Enrichment										
XX	<- # of BP										
XX.XXX	<- BOC Burnup										
2 2.80 13.980	5 3.60 4 0.000										
1 2.70 20.730	2 2.80 15.200	2 2.80 14.670									
4 3.20 4 0.000	4 3.20 17.960	4 3.20 8 0.000	1 2.70 24.600								
2 2.80 10.400	4 3.20 13.060	4 3.20 14.320	4 3.20 8 0.000	1 2.70 24.600							
4 3.20 14.990	4 3.20 10.110	2 2.80 14.950	4 3.20 10.200	5 3.60 4 0.000	4 3.20 0.000						
4 3.20 14.990	5 3.60 8 0.000	5 3.60 4 0.000	4 3.20 0.000	4 3.20 21.570	1 2.70 25.060						
6 2.80 27.200	4 3.20 22.080	4 3.20 25.290	1 2.70 23.160								

Figure 11. Optimal Design Pattern for Cycle 10

Comparisons of three core loading alternatives: the reference (actual) design, the previous 'optimal' design, and the improved design are given in Tables 2 and 3. The superiority of both of the optimal designs as compared to the reference case is apparent. The length of cycle 9 was improved by over 1000 MWD/MT while cycle 10 saw an increase of over 1500 MWD/MT. Also, the number of BP rods necessary for peaking control was decreased. This decrease was quite drastic in the cycle 10 design, where the number of BP's was more than cut in half. The focus of this work, however, is the improvement of the optimization procedure, which is evident by comparison of the new and old 'optimal' results.

The new optimal design of cycle 9 gave a slight improvement in both the achievable cycle length and the number of BP's required as compared to the old optimal design. The cycle was lengthened by about 15 MWD/MT while using 32 fewer BP rods. This improvement also coincided with a small reduction in the cycle maximum power peaking. In fact, the old optimal design would not have been considered acceptable by the present design method due to peaking violations.

At first it may appear that the old optimal design for cycle 10 is superior to the newer design in that it results in an identical cycle length while using slightly fewer BP rods. However, while the old design does give an equal cycle length, the power peaking is significantly higher than the new optimal design.

**TABLE 2. Comparison of Loading Patterns for Cycle 9**

	<b>Loading Pattern</b>		
	<b>Reference</b>	<b>OLD</b>	<b>NEW</b>
<b>Cycle Length (GWD/MT)</b>	11.944	12.964	12.980
<b>Total Number of BP's</b>	496	480	448
<b>Max. Power Peaking</b>	1.3146	1.3342	1.3253

TABLE 3. Comparison of Loading Patterns for Cycle 10

	Loading Pattern		
	Reference	OLD	NEW
Cycle Length (GWD/MT)	11.638	13.241	13.241
Total Number of BP's	544	256	288
Max. Power Peaking	1.3420	1.3457	1.3295



A more detailed comparison of the power peaking for both cycles is given in Tables 4 and 5. The cycle 9 results show that the only peaking violation of the old optimal design occurs at the beginning-of-cycle. As was mentioned earlier, peaking is not an important consideration at BOC. This, coupled with the fact that the power peaking does not greatly exceed the limit, would lead to the conclusion that the old optimal design would, in fact, be acceptable. Even so, the peaking in the new design is superior to both the old optimal design and even the reference design.

From the cycle 10 results, it appears that the old optimal design exceeded power peaking limits for nearly half of the cycle. Clearly, this is an unacceptable design. Even the reference design violates the peaking limit of 1.33, although this violation occurs at BOC (similar to the old optimal design of cycle 9). Just as in the cycle 9 design, the new optimal design for cycle 10 is far superior to the old optimal design, with a lower power peaking factor at nearly every burnup step.

TABLE 4. Maximum Power Peaking - Cycle 9

Cycle Burnup (GWD/MT)	Loading Pattern		
	Reference	OLD	NEW
0.0	1.314	1.334	1.309
1.0	1.315	1.309	1.291
2.0	1.309	1.299	1.284
3.0	1.300	1.294	1.297
4.0	1.283	1.296	1.301
5.0	1.283	1.300	1.309
6.0	1.293	1.321	1.317
8.0	1.284	1.318	1.325
10.0	1.296	1.305	1.314

TABLE 5. Maximum Power Peaking - Cycle 10

Cycle Burnup (GWD/MT)	Loading Pattern		
	Reference	OLD	NEW
0.0	1.342	1.346	1.330
1.0	1.297	1.341	1.323
2.0	1.273	1.336	1.325
3.0	1.256	1.332	1.324
4.0	1.249	1.325	1.319
5.0	1.251	1.320	1.313
6.0	1.254	1.316	1.315
8.0	1.258	1.310	1.312
10.0	1.257	1.300	1.305

#### 4. CONCLUSIONS AND RECOMMENDATIONS

##### Conclusions

With the completion of the present work, there now exists a completely automated method for optimizing the burnable poison loading for a core reload design pattern. It should be stressed that the method is capable of performing reload design only. The code is not designed to be used for startup cycle design. This is due to the fact that in the startup core, all of the assemblies are 'fresh' and available for BP loading. The code is simply not capable of handling such a large number of poisonable assemblies.

The improved method will save the core designer a great deal of time and effort. Perhaps even more significantly, the design work can now be performed by individuals not having a tremendous amount of insight or experience in the core reload design problem. This makes the code ideal for utilities desiring to become vendor-independent in the designing of their own reload cores.

The major benefit of the improved BP assignment method is not in lengthening the cycle, for it is becoming more and more evident that the BP loading has very little effect on the achievable cycle length. Its

greatest worth is in determining a burnable poison loading that makes possible the safe operation of the optimal fuel loading pattern. Also, the improved method investigates a far greater number of alternative loading strategies and therefore increases the likelihood that the safest and most economical scenario is chosen.

### Recommendations

It has been shown in previous work<sup>1</sup> that the control strategy has but a minor effect on the achievable cycle length. This theory is further supported by the current work. The difference between the longest and the shortest cycles that were deemed as acceptable by the code is less than 40 MWD/MT for both cycles 9 and 10. This would seem to lead to the conclusion that the selection of the optimal BP loading on the basis of cycle length is not such a viable alternative. Other possible alternatives for the objective function include: the minimization of the total number of BP rods and the minimization of the power peaking. The minimization of the total number of BP rods in the core would obviously lower fuel cycle costs by reducing the cost of fuel fabrication. This would also lead to an lessening of the reactivity penalty due to BP residue. The minimization of the core power peaking gives no explicit benefit in terms of cycle length or fuel costs. It does, however, give rise to a larger margin of safety, which could ultimately outweigh either of the above benefits.

The entire foundation of the burnable absorber assignment method is based on the attempt to achieve some desired power distribution. Theoretically, this target power shape is the optimum power shape. Since this optimum is currently unknown, a Haling power distribution was used as a target in the present method. Unfortunately, however, this is most probably not the optimum power shape. The method would be improved

if the true optimum power distribution could be determined. This task is currently being undertaken here at Purdue.

The current method provides an adequate tool for optimized reload design utilizing discrete burnable absorber rods. This is by far the most common means of power peaking control in today's reactors. However, the inherent advantages of other types of absorbers, such as gadolinia and IFBA may soon lead to the replacement of the discrete absorber rod. Therefore, it would be advisable to examine the possibility of incorporating the use of these other types of absorbers into the core reload design method.

**LIST OF REFERENCES**

1. Kim, Y./J., 'Optimization of Core Reload Design for Low-Leakage Fuel Management in Pressurized Water Reactors,' Ph. D. Thesis, Purdue University, 1986.
2. Ver Planck, D., 'SIMULATE-E: A Nodal Core Analysis Program for LWR's; Computer Code Users's Manual,' EPRI-NP-2792-CCM, 1983.
3. Haling, R., 'Operating Strategy for Maintaining An Optimum Power Distribution Throughout Life,' Proc. ANS Topl. Mtg.; Nuclear Performance Power Reactor Cores, TID-7672, 1964.
4. Reklaitis, G. V. et. al., 'Engineering Optimization: Methods and Applications,' John Wiley and Sons, New York, 1983.



**APPENDICES**

**APPENDIX A: SOURCE LISTING**

The following partial source listing gives all of the major routines that were either created or altered in this work. Excluded are routines with minor or cosmetic changes.

00001		SUBROUTINE PARTB	PARTB	2
00004		CHARACTER*8 ADATE, BDATE, IHTIME, JHTIME	F200CPU	3
00011		DIMENSION JOBNAM(2)	BLOK4	6
0017		LOGICAL IUOID, ICNTRL, IFIRST, LEVELU, LEVELC		
00018		LOGICAL DEPL(50000), KPFRST	OPUS	20
00019		DIMENSION IBP(50000,20), IROW(20), JCOL(20), EXP(8,8), CONC(6,8,8)	OPUS	21
00020		TEMP=0.0	OPUS	22
00021		ITEMP=0	OPUS	23
00022		NBPERM=1	OPUS	24
00023		DEPL(1)=.TRUE.	OPUS	25
00024		MOPTIM=A1(46)	OPUS	26
00025		IMABWR=MODEB(2)	PARTB	15
00026		IDEBUG=0	PARTB	16
00027		IFIRST=.TRUE.	PARTB	17
00028		CALL QSSNDMDF('LEN=',16,'MSG=', '** IN PARTB **')	SMCY1	102
00029		CALL TIMER (SHPARTB,2)	PARTB	19
00030		LAUSED(2)=LAUSED(1)	PARTB	21
			PARTB	22
			PARTB	23
			PARTB	24
			PARTB	25
			PARTB	26
			PARTB	27
			PARTB	28
			PARTB	29
			PARTB	30
			PARTB	31
			PARTB	32
			PARTB	33
			PARTB	34
			PARTB	35
			PARTB	36
			PARTB	37
			PARTB	38
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			PARTB	46
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			PARTB	48
			PARTB	49
			PARTB	50
			PARTB	51
			PARTB	52
			PARTB	53
			PARTB	54
			PARTB	55
			PARTB	56
			PARTB	57
			PARTB	58
			PARTB	59
			PARTB	60
			PARTB	61
			OPUS	27
			XXX2	1
			PARTB	63
			PARTB	64
			PARTB	65
			PARTB	66
			PARTB	67
			XXX2	2
			XXX2	3
			PARTB	68
			OPUS	28
			OPUS	29
			XXX2	4
			XXX2	5
			OPUS	30
			OPUS	31
			OPUS	32
			OPUS	33
			OPUS	34
			OPUS	35
			XXX2	8
			XXX2	9
			XXX2	10

```

00080      * NFLOW,NPOWER,NSOURC,NUOID,XPOLST,XT,IFIRST,
00081      * KBPOS,NBPERM,IBP,DEPL,IROW,JCOL
00082      IF(IDEBUG.GT.0) GO TO 660
00083      IF(IBPAS.EQ.0) THEN
00084          NU=0
00085          NC=0
00086          GO TO 30
00087      ELSE
00088      ENDIF
00089      CONTINUE
00090      NODEP=0
00091      NODEP1=0
00092      IF(MOPTIM.EQ.2) THEN
00093          WRITE(6,50)
00094          FORMAT(1H1,1X,62H PERFORM DEPLETION OF ALL ACCEPTABLE INTEGRAL BP
00095          *COMBINATIONS:./,3X,61(1H*))
00096      ENDIF
00097      DO 665 KPOIS=1,NBPERM
00098          PTEMP=0.0
00099          XTEMP=0.0
00100          IPEAK=0
00101          JPEAK=0
00102          IF(.NOT. DEPL(KPOIS)) THEN
00103              NODEP=NODEP+1
00104              GO TO 665
00105          ENDIF
00106
00107      C
00108      C
00109      C
00110      C
00111      C
00112      C
00113      C
00114      C
00115      C
00116      C
00117      C
00118      C
00119      C
00120      C
00121      C
00122      C
00123      C
00124      C
00125      C
00126      C
00127      C
00128      C
00129      C
00130      C
00131      C
00132      C
00133      C
00134      C
00135      C
00136      C
00137      C
00138      C
00139      C
00140      C
00141      C
00142      C
00143      C
00144      C
00145      C
00146      C
00147      C
INITIALIZATION
IEOL=0
DEP=DE
TIMXE=0.0
NEC=0
ISTEP=0
FIRST=1.0
IBRN=IBURM
DETERMINE THE ITERATION LEVELS INVOLVED IN THIS CASE.
IUOID=.FALSE.
IF (NUMAX.GT.0) IUOID=.TRUE.
ICNTRL=.FALSE.
IF(MODEB(5).GE.3) ICNTRL=.TRUE.
IF(MODEB(5).GT.5) ICNTRL=.FALSE.
THE HALTING ITERATION USES THE CONTROL ITERATION LEVEL DURING
THE END OF LIFE EXPOSURE SEARCH.
IF (MODEB(5).EQ.6.AND.MODEB(10).GT.0) ICNTRL=.TRUE.
NC=0
NU=0
NS=0
XL=1.0
ITERS=0
ISRCH=0
LPMEM = NPARTA
LNEED=MMACRO*NDIM3D
IF (MODEB(1).EQ.0) LNEED=NDIM3D
PUT ONLY THE SOURCE TO POWER CONVERSION FACTOR ON I/O FOR 1 GROUP
CALL NEXTPT (LPMEM,10,LNEED,20HCROSS SECTION )
LNEED = NSTATE*NDIM3D
CALL NEXTPT(LPMEM,18,LNEED,15HSTATE VARIABLES)
CALL NEXTPT(LPMEM,24,NDIM3D,20HTHERMAL LEAKAGE )
LPW=LPMEM
IF(MOPTIM.EQ.2) THEN
XPO=0.0
A1(1)=0.0
XT=0.0
XPOLST=0.0
NTOTBP=0
DO 1 K=1,KBPOS
I=IROW(K)
J=JCOL(K)
IF(J.EQ.1) THEN
IF(I.EQ.J) THEN
NYV=1
ELSE
NYV=4
ENDIF
ELSEIF(I.EQ.J) THEN
NYV=4
ELSE

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OPUS 106
OPUS 107

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00148		NYV=8	OPUS	108
00149		ENDIF	OPUS	109
00150		NBP(I,J)=IBP(KPOIS,K)*10+(NBP(I,J)/1000)*1000	OPUS	110
00151		NTOTBP=NTOTBP+NYV*IBP(KPOIS,K)	OPUS	111
00152	1	CONTINUE	OPUS	112
00153		ENDIF	OPUS	113
00154	20	CONTINUE	OPUS	114
	C	CONTROL ITERATION STARTS HERE	OPUS	115
00155		IOSC=0	OPUS	116
00156		XLC=XKEFF	OPUS	117
00157		LEVELC=.FALSE.	OPUS	118
00158		NC=NC+1	OPUS	119
00159		CALL POWDIS(DEP, ICNTRL, IDEBUG, IOPTX, IOSC, ISRCH, IUOID,	OPUS	120
		* LEVELC, LEVELU, LPAXLS, LPBU, NCOEFS, NCROSS, NEDITS, NEDTS,	OPUS	121
		* NEXPOS, NFLOW, NPOWER, NSOURC, NUOID, XPOLST, XT, IFIRST, KPFRST)	OPUS	122
00160		IF(MOPTIM .EQ. 2) THEN	OPUS	123
00161		PPLIM=A1(39)	OPUS	124
00162		PPMULT2=A2(52)	OPUS	125
00163		IF(SMAX .GT. PPLIM*PPMULT2) THEN	OPUS	126
00164		NODEP1=NODEP1+1	OPUS	127
00165		KPFRST=.FALSE.	OPUS	128
00166		GO TO 665	OPUS	129
00167		ENDIF	OPUS	130
00168		IF(SMAX .GT. PTEMP) THEN	OPUS	131
00169		PTEMP=SMAX	OPUS	132
00170		IPEAK=MAXSI	OPUS	133
00171		JPEAK=MAXSJ	OPUS	134
00172		XTEMP=XPO	OPUS	135
00173		ENDIF	OPUS	136
00174		ENDIF	OPUS	137
	C		PARTB	503
	C	END BWR LIMITS CALCULATION	PARTB	504
	C		PARTB	505
	C		PARTB	506
	C		PARTB	507
	C		PARTB	508
	C		OPUS	138
00175		IF(MOPTIM .LT. 2) THEN	PARTB	509
00176		CALL EDT1 (8,MZ,MZ,MZ,MZ,MZ)	PARTB	510
00177		IF(MODEB(5).EQ.1.OR.MODEB(5).EQ.2) CALL HTBAL(2)	PARTB	511
00178		CALL PAGES(15+KMAX)	PARTB	512
00179		WRITE (ITAPOT,300)	PARTB	513
00180		WRITE (ITAPN,300)	PARTB	514
00181	300	FORMAT (/54X,27HAVERAGE AXIAL DISTRIBUTIONS,/,	PARTB	515
		1 11X,4HNODE,8X,5HPPOWER,8X,5HSHWATER,7X,8H SOURCE ,7X,7HTHERMAL,7X,	PARTB	516
		2 7HCONTROL,11X,1HK,10X,6HBYPASS/	PARTB	517
		3 35X,7HDENSITY,21X,7HLEAKAGE,20X,8HINFINITY,9X,4HNUOID/	PARTB	518
		4 25X,1HP,13X,1HU,13X,1HS,13X,1HT,12X,2HCT,13X,1HK,12X,2HBU/)	PARTB	519
00182		CALL AXIALP (A(LPAXLS),KD,7,22H(12X,12,7F14.4,12X,12),ITAPOT)	PARTB	520
00183		CALL AXIALP (A(LPAXLS),KD,7,22H(12X,12,7F14.4,12X,12),ITAPN)	OPUS	139
00184		ENDIF	PARTB	521
	C		PARTB	522
	C	APPLY THE CONTROL SEARCH CALLED FOR BY MODEB(5)	PARTB	523
	C		PARTB	524
00185		NU=0	PARTB	525
00186		IF (.NOT.ICNTRL) GO TO 390	PARTB	526
00187		IF (NC.GE.NCMAX) GO TO 390	PARTB	527
00188		IF (NC.LT.2) GO TO 310	PARTB	528
00189		IF (LEVELC) GO TO 390	PARTB	529
00190	310	CONTINUE	PARTB	530
00191		M=MODEB(5)-3	PARTB	531
00192		GO TO (340,350,370),M	SMCY1	113
00193	340	CALL NHWORD(NAM,8HSEQUENCE,8)	PARTB	533
00194		GO TO 360	SMCY1	114
00195	350	CALL NHWORD(NAM,8HPATTERN,8)	PARTB	535
00196	360	ANOTWT=NOTWT	SMCY1	115
00197		CALL NHWORD(NAM(13),8HPOSITION,8)	SMCY1	116
00198		CALL SEARCH(SR(4),ANOTWT,NAM,NAM(13),XKEFF,XLMBDA,ISRCH)	PARTB	538
00199		NOTWT=ANOTWT	PARTB	539
00200		IF (NOTWT.LE.0) NOTWT=0	PARTB	540
00201		CALL RODMUU (A(LPR),A(LPMPCR),A(LPARAY),A(LPNTCH),ID,JD,IRMX,LIMRA	PARTB	541
		1Y,LIMPAS)	PARTB	542
00202		GO TO 380	PARTB	543
00203	370	CONTINUE	PARTB	544
	C	HALING SEARCH (SEE SIGCAL/SIGDAT)	PARTB	545
00204		IF (MODEB(10).EQ.0) GO TO 380	SMCY1	117
00205		CALL NHWORD(NAM,8HHALING D,8)	SMCY1	118
00206		CALL NHWORD(NAM(13),8HEXPOSURE,8)	SMCY1	119
00207		CALL SEARCH(SR(5),DEP,NAM,NAM(13),XKEFF,XLMBDA,ISRCH)	PARTB	549
00208	380	CONTINUE	PARTB	550
00209		GO TO 30		

00210	390	CONTINUE	PARTB	551
00211		NC=0	PARTB	552
	C		PARTB	553
	C	END OF CONTROL LEVEL ITERATION	PARTB	554
	C		PARTB	555
00212		LPSTAT=NFLOW	PARTB	556
00213		LPMEM = LPSTAT + NSTATE*JD*KD	PARTB	557
00214		LNEED = LIST*NDAT*LIMNFT	PARTB	558
00215		CALL FINDPT(LPMEM,LPDAT,16,LNEED)	PARTB	559
00216		LNEED = KDIM*LKNFT	PARTB	560
00217		CALL FINDPT(LPMEM,LPKTAB,17,LNEED)	PARTB	561
00218		LPSIG = LPMEM	PARTB	562
00219		LPCON = LPSIG + NMACRO*JD*KD	PARTB	563
00220		LPCNLD = LPCON + ND*JD*KD	PARTB	564
00221		LPWF = LPCNLD + ND*JD*KD	PARTB	565
00222		LPTKN = LPWF + KD*LIMNFT	PARTB	566
00223		LPTKN2 = LPTKN + KD*LIMNFT	PARTB	567
00224		NEXPOS = LPTKN2 + KD*LIMNFT	PARTB	568
00225		CALL CLSIZE(NEXPOS,LAUSED(2),20HEXPOSURE/EDITS )	PARTB	569
			PARTB	570
		FUEL DEPLETION CONTROL	PARTB	571
			PARTB	572
			PARTB	573
			PARTB	574
		USAGE OF IEOL AND IBURN	PARTB	575
		INPUT 0 LT IEOL LT 99 TO INITIATE SEARCH FOR EOL EXPOSURE BY EOLEX	PARTB	576
		EOLEXP SETS IEOL=99 WHEN IT SEES EOFPL ON NEXT STEP.	PARTB	577
		IEOL=999 WHEN NO MORE EXPOSURE CALCULATIONS ARE TO BE MADE	PARTB	578
		IBURN=0 WHEN AFINAL SOURCE CALCULATION IS TO BE MADE.	PARTB	579
		IBURN=1 IF NO MORE SOURCE CALCULATIONS ARE TO BE MADE	PARTB	580
			PARTB	581
00226		IF (DE.EQ.0.0) IEOL=999	PARTB	582
00227		IF (DE.GT.0.0.AND.XPO.GE.XPOMAX) IEOL=999	PARTB	583
00228		IF (DE.LT.0.0.AND.XPO.LE.XPOMAX) IEOL=999	PARTB	584
00229		IF (IEOL.EQ.999) IBRN=1	PARTB	585
00230		IF (MODEB(5).NE.6) GO TO 400	PARTB	586
00231		IBRN=1	PARTB	587
00232		IEOL=999	PARTB	588
00233	400	CONTINUE	PARTB	589
00234		DEP=DE	PARTB	590
00235		IF (IEOL.EQ.999) GO TO 600	OPUS	140
00236		IF(MOPTIM.LT.2) THEN	PARTB	591
00237		CALL PAGES (MAXLIN)	OPUS	141
00238		ENDIF	PARTB	592
00239		LINES=4	PARTB	593
00240		IOPTEX=1	PARTB	594
00241		ISTEP=ISTEP+1	PARTB	595
00242		MD12=ISTEP/2	PARTB	596
00243		IF (MODEB(12).NE.0.AND.MD12*2.NE.ISTEP) GO TO 490	PARTB	597
00244		IEOL=MODEB(10)	PARTB	598
00245		IF (MODEB(5).GT.5) GO TO 470	PARTB	599
00246		IF (IEOL.LE.0) GO TO 470	PARTB	600
			PARTB	601
		FOLLOWING LOGIC CONTROLS THE END OF LIFE EXPOSURE SEARCH	PARTB	602
		FOR XKEFF=XLMBDA	PARTB	603
			PARTB	604
		ASSIGN CONTROL SEARCH PARAMETER	PARTB	605
00247		MM=MODEB(5)	PARTB	606
00248		IF (MM.EQ.0) GO TO 410	PARTB	607
00249		GO TO (420,430,440,450,450,460). MM	PARTB	608
00250	410	DUMY=0.0	PARTB	609
00251		CALL EOLEXP (XKEFF,DE,XPO,XLMBDA,DEP,DUMY,IEOL,FIRST)	PARTB	610
00252		GO TO 470	PARTB	611
00253	420	CALL EOLEXP (XKEFF,DE,XPO,XLMBDA,DEP,POI,IEOL,FIRST)	PARTB	612
00254		GO TO 470	PARTB	613
00255	430	CONTINUE	PARTB	614
00256		GO TO 470	PARTB	615
00257	440	CONTINUE	PARTB	616
00258		GO TO 470	PARTB	617
00259	450	SRCH=NOTWT	PARTB	618
00260		CALL EOLEXP (XKEFF,DE,XPO,XLMBDA,DEP,SRCH,IEOL,FIRST)	PARTB	619
00261		NOTWT=SRCH	PARTB	620
00262		GO TO 470	PARTB	621
00263	460	CONTINUE	PARTB	622
00264	470	CONTINUE	PARTB	623
00265		IF (DEP.LT.0.0.AND.XPO+DEP.LT.XPOMAX) DEP=XPOMAX-XPO	PARTB	624
00266		IF (DEP.GT.0.0.AND.XPO+DEP.GT.XPOMAX) DEP=XPOMAX-XPO	PARTB	625
			PARTB	626
		SIGCAL DOES THE FUEL DEPLETION WHEN IOPTEX=1	PARTB	627
			PARTB	628
00267		XPO=A1(1)+DEP	PARTB	628

00268		IF (MODEB(12).GT.0.AND.MD12*2.EQ.ISTEP) GO TO 480	PARTB	629
00269		GO TO 510	PARTB	630
00270	480	IOPTEX=3	PARTB	631
00271		GO TO 510	PARTB	632
00272	490	CONTINUE	PARTB	633
00273		WRITE (ITAPOT,500)	PARTB	634
00274	500	FORMAT (/10X,87HTHIS IS THE FIRST STEP IN THE STEP AVERAGE DEPLETI ION PROCEDURE CONTROLLED BY MODEB(12).)	PARTB	635
00275	510	CONTINUE	PARTB	636
00276		CALL SIGCAL(A(LPSTAT),A(LPDAT),A(LPCON),A(LPCNLD),A(LPSIG), 1 A(LPNFT),A(LPTKN),A(LPTKN2),A(LPWF),A(LPKTAB),A(LPKX), 2 A(LPYY),A(LPCNLD),ID,JD,KD,ND,NMACRO,NSTATE,LIMNFT,KDIM,LIST, 3 NDAT,IOPTEX,DEP,XT,XPOLST,KPFRST)	PARTB	637
00277		A1(1)=XPO	PARTB	638
00278		IF (IOPTEX.NE.3) GO TO 530	PARTB	639
00279		LINES=LINES+2	PARTB	640
00280		WRITE (ITAPOT,520) EBAR	PARTB	142
00281		WRITE (ITAPN,520) EBAR	PARTB	642
00282	520	FORMAT (/10X,57HTHIS IS THE REDEPLETION TIMESTEP CONTROLLED BY MOD 1EB(12)./10X,34HTHE CORE AVERAGE EXPOSURE REMAINS ,F10.4)	PARTB	643
00283		GO TO 550	PARTB	644
00284	530	CONTINUE	PARTB	645
00285		IF(MOPTIM .LT. 2) THEN	OPUS	646
00286		WRITE (ITAPOT,540) EBAR	PARTB	647
00287		WRITE (ITAPN,540) EBAR	PARTB	648
00288	540	FORMAT (/10X,63HTHE AVERAGE EXPOSURE OF THE FUEL IN CORE HAS BEEN 1 INCREASED TO ,F10.4)	PARTB	649
00289		ENDIF	OPUS	650
00290	550	CONTINUE	PARTB	143
00291		IF(MOPTIM .LT. 2) THEN	OPUS	651
00292		CALL PAGES (2)	PARTB	652
00293		WRITE (ITAPOT,560) DEP	PARTB	653
00294		WRITE (ITAPN,560) DEP	PARTB	654
00295	560	FORMAT (/10X,42HTHE (ADJUSTED) EXPOSURE INTERVAL USED WAS ,F10.4)	OPUS	144
00296		ENDIF	PARTB	655
00297		IF (MODEB(3).EQ.2.OR.MODEB(3).EQ.4) GO TO 570	PARTB	145
00298		GO TO 590	OPUS	656
00299	570	CONTINUE	PARTB	657
00300		CALL PAGES (1)	PARTB	658
00301		IF(MODEB(12).GT.0.AND.MD12*2.EQ.ISTEP) GO TO 575	PARTB	659
00302		TIMXE1=TIMXE	PARTB	146
00303		TIMXE=TIMXE+XT/3600.	PARTB	660
00304		XETIME=XETIME+XT/3600.	PARTB	661
00305	575	CONTINUE	PARTB	662
00306		WRITE (ITAPOT,580) ISTEP,TIMXE1,TIMXE,XETIME	PARTB	663
00307		WRITE (ITAPN,580) ISTEP,TIMXE1,TIMXE,XETIME	PARTB	664
00308	580	FORMAT (10X,10HTIME STEP ,I4,10H BEGAN AT ,F10.4,20H HOURS AND END 1ED AT ,F10.4,7H HOURS.,/10X,32HTOTAL TIME SINCE INITIALIZATION , 2 18H OF CONCENTRATIONS,8H IS NOW ,F10.4,7H HOURS.)	PARTB	665
00309	590	CONTINUE	PARTB	666
00310	600	CONTINUE	PARTB	667
			PARTB	668
		EDIT AXIAL AVERAGE EXPOSURE, VOID HISTORY AND CONCENTRATIONS	PARTB	669
			PARTB	670
			OPUS	671
00311		IF(MOPTIM .LT. 2) THEN	PARTB	8
00312		CALL CLEAR(A(LPAXLS),7*KD,0.0)	PARTB	673
00313		CALL REWIND (ITAPH,12)	PARTB	674
00314		DO 620 I=1,IMAX	PARTB	675
00315		CALL READX (A(LPCON),LP12,ITAPH,I,ND,JD,KD,LPXTR(12))	PARTB	676
00316		CALL AXIALC(A(LPCON),A(LPYY),ID,JD,KD,ND,A(LPAXLS),NC4,1,7,I)	PARTB	677
00317		IF(ISET1.EQ.0)	PARTB	678
		1CALL AXIALC(A(LPCON),A(LPYY),ID,JD,KD,ND,A(LPAXLS),NCS,2,7,I)	PARTB	147
		IF(ISET2.EQ.0)	PARTB	679
00318		1CALL AXIALC(A(LPCON),A(LPYY),ID,JD,KD,ND,A(LPAXLS),NCS,3,7,I)	PARTB	680
		NCSM=NCS-1	PARTB	681
00319		IF (NCSM.LT.1) GO TO 620	PARTB	682
00320		DO 610 NZ=1,NCSM	PARTB	683
00321		CALL AXIALC(A(LPCON),A(LPYY),ID,JD,KD,ND,A(LPAXLS),NZ,NZ+3,7,I)	PARTB	684
00322			PARTB	685
00323	610	CONTINUE	PARTB	686
00324	620	CONTINUE	PARTB	687
00325		CALL PAGES(15+KMAX)	PARTB	688
00326		WRITE (ITAPOT,630)	PARTB	689
00327		WRITE (ITAPN,630)	PARTB	690
00328	630	FORMAT (/54X,27HAVERAGE AXIAL DISTRIBUTIONS,/, 1 11X,4HNODE,5X,8HEXPOSURE,5X,4HVOID, 8X,7HCONTROL,5X, 2 51H- - - - - C O N C E N T R A T I O N S - - - - - ,/ 3 35X,7HHISTORY, 7X,7HHISTORY,7X,6HIODINE,8X,5HXENON, 47X,10HPROMETHIUM,5X,8HSAMARIUM,/) )	PARTB	691
00329		CALL AXIALP(A(LPAXLS),KD,7,29H(12X,12,3F14.4,4E14.4,12X,12),ITAPOT	PARTB	692
		1)	PARTB	693

00330		CALL AXIALP(A(LPAXLS),KD,7,29H(12X,I2,3F14.4,4E14.4,12X,I2),ITAPN)	PARTB	704
	C	GET 3D EDITS OF CONCENTRATION FILE ARRAYS, E,U,I,XE,PM,SM	PARTB	705
00331		LPMEM=NFLOW	PARTB	706
00332		CALL FINDPT(LPMEM,LPSRC,11,NDIM3D)	PARTB	707
00333		LPE=LPMEM	PARTB	708
00334		LPXE = LPE + NDIM3D	PARTB	709
00335		LPU = LPXE + NDIM3D	PARTB	710
00336		LPCON = LPU + NDIM3D	PARTB	711
00337		LPTKN = LPCON + ND*JD*KD	PARTB	712
00338		LPTKN2 = LPTKN + KD*LIMNFT	PARTB	713
00339		LPWF = LPTKN2 + KD*LIMNFT	PARTB	714
00340		LPMAP1 = LPWF + KD*LIMNFT	PARTB	715
00341		LPMAP2=LPMAP1+MAX0(JD*KD,576)	PARTB	716
00342		NEDTS=LPMAP2*KD*16	PARTB	717
00343		CALL CLSIZE (NEDTS,LAUSED(2),17HTARGET FILE EDITS)	PARTB	718
00344		CALL FZONES(A(LPTKN),A(LPTKN2),A(LPWF),ID,JD,KD,	PARTB	719
		1 LIMNFT)	PARTB	720
00345		CALL EDITB3 (A(LPSRC),A(LPE),A(LPU),A(LPXE),A(LPCON),A(LPMNFT),A(LP	PARTB	721
		INFID),A(LPMAP1),A(LPMAP2),A(LPBATF),A(LPY),ID,JD,KD,ND,	PARTB	722
		2A(LPTKN),A(LPTKN2),A(LPWF),LIMNFT)	PARTB	723
		ENDIF	OPUS	148
00346		IF (DE.GT.0.0.AND.XPO.GE.0.99999*XPOMAX) IEOL=999	PARTB	724
00347		IF (DE.LT.0.0.AND.XPO.LE.1.00001*XPOMAX) IEOL=999	PARTB	725
00348			PARTB	726
	C	SHOULD ANOTHER SOURCE CALCULATION BE INITIATED	PARTB	727
	C		PARTB	728
	C		PARTB	729
00349		IF (IBRN.EQ.0.AND.IEOL.GE.99) GO TO 640	PARTB	730
00350		IF (IEOL.EQ.999) GO TO 650	PARTB	731
00351		GO TO 20	PARTB	732
00352	640	IEOL=999	PARTB	733
00353		IBRN=1	PARTB	734
00354		GO TO 20	PARTB	735
00355	650	IEOL=999	PARTB	736
00356		IF (MODEB(15).EQ.0) GO TO 660	PARTB	737
00357		A1(4)=PTH	PARTB	738
00358		A1(5)=WT	PARTB	739
00359		A1(10)=POI	PARTB	740
00360		A1(11)=NOTWT	PARTB	741
00361		IF(MODEB(5).NE.6) GO TO 660	PARTB	742
	C		PARTB	743
	C	CONTINUATION MODE FOR COMPLETED HALING DEPLETIONS	PARTB	744
	C		PARTB	745
00362		XPO=XPOMAX	PARTB	746
00363		A1(1)=XPO	PARTB	747
00364	660	CONTINUE	OPUS	149
	C		OPUS	150
	C	DEPLETION COMPLETED - PRINT OUT RESULTS	OPUS	151
	C		OPUS	152
00365		IF(MOPTIM.EQ.2) THEN	OPUS	153
00366		WRITE(6,664) KPOIS	OPUS	154
00367	664	FORMAT(10(/),1X,5HCASE,IS,10H COMPLETED,/,1X,20(1H=))	OPUS	155
00368		WRITE(6,667) (IBP(KPOIS,J),J=1,KBPOS)	OPUS	156
00369	667	FORMAT(/,10X,11HBP LOADING,20I4)	OPUS	157
00370		WRITE(6,668) NTOTBP	OPUS	158
00371	668	FORMAT(10X,25HTOTAL NUMBER OF BP RODS, I4)	OPUS	159
00372		WRITE(6,654) PTEMP,XTEMP,IPEAK,JPEAK	OPUS	160
00373	654	FORMAT(/,1X,21HMAX POWER PEAKING OF, F7.5,4H AT, F6.3,	OPUS	161
		*21H GWD/MT AT POSITION (,I1,1H, ,I1,1H))	OPUS	162
00374		WRITE(6,655) XPO	OPUS	163
00375	655	FORMAT(/,1X,25HCYCLE LENGTH ACHIEVED, F7.3,8H GWD/MT)	OPUS	164
00376		IF(XPO.GT. TEMP) THEN	OPUS	165
00377		TEMP=XPO	OPUS	166
00378		ITEMP=KPOIS	OPUS	167
00379		ENDIF	OPUS	168
00380		ENDIF	OPUS	169
00381		KPRST=.FALSE.	PARTB	748
	C		PARTB	749
	C	SET SEARCH EIGENVALUE FOR THE NEXT (SERIES OF )CASES	PARTB	750
	C		PARTB	751
00382		IF(MODEB(21).EQ.0) GO TO 665	PARTB	752
00383		A1(19)=XKEFF	PARTB	753
00384		XLMBDA=A1(19)	PARTB	754
00385		MODEB(21)=0	PARTB	755
00386	665	CONTINUE	PARTB	756
00387		ENDPTH=PTH	PARTB	757
00388		NORIGN=LWA-IBIAS	PARTB	758
00389		CALL PAGES(18)	OPUS	170
00390		IF(MOPTIM.EQ.2) THEN	OPUS	171
00391		DO 675 K=1,KBPOS	OPUS	171
00392		I=IROW(K)	OPUS	172



00393		J=JCOL(K)	OPUS	173
00394		NBP(I,J)=IBP(ITEMP,K)*10+(NBP(I,J)/1000)*1000	OPUS	174
00395	675	CONTINUE	OPUS	175
00396		NBP1=NBPERM-NODEP	OPUS	176
00397		NBP2=NBP1-NODEP1	OPUS	177
00398		CALL OPRES(NBP,NBPERM,NBP1,NBP2,TEMP,ITEMP,EXP,JMX)	OPUS	178
00399		ENDIF	OPUS	179
00400		WRITE(ITAPOT,670) LWA,NORIGN,NPARTA,NFLOW,NUOID,NCROSS,NCOEFS,	PARTB	759
		1 NSOURC,NPOWER,NEDITS,NEXPOS,NEDTS	PARTB	760
00401	670	FORMAT(1H1,10X,26HMEMORY BOUNDARIES IN PARTB,	OPUS	180
		A/10X,18,30H INSTRUCTIONS (LWA)	PARTB	762
		B/10X,18,30H ORIGIN OF CONTAINER ARRAY	PARTB	763
		1/10X,18,30H INPUT DATA BOUNDARY	PARTB	764
		2/10X,18,30H FLOW ITERATION BOUNDARY	PARTB	765
		3//10X,18,30H VOID CALCULATION	PARTB	766
		4/10X,18,30H CROSS SECTION CALCULATION	PARTB	767
		A/10X,18,30H COEFFICIENTS CALCULATION	PARTB	768
		S/10X,18,30H SOURCE ITERATION	PARTB	769
		6/10X,18,30H POWER CALCULATION	PARTB	770
		7/10X,18,30H EDIT SECTION 1	PARTB	771
		8/10X,18,30H EXPOSURE CALCULATION	PARTB	772
		9/10X,18,30H EDIT SECTION 2	PARTB	773
		CALL TIMER (SHPARTB,3)	PARTB	774
00402		RETURN	PARTB	775
00403		END	PARTB	776
00404				

00001		SUBROUTINE POISPUT(POW,XBR,KSTEP,IU,JU,NITERS,IBPAS,	POISPUT	2
	*	DEP,ICNTRL,IDEBUG,IOPTX,IOSC,ISRCH,IUOID,LEVELC,LEVELU,	POISPUT	3
	*	LPAXLS,LPBU,NCOEFS,NCROSS,NEDITS,NEDTS,NEXPOS,NFLOW,NPOWER,	POISPUT	4
	*	NSOURC,NVOID,XPOLST,XT,IFIRST,KBPOS,NBPERM,IBP,DEPL,IROW,JCOL)	POISPUT	5
	CC		POISPUT	6
00003		COMMON/BLOK1/IBLOK1(1),ID,JD,KD,IX1,JX1,KX1,IX2,JX2,KX2,ISET1,	BLOK1	2
00004		CHARACTER*8 ADATE,BDATE,IHTIME,JMTIME	F200CPU	3
00011		DIMENSION JOBNAM(2)	BLOK4	6
	C		POISPUT	14
00015		DIMENSION POW(IU,JU),XBR(IU,JU)	POISPUT	15
00016		DIMENSION IROW(20),JCOL(20),NBPBAS(15,15),PDERIU(15,15,20),	POISPUT	16
	*	NUMBP(20),RHS(120),ACY(120,50),T(100),DET(2)	POISPUT	17
	*	PBASE(15,15),COST(50),PSOL(120),DSOL(120),RW(15000),	POISPUT	18
	*	IW(400),DUM(120),BPLIM(20),IBP(50000,20)	POISPUT	19
00017		LOGICAL IUOID,ICNTRL,IFIRST,LEVELU,LEVELC,DEPL(50000),KPFIRST	POISPUT	20
	C		POISPUT	21
	C		POISPUT	22
	C	AT FIRST STEP OF DEPLETION, FIND FRESH FUEL POSITIONS IN WHICH	POISPUT	23
	C	BP REQUIREMENTS WILL BE SEARCHED FOR	POISPUT	24
	C		POISPUT	25
00018		JBPSOL=IROUND(A1(47))	POISPUT	26
00019		PPLIM=A1(39)	POISPUT	27
00020		BPDEL=A2(50)	POISPUT	28
00021		PPMULT1=A2(51)	POISPUT	29
00022		PPMULT2=A2(52)	POISPUT	30
00023		PPLIM=PPLIM*PPMULT2	POISPUT	31
00024		JBPEDIT=IROUND(A1(48))	POISPUT	32
00025		IF(JBPEDIT.EQ.0) NSTEP=1	POISPUT	33
00026		K=1	POISPUT	34
00027		NODES=0	POISPUT	35
00028		DO 20 I=1,IMAX	POISPUT	36
00029		JL=JMN(I)	POISPUT	37
00030		JR=JMX(I)	POISPUT	38
00031		DO 20 J=JL,JR	POISPUT	39
00032		INBP=NBP(I,J)-(NBP(I,J)/10000)*10000	POISPUT	40
00033		IF(INBP.LE.0) GO TO 30	POISPUT	41
00034		IROW(K)=I	POISPUT	42
00035		JCOL(K)=J	POISPUT	43
00036		K=K+1	POISPUT	44
00037	30	CONTINUE	POISPUT	45
00038		NODES=NODES+1	POISPUT	46
00039	20	CONTINUE	POISPUT	47
00040		KBPOS=K-1	POISPUT	48
	C		POISPUT	49
	C	AT EACH OUTER STEP, KEEP BASE POWER DISTRIBUTION IN ORDER TO GET	POISPUT	50
	C	FIRST ORDER DERIVATIVE, D(POW(I,J))/D(BPNUM)	POISPUT	51
	C		POISPUT	52
00041		DO 10 I=1,IMAX	POISPUT	53
00042		JL=JMN(I)	POISPUT	54
00043		JR=JMX(I)	POISPUT	55
00044		DO 10 J=JL,JR	POISPUT	56
00045		PBASE(I,J)=POW(I,J)	POISPUT	57
00046		NBPBAS(I,J)=NBP(I,J)	POISPUT	58
00047	10	CONTINUE	POISPUT	59
00048		KLCON=0	POISPUT	60
00049		DO 180 KPP=1,KBPOS	POISPUT	61
00050	170	I=IROW(KPP)	POISPUT	62
00051		J=JCOL(KPP)	POISPUT	63
00052		TTNBP=NBPBAS(I,J)-(NBPBAS(I,J)/1000)*1000	POISPUT	64
00053		BPLIM(KPP)=TTNBP/10.	POISPUT	65
00054		LDIFF=ABS(NUMBP(KPP))-KLCON	POISPUT	66
00055		IF(LDIFF.GT.0) KLCON=ABS(NUMBP(KPP))	POISPUT	67
00056	180	CONTINUE	POISPUT	68
	C		POISPUT	69
	C	COMPUTE THE LEAST SQUARE DIFFERENCE, (XBR(I,J)-PBASE(I,J))**2	POISPUT	70
	C		POISPUT	71
00057		TOTAL=0.	POISPUT	72
00058		TOTAL1=0.0	POISPUT	73
00059		IF(JBPSOL.EQ.0) THEN	POISPUT	74
00060		DO 120 KK=1,KBPOS	POISPUT	75
00061		I=IROW(KK)	POISPUT	76
00062		J=JCOL(KK)	POISPUT	77
00063		TOTAL=TOTAL+ABS(XBR(I,J)-PBASE(I,J))/XBR(I,J)	POISPUT	78
00064	120	CONTINUE	POISPUT	79
00065		NTOTL=KBPOS	POISPUT	80
00066		DO 65 I=1,IMAX	POISPUT	81
00067		JL=JMN(I)	POISPUT	82
00068		JR=JMX(I)	POISPUT	83
00069		DO 65 J=JL,JR	POISPUT	84
00070		TOTAL1=TOTAL1+ABS(XBR(I,J)-PBASE(I,J))/XBR(I,J)	POISPUT	85

00071	65	CONTINUE	POISPUT	86
00072		AUGDIFF=100*TOTAL1/NODES	POISPUT	87
00073		ELSE	POISPUT	88
00074		DO 85 I=1,IMAX	POISPUT	89
00075		JL=JMN(I)	POISPUT	90
00076		JR=JMX(I)	POISPUT	91
00077		DO 85 J=JL,JR	POISPUT	92
00078		TOTAL=TOTAL+ABS(XBR(I,J)-PBASE(I,J))/XBR(I,J)	POISPUT	93
00079	85	CONTINUE	POISPUT	94
00080		NTOTL=NODES	POISPUT	95
00081		ENDIF	POISPUT	96
	C		POISPUT	97
	C	TEST CONVERGENCE	POISPUT	98
	C		POISPUT	99
00082		OBSQ=100*TOTAL/NTOTL	POISPUT	100
00083		WRITE(6,3000) OBSQ	POISPUT	101
00084	3000	FORMAT(//,3X,'*** AVERAGE ABSOLUTE PERCENT DIFFERENCE =',F10.5,	POISPUT	102
		* ' PERCENT')	POISPUT	103
00085		IF(JBPSOL.EQ.0) THEN	POISPUT	104
00086		WRITE(6,3060) AUGDIFF	POISPUT	105
00087		ENDIF	POISPUT	106
00088	3060	FORMAT(//,3X,'=== CORE AVERAGE PERCENT DIFFERENCE =',F10.5,	POISPUT	107
		* ' PERCENT')	POISPUT	108
00089		IBPAS=0	POISPUT	109
	C		POISPUT	110
	C		POISPUT	111
00090		IF(NITERS.EQ.1) GO TO 130	POISPUT	112
00091		IF( NITERS .GT. 7 .OR. OBSQ .LE. 0.2 .OR.	POISPUT	113
		* (ABS(OBSQ-OLDSQ) .LE. 0.1 .AND. KLCOM .LT. 5) ) THEN	POISPUT	114
00092		IBPAS=1	POISPUT	115
00093		GO TO 999	POISPUT	116
00094		ENDIF	POISPUT	117
00095	130	OLDSQ=OBSQ	POISPUT	118
	C		POISPUT	119
	C	GET SENSITIVITY COEFFICIENT HERE, INCREMENT OF BPNUM IS 1	POISPUT	120
	C		POISPUT	121
00096		WRITE(6,4545)	POISPUT	122
00097	4545	FORMAT(//,120(1H*),//)	POISPUT	123
00098		WRITE(6,4550) NITERS	POISPUT	124
00099	4550	FORMAT(///,2X,101(1H=),/,2X,'(',I2,')=' , ' OUTER ITERATION STEP'	POISPUT	125
		* ,/,2X,101(1H=),/)	POISPUT	126
00100		IF(OBSQ.LE.3.0) THEN	POISPUT	127
00101		IF(NITERS.EQ.1) GO TO 600	POISPUT	128
00102		WRITE(6,3050)	POISPUT	129
00103	3050	FORMAT(//,3X,55(1H-),/,3X,' NO MORE EVALUATION OF POWER ',	POISPUT	130
		* 'COEFFICIENTS HEREAFTER ',/,3X,55(1H-),/)	POISPUT	131
00104		GO TO 900	POISPUT	132
00105		ENDIF	POISPUT	133
00106	600	CONTINUE	POISPUT	134
00107		MDELBP=2	POISPUT	135
00108		IF(JBPEDIT.EQ.0) THEN	POISPUT	136
00109		IEDIT(2)=1	POISPUT	137
00110		IEDIT(3)=1	POISPUT	138
00111		IEDIT(5)=1	POISPUT	139
00112		ENDIF	POISPUT	140
00113		WRITE(6,4560)	POISPUT	141
00114	4560	FORMAT(///,3X,'=== GET SENSITIVITY COEFFICIENTS ===',//)	POISPUT	142
00115		DO 50 K=1,KBPOS	POISPUT	143
00116		DO 40 I=1,IMAX	POISPUT	144
00117		JL=JMN(I)	POISPUT	145
00118		JR=JMX(I)	POISPUT	146
00119		DO 40 J=JL,JR	POISPUT	147
00120	40	NBP(I,J)=NBPBAS(I,J)	POISPUT	148
00121		IP=IROW(K)	POISPUT	149
00122		JP=JCOL(K)	POISPUT	150
00123		WRITE(6,4570) K,IP,JP	POISPUT	151
00124	4570	FORMAT(//,3X,'(KSTEP)=' ,I3, ' ** POWER SENSITIVITY ARRAY DUE TO',	POISPUT	152
		* ' INSERTION OF 1 BP IN (' ,I2,',' ,I2,') POSITION **')	POISPUT	153
00125		NBP(IP,JP)=NBP(IP,JP)+10*MDELBP	POISPUT	154
00126		NU=0	POISPUT	155
00127		KPFRST=.TRUE.	POISPUT	156
00128		CALL POWDIS(DEP,ICNTRL,IDEBUG,IOPTEX,IOSC,ISRCH,IUOID,	POISPUT	157
		* LEVELC,LEVELU,LPAXLS,LPBU,NCOEFS,NCROSS,NEDITS,NEDTS,NEXPOS,	POISPUT	158
		* NFLOW,NPOWER,NSOURC,NUOID,XPOLST,XT,IFIRST,KPFRST)	POISPUT	159
00129		CALL REWIND(ITAPU,11)	POISPUT	160
00130		CALL READX(POW,LP11,ITAPU,1,IU,JU,KD,LPXTR(11))	POISPUT	161
00131		DO 60 II=1,IMAX	POISPUT	162
00132		JL=JMN(II)	POISPUT	163
00133		JR=JMX(II)	POISPUT	164
00134		DO 61 JJ=JL,JR	POISPUT	165

00135		PDERIU(II, JJ, K)=(POW(II, JJ)-PBASE(II, JJ))/MDELBP	POISPUT	166
00136	61	CONTINUE	POISPUT	167
00137		WRITE(6, 1000) (PDERIU(II, JH, K), JH=1, JR)	POISPUT	168
00138	1000	FORMAT(3X, 15F8.5)	POISPUT	169
00139	60	CONTINUE	POISPUT	170
00140	50	CONTINUE	POISPUT	171
00141	900	CONTINUE	POISPUT	172
	C		POISPUT	173
	C	SOLVE LINEAR EQUATION, AC(I, J)*BPNUM(J)=RHS(J) WHERE	POISPUT	174
	C	I=J=1, KBPOS, RHS(J)=XBR-PBASE	POISPUT	175
	C	HERE SUBROUTINE LEQIF IN IMSL LIBRARY IS USED	POISPUT	176
	C		POISPUT	177
	C	OR.	POISPUT	178
	C	APPLY LINEAR PROGRAMMING TECHNIQUE HERE,	POISPUT	179
	C	F = MAX(- SUM(D(J)))	POISPUT	180
	C		POISPUT	181
	C	SUBJECT TO	POISPUT	182
	C		POISPUT	183
	C	AC(I, J)*X(J) .LE. RHS(J)	POISPUT	184
	C		POISPUT	185
	C	WHERE I=1, (3*NODES)	POISPUT	186
	C	J=1, KBPOS+NODES	POISPUT	187
	C		POISPUT	188
	C	(X(J), J=1, KBPOS+NODES)=(NUMBP(1)..NUMBP(KBPOS), D(1)..D(NODES))	POISPUT	189
	C		POISPUT	190
	C	SUBROUTINE ZX3LP IN IMSL LIBRARY IS UTILIZED	POISPUT	191
	C		POISPUT	192
	C		POISPUT	193
	C		POISPUT	194
00142		IF(JBPSOL.EQ.0) GO TO 700	POISPUT	195
00143		N=KBPOS+NODES	POISPUT	196
00144		NODE2=2*NODES	POISPUT	197
00145		NODE3=3*NODES	POISPUT	198
00146		M2=0	POISPUT	199
00147		WRITE(6, 5000) (BPLIM(L), L=1, KBPOS)	POISPUT	200
00148	5000	FORMAT(//, 4X, 'BPLIM(K) =', 10F8.3)	POISPUT	201
	C		POISPUT	202
00149		DO 76 I=1, NODES	POISPUT	203
00150	76	DUM(I)=0.0	POISPUT	204
00151		KY=1	POISPUT	205
00152		DO 77 I=1, IMAX	POISPUT	206
00153		JL=JMN(I)	POISPUT	207
00154		JR=JMX(I)	POISPUT	208
00155		DO 78 J=JL, JR	POISPUT	209
00156		DO 79 K=1, KBPOS	POISPUT	210
00157	79	DUM(KY)=DUM(KY)+PDERIU(I, J, K)*BPLIM(K)	POISPUT	211
00158	78	KY=KY+1	POISPUT	212
00159	77	CONTINUE	POISPUT	213
00160		KX=1	POISPUT	214
00161		DO 70 I=1, IMAX	POISPUT	215
00162		JL=JMN(I)	POISPUT	216
00163		JR=JMX(I)	POISPUT	217
00164		DO 70 J=JL, JR	POISPUT	218
00165		DO 71 K=1, KBPOS	POISPUT	219
00166		ACY(KX, K)=PDERIU(I, J, K)/XBR(I, J)	POISPUT	220
00167		ACY(KX+NODES, K)=-PDERIU(I, J, K)/XBR(I, J)	POISPUT	221
00168	71	ACY(KX+NODE2, K)=PDERIU(I, J, K)	POISPUT	222
00169		ACY(KX, KX+KBPOS)=-1.0	POISPUT	223
00170		ACY(KX+NODES, KX+KBPOS)=-1.0	POISPUT	224
00171		RHS(KX)=(XBR(I, J)-PBASE(I, J)+DUM(KX))/XBR(I, J)	POISPUT	225
00172		RHS(KX+NODES)=(PBASE(I, J)-XBR(I, J)-DUM(KX))/XBR(I, J)	POISPUT	226
00173		RHS(KX+NODE2)=PPLIM-PBASE(I, J)+DUM(KX)	POISPUT	227
00174		COST(KX+KBPOS)=-1.0	POISPUT	228
00175		KX=KX+1	POISPUT	229
00176	70	CONTINUE	POISPUT	230
	C		POISPUT	231
00177		DO 75 J=1, KBPOS	POISPUT	232
00178	75	COST(J)=0.0	POISPUT	233
	C		POISPUT	234
00179		WRITE(6, 4440)	POISPUT	235
00180	4440	FORMAT(///, 3X, 65(1H=), //, 3X, '??? LINEAR PROGRAMMING SOLVER:'	POISPUT	236
	*	, ' SUBROUTINE ZX3LP IS CALLED ???', //, 3X, 65(1H=), //)	POISPUT	237
00181		TIME1=SECOND()	POISPUT	238
00182		CALL ZX3LP(ACY, 120, RHS, COST, N, NODE3, M2, S, PSOL, DSOL, RW,	POISPUT	239
	\$	IW, IER)	POISPUT	240
00183		TIME2=SECOND()	POISPUT	241
00184		TOTT=TIME2-TIME1	POISPUT	242
00185		DO 500 I=1, KBPOS	POISPUT	243
00186	500	NUMBP(I)=(PSOL(I)-BPLIM(I))*10.	POISPUT	244
00187		SQBQ=S/NODES*100	POISPUT	245
00188		GO TO 800	POISPUT	245

00189	700	CONTINUE	POISPUT	246
	C		POISPUT	247
00190		KK=1	POISPUT	248
00191	80	CONTINUE	POISPUT	249
00192		II=IROW(KK)	POISPUT	250
00193		JJ=JCOL(KK)	POISPUT	251
00194		RHS(KK)=XBR(II,JJ)-PBASE(II,JJ)	POISPUT	252
00195		DO 90 K=1,KBPOS	POISPUT	253
00196		ACY(KK,K)=PDERIU(II,JJ,K)	POISPUT	254
00197	90	CONTINUE	POISPUT	255
00198		IF(KK.EQ.KBPOS) GO TO 95	POISPUT	256
00199		KK=KK+1	POISPUT	257
00200		GO TO 80	POISPUT	258
00201	95	CONTINUE	POISPUT	259
00202		WRITE(6,4449)	POISPUT	260
00203	4449	FORMAT(///,3X,65(1H=),/,3X,'??? LINEAR EQUATION SOLVER:'	POISPUT	261
	*	,' SUBROUTINE LEQIF IS CALLED ???',/,3X,65(1H=),/)	POISPUT	262
00204		TIME1=SECOND()	POISPUT	263
00205		CALL LEQIF(ACY,120,KBPOS,1,RHS,120,1,0,T,IER)	POISPUT	264
00206		TIME2=SECOND()	POISPUT	265
00207		TOTT=TIME2-TIME1	POISPUT	266
00208		DO 100 KI=1,KBPOS	POISPUT	267
00209	100	NUMBP(KI)=RHS(KI)*10	POISPUT	268
00210	800	CONTINUE	POISPUT	269
00211		WRITE(6,4011) TOTT	POISPUT	270
00212	4011	FORMAT(/,3X,'(CPU TIME CONSUMED IN EITHER LP OR LINEAR '	POISPUT	271
	*	' EQUATION SOLVER) =',F10.6,' SECONDS',/)	POISPUT	272
00213		IF(JBPSQL.NE.0) THEN	POISPUT	273
00214		WRITE(6,4000) SOBQ	POISPUT	274
00215	4000	FORMAT(///,3X,'OBJECTIVE FUNCTION =',F10.5,' PERCENT',/)	POISPUT	275
00216		ENDIF	POISPUT	276
00217		WRITE(6,2000) (NUMBP(KM),KM=1,KBPOS)	POISPUT	277
00218	2000	FORMAT(/4X,' ( INCREMENT OF BP DISTRIBUTION ) *10',/,4X,15I5)	POISPUT	278
	C		POISPUT	279
	C	ASSIGN SEARCHED BP DISTRIBUTION AT THIS PLACE	POISPUT	280
	C		POISPUT	281
00219		DO 160 KL=1,KBPOS	POISPUT	282
00220	150	I=IROW(KL)	POISPUT	283
00221		J=JCOL(KL)	POISPUT	284
00222		LNBP=NBPBAS(I,J)-(NBPBAS(I,J)/1000)*1000	POISPUT	285
00223		MTBP=LNBP+NUMBP(KL)	POISPUT	286
00224		IF(MTBP.LE.0) NUMBP(KL)=-LNBP	POISPUT	287
00225		NBP(I,J)=NBPBAS(I,J)+NUMBP(KL)	POISPUT	288
00226	160	CONTINUE	POISPUT	289
00227		WRITE(6,4120)	POISPUT	290
00228	4120	FORMAT(///,3X,'*** CONSTRUCTED POWER USING LINEAR'	POISPUT	291
	*	' APPROXIMATION IN SEARCHED BP DISTRIBUTION ***',/)	POISPUT	292
00229		DO 400 I=1,IMAX	POISPUT	293
00230		JL=JMN(I)	POISPUT	294
00231		JR=JMX(I)	POISPUT	295
00232		DO 410 J=JL,JR	POISPUT	296
00233		PPBAS=PBASE(I,J)	POISPUT	297
00234		DO 420 K=1,KBPOS	POISPUT	298
00235	420	PPBAS=PPBAS+PDERIU(I,J,K)*NUMBP(K)/10	POISPUT	299
00236		POW(I,J)=PPBAS	POISPUT	300
00237	410	CONTINUE	POISPUT	301
00238		WRITE(6,3020) (POW(I,JK),JK=1,JR)	POISPUT	302
00239	3020	FORMAT(3X,8F8.3)	POISPUT	303
00240	400	CONTINUE	POISPUT	304
00241		WRITE(6,3030)	POISPUT	305
00242	3030	FORMAT(///,3X,65(1H=),/,3X,' NORMAL POWER CALCULATION USING'	POISPUT	306
	*	' REESTIMATED BP DISTRIBUTION',/,3X,65(1H=),/)	POISPUT	307
00243		KSTEP=KSTEP+1	POISPUT	308
00244		NITERS=NITERS+1	POISPUT	309
00245	999	CONTINUE	POISPUT	310
00246		IF(JBPEDIT.EQ.0) THEN	POISPUT	311
00247		IEDIT(2)=5	POISPUT	312
00248		IEDIT(3)=5	POISPUT	313
00249		IEDIT(5)=5	POISPUT	314
00250		ENDIF	POISPUT	315
00251		IF (IBFAS .EQ. 1) THEN	POISPUT	316
00252		CALL INTBP(IROW,JCOL,NBPBAS,PBASE,PDERIU,IMAX,JMN,JMX,KBPOS,	POISPUT	317
	*	XBR,IU,JU,NBPERM,IBP,DEPL,PPLIM,BPDEL,PPMULT1)	POISPUT	318
00253		IEDIT(2)=1	POISPUT	319
00254		IEDIT(3)=1	POISPUT	320
00255		IEDIT(5)=1	POISPUT	321
00256		ENDIF	POISPUT	322
00257		RETURN	POISPUT	323
00258		END	POISPUT	324

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00001      SUBROUTINE INTBP(IROW,JCOL,NBPBAS,PBASE,PDERIU,IMAX,JMN,JMX,
*          KBPOS,XBR,IU,JU,NBPERM,IBP,DEPL,PPLIM,BPDEL,PPMULT1)
00002      LOGICAL DEPL(50000)
00003      DIMENSION IBPERM(50000),NDISBP(20),NLOBP(20),NHIBP(20),
*          NXBP(20),NXBP(20),IROW(20),JCOL(20),NBPBAS(15,15),
*          JMN(34),JMX(34),PBASE(15,15),PDERIU(15,15,20),
*          IBPX0(20),LOCX(5),IPOINT(50000,20),
*          PPQW(15,15),XBR(IU,JU),IBPEX(50000),IBP(50000,20)
C
00004      PARAMETER(INTMAX=6)
00005      DATA (NDISBP(I),I=1,7) /0,4,8,12,16,20,24/
C
C          DETERMINE THE HIGH AND LOW INTEGRAL NUMBER OF BP'S
C          CORRESPONDING TO THE POISPUT RESULTS
C
00006      DO 100 K=1,KBPOS
00007      I=IROW(K)
00008      J=JCOL(K)
00009      RNBP(K)=FLOAT(NBPBAS(I,J) - (NBPBAS(I,J)/1000)*1000)/10.0
00010      IF(RNBP(K).GT.NDISBP(INTMAX)) THEN
00011      WRITE(6,150) I,J
00012      150  FORMAT(/,3X,'BP LIMIT EXCEEDED AT POSITION (',I2,',',I2,')')
00013      NHIBP(K)=NDISBP(INTMAX)
00014      NLOBP(K)=NDISBP(INTMAX-1)
00015      ENDIF
00016      DO 200 L=2,INTMAX
00017      IF(RNBP(K).LE.NDISBP(L).AND.RNBP(K).GE.NDISBP(L-1)) THEN
00018      NHIBP(K)=NDISBP(L)
00019      NLOBP(K)=NDISBP(L-1)
00020      ENDIF
00021      200  CONTINUE
00022      NMAX=NDISBP(INTMAX)
00023      IF((NHIBP(K)-RNBP(K)).LE.BPDEL.AND.NHIBP(K).NE.NMAX) THEN
00024      NXBP(K)=NHIBP(K)+4
00025      ELSEIF((RNBP(K)-NLOBP(K)).LE.BPDEL.AND.NLOBP(K).NE.0) THEN
00026      NXBP(K)=NLOBP(K)-4
00027      ELSE
00028      NXBP(K)=-1
00029      ENDIF
00030      100  CONTINUE
C
C          PRODUCE ALL POSSIBLE COMBINATIONS OF INTEGRAL BP LOADINGS
C
00031      NBPERM=2**KBPOS
00032      CALL COMBO(KBPOS,NBPERM,IBPERM)
00033      DO 205 K=1,KBPOS
00034      IPOINT(1,K)=K
00035      205  CONTINUE
C          INCLUDE COMBINATIONS WITH 'EXTRA' BP POSSIBILITIES
00036      NUMX=0
00037      DO 210 K=1,KBPOS
00038      IF(NXBP(K).GE.0) THEN
00039      NUMX=NUMX+1
00040      LOCX(NUMX)=K
00041      ENDIF
00042      210  CONTINUE
00043      NBPX=2**NUMX
00044      CALL COMBO(NUMX,NBPX,IBPX0)
00045      NUPERM=0
00046      DO 220 I=2,NBPX
00047      DO 225 II=1,KBPOS
00048      IPOINT(I,II)=II
00049      225  CONTINUE
00050      NUM=IBPX0(I)
00051      NSUM=0
00052      DO 230 J=1,NUMX
00053      JJ=NUMX-J
00054      JJJ=NUM/10**JJ
00055      NUM=NUM-JJJ*10**JJ
00056      IF(JJJ.EQ.1) THEN
00057      NSUM=NSUM+1
00058      IITEM=IPOINT(I,NSUM)
00059      JPOS=LOCX(J)
00060      IPOINT(I,NSUM)=IPOINT(I,JPOS)
00061      IPOINT(I,JPOS)=IITEM
00062      ENDIF
00063      230  CONTINUE
00064      NUMEX=KBPOS-NSUM
00065      NBPEX=2**NUMEX
00066      CALL COMBO(NUMEX,NBPEX,IBPEX)

```

```

INTBP 2
INTBP 3
INTBP 4
INTBP 5
INTBP 6
INTBP 7
INTBP 8
INTBP 9
INTBP 10
INTBP 11
INTBP 12
INTBP 13
INTBP 14
INTBP 15
INTBP 16
INTBP 17
INTBP 18
INTBP 19
INTBP 20
INTBP 21
INTBP 22
INTBP 23
INTBP 24
INTBP 25
INTBP 26
INTBP 27
INTBP 28
INTBP 29
INTBP 30
INTBP 31
INTBP 32
INTBP 33
INTBP 34
INTBP 35
INTBP 36
INTBP 37
INTBP 38
INTBP 39
INTBP 40
INTBP 41
INTBP 42
INTBP 43
INTBP 44
INTBP 45
INTBP 46
INTBP 47
INTBP 48
INTBP 49
INTBP 50
INTBP 51
INTBP 52
INTBP 53
INTBP 54
INTBP 55
INTBP 56
INTBP 57
INTBP 58
INTBP 59
INTBP 60
INTBP 61
INTBP 62
INTBP 63
INTBP 64
INTBP 65
INTBP 66
INTBP 67
INTBP 68
INTBP 69
INTBP 70
INTBP 71
INTBP 72
INTBP 73
INTBP 74
INTBP 75
INTBP 76
INTBP 77
INTBP 78
INTBP 79
INTBP 80
INTBP 81

```

00067		DO 250 M=1,NBPEX	INTBP	82
00068		NUPERM=NUPERM+1	INTBP	83
00069		IBPERM(NBPERM+NUPERM)=INT((2.0/9.0)*10**NSUM)*10**(KBPOS-NSUM)	INTBP	84
	1	+IBPEX(M)	INTBP	85
00070		DO 255 K=1,KBPOS	INTBP	86
00071		IPOINT(NBPERM+NUPERM,K)=IPOINT(I,K)	INTBP	87
00072	255	CONTINUE	INTBP	88
00073	250	CONTINUE	INTBP	89
00074	220	CONTINUE	INTBP	90
00075		DO 245 M=1,NBPERM	INTBP	91
00076		DO 245 K=1,KBPOS	INTBP	92
00077		IPOINT(M,K)=IPOINT(I,K)	INTBP	93
00078	245	CONTINUE	INTBP	94
00079		NBPERM=NBPERM+NUPERM	INTBP	95
00080	270	CONTINUE	INTBP	96
	C		INTBP	97
00081		DO 300 M=1,NBPERM	INTBP	98
00082		DEPL(M)=.TRUE.	INTBP	99
	C		INTBP	100
	C	DETERMINE THE ACTUAL BP LOADING FROM THE BP IDENTIFIER 'IBPERM'	INTBP	101
	C		INTBP	102
00083		NOLDBP=IBPERM(M)	INTBP	103
00084		DO 350 KK=KBPOS-1,0,-1	INTBP	104
00085		K=IPOINT(M,KBPOS-KK)	INTBP	105
00086		IBPID=NOLDBP/10**KK	INTBP	106
00087		NOLDBP=NOLDBP-IBPID*10**KK	INTBP	107
00088		IF(IBPID .EQ. 0) THEN	INTBP	108
00089		IBP(M,K)=NLOBP(K)	INTBP	109
00090		ELSEIF(IBPID .EQ. 1) THEN	INTBP	110
00091		IBP(M,K)=NHIBP(K)	INTBP	111
00092		ELSEIF(IBPID .EQ. 2) THEN	INTBP	112
00093		IBP(M,K)=NXBP(K)	INTBP	113
00094		ENDIF	INTBP	114
00095	350	CONTINUE	INTBP	115
	C		INTBP	116
	C	CONSTRUCT POWER DISTRIBUTION USING LINEAR APPROXIMATION	INTBP	117
	C	COMPUTE THE DIFFERENCE, (XBR(I,J)-PBASE(I,J))	INTBP	118
	C		INTBP	119
00096		TOTAL=0.	INTBP	120
00097		DO 400 I=1,IMAX	INTBP	121
00098		JL=JMN(I)	INTBP	122
00099		JR=JMX(I)	INTBP	123
00100		DO 410 J=JL,JR	INTBP	124
00101		PPBAS=PBASE(I,J)	INTBP	125
00102		DO 420 K=1,KBPOS	INTBP	126
00103		PPBAS=PPBAS+PDERIU(I,J,K)*(IBP(M,K)-RNBP(K))	INTBP	127
00104	420	CONTINUE	INTBP	128
00105		PPOW(I,J)=PPBAS	INTBP	129
00106		IF(PPBAS .GT. (PPLIM*PPMULT1)) DEPL(M)=.FALSE.	INTBP	130
00107	410	CONTINUE	INTBP	131
00108	400	CONTINUE	INTBP	132
00109	300	CONTINUE	INTBP	133
00110		RETURN	INTBP	134
00111		END	INTBP	135

00001		SUBROUTINE OPRES(NBP,NBPERM,NBP1,NBP2,TEMP,ITEMP,EXP,JMX)	OPRES	2
00002		DIMENSION NBP(15,15),EN(8),NOBP(8),EXP(8,8),JMX(34)	OPRES	3
00003		CALL FACE	OPRES	4
00004		WRITE(6,50)	OPRES	5
00005	50	FORMAT(1H1,15H OPUS RESULTS ,/,15(1H*))	OPRES	6
00006		WRITE(6,100) NBPERM, NBP1, NBP2	OPRES	7
00007	100	FORMAT(///,5X,45H TOTAL NUMBER OF INTEGRAL BP COMBINATIONS	OPRES	8
		*15,/,5X,45H NUMBER OF DEPLETABLE BP COMBINATIONS ,15,	OPRES	9
		*/,5X,45H NUMBER OF ACCEPTABLE BP COMBINATIONS ,15,///)	OPRES	10
00008		WRITE(6,150) ITEMP,TEMP	OPRES	11
00009	150	FORMAT(///,1X,5HCASE ,15,23H GIVES LONGEST CYCLE AT,F7.3,	OPRES	12
		*7H GWD/MT,/,1X,15H FUEL LOADING: ,/,1X,15(1H-))	OPRES	13
00010		WRITE(6,200)	OPRES	14
00011	200	FORMAT(/,25X,16(1H*),/,25X,16H* ENRICHMENT *,/,25X,	OPRES	15
		*16H* BP LOADING *,/,25X,16H* BOC EXPOSURE *,/,25X,16(1H*))	OPRES	16
00012		DO 300 I=1,8	OPRES	17
00013		DO 350 J=1,JMX(I)	OPRES	18
00014		NOBP(J)=MOD(NBP(I,J),1000)/10	OPRES	19
00015		EN(J)=MOD(NBP(I,J),1000000)/10000/100.0	OPRES	20
00016	350	CONTINUE	OPRES	21
00017		GO TO (355,360,365,370,375,380,385,390) I	OPRES	22
00018	355	WRITE(6,356)	OPRES	23
00019	356	FORMAT(1X,10(1H*))	OPRES	24
00020		GO TO 395	OPRES	25
00021	360	WRITE(6,361)	OPRES	26
00022	361	FORMAT(1X,19(1H*))	OPRES	27
00023		GO TO 395	OPRES	28
00024	365	WRITE(6,366)	OPRES	29
00025	366	FORMAT(1X,28(1H*))	OPRES	30
00026		GO TO 395	OPRES	31
00027	370	WRITE(6,371)	OPRES	32
00028	371	FORMAT(1X,37(1H*))	OPRES	33
00029		GO TO 395	OPRES	34
00030	375	WRITE(6,376)	OPRES	35
00031	376	FORMAT(1X,46(1H*))	OPRES	36
00032		GO TO 395	OPRES	37
00033	380	WRITE(6,381)	OPRES	38
00034	381	FORMAT(1X,55(1H*))	OPRES	39
00035		GO TO 395	OPRES	40
00036	385	WRITE(6,386)	OPRES	41
00037	386	FORMAT(1X,55(1H*))	OPRES	42
00038		GO TO 395	OPRES	43
00039	390	WRITE(6,391)	OPRES	44
00040	391	FORMAT(1X,55(1H*))	OPRES	45
00041	395	CONTINUE	OPRES	46
00042		WRITE(6,400) (EN(J),J=1,JMX(I))	OPRES	47
00043	400	FORMAT(1X,1H*,8(F7.2,2H*))	OPRES	48
00044		GO TO (405,410,415,420,425,430,435,440) I	OPRES	49
00045	405	WRITE(6,406)	OPRES	50
00046	406	FORMAT(1X,1H*,8X,1H*)	OPRES	51
00047		GO TO 445	OPRES	52
00048	410	WRITE(6,411)	OPRES	53
00049	411	FORMAT(1X,1H*,2(8X,1H*))	OPRES	54
00050		GO TO 445	OPRES	55
00051	415	WRITE(6,416)	OPRES	56
00052	416	FORMAT(1X,1H*,3(8X,1H*))	OPRES	57
00053		GO TO 445	OPRES	58
00054	420	WRITE(6,421)	OPRES	59
00055	421	FORMAT(1X,1H*,4(8X,1H*))	OPRES	60
00056		GO TO 445	OPRES	61
00057	425	WRITE(6,426)	OPRES	62
00058	426	FORMAT(1X,1H*,5(8X,1H*))	OPRES	63
00059		GO TO 445	OPRES	64
00060	430	WRITE(6,431)	OPRES	65
00061	431	FORMAT(1X,1H*,6(8X,1H*))	OPRES	66
00062		GO TO 445	OPRES	67
00063	435	WRITE(6,436)	OPRES	68
00064	436	FORMAT(1X,1H*,6(8X,1H*))	OPRES	69
00065		GO TO 445	OPRES	70
00066	440	WRITE(6,441)	OPRES	71
00067	441	FORMAT(1X,1H*,4(8X,1H*))	OPRES	72
00068	445	CONTINUE	OPRES	73
00069		DO 450 J=1,JMX(I)	OPRES	74
00070		IF(NOBP(J) .LE. 0) GO TO 450	OPRES	75
00071		GO TO (500,550,600,650,700,750) J	OPRES	76
00072	500	WRITE(6,525) NOBP(J)	OPRES	77
00073	525	FORMAT(1H+,1X,I7)	OPRES	78
00074		GO TO 450	OPRES	79
00075	550	WRITE(6,575) NOBP(J)	OPRES	80
00076	575	FORMAT(1H+,10X,I7)	OPRES	81



```

00077      GO TO 450
00078      600  WRITE(6,625) NOBP(J)
00079      625  FORMAT(1H+,19X,I7)
00080      GO TO 450
00081      650  WRITE(6,675) NOBP(J)
00082      675  FORMAT(1H+,29X,I7)
00083      GO TO 450
00084      700  WRITE(6,725) NOBP(J)
00085      725  FORMAT(1H+,37X,I7)
00086      GO TO 450
00087      750  WRITE(6,775) NOBP(J)
00088      775  FORMAT(1H+,45X,I7)
00089      450  CONTINUE
00090      WRITE(6,800) (EXP(I,J),J=1,JMX(I))
00091      800  FORMAT(1X,1H*,8(F7.3,2H*))
00092      300  CONTINUE
00093      WRITE(6,850)
00094      850  FORMAT(1X,37(1H*))
00095      RETURN
00096      END

```

```

OPRES      82
OPRES      83
OPRES      84
OPRES      85
OPRES      86
OPRES      87
OPRES      88
OPRES      89
OPRES      90
OPRES      91
OPRES      92
OPRES      93
OPRES      94
OPRES      95
OPRES      96
OPRES      97
OPRES      98
OPRES      99
OPRES     100
OPRES     101

```

```
00001 SUBROUTINE COMBO(X,Y,A)
00002 INTEGER X,Y,A(Y)
00003 DO 50 M=1,Y
00004 NUMBER=M-1
00005 IX=0
00006 DO 60 KK=1,X
00007 K=X-KK
00008 IDELX=NUMBER/2**K
00009 IX=IX+IDELX*10**K
00010 NUMBER=NUMBER-IDELX*2**K
00011 60 CONTINUE
00012 A(M)=IX
00013 50 CONTINUE
00014 RETURN
00015 END
```

```
COMBO 2
COMBO 3
COMBO 4
COMBO 5
COMBO 6
COMBO 7
COMBO 8
COMBO 9
COMBO 10
COMBO 11
COMBO 12
COMBO 13
COMBO 14
COMBO 15
COMBO 16
```

APPENDIX B: ADDITIONAL INPUT REQUIREMENTS

A. Card Type 1

- 1(39) PPLIM Nodal power peaking limit
- 1(46) MOPTIM Control of optimization flow:
  - =1 - direct search only
  - =2 - BP search only
  - =3 - combined (currently ~~un~~available)
- 1(47) JBPSOL BP search method.
  - =0 fresh fuel search only
  - =1 linear programming
- 1(48) JBPEDIT Output edit control
  - =0 normal operation
  - =1 for debugging purposes

B. Card Type 2

- 2(50) BPDEL Burnable poison tolerance (0.5 is recommended)
- 2(51) PPMULT1 Reconstructed power peaking tolerance (1.02 is recommended)
- 2(52) PPMULT2 SIMULATE power peaking tolerance (1.00 is recommended)

C. Card Type 17

17BR(i,j) Target power distribution

D. Card Type 32

32NBP(i,j) Burnable poison identifier array, of the form ABBBCDDD.

where,

- A = fuel type (1 - 15x15 std,  
2 - 15x15 ofa)
- BBB= enrichment (e.g. 2.0 w/o - 270)
- C = BP type (1 - glass, 2 - WABA)
- DDD= number of BP's (e.g. 12 BP - 120)

In the highly unlikely situation that none of the candidate BP loadings successfully depletes to end-of-cycle, several courses of action are available. Corrective actions should be pursued in the following order:

1. increasing the BP tolerance, BPDEL, to create more candidate BP loadings to be evaluated
2. increasing the reconstructed power peaking tolerance, PPMULT1, to allow more cases to be passed on to the normal power calculation / depletion
3. increasing the SIMULATE power peaking tolerance, PPMULT2, to allow a higher power peaking to be acceptable for a final design pattern

It is much more likely that the opposite situation would occur, that too many cases reach the depletion stage such that the computation costs become excessive. In this case, the above actions should be reversed (i.e. instead of increasing tolerances, decrease them). However, the same order still applies to these alternatives.

**APPENDIX C: SAMPLE INPUT LISTING**









43  
43  
99  
678  
222  
333  
333  
333  
333  
13  
3

ILAST

15411,BBL,CY,L10000.  
RESOURCE(JCAT=S3)  
ATTACH,ZZSYM.  
ATTACH,SIMTRN.  
SIMTRN.

47 23 6

I \*\*\* ZION-1 CYCLE 9 SIMULATION (2-D MODEL, FLARE OPTION)

```
*****  
* ZION-1 CYCLE-9 2-D SIMULATION, MODERATOR FEEDBACK*  
* INCLUDED, FLARE OPTION USED WITH INODE/1 ASSEM *  
* BUILT-IN BPMODEL USED, SO EVERY XS IS FOR 0 BP *  
* FUEL TYPE DEC.14.1986 *  
* BP SEARCH - PATTERN 2B *  
*****  
0.0 18.0 2.0 12.01 16.875 178.6 2250. 0  
S1 1000. S3 82.17 S2  
.981 .981 1.0 21.608 10.804  
S1 -1 -1 0 0 3 12.01 16.875 99.56 0  
0 1 1 13.563 0 2250. .73634 1.33 3 4  
2 5 6 0 2 1 0 5 1.33  
0 3 1 0 20 0 2 .0005 4 .00005 12  
.00005 4 .0001 1 0 0 0 0 1.0  
1.0 0.6 0 1.0 4 1.2 .6 S2 .35 1.0 0  
5 7 18 1 0 0 0 0 4 6 15 0 0 0 0 20.0  
0.50 1.020 1.000  
1 S6 0.455 S2 1 30.48 S13 S9 6.241E+18 2.679 S2 4 0.0 .06404  
1 S2 .288E-4 .00633 S2 .211E-4 .01255 S2 .385E-5 0.0 S2 0.0  
2 S6 0.455 S2 2 30.48 S13 S9 6.241E+18 2.672 S2 4 0.0 .06405  
2 S2 .288E-4 .00626 S2 .211E-4 .01255 S2 .385E-5 0.0 S2 0.0  
3 S6 0.455 S2 3 30.48 S13 S9 6.241E+18 2.721 S2 4 0.0 .06420  
3 S2 .288E-4 .00670 S2 .211E-4 .01277 S2 .385E-5 0.0 S2 0.0  
4 S6 .455 S2 4 30.48 S13 S9 6.241E+18  
4 2.4428E+00 1.4184E-02 -1.3229E-04 4 0.0  
4 6.3547E-02 1.2444E-05 2.4107E-07 .288E-4  
4 2.7119E-03 2.2301E-04 -2.6877E-06 .211E-4  
4 1.1340E-02 6.2331E-05 -4.5984E-07 .385E-5 0.0 S2 0.0  
5 S6 .455 S2 5 30.48 S13 S9 6.241E+18  
5 2.4421E+00 1.3728E-02 -1.2446E-04 4 0.0  
5 6.3550E-02 1.1127E-05 2.5762E-07 .288E-4  
5 2.6939E-03 2.1710E-04 -2.5757E-06 .211E-4  
5 1.1336E-02 6.0015E-05 -4.2172E-07 .385E-5 0.0 S2 0.0  
6 S6 .455 S2 6 30.48 S13 S9 6.241E+18  
6 2.4396E+00 1.2124E-02 -9.8852E-05 4 0.0  
6 6.3561E-02 6.8125E-06 3.0254E-07 .288E-4  
6 2.6363E-03 1.9567E-04 -2.1907E-06 .211E-4  
6 1.1321E-02 5.1940E-05 -2.9976E-07 .385E-5 0.0 S2 0.0  
4 1 1  
4 2 1 5  
4 3 1 6 2  
4 4 1 5 6 5 4  
4 5 1 3 6 4 6 2  
4 6 1 3 4 6 2 6 5  
4 7 1 4 6 5 6 6 2  
4 8 1 3 3 2 3  
5 1 1 1 1 0 0 1 20 3  
5 1 1 0 1.5000E-01 1.0000E+00 2.0000E+00 4.0000E+00  
5 1 1 6.0000E+00 8.0000E+00 1.0000E+01 1.2000E+01 1.4000E+01  
5 1 1 1.6000E+01 1.8000E+01 2.0000E+01 2.4000E+01 2.8000E+01  
5 1 1 3.2000E+01 3.6000E+01 4.0000E+01 4.5000E+01 5.0000E+01  
5 1 1 2.3604E-01 2.3593E-01 2.3584E-01 2.3561E-01 2.3534E-01  
5 1 1 2.3499E-01 2.3467E-01 2.3443E-01 2.3396E-01 2.3346E-01  
5 1 1 2.3297E-01 2.3253E-01 2.3217E-01 2.3222E-01 2.3322E-01  
5 1 1 2.3408E-01 2.3489E-01 2.3563E-01 2.3640E-01 2.3704E-01  
5 1 1 2 1 0 0 1 20 3  
5 1 1 0 1.5000E-01 1.0000E+00 2.0000E+00 4.0000E+00  
5 1 1 6.0000E+00 8.0000E+00 1.0000E+01 1.2000E+01 1.4000E+01  
5 1 1 1.6000E+01 1.8000E+01 2.0000E+01 2.4000E+01 2.8000E+01  
5 2 1 3.2000E+01 3.6000E+01 4.0000E+01 4.5000E+01 5.0000E+01  
5 2 1 8.9364E-03 8.9471E-03 8.9775E-03 9.0403E-03 9.1868E-03  
5 2 1 9.3689E-03 9.5361E-03 9.6798E-03 9.8195E-03 9.9477E-03  
5 2 1 1.0064E-02 1.0172E-02 1.0272E-02 1.0454E-02 1.0619E-02  
5 2 1 1.0769E-02 1.0905E-02 1.1041E-02 1.1189E-02 1.1326E-02  
5 3 1 3 1 0 0 1 20 3  
5 3 1 0 1.5000E-01 1.0000E+00 2.0000E+00 4.0000E+00  
5 3 1 6.0000E+00 8.0000E+00 1.0000E+01 1.2000E+01 1.4000E+01  
5 3 1 1.6000E+01 1.8000E+01 2.0000E+01 2.4000E+01 2.8000E+01  
5 3 1 3.2000E+01 3.6000E+01 4.0000E+01 4.5000E+01 5.0000E+01  
5 3 1 1.7108E-02 1.7157E-02 1.7138E-02 1.7117E-02 1.7035E-02  
5 3 1 1.6924E-02 1.6822E-02 1.6724E-02 1.6682E-02 1.6661E-02
```



10 4 1	0
10 4 2	9.410
10 4 3	0
10 4 4	13.100
10 4 5	8.270
10 4 6	22.750
10 4 7	0
10 4 8	24.040
10 5 1	25.890
10 5 2	7.910
10 5 3	12.770
10 5 4	8.270
10 5 5	21.220
10 5 6	0
10 5 7	0
10 6 1	24.940
10 6 2	12.240
10 6 3	0
10 6 4	22.750
10 6 5	0
10 6 6	0
10 6 7	31.420
10 7 1	12.960
10 7 2	0
10 7 3	0
10 7 4	0
10 7 5	0
10 7 6	31.420
10 8 1	25.880
10 8 2	22.630
10 8 3	23.770
10 8 4	24.040

12	-0.0001	S3	.008		
13	0.59582	-0.49396	-0.20215		

14	1	1.0	0.0				
15	1	1	1	1	1	1	1
15	2	1	1	1	1	1	1
15	3	1	1	1	1	1	1
15	4	1	1	1	1	1	1
15	5	1	1	1	1	1	1
15	6	1	1	1	1	1	1
15	7	1	1	1	1	1	1
15	8	1	1	1	1	1	1

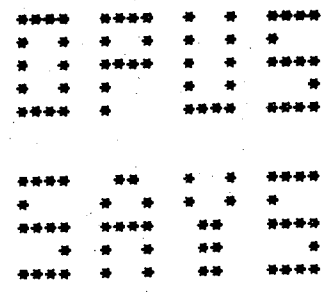
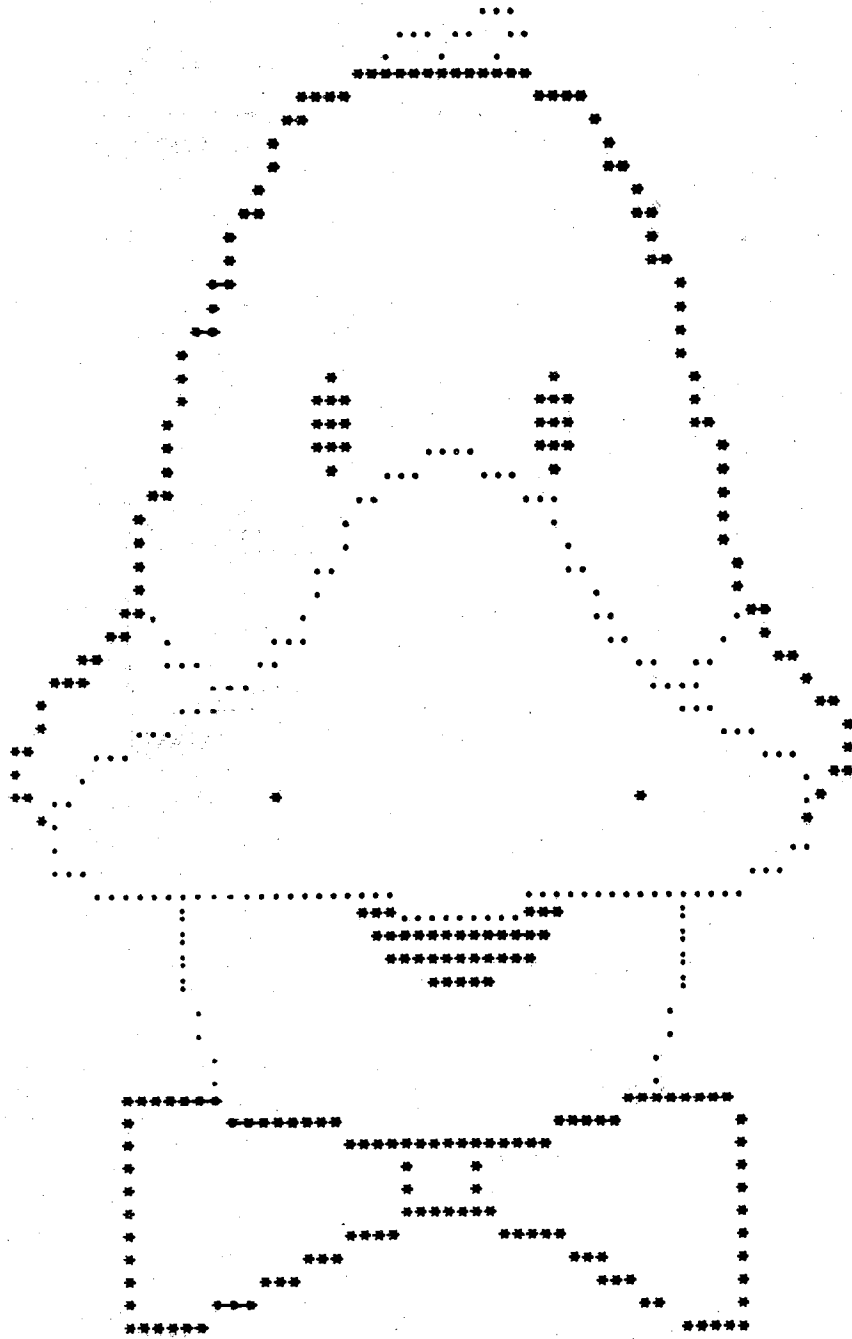
17	1	0.983					
17	2	1.115	1.296				
17	3	1.032	1.176	1.081			
17	4	1.261	1.264	1.299	1.143		
17	5	0.899	1.172	1.149	1.175	1.081	
17	6	0.806	1.051	1.326	1.067	1.241	1.025
17	7	0.849	1.133	1.188	1.124	0.921	0.416
17	8	0.387	0.461	0.486	0.364		

19	1	5	5	1	5	1	R60
20	0	0	S6	0	1	S1	1
21	1	1	3	0	1	0	1
					S2	1	1
					S2	0	1
					1	1	S6
					0		

32	1	22900000					
32	2	22700000	22802080				
32	3	22900000	23200000	23200000			
32	4	22802080	23200000	22802080	22700000		
32	5	22800000	23200000	22700000	23200000	23200000	
32	6	22800000	22700000	23202080	23200000	23202080	22802080
32	7	22700000	23202080	22802080	23202080	23202080	23200000
32	8	22800000	22800000	23200000	22800000		

LAST

**APPENDIX D: SAMPLE OUTPUT**



OPUS RESULTS  
\*\*\*\*\*

TOTAL NUMBER OF INTEGRAL BP COMBINATIONS	1152
NUMBER OF DEPLETABLE BP COMBINATIONS	39
NUMBER OF ACCEPTABLE BP COMBINATIONS	2

CASE 373 GIVES LONGEST CYCLE AT 13.240 GWD/MT

FUEL LOADING:

\*\*\*\*\*  
\* ENRICHMENT \*  
\* BP LOADING \*  
\* BOC EXPOSURE \*  
\*\*\*\*\*

```
*****
* 2.90 *
* *
* 27.530 *
*****
* 2.80 * 3.60 *
* * 4 *
* 13.980 * 0.000 *
*****
* 2.70 * 2.80 * 2.80 *
* * * *
* 20.730 * 15.200 * 14.670 *
*****
* 3.20 * 3.20 * 3.20 * 2.70 *
* 4 * * 8 * *
* 0.000 * 17.960 * 0.000 * 26.100 *
*****
* 2.80 * 3.20 * 3.20 * 3.20 * 2.70 *
* * * * 8 * *
* 10.400 * 13.060 * 14.320 * 0.000 * 24.600 *
*****
* 3.20 * 3.20 * 2.80 * 3.20 * 3.60 * 3.20 *
* * * * * 4 * *
* 14.990 * 10.110 * 14.950 * 10.200 * 0.000 * 0.000 *
*****
* 3.20 * 3.60 * 3.60 * 3.20 * 3.20 * 2.70 *
* * 8 * 4 * * * *
* 14.990 * 0.000 * 0.000 * 0.000 * 21.570 * 25.060 *
*****
* 2.80 * 3.20 * 3.20 * 2.70 *
* * * * *
* 27.200 * 22.080 * 25.290 * 23.160 *
*****
```