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# Classes of Filters in Generalizations of Commutative Fuzzy Structures* 

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#### Abstract

Bounded commutative residuated lattice ordered monoids ( $R \ell$-monoids) are a common generalization of $B L$-algebras and Heyting algebras, i.e. algebras of basic fuzzy logic and intuitionistic logic, respectively. In the paper we develop the theory of filters of bounded commutative $R \ell$ monoids.


Key words: Residuated $\ell$-monoid, deductive system, $B L$-algebra, $M V$-algebra, Heyting algebra, filter.
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## 1 Introduction

$B L$-algebras have been introduced by $P$. Hájek as an algebraic counterpart of the basic fuzzy logic $B L[5]$. Omitting the requirement of pre-linearity in the definition of a $B L$-algebra, one obtains the definition of a bounded commutative residuated lattice ordered monoid ( $R \ell$-monoid). Nevertheless, bounded commutative $R \ell$-monoids are a generalization not only of $B L$-algebras but also of Heyting algebras which are an algebraic counterpart of the intuitionistic propositional logic. Therefore, bounded commutative $R \ell$-monoids could be taken as an algebraic semantics of a more general logic than Hájek's fuzzy logic. It is

[^0]known that every $B L$-algebra (and consequently every $M V$-algebra [2], or equivalently, every Wajsberg algebra [4]) is a subdirect product of linearly ordered $B L$-algebras. Moreover, a bounded commutative $R \ell$-monoid is a subdirect product of linearly ordered $R \ell$-monoids if and only if it is a $B L$-algebra [13]. On the other side, bounded commutative $R \ell$-monoids which need not be $B L$-algebras can be constructed from $B L$-algebras by means of other natural operations, e.g. by means of pasting, i.e. ordinal sums. For example, the pasting of Wajsberg algebras which are not linearly ordered gives bounded commutative $R \ell$-monoids which are not $B L$-algebras $[8,9]$.

In both $B L$-algebras and bounded commutative $R \ell$-monoids, filters coincide with deductive systems of those algebras and are exactly the kernels of their congruences. Various types of filters of $B L$-algebras were studied in [19], [7] and [11]. Boolean filters of bounded commutative $R \ell$-monoids were investigated in [14].

In this paper we further develop the theory of filters of bounded commutative $R \ell$-monoids and among others, we generalize some results of [7] and [11].

For concepts and results concerning $M V$-algebras, $B L$-algebras and Heyting algebras see for instance [2], [5], [1].

## 2 Preliminaries

A bounded commutative $R \ell$-monoid is an algebra $M=(M ; \odot, \vee, \wedge, \rightarrow, 0,1)$ of type $\langle 2,2,2,2,0,0\rangle$ satisfying the following conditions:

| $(R \ell 1)$ | $(M ; \odot, 1)$ is a commutative monoid. |
| :--- | :--- |
| $(R \ell 2)$ | $(M ; \vee, \wedge, 0,1)$ is a bounded lattice. |
| $(R \ell 3)$ | $x \odot y \leq z$ if and only if $x \leq y \rightarrow z$, for any $x, y, z \in M$. |
| $(R \ell 4)$ | $x \odot(x \rightarrow y)=x \wedge y$, for any $x, y \in M$. |

In the sequel, by an $R \ell$-monoid we will mean a bounded commutative $R \ell$ monoid.

On any $R \ell$-monoid $M$ let us define a unary operation negation ${ }^{-}$by $x^{-}:=$ $x \rightarrow 0$ for any $x \in M$.

Bounded commutative $R \ell$-monoids are special cases of residuated lattices, more precisely (see for instance [3]), they are exactly commutative integral generalized $B L$-algebras in the sense of [10].

The above mentioned algebras can be characterized in the class of all $R \ell$ monoids as follows: An $R \ell$-monoid $M$ is
a) a $B L$-algebra if and only if $M$ satisfies the identity of pre-linearity

$$
(x \rightarrow y) \vee(y \rightarrow x)=1 ;
$$

b) an $M V$-algebra if and only if $M$ fulfills the double negation law

$$
x^{--}=x
$$

c) a Heyting algebra if and only if the operation " $\odot$ " is idempotent.

Lemma 2.1 See [15] and [16]. In any bounded commutative $R \ell$-monoid $M$ we have for any $x, y, z \in M$ :

$$
\text { (1) } 1 \rightarrow x=x \text {. }
$$

(2) $\quad x \leq y \Longleftrightarrow x \rightarrow y=1$.
(3) $x \odot y \leq x \wedge y$.
(4) $x \leq y \rightarrow x$.
(5) $(x \odot y) \rightarrow z=x \rightarrow(y \rightarrow z)=y \rightarrow(x \rightarrow z)$.
(6) $\quad(x \vee y) \rightarrow z=(x \rightarrow z) \wedge(y \rightarrow z)$.
(7) $\quad x \rightarrow(y \wedge z)=(x \rightarrow y) \wedge(x \rightarrow z)$.
(8) $\quad x \leq x^{--}, x^{-}=x^{---}$.
(9) $\quad x \leq y \Longrightarrow y^{-} \leq x^{-}$.
(10) $(x \odot y)^{-}=y \rightarrow x^{-}=y^{--} \rightarrow x^{-}=x \rightarrow y^{-}=x^{--} \rightarrow y^{-}$.
(11) $x \leq y \Longrightarrow z \rightarrow x \leq z \rightarrow y, y \rightarrow z \leq x \rightarrow z$.
(12) $x \rightarrow y \leq y^{-} \rightarrow x^{-}$.
(13) $\quad x \vee y \leq((x \rightarrow y) \rightarrow y) \wedge((y \rightarrow x) \rightarrow x)$.
(14) $x \rightarrow y \leq(y \rightarrow z) \rightarrow(x \rightarrow z)$.
(15) $x \rightarrow y \leq(z \rightarrow x) \rightarrow(z \rightarrow y)$.

A non-empty subset $F$ of an $R \ell$-monoid $M$ is called a filter of $M$ if
(F1) $\quad x, y \in F$ imply $x \odot y \in F$;
(F2) $\quad x \in F, y \in M, x \leq y$ imply $y \in F$.
A subset $D$ of an $R \ell$-monoid $M$ is called a deductive system of $M$ if
(i) $1 \in D$;
(ii) $\quad x \in D, x \rightarrow y \in D$ imply $y \in D$.

Proposition 2.2 [3]. Let $H$ be a non-empty subset of $M$. Then $H$ is a filter of $M$ if and only if $H$ is a deductive system of $M$.

By [18], filters of commutative $R \ell$-monoids are exactly the kernels of their congruences. If $F$ is a filter of $M$, then $F$ is the kernel of the unique congruence $\Theta(F)$ such that $\langle x, y\rangle \in \Theta(F)$ if and only if $(x \rightarrow y) \wedge(y \rightarrow x) \in F$, for any $x, y \in M$. Hence we will consider quotient $R \ell$-monoids $M / F$ of $R \ell$-monoids $M$ by their filters $F$.

A filter $F$ of $M$ is called maximal if $F$ is a proper filter of $M$ and is not a proper subset of any proper filter of $M$.

## 3 Implicative filters

Let $M$ be an $R \ell$-monoid and $F$ a subset of $M$. Then $F$ is called an implicative filter of $M$ if
(1) $1 \in F$;
(2) $\quad x \rightarrow(y \rightarrow z) \in F, x \rightarrow y \in F$ imply $x \rightarrow z \in F$.

Proposition 3.1 Every implicative filter of an $R \ell$-monoid $M$ is a filter of $M$.
Proof Let $\emptyset \neq F \subseteq M$ satisfy conditions (1) and (2) and let $x, y \in M$ be such that $x, x \rightarrow y \in F$. Then $1 \rightarrow(x \rightarrow y) \in F, 1 \rightarrow x \in F$, hence $y=1 \rightarrow y \in F$.

If $F$ is a filter of an $R \ell$-monoid $M$ and $a \in M$, put

$$
M_{a}:=\{x \in M: a \rightarrow x \in F\} .
$$

Theorem 3.2 Let $M$ be an $R \ell$-monoid and $F$ be a filter of $M$. Then $F$ is an implicative filter of $M$ if and only if $M_{a}$ is a filter of $M$ for every $a \in M$.

Proof Let $F$ be an implicative filter of $M$ and $a \in M$. Then $1=a \rightarrow 1 \in M$, thus $1 \in M_{a}$. Further, suppose that $x, x \rightarrow y \in M_{a}$, i.e. $a \rightarrow x \in F$ and $a \rightarrow(x \rightarrow y) \in F$. Then we get $a \rightarrow y \in F$, and hence $y \in M_{a}$. That means, $M_{a}$ is a filter of $M$ for arbitrary $a \in M$.

Conversely, let $M_{a}$ be a filter of $M$ for each $a \in M$. Suppose that $x \rightarrow(y \rightarrow$ $z) \in F$ and $x \rightarrow y \in F$. Then $y \rightarrow z \in M_{x}$ and $y \in M_{x}$, hence $z \in M_{x}$ and therefore $x \rightarrow z \in F$. That means, $F$ is implicative.

Theorem 3.3 Let $F$ be a filter of an $R \ell$-monoid $M$. Then the following conditions are equivalent:
(a) $F$ is an implicative filter of $M$.
(b) $y \rightarrow(y \rightarrow x) \in F$ implies $y \rightarrow x \in F$, for any $x, y \in M$.
(c) $z \rightarrow(y \rightarrow x) \in F$ implies $(z \rightarrow y) \rightarrow(z \rightarrow x) \in F$, for any $x, y, z \in M$.
(d) $z \rightarrow(y \rightarrow(y \rightarrow x)) \in F$ and $z \in F$ imply $y \rightarrow x \in F$, for any $x, y, z \in M$.
(e) $x \rightarrow(x \odot x) \in F$, for any $x \in M$.

Proof $(\mathrm{a}) \Rightarrow(\mathrm{b})$ : Suppose that $F$ is an implicative filter of $M, x, y \in M$ and $y \rightarrow(y \rightarrow x) \in F$. Then since $y \rightarrow y=1 \in F$, we obtain $y \rightarrow x \in F$.
(b) $\Rightarrow(\mathrm{c})$ : Let $F$ be a filter of $M$ satisfying the condition (b), $x, y, z \in M$ and $z \rightarrow(y \rightarrow x) \in F$. Then $z \rightarrow(z \rightarrow((z \rightarrow y) \rightarrow x))=z \rightarrow((z \rightarrow y) \rightarrow$ $(z \rightarrow x)) \geq z \rightarrow(y \rightarrow x) \in F$, thus $z \rightarrow(z \rightarrow((z \rightarrow y) \rightarrow x)) \in F$. From this we have $z \rightarrow((z \rightarrow y) \rightarrow x) \in F$, that means $(z \rightarrow y) \rightarrow(z \rightarrow x) \in F$.
$(\mathrm{c}) \Rightarrow(\mathrm{d})$ : Suppose that a filter $F$ satisfies the condition (c). Let $z \rightarrow(y \rightarrow$ $(y \rightarrow x)) \in F$ and $z \in F$. Then also $y \rightarrow(y \rightarrow x) \in F$. At the same time, $y \rightarrow x=(y \rightarrow y) \rightarrow(y \rightarrow x)$, thus $y \rightarrow x \in F$.
(d) $\Rightarrow$ (a): Let a filter $F$ fulfill the condition (d). Let $x \rightarrow(y \rightarrow z) \in F$ and $x \rightarrow y \in F$. Then $x \rightarrow(y \rightarrow z)=y \rightarrow(x \rightarrow z) \leq(x \rightarrow y) \rightarrow(x \rightarrow(x \rightarrow z))$, hence $(x \rightarrow y) \rightarrow(x \rightarrow(x \rightarrow z)) \in F$, and therefore $x \rightarrow z \in F$.
(a) $\Rightarrow(\mathrm{e})$ : Let $F$ be an implicative filter of $M$. Then $x \rightarrow(x \rightarrow(x \odot x))=$ $(x \odot x) \rightarrow(x \odot x)=1 \in F$. Further, $x \rightarrow x=1 \in F$, and hence we obtain $x \rightarrow(x \odot x) \in F$.
(e) $\Rightarrow(\mathrm{a})$ : Let a filter $F$ satisfy the condition (e) and let $x \rightarrow(y \rightarrow z) \in F$ and $x \rightarrow y \in F$. Then $(x \rightarrow(y \rightarrow z)) \odot(x \rightarrow y) \odot x \odot x \leq(y \rightarrow z) \odot y \leq z$, hence $(x \rightarrow(y \rightarrow z)) \odot(x \rightarrow y) \leq(x \odot x) \rightarrow z$, and thus $(x \odot x) \rightarrow z \in F$. Further, $x \rightarrow(x \odot x) \in F,(x \odot x) \rightarrow x=1 \in F$, therefore from $(x \odot x) \rightarrow z \in F$, we obtain $x \rightarrow z \in F$.

Using the proof $(\mathrm{a}) \Rightarrow(\mathrm{e})$ in the preceding theorem, we have as an immediate consequence:

Theorem 3.4 If $F$ is a filter of an $R \ell$-monoid $M$, then $F$ is an implicative filter if and only if the quotient $R \ell$-monoid $M / F$ is a Heyting algebra.

Proposition 3.5 If $F_{1}$ and $F_{2}$ are filters of an $R \ell$-monoid $M, F_{1} \subseteq F_{2}$ and $F_{1}$ is an implicative filter of $M$, then $F_{2}$ is also an implicative filter of $M$.

Proof Suppose that $F_{1}$ and $F_{2}$ are filters of an $R \ell$-monoid $M, F_{1} \subseteq F_{2}$ and $F_{1}$ is implicative. Then, by Theorem 3.3, $x \rightarrow x \odot x \in F_{1} \subseteq F_{2}$ for any $x \in M$, and therefore $F_{2}$ is also implicative.

Let $M$ be an $R \ell$-monoid and $F$ a subset of $M$. Then $F$ is called a positive implicative filter of $M$ if
(1) $1 \in F$;
(3) $\quad x \rightarrow((y \rightarrow z) \rightarrow y) \in F$ and $x \in F$ imply $y \in F$, for any $x, y, z \in M$.

Proposition 3.6 Every positive implicative filter of an $R \ell$-monoid $M$ is a filter of $M$.

Proof Let $x \in F$ and $x \rightarrow y \in F$. Then $x \rightarrow((y \rightarrow 1) \rightarrow y)=x \rightarrow(1 \rightarrow y)=$ $x \rightarrow y$, hence $x \rightarrow((y \rightarrow 1) \rightarrow y) \in F$, and thus $y \in F$.

Proposition 3.7 Every positive implicative filter of $M$ is an implicative filter of $M$.

Proof Let $F$ be a positive implicative filter of $M, x, y, z \in M, x \rightarrow(y \rightarrow z) \in$ $F$ and $x \rightarrow y \in F$. We have $(x \rightarrow y) \rightarrow(x \rightarrow(x \rightarrow z)) \geq y \rightarrow(x \rightarrow z)=x \rightarrow$ $(y \rightarrow z)$, hence $(x \rightarrow y) \rightarrow(x \rightarrow(x \rightarrow z)) \in F$, and thus also $x \rightarrow(x \rightarrow z) \in F$.

Since $((x \rightarrow z) \rightarrow z) \rightarrow(x \rightarrow z) \geq x \rightarrow(x \rightarrow z)$, then we get $((x \rightarrow z) \rightarrow$ $z) \rightarrow(x \rightarrow z) \in F$. Further, $1 \rightarrow(((x \rightarrow z) \rightarrow z) \rightarrow(x \rightarrow z))=((x \rightarrow z) \rightarrow$ $z) \rightarrow(x \rightarrow z)$, and since $1 \rightarrow(((x \rightarrow z) \rightarrow z) \rightarrow(x \rightarrow z)) \in F$ and $1 \in F$, we obtain $x \rightarrow z \in F$.

Therefore $F$ is an implicative filter.

Theorem 3.8 Let $F$ be a filter of an $R \ell$-monoid $M$. Then the following conditions are equivalent:
(a) $F$ is a positive implicative filter of $M$.
(b) $(x \rightarrow y) \rightarrow x \in F$ implies $x \in F$, for any $x, y \in M$.
(c) $\left(x^{-} \rightarrow x\right) \rightarrow x \in F$, for any $x \in M$.

Proof $\quad(\mathrm{a}) \Rightarrow(\mathrm{b})$ : Let $F$ be a positive implicative filter of $M$ and $(x \rightarrow y) \rightarrow$ $x \in F$. Then since $1 \rightarrow((x \rightarrow y) \rightarrow x)=(x \rightarrow y) \rightarrow x \in F$ and $1 \in F$, we get $x \in F$.
(b) $\Rightarrow$ (a): Let a filter $F$ satisfy the condition (b) and let $x \rightarrow((y \rightarrow z) \rightarrow$ $y) \in F$ and $x \in F$. Then $(y \rightarrow z) \rightarrow y \in F$, and therefore $y \in F$. Hence $F$ is a positive implicative filter of $M$.
(b) $\Rightarrow$ (c): Let $F$ be a filter of $M$ and $x \in M$. Then $\left(\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right) \rightarrow\right.$ $0) \rightarrow\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right)=\left(x^{-} \rightarrow x\right) \rightarrow\left(\left(\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right) \rightarrow 0\right) \rightarrow x\right) \geq\left(\left(\left(x^{-} \rightarrow\right.\right.\right.$ $\left.x) \rightarrow x) \rightarrow 0) \rightarrow x^{-}=\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right) \rightarrow 0\right) \rightarrow(x \rightarrow 0) \geq x \rightarrow\left(\left(x^{-} \rightarrow x\right) \rightarrow\right.$ $x)=1 \in F$, thus $\left(\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right) \rightarrow 0\right) \rightarrow\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right) \in F$, and hence $\left(x^{-} \rightarrow x\right) \rightarrow x \in F$.
$(\mathrm{c}) \Rightarrow(\mathrm{b})$ : Let a filter $F$ satisfy condition (c). Let $(x \rightarrow y) \rightarrow x \in F$. We have $(x \rightarrow y) \rightarrow x \leq(x \rightarrow 0) \rightarrow x=x^{-} \rightarrow x$, hence $x^{-} \rightarrow x \in F$. By the assumption, $\left(x^{-} \rightarrow x\right) \rightarrow x \in F$, thus $x \in F$. Therefore $F$ satisfies the condition (b).

Proposition 3.9 If $F_{1}$ and $F_{2}$ are filters of an $R \ell$-monoid $M, F_{1}$ is a positive implicative filter and $F_{1} \subseteq F_{2}$, then $F_{2}$ is also a positive implicative filter of $M$.

Proof Let $F_{1} \subseteq F_{2}$ and $F_{1}$ be positive implicative. Then for any $x \in M$ we get $\left(x^{-} \rightarrow x\right) \rightarrow x \in F_{1}$, thus $\left(x^{-} \rightarrow x\right) \rightarrow x \in F_{2}$. Therefore, by Theorem 3.8, $F_{2}$ is a positive implicative filter of $M$.

Theorem 3.10 Let $M$ be an $R \ell$-monoid. Then the following conditions are equivalent:
(a) $M$ is a Heyting algebra.
(b) Every filter of $M$ is implicative.
(c) $\{1\}$ is an implicative filter of $M$.

Proof $\quad(\mathrm{a}) \Rightarrow(\mathrm{c})$ : It follows from Theorem 3.4.
(a) $\Rightarrow(\mathrm{b})$ : Let $M$ be an idempotent $R \ell$-monoid, $F$ be a filter of $M$, and $x \in M$. Then $x \rightarrow(x \odot x)=x \rightarrow x=1 \in F$, hence by Theorem 3.3, $F$ is an implicative filter.
$(\mathrm{b}) \Rightarrow(\mathrm{c})$ : It is obvious.

Proposition 3.11 Let $F$ be an implicative filter of an $R \ell$-monoid $M$. Then the following conditions are equivalent:
(a) $F$ is a positive implicative filter of $M$.
(b) $(x \rightarrow y) \rightarrow y \in F$ implies $(y \rightarrow x) \rightarrow x \in F$, for any $x, y \in M$.

Proof $(\mathrm{a}) \Rightarrow(\mathrm{b})$ : Let $F$ be a positive implicative filter of $M$ and $(x \rightarrow y) \rightarrow y \in$ $F$. Since $x \leq(y \rightarrow x) \rightarrow x$, we get $((y \rightarrow x) \rightarrow x) \rightarrow y \leq x \rightarrow y$. Hence $(x \rightarrow$ $y) \rightarrow y \leq(y \rightarrow x) \rightarrow((x \rightarrow y) \rightarrow x)=(x \rightarrow y) \rightarrow((y \rightarrow x) \rightarrow x) \leq(((y \rightarrow$ $x) \rightarrow x) \rightarrow y) \rightarrow((y \rightarrow x) \rightarrow x)$, and thus $(((y \rightarrow x) \rightarrow x) \rightarrow y) \rightarrow((y \rightarrow x) \rightarrow$ $x) \in F$. Consequently, also $1 \rightarrow((((y \rightarrow x) \rightarrow x) \rightarrow y) \rightarrow((y \rightarrow x) \rightarrow x)) \in F$, and since $F$ is a positive implicative filter, we get $(y \rightarrow x) \rightarrow x \in F$.
(b) $\Rightarrow$ (a): Let an implicative filter $F$ satisfy the condition (b) and let $x \in F$ and $x \rightarrow((y \rightarrow z) \rightarrow y) \in F$. Then also $(y \rightarrow z) \rightarrow y \in F$. Further, $(y \rightarrow z) \rightarrow$ $y \leq(y \rightarrow z) \rightarrow((y \rightarrow z) \rightarrow z)$, hence $(y \rightarrow z) \rightarrow((y \rightarrow z) \rightarrow z) \in F$. Since $F$ is implicative, $(y \rightarrow z) \rightarrow z \in F$. Then, by the assumption, also $(z \rightarrow y) \rightarrow y \in F$. Further, $z \leq y \rightarrow z$, hence $(y \rightarrow z) \rightarrow y \leq z \rightarrow y$, thus $z \rightarrow y \in F$. We have shown $(z \rightarrow y) \rightarrow y \in F$, therefore $y \in F$.

Theorem 3.12 Let $M$ be an $R \ell$-monoid. Then the following conditions are equivalent:
(a) $\{1\}$ is a positive implicative filter.
(b) Every filter of $M$ is positive implicative.
(c) $M(a):=\{x \in M: a \leq x\}$ is a positive implicative filter of $M$, for every $a \in M$.
(d) $(x \rightarrow y) \rightarrow x=x$, for any $x, y \in M$.
(e) $M$ is a Boolean algebra.

Proof $(\mathrm{a}) \Rightarrow(\mathrm{b})$ : It follows from Proposition 3.9.
(b) $\Rightarrow(\mathrm{c})$ : Let $a \in M$. Then $1 \in M(a)$. Assume that $x, x \rightarrow y \in M(a)$, i.e. $a \rightarrow x=1, a \rightarrow(x \rightarrow y)=1$. Since by the assumption, $\{1\}$ is a positive implicative filter of $M$, we obtain $a \rightarrow y=1$, hence $y \in M(a)$. That means $M(a)$ is a filter of $M$ which is also positive implicative.
(c) $\Rightarrow$ (d): If $x, y \in M$, then $(x \rightarrow y) \rightarrow x \in M((x \rightarrow y) \rightarrow x)$, therefore $(x \rightarrow y) \rightarrow x \leq x$ by Theorem 3.8. Moreover, $x \leq(x \rightarrow y) \rightarrow x$, i.e. $(x \rightarrow y) \rightarrow$ $x=x$.
(d) $\Rightarrow$ (a): It follows from Theorem 3.8.
$(\mathrm{d}) \Rightarrow(\mathrm{e}):$ Since $(x \rightarrow y) \rightarrow x=x$, we obtain $(y \rightarrow x) \rightarrow x=(y \rightarrow x) \rightarrow$ $((x \rightarrow y) \rightarrow x) \geq(x \rightarrow y) \rightarrow y$, and similarly, $(x \rightarrow y) \rightarrow y \geq(y \rightarrow x) \rightarrow x$. Hence $x^{--}=(x \rightarrow 0) \rightarrow 0=(0 \rightarrow x) \rightarrow x=1 \rightarrow x=x$ and therefore by [12], $M$ is an $M V$-algebra. Then by [7, Lemma 3.16], furthermore $M$ is a Boolean algebra.
$(\mathrm{e}) \Rightarrow(\mathrm{d})$ : Since $M$ is a Boolean algebra, $x^{-}$is the lattice complement of $x$ in $M$, and so $x \vee x^{-}=1$. This implies, by [7, Lemma 3.16], $(x \rightarrow y) \rightarrow x=x$ for any $x, y \in M$.

Theorem 3.13 If $F$ is a filter of an $R \ell$-monoid $M$, then the following conditions are equivalent:
(a) $F$ is a maximal and positive implicative filter of $M$.
(b) $F$ is a maximal and implicative filter of $M$.
(c) If $x, y \in M \backslash F$, then $x \rightarrow y \in F$ and $y \rightarrow x \in F$.
(d) $M / F$ is a two-element Boolean algebra.

Proof $(\mathrm{a}) \Rightarrow(\mathrm{b})$ : It is obvious.
(b) $\Rightarrow$ (c): Let $F$ be a maximal and implicative filter of $M$. By Theorem $3.2, M_{y}=\{a \in M: y \rightarrow a \in F\}$ is a filter of $M$. If $b \in F$, then from $b \leq y \rightarrow b$ it follows that $y \rightarrow b \in F$, thus $b \in M_{y}$. Hence $F \subseteq M_{y}$. Since $F$ is a maximal filter of $M$ and $y \notin F$, we have $M_{y}=M$. Therefore $y \rightarrow x \in F$. The assumption $x \notin F$ analogously implies $x \rightarrow y \in F$.
(c) $\Rightarrow$ (a): Let a filter $F$ satisfy the condition (c). Suppose that $F$ is not positive implicative. Then by Theorem 3.8 , there are $x, y \in M$ such that $x \notin F$ and $(x \rightarrow y) \rightarrow x \in F$. If $y \in F$, then $x \rightarrow y \in F$, and hence $x \in F$, a contradiction. If $y \notin F$, then by (c), $x \rightarrow y \in F$, a contradiction. Hence $F$ is a positive implicative filter of $M$. We will prove that $F$ is also a maximal filter of $M$. If $a \notin F$, then by the preceding part of the proof, $F \cup\{a\} \subseteq M_{a}$. We will show that $M_{a}$ is the least filter of $M$ containing $F \cup\{a\}$. Let $G$ be a filter of $M$ such that $F \cup\{a\} \subseteq G$. If $x \in M_{a}$, then $a \rightarrow x \in F \subseteq G$, and since $a \in G$, we have $x \in G$. Therefore $M_{a} \subseteq G$. Consider any element $z \in M$. If $z \in F$, then $z \in M_{a}$. If $z \notin F$, then since also $a \notin F$, the assumption (c) gives $a \rightarrow z \in M_{a}$. Hence $M_{a}=M$, and therefore $F$ is a maximal filter of $M$.
$(\mathrm{c}) \Rightarrow(\mathrm{d})$ : It is obvious.
A filter $F$ of an $R \ell$-monoid $M$ is called
a) Boolean if $x \vee x^{-} \in F$ for every $x \in M$;
b) semi-Boolean if $\left(x \wedge x^{-}\right)^{-} \in F$ for every $x \in M$.

Proposition 3.14 [14, Theorem 3.2]. If $F$ is a filter of an $R \ell$-monoid $M$, then $F$ is Boolean if and only if $M / F$ is a Boolean algebra.

Proposition 3.15 Every Boolean filter of $M$ is semi-Boolean.
Proof Let $x \in M$. Then $x^{-} \leq\left(x \wedge x^{-}\right)^{-}$and $x \leq x^{--} \leq\left(x \wedge x^{-}\right)^{-}$, hence $x \vee x^{-} \leq\left(x \wedge x^{-}\right)^{-}$.

Example 3.16 Let $M=\{0, a, b, c, 1\}$ be the lattice with the diagram in Fig. 1, and let $\odot=\wedge$ and $\rightarrow$ be defined in the corresponding table in Fig. 1.

| $\rightarrow$ | 0 | a | b | c | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 1 | 1 | 1 | 1 |
| a | b | 1 | b | 1 | 1 |
| b | a | a | 1 | 1 | 1 |
| c | 0 | a | b | 1 | 1 |
| 1 | 0 | a | b | c | 1 |



Fig. 1
Then $M=(M ; \vee, \wedge, \odot, \rightarrow, 0,1)$ is an $R \ell$-monoid (which is not a $B L$-algebra). The filter $F=\{1\}$ is semi-Boolean, but it is not Boolean.

Theorem 3.17 a) Let $M$ be an $R \ell$-monoid. Then every Boolean filter of $M$ is positive implicative and every positive implicative filter of $M$ is semi-Boolean.
b) If an $R \ell$-monoid $M$ satisfies condition

$$
\begin{equation*}
\left(\left(x \rightarrow x^{-}\right) \rightarrow x^{-}\right) \wedge\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right)=x \vee x^{-}, \text {for any } x \in M \tag{*}
\end{equation*}
$$

then Boolean and positive implicative filters of $M$ coincide.
Proof a) Let $M$ be an $R \ell$-monoid, let $F$ be a Boolean filter of $M$ and let $x \in M$. Then by Lemma 2.1, $x \vee x^{-} \leq\left(\left(x \rightarrow x^{-}\right) \rightarrow x^{-}\right) \wedge\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right)$, hence $\left(\left(x \rightarrow x^{-}\right) \rightarrow x^{-}\right) \wedge\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right) \in F$, and therefore $\left(x^{-} \rightarrow x\right) \rightarrow x \in F$. That means $F$ is positive implicative.

Let now $F$ be an arbitrary positive implicative filter of $M$ and $x \in M$. Then $\left(x^{--} \rightarrow x^{-}\right) \rightarrow x^{-} \in F$ and by Lemma 2.1, $\left(x^{--} \rightarrow x^{-}\right) \rightarrow x^{-}=\left(x \rightarrow x^{-}\right) \rightarrow$ $x^{-}=\left(\left(x \rightarrow x^{-}\right) \odot x\right)^{-}=\left(x \wedge x^{-}\right)^{-}$. Thus $F$ is a semi-Boolean filter.
b) Let an $R \ell$-monoid $M$ satisfy condition (*) and let $F$ be a positive implicative filter of $M$. Then a fortiori $F$ is also implicative, hence $x \rightarrow(x \odot x) \in F$ for every $x \in M$. We have $\left(x \rightarrow x^{-}\right) \rightarrow x^{-}=(x \rightarrow(x \rightarrow 0)) \rightarrow(x \rightarrow 0)=$ $((x \odot x) \rightarrow 0) \rightarrow(x \rightarrow 0) \geq x \rightarrow(x \odot x)$, hence $\left(x \rightarrow x^{-}\right) \rightarrow x^{-} \in F$, and thus also $x \vee x^{-}=\left(\left(x \rightarrow x^{-}\right) \rightarrow x^{-}\right) \wedge\left(\left(x^{-} \rightarrow x\right) \rightarrow x\right) \in F$. Therefore $F$ is a Boolean filter.

As an immediate consequence we get the following theorem.
Theorem 3.18 [11, Theorem 2]. Boolean and positive implicative filters of any BL-algebra coincide.

Proof If $M$ is a $B L$-algebra, then by [5, Lemma 2.3.4(8)], $((x \rightarrow y) \rightarrow y) \wedge((y \rightarrow$ $x) \rightarrow x)=x \vee y$, for every $x, y \in M$.

Let $F$ be a filter of an $R \ell$-monoid $M$. Then $F$ is called an implicative deductive system if $x \rightarrow\left(z^{-} \rightarrow y\right) \in F$ and $y \rightarrow z \in F$ imply $x \rightarrow z \in F$, for any $x, y, z \in M$.

Theorem 3.19 [14, Theorem 3.2]. Let $F$ be a filter of an $R \ell$-monoid $M$. Then $F$ is an implicative deductive system if and only if $F$ is a Boolean filter.

Remark 3.20 Now we can rephrase Theorem 3.17 in this way. Let $M$ be an $R \ell$-monoid. Then every implicative deductive system of $M$ is a positive implicative filter and every positive implicative filter of $M$ is semi-Boolean. If $M$ satisfies the condition $(*)$, then implicative deductive systems and positive implicative filters of $M$ coincide.

Theorem 3.21 If $F$ is a maximal and (positive) implicative filter of an $R \ell$ monoid $M$, then $F$ is Boolean.

Proof Let $F$ be a maximal and (positive) implicative filter of $M$. Then by Theorem 3.13, $M / F$ is a two element $R \ell$-monoid, hence a two element Boolean algebra. Consequently, by Proposition 3.14, $F$ is a Boolean filter.

Theorem 3.22 If $F$ is a maximal filter of an $R \ell$-monoid $M$, then the following conditions are equivalent:
(a) $F$ is a Boolean filter.
(b) $F$ is a positive implicative filter.
(c) $F$ is an implicative filter.
(d) $F$ is an implicative deductive system.

Proof It follows from Theorems 3.17 and 3.21 and from Remark 3.20.
Let $M$ be an $R \ell$-monoid. If $F$ is a proper filter of $M$, denote

$$
F^{-}:=\left\{x \in M: x \leq y^{-} \text {for some } y \in F\right\} .
$$

By [14, Proposition 3.4], $F \cup F^{-}$is a subalgebra of $M$ for every proper filter $F$ of $M$.

An $R \ell$-monoid $M$ is called bipartite if $M=F \cup F^{-}$for some maximal filter $F$ of $M$.

By [14, Theorem 3.6], $M$ is bipartite if and only if $M$ contains a proper Boolean filter.

An $R \ell$-monoid $M$ is said to be strongly bipartite if $M=F \cup F^{-}$for every maximal filter $F$ of $M$.

If $M$ is an $R \ell$-monoid, denote by $B(M)$ the intersection of all Boolean filters of $M$. Obviously $B(M)$ is the least Boolean filter of $M$.

Further, denote by $\operatorname{Rad}(M)$ the radical of $M$, i.e. the intersection of all maximal filters of $M$.

Theorem 3.23 [14, Theorem 3.8]. If $M$ is an $R \ell$-monoid, then the following conditions are equivalent:
(a) $M$ is strongly bipartite.
(b) Every maximal filter of $M$ is Boolean.
(c) $B(M) \subseteq \operatorname{Rad}(M)$.

The following theorem is an immediate consequence of Theorems 3.22 and 3.23 .

Theorem 3.24 If $M$ is an $R \ell$-monoid, then the following conditions are equivalent:
(a) $M$ is strongly bipartite.
(b) $B(M) \subseteq \operatorname{Rad}(M)$.
(c) Every maximal filter of $M$ is Boolean.
(d) Every maximal filter of $M$ is positive implicative.
(e) Every maximal filter of $M$ is implicative.

## 4 Fantastic filters

Let $M$ be an $R \ell$-monoid and $F$ a subset of $M$. Then $F$ is called a fantastic filter of $M$ if
(1) $1 \in F$;
(4) $z \rightarrow(y \rightarrow x) \in F$ and $z \in F$ imply $((x \rightarrow y) \rightarrow y) \rightarrow x \in F$, for any $x, y, z \in M$.
Proposition 4.1 Every fantastic filter of $M$ is a filter of $M$.
Proof Let $F$ be a fantastic filter of $M$ and $x, y \in M$. If $x, x \rightarrow y \in F$, then also $x \in F$ and $x \rightarrow(1 \rightarrow y)=x \rightarrow y \in F$, and thus by (4), $y \in F$.

Theorem 4.2 A filter $F$ of an $R \ell$-monoid $M$ is fantastic if and only if
(5) $y \rightarrow x \in F$ implies $((x \rightarrow y) \rightarrow y) \rightarrow x \in F$, for every $x, y \in M$.

Proof Let $F$ be a fantastic filter of $M, x, y \in M$ and $y \rightarrow x \in F$. Then $1 \rightarrow(y \rightarrow x)=y \rightarrow x \in F$ and $1 \in F$, hence $((x \rightarrow y) \rightarrow y) \rightarrow x \in F$.

Conversely, let a filter $F$ satisfy the condition (5) and let $z \rightarrow(y \rightarrow x) \in F$ and $z \in F$. Then $y \rightarrow x \in F$, therefore also $((x \rightarrow y) \rightarrow y) \rightarrow x \in F$.

Theorem 4.3 Every positive implicative filter of an $R \ell$-monoid $M$ is a fantastic filter of $M$.

Proof Suppose $F$ is a positive implicative filter of $M$ and $x, y \in M$ are such that $y \rightarrow x \in F$. We have $x \leq((x \rightarrow y) \rightarrow y) \rightarrow x$, thus

$$
(((x \rightarrow y) \rightarrow y) \rightarrow x) \rightarrow y \leq x \rightarrow y
$$

Further, $((((x \rightarrow y) \rightarrow y) \rightarrow x) \rightarrow y) \rightarrow(((x \rightarrow y) \rightarrow y) \rightarrow x) \geq(x \rightarrow y) \rightarrow$ $(((x \rightarrow y) \rightarrow y) \rightarrow x)=((x \rightarrow y) \rightarrow y) \rightarrow((x \rightarrow y) \rightarrow x) \geq y \rightarrow x$.

By the assumption $y \rightarrow x \in F$, hence also

$$
((((x \rightarrow y) \rightarrow y) \rightarrow x) \rightarrow y) \rightarrow(((x \rightarrow y) \rightarrow y) \rightarrow x) \in F
$$

Since $F$ is positive implicative, we get $((x \rightarrow y) \rightarrow y) \rightarrow x \in F$, and hence $F$ is a fantastic filter.

Theorem 4.4 If $F$ is a filter of an $R \ell$-monoid $M$, then the following conditions are equivalent:
(a) $F$ is a fantastic filter of $M$.
(b) $x^{--} \rightarrow x \in F$, for every $x \in M$.
(c) $x \rightarrow u \in F$ and $y \rightarrow u \in F$ imply $((x \rightarrow y) \rightarrow y) \rightarrow u \in F$, for every $x, y, u \in M$.

Proof $(\mathrm{a}) \Rightarrow(\mathrm{b})$ : Let $F$ be a fantastic filter of $M$ and $x \in M$. Since $0 \rightarrow x=$ $1 \in F$, we obtain from (5) that $x^{--} \rightarrow x=((x \rightarrow 0) \rightarrow 0) \rightarrow x \in F$.
(b) $\Rightarrow$ (c): Suppose that $F$ is a filter of $M$ such that $x^{--} \rightarrow x \in F$ for every $x \in M$. Let $x, y, u \in M, x \rightarrow u \in F$ and $y \rightarrow u \in F$. Since $x \rightarrow u \leq u^{-} \rightarrow x^{-}$ and $y \rightarrow u \leq u^{-} \rightarrow y^{-}$, we get $u^{-} \rightarrow x^{-} \in F$ and $u^{-} \rightarrow y^{-} \in F$, and thus $\left(u^{-} \rightarrow x^{-}\right) \wedge\left(u^{-} \rightarrow y^{-}\right) \in F$.

Moreover,

$$
\begin{aligned}
& \left(u^{-} \rightarrow x^{-}\right) \wedge\left(u^{-} \rightarrow y^{-}\right)=u^{-} \rightarrow\left(x^{-} \wedge y^{-}\right) \\
& \quad=u^{-} \rightarrow\left(y^{-} \odot\left(y^{-} \rightarrow x^{-}\right)\right)=u^{-} \rightarrow\left(y^{-} \odot\left(y^{-} \rightarrow(x \rightarrow 0)\right)\right. \\
& \quad=u^{-} \rightarrow\left(y^{-} \odot\left(x \rightarrow\left(y^{-} \rightarrow 0\right)\right)=u^{-} \rightarrow\left(y^{-} \odot\left(x \rightarrow y^{--}\right)\right)\right.
\end{aligned}
$$

Further,

$$
\begin{aligned}
& \left(u^{-} \rightarrow\left(y^{-} \odot\left(x \rightarrow y^{--}\right)\right)\right) \rightarrow\left(u^{-} \rightarrow\left(y^{-} \odot(x \rightarrow y)\right)\right) \\
& \left.\quad \geq\left(y^{-} \odot\left(x \rightarrow y^{--}\right)\right) \rightarrow\left(y^{-} \odot(x \rightarrow y)\right)\right) \\
& \quad \geq\left(x \rightarrow y^{--}\right) \rightarrow(x \rightarrow y) \geq y^{--} \rightarrow y \in F
\end{aligned}
$$

therefore also $u^{-} \rightarrow\left(y^{-} \odot(x \rightarrow y)\right) \in F$.
Moreover,

$$
u^{-} \rightarrow\left(y^{-} \odot(x \rightarrow y)\right) \leq\left(y^{-} \odot(x \rightarrow y)\right)^{-} \rightarrow u^{--}=\left((x \rightarrow y) \rightarrow y^{--}\right) \rightarrow u^{--}
$$

hence $\left((x \rightarrow y) \rightarrow y^{--} \rightarrow u^{--} \in F\right.$. Further we have

$$
\begin{aligned}
& \left(\left((x \rightarrow y) \rightarrow y^{--}\right) \rightarrow u^{--}\right) \rightarrow\left(((x \rightarrow y) \rightarrow y) \rightarrow u^{--}\right) \\
& \quad \geq((x \rightarrow y) \rightarrow y) \rightarrow\left((x \rightarrow y) \rightarrow y^{--}\right) \geq y \rightarrow y^{--}=1 \in F
\end{aligned}
$$

thus $((x \rightarrow y) \rightarrow y) \rightarrow u^{--} \in F$.
Moreover,

$$
\left(((x \rightarrow y) \rightarrow y) \rightarrow u^{--}\right) \rightarrow(((x \rightarrow y) \rightarrow y) \rightarrow u) \geq u^{--} \rightarrow u \in F
$$

therefore also $((x \rightarrow y) \rightarrow y) \rightarrow u \in F$.
(c) $\Rightarrow(\mathrm{a})$ : If $F$ satisfies the condition (c), then for $u=x$ we get that whether $y \rightarrow x \in F$ then $((x \rightarrow y) \rightarrow y) \rightarrow x \in F$, for every $x, y \in M$, hence $F$ is a fantastic filter of $M$.

Theorem 4.5 If $F_{1}$ and $F_{2}$ are filters of an $R \ell$-monoid $M, F_{1} \subseteq F_{2}$ and $F_{1}$ is fantastic in $M$, then $F_{2}$ is also a fantastic filter of $M$.

Proof Let $F_{1}$ and $F_{2}$ be filters of $M, F_{1} \subseteq F_{2}$, and let $F_{1}$ be fantastic. Then by Theorem 4.4, $x^{--} \rightarrow x \in F_{1} \subseteq F_{2}$, for every $x \in M$, hence $F_{2}$ is also fantastic.

Theorem 4.6 $A$ filter $F$ of an $R \ell$-monoid $M$ is fantastic if and only if $M / F$ is an $M V$-algebra.

Proof Let $F$ be a filter of $M$. Then $F$ is fantastic if and only if $x^{--} \rightarrow x \in F$ for every $x \in M$, which is equivalent to the following conditions in $M / F$ :

$$
x^{--} / F \rightarrow x / F=F, \quad x^{--} / F \leq x / F \text { and } x^{--} / F=x / F
$$

for every $x / F \in M / F$, and this is equivalent to $M / F$ is an $M V$-algebra.
Proposition 4.7 If $F$ is a maximal filter of an $R \ell$-monoid $M$, then $F$ is fantastic.

Proof It follows from [3, Proposition 3.5], where it is proved that $M / F$ is an $M V$-algebra for every maximal filter $F$ of $M$.

Remark 4.8 The $M V$-filters of $R \ell$-monoids, i.e. filters such that the corresponding quotient $R \ell$-monoids are $M V$-algebras, were investigated in [16], [17] and [3]. By Theorem 4.6, $M V$-filters of $R \ell$-monoids are exactly their fantastic filters. If $M$ is an $R \ell$-monoid, denote by $D(M):=\left\{x \in M: x^{--}=1\right\}$ the set of all dense elements in $M$. Then $D(M)$ is a proper filter of $M$ and a filter $F$ of $M$ is an $M V$-filter if and only if $D(M) \subseteq F$. Therefore we get as a consequence the following proposition.

Proposition 4.9 A filter $F$ of an $R \ell$-monoid $M$ is fantastic if and only if $D(M) \subseteq F$.
Proposition 4.10 Let $M$ be an $R \ell$-monoid. Then the following conditions are equivalent:
(1) $M$ is an $M V$-algebra.
(2) Every filter of $M$ is fantastic.
(3) $\{1\}$ is a fantastic filter of $M$.

Proof $(1) \Rightarrow(2)$ : Let $M$ be an $M V$-algebra and $F$ be a filter of $M$. Since the class of $M V$-algebras is a subvariety of the variety of $R \ell$-monoids, the quotient $R \ell$-monoid $M / F$ is also an $M V$-algebra. Therefore by Theorem 4.6, $F$ is a fantastic filter.
$(2) \Rightarrow(3)$ : It is obvious.
$(3) \Rightarrow(1)$ : Let $\{1\}$ be a fantastic filter of $M$. Then $M \cong M /\{1\}$ is an $M V$-algebra.

Theorem 4.11 If $F$ is a filter of an $R \ell$-monoid $M$, then the following conditions are equivalent.
(a) $F$ is a Boolean filter.
(b) $F$ is an implicative and fantastic filter.

Proof By Proposition 3.14, a filter $F$ is Boolean if and only if $M / F$ is a Boolean algebra. Moreover, an $R \ell$-monoid $M / F$ is a Boolean algebra if and only if $M / F$ is an $M V$-algebra and $(x / F) \odot(x / F)=x / F$ for every $x / F \in M / F$. This is equivalent to $(x / F)^{--}=x / F$ and $(x / F) \odot(x / F)=x / F$, and it holds, by Theorems 4.6 and 3.4, if and only if $F$ is a fantastic and implicative filter of $M$.

We have characterized filters of $R \ell$-monoids such that the corresponding quotient $R \ell$-monoids are Heyting algebras, Boolean algebras and $M V$-algebras, respectively. (See e.g. Theorem 3.4, Proposition 3.14 and Theorem 4.6.) Now we will complete it for the case when the quotient $R \ell$-monoid is a $B L$-algebra.

A filter $F$ of an $R \ell$-monoid $M$ is called a $B L$-filter of $M$ if

$$
(x \rightarrow y) \vee(y \rightarrow x) \in F
$$

for every $x, y \in M$.
Theorem 4.12 A filter $F$ of an $R \ell$-monoid $M$ is a $B L$-filter of $M$ if and only if $M / F$ is a $B L$-algebra.

Proof We know that an $R \ell$-monoid is a $B L$-algebra if and only if it satisfies the identity of pre-linearity.

Let $M$ be an $R \ell$-monoid and $F$ be a filter of $M$. If $x, y \in M$, then

$$
(x / F \rightarrow y / F) \vee(y / F \rightarrow x / F)=((x \rightarrow y) \vee(y \rightarrow x)) / F
$$

Hence $(x / F \rightarrow y / F) \vee(y / F \rightarrow x / F)=F$ if and only if $(x \rightarrow y) \vee(y \rightarrow x) \in F$.

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