

Herbert Fleischner; Emanuel Wenger  
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## $D_0$ -FAVOURING EULERIAN TRAILS IN DIGRAPHS

HERBERT FLEISCHNER, EMANUEL WENGER

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*Dedicated to the memory of Milan Sekanina*

**Abstract.** A characterization for a special class of Eulerian trails in digraphs which traverse a set of arcs of a subdigraph  $D_0$  before any arc of  $D_1 = D - D_0$  is traversed, is proved. The most general structure of a subdigraph  $D_1$  to allow such a restricted Eulerian trail is given.

**Key words.** Directed graph, Eulerian trail, restricted Eulerian trail, spanning in-tree.

**MS Classification.** 05 C 139

### PRELIMINARIES

For notation and terminology, see [2, 4]. Let  $D$  be a digraph with vertex set  $V(D)$  and  $A(D)$ . In particular,  $V(D)$  and  $A(D)$  are always assumed to be finite.  $A_v^+ \subset A(D)$  denotes the set of arcs, incident from  $v$ , for  $v \in V(D)$ . For a digraph  $D$  and a subdigraph  $D_1$  let  $D - D_1 \subseteq D - A(D_1)$  denote the uniquely determined digraph without isolated vertices. The following lemma is folklore.

**Lemma 1.** *Let  $D$  be a digraph and  $\text{od}_D(v) \geq 1$  for all  $v \in V(D)$ . Then there exists at least one non-trivial strongly connected component  $C$  with no arc of  $D$  incident from  $C$  (that is,  $(a, b) \in A(D)$  implies either  $a \notin V(C)$ , or  $b \notin V(D) - V(C)$ ).*

**Lemma 2.** *Let  $D$  be a digraph satisfying  $\text{od}_D(v) \geq 1$  for all  $v \in V(D)$ . Suppose  $D$  has precisely one (nontrivial) strongly connected component  $C$  with no arc of  $D$  incident from  $C$ . Then there exists a spanning in-tree with root  $v_0$ , where  $v_0$  is an arbitrary vertex of  $C$ .*

**Proof.** Let  $v_0$  be an arbitrary vertex of  $C$ , and let  $B_0$  be an in-tree with root  $v_0$  containing a maximum number of vertices. If  $V(B_0) \neq V(D)$  then we consider  $D_0 = \langle V(B_0) \rangle$ , the digraph induced by  $V(B_0)$ . Because of the maximality of  $B_0$  there does not exist an arc  $(x, y)$  with  $x \in V(D) - V(D_0)$  and  $y \in V(D_0)$ ; furthermore, one easily concludes that  $C \subseteq D_0$ .  $D_1 = D - V(D_0)$  fulfills the assumptions of Lemma 1. Because of Lemma 1 there exists a strongly connected component  $C' \subset D_1$  such that no arc of  $D_1$  is incident from  $C'$ . By construction it follows that  $C' \cap C = \emptyset$  which contradicts the uniqueness of  $C$ .

**Definition.** Let  $D$  be a weakly connected eulerian digraph, and let  $D_0$  be a subdigraph of  $D$ . An eulerian trail  $T$  of  $D$  is called  $D_0$ -favouring if and only if for every  $v \in V(D)$ ,  $T$  traverses every arc of  $D_0$  incident from  $v$  before it traverses any arc of  $D_1 = D - D_0$  incident from  $v$ .

Of course, every eulerian trail of  $D$  is a  $D_0$ -favouring eulerian trail for some  $D_0$  (just take  $D_0 = D$ ). For which subdigraph  $D_0$  of  $D$  does there exist a  $D_0$ -favouring eulerian trail? There are two known results on the existence of  $D_0$ -favouring eulerian trails depending on the structure of  $D_1 = D - D_0$ .

**Theorem 1.** Let  $D$  be a weakly connected eulerian digraph, and for given  $v \in V(D)$  let  $D_0 \subset D$  be chosen such that  $D_1 = D - D_0$  is a spanning in-tree of  $D$  with root  $v$ . Then there exists a  $D_0$ -favouring eulerian trail starting (and ending) at  $v$ . Conversely, if  $T$  is an eulerian trail of  $D$  starting (and ending) at  $v$ , and if we mark at every  $w \in V(D)$ ,  $w \neq v$ , the last arc of  $T$  incident from  $w$ , then  $D_1$ , the subgraph of  $D$  induced by the marked arcs, is a spanning in-tree with root  $v$  (and hence  $T$  is a  $(D - D_1)$ -favouring eulerian trail of  $D$ ).

Theorem 1 plays an essential role in establishing the BEST-Theorem which gives a formula for the number of eulerian trails in an eulerian digraph. A proof of Theorem 1 can be found in [1].

**Theorem 2.** Let  $D$  be an eulerian digraph. Let  $D_1 \subseteq D$  be chosen such that  $\text{od}_{D_1}(v) \geq 1$  for every  $v \in V(D_1) \subset V(D)$ , and let  $D_0 = D - D_1$ .  $D$  has a  $D_0$ -favouring eulerian trail if and only if  $D_1$  has precisely one (nontrivial) strongly connected component  $C_1$  with the property that no arc of  $D_1$  is incident from  $C_1$ . Moreover, every  $D_0$ -favouring eulerian trail of  $D$  must start at some vertex of  $C_1$ , and for any vertex of  $C_1$  there is a  $D_0$ -favouring eulerian trail of  $D$  starting at that vertex.

Theorem 2 was proved by Berkowitz [3].

## A GENERAL THEOREM

In view of Theorems 1 and 2, we ask the following question: What is the most general structure a subdigraph  $D_1$  of an eulerian digraph  $D$  can have in order to imply the existence of a  $(D - D_1)$ -favouring eulerian trail  $T$ ?

Theorem 2 implies that  $D_1$  must not contain more than one nontrivial strongly connected component  $C_1$  with the property that no arc of  $D_1$  is incident from  $C_1$ . But this condition is not sufficient even if  $D_1$  is weakly connected; this can be seen from the digraph  $D^*$  of Figure 1.

What if we go the other way round? That is, given an eulerian digraph  $D$  and  $D_1 \subseteq D$ , can we find  $D_1^+ \subseteq D$  with  $D_1 \subseteq D_1^+$  such that  $D$  has a  $(D - D_1^+)$ -favouring eulerian trail  $T^+$  which induces a  $(D - D_1)$ -favouring eulerian trail  $T$ ?

This approach and Theorem 1 and Theorem 2 lead to the following theorem which answers our original question.

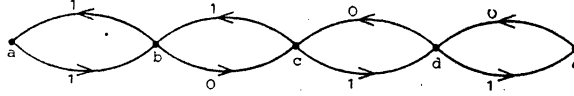


Figure 1

Figure 1. An eulerian digraph  $D^*$  having no  $D_0$ -favouring eulerian trail (the arcs of  $D_i$  are marked with  $i, i = 0, 1$ ).

**Theorem 3.** Let  $D$  be an eulerian digraph, and let  $D_1$  be a subdigraph of  $D$ . Any two of the following statements are equivalent:

1.  $D$  has a  $(D - D_1)$ -favouring eulerian trail.
2. There exists a digraph  $D_1^+$  with  $D_1 \subseteq D_1^+ \subseteq D$  such that for every  $v \in V(D)$ 
  - a)  $\text{od}_{D_1^+}(v) = \text{od}_{D_1}(v)$  if and only if  $\text{od}_{D_1}(v) \neq 0$ ;
  - b)  $\text{od}_{D_1^+}(v) = 1$  otherwise.
  - c)  $D_1^+$  has precisely one non-trivial strongly connected component  $C_1$  with no arc of  $D_1^+$  incident from  $C_1$ .
3. There exists a digraph  $D_1^+$  with  $D_1 \subseteq D_1^+ \subseteq D$  such that
  - a)  $D$  has a  $(D - D_1^+)$ -favouring eulerian trail;
  - b) for every  $D_1'$  with  $D_1 \subseteq D_1' \subseteq D_1^+$ , if  $(x, y) \in A(D_1' - D_1)$ , then  $\text{od}_{D_1}(x) = 0$ .
4.  $D_1$  contains a spanning in-forest  $D_1^-$  such that
  - a) for some  $v_0$  and for every  $x \in V(D_1) - v_0$ ,  $\text{od}_{D_1^-}(x) = 0$  if and only if  $\text{od}_{D_1}(x) = 0$ , and  $\text{od}_{D_1^-}(v_0) = 0$ ;
  - b)  $D$  has an in-tree  $B$  with root  $v_0$  and  $D_1^- \subseteq B$ .

**Proof.** 1. implies 2. Let  $T$  be a  $(D - D_1)$ -favouring eulerian trail starting at  $v_0$ . Define  $D_1^+$  by  $D_1^+ = D_1$  if  $\text{od}_{D_1}(v) \geq 1$  for every  $v \in V(D)$ ; otherwise, for every  $v$  with  $\text{od}_{D_1}(v) = 0$ , mark the last arc of  $T$  which is incident from  $v$ , and let  $D_1^+$  consist of  $D_1$  plus the marked arcs. In any case,  $D_1 \subseteq D_1^+$  and  $D_1^+$  satisfies 2. a), 2. b). Moreover,  $T$  is a  $(D - D_1^+)$ -favouring eulerian trail because of the choice of the elements of  $A(D_1^+) - A(D_1)$ . It remains to show that  $D_1^+$  has precisely one non-trivial strongly connected component  $C_1$  with no arc of  $D_1^+$  incident from  $C_1$ . Because of  $\text{od}_{D_1^+}(v) \geq 1$  for every  $v \in V(D_1^+)$  and the finiteness of  $D_1^+$ ,  $D_1^+$  has at least one non-trivial strongly connected component and, in particular, by Lemma 1 at least one non-trivial strongly connected component  $C_1^+$  with no arc of  $D_1^+$  incident from  $C_1^+$ .

$T$  must start and end in a vertex of  $C_1^+$ . Otherwise, there exist one or more arcs  $(v, w)$  of  $D$  such that  $v \in V(C_1^+)$  and  $w \notin V(C_1^+)$ ; among these arcs let  $(v_1, w_1)$  be

the last arc in  $T$ , such that  $v_1, (v_1, w_1), w_1$  is a section of  $T$ . By definition of  $C_1^+$ ,  $(v_1, w_1) \notin A(D_1^+)$ , and because of  $\text{od}_{D_1^+}(v_1) \geq 1$  we get a contradiction to the fact that  $T$  is a  $(D - D_1^+)$ -favouring eulerian trail. It's clear now that there can be only one component  $C_1^+$  with the desired property. The implication now follows.

2. *implies* 3. Take  $D_1^+$  and  $C_1$  as defined by 2 a), b), and c). At first it will be proved that  $D$  has a  $(D - D_1^+)$ -favouring eulerian trail.

Properties 2. a), b), imply that  $D_1^+$  is a spanning subdigraph of  $D$ . Therefore and because of Lemma 1, and property 2. c) there exists in  $D$  a spanning in-tree  $B_1^+ \subset D_1$  with root  $v_0 \in V(C_1)$  (see Lemma 2).

Mark all the arcs of  $B^+$ . Construct  $T$  by starting at vertex  $v_0$  with any arc  $(v_0, x)$ , choose any unmarked arc incident from  $x$ , if such arc exists; otherwise, choose among the marked arcs one which does not belong to  $B_1^+$  if such arc exists; otherwise, choose the arc of  $B_1^+$ . Continue this way until this procedure terminates at some  $y \in V(D)$ . Then  $y = v_0$ ; otherwise,  $T$  contains more arcs incident to  $y$  than it contains arcs incident from  $y$  contradicting  $D$  being eulerian. Suppose  $T$  does not contain all arcs of  $D$ . Then let  $z$  be a vertex incident with arcs not contained in  $T$ . Since  $D$  is eulerian and  $T$  is a closed trail,  $\text{id}_{D-T}(z) = \text{od}_{D-T}(z) \neq 0$ . Moreover,  $z \neq v_0$  by the very construction of  $T$ . By definition of  $B_1^+$ , there is a path  $P(z, v_0) \subset B_1^+$  joining  $z$  to  $v_0$ . Write

$$P(z, v_0) = z, (z, u_1), u_1, \dots, u_k, (u_k, v_0), v_0;$$

possibly  $z = u_k$  and  $u_1 = v_0$  (i.e.  $P(z, v_0)$  may contain just one arc). By the construction of  $T$  it follows that  $(z, u_1)$  is not contained in  $T$ ; therefore, also  $(u_1, u_2)$  is not contained in  $T$  (note that  $(u_1, u_2)$  can be contained in  $T$  only if all arcs incident to  $u_1$  are contained in  $T$ ); a.s.o. In particular,  $(u_k, v_0)$  is not contained in  $T$ , contradicting the fact that  $\text{id}_T(v_0) = \text{od}_T(v_0) = \text{id}_D(v_0) = \text{od}_D(v_0)$ . Thus,  $T$  contains all arcs of  $D$ . This and the construction of  $T$  imply that  $T$  is a  $(D - D_1^+)$ -favouring eulerian trail of  $D$ .

Now consider any  $D_1'$  with  $D_1 \subseteq D_1' \subseteq D_1^+$  and suppose  $A(D_1' - D_1) \neq \emptyset$ ; let  $(x, y) \in A(D_1' - D_1)$ . By definition of  $D_1^+$  in 2. a), b), an arc of  $D_1^+ - D_1$  is necessarily incident from a vertex  $z$  with  $\text{od}_{D_1}(z) = 0$ . Hence  $(x, y) \in A(D_1' - D_1)$  implies  $\text{od}_{D_1}(x) = 0$ ; thus 3. b) holds as well.

3. *implies* 4. Start with  $D_1^+$  as described in 3., and consider a  $(D - D_1^+)$ -favouring eulerian trail  $T^+$  of  $D$ . If there is  $w \in V(D)$  different from the initial vertex  $v_0$  of  $T^+$  such that the last arc of  $T^+$  incident from  $w$  is not in  $D_1^+$ , then mark this arc. Note that in this case none of the arcs incident from  $w$  lies in  $D_1^+$ .

We define

$$D_1^{++} = D_1^+ \quad \text{if no such } w \text{ exists;}$$

otherwise,

$$D_1^{++} = \langle A(D_1^+) \cup \{a \in A_w^+ / \text{od}_{D_1^+}(w) = 0 \text{ and } a \text{ has been marked}\} \rangle.$$

In any case, by definition of  $D_1^{++}$ ,  $T^+$  is even a  $(D - D_1^{++})$ -favouring eulerian trail of  $D$ , and  $D_1^{++}$  satisfies 3. b) as well. Moreover,  $V(D_1^{++}) = V(D)$ .

Marking for every  $v \neq v_0$  the last arc of  $T^+$  incident from  $v$  yields a spanning subdigraph  $B$  and  $B \subset D_1^{++}$  follows from the very definition of  $D_1^{++}$ . Furthermore,  $\text{od}_B(v) = 1$  for all  $v \neq v_0$  and  $\text{od}_B(v_0) = 0$ . Suppose  $B$  is not connected; then there exists a weakly connected component  $B_1$  of  $B$  which does not contain  $v_0$  and  $\text{od}_{B_1}(w) = \text{od}_B(w) = 1$  for all  $w \in V(B_1)$ . By Lemma 1 there exists at least one nontrivial strongly connected component  $C_1 \subseteq B_1$  with no arc of  $B_1$  incident from  $C_1$ . Now, if  $r$  is the last vertex of  $T$  in  $C_1$ , such that  $r, (r, s), s$  is a section of  $T$ , then it follows from the construction of  $B$  that  $(r, s) \in A(B)$ ; furthermore  $s \in V(C_1)$  because of the definition of  $C_1$ . By the choice of  $r$ ,  $T$  terminates in  $C_1$  contradicting the fact, that  $T$  is an eulerian trail starting in  $v_0 \notin V(B_1) \supset V(C_1)$ . Thus  $B$  is connected, and  $\text{od}_B(v) = 1$  for all  $v \neq v_0$ ,  $\text{od}(v_0) = 0$ . This implies that  $B$  is a spanning in-tree of  $D_1^{++} \subseteq D$  rooted at  $v_0$ .

Define  $D_1^-$  by  $V(D_1^-) = V(D_1)$  and  $A(D_1^-) = A(B) \cap A(D_1)$ ; thus  $D_1^-$  is a spanning in-forest of  $D_1$  which satisfies 4. b). Let  $(x, y)$  be any arc of  $B$  not in  $D_1^-$ ; then  $x \neq v_0$ . If  $(x, y) \notin A(D_1^+)$ , then it follows from the definition of  $D_1^{++}$  and  $D_1^{++} \supset D_1$  that  $\text{od}_{D_1^+}(x) = 0 = \text{od}_{D_1}(x)$ . If  $(x, y) \in A(D_1^+)$ , then  $(x, y) \notin A(D_1)$  by definition of  $D_1^-$ ; and by 3. b) with  $D_1' = D_1^+$ ,  $\text{od}_{D_1}(x) = 0$  follows.

We summarize:  $D_1^-$  is a spanning in-forest of  $D_1$ , and if  $x \neq v_0$  for some  $v_0 \in V(D_1^-)$  (which is the root of  $B$  indeed) satisfies  $\text{od}_{D_1^-}(x) = 0$  then  $\text{od}_{D_1}(x) = 0$  (for,  $x$  not being the root of  $B$  implies  $(x, y) \in A(B - D_1^-)$  for some  $y$ ). Since  $\text{od}_{D_1}(x) = 0$  implies  $\text{od}_{D_1^-}(x) = 0$  anyway and  $\text{od}_{D_1^-}(v_0) = \text{od}_B(v_0) = 0$ , and because  $D_1^- \subseteq B$  with  $V(B) = V(D)$ , the proof of the implication is finished.

4. implies 1. Let  $D_1^- \subseteq D_1$  be chosen as described in 4. a) and let  $B$  be a spanning in-forest of  $D$  with root  $v_0$  and  $D_1^- \subseteq B$ . Marking all arcs of  $B$  we construct a trail  $T$  by starting at vertex  $v_0$  with any arc  $(v_0, x)$ . Choose any unmarked arc incident from  $x$ , if such arc exists; choose the marked arc incident from  $x$ , otherwise.

Continuing this way until this procedure terminates we get a  $(D - B)$ -favouring eulerian trail (for arguments see 2. implies 3.).

Because of the freedom to choose the order in which the arcs of  $A_v^+ - A(B)$  appear in  $T$  for every  $v \in V(D)$  we are even able to construct  $T$  in such a way that the arcs of  $A_v^+ \cap (D - D_1)$  appear in  $T$  before any of the arcs of  $A_v^+ \cap D_1$  are used. This is true even in the case where an arc  $(x, y) \in B$  does not belong to  $D_1$ ; for, in this case  $\text{od}_{D_1^-}(x) = \text{od}_{D_1}(x) = 0$  by 4. a), i.e.  $A_x^+ \cap A(D_1) = \emptyset$ , i.e.,  $A_x^+ \subseteq D - D_1$ . In the case of  $v_0$ , if  $A_{v_0}^+ \cap A(D_1) \neq \emptyset$ , then we proceed in the construction of  $T$  by starting along an arc of  $A_{v_0}^+ \cap A(D - D_1)$ , and each time we arrive in  $v_0$  we continue along an arc of  $A_{v_0}^+ \cap A(D - D_1)$  not traversed before, as long as there is such an arc. Consequently,  $T$  is a  $(D - D_1)$ -favouring eulerian trail of  $D$ . This finishes the proof of the implication. Theorem 3 now follows.

It is easy to see that Theorem 3 is a generalization of Theorem 1 and Theorem 2. Both Theorems can be derived by using the equivalent statements of Theorem 3 and some details of their proof. We also note that in proving Theorem 3 we used ideas developed originally for the proofs of Theorem 1 and Theorem 2.

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*Prof. Herbert Fleischner, Emanuel Wenger*  
*Austrian Academy of Sciences*  
*Institute for Information Processing*  
*Sonnenfelsgasse 19/2*  
*1010-Vienna, Austria*