A formal model for explicit knowledge as awareness of plus awareness that

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Explicit knowledge as awareness of + awareness that

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Explicit knowledge as awareness of + awareness that

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- System of Explicit Knowledge
- The model
- The concepts
- **PROPERTIES AND RELATIONSHIPS**
- Awareness-of and Awareness-that
- Effects of the closure operation
- Moorean Phenomena
- Other Alternatives for the Concepts

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IN A NUTSHELL

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- *Purpose:* reconsider what constitutes *explicit knowledge*.



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- *Purpose:* reconsider what constitutes *explicit knowledge*.
- *Here:* a formal model capturing the theoretical ideas.

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 - *explicit* and *implicit* w.r.t. *deduction* (e.g., Konolige 1984, Levesque 1984);
 - explicit and implicit w.r.t. awareness (Fagin and Halpern 1988);
- Note.
 - *Explicit knowledge:* what the agent *actually has*.
 - 'Implicit' knowledge: what she can reach via some given action.

• Awareness-of and awareness-that.

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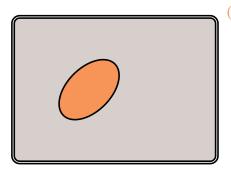
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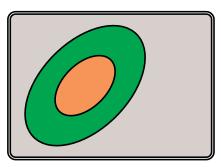
- Fagin and Halpern (1988): awareness has *different interpretations*.
- Dretske (1993): Awareness of things vs awareness of facts.
- Here:
 - *Awareness-of* as *entertaining* (*'working memory'*), not implying any attitude in favour or against.
 - Awareness-that as acknowledgement or acceptance.

Combined Diagram



(5) Awareness-that

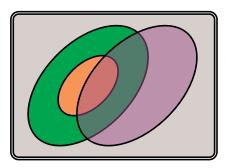
COMBINED DIAGRAM



(5) Awareness-that

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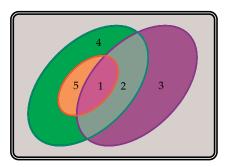


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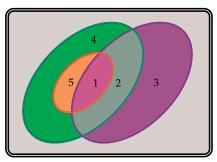
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(5) Awareness-that (not in working memory)
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(2) Aware-of not aware-that, but deducible
(1) Explicit knowledge
(aware-of and aware-that)

Explicit knowledge as awareness of + awareness that

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Awareness neighbourhood model (ANM)

Definition (Awareness neighbourhood model (ANM))

Let **P** be a set of atoms. An *ANM* is a tuple $M = \langle W, N, V, A \rangle$ where

- $W \neq \emptyset$ • $V: \mathbb{P} \to \wp(W)$
- $N: W \to \wp(\wp(W))$ $A \subseteq \mathbf{P}$

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• *Awareness-of*: (global) set of atoms *A*.

Language and semantic interpretation (1)

Definition (Language \mathcal{L})

 $\varphi, \psi ::= \top | p | \neg \varphi | \varphi \land \psi | A^{\circ} \varphi | A^{t} \varphi | [*] \varphi$

- $\llbracket \mathsf{T} \rrbracket^M := W$,
- $\llbracket p \rrbracket^M := V(p),$

- $\llbracket \neg \varphi \rrbracket^M := W \setminus \llbracket \varphi \rrbracket^M$,
- $\llbracket \boldsymbol{\varphi} \wedge \boldsymbol{\psi} \rrbracket^M := \llbracket \boldsymbol{\varphi} \rrbracket^M \cap \llbracket \boldsymbol{\psi} \rrbracket^M.$

Language and semantic interpretation (2)

Definition (Language \mathcal{L})

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,
• $\llbracket \mathbf{A}^{\mathbf{t}} \boldsymbol{\varphi} \rrbracket^{M} := \{ w \in W \mid \llbracket \boldsymbol{\varphi} \rrbracket^{M} \in N(w) \}.$

LANGUAGE AND SEMANTIC INTERPRETATION (3)

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Given $M = \langle W, N, V, A \rangle$, define $M^* = \langle W, N^*, V, A \rangle$ with

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The concepts of *satisfiability* and *validity* are defined as usual.

Explicit knowledge as awareness of + awareness that

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Aware knowledge

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PROPERTIES OF AWARENESS-OF (A^o)

• The agent is aware-of the **concept of truth**: ⊩ A^o ⊤

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- Since A^o is defined as a set of atomic propositions, it is **closed under subformulas and superformulas**:
 - $$\begin{split} \Vdash A^{\circ} \neg \varphi \leftrightarrow A^{\circ} \varphi & \qquad \qquad \Vdash A^{\circ} A^{\circ} \varphi \leftrightarrow A^{\circ} \varphi \\ \Vdash A^{\circ} (\varphi \land \psi) \leftrightarrow (A^{\circ} \varphi \land A^{\circ} \psi) & \qquad \qquad \qquad \Vdash A^{\circ} A^{t} \varphi \leftrightarrow A^{\circ} \varphi \\ \Vdash A^{\circ} [*] \varphi \leftrightarrow A^{\circ} \varphi \end{split}$$

Properties of Awareness-that (A^t)

A^t is what appears in N(w). This pureley semantic concept is closed under logical equivalence (some kind of omniscience):
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- But it is the only closure property, since **I** φ does not imply **I** $A^t \varphi$ **J** $(A^t \varphi \land A^t \psi) \rightarrow A^t(\varphi \land \psi)$ **J** $A^t(\varphi \land \psi) \rightarrow A^t \varphi$ and **J** $A^t(\varphi \land \psi) \rightarrow A^t \psi$
- Hence, A^t is not closed under logical consequence: $\# A^t(\varphi \to \psi) \to (A^t \varphi \to A^t \psi)$

Awareness-of and Awareness-that

In contrast to what happens in *Awareness Logic* by Fagin and Halpern, where $\Vdash A\varphi \rightarrow \Box A\varphi$, with a global awareness set, we do not obtain this result, thanks to the different concepts of awareness we defined.

Recall that **awareness-of** is a global notion and **awareness-that** is locally defined.

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Thus, analogous properties do not hold:

• J $A^{o} \varphi \rightarrow A^{t} A^{o} \varphi$ • J $A^{o} \varphi \rightarrow A^{t} \neg A^{o} \varphi$

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- About logical equivalence: $\mu \phi \leftrightarrow \psi$ does not imply $\mu K_{Ex} \phi \leftrightarrow K_{Ex} \psi$

But, $\Vdash \varphi \leftrightarrow \psi$ implies $\Vdash (K_{Ex} \varphi \land A^{\circ} \psi) \rightarrow K_{Ex} \psi$

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- About logical equivalence:

$$\label{eq:constraint} \begin{split} & \Vdash \varphi \leftrightarrow \psi \text{ does not imply } \Vdash \mathsf{K}_{Ex} \, \varphi \leftrightarrow \mathsf{K}_{Ex} \, \psi \\ & \text{But, } \Vdash \varphi \leftrightarrow \psi \text{ implies } \Vdash (\mathsf{K}_{Ex} \, \varphi \wedge \mathsf{A}^{\mathsf{o}} \, \psi) \rightarrow \mathsf{K}_{Ex} \, \psi \end{split}$$

• About Modus Ponens:

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Recall that $K_{Im} \varphi := A^{\circ} \varphi \wedge [*] A^{t} \varphi$. This has the following consequences:

PROPERTIES OF IMPLICIT KNOWLEDGE (K_{Im})

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Explicit knowledge as awareness of + awareness that

- $\mathbb{H} A^t \varphi \rightarrow [*] A^t \varphi$.

- 'Implicit is not always Explicit': $\mathbb{K} \mathbb{K}_{Ex} \varphi \to \mathbb{K}_{Im} \varphi$ What the agent has acknowledged as true does not need to hold after the closure operation. Thus,
- J* $A^t \varphi \rightarrow [*] A^t \varphi$. Take $\varphi := \neg A^t q$, then $A^t \neg A^t q$ has a similar effect as a *Moore sentence*, stating "the agent is aware that it is the case that she is not aware that q is the case".

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- While 'A^t ¬ A^t q' is true at *M*, it will not be the case at *M**, since its truthset has shrunk after the operation.

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- While 'A^t ¬ A^t q' is true at *M*, it will not be the case at *M**, since its truthset has shrunk after the operation.
- Though, $\Vdash \varphi \rightarrow [*] \varphi$ implies $\Vdash K_{Ex} \varphi \rightarrow K_{Im} \varphi$

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OTHER ALTERNATIVES FOR REPRESENTING OUR BASIC CONCEPTS

For representing Awareness-of:

- Concept of *topics* in Berto and Hawke (2018) (cf. Berto 2018). (A *topic* being what the sentence is *about*.)
- The *issue relation* in, e.g., Grossi (2009), van Benthem and Minică (2012), Baltag et al. (2018). (Equivalence relation that creates partitions of the domain in relational model.)

For representing Awareness-that:

- *Explicit knowledge* in proposals not incorporating the notion of *awareness*, e.g., Konolige 1984, Levesque 1984, Duc 1997, Artemov and Nogina 2005, Jago 2009, Velázquez-Quesada 2013.
- Alternatives where the knowledge/belief relies on *evidences* (van Benthem and Pacuit 2011, Özgün 2017) and *arguments* (Shi et al. 2018a, 2017, 2018b).

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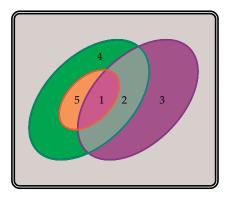
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Recall: Combined Diagram



(5) Awareness-that (not in working memory)
(4) 'Implicit' Awareness-that (not in working memory)
(3) Awareness-of
(2) Aware-of not aware-that, but deducible
(1) Explicit knowledge
(aware-of and aware-that)

Action: Becoming Aware-of $[+\chi]$ (1)

Definition (The becoming aware-of operation)

Given $M = \langle W, N, V, A \rangle$ and $M^{+\chi} = \langle W, N, V, A^{+\chi} \rangle$, we have

 $A^{+\chi} = A \cup \operatorname{atm}(\chi)$

Action: Becoming Aware-of $[+\chi]$ (1)

Definition (The becoming aware-of operation)

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Then, we define $[[+\chi] \varphi]^M = [\varphi]^{M^{+\chi}}$ and extend the language \mathcal{L} with $[+\chi] \varphi$, read as after the agent becomes aware-of χ , φ is the case.

Action: Becoming Unaware-of $[-\chi]$

Definition (The becoming unaware-of operation)

Given $M = \langle W, N, V, A \rangle$ and $M^{-\chi} = \langle W, N, V, A^{-\chi} \rangle$, we have

 $A^{-\chi} = A \setminus \operatorname{atm}(\chi)$

Action: Becoming Unaware-of $[-\chi]$

Definition (The becoming unaware-of operation)

Given $M = \langle W, N, V, A \rangle$ and $M^{-\chi} = \langle W, N, V, A^{-\chi} \rangle$, we have

 $A^{-\chi} = A \setminus \operatorname{atm}(\chi)$

Then, we define $[\![-\chi] \varphi]^M = [\![\varphi]\!]^{M^{-\chi}}$ and extend the language \mathcal{L} with $[-\chi] \varphi$, read as after the agent becomes unaware-of χ , φ will be the case.

Action: Becoming Unaware-of $[-\chi]$

Definition (The becoming unaware-of operation)

Given $M = \langle W, N, V, A \rangle$ and $M^{-\chi} = \langle W, N, V, A^{-\chi} \rangle$, we have

 $A^{-\chi} = A \setminus \operatorname{atm}(\chi)$

Then, we define $\llbracket [-\chi] \varphi \rrbracket^M = \llbracket \varphi \rrbracket^{M^{-\chi}}$ and extend the language \mathcal{L} with $[-\chi] \varphi$, read as after the agent becomes unaware-of χ , φ will be the case.

Alternative definition: weak becoming unawarene-of

- $\llbracket [-'Q] \varphi \rrbracket^M = \llbracket \varphi \rrbracket^{M^{-Q}}$
- $\llbracket [-'\chi] \varphi \rrbracket^M = \llbracket \bigwedge_{\{Q \subseteq \operatorname{atm}(\chi) \mid Q \neq \emptyset\}} [-Q] \varphi \rrbracket^M$
- $[\langle -\chi \rangle \varphi]^M = [[\bigvee_{\{Q \subseteq \operatorname{atm}(\chi) | Q \neq \emptyset\}} [-Q] \varphi]^M$

Definition (The deductive inference operation)

For $\eta, \chi, \varphi \in \mathcal{L}$, $\operatorname{atm}(\eta \to \chi) \subseteq A$, and $M = \langle W, N, V, A \rangle$, we have $M^{\eta \to \chi} = \langle W, N^{\eta \to \chi}, V, A \rangle$ where for any $w \in W$: $N^{\eta \to \chi}(w) = \begin{cases} N(w) \cup [\chi]^M & \text{if } \{[(\eta \to \chi)]^M, [\eta]^M\} \subseteq N(w) \\ N(w) & \text{otherwise} \end{cases}$

 $\eta \rightarrow \chi$

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Action: Deductive Inference (Modus Ponens step) $[\eta \rightarrow \chi]$

Definition (The deductive inference operation)

For $\eta, \chi, \varphi \in \mathcal{L}$, $\operatorname{atm}(\eta \to \chi) \subseteq A$, and $M = \langle W, N, V, A \rangle$, we have $M^{\eta \to \chi} = \langle W, N^{\eta \to \chi}, V, A \rangle$ where for any $w \in W$: $N^{\eta \to \chi}(w) = \begin{cases} N(w) \cup \llbracket \chi \rrbracket^M & if \{\llbracket (\eta \to \chi) \rrbracket^M, \llbracket \eta \rrbracket^M \} \subseteq N(w) \\ N(w) & otherwise \end{cases}$ Then, we define $\llbracket [\eta \to \chi] \varphi \rrbracket^M = \llbracket \varphi \rrbracket^{M^{\eta \to \chi}}$ and extend the language \mathcal{L} with $\llbracket \eta \to \chi \rrbracket \varphi$, read as after the agent performs a deductive inference from $\eta \to \chi$ and η holds, φ is the case.

Action: Forgetting $[\chi]$

DEFINITION (THE FORGETTING OPERATION)

For $\chi \in \mathcal{L}$ such that $\operatorname{atm}(\chi) \subseteq A$, we have $M = \langle W, N, V, A \rangle$ and $M^{\setminus \chi} = \langle W, N^{\setminus \chi}, V, A \rangle$ where for $w \in W$:

 $N^{\setminus \chi}(w) = N(w) \setminus \llbracket \chi \rrbracket^M$

Action: Forgetting $[\chi]$

Definition (The forgetting operation)

For $\chi \in \mathcal{L}$ such that $\operatorname{atm}(\chi) \subseteq A$, we have $M = \langle W, N, V, A \rangle$ and $M^{\setminus \chi} = \langle W, N^{\setminus \chi}, V, A \rangle$ where for $w \in W$:

 $N^{\setminus \chi}(w) = N(w) \setminus \llbracket \chi \rrbracket^M$

Then, we define $\llbracket [\chi] \varphi \rrbracket^M = \llbracket \varphi \rrbracket^{M^{\chi}}$ and extend the language \mathcal{L} with $[\chi] \varphi$, read as after the agent forgets χ, φ is the case.

Explicit knowledge as awareness of + awareness that

INTRODUCTION

- System of Explicit Knowledge • The model
- The model
- The concepts
- **PROPERTIES AND RELATIONSHIPS**
- Awareness-of and Awareness-that
- Effects of the closure operation
- Moorean Phenomena
- Other Alternatives for the Concepts

EPISTEMIC ACTIONS



SUMMARIZING

- *Awareness-of* and *awareness-that* as primitive concepts defining *explicit knowledge*.
- A *semantic model*; defined the involved notions.
- *Properties* as compared with related approaches (e.g., Hintikka 1962, Konolige 1984, Fagin and Halpern 1988).

CURRENT AND FUTURE WORK

- More precise comparison with other *semantic alternatives*
- Axiom system.
- Further *epistemic actions* like *observation or communication*

System of Explicit Knowledge	PROPERTIES AND RELATIONSHIPS		

Thank you! ¡Muchas gracias!

Explicit knowledge as awareness of + awareness that

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