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PATHS AND CIRCUITS IN \mathbb{G} -GRAPHS OF CERTAIN NON-ABELIAN GROUPS

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ABSTRACT. In [BJRTD08], necessary and sufficient conditions were given for the existence of Eulerian and Hamiltonian paths and circuits in the \mathbb{G} -graph of the dihedral group D_n . In this paper, we consider the \mathbb{G} -graphs of the quasihedral, modular, and generalized quaternion group. These groups are of rank 2 and we consider only the graphs $\Gamma(G, S)$ where $|S| = 2$.

1. INTRODUCTION

Let G be a finitely generated group with generating set $S = \{s_1, \dots, s_k\}$. For a subgroup H of G , define the subset T_H of G to be a *left transversal* for H if $\{xH \mid x \in T_H\}$ is precisely the set of all left cosets of H in G . For each $s_i \in S$ let $H_i = \langle s_i \rangle$. Associate a simple graph $\Gamma(G, S)$ to (G, S) with vertex set $V = \{x_j H_i \mid x_j \in T_{H_i}\}$. Two distinct vertices $x_j H_i$ and $x_l H_k$ in V are joined by an edge if $x_j \langle s_i \rangle \cap x_l \langle s_k \rangle$ is nonempty. The edge set E consists of pairs $(x_j H_i, x_l H_k)$. $\Gamma(G, S)$ defined this way has no multiedge or loop. A multiedge graph was defined similarly in 2004. Many of the results about this graph [[BG04], [BGL05], [BG05], and [BG07]] can be modified for the simple graph, $\Gamma(G, S)$, [D08]. The main object of this paper is to study the existence of Eulerian and Hamiltonian paths and circuits in the \mathbb{G} -graphs of the quasihedral, modular, and generalized quaternion group. To explore the existence of Eulerian paths and circuits in $\Gamma(G, S)$, we recall a few theorems of Euler and a result from [BJRTD08].

Theorem 1. (Euler) *Let Γ be a nontrivial connected graph. Then Γ has an Eulerian circuit if and only if every vertex is of even degree.*

Theorem 2. (Euler) *Let Γ be a nontrivial connected graph. Then Γ has an Eulerian path if and only if Γ has exactly two vertices of odd degree. Furthermore, the path begins at one of the vertices of odd degree and terminates at the other.*

Lemma 3. [BJRTD08] *If G is a group with generating set $S = \{s_1, \dots, s_n\}$ and $S_{i,j} = |\langle s_i \rangle \cap \langle s_j \rangle|$, then the degree of the vertex $\langle s_i \rangle$, denoted $\deg(\langle s_i \rangle)$, is*

$$\deg(\langle s_i \rangle) = \left(\sum_{j=1}^n |s_i|/S_{i,j} \right) - 1.$$

Remark 1. Notice that $\deg(\langle s_i \rangle) = \deg(x_j \langle s_i \rangle)$ for all $x_j \langle s_i \rangle$ in V_i .

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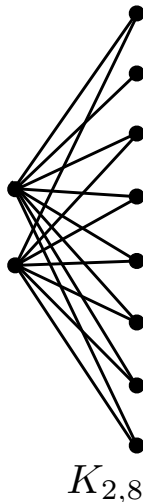
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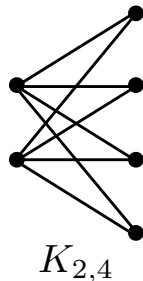
We consider the \mathbb{G} -graphs of the quasihedral, modular, and generalized quaternion group. We start with a few examples of the graphs.

Example 1.

- (i) The modular group, M , has presentation $\langle s, t \mid s^8 = t^2 = e, st = ts^5 \rangle$. Letting $S = \{s, t\}$, the \mathbb{G} graph of this group is $\Gamma(M, S)$.



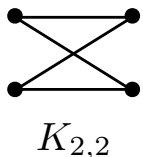
- (ii) The quasihedral group, QS , has presentation $\langle s, t \mid s^8 = t^2 = e, st = ts^3 \rangle$. Letting $S = \{s, ts\}$, the \mathbb{G} graph of this group is $\Gamma(QS, S)$.



- (iii) The generalized quaternion group, Q_{2^n} , has presentation

$$\langle s, t \mid s^{2^{n-1}} = e, s^{2^{n-2}} = t^2, tst^{-1} = s^{-1} \rangle.$$

Letting $n = 3$, $S = \{s, t\}$, the \mathbb{G} graph of this group is $\Gamma(Q_{2^3}, S)$.



The next lemma pertains to all of the groups in question.

Lemma 4. *Let $G = M, QS$, or Q_{2^n} and let j be an odd integer then*

$$\langle s^j \rangle = \langle s \rangle = \{s, s^2, \dots, s^{|s|-1}, e\}.$$

Proof. For each of the above groups, $|s|$ is even. So $\gcd(j, |s|) = 1$ and there exist $x, y \in \mathbb{Z}$ such that $jx + |s|y = 1$. So

$$\begin{aligned} s^1 &= s^{jx+|s|y} \\ s^1 &= s^{jx} s^{|s|y} \\ s^1 &= (s^j)^x (s^{|s|})^y \\ s^1 &= (s^j)^x (e)^y \end{aligned}$$

Therefore $s^1 = (s^j)^x$ and $\langle s \rangle = \langle s^j \rangle$. □

2. THE MODULAR GROUP

Recall that the modular group, M , has presentation $\langle s, t \mid s^8 = t^2 = e, st = ts^5 \rangle$. Next we determine the existence of Eulerian and Hamiltonian circuits and paths.

Lemma 5. *If G is the modular group and n is odd, then*

$$\langle ts^n \rangle = \langle ts \rangle = \{ts, s^6, ts^7, s^4, ts^5, s^2, ts^3, e\}.$$

Lemma 6. *If G is the modular group, then $\langle ts^2 \rangle = \langle ts^6 \rangle = \{ts^2, s^4, ts^6, e\}$.*

Lemma 7. *If G is the modular group and $n = 2$ or 6 , then $|\langle s \rangle \cap \langle ts^n \rangle| = 2$.*

Lemma 8. *If G is the modular group and n is odd, then $|\langle s \rangle \cap \langle ts^n \rangle| = 4$.*

Theorem 9. *If G is the modular group, and S is a minimal generating set, then $\Gamma(G, S)$ contains an Eulerian circuit.*

Proof. Let G be the modular group and S be a minimal generating set. Then $S = \{s^n, ts^k\}$, where n is odd, $1 \leq n \leq 7$, and $0 \leq k \leq 7$ or $S = \{ts^n, ts^m\}$, where n is odd and m is even. By using the lemmas above there exists three distinct graphs.

case i) Let $S = \{s^n, t\}$ where n is odd, then $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle t \rangle| = 1$ and $deg(\langle s^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$, which is even.

Similarly $deg(\langle t \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$, which is even. This graph is $K_{2,8}$ and contains an Eulerian circuit.

case ii) Let $S = \{s^n, ts^m\}$ where n, m are odd, then $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^m \rangle| = 4$ and $deg(\langle s^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{4} - 1 = 2$, which is

even. Similarly $deg(\langle ts^m \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{8}{S_{2,1}} + \frac{8}{S_{2,2}} - 1 = \frac{8}{4} + \frac{8}{8} - 1 = 2$, which is even. This graph is $K_{2,2}$ and contains an Eulerian circuit.

case iii) Let $S = \{s^n, ts^k\}$ where n is odd and $k = 2$ or 6 , then $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^k \rangle| = 2$ and $deg(\langle s^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{2} - 1 = 4$,

which is even. Similarly $\deg(\langle ts^k \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{4}{2} + \frac{4}{4} - 1 =$

2, which is even. This graph is $K_{2,4}$ and contains an Eulerian circuit.

case iv) Let $S = \{s^n, ts^4\}$ where n is odd, then $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^4 \rangle| = 1$ and $\deg(\langle s^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$, which is even.

Similarly $\deg(\langle ts^4 \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$, which is even. This graph is $K_{2,8}$ and contains an Eulerian circuit.

case v) Let $S = \{ts^n, t\}$ where n is odd, then $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle t \rangle| = 1$ and $\deg(\langle ts^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$, which is even.

Similarly $\deg(\langle t \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$, which is even. This graph is $K_{2,8}$ and contains an Eulerian circuit.

case vi) Let $S = \{ts^n, ts^k\}$ where n is odd and $k = 2$ or 6 , then $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle ts^k \rangle| = 2$ and $\deg(\langle ts^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{2} - 1 = 4$, which is even. Similarly $\deg(\langle ts^k \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{4}{2} + \frac{4}{4} - 1 = 2$, which is even. This graph is $K_{2,4}$ and contains an Eulerian circuit.

case vii) Let $S = \{ts^n, ts^4\}$ where n is odd, then $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle ts^4 \rangle| = 1$ and $\deg(\langle ts^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$, which is

even. Similarly $\deg(\langle ts^4 \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 = \frac{2}{1} + \frac{2}{2} - 1 = 2$, which is even. This graph is $K_{2,8}$ and contains an Eulerian circuit. \square

Remark 2. For all minimal generating sets, $\Gamma(M, S)$ does not contain an Eulerian path.

Theorem 10. *If G is the modular group, and $S = \{s^n, ts^m\}$ where n, m are odd, then $\Gamma(G, S)$ contains a Hamiltonian circuit and a Hamiltonian path.*

Proof. The vertex set of $\Gamma(M, S)$ is $V(\Gamma(M, S)) = \{\langle s^n \rangle, t\langle s^n \rangle, \langle ts^m \rangle, t\langle ts^m \rangle\}$. A Hamiltonian circuit is given by

$$\langle s^n \rangle, \langle ts^m \rangle, t\langle s^n \rangle, t\langle ts^m \rangle, \langle s^n \rangle.$$

A Hamiltonian path is given by

$$\langle s^n \rangle, \langle ts^m \rangle, t\langle s^n \rangle, t\langle ts^m \rangle.$$

\square

Remark 3. $S = \{s^n, ts^m\}$ where n, m are odd is the only minimal generating set of M that yields a graph that contains a Hamiltonian circuit (path).

3. THE QUASIHEDRAL GROUP

Recall that the quasihedral group, QS , has presentation $\langle s, t \mid s^8 = t^2 = e, st = ts^3 \rangle$. Next we determine the existence of Eulerian and Hamiltonian circuits and paths.

Lemma 11. *If G is the quasihedral group and n is 1 or 5, then $\langle ts^n \rangle = \{ts, s^4, ts^5, e\}$.*

Lemma 12. *If G is the quasihedral group and n is 3 or 7, then $\langle ts^n \rangle = \{ts^3, s^4, ts^7, e\}$.*

Lemma 13. *If G is the quasihedral group and n is even, then $\langle ts^n \rangle = \{ts^n, e\}$.*

Lemma 14. *If G is the quasihedral group and n is even, then $|\langle s \rangle \cap \langle ts^n \rangle| = 1$.*

Lemma 15. *If G is the quasihedral group and n is odd, then $|\langle s \rangle \cap \langle ts^n \rangle| = 2$.*

Theorem 16. *If G is the quasihedral group, and S is a minimal generating set, then $\Gamma(G, S)$ contains a Eulerian circuit.*

Proof. Let G be the quasihedral group and S be a minimal generating set. Then $S = \{s^n, ts^k\}$, where n is odd and $1 \leq n \leq 7$ and $1 \leq k \leq 3$ or $S = \{ts^n, ts^m\}$, where n is odd and m is even. By using the above lemmas, there exists three distinct graphs.

case i) Let $S = \{s^n, ts^m\}$, where n, m are odd, then $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^m \rangle| = 2$ and $deg(\langle s^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{2} - 1 = 4$, which is

even. Similarly, $deg(\langle ts^m \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{4}{S_{2,1}} + \frac{4}{S_{2,2}} - 1 = \frac{4}{2} + \frac{4}{4} - 1 = 2$,

which is even. This graph is $K_{2,4}$ and contains an Eulerian circuit.

case ii) Let $S = \{s^n, ts^m\}$, where n is odd and m is even, then $S_{1,2} = S_{2,1} = |\langle s^n \rangle \cap \langle ts^m \rangle| = 1$ and $deg(\langle s^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{8}{S_{1,1}} + \frac{8}{S_{1,2}} - 1 = \frac{8}{8} + \frac{8}{1} - 1 = 8$,

which is even. Similarly, $deg(\langle ts^m \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 =$

$\frac{2}{1} + \frac{2}{2} - 1 = 2$, which is even. This graph is $K_{2,8}$ and contains an Eulerian circuit.

case iii) Let $S = \{ts^n, ts^m\}$ where n is odd and m is even, then $S_{1,2} = S_{2,1} = |\langle ts^n \rangle \cap \langle ts^m \rangle| = 1$ and $deg(\langle ts^n \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{4}{S_{1,1}} + \frac{4}{S_{1,2}} - 1 = \frac{4}{4} + \frac{4}{1} - 1 =$

4 , which is even. Similarly $deg(\langle ts^m \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{2}{S_{2,1}} + \frac{2}{S_{2,2}} - 1 =$

$\frac{2}{1} + \frac{2}{2} - 1 = 2$, which is even. By applying Euler's theorem, this graph contains an Eulerian circuit. \square

Remark 4. For all minimal generating sets, $\Gamma(QS, S)$ does not contain an Eulerian path, a Hamiltonian path, or a Hamiltonian circuit.

4. GENERALIZED QUATERNION GROUP

Recall that the generalized quaternion group, Q_{2^n} , has presentation $\langle s, t \mid s^{2^{n-1}} = e, s^{2^{n-2}} = t^2, tst^{-1} = s^{-1} \rangle$. Next we determine the existence of Eulerian and Hamiltonian circuits and paths.

Lemma 17. *If G is the generalized quaternion group, then $t^4 = e$.*

Proof. Let G be the generalized quaternion group. Recall that $t^2 = s^{2^{n-2}}$. Squaring both sides,

$$\begin{aligned} (t^2 = s^{2^{n-2}})^2 \\ t^4 = s^{2^{n-1}} = e. \end{aligned}$$

□

Lemma 18. *Let G be the generalized quaternion group, then $(ts^j)^2 = t^2$ for all j .*

Proof. We proceed with induction on j . Let $j = 1$, then $(ts^1)^2 = tsts = ts(s^{-1}t) = t^2$ and the theorem holds for $j = 1$. Assume that the theorem holds for $j = k$, i.e., $(ts^k)^2 = t^2$.

Now let $j = k + 1$, then $(ts^{k+1})^2 = ts^{k+1}ts^{k+1} = ts^{k+1}ts^k = ts^{k+1}s^{-1}ts^k = ts^kts^k = (ts^k)^2 = t^2$. Therefore $(ts^j)^2 = t^2$ for all j .

□

Lemma 19. *Let G be the generalized quaternion group, then $\langle ts^j \rangle = \{ts^j, t^2, t^3s^j, e\}$ for all j .*

Lemma 20. *If G is the generalized quaternion group and $\langle ts^j \rangle \neq \langle ts^k \rangle$, then $\langle ts^j \rangle \cap \langle ts^k \rangle = \{t^2, e\}$ and $|\langle ts^j \rangle \cap \langle ts^k \rangle| = 2$.*

Theorem 21. *If G is the generalized quaternion group, and S is a minimal generating set, then $\Gamma(G, S)$ contains an Eulerian circuit.*

Proof. Let G be the generalized quaternion group and S be a minimal generating set. Then, $S = \{s^k, ts^j\}$ where k is odd or $S = \{ts^k, ts^m\}$, where k is odd and m is even.

case i) Let $S = \{s^k, ts^j\}$ where k is odd, then $S_{1,2} = S_{2,1} = |\langle s^k \rangle \cap \langle ts^j \rangle| = 2$ and $deg(\langle s^k \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{2^{n-1}}{S_{1,1}} + \frac{2^{n-1}}{S_{1,2}} - 1 = \frac{2^{n-1}}{2^{n-1}} + \frac{2^{n-1}}{2} - 1 = \frac{2^{n-1}}{2} = 2^{n-2}$,

which is even. Similarly $deg(\langle ts^j \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{4}{S_{2,1}} + \frac{4}{S_{2,2}} - 1 = \frac{4}{4} + \frac{4}{2} - 1 = 2$,

which is even. This graph is $K_{2,2^{n-2}}$ and contains an Eulerian circuit.

case ii) Let $S = \{ts^k, ts^m\}$, where k is odd and m is even, then $S_{1,2} = S_{2,1} = |\langle ts^k \rangle \cap \langle ts^m \rangle| = 2$ and $deg(\langle ts^k \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_1 \rangle|}{S_{1,j}} \right) - 1 = \frac{4}{S_{1,1}} + \frac{4}{S_{1,2}} - 1 = \frac{4}{4} + \frac{4}{2} - 1 = 2$,

which is even. Similarly $deg(\langle ts^m \rangle) = \left(\sum_{j=1}^2 \frac{|\langle s_2 \rangle|}{S_{2,j}} \right) - 1 = \frac{4}{S_{2,1}} + \frac{4}{S_{2,2}} - 1 =$

$\frac{4}{4} + \frac{4}{2} - 1 = 2$, which is even. By applying Euler's theorem, this graph contains an Eulerian circuit.

□

Remark 5. For all minimal generating sets, $\Gamma(Q_{2^n}, S)$ does not contain an Eulerian path.

Theorem 22. *If G is the generalized quaternion group, Q_{2^n} , and $S = \{s^k, ts^m\}$ where k is odd, then $\Gamma(G, S)$ contains a Hamiltonian circuit and a Hamiltonian path for $n = 3$.*

Proof. The vertex set of $\Gamma(Q_{2^3}, S)$ is $V(\Gamma(Q_{2^3}, S)) = \{\langle s^k \rangle, t\langle s^k \rangle, \langle ts^m \rangle, t\langle ts^m \rangle\}$. A Hamiltonian circuit is given by

$$\langle s^k \rangle, \langle ts^m \rangle, t\langle s^k \rangle, t\langle ts^m \rangle, \langle s^k \rangle.$$

A Hamiltonian path is given by

$$\langle s^k \rangle, \langle ts^m \rangle, t\langle s^k \rangle, t\langle ts^m \rangle.$$

□

Theorem 23. *If G is the generalized quaternion group, Q_{2^n} , and $S = \{ts^k, ts^m\}$, where k is odd and m is even, then $\Gamma(G, S)$ contains a Hamiltonian circuit and a Hamiltonian path.*

Proof. The vertex set of $\Gamma(Q_{2^n}, S)$ is

$$V(\Gamma(Q_{2^n}, S)) = \{\langle ts^k \rangle, s\langle ts^k \rangle, \dots, s^{2^{n-2}-1}\langle ts^k \rangle, \langle ts^m \rangle, s\langle ts^m \rangle, \dots, s^{2^{n-2}-1}\langle ts^m \rangle\}.$$

A Hamiltonian circuit is given by

$$\langle ts^k \rangle, \langle ts^m \rangle, s^{k-m}\langle ts^k \rangle, s^{k-m}\langle ts^m \rangle, \dots, s^{k-(2^{n-2}-1)m}\langle ts^k \rangle, s^{k-(2^{n-2}-1)m}\langle ts^m \rangle, \langle ts^k \rangle.$$

A Hamiltonian path is given by

$$\langle ts^k \rangle, \langle ts^m \rangle, s^{k-m}\langle ts^k \rangle, s^{k-m}\langle ts^m \rangle, s^{k-2m}\langle ts^k \rangle, s^{k-2m}\langle ts^m \rangle, \dots, s^{k-(2^{n-2}-1)m}\langle ts^k \rangle, s^{k-(2^{n-2}-1)m}\langle ts^m \rangle.$$

□

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