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<b>Auteurs:</b> Authors:	Mohamad Gebai et Michel R. Dagenais
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# Survey and Analysis of Kernel and Userspace Tracers on Linux: Design, Implementation, and Overhead

Mohamad Gebai, Michel R. Dagenais Polytechnique Montreal

**Abstract**—As applications and operating systems are becoming more complex, the last decade has seen the rise of many tracing tools all across the software stack. This paper presents a hands-on comparison of modern tracers on Linux systems, both in user space and kernel space. The authors implement microbenchmarks that not only quantify the overhead of different tracers, but also sample fine-grained metrics that unveil insights into the tracers’ internals and show the cause of each tracer’s overhead. Internal design choices and implementation particularities are discussed, which helps to understand the challenges of developing tracers. Furthermore, this analysis aims to help users choose and configure their tracers based on their specific requirements in order to reduce their overhead and get the most of out of them.

## I. INTRODUCTION

Tracing has proved itself to be a robust and efficient approach to debugging and reverse-engineering complex systems. The past decade has seen the rise of many tracers across all layers of the software stack, and even at the hardware level [Intel CorporationIntel Corporation2016] [Sharma and DagenaisSharma and Dagenais2016a]. Some applications, such as Google Chrome, even provide tracers natively integrated within the product itself. Fundamentally, tracing is a sophisticated form of logging, where a software component, called the tracer, provides a framework that implements efficient and configurable logging. The most common use of logging by developers is done via the `printf()` function (or an equivalent), although this method is largely inefficient and limited. Tracers provide more flexible and robust approaches that can be easily maintained over time and usually add little overhead. Tracing is common in user applications, but is also widely used in the Linux kernel, which provides multiple tracing infrastructures. With complex online distributed systems, tracing becomes an efficient way of debugging problems whenever they arise. Although it is a problem that is often underestimated, the need for efficient and low-impact tracers is increasing, especially with modern parallel heterogeneous systems of ever increasing complexity. The Linux kernel contains over 1000 tracepoints and the volume of events that can be generated at runtime reinforces the need for low-impact tracers. In this article, we focus on the overhead that different tracers add to the traced applications, both at the user and kernel levels, on Linux systems. We start by establishing the nomenclature used in this work, and categorizing the many tools that have been gathered under the term “tracer”. We explain the differences

between them from the user point of view, we summarize the mechanisms used by each to perform its tasks, and we show key design and implementation decisions for each when relevant. We propose a microbenchmark that aims at providing reliable low-level and fine-grained metrics for an in-depth analysis and comparison.

In this work, we highlight the performance and the footprint of multiple tracers, as well as their underlying infrastructure. Many commercial and broadly-known tools rely on the tracing infrastructure variants studied here, and thus the overhead measured directly applies. The contribution of this paper is a deep dive into the design principles of modern tracers on Linux. This work tackles the problem of comparing tracers by measuring fine-grained and low-level performance metrics, so that the design choices made by the tracer developers, as well as the implementation and coding decisions, are taken into consideration when assessing the impact of a tracing tool on the traced system. Moreover, the contribution also encompasses a methodology for low-level benchmarking that unravels the real behavior of tracers, instead of using platform and micro-architecture emulators.

The rest of the paper is structured as follows: section II goes over previous work on tracers benchmarks, section III includes a reminder of key concepts required to understand the work presented and sets up the nomenclature used in this paper, section IV explains the mechanisms used in tracing, section V introduces the tracers and explains their internals when relevant, section VI explains the benchmarks, section VII shows the results of this work, section VIII concludes.

## II. PREVIOUS WORK

Bitzes et al. [Bitzes and NowakBitzes and Nowak2014] studied the overhead of sampling using performance counters. However, their studies don’t address tracing in general. It also only focuses on the approaches for collecting hardware counters data and their performance overhead, rather than covering the internals of the tracers, or design and implementation choices. Sivakumar et al. [Sivakumar and Sundar RajanSivakumar and Sundar Rajan2010] measured the impact of the LTTng tracer, both in user and kernel space. The authors ran multiple known general benchmarks and reported the overhead that the tracer added. This approach helps to estimate the impact that LTTng may have on specific workloads but doesn’t quantify in detail the cost of the instrumentation,

or the cause of the overhead. Mihajlovic et al. [Mihajlović, Žilić, and GrossMihajlović et al.2014] discuss their work to enable hardware tracing in virtual environments by modifying the ARM ISA emulation and show the overhead added by their approach. While the approach it presents is interesting, their approach relies on dynamic tracing which is a specific approach to tracing. Furthermore, it doesn't cover or benchmark the fundamental work of tracers. Furthermore, no detailed comparison with other tracers is presented. Moore et al. [Moore, Cronk, London, and DongarraMoore et al.2001] reviewed performance analysis tools for MPI application. While they cover two of the tracers of this work, the work focuses primarily on the MPI framework, and doesn't address kernel space tracing. The objective of the work also differs in that fundamental design choices of tracers with different scopes aren't covered. In [GhodsGhods2016], the author explains and analyzes the internals of the Perf tool, mainly for sampling performance hardware counters. This work doesn't cover comparisons with other tracing tools. The work of Desnoyers [DesnoyersDesnoyers2009] reports benchmarking results for LTTng and other tracers, albeit only showing the latency of recording an event, without presenting a detailed comparison with other tracers.

The infrastructures and frameworks presented in this article are often the basis for other commercial and more broadly-known monitoring and performance tools. For instance, the work by B. Gregg [Gregg2017] makes extensive use of Perf and eBPF. Flamegraphs are often generated by profiling applications using Perf, although it is a profiler rather than a tracer. Another example is SysDig, which uses the tracepoint infrastructure to extract information from the kernel.

### III. DEFINITIONS AND NOMENCLATURE

#### A. Definitions

This section serves as a reminder of some common terms in the world of tracing that are essential for understanding the rest of the paper.

**Tracepoint:** A tracepoint is a statement placed directly in the code of an application that provides a hook to invoke a probe. A tracepoint generally provides a way in which it can be enabled or disabled dynamically.

**Probe:** A probe is a function that is hooked to a tracepoint and is called whenever the tracepoint is encountered at run-time (if enabled). A probe usually performs a custom task and is either implemented by the tracer or by the user. Typically, a probe needs to be as small and fast as possible, to add as little overhead as possible and reduce the perturbation to the system.

**Event:** An event marks the encounter of a tracepoint at run-time. Depending on the kind of tracing, the event can have a logical meaning, such as a context switch, or can simply represent a location in the code, such as a function entry or exit. An event is punctual and has no duration, and is usually annotated with a timestamp.

**Payload:** An event typically holds a payload, which is additional information related to the event. For instance, the payload of a context switch may contain the identifiers of the

two tasks involved.

**Ring buffer:** A data structure that serves as a placeholder for events. When an event is recorded at run-time, the tracer's probe is invoked. The probe records the encountered event by placing it in memory in the ring buffer (producer). At a later time, a consumer can read the contents of the ring buffer and report them to the user. When the data structure is full, the incoming events may either replace the oldest events (in a ring-like fashion) or they may be discarded until some events have been consumed from the buffer.

**Atomic operation** An atomic operation has the characteristic of being indivisible, which implies that intermediate values or intermediate states are invisible to concurrent operations. Atomic operations usually require support from the hardware or the operating system, and great care must be taken by the developer to guarantee atomicity. For instance, on x86 architectures, a `mov` instruction isn't guaranteed to be atomic unless its operands are cache-aligned. Consider the case where one of the operands is stored across two pages: the `mov` instruction will require access to different pages (and potentially cause virtual address translations), making the operation divisible and non-atomic, as an intermediate unstable value can be visible if the operand is accessed by another instruction between these steps.

#### B. Nomenclature

We previously defined a tracepoint as a location in an application where a probe can be hooked. This section starts by introducing the different mechanisms used for probe callback, as well as their implementations. A mechanism is a known theoretical approach as to how a callback can be implemented, but the actual implementation is left to the tracing infrastructure. For instance, a trampoline is a mechanism that allows instrumentation at run-time, but the actual implementation of the trampoline is left to a tracing infrastructure such as `DynInst` or `Kprobes`. Similarly, a tool can support multiple mechanisms and allow its users to configure the mechanism to be used, depending on their needs. Tracers can then be built atop of these technologies to leverage their callback mechanisms, thus outsourcing this crucial part. Tracers can be built to support multiple callback mechanisms, for better flexibility and feature offerings. In summary, a tracer can use one or many callback implementations, which in turn implement one or many mechanisms. For instance, LTTng can use either `TRACE_EVENT` or `Kprobes`, and `Kprobes` can use either a trap or a trampoline.

We define a tracer as a tool that implements the following pattern: *callback, serialize, write*<sup>1</sup>. The output of a tracer is a trace, and efforts are dedicated to reducing as much as possible its overhead. On the other hand, tools such as eBPF and SystemTap fundamentally follow a different pattern: *callback, compute, update*. We refer to them as aggregators, since their work is often to collect and aggregate metrics in real time or in a live fashion, on the critical path of the applications, contrary to the post-mortem nature of trace analysis. To this end, they provide users with scripting capabilities and advanced

<sup>1</sup>The write step here refers to a ring buffer rather than to a file

data structures (such as hashmaps) to implement aggregation methods to be executed upon certain events. As opposed to tracers, the output of aggregators is the result of the user-defined probe, which is typically a collection of metrics, an alert when a threshold is exceeded at runtime, and so on. They often neglect the timing aspect and don't implicitly perform a clock read on each callback.

#### IV. CALLBACK MECHANISMS

This section introduces the different mechanisms used to instrument applications. The instrumentation can be static or dynamic, where the former is built into the binary at compile-time and tracepoint location is known in advance, and the latter is inserted at run-time at user-defined locations. This is not to be confused with dynamic tracing, which means that tracing can be turned on or off at run-time. Dynamic tracing can be supported for either static or dynamic instrumentation. To give better insights on how the mechanisms work, we cover their implementation in various technologies to show how they are effectively implemented and used.

##### A. Function instrumentation

Function instrumentation is a static instrumentation method that requires support from the compiler. The approach is to have each function call of an application prefaced by a call to a common tracing probe. In other words, the binary contains explicit calls to a specific routine upon each function entry (and exit in some cases). The implementation of this routine is left to the developer, or the tracer, and can implement tracing, profiling or any other monitoring feature.

GCC implements this callback mechanism in various ways. For instance, the `-pg` flag will generate a binary where each function has the `mcount` routine call as a preamble [FrysingerFrysinger2016]:

```
$> echo "int main() {}" | gcc -pg -S -x c \
  - -o /dev/stdout | grep mcount

    call    mcount
```

Since `mcount` is called at each function entry, additional efforts need to be put into its implementation to provide tracing at the lowest cost possible. `ftrace` uses the `mcount` implementation to trace kernel functions entries, and implements the `mcount` routine in platform-specific assembly. `LTTng UST` also uses this method for userspace tracing, albeit with the `-finstrument-functions` flag of `gcc`. Similarly to `-pg`, calls to specific routines are built into the binary not only at each function call entry, but also at function exit:

```
$> echo "int main() {}" | gcc
  -finstrument-functions -S -x c - -o \
  /dev/stdout | grep cyg_profile

    call    __cyg_profile_func_enter
    call    __cyg_profile_func_exit
```

With the `-finstrument-functions` flag, the instrumentation routines are called

```
__cyg_profile_func_enter() and
__cyg_profile_func_exit() for function entry
and exit respectively.
```

##### B. Static Tracepoints

A tracepoint is a static instrumentation approach manually inserted directly in the application code by the developers.

Unlike regular function calls, tracepoint statements in the Linux kernel are optimized to have a minimal impact on performance. As the instrumentation is directly in the code and always built into the binary (unless the kernel is configured otherwise), great care must be taken to reduce the added overhead, especially when tracing is disabled, as is the case most of the time. The rest of this subsection discusses how this goal is achieved in the Linux kernel.

A disabled tracepoint has no effect and translates to a simple condition check for a branch [DesnoyersDesnoyers2016c] (in case it is enabled). To reduce the overhead for a disabled tracepoint, a hint is given to the compiler to make the tracepoint instructions far from the cache lines of the fast path (which is the regular code). In that way, the cache-friendliness of the fast path isn't affected by the unexecuted code of the tracepoint. Furthermore, for kernel tracing, the tracepoint call is implemented as a C macro that translates to a branch over a function call. In that manner, the overhead of the function call and stack setup is avoided altogether.

Although the overhead impact of this approach is minimal in theory, it still requires reading from main memory the operand of the condition, to avoid the branch when the tracepoint is off. This adds non-negligible overhead as reading from memory not only is a slow process, but ultimately affects the efficiency of the pipeline. To overcome this issue, the Immediate Value infrastructure was set in place [Desnoyers and DagenaisDesnoyers and Dagenais2008] by the `LTTng` project. This mechanism uses a constant value directly into the instruction's operand. In that manner, no read from memory is required for the condition's operand. A disassembly of the generated tracepoint code clearly shows the use of immediate values:

```
test    %edx,%edx
jne     [tracepoint tag]
```

When a tracepoint is turned on, the code is safely modified on the fly to change the value of the constant test check. Synchronization implications need to be taken into account for this process as the code resides in memory, and is shared amongst multiple CPUs, but copies may exist in instruction caches.

In the Linux kernel, static tracepoints are implemented as the `TRACE_EVENT` macro [RostedtRostedt2010], which allows developers to easily define a tracepoint that can be inserted using the `trace_tracepoint_name()` function directly in the kernel code. Many of the current kernel tracers can interface to the `TRACE_EVENT` infrastructure by providing their own probes. This infrastructure also makes for an easy mechanism for a developer to implement their own tracer or aggregator.

For static tracepoints in user space, the LTTng tracer implements this mechanism with the same optimizations as for kernel space, reducing as much as possible the overhead added by the instrumentation.

### C. Trap

Trap-based instrumentation is a mechanism to dynamically instrument an application at run-time. It relies on the operating system support for traps, which it exploits to insert and execute custom probes at virtually any location in the kernel or application code. In the Linux kernel, this mechanism is implemented by the Kprobe infrastructure, which uses a trap-based approach to dynamically hook into kernel code [Keniston J. Keniston J. 2016] [Mavinakayanahalli, Panchamukhi, Keniston, Keshavamurthy, and Hiramatsu Mavinakayanahalli et al. 2006]. When a Kprobe is loaded and registered at a given instruction, that instruction is copied and replaced with a breakpoint instruction (`int3` on x86). When the breakpoint instruction is later executed by the CPU, the kernel's breakpoint handler is invoked. It saves the state of the application (registers, stack, etc.) and gives control to the Kprobe infrastructure using the Linux notifier call chain, which ends up calling the tracing probe. Once this process completes and the trap has been handled, the copied instruction (that was replaced by the breakpoint) is finally executed and control continues normally at the call site. Image 1 shows this entire process. Kprobes also offer support for pre-handler and post-handler probes, as well as function return instrumentation (Kretprobe) which aren't covered in this study.

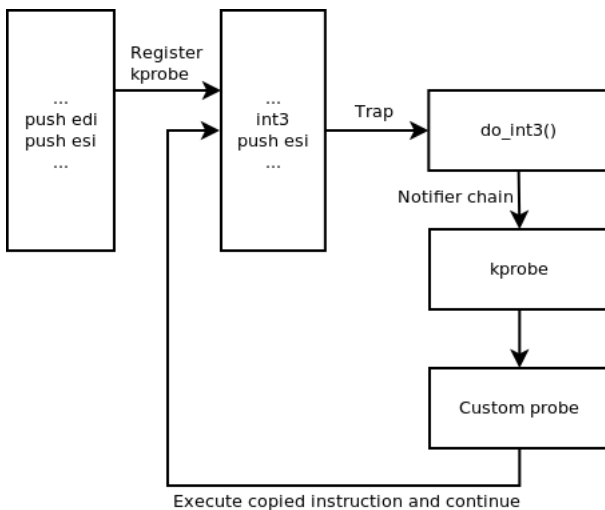


Fig. 1: Trap-based callback mechanism with Kprobes

A tracing infrastructure can be built atop Kprobes, where instead of manually inserting the `trace_tracepoint_name()` statement in the code at the call site, a Kprobe can be registered at the desired location at run-time. Tracers leverage this approach to connect their probe to a Kprobe instead of the `TRACE_EVENT` macro. In that manner, the callback mechanism is abstracted and only connecting the tracer's probe to different backends provides more flexibility to the user. The resulting trace is identical to one that is generated using the

`TRACE_EVENT` macro, but the callback mechanisms used to invoke the probes are different, which can have an impact on performance. When a Kprobe is unloaded, the breakpoint instruction is replaced with the original one, thus removing completely any trace of the instrumentation.

The Ptrace infrastructure in the kernel [Haardt and Coleman Haardt and Coleman 1999] [Padala Padala 2002] also uses traps to offer to userspace applications a mechanism to hook onto processes. It is important to note that, contrary to what its name suggests, Ptrace is not a tracer in itself, but rather an infrastructure provided by the Linux kernel for processes to monitor other processes. It allows a process to “hook” into another one and interrupt its execution, inspect its internal data, access its registers, etc. Many debuggers use Ptrace as a backend, including GDB.

### D. Trampoline

Trampolines are a jump-based approach to dynamically patch or instrument an application at runtime. They provide a lower overhead alternative to trap-based mechanisms for the price of a more complex implementation. In more recent versions, the Linux kernel tries to optimize registered Kprobes using a jump-based trampoline instead of a costly breakpoint. The core of the optimization is to have a “detour” buffer (called the optimized region) to mimic the breakpoint approach [Keniston J. Keniston J. 2016]. Instead of patching an instruction with a breakpoint instruction that triggers a trap, it is replaced by a simple jump to the optimized region. The jump-based approach starts by pushing the CPU's registers onto the stack, jumps to a trampoline that acts as an intermediate, which in turn jumps to the user-defined probe [Hiramatsu Hiramatsu 2010]. When it completes execution, the process is reversed: the code jumps out of the optimized region, the registers are restored from the stack and execution of the original path continues. Note that not all loaded Kprobes use the trampoline approach, as it requires a set of conditions to be met (e.g., length of the instruction at the target location). If they aren't, the kernel falls back to the breakpoint-based approach described previously.

## V. THE TRACERS

This section introduces the tracers studied and benchmarked in this work. When relevant, details of the design and implementation are provided for each tracer, and are later correlated with the results. Table I, presented at the end of this section, shows a summary of the tracers, as well as the mechanisms used by each (see section IV for an explanation of the mechanisms).

This work does not evaluate Dtrace [Gregg and Mauro Gregg and Mauro 2011] and Ktap [Kata 2017]. The former is a proven tracer on Solaris and considered one of the pioneers in the field of tracing. However, it does not appear to be actively developed any more, with a total of 13 mailing list postings in the first half of 2017 on dtrace.org, its Linux port never reached the stability or broad usage of the Solaris port, and its strength lied more

in its flexibility and ease of use rather than its optimized performance and scalability [BrousseauBrousseau2017]. Ktap was an interesting lightweight dynamic tracing tool experiment based on bytecode, but was quickly superseded by eBPF which offers similar functionality but sharing a core infrastructure with other kernel subsystems.

### A. Kernel tracers

1) *None*: The samples marked as 'None' designate the baseline, which represents the system with all tracing disabled, effectively only benchmarking the instrumentation itself which is negligible (translates to a constant check as explained earlier).

2) *Ftrace*: Ftrace is a tracer included in the Linux kernel and shows insights into its internal behavior [RostedtRostedt2009a] [RostedtRostedt2016a] [RostedtRostedt2009b] by tracing all kernel function entries. It is controlled by a set of files in the `debugfs` pseudo-filesystem. Some of the main configurations include which "subtracer" to use, the size of the trace buffers and which clock source to use to timestamp the events. It is also possible to enable and disable specific events. Ftrace can be used in many ways: function tracing, tracepoints, system calls tracing, dynamic instrumentation and so on.

Function (and function graph) tracing reports the entry and exit of all functions at the kernel level. Ftrace can use the `TRACE_EVENT` infrastructure for static instrumentation tracing or the `Kprobe` infrastructure to dynamically hook into various parts of the kernel.

When tracing is enabled, the callback mechanism calls Ftrace's probe which stores the events in a ring buffer. It is possible to configure Ftrace to either overwrite the oldest events or drop incoming events, once the ring buffer is full. It is interesting to note that the trace is kept in memory and not flushed to disk, and only made available upon reading the contents of the `trace` memory-backed file (it is possible to manually dump the contents of the `trace` file to disk). It is also possible to consume the ring buffer as it is written, through the `trace_pipe` file.

By default, Ftrace uses the local clock to timestamp the events it records. The local clock is a clock source that is CPU-local and is thus faster to read, but doesn't provide any guarantee in terms of monotonicity and synchronization with the other CPUs' clocks. It is, however, possible to configure Ftrace to use other clock sources, such as a global clock (which is system-wide), a logical counter or even architecture-specific clocks such as the TSC<sup>2</sup> on x86.

Ftrace limits the size of an event, including its payload, to that of a page. It uses per-CPU buffers, which avoids the need for synchronization of the buffer when tracing on multiple cores. Ftrace segments its ring buffers into pages and manipulates them individually. The ring buffer itself is a linked list of pages [RostedtRostedt2016b], and internal references are kept for bookkeeping. For instance, the `tail_page` is a reference to the page into which the next event should be written, and the `commit_page` is a reference to the page that

last finished a write. Although there can't be simultaneous writers to the same page (per-CPU buffers), a writer can still be interrupted by another writer by the means of interrupts and NMIs. An example of that can be an event to be written into the buffer, from an interrupt handler context that was invoked while a write to the ring buffer was already happening, as shown in the following sequence:

```
Start writing
-> Interrupt raised
Enter interrupt handler
Start writing
Finish writing
Exit interrupt handler
Finish writing
```

These nested writes require the tracer and its ring buffer to guarantee reentrancy in order to avoid data corruption and misbehavior. Note that there can be more than two levels of preemption (normal execution, interrupts, non-maskable interrupts, machine check exceptions, etc.). The way Ftrace ensures reentrancy is by dividing the writing process into three parts: reserving, writing, committing. The writing process starts by reserving a slot in memory in an atomic fashion, making it indivisible and thus guaranteeing reentrancy for this step. Only then can the writing into the reserved slot step begin. If a nested write occurs, it has to follow the same pattern, starting with a slot reservation. Since this can only happen after the preempted write has already completed its slot reservation (since it is indivisible), there can be no contention over the writing area, making the writing process safe. Once the nested write completes, it can commit which seals the writing transaction. The interrupted write can then complete its write and commit:

```
Reserve slot
Start writing
-> Interrupt raised
Enter interrupt handler
Reserve slot
Start writing
Finish writing
Exit interrupt handler
Finish writing
```

By following this scheme, writing transactions appear as being atomic, in the sense that no two nested writes can write to the same slot in the buffer. When nested writes occur, some subtleties need to be implemented. For instance, the nested write cannot commit before the write it preempted. Until then, it is in the "pending commit" state. This is required since all events prior to the `commit_page` have to be committed (the `commit_page` actually points to the latest *event* that was committed, and masking its least significant bits gives the address of the page).

Furthermore, in order to consume tracing data, Ftrace keeps an additional page, called the `reader_page`, which is not a part of the ring buffer. Rather, it is used as an interim to extract the trace data from the buffer. When a reader wants to consume trace data, the `reader_page` is atomically swapped with the

<sup>2</sup>TimeStamp Counter

*head\_page*. The *head\_page* is simply an internal reference to the page that should be read next. As this swap happens, the old *reader\_page* becomes part of the ring buffer, and the old *head\_page* can be safely consumed. This swapping happens in an atomic fashion using CAS (compare-and-swap) operations. After the swap happens, the *head\_page* pointer is updated to reference the next page to be read.

We explained earlier in this section that Ftrace manipulates pages individually in its ring buffer. Although this implementation choice has benefits such as avoiding lazy memory allocations (and often a lazy TLB update), it also results in two main limitations: the size of an event (including its payload) is limited to the size of a page, and memory barriers are required at page boundaries. The former limitation is due to the fact that single pages are consumable. Consumption of the buffer can be done at page granularity, which implies that a single event cannot be stored across page boundaries. The latter limitation is less obvious; a page of the buffer can only be read once it is full. Thus, page management variables are required for internal bookkeeping, such as flagging a page as ready to be consumed. These variables are shared by the writer and any reader. To guarantee coherent ordering between buffer data and buffer management variables, memory barriers are required to ensure that a page is not flagged as full before its data is actually propagated to main memory.

3) *LTTng*: *LTTng* was created in 2006 [Desnoyers and DagenaisDesnoyers and Dagenais2006b] [DesnoyersDesnoyers2009] [Desnoyers and DagenaisDesnoyers and Dagenais2009] around the same time as Ftrace and thus both tracers share many similarities design-wise. However, *LTTng* isn't part of the mainline Linux kernel and is rather deployed as a group of loadable kernel modules for the kernel tracing part, and a userspace component for tracing management and trace consumption (as opposed to Ftrace's `debugfs` interface). *LTTng* was designed and implemented with the objective of minimal performance impact [Desnoyers and DagenaisDesnoyers and Dagenais2006a] while being fully-reentrant, interrupt-safe and NMI-safe. Similarly to Ftrace, *LTTng* uses per-CPU variables and buffers to avoid concurrent writes and the need for synchronization. Reentrancy is guaranteed by the means of atomic slot reservation using local CAS (Compare-And-Swap) to permit nested writes [Desnoyers and DagenaisDesnoyers and Dagenais2012], similarly to what was explained in section V-A2. Image 2 shows how having (local) atomic slot reservation guarantees reentrancy. As the figure shows, in sub-buffer 1 of buffer 0 (on CPU 0), a first event was written into the buffer. A slot was reserved after it for writing, but the process was interrupted midway through by another write. We see that this nested write completes successfully, and doesn't affect the end result of the interrupted write, as its slot is already reserved.

The low-overhead requirements of the *LTTng* project have lead to the creation of *local atomic operations* in the Linux kernel [DesnoyersDesnoyers2016b], which aim to provide a lower performance cost than regular atomic operations, by leveraging the fact that some data is CPU-local (that is, data

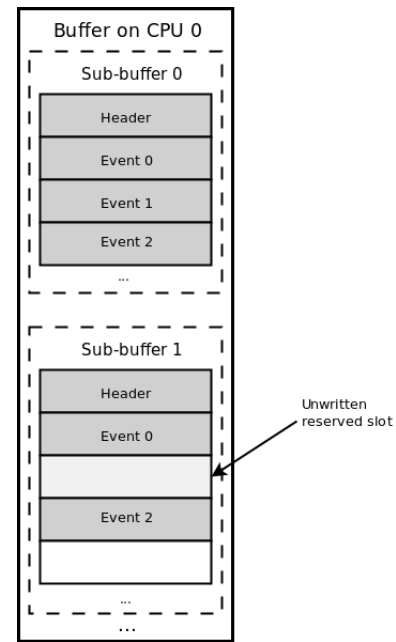


Fig. 2: Anatomy of *LTTng*'s sub-buffers

that won't be accessed by another CPU). When that is the case, atomic operations accessing local-only data don't require the use of the `LOCK` prefix (which locks the memory bus) or memory barriers to ensure read coherency between CPUs, which leads to *local atomic operations*. This performance improvement comes at the cost of a higher usage complexity, and care is needed when accessing the data from other CPUs due to weak ordering. It is worth mentioning that local atomic operations disable preemption around the call site to avoid CPU migration when accessing local variables from a process context (as opposed to interrupt context). Ftrace also makes use of local atomic operations.

To guarantee wait-free reads of tracing management variables (such as enabling/disabling tracing, filters to be applied, etc.), *LTTng* uses RCU<sup>3</sup> for the synchronization of concurrent accesses [McKenney and SlingwineMcKenney and Slingwine1998] [Desnoyers and DagenaisDesnoyers and Dagenais2010]. Since writes to these variables are rare, but reads are abundant and concurrent, RCU is ideal for such an access pattern since it avoids all waiting on the reader side.

Contrary to Ftrace, *LTTng* doesn't use pages as the finest-grained entity for ring buffer management. Instead, it uses per-CPU sub-buffers. The advantage of this approach is that the size of an event can be greater than the size of a page, and the performance hit of memory barriers is better amortized (barriers are only required at sub-buffer boundaries instead of page boundaries). Sub-buffer boundaries require memory barriers mainly to guarantee in-order memory writes to the sub-buffer data and its management variables. In other words, out-of-order memory accesses to sub-buffer data and the sub-buffer management variables can result in incoherent perception from the reader's side, where a sub-buffer can be

<sup>3</sup>Read-Copy-Update

flagged as ready to be read while the tracing data hasn't been propagated to memory yet. It is interesting to note that the impact of memory barriers is negligible on x86\_64. Since the architecture guarantees Total Store Ordering (TSO) between CPUs [Hennessy and PattersonHennessy and Patterson2011], write memory barriers do not actually create a fence but resolve to nops and simply avoid instructions reordering at compile-time. On architectures with other memory orderings such as ARM, the use of memory barriers at sub-buffer boundaries (page boundaries for Ftrace) has a larger impact on the overhead.

4) *LTTng-kprobe*: We introduce LTTng-kprobe as a standalone tracer for simplicity purposes. In reality, it is simply the LTTng tracer configured to use Kprobes instead of TRACE\_EVENT for probe callback. Other than the callback mechanism, there are no differences for the probes or the ring buffer, and the same internals and guarantees as LTTng are valid for this tracer. Comparing the benchmark results between LTTng-kprobe and LTTng will highlight the impact of using Kprobes instead of TRACE\_EVENT. This comparison can help developers assess the overhead of Kprobes-based tracers in their applications when low overhead is key.

5) *Perf*: Perf is a performance monitoring tool integrated into the mainline kernel. Its use case is typically different than that of Ftrace or LTTng. Perf is targeted for sampling and profiling applications, although it can interface with the tracepoint infrastructure within the kernel and record tracepoints (including system calls). Perf can also gather hardware-level PMU<sup>4</sup> such as different level of cache misses, TLB misses, CPU cycles, missed branch predictions, and so on. Contrary to Ftrace, Perf's monitoring scope is restricted to a single process. The events and counters reported by Perf are those which occurred within the context of the traced process and thus have been accounted for it. This property makes Perf more suited to analyze the behavior of a given program and can help answer practical questions such as cache-friendliness of the code, or the amount of time spent in each function.

The Perf tool itself is built on top of the kernel's perf\_events subsystem [GhodsGhods2016] which is the part that actually implements the tracing, profiling and sampling functionalities. Perf can also use Ftrace's architecture to hook to tracepoints, and trace similarly to Ftrace and LTTng. Perf\_events internally uses a ring buffer which can be mapped to userspace to be consumed. Through the perf\_event\_open() system call, a userspace process can obtain a file descriptor on the metric/counter it wants to measure. The file descriptor can then be mmap()'d and accessed from the userspace process's memory space.

6) *eBPF*: eBPF has evolved from the Berkeley Packet Filter [Schulist J.Schulist J.2016] [McCanne and JacobsonMcCanne and Jacobson1993] to a standalone monitoring tool included in the Linux kernel. It allows users to write programs (similarly to probes) which can be dynamically inserted at any location in the kernel using Kprobes (section IV-C). eBPF programs are compiled to bytecode which is interpreted and executed by the kernel

in a limited context to ensure security. The kernel also supports just-in-time compilation for sections of the generated bytecode [Sharma and DagenaisSharma and Dagenais2016b]. It is important to note that eBPF isn't a tracer, as it doesn't follow the *callback, serialize, write* scheme, but rather an aggregator. The eBPF interpreter provides data structures to its users, such as simple arrays and hashmaps, which makes it a great tool for aggregation, bookkeeping, and live monitoring. As eBPF provides arrays that can be shared from the kernel space to user space, a tracer-like behavior can be implemented. For the purpose of this study, we wrote a minimal eBPF program that samples the clock and writes a data structure holding a timestamp and a payload to an eBPF array. We hook this program to the same static tracepoint used for benchmarking other tracers using Kprobes. Although the data is never read, this program simulates the behavior of a tracer, making its benchmarking relevant for this study. As of version 4.7, the Linux kernel supports eBPF tracing, which hooks directly into the TRACEPOINT infrastructure. The same eBPF program was then ran on a 4.12 Linux kernel, with the only difference being a direct hook onto kernel tracepoints instead of going through Kprobes.

7) *SystemTap*: SystemTap is similar to eBPF (section V-A6) as it provides a simple language to write probes for aggregation and live monitoring [Prasad, Cohen, Eigler, Hunt, Keniston, and ChenPrasad et al.2005]. SystemTap provides an easy scripting language for users to create custom probes to monitor their systems. The users can provide kernel or userspace symbols to hook on. SystemTap programs are then translated to C code and compiled into a loadable kernel module [EiglerEigler2006]. Once loaded, the module dynamically inserts the probe into the kernel's code using Kprobes. Similarly to eBPF, since SystemTap follows the *callback, compute, update* pattern, it is, in fact, an aggregator rather than a tracer. However, for the purposes of our work, a tracing behavior can be simulated by implementing a probe that samples the time and writes the value along with a constant payload to an internal array.

8) *Strace*: Strace is a tool for system calls tracing [KerriskKerrisk2010] [Johnson and TroanJohnson and Troan2004]. Using Ptrace (section IV-C), it hooks into a process and intercepts all its system calls, along with their arguments. The result is written to a file descriptor for later analysis. Due to the heavy trap mechanism, along with the scheduling costs (as multiple processes are involved, the monitored and the monitoring), Strace typically adds a large overhead which doesn't usually suit production environments. Other tracers, such as LTTng and Ftrace, can provide the same information using different tracing mechanisms.

9) *SysDig*: SysDig [Selij and van den HaakSelij and van den Haak201] is a modern commercial tool used for monitoring systems. It covers a wide range of applications, from containers, to web services, etc. SysDig also allows users to write Chisels [chi2017], which are lua scripts capable of custom analysis similar to eBPF and SystemTap. SysDig leverages the TRACEPOINT infrastructure in the kernel, which was introduced by the LTTng project, to hook onto available

<sup>4</sup>Performance Monitoring Units



tracepoints and provide system and application monitoring. Upon loading, the SysDig kernel modules register to context switch and system calls events. The work presented in this article does not cover SysDig for two main reasons. Primarily, SysDig doesn't allow the users to control the kernel tracing part of the product, and this aspect is only used to gather a few metrics about the monitored applications, such as reads, writes and CPU usage. Moreover, SysDig doesn't follow the *callback, serialize, write* pattern, but rather uses tracepoints for on-the-fly metric gathering. Finally, as SysDig uses the same underlying infrastructure as LTTng and Ftrace, the overhead of SysDig is implicitly covered when studying the TRACEPOINT infrastructure.

## B. Userspace tracers

1) *Printf*: `printf()` is a rudimentary form of tracing, but is the easiest to use. It uses a string as input as well as some variables for pretty-printing. `printf()` uses an internal buffer to store the data before it is flushed to a file descriptor. Thus, we can implement a basic tracer using `printf()` by sampling the time and printing it along with a payload into `printf()`'s internal buffer. This satisfies the definition of a tracer given in section III-A.

Glibc's implementation of `printf()` is thread-safe [PeekPeek1996], although multiple threads within the same application share the same global buffer for the output stream. `printf()` uses an internal lock to protect the output buffer and avoid corruption on contention. However, reentrancy is not guaranteed, and calling `printf()` from different contexts (such as from a signal handler that interrupts `printf()`) might have unexpected results.

2) *LTTng UST*: LTTng UST (UserSpace Tracer) [Blunck, Desnoyers, and FournierBlunck et al.2009] is a port to userspace of LTTng (section V-A3). Although they are independent, both tracers share the same design, using a ring buffer to store the trace [Fournier, Desnoyers, and DagenaisFournier et al.2009], RCU mechanism for data structure synchronization and atomic operations for reentrancy. To this effect, the RCU mechanism was ported to userspace as well, creating the URCU (Userspace Read-Copy-Update) project [Desnoyers, McKenney, Stern, Dagenais, and WalpoleDesnoyers et al.2012]. All programs to be traced with LTTng UST should be linked against the library, as well as `libust` (the tracing library). Similarly to its kernel counterpart, LTTng UST was designed and implemented to perform tracing at low cost, while guaranteeing reentrancy, interrupt-safety, signal-safety, etc. No tracing data is copied and no system calls are invoked by the tracer, removing the typical sources of overhead.

3) *LTTng using tracef()*: Since part of `printf()`'s latency is split between pretty printing the input and storing to an internal buffer, we added LTTng's `tracef()` function for a more equitable/fair comparison. `Tracef()` is a function that combines, from the developer's point of view, LTTng's tracer and `printf()`. Instead of implementing actual tracepoints that can be called within the code, `tracef()` generates a `ltnng_ust_tracef` event which holds as payload a

pretty-printed string similarly to `printf()`. In that manner, LTTng's internal ring buffer mechanism is used, while including the cost of pretty printing for a more equitable comparison with `printf()`. In that way, we compare more accurately the actual serialization to a ring buffer between LTTng and `printf()`.

Since `tracef()` uses LTTng's internals and only affects the serialization part, the same guarantees as tracing using regular tracepoints apply, such as reentrancy, thread-safety, interrupt-safety and so on.

4) *Extrac*: Extrac is a tracer developed at the Barcelona Supercomputing Center (BSC) as part of a tracing and analysis infrastructure for high-performance computing. It has complementary software such as Paraver [Pillet, Labarta, Cortes, and GironaPillet et al.1995] for trace visualization and analysis. In this paper, we focus exclusively on the tracing part. The tracer supports many mechanisms for both static and dynamic instrumentation, as shown in Table I. It also uses `LD_PRELOAD` to intercept MPI library calls at runtime and instrument them. In addition, Extrac supports sampling hardware counters through the PAPI interface [Terpstra, Jagode, You, and DongarraTerpstra et al.2010], and other features which aren't covered in this study as they are beyond the scope of tracing [??ext2016]. Extrac also uses internal per-thread buffers to store the data on the fly. Static tracepoints have limited features: only one type of tracepoint exists which is called `Extrac_event`, having two fields as a payload. The fields are a pair of integers, the first one being a number representing the type of event that occurred, and the other flags either a function entry or exit. This provides less flexibility to the user to create and use custom tracepoints with variable payloads.

5) *VampirTrace*: VampirTrace is a high-performance computing tracer [Knüpfer, Brunst, Doleschal, Jurenz, Lieber, Mickler, Müller, and NagelKnüpfer et al.2010], [Müller, Knüpfer, Jurenz, Lieber, Brunst, Mix, and NagelMüller et al.2009], [Schöne, Tschüter, Ilsche, and HackenbergSchöne et al.2010], with the ability to interface with large-scale computing frameworks such as MPI. Similarly to Extrac, it offers many ways to instrument applications, either statically or dynamically. For static tracing, VampirTrace uses a different approach and relies on the developer to define the sections of code to be analyzed. When tracing of a section is disabled through a delimiter, an event is generated and written to the buffer. VampirTrace uses per-thread buffers to avoid synchronization and maintain scalability. The microbenchmark of this paper for VampirTrace consists of a tight loop that enables and disables tracing. This implementation defines an empty section to trace, but generates an event on each loop, due to disabling tracing on each loop so that a behavior similar to other tracers is achieved.

6) *Lightweight UST*: Lightweight UST (LW-ust) is an in-house minimal tracer built for the purpose of this study. It targets the fastest possible naive tracing implementation, at the cost of actual usability, reentrancy, and thread-safeness. The purpose of this tracer is to show a baseline of how fast tracing can be achieved, which is a simple clock read and a copy into an internal buffer. LW-ust uses per-CPU circular buffers to

write data, which are cache line-aligned to avoid false sharing. The data copied into the buffer is simply an integer referring to the tracepoint type, a timestamp (using `clock_gettime()`, similarly to other tracers) as well as another integer as the tracepoint payload. No effort is made towards reentrancy or thread-safety, and thus data integrity is not guaranteed. Although LW-ust can not be used on production systems, it is still interesting to benchmark it as it provides a lower bound on the impact of a tracer.

## VI. BENCHMARKING

### A. Objectives

The objective of this study is to quantify the overhead added by tracing, taking into consideration different callback mechanisms, internal tracers' architectures, and the guarantees they provide. We focus on the cost of individual tracepoints, rather than analyzing the impact on more general workloads. By analyzing the results, we hope to help the users configure the tracers to adapt them for their specific use cases when low overhead is a requirement in production systems.

### B. Approach

1) *Overview:* We start by explaining our benchmark related to kernel tracing as well as the metrics we measured. Our objective is to measure the cost of tracing an event using the different tracers. To be able to get this metric with the best possible precision, two conditions have to be met: first, all tracers must be running in the overwrite mode. That way, we guarantee that all tracepoints are executed, even when the internal ring buffers are full, as opposed to benchmarking a test and nop operation (testing the buffer size and jumping over the tracepoint if the buffer is full). Secondly, when possible, we launch tracers in a producer-only mode. This requirement helps to cancel outside noise that might interfere with the benchmarking. In our case, the consumer process might preempt the traced process, or the action of writing to the disk or the network might also interfere in some way with the producer. Although we are only measuring isolated tracepoints, and preemption doesn't directly affect the duration of executing a tracepoint, it still, however, impacts the duration of a tracepoint since it might invalidate caches (including the TLB for which misses are particularly expensive).

The actual measurement is done in a microbenchmark which is simply a tight loop executing a tracepoint call. The payload of the event is 4 bytes. The microbenchmark is implemented in a kernel module to avoid the overhead of switching between user and kernel spaces. A single `ioctl()` call to the module (via the `sysfs` pseudo-filesystem) triggers the benchmark and blocks until its completion.

The time is measured prior and after the call at the nanosecond granularity. We point out that there is no system call for reading the time since the benchmark is already running in kernel space (although some clock access functions can be used in user space to avoid a system call, such as the monotonic clock). Hardware performance counters are also read before and after the tracepoint, allowing us to track specific metrics that give insight into each kernel's implementation.

The benefit of sampling many hardware counters in a single run, instead of sampling a single counter across many runs (to reduce the impact of these samplings) is that it is possible a posteriori to make correlations between many metrics for a single recorded event. For instance, we can verify if the slow path (given by sampling the number of instructions) can be further optimized by reducing cache misses (given by sampling the number of cache misses). Due to hardware limitations, only four hardware counters can be sampled at a time. Thus, two runs of the benchmark are executed, each sampling different counters. Furthermore, in order to avoid interruption while the tracepoint is executing, which would interfere with the measurements, we disable interrupts before each call and enable them after. Since this approach doesn't disable NMIs (which are, by definition, non-maskable), we read the NMI count prior to and after each tracepoint call, discarding the result values in case an NMI is detected. Algorithm 1 shows the pseudo-code of the benchmark. The results shown in this paper are gathered by running the tight loop 5000 times in a steady state for each benchmark run. For scalability tests, the tight loop iterates 5000 times per core (in parallel) in a steady state for each benchmark run.

---

#### Algorithm 1: Pseudo-code of the kernel tracer benchmark

---

**Input:** *numberOfLoops*

**Output:** *ArrayOfResults*

```

payload = (int32_t) 0;
results = allocateArray();
i = 0;
while i != numberOfLoops do
  disableInterrupts();
  numberNMI = getNumberNMI();
  readPMU(metric1_start);
  readPMU(metric2_start);
  readPMU(metric3_start);
  readPMU(metric4_start);
  start_time = gettimeofday();
  tracepoint(payload);
  end_time = gettimeofday();
  readPMU(metric1_end);
  readPMU(metric2_end);
  readPMU(metric3_end);
  readPMU(metric4_end);
  enableInterrupts();
  if numberNMI == getNumberNMI() then
    results[i].duration = end_time - start_time;
    results[i].metric1 = metric1_end - metric1_start;
    results[i].metric2 = metric2_end - metric2_start;
    results[i].metric3 = metric3_end - metric3_start;
    results[i].metric4 = metric4_end - metric4_start;
    i++;
  end
end
return results;

```

---

The final `results` array is then dumped in a CSV (Comma-Separated Value) file, which is used to generate the graphs shown in section VII. The same algorithm is used for both kernel and user space.

Benchmarking system calls and analyzing the results require a different methodology than regular tracepoints. The added cost of the actual system call, as well as the fact that the Linux

TABLE I: Callback mechanisms for kernel and userspace tracers

'X' indicates that a callback mechanism is supported by the tracer and is covered in this work

'O' indicates that the mechanism is supported by the tracer but wasn't benchmarked in this work

Mechanism	Kernel space				User space				
	Static		Dynamic		Static		Dynamic		
	Static tracepoint	Function tracing	Trap	Trampoline	Static tracepoint	Function tracing	Trampoline	Uprobe	Ptrace
Implementation	TRACE_EVENT	Compiler	Kprobe	Optimized kprobe	Function call	Compiler	DynInst	Uprobe	Ptrace
LTTng	X		X	X	X	O		O	
Ftrace	X	O	O	O					
Perf	X		O	O					
eBPF	O		X	X					
SystemTap			X	X			O	X	
Strace									X
Printf					X				
Extrace					X		O		
LW-ust					X				

kernel automatically instruments system call entry and exit, make getting fine-grained metrics about the tracing mechanism only more difficult. Since actual work is done by the system call, PMU counters values aren't representative of the tracing mechanism solely, but also include the system call's execution. In addition, the CPU transitions from user mode to kernel mode through a trap that represents the actual kernel invocation. This is a relatively costly process, since it requires saving the state of the stack, registers, and switching CPU rings [Tanenbaum and BosTanenbaum and Bos2014]. This procedure is therefore implicitly accounted for the tracepoint in the benchmark. The analysis and benchmarking of system calls are thus done at a higher level, with coarse-grained results. It is worth mentioning that system calls tracing is a subset of kernel tracing, which we presented in the previous section. In other words, writing to the ring buffer is done in the same fashion, and only serializing the payload differs from a regular tracepoint.

In order to reduce as much as possible the duration of the system call's actual work and get samples that are as close as possible to the cost of instrumentation solely, we benchmarked an empty `ioctl()` system call. For the purpose of this work, we wrote a kernel module that exports an entry in `procfs`. We implemented its `ioctl()` function as an empty function, such that calling `ioctl()` on the file descriptor returns as quickly as possible, and ends up being only the tracing instructions and a CPU context switch to kernel space and back.

2) *Test setup*: We ran the benchmarks on an Intel i7-3770 CPU running at a 3.40 GHz frequency with 16 GB of RAM, running a 4.5.0 vanilla Linux kernel, and a 4.7.0 Linux kernel for eBPF tracing (as it is the first version that introduced built-in eBPF tracing). We built manually the following software: LTTng from its 2.9-rc1 branch, SystemTap from its master branch at version 3.0/0.163, and Extrace at version 3.3.0, and VampirTrace at version 5.14.4.

3) *Configuration*: Benchmarking at such a low granularity proves to be a tedious task relatively to higher level

benchmarks [Kalibera and JonesKalibera and Jones2013]. Optimizations, both at the operating system level as well as the hardware level, can interfere greatly with the results. Not only are the results biased, but reproducibility of the results is also affected. For instance, dynamically adjustable CPU frequency has a direct impact on the latency of a tracepoint. In order to have the most control over the environment, we disabled the following features at the hardware level: hyperthreading, C-states, dynamic frequency and Turbo Boost.

Furthermore, we configured the tracers to have as much of the same behavior as possible to obtain relatively fair results. For instance, as shown in section V-A2, Ftrace uses a CPU-local clock source to sample the time, whereas LTTng uses a global clock. At the nanosecond scale, as we are benchmarking, this behavioral difference can have a major impact on the results. We configured Ftrace to use the global clock so that the same time sampling mechanism is used by all tracers. We configured LTTng to use 2 sub-buffers per CPU buffer, each of 8 KB, and Ftrace to use a 16 KB buffer per CPU (same ring buffer size). We also configured LTTng to trace in flight recorder mode, in which the contents of the ring buffer are never flushed to disk, and new tracing data overwrites the older one. This guarantees that flushing to disk doesn't interfere with the benchmark, which is also the behavior of Ftrace.

4) *Steady state*: Another important factor that impacts the accuracy of the results is benchmarking in a steady state. When tracing is first enabled, the tracer typically allocates memory for the internal ring buffer. Running in the steady state means that all code paths have been covered at least once (and hence are in some cache level), all memory locations have been touched at least once, virtual addresses have been resolved and stored in the TLB, and all other initialization routines have been covered. For instance, when memory is allocated, it typically isn't physically reserved until it is actually touched by the owner process. This mechanism happens by the means of a page fault, which implies a frame allocation by the operating system. The overhead of this procedure is accounted to the

process, and particularly during the call to `tracepoint()` (or its equivalent), which might be misleading and not representative of the actual cost of the tracepoint. Note that this particular example doesn't apply for LTTng and Ftrace as they make sure memory frame allocation is done at allocation time, avoiding postponed page faults and lazy TLB filling. Thus, running in a transient state may result in biased numbers for the `tracepoint()` call, since a costly TLB miss due to an initial virtual address resolution will get accounted to the tracepoint. On the other hand, earlier runs have shown that Perf varies greatly between the transient state and the steady state, due to internal initialization being done on the critical path of tracepoints. In the remainder of this paper, all benchmarks results are done in a steady state.

The steady state is reached once the amount of trace data is equal to the size of the tracers' internal buffers to guarantee that all locations have been accessed at least once. We make sure we that the entire ring buffer has been filled at least once before results are recorded. Since tracers are configured to overwrite mode, all calls to tracepoint should happen in optimal conditions, with all memory allocated, hot code in the cache, and data and instruction addresses in the TLB.

5) *Metrics*: Table II shows the collected metrics for each executed tracepoint. The raw data generated by the benchmarks contains the value of these metrics for each recorded event. These values are sampled using the `perf_events` infrastructure from the kernel space, for kernel tracing, and the `perf_event_open()` system call, for user space tracing.

TABLE II: Metrics

Metric	Meaning
Latency	Time to execute a tracepoint in nanoseconds
L1 misses	Number of first level cache misses
Cache misses	Number of all levels cache misses
CPU cycles	Number of CPU cycles
Instructions	Number of instructions
TLB misses	Number of TLB misses
Bus cycles	Number of memory bus cycles
Branch misses	Number of mispredicted branch instructions
Branch instructions	Number of branch instructions

Some of these metrics are more significant than others, as the results section will show. We start by explaining the runtime penalty of some of these metrics.

A TLB (Translation lookaside buffer) miss happens when a virtual address to physical address mapping isn't stored in the TLB. The penalty is the page walk process, which needs to perform the translation from virtual to physical addresses. Depending on the architecture, this process might be more or less costly. On `x86_64`, a page walk results in 4 memory accesses (one for each level of the page tables) which is a major performance setback, given the low frequency of the memory bus relative to the frequency of modern CPUs.

A branch miss indicates a branch misprediction from the compiler or the hardware itself. The compiler can optimize the fast path of an application by predicting branches based on hints given by the developer. The CPU can also predict branches at run-time, depending on the path that is most often taken. Instruction and data caches, as well as overall

pipeline efficiency, are optimized for the fast path by fetching instructions prematurely, before the branching instruction has been evaluated. The penalty of a branch miss is having to load the instructions that were assumed unused, as well as any data related to them. This fetching might occur from either main memory or from higher levels of the cache (L2 or L3). Consequently, this process reduces the efficiency of the pipeline as it stalls while the data and instructions are fetched.

## VII. RESULTS

### A. Kernel tracing

We now present the results of the benchmark for the kernel tracers. Unless stated otherwise, all results are gathered from single-core benchmark runs. We start by showing in Table III an overview of the results which we will try to explain using the more detailed graphs. The standard deviation is provided when the results indicate an average value. The 90th percentile helps filtering out corner cases where tracepoint latency is exceptionally high.

TABLE III: Average latency of a kernel space tracepoint

	Average (ns)	STD	90th percentile
None	17	0	17
LTTng	92	5	89
Ftrace	116	7	114
Perf	121	1	118
LTTng-Kprobe	123	5	121
eBPF-tracing	128	5	124
SystemTap	130	5	125
eBPF	140	10	126

In the graphs that follow, we show, for each run, the value of different metrics for each tracepoint call. By showing the results of all tracepoint calls for each tracer, it is much easier to find trends in the usage of different resources. Furthermore, this approach helps explaining corner cases, where a given set of circumstances influences the cost of a single tracepoint and helps setting an upper bound for the cost of an individual tracepoint, once the steady state is reached.

As explained in section VI-B, we sample performance counters at the beginning and end of each tracepoint invocation in order to precisely measure different metrics, such as the number of instructions executed, the number of cycles, or the number of cache misses. The number of instructions is directly related to the code logic and should not bring any surprise. Furthermore, although the number of instructions as an absolute value is not significant in isolation (compiler versions, options, etc. directly influence this value), it still is an interesting metric to capture, as it uncovers the different code paths that are taken by the program at runtime. Other metrics such as the number of cycles and the number of cache misses are interesting to show any erratic behavior. For instance, a larger number of cycles per instruction<sup>5</sup> (much

<sup>5</sup>This value depends on the architecture of the CPU. Superscalar processors can have multiple pipelines to achieve parallelism within a single processor. When that is the case, we can see a number of cycles per instructions that is lower than 1 (multiple instructions per cycle per processor). This paper focuses on multi-core scalar processors, which typically have one pipeline per processor.

higher than 1) might be due to contention over the memory bus, or to inefficient instructions. A lower number of cycles per instruction (closer to 1) could be a good sign, but it can often hint at potential optimizations, such as doing calculations in advance and caching them for future use. Added latency might be caused by an abnormally high number of page faults, and might help discover hard-to-detect and unexpected issues. On the other hand, these low-level metrics can confirm that everything is running smoothly, for instance, when most of the tracepoints seem to indicate a ratio of instructions per cycle close to 1.

For all of the following graphs, the y-axis always shows the tracepoint latencies whereas the x-axis shows the values of different counters.

Looking at Figure 3 a) and b), it is interesting to note the lack of direct effect of the cache misses on the tracepoint latency. Contrary to what would have been expected, more L1 misses do not impact in any significant manner the cost of a tracepoint. We can make this assessment since the tracepoints that triggered 0 L1 misses range widely in the latency spectrum. Furthermore, tracepoints with higher L1 misses record similar latency as the ones with no misses. Of course, cache misses do affect latency, but their impact is diluted by other factors. We also notice that benchmarking in a steady state helps to keep the number of cache misses to a minimum, as shown in b).

On the other hand, CPU cycles have a direct linear relation with the latency of a tracepoint. However, this is more a consequence than a cause, since tracepoints are uninterruptible in our setup (interrupts disabled and the ones interrupted by NMIs are ignored) and the CPU frequency is maintained at a maximum, it is only natural that the more costly tracepoints require more CPU cycles. The interesting thing to note is that the nearly-perfect linear relation between CPU cycles and latency doesn't exist with dynamic CPU frequency enabled. Since the CPU frequency can change dynamically, a high latency tracepoint can actually record a low number of cycles, changing the trend between latency and CPU cycles. A few outliers exist, potentially caused by an imprecision by sampling the pipeline to extract CPU cycles values [Wilhelm, Grund, Reineke, Schlickling, Pister, and Ferdinand-Wilhelm, 2009].

Figure 3 d) shows a metric that directly impacts latency: the number of instructions. As the number of instructions grows, the groups of samples are higher on the latency axis. As we might have expected, more instructions per tracepoint usually implies more time to complete and thus record a higher latency. With that said, the graph suggests that other factors impact the latency, as even for tracepoints that require the same number of instructions, their distribution on the latency spectrum is quite wide (tracepoints recorded using Ftrace that required 1052 instructions range between 110ns and 225ns).

Another interesting point that we can take from Figure 3 d) is the code path for each tracer. Looking at tracepoints sampled for LTTng, we can easily guess the three internal code paths of the tracer: the samples are grouped into three possible number of instructions: 870 instructions, 927 instructions, or 1638 instructions (for this particular build of LTTng and kernel). As explained in section V-A3, LTTng uses internal

sub-buffers with their size being a multiple of a page for its ring buffer. We might guess that these three code paths shown in the graph represent the tracepoints that cross boundaries: the tracepoints requiring 870 instructions to complete are the most frequent ones and execute the fast path. Storing the tracepoint into a sub-buffer is straightforward and translates to a simple `memcpy()`. The middle path, requiring 927 instructions, is covered when storing a tracepoint is still within the same sub-buffer but crosses page boundaries. The memory area backing a sub-buffer is manually managed by LTTng, and thus pages that make up a sub-buffer aren't contiguous: LTTng doesn't use virtual addresses but rather uses physical memory frames, which requires page stitching when data needs to be stored across (or read from) more than a single page. Finally, the slow path, requiring 1638 instructions, is covered by the tracepoints that cross sub-buffer boundaries (and implicitly page boundaries, as sub-buffers are page-aligned). LTTng then requires internal bookkeeping, such as writing some header data into the sub-buffer, which adds instructions to the critical path and further latency. Thus, if tracepoint latency is an issue, avoiding the slow path is possible by allocating larger sub-buffers and reducing the frequency of the slow path (although this might lead to other problems, such as events loss). With this information, it is possible to predict how many instructions an event might require, depending on the empty space left in the sub-buffer at the moment the tracepoint is hit. In other words, the middle and slow paths are taken at regular intervals. The middle path is taken every  $(\text{PAGE\_SIZE} / \text{EVENT\_SIZE})$  events, and the slow path is taken every  $(\text{SUB-BUFFER\_SIZE} / \text{EVENT\_SIZE})$  events, assuming all events are the same size.

This analysis can help the users choose the right buffer size to configure their tracer. In order to reduce the average tracepoint latency, the slow path should be avoided as much as possible. This is possible by setting larger sub-buffers. On the other hand, since sub-buffers can only be flushed once they are full, having larger sub-buffers usually implies a higher probability of event loss. Flushing large buffers requires more time, and if the events are recorded at a high rate, the ring buffer has time to loop and start overwriting unread sub-buffers. For a fixed buffer size, a trade-off has to be made between the number of sub-buffers and their size. Smaller sub-buffers reduce the risk of lost events, but larger sub-buffers result in faster tracepoints on average. We conclude this by mentioning that comparing the number of instructions per event between tracers isn't necessarily a relevant metric, as the number of cycles an instruction requires may vary greatly and isn't a direct indicator of latency (unless the numbers of instructions differ greatly). However, it is still interesting to analyze the number of instructions for the same tracer to deconstruct different code paths that have been taken, and get a deeper understanding of a tracer's internals.

A similar observation can be made for Ftrace. However, we can only group the samples as executing one of two paths: 1052 instructions or 1341 instructions, which we reference respectively as the fast path and the slow path. The reason Ftrace samples only show two code paths, instead of 3 like

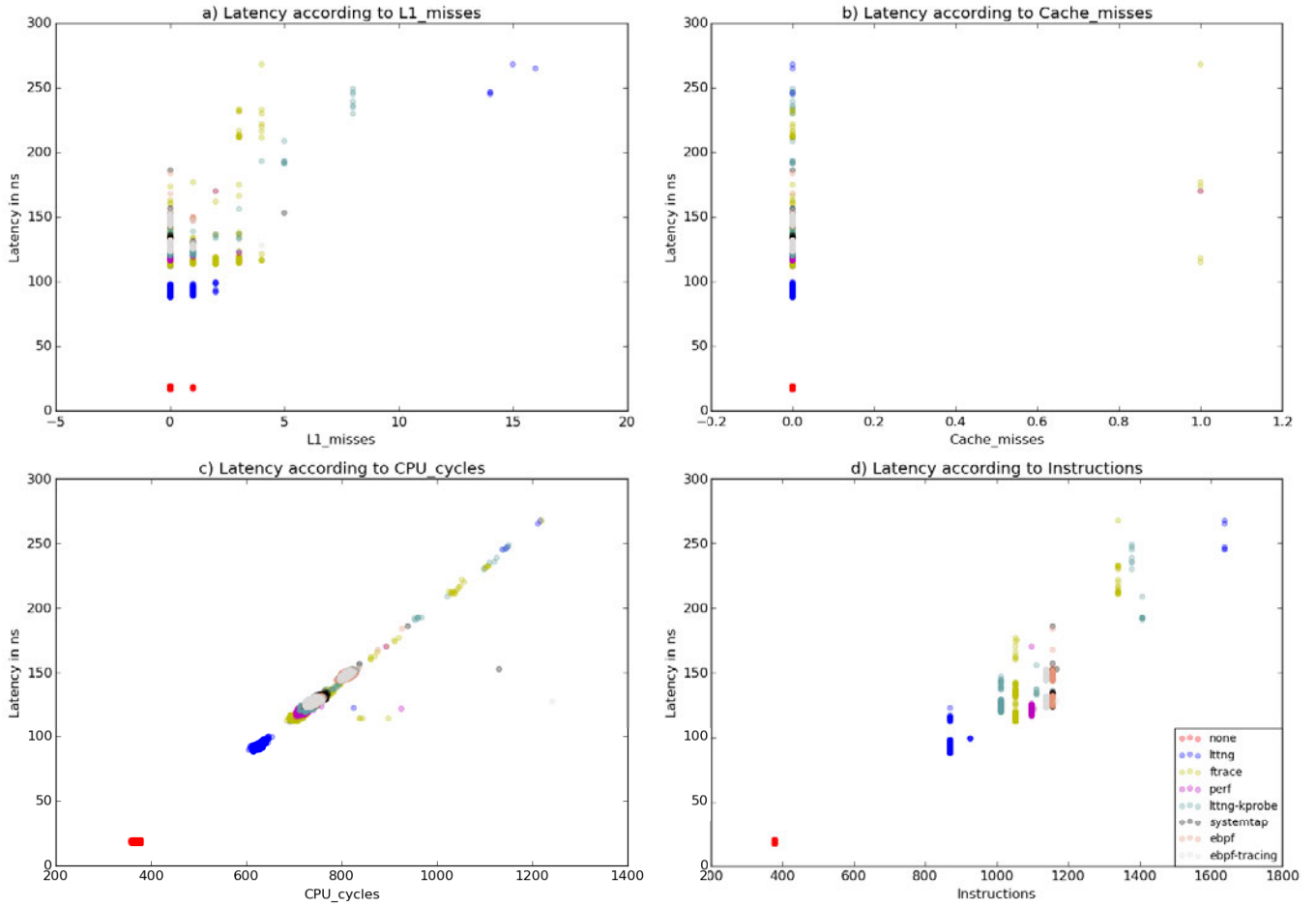


Fig. 3: All tracepoints latencies in kernel space against different hardware counters (Part 1)

LTTng, is the fact that Ftrace doesn't use the notion of sub-buffers and only manipulates pages, albeit manually similarly to LTTng. Thus, a tracepoint recorded with Ftrace can only cover one of two cases: it either fits directly into a page, or it crosses page boundaries. This behavior can be achieved with LTTng if the size of the sub-buffers is set to the page size.

Figure 4 a), indicating the number of bus cycles per tracepoint, shows a result similar to Figure 3 c), which is naturally expected. Figure 4 b) shows that all tracepoints for all tracers are recorded without triggering any TLB misses. This is due to the fact that the events are sampled in the steady state and all TLB misses have already gradually been fulfilled. This is also to be expected as there are no outliers in the samples, since a TLB miss is costly relatively to the average tracepoint latency (memory accesses are orders of magnitude slower than cache accesses [Hennessy and PattersonHennessy and Patterson2011]).

Finally, Figure 4 c) showing the branch misses can support the theory about the slow and fast paths we discussed for Ftrace and LTTng. The binary is typically optimized for the fast path by the compiler (and by the pipeline at run-time), and thus should trigger no branch misses [SmithSmith1981]. Figure 4 c) validates this theory as most of the samples have 0 branch misses. When the tracepoint data to be written into

the ring buffer crosses a page boundary, a branch miss should occur when the remaining free size in the current page of the buffer is tested against the size of the event. This process explains sample where 1 branch miss happens. Looking at the raw data, we can confirm that all cases going through the slow (and middle) paths for Ftrace and LTTng trigger exactly 1 branch miss. The reciprocal of this hypothesis is also valid: all samples that trigger at least one branch miss are executing either the middle or slow path.

The difference in tracepoint latency between LTTng and LTTng-kprobe highlights the impact of the callback mechanism used by the tracer. We see that using Kprobes increases the frequency of L1 misses (Figure 3 a)) as well as the number of instructions per tracepoint (Figure 3 d)) which contribute to a higher overall latency. Table III shows that changing the callback mechanism from `TRACE_EVENT` to `Kprobes` adds around 30 ns of overhead.

Figure 5 shows how the kernel tracers scale as the number of cores involved in the tracing effort grows. Notice that the latency axis is in logarithmic scale. Ftrace, eBPF and SystemTap show poor scalability, while LTTng and Perf scale with almost no added overhead. Such a significant performance impact on parallel systems usually indicates the use of an internal lock. It is indeed the case for Ftrace, as the global clock is protected

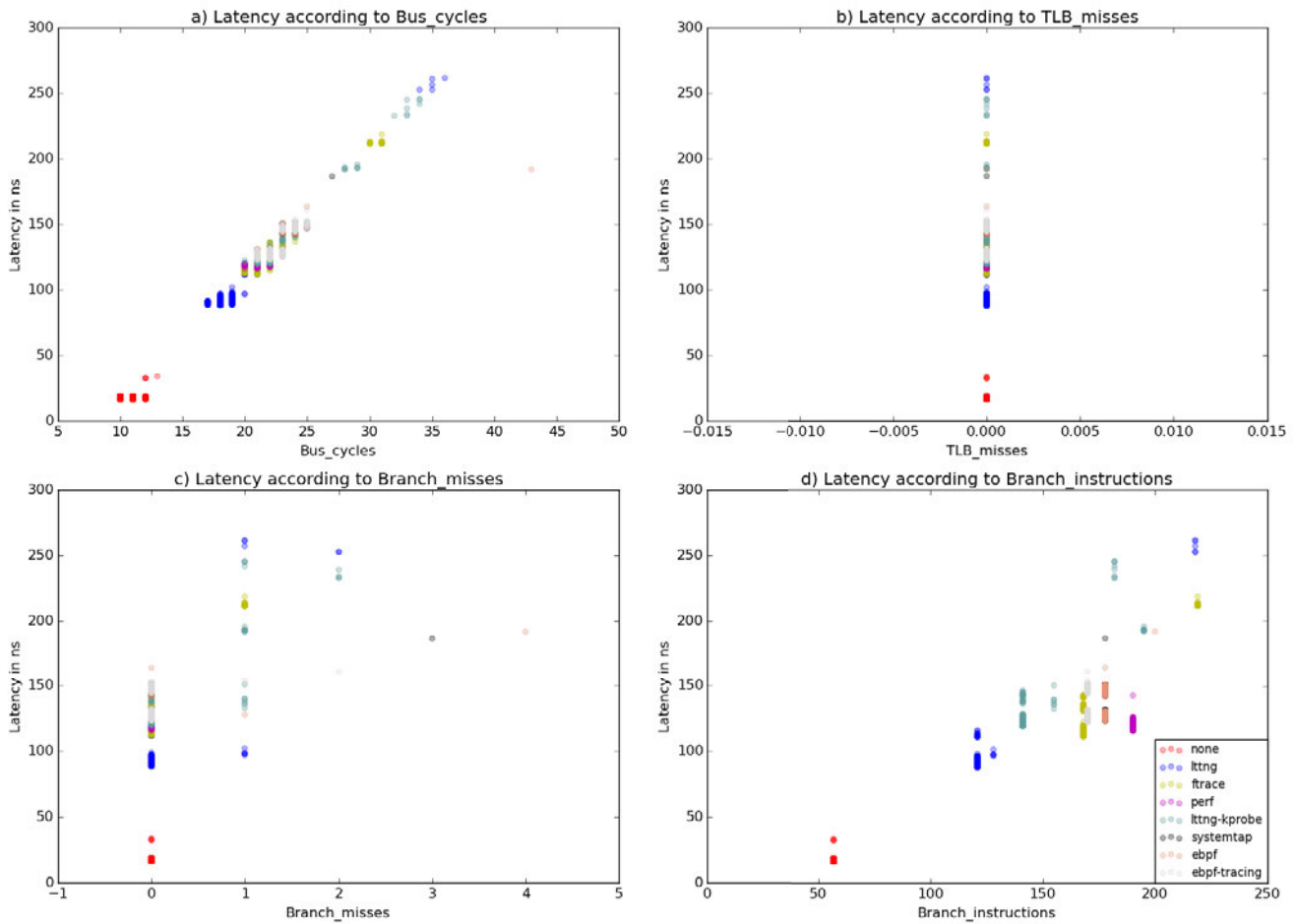


Fig. 4: All tracepoints latencies in kernel space against different hardware counters (Part 2)

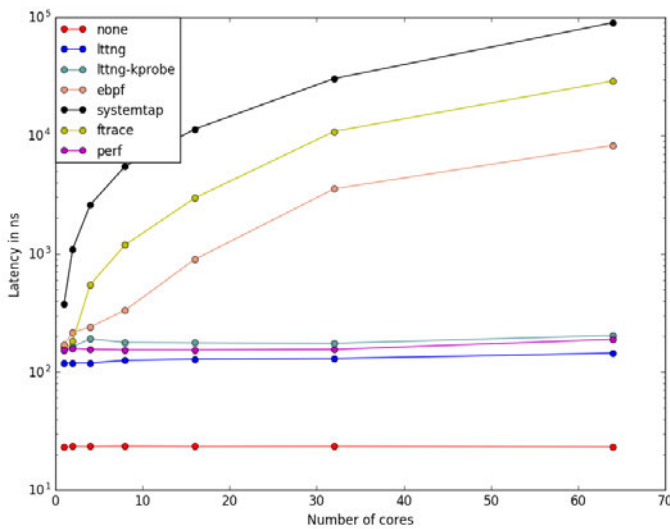


Fig. 5: Average tracepoint latency in kernel space against the number of cores (Log scale)

using a spin lock [LoveLove2005]. When configured to use a global clock, Ftrace internally manages a data structure used as a clock source to timestamp all events. This data structure

simply holds the timestamp of the last clock value at the last timestamp, so that consecutive clock reads perceive the time as strictly monotonically increasing. However, as this clock is global to the system and shared amongst CPUs, proper synchronization is required to avoid concurrent writes. The following code, taken from the kernel source tree, shows a snippet of the global clock read function in Ftrace. In the following snippet, function `sched_clock_cpu()` reads the local CPU clock.

```
u64 notrace trace_clock_global(void)
{
    unsigned long flags;
    int this_cpu;
    u64 now;
    ...
}
```

```
this_cpu = raw_smp_processor_id();
now = sched_clock_cpu(this_cpu);
/*
 * If in an NMI context then dont risk lockups
 * and return the cpu_clock() time:
 */
if (unlikely(in_nmi()))
```

```

goto out;

arch_spin_lock(&trace_clock_struct.lock);

if ((s64)(now - trace_clock_struct.prev_t
now = trace_clock_struct.prev_time +

trace_clock_struct.prev_time = now;

arch_spin_unlock(&trace_clock_struct.lock

...

out:
return now;
}

```

This implementation choice might be eligible for optimization as the global clock is, in fact, most often used for highly parallel workloads, making the use of an internal lock counter-intuitive. On the other hand, our micro-benchmark implements an extreme case, where a CPU has a hold of the lock at any given time, which might not be the case in real-life workloads unless the event rate is unusually high. Once again, discretion is given to the user on the choice of the clock. Figure 6 shows that when choosing the CPU local clock, the scalability of Ftrace is greatly improved and the added overhead is acceptable (*ftrace-local*). The Figure also shows the scalability of a global counter as a clock source, which doesn't require a lock but rather uses atomic operations to increment the variable at each clock read. The use of a global counter also shows somewhat poor scalability and makes its choice less than ideal for parallel workloads with a high event density. It is worth noting that atomic operations use the `LOCK` prefix to lock the memory bus, which explains the poor scalability.

Although great care was taken to use per-CPU arrays for eBPF and SystemTap programs, both tracers show poor scalability and the results suggest that they use internal locks either to access per-CPU data or the system clock.

### B. Userspace tracing

As with kernel tracing, we start by showing an overview of the results for userspace tracing in Table IV.

TABLE IV: Average latency of a userspace tracepoint

	Average (ns)	STD	90th percentile
None	17	0	17
LW-UST	66	2	64
VampirTrace	84	87	72
Extrae	120	7	118
LTTng-UST	158	27	150
Printf	250	65	242
LTTng-printf	446	73	433
SystemTap	1039	52	1020

Figure 7 shows the same metrics as the ones for kernel tracing. For simplicity purposes, as the results are more scarce

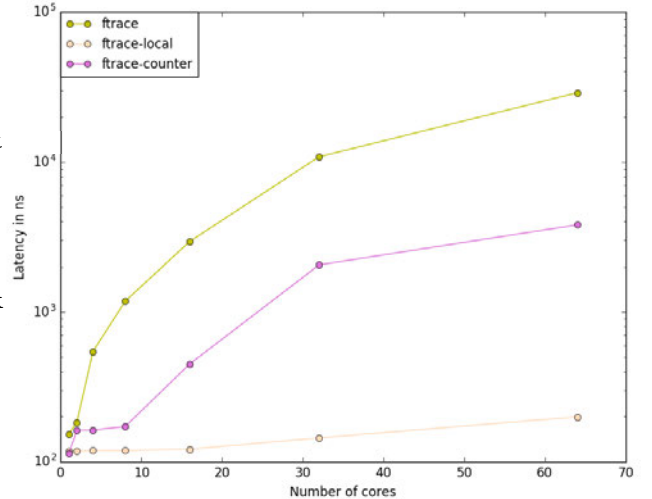


Fig. 6: Average tracepoint latency in kernel space for Ftrace using different clock sources against the number of cores (Log scale)

than for kernel space, the Figure only shows 1000 sample for each tracer, instead of 5000 for the previous Figures 3 and 4. The user space results can be interpreted in the same manner as for kernel space: L1 misses aren't a major factor of tracepoint latency since events that caused between 0 and 20 L1 miss range in the same latency spectrum. However, one of the main differences with kernel tracing is that the order of magnitude of the metrics is much higher. Looking at the number of instructions of individual tracepoints in Figure 7 d), we notice that LTTng has two main code path (for the regular and the `printf()` variants). We've explained in section V-A3 that LTTng uses internal sub-buffers to write tracing data. The slow path in the Figure indicates events write that cross sub-buffer boundaries. Contrary to kernel space, LTTng UST doesn't handle individual physical frames of memory that make up the sub-buffer, which explains why the middle path seen for the kernel tracer that covers page boundaries doesn't exist in user space. Extrae, the high-performance computing tracer, only seems to have a single code path in the steady state for all events. As explained earlier, this performance gain comes at the price of lesser flexibility for the user. Furthermore, it is interesting to mention that the `printf()` average tracepoint cost is almost equal to the sum of LTTng's and `printf()`'s costs. This was to be expected as the `printf()` function simply combines LTTng's regular tracepoint mechanisms to `printf()`'s pretty-printing functionality.

LTTng UST version 2.9, used in this work, brings major performance improvements. Previous benchmarks ran on slightly older versions of the tracer had suggested that tracepoints were 72% slower. Version 2.9 introduces a faster inline `memcpy`, inlining some ring-buffer access functions, as well as other internal improvements regarding shared memory management. Additionally, defining the `_LGPL_SOURCE` macro in the benchmark code, which causes `liburcu`<sup>6</sup> to inline functions, further improves performance by 5 nanoseconds per tracepoint

<sup>6</sup>Userspace RCU library, required by LTTng UST



on average. On the other hand, a major performance setback for LTTng UST is the usage of the `getcpu()` system call on each tracepoint, as this information is part of the payload of all events. Furthermore, scheduling is disabled around the call-site to avoid CPU migration while the CPU number is being read which further adds overhead. It is eligible for improvement in the near future as a new mechanism might be integrated into the Linux kernel for a faster access to kernel data from user space. The patch which introduces the *restartable sequences* as an alternative to `getcpu()` [DesnoyersDesnoyers2016a], is submitted by the creator of the LTTng project, and implements the *Prepare-Write-Restart* pattern. Instead of disabling scheduling, the user space process alerts the kernel of the operation it wants to perform (the *prepare* step). It then tries to read the CPU number and is forced to restart if the information is invalid (the thread was migrated to another CPU) by the time it has been read, thus the *restart* step of the pattern. A cleanup function has to be provided to handle the case where the sequence is aborted and has to be restarted. Restartable sequences promise major improvements for user space tracepoints as they directly shorten their critical path. Preliminary benchmarks have shown the latency of tracepoints to drop to a little over 100ns per event (instead of 158ns) with a kernel supporting restartable sequences and a version of LTTng which takes advantage of them. Reading the CPU number and shared memory abstraction in user space are the two major causes of the performance difference between LTTng in kernel space and user space.

Although the Extrae tracer shows low overhead, it is still eligible for improvement. Sampling the clock either uses the TSC<sup>7</sup> (x86 only), or `rusage`<sup>8</sup> to query the operating system about user and kernel CPU usage (respectively `rusage.utime` and `rusage.stime`). The latter case, although not used in this work, adds significant overhead to the critical path of each event write, as interaction with the operating system via a system call is required. The tracer also keeps internal states between which it switches when entering and leaving a tracepoint. When entering a tracepoint, the tracer checks for the amount of free memory in the buffer for the event to be written. When the amount of the remaining free memory is too low, the buffer is flushed synchronously before the event can be recorded. Although this approach has the advantage of not dropping any events, it may significantly alter the behavior of the traced application. Moreover, Extrae samples hardware counters using PAPI<sup>9</sup> at each event write. Extrae uses arrays to store per-thread data, where each element of the array is specific to a thread. However, because the array isn't cache-aligned, false sharing might occur when the traced application is multi-threaded, which might slightly impact the pipeline efficiency. Additional experimental runs did in fact show that L1 Data cache misses increase when Extrae is tracing a multi-threaded software. Finally, browsing the Extrae source code shows the use of many levels of unnecessary function calls. For instance, some functions have the sole role

of assigning values to global variables and, in some cases, two functions only differ by the value they assign to the same global variable. Although some of these function calls might be inlined by the compiler, they can contribute to the performance overhead in case they are not. In all cases, the critical path of the tracepoint code goes through many layers of function calls that are not eligible for inlining. Having the instruction pointer move across a large memory region might increase the number of cache misses (L1 Instruction cache), in addition to the regular procedure of setting up the stack for each function call and exit. Moreover, some functions in the Extrae code do not contribute to time nor space optimization, and can potentially be replaced by C macros to maintain code readability. It is worth mentioning that Extrae isn't reentrant, and thus can have a much more straightforward probe with regard to writing the tracing data, since there is no need for synchronization nested writes.

VampirTrace is more optimized than Extrae. First, data accessed by the tracing probe is cache-aligned to avoid false sharing. Furthermore, it makes extensive use of macros along the critical path of the probe, to avoid costly function calls. Similarly to Extrae, a check for the free size of the buffer is made on the critical path and causes a flush to disk when not enough free memory is left. The implementation of VampirTrace is straightforward and is highly optimized, but doesn't guarantee reentrancy. Finally, as shown in Table IV, the standard deviation shown in VampirTrace numbers is due to the slow path that causes trace data to be flushed, which has a small period compared to other tracers. In our benchmark, the flush rate was exactly 85 events, which means that every 85 events a flush to the filesystem had to be made. Lastly, Table IV shows that the 90th percentile of latency values for VampirTrace is much lower (17% lower) than the average, comparatively to other tracers, and is actually very close to LW-ust's. This information is valuable as it shows how much VampirTrace is optimized for the fast path (shown by the 90th percentile). The 8 nanoseconds difference between LW-ust and VampirTrace fast paths values is quite negligible and is a low price to pay on top of LW-ust to get a usable tracer. This slight overhead is due to verifying the remaining space in the buffer, on top of other internal minor procedures. VampirTrace is very close to the lower bound we set for userspace tracing.

1) *Printf*: Figure 8 shows how userspace tracers scale with the number of cores when multiple threads simultaneously trace as part of the same application. Other than `printf()` and `SystemTap`, all tracers show good scalability as they leverage per-CPU or per-thread data structures to guarantee a lockless buffering scheme. `Printf()` shows poor scalability: the average tracepoint latency more than doubles going from a single-threaded application to a two-threaded one and again when four threads are spawned. These results confirm that `printf()` uses an internal lock to protect the buffer where pretty-printed strings are stored.

### C. System calls

We explained in section VI-B that we benchmarked system calls through an empty `ioctl()` function to a `procfs`

<sup>7</sup>TimeStamp Counter

<sup>8</sup>Resource Usage

<sup>9</sup>Performance API, a Linux infrastructure to access hardware counters

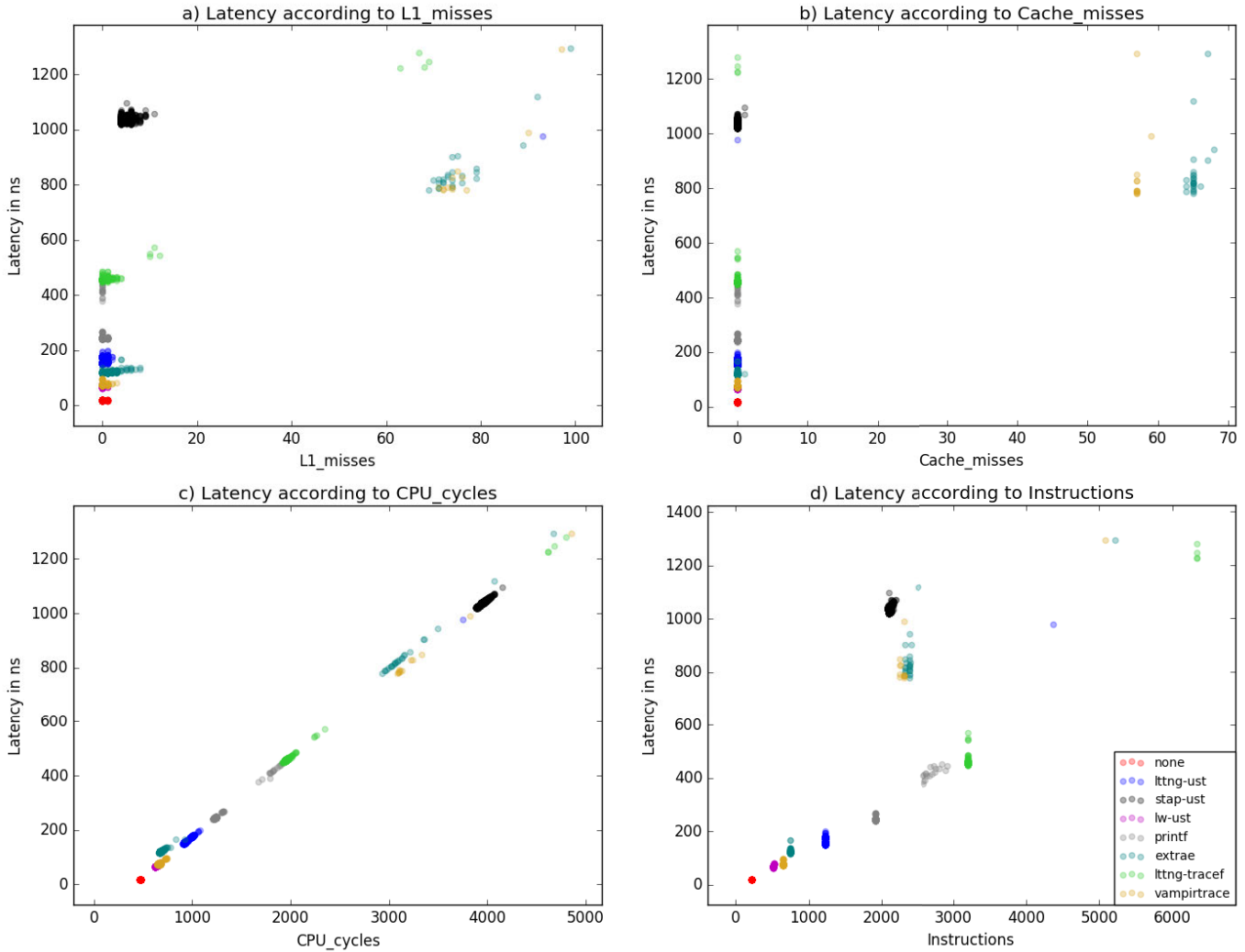


Fig. 7: All tracepoints latencies in userspace against different hardware counters

file exposed by a custom kernel module. Table V shows the results for system call tracing with different tracers. We notice Strace’s overwhelming performance overhead, making an empty `ioctl()` call almost 19 times slower. We mention that our benchmark redirects Strace’s output to `/dev/null` to avoid all overhead due to writing to a file or the terminal. We also used the `-t` flag to assign a timestamp to each system call, although the performance impact is similar even without timestamping. The large performance overhead added by Strace is due to the Ptrace infrastructure and mechanism, where a trap has to be generated by the operating system on each system call entry and exit. Along with this trap, the traced process needs to be put in the blocked state, and the tracer is woken up from the blocked state. Context switches and moving processes from the blocked to the running state by the operating system also contribute to the overhead. Furthermore, Strace doesn’t implement its own low-cost internal buffer for storing collected data, and simply outputs the results to a file descriptor.

Although Ftrace and LTTng use a different mechanism than Strace to trace system calls, they provide the same information

at a much lower cost, with even more context information than Strace. However, it is interesting to note that for system calls tracing, Ftrace shows limited information compared to Strace and LTTng. Often times, the parameters of a system call are required for investigation, rather than the system call itself (for instance, when looking at `open()` system calls to investigate files opened by a certain process). Ftrace records system calls, but data structures and literal strings received as parameters aren’t resolved and presented in the trace. Instead, the value of the pointers to those arguments is recorded, which usually has no significance a posteriori. On the other hand, resolving string literals as LTTng does has the effect of adding some overhead, as copying of internal data will be required. This added information explains the difference in system call tracing performance between LTTng and Ftrace. Although LTTng showed better performance for regular tracepoints, Ftrace has a lower latency for system calls. To prove this artifact, we benchmarked (latency only) the `open()` system call, depending on the length of the file name to open. The results are shown in Figure 9. Except for an empty file name, which we assume causes `open()` to return early without

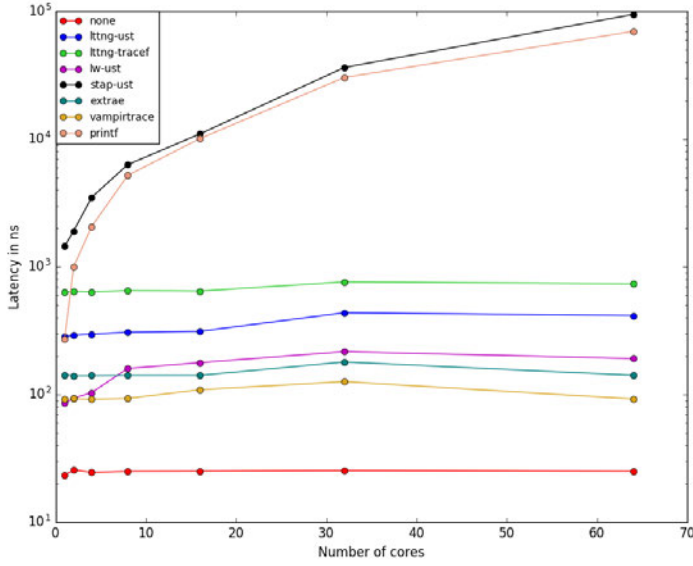


Fig. 8: Average tracepoint latency in userspace against the number of threads/cores

reaching the Virtual File System, tracing the system call using Ftrace takes 700 ns regardless of the length of the file’s name. However, tracing with LTTng adds overhead that is in linear relation to the file name length. Once again, a trade-off has to be made between overhead and usability. The requirements and needs of a specific use case can lead the users to choose either one of those tracers. Finally, Strace provides the most ease of use at a much higher cost.

TABLE V: Average latency of an empty `ioctl()` system call for multiple tracers

	Average (ns)	STD	90th percentile
None	63	0	61
Ftrace	300	13	296
LTTng	327	13	320
Strace	18991	351	18288

#### D. Additional analysis

To dig a little further into the reasons for tracepoint latency, we decided to compare the amount of time spent in the different steps of recording an event for the two most efficient tracers. As was presented in section III-B, recording an event for an enabled tracepoint is summed up to the callback mechanism, the probe, and the write to a ring buffer. The serialization part can be further divided into sampling the clock, and formatting the event and its payload into data to be written in the buffer (LTTng uses CTF [DesnoyersDesnoyers2012] as a format and Ftrace can be configured to use a binary format as well). Table III showed that recording an event requires about 92 ns with LTTng and 116 ns with Ftrace. Figure 10 shows the profile of these latencies for each tracer and breaks down the time spent at each step.

We showed in Table I that we benchmark both LTTng and Ftrace for tracing static tracepoints using the `TRACE_EVENT`

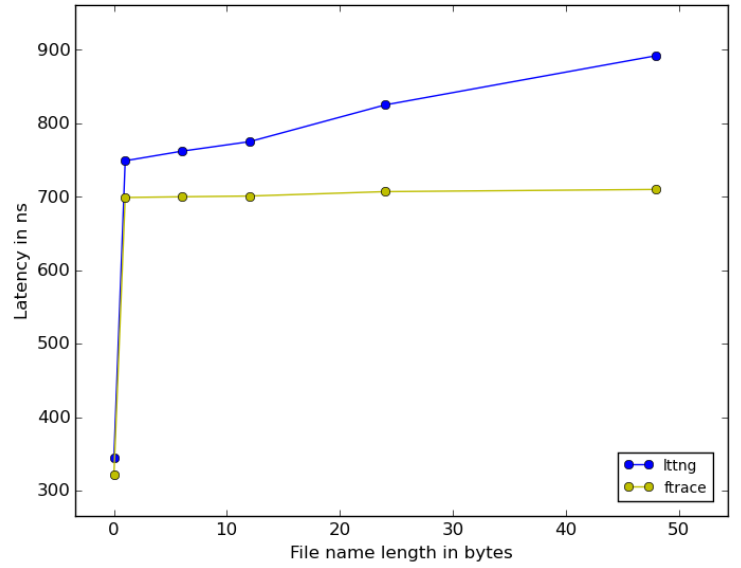


Fig. 9: Latency to trace the `open()` system call according to the file name length

macro. This callback mechanism is optimized to have the lowest cost possible and ends up requiring only 4 ns. As presented in section VI-B3, we configured Ftrace to use a global clock (instead of the default local clock) for time sampling, similarly to LTTng. Reading a global clock requires 32 ns which accounts for 34% and 27% of the total time for LTTng and Ftrace respectively. Serialization is the part that differentiates the most the overhead between the tracers. LTTng’s serialization step takes 29 ns, which accounts for 31% of its total time, while Ftrace requires 56 ns, which accounts for 47% of the entire process. Writing to the ring buffer, as they share similar designs, shows very similar overhead, although LTTng’s ring buffer is slightly slower due to the sub-buffer granularity and additional checks for boundaries. This analysis shows where Ftrace requires more time than LTTng, at the serialization step, which could be eligible for improvement.

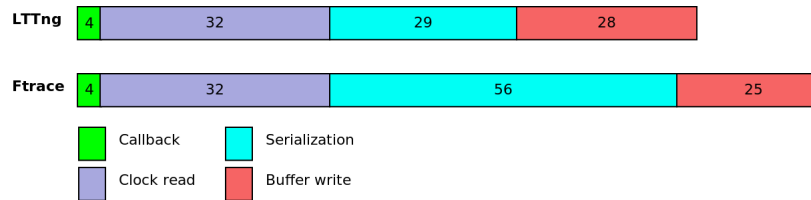


Fig. 10: Time distribution for different steps of a tracepoint for LTTng and Ftrace

#### E. Summary

As a reference, Table VI consolidates all the results and shows a summary of all tracers. The table is sorted by ascending values of the average latency. Other metrics that were only shown as graphs previously are also included.

The results of table VI allow us not only to compare tools among themselves, but also different underlying mechanisms.

TABLE VI: Summary of the overhead per tracepoint for all tracers (Average and Standard Deviation) sorted by latency average

	Latency (in ns)		L1 misses		Cache misses		Instructions		CPU cycles	
	Average	STD	Average	STD	Average	STD	Average	STD	Average	STD
VampirTrace (Userspace)	84	87	0	8	0	6	670	233	695	299
LTTng	92	5	0	0	0	0	870	24	629	18
Ftrace	116	7	0	0	0	0	1053	18	709	26
Extrac (Userspace)	117	7	1	2	0	0	750	0	672	25
Perf	121	1	0	0	0	0	1099	0	720	6
LTTng-Kprobe	123	5	0	0	0	0	1012	20	731	18
eBPF-tracing	128	5	0	0	0	0	1138	0	746	18
SystemTap	130	5	0	0	0	0	1157	0	753	20
eBPF	140	10	0	0	0	0	1157	0	782	32
LTTng-UST (Userspace)	158	27	0	3	0	0	1228	117	944	161
Printf (Userspace)	245	144	0	6	0	3	1916	168	1250	492
LTTng-tracef (Userspace)	446	73	0	4	0	0	3211	205	1917	253
SystemTap (Userspace)	1046	45	4	2	0	0	2111	68	3958	155
Strace	18224	252	336	14	0	0	10666	362	11984	447

The TRACEPOINT infrastructure is the most efficient callback mechanism, compared to Kprobes and other mechanisms. The choice of the tracer to use can be difficult depending on the needs of the developer. For kernel tracing, LTTng and Ftrace serve similar purposes and both add little overhead to the system when tracing. Ftrace is more configurable than LTTng, for low-level tweaking such as the choice of the clock. Ftrace is integrated into the Linux kernel in most distributions, and as a result is more easily accessible. On the other hand, LTTng is more feature-rich in terms of usability. For example, the CTF output allows LTTng to be compatible with visualization and analysis tools that understand this format. Tracing using Perf is the best choice when quick statistics are needed. Its ease of use is overwhelming compared to Ftrace and LTTng, but its output (as a tracer) is a short summary and hides the details that are provided by other kernel tracers. SystemTap and eBPF serve similar roles, although eBPF is gaining more traction in recent kernel versions. eBPF adds slightly more overhead than SystemTap due to the need of an in-kernel virtual machine, while SystemTap translates user scripts into loadable kernel modules. The wide availability of targeted tools for kernel subsystems make eBPF, particularly within the BCC project [??bcc2017], a more user-friendly tool than SystemTap. As for high performance applications, VampirTrace performs better than Extrac, and provides similar features as well, making it a better choice in terms of added overhead.

### VIII. CONCLUSION

In this paper, we explained the designs of different kernel and user space tracers. The benchmark results not only helped to quantify the overhead that each tracer adds to the traced application, but also point out specific details about the internals of each tracer. By sampling low-level performance counters, this analysis helps users determine where the overhead usually originates from, and how to configure tracers to reduce their footprint depending on their requirements. Lastly, this research also helps tracer developers, as we were able to point out potential optimizations in some tracers, and showed design limitations in others. In conclusion, designing a tracer is a trade-off between usability, features and overhead. Fundamental design choices have to be made that dictate how the tracer behaves under critical conditions.

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