# Monogamy of entanglement without inequalities 

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We provide a fine-grained definition of monogamous measure of entanglement that does not invoke any particular monogamy relation. Our definition is given in terms of an equality, as opposed to inequality, that we call the "disentangling condition". We relate our definition to the more traditional one, by showing that it generates standard monogamy relations. We then show that all quantum Markov states satisfy the disentangling condition for any entanglement monotone. In addition, we demonstrate that entanglement monotones that are given in terms of a convex roof extension are monogamous if they are monogamous on pure states, and show that for any quantum state that satisfies the disentangling condition, its entanglement of formation equals the entanglement of assistance. We characterize all bipartite mixed states with this property, and use it to show that the G-concurrence is monogamous. In the case of two qubits, we show that the equality between entanglement of formation and assistance holds if and only if the state is a rank 2 bipartite state that can be expressed as the marginal of a pure 3 -qubit state in the $W$ class.

Monogamy of entanglement is one of the nonintuitive phenomena of quantum physics that distinguish it from classical physics. Classically, three random bits can be maximally correlated. For example, three coins can be prepared in a state in which with $50 \%$ chance all three coins show "head", and with the other $50 \%$ chance they all show "tail".

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In such a preparation, any two coins are maximally correlated. In contrast with the classical world, it is not possible to prepare three qubits $A, B, C$ in a way that any two qubits are maximally entangled [1]. In fact, if qubit $A$ is maximally entangled with qubit $B$, then it must be uncorrelated (not even classically) with qubit $C$. This phenomenon of monogamy of entanglement was first quantified in a seminal paper by Coffman, Kundu, and Wootters (CKW) [1] for three qubits, and later on studied intensively in more general settings $[2,3,4,5,14,15,6,16,17,18,7,19,20,21$, $22,23,24,25,8,9,26,27,10,11,28,29,30,12$, $13,31,32,33,34,35,36,37,38,39,40]$.

Qualitatively, monogamy of entanglement measures the shareability of entanglement in a composite quantum system, i.e., the more two subsystems are entangled the less this pair has entanglement with the rest of the system. This feature of entanglement has found potential applications in many quantum information tasks and other areas of physics, such as quantum key distribution [41, $42,43]$, classification of quantum states [44, 45, 46], condensed-matter physics [47, 48, 49], frustrated spin systems [50], statistical physics [51], and even black-hole physics [52, 53].

A monogamy relation is quantitatively displayed as an inequality of the following form:

$$
\begin{equation*}
E(A \mid B C) \geq E(A \mid B)+E(A \mid C) \tag{1}
\end{equation*}
$$

where $E$ is a measure of bipartite entanglement and $A, B, C$ are three subsystems of a composite quantum system. It states that the sum of the entanglement between $A$ and $B$, and between $A$ and $C$, can not exceed the entanglement between $A$ and the joint system $B C$. While not all measures of entanglement satisfy this relation, some do. Con-
sequently, any measure of entanglement that does satisfy (1) was called in literature monogamous.

Here we argue that this definition of monogamous measure of entanglement captures only partially the property that entanglement is monogamous. This is evident from the fact that many important measures of entanglement do not satisfy the relation (1). Some of these measures are not even additive under tensor product [54, 55, 56] (in fact, some measures [57] are multiplicative under tensor product). Therefore, the summation in the RHS of (1) is clearly only a convenient choice and not a necessity. For example, it is well known that if $E$ does not satisfy this relation, it is still possible to find a positive exponent $\alpha \in \mathbb{R}_{+}$, such that $E^{\alpha}$ satisfies the relation. This was already apparent in the seminal work of [1] in which $E$ was taken to be the square of the concurrence and not the concurrence itself. More recently, it was shown that many other measures of entanglement satisfy the monogamy relation (1) if $E$ is replaced by $E^{\alpha}$ for some $\alpha>1[13,8,12,5,7,11,9,6,10]$.

One attempt to address these issues with the current definition of a monogamous measure of entanglement, is to replace the relation (1) with a family of monogamy relations of the form, $E(A \mid B C) \geq$ $f(E(A \mid B), E(A \mid C))$, where $f$ is some function of two variables that satisfy certain conditions [31]. However, such an approach is somewhat artificial in the sense that the monogamy relations are not derived from more basic principles.

In this paper we take another approach to monogamy of entanglement which is more "finegrained" in nature, and avoid the introduction of such a function $f$. Therefore, our definition of monogamous measure of entanglement (see Definition 1 below) does not involve a monogamy relation, but instead a condition on the measure of entanglement that we call the disentangling condition (following the terminology of [28]). Yet, we show that our definition is consistent with the more traditional notion of monogamy of entanglement, by showing that $E$ is monogamous according to our definition if and only if there exists an $\alpha>0$ such that $E^{\alpha}$ satisfies (1). Consequently, many more measures of entanglement are monogamous according to our definition. We then provide a character-
ization for the disentangling condition in the form of an equality between the entanglement of formation (EoF) associated with the given entanglement measure (see Eq. (7) below) and the entanglement of assistance (EoA) [58], and discuss its relation to quantum Markov chains [59]. In addition, we characterize all states for which EoF equals to EoA when the measure of entanglement is taken to be the G-concurrence, and use that to show that the G-concurrence is monogamous. Finally, we show that in the 2-dimensional case, the bipartite 2-qubit mixed states that can be expressed as the marginal of the 3 -qubit $W$-state, are the only 2 -qubit entangled states for which the EoF equals the EoA with respect to any measure of entanglement.

Let $\mathcal{H}^{A} \otimes \mathcal{H}^{B} \equiv \mathcal{H}^{A B}$ be a bipartite Hilbert space, and $\mathcal{S}\left(\mathcal{H}^{A B}\right) \equiv \mathcal{S}^{A B}$ be the set of density matrices acting on $\mathcal{H}^{A B}$. A function $E$ : $\mathcal{S}^{A B} \rightarrow \mathbb{R}_{+}$is called a measure of entanglement if (1) $E\left(\sigma^{A B}\right)=0$ for any separable density matrix $\sigma^{A B} \in \mathcal{S}^{A B}$, and (2) $E$ behaves monotonically under local operations and classical communications (LOCC). That is, given an LOCC map $\Phi$ we have

$$
\begin{equation*}
E\left(\Phi\left(\rho^{A B}\right)\right) \leq E\left(\rho^{A B}\right), \quad \forall \rho^{A B} \in \mathcal{S}^{A B} \tag{2}
\end{equation*}
$$

The measure is said to be faithful if it is zero only on separable states.

The map $\Phi$ is completely positive and trace preserving (CPTP). In general, LOCC can be stochastic, in the sense that $\rho^{A B}$ can be converted to $\sigma_{j}^{A B}$ with some probability $p_{j}$. In this case, the map from $\rho^{A B}$ to $\sigma_{j}^{A B}$ can not be described in general by a CPTP map. However, by introducing a 'flag' system $A^{\prime}$, we can view the ensemble $\left\{\sigma_{j}^{A B}, p_{j}\right\}$ as a classical quantum state $\sigma^{A^{\prime} A B} \equiv$ $\sum_{j} p_{j}|j\rangle\left\langle\left. j\right|^{A^{\prime}} \otimes \sigma_{j}^{A B}\right.$. Hence, if $\rho^{A B}$ can be converted by LOCC to $\sigma_{j}^{A B}$ with probability $p_{j}$, then there exists a CPTP LOCC map $\Phi$ such that $\Phi\left(\rho^{A B}\right)=\sigma^{A^{\prime} A B}$. Therefore, the definition above of a measure of entanglement captures also probabilistic transformations. Particularly, $E$ must satisfy $E\left(\sigma^{A^{\prime} A B}\right) \leq E\left(\rho^{A B}\right)$.

Almost all measures of entanglement studied in literature (although not all [60]) satisfy

$$
\begin{equation*}
E\left(\sigma^{A^{\prime} A B}\right)=\sum_{j} p_{j} E\left(\sigma_{j}^{A B}\right) \tag{3}
\end{equation*}
$$

which is very intuitive since $A^{\prime}$ is just a classical system encoding the value of $j$. In this case the condition $E\left(\sigma^{A^{\prime} A B}\right) \leq E\left(\rho^{A B}\right)$ becomes $\sum_{j} p_{j} E\left(\sigma_{j}^{A B}\right) \leq E\left(\rho^{A B}\right)$. That is, LOCC can not increase entanglement on average. Measures of entanglement that satisfy this property are called entanglement monotones, and they can also be shown to be convex [61]. In the following definition we denote density matrices acting on a finite dimensional tripartite Hilbert space $\mathcal{H}^{A B C}$ by $\rho^{A \mid B C}$, where the vertical bar indicates the bipartite split across which we will measure the (bipartite) entanglement.

Definition 1. Let $E$ be a measure of entanglement. $E$ is said to be monogamous if for any $\rho^{A \mid B C} \in \mathcal{S}^{A B C}$ that satisfies

$$
\begin{equation*}
E\left(\rho^{A \mid B C}\right)=E\left(\rho^{A B}\right) \tag{4}
\end{equation*}
$$

we have that $E\left(\rho^{A C}\right)=0$.
The condition in (4) is a very strong one and typically is not satisfied by most tripartite states $\rho^{A B C}$. Following the terminology of [28] we call it here the "disentangling condition". We will see below that states that saturate the strong subadditivity inequality for the von Neumann entropy (i.e. quantum Markov states) always satisfy this equality, for any entanglement monotone $E$.
Note that the condition in (1) is stronger than the one given in Definition 1. Indeed, if $E$ satisfies (1), then any $\rho^{A \mid B C}$ that satisfies (4) must have $E\left(\rho^{A C}\right)=0$. At the same time, Definition 1 captures the essence of monogamy: that is, if system $A$ shares the maximum amount of entanglement with subsystem $B$, it is left with no entanglement to share with $C$.

In Definition 1 we do not invoke a particular monogamy relation such as (1). Instead, we propose a minimalist approach which is not quantitative, in which we only require what is essential from a measure of entanglement to be monogamous. Yet, this requirement is sufficient to generate a more quantitative monogamy relation:

Theorem 1. Let $E$ be a continuous measure of entanglement. Then, $E$ is monogamous according
to Definition 1 if and only if there exists $0<\alpha<$ $\infty$ such that

$$
\begin{equation*}
E^{\alpha}\left(\rho^{A \mid B C}\right) \geq E^{\alpha}\left(\rho^{A B}\right)+E^{\alpha}\left(\rho^{A C}\right) \tag{5}
\end{equation*}
$$

for all $\rho^{A B C} \in \mathcal{S}^{A B C}$ with fixed $\operatorname{dim} \mathcal{H}^{A B C}=d<$ $\infty$.

We call the smallest possible value for $\alpha$ that satisfies Eq. (5) in a given dimension $d=\operatorname{dim}\left(\mathcal{H}^{A B C}\right)$, the monogamy exponent associated with a measure $E$, and denote it simply as $\alpha(E)$. In general, the monogamy exponent is hard to compute, and in the supplemental material we provide (along with the proofs of the theorems in this paper) a comprehensive list of known bounds for the monogamy exponent when $E$ is one of the measures of entanglement that were studied extensively in literature.

It is important to note that the relation given in (5) is not of the form $E(A \mid B C) \geq$ $f(E(A \mid B), E(A \mid C))$, where $f$ is some function of two variables that satisfies certain conditions and is independent of $d$ [31]. This is because the monogamy exponent in (5) depends on the dimension $d$, whereas $f$ as was used in [31] is universal in the sense that it does not depend on the dimension. Therefore, if a measure of entanglement such as the entanglement of formation is not monogamous according to the class of relations given in [31], it does not necessarily mean that it is not monogamous according to our definition. Moreover, since our approach allows for dependence in the dimension, it avoids the issues raised in [31] that measures of entanglement cannot be simultaneously monogamous and (geometrically) faithful (see [31] for their definition of faithfulness).

In general, the class of all states $\rho^{A B C}$ that satisfy the disentangling condition (4) depends on the choice of the entanglement measure $E$. However, there is a class of states that satisfy this condition for any choice of an entanglement monotone $E$. These are precisely the states that saturate the strong subadditivity of the von-Neumann entropy [59]. For such states, the system $B$ Hilbert space $\mathcal{H}^{B}$ must have a decomposition into a direct sum of tensor products $\mathcal{H}^{B}=\bigoplus_{j} \mathcal{H}^{B_{j}^{L}} \otimes \mathcal{H}^{B_{j}^{R}}$, such that the state $\rho^{A B C}$ has the form

$$
\begin{equation*}
\rho^{A B C}=\bigoplus_{j} q_{j} \rho^{A B_{j}^{L}} \otimes \rho^{B_{j}^{R} C}, \tag{6}
\end{equation*}
$$

where $q_{j}$ is a probability distribution [59].
Theorem 2. Let $E$ be an entanglement monotone. Then, E satisfies the disentangling condition (4) for all Markov quantum states $\rho^{A B C}$ of the form (6).

Note that a state of the form (6) has a separable marginal state $\rho^{A C}$, and therefore the above theorem is consistent with Definition 1. Consequently, to check if an entanglement monotone is monogamous, one has to consider tripartite states that satisfy (4) but have a different form than (6). Perhaps such states do not exist for certain entanglement monotones. Indeed, partial results in this direction were proved recently in [28]. Particularly, it was shown that any pure tripartite entangled state $\psi^{A B C}$ with bipartite marginal state $\rho^{A B}$, that satisfies $N\left(\psi^{A \mid B C}\right)=N\left(\rho^{A B}\right)$, where $N$ is the negativity measure, must have the form $\psi^{A B C}=\psi^{A B^{L}} \otimes \psi^{B^{R} C}$. Note that this is precisely the form (6) when $\rho^{A B C}$ is pure. Like the Negativity, we will see later that also the G-concurrence has this property for pure tripartite states.

Typically it is hard to check if a measure of entanglement is monogamous since the conditions in (4) involves mixed tripartite states. However, as we show below, it is sufficient to consider only pure tripartite states in (4), if the entanglement measure $E$ is defined on mixed states by a convex roof extension; that is,

$$
\begin{equation*}
E_{f}\left(\rho^{A B}\right) \equiv \min \sum_{j=1}^{n} p_{j} E\left(\left|\psi_{j}\right\rangle\left\langle\left.\psi_{j}\right|^{A B}\right)\right. \tag{7}
\end{equation*}
$$

where the minimum is taken over all pure state decompositions of $\rho^{A B}=\sum_{j=1}^{n} p_{j}\left|\psi_{j}\right\rangle\left\langle\left.\psi_{j}\right|^{A B}\right.$. We call $E_{f}$ the entanglement of formation (EoF) associated with $E$. In general, it could be that $E_{f} \neq E$ on mixed states like the convex roof extended Negativity is different than the Negativity itself.
Theorem 3. Let $E_{f}$ be an entanglement monotone as above. If $E_{f}$ is monogamous (according to Definition 1) on pure tripartite states in $\mathcal{H}^{A B C}$, then it is also monogamous on tripartite mixed states acting on $\mathcal{H}^{A B C}$.

From the above theorem and Corollary 5 below it follows that the convex roof extended negativity is monogamous, since it is monogamous for
pure states [28]. Similarly, we will use it to show that the G-concurrence is monogamous. However, monogamy alone does not necessarily imply that a tripartite state is a Markov state if it satisfies the disentangling condition. In the following theorem, we provide yet another property of a tripartite state $\rho^{A B C}$ that satisfies the disentangling condition. The property is that any LOCC protocol on such a tripartite state can not increase on average the initial bipartite entanglement between $A$ and $B$. The maximum such average of bipartite entanglement is known as entanglement of collaboration [62]. It is defined by:

$$
\begin{equation*}
E_{c}^{A B \mid C}\left(\rho^{A B C}\right)=\max \sum_{j=1}^{n} p_{j} E\left(\rho_{j}^{A B}\right) \tag{8}
\end{equation*}
$$

where $E$ is a given measure of bipartite entanglement, and the maximum is taken over all tripartite LOCC protocols yielding the bipartite state $\rho_{j}^{A B}$ with probability $p_{j}$. This measure of entanglement is closely related to EoA, denoted here by $E_{a}$, in which the optimization above is restricted to LOCC with one way classical communication from system $C$ to systems $A$ and $B$. Therefore, $E_{c}^{A B \mid C} \geq E_{a}$ and in [62] it was shown that this inequality can be strict.
Theorem 4. Let $E$ be an entanglement monotone on bipartite mixed states, and let $\rho^{A B C}$ be a (possibly mixed) tripartite state satisfying the disentangling condition (4). Then,

$$
\begin{equation*}
E\left(\rho^{A B}\right)=E_{c}^{A B \mid C}\left(\rho^{A B C}\right) \tag{9}
\end{equation*}
$$

The condition in (9) is a very strong one. Particularly, it states that all measurements on system $C$ yield the same average of bipartite entanglement between $A$ and $B$. Therefore, the entanglement of the marginal state $\rho^{A B}$ is resilient to any quantum process or measurement on system $C$. In the case that $\rho^{A B C}$ is pure, the EoA, $E_{a}$, depends only on the marginal $\rho^{A B}$, and we get the following corollary:
Corollary 5. Let $E$ be an entanglement monotone on bipartite states, and let $\rho^{A B C}$ be a pure tripartite state satisfying the disentangling condition (4). Then,

$$
\begin{equation*}
E\left(\rho^{A B}\right)=E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right) \tag{10}
\end{equation*}
$$

where the EoF, $E_{f}$, is defined in (7), and EoA, $E_{a}$, is also defined as in (7) but with a maximum replacing the minimum.

The equality of $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ in corollary 5 implies that any pure state decomposition of $\rho^{A B}$ has the same average entanglement when measured by $E$. Unless $\rho^{A B}$ is pure, it is almost never satisfied. Nonetheless, there are non-trivial states (of measure zero) that do satisfy it. Such states have the following property:

Theorem 6. Let $E$ be a measure of bipartite entanglement, and let $\rho^{A B} \in \mathcal{S}^{A B}$ be a bipartite state with support subspace $\operatorname{supp}\left(\rho^{A B}\right)$. If $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ then for any $\sigma^{A B} \in \mathcal{S}^{A B}$ with $\operatorname{supp}\left(\sigma^{A B}\right) \subseteq \operatorname{supp}\left(\rho^{A B}\right)$ we have $E_{f}\left(\sigma^{A B}\right)=$ $E_{a}\left(\sigma^{A B}\right)$.

Remark 1. The theorem above follows from some of the results presented in [33], and for completeness we provide its full proof in the appendix.

Theorem 6 demonstrates that the equality between the EoF and EoA corresponds to a property of the support of $\rho^{A B}$ rather than $\rho^{A B}$ itself. In the following theorem, we characterize the precise form of the support of $\rho^{A B}$ that yields such an equality. The entanglement monotone we are using is a generalization of the concurrence measure known as the G-concurrence [57].

The G-concurrence is an entanglement monotone that on pure bipartite states is equal to the geometric-mean of the Schmidt coefficients. On a (possibly unnormalized) vector $|x\rangle \in \mathbb{C}^{d} \otimes \mathbb{C}^{d}$, it can be expressed as:

$$
\begin{equation*}
G(|x\rangle)=|\operatorname{det}(X)|^{2 / d} ; \quad|x\rangle=X \otimes I\left|\phi_{+}\right\rangle \tag{11}
\end{equation*}
$$

where $\left|\phi_{+}\right\rangle=\sum_{i=1}^{d}|i\rangle^{A}|i\rangle^{B}$ is the maximally entangled state, and $X \in \mathcal{M}_{d}(\mathbb{C})$ is a $d \times d$ complex matrix. For mixed states the G-concurrence is defined in terms of the convex roof extension; that is, $G\left(\rho^{A B}\right):=G_{f}\left(\rho^{A B}\right)$.

In the following theorem we denote by $\mathcal{N} \subset$ $\mathcal{M}_{d}(\mathbb{C})$, the nilpotent cone, consisting of all $d \times d$ nilpotent complex matrices (i.e. $X \in \mathcal{N}$ iff $X^{k}=0$ for some $1 \leq k \leq d)$. While the set $\mathcal{N}$ is not a vector space, it contains subspaces. For example, the set $\mathcal{T}$ of all strictly upper triangular matrices
(i.e. upper triangular matrices with zeros on the diagonal) is a $d(d-1) / 2$-dimensional subspace in $\mathcal{N}$. From Gerstenhaber's theorem [63] it follows that the largest dimension of a subspace in $\mathcal{N}$ is $d(d-1) / 2$, and if a subspace $\mathcal{N}_{0} \subset \mathcal{N}$ has this maximal dimension, then it must be similar to $\mathcal{T}$ (i.e. their exists an invertible matrix $S$ such that $\left.\mathcal{N}_{0}=S \mathcal{T} S^{-1}\right)$.
Theorem 7. Let $\rho^{A B}$ be a bipartite density matrix acting on $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$ with rank $r$. Then,

$$
\begin{equation*}
G\left(\rho^{A B}\right)=G_{a}\left(\rho^{A B}\right)>0, \tag{12}
\end{equation*}
$$

if and only if $r \leq 1+d(d-1) / 2$, and there exists a full Schmidt rank state $|x\rangle^{A B} \in \mathbb{C}^{d} \otimes \mathbb{C}^{d}$, and a subspace $\mathcal{N}_{0} \subset \mathcal{N}$, with $\operatorname{dim}\left(\mathcal{N}_{0}\right)=r-1$, such that

$$
\begin{equation*}
\operatorname{supp}\left(\rho^{A B}\right)=\left\{I \otimes Y|x\rangle^{A B} \mid Y \in \mathcal{K}\right\} \tag{13}
\end{equation*}
$$

where $\mathcal{K} \equiv\{I\} \oplus \mathcal{N}_{0} \equiv\left\{c I+N \mid c \in \mathbb{C} ; N \in \mathcal{N}_{0}\right\}$.
The direct sum above is consistent with the fact that a nilpotent matrix has a zero trace, so that it is orthogonal to the identity matrix in the HilbertSchmidt inner product. The theorem above implies that the G-concurrence is monogamous on pure tripartite states:
Corollary 8. Let $\left|\psi^{A B C}\right\rangle \in \mathbb{C}^{d} \otimes \mathbb{C}^{d} \otimes \mathbb{C}^{n}$ be a pure tripartite state, with bipartite marginal $\rho^{A B}$. If $G\left(\psi^{A \mid B C}\right)=G\left(\rho^{A B}\right)>0$ then $|\psi\rangle^{A B C}=$ $|\chi\rangle^{A B}|\varphi\rangle^{C}$ for some bipartite state $|\chi\rangle^{A B} \in \mathbb{C}^{d} \otimes \mathbb{C}^{d}$ and a vector $|\varphi\rangle^{C} \in \mathbb{C}^{n}$.

The above result is somewhat surprising since the G-concurrence is not a faithful measure of entanglement. Yet, it states that the disentangling condition forces the marginal state $\rho^{A C}$ to be a product state. This is much stronger than $G\left(\rho^{A C}\right)=0$ (which can even hold for some entangled $\rho^{A C}$ ), and in particular, it states that $A$ and $C$ can not even share classical correlation. Combining the above corollary with Theorem 3 implies that the G-concurrence is monogamous on any mixed state that is acting on $\mathbb{C}^{d} \otimes \mathbb{C}^{d} \otimes \mathbb{C}^{n}$. Finally, in the qubit case, Theorem 7 takes the following form:

Corollary 9. Let $\rho^{A B}$ be an entangled two qubit state with rank $r>1$, and let $E$ be any injective (up to local unitaries) measure of pure two qubit entanglement. Then,

1. If $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ then $r=2$, and $\rho^{A B}=\operatorname{Tr}_{C}|W\rangle\left\langle\left. W\right|^{A B C}\right.$ is the 2-qubit marginal of $|W\rangle=\lambda_{1}|100\rangle+\lambda_{2}|010\rangle+\lambda_{3}|001\rangle+\lambda_{4}|000\rangle$, $\lambda_{i} \in \mathbb{C}$.
2. Conversely, if $\rho^{A B}$ is a marginal of a state in the $W$-class then $C_{f}\left(\rho^{A B}\right)=C_{a}\left(\rho^{A B}\right)$, where $C$ is the concurrence.

Remark 2. The second part of the theorem implies that also $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$, for any $\rho^{A B}$ that is a marginal of a W -state, and any measure $E$, that can be expressed as a convex function of the concurrence $C$ [64].

In conclusions, we introduced a new definition for a monogamous measure of entanglement. Our definition involves an equality (the disentangling condition (4)) rather than the inequality (1). Yet, we showed that our notion of monogamy can reproduce monogamy relations like in (1) with a small change that the measure $E$ is replaced by $E^{\alpha}$ for some exponent $\alpha>0$. We then showed that convex roof based entanglement monotones of the form (7) are monogamous iff they are monogamous on pure tripartite states, and showed further that the disentangling condition in (4) holds for any entanglement monotone if $\rho^{A B C}$ is a quantum Markov state. While it is left open if the converse is also true (at least for some entanglement monotones), we were able to show that for the G-concurrence, the only pure tripartite states that satisfy the disentangling condition (4) are Markov states. In addition, we related the disentangling condition to states that have the same average entanglement for all convex pure state decompositions, and found a characterization of such states in Theorem 6 (for general measures of entanglement), and a complete characterization in Theorem 7 (for the G-concurrence). Clearly, much more is left to investigate along these lines.

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## Supplementary Material <br> Monogamy of entanglement without inequalities

## A The Monogamy exponent

Theorem 1. Let $E$ be a continuous measure of entanglement. Then, $E$ is monogamous according to Definition 1 if and only if there exists $0<\alpha<\infty$ such that

$$
\begin{equation*}
E^{\alpha}\left(\rho^{A \mid B C}\right) \geq E^{\alpha}\left(\rho^{A B}\right)+E^{\alpha}\left(\rho^{A C}\right) \tag{14}
\end{equation*}
$$

for all $\rho^{A B C} \in \mathcal{S}^{A B C}$ with fixed $\operatorname{dim} \mathcal{H}^{A B C}=d<\infty$.
Proof. Let $E$ be a monogamous measure of entanglement according to Def.1. Since $E$ is a measure of entanglement, it is non-increasing under partial traces, and therefore $E\left(\rho^{A \mid B C}\right) \geq \max \left\{E\left(\rho^{A B}\right), E\left(\rho^{A C}\right)\right\}$ for any state $\rho^{A \mid B C} \in \mathcal{S}^{A B C}$. We assume $E\left(\rho^{A \mid B C}\right)>0$ and set $x_{1} \equiv E\left(\rho^{A B}\right) / E\left(\rho^{A \mid B C}\right)$ and $x_{2} \equiv$ $E\left(\rho^{A C}\right) / E\left(\rho^{A \mid B C}\right)$. Clearly, there exists $\gamma>0$ such that

$$
\begin{equation*}
x_{1}^{\gamma}+x_{2}^{\gamma} \leq 1, \tag{15}
\end{equation*}
$$

since either $x_{j}^{\gamma} \rightarrow 0$ when $\gamma$ increases, or if $x_{1}=1$ then by assumption $x_{2}=0$ and vise versa. We denote by $f\left(\rho^{A B C}\right)$ the smallest value of $\gamma$ that achieves equality in (15). Since $E$ is continuous, so is $f$, and the compactness of $\mathcal{S}^{A B C}$ gives:

$$
\begin{equation*}
\alpha \equiv \max _{\rho^{A B C} \in \mathcal{S}^{A B C}} f\left(\rho^{A B C}\right)<\infty . \tag{16}
\end{equation*}
$$

By definition, $\alpha$ satisfies the condition in (5).
As discussed in the paper, the expression for $\alpha$ in (16) is optimal in the sense that it provides the smallest possible value for $\alpha$ that satisfies Eq. (5). This monogamy exponent is a function of the measure $E$, and we denote it by $\alpha(E)$. It may depend also on the dimension $d \equiv \operatorname{dim}\left(\mathcal{H}^{A B C}\right)$ (see in Table 1 ), and, in general, is hard to compute especially in higher dimensions $[1,4,6,2,3,5]$. By definition, $\alpha(E)$ can only increase with $d$ (e.g. Table 1). Table 1 indicates that almost any entanglement measure is monogamous at least for multi-qubit systems. In addition, almost all the entanglement measures studied in the literature are continuous, and in particular $C, N, N_{c r}, E_{f}, \tau, T_{q}, R_{\alpha}$ and $E_{r}$ are all continuous [27, 28, 29, 30].

## B Quantum Markov States and Monogamy of Entanglement

Recall that quantum Markov states are states that saturate the strong subadditivity of the von-Neumann entropy. That is, they saturate the inequality:

$$
\begin{equation*}
S\left(\rho^{A B}\right)+S\left(\rho^{B C}\right) \geq S\left(\rho^{A B C}\right)+S\left(\rho^{B}\right), \tag{17}
\end{equation*}
$$

where $S(\rho)=-\operatorname{Tr}[\rho \log \rho]$ is the von-Neumann entropy. In [31] it was shown that the inequality above is saturated if and only if the Hilbert space of system $B, \mathcal{H}^{B}$, can be decomposed into a direct sum of tensor products

$$
\begin{equation*}
\mathcal{H}^{B}=\bigoplus_{j} \mathcal{H}^{B_{j}^{L}} \otimes \mathcal{H}^{B_{j}^{R}} \tag{18}
\end{equation*}
$$

Table 1: A comparison of the monogamy exponent of several entanglement measures. We denote the one-way distillable entanglement, concurrence, negativity, convex roof extended negativity, entanglement of formation (the original one defined in [7]), tangle, squashed entanglement, Tsallis- $q$ entanglement and Rényi- $\alpha$ entanglement by $E_{d}, C, N, N_{c r}, E_{f}$, $\tau, E_{s q}, T_{q}$ and $R_{\alpha}$, respectively.

| $E$ | $\alpha(E)$ | System | Reference |
| :---: | :---: | :---: | :---: |
| $E_{d}$ | $\leq 1$ | any system | $[8]$ |
| $C$ | 2 | $2^{\otimes 3}$ | $[9]^{a}$ |
|  | $\leq \sqrt{2}$ | $2^{\otimes n}$ | $[10]$ |
|  | $\leq 2$ | $2 \otimes 2 \otimes 2^{m}$ | $[11]$ |
|  | $\leq 2$ | $2^{\otimes n}$ | $[11]$ |
|  | $>2$ | $3^{\otimes 3}$ | $[12]$ |
|  | $\leq 2$ | $2 \otimes 2 \otimes 4$ | $[13]$ |
| $N$ | $>3$ | $3 \otimes 2 \otimes 2$ | $[14]$ |
|  | $\leq 2$ | $2^{\otimes n}$ | $[15,16]^{b}$ |
|  | $\leq 2$ | $2 \otimes 2 \otimes 2^{m}$ | $[16]^{b}$ |
| $N_{c r}$ | $\leq 2$ | $d \otimes d \otimes d, d=2,3,4$ | $[1]$ |
| $E_{f}$ | $\leq 2$ | $2^{\otimes n}$ | $[17,16,18]$ |
|  | $\leq \sqrt{2}$ | $2^{\otimes n}$ | $[19,20]$ |
| $\tau$ | $>1$ | $2 \otimes 2 \otimes 2$ | $[10]$ |
| $E_{s q}$ | $\leq 1$ | $2 \otimes 2 \otimes 4$ | $[20,21]$ |
| $T_{q}, 2 \leq q \leq 3$ | $\leq 1$ | $a n y s^{\otimes y s t e m}$ | $[8]$ |
| $T_{q}$ | $\leq 2$ | $2 \otimes 2 \otimes 2^{m}$ | $[22]$ |
| $R_{\alpha}, \alpha \geq 2$ | $\leq 1$ | $2^{\otimes n}$ | $[23]^{c}$ |
| $R_{\alpha}, \alpha \geq \frac{\sqrt{7}-1}{2}$ | $\leq 2$ | $2^{\otimes n}$ | $[24,25]$ |

${ }^{a} \alpha(C) \leq 2$ was shown in [9], and the equality $\alpha(C)=2$ follows from the saturation by $W$ states (see, for example, Corollary 9).
${ }^{b}$ For pure states.
${ }^{c}$ For mixed states, and $q \in\left[\frac{5-\sqrt{13}}{2}, 2\right] \cup\left[3, \frac{5+\sqrt{13}}{2}\right]$.
${ }^{d}$ For mixed states.
such that the state $\rho^{A B C}$ has the form

$$
\begin{equation*}
\rho^{A B C}=\bigoplus_{j} q_{j} \rho^{A B_{j}^{L}} \otimes \rho^{B_{j}^{R} C}, \tag{19}
\end{equation*}
$$

where $q_{j}$ is a probability distribution.
Theorem 2. Let $E$ be an entanglement monotone. Then, E satisfies the disentangling condition (4) for all entangled Markov quantum states $\rho^{A B C}$ of the form (19).

Proof. Since local ancillary systems are free in entanglement theory, one can append an ancillary system $B^{\prime}$ that encodes the orthogonality of the subspaces $\mathcal{H}^{B_{j}^{L}} \otimes \mathcal{H}^{B_{j}^{R}}$. This can be done with an isometry that maps states in $\mathcal{H}^{B_{j}^{L}} \otimes \mathcal{H}^{B_{j}^{R}}$ to states in $\mathcal{H}^{B^{L}} \otimes \mathcal{H}^{B^{R}} \otimes|j\rangle\left\langle\left. j\right|^{B^{\prime}}\right.$, where systems $B^{L}$ and $B^{R}$ have dimensions $\max _{j} \operatorname{dim}\left(\mathcal{H}^{B_{j}^{L}}\right)$ and $\max _{j} \operatorname{dim}\left(\mathcal{H}^{B_{j}^{R}}\right)$, respectively. Therefore, w.l.o.g. we can write the above state as

$$
\begin{equation*}
\rho^{A B C}=\sum_{j} q_{j} \rho_{j}^{A B^{L}} \otimes|j\rangle\left\langle\left. j\right|^{B^{\prime}} \otimes \rho_{j}^{B^{R} C} .\right. \tag{20}
\end{equation*}
$$

Now, note that with any entanglement monotone $E$, the entanglement between $A$ and $B C$ is measured by:

$$
\begin{equation*}
E\left(\rho^{A \mid B C}\right)=\sum_{j} q_{j} E\left(\rho^{A \mid B_{j}^{L}} \otimes \rho^{B_{j}^{R} C}\right)=\sum_{j} q_{j} E\left(\rho^{A B_{j}^{L}}\right), \tag{21}
\end{equation*}
$$

where in the first equality we used the property (3) of entanglement monotones. Similarly, the entanglement between $A$ and $B$ is measured by

$$
\begin{equation*}
E\left(\rho^{A B}\right)=\sum_{j} q_{j} E\left(\rho^{A B_{j}^{L}} \otimes \rho^{B_{j}^{R}}\right)=\sum_{j} q_{j} E\left(\rho^{A B_{j}^{L}}\right) \tag{22}
\end{equation*}
$$

We therefore obtain (4) as long as $E\left(\rho^{A \mid B_{j}^{L}}\right)>0$ for some $j$ for which $q_{j}>0$. This completes the proof.

## C Monogamy of entanglement: pure vs mixed tripartite states

As discussed in the paper, it is typically hard to check if a measure of entanglement is monogamous since the condition in (4) involves mixed tripartite states. On the other hand, it is significantly simpler to check the disentangling condition if $\rho^{A B C}$ that appears in the disentangling condition (4) is pure. We say that $E$ is monogamous on pure states if for any pure tripartite state $\rho^{A B C}$ that satisfies (4), $E\left(\rho^{A C}\right)=0$. In the theorem below we shown that sometimes if $E$ is monogamous on pure states it is also monogamous on mixed states (that is, it is monogamous according to Def. 1).

For any entanglement monotone $E$ on the set of bipartite density matrices, $\mathcal{S}^{A B}$, we define a corresponding entanglement of formation measure, $E_{f}$ which is defined by the following convex roof extension:

$$
\begin{equation*}
E_{f}\left(\rho^{A B}\right) \equiv \min \sum_{j=1}^{n} p_{j} E\left(\left|\psi_{j}\right\rangle\left\langle\left.\psi_{j}\right|^{A B}\right),\right. \tag{23}
\end{equation*}
$$

where the minimum is taken over all pure state decompositions of $\rho^{A B}=\sum_{j=1}^{n} p_{j}\left|\psi_{j}\right\rangle\left\langle\left.\psi_{j}\right|^{A B}\right.$. Clearly, $E=E_{f}$ on pure bipartite states, but on mixed states they can be different, like the convex roof extended Negativity is different from the Negativity itself. Since we assume that $E$ is entanglement monotone it is convex so $E\left(\rho^{A B}\right) \leq E_{f}\left(\rho^{A B}\right)$ for all $\rho^{A B} \in \mathcal{S}^{A B}$. The corresponding entanglement of formation of
a given entanglement monotone, is itself an entanglement monotone, and has the following remarkable property.

Theorem 3. Let $E$ be an entanglement monotone, and let $E_{f}$ be its corresponding entanglement of formation (23). If $E_{f}$ is monogamous (according to Definition 1) on pure tripartite states in $\mathcal{H}^{A B C}$, then it is also monogamous on tripartite mixed states acting on $\mathcal{H}^{A B C}$.

Proof. Let $\rho^{A \mid B C}=\sum_{j} p_{j}\left|\psi_{j}\right\rangle\left\langle\left.\psi_{j}\right|^{A B C}\right.$ be a tripartite state acting on $\mathcal{H}^{A B C}$ with $\left\{p_{j},\left|\psi_{j}\right\rangle^{A B C}\right\}$ being the optimal decomposition such that

$$
\begin{equation*}
E_{f}\left(\rho^{A \mid B C}\right)=\sum_{j} p_{j} E_{f}\left(\left|\psi_{j}\right\rangle^{A \mid B C}\right) \tag{24}
\end{equation*}
$$

We also assume w.l.o.g. that $p_{j}>0$. Now, suppose $E_{f}\left(\rho^{A \mid B C}\right)=E_{f}\left(\rho^{A B}\right)$, and denote $\rho_{j}^{A B} \equiv$ $\operatorname{Tr}_{C}\left|\psi_{j}\right\rangle\left\langle\left.\psi_{j}\right|^{A B C}\right.$. Since discarding a subsystem can only decrease the entanglement, we get

$$
\begin{equation*}
\sum_{j} p_{j} E_{f}\left(\left|\psi_{j}\right\rangle^{A \mid B C}\right) \geq \sum_{j} p_{j} E_{f}\left(\rho_{j}^{A B}\right) \geq E_{f}\left(\rho^{A B}\right) \tag{25}
\end{equation*}
$$

where the last inequality follows from the convexity of $E_{f}$ and the fact that $\rho^{A B}=\sum_{j} p_{j} \rho_{j}^{A B}$. However, all the inequalities above must be equalities since $E_{f}\left(\rho^{A \mid B C}\right)=E_{f}\left(\rho^{A B}\right)$. In particular, we get

$$
\begin{equation*}
\sum_{j} p_{j} E_{f}\left(\left|\psi_{j}\right\rangle^{A \mid B C}\right)=\sum_{j} p_{j} E_{f}\left(\rho_{j}^{A B}\right) \tag{26}
\end{equation*}
$$

This in turn implies that $E_{f}\left(\left|\psi_{j}\right\rangle^{A \mid B C}\right)=E_{f}\left(\rho_{j}^{A B}\right)$ for each $j$ since $E_{f}\left(\left|\psi_{j}\right\rangle^{A \mid B C}\right) \geq E_{f}\left(\rho_{j}^{A B}\right)$ for each $j$ (i.e. tracing out subsystem cannot increase entanglement). Since we assume that $E_{f}$ is monogamous on pure tripartite states, we conclude that for each $j, E_{f}\left(\rho_{j}^{A C}\right)=0$, where $\rho_{j}^{A C} \equiv \operatorname{Tr}_{B}\left|\psi_{j}\right\rangle\left\langle\left.\psi_{j}\right|^{A B C}\right.$. Hence, $E_{f}\left(\rho^{A C}\right)=0$ since $\rho^{A C}=\sum_{j} p_{j} \rho_{j}^{A C}$ and $E_{f}$ is convex.

## D Entanglement of Collaboration and Monogamy of Entanglement

Monogamy of entanglement is closely related to entanglement of collaboration. Given a measure of bipartite entanglement $E$, its corresponding entanglement of collaboration, $E_{c}^{A B \mid C}$, is a measure of entanglement on tripartite mixed states, $\rho^{A B C}$, given by [32]:

$$
\begin{equation*}
E_{c}^{A B \mid C}\left(\rho^{A B C}\right)=\max \sum_{j=1}^{n} p_{j} E\left(\rho_{j}^{A B}\right) \tag{27}
\end{equation*}
$$

where the maximum is taken over all tripartite LOCC protocols yielding the bipartite state $\rho_{j}^{A B}$ with probability $p_{j}$. The following theorem demonstrates the connection between the disentangling condition and entanglement of collaboration.

Theorem 4. Let $E$ be an entanglement monotone on bipartite mixed states, and let $\rho^{A B C}$ be a (possibly mixed) tripartite state satisfying the disentangling condition (4). Then,

$$
\begin{equation*}
E\left(\rho^{A B}\right)=E_{c}^{A B \mid C}\left(\rho^{A B C}\right) \tag{28}
\end{equation*}
$$

Proof. Let $\left\{\rho_{j}^{A B}, p_{j}\right\}$ be the optimal ensemble in (8) obtained by LOCC on $\rho^{A B C}$. Since $E$ is a bipartite entanglement monotone, it does not increase on average:

$$
\begin{equation*}
E\left(\rho^{A \mid B C}\right) \geq \sum_{j} p_{j} E\left(\rho_{j}^{A B}\right)=E_{c}^{A B \mid C}\left(\rho^{A B C}\right) . \tag{29}
\end{equation*}
$$

On the other hand, by definition $E\left(\rho^{A B}\right) \leq E_{c}^{A B \mid C}\left(\rho^{A B C}\right)$, so that together with (4) we get (9).
Entanglement of collaboration is different from entanglement of assistance, $E_{a}$, in which the optimization in (27) is restricted to LOCC of the following form: Charlie (system C) performs a measurement, and communicates the outcome $j$ to Alice and Bob. In [32] the following LOCC protocol was considered: Alice performs a measurement, then sending the outcome to Charlie, and then Charlie performs his measurement, and sends back the result to Alice and Bob. It was shown that in such a scenario it is possible to increase the average entanglement between systems $A$ and $B$ to a value beyond the average entanglement that can be achieved if only Charlie performed a measurement. Therefore, $E_{c}^{A B \mid C}$ can be strictly larger than $E_{a}$, and in general, $E_{c}^{A B \mid C} \geq E_{a}$. However, if $\rho^{A B C}$ satisfies the disentangling condition then we must have $E_{c}^{A B \mid C}=E_{a}$. Indeed, if $\rho^{A B C}$ satisfies (4) we get

$$
\begin{equation*}
E_{a}\left(\rho^{A B C}\right) \geq E\left(\rho^{A B}\right)=E\left(\rho^{A \mid B C}\right)=E_{c}^{A B \mid C}\left(\rho^{A B C}\right) \tag{30}
\end{equation*}
$$

Therefore, one can replace $E_{c}^{A B \mid C}$ in (28) with $E_{a}$, which may be convenient since $E_{a}$ is somewhat a simpler measure than $E_{c}^{A B \mid C}$. Note however that we left $E_{c}^{A B \mid C}$ in (28) since $E\left(\rho^{A \mid B C}\right)=E_{c}^{A B \mid C}\left(\rho^{A B C}\right)$ implies $E\left(\rho^{A \mid B C}\right)=E_{a}\left(\rho^{A B C}\right)$ but not vice versa.

## D. 1 When Entanglement of Formation equals Entanglement of Assistance?

An immediate consequence of Theorem 4 above is that if $\rho^{A B C}$ is a pure state that satisfies the disentangling condition then the entanglement of formation of $\rho^{A B}$ must be equal to its entanglement of assistance.

Corollary 5. Let $E$ be an entanglement monotone on bipartite states, and let $\rho^{A B C}$ be a pure tripartite state satisfying the disentangling condition (4). Then,

$$
\begin{equation*}
E\left(\rho^{A B}\right)=E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right), \tag{31}
\end{equation*}
$$

where the entanglement of formation, $E_{f}$, is defined in (23), and the entanglement of assistance, $E_{a}$, is also defined as in (23) but with a maximum replacing the minimum.

Proof. The proof follows straightforwardly from Theorem (4) recalling that

$$
\begin{equation*}
E\left(\rho^{A B}\right) \leq E_{f}\left(\rho^{A B}\right) \leq E_{a}\left(\rho^{A B}\right) \leq E_{c}^{A B \mid C}\left(\rho^{A B C}\right)=E\left(\rho^{A \mid B C}\right) \tag{32}
\end{equation*}
$$

where the first inequality follows from the fact that $E$ is an entanglement monotone, and the last equality from Theorem (4). Therefore, all the inequalities above are equalities since we assume the disentangling condition $E\left(\rho^{A \mid B C}\right)=E\left(\rho^{A B}\right)$. This completes the proof.

In the following theorem we show that the equality between the entanglement of formation and entanglement of assistance is a property of the support space of the the bipartite state in question.

Theorem 6. Let $E$ be a measure of bipartite entanglement, and let $\rho^{A B} \in \mathcal{S}^{A B}$ be a bipartite state with support subspace $\operatorname{supp}\left(\rho^{A B}\right)$. If $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ then for any $\sigma^{A B} \in \mathcal{S}^{A B}$ with $\operatorname{supp}\left(\sigma^{A B}\right) \subseteq$ $\operatorname{supp}\left(\rho^{A B}\right)$ we have $E_{f}\left(\sigma^{A B}\right)=E_{a}\left(\sigma^{A B}\right)$.

Proof. In the proof we use some of the ideas introduced in [33]. Suppose $\rho^{A B}=\sum_{j} p_{j} \rho_{j}^{A B}$, where $\left\{p_{j}\right\}$ are (non-zero) probabilities and $\rho_{j}^{A B} \in \mathcal{S}^{A B}$. Then, the condition $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ together with the convexity (concavity) of $E_{f}\left(E_{a}\right)$ gives

$$
\begin{equation*}
\sum_{j} p_{j} E_{f}\left(\rho_{j}^{A B}\right) \geq E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right) \geq \sum_{j} p_{j} E_{a}\left(\rho_{j}^{A B}\right) . \tag{33}
\end{equation*}
$$

But since for all $j$ we also have $E_{f}\left(\rho_{j}^{A B}\right) \leq E_{a}\left(\rho_{j}^{A B}\right)$, we get that $E_{f}\left(\rho_{j}^{A B}\right)=E_{a}\left(\rho_{j}^{A B}\right)$ for all $j$. Let $\mathcal{F}\left(\rho^{A B}\right)$ be the set of all density matrices in $\mathcal{S}^{A B}$ that appear in a convex decomposition of $\rho^{A B}$. In convex analysis, $\mathcal{F}\left(\rho^{A B}\right)$ is called a face of $\mathcal{S}^{A B}$. Now, from the argument above we have that if $\sigma^{A B} \in$ $\mathcal{F}\left(\rho^{A B}\right)$ then $E_{f}\left(\sigma^{A B}\right)=E_{a}\left(\sigma^{A B}\right)$. On the other hand, for any $\sigma^{A B} \in \mathcal{S}^{A B}$ with the property that $\operatorname{supp}\left(\sigma^{A B}\right) \subset \operatorname{supp}\left(\rho^{A B}\right)$ there exists a small enough $\epsilon>0$ such that $\rho^{A B}-\epsilon \sigma^{A B} \geq 0$. Denoting by $\gamma^{A B} \equiv\left(\rho^{A B}-\epsilon \sigma^{A B}\right) /(1-\epsilon)$ we get that $\gamma^{A B} \in \mathcal{S}^{A B}$ and $\rho^{A B}$ can be expressed as the convex combination $\rho^{A B}=\epsilon \sigma^{A B}+(1-\epsilon) \gamma^{A B}$. We therefore must have $E_{f}\left(\sigma^{A B}\right)=E_{a}\left(\sigma^{A B}\right)$. This completes the proof.

In the proof above we called $\mathcal{F}\left(\rho^{A B}\right)$ a face. A face $\mathcal{F}$ of $\mathcal{S}^{A B}$ is a convex subset of $\mathcal{S}^{A B}$ that satisfies the following property: if $t \rho_{1}+(1-t) \rho_{2} \in \mathcal{F}$ for some $\rho_{1}, \rho_{2} \in \mathcal{S}^{A B}$ and $0<t<1$, then $\rho_{1}, \rho_{2} \in \mathcal{F}$. To see that $\mathcal{F}\left(\rho^{A B}\right)$ is a face of $\mathcal{S}^{A B}$, we first show that it is convex. Indeed, let $\sigma \equiv t \sigma_{1}+(1-t) \sigma_{2}$ for some $t \in[0,1]$ and $\sigma_{1}, \sigma_{2} \in \mathcal{F}\left(\rho^{A B}\right)$. Since $\sigma_{1}, \sigma_{2} \in \mathcal{F}\left(\rho^{A B}\right)$ there exists $p, q \in(0,1]$ and $\gamma_{1}, \gamma_{2} \in \mathcal{S}^{A B}$ such that

$$
\begin{equation*}
\rho^{A B}=p \sigma_{1}+(1-p) \gamma_{1}=q \sigma_{2}+(1-q) \gamma_{2} . \tag{34}
\end{equation*}
$$

The first equality implies that $\frac{t}{p} \rho^{A B}=t \sigma_{1}+\frac{t(1-p)}{p} \gamma_{1}$, and the second equality gives $\frac{1-t}{q} \rho^{A B}=(1-t) \sigma_{2}+$ $\frac{(1-t)(1-q)}{q} \gamma_{2}$. Therefore,

$$
\begin{equation*}
\left(\frac{t}{p}+\frac{1-t}{q}\right) \rho^{A B}=\sigma^{A B}+\frac{t(1-p)}{p} \gamma_{1}+\frac{(1-t)(1-q)}{q} \gamma_{2} . \tag{35}
\end{equation*}
$$

After dividing by $\frac{t}{p}+\frac{1-t}{q}$, we can see that $\sigma^{A B}$ appears in a convex combination of $\rho^{A B}$. Therefore, $\mathcal{F}\left(\rho^{A B}\right)$ is convex. To complete the proof that $\mathcal{F}\left(\rho^{A B}\right)$ is a face, note that if $\tau \equiv t \rho_{1}+(1-t) \rho_{2} \in \mathcal{F}$ for some $\rho_{1}, \rho_{2} \in \mathcal{S}^{A B}$ and $0<t<1$ then clearly $\rho_{1}, \rho_{2} \in \mathcal{F}\left(\rho^{A B}\right)$.

Note that the condition $\operatorname{supp}\left(\sigma^{A B}\right) \subseteq \operatorname{supp}\left(\rho^{A B}\right)$ is equivalent to $\sigma^{A B} \in \mathcal{F}\left(\rho^{A B}\right)$. The precise form of the support space of a bipartite state $\rho^{A B}$ that satisfies $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ depends on the measure of entanglement $E$. In the following sections we find it precisely for the case where $E$ is the G-concurrence, and we use it to show that the G-concurrence is monogamous.

## D. 2 Monogamy of the G-concurrence

Any pure bipartite state, $|x\rangle \in \mathbb{C}^{d} \otimes \mathbb{C}^{d}$, can be written as:

$$
\begin{equation*}
|x\rangle=X \otimes I\left|\phi_{+}\right\rangle \quad \text { where } \quad\left|\phi_{+}\right\rangle=\sum_{i=1}^{d}|i\rangle^{A}|i\rangle^{B}, \tag{36}
\end{equation*}
$$

and $X$ is a $d \times d$ complex matrix. The relation above between a complex matrix $X \in \mathcal{M}_{d}(\mathbb{C})$ and a bipartite vector $|x\rangle \in \mathbb{C}^{d} \otimes \mathbb{C}^{d}$ defines an isomorphism between $\mathcal{M}_{d}(\mathbb{C})$ and $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$. Using this isomorphism, in the remaining of this section we will view interchangeably the support of a density matrix both as a subspace of $\mathcal{M}_{d}(\mathbb{C})$ or as a subspace of $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$, depending on the context. We will use capital letters $X, Y, Z, W$
for matrices in $\mathcal{M}_{d}(\mathbb{C})$, and use lower case letters $|x\rangle,|y\rangle,|z\rangle,|w\rangle$ to denote their corresponding bipartite vectors.

With these notations, the G-concurrence of $|x\rangle$, which is the geometric mean of the Schmidt coefficients of $|x\rangle$, can be expressed as:

$$
\begin{equation*}
G(|x\rangle)=|\operatorname{det}(X)|^{2 / d} \tag{37}
\end{equation*}
$$

and for mixed states it is defined in terms of the convex roof extension; that is, $G\left(\rho^{A B}\right):=G_{f}\left(\rho^{A B}\right)$ for any $\rho^{A B} \in \mathcal{S}^{A B}$. Note that the G-concurrence is homogeneous, and in particular, $G(c|x\rangle)=|c|^{2} G(|x\rangle)$. We start in proving the following Lemma:
Lemma: Let $\rho^{A B}$ be a bipartite density matrix acting on $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$ with $G\left(\rho^{A B}\right)>0$. Then, there exists a pure state decomposition of $\rho^{A B}$ with the following properties:

$$
\begin{equation*}
\rho^{A B}=\sum_{j=1}^{r}\left|w_{j}\right\rangle\left\langle w_{j}\right|, \quad G\left(\left|w_{j}\right\rangle\right)=0, \quad \forall j=2, \ldots, r, \tag{38}
\end{equation*}
$$

where $r$ is the rank of $\rho^{A B}$ and $\left|w_{j}\right\rangle$ are sub-normalized vectors in $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$ (i.e. vectors with norm no greater than 1).
Proof. Let $\rho^{A B}=\sum_{j=1}^{r}\left|x_{j}\right\rangle\left\langle x_{j}\right|$ be the spectral decomposition of $\rho^{A B}$ with $\left|x_{j}\right\rangle$ being the sub-normalized eigenvectors of $\rho^{A B}$. Clearly, $G\left(\left|x_{j}\right\rangle\right)>0$ for at least one $j$. Therefore, w.l.o.g. we assume $G\left(\left|x_{1}\right\rangle\right)>0$. Let $\mathcal{K}_{2} \subset \operatorname{supp}\left(\rho^{A B}\right)$ be the two dimensional subspace spanned by $X_{1}$ and $X_{2}$. We first show that $\mathcal{K}_{2}$ contains a matrix with zero determinant. Indeed, if $\operatorname{det}\left(X_{2}\right)=0$ we are done. Otherwise, for any $\lambda \in \mathbb{C}$ we get

$$
\begin{align*}
G\left(\left|x_{1}\right\rangle+\lambda\left|x_{2}\right\rangle\right) & =\left|\operatorname{det}\left(X_{1}+\lambda X_{2}\right)\right|^{2 / d} \\
& =\left|\operatorname{det}\left(X_{2}\right)\right|^{2 / d}\left|\operatorname{det}\left(X_{1} X_{2}^{-1}+\lambda I\right)\right|^{2 / d} \tag{39}
\end{align*}
$$

Note that $\operatorname{det}\left(X_{1} X_{2}^{-1}+\lambda I\right)$ is a polynomial of degree $d$ in $\lambda$ and must have at least one complex root. Therefore, there exists $\lambda=\lambda_{0} \neq 0$ such that $\operatorname{det}\left(X_{1}+\lambda_{0} X_{2}\right)=0$. This completes the assertion that $\mathcal{K}_{2}$ contains a matrix with zero determinant.

Next, denote $a \equiv \frac{1}{\sqrt{1+\left|\lambda_{0}\right|^{2}}}$ and $b \equiv \frac{\lambda_{0}}{\sqrt{1+\left|\lambda_{0}\right|^{2}}}$, so that $|a|^{2}+|b|^{2}=1$, and the vectors $\left|w_{2}\right\rangle \equiv a\left|x_{1}\right\rangle+b\left|x_{2}\right\rangle$ and $|y\rangle \equiv \bar{b}\left|x_{1}\right\rangle-\bar{a}\left|x_{2}\right\rangle$ satisfy

$$
\begin{equation*}
\left|x_{1}\right\rangle\left\langle x_{1}\right|+\left|x_{2}\right\rangle\left\langle x_{2}\right|=|y\rangle\langle y|+\left|w_{2}\right\rangle\left\langle w_{2}\right| \tag{40}
\end{equation*}
$$

with $G\left(\left|w_{2}\right\rangle\right)=0$ (by construction, $W_{2}$ is proportional to $X_{1}+\lambda_{0} X_{2}$ and therefore has zero determinant). We can therefore write

$$
\begin{equation*}
\rho^{A B}=\left|w_{2}\right\rangle\left\langle w_{2}\right|+|y\rangle\langle y|+\sum_{j=3}^{r}\left|x_{j}\right\rangle\left\langle x_{j}\right| . \tag{41}
\end{equation*}
$$

Now, if $G(|y\rangle)=0$ then we denote it as $\left|w_{3}\right\rangle$ and pick from $\left\{\left|x_{j}\right\rangle\right\}_{j=3}^{r}$ another state that does not have a vanishing G-concurrence. We therefore assume $G(|y\rangle) \neq 0$ and from the same arguments as above we conclude that $|y\rangle\langle y|+\left|x_{3}\right\rangle\left\langle x_{3}\right|=|\tilde{y}\rangle\langle\tilde{y}|+\left|w_{3}\right\rangle\left\langle w_{3}\right|$ for some vectors $|\tilde{y}\rangle$ and $\left|w_{3}\right\rangle$, with $G\left(\left|w_{3}\right\rangle\right)=0$. By replacing $|y\rangle$ and $\left|x_{3}\right\rangle$ with $|\tilde{y}\rangle$ and $\left|w_{3}\right\rangle$, and repeating the process we arrive at the desired decomposition of $\rho^{A B}$.

In the following theorem we denote by $\mathcal{N} \subset \mathcal{M}_{d}(\mathbb{C})$, the nilpotent cone, consisting of all $d \times d$ nilpotent complex matrices (i.e. $X \in \mathcal{N}$ iff $X^{k}=0$ for some $1 \leq k \leq d$ ). While the set $\mathcal{N}$ is not a vector space, it contains subspaces. For example, the set $\mathcal{T}$ of all strictly upper triangular matrices (i.e. upper triangular
matrices with zeros on the diagonal) is a $d(d-1) / 2$-dimensional subspace in $\mathcal{N}$. From Gerstenhaber's theorem [34] it follows that the largest dimension of a subspace in $\mathcal{N}$ is $d(d-1) / 2$, and if a subspace $\mathcal{N}_{0} \subset \mathcal{N}$ has this maximal dimension, then it must be similar to $\mathcal{T}$ (i.e. their exists an invertible matrix $S$ such that $\mathcal{N}_{0}=S \mathcal{T} S^{-1}$ ).

Theorem 7. Let $\rho^{A B}$ be a bipartite density matrix acting on $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$, and suppose it has rank $r>1$. Then,

$$
\begin{equation*}
G\left(\rho^{A B}\right)=G_{a}\left(\rho^{A B}\right)>0 \tag{42}
\end{equation*}
$$

if and only if $r \leq 1+d(d-1) / 2$, and there exists a full rank matrix $X \in \mathcal{M}_{d}(\mathbb{C})$ and a subspace $\mathcal{N}_{0} \subset \mathcal{N}$ with $\operatorname{dim}\left(\mathcal{N}_{0}\right)=r-1$ such that

$$
\begin{equation*}
\operatorname{supp}\left(\rho^{A B}\right)=X \mathcal{K} \equiv\{X Y \mid Y \in \mathcal{K}\} \tag{43}
\end{equation*}
$$

where $\mathcal{K} \equiv\{I\} \oplus \mathcal{N}_{0} \equiv\left\{c I+N \mid c \in \mathbb{C} ; N \in \mathcal{N}_{0}\right\}$.
Remark 3. In the statement of Theorem 7 of the main text, we viewed $\operatorname{supp}\left(\rho^{A B}\right)$ as a subspace of $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$. Here, for convenience of the proof, we use the isomorphism (36) between $\mathcal{M}_{d}(\mathbb{C})$ and $\mathbb{C}^{d} \otimes \mathbb{C}^{d}$, and view the $\operatorname{supp}\left(\rho^{A B}\right)$ as a subspace of $\mathcal{M}_{d}(\mathbb{C})$. The matrix $X$ in (43) is related to the bipartite state $|x\rangle^{A B}$ in (13) via (36). Recall also that a nilpotent matrix has a zero trace, and is orthogonal to the identity matrix in the Hilbert-Schmidt inner product. This is consistent with the direct sum in the definition of $\mathcal{K}$.

Proof. Suppose there exist $X$ and $\mathcal{N}_{0}$ as above such that $\operatorname{supp}\left(\rho^{A B}\right) \subseteq X \mathcal{K}$, and let

$$
\begin{equation*}
\rho^{A B}=\sum_{j=1}^{r}\left|w_{j}\right\rangle\left\langle w_{j}\right| \tag{44}
\end{equation*}
$$

where $\left|w_{j}\right\rangle$ are the sub-normalized vectors given in (38); i.e. $G\left(\left|w_{j}\right\rangle\right)=0$ for $j>1$. Since $W_{j} \in X \mathcal{K}$, there exist constants $c_{j} \in \mathbb{C}$ and matrices $Z_{j} \in \mathcal{N}_{0}$ (corresponding to some sub-normalized vectors $\left|z_{j}\right\rangle$ ) such that

$$
\begin{equation*}
W_{j}=c_{j} X+X Z_{j} \tag{45}
\end{equation*}
$$

For $j>1$ we have

$$
\begin{equation*}
0=G\left(\left|w_{j}\right\rangle\right)=\left|\operatorname{det}\left(c_{j} X+X Z_{j}\right)\right|^{2 / d}=G(|x\rangle)\left|\operatorname{det}\left(c_{j} I+Z_{j}\right)\right|^{2 / d} \tag{46}
\end{equation*}
$$

so that $\operatorname{det}\left(c_{j} I+Z_{j}\right)=0$ since $G(|x\rangle)>0$. But since $Z_{j}$ is a nilpotent matrix $\operatorname{det}\left(c_{j} I+Z_{j}\right)=\left(c_{j}\right)^{d}$. Hence, $c_{j}=0$ for $j>1$ and we denote by $c \equiv c_{1}$. Note that $c \neq 0$ since otherwise we will get $G\left(\rho^{A B}\right)=0$.

Therefore, the average G-concurrence of decomposition (44) is given by

$$
\begin{equation*}
\sum_{j=1}^{r} G\left(\left|w_{j}\right\rangle\right)=G\left(\left|w_{1}\right\rangle\right)=\left|\operatorname{det}\left(c X+X Z_{1}\right)\right|^{2 / d}=G(|x\rangle)|c|^{2} \tag{47}
\end{equation*}
$$

Next, let

$$
\begin{equation*}
\rho^{A B}=\sum_{k=1}^{m}\left|y_{k}\right\rangle\left\langle y_{k}\right| \tag{48}
\end{equation*}
$$

be another pure state decomposition of $\rho^{A B}$ with $m \geq r$ and $\left\{\left|y_{k}\right\rangle\right\}$ some sub-normalized states. Then, there exists an $m \times r$ isometry $U=\left(u_{k j}\right)$, i.e. $U^{\dagger} U=I_{r}$, such that

$$
\begin{equation*}
\left|y_{k}\right\rangle=\sum_{j=1}^{r} u_{k j}\left|w_{j}\right\rangle . \tag{49}
\end{equation*}
$$

Using the form (45) with $c_{1} \equiv c$ and $c_{2}=c_{3}=\ldots=c_{r}=0$, we get

$$
\begin{equation*}
Y_{k}=\sum_{j=1}^{r} u_{k j} W_{j}=X\left(\sum_{j=1}^{r} u_{k j}\left(c \delta_{1 j} I+Z_{j}\right)\right) \equiv X\left(u_{k 1} c I+N_{k}\right), \tag{50}
\end{equation*}
$$

where $N_{k} \equiv \sum_{j=1}^{r} u_{k j} Z_{j} \in \mathcal{N}_{0}$ are nilpotent matrices. Consequently,

$$
\begin{equation*}
G\left(\left|y_{k}\right\rangle\right)=\left|\operatorname{det}\left(X\left(u_{k 1} c I+N_{k}\right)\right)\right|^{2 / d}=\left|u_{k 1}\right|^{2}|c|^{2} G(|x\rangle), \tag{51}
\end{equation*}
$$

where we used the fact that $\operatorname{det}\left(u_{k 1} c I+N_{k}\right)=\left(u_{k 1} c\right)^{d}$ since $N_{k}$ is nilpotent. Note that since the matrix $U$ is an isometry, $\sum_{k=1}^{m}\left|u_{k 1}\right|^{2}=1$. Therefore,

$$
\begin{equation*}
\sum_{k=1}^{m} G\left(\left|y_{k}\right\rangle\right)=|c|^{2} G(|x\rangle)=\sum_{j=1}^{r} G\left(\left|w_{j}\right\rangle\right), \tag{52}
\end{equation*}
$$

where the last equality follows from (47). Therefore, all pure states decomposition of $\rho^{A B}$ have the same average G-concurrence so that $G\left(\rho^{A B}\right)=G_{a}\left(\rho^{A B}\right)$.

To prove the converse, suppose $G\left(\rho^{A B}\right)=G_{a}\left(\rho^{A B}\right)>0$, and consider the decomposition (38) of $\rho^{A B}$ as in the Lemma above. Consider the unnormalized state $\sigma^{A B} \equiv \rho^{A B}-\left|w_{1}\right\rangle\left\langle w_{1}\right|$. The state $\sigma^{A B}=\sum_{j=2}^{r}\left|w_{j}\right\rangle\left\langle w_{j}\right|$ has zero $G$ concurrence since $G\left(\left|w_{j}\right\rangle\right)=0$ for all $j=2, \ldots, r$. Consider another decomposition of $\sigma^{A B}=\sum_{j=2}^{r}\left|y_{j}\right\rangle\left\langle y_{j}\right|$. Since $G\left(\rho^{A B}\right)=G_{a}\left(\rho^{A B}\right)$ we must have that $G\left(\left|y_{j}\right\rangle\right)=0$ for all $j=2, \ldots, r$. Otherwise, the decomposition $\rho^{A B}=\left|w_{1}\right\rangle\left\langle w_{1}\right|+\sum_{k=2}^{r}\left|y_{k}\right\rangle\left\langle y_{k}\right|$ will have a higher average G-concurrence than the decomposition in (38). But since the $\left\{\left|y_{j}\right\rangle\right\}$ decomposition was arbitrary, we conclude that all the states in the subspace $\mathcal{W} \equiv \operatorname{span}\left\{\left|w_{2}\right\rangle, \ldots,\left|w_{r}\right\rangle\right\}$ have zero G-concurrence.

We now denote $W_{1} \equiv X$ and for $j=2, \ldots, r$ we set $N_{j} \equiv X^{-1} W_{j}$. Note that with these notations, $\operatorname{supp}\left(\rho^{A B}\right)=\operatorname{span}\left\{X, X N_{2}, X N_{3}, \ldots, X N_{r}\right\}$. Since $X$ is invertible, and $\operatorname{det}(W)=0$ for any matrix $W$ in the span of $W_{2}, \ldots, W_{r}$, we also have $\operatorname{det}(N)=0$ for any $N$ in the span of $N_{2}, \ldots, N_{r}$. Let $N \in \operatorname{span}\left\{N_{2}, \ldots, N_{r}\right\}$ be a fixed matrix, and consider the two dimensional subspace $\mathcal{W}_{2} \equiv\{X, X N\} \subset$ $\operatorname{supp}\left(\rho^{A B}\right)$. We first show that $\mathcal{W}_{2}$ does not contain a matrix with zero determinant that is linearly independent of $W \equiv X N$. Indeed, if there exists a normalized matrix $Z \in \mathcal{W}$ such that $\operatorname{det}(Z)=\operatorname{det}(W)=0$ with $W, Z$ being linearly independent, then $\mathcal{W}_{2}=\operatorname{span}\{W, Z\}$, and the rank 2 density matrix

$$
\begin{equation*}
\sigma^{A B}=|w\rangle\langle w|+|z\rangle\langle z| \tag{53}
\end{equation*}
$$

must have zero G-concurrence (recall that $G(|w\rangle)=G(|z\rangle)=0$ ). On the other hand, since $|x\rangle \in \mathcal{W}_{2}=$ $\operatorname{supp}\left(\sigma^{A B}\right)$, the density matrix $\sigma^{A B}$ must have a pure state decomposition containing $|x\rangle$. Since $G(|x\rangle)>0$ this decomposition does not have a zero average G-concurrence. Therefore, $0=G\left(\sigma^{A B}\right)<G_{a}\left(\sigma^{A B}\right)$. However, from Theorem 6 any density matrix $\sigma^{A B}$ with a support $\operatorname{supp}\left(\sigma^{A B}\right)=\mathcal{W}_{2} \subset \operatorname{supp}\left(\rho^{A B}\right)$ has the same average G-concurrence for all pure state decompositions. We therefore get a contradiction with Theorem 6, and thereby prove the assertion that $\mathcal{W}_{2}$ does not contain another matrix with zero determinant that is linearly independent of $W$.

Since $W=X N$ is the only matrix in $\mathcal{W}_{2}$ with zero determinant (up to multiplication by a constant), we must have that for any $\lambda \neq 0$

$$
\begin{equation*}
0 \neq \operatorname{det}(X+\lambda W)=\operatorname{det}\left(X^{-1}\right) \operatorname{det}\left(I+\lambda X^{-1} W\right)=\lambda^{d} \operatorname{det}\left(X^{-1}\right) \operatorname{det}\left(\frac{1}{\lambda} I+N\right) . \tag{54}
\end{equation*}
$$

Setting $t \equiv \frac{1}{\lambda}$ we conclude that the polynomial

$$
\begin{equation*}
f(t) \equiv \operatorname{det}(t I+N) \tag{55}
\end{equation*}
$$

is never zero for $t \neq 0$. On the other hand, $f(t)$ is a polynomial of degree $d$ and the coefficient of $t^{d}$ is one (note that it is the characteristic polynomial of $-N$ ). But since $t=0$ is the only root of $f(t)$, we must have $f(t)=t^{d}$; that is,

$$
\begin{equation*}
\operatorname{det}(t I+N)=t^{n}, \quad \forall t \in \mathbb{C} \tag{56}
\end{equation*}
$$

Therefore, $N$ must be a nilpotent matrix. Since $N$ was arbitrary, it follows that the subspace $\mathcal{N}_{0} \equiv$ $\operatorname{span}\left\{N_{2}, \ldots, N_{r}\right\}$ is a subspace of nilpotent matrices. This completes the proof.

The G-concurrence is defined in (37) on pure bipartite states with the same local dimension. For a pure bipartite state $|\psi\rangle^{A B} \in \mathcal{H}^{A} \otimes \mathcal{H}^{B}$ with $\operatorname{dim}\left(\mathcal{H}^{A}\right)<\operatorname{dim}\left(\mathcal{H}^{B}\right)$ it is defined by:

$$
\begin{equation*}
G\left(|\psi\rangle^{A B}\right)=\left(\operatorname{det}\left(\rho^{A}\right)\right)^{1 / d_{A}} \tag{57}
\end{equation*}
$$

where $\rho^{A} \equiv \operatorname{Tr}_{B}\left(|\psi\rangle\left\langle\left.\psi\right|^{A B}\right)\right.$ is the reduced density matrix, and $d_{A} \equiv \operatorname{dim}\left(\mathcal{H}^{A}\right)$. With this extended definition, we have the following result:

Corollary 8. Let $\left|\psi^{A B C}\right\rangle \in \mathbb{C}^{d} \otimes \mathbb{C}^{d} \otimes \mathbb{C}^{n}$ be a pure tripartite state, with bipartite marginal $\rho^{A B}$. If

$$
\begin{equation*}
G\left(\psi^{A \mid B C}\right)=G\left(\rho^{A B}\right)>0, \tag{58}
\end{equation*}
$$

then $|\psi\rangle^{A B C}=|\chi\rangle^{A B}|\varphi\rangle^{C}$ for some bipartite state $|\chi\rangle^{A B} \in \mathbb{C}^{d} \otimes \mathbb{C}^{d}$ and a vector $|\varphi\rangle^{C} \in \mathbb{C}^{n}$.
Proof. Since $|\psi\rangle^{A B C}$ satisfies the disentangling condition, we get from Corollary 5 that $G\left(\rho^{A B}\right)=$ $G_{a}\left(\rho^{A B}\right)>0$. Therefore, from the theorem above there exists a pure state decomposition of the marginal state $\rho^{A B}$ consisting of sub-normalized bipartite states $\left\{\left|w_{j}\right\rangle\right\}$ as in (38), with the form

$$
\begin{align*}
& \left|w_{1}\right\rangle=\left(c X+X N_{1}\right) \otimes I^{B}\left|\phi_{+}\right\rangle^{A B} \\
& \left|w_{j}\right\rangle=X N_{j} \otimes I^{B}\left|\phi_{+}\right\rangle^{A B}, \quad \forall j=2, \ldots, r, \tag{59}
\end{align*}
$$

where $X$ is a full rank matrix, $N_{j} \in \mathcal{N}_{0}$, and $c \in \mathbb{C}$. Since all decompositions have the same average G-concurrence, we have

$$
\begin{equation*}
G\left(\rho^{A B}\right)=\sum_{j=1}^{r} G\left(\left|w_{j}\right\rangle\right)=|\operatorname{det}(X)|^{2 / d} c . \tag{60}
\end{equation*}
$$

On the other hand, using the property that $\operatorname{Tr}_{B}\left|\phi_{+}\right\rangle\left\langle\left.\phi_{+}\right|^{A B}=I^{A}\right.$, we get from (59) that the marginal of $\rho^{A B}=\sum_{j=1}^{r}\left|w_{j}\right\rangle\left\langle\left. w_{j}\right|^{A B}\right.$ is given by:

$$
\begin{equation*}
\rho^{A}=X\left(c I+N_{1}\right)\left(\bar{c} I+N_{1}^{\dagger}\right) X^{\dagger}+X \sum_{j=2}^{r} N_{j} N_{j}^{\dagger} X^{\dagger} \tag{61}
\end{equation*}
$$

Therefore,

$$
\begin{equation*}
G\left(\left|\psi^{A \mid B C}\right\rangle\right)=\left(\operatorname{det}\left(\rho^{A}\right)\right)^{1 / d}=|\operatorname{det}(X)|^{2 / d}\left(\operatorname{det}\left(\left|c I+N_{1}\right|^{2}+\sum_{j=2}^{r}\left|N_{j}\right|^{2}\right)\right)^{1 / d} \tag{62}
\end{equation*}
$$

where we used the notation $|A| \equiv \sqrt{A A^{\dagger}}$ for any $d \times d$ matrix $A$. Hence, the condition $G\left(\left|\psi^{A \mid B C}\right\rangle\right)=$ $G\left(\rho^{A B}\right)=|\operatorname{det}(X)|^{2 / d}|c|^{2}$ gives

$$
\begin{equation*}
|c|^{2}=\left(\operatorname{det}\left(\left|c I+N_{1}\right|^{2}+\sum_{j=2}^{r}\left|N_{j}\right|^{2}\right)\right)^{1 / d} \tag{63}
\end{equation*}
$$

Since $c \neq 0$ the matrix $c I+N_{1}$ is invertible. Denoting

$$
\begin{equation*}
A \equiv\left|c I+N_{1}\right| \text { and } B \equiv A^{-1}\left(\sum_{j=2}^{r}\left|N_{j}\right|^{2}\right) A^{-1} \tag{64}
\end{equation*}
$$

we get that (63) can be expressed as

$$
\begin{equation*}
|c|^{2}=[\operatorname{det}(A(I+B) A)]^{1 / d}=\left[\operatorname{det}\left(A^{2}\right)\right]^{1 / d}[\operatorname{det}(I+B)]^{1 / d} \tag{65}
\end{equation*}
$$

But since $N_{1}$ is nilpotent, we have

$$
\begin{equation*}
\left[\operatorname{det}\left(A^{2}\right)\right]^{1 / d}=\left[\operatorname{det}\left(\left|c I+N_{1}\right|\right)\right]^{2 / d}=\left|\operatorname{det}\left(c I+N_{1}\right)\right|^{2 / d}=|c|^{2} \tag{66}
\end{equation*}
$$

We therefore conclude that

$$
\begin{equation*}
\operatorname{det}(I+B)=1 \tag{67}
\end{equation*}
$$

However, since $B \geq 0$, Eq. (67) can hold only if $B=0$. This in turn is possible only if $\sum_{j=2}^{r}\left|N_{j}\right|^{2}=0$; i.e. $\quad N_{j}=0$ for all $j=2, \ldots, d$. We therefore conclude that $\rho^{A B}$ is a pure state which implies that $|\psi\rangle^{A B C}=\left|w_{1}\right\rangle^{A B}|\varphi\rangle^{C}$, where $|\varphi\rangle^{C}$ is some pure state.

Note that by combining the above corollary with Theorem 3 we get that the G-concurrence is fully monogamous (even for mixed tripartite states). This is to our knowledge the fist example of a monogamous measure of entanglement that is highly non-faithful. We conclude now with the 2-dimensional version of Theorem 7. Recall that the concurrence is the two dimensional G-concurrence.

Corollary 9. Let $\rho^{A B}$ be an entangled two qubit state with rank $r>1$, and let $E$ be any injective (up to local unitaries) measure of pure two qubit entanglement. Then,

1. If $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ then $r=2$, and

$$
\rho^{A B}=\operatorname{Tr}_{C}|W\rangle\left\langle\left. W\right|^{A B C}\right.
$$

is the 2-qubit marginal of the $W$-class state

$$
\begin{equation*}
|W\rangle=\lambda_{1}|100\rangle+\lambda_{2}|010\rangle+\lambda_{3}|001\rangle+\lambda_{4}|000\rangle, \tag{68}
\end{equation*}
$$

where $\lambda_{j} \in \mathbb{C}$ and $\sum_{j=1}^{4}\left|\lambda_{j}\right|^{2}=1$.
2. Conversely, if $\rho^{A B}$ is a marginal of a state in the $W$-class then $C_{f}\left(\rho^{A B}\right)=C_{a}\left(\rho^{A B}\right)$, where $C$ is the concurrence.

Proof. Part 1: In the two qubit case, we can also write $E\left(|\psi\rangle^{A B}\right)=g\left(C\left(|\psi\rangle^{A B}\right)\right)$, where $C\left(|\psi\rangle^{A B}\right)$ is the concurrence measure of entanglement, which is itself an injective measure. Moreover, since $E$ is injective so is the function $g$. Now, in [35], Wootters showed that there always exists a pure state decomposition of $\rho^{A B}$ with each element in the decomposition being equal to $C_{f}\left(\rho^{A B}\right)$. Denoting by $\left\{p_{j},\left|\psi_{j}\right\rangle^{A B}\right\}$ this decomposition, we have that

$$
\begin{equation*}
E_{f}\left(\rho^{A B}\right) \leq \sum_{j} p_{j} g\left(C\left(\left|\psi_{j}\right\rangle^{A B}\right)\right)=g\left(C_{f}\left(\rho^{A B}\right)\right) \tag{69}
\end{equation*}
$$

Similarly, there exists a pure state decomposition $\left\{q_{k},\left|\phi_{k}\right\rangle^{A B}\right\}$ of $\rho^{A B}$ such that for each $k, C\left(\left|\phi_{k}\right\rangle^{A B}\right)=$ $C_{a}\left(\rho^{A B}\right)$ [36]. Hence,

$$
\begin{equation*}
E_{a}\left(\rho^{A B}\right) \geq \sum_{k} q_{k} g\left(C\left(\left|\phi_{k}\right\rangle^{A B}\right)\right)=g\left(C_{a}\left(\rho^{A B}\right)\right) \tag{70}
\end{equation*}
$$

Therefore, the equality $E_{f}\left(\rho^{A B}\right)=E_{a}\left(\rho^{A B}\right)$ implies that $C_{f}\left(\rho^{A B}\right)=C_{a}\left(\rho^{A B}\right)$. so that we can apply Theorem 7. There is only one 1-dimensional subspace of $\mathcal{N}$ (up to similarity), given by $\mathcal{N}_{0}=$ $\operatorname{span}\left\{S\left(\begin{array}{ll}0 & 1 \\ 0 & 0\end{array}\right) S^{-1}\right\}$, where $S$ is a $2 \times 2$ invertible matrix. Now, from (13) together with the decomposition (38) we conclude that $\rho^{A B}$ can be expressed as:

$$
\begin{equation*}
\rho^{A B}=|x\rangle\left\langle\left. x\right|^{A B}+\mid y\right\rangle\left\langle\left. y\right|^{A B},\right. \tag{71}
\end{equation*}
$$

where $|x\rangle^{A B}=X \otimes I^{B}\left|\phi_{+}\right\rangle$for some invertible matrix $X$, and

$$
|y\rangle \equiv I \otimes S\left(\begin{array}{ll}
0 & 1  \tag{72}\\
0 & 0
\end{array}\right) S^{-1}|x\rangle^{A B}
$$

Therefore, $\rho^{A B}$ is a marginal of the tripartite state

$$
\begin{equation*}
|\psi\rangle^{A B C}=|x\rangle^{A B}|0\rangle^{C}+|y\rangle^{A B}|1\rangle^{C} . \tag{73}
\end{equation*}
$$

Multiplying the above 3-qubit state by the SLOCC element $S^{T} X^{-1} \otimes S^{-1} \otimes I^{C}$, and using the symmetry $S^{T} \otimes S^{-1}\left|\phi_{+}\right\rangle=\left|\phi_{+}\right\rangle$for any invertible matrix $S$, gives

$$
\begin{align*}
S^{T} X^{-1} \otimes S^{-1} \otimes I^{C}|\psi\rangle^{A B C} & =\left(S^{T} \otimes S^{-1}\left|\phi_{+}\right\rangle^{A B}\right)|0\rangle^{C}+\left(S^{T} \otimes\left(\begin{array}{ll}
0 & 1 \\
0 & 0
\end{array}\right) S\left|\phi_{+}\right\rangle^{A B}\right)|1\rangle^{C} \\
& =\left|\phi_{+}\right\rangle^{A B}|0\rangle^{C}+\left(I^{A} \otimes\left(\begin{array}{ll}
0 & 1 \\
0 & 0
\end{array}\right)\left|\phi_{+}\right\rangle^{A B}\right)|1\rangle^{C} \\
& =|000\rangle+|110\rangle+|101\rangle \tag{74}
\end{align*}
$$

note that the state above is the W state after a local flip of the first qubit. Therefore, $|\psi\rangle^{A B C}$ is in the $W$ class.
Proof of Part 2: We use for this part Wootters' formula for the concurrence. Let $R=\sqrt{\rho^{1 / 2} \tilde{\rho} \rho^{1 / 2}}$ be Wootters matrix [35], with $\rho \equiv \rho^{A B}$ and $\tilde{\rho} \equiv \sigma_{y} \otimes \sigma_{y} \rho^{*} \sigma_{y} \otimes \sigma_{y}$. Wootters showed in [35] that for entangled states $C_{f}\left(\rho^{A B}\right)=\lambda_{1}-\lambda_{2}-\lambda_{3}-\lambda_{4}$, where $\lambda_{1}, \ldots, \lambda_{4}$ are the eigenvalues of $R$ in decreasing order. Furthermore, in $[37,36]$ it was shown that $C_{a}\left(\rho^{A B}\right)=\operatorname{Tr}[R]=\lambda_{1}+\lambda_{2}+\lambda_{3}+\lambda_{4}$. Therefore, $C_{f}\left(\rho^{A B}\right)=C_{a}\left(\rho^{A B}\right)$ if and only if $R$ is a rank one matrix. A straightforward calculation shows that for the bipartite marginals, of the W-class states (68), $\operatorname{Tr}\left[R^{2}\right]=(\operatorname{Tr}[R])^{2}$ so that the rank of $R$ is one (recall that $R$ is positive semidefinite). This completes the proof.

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