# A methodology for creating feeding routes in mass transit sustems 

## Metodología para crear rutas alimentadoras en sistemas de transporte masivo

Metodologia para criar rotas alimentadoras em sistemas
de transporte massivo

Fecha de recepción: 2 de agosto de 2016
Fecha de aprobación: 26 de marzo de 2017

Daniela Ospina-Toro*<br>Eliana Mirledy Toro-Ocampo**<br>Ramón Alfonso Gallego-Rendón***


#### Abstract

This paper proposes a methodology to identify feeding routes for areas disconnected to the Mass Transit System (MTS), in order to propose an alternative solution to the deficit in the number of carried passengers. The proposed methodology consists of two steps: (1) structuring scenarios for areas disconnected from the transport system, and (2) combining heuristic and exact techniques to solve the feeding routes problem. The methodology considers among its restrictions the path length and the passenger vehicle capacity. To model the problem, a comparison with the Location Routing Problem (LRP), which is usually applied to freight transport problems, was established. The proposed methodology is a math-heuristic that combines the Lin-Kernighan-Helsgaun algorithm (LKH), and the Clark and Wright's Savings heuristic with the Branch-and-Cut exact algorithm, which is applied into a Mixed Integer Linear Programming model (MILP), also known as a Set Partitioning model (SP) for LRP. This methodological approach is validated with real instances in the massive transport system of Pereira (Megabús), considering some areas disconnected from the Central-Occidental Metropolitan Area System (AMCO) of Pereira, located in the Colombia's Coffee Axis.


Keywords: Feeding Routes; LKH algorithm; Location Routing Problem; Math-Heuristic; Savings Algorithm; Set-Partitioning model.

## Resumen

Se propone una metodología para identificar rutas alimentadoras en zonas no conectadas para un sistema de transporte masivo, con el fin de aumentar la cobertura del servicio y mejorar el nivel de ocupación del sistema. La metodología propuesta consta de dos etapas: 1) estructurar escenarios de áreas no conectadas al sistema de transporte y 2 ) combinar técnicas heurísticas y exactas para resolver el problema de rutas alimentadoras. La metodología considera dentro de sus restricciones la duración de la ruta y la capacidad del vehículo alimentador. Para su modelamiento se establece una analogía entre los problemas del transporte de pasajeros y el problema de localización y ruteo, Location Routing Problem (LRP), que usualmente es aplicado a problemas de transporte

[^0]de mercancías. La metodología de solución propuesta es una matheurística que combina las heurísticas Lin-Kernighan-Helsgaun (LKH) y ahorros con el algoritmo de ramificación y corte, Branch-and-Cut, aplicado sobre un modelo lineal entero mixto de partición de conjuntos (Set Partitioning) para LRP. Esta propuesta metodológica es validada con casos de prueba reales del sistema de transporte masivo de la ciudad de Pereira (Megabús), donde se consideran algunas zonas no conectadas del Área Metropolitana Centro Occidente, localizada en el eje cafetero colombiano.

Palabras clave: Heurística de ahorros; Heurística LKH; Math-heurística; Modelo de partición de conjuntos; Problema de localización y ruteo; Rutas alimentadoras.

## Resumo

Propõe-se uma metodologia para identificar rotas alimentadoras em áreas não conectadas para um sistema de transporte massivo, com o fim de aumentar a cobertura do serviço e melhorar o nível de ocupação do sistema. A metodologia proposta consta de duas etapas: 1) estruturar cenários de áreas não conectadas ao sistema de transporte e 2) combinar técnicas heurísticas e exatas para resolver o problema de rotas alimentadoras. A metodologia considera, dentro de suas restrições, a duração da rota e a capacidade do veículo alimentador. Para seu modelamento estabelece-se uma analogia entre os problemas do transporte de passageiros e o problema de localização e roteamento, Location Routing Problem (LRP), que usualmente é aplicado a problemas de transporte de mercadorias. A metodologia de solução proposta é uma matheurística que combina as heurísticas Lin-Kernighan-Helsgaun (LKH) e economias com o algoritmo de ramificação e corte Branch-and-Cut, aplicado sobre um modelo lineal inteiro misto de partição de conjuntos (Set Partitioning) para LRP. Esta proposta metodológica é validada com casos de prova reais do sistema de transporte massivo da cidade de Pereira (Megabús), onde se consideram algumas regiões não conectadas da Área Metropolitana Centro Ocidente, localizada no Eje Cafetero colombiano.

Palavras chave: Heurística de economia; Heurística LKH; Matheurística; Modelo de partição de conjuntos; Problema de localização e roteamento; Rotas alimentadoras.

Cómo citar este artículo:
D. Ospina-Toro, E. M. Toro-Ocampo, and R. A. Gallego-Rendón, "A methodology for creating feeding routes in mass transit systems," Rev. Fac. Ing., vol. 26 (45), pp. 9-21, May. 2017

## I. Introduction

The Colombian automotive market is more aggressive everyday regarding sale strategies, proposing a big quantity of opportunities for buying vehicles, discouraging public transportation use, and causing an unstoppable increase in traffic that has led to a great urban congestion in many cities. The fleet of vehicles and motorcycles is bigger every day, the parking zones in the city centers are insufficient, and the road infrastructure collapses faced with the traffic flow in the rush hours. This is why nowadays, proposals that include a complete, effective, and environmentally respectful public transportation, and that integrate an important part of the city areas have gained great interest in the community.

The urban transport, as a whole, is an activity of important magnitude for a country. The operation of vehicles that circulate in roads of Latin American and the Caribbean cities with populations higher than 100,000 inhabitants consumes close to $3.5 \%$ of the GDP [1]. In the case of Colombia, the National Government implemented a policy known as CONPES 3197/2002 [2] for improving the service of urban public transportation for passengers, followed by the CONPES 3197/2002 [2] for improving the urban and public transportation service for passengers, and the CONPES 3260/2003 [3] for urban and massive transportation. These policies were implemented after analyzing the transportation deficit in populations with more than 500,000 inhabitants, and high level of development and growth, which led to the conclusion that the suitable solution was to implement integrated massive transportation systems. So far, these systems have been implemented in several capital cities as a solution to the urban mobility issue, but five of them have not achieved the projected goals, due to the competition with other means of transportation (formal and informal).

Pereira is one of those five cities that have not reached the goal number of transported passengers, which should be around 140,000 per day. Nowadays, 90,000 passengers are being transported, indicating that there is still a deficit of 50,000 daily passengers to reach the break-even point. Hence, the challenge is to expedite the integration process into the metropolitan area [4], and to achieve that, a methodology from a mathematical-optimization perspective is presented in this paper. The methodology proposes a design of
feeding routes for zones disconnected from the mass transportation system (MTS). The feeding routes are integrated to the interchange station or stop points, taking as reference the record available in the origindestination matrix, and projecting them to the date of study.

The contributions of this paper are the following: 1) a methodology for generating feeding routes in unintegrated sectors of the MTS that introduces an analogy to the Location-Routing Problem (LRP), and 2) a strategy for integrating disconnected zones from the MTS, in order to reach a break-even point regarding the number of passengers, and thus achieve its economic viability, so wider service coverage is attained.

## II. Demand Forecasting

The compound annual growth rate (CAGR) of the population growth is the traditionally used measurement in the business and investment fields for estimating the profits in a time lapse [5]. The computation of the population growth rate for cities and communities has been applied to the U.S. prison population [6], to the forecast of methane emissions in Delhi, as a consequence of its population growth, and to scientific papers written by mathematicians on global population growth [7]. The demand forecast was calculated using (1) and (2), as proposed in [5].

$$
\begin{align*}
& r_{p}=\left(\frac{P_{t}}{P_{0}}\right)^{\frac{1}{n}}-1  \tag{1}\\
& r_{u}=\left(\frac{D_{F}}{D_{A}}\right)-1 \tag{2}
\end{align*}
$$

Where $r_{p}=$ population growth rate per time unit, $P_{t}=$ population in the year $\mathrm{t}, \mathrm{P}_{0}=$ population in the year $0, \mathrm{n}$ $=$ number of years, $r_{u}=$ user's growth rate based on the forecast demand, $D_{f}=$ future customers demand, $D_{A}=$ current customers demand.

## III. Location-Routing Problem

The Location-Routing Problem (LRP) can be defined as a finite set of potential depots $I$ with a fixed opening $\operatorname{cost} f_{i}$ and capacity $b_{i}$, and a set of costumers $J$, where every costumer has a demand $d_{j}$. A homogeneous vehicle fleet with a limited size must be routed from
the depots for attending customers demand. Every vehicle has an associated capacity $Q$, and every pair of nodes, $i$ and $j$, has an associated traveling $\operatorname{cost} c_{i j}$ [8].

The objective is to determine the set of depots that must be open, the order for visiting the customers, and the allocation of routes to the opened depots. All these actions must be performed minimizing the total operative costs, which represent the opening depots costs, the opening routes costs, and the costs associated to the utilization of the active arcs [9].

Several authors have focused on developing heuristic methods for finding solutions in computationally reasonable times. Nagy and Salhi [10] classified these algorithms as sequential, iterative, hierarchic, and clustering-oriented. The sequential methods initially solve the depots location, and then the routes associated to every depot; however, this type of methodology reduces the feedback between both problems, and thereupon is inadvisable [11]. The iterative methods solve both sub-problems simultaneously, and give permanent feedback through the obtained outputs from every sub-problem separately; these strategies applied to the LRP solve the problem without considering the locating variables, allocating one depot for every set of customers, and at least one route to the open depots [12]. The hierarchical methods consider depots location as the main problem, and the routing problem as a subordinated problem; first, the location problem is solved in an approximated way, and then, an associated subroutine to the corresponding routing problem is added [13, 14]. The methods based on clustering customers were proposed by Barreto et al. [15].

In the first phase, the set of customers is divided into groups depending on the vehicle capacity. In the second phase, a Traveling Salesman Problem (TSP) is solved for every cluster. In the last phase, every TSP tour is clustered in super nodes to solve the capacitated depots location problem.

To solve big-sized instances, Prins et al. [16], Prodhon [9, 17, 18], Duhamel et al. [19] and Escobar et al. [20] proposed several metaheuristics; among them, the techniques based on Granular Tabu Search (GTS) have been the most efficient for solving instances of up to 200 customers. The methods combine Integer Lineal Programming (ILP) models using Lagrangean relaxation for setting up the routes. Pirkwieser and Raidl [21] proposed the Variable Neighborhood

Search (VNS) for solving the Periodic Location Routing Problem (PLRP), and the CLRP approached through a combination of a VNS with the sequential solution of a ILP. The considered ILPs are two-index models aimed to solve the location problem, and the reallocation by means of a set-partitioning model (SP).

Yu et al. [22] proposed the Simulated Annealing algorithm, where diversification is controlled according to the temperature parameter, in order to escape from local optima. Better gaps were found using the proposed methodology, compared to the ones found in [16, 19]; also, the location-allocation concept is introduced, where the location and routing decisions are made simultaneously. In Hemmelmayr et al. [23], the Adaptive Large Neighborhood Search heuristic (ALS) was developed for the LRP. In this technique, different neighborhoods are applied and classified according to the obtained results; for the following iterations, the best scored neighborhoods have the higher possibilities to be selected.

Furthermore, the LRP has been solved throughout exact techniques, including the Branch-and-Bound method [24], column generation, Branch and Price [25], and Branch and Cut [26]. These methods solve instances of up to 200 customers, but some instances of 100 customers are still unsolved. Prodhon and Prins [27] presented a complete study on problems associated to the LRP, where models, solution methodologies, and a comparison of the results obtained using the best algorithms are encompassed, analyzed, discussed, and classified.

## IV. Proposed Mathematical Model for Defining Feeding Routes

The solved model is based on the formulation presented by Akca et al. [25] that corresponds to a SP model for
the LRP. $\Omega$ is defined as the set of all possible routes starting and finishing in the depot (interchange station)
$i \in I$, subserving a subset of clients (bus stations) of at least two customers with a demand $\mathrm{q}_{1}$. For every $l \in \Omega$, a binary variable $\theta_{l}$ equal to 1 is associated if $l$ appears in the optimal solution of the LRP, and 0 otherwise, and a $C_{l}$ cost for using a route. For every client (bus station) $j \in J$, where $J$ corresponds to the
set of bus stations, and the route $l \in \Omega$ is assigned to the variable $a_{j}^{l}$, which corresponds to the number of times the customer $j$ appears in the route $l, z_{i}$ is a binary variable that is 1 if the depot (interchange station) $i$ is used as the starting and ending point of the feeding route, and 0 otherwise. $N_{\text {min }}$ represents the number of closed depots (interchange stations).
$\min \sum_{l \in \Omega} c_{l} \theta_{l}$
Subject to

$$
\begin{align*}
& \sum_{\substack{l \in \Omega \\
j \in J}} a_{j}^{l} \theta_{l}=1, \forall j \in J \\
& \sum_{l \in \Omega} q_{l} \theta_{l} \leq b_{i}\left(1-z_{i}\right), \forall i \in I \\
& \sum_{i=1}^{H} z_{i}=N_{\min }  \tag{5}\\
& \theta_{l} \in[0,1]  \tag{6}\\
& a_{j \geq 0}^{l} \geq \text { and integer }  \tag{7}\\
& z_{i} \in[0,1] \tag{8}
\end{align*}
$$

The objective function (3) formulates the feedingrouting cost minimization, and the interchange stations use. Constraints presented in (4) indicate that a bus station can only be allocated to one route, so there are as many constraints as existing bus stations. Constraints (5) indicate that the accumulated demand of bus stations included into the feeding routes do not exceed the interchange stations capacity; there are as many constraints of this type, as considered interchange stations. Constraint (6) assures that the interchange stations that remain closed must be
equal to $N_{\min }$; this constraint has been added to the original model, and represents a valid cut for the model, guaranteeing that only a set of suitable interchange stations will be added, the ones strictly necessary for attending the feeding routes demand. Finally, (7), (8), and (9) define the variables nature.

## V. Proposed Methodology

According to the bibliographic review, we can make an analogy between the LRP and the feeding routing problem, where the depots are equivalent to the interchange stations on the backbone, the customers are associated to the interchange stations, and the demand is associated to the number of passengers required to board the vehicle at a determined bus station. Like in the LRP, the route is designed and connected to the closer interchange station in the MTS backbone.

The proposed methodology includes demand forecast and a math-heuristic developed in three steps: 1) a big tour construction over a network of possible bus stations in a disconnected zone; 2) generation of $n$ possible route alternatives that are subsequently connected to one or more points on the backbone, and identification of $n$ possible circular routes that will be sent as input of the third step; 3) use of a Mixed Integer Linear Programming problem (MILP) to identify the set of routes attending the bus station demands, considering as objective function the travel time minimization.

## A. Math-heuristic for solving the problem

The problem solution is described in three steps:
Step 1. Grand tours attainment. This step consists in generating three TSP or grand tours with the bus stations of the route to be designed, using three different techniques: 1) Clark \& Wright Savings [28], 2) Walk, and 3) Nearest Neighborhood. All three tours are improved using them as the Lin-KernighanHelsgaun algorithm (LKH) input, an improved version of the traditional Lin-Kernighan algorithm developed by Helsgaun [29].

Step 2. Route Segmentation. A tour-partitioning is performed to the three grand tours, using the Prins algorithm [30]; then $n$ routes are obtained and connected to the interchange stations considered in the study. Infeasible solutions in terms of vehicle capacity and time duration are discarded.

Step 3. Mixed Integer Linear Programming (MILP) solution. The routes generated in the former step are defined as decision variables of the problem that
is solved using the Branch-and-Cut algorithm. The MILP solution is defined in equations (3) to (7).

1) Proposed algorithm: The algorithm is presented in

Algorithm 1. The required input parameters are added in lines 2 to 7 .

Step 1: To obtain grand tours, the LKH and the savings algorithms were used (lines 8 to 20).

Step 2: To generate feasible routes, they are built starting with a single client up to $k$ clients, considering the vehicle capacity and time constraints. This step is performed for every grand tour previously obtained (lines 21 to 26).

Step 3: To solve the PLEB Model using the CPLEX software, the inputs for the model are the set of routes $\Omega_{l}$ the set of demands $q_{l \text {, and the set of times }} c_{l}$.

```
1: Algorithm
\(c_{i j} \leftarrow\) Time_Matrix[][]
\(q_{i} \leftarrow\) Demand[]
\(G T_{k} \leftarrow\) Grand Tour []
\(\Omega_{l} \leftarrow\) Set of routes[][]
\(q_{l} \leftarrow\left\{q_{i}\right\}\)
\(c_{l} \leftarrow\left\{c_{i j}\right\}\)
for \(t \leftarrow 1 ; t<=3 ; t++\)
    switch (t)
                case 1:
                \(G T_{t}=L K H\left(c_{i j}\right.\) walk \()\)
                break
                case 2:
                \(G T_{t}=L K H\left(c_{i j}\right.\), nearest_neighborhood \()\)
                break
                case 3:
                \(G T_{t}=\operatorname{savings}(c i j)\)
                break
            End
            end
            for \(t \leftarrow 1 ; t<=3 ; t++\)
            for \(i \in I\)
            if \(q_{l} \leftarrow q_{l}+q_{i}>b_{i}\)
                                    \(\left[c_{b}, q_{l}, \Omega_{l}\right]=\) Segment_route \(\left(G T_{t}\right)\)
            End
            End
            end
            \(s=\operatorname{MILPB}\left(c_{l}, q_{b}, \Omega_{l}\right)\)
```

Algorithm 1. Algorithm for adding feeding routes into a MTS by using the proposed methodology.
2) Applied example: We used a system of five bus stations (3, 4, 5, 6, and 7) located in a disconnected zone, which must be attended from the interchange stations D1 and/or D2 that are placed on the backbone.

Table 1 and 2 present the input data.
Table 1
Time matrix

|  | D1 | D2 | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{D 1}$ | - | - | 4 | 1 | 8 | 3 | 9 |
| $\mathbf{D 2}$ | - | - | 6 | 2 | 3 | 2 | 7 |
| $\mathbf{3}$ | 4 | 6 | - | 5 | 6 | 7 | 4 |
| $\mathbf{4}$ | 1 | 2 | 5 | - | 3 | 9 | 5 |
| $\mathbf{5}$ | 8 | 3 | 6 | 3 | - | 8 | 9 |
| $\mathbf{6}$ | 3 | 2 | 7 | 9 | 8 | - | 3 |
| $\mathbf{7}$ | 9 | 7 | 4 | 5 | 9 | 3 | - |

Table 2
Coordinates and demand/supply of passengers per bus/interchange station

|  | Coordinate X | Coordinate Y | Demand/Supply |
| :---: | :---: | :---: | :---: |
| D1 | 10 | 7 | 22 |
| D2 | 5 | 3 | 22 |
| $\mathbf{3}$ | 0 | 5 | 5 |
| $\mathbf{4}$ | 4 | 6 | 2 |
| $\mathbf{5}$ | 6 | 5 | 7 |
| $\mathbf{6}$ | 8 | 5 | 3 |
| $\mathbf{7}$ | 7 | 9 | 4 |

Step 1. The grand tours were generated by using heuristic techniques, considering the bus stations described in

Table 1 . We obtained three sequences where depots D1 and D2 are not included: a) 3-4-5-6-7; b) 6-7-5-34; c) 5-7-4-3-6.

Step 2. $\mathrm{k}_{\mathrm{n}}$ possible feeding routes were generated for every sequence, where k represented the number
of depots. Every feeding route corresponded to a decision variable for the BILP. In this step, we verified the alternative to quantify time and demand for determining the feeding routes feasibility. We considered a homogeneous fleet with a capacity of 15 passengers per vehicle, and a maximum time per route of 25 minutes.

Table 3 shows some of the obtained feeding routes with their respective time, capacity, and feasibility. Figure 1 shows the optimal solution of the example.

Table 3
Set of feeding routes generated with their respective time, capacity, and feasibility

| Sequence | Time | Capacity | Feasibility | Route |
| :---: | :---: | :---: | :---: | :---: |
| 131 | $4+4=8$ | 5 | Yes | $\# 1$ |
| 232 | $6+6=12$ | $4+5+1=10$ | 5 | Yes |
| 1341 | $6+5+2=13$ | $5+2=7$ | Yes | $\# 2$ |
| 2342 | $4+5+3+8=20$ | $5+2+7=14$ | Yes | $\# 3$ |
| 13451 | $6+5+3+3=17$ | $5+2+7=14$ | Yes | $\# 4$ |
| 23452 | $4+5+3+8+3=23$ | $5+2+7+3=17$ | Ynly time | $\# 5$ |
| 134561 | $6+5+3+8+3=24$ | $5+2+7+3=17$ | Only time | No |
| 234562 | $4+5+3+8+3+9=32$ | $5+2+7+3+4=21$ | No | No |
| 1345671 |  |  |  | No |
| 16751 | $3+3+9+8=23$ | $3+4+7=14$ | Yes | \#m |
| 26752 | $2+3+9+3=17$ | $3+4+7=14$ | Yes | $\# m+1$ |
| $\ldots$ |  |  |  | No |
| 257432 | $3+9+5+5+4=26$ | $7+4+2+5=18$ | No |  |



Fig. 1. Optimal solution for the applied example.

## VI. Study case of Megabús: Central-Occidental Metropolitan Area System (AMCO)

The study case started by updating the information for 2015 by incorporating the demand forecast methodology discussed in section 2. The basic information corresponding to Pereira and La Virginia municipalities can be found in [31]. On the other hand, the information about the Technological University of Pereira (UTP) zone was taken from [32], and it was divided into zones in order to identify the user's origin points. The data forecast for 2015 was performed by using the growth rates obtained in (8), (9), and (10).

## Growth rate for Pereira, La Virginia, and the UTP

Pereira (2009 to 2012)

La Virginia (2009 to 2012)

UTP (2005 to 2012)

$$
\begin{align*}
& r_{p}=\left(\frac{459690}{451814}\right)^{\frac{1}{3}}-1=0.0058  \tag{8}\\
& r_{p}=\left(\frac{31734}{310504}\right)^{\frac{1}{3}}-1=0.0024  \tag{9}\\
& r_{p}=\left(\frac{15042}{9193}\right)^{\frac{1}{7}}-1=0.0729 \tag{10}
\end{align*}
$$

The Megabus system currently mobilizes 90.000 passengers, but its objective is to reach 140.000 ; therefore, the demand must be projected considering the users growth rate factor denominated $r_{u}$.

$$
\begin{equation*}
r_{u}=\left(\frac{140000}{97000}\right)-1=0.4433 \tag{12}
\end{equation*}
$$

## A. Construction of the test cases

The AMCO zonification was applied using the transport zones analysis (ZAT). The bus stations coordinates, the origin-destination matrix, and the Pereira city zonification can be consulted at http:// blog.utp.edu.co/problemasderuteo/? $\mathrm{p}=331$.

The considered study cases were taken from every disconnected zone of the Megabús backbone, which are structured considering the proximity of the bus stations, the topography, and the city road design. Based on this information, the demand matrix was generated using Google Maps, and the traveling times among nodes were computed. The same value of $r_{u}$ was taken for all zones to forecast the demand.

The $r_{p}$ value used for the ZAT of 10 bus stations was the one found for the UTP. The current Megabús stations,
hereafter called interchange stations, were codified with values between 1 and 100. For the Samaria zone, the bus stations are considered from 101 to 199 , for the UTP zone, the bus stations are codified from 201 to 299, and for La Virginia zone from 301 to 399.

In total, we considered five study cases; every case had an associated bus station number that corresponded to the connection of the route to a station in the backbone, and the bus stations number. Table 4 presents the different scenarios.

- 1V: Scenario for La Virginia municipality. Time limit: 2 hours.
- 1 S and 2 S : Scenarios for the Samaria neighborhood. Time limit: 1 hour 10 minutes.
- 1UTP, 2UTP, and 3UTP: Scenarios for the Technological University of Pereira. Time limit: 1 hour 10 minutes.

The shown scenarios were disconnected from the MTS Megabús. For the circular routes generation, in Samaria and UTP zones, an infinite demand, and a limit time of 1 hour 10 minutes were considered. For the Virginia zone, a limit time of 2 hours was considered because it is an intermunicipal route.

Table 4
Scenarios. V: La Virginia. S: Samaria Neighborhood. UTP: Technological University of Pereira

| Study case | Interchange <br> Stations | Bus Stations |
| :--- | :--- | :--- |
| 1 V | $1,2,3$ | $301,302,303,304,305,306,307,308$ |
| 1 S | $28,29,36$ | $101,102,104,106,107,108,109,110,111,115,116$, |
|  |  | $117,118,119,120,127,128,132$ |
| 2 S | $28,29,36$ | $109,110,111,112,113,114,115,116,117$, |
|  | $118,119,120,122,123,124,127,128,129,132$ |  |
| 1 UTP | $6,7,8,20,21$ | $204,205,208,216,217,220,221,228,229,230,231,232$, |
|  |  | $233,234,235,236,238,239,240,241,242,243,245,247,248$ |
| 2 UTP | $5,6,7,8,9,20,21$ | $205,208,215,220,221,228,229,230,231,232,23,234,235$ |
|  |  | $236,238,239,240,241,242,245,247,248,249,250$ |
| 3 UTP | $5,6,7,8,9,20,21$, | $206,207,208,209,211,212,213,214,215,25,239,240$, |
|  | $22,23,29$ | $241,242,243,244,249,250,251,252,253,254$ |

Step 1: The grand tours or TSP were generated using as input the data presented in column "Interchange Stations" in Table 5.

Step 2: The set of feeding routes was generated from the TSP obtained in Step 1. The Megabús interchanging stations are considered depots for generating circular
routes between interchange stations and bus stations. Taking into account demand and time restrictions, the feeding routes were categorized as feasible or unfeasible. Thus, only the feasible routes were considered as decision variables for step 3.

Step 3: The MILP presented in equations (3) to (7) was solved to obtain the results shown in Table 5, where TEA represents the algorithm running time.

## Table 5

Solutions found for the proposed scenarios

| Study case | Generated Sequence | Feeding route duration | TEA (seconds) |
| :---: | :---: | :---: | :---: |
| 1V | Route 1: 2-308-304-305-306-307-301-302-2 | 105 min | 0.13 |
| 1S | Route 1: 36-132-120-127-128-116-36-36 | 33 min | 2.33 |
|  | Route 2: 36-105-106-107-104-102-101-108-109-110-111-115-117- |  |  |
| 1S | 118-119-36 | 41 min |  |
|  | 36-119-118-117-115-111-114-110-124-109-112-123-122-113-129-116- |  | 1.9 |
| 2S | $128-127-120-132-36$ | 64 min |  |
| 1 UTP | Route 1: 21-239-253-240-241-249-250-254-242-243-21 | 33 min |  |
|  | Route 2: 21-236-235-248-221-216-217-204-205-220-232-233-229- |  | 5.9 |
| 1 UTP | 228-230-247-245-21 | 66 min |  |
| $2 \text { UTP }$ | Route 1: 21-239-253-240-241-249-250-254-242-243-21 | 33 min | 9.71 |
|  | Route 1: 21-208-238-245-247-220-215-205-221-232-233-229-228- |  | 7.01 |
| 3 UTP | 230-234-231-248-235-236-21 | 61 min |  |

## B. Computational implementation

Step 1 was performed using the LHK executable file (available at http://www.akira.ruc.dk/~keld/ research/LKH/); additionally, the savings algorithm was built on $\mathrm{C}++$. For step 2, a code was developed on $\mathrm{C}++$, and step 3 was run by using the IBM ILO CPLEX (available at http://www-03.ibm.com/ software/products/es/ibmilogcpleoptistud) to solve the optimization model. The methodology is a whole routine, where LKH, C++, and IBM ILO CLPLEX are integrated. The algorithm was executed on a Lenovo

Z470 with an Intel Core $15-2410 \mathrm{M}$, processor of 2.3 GHz , and a RAM memory of 4GB.

## C. Feeding routes tracing

Feeding routes were identified for attending the disconnected zones, and thus integrating new users to the MTS, so the break-even point could be reached. Thereby, a viable system from an economic perspective may be generated. Figure 2 corresponds to the 1UTP solution.


Fig. 2. Solution route: Scenario UTP, feeding route 1.

## VII. Conclusions

An application to the Mass Transport System problem was developed for reaching the operational breakeven point. To solve the problem, the passengers forecast concept was incorporated, as well as the math-heuristic optimization method, combining heuristic optimization techniques to determine possible solutions, and an exact method to choose the best proposed alternatives. The feasibility criteria were met regarding the feeding route duration and the vehicle capacity.

The methodology described apply LRP concepts to the MTS, where the depots are equivalent to the interchange stations placed on the backbone. The set of costumers is associated to the set of bus stations, and the demand is associated to the amount of passengers who are required to board the route at a determined bus station. Like in the LRP, the route is designed and connected to the closer backbone interchange station of Megabús.

Many MTS operators still must reach their operative goals (number of users), due to the lack of feeding routes that would allow access to different locations of the city that represent potential sources of users. This is why this methodology attempts to meet the users demand and, from this information, establish possible feeding routes for accessing the system at a minimum
cost, and satisfying the feasibility constraints related to vehicle capacity and route time duration.

## References

[1] I. Thomson, Impacto de las tendencias sociales, económicas y tecnológicas sobre el transporte público: investigación preliminar en ciudades de América Latina; United Nations Publications, 2002.
[2] Ministerio de Transporte de Colombia, Política para mejorar el servicio de transporte Público urbano de pasajeros; Consejo Nacional de Política Económica y Social, CONPES 3167, Departamento Nacional de Planeación, 2002.
[3] Ministerio de Transporte de Colombia, Política Nacional de Transporte Urbano y Masivo; Consejo Nacional de Política Económica y Social, CONPES 3260, DEPARTAMENTO NACIONAL DE PLANEACIÓN2003.
[4] El Tiempo. Redacción Barranquilla, Medellín, Bucaramanga y Pereira, "Ninguno de los 'Transmilenios' del país se salva de líos financieros," ed: eltiempo.com, 2013.
[5] M. J. Anson, F. J. Fabozzi, and F. J. Jones, The handbook of traditional and alternative investment vehicles: investment characteristics and strategies, vol. 194: John Wiley \& Sons, 2010. DOI: http://doi. org/10.1002/9781118258248.
[6] A. Blumstein, and A. J. Beck, "Population growth in US prisons, 1980-1996," Crime. \& Just., vol. 26, pp. 17-61, Jan. 1999. DOI: http://doi. org/10.1086/449294.
[7] J. W. Grossman, "Patterns of collaboration in mathematical research," SIAM News, vol. 35, pp. 8-9, 2002.
[8] G. Laporte, Y. Nobert, and D. Arpin, "An exact algorithm for solving a capacitated location-routing problem," Annals of Operations Research, vol. 6 (9), pp. 291-310, Sep. 1986. DOI: http://doi.org/10.1007/ BF02023807.
[9] C. Prodhon, and C. Prins, "A memetic algorithm with population management (MA| PM) for the periodic location-routing problem," Lecture Notes in Computer Science, vol. 5296, pp. 43-57, 2008. DOI: http://doi.org/10.1007/978-3-540-88439-2_4.
[10] G. Nagy, and S. Salhi, "Location-routing: Issues, models and methods," European Journal of Operational Research, vol. 177 (2), pp. 649672, Mar. 2007. DOI: http://doi.org/10.1016/j. ejor.2006.04.004.
[11] S. Salhi, and G. K. Rand, "The effect of ignoring routes when locating depots," European Journal of Operational Research, vol. 39 (2), pp. 150-156, Mar. 1989. DOI: http://doi.org/10.1016/0377-2217(89)90188-4.
[12] D. Tuzun, and L. I. Burke, "A two-phase tabu search approach to the location routing problem," European Journal of Operational Research, vol. 116 (1), pp. 87-99, Jul. 1999. DOI: http://doi.org/10.1016/ S0377-2217(98)00107-6.
[13] M. Albareda-Sambola, J. A. Díaz, and E. Fernández, "A compact model and tight bounds for a combined location-routing problem," Computers \& Operations Research, vol. 32 (3), pp. 407-428, Mar. 2005. DOI: http://doi.org/10.1016/S0305-0548(03)00245-4.
[14] J. Melechovský, C. Prins, and R. W. Calvo, "A metaheuristic to solve a location-routing problem with non-linear costs," Journal of Heuristics, vol. 11 (5-6), pp. 375-391, Dec. 2005. DOI: http://doi. org/10.1007/s10732-005-3601-1.
[15] S. Barreto, C. Ferreira, J. Paixao, and B. S. Santos, "Using clustering analysis in a capacitated locationrouting problem," European Journal of Operational Research, vol. 179 (3), pp. 968-977, Jun. 2007. DOI: http://doi.org/10.1016/j.ejor.2005.06.074.
[16] C. Prins, C. Prodhon, A. Ruiz, P. Soriano, and R. Wolfler Calvo, "Solving the capacitated locationrouting problem by a cooperative Lagrangean relaxation-granular tabu search heuristic," Transportation Science, vol. 41 (4), pp.470-483, Nov. 2007. DOI: http://doi.org/10.1287/trsc.1060.0187.
[17] C. Prodhon, "An elsxpath relinking hybrid for the periodic location-routing problem," Lecture Notes in Computer Science, vol. 5818, pp. 15-29, 2009. DOI: http://doi.org/10.1007/978-3-642-04918-7_2.
[18] C. Prodhon, "A hybrid evolutionary algorithm for the periodic location-routing problem," European Journal of Operational Research, vol. 210 (2), pp.

204-212, Apr. 2011. DOI: http://doi.org/10.1016/j. ejor.2010.09.021.
[19] C. Duhamel, P. Lacomme, C. Prins, and C. Prodhon, "A GRASP $\times$ ELS approach for the capacitated location-routing problem," Computers \& Operations Research, vol. 37 (11), pp. 1912-1923, Nov. 2010. DOI: http://doi.org/10.1016/j.cor.2009.07.004.
[20] J. W. Escobar, R. Linfati, and P. Toth, "A twophase hybrid heuristic algorithm for the capacitated location-routing problem," Computers \& Operations Research, vol. 40 (1), pp. 70-79, Jan. 2013. DOI: http://doi.org/10.1016/j.cor.2012.05.008.
[21] S. Pirkwieser, and G. R. Raidl, Variable neighborhood search coupled with ILP-based very large neighborhood searches for the (periodic) location-routing problem: Springer, 2010. DOI: http://doi.org/10.1007/978-3-642-16054-7_13.
[22] B. Yu, Z.-Z. Yang, and B. Yao, "An improved ant colony optimization for vehicle routing problem," European Journal of Operational Research, vol. 196 (1), pp. 171-176, Jul. 2009. DOI: http://doi. org/10.1016/j.ejor.2008.02.028.
[23] V. C. Hemmelmayr, J.-F. Cordeau, and T. G. Crainic, "An adaptive large neighborhood search heuristic for two-echelon vehicle routing problems arising in city logistics," Computers \& Operations Research, vol. 39 (12), pp. 3215-3228, Dec. 2012. DOI: http://doi. org/10.1016/j.cor.2012.04.007.
[24] R. T. Berger, C. R. Coullard, and M. S. Daskin, "Location-routing problems with distance constraints," Transportation Science, vol. 41 (1), pp. 29-43, Feb. 2007. DOI: http://doi.org/10.1287/ trsc.1060.0156.
[25] Z. Akca, R. Berger, and T. Ralphs, "A branch-andprice algorithm for combined location and routing problems under capacity restrictions," in Operations Research and Cyber-Infrastructure, ed: Springer, pp. 309-330, 2009. DOI: http://doi.org/10.1007/978-0-387-88843-9_16.
[26] J.-M. Belenguer, E. Benavent, C. Prins, C. Prodhon, and R. W. Calvo, "A branch-and-cut method for the capacitated location-routing problem," Computers \& Operations Research, vol. 38 (6), pp. 931941, Jun. 2011. DOI: http://doi.org/10.1016/j. cor.2010.09.019.
[27] C. Prodhon and C. Prins, "A survey of recent research on location-routing problems," European Journal of Operational Research, vol. 238 (1), pp. 1-17, Oct. 2014. DOI: http://doi.org/10.1016/j. ejor.2014.01.005.
[28] G. U. Clarke and J. W. Wright, "Scheduling of vehicles from a central depot to a number of delivery points," Operations Research, vol. 12, pp. 568-581, 1964. DOI: http://doi.org/10.1287/opre.12.4.568.
[29] K. Helsgaun, "An effective implementation of the Lin-Kernighan traveling salesman heuristic," European Journal of Operational Research, vol. 126,
pp. 106-130, 2000. DOI: http://doi.org/10.1016/ S0377-2217(99)00284-2.
[30] S. Liu, W. Huang, and H. Ma, "An effective genetic algorithm for the fleet size and mix vehicle routing problems," Transportation Research Part E: Logistics and Transportation Review, vol. 45 (3), pp. 434-445, May. 2009. DOI: http://doi.org/10.1016/j. tre.2008.10.003.
[31] A. Giraldo and J. Caicedo, Construcción de matriz origen destino de transporte para amco (2008-
2012), Facultad de Ingeniería Industrial. Universidad Tecnológica de Pereira, Pereira 2010.
[32] F. Carmona, E. Guillen, Estudio de factibilidad para la creación de rutas exclusivas de transporte en los municipios de Pereira y Dosquebradas para la comunidad estudiantil de jornada diurna de la Universidad Tecnológica de Pereira, Ingeniería Industrial, Universidad Tecnológica de Pereira, Pereira 2014.


[^0]:    * M.Sc. Universidad Tecnológica de Pereira (Pereira-Risaralda, Colombia). dospina@utp.edu.co.
    ** Ph.D. Universidad Tecnológica de Pereira (Pereira-Risaralda, Colombia). elianam@utp.edu.co.
    *** Ph.D. Universidad Tecnológica de Pereira (Pereira-Risaralda, Colombia). ragr@utp.edu.co.

