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# A programming language view to model-driven engineering

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# Outline

- > **Model-Based System/Software Engineering vs. the real world**
- > **AADL, an overview**
- > **MBSE as an extension to programming in the large**

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# Engineering cycles

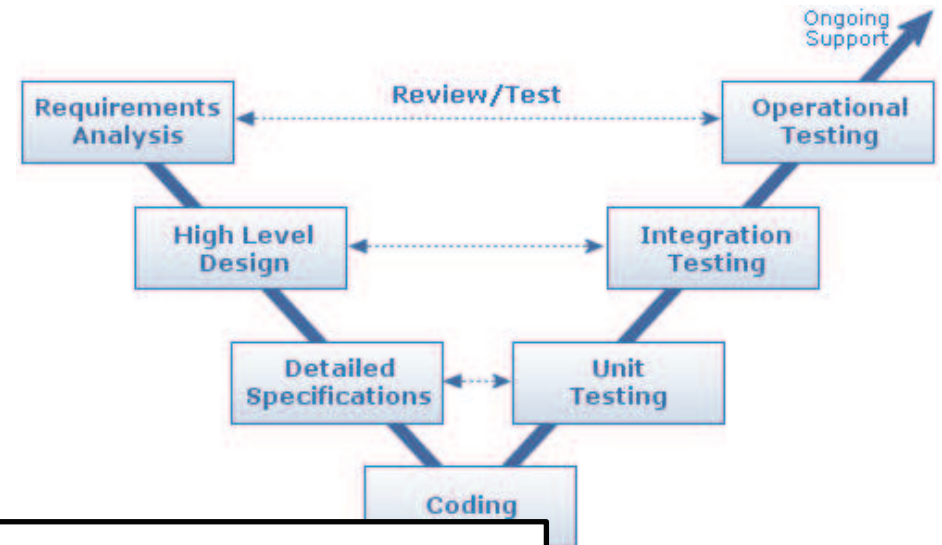
## > Now typical in SW engineering

» Various refinements

+ traceability across layers

+ split across teams (HW, SW, ...)

and consolidation



**Question:** where are the difficulties?

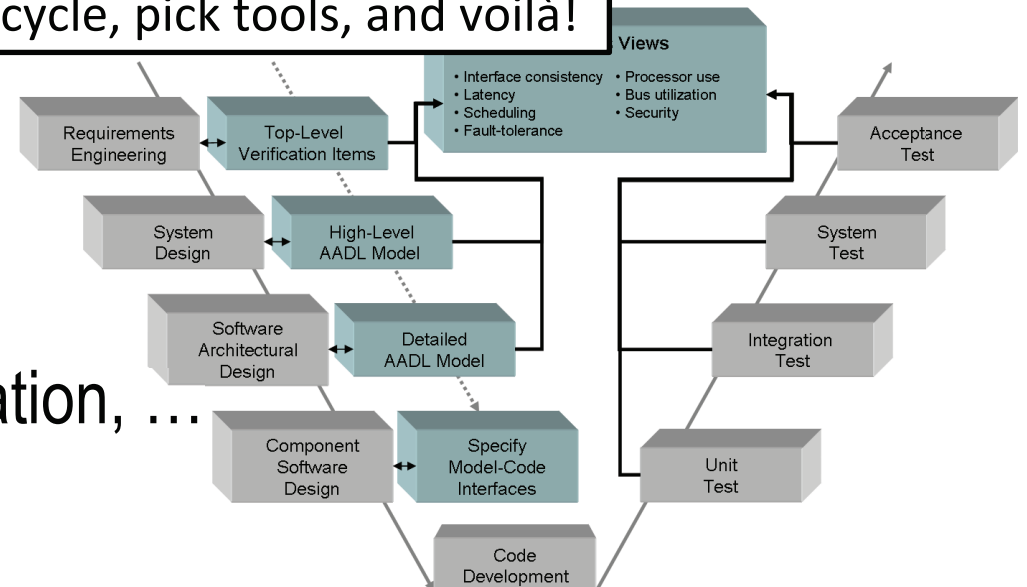
One could simply apply V-cycle, pick tools, and voilà!

## > Supported by AADL

» Scheduling, safety analysis

» Model checking, Code generation, ...

» Single architectural model



# Why is analysis in a V-cycle so difficult? (System Engineers 1 – 0 Software Engineers)

## SECURITY

Intrusion

Increased confidentiality requirement

- change of encryption policy

**Lesson#1:** Iterative V-cycle is ineffective: any design choice could jeopardize previous results.

Need to reconcile many domains, many approaches with pre-existent different tool supports (or lack of – Excel/Visio ..)

Also, SW is a **minor** part of the model compared to properties, behavior, interactions, etc.

## REAL-TIME PERFORMANCE

Deadlock/Starvation

Latency

Execution Time/Deadline

- increases CPU utilization
- increases power consumption
- may increase latency

Confidence

changes

es

on

n

plexity

By Feiler and Lewis

# Why is model-based so difficult?

## > Order of complexity (gratuitous comparison)

» Mathematics: axioms + proof, no interpretation

» ~~Electronics physics, structural physics, human factors, ergonomics~~

**Lesson:** Under-specified MDE process incurs 4 out of 7 « wastes »

» depicted in Lean Management

1. **Waiting:** some analysis are time-consuming, e.g. model checking

> **M** 2. **Over-processing:** when to trigger analysis?

» 3. **Over-production:** model too verbose/detailed

» 4. **Defects:** late discovery of model inaccuracies

» ⇒ Common root: analysis is a button in the GUI, not part of the model

> **A** Need to reflect on analysis process: prog. language can help!

» E.g. scheduling analysis: true/false

» Safety: error rate, stop when below threshold

» Also, Analysis part of the GUI space, not the modeling space!

ation”

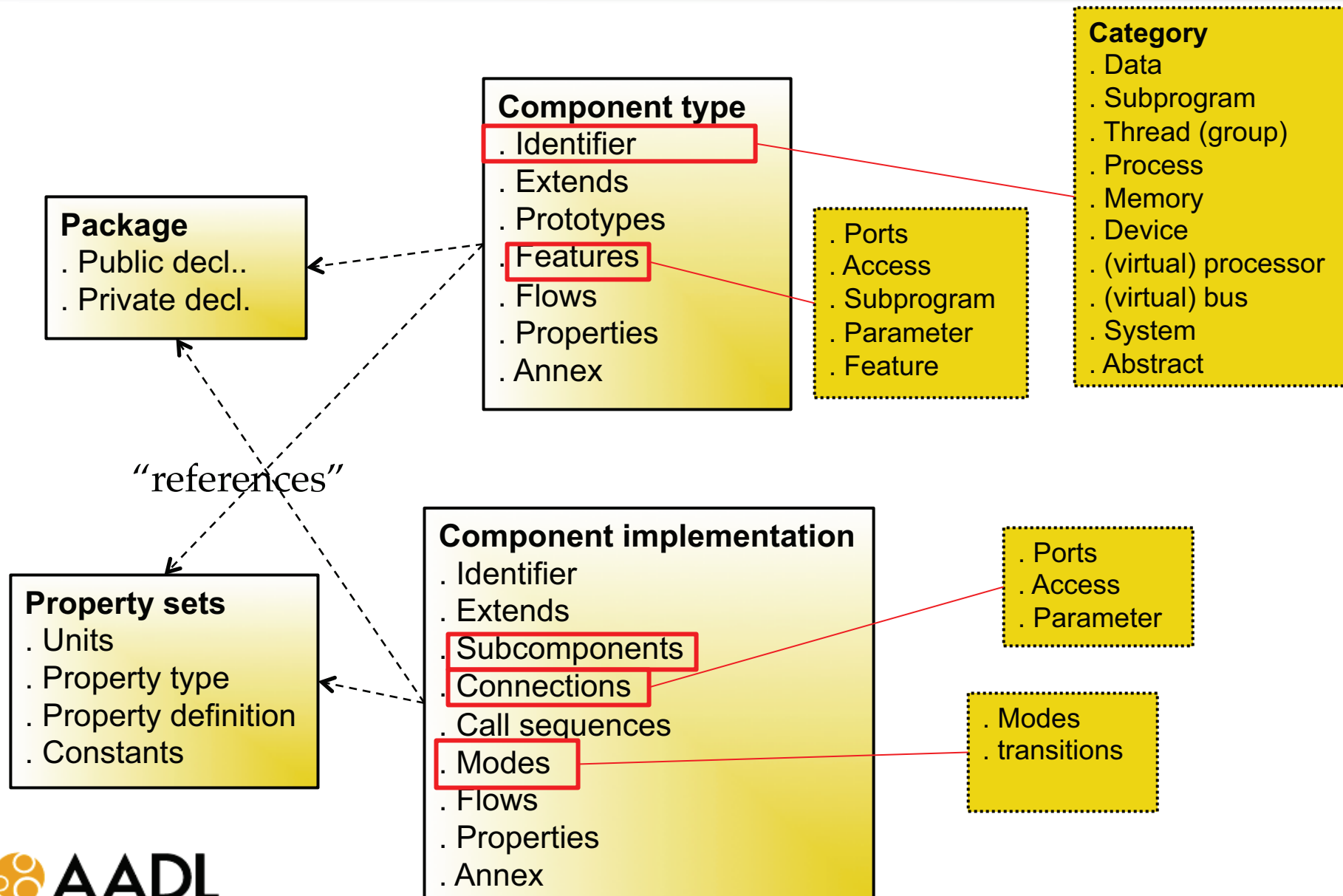
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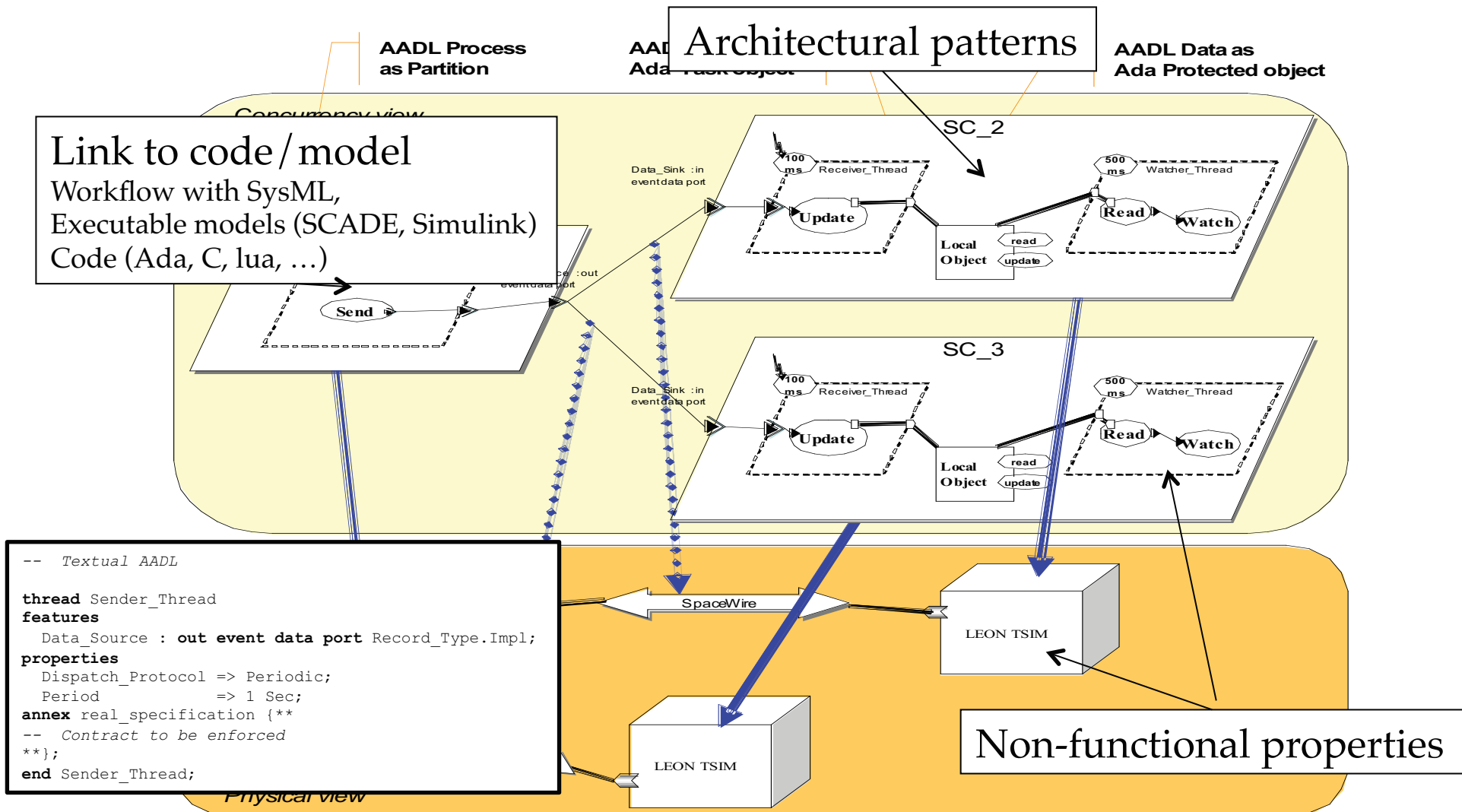
- > **International standard promoted by SAE International, AS-2C committee, released as AS-5506A**
- > **Version 1.0 published in 2004, version 2 in 2009**
  - » Committee driven by inputs from the avionics and space industry
  - » Academics drive analysis capability, to ensure they match with modeling patterns
- > **<http://aadl.info> list all resources around AADL**
  - » Public wiki with lot of resources: [https://wiki.sei.cmu.edu/aadl/index.php/Main\\_Page](https://wiki.sei.cmu.edu/aadl/index.php/Main_Page)
  - » Include link to most research activities around AADL
- > **AADL is dedicated to real-time embedded domain**
  - Modeling software and hardware resources for V&V
  - Extension & refinements concept to iterate down to generation
- > **Different representations**
  - » Graphical: high-level view of the system
  - » Textual: to view all details
  - » XML: to ease processing by 3rd party tool
- > **Some interactions with SysML (higher-level design)**

# AADL model elements



# AADL in one slide (!)

Architecture helps you focusing on the actual system



# AADL Extensions

- > **AADL is meant to be extensible**
- > **New property sets for specific concerns: e.g. ARINC653**
- > **Additional language to extend semantics**
  - » Behavioral specifications: AADL-BA
  - » Error modeling, propagation in a system: AADL-EMV2
  - » Constraints on model (on going)
    - Algebraic specifications for contracts, patterns, ...
  - » Requirement engineering (on going)
- > **Each extension has to remain compatible with core**
  - » Can be safely ignored if not relevant for a particular objective

# Some examples of AADL tool support

- > **AADL as a backbone, federating multiple activities**
  - » analysis through generation of intermediate models + external tools
- > **Common tool IDE: OSATE2 from SEI (FLOSS)**
  - » AADL core (SEI) + Behavioral (TPT) + Error (SEI) annexes
- > **Non exhaustive list of tools, European-centric (see <http://www.aadl.info>)**
  - » Integration to a process: with SysML, Simulink, SCADE
  - » Architectural pattern checks: MILS, ARINC, Ravenscar, Synchronous
  - » Model checking:
    - Timed/Stochastic/Colored Petri Nets
    - Timed automata et al.: UPPAAL, Versa, TASM
  - » Scheduling: MAST, Cheddar, CARTS
  - » Performance evaluation: real-time and network calculus
  - » Fault analysis: COMPASS, Stochastic Petri Nets, PRISM, FHA
  - » Simulation: ADeS, Marzhin
  - » Energy consumption of SoC: OpenPeople project
  - » Code generation: SystemC, C, Ada, RTSJ, Lustre
  - » WCET analysis: mapping to Bound-T

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# Moving back to programming language

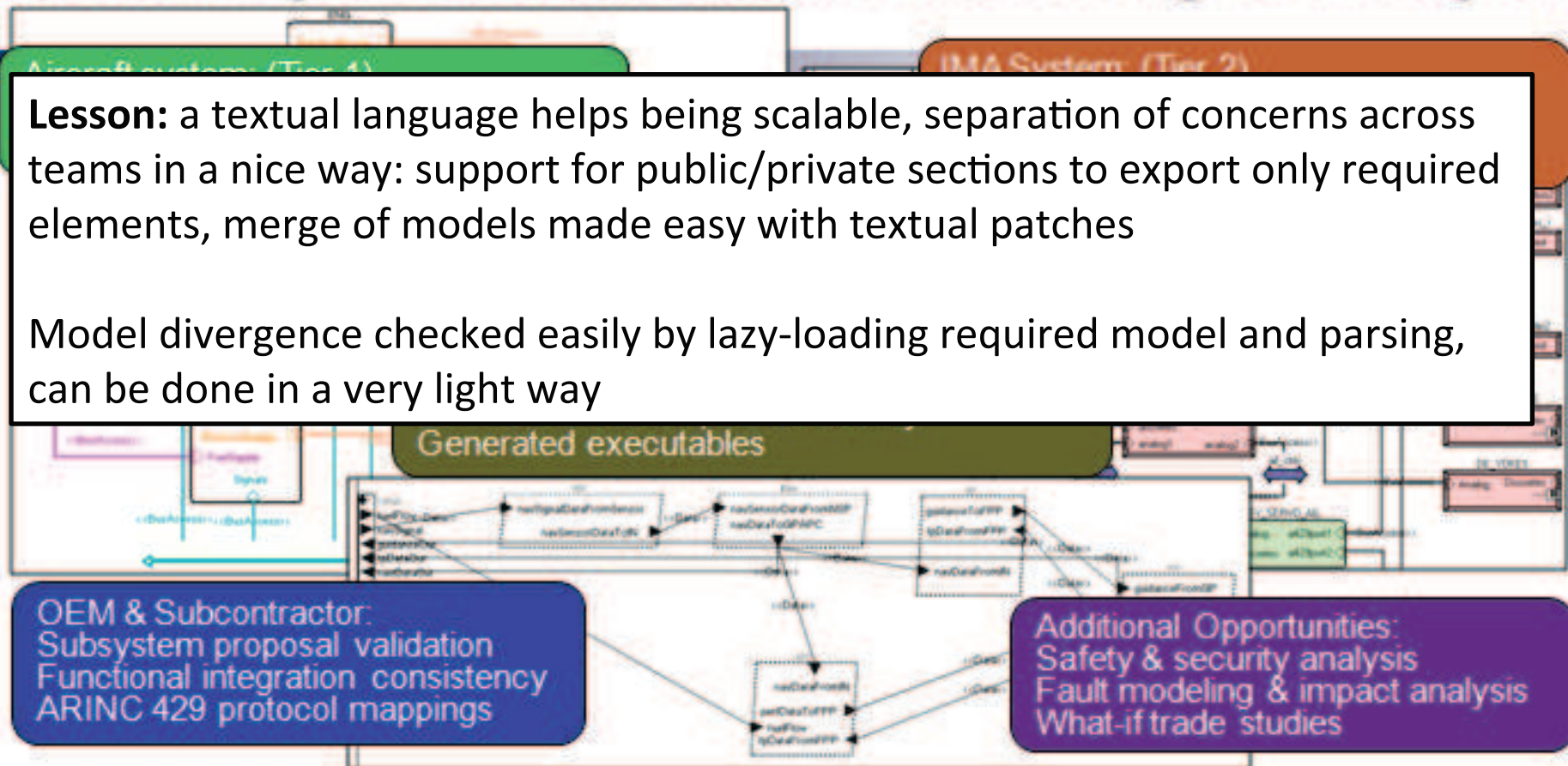
- > **AADL has a concrete syntax**
  - » Concrete means also rock-solid to build foundation
- > **Scalable: AADL package system close to Ada one**
  - » Potential for modular processing
  - » Optimizations in representation/processing of the AST
    - OSATE2: EMF, issues with object ids and cache
    - Ocarina: GNAT-like tree: faster, leaner
- > **Text also means potential for detailed syntactic constructs**
  - » Liskov principle, multiple bindings, formal specs, etc
  - » Cannot be (easily) represented graphically !

# Example#1: SAVI <http://www.avsi.aero>

## Incremental Multi-Fidelity Multi-dimensional Multi-Layered Architecture Modeling & Analysis

**Lesson:** a textual language helps being scalable, separation of concerns across teams in a nice way: support for public/private sections to export only required elements, merge of models made easy with textual patches

Model divergence checked easily by lazy-loading required model and parsing, can be done in a very light way



Use of AADL to cover a whole modeling cycle, focusing on validation of high level budgets (mass, energy), interface consistencies, etc.  
Modeling teams scattered across multiple teams and companies



## Example#2: TASTE <http://assert-project.net/-TASTE->

- > **Code generation and analysis for Space-critical systems**
  - » Subset of AADL as input language + model transformations

**Lesson:** a textual language, free from meta-model management issues is a must to avoid maintenance issues.

TASTE is 7 years old (!), each layer evolves independently, coordinated by an orchestrator

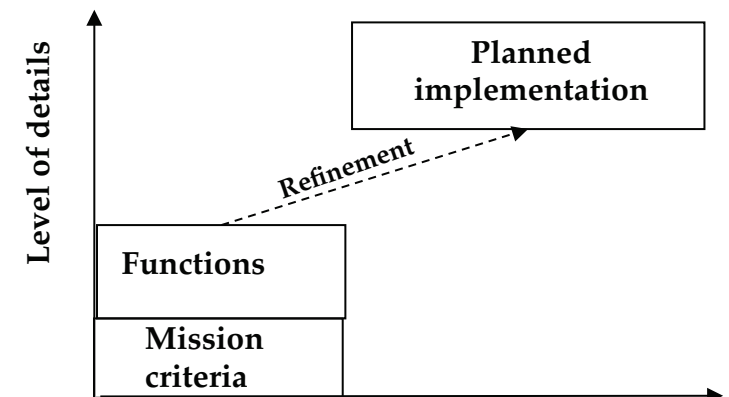
- ⇒ Each tool either reuses one existing parser (from Ocarina or ANTLR);
- ⇒ Or simple regexp to find the information it needs.

Simply follow Unix philosophy to address a complex transformation issue

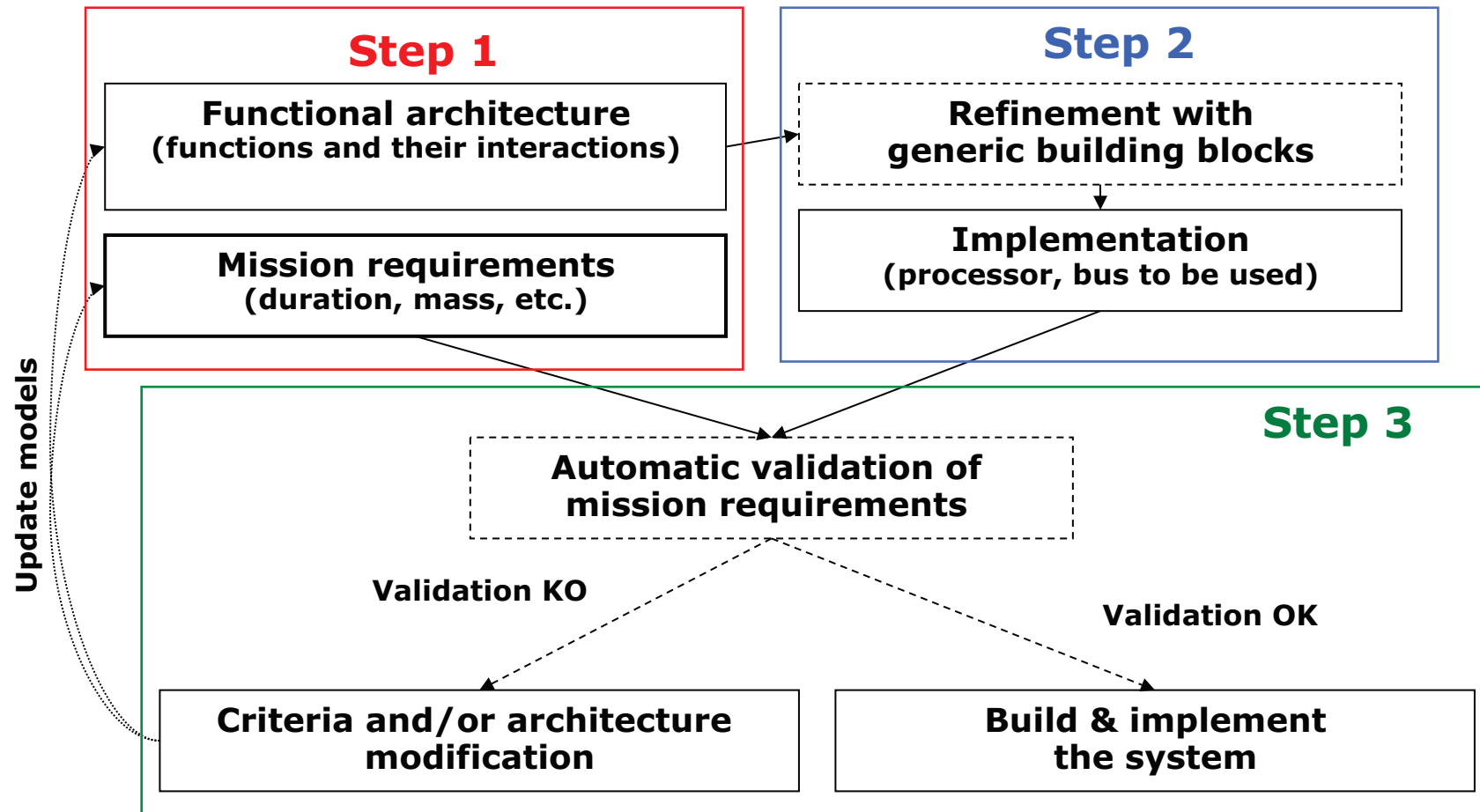
```
safe-mode,  
switch-to-redundant,  
...  
}  
  
AOCS-tm ::= SEQUENCE {  
attitude Attitude-ty,  
orbit Orbit-ty,  
...  
}
```

# Example#3: ARAM (joint project with ESA, 2011)

- > **Based on current practice for space projects at ESA**
- > **Define mission criteria**
  - » Max weight, orbit position, duration, etc.
- > **Specify functional aspects**
  - » What will be provided by the platform
  - » Specify requirements & constraints
- > **Refine the architecture**
  - » Replace functions by implementation
  - » Reuse existing components
- > **Validate planned implementation**
  - » Implementation properties vs. Function requirements
  - » Automate system integration verification



# ARAM Proposed approach, cont'd



# System exploration, design, integration

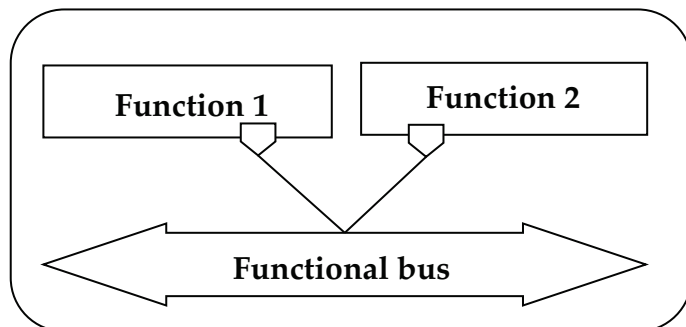
```
abstract function1
features
  ba : requires bus access
      genericbus;
end function1;
```

```
system implementation mission.i
subcomponents
  f1 : abstract function1;
  f2 : abstract function2;
  b  : bus genericbus;
connections
  bus access f1.ba -> b;
  bus access f2.ba -> b;
annex Constraints {**
  -- list of contracts to be met
**};
end mission.i;
```

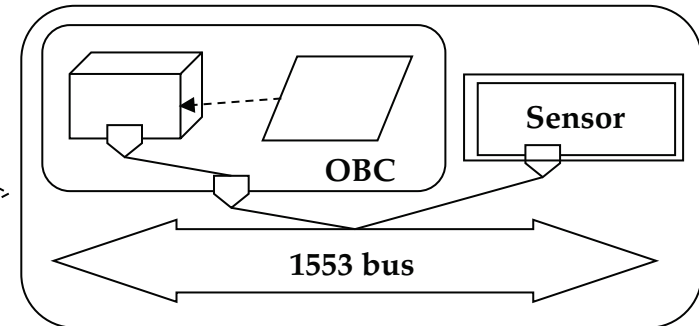
```
system obc
features
  ba : requires bus
      access bus1553;
end obc;
```

```
device sensor
features
  ba : requires bus
      access bus1553;
end sensor;
```

```
system implementation mission.planned
  extends mission.i
subcomponents
  f1 : refined to system obc;
  f2 : refined to device sensor;
  bus : refined to bus1553;
-- contracts inherited from mission.i
end mission.i;
```



Architecture refinement



# Contract example

```
-- gaia::functions
abstract fpa_data_get
features
  dataout      : out data port Data_Types::fpa_data;
  ctrlout      : out data port Data_Types::fpa_ctrl;
end fpa_data_get;
-- blocks
device FPA
features
  dataout : out data port Data_Types::FPA_data;
  ctrlout : out data port Data_Types::FPA_ctrl;
properties
  ARAM_Properties::Realizes =>
    (classifier (GAIA::Functions::FPA_data_get));
end FPA;
-- gaia::validation
system implementation Gaia.Validation
subcomponents
  Functional : system GAIA::Functions::Gaia.Functional;
  Impl       : system GAIA::Implementations::Gaia.First_Architecture;
properties
  ARAM_Properties::Actual_Function_Binding =>
    (reference (Functional.get1)) applies to Impl.fp1.datapart;
```

« Interface » of a function

One implementation

Function/implementation  
mapping

# Function coverage

**Lesson:** use a common language to model the architecture, **shared** by system engineer and software/hardware engineers

=> Each « role » can model its facet of the model

Syntax and semantics of AADL to bind them all, (like a programming language !)

External model bound to architecture (SysML, Simulink, DOORS, ...) for detailed info

=> Each level is attached its own set of verification (constraints, checks, computation, ..) and associated evaluation tool

Refined entities may « inherit » constraints from parent (à-la Liskov)

⇒ Verification rules using DSL evaluate specific architecture patterns,

⇒ “if A is connected to B, then the bandwidth of the bus used is less than ..”

⇒ Part of SAE AS2-C standardization effort

⇒ Nice side-effect: can be used to enforce requirements, subsets, contracts

⇒ Used for ARINC, MILS, Ravenscar architectural styles using Ocarina

|                     |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |
|---------------------|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| Impl.U3_1.runtime   | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x |
| Impl.transportlayer | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x |

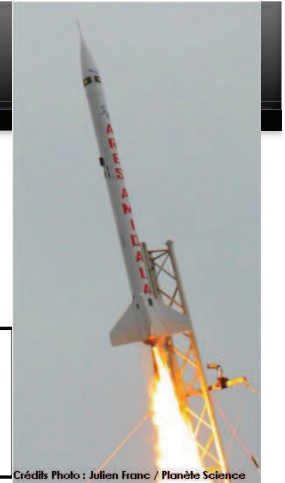
# AADL Constraint Language

- > **Work in progress as part of SAE AS2-C committee work**
  - » Defines accessors and computation rules on model elements
- > **E.g. AADLv2 and ARINC653 annex support IMA concepts**
  - » Needs to constraint models to respect some invariants

```
theorem scheduling_major_frame -- Check configuration of partition scheduling  
  foreach cpu in processor_set do  
    check ((float (property (cpu, "ARINC653::Module_Major_Frame")) =  
            sum (property (cpu, "ARINC653::Partition_Slots"))));  
end scheduling_major_frame;
```

```
system IMA_System extends AADL_System - implementation/extension must respect profile  
annex real_specification {**  
  theorem check_IMA  
    foreach s in local_set do  
      requires (check_IMA_profile); -- logical conjunction of theorems  
end check_IMA; **};  
end IMA_System;
```

# Example#4: PERSEUS supersonic rocket

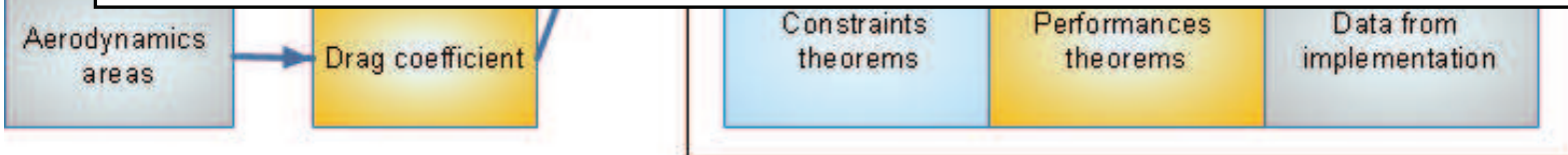


## > Analysis of rocket kinematics performance

System-level analysis done by combining atomic computations  
Each defined separately

**Lesson# 2:** verification should be driven by domain expert  
Expert knows what to compute, the dependencies between parameters  
Architects will attach analysis rule to model entities they apply to  
⇒ use of AADL annex subclause mechanism

**Lesson# 2bis:** notion of ordering of rules (*makefile-like*)  
Some properties deduced from analysis, reused in another analysis  
⇒ Resolution in a compiler AST  
⇒ Need also caching and “semantic timestamping”







## Wrap-up (System Engineers 1 – 1 Software Engineers)

- > **Equating Model-Based Analysis and Compiling is appealing**
  - » Text-based allows for optimization and more precise semantics
    - Fast evaluation for static/simple contracts, proof for complex one (BLESS)
    - Integration of IEEE PSL (dynamic traces) for observers
  - » Links with analysis tools made easy
- > **Integrating analysis contract to models helps solving**
  - » Waiting, Over-processing, Over-production, Defects
  - » A compiler/makefile-like approach would optimize analysis effort
    - Run only when required (i.e. model changed “significantly”)
- > **Integrating contracts as model elements, and analysis as compilation steps allow for better usage of designer time, and split: analysis designer vs. system designer**