



Coloring Curves that Cross a Fixed Curve

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Abstract

We prove that for every integer $t \geq 1$, the class of intersection graphs of curves in the plane each of which crosses a fixed curve in at least one and at most t points is χ -bounded. This is essentially the strongest χ -boundedness result one can get for those kind of graph classes. As a corollary, we prove that for any fixed integers $k \geq 2$ and $t \geq 1$, every k -quasi-planar topological graph on n vertices with any two edges crossing at most t times has $O(n \log n)$ edges.

Keywords String graphs · χ -Boundedness · k -Quasi-planar graphs

Mathematics Subject Classification 05C62 · 05C15

1 Introduction

Overview

A *curve* is a homeomorphic image of the real interval $[0, 1]$ in the plane. The *intersection graph* of a family of curves has these curves as vertices and the intersecting

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pairs of curves as edges. Combinatorial and algorithmic aspects of intersection graphs of curves, known as *string graphs*, have been attracting researchers for decades. A significant part of this research has been devoted to understanding classes of string graphs that are χ -bounded, which means that every graph G in the class satisfies $\chi(G) \leq f(\omega(G))$ for some function $f: \mathbb{N} \rightarrow \mathbb{N}$. Here, $\chi(G)$ and $\omega(G)$ denote the chromatic number and the clique number (the maximum size of a clique) of G , respectively. Although the class of all string graphs is not χ -bounded [24,25], known constructions of string graphs with small clique number and large chromatic number are very specific. One can expect that an appropriate restriction on the arrangement of the curves will prevent such constructions and thus lead to a χ -bounded subclass of the class of string graphs.

McGuinness [19,20] proposed studying families of curves that cross a fixed curve *exactly once*. This initiated a series of results culminating in the proof that the class of intersection graphs of such families is indeed χ -bounded [26]. By contrast, the class of intersection graphs of curves each crossing a fixed curve *at least once* is equal to the class of all string graphs and therefore is not χ -bounded. We prove an essentially strongest possible generalization of the former result, allowing curves to cross the fixed curve *at least once and at most t times*, for any bound t .

Theorem 1 *For every integer $t \geq 1$, the class of intersection graphs of curves each crossing a fixed curve in at least one and at most t points is χ -bounded.*

Additional motivation for Theorem 1 comes from its application to bounding the number of edges in so-called k -quasi-planar graphs, which we discuss at the end of this introduction.

Context

Colorings of intersection graphs of geometric objects have been investigated since the 1960s, when Asplund and Grünbaum [3] proved that intersection graphs of axis-parallel rectangles in the plane satisfy $\chi = O(\omega^2)$ and conjectured that the class of intersection graphs of axis-parallel boxes in \mathbb{R}^d is χ -bounded for every integer $d \geq 1$. A few years later Burling [5] discovered a surprising construction of triangle-free intersection graphs of axis-parallel boxes in \mathbb{R}^3 with arbitrarily large chromatic number. Since then, the upper bound of $O(\omega^2)$ and the trivial lower bound of $\Omega(\omega)$ on the maximum possible chromatic number of a rectangle intersection graph have been improved only in terms of multiplicative constants [11,13].

Another classical example of a χ -bounded class of geometric intersection graphs is provided by circle graphs—intersection graphs of chords of a fixed circle. Gyárfás [10] proved that circle graphs satisfy $\chi = O(\omega^2 4^\omega)$. The best known upper and lower bounds on the maximum possible chromatic number of a circle graph are $O(2^\omega)$ [14] and $\Omega(\omega \log \omega)$ [12,13].

McGuinness [19,20] proposed investigating the problem in a setting that allows much more general geometric shapes but restricts the way how they are arranged in the plane. In [19], he proved that the class of intersection graphs of L-shapes crossing a fixed horizontal line is χ -bounded. Families of L-shapes in the plane are *simple*, which means that any two members of the family intersect in at most one point. McGuinness

[20] also showed that triangle-free intersection graphs of simple families of curves each crossing a fixed line in exactly one point have bounded chromatic number. Further progress in this direction was made by Suk [27], who proved that simple families of x -monotone curves crossing a fixed vertical line give rise to a χ -bounded class of intersection graphs, and by Lasoń et al. [17], who reached the same conclusion without assuming that the curves are x -monotone. Finally, in [26], we proved that the class of intersection graphs of curves each crossing a fixed line in exactly one point is χ -bounded. These results remain valid if the fixed straight line is replaced by a fixed curve [28].

The class of string graphs is not χ -bounded. Pawlik et al. [24,25] showed that Burling's construction for boxes in \mathbb{R}^3 can be adapted to provide a construction of triangle-free intersection graphs of straight-line segments (or geometric shapes of various other kinds) with chromatic number growing as fast as $\Theta(\log \log n)$ with the number of vertices n . It was further generalized to a construction of string graphs with clique number ω and chromatic number $\Theta_\omega((\log \log n)^{\omega-1})$ [16]. The best known upper bound on the chromatic number of string graphs in terms of the number of vertices is $(\log n)^{O(\log \omega)}$, proved by Fox and Pach [8] using a separator theorem for string graphs due to Matoušek [18]. For intersection graphs of segments and, more generally, x -monotone curves, upper bounds of the form $\chi = O_\omega(\log n)$ follow from the above-mentioned results in [27] and [26] via recursive halving. Upper bounds of the form $\chi = O_\omega((\log \log n)^{f(\omega)})$ (for some function $f: \mathbb{N} \rightarrow \mathbb{N}$) are known for very special classes of string graphs: rectangle overlap graphs [15,16] and subtree overlap graphs [16]. The former still allow the triangle-free construction with $\chi = \Theta(\log \log n)$ and the latter the construction with $\chi = \Theta_\omega((\log \log n)^{\omega-1})$.

Quasi-Planarity

A *topological graph* is a graph with a fixed curvilinear drawing in the plane. For $k \geq 2$, a *k -quasi-planar graph* is a topological graph with no k pairwise crossing edges. In particular, a 2-quasi-planar graph is just a planar graph. It is conjectured that k -quasi-planar graphs with n vertices have $O_k(n)$ edges [4,23]. For $k = 2$, this asserts a well-known property of planar graphs. The conjecture is also verified for $k = 3$ [2,22] and $k = 4$ [1], but it remains open for $k \geq 5$. The best known upper bounds on the number of edges in a k -quasi-planar graph are $n(\log n)^{O(\log k)}$ in general [7,8], $O_k(n \log n)$ for the case of x -monotone edges [9,29], $O_k(n \log n)$ for the case that any two edges intersect at most once [28], and $2^{\alpha(n)^v} n \log n$ for the case that any two edges intersect in at most t points, where α is the inverse Ackermann function and v depends on k and t [28]. We apply Theorem 1 to improve the last bound to $O_{k,t}(n \log n)$.

Theorem 2 *Every k -quasi-planar topological graph G on n vertices such that any two edges of G intersect in at most t points has at most $\mu_{k,t} n \log n$ edges, where $\mu_{k,t}$ depends only on k and t .*

The proof follows the same lines as the proof in [28] for the case $t = 1$ (see Sect. 3).

2 Proof of Theorem 1

Setup

We let \mathbb{N} denote the set of positive integers. Graph-theoretic terms applied to a family of curves \mathcal{F} have the same meaning as applied to the intersection graph of \mathcal{F} . In particular, the *chromatic number* of \mathcal{F} , denoted by $\chi(\mathcal{F})$, is the minimum number of colors in a *proper coloring* of \mathcal{F} (a coloring that distinguishes pairs of intersecting curves), and the *clique number* of \mathcal{F} , denoted by $\omega(\mathcal{F})$, is the maximum size of a *clique* in \mathcal{F} (a set of pairwise intersecting curves in \mathcal{F}).

Theorem 1 (Rephrased) *For every $t \in \mathbb{N}$, there is a non-decreasing function $f_t: \mathbb{N} \rightarrow \mathbb{N}$ with the following property: for any fixed curve c_0 , every family \mathcal{F} of curves each intersecting c_0 in at least one and at most t points satisfies $\chi(\mathcal{F}) \leq f_t(\omega(\mathcal{F}))$.*

We do not state any explicit bound on the function f_t above, because it depends highly on the bound on the function f in Theorem 13 (one of our main tools), and no explicit bound on f is provided in [6]. We assume (implicitly) that the intersection points of all curves $c \in \mathcal{F}$ with c_0 considered in Theorem 1 are distinct and each of them is a *proper crossing*, which means that c passes from one to the other side of c_0 in a sufficiently small neighborhood of the intersection point. We can make this assumption without loss of generality, as it can be guaranteed by appropriate small perturbations of the curves that do not influence the intersection graph.

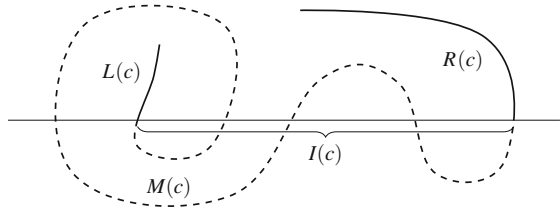
Initial Reduction

We start by reducing Theorem 1 to a somewhat simpler and more convenient setting. We fix a horizontal line in the plane and call it the *baseline*. The upper closed half-plane determined by the baseline is denoted by H^+ . A *1-curve* is a curve in H^+ that has one endpoint (called the *basepoint* of the 1-curve) on the baseline and does not intersect the baseline in any other point. Intersection graphs of 1-curves are known as *outerstring graphs* and form a χ -bounded class of graphs—this result, due to the authors, is the starting point of the proof of Theorem 1.

Theorem 3 ([26]) *There is a non-decreasing function $f_0: \mathbb{N} \rightarrow \mathbb{N}$ such that every family \mathcal{F} of 1-curves satisfies $\chi(\mathcal{F}) \leq f_0(\omega(\mathcal{F}))$.*

An *even-curve* is a curve that has both endpoints above the baseline and a positive even number of intersection points with the baseline, each of which is a proper crossing. For $t \in \mathbb{N}$, a *2t-curve* is an even-curve that intersects the baseline in exactly $2t$ points. A *basepoint* of an even-curve c is an intersection point of c with the baseline. Like above, we assume (implicitly, without loss of generality) that the basepoints of all even-curves in any family that we consider are distinct. Every even-curve c determines two 1-curves—the two parts of c from an endpoint to the closest basepoint along c . They are called the *1-curves of c* and denoted by $L(c)$ and $R(c)$ so that the basepoint of $L(c)$ lies to the left of the basepoint of $R(c)$ on the baseline (see Fig. 1). A family \mathcal{F} of even-curves is an *LR-family* if every intersection between two curves $c_1, c_2 \in \mathcal{F}$ is

Fig. 1 $L(c)$, $R(c)$, $M(c)$ (all the dashed part), and $I(c)$ for a 6-curve c



an intersection between $L(c_1)$ and $R(c_2)$ or between $L(c_2)$ and $R(c_1)$. The main effort in this paper goes to proving the following statement on LR -families of even-curves.

Theorem 4 *There is a non-decreasing function $f : \mathbb{N} \rightarrow \mathbb{N}$ such that every LR -family \mathcal{F} of even-curves satisfies $\chi(\mathcal{F}) \leq f(\omega(\mathcal{F}))$.*

Theorem 4 makes no assumption on the maximum number of intersection points of an even-curve with the baseline. We derive Theorem 1 from Theorem 4 in two steps, first proving the following lemma, and then showing that Theorem 1 is essentially a special case of it.

Lemma 5 *For every $t \in \mathbb{N}$, there is a non-decreasing function $f_t : \mathbb{N} \rightarrow \mathbb{N}$ such that every family \mathcal{F} of $2t$ -curves no two of which intersect below the baseline satisfies $\chi(\mathcal{F}) \leq f_t(\omega(\mathcal{F}))$.*

Proof of Lemma 5 from Theorem 4 The proof goes by induction on t . Let f_0 and f be the functions claimed by Theorems 3 and 4, respectively, and let $f_t(k) = f_{t-1}^2(k) f(k)$ for $t \geq 1$ and $k \in \mathbb{N}$. We establish the base case for $t = 1$ and the induction step for $t \geq 2$ simultaneously. Namely, fix an integer $t \geq 1$, and let \mathcal{F} be as in the statement of the lemma. For every $2t$ -curve $c \in \mathcal{F}$, enumerate the endpoints and basepoints of c as $p_0(c), \dots, p_{2t+1}(c)$ in their order along c so that $p_0(c)$ and $p_1(c)$ are the endpoints of $L(c)$ while $p_{2t}(c)$ and $p_{2t+1}(c)$ are the endpoints of $R(c)$. Build two families of curves \mathcal{F}_1 and \mathcal{F}_2 putting the part of c from $p_0(c)$ to $p_{2t-1}(c)$ to \mathcal{F}_1 and the part of c from $p_2(c)$ to $p_{2t+1}(c)$ to \mathcal{F}_2 for every $c \in \mathcal{F}$. If $t = 1$, then \mathcal{F}_1 and \mathcal{F}_2 are families of 1-curves. If $t \geq 2$, then \mathcal{F}_1 and \mathcal{F}_2 are equivalent to families of $2(t - 1)$ -curves, because the curve in \mathcal{F}_1 or \mathcal{F}_2 obtained from a $2t$ -curve $c \in \mathcal{F}$ can be shortened a little at $p_{2t-1}(c)$ or $p_2(c)$, respectively, losing that basepoint but no intersection points with other curves. Therefore, by Theorem 3 or the induction hypothesis, we have $\chi(\mathcal{F}_k) \leq f_{t-1}(\omega(\mathcal{F}_k)) \leq f_{t-1}(\omega(\mathcal{F}))$ for $k \in \{1, 2\}$. For $c \in \mathcal{F}$ and $k \in \{1, 2\}$, let $\phi_k(c)$ be the color of the curve obtained from c in an optimal proper coloring of \mathcal{F}_k . Every subfamily of \mathcal{F} on which the pair $(\phi_1(c), \phi_2(c))$ is constant is an LR -family and therefore, by Theorem 4 and monotonicity of f , has chromatic number at most $f(\omega(\mathcal{F}))$. We conclude that $\chi(\mathcal{F}) \leq \chi(\mathcal{F}_1)\chi(\mathcal{F}_2)f(\omega(\mathcal{F})) \leq f_{t-1}^2(\omega(\mathcal{F}))f(\omega(\mathcal{F})) = f_t(\omega(\mathcal{F}))$. \square

A closed curve is a homeomorphic image of a unit circle in the plane. For a closed curve γ , the Jordan curve theorem asserts that the set $\mathbb{R}^2 \setminus \gamma$ consists of two arc-connected components, one of which is bounded and denoted by $\text{int } \gamma$ and the other is unbounded and denoted by $\text{ext } \gamma$.

Proof of Theorem 1 from Theorem 4 We elect to present this proof in an intuitive rather than rigorous way. Let \mathcal{F} be a family of curves each intersecting c_0 in at least one and at most t points. Let γ_0 be a closed curve surrounding c_0 very closely so that γ_0 intersects every curve in \mathcal{F} in exactly $2t$ points (winding if necessary to increase the number of intersections) and all endpoints of curves in \mathcal{F} and intersection points of pairs of curves in \mathcal{F} lie in $\text{ext } \gamma_0$. We apply geometric inversion to obtain an equivalent family of curves \mathcal{F}' and a closed curve γ'_0 with the same properties except that all endpoints of curves in \mathcal{F}' and intersection points of pairs of curves in \mathcal{F}' lie in $\text{int } \gamma'_0$. It follows that some part of γ'_0 lies in the unbounded component of $\mathbb{R}^2 \setminus \bigcup \mathcal{F}'$. We “cut” γ'_0 there and “unfold” it into the baseline, transforming \mathcal{F}' into an equivalent family \mathcal{F}'' of $2t$ -curves all endpoints of which and intersection points of pairs of which lie above the baseline. The “equivalence” of \mathcal{F} , \mathcal{F}' , and \mathcal{F}'' means in particular that the intersection graphs of \mathcal{F} , \mathcal{F}' , and \mathcal{F}'' are isomorphic, so the theorem follows from Lemma 5 (and thus Theorem 4). \square

A statement analogous to Theorem 4 fails for families of objects each consisting of two 1-curves only, without the “middle part” connecting them. Specifically, we define a *double-curve* as a set $X \subset H^+$ that is a union of two disjoint 1-curves, denoted by $L(X)$ and $R(X)$ so that the basepoint of $L(X)$ lies to the left of the basepoint of $R(X)$, and we call a family \mathcal{X} of double-curves an *LR-family* if every intersection between two double-curves $X_1, X_2 \in \mathcal{X}$ is an intersection between $L(X_1)$ and $R(X_2)$ or between $L(X_2)$ and $R(X_1)$.

Theorem 6 *There exist triangle-free LR-families of double-curves with arbitrarily large chromatic number.*

The proof of Theorem 6 is an easy adaptation of the construction from [24,25] and is presented in detail in Sect. 4. The rest of this section is devoted to the proof of Theorem 4.

Overview of the Proof of Theorem 4

Recall the assertion of Theorem 4: the *LR*-families of even-curves are χ -bounded. The proof is quite long and technical, so we find it useful to provide a high-level overview of its structure. The proof will be presented via a series of reductions. First, we will reduce Theorem 4 to the following statement (Lemma 7): the *LR*-families of 2-curves are χ -bounded. This statement will be proved by induction on the clique number. Specifically, we will prove the following as the induction step: if every *LR*-family of 2-curves \mathcal{F} with $\omega(\mathcal{F}) \leq k - 1$ satisfies $\chi(\mathcal{F}) \leq \xi$, then every *LR*-family of 2-curves \mathcal{F} with $\omega(\mathcal{F}) \leq k$ satisfies $\chi(\mathcal{F}) \leq \zeta$, where ζ is a constant depending only on k and ξ . The only purpose of the induction hypothesis is to guarantee that if $\omega(\mathcal{F}) \leq k$ and $c \in \mathcal{F}$, then the family of 2-curves in $\mathcal{F} \setminus \{c\}$ that intersect c has chromatic number at most ξ . For convenience, *LR*-families of 2-curves with the latter property will be called ξ -families. We will thus reduce the problem to the following statement (Lemma 9): the ξ -families are χ -bounded, where the χ -bounding function depends on ξ .

We will deal with ξ -families via a series of technical lemmas of the following general form: every ξ -family with sufficiently large chromatic number contains a specific configuration of curves. Two kinds of such configurations are particularly important: (a) a large clique, and (b) a 2-curve c and a subfamily \mathcal{F}' with large chromatic number such that all basepoints of 2-curves in \mathcal{F}' lie between the basepoints of c . At the core of the argument are the proofs that

- every ξ -family with sufficiently large chromatic number contains (a) or (b) (Lemma 16),
- assuming the above, every ξ -family with sufficiently large chromatic number contains (a).

Combined, they complete the argument. Since the two proofs are almost identical, we introduce one more reduction—to (ξ, h) -families (Lemma 15). A (ξ, h) -family is just a ξ -family that satisfies an additional technical condition that allows us to deliver both proofs at once.

More Notation and Terminology

Let \prec denote the left-to-right order of points on the baseline ($p_1 \prec p_2$ means that p_1 is to the left of p_2). For convenience, we also use the notation \prec for curves intersecting the baseline ($c_1 \prec c_2$ means that every basepoint of c_1 is to the left of every basepoint of c_2) and for families of such curves ($\mathcal{C}_1 \prec \mathcal{C}_2$ means that $c_1 \prec c_2$ for any $c_1 \in \mathcal{C}_1$ and $c_2 \in \mathcal{C}_2$). For a family \mathcal{C} of curves intersecting the baseline (even-curves or 1-curves) and two 1-curves x and y , let $\mathcal{C}(x, y) = \{c \in \mathcal{C} : x \prec c \prec y\}$ or $\mathcal{C}(x, y) = \{c \in \mathcal{C} : y \prec c \prec x\}$ depending on whether $x \prec y$ or $y \prec x$. For a family \mathcal{C} of curves intersecting the baseline and a segment I on the baseline, let $\mathcal{C}(I)$ denote the family of curves in \mathcal{C} with all basepoints on I .

For an even-curve c , let $M(c)$ denote the subcurve of c connecting the basepoints of $L(c)$ and $R(c)$, and let $I(c)$ denote the segment on the baseline connecting the basepoints of $L(c)$ and $R(c)$ (see Fig. 1). For a family \mathcal{F} of even-curves, let $L(\mathcal{F}) = \{L(c) : c \in \mathcal{F}\}$, $R(\mathcal{F}) = \{R(c) : c \in \mathcal{F}\}$, and $I(\mathcal{F})$ denote the minimal segment on the baseline that contains $I(c)$ for every $c \in \mathcal{F}$.

A *cap-curve* is a curve in H^+ that has both endpoints on the baseline and does not intersect the baseline in any other point. It follows from the Jordan curve theorem that for every cap-curve γ , the set $H^+ \setminus \gamma$ consists of two arc-connected components, one of which is bounded and denoted by $\text{int } \gamma$ and the other is unbounded and denoted by $\text{ext } \gamma$.

Reduction to LR-Families of 2-Curves

We will reduce Theorem 4 to the following statement on LR-families of 2-curves, which is essentially a special case of Theorem 4.

Lemma 7 *There is a non-decreasing function $f : \mathbb{N} \rightarrow \mathbb{N}$ such that every LR-family \mathcal{F} of 2-curves satisfies $\chi(\mathcal{F}) \leq f(\omega(\mathcal{F}))$.*

A *component* of a family of 1-curves \mathcal{S} is an arc-connected component of $\bigcup \mathcal{S}$ (the union of all curves in \mathcal{S}). The following easy but powerful observation reuses an idea from [17,20,27].

Lemma 8 *For every LR-family of even-curves \mathcal{F} , if \mathcal{F}^* is the family of curves $c \in \mathcal{F}$ such that $L(c)$ and $R(c)$ lie in different components of $L(\mathcal{F}) \cup R(\mathcal{F})$, then $\chi(\mathcal{F}^*) \leq 4$.*

Proof Let G be an auxiliary graph where the vertices are the components of $L(\mathcal{F}) \cup R(\mathcal{F})$ and the edges are the pairs $V_1 V_2$ of components such that there is a curve $c \in \mathcal{F}^*$ with $L(c) \subseteq V_1$ and $R(c) \subseteq V_2$ or $L(c) \subseteq V_2$ and $R(c) \subseteq V_1$. Since \mathcal{F} is an LR-family, the curves in \mathcal{F}^* can intersect only within the components of $L(\mathcal{F}) \cup R(\mathcal{F})$. It follows that G is planar and thus 4-colorable. Fix a proper 4-coloring of G , and assign the color of a component V to every curve $c \in \mathcal{F}^*$ with $L(c) \subseteq V$. For any $c_1, c_2 \in \mathcal{F}^*$, if $L(c_1)$ and $R(c_2)$ intersect, then $L(c_1)$ and $R(c_2)$ lie in the same component V_1 while $L(c_2)$ lies in a component V_2 such that $V_1 V_2$ is an edge of G , so c_1 and c_2 are assigned different colors. The coloring of \mathcal{F}^* is therefore proper. \square

Proof of Theorem 4 from Lemma 7 We show that $\chi(\mathcal{F}) \leq f(\omega(\mathcal{F})) + 4$, where f is the function claimed by Lemma 7. We have $\mathcal{F} = \mathcal{F}_1 \cup \mathcal{F}_2$, where $\mathcal{F}_1 = \{c \in \mathcal{F} : L(c) \text{ and } R(c) \text{ lie in the same component of } L(\mathcal{F}) \cup R(\mathcal{F})\}$ and $\mathcal{F}_2 = \{c \in \mathcal{F} : L(c) \text{ and } R(c) \text{ lie in different components of } L(\mathcal{F}) \cup R(\mathcal{F})\}$. Lemma 8 yields $\chi(\mathcal{F}_2) \leq 4$. It remains to show that $\chi(\mathcal{F}_1) \leq f(\omega(\mathcal{F}))$.

Let $c_1, c_2 \in \mathcal{F}_1$. We claim that the intervals $I(c_1)$ and $I(c_2)$ are nested or disjoint. Suppose they are neither nested nor disjoint. The components of $L(\mathcal{F}) \cup R(\mathcal{F})$ are disjoint from the curves of the form $M(c)$ with $c \in \mathcal{F}$ except at common basepoints. For $k \in \{1, 2\}$, since $L(c_k)$ and $R(c_k)$ belong to the same component of $L(\mathcal{F}) \cup R(\mathcal{F})$, the basepoints of $L(c_k)$ and $R(c_k)$ can be connected by a cap-curve γ_k disjoint from $M(c)$ for every $c \in \mathcal{F}$ except at the endpoints of $M(c)$ when $c = c_k$. We assume (without loss of generality) that γ_1 and γ_2 intersect in a finite number of points and each of their intersection points is a proper crossing. Since the intervals $I(c_1)$ and $I(c_2)$ are neither nested nor disjoint, the basepoints of $L(c_2)$ and $R(c_2)$ lie one in $\text{int } \gamma_1$ and the other in $\text{ext } \gamma_1$. This implies that γ_1 and γ_2 intersect in an odd number of points, by the Jordan curve theorem. For $k \in \{1, 2\}$, let $\tilde{\gamma}_k$ be the closed curve obtained as the union of γ_k and $M(c_k)$. It follows that $\tilde{\gamma}_1$ and $\tilde{\gamma}_2$ intersect in an odd number of points and each of their intersection points is a proper crossing, which is a contradiction to the Jordan curve theorem.

Transform \mathcal{F}_1 into a family of 2-curves \mathcal{F}'_1 replacing the part $M(c)$ of every 2-curve $c \in \mathcal{F}_1$ by the lower semicircle connecting the endpoints of $M(c)$. Since the intervals $I(c)$ with $c \in \mathcal{F}_1$ are pairwise nested or disjoint, these semicircles are pairwise disjoint. Consequently, \mathcal{F}'_1 is an LR-family. Since the intersection graphs of \mathcal{F}_1 and \mathcal{F}'_1 are isomorphic, Lemma 7 implies $\chi(\mathcal{F}_1) = \chi(\mathcal{F}'_1) \leq f(\omega(\mathcal{F}'_1)) \leq f(\omega(\mathcal{F}))$. \square

Reduction to ξ -Families

For $\xi \in \mathbb{N}$, a ξ -family is an LR -family of 2-curves \mathcal{F} with the following property: for every 2-curve $c \in \mathcal{F}$, the family of 2-curves in $\mathcal{F} \setminus \{c\}$ that intersect c has chromatic number at most ξ . We reduce Lemma 7 to the following statement on ξ -families.

Lemma 9 *For any $\xi, k \in \mathbb{N}$, there is a constant $\zeta \in \mathbb{N}$ such that every ξ -family \mathcal{F} with $\omega(\mathcal{F}) \leq k$ satisfies $\chi(\mathcal{F}) \leq \zeta$.*

Proof of Lemma 7 from Lemma 9 Let $f(1) = 1$. For $k \geq 2$, let $f(k)$ be the constant claimed by Lemma 9 such that every $f(k-1)$ -family \mathcal{F} with $\omega(\mathcal{F}) \leq k$ satisfies $\chi(\mathcal{F}) \leq f(k)$. Let $k = \omega(\mathcal{F})$, and proceed by induction on k to prove $\chi(\mathcal{F}) \leq f(k)$. Clearly, if $k = 1$, then $\chi(\mathcal{F}) = 1$. For the induction step, assume $k \geq 2$. For every $c \in \mathcal{F}$, the family of 2-curves in $\mathcal{F} \setminus \{c\}$ that intersect c has clique number at most $k-1$ and therefore, by the induction hypothesis, has chromatic number at most $f(k-1)$. That is, \mathcal{F} is an $f(k-1)$ -family, and the definition of f yields $\chi(\mathcal{F}) \leq f(k)$. \square

Dealing with ξ -Families

First, we establish the following special case of Lemma 9.

Lemma 10 *For every $\xi \in \mathbb{N}$, every ξ -family \mathcal{F} with $\bigcap_{c \in \mathcal{F}} I(c) \neq \emptyset$ satisfies $\chi(\mathcal{F}) \leq 4\xi + 4$.*

The proof of Lemma 10 is essentially the same as the proof of Lemma 19 in [28]. We need the following elementary lemma, which was also used in various forms in [17,19,20,26,27]. We include its proof, as we will later extend it when proving Lemma 12.

Lemma 11 (McGuinness [19, Lem. 2.1]) *Let G be a graph, $<$ be a total order on the vertices of G , and $\alpha, \beta \in \mathbb{N}$. Let $G(u, v)$ denote the subgraph of G induced on the vertices strictly between u and v in the order $<$. If $\chi(G) > (2\beta + 2)\alpha$, then G has an induced subgraph H such that $\chi(H) > \alpha$ and $\chi(G(u, v)) > \beta$ for every edge uv of H . In particular, if $\chi(G) > 2\beta + 2$, then G has an edge uv with $\chi(G(u, v)) > \beta$.*

Proof Let $G[U]$ denote the subgraph of G induced on a set of vertices U . Partition the vertices of G into subsets $V_0 < \dots < V_n$ so that $\chi(G[V_i]) = \beta + 1$ for $0 \leq i < n$ and $\chi(G[V_n]) \leq \beta + 1$. This is done greedily, by processing the vertices of G in the order $<$, adding them to V_0 until $\chi(G[V_0]) = \beta + 1$, then adding them to V_1 until $\chi(G[V_1]) = \beta + 1$, and so on. For $0 \leq i \leq n$, a proper $(\beta + 1)$ -coloring of $G[V_i]$ yields a partition of V_i into color classes $V_i^1, \dots, V_i^{\beta+1}$ that are independent sets in G . Let $r \in \{1, \dots, \beta + 1\}$ be such that $\chi(G[\bigcup_{i=0}^n V_i^r])$ is maximized. It follows that $\chi(G[\bigcup_{i=0}^n V_i^r]) \geq \chi(G)/(\beta + 1) > 2\alpha$ and thus $\chi(G[\bigcup_{i \text{ even}} V_i^r]) > \alpha$ or $\chi(G[\bigcup_{i \text{ odd}} V_i^r]) > \alpha$. Let $H = G[\bigcup_{i \text{ even}} V_i^r]$ or $H = G[\bigcup_{i \text{ odd}} V_i^r]$ accordingly, so that $\chi(H) > \alpha$. Now, if uv is an edge of H , then $u \in V_k^r$ and $v \in V_\ell^r$ for two distinct indices $k, \ell \in \{0, \dots, n\}$ of the same parity (because each V_i^r is an independent set in G), and therefore $G[V_i]$ is a subgraph of $G(u, v)$ for every (at least one) index $i \in \{1, \dots, n-1\}$ strictly between k and ℓ , witnessing $\chi(G(u, v)) > \beta$. \square

Proof of Lemma 10 Suppose $\chi(\mathcal{F}) > 4\xi + 4$. Since $\bigcap_{c \in \mathcal{F}} I(c) \neq \emptyset$, the 2-curves in \mathcal{F} can be enumerated as c_1, \dots, c_n so that $L(c_1) < \dots < L(c_n) < R(c_n) < \dots < R(c_1)$, where $n = |\mathcal{F}|$. Lemma 11 applied to the intersection graph of \mathcal{F} and the order c_1, \dots, c_n provides two indices $i, j \in \{1, \dots, n\}$ such that the 2-curves c_i and c_j intersect and $\chi(\{c_{i+1}, \dots, c_{j-1}\}) > 2\xi + 1$. Assume $L(c_i)$ and $R(c_j)$ intersect; the argument for the other case is analogous. There is a cap-curve $v \subseteq L(c_i) \cup R(c_j)$ connecting the basepoints of $L(c_i)$ and $R(c_j)$. Every curve intersecting v intersects c_i or c_j . Since \mathcal{F} is a ξ -family, the 2-curves in $\{c_{i+1}, \dots, c_{j-1}\}$ that intersect c_i have chromatic number at most ξ , and so do those that intersect c_j . Every 2-curve $c_k \in \{c_{i+1}, \dots, c_{j-1}\}$ not intersecting v satisfies $L(c_k) \subset \text{int } v$ and $R(c_k) \subset \text{ext } v$, so these 2-curves are pairwise disjoint. We conclude that $\chi(\{c_{i+1}, \dots, c_{j-1}\}) \leq 2\xi + 1$, which is a contradiction. \square

It easily follows from Lemma 11 that every family of 2-curves \mathcal{F} with $\chi(\mathcal{F}) > (2\beta + 2)^2\alpha$ contains a subfamily \mathcal{H} with $\chi(\mathcal{H}) > \alpha$ such that any two intersecting 2-curves $c_1, c_2 \in \mathcal{H}$ satisfy $\chi(\mathcal{F}(L(c_1), L(c_2))) > \beta$ and $\chi(\mathcal{F}(R(c_1), R(c_2))) > \beta$. This is considerably strengthened by the following lemma. Its proof is based on the same general idea as the proof of Lemma 11 presented above.

Lemma 12 *For every $\xi \in \mathbb{N}$, there is a function $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ with the following property: for any $\alpha, \beta \in \mathbb{N}$ and every ξ -family \mathcal{F} with $\chi(\mathcal{F}) > f(\alpha, \beta)$, there is a subfamily $\mathcal{H} \subseteq \mathcal{F}$ such that $\chi(\mathcal{H}) > \alpha$ and $\chi(\mathcal{F}(x, y)) > \beta$ for any two intersecting 1-curves $x, y \in L(\mathcal{H}) \cup R(\mathcal{H})$.*

Proof Let $f(\alpha, \beta) = (2\beta + 12\xi + 20)\alpha$. Let \mathcal{F} be a ξ -family with $\chi(\mathcal{F}) > f(\alpha, \beta)$. Construct a sequence of points $p_0 < \dots < p_{n+1}$ on the baseline with the following properties:

- the points p_0, \dots, p_{n+1} are distinct from all basepoints of 2-curves in \mathcal{F} ,
- p_0 lies to the left of and p_{n+1} lies to the right of all basepoints of 2-curves in \mathcal{F} ,
- $\chi(\mathcal{F}(p_i p_{i+1})) = \beta + 1$ for $0 \leq i < n$ and $\chi(\mathcal{F}(p_n p_{n+1})) \leq \beta + 1$.

This is done greedily, by first choosing p_1 so that $\chi(\mathcal{F}(p_0 p_1)) = \beta + 1$, then choosing p_2 so that $\chi(\mathcal{F}(p_1 p_2)) = \beta + 1$, and so on. For $0 \leq i \leq j \leq n$, let $\mathcal{F}_{i,j} = \{c \in \mathcal{F} : p_i < L(c) < p_{i+1} \text{ and } p_j < R(c) < p_{j+1}\}$. In particular, $\mathcal{F}_{i,i} = \mathcal{F}(p_i p_{i+1})$ for $0 \leq i \leq n$. Since $\mathcal{F} = \bigcup_{0 \leq i \leq j \leq n} \mathcal{F}_{i,j}$ and $\chi(\mathcal{F}) \geq (2\beta + 12\xi + 20)\alpha$, at least one of the following inequalities holds:

$$\chi\left(\bigcup_{i=0}^n \mathcal{F}_{i,i}\right) > (2\beta + 2)\alpha, \quad \chi\left(\bigcup_{i=0}^{n-1} \mathcal{F}_{i,i+1}\right) > (12\xi + 12)\alpha,$$

$$\chi\left(\bigcup_{i=0}^{n-2} \bigcup_{j=i+2}^n \mathcal{F}_{i,j}\right) > 6\alpha.$$

In each case, we claim we can find a subfamily $\mathcal{H} \subseteq \mathcal{F}$ such that any two intersecting 1-curves $x \in R(\mathcal{H})$ and $y \in L(\mathcal{H})$ satisfy $x \in R(\mathcal{F}_{i,j})$ and $y \in L(\mathcal{F}_{r,s})$, where $0 \leq i \leq j \leq n, 0 \leq r \leq s \leq n$, and $|j - r| \geq 2$. Then, we have $\chi(\mathcal{F}(x, y)) \geq \chi(\mathcal{F}(p_{\min(j,r)+1} p_{\max(j,r)})) \geq \beta + 1$, as required.

Suppose $\chi(\bigcup_{i=0}^n \mathcal{F}_{i,i}) > (2\beta + 2)\alpha$. We have $\chi(\mathcal{F}_{i,i}) \leq \beta + 1$ for $0 \leq i \leq n$. Color the 2-curves in each $\mathcal{F}_{i,i}$ properly using the same set of $\beta + 1$ colors on $\mathcal{F}_{i,i}$ and $\mathcal{F}_{r,r}$ whenever $i \equiv r \pmod{2}$, thus partitioning the family $\bigcup_{i=0}^n \mathcal{F}_{i,i}$ into $2\beta + 2$ color classes. Since $\chi(\bigcup_{i=0}^n \mathcal{F}_{i,i}) > (2\beta + 2)\alpha$, at least one such color class $\mathcal{H} \subseteq \bigcup_{i=0}^n \mathcal{F}_{i,i}$ satisfies $\chi(\mathcal{H}) > \alpha$. To conclude, for any two intersecting 1-curves $x \in R(\mathcal{H})$ and $y \in L(\mathcal{H})$, we have $x \in R(\mathcal{F}_{i,i})$ and $y \in L(\mathcal{F}_{r,r})$ for some distinct indices $i, r \in \{0, \dots, n\}$ such that $i \equiv r \pmod{2}$ and thus $|i - r| \geq 2$.

Now, suppose $\chi(\bigcup_{i=0}^{n-1} \mathcal{F}_{i,i+1}) > (12\xi + 12)\alpha$. Lemma 10 yields $\chi(\mathcal{F}_{i,i+1}) \leq 4\xi + 4$ for $0 \leq i \leq n - 1$. Color the 2-curves in every $\mathcal{F}_{i,i+1}$ properly using the same set of $4\xi + 4$ colors on $\mathcal{F}_{i,i+1}$ and $\mathcal{F}_{r,r+1}$ whenever $i \equiv r \pmod{3}$, thus partitioning the family $\bigcup_{i=0}^{n-1} \mathcal{F}_{i,i+1}$ into $12\xi + 12$ color classes. At least one such color class \mathcal{H} satisfies $\chi(\mathcal{H}) > \alpha$. To conclude, for any two intersecting 1-curves $x \in R(\mathcal{H})$ and $y \in L(\mathcal{H})$, we have $x \in R(\mathcal{F}_{i,i+1})$ and $y \in L(\mathcal{F}_{r,r+1})$ for some distinct indices $i, r \in \{0, \dots, n - 1\}$ such that $i \equiv r \pmod{3}$ and thus $|i + 1 - r| \geq 2$.

Finally, suppose $\chi(\bigcup_{i=0}^{n-2} \bigcup_{j=i+2}^n \mathcal{F}_{i,j}) > 6\alpha$. This implies $\chi(\bigcup_{i \in I} \bigcup_{j=i+2}^n \mathcal{F}_{i,j}) > 3\alpha$, where $I = \{i \in \{0, \dots, n - 2\} : i \equiv 0 \pmod{2}\}$ or $I = \{i \in \{0, \dots, n - 2\} : i \equiv 1 \pmod{2}\}$. Consider an auxiliary graph G with vertex set I and edge set $\{ij : i, j \in I, i < j, \text{ and } \mathcal{F}_{i,j-1} \cup \mathcal{F}_{i,j} \neq \emptyset\}$. If there were two edges i_1j_1 and i_2j_2 in G with $i_1 < i_2 < j_1 < j_2$, then the 2-curves witnessing these edges, one from $\mathcal{F}_{i_1,j_1-1} \cup \mathcal{F}_{i_1,j_1}$ and the other from $\mathcal{F}_{i_2,j_2-1} \cup \mathcal{F}_{i_2,j_2}$, would intersect below the baseline, which is impossible. This shows that G is an outerplanar graph, and thus $\chi(G) \leq 3$. Fix a proper 3-coloring of G , and use the color of i on every 2-curve in $\bigcup_{j=i+2}^n \mathcal{F}_{i,j}$ for every $i \in I$, partitioning the family $\bigcup_{i \in I} \bigcup_{j=i+2}^n \mathcal{F}_{i,j}$ into 3 color classes. At least one such color class \mathcal{H} satisfies $\chi(\mathcal{H}) > \alpha$. To conclude, for any two intersecting 1-curves $x \in R(\mathcal{H})$ and $y \in L(\mathcal{H})$, we have $x \in R(\mathcal{F}_{i,j})$ and $y \in L(\mathcal{F}_{r,s})$ for some indices $i, r \in I, j \in \{i + 2, \dots, n\}$, and $s \in \{r + 2, \dots, n\}$ such that $j \notin \{r - 1, r\}$ (otherwise ir would be an edge of G), $j \neq r + 1$ (otherwise two 2-curves, one from $\mathcal{F}_{i,r+1}$ and one from $\mathcal{F}_{r,s}$, would intersect below the baseline), and thus $|j - r| \geq 2$. □

Lemma 2 in [26] asserts that for every family of 1-curves \mathcal{S} with at least one intersecting pair, there are a cap-curve γ and a subfamily $\mathcal{T} \subseteq \mathcal{S}$ with $\chi(\mathcal{T}) \geq \chi(\mathcal{S})/2$ such that every 1-curve in \mathcal{T} is entirely contained in $\text{int } \gamma$ and intersects some 1-curve in \mathcal{S} that intersects γ (equivalently, $\text{ext } \gamma$). The proof follows a standard idea, originally due to Gyárfás [10], to choose \mathcal{T} as one of the sets of 1-curves at a fixed distance from an appropriately chosen 1-curve in the intersection graph of \mathcal{S} . However, this method fails to imply an analogous statement for 2-curves. We will need a more powerful tool—part of the recent series of works on induced subgraphs that must be present in graphs with sufficiently large chromatic number.

Theorem 13 (Chudnovsky et al. [6]) *There is a function $f : \mathbb{N} \rightarrow \mathbb{N}$ with the following property: for every $\alpha \in \mathbb{N}$, every string graph G with $\chi(G) > f(\alpha)$ contains a vertex v such that the subgraph of G induced on the vertices within distance at most 2 from v has chromatic number greater than α .*

The special case of Theorem 13 for triangle-free intersection graphs of curves any two of which intersect in at most one point was proved earlier by McGuinness [21, Thm. 5.3].

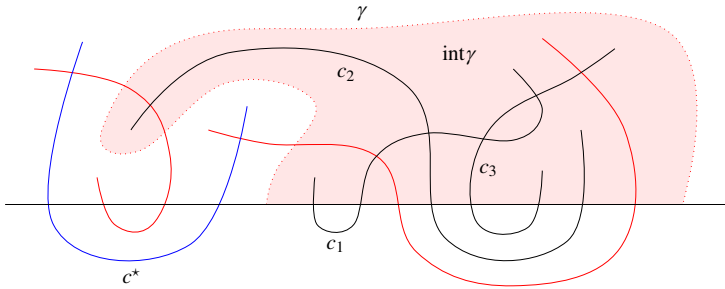


Fig. 2 Illustration for Lemma 14: $\mathcal{G} = \{c_1, c_2, c_3\}$

Lemma 14 For every $\xi \in \mathbb{N}$, there is a function $f: \mathbb{N} \rightarrow \mathbb{N}$ with the following property: for every $\alpha \in \mathbb{N}$ and every ξ -family \mathcal{F} with $\chi(\mathcal{F}) > f(\alpha)$, there are a cap-curve γ and a subfamily $\mathcal{G} \subseteq \mathcal{F}$ with $\chi(\mathcal{G}) > \alpha$ such that every 2-curve $c \in \mathcal{G}$ satisfies $L(c) \cup R(c) \subset \text{int } \gamma$ and intersects some 2-curve in \mathcal{F} that intersects $\text{ext } \gamma$.

Proof Let $f(\alpha) = f_1(3\alpha + 5\xi + 5)$, where f_1 is the function claimed by Theorem 13. Let \mathcal{F} be a ξ -family with $\chi(\mathcal{F}) > f(\alpha)$. It follows that there is a 2-curve $c^* \in \mathcal{F}$ such that the family of 2-curves within distance at most 2 from c^* in the intersection graph of \mathcal{F} has chromatic number greater than $3\alpha + 5\xi + 5$. For $k \in \{1, 2\}$, let \mathcal{F}_k be the 2-curves in \mathcal{F} at distance exactly k from c^* in the intersection graph of \mathcal{F} . We have $\chi(\{c^*\} \cup \mathcal{F}_1 \cup \mathcal{F}_2) > 3\alpha + 5\xi + 5$ (by Theorem 13) and $\chi(\mathcal{F}_1) \leq \xi$ (because \mathcal{F} is a ξ -family), so $\chi(\mathcal{F}_2) > 3\alpha + 4\xi + 4$. We have $\mathcal{F}_2 = \mathcal{G}_1 \cup \mathcal{G}_2 \cup \mathcal{G}_3 \cup \mathcal{G}_4$, where

$$\begin{aligned} \mathcal{G}_1 &= \{c \in \mathcal{F}_2 : L(c) < R(c) < L(c^*) < R(c^*)\}, \\ \mathcal{G}_2 &= \{c \in \mathcal{F}_2 : L(c^*) < L(c) < R(c) < R(c^*)\}, \\ \mathcal{G}_3 &= \{c \in \mathcal{F}_2 : L(c^*) < R(c^*) < L(c) < R(c)\}, \\ \mathcal{G}_4 &= \{c \in \mathcal{F}_2 : L(c) < L(c^*) < R(c^*) < R(c)\}. \end{aligned}$$

Since $\chi(\mathcal{F}_2) > 3\alpha + 4\xi + 4$ and $\chi(\mathcal{G}_k) \leq 4\xi + 4$ (by Lemma 10), we have $\chi(\mathcal{G}_k) > \alpha$ for some $k \in \{1, 2, 3\}$. Since no 2-curve in \mathcal{G}_k intersects c^* and neither basepoint of c^* lies on the segment $I(\mathcal{G}_k)$, the 2-curve c^* and the family \mathcal{G}_k can be separated in H^+ by a cap-curve γ such that $L(c^*) \cup R(c^*) \subset \text{ext } \gamma$ and $L(c) \cup R(c) \subset \text{int } \gamma$ for all $c \in \mathcal{G}_k$ (see Fig. 2). The lemma follows with $\mathcal{G} = \mathcal{G}_k$. \square

Reduction to (ξ, h) -Families

For $\xi \in \mathbb{N}$ and a function $h: \mathbb{N} \rightarrow \mathbb{N}$, a (ξ, h) -family is a ξ -family \mathcal{F} with the following additional property: for every $\alpha \in \mathbb{N}$ and every subfamily $\mathcal{G} \subseteq \mathcal{F}$ with $\chi(\mathcal{G}) > h(\alpha)$, there is a subfamily $\mathcal{H} \subseteq \mathcal{G}$ with $\chi(\mathcal{H}) > \alpha$ such that every 2-curve in \mathcal{F} with a basepoint on $I(\mathcal{H})$ has both basepoints on $I(\mathcal{G})$. We will prove the following lemma.

Lemma 15 For any $\xi, k \in \mathbb{N}$ and any function $h: \mathbb{N} \rightarrow \mathbb{N}$, there is a constant $\zeta \in \mathbb{N}$ such that every (ξ, h) -family \mathcal{F} with $\omega(\mathcal{F}) \leq k$ satisfies $\chi(\mathcal{F}) \leq \zeta$.

The notion of a (ξ, h) -family and Lemma 15 provide a convenient abstraction of what is needed to prove the next lemma and then to derive Lemma 9 from the next lemma.

Lemma 16 *For any $\xi, k \in \mathbb{N}$, there is a function $f: \mathbb{N} \rightarrow \mathbb{N}$ such that for every $\alpha \in \mathbb{N}$, every ξ -family \mathcal{F} with $\omega(\mathcal{F}) \leq k$ and $\chi(\mathcal{F}) > f(\alpha)$ contains a 2-curve c with $\chi(\mathcal{F}(I(c))) > \alpha$.*

Proof of Lemma 16 from Lemma 15 For $\alpha \in \mathbb{N}$, let $h_\alpha: \mathbb{N} \ni \beta \mapsto \beta + 2\alpha + 2 \in \mathbb{N}$, and let $f(\alpha)$ be the constant claimed by Lemma 15 such that every (ξ, h_α) -family \mathcal{F} with $\omega(\mathcal{F}) \leq k$ satisfies $\chi(\mathcal{F}) \leq f(\alpha)$. Let \mathcal{F} be a ξ -family with $\omega(\mathcal{F}) \leq k$ and $\chi(\mathcal{F}(I(c))) \leq \alpha$ for every $c \in \mathcal{F}$. We show that \mathcal{F} is a (ξ, h_α) -family, which then implies $\chi(\mathcal{F}) \leq f(\alpha)$. To this end, consider a subfamily $\mathcal{G} \subseteq \mathcal{F}$ with $\chi(\mathcal{G}) > h_\alpha(\beta)$ for some $\beta \in \mathbb{N}$. Take $\mathcal{G}_L, \mathcal{G}_R \subseteq \mathcal{G}$ greedily so that $L(\mathcal{G}_L) < L(\mathcal{G} \setminus \mathcal{G}_L)$, $\chi(\mathcal{G}_L) = \alpha + 1$, $R(\mathcal{G} \setminus \mathcal{G}_R) < R(\mathcal{G}_R)$, and $\chi(\mathcal{G}_R) = \alpha + 1$. Let $\mathcal{H} = \mathcal{G} \setminus (\mathcal{G}_L \cup \mathcal{G}_R)$. It follows that $\chi(\mathcal{H}) \geq \chi(\mathcal{G}) - \chi(\mathcal{G}_L) - \chi(\mathcal{G}_R) > h_\alpha(\beta) - 2\alpha - 2 = \beta$. If there is a 2-curve $c \in \mathcal{F}$ with one basepoint on $I(\mathcal{H})$ and the other basepoint not on $I(\mathcal{G})$, then $\mathcal{G}_L \subseteq \mathcal{F}(I(c))$ or $\mathcal{G}_R \subseteq \mathcal{F}(I(c))$, so $\chi(\mathcal{F}(I(c))) \geq \alpha + 1$, which is a contradiction. Therefore, every 2-curve in \mathcal{F} with a basepoint on $I(\mathcal{H})$ has both basepoints on $I(\mathcal{G})$. This shows that \mathcal{F} is a (ξ, h_α) -family. \square

Proof of Lemma 9 from Lemma 15 Let h be the function claimed by Lemma 16 for ξ and k . Let ζ be the constant claimed by Lemma 15 for ξ, k , and h . Let \mathcal{F} be a ξ -family with $\omega(\mathcal{F}) \leq k$. We show that \mathcal{F} is a (ξ, h) -family, which then implies $\chi(\mathcal{F}) \leq \zeta$. To this end, consider a subfamily $\mathcal{G} \subseteq \mathcal{F}$ with $\chi(\mathcal{G}) > h(\alpha)$ for some $\alpha \in \mathbb{N}$. Lemma 16 yields a 2-curve $c \in \mathcal{G}$ such that $\chi(\mathcal{G}(I(c))) > \alpha$. Every 2-curve in \mathcal{F} with a basepoint on $I(c)$ has both basepoints on $I(c)$, otherwise it would intersect c below the baseline. Therefore, the condition on \mathcal{F} being a (ξ, h) -family is satisfied with $\mathcal{H} = \mathcal{G}(I(c))$. \square

Dealing with (ξ, h) -Families

The rest of this section is devoted to the proof of Lemma 15. Its structure and principal ideas are based on those of the proof of Theorem 3 presented in [26]. For each forthcoming lemma, we provide a reference to its counterpart in [26].

A *skeleton* is a pair (γ, \mathcal{U}) such that γ is a cap-curve and \mathcal{U} is a family of pairwise disjoint 1-curves each of which has one endpoint (other than the basepoint) on γ and all the remaining part in $\text{int } \gamma$ (see Fig. 3). For a family of 1-curves \mathcal{S} , a skeleton (γ, \mathcal{U}) is an \mathcal{S} -*skeleton* if every 1-curve in \mathcal{U} is a subcurve of some 1-curve in \mathcal{S} . A family of 2-curves \mathcal{G} is *supported* by a skeleton (γ, \mathcal{U}) if every 2-curve $c \in \mathcal{G}$ satisfies $L(c) \cup R(c) \subset \text{int } \gamma$ and intersects some 1-curve in \mathcal{U} . A family of 2-curves \mathcal{H} is *supported from outside* by a family of 1-curves \mathcal{S} if every 2-curve in \mathcal{H} intersects some 1-curve in \mathcal{S} and every 1-curve $s \in \mathcal{S}$ satisfies $s < \mathcal{H}$ or $\mathcal{H} < s$.

Lemma 17 (cf. [26, Lem. 5]) *For every $\xi \in \mathbb{N}$ and every function $h: \mathbb{N} \rightarrow \mathbb{N}$, there is a function $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ such that for any $\alpha, \beta \in \mathbb{N}$, every (ξ, h) -family \mathcal{F} with $\chi(\mathcal{F}) > f(\alpha, \beta)$ contains at least one of the following configurations:*

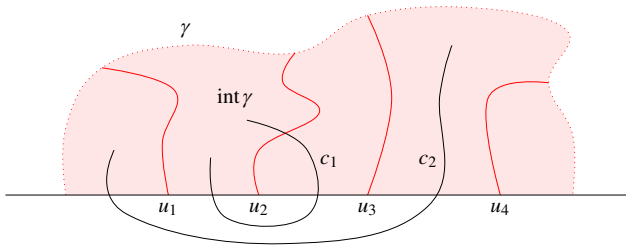


Fig. 3 A skeleton $(\gamma, \{u_1, u_2, u_3, u_4\})$, which supports c_1 but not c_2

- a subfamily $\mathcal{G} \subseteq \mathcal{F}$ with $\chi(\mathcal{G}) > \alpha$ supported by an $L(\mathcal{F})$ -skeleton or an $R(\mathcal{F})$ -skeleton,
- a subfamily $\mathcal{H} \subseteq \mathcal{F}$ with $\chi(\mathcal{H}) > \beta$ supported from outside by a family of 1-curves \mathcal{S} such that $\mathcal{S} \subseteq L(\mathcal{F})$ or $\mathcal{S} \subseteq R(\mathcal{F})$.

Proof Let $f(\alpha, \beta) = f_1(2\alpha + h(2\beta) + 4)$, where f_1 is the function claimed by Lemma 14. Let \mathcal{F} be a (ξ, h) -family with $\chi(\mathcal{F}) > f(\alpha, \beta)$. Apply Lemma 14 to obtain a cap-curve γ and a subfamily $\mathcal{G} \subseteq \mathcal{F}$ with $\chi(\mathcal{G}) > 2\alpha + h(2\beta) + 4$ such that every 2-curve $c \in \mathcal{G}$ satisfies $L(c) \cup R(c) \subset \text{int } \gamma$ and intersects some 2-curve in \mathcal{F}_{ext} . Here and further on, \mathcal{F}_{ext} denotes the family of 2-curves in \mathcal{F} that intersect $\text{ext } \gamma$. Let \mathcal{U}_L be the 1-curves that are subcurves of 1-curves in $L(\mathcal{F})$, have one endpoint (other than the basepoint) on γ , and have all the remaining part in $\text{int } \gamma$. Let \mathcal{U}_R be the analogous subcurves of 1-curves in $R(\mathcal{F})$. Thus (γ, \mathcal{U}_L) is an $L(\mathcal{F})$ -skeleton, and (γ, \mathcal{U}_R) is an $R(\mathcal{F})$ -skeleton. Let \mathcal{G}_L be the 2-curves in \mathcal{G} that intersect some 1-curve in \mathcal{U}_L , and let \mathcal{G}_R be those that intersect some 1-curve in \mathcal{U}_R . If $\chi(\mathcal{G}_L) > \alpha$ or $\chi(\mathcal{G}_R) > \alpha$, then the first conclusion of the lemma holds. Thus assume $\chi(\mathcal{G}_L) \leq \alpha$ and $\chi(\mathcal{G}_R) \leq \alpha$. Let $\mathcal{G}' = \mathcal{G} \setminus (\mathcal{G}_L \cup \mathcal{G}_R)$. It follows that $\chi(\mathcal{G}') \geq \chi(\mathcal{G}) - 2\alpha > h(2\beta) + 4$.

The rest of the argument is illustrated in Fig. 4. By Lemma 8, the 2-curves $c \in \mathcal{G}'$ such that $L(c)$ and $R(c)$ lie in different components of $L(\mathcal{G}') \cup R(\mathcal{G}')$ have chromatic number at most 4. Therefore, there is a component V of $L(\mathcal{G}') \cup R(\mathcal{G}')$ such that $\chi(\mathcal{G}'_V) \geq \chi(\mathcal{G}') - 4 > h(2\beta)$, where $\mathcal{G}'_V = \{c \in \mathcal{G}' : L(c) \cup R(c) \subseteq V\}$. There is a cap-curve $v \subseteq V$ connecting the two endpoints of the segment $I(\mathcal{G}'_V)$. Suppose there is a 2-curve $c \in \mathcal{F}_{\text{ext}}$ with both basepoints on $I(\mathcal{G}'_V)$. If $L(c)$ intersects $\text{ext } \gamma$, then the part of $L(c)$ from the basepoint to the first intersection point with γ , which is a 1-curve in \mathcal{U}_L , intersects v (as $v \subseteq V \subset \text{int } \gamma$) and thus intersects a 2-curve in \mathcal{G}' ; this implies $\mathcal{G}' \cap \mathcal{G}_L \neq \emptyset$, which is a contradiction. An analogous contradiction is reached if $R(c)$ intersects $\text{ext } \gamma$. This shows that no 2-curve in \mathcal{F}_{ext} has both basepoints on $I(\mathcal{G}'_V)$.

Since \mathcal{F} is a (ξ, h) -family and $\chi(\mathcal{G}'_V) > h(2\beta)$, there is a subfamily $\mathcal{H}' \subseteq \mathcal{G}'_V$ such that $\chi(\mathcal{H}') > 2\beta$ and every 2-curve in \mathcal{F} with a basepoint on $I(\mathcal{H}')$ has the other basepoint on $I(\mathcal{G}'_V)$. This and the above imply that no 2-curve in \mathcal{F}_{ext} has a basepoint on $I(\mathcal{H}')$. Since every 2-curve in \mathcal{H}' intersects some 2-curve in \mathcal{F}_{ext} , the family \mathcal{H}' is supported from outside by the family of 1-curves $L(\mathcal{F}_{\text{ext}}) \cup R(\mathcal{F}_{\text{ext}})$. We have $\mathcal{H}' = \mathcal{H}'_L \cup \mathcal{H}'_R$, where \mathcal{H}'_L are the 2-curves in \mathcal{H}' that intersect some 1-curve in $L(\mathcal{F}_{\text{ext}})$ and \mathcal{H}'_R are those that intersect some 1-curve in $R(\mathcal{F}_{\text{ext}})$. Since $\chi(\mathcal{H}') > 2\beta$,

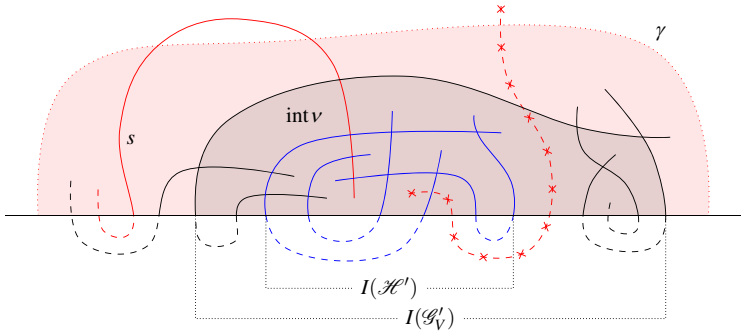


Fig. 4 Illustration for the proof of Lemma 17: the component V is drawn by black and blue solid lines; the family \mathcal{H}' (drawn blue) is supported from outside by the single red 1-curve s

we conclude that $\chi(\mathcal{H}_L) > \beta$ or $\chi(\mathcal{H}_R) > \beta$ and thus the second conclusion of the lemma holds with $(\mathcal{H}, \mathcal{S}) = (\mathcal{H}_L, L(\mathcal{F}_{\text{ext}}))$ or $(\mathcal{H}, \mathcal{S}) = (\mathcal{H}_R, R(\mathcal{F}_{\text{ext}}))$. \square

Lemma 18 (cf. [26, Lem. 8]) *For every $\xi \in \mathbb{N}$ and every function $h: \mathbb{N} \rightarrow \mathbb{N}$, there is a function $f: \mathbb{N} \rightarrow \mathbb{N}$ such that for every $\alpha \in \mathbb{N}$, every (ξ, h) -family \mathcal{F} with $\chi(\mathcal{F}) > f(\alpha)$ contains a subfamily $\mathcal{G} \subseteq \mathcal{F}$ with $\chi(\mathcal{G}) > \alpha$ supported by an $L(\mathcal{F})$ -skeleton or an $R(\mathcal{F})$ -skeleton.*

Proof Let $f(\alpha) = f_1(\alpha, f_1(\alpha, f_1(\alpha, 4\xi)))$, where f_1 is the function claimed by Lemma 17. Let \mathcal{F} be a (ξ, h) -family with $\chi(\mathcal{F}) > f(\alpha)$. Suppose for the sake of contradiction that every subfamily of \mathcal{F} supported by an $L(\mathcal{F})$ -skeleton or an $R(\mathcal{F})$ -skeleton has chromatic number at most α . Let $\mathcal{F}_0 = \mathcal{F}$. Apply Lemma 17 (and the second conclusion thereof) three times to find families $\mathcal{F}_1, \mathcal{F}_2, \mathcal{F}_3, \mathcal{S}_1, \mathcal{S}_2$, and \mathcal{S}_3 with the following properties:

- $\mathcal{F} = \mathcal{F}_0 \supseteq \mathcal{F}_1 \supseteq \mathcal{F}_2 \supseteq \mathcal{F}_3$,
- for $1 \leq i \leq 3$, we have $\mathcal{S}_i \subseteq L(\mathcal{F}_{i-1})$ or $\mathcal{S}_i \subseteq R(\mathcal{F}_{i-1})$, and \mathcal{F}_i is supported from outside by \mathcal{S}_i .
- $\chi(\mathcal{F}_1) > f_1(\alpha, f_1(\alpha, 4\xi))$, $\chi(\mathcal{F}_2) > f_1(\alpha, 4\xi)$ and $\chi(\mathcal{F}_3) > 4\xi$.

There are indices i and j with $1 \leq i < j \leq 3$ such that \mathcal{S}_i and \mathcal{S}_j are of the same “type”: either $\mathcal{S}_i \subseteq L(\mathcal{F}_{i-1})$ and $\mathcal{S}_j \subseteq L(\mathcal{F}_{j-1})$ or $\mathcal{S}_i \subseteq R(\mathcal{F}_{i-1})$ and $\mathcal{S}_j \subseteq R(\mathcal{F}_{j-1})$. Assume for the rest of the proof that $\mathcal{S}_i \subseteq R(\mathcal{F}_{i-1})$ and $\mathcal{S}_j \subseteq R(\mathcal{F}_{j-1})$; the argument for the other case is analogous. Let $\mathcal{S}_L = \{s \in \mathcal{S}_j : s < \mathcal{F}_j\}$, $\mathcal{S}_R = \{s \in \mathcal{S}_j : \mathcal{F}_j < s\}$, \mathcal{F}_L be the 2-curves in \mathcal{F}_j that intersect some 1-curve in \mathcal{S}_L , and \mathcal{F}_R be those that intersect some 1-curve in \mathcal{S}_R . Thus $\mathcal{F}_L \cup \mathcal{F}_R = \mathcal{F}_j$. This and $\chi(\mathcal{F}_j) \geq \chi(\mathcal{F}_3) > 4\xi$ yield $\chi(\mathcal{F}_L) > 2\xi$ or $\chi(\mathcal{F}_R) > 2\xi$. Assume for the rest of the proof that $\chi(\mathcal{F}_L) > 2\xi$; the argument for the other case is analogous.

The rest of the argument is illustrated in Fig. 5. Let \mathcal{S}_L^{\min} be an inclusion-minimal subfamily of \mathcal{S}_L subject to the condition that $L(c)$ intersects some 1-curve in \mathcal{S}_L^{\min} for every 2-curve $c \in \mathcal{F}_L$. Let s^* be the 1-curve in \mathcal{S}_L^{\min} with rightmost basepoint, and let $\mathcal{F}_L^* = \{c \in \mathcal{F}_L : L(c) \text{ intersects } s^*\}$. Since \mathcal{F} is a ξ -family, we have $\chi(\mathcal{F}_L^*) \leq \xi$. By minimality of \mathcal{S}_L^{\min} , the family \mathcal{F}_L^* contains a 2-curve c^* disjoint from every 1-curve

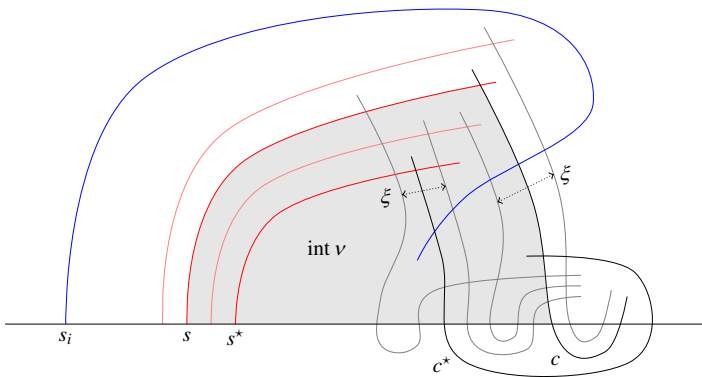


Fig. 5 Illustration for the proof of Lemma 18: members of \mathcal{S}_L^{\min} are drawn red

in \mathcal{S}_L^{\min} other than s^* . Since $c^* \in \mathcal{F}_j \subseteq \mathcal{F}_i$ and \mathcal{F}_i is supported from outside by \mathcal{S}_i , there is a 1-curve $s_i \in \mathcal{S}_i$ that intersects $L(c^*)$. We show that every 2-curve in $\mathcal{F}_L \setminus \mathcal{F}_L^*$ intersects s_i .

Let $c \in \mathcal{F}_L \setminus \mathcal{F}_L^*$, and let s be a 1-curve in \mathcal{S}_L^{\min} that intersects $L(c)$. We have $s \neq s^*$, as $c \notin \mathcal{F}_L^*$. There is a cap-curve $v \subseteq s \cup L(c)$. Since $s < s^* < L(c)$ and s^* intersects neither s nor $L(c)$, we have $s^* \subset \text{int } v$. Since $L(c^*)$ intersects s^* but neither s nor $L(c)$, we also have $L(c^*) \subset \text{int } v$. Since $s \in \mathcal{S}_j \subseteq R(\mathcal{F}_i)$ and $s_i < \mathcal{F}_i$ or $\mathcal{F}_i < s_i$, the basepoint of s_i lies in $\text{ext } v$. Since s_i intersects $L(c^*)$ and $L(c^*) \subset \text{int } v$, the 1-curve s_i intersects v and thus $L(c)$. This shows that every 2-curve in $\mathcal{F}_L \setminus \mathcal{F}_L^*$ intersects s_i . This and the assumption that \mathcal{F} is a ξ -family yield $\chi(\mathcal{F}_L \setminus \mathcal{F}_L^*) \leq \xi$. We conclude that $\chi(\mathcal{F}_L) \leq \chi(\mathcal{F}_L^*) + \chi(\mathcal{F}_L \setminus \mathcal{F}_L^*) \leq 2\xi$, which is a contradiction. \square

A chain of length n is a sequence $((a_1, b_1), \dots, (a_n, b_n))$ of pairs of 2-curves with the following properties:

- for $1 \leq i \leq n$, the 1-curves $R(a_i)$ and $L(b_i)$ intersect,
- for $2 \leq i \leq n$, the basepoints of $R(a_i)$ and $L(b_i)$ lie between the basepoints of $R(a_{i-1})$ and $L(b_{i-1})$, and $L(a_i)$ intersects $R(a_1), \dots, R(a_{i-1})$ or $R(b_i)$ intersects $L(b_1), \dots, L(b_{i-1})$.

Lemma 19 (cf. [26, Lem. 11]) *For every $\xi \in \mathbb{N}$ and every function $h : \mathbb{N} \rightarrow \mathbb{N}$, there is a function $f : \mathbb{N} \rightarrow \mathbb{N}$ such that for every $n \in \mathbb{N}$, every (ξ, h) -family \mathcal{F} with $\chi(\mathcal{F}) > f(n)$ contains a chain of length n .*

Proof of Lemma 15 from Lemma 19 Let $\zeta = f(2k + 1)$, where f is the function claimed by Lemma 19 for ξ and h . Let \mathcal{F} be a (ξ, h) -family with $\chi(\mathcal{F}) > \zeta$. By Lemma 19, \mathcal{F} contains a chain of length $2k + 1$. This chain contains a subchain $((a_1, b_1), \dots, (a_{k+1}, b_{k+1}))$ of pairs of the same “type”—such that $L(a_i)$ intersects $R(a_1), \dots, R(a_{i-1})$ for $2 \leq i \leq k + 1$ or $R(b_i)$ intersects $L(b_1), \dots, L(b_{i-1})$ for $2 \leq i \leq k + 1$. This subchain contains a clique $\{a_1, \dots, a_{k+1}\}$ or $\{b_1, \dots, b_{k+1}\}$, respectively, which is not possible when $\omega(\mathcal{F}) \leq k$. \square

Proof of Lemma 19 We define the function f by induction. We set $f(1) = 1$; if $\chi(\mathcal{F}) > 1$, then \mathcal{F} contains two intersecting 2-curves, which form a chain of length 1. For the induction step, fix $n \geq 1$, and assume that $f(n)$ is defined so that every (ξ, h) -family \mathcal{H} with $\chi(\mathcal{H}) > f(n)$ contains a chain of length n . Let f_1 be the function claimed by Lemma 12 and f_2 be the function claimed by Lemma 18. Let

$$\beta = f_1(f(n), h(2\xi) + 4\xi + 2), \quad f(n + 1) = f_2(f_2(f_2(\beta))).$$

Let \mathcal{F} be a (ξ, h) -family with $\chi(\mathcal{F}) > f(n + 1)$. We claim that \mathcal{F} contains a chain of length $n + 1$.

Let $\mathcal{F}_0 = \mathcal{F}$. Lemma 18 applied three times provides families of 2-curves $\mathcal{F}_1, \mathcal{F}_2, \mathcal{F}_3$ and skeletons $(\gamma_1, \mathcal{U}_1), (\gamma_2, \mathcal{U}_2), (\gamma_3, \mathcal{U}_3)$ with the following properties:

- $\mathcal{F} = \mathcal{F}_0 \supseteq \mathcal{F}_1 \supseteq \mathcal{F}_2 \supseteq \mathcal{F}_3$,
- for $1 \leq i \leq 3$, $(\gamma_i, \mathcal{U}_i)$ is an $L(\mathcal{F}_{i-1})$ -skeleton or an $R(\mathcal{F}_{i-1})$ -skeleton supporting \mathcal{F}_i ,
- $\chi(\mathcal{F}_1) > f_2(f_2(\beta)), \chi(\mathcal{F}_2) > f_2(\beta)$, and $\chi(\mathcal{F}_3) > \beta$.

There are indices i and j with $1 \leq i < j \leq 3$ such that the skeletons $(\gamma_i, \mathcal{U}_i)$ and $(\gamma_j, \mathcal{U}_j)$ are of the same “type”: either an $L(\mathcal{F}_{i-1})$ -skeleton and an $L(\mathcal{F}_{j-1})$ -skeleton or an $R(\mathcal{F}_{i-1})$ -skeleton and an $R(\mathcal{F}_{j-1})$ -skeleton. Assume for the rest of the proof that $(\gamma_i, \mathcal{U}_i)$ is an $L(\mathcal{F}_{i-1})$ -skeleton and $(\gamma_j, \mathcal{U}_j)$ is an $L(\mathcal{F}_{j-1})$ -skeleton; the argument for the other case is analogous.

By Lemma 12, since $\chi(\mathcal{F}_j) \geq \chi(\mathcal{F}_3) > \beta$, there is a subfamily $\mathcal{H} \subseteq \mathcal{F}_j$ such that $\chi(\mathcal{H}) > f(n)$ and $\chi(\mathcal{F}_j(x, y)) > h(2\xi) + 4\xi + 2$ for any two intersecting 1-curves $x, y \in L(\mathcal{H}) \cup R(\mathcal{H})$. Since $\chi(\mathcal{H}) > f(n)$, there is a chain $((a_1, b_1), \dots, (a_n, b_n))$ of length n in \mathcal{H} . Let x and y be the 1-curves $R(a_n)$ and $L(b_n)$ ordered so that $x < y$. Since they intersect, we have $\chi(\mathcal{F}_j(x, y)) > h(2\xi) + 4\xi + 2$.

The rest of the argument is illustrated in Fig. 6. Since $\mathcal{F}_j \subseteq \mathcal{F}_i$ and \mathcal{F}_i is supported by the skeleton $(\gamma_i, \mathcal{U}_i)$, every 2-curve in $\mathcal{F}_j(x, y)$ intersects some 1-curve in \mathcal{U}_i . Let \mathcal{G} be the family of curves in $\mathcal{F}_j(x, y)$ that intersect at least one 1-curve in $\mathcal{U}_i(x, y)$. If a 2-curve $c \in \mathcal{F}_j(x, y)$ intersects no 1-curve in $\mathcal{U}_i(x, y)$, then c intersects the 1-curve in \mathcal{U}_i with rightmost basepoint to the left of the basepoint of x (if such a 1-curve exists) or the 1-curve in \mathcal{U}_i with leftmost basepoint to the right of the basepoint of y (if such a 1-curve exists). This and the fact that \mathcal{F} is a ξ -family imply $\chi(\mathcal{F}_j(x, y) \setminus \mathcal{G}) \leq 2\xi$ and thus $\chi(\mathcal{G}) \geq \chi(\mathcal{F}_j(x, y)) - 2\xi > h(2\xi) + 2\xi + 2$.

Let u_L be the 1-curve in $\mathcal{U}_i(x, y)$ with leftmost basepoint and u_R be the 1-curve in $\mathcal{U}_i(x, y)$ with rightmost basepoint. Every 1-curve in $\mathcal{U}_i(x, y)$ lies in the closed region K bounded by u_L, u_R , the segment of the baseline between the basepoints of u_L and u_R , and the part of γ_i between its intersection points with u_L and u_R . Since \mathcal{F} is a ξ -family, the 2-curves in \mathcal{G} intersecting u_L or u_R have chromatic number at most 2ξ . Every other 2-curve $c \in \mathcal{G}$ satisfies $L(c) \subset K$ or $R(c) \subset K$. Those for which $L(c) \subset K$ but $R(c) \not\subset K$ satisfy $R(c) \cap K = \emptyset$ and therefore are disjoint from each other. Similarly, those for which $R(c) \subset K$ but $L(c) \not\subset K$ are disjoint from each other. Let $\mathcal{G}' = \{c \in \mathcal{G} : L(c) \subset K \text{ and } R(c) \subset K\}$. It follows that $\chi(\mathcal{G} \setminus \mathcal{G}') \leq 2\xi + 2$ and thus $\chi(\mathcal{G}') \geq \chi(\mathcal{G}) - 2\xi - 2 > h(2\xi)$.

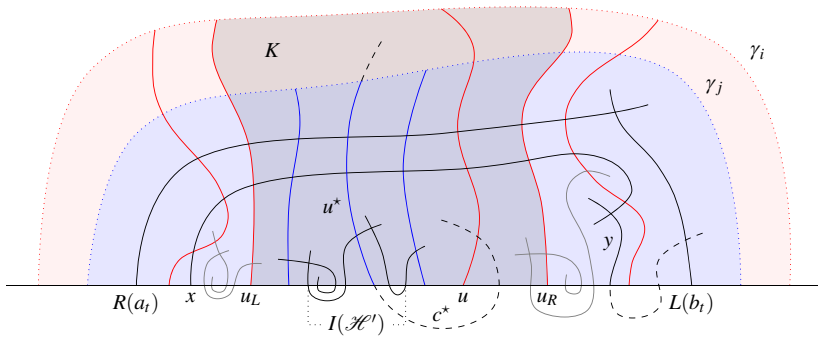


Fig. 6 Illustration for the proof of Lemma 19

Since \mathcal{F} is a (ξ, h) -family and $u_L < \mathcal{G}' < u_R$, there is a subfamily $\mathcal{H}' \subseteq \mathcal{G}'$ with $\chi(\mathcal{H}') > 2\xi$ such that every 2-curve $c \in \mathcal{F}$ with a basepoint on $I(\mathcal{H}')$ satisfies $u_L < c < u_R$. Since $\mathcal{H}' \subseteq \mathcal{F}_j$ and \mathcal{F}_j is supported by $(\gamma_j, \mathcal{U}_j)$, every 2-curve in \mathcal{H}' intersects some 1-curve in \mathcal{U}_j . If a 2-curve $c \in \mathcal{H}'$ intersects no 1-curve in $\mathcal{U}_j(I(\mathcal{H}'))$, then c intersects the 1-curve in \mathcal{U}_j with rightmost basepoint to the left of $I(\mathcal{H}')$ (if such a 1-curve exists) or the 1-curve in \mathcal{U}_j with leftmost basepoint to the right of $I(\mathcal{H}')$ (if such a 1-curve exists). Since \mathcal{F} is a ξ -family, the 2-curves in \mathcal{H}' intersecting at least one of these two 1-curves have chromatic number at most 2ξ . Therefore, since $\chi(\mathcal{H}') > 2\xi$, some 2-curve in \mathcal{H}' intersects a 1-curve in $\mathcal{U}_j(I(\mathcal{H}'))$. In particular, the family $\mathcal{U}_j(I(\mathcal{H}'))$ is non-empty.

Let $u^* \in \mathcal{U}_j(I(\mathcal{H}'))$. The 1-curve u^* is a subcurve of $L(c^*)$ for some 2-curve $c^* \in \mathcal{F}_{j-1}$. The fact that the basepoint of $L(c^*)$ lies on $I(\mathcal{H}')$ and the property of \mathcal{H}' imply $u_L < c^* < u_R$. Since $c^* \in \mathcal{F}_{j-1} \subseteq \mathcal{F}_i$ and \mathcal{F}_i is supported by $(\gamma_i, \mathcal{U}_i)$, the 1-curve $R(c^*)$ intersects a 1-curve $u \in \mathcal{U}_i$, which can be chosen so that $u_L \leq u \leq u_R$, because $c^* \subset \text{int } \gamma_i$ and both basepoints of c^* lie in K . Let $a_{n+1} = c^*$ and b_{n+1} be the 2-curve in \mathcal{F}_{i-1} such that u is a subcurve of $L(b_{n+1})$. Thus $x < \{R(a_{n+1}), L(b_{n+1})\} < y$. For $1 \leq t \leq n$, the facts that the 1-curves $R(a_t)$ and $L(b_t)$ intersect, they are both contained in $\text{int } \gamma_j$ (as $a_t, b_t \in \mathcal{F}_j$), the basepoint of u^* lies between the basepoints of $R(a_t)$ and $L(b_t)$, and u^* intersects γ_j imply that u^* and therefore $L(a_{n+1})$ intersects $R(a_t)$. We conclude that $((a_1, b_1), \dots, (a_{n+1}, b_{n+1}))$ is a chain of length $n + 1$. \square

3 Proof of Theorem 2

Lemma 20 (Fox et al. [9, Lem. 3.2]) *For every $t \in \mathbb{N}$, there is a constant $v_t > 0$ such that every family of curves \mathcal{F} any two of which intersect in at most t points has subfamilies $\mathcal{F}_1, \dots, \mathcal{F}_d \subseteq \mathcal{F}$ (where d is arbitrary) with the following properties:*

- for $1 \leq i \leq d$, there is a curve $c_i \in \mathcal{F}_i$ intersecting all curves in $\mathcal{F}_i \setminus \{c_i\}$,
- for $1 \leq i < j \leq d$, every curve in \mathcal{F}_i is disjoint from every curve in \mathcal{F}_j ,
- $|\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d| \geq v_t |\mathcal{F}| / \log |\mathcal{F}|$.

Proof of Theorem 2 Let \mathcal{F} be a family of curves obtained from the edges of G by shortening them slightly so that they do not intersect at the endpoints but all other intersection points are preserved. It follows that $\omega(\mathcal{F}) \leq k - 1$ (as G is k -quasi-planar) and any two curves in \mathcal{F} intersect in at most t points. Let $v_t, \mathcal{F}_1, \dots, \mathcal{F}_d$, and c_1, \dots, c_d be as claimed by Lemma 20. For $1 \leq i \leq d$, since $\omega(\mathcal{F}_i \setminus \{c_i\}) \leq \omega(\mathcal{F}) - 1 \leq k - 2$, Theorem 1 yields $\chi(\mathcal{F}_i \setminus \{c_i\}) \leq f_t(k - 2)$. Thus $\chi(\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d) \leq f_t(k - 2) + 1$. For every color class \mathcal{C} in a proper coloring of $\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d$ with $f_t(k - 2) + 1$ colors, the vertices of G and the curves in \mathcal{C} form a planar topological graph, and thus $|\mathcal{C}| < 3n$. Thus $|\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d| < 3(f_t(k - 2) + 1)n$. This, the third property in Lemma 20, and the fact that $|\mathcal{F}| < n^2$ yield $|\mathcal{F}| < 3v_t^{-1}(f_t(k - 2) + 1)n \log |\mathcal{F}| < 6v_t^{-1}(f_t(k - 2) + 1)n \log n$. \square

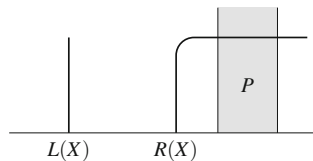
4 Proof of Theorem 6

Proof of Theorem 6 A *probe* is a section of H^+ bounded by two vertical rays starting at the baseline. We use induction to construct, for every positive integer k , an LR -family \mathcal{X}_k of double-curves and a family \mathcal{P}_k of pairwise disjoint probes with the following properties:

- (1) every probe in \mathcal{P}_k is disjoint from $L(X)$ for every double-curve $X \in \mathcal{X}_k$,
- (2) for every probe $P \in \mathcal{P}_k$, the double-curves in \mathcal{X}_k intersecting P are pairwise disjoint,
- (3) \mathcal{X}_k is triangle-free, that is, $\omega(\mathcal{X}_k) \leq 2$,
- (4) for every proper coloring of \mathcal{X}_k , there is a probe $P \in \mathcal{P}_k$ such that at least k distinct colors are used on the double-curves in \mathcal{X}_k intersecting P .

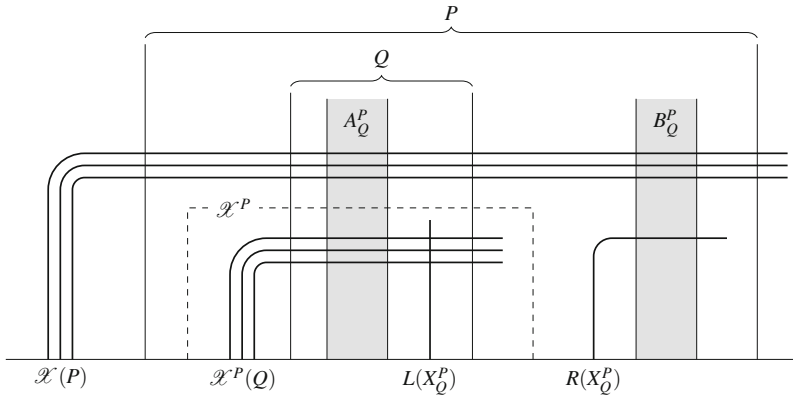
This is enough for the proof of theorem, because the last property implies $\chi(\mathcal{X}_k) \geq k$. For a pair $(\mathcal{X}_k, \mathcal{P}_k)$ satisfying the conditions above and a probe $P \in \mathcal{P}_k$, let $\mathcal{X}_k(P)$ denote the set of double-curves in \mathcal{X}_k intersecting P .

For the base case $k = 1$, we let $\mathcal{X}_1 = \{X\}$ and $\mathcal{P}_1 = \{P\}$, where X and P look as follows:



It is clear that conditions (1)–(4) are satisfied.

For the induction step, we assume $k \geq 1$ and construct the pair $(\mathcal{X}_{k+1}, \mathcal{P}_{k+1})$ from $(\mathcal{X}_k, \mathcal{P}_k)$. Let $(\mathcal{X}, \mathcal{P})$ be a copy of $(\mathcal{X}_k, \mathcal{P}_k)$. For every probe $P \in \mathcal{P}$, put another copy $(\mathcal{X}^P, \mathcal{P}^P)$ of $(\mathcal{X}_k, \mathcal{P}_k)$ inside P below the intersections of P with the double-curves in $\mathcal{X}(P)$. Then, for every probe $P \in \mathcal{P}$ and every probe $Q \in \mathcal{P}^P$, let a double-curve X_Q^P and probes A_Q^P and B_Q^P look as follows:



In particular, X_Q^P intersects the double-curves in $\mathcal{X}^P(Q)$, A_Q^P intersects the double-curves in $\mathcal{X}(P) \cup \mathcal{X}^P(Q)$, and B_Q^P intersects the double-curves in $\mathcal{X}(P) \cup \{X_Q^P\}$. Let

$$\mathcal{X}_{k+1} = \mathcal{X} \cup \bigcup_{P \in \mathcal{P}} \mathcal{X}^P \cup \bigcup_{P \in \mathcal{P}} \{X_Q^P : Q \in \mathcal{P}^P\}, \quad \mathcal{P}_{k+1} = \bigcup_{P \in \mathcal{P}} \{A_Q^P, B_Q^P : Q \in \mathcal{P}^P\}.$$

Conditions (1) and (2) clearly hold for $(\mathcal{X}_{k+1}, \mathcal{P}_{k+1})$, and (2) for $(\mathcal{X}_k, \mathcal{P}_k)$ implies (3) for $(\mathcal{X}_{k+1}, \mathcal{P}_{k+1})$. To see that (4) holds for $(\mathcal{X}_{k+1}, \mathcal{P}_{k+1})$ and $k + 1$, consider a proper coloring ϕ of \mathcal{X}_{k+1} . Let $\phi(X)$ denote the color of a double-curve $X \in \mathcal{X}_{k+1}$ and $\phi(\mathcal{Y})$ denote the set of colors used on a subset $\mathcal{Y} \subseteq \mathcal{X}_{k+1}$. By (4) applied to $(\mathcal{X}, \mathcal{P})$, there is a probe $P \in \mathcal{P}$ such that $|\phi(\mathcal{X}(P))| \geq k$. By (4) applied to $(\mathcal{X}^P, \mathcal{P}^P)$, there is a probe $Q \in \mathcal{P}^P$ such that $|\phi(\mathcal{X}^P(Q))| \geq k$. Since X_Q^P intersects the double-curves in $\mathcal{X}^P(Q)$, we have $\phi(X_Q^P) \notin \phi(\mathcal{X}^P(Q))$. If $\phi(\mathcal{X}(P)) \neq \phi(\mathcal{X}^P(Q))$, then $\mathcal{X}_{k+1}(A_Q^P) = \mathcal{X}(P) \cup \mathcal{X}^P(Q)$ yields $|\phi(\mathcal{X}_{k+1}(A_Q^P))| = |\phi(\mathcal{X}(P)) \cup \phi(\mathcal{X}^P(Q))| \geq k + 1$. If $\phi(\mathcal{X}(P)) = \phi(\mathcal{X}^P(Q))$, then $\mathcal{X}_{k+1}(B_Q^P) = \mathcal{X}(P) \cup \{X_Q^P\}$ and $\phi(X_Q^P) \notin \phi(\mathcal{X}(P))$ yield $|\phi(\mathcal{X}_{k+1}(B_Q^P))| = |\phi(\mathcal{X}(P)) + 1| \geq k + 1$. This shows that (4) holds for $(\mathcal{X}_{k+1}, \mathcal{P}_{k+1})$ and $k + 1$. \square

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