

On Calculation of Monomial Automorphisms of Linear Cyclic Codes

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Abstract—A description of the monomial automorphisms group of an arbitrary linear cyclic code in term of polynomials is presented. This allows us to reduce a task of code's monomial automorphisms calculation to a task of solving some system of equations (in general, nonlinear) over a finite field. The results are illustrated with examples of calculating the full monomial automorphisms groups for two codes.

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1. INTRODUCTION

This paper is devoted to calculation of the full monomial automorphisms groups of linear cyclic codes. The knowledge of code's automorphisms allows us to refine a code's structure and can be used when designing error-correcting decoding algorithms. We start with needed definitions and notions. Let \mathbb{F}_q be a finite field of q elements, \mathbb{F}_q^* be its multiplicative group, \mathbb{F}_q^n be a set of vectors of length n over \mathbb{F}_q . Any non-empty set $C \subseteq \mathbb{F}_q^n$ is called a code over \mathbb{F}_q of length n ; if C is a linear subspace over \mathbb{F}_q , then the code C is called a *linear code*. For short, we call a linear code V of length ν over \mathbb{F}_q , which is invariant under the cyclic shift of vectors, as $C(\nu, q)$ -code.

A component-wise product of vectors $\mathbf{v} = (v_1, v_2, \dots, v_n)$ and $\mathbf{w} = (w_1, w_2, \dots, w_n)$ is determined as $\mathbf{v} \circ \mathbf{w} = (v_1 w_1, v_2 w_2, \dots, v_n w_n)$. For vectors $\mathbf{v} = (v_1, v_2, \dots, v_n) \in \mathbb{F}_q^n$, $\mathbf{c} = (c_1, c_2, \dots, c_n) \in$

$(\mathbb{F}_q^*)^n$, and a permutation $\pi = \begin{pmatrix} 1 & 2 & \dots & n \\ \pi(1) & \pi(2) & \dots & \pi(n) \end{pmatrix}$ on the set of indexes $I_n = \{1, 2, \dots, n\}$

(π is an element of symmetric group S_n on n points) we determine the following vectors: $\mathbf{v}^\pi = (v_{\pi(1)}, v_{\pi(2)}, \dots, v_{\pi(n)})$ is a permutation of components of vector \mathbf{v} by the action π ; $\mathbf{v}^{(\pi, \mathbf{c})} = (c_1 v_{\pi(1)}, c_2 v_{\pi(2)}, \dots, c_n v_{\pi(n)})$ is a result of transformation of the vector \mathbf{v} under the action (π, \mathbf{c}) .

Notice that $(\mathbf{v}^\pi)^\rho = \mathbf{v}^{\pi\rho}$ and $(\mathbf{v}^{(\pi, \mathbf{a})})^{(\rho, \mathbf{b})} = (\mathbf{v}^\pi \circ \mathbf{a})^{(\rho, \mathbf{b})} = \mathbf{v}^{\pi\rho} \circ \mathbf{a}^\rho \circ \mathbf{b} = \mathbf{v}^{\pi\rho} \circ \mathbf{a}^{(\rho, \mathbf{b})}$, where $\pi\rho$ is a product of permutations π and ρ with multiplication on the left, i.e. $\pi\rho(x) = \pi(\rho(x)) \forall x \in I_n$.

For a code $V \subseteq \mathbb{F}_q^n$, $\pi \in S_n$, $\mathbf{c} \in (\mathbb{F}_q^*)^n$ we assume $V^\pi = \{\mathbf{v}^\pi | \mathbf{v} \in V\}$, $V^{(\pi, \mathbf{c})} = V^\pi \circ \mathbf{c} = \{\mathbf{v}^{(\pi, \mathbf{c})} | \mathbf{v} \in V\}$. A transformation (π, \mathbf{c}) , preserving the code V , i.e. $V^{(\pi, \mathbf{c})} = V$, is called a *monomial automorphism* of the code V . A set $\text{MAut}(V)$ of all such automorphisms forms a group with multiplication (on the left): $(\rho, \mathbf{b}) \times (\pi, \mathbf{a}) = (\pi\rho, \mathbf{a}^{(\rho, \mathbf{b})})$.

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