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# Formula Layout 

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#### Abstract

Both the quality of the results of TEX's formula layout algorithm and the complexity of its description in the $\mathrm{T}_{\mathrm{E}} \mathrm{Xb}$ book [1] are hard to beat. The algorithm is (verbally) described as an imperative program with very complex control flow and complicated manipulations of the data structures representing formulae. In a forthcoming textbook [3], we describe TEX's formula layout algorithm as a functional program transforming mlist-terms into box-terms. This transformation is given in this paper.


## 1 Introduction

The quality of the results of $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ 's formula layout algorithm are convincing. However, any attempt to understand the reasons for that leads to deep frustration when Knuth's description of the algorithm from the $T_{E} X b o o k$ [1] is used. In an attempt to understand this problem, one has to cleanly separate the reasons for the lack of understandability.

1. The problem may have a nature that does not allow for a solution which is easily described in some readable way. Not much can be done about that.
2. The algorithm used to solve the problem may not be the simplest possible, but may be tuned for efficiency or optimality of the result. Here, a clean separation between principles of a space of solutions and the optimizations applied would help the interested reader.
3. The context of the chosen algorithm may enforce a bad design. Here, a new describer may take the freedom to abstract from this context.
4. The description may not be the best possible for the given algorithm. This is a particularly favorable situation for an attempt to explain an interesting subject better.

Knuth's description of formula layout is an imperative program with very complex control flow and complicated manipulations of the data structures representing formulae. The 'programming language' is English prose with some formal fragments. In this paper, we present a new description using the pure functional language Miranda ${ }^{1}$ [2]. The use of a functional language gives a completely new flavor to the description. Of course, the mere fact that we use a concrete programming language instead of English phrases adds rigor and exactness to the exposition.

Our criticism and attempt to improve the presentation of the formula layout algorithm of $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ mainly touches points $2-4$ above. In Section 2, we consider the input of the layout algorithm, i.e., the internal representation of formulae. In 2.1, we discuss Knuth's original data structure. In our

[^0]opinion, it is misconceived. Many difficulties in Knuth's description result from the design of this data structure. In 2.2, we propose a new data structure for formulae with a clean and simple design.

In Section 3 we present some more details which influence formula layout: the styles of formulae and subformulae (which communicate information about their context), the representation of characters, and the layout parameters which control the positions of subformulae. Knuth's account of these things is very concrete. In contrast, we present an abstract interface which hides the details of font table organization, and makes clear how the information is used.

In Section 4, we consider the output of the layout algorithm, box terms. Knuth's description of this data structure and its operations is particularly vague. We try to model Knuth's intentions by a Miranda data type and functions defined in Miranda.

In Section 5, we present a bunch of specialized functions which translate subformulae of various kinds into box terms. In Section 6, we deal with the translation of whole formulae, i.e., the recursive descent to subformulae and the selection of the appropriate specialized subformula functions.

We give an honest estimation of the improvements in the conclusion (Section 7). Of course, our functional solution is not simpler than the problem admits. Formula layout is an inherently difficult problem; not in terms of computational, but of algorithmic complexity. There are many different kinds of mathematical formulae, whose layout is governed by tradition and aesthetics. Algorithms for formula layout have to distinguish many cases and pay attention to lots of little details.

## 2 Internal Representation of Formulae

### 2.1 The Original $\mathrm{T}_{\mathrm{E}} \mathrm{X}$-Representation

$\mathrm{T}_{\mathrm{E}} \mathrm{X}$ reads a formula specification from the input and converts it into an internal representation, a math list. A math list is a sequence of math items.

According to the description in the $\mathrm{T}_{\mathrm{E}} \mathrm{Xbook}$ [1, page 157], a math item is an atom, a horizontal space, a style command (e.g., \textstyle), a generalized fraction, or some other material which we do not consider here for simplification.

Atoms have (at least) three parts: a nucleus, a superscript, and a subscript. Each of these fields may be empty, a math symbol, or a math list. There are thirteen kinds of atoms, some of which with additional parts. Eight atom kinds mainly regulate the spacing between two adjacent atoms: a relation atom such as ' $=$ ' is surrounded by some amount of space, a binary atom such as ' + ' by less space, and an ordinary atom such as ' $x$ ' by no extra space at all. The remaining five kinds of atoms have a more serious semantics. An overline atom for instance is an overlined subformula.

The formula $\left(x_{i}+y\right)^{\overline{n+1}}$ for instance may be specified as $\$\left(x_{-} i+y\right) \wedge\{\backslash o v e r l i n e\{n+1\}\} \$$. In internal form, it is represented by a math list consisting of five atoms: an 'Open'-atom with nucleus '(' (and empty superscript and subscript); an 'Ord'-atom with nucleus ' $x$ ', empty superscript, and subscript ' $i$ '; a 'Bin'-atom with nucleus ' + '; an 'Ord'-atom with nucleus ' $y$ '; and finally a 'Close'-atom with nucleus ')', whose superscript is a math list consisting of a single 'Over'-atom, whose nucleus is a math list of three atoms corresponding to $n+1$.

This internal representation deserves some criticism. The superscript and subscript fields are empty in most cases; there should really be superscript and subscript constructors. The thirteen kinds of atoms combine two completely different aspects: a classification needed to control spacing, and the adjunction of meaningful constructors. These two aspects should not be mixed into a single concept. Interestingly, $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ 's layout algorithm internally tries hard to distinguish these aspects, as we explain by two examples.

Overline atoms are handled during a first pass through the formula. The overline rule is added to the corresponding subformula, and afterwards, it is transformed into an 'Ord' atom since the spacing
of overline atoms and 'Ord' atoms is identical. The actual inter-atom spaces are added in a second pass through the formula.

Fractions are math items, but not atoms. Their layout is computed during the first pass of the algorithm, and afterwards, they are transformed into 'Inner' atoms. The kind 'Inner' controls the spacing around fractions in the second pass of the algorithm.

Thus, we see that the mixture of different concepts into the same notion leads to the need to destructively transform the data structure of formulae which makes $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ 's layout algorithm hard to understand.

### 2.2 An Alternative Representation Defined in Miranda

To avoid the problems mentioned above, we completely redesigned the internal representation of formulae. The following definition is given in Miranda.

As in the original representation, formulae are math lists (mlist) consisting of math items (mitem).
mlist $==$ [mitem]
Mitems are defined as the elements of a constructor type. We do not distinguish between atoms and non-atoms, and restrict ourselves to semantically meaningful constructors.

```
mitem ::=
    Sym class mathchar | || a single symbol with its class
    MathSpace num | || space (in relative math units)
    Over mlist | || overlined subformula
    Under mlist | || underlined subformula
    Frac mlist mlist | || fraction with numerator and denominator
    Sup mitem mlist | || formula with superscript
    Sub mitem mlist | || formula with subscript
    SupSub mitem mlist mlist | || with superscript and subscript
    Class_cmd class mlist | || from \mathord{ }, \mathop{ } etc
    Style_cmd style | || from \displaystyle etc
    Group mlist || a nested group, indicated by {...}
```

For reasons of simplicity, we omitted some of $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ 's possibilities. To cover the full power of $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ formulae, additional constructors would be needed for left and right big delimiters, for accented characters, for roots $(\sqrt{3}+\sqrt[3]{2})$, for vertically centered subformulae, etc. They don't offer principally new problems, although the treatment of accented characters and roots in [1] is particularly hard to grasp. The constructor Frac represents a special case of $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ 's generalized fractions; for a full treatment, more argument fields would be needed.

To complete our description, we have to define the types class, mathchar, and style. The type mathchar is defined in Section 3.2, and style in Section 3.1. The type class enumerates nine constructors:

```
class ::= Ord | Op | Bin | Rel | Open | Close | Punct | Inner | None
```

The first eight classes correspond to those atom kinds which control spacing. The ninth class None is used for mitems which are not atoms in the original $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ representation, and are never transformed into atoms.

Using our representation, the formula ( $x_{-} i+y$ ) ^\{\overline\{n+1\}\} is internally described as the following term:

```
[ Sym Open '(',
    Sub (Sym Ord 'x') [Sym Ord 'i'],
    Sym Bin '+',
    Sym Ord 'y',
```

Sup (Sym Close ')') [Over [Sym Ord 'n', Sym Bin '+', Sym Ord '1']]

The internal representation is created by a parser starting from the external formula description. We have to assume that this parser is a bit more powerful than the one employed in the $\mathrm{T}_{\mathrm{E}} \mathrm{Xbook}$. It has to correctly transform the input string into our data structure obeying the subformula structure. Note that the nucleus of Sup etc. needs grouping if it is not a single symbol. (The nucleus is an mitem instead of an mlist, since otherwise, class computation and spacing would fail.)

When reading a character or mathematical symbol, the parser knows about the pre-assigned class of this symbol, e.g., Rel for ' $=$ ' and Open for ' ('. This class is stored in the internal representation as the first argument of the Sym constructor. As the biggest difference to the original $\mathrm{T}_{\mathrm{E}} \mathrm{X}$-situation, we assume that the parser is able to recognize binary symbols which are used in non-binary contexts, e.g., the plus symbol in $f^{+}$. The class of these symbols should be Ord instead of Bin. In the original $\mathrm{T}_{\mathrm{E}} \mathrm{X}-$ algorithm, atoms of kind 'Bin' change their kind into 'Ord' depending on the kinds of neighboring atoms during the computation of inter-atom spaces. This solution could also be programmed in Miranda, but would make function do_mlist in Section 6.3 overly complex.

In contrast to the original description, classes are not stored with all mitems. The reason is that in almost all cases, the class of an mitem can be derived mechanically from its structure. The only exceptions are symbols which are classified by some external declarations, and the Class_cmd items which come from explicit class assertions in the formula description (by the commands \mathord, $\backslash$ mathop etc.).

The following function computes the class of every subformula:

```
class_ :: mitem -> class
class_ (Sym cl mc) = cl || the class is a symbol property
class_ (MathSpace w) = None || spaces are not atoms, and never will be
class_ (Over ml) = Ord || Over-atoms are changed into Ord-atoms
class_ (Under ml) = Ord || Under-atoms are changed into Ord-atoms
class_ (Frac num den) = Inner || fractions become Inner-atoms
class_ (Sup mi sup) = class_ mi
class_ (Sub mi sub) = class_ mi
class_ (SupSub mi sup sub) = class_ mi
class_ (Class_cmd cl ml) = cl || class is explicitly set
class_ (Style_cmd st) = None || style commands are not atoms
class_ (Group ml) = Ord || this is an Ord-atom in TeX
```

In the comments, we tried to explain the reason for this rule. For instance, Over-items are classified as Ord because they are transformed into 'Ord'-atoms in the course of $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ 's layout algorithm.

## 3 Additional Details

In this section, we present some additional detail information needed for the formula layout: the styles of formulae and subformulae, the representation of characters, and the layout parameters controlling the positions of subformulae.

### 3.1 Formula Styles

The layout of formulae and subformulae in $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ documents depends on a style parameter. There are two kinds of basic styles: formulae may appear on a separate line by their own (display style) or as
part of a line of text (text style). Consider the following displayed formula

$$
\sum_{j=1}^{n} A^{i^{j}}+\frac{y^{2}}{y^{2}+z^{2}}
$$

and its inline counterpart $\sum_{j=1}^{n} A^{i^{j}}+\frac{y^{2}}{y^{2}+z^{2}}$. We observe that in display style, the sum symbol is bigger, and the limits of the summation are placed vertically below and above it (this is called limit position). In text style however, the position of the limits is to the right of the symbol. All superscripts are set in styles with smaller characters and spaces. The same is true for the constituents of the fraction in text style. Notice also how the position of superscripts depends on their context, i.e., on the style of the corresponding subformula. In the denominator, their position is lower than in the numerator.

In the $\mathrm{T}_{\mathrm{E}} \mathrm{Xbook}$ [1], there are eight styles altogether: display style $D$, text style $T$, script style $S$, script-script style $S S$, and four 'cramped' styles $D^{\prime}, T^{\prime}, S^{\prime}$, and $S S^{\prime}$. In cramped styles, which are used for denominators, superscripts are placed in a lower position than in the corresponding uncramped styles. In analyzing the usage of these styles, it turned out that they may be regarded as pairs of a main style and a Boolean value 'cramped'. The two components of the pairs are independently calculated and used, so that it is easy to separate them completely. This is done in our Miranda program. Hence, we have only four styles:

```
style ::= D | T | S | SS
```

Function script computes the style for subscripts and superscripts from the current style, and fract calculates the styles of numerators and denominators.

```
script, fract :: style -> style
script D = S; script T = S; script S = SS; script SS = SS
fract D = T; fract T = S; fract S = SS; fract SS = SS
```


### 3.2 Math Characters and Output Characters

Characters from a formula description do not yet completely determine the characters which appear in the printed document. The formula description $\mathrm{x}^{\wedge} \mathrm{x}$ for instance yields the printed formula $x^{x}$, where the two occurrences of $x$ appear in different sizes. The reason is that the first $x$ is set in text style T, whereas for the second one, script style S is used.

In our description, we model this behavior by using two different types of characters and a styledependent transfer function. For characters in the internal representation, whose appearance is not yet determined, we use type mathchar, whereas type outchar is used for characters in the result of the formula layout. The transfer function is setchar : : style $\rightarrow$ mathchar $\rightarrow$ outchar. Although the two types and setchar could be specified further following the hints in the $\mathrm{T}_{\mathrm{E}} \mathrm{Xb}$ book, we refrain from doing it since a complete definition would be difficult and hardly interesting.

For formula layout, we need some information about the size and form of characters. The height of a character is the distance from its top end to the base line; e.g., ' $a$ ' and ' $g$ ' have the same height, and ' $f$ ' has a bigger one. The depth is the distance from the base line to the bottom end; e.g., ' $a$ ' has depth 0 , whereas ' $g$ ' has non-zero depth. The width is the horizontal size, and the slant gives information how far the character is slanted to the right.

These character informations are given by the four functions char_height, char_depth, char_width, and char_slant, all with type outchar $\rightarrow$ dim, where dim is the type of dimensions, i.e., amounts of length, measured in basic units. We may simply assume dim == num. The four functions are left unspecified here; in practice, their values are read off from the appropriate font tables.

### 3.3 Layout Parameters

The exact layout of a formula depends on some layout parameters. They control the position of superscripts, the distance between numerator and fraction stroke, the thickness of the stroke, etc.

In the $\mathrm{T}_{\mathrm{E}} \mathrm{Xbook}$, the layout parameters are attached to the fonts used to make formulae. Since the choice of the font depends on the style, we incorporate the layout parameters as functions of type style -> dim. The function names are (abbreviations of) the symbolical names given in the table in [1, page 447].
Height of ' $x$ ' in current font:

```
x_height
quad
num1 num2
denom1 denom2
sub_drop sub1 sub2
big_op1 through big_op5
rule_thickness
axis_height
```

for superscripts: sup_drop sup1 sup2 sup3
Width of ' $M$ ' in current font:
Parameters for numerators:
for denominators:
for subscripts:
for limits at large operators:
Default thickness of rules:

Distance from 'axis' to base line: axis_height
The axis is the line where fraction strokes sit on. Consider e.g., $x+\frac{y}{z}$. The base line is at the bottom end of the ' $x$ ' and the ' + '.
Some font parameters are used in special contexts only. This is realised by three auxiliary functions.

```
num_level, den_level :: style -> dim
num_level D = num1 D; num_level st = num2 st
den_level D = denom1 D; den_level st = denom2 st
sup_level :: style -> bool -> dim
sup_level D False = sup1 D || Display style, not cramped
sup_level st True = sup3 st || all cramped styles
sup_level st cr = sup2 st || style T, S, or SS; not cramped
```

In addition to the style-dependent layout parameters, there is a constant scriptspace of type dim.

## 4 The Target Representation: Box Terms

During formula layout, an input term of type mlist is translated into a term of type box. Boxes are rectangles whose edges are parallel to the page edges. Compound boxes are built from smaller boxes, and atomic boxes contain symbols or are filled with black. Each box has a horizontal base line. It has the reference point of the box at its left end. Boxes have heights, h, depths, $d$, and widths, $w$. These dimensions may be negative. This is the case for boxes which are shifted upwards or downwards beyond their base line and for boxes which represent negative distances.


In the $\mathrm{T}_{\mathrm{E}} \mathrm{Xbook}$, boxes and their properties are described verbally. At first glance, the size attributes of a compound box seem to be totally determined by the sizes of its constituents. Later however, it seems as if the size dimensions of a box may be arbitrarily changed. For, the description of the formula layout contains phrases such as "increase the depth of the box by", "add ... to the width of the box", or "construct a box with depth ... and height ...".

Here, we represent boxes as a Miranda data type. The operations on boxes are formalized. The size dimensions of our boxes are determined by their structure. We tried to catch the intended meaning of the size manipulations in [1] by adding space boxes without visible content.

```
box ::= HSpace dim | || horizontal space with width
    VSpace dim dim | || vertical space with height and depth
    Rule dim dim dim | || black box with height, depth, and width
    Chr outchar | || character box
    HBox [box] | || horizontal list of boxes
    Vdn [box] | || vertical list of boxes, downward
    Vup [box] || vertical list of boxes, upward
```

Their are four kinds of atomic boxes and three kinds of compound boxes. An HBox is the horizontal concatenation of a list of boxes, ordered from left to right. The boxes are concatenated such that their base lines become adjacent. The reference point of an HBox is the one of its leftmost constituent. An HBox may be empty.

Both Vdn and Vup boxes represent vertical concatenations of boxes. In both cases, the concatenation is done so that the reference points of the constituent boxes are vertically aligned. In Vdn boxes, the constituents are ordered from top to bottom. The reference point of a Vdn box is the reference point of its topmost component. In contrast, the components of a Vup list are ordered from bottom to top. The reference point of a Vup box is the one of its lowest component. Thus, in both cases, the reference point of the compound box is the one of the head of its list of components. Both Vdn and Vup lists should never be empty.

## The Dimensions of a Box

Height, depth, and width of a box are uniquely defined from its structure. We call the sum of height and depth vsize.

```
height, depth, width, vsize :: box -> dim
vsize box = height box + depth box
height (HSpace w) = 0; height (VSpace h d) = h; height (Rule h d w) = h
height (Chr ch) = char_height ch
height (HBox boxl) = max0 (map height boxl) || as max, but max0 [] = 0
height (Vdn (top : rest)) = height top
height (Vup (bot : rest)) = height bot + sum (map vsize rest)
depth (HSpace w) = 0; depth (VSpace h d) = d; depth (Rule h d w) = d
depth (Chr ch) = char_depth ch
depth (HBox boxl) = max0 (map depth boxl)
depth (Vdn (top : rest)) = depth top + sum (map vsize rest)
depth (Vup (bot : rest)) = depth bot
width (HSpace w) = w; width (VSpace h d) = 0; width (Rule h d w) = w
width (Chr ch) = char_width ch
width (HBox boxl) = sum (map width boxl)
width (Vdn boxl) = max (map width boxl)
width (Vup boxl) = max (map width boxl)
```

For the sake of efficiency, all three dimensions could be stored at HBox, Vdn, and Vup constructors, in order to avoid costly recomputations (memoization).

## Some Operations on Boxes

hconc concatenates two boxes to form an HBox. If one of them is an HBox already, nesting of HBoxes is avoided.

```
hconc :: box -> box -> box
hconc (HBox boxl1) (HBox boxl2) = HBox (boxl1 ++ boxl2) || list concatenation
hconc box1 (HBox boxl2) = HBox (box1 : boxl2)
hconc (HBox boxl1) box2 = HBox (boxl1 ++ [box2])
hconc box1 box2 = HBox [box1 , box2]
```

right moves a box to the right by putting an HSpace box in front of it.
right : : dim -> box -> box
right 0 box $=$ box
right 1 box $=$ (HSpace 1) \$hconc box || \$hconc = hconc as infix operator center centers a given box inside a space of given width. It uses right.
center : : dim -> box $\rightarrow$ box; center w box $=$ right ( $($ w - width box)/2) box center is only called with $w \geq$ width box. It does not matter that there is no HSpace to the right of the box since centered boxes are placed in vertical lists where widths are maximized.
The next operation extends a box to the right ("increases its width").

```
extend :: dim -> box -> box
extend 0 box = box; extend l box = box $hconc (HSpace l)
```

A box is raised by increasing its height and decreasing its depth; the vsize does not change. This is done by vertically adjoining an empty box of vsize 0 , but non-zero height and depth (one of these must be negative).

```
raise :: dim -> box -> box
raise 0 box = box
raise l box = Vup [VSpace (l - d) (d - l), box] where d = depth box
To verify raise, show the two equations
```

```
height (raise l box) = height box + l depth (raise l box) = depth box - l.
```

```
height (raise l box) = height box + l depth (raise l box) = depth box - l.
```

Instead of Vup, Vdn could be used equally well (with a different argument).
Finally, we define an operation vlist which takes three arguments: a box $B$, a list of boxes in upward order which goes above $B$, and a list of boxes in downward order which goes below $B$. The reference point of the whole thing is that of $B$.

```
vlist :: box -> [box] -> [box] -> box
vlist box up_list dn_list = Vdn ( Vup (box : up_list) : dn_list )
vlist could equally well be specified the other way round: ...Vup (Vdn (box:dn_list) : up_list).
```


## 5 Setting of Subformulae

In the sequel, we show how subformulae of the various kinds are translated into box terms. Later, we combine these functions to a function that computes the layout of arbitrary mitems.

### 5.1 Symbols and Spaces

Symbols (mathchars) are transformed into character boxes by choosing the appropriate output character (function setchar of Section 3.2) and putting it into a box ( Chr ) which is vertically centered around the axis in some cases (vcenter). The result is not only the box, but also the slant ('italic correction') of the produced character. This information is needed later.

```
set_sym :: style -> class -> mathchar -> (box, dim)
set_sym st cl mc = (vcenter (Chr ch), char_slant ch), if cl = Op
    = ( Chr ch , char_slant ch), otherwise
    where ch = setchar st mc
vcenter :: style -> box -> box
vcenter st box = raise (axis_height st - (height box - depth box)/2) box
```

When spaces are set, their size has to be transformed from style-dependent mathematical units into an absolute dimension.
set_space : : style -> num -> box; set_space st ml = HSpace (ml * quad st/18)

### 5.2 Setting Overlined and Underlined Subformulae

We assume that the subformula is already translated into a box. The thickness of the line, th, depends on the style. Between the line and the formula, there is a gap of size 3 th, and above the overline / below the underline, there is white space of size $t$. Since the reference point of the whole thing should be that of the subformula, we use Vup for overlines and Vdn for underlines. In both cases, the list of constituent boxes starts with the subformula box, followed by the distance to the line, the line itself, and the white space beyond it.

```
set_over, set_under :: style -> box -> box
set_over st box = Vup [box, VSpace ( 3 * th) 0, Rule th 0 w, VSpace th 0]
    where w = width box; th = rule_thickness st
set_under st box = Vdn [box, VSpace ( 3 * th) 0, Rule th 0 w, VSpace th 0]
    where w = width box; th = rule_thickness st
```


### 5.3 Setting of Fractions

Numerator and denominator are already given as boxes. The desired vertical position of the numerator is given by num_level relative to the base line. However, the fraction stroke will be positioned at the axis, an invisible line somewhere above the base line. Thus, we compute the position num_pos of the reference point of the numerator relative to the axis. From this, the actual distance num_dist between the bottom edge of the numerator and the top edge of the stroke is calculated. There is a style-dependent minimal distance min_dist. If num_dist is too small, it is increased up to min_dist. The denominator is handled analogously. Next, both numerator and denominator are centered to the maximum of their width. Then, we form a vertical list whose reference point is at the middle of the fraction stroke using vlist, and finally raise the resulting box to the level of the axis.

```
set_frac :: style -> box -> box -> box
set_frac st num den =
    raise ax fracbox
    where ax = axis_height st; th = rule_thickness st
                num_pos = num_level st - ax; den_pos = den_level st + ax
                num_dist = num_pos - depth num - th/2
                den_dist = den_pos - height den - th/2
                min_dist = 3* th, if st = D
            = th, otherwise
                num_dist' = num_dist $max2 min_dist || maximum as infix operator
                den_dist' = den_dist $max2 min_dist
                w = (width num) $max2 (width den)
                num_list = [VSpace num_dist' 0, center w num]
                den_list = [VSpace den_dist' 0, center w den]
                fracbox = vlist (Rule (th/2) (th/2) w) num_list den_list
```


### 5.4 Superscripts and Subscripts in Limit Position

The following functions deal with superscripts and subscripts in the limit position, i.e., vertically above and below the nucleus as in $\sum_{i=1}^{\infty}$. (In the nolimit position, they are to the right of the nucleus.) Function lim_sup deals with the case of superscripts only, lim_sub is called if there are only subscripts, and lim_supsub is for joint superscript / subscript combinations. We assume that nucleus, superscript, and subscript are already given as boxes.

In limit position, superscripts are placed above the nucleus in some distance, and white space is added above them. Superscript and nucleus are first centered to their maximum width. Afterwards,
the superscript is shifted to the right by some amount shift which depends on the slant of the nucleus. This is visible in e.g., $\int_{1}^{1}$. Subscripts are handled symmetrically.

To partially reduce the three functions to two, we use two auxiliary functions mksup and mksub which transform superscripts and subscripts into a list of boxes. For superscripts, the list is ordered upward, and for subscripts downward.

```
mksup :: style -> dim -> box -> [box]
mksup st shift sup =
    [VSpace dist 0 , right shift sup, VSpace space 0] || upward list
    where dist \(=\) (big_op1 st) \$max2 (big_op3 st - depth sup)
            space \(=\) big_op5 st
mksub :: style -> dim -> box -> [box]
mksub st shift sub =
    [VSpace dist 0, right (-shift) sub, VSpace space 0] || downward list
    where dist \(=\left(b i g \_o p 2\right.\) st) \(\$\) max 2 (big_op4 st - height sub)
                space \(=\) big_op5 st
```

In the actual functions, the appropriate auxiliary functions are called and their results are vertically combined.

```
lim_sup :: style -> dim -> box -> box -> box
lim_sup st shift nuc sup =
    Vup (nuc' : sup_list)
    where w = (width sup) $max2 (width nuc)
                sup' = center w sup; nuc' = center w nuc
                sup_list = mksup st shift sup'
lim_sub :: style -> dim -> box -> box -> box
lim_sub st shift nuc sub =
    Vdn (nuc' : sub_list)
    where w = (width nuc) $max2 (width sub)
            nuc' = center w nuc; sub' = center w sub
            sub_list = mksub st shift sub'
lim_supsub :: style -> dim -> box -> box -> box -> box
lim_supsub st shift nuc sup sub =
    vlist nuc' sup_list sub_list
    where w = (width sup) $max2 (width nuc) $max2 (width sub)
        sup' = center w sup; nuc' = center w nuc; sub' = center w sub
        sup_list = mksup st shift sup'
        sub_list = mksub st shift sub'
```


### 5.5 Superscripts and Subscripts in Nolimit Position

Here, the superscripts and subscripts are put to the right of the nucleus as in $\sum_{i=1}^{\infty}$. Their exact position depends on the fact whether the nucleus is a "character box, possibly followed by a kern". This information is passed as a Boolean to the functions nolim_sup, nolim_sub, and nolim_supsub. The third function has an additional argument: the slant of the nucleus, which is used to move the superscript to the right. This is visible in e.g., $P_{2}^{2}$. The functions involving superscripts need the information whether the style is 'cramped'.

There are several auxiliary functions to compute the positions of superscripts and subscripts. For a (partial) motivation for the formulae appearing in these functions, we refer to the $\mathrm{T}_{\mathrm{E}} \mathrm{Xbook}$ [1].

```
sup_position :: style -> bool -> bool \(->\) dim \(->\) dim -> dim
sup_position st cramped is_char hnuc dsup =
    sup_pos \$max2 sup_level st cramped || hnuc = height nuc
    \$max2 dsup + abs (x_height st)/4 || dsup = depth sup
    where sup_pos \(=0\), if is_char
    \(=\) hnuc + sup_drop (script st), otherwise
sub_position0 :: style -> bool -> dim -> dim
sub_position0 st True dnuc \(=0 \quad| |\) dnuc \(=\) depth nuc
sub_position0 st False dnuc = dnuc + sub_drop (script st)
sub_position1 :: style -> bool -> dim -> dim -> dim
sub_position1 st is_char dnuc hsub \(=\quad| | ~ h s u b=h e i g h t ~ s u b ~\)
    sub_position0 st is_char dnuc
        \$max2 sub1 st
        \$max2 hsub - 4 * abs (x_height st) / 5
sub_position2 :: style -> bool -> dim -> dim
sub_position2 st is_char dnuc =
    sub_position0 st is_char dnuc \$max2 sub2 st
```

The cases where there are only superscripts or only subscripts are relatively simple. Note that every 'script' is extended to the right by scriptspace.

```
nolim_sup :: style -> bool -> bool -> box -> box -> box
nolim_sup st cramped is_char nuc sup =
    nuc $hconc (raise sup_pos sup')
    where dsup = depth sup; hnuc = height nuc
                sup_pos = sup_position st cramped is_char hnuc dsup
        sup' = extend scriptspace sup
nolim_sub :: style -> bool -> box -> box -> box
nolim_sub st is_char nuc sub =
    nuc $hconc (raise (-sub_pos) sub')
    where dnuc = depth nuc; hsub = height sub
        sub_pos = sub_position1 st is_char dnuc hsub
        sub' = extend scriptspace sub
```

The case of a joint superscript / subscript combination is much more difficult. First, the desired position sup-pos of the superscript is computed. It is the distance from the reference point of the superscript to the base line. From it, the distance sup_dist between the bottom edge of the superscript and the base line is derived. The subscript is handled analogously. There is a minimum value min_sup for sup_dist, and a minimum value min_dist for the total distance sup_dist + sub_dist between superscript and subscript. There are correction values corr_sup and corr_dist if these minimum values are not reached. The correction of sup_dist is done by raising both superscript and subscript, i.e., adding corr_dist to sup_dist and subtracting it from sub_dist. The correction of sup_dist + sub_dist is done by lowering the subscript, i.e., adding corr_dist to sub_dist.

```
nolim_supsub :: style -> bool -> bool -> dim -> box -> box -> box -> box
nolim_supsub st cramped is_char slant nuc sup sub =
    nuc $hconc sup_sub
    where dsup = depth sup; hsub = height sub
        hnuc = height nuc; dnuc = depth nuc
        sup_pos = sup_position st cramped is_char hnuc dsup
        sub_pos = sub_position2 st is_char dnuc
        sup_dist = sup_pos - dsup; sub_dist = sub_pos - hsub
```

```
min_dist = 4 * rule_thickness st
corr_dist = correction (sup_dist + sub_dist) min_dist
min_sup = 4 * abs (x_height st) / 5
corr_sup = correction sup_dist min_sup
sup_dist' = sup_dist + corr_sup
sub_dist' = sub_dist + corr_dist - corr_sup
sup' = right slant (extend scriptspace sup)
sub' = extend scriptspace sub
sup_sub = vlist (VSpace sup_dist' sub_dist') [sup'] [sub']
```

The amount of the necessary correction values is computed by the following function:

```
correction :: num -> num -> num
correction value min_value = min_value - value, if value < min_value
    = 0, otherwise
```

The definitions of sup_dist' and sub_dist' can be algebraically simplified to

```
sup_dist' = sup_dist $max2 min_sup
sub_dist' = ((sup_dist + sub_dist) $max2 min_dist) - sup_dist'
```

After that, corr_sup, corr_dist, and the function correction are no longer needed. We did not directly introduce the simplified definitions since they are hard to explain by themselves.

## 6 From Subformulae to Whole Formulae

### 6.1 Some Auxiliary Functions

In some cases, white space is appended to a symbol to compensate for its slant (italic correction).

```
it_corr :: (box, dim) -> box; it_corr (box, slant) = extend slant box
```

The function lim_ computes whether superscripts and subscripts are placed in limit position.

```
lim_ :: style -> class -> bool
lim_ D Op = True; lim_ st cl = False
```

Actually, this is a bit simplified since it only realizes $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ 's default rule. In full $\mathrm{T}_{\mathrm{E}} \mathrm{X}$, there are commands \limits and \nolimits which may be appended to large operators. The default rule is only applied if none of these commands is issued.

### 6.2 Translation of Math Items

Math items are translated by the function set_mitem. For the nucleus of formulae with superscripts or subscripts, we need a special version of set_mitem, called set_nuc, which not only returns a box, but also the information whether its argument was a single character, and if so, its slant (the value needed for italic correction).

```
set_nuc :: style -> bool -> mitem -> ((box, dim), bool)
set_nuc st cr (Sym cl mc) = ( set_sym st cl mc, True )
set_nuc st cr mitem = ((set_mitem st cr mitem, 0), False)
```

Function set_mitem deals with the various cases of mitems. It handles the recursive setting of subformulae and then passes control to specialized functions. Its Boolean parameter is the bit indicating cramped styles. Notice that denominators, overlined formulae, and subscripts are always cramped. Other subformulae inherit the cramp status of their context.

```
set_mitem :: style -> bool -> mitem -> box
set_mitem st cr (Sym cl mc) = it_corr (set_sym st cl mc)
set_mitem st cr (MathSpace mw) = set_space st mw
set_mitem st cr (Over ml) = set_over st (set_mlist st True ml)
set_mitem st cr (Under ml) = set_under st (set_mlist st cr ml)
set_mitem st cr (Frac num den)
    = set_frac st numbox denbox
        where st' = fract st; numbox = set_mlist st' cr num
                                denbox = set_mlist st' True den
set_mitem st cr (Sup nuc sup)
    = lim_sup st (slant/2) boxnuc' boxsup, if lim
    = nolim_sup st cr is_char boxnuc' boxsup, otherwise
        where lim = lim_ st (class_ nuc)
                    ((boxnuc, slant), is_char) = set_nuc st cr nuc
                    boxnuc' = it_corr (boxnuc, slant)
                    boxsup = set_mlist (script st) cr sup
set_mitem st cr (Sub nuc sub)
    = lim_sub st (slant/2) boxnuc' boxsub, if lim
    = nolim_sub st is_char boxnuc boxsub, otherwise
        where lim = lim_ st (class_ nuc)
                            ((boxnuc, slant), is_char) = set_nuc st cr nuc
            boxnuc' = it_corr (boxnuc, slant)
            boxsub = set_mlist (script st) True sub
set_mitem st cr (SupSub nuc sup sub)
    = lim_supsub st (slant/2) boxnuc' boxsup boxsub, if lim
    = nolim_supsub st cr is_char slant boxnuc boxsup boxsub, otherwise
        where lim = lim_ st (class_ nuc)
            st' = script st
            ((boxnuc, slant), is_char) = set_nuc st cr nuc
            boxnuc' = it_corr (boxnuc, slant)
            boxsup = set_mlist st' cr sup
            boxsub = set_mlist st' True sub
set_mitem st cr (Class_cmd cl ml) = set_mlist st cr ml
set_mitem st cr (Group ml) = set_mlist st cr ml
```

The case of the Style_cmd constructor is missing since it is handled by set_list presented below.

### 6.3 Translation of Math Lists

In the $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ book, there is a second pass during formula layout where appropriate spaces are inserted between adjacent atoms, ignoring any non-atoms in between.

Here, insertion of inter-atom spaces is done by the function set_mlist which translates math lists into boxes. It does its job by calling an auxiliary function do_mlist with an additional class argument. This argument remembers the class of the previously set item, ignoring items of class None. At the beginning of the math list, the remembered class is None.

```
set_mlist :: style -> bool -> mlist -> box
set_mlist st cr ml = do_mlist st cr None ml
```

Function do_mlist handles style commands, inserts inter-atom spaces (set_space), and calls set_mitem to translate items into boxes.
do_mlist $::$ style $->$ bool $\rightarrow$ class $\rightarrow$ mlist $->$ box

```
do_mlist st cr old [] = HBox []
do_mlist st cr old (Style_cmd st' : ml) = do_mlist st' cr old ml
do_mlist st cr old (mi : ml) =
    boxmi $hconc rest
    where boxmi = set_mitem st cr mi; new = class_ mi
        rest = do_mlist st cr old ml, if new = None
            = set_space st (space st old new) $hconc
                do_mlist st cr new ml, otherwise
```

The auxiliary function space : : style -> class $->$ class $->$ num computes the space between two 'atoms' according to the table in the $\mathrm{T}_{\mathrm{E}} \mathrm{Xbook}$ [1, page 170]. The output is assumed to be in relative mathematical units. It depends on the style, and the conversion to an absolute dimension by set_space is again style dependent.

### 6.4 Setting of Whole Formulae

Function set_display handles displayed formulae, and set_inline formulae within text lines.

```
set_display, set_inline :: mlist -> box
set_display ml = set_mlist D False ml || Display style, not cramped
set_inline ml = set_mlist T False ml || Text style, not cramped
```

Actually, the difference between these two functions should be bigger. In $\mathrm{T}_{\mathrm{E}} \mathrm{X}$, displayed formulae require some postprocessing for positioning, and potential line breaks are computed within inline formulae.

## 7 Conclusion

Let us summarize what we achieved by our description.

- Defining an adequate data type separates concerns, i.e., spacing aspects from structure aspects. This is of great help for a better understanding of the algorithm.
- Using a powerful parser instead of a macro expansion mechanism avoids some postprocessing of the input on the data structure representing math formulae. Hence, we can translate mlists to box terms in a single pass, whereas Knuth needs two passes.
- Defining the result data type (box terms) makes many aspects of the algorithm explicit, which are implicit or at most verbally described in Knuth's description.
- Using a functional description language forced us to transform the updatable global variables of Knuth's description into explicit function parameters. On the one hand, this adds complexity to the description, but on the other hand, the flow of information becomes visible: it can be seen where information comes from, where it is updated, and where it is used. Thus, it becomes apparent which subtasks depend on others, and which are independent from each other.
- Some constructs, e.g., roots, were not treated here for space reasons. They don't offer principally new problems, although their treatment in [1] is particularly hard to grasp.
- Some postprocessing parts of the algorithm look somewhat 'imperative'. These are those, where some subformulae are set independently of each other only to detect afterwards, that certain minimal distances between them are not satisfied (set_frac and nolim_supsub). Miranda is not the best language to describe this, but it is possible as you can see.
- In our presentation, we preferred clarity over efficiency. The main problem is the computation of the size attributes height, depth, and width of box terms which is repeated many terms. The attributes should already be computed when box terms are constructed, and stored at
their constructors so that e.g., HBox [box] becomes HBox dim dim dim [box]. The class_ attribute of mitems should be handled similarly. The necessary program transformations are not difficult, but afterwards, it is no longer obvious that the size attributes of a compound box are fully determined by its constituent boxes.


## References

[1] D.E. Knuth. The $T_{E} X b o o k$. Addison Wesley, 1986.
[2] D. Turner. An overview of Miranda. SIGPLAN Notices, December 1986.
[3] R. Wilhelm and R. Heckmann. Dokumentenverarbeitung. Addison Wesley, 1996.


[^0]:    ${ }^{1}$ Miranda is a trademark of Research Software Ltd.

