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Augmented monomials in terms of power sums

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Abstract

The problem of base changes for the classical symmetric functions has been solved a long time ago and has been incorporated into most computer software packages for symmetric functions. In this paper, we develop a simple recursive formula for the expansion of the augmented monomial symmetric functions into power sum symmetric functions. As corollaries, we present two algorithms that can be used to expressing the augmented monomial symmetric functions in terms of the power sum symmetric functions.

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Background

Any positive integer n can be written as a sum of one or more positive integers, i.e.,

$$n = \lambda_1 + \lambda_2 + \cdots + \lambda_r. \tag{1}$$

When the order of integers λ_i does not matter, this representation is known as an integer partition Andrews ([1976](#)) and can be rewritten as

$$n = t_1 + 2t_2 + \cdots + nt_n$$

where each positive integer i appears t_i times. If the order of integers λ_i is important, then the representation ([1](#)) is known as a composition. For

$$\lambda_1 \geqslant \lambda_2 \geqslant \cdots \geqslant \lambda_r$$

we have a descending composition. We notice that more often than not there appears the tendency of defining partitions as descending compositions and this is also the convention used in this paper. In order to indicate that

$$\lambda = [\lambda_1, \lambda_2, \ldots, \lambda_r] \qquad \text{or} \qquad \lambda = [1^{t_1} 2^{t_2} \ldots n^{t_n}]$$

is a partition of n, we use the notation $\lambda \vdash n$. We denote by $l(\lambda)$ the number of parts of λ, i.e.,

$$l(\lambda) = r \qquad \text{or} \qquad l(\lambda) = t_1 + t_2 + \cdots + t_n.$$

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If $\alpha, \beta \vdash n$, then α precedes β in the dominance order if and only if for any $k \geqslant 1$, the sum of the k largest parts of α is less than the sum of the k largest parts of β, i.e.,

$$\alpha_1 + \cdots + \alpha_k < \beta_1 + \cdots + \beta_k$$

for all $k \geqslant 1$. In this definition, partitions are extended by appending zero parts at the end as necessary. If $\alpha = [\alpha_1, \alpha_2, \ldots, \alpha_r]$ and $\beta = [\beta_1, \beta_2, \ldots, \beta_s]$ are partitions of the same positive integer, then α precedes β in the lexicographic order if there is a positive integer t with the following properties:

1. $t \leqslant r$ and $t \leqslant s$;
2. for ever positive integer $i \leqslant t, \alpha_i = \beta_i$; and
3. either $\alpha_{t+1} < \beta_{t+1}$ or $t = r$ and $t < s$.

When α precedes β in lexicographic order, we use the notation $\alpha \prec \beta$. If $\alpha \prec \beta$ or $\alpha = \beta$, then we use the notation $\alpha \preceq \beta$. It is clear that the dominance order implies lexicographical order.

We recall some basic facts about monomial symmetric functions. Proofs and details can be found in Macdonald's book (Macdonald 1995). Let $\lambda = [\lambda_1, \lambda_2, \ldots, \lambda_k]$ be a partition with $k \leqslant n$. Being given a set of variables $\{x_1, x_2, \ldots, x_n\}$, the monomial symmetric function

$$m_\lambda = m_{[\lambda_1, \lambda_2, \ldots, \lambda_k]}(x_1, x_2, \ldots, x_n)$$

on these variables is the sum of monomial $x_1^{\lambda_1} x_2^{\lambda_2}, \ldots, x_k^{\lambda_k}$ and all distinct monomials obtained from it by a permutation of variables. For instance, with $\lambda = [2, 1, 1]$ and $n = 4$, we have:

$$\begin{aligned}
m_{[2,1,1]} = {}& x_1^2 x_2 x_3 + x_1 x_2^2 x_3 + x_1 x_2 x_3^2 + x_1^2 x_2 x_4 \\
& + x_1 x_2^2 x_4 + x_1 x_2 x_4^2 + x_1^2 x_3 x_4 + x_1 x_3^2 x_4 \\
& + x_1 x_3 x_4^2 + x_2^2 x_3 x_4 + x_2 x_3^2 x_4 + x_2 x_3 x_4^2
\end{aligned}$$

In particular, when $\lambda = [k]$, we have the kth power sum symmetric function $p_k = p_k(x_1, x_2, \ldots, x_n)$, i.e.,

$$m_{[k]} = p_k = \sum_{i=1}^{n} x_i^k \,.$$

In every case $p_0(x_1, x_2, \ldots, x_n) = n$.

If $\lambda \vdash k$ then m_λ is a symmetric function of degree k. It is well-known that the set

$$\{m_\lambda(x_1, x_2, \ldots, x_n) \mid \lambda \vdash k \text{ and } l(\lambda) \leqslant n\}$$

is a basis for the vector space Λ_n^k of symmetric functions of degree k of n variables. The dimension of Λ_n^k is the number of partitions of k. The power sum symmetric functions p_k do not have enough elements to form a basis for Λ_n^k, there must be one function for every partition $\lambda \vdash k$. To that end in each case we form multiplicative function $p_\lambda = p_\lambda(x_1, x_2, \ldots, x_n)$ so that for

$$\lambda = [\lambda_1, \lambda_2, \ldots, \lambda_{l(\lambda)}]$$

we note

$$p_\lambda = p_{\lambda_1} p_{\lambda_2} \cdots p_{\lambda_{l(\lambda)}}.$$

Also, it is known that the set

$$\{p_\lambda(x_1, x_2, \ldots, x_n) \mid \lambda \vdash k \text{ and } l(\lambda) \leqslant n\}$$

is another basis for Λ_n^k.

For each partition

$$\lambda = [\lambda_1, \lambda_2, \ldots, \lambda_k] = [1^{t_1} 2^{t_2} \cdots r^{t_r}]$$

with $k \leqslant n$, the augmented monomial symmetric function

$$\tilde{m}_\lambda = \tilde{m}_{[\lambda_1, \lambda_2, \ldots, \lambda_k]}(x_1, x_2, \ldots, x_n)$$

is defined by

$$\tilde{m}_\lambda = t_1! t_2! \cdots t_r! \cdot m_\lambda.$$

In this paper, we develop a simple recursive formula for the expansion of the augmented monomial symmetric functions into power sum symmetric functions. As corollaries, we present two algorithms that can be used to expressing the augmented monomial symmetric functions in terms of the power sum symmetric functions.

Two theorems for expanding augmented monomials

The cardinality of a set A is usually denoted $|A|$. Recall that a partition of the set A is a collection of non-empty, pairwise disjoint subsets of A whose union is A.

A simple way to express the augmented monomial symmetric function $\tilde{m}_\lambda$ in terms of the power sum is given by

Theorem 1 *Let $[\lambda_1, \lambda_2, \ldots, \lambda_k]$ be an integer partition. Then*

$$\tilde{m}_{[\lambda_1, \lambda_2, \ldots, \lambda_k]} = p_{\lambda_k} \cdot \tilde{m}_{[\lambda_1, \lambda_2, \ldots, \lambda_{k-1}]}$$
$$- \sum_{i=1}^{k-1} \tilde{m}_{[\lambda_1, \ldots, \lambda_{i-1}, \lambda_i + \lambda_k, \lambda_{i+1}, \ldots, \lambda_{k-1}]}$$

where $\tilde{m}$ and p are functions of n variables, $n \geqslant k$.

Proof We denote by M the set of terms in the expression $p_{\lambda_k} \cdot \tilde{m}_{[\lambda_1, \lambda_2, \ldots, \lambda_{k-1}]}$, by M_k the set of terms in the expression $\tilde{m}_{[\lambda_1, \lambda_2, \ldots, \lambda_k]}$ and by M_i the set of terms in the expression $\tilde{m}_{[\lambda_1, \ldots, \lambda_{i-1}, \lambda_i + \lambda_k, \lambda_{i+1}, \ldots, \lambda_{k-1}]}$, for $i = 1, 2, \ldots, k-1$. According to

$$|M| = \frac{n \cdot n!}{(n-k+1)!}, \quad |M_k| = \frac{n!}{(n-k)!}$$

and

$$|M_i| = \frac{n!}{(n-k+1)!},$$

we get

$$|M| = |M_k| + \sum_{i=1}^{k-1} |M_i|.$$

Taking into account that $\{M_i\}_{1 \leqslant i \leqslant k}$ is pairwise disjoint, we deduce that $\{M_i\}_{1 \leqslant i \leqslant k}$ is a set partition of M. Therefore, the theorem is proved. □

Example 1 Replacing k by 2 in Theorem 1, we get

$$\tilde{m}_{[\lambda_1, \lambda_2]} = p_{\lambda_1} p_{\lambda_2} - p_{\lambda_1 + \lambda_2}. \tag{2}$$

Then, for $k = 3$, we obtain

$$\begin{aligned}\tilde{m}_{[\lambda_1, \lambda_2, \lambda_3]} = {}& p_{\lambda_3} \cdot \tilde{m}_{[\lambda_1, \lambda_2]} \\ & - \tilde{m}_{[\lambda_1 + \lambda_3, \lambda_2]} - \tilde{m}_{[\lambda_1, \lambda_2 + \lambda_3]}.\end{aligned} \tag{3}$$

By (2) and (3), we deduce that

$$\begin{aligned}\tilde{m}_{[\lambda_1, \lambda_2, \lambda_3]} = {}& p_{\lambda_1} p_{\lambda_2} p_{\lambda_3} - p_{\lambda_1} p_{\lambda_2 + \lambda_3} - p_{\lambda_2} p_{\lambda_1 + \lambda_3} \\ & - p_{\lambda_3} p_{\lambda_1 + \lambda_2} + 2 p_{\lambda_1 + \lambda_2 + \lambda_3}.\end{aligned}$$

It is clear that in the expansion of the augmented monomial $\tilde{m}_\lambda$ generated by Theorem 1, the number of terms is equal to the number of parts of λ.

The following result is immediate from Theorem 1.

Corollary 1 *Let $\lambda = [1^{t_1} 2^{t_2} \cdots]$ be an integer partition and let j be a positive integer such that $t_j > 0$. Then*

$$\tilde{m}_\lambda = p_j \cdot \tilde{m}_{\lambda^0} - \sum_{i>0} (t_i - \delta_{ij}) \tilde{m}_{\lambda^i}$$

where δ_{ij} is the Kronecker delta and

$$\lambda^i = [1^{t_1(i)} 2^{t_2(i)} \cdots],$$

with

$$t_r(0) = \begin{cases} t_r - 1, & \text{if } r = j, \\ t_r, & \text{otherwise,} \end{cases}$$

and

$$t_r(i) = \begin{cases} t_r - 1 - \delta_{ij}, & \text{if } r \in \{i, j\}, \\ t_r + 1, & \text{if } r = i + j, \\ t_r, & \text{otherwise} \end{cases}$$

for all $i > 0$.

In this corollary, if $\lambda \vdash k$ then we remark that $\lambda^0 \vdash k - j$ and $\lambda \prec \lambda^i$ for all $i > 0$ with $t_i > \delta_{ij}$. If $t_j = 1$ then we have $t_j(j) = -1$. This drawback is eliminated by the fact that $t_j - \delta_{jj} = 0$.

Example 2 For $\lambda = [1^3 2^1 3^1]$ and $j = 3$, by Corollary 1, we have

$$\lambda^0 = [1^3 2^1], \quad \lambda^1 = [1^2 2^1 4^1] \quad \text{and} \quad \lambda^2 = [1^3 5^1].$$

Clearly, the coefficient of $\tilde{m}_{\lambda^0}$ is p_3, the coefficient of $\tilde{m}_{\lambda^1}$ is -3, the coefficient of $\tilde{m}_{\lambda^2}$ is -1, and for $i > 2$ all the coefficients are 0. Thus, we obtain

$$\tilde{m}_{[1^3 2^1 3^1]} = p_3 \cdot \tilde{m}_{[1^3 2^1]} - 3\tilde{m}_{[1^2 2^1 4^1]} - \tilde{m}_{[1^3 5^1]}.$$

We remark that in the expansion of $\tilde{m}_\lambda$ generated by Corollary 1, the number of terms is equal to

$$\text{the number of distinct parts of } \lambda + \begin{cases} 1, & \text{for } t_j > 1, \\ 0, & \text{for } t_j = 1. \end{cases}$$

So, we can say that this corollary is a refined form of Theorem 1.

We denote by $\mathcal{P}_n$ the set of all partitions of $\{1, 2, \ldots, n\}$. The cardinality of the set $\mathcal{P}_n$ is well-known as the nth Bell number, B_n (see Sloane 2012, A000110). The Möbius function of $\mathcal{P}_n$ (Bender and Goldman 1975; Rota 1964), namely

$$\mu(v) = \prod_{i=1}^{|v|} (-1)^{|v_i| - 1}(|v_i| - 1)!, \tag{4}$$

can be used to express the augmented monomial symmetric functions in terms of the power sum symmetric functions.

Theorem 2 *Let λ be an integer partition. Then*

$$\tilde{m}_\lambda = \sum_{v \in \mathcal{P}_{l(\lambda)}} \mu(v) p_{s(v)}$$

where $s(v) = [s_1, s_2, \ldots, s_{|v|}]$ is an integer partition with

$$s_i = \sum_{j \in v_i} \lambda_j, \quad i = 1, \ldots, |v|,$$

$\tilde{m}$ and p are functions of n variables, $n \geqslant l(\lambda)$.

Proof Let $\lambda = [\lambda_1, \lambda_2, \ldots, \lambda_k]$ be an integer partition. For $v = (v_1, v_2, \ldots, v_r) \in \mathcal{P}_{k-1}$ and $1 \leqslant i \leqslant r$ let us consider $f(v), f_i(v) \in \mathcal{P}_k$ defined by

$$f(v) = (v_1, v_2, \ldots, v_r, \{k\})$$

and

$$f_i(v) = (v_1, \ldots, v_{i-1}, v_i \cup \{k\}, v_{i+1}, \ldots, v_r).$$

By (4), we deduce that

$$\mu(f(v)) = \mu(v) \quad \text{and} \quad \mu(f_i(v)) = -|v_i|\mu(v). \tag{5}$$

Let $\mathcal{P'}_k$ be a subset of $\mathcal{P}_k$ defined by

$$\mathcal{P'}_k = \left\{ v \in \mathcal{P}_k : \{k\} \notin v \right\}. \tag{6}$$

We are to prove the theorem by induction on k. For $k = 1$, we have $\mu(\{1\}) = 1$ and $s(\{1\}) = [\lambda_1]$. Considering that $\tilde{m}_{[\lambda_1]} = \mu(\{1\})p_{s(\{1\})}$, the base case of induction is finished. We suppose that the relation

$$\tilde{m}_{[\lambda_1, \lambda_2, \ldots, \lambda_{k'}]} = \sum_{v \in \mathcal{P}_{k'}} \mu(v) p_{s(v)}$$

is true for any integer $k', 1 \leqslant k' < k$. By (5), (6) and Theorem 1, we can write

$$\tilde{m}_{[\lambda_1, \lambda_2, \ldots, \lambda_k]}$$

$$= p_{\lambda_k} \cdot \sum_{v \in \mathcal{P}_{k-1}} \mu(v) p_{s(v)} - \sum_{i=1}^{k-1} \sum_{v \in \mathcal{P}_{k-1}} \mu(f_i(v)) p_{s(f_i(v))}$$

$$= \sum_{v \in \mathcal{P}_k - \mathcal{P'}_k} \mu(v) p_{s(v)} + \sum_{v \in \mathcal{P'}_k} \mu(v) p_{s(v)}.$$

Thus, the proof is finished. □

Example 3 For $\{1, 2, 3\}$, we have $\mathcal{P}_3 = \{a, b, c, d, e\}$ with

$a = \{\{1\}, \{2\}, \{3\}\},$
$b = \{\{1\}, \{2, 3\}\},$
$c = \{\{2\}, \{1, 3\}\},$
$d = \{\{3\}, \{1, 2\}\}$ and
$e = \{\{1, 2, 3\}\}.$

According to (4), we have

$\mu(a) = (-1)^{3-3} 0!0!0! = 1$
$\mu(b) = \mu(c) = \mu(d) = (-1)^{3-2} 0!1! = -1$ and
$\mu(e) = (-1)^{3-1} 2! = 2.$

Taking into account Theorem 2, we get

$$\tilde{m}_{[2,1,1]} = p_{[2,1,1]} - p_{[2,2]} - 2p_{[3,1]} + 2p_{[4]}.$$

Iterative algorithm for computing transition matrix

If $\lambda \vdash k$, then it is immediate from Theorem 1 or Theorem 2 the fact that the augmented monomial symmetric function $\tilde{m}_\lambda$ is a sum over integer partitions of k.

Corollary 2 *Let λ be an integer partition. Then*

$$\tilde{m}_\lambda = \sum_{\lambda \preceq \beta} T_{\lambda\beta} \cdot p_\beta$$

where $T_{\lambda\beta}$ is an integer such that

$$(-1)^{l(\lambda)-l(\beta)} T_{\lambda\beta} \geqslant 0,$$

m̃ and p are functions of n variables, $n \geqslant l(\lambda)$.

We observe that the transition matrix expanding the augmented monomial symmetric functions in p_λ is lower triangular (with respect to any extension of the dominance ordering on partitions to a total order on the partitions $\lambda \vdash k$), i.e.,

$$\begin{bmatrix} \tilde{m}_{[k^1]} \\ \vdots \\ \tilde{m}_{[1^k]} \end{bmatrix} = T^{(k)} \cdot \begin{bmatrix} p_{[k^1]} \\ \vdots \\ p_{[1^k]} \end{bmatrix},$$

where

$$T^{(k)} = \left[T_{\lambda\beta} \right]_{\lambda,\beta \vdash k},$$

with

$$T_{\lambda\beta} = \begin{cases} 0, & \text{for } \lambda \not\preceq \beta, \\ 1, & \text{for } \lambda = \beta. \end{cases} \tag{7}$$

Example 4 For $k = 4$, according to Theorems 1 or 2, we obtain

$$\begin{bmatrix} \tilde{m}_{[4]} \\ \tilde{m}_{[3,1]} \\ \tilde{m}_{[2^2]} \\ \tilde{m}_{[2,1^2]} \\ \tilde{m}_{[1^4]} \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ -1 & 1 & 0 & 0 & 0 \\ -1 & 0 & 1 & 0 & 0 \\ 2 & -2 & -1 & 1 & 0 \\ -6 & 8 & 3 & -6 & 1 \end{bmatrix} \begin{bmatrix} p_{[4]} \\ p_{[3,1]} \\ p_{[2^2]} \\ p_{[2,1^2]} \\ p_{[1^4]} \end{bmatrix}.$$

We remark that

$$\tilde{m}_{[1^k]} = k! \cdot m_{[1^k]} = k! \cdot e_k,$$

where e_k is the kth elementary symmetric function. For $k = t_1 + 2t_2 + \cdots + kt_k$, the number of ways of partitioning a set of k different objects into t_i subsets containing i objects, $i = 1, 2, \ldots, k$ is

$$\frac{k!}{\prod_{i=1}^{k} t_i! \cdot (i!)^{t_i}}$$

[see (s.24.1.2 Abramovitz and Stegun 1972)]. Thus, the formula

$$T_{[1^k][1^{t_1} 2^{t_2} \cdots k^{t_k}]} = (-1)^{k-t_1-t_2-\cdots-t_k} \frac{k!}{\prod_{i=1}^{k} t_i! i^{t_i}},$$

where $k = t_1 + 2t_2 + \cdots + kt_k$, can be easily derived from Theorem 2. Unfortunately, for $T_{\lambda\beta}$ with $[1^k] \prec \lambda$ and $\lambda \prec \beta$ such formulas are not known.

The following result is immediate from Corollaries 1 and 2.

Corollary 3 *Let k be a positive integer. If $\lambda = [1^{t_1}2^{t_2}\cdots]$ and $\beta = [1^{v_1}2^{v_2}\cdots]$ are two integer partitions of k such that $\lambda \prec \beta$ then*

$$T_{\lambda\beta} = \left(1 - \delta_{0,v_j}\right)T_{\lambda^0\beta^0} - \sum_{i>0}(t_i - \delta_{ij})T_{\lambda^i\beta},$$

where j is a positive integer such that $t_j > 0$, δ_{ij} is the Kronecker delta,

$$\beta^0 = [1^{v_1(0)}2^{v_2(0)}\cdots] \quad \text{and} \quad \lambda^i = [1^{t_1(i)}2^{t_2(i)}\cdots],$$

with

$$v_r(0) = v_r - \delta_{rj}, \quad t_r(0) = t_r - \delta_{rj}$$

and

$$t_r(i) = \begin{cases} t_r - 1 - \delta_{ij}, & \text{if } r \in \{i,j\}, \\ t_r + 1, & \text{if } r = i + j, \\ t_r, & \text{otherwise} \end{cases}$$

for all $i > 0$.

In this corollary, for $v_j = 0$ we have $v_j(0) = -1$. Fortunately, this drawback is eliminated by the fact that $1 - \delta_{0,v_j} = 0$. Recall that λ^0 is an integer partition of $k - j$ and $\lambda \prec \lambda^i$ for all $i > 0$ with $t_i > \delta_{ij}$. We remark that $\beta^0 \vdash k - j$ for $v_j > 0$.

Example 5 By Corollary 3, for $\lambda = [1^4]$ and $\beta = [1^13^1]$, we have

$$\begin{aligned} T_{[1^4][1^13^1]} &= T_{[1^3][3^1]} - 3T_{[1^22^1][1^13^1]} \\ &= -2T_{[1^12^1][3^1]} - 3\left(-2T_{[1^13^1][1^13^1]}\right) \\ &= -2\left(-T_{[3^1][3^1]}\right) + 6 \\ &= 8. \end{aligned}$$

According to (7) and Corollary 3, we obtain Algorithm 1 for computing the transition matrix $T^{(k)}$. We can see that in order to compute the transition matrix $T^{(k)}$, Algorithm 1 is based on generating the immediate lexicographic predecessor of an integer partition (see lines 10 and 22). The problem of generating the immediate lexicographic predecessor of an integer partition is well-known in literature. For more details, one can refer to (Kelleher and O'Sullivan 2009) and the references therein.

Algorithm 1 Computing transition matrix

Require: k

1:   $T_{[1][1]} \leftarrow 1$
2:   $r \leftarrow 2$
3:   **while** $r \leq k$ **do**
4:      $\lambda \leftarrow [r]$
5:      $T_{\lambda\lambda} \leftarrow 1$
6:      **for all** $\beta \prec \lambda$ **do**
7:         $T_{\lambda\beta} \leftarrow 0$
8:      **end for**
9:      **while** $[1^r] \prec \lambda$ **do**
10:        $\lambda \leftarrow$ predecessor of λ
11:        $\beta \leftarrow [r]$
12:        **while** $\lambda \prec \beta$ **do**
13:          $j \leftarrow \lambda_1$
14:          **if** $v_j > 0$ **then**
15:            $T_{\lambda\beta} \leftarrow T_{\lambda^0\beta^0}$
16:          **else**
17:            $T_{\lambda\beta} \leftarrow 0$
18:          **end if**
19:          **for all** $i > 0$ with $t_i > \delta_{ij}$ **do**
20:            $T_{\lambda\beta} \leftarrow T_{\lambda\beta} - (t_i - \delta_{ij})T_{\lambda^i\beta}$
21:          **end for**
22:          $\beta \leftarrow$ predecessor of β
23:        **end while**
24:        $T_{\lambda\lambda} \leftarrow 1$
25:        **for all** $\beta \prec \lambda$ **do**
26:          $T_{\lambda\beta} \leftarrow 0$
27:        **end for**
28:      **end while**
29:      **visit** $T^{(r)}$
30:      $r \leftarrow r + 1$
31:   **end while**

Example 6 Applying Algorithm 1 for $k = 5$, we get successively:

$$T^{(2)} = \begin{bmatrix} 1 & 0 \\ -1 & 1 \end{bmatrix},$$

$$T^{(3)} = \begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 2 & -3 & 1 \end{bmatrix},$$

$$T^{(4)} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ -1 & 1 & 0 & 0 & 0 \\ -1 & 0 & 1 & 0 & 0 \\ 2 & -2 & -1 & 1 & 0 \\ -6 & 8 & 3 & -6 & 1 \end{bmatrix},$$

$$T^{(5)} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 1 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 2 & -2 & -1 & 1 & 0 & 0 & 0 \\ 2 & -1 & -2 & 0 & 1 & 0 & 0 \\ -6 & 6 & 5 & -3 & -3 & 1 & 0 \\ 24 & -30 & -20 & 20 & 15 & -10 & 1 \end{bmatrix}.$$

At the end of this section, we remark the following

Conjecture 1 *Let k be a positive integer. The identities*

$$\sum_{\lambda \preceq \beta} T_{\lambda\beta} = 0 \qquad \text{and} \qquad \sum_{\nu \in \mathcal{P}_{l(\lambda)}} \mu(\nu) = 0$$

are true for all $\lambda \prec [k]$.

Recursive algorithm for computing an element of the transition matrix

A specific augmented monomial function $\tilde{m}_\lambda$ can be expressed in terms of power sums

$$\tilde{m}_\lambda = \sum_{\lambda \preceq \beta} T_{\lambda\beta} \cdot p_\beta$$

without computing the transition matrices

$$T^{(r)} = \left[T_{\lambda\beta} \right]_{\lambda,\beta \vdash r}, \quad r \leqslant l(\lambda).$$

According to (7) and Corollary 3, we obtain Algorithm 2 for computing the coefficient $T_{\lambda\beta}$ of the power sums p_β.

In Algorithm 2, λ and β are two integer partitions of k such that

$$\lambda = [1^{t_1} 2^{t_2} \cdots k^{t_k}] \quad \text{and} \quad \beta = [1^{v_1} 2^{v_2} \cdots k^{v_k}].$$

The recursive function TLB(k) is presented in a form that allows fast identification of the correlation between Corollary 3 and the operations executed with the arrays $(t_1, t_2, \ldots, t_k)$ and $(v_1, v_2, \ldots, v_k)$. Thus, the lines 2–9 are useful to determine whether $\beta = \lambda$ or $\beta \prec \lambda$. The value of j is selected in the lines 16–24 such that j is the largest positive integer with

$$t_j = \min\{t_i | t_i > 0\}.$$

This selection of j allows us to reduce the number of recursive calls from the lines 30 and 39.

The arrays $(t_1, t_2, \ldots, t_k)$ and $(v_1, v_2, \ldots, v_k)$ are the global variables of the recursive function TLB(k). These global variables are very important because help us save memory. The integer partitions λ and λ^i with $i \geqslant 0$ are alternatively stored in the same array $(t_1, t_2, \ldots, t_k)$. The integer partition λ^0 is immediately derived from the integer partition λ in the line 28. Then λ is derived from λ^0 in the line 31. The integer partition λ^i with $i > 0$ is derived from the integer partition λ in the lines 36–38. Then λ is derived from λ^i in the lines 40–42. The integer partitions β and β^0 are alternatively stored in the same array $(v_1, v_2, \ldots, v_k)$. The integer partition β^0 is immediately derived from the integer partition β in the line 29. Then β is derived from β^0 in the line 32.

The function TLB(k) can be integrated into any algorithm for generating integer partitions to get the expression of the augmented monomial $\tilde{m}_\lambda$ in terms of power sums.

Algorithm 2 Computing $T_{\lambda\beta}$

Require: $k, t_1, t_2, \ldots, t_k, v_1, v_2, \ldots, v_k$
 1: **function** $\text{TLB}(k)$
 2:     $i \leftarrow k$
 3:     $L \leftarrow i \cdot t_i$
 4:     $B \leftarrow i \cdot v_i$
 5:     **while** $L = B$ and $i > 1$ **do**
 6:         $i \leftarrow i - 1$
 7:         $L \leftarrow L + i \cdot t_i$
 8:         $B \leftarrow B + i \cdot v_i$
 9:     **end while**
10:     **if** $B = L$ **then**
11:         $S \leftarrow 1$
12:     **else**
13:         **if** $B < L$ **then**
14:             $S \leftarrow 0$
15:         **else**
16:             $j \leftarrow 1$
17:             **while** $t_j = 0$ **do**
18:                 $j \leftarrow j + 1$
19:             **end while**
20:             **for** $i = j + 1$ **to** k **do**
21:                 **if** $t_i > 0$ and $t_i < t_j$ **then**
22:                     $j \leftarrow i$
23:                 **end if**
24:             **end for**
25:             **if** $v_j = 0$ **then**
26:                 $S \leftarrow 0$
27:             **else**
28:                 $t_j \leftarrow t_j - 1$
29:                 $v_j \leftarrow v_j - 1$
30:                 $S \leftarrow \text{TLB}(k - j)$
31:                 $t_j \leftarrow t_j + 1$
32:                 $v_j \leftarrow v_j + 1$
33:             **end if**
34:             **for** $i = 1$ **to** k **do**
35:                 **if** $t_i > \delta_{ij}$ **then**
36:                     $t_i \leftarrow t_i - 1$
37:                     $t_j \leftarrow t_j - 1$
38:                     $t_{i+j} \leftarrow t_{i+j} + 1$
39:                     $S \leftarrow S - (t_i + 1) \cdot \text{TLB}(k)$
40:                     $t_i \leftarrow t_i + 1$
41:                     $t_j \leftarrow t_j + 1$
42:                     $t_{i+j} \leftarrow t_{i+j} - 1$
43:                 **end if**
44:             **end for**
45:         **end if**
46:     **end if**
47:     **return** S
48: **end function**

Concluding remarks

An iterative algorithm for computing transition matrix expanding the augmented monomial symmetric functions in terms of the power sums symmetric functions has been derived in this paper. It is clear that the efficiency of this algorithm is directly influenced by the efficiency of the algorithm used for generating integer partitions in reverse lexicographic order. To express a specific augmented monomials in terms of power sums, we need a single line of the transition matrix. In this case, the computation of all transition matrix elements is not justified. Thus, a recursive function that computes the value of a single element of the transition matrix has been derived. Clearly, behind these algorithms is Theorem 1.

A recursive algorithm that requires algebraic symbol manipulation for expressing the augmented monomial $\tilde{m}_\lambda$ in terms of power sums can be easily derived from Theorem 1. For instance, in Maple this algorithm can be written as

```
monom := proc(a :: list, n :: integer);
local i, s, b;
if n = 1 then
    s := p_{a[1]}
else
    if n = 2 then
        s := p_{a[1]} · p_{a[2]} − p_{a[1]+a[2]}
    else
        s := p_{a[n]} · monom(a, n − 1);
        for i from 1 by 1 while i < n do
            b := a;
            b[i] := b[i] + a[n];
            s := s − monom(b, n − 1);
        end do;
    end if;
end if;
sort(expand(s));
end proc;
```

and the command

$$monom([3, 2, 1, 1], 4)$$

generates the following expression

$$p_1^2 p_2 p_3 - p_1^2 p_5 - 2 p_1 p_2 p_4 - 2 p_1 p_3^2$$
$$- p_2^2 p_3 + 4 p_1 p_6 + 3 p_2 p_5 + 4 p_3 p_4 - 6 p_7.$$

Such a recursive algorithm is very simple but its effectiveness can not be called into question because of the large number of recursive calls (for the augmented monomial $\tilde{m}_\lambda$ the number of recursive calls is the factorial of $l(\lambda) - 1$).

Unfortunately, Theorem 2 is more difficult to exploit in order to give similar results. However, a special case can be considered.

Corollary 4   *Let $\lambda = [\lambda_1, \lambda_2, \ldots, \lambda_r]$ be an integer partition of k such that*

$$\lambda_i > \lambda_{i+1} + \cdots + \lambda_r$$

for all $i < r$.

1. The number of integer partition β with $\lambda \preceq \beta$ is greater than or equal to B_r.
2. The number of integer partition β with $T_{\lambda\beta} = 0$ is equal to

   $$p(k) - B_r,$$

   where $p(k)$ is the Euler partition function.
3. For all $\nu \in \mathcal{P}_r$ the following formula holds:

   $$T_{\lambda,s(\nu)} = \mu(\nu),$$

   where $s(\nu) = [s_1, s_2, \ldots, s_{|\nu|}]$ is an integer partition with

   $$s_i = \sum_{j \in \nu_i} \lambda_j, \quad i = 1, \ldots, |\nu|.$$

This corollary is immediate from Theorem 2 and Corollary 3.

Competing interests

The authors declare that they have no competing interests.

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