# PRIORITIZATION METHOD FOR FRONTIER DMUs: A DISTANCE-BASED APPROACH

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In nonparametric methods, if the number of observations is relatively small as compared to the sum of number of inputs and outputs, many units are evaluated as efficient. Several methods for prioritizing these efficient units are reported in literature. Andersen et al. and Mehrabian et al. proposed two methods for ranking efficient units, but both methods break down in some cases. This paper describes a new DEA ranking approach that uses  $L_2$ -norm.

# 1. Introduction

Data envelopment analysis (DEA) originally from Farell's [7] seminal work and popularized by Charnes et al. [5] has gained a wide range of applications measuring comparative efficiency. The main point of DEA is to assign an efficiency score to each of the decisionmaking units (DMUs). This efficiency score depends on the orientation of the problem. The data set in DEA that consists of *n* DMUs defines a production set with a subset consisting of boundary points that form an efficient frontier. The main object of DEA is to determine the relative position of a unit in a production set. Banker et al. [3] and Charnes, Cooper, and Rhodes (CCR) [5] used the term *Production possibility set* (PPS) for production set. In formal terms, we define *T* as

$$T = \{(X, Y) \mid \text{input } X \text{ can produce output } Y\}.$$
(1.1)

The usefulness of the set notion for our purposes comes from the fact that it brings along two further notions, namely, (i) that of the boundary (or frontier) of the set, and (ii) that of the interior of the set. This permits one to distinguish between DMUs that belong to the interior of the production set, which are called inefficient, and those that do belong to a particular part of the frontier and are called efficient. DEA assigns 1 to all efficient DMUs. This approach ranks all DMUs based upon their efficiency score. In cases where several DMUs have the same efficiency score 1, DEA approach cannot rank these

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DMUs. However, decision makers are often interested in a complete ranking, beyond the dichotomized classification, in order to refine the evaluation of the units. Discriminating between these efficient DMUs is an interesting subject. This problem has recently been studied by a number of DEA researchers (see [1, 2, 6, 8] etc.). Andersen and Petersen [1] proposed a method for ranking efficient DMUs (AP model). In their method the column corresponding to the unit under consideration is omitted from technological matrix. However, the feasibility for the problem is not guaranteed and for small variation in data, some problems are not stable. Another method was suggested by Mehrabian et al. [8] (MAJ model). Although in MAJ model the basic idea is same as AP model, the movement strategy towards frontier is different from AP model. However, in some cases, the infeasibility for the problem may occur; we will prove that too. The method proposed by the authors removes the existing difficulties in both methods. The paper is structured as follows. Section 2 provides basic DEA models. An introduction to AP and MAJ models is given in Section 3. The next section of the paper addresses a method for ranking efficient DMUs. Some numerical examples are solved in Section 5. Conclusions appear in Section 6.

# 2. Background

To describe the DEA efficiency measurement, let there be *n* DMUs and let the performance of each DMU be characterized by a production process of *m* inputs ( $x_{ij}$ , i = 1,...,m) to yield *s* outputs ( $y_{rj}$ , r = 1,...,s).

*Definition 2.1.*  $DMU_p$  is called inefficient if and only if it is dominated within the set of the *n* DMUs.

The *p*th DMU is dominated when

$$\sum_{j=1}^{n} \lambda_j x_{ij} \le x_{ip} \quad \text{for } i = 1, 2, \dots, m,$$

$$\sum_{j=1}^{n} \lambda_j y_{rj} \ge y_{rp} \quad \text{for } r = 1, 2, \dots, s,$$

$$\lambda_j \ge 0, \quad \text{for } j = 1, 2, \dots, n,$$

$$(2.1)$$

and at least one of the above conditions is a strict inequality. Every nondominant DMU is called efficient. To estimate a DEA efficiency score of the specific *p*th DMU, we use the following original DEA model (Charnes et al. [5]):

Max 
$$e_p$$
 subject to  $e_q = \frac{\sum_{r=1}^{s} u_r y_{rq}}{\sum_{i=1}^{m} v_i x_{iq}}, \quad q = 1, \dots, n,$   
 $0 \le e_q \le 1, \quad q = 1, \dots, n,$   
 $u_r \ge \epsilon, \quad r = 1, \dots, s,$   
 $v_i \ge \epsilon, \quad i = 1, \dots, m,$ 

$$(2.2)$$

where  $e_p$ , p = 1,...,n, are the efficiency of DMU<sub>p</sub>. This model can be reduced to a nonratio format in the usual manner of Charnes and Cooper [4] as follows:

$$\operatorname{Max} \bar{e}_{p} = \sum_{r=1}^{s} u_{r} y_{rp}, \quad \text{subject to } \sum_{i=1}^{m} v_{i} x_{ip} = 1,$$
$$\sum_{r=1}^{s} u_{r} y_{rj} - \sum_{i=1}^{m} v_{i} x_{ij} \leq 0, \quad j = 1, \dots, n,$$
$$u_{r} \geq \epsilon, \quad r = 1, \dots, s,$$
$$v_{i} \geq \epsilon, \quad i = 1, \dots, m.$$
$$(2.3)$$

Hence, to determine all DMU efficiencies, the above program is solved *n* times, once for each DMU<sub>*p*</sub>, p = 1,...,n.

#### 3. AP model and MAJ model

**3.1. AP model.** Andersen and Petersen [1] developed a procedure for ranking efficient units. Their method enables an efficient  $DMU_p$  to achieve an efficiency score greater than or equal to one by removing the *p*th constraint in (2.3). The mathematical formulation of this model is as follows:

$$\operatorname{Max} \sum_{r=1}^{s} u_r y_{rp}, \quad \text{subject to} \quad \sum_{i=1}^{m} v_i x_{ip} = 1,$$

$$\sum_{r=1}^{s} u_r y_{rj} - \sum_{i=1}^{m} v_i x_{ij} \le 0, \quad j = 1, \dots, n, \ j \ne p,$$

$$u_r \ge \epsilon, \quad r = 1, \dots, s,$$

$$v_i \ge \epsilon, \quad i = 1, \dots, m.$$
(3.1)

The dual formulation of (3.1) is as follows:

$$\operatorname{Min} \theta_{p} - \epsilon \left[ \sum_{i=1}^{m} s_{i}^{-} + \sum_{r=1}^{s} s_{r}^{+} \right] \quad \text{subject to} \quad \sum_{\substack{j=1\\j\neq p}}^{n} \lambda_{j} x_{ij} + s_{i}^{-} = \theta_{p} x_{ip}, \quad i = 1, \dots, m,$$

$$\sum_{\substack{j=1\\j\neq p}}^{n} \lambda_{j} y_{rj} - s_{r}^{+} = y_{rp}, \quad r = 1, \dots, s,$$

$$\lambda_{j}, s_{i}^{-}, s_{r}^{+} \ge 0, \quad \forall i, j, r.$$

$$(3.2)$$

However, there are some problematic areas with this methodology that is shown by the following example. There are 5 DMUs with two inputs and two outputs. We consider three cases of DMU1, denoted by DMU1<sub>1</sub>, DMU1<sub>2</sub> and DMU1<sub>3</sub> (see Table 3.1).

The AP model for evaluating  $DMU1_2$  is infeasible, this model for evaluating  $DMU1_3$  leads to a large score. Table 3.2 presents the results.

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	$DMU1_1$	DMU12	DMU1 <sub>3</sub>	DMU2	DMU3	DMU4	DMU5
$I_1$	2	0	0.09	5	10	10	2
$I_2$	10	10	10	7	5	7	13
$O_1$	2	2	2	2	3	2	3
$O_2$	3	3	3	3	1	1	2

Table 3.1. The data set for infeasibility and unstability of AP model.

Table 3.2. The AP and CCR efficiency indexes.

Efficiency index	$DMU1_1$	DMU1 <sub>2</sub>	DMU1 <sub>3</sub>
CCR	1.000	1.000	1.000
AP	1.700	Infeasible	33.300

**3.2. MAJ model.** An alternative approach suggested in the literature, is MAJ model (Mehrabian and et al. [8]). The mathematical formulation of MAJ model is as follows:

$$\begin{array}{ll} \text{Min } 1 + W_p & \text{subject to } \sum_{\substack{j=1\\j \neq p}}^n \lambda_j x_{ij} + s_i^- = W_p + x_{ip}, \quad i = 1, \dots, m, l \\ & \sum_{\substack{j=1\\j \neq p}}^n \lambda_j y_{rj} - s_r^+ = y_{rp}, \quad r = 1, \dots, s, \\ & \lambda_j, s_i^-, s_r^+ \ge 0, \quad \forall i, j, r. \end{array}$$

$$(3.3)$$

In what follows, we show that MAJ model fails to rank DMUs in some cases.

(i) *First deficiency*. For special data, MAJ model is infeasible. To show this, consider the dual formulation of MAJ model as follows:

$$\text{Max} \sum_{r=1}^{s} u_r y_{rp} - \sum_{i=1}^{m} v_i x_{ip}, \quad \text{subject to} \ \sum_{r=1}^{s} u_r y_{rj} - \sum_{i=1}^{m} v_i x_{ij} \le 0, \quad j = 1, \dots, n, \ j \ne p,$$

$$\sum_{i=1}^{m} v_i x_{ip} = 1,$$

$$u_r, v_i \ge 0, \quad \forall i, r.$$

$$(3.4)$$

Consider the system of linear equations

$$Y^{p}d_{u} = 0, \qquad Y_{p}d_{u} - t = 0,$$
 (3.5)

	DMU1	DMU2	DMU3	DMU4	DMU5
$I_1$	2	0	0.1	5	10
$I_2$	8	8	8	5	4
$O_1$	1	1	1	1	2
$O_2$	2	2	2	1	1

Table 3.3. The data set.

Table 3.4. The results for MAJ model.

	DMU1	DMU2	DMU3	DMU4	DMU5
MAJ input orientation	4	2	3	6	1
MAJ output orientation	4	1	3	6	2

where t > 0;  $Y^p$  is an (n - 1)s-matrix of outputs of  $DMU_1, ..., DMU_{p-1}, DMU_{p+1}, ..., DMU_n$ ;  $Y_p$  is the output vector of  $DMU_p$ ; and  $d_u \ge 0$  is a nonzero *s*-vector. For some  $Y^p$  and  $Y_p$ , (3.5) has a solution. For example

$$Y_{p} = (0, \dots, 0, y_{kp} > 0, 0, \dots, 0),$$

$$d_{u} = \left(0, \dots, 0, \frac{t}{y_{kp}}, 0, \dots, 0\right),$$

$$Y^{p} = \begin{bmatrix} y_{11} & y_{12} & \cdots & y_{1p} = 0 & \cdots & y_{1s} \\ y_{21} & y_{22} & \cdots & y_{2p} = 0 & \cdots & y_{2s} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ y_{n-1,1} & y_{n-1,2} & \cdots & y_{n-1,p} = 0 & \cdots & y_{n-1,s} \end{bmatrix}$$
(3.6)

is a solution to (3.5). Now, consider  $d = [d_u, 0]^T$ . It can be seen that  $Ad = 0, d \ge 0, d \ne 0, Cd > 0$ , where A is the technological matrix and C is the cost coefficient in the objective function in (3.4). This means that for special data, (3.4) is unbounded and by duality theorem, MAJ model is infeasible. This completes the proof.

(ii) *Second deficiency.* The following example shows that in some cases of MAJ model, ranking in input orientation is different from ranking in output orientation. Our illustrative example entails 5 DMUs (see Table 3.3). Each unit consumes two inputs to produce two outputs.

Running MAJ model yields the results presented in Table 3.4.

It is evident that ranking in input orientation is different from ranking in output orientation (see DMU2 and DMU5).

	DMU1	DMU2	DMU3
$I_1$	5	2	1
$I_2$	1	2	4
0	1	1	1

Table 4.1. The data set for the simple example.

#### 4. A distance-based approach

This section describes a new DEA-based ranking approach that uses  $L_2$ -norm. In the proposed method, to rank an efficient unit, this unit is omitted from PPS and the new frontier is constructed. Then, the shortest distance ( $L_2$ -norm) from this unit to the new frontier is determined. This is done for all DEA efficient units.

Definition 4.1. A surface  $H = \{(X, Y) : -\alpha^T X + \beta^T Y = 0, \alpha \ge 0, \beta \ge 0\} \cap T_c$  is called a supporting efficient surface if for each observation  $j = 1, ..., n, -\alpha^T X_j + \beta^T Y_j \le 0$  and at least m + s - 1 inequalities are binding.

To have at least m + s - 1 inequalities  $-\alpha^T X_j + \beta^T Y_j \le 0$  as equalities, we use the slack variables  $s_j$ . The constraints  $s_j \le (1 - \gamma_j)M$ ,  $s_j \ge 0$ ,  $\gamma_j \in \{0,1\}$ ,  $\sum_{j \in E} \gamma_j \ge m + s - 1$  force some of the  $s_j$  at zero level (*M* is a large positive constant). Clearly, selecting  $\gamma_t = 1$  forces the  $s_t = 0$ . (*E* denotes the set of all CCR-efficient DMUs.) Consider the following program:

$$D_{p} = \operatorname{Max} - \alpha^{T} X_{p} + \beta^{T} Y_{p} \quad \text{subject to} \quad -\alpha^{T} X_{j} + \beta^{T} Y_{j} + s_{j} = 0, \quad j \in E, \ j \neq p,$$

$$\alpha^{T} 1_{m} + \beta^{T} 1_{s} = 1,$$

$$s_{j} \leq (1 - \gamma_{j}) M, \quad j \in E, \ j \neq p,$$

$$\sum_{j \in E, \ j \neq p} \gamma_{j} \geq m + s - 1,$$

$$\alpha \geq \epsilon \cdot 1_{m},$$

$$\beta \geq \epsilon \cdot 1_{s},$$

$$\gamma_{j} \in \{0, 1\}, \quad j \in E, \ j \neq p,$$

$$s_{j} \geq 0, \quad j \in E, \ j \neq p,$$

$$(4.1)$$

where  $\alpha$  is an *m*-vector,  $\beta$  is an *s*-vector,  $\epsilon > 0$ ,  $X_j = (x_{1j}, \dots, x_{mj})^T$ ,  $Y_j = (y_{1j}, \dots, y_{sj})^T$ and  $1 = (1, 1, \dots, 1)^T$ . It is inherent in this formulation of the DEA model that not all inefficiencies will be captured by the radial efficiency measures, but that slacks also have to be considered when judging the extent of inefficiency. In this model, efficiency relates to the observed ability to generate more output consuming less input (see the dual version of (4.1)). This model computes the distance between the efficient frontier, evaluated without unit *p*, and the unit itself. To highlight the practical implication of this model, consider the data set that consists of two inputs and one output. The data are listed in Table 4.1. Model 8 for evaluating  $DMU_2$  is as follows:

$$D_{2} = Max - 2\alpha_{1} - 2\alpha_{2} + \beta \quad \text{subject to} \quad -5\alpha_{1} - \alpha_{2} + \beta + s_{1} = 0,$$
  

$$-\alpha_{1} - 4\alpha_{2} + \beta + s_{2} = 0,$$
  

$$\alpha_{1} + \alpha_{2} + \beta = 1,$$
  

$$s_{1} \leq (1 - \gamma_{1})M,$$
  

$$s_{2} \leq (1 - \gamma_{2})M,$$
  

$$\gamma_{1} + \gamma_{2} \geq 2,$$
  

$$\gamma_{1}, \gamma_{2} \in \{0, 1\} \cdot \alpha_{1}, \alpha_{2}, \beta \geq 0.$$
  
(4.2)

The optimal solution to this program is  $(\alpha_1, \alpha_2, \beta) = (3/26, 4/26, 19/26), D_2 = 5/26$ , and the supporting efficient surface is

$$F' = \left\{ (x_1, x_2, y) \mid -\frac{3}{26} x_1 - \frac{4}{26} x_2 + \frac{19}{26} y = 0 \right\} \cup T'_c.$$
(4.3)

(Note that  $T'_c$  is the new PPS obtained from omission of  $DMU_p$ .) From (4.2) it is easy to show that  $F' = \{(X, Y) : -\alpha^{*^T}X + \beta^{*^T}Y = 0\} \cap T'_c$  is an efficient supporting surface to  $T'_c$ , in which  $(\alpha^*, \beta^*)$  is the optimal solution to (4.1) in evaluating  $DMU_2$ . Consider the following two cases:

(i) *F*' is unique and hence  $(\alpha^*, \beta^*)$  is unique. In this case, set

$$EF_p = \frac{D_p}{\sqrt{\alpha^{*^T} \alpha^* + \beta^{*^T} \beta^*}},\tag{4.4}$$

(ii) *F*' is not unique. In this case, let  $X_{\lambda} = \sum_{j \neq p} \lambda_j X_j$ ,  $Y_{\lambda} = \sum_{j \neq p} \lambda_j Y_j$  in which  $\lambda_j \ge 0$  for all *j*. Set

 $EF_p = Min_{\lambda \ge 0}$  {the distance between  $(X_p, Y_p)$  and  $(X_\lambda, Y_\lambda)(L_2\text{-norm})$ }. (4.5)

The proposed efficiency index is now defined as

$$\widehat{EF}_p = \frac{EF_p}{\operatorname{Max}_{1 \le j \le n} \{EF_j\}} \le 1.$$
(4.6)

It is to be noted that in a real data set, the second case is of seldom occurrence.

Definition 4.2.  $(X_p, Y_p) \in T_c$  is nondominant if and only if there exist no  $\lambda = (\lambda_1, ..., \lambda_n) \ge 0$  such that  $\sum_{j=1}^n \lambda_j X_j \le X_p$  and  $\sum_{j=1}^n \lambda_j Y_j \ge Y_p$  and strict inequality holds true for at least one component.

The following theorem gives a necessary and sufficient condition for efficiency of  $DMU_p$ .

THEOREM 4.3. Considering (4.1),  $D_p \ge 0$  if and only if  $DMU_p$  is efficient in CCR model.

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*Proof.* We will first prove that if  $D_p \ge 0$  then  $DMU_p$  is nondominant in  $T_c$ . Suppose, on the contrary, that  $DMU_p$  is dominant. Then, there exist some  $\lambda = (\lambda_1, ..., \lambda_n)$  such that  $\lambda_j \ge 0, j = 1, ..., n$ , and

$$X_p \ge \sum_{j \ne p} \lambda_j X_j = X_\lambda, \qquad Y_p \le \sum_{j \ne p} \lambda_j Y_j = Y_\lambda,$$

$$(4.7)$$

and strict inequality holds true for at least one component. Now,

$$0 \le -\alpha^{*^{T}} X_{p} + \beta^{*^{T}} Y_{p} < -\alpha^{*^{T}} X_{\lambda} + \beta^{*^{T}} Y_{\lambda},$$
(4.8)

where  $(\alpha^*, \beta^*)$  is the optimal solution to (4.1). But we note that  $(X_{\lambda}, Y_{\lambda}) \in T_c$  and

$$H^{-} = \{ (X, Y) : -\alpha^{T} X + \beta^{T} Y \le 0 \}$$
(4.9)

is a convex set, hence we must have

$$-\alpha^{*^{T}}X_{\lambda} + \beta^{*^{T}}Y_{\lambda} \le 0 \tag{4.10}$$

which contradicts (4.2). To show the converse, suppose that  $DMU_p$  is CCR-efficient; this means that  $Z_p^* \ge 1$ , where

$$Z_p^* = \operatorname{Max} U^T Y_p,$$
  
s.t.  $V^T X_p = 1,$   
 $U^T Y_j - V^T X_j \le 0, \quad j = 1, \dots, n, \ j \ne p,$   
 $U \ge \epsilon \cdot 1,$   
 $V \ge \epsilon \cdot 1.$   
(4.11)

Hence, if  $(U^*, V^*)$  is the optimal solution of this problem, we must have  $U^T Y_p - V^T X_p \ge 0$ . It is easy to show that  $(U^*, V^*)$  is the optimal solution to this problem if and only if  $(U^*/(1U^* + 1V^*), V^*/(1U^* + 1V^*)) = (\alpha^*, \beta^*)$  is the optimal solution to (3.3). Hence, we have

$$\sum_{r=1}^{s} \frac{u_{r}^{*}}{1U^{*} + 1V^{*}} y_{rp} - \sum_{i=1}^{m} \frac{v_{i}^{*}}{1U^{*} + 1V^{*}} x_{ip} \ge 0$$

$$\implies \beta^{*T} Y_{p} - \alpha^{*T} X_{p} = D_{p} \ge 0.$$
(4.12)

This completes the proof.

Hence, to rank DMUs, DMU<sub>j</sub> is better than DMU<sub>i</sub> if and only if  $\widehat{EF}_j > \widehat{EF}_i$ . In what follows, we show that the suggested efficiency index  $\widehat{EF}_j$  is units-invariant. Suppose that  $(X_j, Y_j), j = 1, ..., n$ , is substituted by  $(\lambda X_j, \lambda Y_j) = DMU_{\hat{j}}$ , where  $\lambda > 0$ . It suffices to show

Method	$DMU1_1$	$DMU1_2$	DMU1 <sub>3</sub>
CCR	1.00	1.00	1.00
Proposed method	0.788	1.50	1.46
AP method	1.70	Infeasible	33.30

Table 4.2. The results for the simple example.

Table 5.1. The data set for Example 5.1.

	DMU1	DMU2	DMU3	DMU4	DMU5	DMU6
$I_1$	13	6	2	1	9	4
$I_2$	1	3	6	10	5	8
Ο	1	1	1	1	1	1

that  $\widehat{EF}_{\hat{p}} = \widehat{EF}_p$ . If  $(\alpha^*, \beta^*)$  be the optimal solution to (4.1) for evaluation DMU<sub>p</sub>, it also is the optimal solution to (4.1) for evaluation DMU<sub> $\hat{p}$ </sub>. Hence,  $D_{\hat{p}} = \lambda D_p$ . On the other hand,

$$EF_{\hat{p}} = \frac{D_{\hat{p}}}{\sqrt{\alpha^{* T} \alpha^{*} + \beta^{* T} \beta^{*}}} = \frac{\lambda D_{p}}{\sqrt{\alpha^{* T} \alpha^{*} + \beta^{* T} \beta^{*}}},$$
  

$$\widehat{EF}_{\hat{p}} = \frac{EF_{\hat{p}}}{\operatorname{Max}_{1 \le j \le n} \{EF_{\hat{j}}\}} = \widehat{EF_{p}}.$$
(4.13)

In this method, a small variation in data leads to a small shift on the frontier. Also, the proposed model is always feasible and bounded. These guarantee the feasibility and stability of the method. Consider the presented example in Section 3, again. Using the proposed method, we have the results presented in Table 4.2.

#### 5. Illustrative examples

In order to provide a numerical illustration of the proposed approach, three examples are given.

*Example 5.1.* Consider the data set in Table 5.1 consisting of six DMUs each consuming two inputs to produce one output.

The results of the proposed approach are summarized in Table 5.2.

*Example 5.2.* Another example consists of a data set on 15 units with four inputs and three outputs. Table 5.3 contains a listing of the original data.

Running the proposed model yields the results listed in Table 5.4.

*Example 5.3.* In order to provide an application of the proposed ranking approach, the Sherman and Gold's [9] data set on 14 bank branches is used. Sherman and Gold utilized

$DMU_j$	$\widehat{EF}_j$	Classification	Proposed rank	AP rank	MAJ rank
1	1.000	Efficient	1	1	1
2	0.339	Efficient	4	4	4
3	0.454	Efficient	3	3	2
4	0.726	Efficient	2	2	3
5	-1.122	Inefficient	6	6	6
6	-0.924	Inefficient	5	5	5

Table 5.2. The results for Example 5.1.

DMUs	$I_1$	$I_2$	$I_3$	$I_4$	$O_1$	$O_2$	<i>O</i> <sub>3</sub>
01	63.46	21.39	0	0	42.23	0	4063.68
02	0	0	0	229.24	13.89	0	11250.82
03	0	22.86	15.54	21.59	55.56	0	0
04	199.62	5.47	0	52.29	56.17	0	9829.4
05	302.82	15.52	0	0	92.59	0	0
06	109.1	0	20.47	0	36.29	0	2410.96
07	0	0	56.21	14.03	0	34.13	0
08	6.38	11.52	43.03	0	15.14	13.04	10574.70
09	33.61	28.34	0	75.24	69.46	0	4971.83
10	167.28	0	55.79	0	0	29.87	12458.02
11	99.64	9.77	6.76	0	0	19.89	10099.37
12	0	23.71	0	0	0	27.78	0
13	194.02	8.24	0	0	46.17	0	6506.39
14	0	6.42	119.09	0	0	74.07	0
15	0	10.83	7.31	10.23	26.32	0	0

DEA to evaluate the technical efficiency of branch operation. The data set consists of three inputs and four outputs over 14 branches. The input and output categories are as follows:

Input 1: rent (thousands of dollars),

Input 2: full-time equivalent personnel,

Input 3: supplies (thousands of dollars),

Output 1: loan applications, new pass-book loans, life-insurance sales,

Output 2: new accounts, closed accounts,

Output 3: travelers checks sold, bonds sold, bonds redeemed,

Output 4: deposits, withdrawals, checks sold, treasury checks issued, checks, loan payments, pass-book loan payments, life insurance payments, mortgage payments.

Table 5.5 contains a listing of the original data.

DMU <sub>j</sub>	$\widehat{EF}_j$	Classification	Proposed rank	AP rank	MAJ rank
01	15.45	Efficient	6	4	6
02	41.95	Efficient	2	Infeasible	3
03	0.004	Efficient	15	10	15
04	5.69	Efficient	12	7	12
05	8.29	Efficient	9	8	10
06	6.39	Efficient	11	Infeasible	9
07	2.95	Efficient	13	Infeasible	13
08	88.19	Efficient	1	1	1
09	11.44	Efficient	7	5	7
10	10.674	Efficient	8	Infeasible	8
11	22.12	Efficient	5	6	5
12	23.39	Efficient	3	Infeasible	2
13	6.72	Efficient	10	3	11
14	23.06	Efficient	4	2	4
15	0.044	Efficient	14	9	14

Table 5.4. The results for Example 5.2 using different algorithms.

Table 5.5. The raw data for Example 5.3.

$\begin{array}{ c c c c c c c c c c c c c c c c c c c$								
0248,80017,40037,900384,0001,685,500139,7803,336,8600336,60014,20029,800209,0001,058,90065,7203,570,0500447,1009,30026,800157,000879,40027,3402,081,3500532,6004,60019,60046,000370,90018,9201,069,1000650,8008,30018,900272,000667,40034,7502,660,0400740,8007,50020,40053,000465,70020,2401,800,2500831,9009,20021,400250,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	DMUs	$I_1$	$I_2$	$I_3$	$O_1$	$O_2$	$O_3$	$O_4$
0336,60014,20029,800209,0001,058,90065,7203,570,0500447,1009,30026,800157,000879,40027,3402,081,3500532,6004,60019,60046,000370,90018,9201,069,1000650,8008,30018,900272,000667,40034,7502,660,0400740,8007,50020,40053,000465,70020,2401,800,2500831,9009,20021,400250,000642,70043,2802,296,7400936,40076,00021,000407,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	01	140,000	42,900	87,500	484,000	4,139,100	59,860	2,951,430
0447,1009,30026,800157,000879,40027,3402,081,3500532,6004,60019,60046,000370,90018,9201,069,1000650,8008,30018,900272,000667,40034,7502,660,0400740,8007,50020,40053,000465,70020,2401,800,2500831,9009,20021,400250,000642,70043,2802,296,7400936,40076,00021,000407,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	02	48,800	17,400	37,900	384,000	1,685,500	139,780	3,336,860
0532,6004,60019,60046,000370,90018,9201,069,1000650,8008,30018,900272,000667,40034,7502,660,0400740,8007,50020,40053,000465,70020,2401,800,2500831,9009,20021,400250,000642,70043,2802,296,7400936,40076,00021,000407,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	03	36,600	14,200	29,800	209,000	1,058,900	65,720	3,570,050
0650,8008,30018,900272,000667,40034,7502,660,0400740,8007,50020,40053,000465,70020,2401,800,2500831,9009,20021,400250,000642,70043,2802,296,7400936,40076,00021,000407,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	04	47,100	9,300	26,800	157,000	879,400	27,340	2,081,350
0740,8007,50020,40053,000465,70020,2401,800,2500831,9009,20021,400250,000642,70043,2802,296,7400936,40076,00021,000407,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	05	32,600	4,600	19,600	46,000	370,900	18,920	1,069,100
0831,9009,20021,400250,000642,70043,2802,296,7400936,40076,00021,000407,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	06	50,800	8,300	18,900	272,000	667,400	34,750	2,660,040
0936,40076,00021,000407,000647,70032,3601,981,9301025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	07	40,800	7,500	20,400	53,000	465,700	20,240	1,800,250
1025,7007,90019,00072,000402,50019,9302,284,9101144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	08	31,900	9,200	21,400	250,000	642,700	43,280	2,296,740
1144,5008,70021,700105,000482,40049,3202,245,1601242,3008,90025,80094,000511,00026,9502,303,0001340,6005,50019,40084,000287,40034,9401,141,750	09	36,400	76,000	21,000	407,000	647,700	32,360	1,981,930
12       42,300       8,900       25,800       94,000       511,000       26,950       2,303,000         13       40,600       5,500       19,400       84,000       287,400       34,940       1,141,750	10	25,700	7,900	19,000	72,000	402,500	19,930	2,284,910
13 40,600 5,500 19,400 84,000 287,400 34,940 1,141,750	11	44,500	8,700	21,700	105,000	482,400	49,320	2,245,160
	12	42,300	8,900	25,800	94,000	511,000	26,950	2,303,000
14 76,100 11,900 32,800 199,000 694,600 67,160 3,338,390	13	40,600	5,500	19,400	84,000	287,400	34,940	1,141,750
	14	76,100	11,900	32,800	199,000	694,600	67,160	3,338,390

Running AP, MAJ, and proposed models yield a ranking pattern which is reported in Table 5.6. Note that DMU1, DMU2, DMU3, DMU4, DMU6, DMU9, DMU10 and DMU14 are efficient in  $T_c$ .

$DMU_j$	$\widehat{EF}_j$	Classification	Proposed rank	AP rank	MAJ rank
01	0.284	Efficient	4	6	4
02	1.000	Efficient	1	2	1
03	0.411	Efficient	3	4	3
04	0.007	Efficient	8	8	8
05	-0.02	Inefficient	11	12	11
06	0.263	Efficient	5	3	5
07	-0.105	Inefficient	14	14	14
08	-0.011	Inefficient	9	9	9
09	0.737	Efficient	2	1	2
10	0.046	Efficient	6	5	6
11	-0.016	Inefficient	10	10	10
12	-0.105	Inefficient	13	13	13
13	-0.026	Inefficient	12	11	12
14	0.011	Efficient	7	7	7

Table 5.6. The results for Example 5.3 using different algorithms.

# 6. Conclusion

In multidimensional cases, or if the number of observations is relatively small as compared to the sum of number of inputs and outputs, running the DEA model will typically result in a relatively large number of efficient units. Often, decision makers are interested in a complete ranking, beyond the dichotomized classification, in order to refine the evaluation of the units. This paper provides an analytical procedure for ranking efficient units. In this method, the difficulties of the existing methods have been removed. The proposed approach can be analytically separated in the following two computational processes: (i) the first stage omits the unit being scored from PPS and determines the new frontier and (ii) the second stage finds the shortest distance from the unit being scored to new frontier.

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