

### Example 1

A Health Agency has to decide about opening new ambulatories to serve a set of  $m$  users ( $i = 1, \dots, m$ ). A preliminary investigation suggests a set of  $n$  potential locations. Each users can be served by an ambulatory on location  $j$  ( $j = 1, \dots, n$ ) if the location is not “too far” from the user’s home. Thus, we are given a matrix  $A$  whose entry  $a_{ij}$  has value 1 if user  $i$  can be served by an ambulatory open in location  $j$  and 0 otherwise. The cost of opening an ambulatory on location  $j$  is equal to  $c_j$ . Moreover, the ambulatory in location  $j$  can serve at most  $M_j$  users. Finally, for planning reasons, the Health Agency wants to open at least  $P$  ambulatories, serve all users and minimize the overall cost.

Integer Linear Programming Model:

$$x_{ij} := \begin{cases} 1 & \text{if user } i \text{ is served by ambulatory } j \\ 0 & \text{otherwise} \end{cases} \quad i = 1, \dots, m; j = 1, \dots, n;$$

$$y_j := \begin{cases} 1 & \text{if ambulatory } j \text{ is open} \\ 0 & \text{otherwise} \end{cases} \quad j = 1, \dots, n;$$

$$\min \sum_{j=1}^n c_j y_j \quad (1a)$$

$$\sum_{j=1}^n a_{ij} x_{ij} = 1 \quad \forall i = 1, \dots, m \quad (1b)$$

$$\sum_{i=1}^m x_{ij} \leq M_j y_j \quad \forall j = 1, \dots, n \quad (1c)$$

$$\sum_{j=1}^n y_j \geq P \quad (1d)$$

$$y_j \in \{0, 1\} \quad \forall j = 1, \dots, n \quad (1e)$$

$$x_{ij} \in \{0, 1\} \quad \forall i = 1, \dots, m; \quad \forall j = 1, \dots, n \quad (1f)$$

## Example 2

The Ministry of Home Security must decide concerning the opening of new Fire Stations so as to serve  $q$  neighborhoods ( $i = 1, \dots, q$ ) in Bologna. A preliminary investigation led to the consider a set of  $s$  potential areas  $j, j = 1, \dots, s$  where a Fire Station can be built. Each neighborhood  $i$  can be served by a Fire Station on area  $j$  if that is not “too far”. Then, we are given a binary matrix  $A$  whose entry  $a_{ij}$  is 1 if neighborhood  $i$  can be served by a Fire Station on area  $j$  and 0 otherwise. The opening cost for the Fire Station in area  $j$  is  $r_j$ . For security reasons, each neighborhood must be potentially served by at least 2 (open) Fire Stations, and at least  $B$  Fire Stations must be open. The constraints must be satisfied by minimizing the overall opening cost.

Integer Linear Programming Model:

$$y_j := \begin{cases} 1 & \text{if Fire Station } j \text{ is open} \\ 0 & \text{otherwise} \end{cases} \quad j = 1, \dots, s;$$

$$\min \sum_{j=1}^s r_j y_j \quad (2a)$$

$$\sum_{j=1}^s a_{ij} y_j \geq 2 \quad \forall i = 1, \dots, q \quad (2b)$$

$$\sum_{j=1}^s y_j \geq B \quad (2c)$$

$$y_j \in \{0, 1\} \quad \forall j = 1, \dots, s \quad (2d)$$

$$(2e)$$

### Example 3

A pet shop has  $n$  water pools to be used to show a set of  $m$  tropical fishes. Each pool  $j$  has a cost of  $c_j$  and a volume  $V_j$ . The shop owner wants to minimize the overall cost of the pools used by considering that, for space reasons, the shop can use at most  $K$  pools simultaneously. It is given a graph  $G = (V, E)$  whose generic edge  $a_{ih}$  defines the incompatibility between fishes  $i$  and  $j$ , respectively, of being put in the same pool (because of the natural nutrition chain). Moreover, we are given a binary matrix  $B$  whose entry  $b_{ij}$  is equal to 1 if fish  $i$  can be assigned to pool  $j$  and 0 otherwise. Finally, each fish  $i$  needs a living space of  $v_i$ , thus the sum of the living space of the fishes assigned to the same pool  $j$  cannot be higher than the volume of the volume  $V_j$ .

Integer Linear Programming Model:

$$y_j := \begin{cases} 1 & \text{if pool } j \text{ is used} \\ 0 & \text{otherwise} \end{cases} \quad j = 1, \dots, n;$$

$$x_{ij} := \begin{cases} 1 & \text{if fish } i \text{ is assigned to pool } j \\ 0 & \text{otherwise} \end{cases} \quad i = 1, \dots, m, \quad j = 1, \dots, n;$$

$$\min \sum_{j=1}^n c_j y_j \quad (3a)$$

$$\sum_{j=1}^n y_j \leq K \quad (3b)$$

$$\sum_{j=1}^n b_{ij} x_{ij} = 1 \quad \forall i = 1, \dots, m \quad (3c)$$

$$\sum_{i=1}^m v_i x_{ij} \leq V_j y_j \quad \forall j = 1, \dots, n \quad (3d)$$

$$x_{ij} + x_{hj} \leq y_j \quad \forall (i, h) \in A, j = 1, \dots, n \quad (3e)$$

$$x_{ij}, y_j \in \{0, 1\} \quad \forall i = 1, \dots, m, j = 1, \dots, n \quad (3f)$$

Discuss alternatives and/or possible improvements to the above model.