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# Differential Cluster Analysis

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September 7, 2009

CHES Workshop 2009, Lausanne, Switzerland



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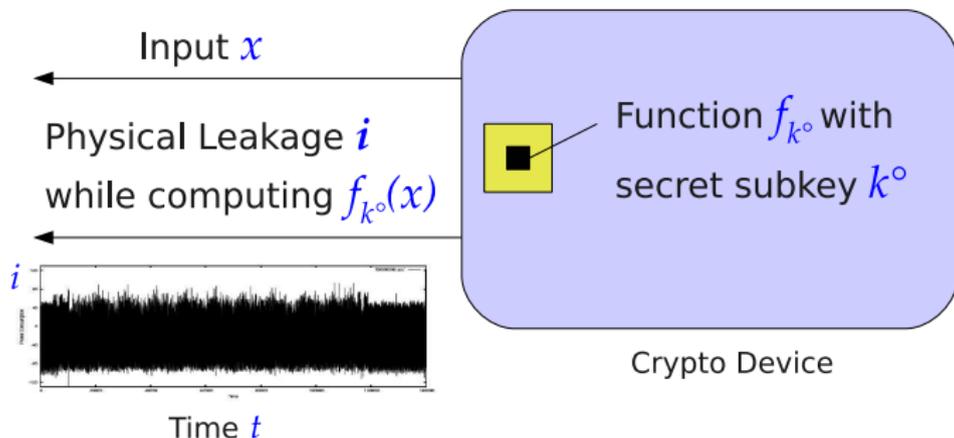


## Presentation Outline

- Introduction
- Differential Cluster Analysis
- Applications
- Experimental Results
- Conclusion

# Introduction: Adversary Model

## Differential Side Channel Adversary



## Adversary Success

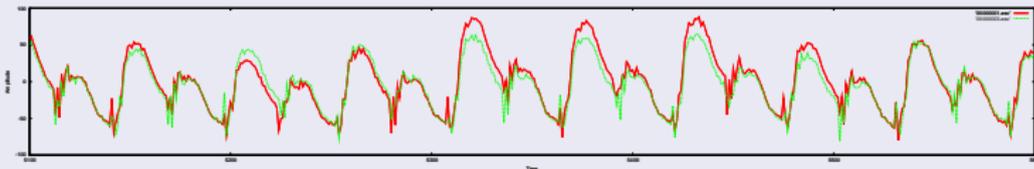
The adversary is successful if she recovers the secret subkey  $k^o$ .

# Introduction: Differential Power Analysis in a Nutshell

## The DPA (Differential Power Analysis) Problem

Given measurements  $i_n$  ( $n \in \{1, \dots, N\}$ ) while the crypto device computes function  $f_{k_0}$  with random inputs  $x_n$ :

- Does a statistics prove significant differences of the measurements according to a partitioning function  $g(f_k(x_n))$ ?



## DPA Partitioning Functions $g$ :

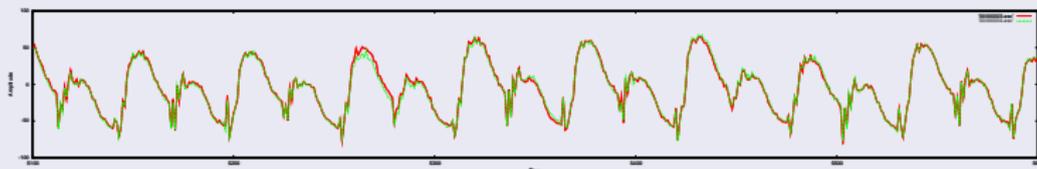
- Single-bit partitioning function (Kocher *et al.*, 1999)
- Multi-bit partitioning/comparison function:
  - “All-or-Nothing” with a leakage model (e.g., Hamming weight) (Messerges *et al.*, 2000)
  - CPA (Correlation Power Analysis) with a leakage model (e.g., Hamming distance) (Brier *et al.*, 2002)

# Introduction: Collision Analysis in a Nutshell

## The CA (Collision Analysis) Problem

Given measurements  $i_n$  ( $n \in \{1, \dots, N\}$ ) while the crypto device computes the many-to-one function  $f_{k^0}$  with random inputs  $x_n$ :

- Does a statistics prove high similarity of measurements with two inputs  $x_i \neq x_j$ ?



## Collision Detection:

- Euclidean distance of measurement vectors over some  $t$ . (Schramm *et al.*, 2003)

# Our Approach

## Objectives

- Combination of leakage detection functions for DPA (*Separation*) and CA (*Cohesion*).
- Sensitivity to general leakage features
- Multi-bit approach
  - using all measurements and
  - without the need for a good power model.
- Multivariate approach

## Idea

- Our basic approach: Cluster Analysis
- Our basic question: Do clusters of measurements exist?

# Differential Cluster Analysis

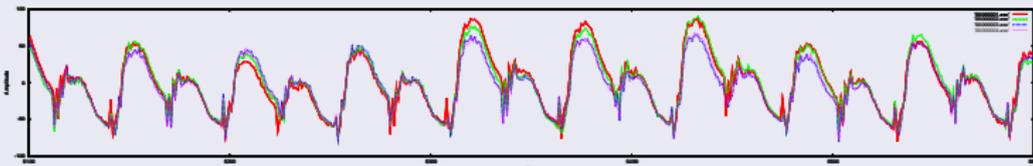
## Requirement

- $f_k : \{0, 1\}^u \mapsto \{0, 1\}^w$  is a many-to-one collision function, i.e., at least two inputs  $x_i, x_{i'} \in \{0, 1\}^u$  with  $x_i \neq x_{i'}$  collide in one state  $\Delta \in \{0, 1\}^w$ .

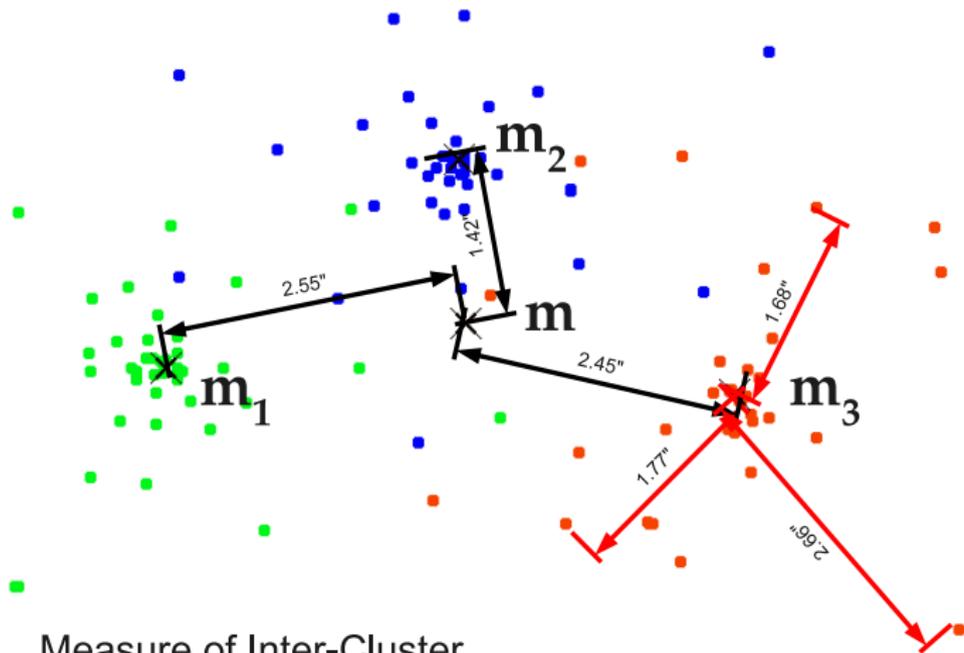
## The DCA (Differential Cluster Analysis) Problem

Given measurements  $i_n$  ( $n \in \{1, \dots, N\}$ ) while the crypto device computes the many-to-one function  $f_{k^0}$  with random inputs  $x_n$ :

- Does a cluster criterion function prove the existence of clusters of measurements according to partitioning function  $f_k(x_n)$ ?



# Measuring Clustering Quality



Measure of Inter-Cluster Separation

Measure of Intra-Cluster Cohesion

# Cluster Criterion Functions

## Sum-of-Squared-Error:

$$J_{SSE} = \sum_{i=1}^c \sum_{\mathbf{x} \in \mathcal{D}_i} \|\mathbf{x} - \mathbf{m}_i\|^2$$

$J_{SSE}$  evaluates intra-cluster cohesion. The optimal partition minimizes  $J_{SSE}$ .

## Sum-of-Squares:

$$J_{SOS} = \sum_{i=1}^c n_i \|\mathbf{m}_i - \mathbf{m}\|^2$$

$J_{SOS}$  evaluates inter-cluster separation. The optimal partition maximizes  $J_{SOS}$ .

The sum of  $J_{SSE}$  and  $J_{SOS}$  is a constant.

**Variance criterion** (Standaert *et al.*: ICISC 2008):

$$J_{VAR} = \frac{\|\mathbf{v}\|^2}{\frac{1}{N} \sum_{i=1}^c n_i \|\mathbf{v}_i\|^2}$$

$J_{VAR}$  evaluates overall variance vs. intra cluster variances. The optimal partition maximizes  $J_{VAR}$ .

**T-test criterion** (Gierlichs *et al.*: CHES 2006):

$$J_{STT} = \sum_{i,j=1;i \neq j}^c \frac{\|\mathbf{m}_i - \mathbf{m}_j\|^2}{\sqrt{\frac{\|\mathbf{v}_i\|^2}{n_i} + \frac{\|\mathbf{v}_j\|^2}{n_j}}}$$

$J_{STT}$  evaluates inter cluster separation, normalized by intra cluster variances and cluster sizes. The optimal partition maximizes  $J_{STT}$ .

# Differential Cluster Analysis (General Approach)

## Differential Cluster Analysis (General Approach)

- 1 For each subkey hypothesis  $k$ :
  - Sort measurements into  $2^w$  clusters  $\mathcal{D}_0, \dots, \mathcal{D}_{2^w-1}$  according to  $\Delta_i = f_k(x_n)$  ( $1 \leq n \leq N$ ).
  - Compute a cluster criterion function:  $J_k$ .
- 2 Rank the pairs  $(k, J_k)$  according to  $J_k$ .
- 3 Output subkey candidate that leads to the best clustering quality.

# Detailed Comparison with CA and DPA

	DPA	CA	DCA
Many-to-one function	not required	required	required
Leakage model	<ul style="list-style-type: none"><li>• none for single-bit DPA</li><li>• required for multi-bit DPA</li></ul>	none	not required, can be integrated
Statistics	based on differences	based on similarity	based on both separation and cohesion
Detected Leakage Features	<ul style="list-style-type: none"><li>• differences of two states for single-bit and "all-or-nothing" multi-bit DPA</li><li>• linearity of differences for CPA</li></ul>	general features	general features
Multivariate Leakage	original approach can be extended	yes	yes

# Comparison with DPA: An Example

## Example

Assume  $f_k : \{0, 1\}^u \mapsto \{0, 1\}^2$  is a many-to-one function.

- Single-bit DPA fails and
- Multi-bit DPA (with Hamming weight model) fails.

Does this assure that there is no leakage at all?

# Comparison with DPA: An Example

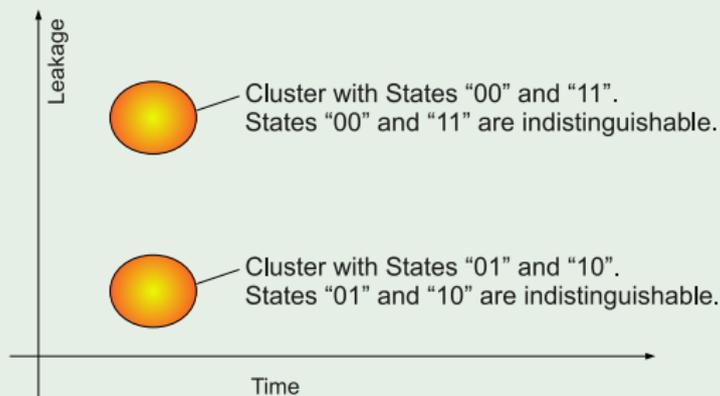
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- Single-bit DPA fails and
- Multi-bit DPA (with Hamming weight model) fails.

Does this assure that there is no leakage at all?

**No. Counter-Example:**



# Applications: Algorithmic Collisions

## Algorithmic Collisions:

- $f_k$  emerges from an abstract concept (e.g., cryptographic standard).
- $f_k$  is implementation independent (despite of masking ...).

### Example: DES

- DES S-box function is 4-to-1:  $f_k : \{0, 1\}^6 \mapsto \{0, 1\}^4$  yields  $2^4$  clusters.

### Example: AES

- AES S-box is not a collision function.
- Targeting only  $r$ -bit ( $1 \leq r < 8$ ) of AES S-box outcome:  
 $f_k : \{0, 1\}^8 \mapsto \{0, 1\}^r$  yields  $2^r$  clusters.
- AES MixColumns transformation is  $2^{24} - to - 1$ :  
 $f_k : \{0, 1\}^{32} \mapsto \{0, 1\}^8$  yields  $2^8$  clusters.

# Applications: Implementation Specific Collisions

Implementation specific collisions:

- $f_k$  emerges from implementation properties.
- $f_k$  is not obvious in the algorithmic description.

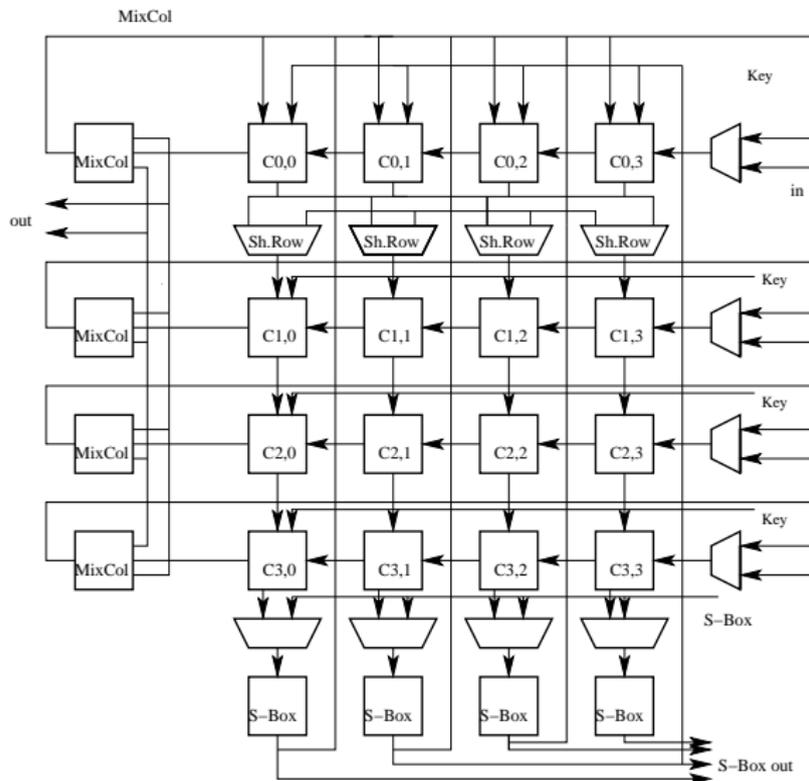
## Example: AES Hardware Module

Differential of two adjacent data cells in the studied AES hardware architecture:  $f_k : \{0, 1\}^{16} \mapsto \{0, 1\}^8$  yields  $2^8$  clusters.

$$f_k(x) = S(x_i \oplus k_i) \oplus S(x_{i'} \oplus k_{i'}) \quad (1)$$

General DCA Approach requires  $2^{16}$  subkey hypotheses.

# Application: AES Hardware Module



# AES Hardware Module: A New Attack Strategy

If  $f_k(x) = 0$  then  $f_k(x) = S(x_i \oplus k_i) \oplus S(x_{i'} \oplus k_{i'})$  simplifies to

$$S(x_i \oplus k_i) = S(x_{i'} \oplus k_{i'}) \Rightarrow x_i \oplus k_i = x_{i'} \oplus k_{i'} \Rightarrow k_i \oplus k_{i'} = x_i \oplus x_{i'}.$$

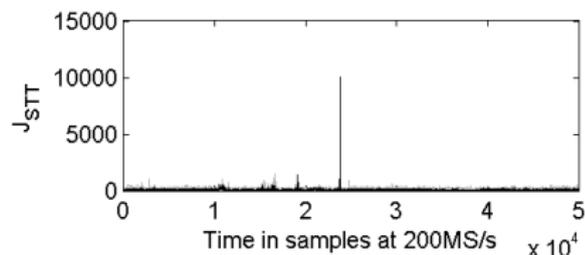
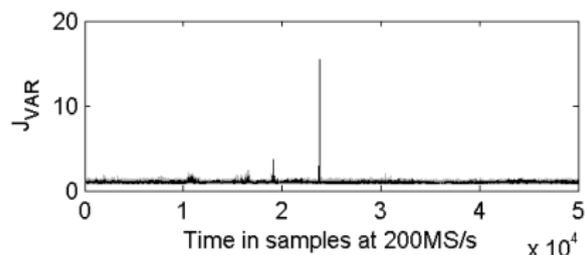
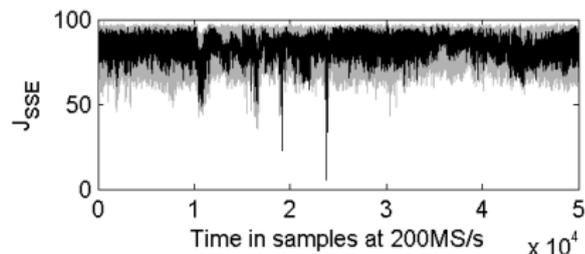
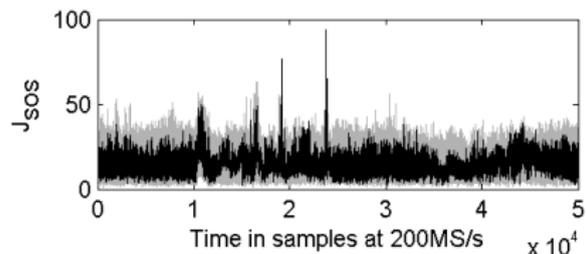
The elements of one cluster are identical if  $k_i \oplus k_{i'} = k_i^{\circ} \oplus k_{i'}^{\circ}$ .

A new two-step key recovery attack:

- 1 Determine the correct xor-difference  $k_i^{\circ} \oplus k_{i'}^{\circ}$  based on  $2^8$  hypotheses.
  - This is the difficult step that checks whether a special (small) cluster for  $f_k(x) = 0$  exists.
- 2 Determine the correct pair  $(k_i^{\circ}, k_{i'}^{\circ})$  based on  $2^8$  hypotheses.
  - This is the easy step that checks whether up to  $2^8$  clusters exist.

- Attack strategy can also be applied with DPA.

## Target Device: AVR Microcontroller



# DES Implementation in Software: Comparison with CPA

Table: Success rates in % for various univariate and multivariate attack scenarios.

Test	Model	Time	N=15	N=20	N=25	N=30	N=40	N=50
CPA	LSB	overall	3	15	37	62	95	98
CPA	LSB	A	42	64	69	77	93	96
CPA	LSB	B	64	77	83	93	98	99
CPA	LSB	C	17	28	29	38	55	65
$J_{SSE}$	LSB	overall	3	15	37	62	95	98
$J_{SSE}$	LSB	A	42	64	70	77	93	96
$J_{SSE}$	LSB	B	64	78	82	93	98	99
$J_{SSE}$	LSB	C	18	28	31	38	56	65
CPA	LSB	AB	70	85	90	96	100	100
$J_{SSE}$	LSB	AB	70	83	91	97	100	100
CPA	LSB	ABC	76	90	96	99	100	100
$J_{SSE}$	LSB	ABC	78	94	96	99	100	100

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**Target Device:** Prototype chip which implements an AES-128 co-processor in  $0.13 \mu\text{m}$  sCMOS technology without countermeasures.

## General Approach with DCA:

- $2^{16}$  key hypotheses
- key recovery with approx. 5000 measurements

## Two-Step Attack Strategy with DCA:

- $2^9$  key hypotheses
- key recovery with approx. 50 000 measurements
  - Step 1: 50 000 measurements
  - Step 2: 5000 measurements

# Conclusion

- Introduction of Differential Cluster Analysis (DCA) – a new technique bridging the gap between Collision Analysis (CA) and Differential Power Analysis (DPA).
- Introduction of implementation specific collisions.
- Confirmation of DCA on both software (DES) and hardware (AES) implementation.
- New two-step attack strategy for an AES hardware module.

*Collision attacks on AES are not constrained to 8-bit software implementations on simple controllers anymore...*