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## Blocking Visibility for Points in General Position

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**Abstract** For a finite set  $P$  in the plane, let  $b(P)$  be the smallest possible size of a set  $Q$ ,  $Q \cap P = \emptyset$ , such that every segment with both endpoints in  $P$  contains at least one point of  $Q$ . We raise the problem of estimating  $b(n)$ , the minimum of  $b(P)$  over all  $n$ -point sets  $P$  with no three points collinear. We review results providing bounds on  $b(n)$  and mention some additional observations.

**Keywords** Visibility · Visibility-blocking set · Behrend's construction

Let  $P$  be an  $n$ -point set in the plane (or, more generally, in  $\mathbf{R}^d$ ). We define a *visibility-blocking set for  $P$*  as a set  $Q$  that is *disjoint* from  $P$  and such that every segment with endpoints in  $P$  contains at least one point of  $Q$ .

If the points of  $P$  are all collinear, then there is a visibility blocking set with  $n - 1$  points. The question raised in this note is what is the smallest possible size of a visibility-blocking set for  $P$  having no three points collinear? That is, we let

$$b(P) := \min\{|Q| : Q \text{ a visibility-blocking set for } P\},$$
$$b(n) := \min\{b(P) : P \subset \mathbf{R}^2 \text{ with no three points collinear, } |P| = n\},$$

and we would like to estimate the asymptotics of  $b(n)$  for large  $n$ .

I arrived at this question when (unsuccessfully) attempting to prove the following nice conjecture of Kára, Pór, and Wood [7, Conjecture 2]: *For all integers  $k, \ell \geq 2$ ,*

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there is an integer  $n$  such that every  $n$ -point set in the plane contains  $\ell$  collinear points or  $k$  pairwise visible points.

I obtained Theorems 1 and 3 below, but after this note was accepted for publication in the Klee Festschrift, it turned out that both of these results had been known earlier in somewhat different contexts. Moreover, I learned that other people considered the asymptotics of  $b(n)$  independently (a group of researchers at the Courant Institute in New York including Andreas Holmsen, János Pach, Radoš Radoičić, and Gábor Tardos [6]) and also rediscovered some of these results. Thus, the present note mostly reviews known results and adds some observations, which have not appeared in print as far as I know and which illustrate some of the difficulties inherent in the problem.

I still believe that the inclusion of this note in the Festschrift is warranted by the beauty of the problem—I think that Vic Klee would like it.

*Midpoints and an Upper Bound* For a point set  $P$ , let  $\mu(P)$  be the cardinality of the set  $\{\frac{1}{2}(p + q) : p, q \in P, p \neq q\}$  of *midpoints* of all pairs of points of  $P$ . Clearly,  $b(P) \leq \mu(P)$ . The problem of estimating  $\mu(n) := \min \mu(P)$ , where the minimum is over all  $n$ -point planar sets  $P$  in general position, was raised, according to Pach [10], by F. Hurtado. Earlier Erdős, Fishburn, and Füredi [4] studied the problem of estimating  $\min \mu(P)$  over all  $n$ -point planar sets  $P$  in *convex position* and established the (surprising) lower and upper bounds of  $0.40n^2$  and  $0.45n^2$ , respectively.

Now we recall an upper bound on  $\mu(n)$  (and  $b(n)$ ) due to Pach [10] (based on Erdős et al. [5]). For an integer  $N$ , let  $\nu(N)$  be the maximum number of elements of a set  $A \subseteq \{1, 2, \dots, N\}$  that contains no three-term arithmetic progression. Behrend [3], improving on a construction by Salem and Spencer, proved that  $\nu(N) \geq N^{1-O(1/\sqrt{\log N})}$ . The following theorem, whose proof we recall, is based on Behrend’s construction.

**Theorem 1** (Pach [10]) *There is a constant  $C$  such that*

$$b(n) \leq \mu(n) \leq ne^{C\sqrt{\log n}}$$

for all sufficiently large  $n$ .

*Proof* Let  $n$  be given and large, and let  $m$  and  $s$  be integer parameters to be specified later. Let  $G := \{0, 1, \dots, s - 1\}^m \subset \mathbf{R}^m$ , and let  $S_k := \{x \in G : \|x\|^2 = k\}$ , where  $\|\cdot\|$  is the Euclidean norm.

As in Behrend’s argument,  $G = \bigcup_{k=0}^{m(s-1)^2} S_k$ , and thus by the pigeonhole principle there exists  $k$  with  $|S_k| \geq s^{m-2}/m$ . We let  $P := S_k$  for some such  $k$ .

Since the points of  $P$  lie on a sphere, no three of them are collinear. The midpoint of every two points  $p, q \in G$  lies in  $G' := \{0, \frac{1}{2}, 1, \frac{3}{2}, \dots, s - 1\}^m$ , and thus  $Q := G' \setminus P$  is a visibility-blocking set for  $P$ .

We let  $m := \lfloor \ln n \rfloor$ , and let  $s$  be the smallest integer with  $s^{m-2}/m \geq n$ . Then  $|P| \geq n$ ,  $(s - 1)^{m-2} < mn$ , and

$$|Q| \leq |G'| = (2s - 1)^m \leq (3(s - 1))^m \leq 3^m (s - 1)^2 mn \leq ne^{O(\sqrt{\log n})},$$

as can easily be calculated (using  $(s - 1)^2 \leq (mn)^{2/(m-2)} \leq (n^2)^{4/m}$ , say). The implicit constant in the exponent can be improved by a more careful choice of  $m$  and by more precise calculations.

The bound in the theorem follows by projecting  $P$  and  $Q$  to a generic two-dimensional subspace of  $\mathbf{R}^m$ . □

*Lower Bounds* I am aware only of the following rather trivial lower bound for  $b(n)$ :

**Observation 2** *If  $P$  is an  $n$ -point planar set with no three points collinear and with  $\text{conv}(P)$  having  $p$  vertices, then  $b(P) \geq 3n - p - 3$ . In particular,  $b(n) \geq 2n - 3$ .*

*Proof* A triangulation of  $P$  has  $3n - p - 3$  edges, and each point of a visibility-blocking set covers at most one of these. □

For point sets in convex position, a slightly superlinear lower bound can be given. The result is implicitly contained in Araujo et al. [1], and the argument goes back (at least) to Kostochka and Kratochvíl [8].

**Theorem 3** *For every  $n$ -point planar set  $P$  in (strictly) convex position, we have*

$$b(P) \geq \begin{cases} n \sum_{k=1}^m 1/k & \text{for } n = 2m + 1 \text{ odd,} \\ 1 + n \sum_{k=1}^{m-1} 1/k & \text{for } n = 2m \text{ even.} \end{cases}$$

Thus,  $b(P) = \Omega(n \log n)$ .

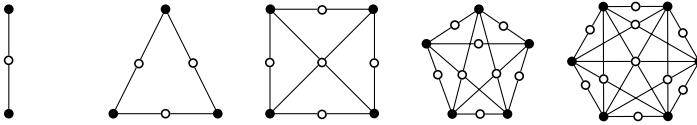
*Proof* Let  $p_1, p_2, \dots, p_n$  be the points of  $P$  numbered along the circumference of  $\text{conv}(P)$ . We define the length  $\ell(p_i p_j)$  of the segment  $p_i p_j$ ,  $i < j$ , as the number of convex hull edges between  $p_i$  and  $p_j$ , i.e.,  $\min(j - i, n + i - j)$ .

Let us suppose that  $Q$  is a visibility-blocking set for  $P$ , and let  $q \in Q$ . The key observation is that if  $\ell$  is the smallest of the lengths of the segments  $p_i p_j$  incident to  $q$ , then  $q$  is incident to at most  $\ell$  segments.

Thus, if we give the segment  $p_i p_j$  weight  $1/\ell(p_i p_j)$ , no  $q \in Q$  is incident to segments of total weight more than 1. The theorem follows by summing the weights of all segments. □

For the case where  $P$  is the vertex set of a regular convex  $n$ -gon, Poonen and Rubinstein [13] show that, apart from the center, no point is the intersection of eight or more of the diagonals of the  $n$ -gon, and thus  $b(P) = \Omega(n^2)$  in this case.

*Small Cases* Now we return to arbitrary sets (not necessarily in convex position). The first few values of  $b(n)$  are  $b(2) = 1$ ,  $b(3) = 3$ ,  $b(4) = 5$ ,  $b(5) = 8$ ,  $b(6) = 10$ . The upper bounds are witnessed by Fig. 1. The lower bounds for  $n \leq 4$  are trivial (or follow from Observation 2). For  $n = 5, 6$ , we distinguish two cases: If all vertices of  $P$  are on the convex hull, then the lower bound follows from Theorem 3, and otherwise, it is given by Observation 2. Determining the exact value gets more complicated for larger  $n$ , and I am not aware of a reasonably clean argument for any



**Fig. 1** Upper bound examples for small  $n$

$n \geq 7$ . Interesting upper bound constructions for several small cases, to be reported elsewhere, were given by Snoeyink and Speckmann (private communication).

*Additional Remarks*

1. A possibly easier version of the problem deals with pseudosegments instead of straight segments. That is, for a given point set  $P$ , we want to construct an arrangement  $\mathcal{A}$  of pseudolines and a subset  $Q$  of its vertices such that  $P \cap Q = \emptyset$ , each  $p \in P$  is a vertex of  $\mathcal{A}$ , no three points of  $P$  lie on a common pseudoline, and every two points of  $P$  lie on a common pseudoline  $\ell$  and have a point of  $Q$  on the segment of  $\ell$  between them.

The lower bound of Theorem 3 still applies in this setting (with “convex position” interpreted appropriately in terms of the arrangement  $\mathcal{A}$ ), and here it is easy to provide an  $O(n \log n)$  upper bound.

Is there a linear upper bound for some  $P$ , not in “convex position”, in the pseudoline setting?

Another variant of the problem, asked by Pach, is partitioning all the  $\binom{n}{2}$  straight segments defined by the points into a small number of *crossing families*, i.e., families in which every two segments cross. Is a superlinear number of such families always necessary? (Also see Pach et al. [11] for some related blocking-type questions.)

2. Here is another variation of the problem: Instead of requiring that each *segment* determined by two points of  $P$  contains a point of  $Q$ , it now suffices that each *line* determined by two points of  $P$  contains a point of  $Q$ . Perhaps surprisingly, there are examples showing that  $O(n)$  points suffice to block all lines.

One such example can be constructed using integer points on the curve  $y = x^3$  (a similar construction was used, e.g., for the *orchard problem*; see [9]). A simpler example, inspired by a relation of the problem to Ungar’s theorem (see [9] again), was communicated to me by Pinchasi [12]: Take the vertex set of a regular  $(2n)$ -gon centered at the origin and apply a projective transform that sends the line at infinity to the  $x$ -axis, obtaining a  $(2n)$ -point set  $P$  (the points of  $P$  lie on a hyperbola). Since the vertices of the regular  $(2n)$ -gon determine only  $2n$  distinct directions, there is a  $(2n)$ -point set  $Q$  on the  $x$ -axis that intersects all the lines determined by  $P$ .

3. The set  $P$  in the example just mentioned can also be partitioned into two  $n$ -point subsets  $P_1$  and  $P_2$  (namely, the points above and below the  $x$ -axis) in such a way that every segment  $p_1 p_2$ ,  $p_1 \in P_1$ ,  $p_2 \in P_2$ , contains a point of  $Q$ . This shows that a natural bipartite version of the original visibility-blocking problem has a linear upper bound.

Pach [10] proves a superlinear lower bound for  $\mu(n)$ , the minimum cardinality of the set of all midpoints for  $n$  points in general position, using Freiman's theorem on set addition. This argument can be adapted to give a superlinear lower bound for the midpoints in the bipartite setting as well. Indeed, let  $P_1, P_2$  be disjoint  $n$ -point sets with  $P := P_1 \cup P_2$  in general position, and suppose that the set  $P_1 + P_2 = \{p_1 + p_2 : p_1 \in P_1, p_2 \in P_2\}$  has cardinality  $O(n)$ . Then  $P$  has  $\Omega(n^3)$  additive four-tuples, i.e., four-tuples  $(p_1, p_2, p_3, p_4)$  with  $p_1 + p_2 = p_3 + p_4$ , and by the Balog–Szemerédi theorem [2] there is a subset  $P' \subseteq P$  of size  $\Omega(n)$  with  $|P' + P'| = O(|P'|)$ . Then, as in Pach's argument, Freiman's theorem implies that for a sufficiently large  $n$ , the set  $P'$  contains three collinear points, a contradiction. This shows that Pach's lower bound idea for  $\mu(n)$  is not directly applicable to  $b(n)$ .

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