# WEIGHT DISTRIBUTION FORMULA FOR SOME CLASS OF CYCLIC CODES 

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## Abstract

Let $h_{1}(X)$ and $h_{2}(X)$ be different irreducible polynomials such that $h_{1}\left(\alpha^{-2^{h}-1}\right)=0$ for some $h(0<h<m)$ and $h_{2}\left(\alpha^{-1}\right)=0$, $\alpha$ being a primitive element of $\mathrm{GF}\left(2^{\mathrm{m}}\right)$. This paper presents the weight distribution formula of the code of length $2^{m}-1$ generated by $\left(x^{2^{m}-1}-1\right) /\left(h_{1}(X) h_{2}(X)\right)$ for any $m$ and h. Some applications to the cross-correlation problem between two different maximum length sequences are presented.

## 1. Introduction

W. W. Peterson [1] calculated a number of weight distributions for BCH codes of lengths 63 to 1023 and their dual codes by digital computation. He observed that some BCH codes with large $t$ for a given $m(5 \leq m \leq 10)$ have a very simple structure of weight distribution. The result presented here is a theoretical development of his observation.

Let $C$ be a cyclic code of length $2^{m}-1$. The extended code of $C$ is the code with an overall parity check added to $C$ as the first digit. The first symbol in a code vector is numbered 0 , and for $i>1$ the $i-t h$ digit is numbered $\alpha^{i-2}, \alpha$ being a primitive element of $G F\left(2^{m}\right)$. Now for $a(\neq 0)$ and $b \in G F\left(2^{m}\right)$ and for a code vector $v$ of the extended code, permute the symbol in position X to position $\mathrm{aX}+\mathrm{b}$. Then, the resulting vector is denoted by $\pi_{a b}$. W. W. Peterson [1] proved that the extended codes of BCH codes are invariant under doubly transitive group of permutations $\pi=\left\{\pi_{a b} \mid a(\neq 0)\right.$, beGF ( $2^{m}$ ) \}. This paper presents the weight distribution formula for a class of cyclic codes of length $2^{m}-1$ whose extended codes are invariant under $\pi$.

Let $g_{1}(X)$ and $g_{2}(X)$ be different irreducible polynomials such that

$$
\begin{align*}
& g_{1}\left(\alpha^{2^{h}+1}\right)=0 \quad \text { for some } h(0<h<m)  \tag{1}\\
& g_{2}(\alpha)=0 \tag{2}
\end{align*}
$$

The degree of $g_{1}(X)$ is a factor $m^{\prime}$ of $m$ and the degree of $g_{2}(X)$ is $m$. Let

$$
h_{1}(x)=x^{m^{\prime}} g_{1}\left(x^{-1}\right), h_{2}(x)=x^{m} g_{2}\left(x^{-1}\right)
$$

Let $C_{0}$, $C$ and $C^{\prime}$ denote binary cyclic codes with length $2^{m}-1$ generated by $g_{1}(X) g_{2}(X),\left(X^{2^{m}-1}-1\right) /\left(h_{1}(X) h_{2}(X)\right)$ and $\left(x^{2^{m}-1}-1\right) /\left[(X-1) h_{1}(X) h_{2}(X)\right]$ respectively. Then $C$ is the dual code of $C_{0}$ and a subcode of $C^{\prime}$. If $h=1$, then $C_{o}$ is a double error correcting $B C H$ code, and if $m$ is odd and $h=(m-1) / 2$, then $C$ is
a BCH code with the second largest $t$ for given $m$.
In what follows, the weight distribution formula of code $C$ for any $m$ and $h$ will be derived. This problem is closely related to the crosscorrelation problem between two different maximum length sequences.* Some applications to the problem will be presented in section 6 .

## 2. Preliminary Lemmas

Lemma 1: The extended code of $C^{\prime}$ or $C_{0}$ is invariant under $\pi$.
This lemma follows from the definition of $C^{\prime}$ or $C_{o}$ and a general theorem [2]. Let

$$
\begin{equation*}
m=m^{\prime} m^{\prime \prime} . \tag{1}
\end{equation*}
$$

Since $\alpha^{2^{m-h}+1}$ is a root of $g_{1}(X)$, it can be assumed that

$$
\begin{equation*}
2 \mathrm{~h} \leq \mathrm{m} . \tag{2}
\end{equation*}
$$

Since $\left(2^{m^{\prime}}-1\right)\left(2^{h}+1\right)$ is divisible by $\left(2^{m}-1\right)=\left(2^{m^{\prime}}-1\right)\left(2^{m^{\prime}\left(m^{\prime \prime}-1\right)}+\right.$ $\left.2^{m^{0}\left(m^{\prime \prime}-2\right)}+\ldots+1\right)$,

$$
h \geq m^{\prime}\left(m^{\prime \prime}-1\right) .
$$

From (1) and (2),

$$
\begin{equation*}
m^{\prime} m^{\prime \prime} \geq 2 h \geq 2 m^{\prime}\left(m^{\prime \prime}-1\right) \tag{3}
\end{equation*}
$$

Hence,

$$
\mathrm{m}^{\prime \prime}=1 \text { or } 2 .
$$

If $m^{\prime \prime}=2$, then it follows from (3) that

$$
m^{\prime}=h .
$$

That is, there are only two cases:

$$
m^{\prime}=m
$$

and

$$
m^{\prime}=m / 2=h .
$$

[^0]The following well-known lemmas will be used later.
Lemma 2: Let $u(\ell)$ denote the smallest positive integer $u$ such that $2^{u} \equiv 1(\bmod \ell)$. Then, $2^{u^{\prime}} \equiv 1(\bmod \ell)$ if and only if $u^{\prime} \equiv 0(\bmod u(\ell))$. Let $\left(\ell, \ell^{\ell}\right)$ denote the greatest common divisor of $\ell$ and $\ell$ :

Coro1lary 3: Let $u=\left(u_{1}, u_{2}\right)$. Then,

$$
2^{\mathrm{u}}-1=\left(2^{\mathrm{u}_{1}}-1,2^{\mathrm{u}_{2}}-1\right)
$$

Lemma 4: $2^{\mathrm{u}}+1$ (or $2^{\mathrm{u}}-1$ ) is divisible by $2^{\mathrm{u}^{\prime}}+1$, if and only if $u$ is divisible by $u^{\prime}$ and $u / u^{\prime}$ is odd (or even).

$$
\begin{align*}
& \text { Let } c^{0}=(m, h), c=(m, 2 h) \text { and } v=\left(2^{m}-1,2^{h}+1\right) . \text { By Corollary } 3 \\
& 2^{c^{\prime}}-1=\left(2^{m}-1,2^{h}-1\right)  \tag{4}\\
& 2^{c}-1=\left(2^{m}-1,2^{2 h}-1\right) \tag{5}
\end{align*}
$$

Since $\left(2^{h}+1,2^{h}-1\right)=1$,

$$
\left(2^{m}-1,2^{2 h}-1\right)=\left(2^{m}-1,2^{h}+1\right)\left(2^{m}-1,2^{h}-1\right)
$$

Thus,

$$
\left(2^{c}-1\right) / v=2^{c^{\prime}}-1
$$

By definition, $c=c^{\prime}$ or $2 c^{\prime}$. Therefore, we have:
Lemma 5: If $c=c^{\prime}$, then $v=1$. Otherwise,

$$
\begin{equation*}
v=2^{c^{\prime}}+1=2^{c / 2}+1 \tag{6}
\end{equation*}
$$

The next lemma is due to Pless [3].
Lemma 6: Let $a_{j}$ and $b_{j}$ denote the number of code vectors of weight $j$ in a code $A$ and the number of code vectors of weight $j$ in the dual code of $A$ respectively. If $b_{1}=b_{2}=0$, then the following power moment identities hold:

$$
\begin{aligned}
& \Sigma a_{j}=2^{k} \\
& \Sigma j a_{j}=2^{k-1} n \\
& \Sigma j^{2} a_{j}=2^{k-2} n(n+1) \\
& \Sigma j^{3} a_{j}=2^{k-3}\left(n^{3}+3 n^{2}\right)-3!2^{k-3} b_{3} \\
& \Sigma j^{4} a_{j}=2^{k-4}\left(n^{4}+6 n^{3}+3 n^{2}-2 n\right)-4!2^{k-4} n b_{3}+4!2^{k-4} b_{4},
\end{aligned}
$$

where $k$ denotes the number of information digits.
Let $C_{1}$ and $C_{2}$ denote binary cyclic codes with length $2^{m}-1$ generated by $\left(x^{2^{m}-1}-1\right) / h_{1}(X)$ and $\left(x^{2^{m}-1}-1\right) / h_{2}(X)$ respectively. Codes $C_{1}$ and $C_{2}$ are subcodes of $C$ and $C^{\prime}$. If the degree of $g_{1}(X)$ is $m$, then the roots of $h_{1}(X)=X^{m} g_{1}\left(X^{-1}\right)$ :are:

$$
\begin{aligned}
& \alpha^{-\left(2^{h}+1\right)}=\alpha^{2^{m}-2^{h}-2}, \alpha^{-2\left(2^{h}+1\right)}=\alpha^{2^{m}-2^{h+1}-2-1}, \cdots, \\
& \alpha^{-2^{m-h-1}\left(2^{h}+1\right)}=\alpha^{2^{m-1}-2^{m-h-1}-1}, \alpha^{-2^{(m-h)}\left(2^{h}+1\right)}=\alpha^{2^{m}-2^{m-h}-2}, \\
& \cdots, \alpha^{-2^{m-1}}\left(2^{h}+1\right)=\alpha^{2^{m-1}-2^{h-1}-1} .
\end{aligned}
$$

There is no $i\left(0 \leq i<2^{m-1}\right)$ with $h_{1}\left(\alpha^{i}\right)=0$ except for

$$
i_{1}=2^{m-1}-2^{m-h-1}-1
$$

and

$$
i_{1}^{\prime}=2^{m-1}-2^{h-1}-1
$$

By (2) , $i_{1}<i_{1}^{p}$ 。
If $m^{p}=m / 2$, then the roots of $h_{1}(X)$ are:

$$
\begin{aligned}
& \alpha^{-\left(2^{m^{p}}+1\right)}=\alpha^{2^{m}-2^{m^{\prime}}-2}, \alpha^{-2\left(2^{m^{\prime}}+1\right)}=\alpha^{2^{m}-2^{m^{\prime}+1}-2-1}, \ldots \\
& \alpha^{-2^{m^{0}-1}\left(2^{m^{\prime}}+1\right)}=\alpha^{2^{m-1}-2^{m^{\prime}-1}-1} .
\end{aligned}
$$

there is no $i\left(0 \leq i<2^{m-1}\right)$ with $h_{1}\left(\alpha^{i}\right)=0$ except for

$$
i_{1}=2^{m-1}-2^{m^{\prime}-1}-1
$$

The roots of $h_{2}(X)$ are $\alpha^{-1}=\alpha^{2^{m}-2}, \alpha^{-2}=\alpha^{2^{m}-3}, \ldots, \alpha^{-2^{m-1}}=\alpha^{2^{m-1}-1}$. As it is done above, here we let

$$
i_{2}=2^{m-1}-1
$$

Let $X_{1}, X_{2}, \ldots, X_{w}$ be the location numbers of code vector* $v(x)$ of $C^{\prime}$. Then,

$$
v\left(\alpha^{i}\right)=\sum_{f=1}^{W} X_{f}^{i} \quad, \quad 0 \leq i<2^{m}-1
$$

For any $B_{0} € G F(2)$, any $B_{1} \in G F\left(2^{m^{\prime}}\right)$ and any $B_{2} \in G F\left(2^{m}\right)$, there exists a unique code vector $v(x)$ of $C^{\prime}$ such that $v(1)=B_{0}, v\left(\alpha^{i_{1}}\right)=B_{1}$ and $v\left(\alpha^{i_{2}}\right)=B_{2}$ (Mattson, Solomon [5]). Let $v\left(B_{o}, \beta_{1}, B_{2} ; x\right)$ denote the code vector specified by $B_{0}, B_{1}$ and $B_{2}$. By definition,

$$
x^{\ell} v\left(B_{0}, B_{1}, B_{2} ; x\right)=v\left(B_{0}, \alpha^{l i} 1_{B_{1}}, \alpha{ }^{l i}{ }_{B_{2}} ; x\right)
$$

If and only if $B_{0}=0, v\left(B_{0}, \beta_{1}, \beta_{2} ; x\right) \in C$. If and only if $B_{0}=\beta_{1}=0$ (or $B_{0}=B_{2}=0$ ), then $v\left(B_{0}, B_{1}, B_{2} ; x\right) \in C_{2}$ (or $C_{1}$ ). The cyclic permutations on code word symbols induce a permutation group on the code vectors of $\mathrm{C}^{\prime}$, which divides $\mathrm{C}-\mathrm{C}_{2}$ into disjoint sets of transitivity. Since $\nu=\left(\mathrm{i}_{1}, 2^{m}-1\right)$, each set consists of $\left(2^{m}-1\right) / \nu$ code vectors. In case of $m^{\prime}=m$, let $v\left(0, \alpha^{i}, B_{2} ; x\right)\left(0 \leq i<\nu, B_{2} \in G F\left(2^{m}\right)\right)$ represent each set. In case of $m^{\prime}=m / 2$, let $\mathrm{v}\left(0,1, \beta_{2} ; x\right)$ represent each set.

[^1]Now, consider the extended code $C_{\text {ex }}$ of code $C^{\prime}$. Let $\bar{v}\left(B_{0}, B_{1}, B_{2}\right)$ denote the vector with an overall parity check added to $v\left(B_{0}, \beta_{1}, B_{2} ; x\right)$ as the first digit. By definition $C e x=\left\{\bar{v}\left(\beta_{0}, B_{1}, \beta_{2}\right) \mid v\left(B_{0}, B_{1}, B_{2} ; x\right) \in C^{\prime}\right\}$. Let $X_{1}, X_{2}, \ldots, X_{w}$ be the location numbers of $\bar{v}\left(B_{o}, B_{1}, B_{2}\right)$ and let

$$
\begin{equation*}
S_{i}=\sum_{f=1}^{W} x_{f}^{i}, \quad 0 \leq i<2^{m}-1 \tag{7}
\end{equation*}
$$

Then, by definition

$$
\begin{align*}
& s_{i_{1}}=B_{1},  \tag{8}\\
& s_{i_{2}}=B_{2},  \tag{9}\\
& S_{i}=0\left(i \neq i_{1} 2^{\ell}, i_{2} 2^{\ell}\left(\bmod 2^{m}-1\right), 0 \leq \ell<m\right)
\end{align*}
$$

Therefore,

$$
\begin{equation*}
S_{i}=0\left(i \neq i_{1}, i_{1}^{\prime}, \quad 1 \leq i<2^{m-1}-1\right) \tag{10}
\end{equation*}
$$

By Lemma $1, \pi_{1 b} \overline{\mathrm{v}}\left(\mathcal{B}_{o}, \beta_{1}, B_{2}\right) \in C$ ex for any $b \in G F\left(2^{m}\right)$. Let

$$
\begin{equation*}
\pi_{1 b} \overline{\mathrm{v}}\left(\beta_{0}, B_{1}, \beta_{2}\right)=\overline{\mathrm{v}}\left(B_{0}^{\prime}, B_{1}^{\prime}, B_{2}^{\prime}\right) . \tag{11}
\end{equation*}
$$

The weights of $\bar{v}\left(B_{0}, \beta_{1}, \beta_{2}\right)$ and $\bar{v}\left(\beta_{0}^{\prime}, \beta_{1}^{\prime}, B_{2}^{\prime}\right)$ are the same. By the definition of $\pi_{1 b}$,

$$
\begin{aligned}
& B_{1}^{i}=\sum_{f=1}^{w}\left(X_{f}+b\right)^{i_{1}} \\
& B_{2}^{\prime}=\sum_{f=1}^{w}\left(X_{f}+b\right)^{i}{ }_{2} .
\end{aligned}
$$

Hence

$$
B_{1}^{p}=\sum_{f=1}^{w} \sum_{i=0}^{i_{1}}\left({ }_{i}{ }^{i_{1}}\right) x_{f}{ }^{i_{b}{ }^{i_{1}-i}}=\sum_{i=0}^{i_{1}}\left({ }_{i}{ }^{i}\right) S_{i} b^{i_{1}-i}
$$

From (8) and (10),

$$
\begin{equation*}
B_{1}^{\prime}=B_{1} \tag{12}
\end{equation*}
$$

Note that $\left(X_{f}+b\right)^{i_{2}}=\left(X_{f}+b\right)^{2^{m-1}-1}=\left(X_{f}{ }^{2^{m-1}}+b^{2^{m-1}}\right) /\left(X_{f}+b\right)=X_{f}^{2^{m-1}-1}+$ $x_{f}^{2^{m-1}-1} b+\cdots+b^{2^{m-1}-1}$. Then,

$$
\begin{equation*}
B_{2}^{0}=\sum_{f=1}^{w} \sum_{i=0}^{2^{m-1}-1} x_{f} i^{i} b^{m-1}-1-i=\sum_{i=0}^{2^{m-1}-1} S_{i} b^{2^{m-1}-1-i} \tag{13}
\end{equation*}
$$

Consider the case of $m^{\prime}=m$. Since $i_{1}^{\prime} \equiv 2^{h} i_{1}\left(\bmod 2^{m}-1\right), S_{i_{1}^{\prime}}=s_{i_{1}}^{2^{h}}$. From (8), (9), (10) and (13),

$$
\begin{align*}
& B_{2}^{0}=B_{1} b^{2^{m-1}-1-i} 1+B_{1}^{2^{h}} b^{2^{m-1}-1-i} 1 \\
&+B_{2}=B_{1} b^{2^{m-h-1}}  \tag{14}\\
&+B_{1}^{2^{h}} b^{2^{h-1}}+B_{2}
\end{align*}
$$

For the case of $m^{9}=m / 2$, it follows from (8), (9), (10) and (13) that

$$
\begin{equation*}
B_{2}^{\prime}=B_{1} b^{2^{h-1}}+B_{2} . \tag{15}
\end{equation*}
$$

Hereafter we shall consider the case of $m^{\prime}=m$ except for section 5 . For each $i(0 \leq i<\nu), V_{i}=\left\{\alpha^{i} b^{2^{m-h-1}}+\alpha^{i 2^{h}} b^{2^{h-1}} \mid b \in G F\left(2^{m}\right)\right\}$ forms $a$ subspace of GF (2 $2^{m}$. Let

$$
\begin{aligned}
F_{i}(X) & =\alpha^{i} x^{2^{m-h-1}}+\alpha^{i 2^{h}} x^{2^{h-1}} \\
& =\alpha^{i} x^{2^{h-1}}\left(x^{2^{h-1}}\left(2^{m-2 h}-1\right)+\alpha^{i\left(2^{h}-1\right)}\right)
\end{aligned}
$$

If $i=0$, the order of a nonzero root in $G F\left(2^{m}\right)$ of $F_{O}(X)$ is a factor of $2^{m-2 h}-1$. Since $c=(m, 2 h)=(m, m-2 h), 2^{c}-1=\left(2^{m}-1,2^{m-2 h}-1\right)$. This implies
that the roots in $G F\left(2^{m}\right)$ of $F_{0}(X)$ are in subfield $G F\left(2^{c}\right)$. Conversely, any element in this subfield is a root of $\mathrm{F}_{0}(\mathrm{X})$. Hence, the dimension of $\mathrm{V}_{0}$ is $\mathrm{m}^{-c}$. Let $\mathrm{V}_{00}\left(=\mathrm{V}_{0}\right), \mathrm{V}_{01}, \mathrm{~V}_{02}, \ldots, \mathrm{~V}_{02}{ }^{\mathrm{c}}-1$ be the cosets of $\mathrm{GF}\left(2^{\mathrm{m}}\right)$ with respect to $\mathrm{V}_{0}$. Each coset has $2^{\mathrm{m}-\mathrm{c}}$ elements.

For $i \neq 0$, assume that $\alpha^{j}$ is a root of $F_{i}(X)$. Then,

$$
\begin{aligned}
& 2^{h-1}\left(2^{m-2 h}-1\right) j \equiv i\left(2^{h}-1\right)\left(\bmod 2^{m}-1\right) \\
&\left(1-2^{2 h}\right) j \equiv i 2^{h+1}\left(2^{h}-1\right) \\
&\left(\bmod 2^{m}-1\right) \\
&-\left(2^{h}+1\right) j \equiv i 2^{h+1}\left(\bmod 2^{m}-1\right)
\end{aligned}
$$

Since $v$ divides both $2^{h}+1$ and $2^{m}-1$, $v$ must divide $i$. However, $0<i<v$. Therefore, there is no root in $G F\left(2^{m}\right)$ of $F_{i}(X)$ except for zero. Consequently, $V_{i}=G F\left(2^{m}\right)$.

Let $B_{0 j}=\left\{\left(\alpha^{\ell i_{1}}, \alpha^{\ell i^{2}} B\right) \mid 0 \leq \ell<2^{m}-1, B \in V_{0 j}\right\} \quad\left(0 \leq j \leq 2^{C}-1\right)$ and $B_{i}=\left\{\left(\alpha^{i+\ell i_{1}}, B\right) \mid 0 \leq \ell<2^{m}-1, \quad B \in G F\left(2^{m}\right)\right\}(0<i<v)$. Then,

$$
\begin{array}{ll}
\left|B_{0 j}\right|=\left(2^{m}-1\right) 2^{m-c} / \nu & \left(0 \leq j<2^{c}\right), \\
\left|B_{i}\right|=\left(2^{m}-1\right) 2^{m} / \nu & (0<i<\nu)^{*} . \tag{17}
\end{array}
$$

It follows from the definition of $B_{0 j}$ or $B_{i}$ that for any $\left(B_{1}, B_{2}\right)$ and $\left(B_{1}^{0}, B_{2}^{0}\right)$ in the same $B_{0}$ or $B_{i}$ and for any $B_{o} \in G F(2)$, there exists permutation $\pi_{a b}$ such that $\pi_{a b} \bar{v}\left(B_{o}, B_{1}, B_{2}\right)=\bar{v}\left(B_{o}^{\prime}, B_{1}^{\prime}, B_{2}^{\prime}\right)$ and $B_{o}^{\prime} \in G F(2)$. Therefore, $\overline{\mathrm{v}}\left(B_{0}, B_{1}, B_{2}\right)$ and $\overline{\mathrm{v}}\left(B_{0}^{\prime}, B_{1}^{\prime}, \beta_{2}^{\prime}\right)$ have the same weight w. If $B_{0}$ (or $B_{0}^{\prime}$ ) is zero,

[^2]then $v\left(0, B_{1}, B_{2} ; x\right)$ (or $\left.v\left(0, B_{1}^{\prime}, B_{2}^{\prime} ; x\right)\right)$ has weight w. If $B_{0}$ (or $B_{0}^{\prime}$ ) is one, then $v\left(1, B_{1}, \beta_{2} ; x\right)$ (or $\left.v\left(1, B_{1}^{\prime}, \beta_{2}^{\prime} ; x\right)\right)$ has weight $w-1$ by definition. Since $C^{0}$ contains all one vector $e=(1,1, \ldots, 1)$ and $e\left(\alpha^{i}\right)=\sum_{f=0}^{2^{m}-2} \alpha^{i f}=\left[\left(\alpha^{i}\right)^{2^{m}-1}-1\right] /$ $\left(\alpha^{i}-1\right)=0\left(0<i<2^{m}-1\right)$,
$$
v\left(1, B_{1}, B_{2} ; x\right)=v\left(0, B_{1}, B_{2} ; x\right)+e(x) .
$$

Hence, if $B_{0}$ (or $B_{0}^{0}$ ) is one, then $v\left(0, B_{1}, \beta_{2} ; x\right)$ (or $v\left(0, B_{1}^{\prime}, \beta_{2}^{\prime} ; x\right)$ ) has weight $2^{m}-w$. Therefore, we have Lemma 8 .

Lemma 8: For each $j$ (or i) $\left(0 \leq j<2^{c}, 0<i<\nu\right)$, there is $w_{o j}$ (or $w_{i}$ ) such that for any $\left(B_{1}, B_{2}\right) \in B_{0 j}$ (or $B_{i}$ ) the weight of $v\left(0, B_{1}, B_{2} ; x\right)$ is either $w_{0 j}\left(\right.$ or $\left.w_{i}\right)$ or $2^{m}-w_{0 j}\left(\right.$ or $\left.2^{m}-w_{i}\right)$.
3. Case I: $\quad(m, h)=(m, 2 h)$

Hereafter $a_{w}$ will denote the number of code vectors of weight $w$ in $C$ and $b_{w}$ will denote the number of code vectors of weight $w$ in $C_{0}$.

Lemma 9: For even w,

$$
\begin{aligned}
& w a_{w}=\left(2^{m}-w\right) a_{2^{m}-w} \\
& w b_{w}=\left(2^{m}-w\right) b_{2^{m}-w}
\end{aligned}
$$

This lemma follows from a theorem due to Peterson [1] and Lemma 1. Consequently, if the values of $w_{0}{ }^{\prime} s\left(0 \leq j<2^{c}\right)$ and $w_{i}^{\prime} s(0<i<\nu)$ are known, the weight distribution of $\mathrm{C}-\mathrm{C}_{2}$ is completely determined. Furthermore, any nonzero vector of $C_{2}$ has weight $2^{m-1}$, because $C_{2}$ is a maximum length sequence code.

Lemma 10: $\quad b_{1}=b_{2}=0, b_{3}=\left(2^{c^{\prime}-1}-1\right)\left(2^{m}-1\right) / 3$.

Proof: Since $C_{0}$ is a subcode of Hamming code,

$$
b_{1}=b_{2}=0
$$

Assume that $\alpha^{j_{1}}, \alpha^{j_{2}}$ and $\alpha^{j_{3}}$ are the location numbers of a code vector of weight 3 in $C_{0}$. Then,

$$
\begin{align*}
& \alpha^{j_{1}}+\alpha^{j_{2}}=\alpha^{j_{3}}  \tag{18}\\
& \alpha^{j_{1}\left(2^{h}+1\right)}+\alpha^{\left.j_{2}^{(2 h}+1\right)}=\alpha^{j_{3}\left(2^{h}+1\right)} \tag{19}
\end{align*}
$$

From (18),

$$
\begin{aligned}
\alpha^{j_{3}\left(2^{h}+1\right)} & =\left(\alpha^{j_{1}}+\alpha^{j_{2}}\right)^{2^{h}+1}=\left(\alpha^{j} 1^{2^{h}}+\alpha^{j_{2} 2^{h}}\right)\left(\alpha^{j_{1}}+\alpha^{j_{2}}\right) \\
& =\alpha^{j_{1}\left(2^{h}+1\right)}+\alpha^{j_{1} 2^{h}} \alpha^{j_{2}}+\alpha^{j_{1}} \alpha^{j_{2} 2^{h}}+\alpha^{j_{2}\left(2^{h}+1\right)}
\end{aligned}
$$

By combining with (19),

$$
\begin{align*}
& \alpha^{j_{1} 2^{h}} \alpha^{j_{2}}+\alpha^{j} 1_{1}{ }^{2^{2^{h}}}=0 \\
& \alpha^{\left(j_{1}-j_{2}\right)\left(2^{h}-1\right)}=1 \tag{20}
\end{align*}
$$

Thus,

$$
\left(j_{1}-j_{2}\right)\left(2^{h}-1\right) \equiv 0 \quad\left(\bmod 2^{m}-1\right)
$$

If $c^{\prime}=(m, h)=1$, then $\left(2^{h}-1,2^{m}-1\right)=1$. Therefore

$$
j_{1}=j_{2}
$$

This is a contradiction, which leads to the conclusion that $b_{3}=0$. If $c^{\prime} \neq 1$, then $\left(2^{m}-1,2^{h}-1\right)=2^{c}-1$. Let

$$
\begin{equation*}
\mu=\left(2^{m}-1\right) /\left(2^{c}-1\right) . \tag{21}
\end{equation*}
$$

Then,

$$
j_{1} \equiv j_{2}(\bmod \mu) .
$$

Let $j_{1}=\ell_{1} \mu+i$ and $j_{2}=\ell_{2} \mu+i(0 \leq i<\mu)$. Since $\alpha^{\ell_{1} \mu}+\alpha^{\ell_{2} \mu}=\alpha^{\ell_{3} \mu}$ for some $\ell_{3}$, it follows from (18) that $j_{3}=\ell_{3} \mu+i$. Conversely, for any $i$ $(0 \leq i<\mu)$ and for $\ell_{1}, \ell_{2}$, and $\ell_{3}$ such that

$$
\begin{align*}
& \alpha^{\ell_{1} \mu}+\alpha^{\ell_{2}^{\mu}}=\alpha^{\ell_{3}^{\mu}}  \tag{22}\\
& 0 \leq \ell_{1}, \ell_{2}, \ell_{3}<2^{c^{\prime}}-1 \tag{23}
\end{align*}
$$

$\alpha^{\ell_{1} \mu+i}, \alpha^{\ell_{2} \mu+i}$, and $\alpha^{\ell_{3} \mu+i}$ satisfy (18) and (19). The number of unordered triplets $\left(l_{1}, l_{2}, l_{3}\right)$ 's satisfying (22) and (23) is equal to $\left(2^{c^{\prime}}-1\right) / 3$. Consequently,

$$
b_{3}=\mu\left(2^{c^{\prime}}-1\right) / 3=\left(2^{m}-1\right)\left(2^{c^{\prime}-1}-1\right) / 3
$$

Lemma 11: Let $I_{2}$ and $I_{4}$ denote $\sum_{j \neq 0}\left(j-2^{m-1}\right)^{2} a_{j}$ and $\sum_{j \neq 0}\left(j-2^{m-1}\right)^{4} a_{j}$ respectively. Then,

$$
\begin{aligned}
& I_{2}=2^{2 m-2}\left(2^{m}-1\right) \\
& I_{4}=2^{3 m+c^{\prime}-4}\left(2^{m}-1\right)
\end{aligned}
$$

Proof: Note that $k=2 \mathrm{~m}$. By using the power moment identities of Lemma 6,

$$
\begin{aligned}
I_{2} & =\Sigma j^{2} a_{j}-2^{m} \Sigma j a_{j}+2^{2 m-2} \sum_{j \neq 0} a_{j} \\
& =2^{2 m-2} n(n+1)-2^{3 m-1} n+2^{2 m-2}\left(2^{2 m}-1\right) \\
& =\left(2^{m}-1\right)\left(2^{3 m-2}-2^{3 m-1}+2^{3 m-2}+2^{2 m-2}\right) \\
& =\left(2^{m}-1\right) 2^{2 m-2}
\end{aligned}
$$

$$
\begin{align*}
I_{4}= & \sum j^{4} a_{j}-2^{m+1} \Sigma j^{3} a_{j}+3 \cdot 2^{2 m-1} \Sigma j^{2} a_{j}-2^{3 m-1} \Sigma j a_{j} \\
+ & 2^{4 m-4} \sum_{j \neq 0} a_{j} \\
= & 2^{2 m-4}\left(n^{4}+6 n^{3}+3 n^{2}-2 n\right)-2^{3 m-2}\left(n^{3}+3 n^{2}\right) \\
+ & 3 \cdot 2^{4 m-3} n(n+1)-2^{5 m-2} n+2^{4 m-4}\left(2^{2 m}-1\right) \\
+ & 3\left(2^{3 m-1}-2^{2 m-1} n\right) b_{3}+3 \cdot 2^{2 m-2} b_{4} \\
= & n\left\{2^{2 m-4}\left[(n+1)^{3}+3(n+1)(n-1)\right]-2^{3 m-2}\left[(n+1)^{2}+\right.\right. \\
& \left.n-1]+3 \cdot 2^{4 m-3}(n+1)-2^{5 m-2}+2^{4 m-4}\left(2^{m}+1\right)\right\} \\
+ & 3 \cdot 2^{2 m-1}\left(b_{3}+b_{4}\right) \\
= & n\left\{2^{5 m-4}-2^{5 m-2}+3 \cdot 2^{5 m-3}-2^{5 m-2}+2^{5 m-4}\right. \\
& \left.+3 \cdot 2^{3 m-4}\left(2^{m}-2\right)-2^{3 m-2}\left(2^{m}-2\right)+2^{4 m-4}\right\} \\
+ & 3 \cdot 2^{2 m-1}\left(b_{3}+b_{4}\right) \\
= & \left(2^{m}-1\right) 2^{3 m-3}+3 \cdot 2^{2 m-1}\left(b 3^{+b}\right) \tag{24}
\end{align*}
$$

Since all one vector $(1,1, \ldots, 1)$ is in $C_{0}$,

$$
\mathrm{b}_{2^{\mathrm{m}}-4}=\mathrm{b}_{3}
$$

By Lemmas 9 and 10 ,

$$
\begin{align*}
b_{3}+b_{4} & =b_{2^{m}-4}+b_{4}=2^{m-2}{ }_{2^{m}-4}=2^{m-2} b_{3} \\
& =2^{m-2}\left(2^{c^{\prime}-1}-1\right)\left(2^{m}-1\right) / 3 \tag{25}
\end{align*}
$$

By substituting the right hand side of (25) into (24),

$$
I_{4}=\left(2^{m}-1\right) 2^{3 m+c^{1}-4}
$$

Let $j_{M}$ be the smallest nonzero integer such that

$$
a_{j_{M}}+a_{2^{m}-j_{M}} \neq 0
$$

By the definition of $I_{2}$ and $I_{4}$,

$$
\begin{equation*}
\left(j_{M}-2^{m-1}\right)^{2} \geq I_{4} / I_{2}=2^{m+c^{\prime}-2} \tag{26}
\end{equation*}
$$

Consider the case where $c=c^{\prime}$. Then, $\nu=1$ by Lemma 5. Since all nonzero vectors in $C_{2}$ are of weight $2^{m-1}$, it follows from Lemma 8 and (16) that $a_{j}+a_{2^{m}-j}\left(j \neq 0,2^{m-1}\right)$ must be divisible by $2^{m-c^{\prime}}\left(2^{m}-1\right)=2^{m-c}\left(2^{m}-1\right)$.
Therefore, from (26)

$$
\begin{equation*}
I_{2} \geq\left(j_{M}-2^{m-1}\right)^{2}\left(a_{j_{M}}+a{ }_{2^{m}-j_{M}}\right) \geq 2^{2 m-2}\left(2^{m}-1\right) \tag{27}
\end{equation*}
$$

By Lemma 11 and (27),

$$
I_{2}=\left(j_{M}-2^{m-1}\right)^{2}\left(a_{j_{M}}+a{ }_{2^{m}-j_{M}}\right)=2^{2 m-2}\left(2^{m}-1\right)
$$

Consequently,

$$
\begin{align*}
& \left(j_{M}-2^{m-1}\right)^{2}=2^{m+c^{\prime}-2}=2^{m+c-2}  \tag{28}\\
& a_{j_{M}}+a_{2^{m}-j_{M}}=2^{m-c}\left(2^{m}-1\right)  \tag{29}\\
& a_{j}=0\left(j \neq 0, \quad j_{M}, 2^{m-1}, 2^{m}-j_{M}\right)
\end{align*}
$$

Hence,

$$
\begin{align*}
j_{M} & =2^{m-1}-2^{(m+c) / 2-1} \\
a_{2^{m-1}} & =2^{2 m}-1-\left(a_{j_{M}}+a^{m}-j_{M}\right)  \tag{30}\\
& =\left(2^{m}-2^{m-c}+1\right)\left(2^{m}-1\right)
\end{align*}
$$

By Lemma 9, (29) and (30),

$$
\begin{aligned}
& a_{j_{M}}=\left(2^{m-c-1}+2^{(m-c) / 2-1}\right)\left(2^{m}-1\right) \\
& a_{2^{m}-j_{M}}=\left(2^{m-c-1}-2^{(m-c) / 2-1}\right)\left(2^{m}-1\right)
\end{aligned}
$$

Thus, we have the following theorem.
Theorem 1: If $(m, h)=(m, 2 h)=c$, then

$$
\begin{aligned}
& a_{0}=1 \\
& a_{2^{m-1}-2^{(m+c) / 2-1}}=\left(2^{m-c-1}+2^{(m-c) / 2-1}\right)\left(2^{m}-1\right) \\
& a_{2^{m-1}}=\left(2^{m}-2^{m-c}+1\right)\left(2^{m}-1\right) \\
& a_{2^{m-1}}+2^{(m+c) / 2-1}=\left(2^{m-c-1}-2^{(m-c) / 2-1}\right)\left(2^{m}-1\right) \\
& a_{j}=0 \text { for other } j .
\end{aligned}
$$

4. Case II: $2(\mathrm{~m}, \mathrm{~h})=(\mathrm{m}, 2 \mathrm{~h}) \neq \mathrm{m}$

Consider the case in which $2(m, h)=(m, 2 h)$. Then, $\nu=2^{c^{0}}+1$ by
Lemma 5. Since $v\left(0, B_{1}, 0 ; x\right)$ is in $C_{1}, w_{00}=w_{0}, w_{1}, \ldots, w_{v-1}$ can be found from the weight distribution of $C_{1}$. Each code vector of $C_{1}$ is a $\nu$ concatenation of a code vector in cyclic code $C_{1}^{\prime}$ of length $\left(2^{m}-1\right) / \nu$ which is generated by $\left(X^{\left(2^{m}-1\right) / \nu}-1\right) / h_{1}(X)$. Let $v^{\prime}(B ; x)$ denote a code vector $v^{0}(x)$
in $C_{1}^{\prime}$ such that $v^{\prime}\left(\alpha^{i}\right)=B$. With the same argument as the one of section 3 , set $C_{1 i}^{\prime}=\left\{v\left(\alpha^{i+\ell \nu} ; x\right) \mid 0 \leq \ell<\left(2^{m}-1\right) / \nu\right\} \quad(0 \leq i<\nu)$ consists of $\left(2^{m}-1\right) / \nu$ vectors of the same weight $w_{i}^{\prime}$. Since $v(0, B, 0 ; x)=\left(v^{\prime}(B ; x)\right.$, $\underbrace{\left.v^{\prime}(B ; x), \ldots, v^{\prime}(B ; x)\right)}_{\nu}, w_{i}$ can be given by the following equation:

$$
w_{i}=\nu w_{i}^{\prime} \cdot \quad(0 \leq i<\nu)
$$

Therefore, by applying Lemma 6 to code $C_{1}^{\prime}$, we have

$$
\begin{aligned}
& \left(2^{m}-1\right) / \nu \sum_{i=0}^{\nu-1} w_{i} / \nu=2^{m-1}\left(2^{m}-1\right) / \nu \\
& \left(2^{m}-1\right) / \nu \sum_{i=0}^{\nu-1}\left(w_{i} / \nu\right)^{2}=2^{m-2}\left(2^{m}-1\right)\left(2^{m}-1+\nu\right) / \nu^{2}
\end{aligned}
$$

Thus,

$$
\begin{align*}
& \sum_{i=0}^{\nu-1} w_{i}=\nu 2^{m-1}  \tag{31}\\
& \sum_{i=0}^{\nu-1} w_{i}^{2}=\nu 2^{m-2}\left(2^{m}-1+\nu\right)=\nu 2^{m-2}\left(2^{m}+2^{c^{\prime}}\right) \tag{32}
\end{align*}
$$

Therefore,

$$
\begin{align*}
\sum_{i=0}^{\nu-1}\left(w_{i}-2^{m-1}\right)^{2} & =\nu 2^{m-2}\left(2^{m}-1+\nu\right)-2^{m} \nu 2^{m-1}+2^{2 m-2} \nu \\
& =\nu 2^{m-2}(\nu-1)=\nu 2^{m+c^{\prime}-2} \tag{33}
\end{align*}
$$

On the other hand, it follows from Lemma 8, (16), (17) and the definition of $I_{2}$ that

$$
\begin{align*}
& I_{2}=I_{20}+I_{21}, \\
& I_{20}=2^{m-c}\left(2^{m}-1\right) \sum_{j=1}^{2^{c}-1}\left(w_{0} j^{-2^{m-1}}\right)^{2} / \nu  \tag{34}\\
& I_{21}=2^{m-c}\left(2^{m}-1\right)\left(w_{0}-2^{m-1}\right)^{2} / \nu+2^{m}\left(2^{m}-1\right) \sum_{i=1}^{\nu-1}\left(w_{i}-2^{m-1}\right)^{2} / \nu .
\end{align*}
$$

By Lemma 11 and (33),

$$
\begin{align*}
I_{2} & =2^{2 m-2}\left(2^{m}-1\right) \\
& \geq I_{21}=2^{2 m+c^{\prime}-2}\left(2^{m}-1\right)-\left(2^{m}-1\right)\left(2^{m}-2^{m-c}\right)\left(w_{0}-2^{m-1}\right)^{2} / \nu . \tag{35}
\end{align*}
$$

By a simple calculation,

$$
\begin{aligned}
\left(w_{0}-2^{m-1}\right)^{2} & \geq 2^{2 m-2}\left(2^{c^{\prime}}-1\right) v /\left(2^{m}-2^{m-c}\right) \\
& =2^{m+c-2}
\end{aligned}
$$

Hence,

$$
\begin{equation*}
\left.w_{0}=2^{m-1} \pm 2^{(m+c) / 2-1}+\delta\right), \quad \delta \geq 0 \tag{36}
\end{equation*}
$$

Now, by (31) and (32)

$$
\begin{align*}
I_{2}^{\prime} & =\sum_{i=0}^{\nu-1}\left(w_{i}-\left(2^{m-1} \mp 2^{m / 2-1}\right)\right)^{2} \\
& =\nu 2^{m-2}\left(2^{m}-1+\nu\right)-2\left(2^{m-1}+2^{m / 2-1}\right) \nu 2^{m-1}+\left(2^{m-1} \mp 2^{m / 2-1}\right)^{2} \nu \\
& =\nu\left\{2^{2 m-2}+2^{m+c^{\prime}-2}-2^{2 m-1} \pm 2^{3 m / 2-1}+2^{2 m-2}+2^{3 m / 2-1}+2^{m-2}\right\} \\
& =2^{m-2}\left(2^{c^{\prime}}+1\right)^{2} \tag{37}
\end{align*}
$$

On the other hand,

$$
\begin{align*}
I_{2}^{\prime} \geq\left(w_{0}-\left(2^{m-1}+2^{m / 2-1}\right)\right)^{2} & =\left[2^{m-1} \pm\left(2^{(m+c) / 2-1}+\delta\right)-2^{m-1} \pm 2^{m / 2-1}\right]^{2} \\
& =2^{m-2}\left(2^{c^{\prime}}+1+\delta 2^{1-m / 2}\right)^{2} \tag{38}
\end{align*}
$$

From (35), (36), and (37), we have that

$$
\begin{align*}
& \delta=0 \\
& w_{0}=2^{m-1} \pm 2^{(m+c) / 2-1}  \tag{39}\\
& w_{i}=2^{m-1} \mp 2^{m / 2-1} \quad(0<i<\nu) \tag{40}
\end{align*}
$$

Since $w_{i}(0 \leq i<\nu)$ is divisible by $\nu$, the $\pm$ sign is determined by Lemma 4 . Thus, we have:

Theorem 2: Let $a_{j \nu}^{\rho}$ denote the number of code vectors of weight $j$ in $C_{1}^{\rho}$. If $\mathrm{m} / \mathrm{c}$ is odd (or even), then

$$
\begin{aligned}
& a_{0}^{\prime}=1 \\
& a_{2^{\prime}}^{\prime m-1}-2^{(m+c) / 2-1}\left(\text { or } a_{2^{\prime}}^{m-1}+2^{(m+c) / 2-1}\right)=\left(2^{m}-1\right) /\left(2^{c / 2}+1\right) \\
& a^{\prime} \\
& 2^{m-1}+2^{m / 2-1}\left(\text { or } a_{2^{\prime}}^{m-1}-2^{m / 2-1}\right)=2^{c / 2}\left(2^{m}-1\right) /\left(2^{c / 2}+1\right) . \\
& a_{j}^{\prime}=0 \quad \text { for other } j .
\end{aligned}
$$

From (35) and (39) it follows that

$$
\begin{aligned}
I_{21} & =2^{2 m+c^{\prime}-2}\left(2^{m}-1\right)-\left(2^{m}-1\right)\left(2^{m}-2^{m-c}\right) 2^{m+c-2} / \nu \\
& =2^{2 m-2}\left(2^{m}-1\right)\left(2^{c^{\prime}}-\left(2^{c}-1\right) / \nu\right) \\
& =2^{2 m-2}\left(2^{m}-1\right) \\
& =I_{2}
\end{aligned}
$$

By (34),

$$
I_{20}=\left(2^{m}-1\right) 2^{m-c} / \nu \sum_{j=1}^{2^{c}-1}\left(w_{0 j}-2^{m-1}\right)^{2}=0
$$

Hence,

$$
\begin{equation*}
w_{0 j}=2^{m-1} \quad\left(1 \leq j<2^{c}\right) \tag{41}
\end{equation*}
$$

By (16), (17), Lemma 8, (39), (40) and (41), we have:

$$
\begin{aligned}
& a_{2^{m-1}} 2^{m / 2-1}+a{ }_{2^{m-1}+2^{m / 2-1}}=\left(2^{m}-1\right) 2^{m+c / 2} /\left(2^{c / 2}+1\right) \\
& a_{2^{m-1}}=\left(2^{m}-1\right) 2^{m-c}\left(2^{c}-1\right) /\left(2^{c / 2}+1\right)+2^{m}-1
\end{aligned}
$$

$$
{ }_{2}^{m-1}-2(m+c) / 2-1+a_{2^{m-1}+2}(m+c) / 2-1=\left(2^{m}-1\right) 2^{m-c} /\left(2^{c / 2}+1\right)
$$

Thus, the next theorem follows from Lemma 9.
Theorem 3: If $2(m, h)=(m, 2 h)=c$ and $c \neq m$, then

$$
\begin{aligned}
& a_{0}=1 \\
& a_{2^{m-1}-2}(m+c) / 2-1=\left(2^{m}-1\right)\left(2^{(m-c) / 2}+1\right) 2^{(m-c) / 2-1} /\left(2^{c / 2}+1\right) \\
& a_{2^{m-1}-2^{m / 2}}=\left(2^{m}-1\right)\left(2^{m / 2}+1\right) 2^{(m+c) / 2-1} /\left(2^{c / 2}+1\right) \\
& a_{2^{m-1}}=\left(2^{m}-1\right)\left(\left(2^{c / 2}-1\right) 2^{m-c}+1\right) \\
& a_{2^{m-1}}{ }^{m} 2^{m / 2}=\left(2^{m}-1\right)\left(2^{m / 2}-1\right) 2^{(m+c) / 2-1} /\left(2^{c / 2}+1\right) \\
& a_{2^{m-1}}^{m-2}(m+c) / 2-1=\left(2^{m}-1\right)\left(2^{(m-c) / 2}-1\right) 2^{(m-c) / 2-1} /\left(2^{c / 2}+1\right) \\
& a_{j}=0 \quad \text { for other } j .
\end{aligned}
$$

## 5. Case III: $2(m, h)=(m, 2 h)=m$

Consider the case of $m=2 h$. For any $B_{1} \neq 0$, $B_{2}$ in $G F\left(2^{m}\right)$, there exists $b \in G F\left(2^{m}\right)$ such that

$$
B_{1} b^{2^{h-1}}+B_{2}=0
$$

because $\left(2^{h-1}, 2^{m}-1\right)=1$. From (15) and a similar argument to the one for the case of $m=m^{\prime}$, it follows that there exists $w$ such that the weight of any code vector in $C-C_{2}$ is either $w$ or $2^{m}-w$. Since $C_{1}^{\prime}$, the cyclic code of length $2^{m / 2}-1$ generated by $\left(x^{2^{m / 2}-1}-1\right) / g_{1}(X)$, is a maximum length sequence code, $C_{1}$ consists of one zero vector and $2^{m / 2}-1$ vectors of weight $2^{m / 2-1}$ $\left(2^{m / 2}+1\right)$. Therefore,

$$
\mathrm{w}=\left(2^{\mathrm{m} / 2}+1\right) 2^{\mathrm{m} / 2-1}
$$

On the other hand, $C_{2}$ is a maximum length sequence code of length $2^{m}-1$. Hence,

$$
a_{2^{m-1}}=2^{m}-1
$$

According to Lemma 9, we have:
Theorem 4: If $m=2 h$, then

$$
\begin{aligned}
& a_{0}=1 \\
& a_{2^{m-1}}-2^{m / 2-1}=\left(2^{m / 2}-1\right)\left(2^{m-1}+2^{m / 2-1}\right) \\
& a_{2^{m-1}}=2^{m}-1 \\
& a_{2^{m-1}}+^{m / 2-1}=\left(2^{m / 2}-1\right)\left(2^{m-1}-2^{m / 2-1}\right) \\
& a_{j}=0 \quad \text { for other } j .
\end{aligned}
$$

## 6. Crosscorrelation Functions of Two Maximum Length Sequences

It follows from Lemma 5 that $\alpha^{2^{h}+1}$ is a primitive element if and only if $c=c^{\prime}$. Assume that $c=c^{\prime}$. Let

$$
\begin{aligned}
& v(0,1,0 ; x)=\sum_{f=0}^{2^{m}-2} v_{1 f^{x}} x^{f} \\
& v(0,0,1 ; x)=\sum_{f=0}^{2^{m}-2} v_{2 f} x^{f}
\end{aligned}
$$

Then, $\mathrm{v}_{1}=\mathrm{v}_{10}, \mathrm{v}_{11}, \ldots, \mathrm{v}_{12^{\mathrm{m}}-2}$ and $\mathrm{v}_{2}=\mathrm{v}_{20}, \mathrm{v}_{21}, \ldots, \mathrm{v}_{22^{\mathrm{m}}-2}$ are maximum length sequences of length $2^{m}-1$. In $v_{1}$ and $v_{2}$, replace 0 by -1 . Let $\mathrm{u}_{1}=\mathrm{u}_{10}, \mathrm{u}_{11}, \ldots, \mathrm{u}_{12^{\mathrm{m}}-2}$ and $\mathrm{u}_{2}=\mathrm{u}_{20}, \mathrm{u}_{21}, \ldots, \mathrm{u}_{22^{\mathrm{m}}-2}$ be the resulting
sequences of real numbers 1 and -1 . Correlation function $\theta(j)$ of $u_{1}$ and $u_{2}$ is defined by

$$
\theta(j)=\sum_{f=0}^{2^{m}-2} u_{1 f} u_{2 f-j}
$$

where suffix $f-j$ is to be taken $\bmod 2^{m}-1$. Note that

$$
v(0,1,0 ; x)+x^{j} v(0,0,1 ; x)=v\left(0,1, \alpha^{j i_{2}} ; x\right)
$$

If $v\left(0,1, \alpha^{j i} ; x\right)$ has weight $w$, then

$$
\begin{equation*}
\theta(j)=2^{m}-1-2 w . \tag{42}
\end{equation*}
$$

Let $s_{i}$ denote the number of $j^{\prime} s\left(0 \leq j<2^{m}-1\right)$ with $\theta(j)=i$. Then, $s_{i}$ is the number of vectors $v(0,1, \beta ; x)$ with weight $\left(2^{m}-1-i\right) / 2$. From sections 2 and 3 , we have theorem 5.

Theorem 5:

$$
\begin{aligned}
& s_{-2}^{(m+c) / 2}-1=2^{m-c-1}-2^{(m-c) / 2-1} \\
& s_{-1}=2^{m}-2^{m-c} \\
& s_{2}^{(m+c) / 2}-1=2^{m-c-1}+2^{(m-c) / 2-1}, \\
& s_{i}=0 \quad \text { for other } i .
\end{aligned}
$$

For $0 \leq j<2^{m}-1$, let

$$
\begin{aligned}
& \theta_{j}=1 \text { if } \theta(j)=-2^{(m+c) / 2}-1 \text { or } 2^{(m+c) / 2}-1, \\
& \theta_{j}=0 \text { if } \theta(j)=-1
\end{aligned}
$$

Sequence $\theta=\theta_{0}, \theta_{1}, \ldots, \theta_{2^{m}-2}$ will be called the correlation sequence of $u_{1}$ and $u_{2}$. We shall characterize the correlation sequence below. Recall that

$$
v_{0}=v_{00}=\left\{b^{2^{m-h-1}}+b^{2^{h-1}} \mid b \in G F\left(2^{m}\right)\right\}
$$

Since $\left(2^{\mathrm{h}-1}, 2^{\mathrm{m}}-1\right)=1$,

$$
v_{00}=\left\{b^{2^{m-2 h}}+b \mid b \in G F\left(2^{m}\right)\right\}
$$

Since $(m-2 h, m)=c$, the Galois group of $G F\left(2^{m}\right)$ over $G F\left(2^{c}\right)$ is generated by the automorphism $x \rightarrow x^{2^{m-2 h}}$ (by Theorem 9, p. 127 of [6]). Therefore, from the trace theorem (p. 121 of [6]) it follows that

$$
\mathrm{V}_{00}=\left\{b \mid \sigma(b)=0, \quad b \in G F\left(2^{m}\right)\right\},
$$

where $\sigma(b)$ denotes the trace of $b$ in $\operatorname{GF}\left(2^{m}\right)$ over $\operatorname{GF}\left(2^{c}\right)$, and

$$
\sigma(b)=b+b^{2^{c}}+b^{2^{2 c}}+\ldots+b^{2^{m-c}}
$$

Hence, any element of each coset $V_{0 j}$ has the same trace $t_{j} \in G F\left(2^{c}\right)$, and if $j \neq j^{r}, t_{j} \neq t_{j r}$ 。

Note that $v^{2}\left(0, B_{1}, B_{2} ; x\right)=v\left(0, B_{1}, B_{2} ; x^{2}\right)=v\left(0, B_{1}{ }^{2}, B_{2}{ }^{2} ; x\right) \in C$ and that $v\left(0, B_{1}, B_{2} ; x\right)$ and $v\left(0, B_{1}, B_{2} ; x^{2}\right)$ have the same weight. Consequently, if $t_{j^{\prime}}=t_{j}{ }^{2}$, then $w_{0 j^{\prime}}=w_{0 j^{\prime}}$. From the proof of Theorem 1 it follows that there is only one $j_{0}$ such that

$$
\begin{aligned}
& { }^{w_{0}} j_{0}=2^{m-1}-2^{(m+c) / 2-1} \\
& { }^{w_{0 j}}=2^{m-1}, \quad j \neq j_{0} .
\end{aligned}
$$

Since $t_{j_{0}}=t_{j_{0}}^{2}, t_{j_{0}}=0$ or 1 . Since $C_{1}$ is a maximum length sequence code, ${ }_{0} 0$ must be $2^{m-1}$. Therefore, $t_{j_{0}}=1$. This implies that the weight of $v(0,1, B ; x)$ is not equal to $2^{m-1}$ if and only if $B+B^{2^{c}}+\beta^{2^{2 c}}+\ldots+$ $B^{2^{m-c}}=1$. From (42), $\theta(\mathrm{j}) \neq-1$, if and only if $\alpha^{j i} 2+\alpha^{j i} 2^{2 c}+\alpha^{j i^{2} 2^{2 c}}+$ $\ldots+\alpha^{j i} 2^{2^{m-c}}=1$. Since $\alpha^{2 i_{2}}=\alpha^{2^{m}-2}=\alpha^{-1}$ and $\alpha^{-2^{m-1}}=\alpha^{i_{2}}$, we have

Theorem 6 。

Theorem 6: $\theta_{j}=1$ if and only if $\alpha^{-j}+\alpha^{-j 2^{c}}+\alpha^{-j 2^{2 c}}+\ldots+\alpha^{-j 2^{m-c}}$ $=1$ 。

Now consider the case of $c=1$. By definition, $v\left(0,0,1 ; \alpha^{i}\right)=1$. Therefore,

$$
\mathrm{v}\left(0,0,1 ; \alpha^{-2^{\ell}}\right)=1 \quad(0 \leq \ell<m)
$$

On the other hand, for any $\alpha^{j} \neq \alpha^{-2^{\ell}}(0 \leq \ell<m)$,

$$
v\left(0,0,1 ; \alpha^{j}\right)=0
$$

By the formula due to Reed and Solomon [7],

$$
\begin{aligned}
v_{2 f} & =\sum_{j=0}^{2^{m}-2} v\left(0,0,1 ; \alpha^{j}\right) \alpha^{-j f} \\
& =\alpha^{f}+\alpha^{2 f}+\alpha^{2^{2} f}+\cdots+\alpha^{2^{m-1}} f
\end{aligned}
$$

Hence, by Theorem 6

$$
\theta_{j}=v 2^{m}-1-j
$$

This implies the following corollary.
Corollary 7: If $(m, h)=(m, 2 h)=1$, then the correlation sequence of the maximum length sequence generated by $h_{1}(X)$ and the one generated by $h_{2}(X)$ is the maximum length sequence generated by $g_{2}(X)=X^{m} h_{2}\left(X^{-1}\right)$.
R. Gold and E. Kopitzka [8] observed that for some pairs of maximum length sequences the correlation sequences are also maximum length sequences and they listed all such pairs of sequences of length 8191 or less. Among 28 listed pairs, 25 cases are covered by Corollary 7.

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[^0]:    * Dr. B. Elspas pointed out this relation.

[^1]:    *A polynomial representation will be used for a code vector [4].

[^2]:    *|B| means the number of elements of $B$.

