WRC RESEARCH REPORT NO. 122

OPTIMAL OPERATION OF FLOOD CONTROL SYSTEMS

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FINAL REPORT: VOLUME II

Project No. A-079-ILL

This project was partially supported by the U. S. Department of the Interior in accordance with the Water Resources Research Act of 1964, P.L. 88-379, Agreement No. 14-31-0001-6014.

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March 1977

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ABSTRACT

OPTIMAL OPERATION OF FLOOD CONTROL SYSTEMS

The management and control of multiple-reservoir flood control systems is studied. Our objective is to devise operating policies which minimize flood damages as determined by the flood peaks. Methodologies are presented that employ dynamic programming and stochastic dynamic programming for the optimal operation of multiple-reservoir flood control systems with deterministic and stochastic inflows, respectively. The methodologies are applied to a number of real-world problems involving river basins in Illinois and elsewhere. We then study the effects of parametrically varying a number of the input parameters. A Markov renewal flood synthesis model and a methodology for determining the optimal capacity of a new flood control reservoir are also presented.

Feng, C-J and Morin, T. L.

OPTIMAL OPERATION OF FLOOD CONTROL SYSTEMS

University of Illinois Water Resources Center Report No. 122

KEYWORDS -- flood control/reservoirs/dynamic programming/operation/water resource system

PREFACE

This report is one of the two volumes comprising the completion report for Project No. A-079-ILL. This volume deals with the operation of flood control systems, whereas the other volume deals with the planning of flood control systems. The authors would like to acknowledge useful discussions with E. Çinlar, R. Gemmell and S. Pliska.

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CHAPTER I

INTRODUCTION

1.1 Reservoir System Management and Flood Control

Despite the substantial expenditures on structural measures for flood protection, flood damages continue to increase. For example, at the local level on the upper Salt Creek Watershed in Illinois the annual damages are projected [41] to increase from \$412,600 in 1970 to \$853,000 in 1995. The increase is even more dramatic at the national level, where the total annual flood damage potential is estimated to increase from \$1.7 billion in 1966 to \$5.0 billion in 2020 on the basis of the current status of flood control improvements and the projected conditions of floodplain use [86]. These dramatic increases have given rise to increased awareness and interest in floodplain management and control. This is precisely the issue to which this study is addressed.

1.2 Review of Relevant Research

Several researchers have applied linear decision rules in discrete time parameter models in order to determine the optimal capacity of a flood control reservoir. ReVelle [75] developed a chance-constrained model for determining the optimal size and operating policy of a multipurpose reservoir. Subsequently, ReVelle and Kirby [74] developed the model more rigorously, modified it to include evaporation loss, and used the basic model with four performance objectives, given the reservoir size. Eastman and ReVelle [25] developed a chance constrained model for minimizing the required capacity subject to chance constraints for water supply, recreation and flood control. Nayak and Arora [62] considered a chance constraint formulation of multiple-reservoir system in order to find the optimal reservoir capacity. Recently, Sobel [80] determined the optimal reservoir capacity by minimizing the maximum stored quantity and specified operating rules for multiple-reservoir systems in the case of independent inflow, assuming that the outflow of upstream reservoirs does not affect the downstream reservoirs.

There are also some papers which propose the use of linear decision, chance-constrained models for reservoir management, such as Morton [59], in which the demand for water is stochastic. Eisel [26] developed a model for a single reservoir which demonstrates the mathematical complexities resulting from the convolution problem in chance-constrained programming. Young and Asce [91] developed recursive algorithm in order to find the optimal operating policies for a single reservoir which minimize a strictly convex quadratic loss function while satisfying pre-specified target releases. They also studied the effects of stochastic inflows using a Monte Carlo approach. Roefs and Bodin [76] considered the multiple-reservoir system operation problem with energy production considerations. For other operating models see [87, 88]. However, none of the foregoing models consider the possibility of overflow which may be very important in flood control. More importantly, most models are for long term operation instead of individual storms.

There is also a body of relevant literature on descriptive mathematical models of reservoir systems. Çinlar and Pinsky [16, 17] considered the case where the release function is assumed to be arbitrary

...2

and the input process to be an additive process. They solved the storage equation and obtained a number of properties. Çinlar [13] generalized the model to the continuous time and state space case with arbitrary release functions. One of the basic features in the input process is the dependence on environmental factors. Çinlar [14] also studied one of the most general cases, in which the release function is an arbitrary continuous nondecreasing function vanishing at the origin, the environment process is a Markov Process and the input process is a nonstationary additive process defined on the environment process. Moran's [50, 51] model considered a constant environment input process and constant release rules. Bather [3] supposed a controlled Brownian motion process as a model of a reservoir. More recently, Pliska [70] considered the water level in a reservoir as a controlled diffusion process on a compact interval of the real line. These control models also did not consider the overflows.

Additional relevant literature will be cited and reviewed in the body of study where appropriate.

1.3 Purposes and Outline of the Study

The purposes of this study are:

- (1) To formulate mathematical programming models for
 - a) the optimal control of both parallel and series multiplereservoir flood control systems for both single-storm and long-term operation with both deterministic and stochastic inflows.
 - b) the determination of the optimal capacity of a new reservoir to add to an existing flood control system.

- (2) To develop methods for sloving these models.
- (3) To implement the solution methodologies on a digital computer.
- (4) To illustrate and assess the practicality and feasibility of the solution methodologies using data from a number of real-world examples.

These purposes are realized in the following manner.

In Chapter II, mathematical models of the optimal operation of both single-reservoir systems and parallel and series multiple-reservoir flood control systems with the deterministic inflows are formulated. Dynamic programming is used to determine the optimal release policies. The dynamic programming algorithms are coded in FORTRAN IV and implemented on a digital computer. Several real-world examples are solved to illustrate the methodology.

In Chapter III, mathematical models of the optimal operation of both single-reservoir flood control systems and parallel and series multiple-reservoir systems with stochastic inflows are formulated. Stochastic Dynamic Programming is used to determine optimal release policies under the assumptions that the input probabilities are time independent and time-dependent respectively. The dynamic programming algorithms are implemented on a digital computer and used to solve several real-world examples. The effectiveness of these policies is compared with both actual policies and standard policies using both historical flow data and simulated flow data. A stochastic flood synthesis model is also presented.

The effect of varying various flood control parameters is studied in Chapter IV. A number of characterizations of optimal flood control policies are derived.

In Chapter V, the optimal expansion of flood control systems is modeled. Dynamic Programming is used to determine the optimal capacity of additional reservoirs in both the parallel and series cases. The dynamic programming algorithms are implemented on a digital computer and several real-world examples are solved to illustrate the solution methodologies.

The study concludes with a summary and recommendations for future research in Chapter VI.

CHAPTER II

DETERMINISTIC INFLOW

2.1 Introduction

In this chapter, the basic reservoir control model is formulated in §2.2 and §2.3, assuming a given input hydrograph. The linear programming and mixed-integer programming solution methodologies are reviewed and critically evaluated in §2.4. Dynamic programming algorithms, which overcome the shortcomings of these other methodologies, are presented in §2.5. The algorithms are implemented on Northwestern University's CDC 6600 and used to solve several real-world examples.

2.2 The Basic Model

Figure 2.1 depicts the physical structure of the hypothetical river basin systems under investigation. The extension of procedure to handle other system topologies should be immediately apparent.

For the series system shown in the Figure 2.1a, the basic assumptions are that:

- the water released from reservoir 1 at time t arrives at reservoir 2 at time t, i.e., instantaneous flow;
- (2) each reservoir spillway is designed in such a way that it allows for any kind of release rate;
- (3) the capacity C_2 of the downstream reservoir exceeds that C_1 of the upstream reservoir;



Figure 2.1 Proposed Multiple-reservoir flood control system. (a)Series (b)Parallel.

(4) the releases and inflows occur simultaneously.

The storage equation for the series system can be written as

$$Z_{t}^{1} = Z_{0}^{1} + \int_{0}^{t} Y^{1}(s) ds - \int_{0}^{t} r^{1}(s) ds \qquad t \ge 0,$$
 (2.1)

and

$$z_{t}^{2} = z_{0}^{2} + \int_{0}^{t} y^{2}(s) ds + \int_{0}^{t} (r^{1}(s) - r^{2}(s)) ds \qquad t \ge 0, \qquad (2.2)$$

in which $r^{i}(s)$ is the release rate of ith reservoir at time s, { $Y^{i}(t):t \ge 0$ } is the deterministic input process of ith reservoir, Z_{t}^{i} is the storage of the ith reservoir at time t, and C_{i} is the storage capacity of the ith reservoir.

The object is to determine $\{r^{i}(t) | i = 1, 2\}$ in order to

{
$$r^{1}$$
 (t), r^{2} (t)} $\sum_{i=1}^{2} U_{i}(r^{i})$ (2.3)

subject to the constraints

(1)
$$0 \leq Z_t^i \leq C_i$$
 $t \geq 0, i = 1, 2,$ (2.4)

(2) $0 \le r^{i}(t)$ $t \ge 0, i = 1, 2.$ (2.5)

Here U(•) is the damage cost function and \mathbf{r}^{i} is the peak release rate of reservoir i, i.e., $\mathbf{r}^{i} = \max_{t \ge 0} \mathbf{r}^{i}(t)$.

For the parallel system shown in Figure 2.1b, the basic assumptions are as follows.

 The water released from reservoirs 1 and 2 arrives at the damage center at the same time.

- (2) Each reservoir spillway is designed in such a way that it allows for any kind of release rate.
- (3) The releases and inflows occur simultaneously.

The storage equations for the parallel system can be written as:

$$Z_{t}^{1} = Z_{0}^{1} + \int_{0}^{t} Y^{1}(s) ds - \int_{0}^{t} r^{1}(s) ds \qquad t \ge 0,$$
 (2.6)

and

$$Z_{t}^{2} = Z_{o}^{2} + \int_{O}^{t} Y^{2}(s)ds - \int_{O}^{t} r^{2}(s)ds \qquad t \ge 0.$$
 (2.7)

The object is to determine $\{r^{i}(t) | i = 1, 2\}$ in order to

$$\{r^{1} Min \\ (t) + r^{2} (t)\}$$
 (2.8)

subject to the constraints:

(1)
$$0 \leq A_{t}^{i} \leq C_{i}$$
 $t \geq 0, i = 1, 2,$ (2.9)

(3) $0 \le r^{i}(t)$ $t \ge 0, i = 1, 2.$ (2.10)

Here $U(\cdot)$ is the damage cost function at the damage center, r is the peak flow rate at the damage center, and

$$r = \max_{t} \{ (r^{1}(t) + r^{2}(t)) \}.$$

Estimating potential flood damages is an important problem in planning water resources projects. Unfortunately, however, there is a paucity of published data for use by engineers and planners in making such damage estimates. However, it is reasonable to assume that flood damage can be modeled as a monotone nondecreasing function of the depth of flood water [42]. On some streams the stage-discharge curve is roughly parabolic [33], yielding an equation of the type

$$r = d(h - a)^b$$
,

where r is the discharge (cfs),

a, b, d, are constants, and

h is the stage (ft).

Here we can see that the damage cost U is a monotone nondecreasing function of discharge rate.

The exact amount of flood damage may also depend on the time of year, the velocity of flow, the depth and duration of inunduation as well as the nature and quantity of silt deposited. It will be assumed in this study, however, that the critical factor is the peak discharge and the effect of all other factors are of secondary importance and may be neglected. We will also assume that we know or can estimate the damage cost function U(r).

2.3 Continuous Time Process Approximated by Discrete Time Process

The continuous time process can be approximated by a discrete time process. From the basic storage equation:

$$Z(t) = Z(0) + \int_{0}^{t} Y(s)ds - \int_{0}^{t} r(s)ds$$
 (2.12)

where Y(s) is the input rate, we have

$$Z(t) = Y(t) - r(t)$$
 (2.13)

(2.11)

Recall that the object is to determine r(t), which satisfies $Z(t) \ge 0$, $Z(t) \le C$ for $0 \le t \le T$, where T is the duartion of flood hydrograph, so as to minimize the return

$$I = U(max r(t)).$$
 (2.14)

The value of the continuous-time state variable Z(t) at the discrete time

$$t = i\Delta$$
 $i = 0, 1, 2, ...$

can be approximated by a discrete-time variable Z_i satisfying the difference equation

$$Z_{i+1} = Z_i + \frac{dZ(i\Delta)}{dt} \Delta.$$
 (2.15)

This approximation is illustrated on Figure 2.2.

The difference equation can also be written as

$$Z_{i+1} = Z_i + (Y(i\Delta) - r(i\Delta))\Delta.$$
 (2.16)

Then, if \triangle is sufficiently small, we have

$$Z_{i} \approx Z(i\Delta). \qquad (2.17)$$

Now, consider a discrete-time decision process with state variable W_i , and decision variable v_i . The transformation of state is

$$W_{i+1} = W_i + (Y_i - V_i)\Delta.$$
 (2.18)





The object is to select $v_0, \ v_1, \ \cdots \ v_{T/\Delta}$ so as to minimize the return function

$$J = U(\max(v_0, v_1, v_2, ..., v_{T/\Delta}))$$
(2.19)

subject to the constraints

$$W_i \ge 0, \quad W_i \le C$$
 $i = 1, 2, ...$ (2.20)

Notice that the discrete-time decision variable v_i takes the same value as the continuous-time decision variable r(t) at $t = i\Delta$. This discrete time process is such that as Δ goes to zero, the state variable W_i approaches the discrete time variable Z_i . But as Δ goes to zero, the approximation (2.17) also become an equality and so

$$W_i \rightarrow Z(i\Delta)$$
 as $\Delta \rightarrow 0$, (2.21)

and we have

$$\lim_{\Delta \to 0} J = \lim_{\Delta \to 0} (\max(v_0, v_1, v_2, \dots, v_{T/\Delta})) = I$$
(2.22)

Thus, the discrete-time process tends to behave like the continuoustime process as the discrete time interval goes to zero.

The procedure is to divide the time interval into N + 1 stages and piecewise linearize the hydrograph as shown in Figure 2.3. We determine the N + 1 cut points, in such a way that all the critical points of the hydrograph are contained. Then in the interval ΔT_i , the input is

$$Q_i = a_i \Delta T_i + \frac{b_i}{2} (\Delta T_i)^2.$$
 (2.23)

If we denote $\sum_{j=0}^{n} \Delta T_j = T_n$, then the total release in time interval

 ΔT_n is

$$R_{n} = \int_{T_{n-1}}^{T_{n}} r(t) dt = r_{n} \Delta T_{n}.$$

The state transformation is

$$Z_{n+1} - Z_n = Q_n - R_n$$
 $n = 0, 1, 2, ... N$ (2.24)

The objective function is

$$\{(r_0, r_1^{Min}, r_n)\}^{\{\max (U(r_0), U(r_1), \dots U(r_N))\}}$$
(2.25)

and the constraints are

$$0 \le Z_n \le C$$
 $n = 0, 1, ..., N$
 $0 \le R_n$ $n = 0, 1, ..., N$
 $0 \le Q_n$ $n = 0, 1, ..., N$

2.4 Evaluation of Previous Solution Approaches

 $r_n = d_1 \lambda_{1,n} + d_2 \lambda_{2,n} + \dots + d_m \lambda_{m,n}$

Windsor [87] suggested a linear programming method to solve a similar problem using the piecewise linearization of the damage cost function. Consider the following three cases.

(a) Convex U(r)

If U is a convex function of r, as in Figure 2.4a, then we have

n = 0, 1, ..., N (2.26)







$$r = d_1 \lambda_1 + d_2 \lambda_2 + \dots + d_m \lambda_m,$$
 (2.27)

$$U(r) = e_1 \lambda_1 + e_2 \lambda_2 + \dots + e_m \lambda_m.$$
 (2.28)

The problem becomes

$$\{(\lambda_1,\lambda_2,\ldots,\lambda_m)\}_{i}^{\{\sum (e_i\lambda_i)\}}$$
(2.29)

subject to

$$Z_{n+1} - Z_n + (d_1\lambda_{1,n} + d_2\lambda_{2,n} + \dots + d_m\lambda_{m,n})\Delta T = Q_n \quad n = 0, 1, \dots N \quad (2.30)$$
$$d_1\lambda_1 + d_2\lambda_2 + \dots + d_m\lambda_m - (d_1\lambda_{1,n} + \dots + d_m\lambda_{m,n}) \ge 0$$

$$n = 0, 1, \dots N$$
 (2.31)

$$0 \le Z_n \le C$$

n = 0, 1, ...N (2.32)

This is a linear programming problem.

(b) U(r) Concave

If U is a concave function of r, then we cannot use the piecewise linear approximation that was employed in the convex case. For example, if $d_3 < r < d_4$ as Figure 2.5, then r can be a convex combination of d_3 and d_4 . But it can also be a convex combination of d_1 and d_5 or d_2 and d_5 . In the linear programming process, we will find a variable which yields minimum the cost value. Hence in this case, we will follow the straight line joining (d_1, e_1) and (d_5, e_5) which obviously introduces a significant error.



Figure 2.5 Illustration of the error incurred for a concave curve approximated by convex piece-wise linearization.



Figure 2.6 Illustration of convex concave mixed cost function and error introduced by piecewise linearization.

Windsor [88] suggested another method of linearization for this case. As shown in Figure 2.4b, define a set of auxiliary variables $h_{1,n}$, $h_{2,n}$, ..., $h_{m,n}$ and h_1 , h_2 , ..., h_m such that

$$0 \le h_{i,n} \le d_i \lambda_{i,n}$$
 $i = 1, 2, ... m$ (2.33)

$$0 \le h_i \le d_i \lambda_i$$
 $i = 1, 2, ... m$ (2.34)

where

$$\lambda_{i} = 0 \quad \text{or} \quad 1 \quad \text{and}$$

$$\lambda_{i,n} = 0 \quad \text{or} \quad 1.$$

$$(2.35)$$

The cost function becomes

or

$$\sum_{i=1}^{m} s_{i}h_{i} + \sum_{i=1}^{m} y_{i}\lambda_{i}$$
(2.36)

In this linearization, every point is a convex combination of the two nearest points only.

The problem becomes

$$\left\{ \begin{pmatrix} \mathbf{h}_{1}, \mathbf{h}_{2}^{\mathrm{Min}}, \dots \mathbf{h}_{m} \\ (\lambda_{1}, \lambda_{2}^{2}, \dots, \lambda_{m}^{m}) \end{pmatrix} \left\{ \begin{matrix} \sum_{i=1}^{m} S_{i} \mathbf{h}_{i} + \sum_{i=1}^{m} y_{i} \lambda_{i} \\ i = 1 \end{matrix} \right\}$$
(2.37)

subject to

$$Z_{n+1} - Z_n + (h_{1,n} + h_{2,n} + \dots + h_{m,n}) \Delta T = Q_n \quad n = 0, 1, \dots, N$$
 (2.38)

$$\lambda_{1,n} + \lambda_{2,n} + \dots + \lambda_{m,n} = 1 \qquad n = 0, 1, \dots, N \qquad (2.39)$$

$$0 \le h_{i,n} \le d_i \lambda_{i,n} \qquad n = 0, 1, \dots, N \qquad i = 1, 2, \dots, m \qquad (2.40)$$

$$0 \le h_i \le d_i \lambda_i$$
 $i = 1, 2, ..., m$ (2.41)

$$h_1 + h_2 + \dots + h_m - (h_{1,n} + h_{2,n} + \dots + h_{m,n}) \ge 0$$

$$n = 0, 1, \ldots, N$$
 (2.42)

The peak release rate is

$$r = h_1 + h_2 + \dots + h_m$$

This is a mixed-integer programming problem [88]. The solution method has a major shortcoming in that it only can find the optimal peak release rate. But we still do not know how to control the release rate. Another flaw is that the method only be employed to solve the pure concave damage function case. Finally, computational limitations will limit the size of the problems that can be solved. (See [28] and [31]).

(c) U(r) Mixed Convex and Concave

Sometimes we have a mixed convex and concave damage cost function. In the simplest case if the release rate is below some value, say L (flood stage), it will not cause a flood. So we will not incur any damage. The cost function is illustrated as Figure 2.6. For example, if we are at a point x, $0 \le x \le L$, then we want to follow the line of the r-axis. But the mixed-integer programming solution will follow a line which yields the minimum value of the objective function. Then it

will follow the line that has negative cost value. This is an obvious error. Furthermore, there is also another difficulty in that for a release rate between 0 and L, the linear and mixed-integer programming solution will not guarantee choice of the biggest value of release rate which has the same return. But in a flood control problem, we need to ensure that the storage space as large as possible to accommodate the flood flow during the next stage.

In order to overcome these difficulties we propose a solution by dynamic programming.

2.5 Dynamic Programming Approach

2.5.1 Single-Reservoir Flood Control Systems

In the dynamic programming approach we divide the time horizon into N equal time intervals ΔT and linearize the cost function, whether convex, concave, or mixed, by piecewise linearization. We also assume that in each small time interval ΔT , the release rate is constant. The storage at stage n after release (at the beginning of stage n + 1) Z_{n+1} is

$$Z_{n+1} = Z_n + Q_n - r_n \Delta T$$
 (2.43)

This is also the transition function for dynamic programming. Since the constraints are $Z_n \ge 0$, $Z_n \le C$, for all n, the upper bound for r_n when at state Z_n with input Q_n is

$$Ub(Z_n, Q_n) = (Z_n + Q_n)/\Delta T,$$
 (2.44)

and the lower bound of r_n is

$$Lb(Z_n, Q_n) = max\{0, (Z_n + Q_n - C)/\Delta T\}.$$
 (2.45)

Define $f_n(Z_n)$ as the minimum cost from stage n (time nAT) to stage N (time T), if the storage level is Z_n at time nAT (stage n). Then, invoking Bellman's principle of optimality yields the following functional equation of dynamic programming

$$f_{n}(Z_{n}) = \min_{\substack{r_{n} \in \mathbb{R}(Z_{n},Q_{n})}} \{\max(U(r_{n}),f_{n+1}(Z_{n} + Q_{n} - r_{n}\Delta T))\}$$

 $n = 0, 1, \dots, N$ (2.46)

with the boundary condition

$$f_{N+1}(Z_{N+1}) = 0,$$
 (2.47)

where

$$R(Z_n, Q_n) = [max \{0, (Z_n + Q_n - C)/\Delta T\}, (Z_n + Q_n)/\Delta T].$$
 (2.48)

The dynamic programming algorithm was coded in FORTRAN IV and implemented on Northwestern University's CDC 6600. A macro-flow diagram of the algorithm is shown on Figure 2.7 and the FORTRAN program appears in Appendix A1. We note that the digital computer provides a powerful tool for executing the tedius computational work. One value of the tool comes through its rapid and low cost determination of the economic consequences of alternative policies. The decision maker can readily obtain the information he or she needs to predict the



Figure 2.7

Macro-Flow Diagram of the Dynamic Programming Algorithm for a Single-Reservoir Flood Control System. consequences of his or her choices.

We also note that with any discrete model of a process in which the state variables vary continuously there is always a problem associated with deciding on the class interval scheme upon which the discrete version of the process is to be based. This is the problem of how the accuracy of the results is related to both the size of the discrete unit and the geometry of the discretization scheme. The smaller the discrete unit is the closer the discrete approximation approaches the exact continuous situation and so the greater the reliability of the results. However, smaller time units require larger computation time and the solution process may be very expensive, if it is possible at all. In actual processes, the hydrological data are discrete time and discrete value over a considerably large interval. Hence, we need to choose the size of discrete unit based on a compromise between the accuracy desired and the quantity of time and money available for the analysis.

2.5.2 An Example

The Crooked Creek Reservoir on Crooked Creek River, Pennsylvania was constructed by the Pittsburgh flood commission [69]. The horizontal area vs. elevation curve is given on Figure 2.10. The capacity of the reservoir is 3,310,000,000 ft³ and the dam is 918 ft. high. A historical flood hydrograph from [85] is shown in Figure 2.9. When the flood is coming, the initial storage in the reservoir was 840 ft. high and the volume was 220,000,000 ft³. The actual release rule is a function of the water level of the reservoir which is shown on Figure 2.11. The maximum allowable release rate which will not cause



Pittsburgh, Pennsylvania.





a flood is 7,500 cfs. The damage cost function is given as Figure 2.12. The "Standard" release rule [60] is shown on Figure 2.13.

The optimal release rule determined by dynamic programming is displayed on Figure 2.9. From the Figure 2.9, we can find that using optimal release rule the reservoir absorbed the peak flow. Furthermore, the optimal release rule also performed better than both the currently used routing release rule which is as a function of the water levels, and the standard release rule, which is as a function of water storage.

By integrating the area to the left of the curve on Figure 2.10, we can obtain the storage volume when the elevation of the reservoir is z. We can then plot the volume vs. elevation curve S, volume with unit of 6 hr.-sec.-ft., as on Figure 2.15, we can also plot the curve of S - $\frac{Q\Delta t}{2}$ and S + $\frac{Q\Delta t}{2}$ on each side of S and calculate the input of each interval Δt , and then plot I Δt on curve S. In this manner we obtain the elevation of each stage and the release rate of each stage. These computations are summarized in Table 2.1.

The optimal release rule, the currently used routing release rule and standard release rule for operating Crooked Creed Reservoir are shown on Figure 2.9. The initial flood cost is \$5.26 million, the routing rule cost is \$0.864 million, the standard release rule cost is \$3.876 million, and the optimal release rule cost is \$0.758 million. From this we can see that the stationary, Standard Policy, did not absorb the flood peak effectively. The computing time to determine the optimal policies was 122 CPU seconds on Northwestern University's CDC 6600.







Figure 2.11 Routing release rate of Crooked Creek Reservoir.







Inflow Initial storage in any stage





Time	IAt	Pool Elevation	Out flow Rate
	(6hr-sec-ft)	(ft)	(CFS)
$\begin{array}{r} 24\\ 30\\ 36\\ 42\\ 48\\ 54\\ 60\\ 66\\ 72\\ 78\\ 84\\ 90\\ 96\\ 102\\ 108\\ 114\\ 120\\ 126\\ 132\\ 138\\ 144\\ 150\\ 156\\ 162\\ 168\\ \end{array}$	6250 9750 13750 18250 23000 27500 30500 28850 26350 26350 26350 22750 18250 14500 14500 11500 8750 6750 5400 3900 2750 2250 1750 1250 1000 1000 1000	840 841.1 845 853.1 862.5 872.5 881.87 900 906.26 911.3 915 916.25 916.25 916.8 917 916.4 916.3 915.63 915.63 915.63 915.913.32 911.87 910 907.5 904.3 902.5	5020 5020 5510 6067 6870 7606 8137 8798 9225 9553 9796 9965 10021 10025 10028 10023 9993 9965 9880 9824 9740 9592 9450 9356

Table 2.1 Inflow Data, Computions and Routing Outflow Results

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for Crooked Creek Reservoir

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2.5.3 Multiple-Reservoir Flood Control Systems.

Since it is the middle and lower reaches of a stream which require flood protection, flood control reservoirs must be located at the headwaters or upon the upper reaches of the stream; or in the case of larger streams upon the tributaries which contribute to the crest of the flood in the main stream. Then, too, it is upon the tributaries that reservoir sites are most likely to be found which will meet the physical and economic requirements. For this reason, a system of reservoirs, rather than a single reservoir, is necessary to control flood flows, except on small watersheds where a single reservoir at the upper end may suffice [65].

We will first analyze the case of two reservoirs in series and then the case of two reservoirs in parallel. Extention to other system topologies should be immediately apparent.

(a) Two Reservoirs in Series

For two series reservoirs the storage equations can be written

- $Z_{1,n+1} = Z_{1,n} + Q_{1,n} r_{1,n} \Delta T$ n = 0, 1, ..., N (2.49)
- $Z_{2,n+1} = Z_{2,n} + Q_{2,n} + r_{1,n}\Delta T r_{2,n}\Delta T$ n = 0, 1, ..., N (2.50)
 - $0 \le Z_{1,n} \le C_1$ n = 0, 1, ..., N (2.51)

$$0 \le Z_{2,n} \le C_2$$
 $n = 0, 1, ..., N$ (2.52)

where C_1 and C_2 are the capacities of reservoirs 1 and 2, respectively.
The objective function is

$$\begin{cases} {\binom{min}{(r_{1,0},\cdots,r_{1,N})}} \\ {\binom{r_{2,0},\cdots,r_{2,N}}{(r_{2,0},\cdots,r_{2,N})}} \end{cases}^{\{\max[(U_1(r_{1,0})+U_2(r_{2,0})),\cdots,(U_1(r_{1,N})+U_2(r_{2,N}))]\}. \end{cases}$$

(2.53)

We can also solve this problem by dynamic programming. The upper bound of $r_{1,n}$ is

$$Ub_{1,n}(z_{1,n}, Q_{1,n}) = (Z_{1,n} + Q_{1,n})/\Delta T,$$
 (2.54)

and the lower bound of $r_{1,n}$ is

$$Lb_{1,n}(Z_{1,n}, Q_{1,n}) = \max\{0, (Z_{1,n} + Q_{1,n} - C_1)/\Delta T\}.$$
 (2.55)

For a given $r_{1,n}$, the upper bound of $r_{2,n}$ is

$$Ub_{2.n}(Z_{2.n}, Q_{2.n}) = (Z_{2.n} + Q_{2.n} + r_{1,n}\Delta T)/\Delta T.$$
 (2.56)

and the lower bound of r_{2.n} is

$$Lb_{2,n}(Z_{2,n}, Q_{2,n}) = \max \{0, (Z_{2,n} + Q_{2,n} + r_{1,n}\Delta T - C_{2})/\Delta T\}.$$
 (2.57)

.....

The transition functions are

$$Z_{1,n+1} = Z_{1,n} + Q_{1,n} - r_{1,n} \Delta T$$
 $n = 0, 1, ..., N$ (2.58)

$$Z_{2,n+1} = Z_{2,n} + Q_{2,n} + r_{1,n} \Delta T - r_{2,n} \Delta T$$
 $n = 0, 1, ..., N$ (2.59)

Define $f_n(Z_{1,n}, Z_{2,n})$ as the minimum cost from stage n (time nAT) to stage N (time T) if the levels of reservoirs 1 and 2 are $Z_{1,n}$ and $Z_{2,n}$, respectively at time nAT (stage n). Then, invoking Bellman's principle of optimality we obtain the following functional equation of dynamic programming

$$f_{n}(Z_{1,n}, Z_{2,n}) = \min_{\substack{r_{1,n} \in \mathbb{R}_{1,n}}} \left\{ \min_{\substack{r_{2,n} \in \mathbb{R}_{2,n}}} [\max(U(r_{1.n}) + U(r_{2.n}), f_{n+1}(Z_{1,n+1}, Z_{2,n+1}))] \right\}, n = 0, 1, \dots N \quad (2.60)$$

with the boundary condition

$$f_{N+1}(Z_{1,N+1}, Z_{2,N+1}) = 0$$
 (2.61)

in which

 $R_{1,n} = [Lb_{1,n}, Ub_{1,n}],$ (2.62) $R_{2,n} = [Lb_{2,n}, Ub_{2,n}].$

and

A macro flow diagram of the dynamic programming is displayed on Figure 2.16 and the FORTRAN program appears in Appendix A2.

(b) Two Reservoirs in Parallel

For two parallel reservoirs the storage equation can be written

as



Figure 2.16 <u>Macro Plow Diagram of the Dynamic Programming</u> <u>Algorithm for a two Reservoir Series Flood</u> <u>Control System</u>

$$Z_{1,n+1} = Z_{1,n} + Q_{1,n} - r_{1,n}\Delta T \qquad n = 0, 1, ..., N, \quad (2.63)$$

$$Z_{2,n+1} = Z_{2,n} + Q_{2,n} - r_{2,n}\Delta T \qquad n = 0, 1, ..., N, \quad (2.64)$$

$$0 \le Z_{1,n} \le C_1 \qquad n = 0, 1, ..., N, \quad (2.65)$$

$$0 \le Z_{2,n} \le C_2 \qquad n = 0, 1, ..., N, \quad (2.66)$$

where C_1 and C_2 are the capacities of reservoirs 1 and 2, respectively. The objective function is

$$\begin{cases} (r_{1,0}, \dots, r_{1,N}) \\ (r_{2,0}, \dots, r_{2,N}) \end{cases} \begin{cases} \max \left[U(r_{1,0} + r_{2,0}), \dots, U(r_{1,N} + r_{2,N}) \right] \end{cases}$$
(2.67)

We can also solve this problem by dynamic programming. The upper bound of $r_{1,n}$ is

$$Ub_{1,n}(Z_{1,n}, Q_{1,n}) = (Z_{1,n} + Q_{1,n})/\Delta T$$
 (2.68)

and the lower bound of $r_{1,n}$ is

$$Lb_{1,n}(Z_{1,n}, Q_{1,n}) = \max \{0, (Z_{1,n} + Q_{1,n} - C_1)/\Delta T\}.$$
 (2.69)

Likewise, the upper bound of r2,n is

$$Ub_{2,n}(Z_{2,n}, Q_{2,n}) = (Z_{2,n} + Q_{2,n})/\Delta T$$
 (2.70)

and the lower bound of r2,n is

$$Lb_{2,n}(Z_{2,n}, Q_{2,n}) = \max\{0, (Z_{2,n} + Q_{2,n} - C_2)/\Delta T\}.$$
 (2.71)

The transition functions are

$$Z_{1,n+1} = Z_{1,n} + Q_{1,n} - r_{1,n} \Delta T \qquad n = 0, 1, ..., N$$

$$Z_{2,n+1} = Z_{2,n} + Q_{2,n} - r_{2,n} \Delta T \qquad n = 0, 1, ..., N$$
(2.72)

Define $f_n(Z_{1,n}, Z_{2,n})$ as the minimum cost from stage n(time nAT) to stage N (time T) if the levels of reservoirs 1 and 2 are $Z_{1,n}$ and $Z_{2,n}$, respectively at time nAT (stage n). Then invoking Bellman's principle of optimality we obtain the following functional equation of dynamic programming

$$f_{n}(Z_{1,n}, Z_{2,n}) = \min_{\substack{r_{1,n} \in \mathbb{R}_{1,n}}} \left\{ \min_{\substack{r_{2,n} \in \mathbb{R}_{2,n}}} [\max(U(r_{1,n} + r_{2,n}), f_{n+1}(Z_{1,n+1}, Z_{2,n+1}))] \right\} = 0, 1, ... N \quad (2.73)$$

with the boundary condition

$$f_{N+1}(Z_{1,N+1}, Z_{2,N+1}) = 0,$$
 (2.74)

in which

$$R_{1,n} = [Lb_{1,n}, Ub_{1,n}]$$

$$R_{2,n} = [Lb_{2,n}, Ub_{2,n}].$$
(2.75)

A macro flow diagram of the dynamic programming algorithm is displayed

on Figure 2.17 and the FORTRAN program appears in Appendix A3.

2.5.4 Multiple-Reservoir Examples

(a) Two Reservoirs in Series

The Los Angeles County Flood Control District was organized in 1915 to control floods on the Los Angeles, Hondo, and San Gabrial River [69]. There are two reservoirs, San Gabriel No. 2 (or Cogswell), and San Gabriel No. 1, in series on San Gabriel River. The basic information is displayed on Table 2.2 [49]. Suppose that we know the input hydrograph of a storm in history [69] as shown on Table 2.3 and plotted on Figure 2.20. The cost functions for damage areas 1 and 2 are plotted on Figure 2.19. The stage time interval used is $\Delta T = 6$ hrs. The flood stage at damage center 1 is 3500 cfs, and the flood stage at damage center 2 is 5000 cfs. The optimal release rule for Cogswell Reservoir and San Gariel No. 1 Reservoir are shown on Table 2.4. Assuming no reservoirs the initial flood damage cost is \$2.357 million. The optimal flood damage cost is \$0. The standard release rule damage cost is also \$0. Reference to Figure 2.21 shows that these two reservoirs when operated by optimal or standard release rules were completely absorbed the flood peak on San Gabriel River. On this particular example, the standard release rule was as good as the optimal release rule. The reason is that the capacities of these two reservoir flood control system is sufficient to absorb the flood peak of this given hydrograph. The computing time for determining the optimal policy was 581 seconds on Northwestern University's CDC 6600.

(b) Two Reservoir in Parallel

McDonough County, Illinois, lies wholly in the Illinois River basin.











Figure 2.19 Damage cost function of Damage Center 1 and 2.

Reservoir Name	Тур е	High (ft)	Length Crest(ft)	Capacity (acre-ft)	Purpose
Cogswell Reservpir	Rockfill	270	620	10915	Flood control
San Gabriel Reservoir NO	.1 Earthfill	l 285	1500	43928	Flood control

Table 2.2 Basic data of reservoirs on San Gabriel River.

	·	1			84	500	1000
60	300				78	1000	2000
84	720		oir.		72	1500	3000
73	1300		eserv		66	2000	4000
72	1950		iel B		60	2750	5000
66	2500		Gabr		54	3250	5000
60	3100		d San		48	3500	5000
54	3700		ir an	•	745	3500	5000
48	001717		servo		36	3500	5000
747	4300		11 Re	-	30	3500	5000
36	5000		осзме		54	3500	5000
30	0024		of C		18	3000	5000
54	3800		rate		12	1750	3500
18	2400		flow		1 9:	750	5000
12	1400		Input		, 0	3500	5000
9	450		۳. •			Rule 1	Rule 2
0	0		(V (V			1989 1989	ese.
Time (hr)	CFS		Tabl		Tine (hr)	Optir Relea of Re	Optim Relea of Re

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Optimal release rate of San Gabriel River Flood Control System Table 2.4





1:Inflow rate to San Gabriel Reservoir if completely unregulated by any reservoir.
2:Inflow rate to San Gabriel Reservoir after regulated. by Cogswell Reservoir with optimal policy.
3:Optimal release rate of San Gabriel Reservoir.
4:Inflow rate to San Gabriel Reservoir after regulated by Cogswell Reservoir with the standard policy.
5:Standard release rate of San Gabriel Reservoir.

Figure 2.21 Inflow Hydrograph, Optimal Release Rate and Standard Release Rate

At least 90% of the county drains southwestward through the East Fork La Moine River. Several potential reservoirs sites exist in McDonough County. In particular there are two sites: a potential reservoir site on the East Fork La Moine River 4.25 miles west and 3.75 miles north of Bushnell and another on Short Fork a tributary of the East Fork La Moine River. There positions are shown on Figure 2.22 and the basic data is displayed in Table 2.5 [21]. The input hydrograph of a typical flood is displayed in Table 2.6 and shown graphically on Figure 2.23 The flood stage of La Moine River is 2000 cfs. If no reservoirs [79]. exist, the initial flood damage is \$3.04 million. If we consider that these two potential reservoirs are independently operated optimally, the damage cost is \$2.26 million. The release rate of each reservoir and flow rate on damage center of La Moine River when reservoirs operated independently is displayed in Table 2.7 and shown graphically on Figure 2.23. With the whole system of parallel reservoirs is operating optimally, the release rates are as displayed in Table 2.8 and shown graphically on Figure 2.23. The optimal damage cost is \$1.96 million. If we operated these systems with standard policy, the damage cost is \$3.04 million. The release rate is shown graphically on Figure 2.23. In this case the standard policy cannot reduce the peak. This is due to the fact that this policy does not consider the shape of hydrograph and the interaction between the releases of these two reservoirs. The computing time to determine the optimal policy was 511 seconds on Northwestern University's CDC 6600.

From these results, we may conclude that the system of flood control reservoirs should not be operated independently.





					. •		18	650	650	1300
							17	750	750	1500
[•	16	750	750	1500
					- - -	1	15	.750	750	1500
ur pos (lood ontro	lood ontro		•			14	800	006	1700
hed Prai	E Ŭ	с о		ł			13	1100	1400	2500
aters] (sq-1	15.2	6.5	ver.		:		12	1300	2100	3400
at _{We}			ne H1				11	1700	2550	4250
Depth Dam (ft)	39	32	a Moi				10	2200	3000	5200
city ft)	368	040	I UO				ه	2700	3100	5800
Capa (ac		·	volrs				8	3250	2950	6200
Area es)	. 9	4	Reser				2	3500	2600	6100
Pool (acr	33	28	ូ ប្រ ប្រ			· · · · ·	9	3000	2600	5600
Ele.			c Dat	1.2			Ś	5400	1800	4200
llway (ft)	200	200	Bas1				4	1700	1400	3100
Spi	1r	ਸ	2.5				e	1250	1250	2500
	servo to 1	servo te 2	Table				2	800	1000	1800
	9 1 0	940					++	500	r 725	1225
							Time Stage	Inflow at East Folk	Inflow on Short Folk	La Moine

Table 2.6 Inflow Data on La Noine River and it's tributary





Inflow rate and Outflow rate at Damage Center of Ia Moine River under different operating conditions.

ရ ပု စ		~	۳ ا	4	N	9	~	8	, , 6	10	ਜ ਜ	12	13	14	15	16	17	18
volr	1100	1000	1400	2000	2000	2000	2000	1900	1800	1800	1300	1000	1000	1000	1000	1000	1000	1000
voir	1000	1100	1300	1600	2100	2100	2100	2100	2100	2100	2100	1600	1000	1000	1000	1000	1000	1000
the	2100	2700	3600	4100	4100	4100	0004	3900	3900	3900	3400	2600	2000	2000	2000	2000	2000	2000
	able	6.0	Optir	nal re	lease	rate	and	flow	rat e	at De	na <i>k</i> e	Cente	r whe	an eac	h res	ervoi	54	-
			opers	ated 1	Indepe	ndent	-1y.				>							
							3					3+ 	ta				• -	
		∾	m	. 	2	9	2	ω	0	1 0	11	12	13	77	15	16	17	18
rvoir	1500	1000	1400	2000	2700	2900	1800	200	2000	1100	1100	800	1300	1600	1700	1700	200	600
rvoir	1300	1100	1300	1600	1000	800	1900	3500	1700	2400	1900	1300	200	001	300	300	1300	1400
oine er	2800	2100	2700	3600	3700	3700	3700	3700	3700	3500	3000	2100	2000	2000	2000	2000	2000	2000
								÷.,				5 . A. F						
	ľable	2.8	Opt1	mal re ated e	eleas(optime	e rate ally.	e and	flow	rate	st De	າມລຸຮູຣ	Cente	ar who	en the	[ouw e	Le sys	sten B	

CHAPTER III

STOCHASTIC INFLOW

3.1 Introduction

In this chapter, we first formulate a Markov Renewal flood synthesis model in §3.2. The rainfall during each storm and the times of occurrence of heavy storms is modeled as a Markov Renewal process. The rainfall quantity of the first storm of a synthesized flood and the successive occurrence times of synthesized floods are also Markov Renewal Processes. In §3.3 we develop stochastic dynamic programming algorithms to determine optimal release policies for both series and parallel multiple-reservoir flood control systems with stochastic inflows. Examples involving real-world data are provided to illustrate and assess the effectiveness of the solution methods. Finally in §3.4 we simulate the real-time use of optimal operating rules obtained in §3.3 with both historical and simulated Markov model inflow data.

3.2 Stochastic Flood Synthesis Model

In the analysis of floods, it is of interest to determine both the distribution of the times of flood occurrences and the distribution of the duration of each flood.

The relationship between precipitation and runoff is usually complex and is influenced by various storm patterns, antecedents, and basin characteristics. Because of the complexities and the frequent paucity of adequate data, many approximate formulas have been developed to relate rainfall and runoff.

Method of predicting flood-peak discharges and discharge hydrographs from rainfall events have been studied intensively. One method receiving considerable use to estimate peak discharge rates is called the unit hydrograph method. In this method it is assumed that for a given duration of rainfall, the hydrograph time base remains constant. The unit hydrograph is defined [85] as a hydrograph of direct runoff resulting from 1 in. of effective rainfall generated uniformly over the basin area at a uniform rate during a specified period of time or duration. An illustration of the derivation of the unit hydrograph is presented on Figure 3.1 [85].

Once a unit hydrograph has been developed for a basin, it can be used to obtain the surface runoff hydrographs for storm events on the basin. The runoff records can be extended then, for periods in which rainfall was measured but runoff was not. If the unit hydrograph is applied to the maximum probable rain storm for the basin, then the maximum probable flood peak may be obtained.

Application of a unit hydrograph to design rainfall excess amounts other than 1 in. is accomplished by simply multiplying the excess rainfall amount by the unit hydrograph ordinate since the runoff ordinate for a given duration is assumed to be directly proportional to rainfall excess. From the characteristic unit hydrograph of each watershed, we can see that there is a time lag between the flood and the rainfall duration, e.g., a 1 time-unit rainfall may take 10 time units to run through the river basin.

In application to storms of longer or smaller durations the





method of "lagging" often used. It is based on the assumption that the linear response of the watershed is not influenced by previous storms-that is, one can superimpose storms and the results are directly additive. This is illustrated on Figure 3.2. In the 2-hr. unit hydrograph, the total rainfall is 1 inch.

In the hydrograph displayed on Figure 3.3, f_c is the maximum allowable streamflow that will not cause a flood. The portion of the hydrograph Y(t) where Y(t) > f_c will cause a flood. If two storms' hydrograph durations overlap, then they may be grouped. The flood duration is a function of rainfall quantity and the rainfall duration. From the unit hydrograph, we can determine the level of total rainfall \overline{X} such that when the total rainfall of a storm exceeds \overline{X} , it causes a flood. We shall call such storms "heavy storms".

Let the rainfall for each storm in [0, t] be a sequence of random variables I₁, I₂, Now consider all those heavy storms, whose total rainfall, I₁, in [0, t] exceeds X. Denote by N(t) the number of heavy storms in the interval of time [0, t]. N(t) is an integervalued variable. Assume that it satisfies the following conditions: a) If two heavy storms occur in a sufficiently small interval, then it is considered to be a single storm.

b) The number of heavy storms over disjoint intervals are independent.c) The probability of two storms over a time set of measure 0 is equal to zero.

Now, condition (a) says that $t \rightarrow N_t$ increases by jumps of size 1 only; and (b) says that N has independent increments. Therefore [15] (N_t) is a nonstationary Poisson process with some mean function $\Lambda(t)$. By condition (c), the expected number of storms over an interval of



Time

Figure 3.2 Unit hydrograph lagging procedure



Figure 3.3 Flood duration of storms

zero Lebesgue measure is 0; and hence Λ must be absolutely continuous, that is,

$$\Lambda(t) = \int_{0}^{t} \lambda(s) ds \qquad (3.1)$$

for some function λ . Then

$$P\{N_{t+u} - N_t = k\} = \frac{e^{-b}b^k}{k!}, k \in N$$
 (3.2)

where

$$b = \Lambda(t + u) - \Lambda(t) = \int_{t}^{t+u} \lambda(s) ds. \qquad (3.3)$$

In particular, if we consider only one flood season, $\lambda(s) = \lambda$ independent of s, then N becomes a stationary Poisson process. But in general, λ should depend on s, and show a periodicity of one year.

Hence the time between each heavy storms are independent if and only if $\lambda(s)$ = constant, otherwise, if λ varies, time between storms are still exponential but not independent of each other [15]. The cumulative distribution function is

$$\Phi(t) = 1 - e^{-\lambda t}$$
 (3.4)

Let X_0 , X_1 , ... be the quantity of the successive heavy rainfalls and T_n be the time of the nth storm. Assume the successive rainfalls form a Markov Chain. Then for each $n \in N$, the X_n is a random variable taking value in a countable set δ (the set of all possible rainfalls) and T_n is a random variable taking values in $R^+ = [0, +\infty)$ with

$$0 \le T_0 \le T_1 \le T_2 \dots$$
 (3.5)

Assume that $\{X_n, n = 0, 1, ...\}$ is a Markov Chainwith transition matrix G(i, j). Then $\{X_n, T_n; n = 0, 1, ...\}$ is a Markov Renewal Process [15] with Semi-Markov Kernel over δ as

For each heavy storm the quantity of rainfall and duration of rainfall are both random variables. Hence from the unit hydrograph, the flood duration of each heavy storm is also a random variable. It is dependent on both the rainfall quantity and the rainfall duration, but is independent of the time which floods occur. For each fixed rainfall quantity k, the distribution of the flood duration of individual storms' rainfalls are connected together, we consider this a single storm. Hence we can assume that the time between storms is longer than the time duration of its rainfall duration. The Stochastic Flood Model is described below.

Let τ_0 , τ_1 , τ_2 , ... be the flood durations associated with the rainfalls occurring at T_0 , T_1 , T_2 , When a heavy storm with rainfall quantity i comes, it causes a flood with a time duration τ_1 , a random variable with distribution $\Psi(i, \cdot)$ independent of the time the flood occurred. If another heavy storm with rainfall quantity j arrives during the previous flood period, then it will cause a flood with τ_2 units of time; τ_2 is random variable with distribution $\Psi(j, \cdot)$. Let Y_t be the quantity of rainfall of a heavy storm whose flood period covers time t. Let L be the length of the first possible synthesized flood. Assume that at time 0, there is a heavy storm of

rainfall i, which causes a flood. The probability that the flood started at 0 is still going on at time t and the last rainfall before t was of amount k is

$$f(i, t) = P_i \{Y_t = k, L > t\}.$$
 (3.7)

Now,

$$f(i, t) = P(Y_t = k, L > t, T_1 > t | X_o = i) + P(Y_t = k, L > t, T_1 \le t | X_o = i)$$
 (3.8)

From Figure 3.4a, we have

$$P(Y_{t} = k, L > t, T_{1} > t | X_{0} = i)$$

= I(i, k) [1 - \sum Q(i, j, t)](1 - \y(i, t)) (3.9)

From Figure 3.4b, we have

$$P(Y_{t} = k, L > t, T_{1} < t | X_{0} = i)$$

$$= E[P(Y_{t} = k, L > t, T_{1} < t | X_{0} = i, X_{1}, T_{1})$$

$$= E[P(Y_{t} = k, \tilde{L} > t - T_{1} | X_{0} = i, X_{1}, T_{1})$$

$$= \sum_{j} \int_{0}^{t} Q(i, j, ds)f(j, t - s). \qquad (3.10)$$



Hence, we have

$$f(i, t) = g(i, t) + Q*f(i, t)$$
 (3.11)

$$g(i, t) = I(i, k) [1 - \sum_{j} Q(i, j, t)] (1 - \Psi(i, t))$$

= I(i, k)e^{- λt} (1 - $\Psi(i, t)$) (3.12)

This is a Markov Renewal equation [15]. From Markov renewal theory, the solution of (3.11) is

$$f = R*g$$
 (3.13)

S.

Where R is the Markov renewal function corresponding to Q. Hence,

$$f(i, t) = \sum_{\substack{\ell \in \mathcal{S} \\ 0}} \int_{0}^{t} R(i, \ell, ds)g(\ell, t - s)$$

$$= \sum_{\substack{\ell \in \mathcal{S} \\ 0}} \int_{0}^{t} R(i, \ell, ds)I(\ell, k) [1 - \sum_{j} Q(\ell, j, t - s)](1 - \Psi(\ell, t - s))$$

$$= \sum_{\substack{\ell \in \mathcal{S} \\ 0}} \int_{0}^{t} R(i, \ell, ds)I(\ell, k)e^{-\lambda(t-s)}(1 - \Psi(\ell, t - s))$$

$$= \int_{0}^{t} R(i, k, ds)e^{-\lambda(t-s)}(1 - \Psi(k, t - s)). \quad (3.14)$$

.).***

Also, since

$$Q(i, j, t) = G(i, j)(1 - e^{-\lambda t}),$$
 (3.15)

R must have a special form; namely,

$$R(i, j, t) = \sum_{n=0}^{\infty} \frac{e^{-\lambda t} (\lambda t)^{n}}{n!} \sum_{m=0}^{n} G^{m}(i, j)$$
$$= \sum_{m=0}^{\infty} G^{m}(i, j) \sum_{n=m}^{\infty} \frac{e^{-\lambda t} (\lambda t)^{n}}{n!} . \qquad (3.16)$$

Hence,

$$R(i, j, dt) = \sum_{m=1}^{\infty} G^{m}(i, j) \frac{\lambda e^{-\lambda t} (\lambda t)^{m-1}}{(m-1)!} + I(i, j) \epsilon_{o}(dt). \quad (3.17)$$

Now,

$$\int_{0}^{t} R(i, k, ds) e^{-\lambda(t-s)} [1 - \Psi(k, t-s)]$$

$$= I(i, k) e^{-\lambda t} [1 - \Psi(k, t)]$$

$$+ \sum_{m=1}^{\infty} G^{m}(i, k) \int_{0}^{t} \frac{\lambda e^{-\lambda s} (\lambda s)^{m-1}}{(m-1)!} e^{-\lambda(t-s)} [1 - \Psi(k, t-s)]. \quad (3.18)$$

We next consider the process of the synthesized flood. Let $\hat{T} = 0$, \hat{T}_1 , \hat{T}_2 , ... be the time of successive occurences of synthesized floods and let \hat{X}_0 , \hat{X}_1 , \hat{X}_2 , ... be the rainfall quantities of the first storms of each successive synthesized flood. Since (X, T) is a Markov renewal process, we have

$$P(X_{n+1} = j, T_{n+1} - T_n \le t | X_0, \dots X_n; T_0, \dots T_n)$$

= $P(X_{n+1} = j, T_{n+1} - T_n \le t | X_n).$ (3.24)

Now

$$P(\hat{X}_{n+1} = j, \hat{T}_{n+1} - \hat{T}_{n} \le t | \hat{X}_{o}, \dots \hat{X}_{n}; \hat{T}_{o}, \dots \hat{T}_{n})$$

= $P(X_{m} = j, T_{m} - T_{\ell} \le t | X_{o}, \dots X_{\ell}; T_{o}, \dots T_{\ell})$ (3.25)

for some m, l, m > l, where m and l are stopping times of the (X, T) process and (X, T) is Markov renewal process, By the strong Markov property, we obtain

$$P(X_{m} = j, T_{m} - T_{\ell} \le t | X_{0}, \dots X_{\ell}; T_{0}, \dots T_{\ell})$$

$$= P(X_{m} = j, T_{m} - T_{\ell} \le t | X_{\ell})$$

$$= P(\hat{X}_{n+1} = j, \hat{T}_{n+1} - \hat{T}_{n} \le t | \hat{X}_{n}) \qquad (3.26)$$

Thus (\hat{X}, T) is also Markov renewal process. Let \hat{Q} be the semi-Markov kernel corresponding to (\hat{X}, T) . Let

$$h(i, t) = \hat{Q}(i, k, t)$$

= $P(\hat{X}_{T_1} = k, \hat{T}_1 - \hat{T}_0 \le t | \hat{X}_{T_0} = i)$ (3.27)

be the probability that the next flood occurs before time t and is caused by a storm with rainfall k given that at time 0 the synthesized flood occurs as a result of storm with rainfall i.

$$h(i, t) = P(\hat{x}_{T_{1}} = k, \hat{T}_{1} - \hat{T}_{o} \le t | \hat{x}_{T_{o}} = i)$$

= $P(\hat{x}_{T_{1}} = k, \hat{T}_{1} - \hat{T}_{o} \le t, \hat{T}_{1} = T_{1} | \hat{x}_{T_{o}} = i) + P(\hat{x}_{T_{1}} = k, \hat{T}_{1} - \hat{T}_{o} \le t, \hat{T}_{1} > T_{1} | \hat{x}_{T_{o}} = i)$
(3.28)

From Figure 3.5a, we obtain

$$P(\hat{X}_{\hat{T}_{1}} - \hat{T}_{1} - \hat{T}_{0} \le t, \hat{T}_{1} = T_{1} | \hat{X}_{\hat{T}_{0}} = i)$$

$$= P(X_{T_{1}} = k, T_{1} - T_{0} \le t | X_{T_{0}} = i)$$

$$= \int_{0}^{t} Q(i, k, ds) \Psi(i, s), \qquad (3.29)$$

from Figure 3.5b, we obtain

$$P(\hat{X}_{T_{1}} = k, \hat{T}_{1} - \hat{T}_{0} \le t, \hat{T}_{1} | \hat{X}_{T_{0}} = i)$$

$$= E[P(\hat{X}_{T_{1}} = k, \hat{T}_{1} - \hat{T}_{0} \le t, \hat{T}_{1} > T_{1} | \hat{X}_{T_{0}} = i, X_{T_{1}}, T_{1})]$$

$$= E[P(\hat{X}_{T_{1}} = k, \hat{T}_{1} - T_{1} \le t - T_{1}, \hat{T}_{1} > T_{1} | \hat{X}_{T_{0}} = i, X_{T_{1}} = j, T_{1})]$$

$$= \sum_{j=0}^{t} \int_{0}^{t} Q(i, j, ds)(1 - \Psi(i, s))h(j, t - s). \qquad (3.30)$$

Hence, we have

h(i, t) = g(i, t) +
$$\sum_{j=0}^{t} \overline{Q}(i, j, ds) f(j, t - s)$$
 (3.31)

where

$$\overline{Q}(i, j, ds) = Q(i, j, ds)(1 - \Psi(i, s))$$
 (3.32)

(3.33)

$$g(i, t) = \int_{0}^{t} Q(i, k, ds) \Psi(i, s)$$

= $-\int_{0}^{t} Q(i, k, ds) (1 - \Psi(i, s)) + \int_{0}^{t} Q(i, k, ds)$
= $Q(i, k, t) - \overline{Q}(i, k, t).$

From Markov renewal theory, the solution of (3.31) is

$$h(i, t) = \overline{R*g(i, t)}$$

where $\overline{R} = \sum_{n} \overline{Q}^{n}$ is the Markov renewal function corresponding to \overline{Q} and n

$$h(i, t) = \sum_{j} \int_{0}^{t} \overline{R}(i, j, ds)g(j, t - s)$$

= $\sum_{j} \int_{0}^{t} \overline{R}(i, j, ds)[Q(j, k, t - s) - \overline{Q}(j, k, t - s)].$ (3.3)

Hence the probability that next flood occurs before time t and is caused by storm with rainfall k given that at time 0, the synthesized flood occurs as a result of storm with rainfall i is

$$\hat{Q}(i,k,t) = \sum_{j=0}^{t} \overline{R}(i,j,ds) [Q(j,k,t-s) - \overline{Q}(j,k,t-s)]. \quad (3.34)$$

3.3 Stochastic Decision Process

A flood control reservoir differs from a storage reservoir in that the goal is to keep the reservoir as nearly empty as possible, rather than as full as possible. Thus there is involved in the operation of a flood control reservoir a trade-off between possible downstream flood damage and the risk of subsequent flood damage due to storing water in the flood control space. Askew [2], Eastman [22], Eisel [23], Morton [46], Nayak [49, 50], have all described methods for optimizing the operation of a storage reservoir minimizing a lose function subject to chance constraints. But in flood control, we can not use the target constraints on the release rate and chance constraints on the storage volume. The constraints on flood control, such as the storage constraints and the release constraints are physical constraints and cannot be violated. The cost function in flood control depends on peak discharge rate, but not on the size of missing the target. We will describe a stochastic dynamic programming method for finding the optimal release rule in the following.

Assume that the water inflow to the reservoir is a stochastic process defined on [0, T]. At time t, the input rate Y(t) is a random variable. Then the problem is to find $\{r(t) | t \in [0, T]\}$ so as to

$$\begin{array}{l} \text{Min} \{ \max E[U(r(t))] \} \\ \{r(t)\} t \end{array}$$
(3.35)

As in Chapter II the storage equation is

$$Z_{t} = Z_{0} + \int_{0}^{t} Y(s) ds - \int_{0}^{t} r(s) ds$$

= $Z_{0} + Q(t) - \int_{0}^{t} r(s) ds$, (3.36)

in which Z_t is the storage at time t, $0 \le Z_t \le C$, where C is capacity of reservoir, and Q_t is the total input from time 0 to time t.

We will use a discrete time procedure to approximate Y(t). Divide the time interval [0, T] of interest into N subdivisions, each stage of time duration ΔT . Then the storage equation becomes

$$Z_n = Z_0 + \sum_{i=0}^{n} Q_i - \sum_{i=1}^{n} r_i \Delta T$$
 $n = 0, 1, ..., N$ (3.37)

$$D \le Z_n \le C$$
 $n = 0, 1, ..., N$ (3.38)

For a given cost function U(r), where r is the peak discharge rate, the objective function can be written as

$$\min \{\max(EU(r_0), EU(r_1), \dots, EU(r_N))\}.$$
(3.39)
$$\{(r_0, r_1, \dots, r_n)\}$$

The transition of state, however, depends not only on the current state Z_n , the decision r_n but also on a random variable Q_n . The state transition function is

$$Z_{n+1} = Z_n + Q_n - r_n \Delta T$$
 $n = 0, 1, ..., N$ (3.40)

3.3.1 Time Dependent Inflow Rate

Let $P_n(q_n)$ be the probability density of the random variable Q_n , n = 0, 1, 2, ..., N. Define $f_n(Z_n)$ as the optimal expected value with respect to Q_n of the total cost from the n^{th} stage to the end of process, given that the process is in state Z_n at stage n. Then, invoking Bellman's principle of optimality we obtain the following functional equation of dynamic programming

$$f_{n}(Z_{n}) = \underset{r_{n} \in \mathbb{R}_{n}(Z_{n},Q_{n})}{\min} \{ \max(\mathrm{EU}(r_{n}), \mathrm{Ef}_{n+1}(Z_{n}+Q_{n}-r_{n}\Delta T)) \}$$

$$n = 0, 1, \ldots, N$$
 (3.41)

with the boundary condition

$$f_{N+1}(Z_{N+1}) = 0$$
 (3.42)

For fixed Q_n , the lower bound of r_n is

$$Lb_n(Z_n, Q_n) = max\{0, (Z_n + Q_n - C)/\Delta T\},$$
 (3.43)

The upper bound of r is

$$Ub_n(Z_n, Q_n) = (Z_n + Q_n)/\Delta T.$$
 (3.44)

Since Q_n is a random variable, the interval [Lb_n, Ub_n] is a random interval on R^+ .

The physical meaning of Ub_n is that all the water is released in period n and there is no more water in the reservoir. Hence, when the decision is r_n where $r_n > Ub_n$ the actual release is only equal to Ub_n . Likewise, the physical meaning of Lb_n is that the release is only the overflow part of water and the reservoir is full. The release rate cannot be less than this quantity. Hence, when the decision is r_n where $r_n < Lb_n$, the actual release is Lb_n .

$$\mathbf{r}_{n} = \begin{cases} \mathbf{Lb}_{n} & \text{if } \mathbf{r}_{n} < \mathbf{Lb}_{n} \\ \mathbf{r}_{n} & \text{if } \mathbf{r}_{n} \in [\mathbf{Lb}_{n}, \mathbf{Ub}_{n}] \\ \mathbf{Ub}_{n} & \text{if } \mathbf{r}_{n} > \mathbf{Ub}_{n} \end{cases}$$
(3.45)

The range of r is

$$R_n(Z_n, Q_n) = [Min \ Lb_n(Z_n, Q_n), \max \ Ub_n(Z_n, Q_n)]. \qquad (3.46)$$

$$Q_n \qquad Q_n$$

A macro-flow diagram of the dynamic programming algorithm is given on Figure 3.6. The algorithm was coded in FORTRAN IV and implemented on Northwestern University's CDC 6600. The program listing appears in Appendix A4.

3.3.2 Individual Storm

Each watershed area of a reservoir has its own characteristic unit hydrograph. Let Y'(t) be the ordinate of unit hydrograph at time t and M be the total rainfall excess of a specific storm. The inflow hydrograph of that storm is

$$Y(t) = Y'(t)M$$
. (3.47)

In practical cases, the total rainfall of a storm is probabilistic. Suppose that we know its probability density function p(m), then in the functional recurrence equation

$$f_n(Z_n) = \min\{\max(EU(r_n), Ef_{n+1}(Z_n + Q_n - r_n \Delta T))\}$$

$$n = 0, 1, ..., N$$
 (3.48)

the distributions of input at each stage are no longer independent of each other; they depend on the previous stage. Hence the functional



Figure 3.6 Macro Flow Diagram of the Stochastic Dynamic Programming Algorithm for Single Reservoir with Stochastic Inputs.
equation of dynamic programming changes to

$$f_{n}(Z_{n}) = \min\{E[\max(U(r_{n}), g_{n+1}(Z_{n}+Q_{n}-r_{n}\Delta T, m))]\} \ n = 0, 1, ... N$$
(3.49)
$$r_{n}^{m}$$

in which $g_{n+1}(Z_{n+1}, m)$ is the optimal return from the deterministic process when N - n - 1 stages are left, and the total input rainfall is m, where m is a random variable.

$$g_n(Z_n, m) = \min\{\max(U(r_n), g_{n+1}(Z_n + Q_n - r_n \Delta T, m))\} n = 0, 1, ... N$$
 (3.50)
 r_n

$$f_{N+1}(Z_{N+1}) = 0 \tag{3.51}$$

 $g_{N+1}(Z_{N+1}) = 0$ (3.52)

For fixed Q_n , the lower bound of r_n is

$$Lb_n(Z_n, Q_n) = max\{0, (Z_n + Q_n - C)/\Delta T\},$$
 (3.53)

and the upper bound is

$$Ub_n(Z_n, Q_n) = (Z_n + Q_n) / \Delta T.$$
 (3.54)

Since $Y_n = Y_n^{\dagger} \cdot m$; m is random variable, hence the interval $[Lb_n, Ub_n]$ is a random interval on R^{\dagger} . We have

$$\mathbf{r}_{n} = \begin{cases} \mathbf{Lb}_{n} & \text{if} & \mathbf{r}_{n} < \mathbf{Lb}_{n} \\ \mathbf{r}_{n} & \text{if} & \mathbf{r}_{n} \in [\mathbf{Lb}_{n}, \mathbf{Ub}_{n}] \\ \mathbf{Ub}_{n} & \text{if} & \mathbf{r}_{n} > \mathbf{Ub}_{n}. \end{cases}$$
(3.55)

3.3.3 Multiple-Reservoir Flood Control Systems

For the case of a system of two series reservoirs, assume that at time t the input rates $Y_1(t)$ and $Y_2(t)$ are both random variables. Then the problem is

$$\underset{\{r_1(t),r_2(t)\}}{\min} \{ \underset{t}{\max(EU_1(r_1(t)) + EU_2(r_2(t)))} \}.$$
(3.56)

The storage equations are

$$Z_{1}(t) = Z_{1}(0) + \int_{0}^{t} Y_{1}(s)ds - \int_{0}^{t} r_{1}(s)ds,$$
 (3.57)

$$Z_{2}(t) = Z_{2}(0) + \int_{0}^{t} Y_{2}(s) ds + \int_{0}^{t} r_{1}(s) ds - \int_{0}^{t} r_{2}(s) ds.$$
(3.58)

Dividing the time interval into N stages, such that each stage has time duration ΔT , the storage equations become

$$z_{1,n} = z_{1,0} + \sum_{j=0}^{n} q_{1,j} - \sum_{j=0}^{n} r_{1,j} \Delta T$$
 $n = 0, 1, ..., N$ (3.59)

$$Z_{2,n} = Z_{2,0} + \sum_{j=0}^{n} Q_{2,j} + \sum_{j=0}^{n} r_{1,j} \Delta T - \sum_{j=0}^{n} r_{2,j} \Delta T \quad n = 0, 1, ..., N$$
 (3.60)

The objective function can be written as

$$\{ \begin{pmatrix} r_{1,0}, \dots, r_{1,N} \\ (r_{2,0}, \dots, r_{2,N}) \end{pmatrix} \}^{ \{ \max_{n} (E[U(r_{1,n}) + U(r_{2,n})]) \}, } (3.61)$$

Define $f_n(Z_{1,n}, Z_{2,n})$ as the optimal expected value with respect

to $Q_{1,n}$ and $Q_{2,n}$ of the total cost from the nth stage to the end of process, given that the process is in states $Z_{1,n}$ and $Z_{2,n}$ at stage n. Then invoking Bellman's principle of optimality we obtain the following functional equation of dynamic programming

$$f_{n}(Z_{1,n}, Z_{2,n}) = \min_{\substack{r_{1,n} \in \mathbb{R}_{1,n} \\ r_{2,n} \in \mathbb{R}_{2,n} \\ Ef_{n+1}(Z_{1,n+1}, Z_{2,n+1}))}} \min_{\substack{r_{2,n} \in \mathbb{R}_{2,n} \\ n = 0, 1, ..., N}} u_{2}(r_{2,n})],$$
(3.62)

with the boundary condition

$$f_{N+1}(Z_{1,N+1}, Z_{2,N+1}) = 0.$$

The transition functions are

$$Z_{1,n+1} = Z_{1,n} + Q_{1,n} - r_{1,n} \Delta T$$
 $n = 0, 1, ..., N$ (3.63)

$$Z_{2,n+1} = Z_{2,n} + Q_{2,n} + r_{1,n}\Delta T - r_{2,n}\Delta T$$
 $n = 0, 1, ..., N$ (3.64)

in which $Q_{i,n}$ is the input to reservoir i at stage n, a random variable with probability density function $p_i(Q)$.

The decision sets are

$$R_{1,n} = [\max\{0, (Z_{1,n} + Q_{1,n} - C_1)/\Delta T\}, (Z_{1,n} + Q_{1,n})/\Delta T],$$

$$R_{2,n} = [\max\{0, (Z_{2,n} + Q_{2,n} + r_{1,n}\Delta T - C_2)/\Delta T\}, (Z_{2,n} + Q_{2,n} + r_{1,n}\Delta T)/\Delta T].$$
(3.65)

In the case of a system of two parallel reservoirs, assume that at time t, the inflow rates $Y_1(t)$ and $Y_2(t)$ are both random variables. The objective function is

$$\{r_{1}(t), r_{2}(t)\} \{\max_{t} (EU(r_{1}(t) + r_{2}(t)))\}.$$
(3.66)

The storage equations are

$$Z_{1}(t) = Z_{1}(0) + \int_{0}^{t} Y_{1}(s)ds - \int_{0}^{t} r_{1}(s)ds, \qquad (3.67)$$

$$Z_{2}(t) = Z_{2}(0) + \int_{0}^{t} Y_{2}(s)ds - \int_{0}^{t} r_{2}(s)ds. \qquad (3.68)$$

Dividing the time interval into N stages, such that each stage has time duration ΔT , the storage equations become

$$Z_{1,n} = Z_{1,0} + \sum_{j=0}^{n} Q_{1,j} - \sum_{j=0}^{n} r_{1,j} \Delta T$$
 $n = 0, 1, ..., N$ (3.69)

$$Z_{2,n} = Z_{2,0} + \sum_{j=0}^{n} Q_{2,j} - \sum_{j=0}^{n} r_{2,j} \Delta T$$
 $n = 0, 1, ..., N$ (3.70)

$$0 \le Z_{i,n} \le C_i$$
 $i = 1, 2; n = 0, 1, ..., N$ (3.71)

The objective function can be written as

$$\{ \begin{pmatrix} r_{1,0}, \dots, r_{1,N} \\ (r_{2,0}, \dots, r_{2,N}) \end{pmatrix} \{ max (E[U(r_{1,n} + r_{2,n})]) \}.$$
(3.72)

Define $f_n(Z_{1,n}, Z_{2,n})$ as the optimal expected value with respect to $Q_{1,n}$ and $Q_{2,n}$ of the total cost from the nth stage to the end of

process, given that the process is in states $Z_{1,n}$ and $Z_{2,n}$ at stage n. Then invoking Bellman's principle of optimality we obtain the following functional equation of dynamic programming

$$f_{n}(Z_{1,n}, Z_{2,n}) = Min \{ Min (max(EU(r_{1,n} + r_{2,n})), r_{1,n} \in R_{1,n}, r_{2,n} \in R_{2,n} \}$$

$$Ef_{n+1}(Z_{1,n+1}, Z_{2,n+1})))$$
 n = 0, 1, ..., N (3.73)

with the boundary condition

$$f_{N+1}(Z_{1,N+1}, Z_{2,N+1}) = 0.$$

The transition functions are

$$Z_{1,n+1} = Z_{1,n} + Q_{1,n} - r_{1,n} \Delta T$$
 $n = 0, 1, ..., N$ (3.74)

$$Z_{2,n+1} = Z_{2,n} + Q_{2,n} - r_{2,n} \Delta T$$
 $n = 0, 1, ..., N$ (3.75)

The decision sets are

$$R_{1,n} = [\max\{0, (Z_{1,n} + Q_{1,n} - C_1)/\Delta T\}, (Z_{1,n} + Q_{1,n})/\Delta T],$$
 (3.76)

$$R_{2,n} = [\max[0, (Z_{2,n} + Q_{2,n} - C_2)/\Delta T], (Z_{2,n} + Q_{2,n})/\Delta T].$$
 (3.77)

3.3.4 Example

Shelbyville Reservoir is situated between Shelby County and Moultrie County, Illinois, as show on Figure 3.7 [21]. The basic data are given on Table 3.1. The daily flow rate of the Kaskaskia

. 2 . * * *



Figure 3.7 Shelbyville Reservoir, Shelby County, Illinois.

River on Shelbyville is given on Appendix B[79]. The cumulative distribution function of the stream flow rate derived from these data are presented on Table 3.2.

In this example, we will consider three kinds of flow process to approximate the actual stream process. They are (1) individual storm, (2) time independent flow rate for each stage, and (3) time dependent flow rate.

A. Individual Storm

Recall that each watershed area has its characteristic unit hydrograph. Hence, we will assume that the unit hydrograph of Shelbyville Reservoir watershed is known. The total runoff is a stochastic process. From the hydrograph of the storm on March 1911, the total runoff is 1.14 inch. The hydrograph and unit hydrograph data is present on Table 3.3 and the initial storage is 50,000,000 ft.³.

Because of our emphasis on flood control, in determining the distribution of total runoff of each storm we only consider the runoff data of some rainy season. For our example, there are December, January, February, March, April, May and June. The summer and fall seasons are relatively dry. There are a total 35 data points 0.85 3.37 2.08 5.29 0.62 0.03 0.05 1.67 1.22 4.60 2.17 1.89 2.46 0.66 2.14 0.69 1.36 0.30 2.06 0.53 0.57 1.84 1.42 1.14 2.63 0.74 0.11 0.98 0.95 1.30 4.59 3.20 2.08 0.48 0.07

The runoff can be modeled as a Gamma distribution with $\alpha = 1.25$, r = 2 (the test of Goodness of Fit by Chi-square Test and Komogorov Test is on Appendix Cl). The expected value is 2/1.25 = 1.6 in. The result of optimal release policy for the individual storm case is

NAME	OWNER	WATERSHED AREA (sq.m1)	HIGHT OF DAM (ft)	POOL AREA (acre)	STORAGE CAPACITY
SHELBYVILLE RESERVOIR	U.S.COAPS OF ENGINEERS	1030	65	10000	(ac-ft) 22960

Table 3.1 Easte data of Shelbyville Reservoir.

於 0	10600
58	4600
10%	3150
20%	1740
30%	1070
×017	720
50%	420
¢0%	238
20%	123
80%	61
×06	26
95%	17
min cfs	0.2
mo.x cfs	10600
Feriod of record	4.8 year

Table 3.2 Cumulative distribution of stream flow rate.

21	666	0
20	702	31
19	738	63
18	810	126
17	918	264
16	1030	319
15	1150	ヤこヤ
14	1450	687
13	1680	889
12	1910	1091
1	2050	1214
10	2300	1443
6	2450	1564
ο ΟΟ	2400	1521
~	1360	608
Kar 6	600	0
(date)	arte of graph	ate of graph
EMIT	Ord1r Eydro	Ordir Unit

Table 3.3 Ordinate of hydrograph of Shelbyville Reservoir.

present on Table 3.4. Reference to Table 3.4 reveals that the optimal release is 1500 cfs (maximum nonflood release rate) at most stages. This is because in most cases, the capacity of the reservoir is sufficient to handle any single storm flood. The initial expected damage cost is \$2.385 million. For standard release rule, the expected damage cost is \$1.92 million. We cannot determine the expected damage cost of the routing policy since we are lacking both the storage vs. elevation and the outflow vs elevation curves. The optimal expected damage cost is \$0.5898 million. The computing time to determine the optimal policy was 158 CPU second on Northwestern University's CDC 6600.

B. Time Independent Inflow Rate

Suppose that at any time interval the flow rate is distributed as in Table 3.3. Here for a two month period (12 stages), with five days for each stage interval, the optimal release rule is as shown on Table 3.5. When the initial storage is half full the optimal expected cost is \$0.56 million.

C. Time Dependent Inflow Rate

A set of historical or synthetic flow data for a stream is a sequence of numbers or values produced by a random process in a succession of time intervals; such a sequence is called a time series [60]. In general, the ith member of a time series which we write X_i, is the sum of two parts [61].

$$X_i = d_i + e_i$$

(3.78)

Here d_i is the deterministic part. Typically, d_i might be a function

14	
13	анананананананана 84 <i>мимимимимимимими</i> 7700000000000000000000000000000000000
12	44444444444444444444444444444444444444
11	44444444444444444444444444444444444444
10	44444444444444444444444444444444444444
σ	44444444444444444444444444444444444444
ω	44444444444444444444444444444444444444
2	44444444444444444444444444444444444444
\$	44444444444444444444444444444444444444
Ń	
4	++++++++++++++++++++++++++++++++++++++
Ϋ́	44444444444444444444444444444444444444
2	++++++++++++++++++++++++++++++++++++++
ᠳ	1500
FS STATE	L 000000000000000000000000000000000000

Optimal release rule of individual storm of Shelbyville Res. Table 3.4

Optimal Expect Cost	00000000000000000000000000000000000000	
12		
1-1 1-1	44444444444444444444444444444444444444	
10	44444444444444444444444444444444444444	
σ	44444444444444444444444444444444444444	
со -	44444444444444444444444444444444444444	
2	××××××××××××××××××××××××××××××××××××××	
9	x x + + + + + + + + + + + + + + + + + +	
Ń	88844444444444444444 4408887770000000000	
4	88844444444444444444444444444444444444	
m :	22222222222222222222222222222222222222	
2	22222222222222222222222222222222222222	
÷	22222222222222222222222222222222222222	
STAGE 10 ⁰ Cfs STATE	н н н н н н н н н н н н н н	

Optimal release rate with time independent inflow rate. Table 3.5

Sec. and

of the mean flow, of the variability of flow, (as measured by their standard deviation), and of previous flow such as X_{i-1}, X_{i-2}, \dots . The random component of the generation scheme is e_i . It is a random number drawn or sampled from a set of random numbers with a certain probability distribution or pattern. In this discussion, we will assume that d_1 depends on stage i, but not on the previous flows. In the Markovian case, we will use the limiting probability of the flow rate distribution.

We assume the average flow rate at each stage of the same month is constant. These values are presented on Table 3.6. The data are not sufficient to test the distribution of e_i ; hence, we will assume the distribution of e_i is a Normal distribution [61] with mean 0 and variance as the sample variance. The problem was solved by the dynamic programming algorithm presented in Figure 3.6. The release policies for the two month periods of January-February, March-April, and May-June are presented on Table 3.7, 3.8, 3.9 respectively. The optimal expected damage cost for these three two-month periods are \$1.93, \$3.6, and \$2.15 million respectively. The computing time for finding the optimal policy of each two-month period was 211 CPU seconds on Northwestern University's CDC 6600.

3.4 Simulation

In this section we simulate the real-time operation of flood control systems using the optimal release policies determined in §3.2 and §3.3. In this manner we can assess the feasibility and practicality of using calculated optimal release rules in real time.

Digital computer simulation is a popular and powerful tool for using

Nov. Dec.	485 415	252 213	
Oct.	839	438	
ុក ខ ខ	しわわ	229	
Aug.	166	83	
.ytul	672	333	
Jun.	669	336.3	
May.	2208	1105.6	
Apr.	2364	1181.7	
Mar.	2038	1043.5	
Feb.	1938	969.1	
Jan,	1114	560.3	
Month	Mean CFS	Stand Deviation	

Wonthly Mean and Stand Deviation of Kaskaskia River. Table 3.6 1:

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	1																						
Optimal Expected	1 03016		1.93016	1.93016	1.93016	1.93016	1,93016	1.93016	1.93016	1.93016	1.93016	1.93016	1.93016	1.93016	1.93016	1.93016	1.93017	1.93017	1.93018	1.93020	1.93022		
12	1500	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1) い) い) い	1400	1200	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1700	11000	1500		
11	1 600	1600	1650	1650	1650	1650	1700	1700	1700	1700	1700	1700	1750	1750	1700	1700	1650	1650	140	1400	1500		
10	1800	1800	1800	1800	1800	1800	1800	1800	1850	1800	1850	1800	1850	1850	1800	1800	1750	1700	1650	1550	1500		
6	1950	1900	1900	1900	1900	1900	1900	1900	1900	1900	1900	1900	1900	1900	1900	1850	1850	1800	1750	1600	1500		
ω	2100	2000	2000	2000	2000	2000	2000	2000	2000	1950	1950	1950	1950	1950	1950	1900	1900	1850	1800	1700	1550		
2	2200	2150	2100	2100	2100	2050	2050	2050	2050	2050	2050	2000	2000	2000	2000	1950	1950	1900	1850	1750	1650		
· vo	2200	2300	2450	2550	2650	2750	2900	2550	2200	2100	2100	2050	2050	2050	2050	2050	2050	2000	2000	2000	2000		
Ś	2200	2300	2450	2550	2650	2750	2900	3000	2300	2150	2100	2050	2050	2050	2050	2050	2050	2050	2050	2050	2050		
4	2200	2300	2450	2550	2650	2750	2900	3000	2650	2250	2150	2100	2100	2050	2050	2050	2050	2050	2050	2050	2050		
Э	2200	2300	2450	2550	2650	2750	2950	3000	3100	2500	2250	2200	2150	2150	2100	2100	2100	2100	2100	2100	2100		
N	2200	2300	2450	2550	2650	2750	2900	3000	3100	3250	2450	2300	2200	2200	2150	2150	2150	2150	2150	2150	2150	-	
1	2200	2300	2450	2550	2650	2750	2900	3000	3100	3250	2850	2500	2350	2300	2250	2250	2200	2200	2200	2200	2200		
Stage State	0	50	100	150	200	250	000 000	350	1000	450	500	0. 2 2	600	650	200	750	800	850	900	950	1000		

Jan-Feb Optimal Release Policy for Shelbyville Reservoir. Table 3.7

Optimal	Expected Cost	3.600	3.600	3.600	3.600	3.600	3.600	3.600	3.600	3,600	3.600	3.600	3,600	3,601	3.603	3.605	3,607	3.609	3.612	3.615	3.617	3.621
	N	1500	1400	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
	1	1850	18.10	1900	1900	1900	1900	1900	1950	1950	1900	1950	1950	1950	1900	1900	1900	1900	1800	1750	1600	1500
	10	2150	2150	2150	2150	2150	2150	2150	2150	2150	2150	2150	2100	2100	2100	2100	2050	2000	1950	1900	1800	1550
	5	2500	2450	2400	2400	2400	2350	2350	2300	2300	2300	2250	2250	2250	2200	2200	2150	2150	2100	2050	1950	1800
c	Ø	2850	2750	2700	2650	2600	2600	2500	2500	2500	2450	2450	2400	2400	2400	2350	2350	2250	2200	2200	2100	2000
2		3400	3150	2950	2850	2800	2750	2750	2700	2650	2650	2600	2600	2550	2550	2500	2550	2450	2400	2350	2250	2150
7	o	4150	4250	0077	4500	3400	3150	3000	2900	2850	2800	2750	2750	2750	2700	2700	2700	2650	2650	2650	2600	2600
Ľ	n	4150	4250	6644	4500	4600	3600	3300	3100	2950	2900	2850	2850	2800	2800	2750	2750	2750	2750	2750	2700	2650
ų	ħ	4150	4250	00 77	4500	4600	4700	4020	3500	3250	3100	3050	2950	2950	2900	2900	2850	2850	2850	2800	2800	2750
c	ſ	4150	4250	0077	4500	4600	4700	4850	4950	3750	3450	3350	3200	3150	3100	3050	3000	3000	2950	2950	2900	2900
c	N.	4150	4250	0077	4500	4600	002行	4850	4950	5050	4450	3850	3600	3450	3350	3300	3250	3200	3150	3150	3100	3100
٣	-1	4150	4250	0044	4500	4600	4700	4850	4950	5050	5200	5300	4550	4050	3800	3650	3600	3500	3500	3450	3400	3350
STAGE	10 ⁶ cfs	0	0 20	00	150	200	250	300	350	007	450	500	57 0	600	650	200	750	800	850	006	950	1000

Mar-Apr Optimal Release Policy for Shelbyville Reservoir. Table 3.8

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May-June Optimal Release Policy for Shelbyville Reservoir.

Table 3.9

												_										
Optimal Expected	Cost	1.8351	1.8667	1.8914	1.9302	1.9596	1.9918	2,0221	2,0488	2.0897	2,1181	2,1508	2,1824	2,2177	2,2421	2,2722	2.2996	2.3346	2.3513	2.3816	2.3988	2,4227
12		1300	1400	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
11		1300	1400	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
10		1300	1400	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
6		1300	1400	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
ω		1300	1400	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
2		1300	1400	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
\$		1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500
Ŋ		1750	1750	1800	1800	1800	1800	1850	1850	1850	1850	1850	1850	1850	1850	1850	1800	1800	1700	1650	1550	1500
4		2000	2000	2000	2000	2000	2000	2000	2000	2000	2000	2000	2000	2000	2000	2000	1950	1900	1900	1800	1700	1500
ŝ		2200	2200	2200	2200	2150	2150	2150	2150	2150	2150	2100	2100	2100	2100	2100	2050	2000	2000	1900	1800	1600
N		2500	2450	2400	2350	2350	2300	2300	2300	2250	2250	2200	2200	2200	2200	2150	2150	2100	2100	2050	1950	1800
Ч		2900	2750	2650	2550	2550	2500	2450	2450	2400	2400	2350	2350	2300	2300	2250	2250	2200	2200	2150	2100	1950
Stage State	10° cfs/	0	50	100	150	200	250	000	920 920	001	450	500 002	520 220	600	650	200	750	800	850	006	950	1000

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in water resources planning. It allows for a rapid evaluation of system performance for a set of trial policies. The performance of a system is evaluated through a quantitative model which is a mathematical representation of the system under study. The operating policy of the system was the major concern of this study.

3.4.1 Effectiveness Parameters for Flood Control

Three efectiveness parameters for flood control which have been proposed [10] are presented below.

1. Absolute Flood Peak Reduction (10)

The absolute flood peak reduction, ΔQ_p is defined as

$$\Delta Q_{p} = Q_{p} - Q_{pR} \qquad (cfs), \qquad (3.79)$$

where

 Q_p = flood peak for non-regulated flood, and

 Q_{pR} = flood peak for flow regulated by reservoirs. The parameter ΔQ_p would be the one used directly in flood damage mitigation calculations.

2. Relative Flood Peak Reduction (RR)

In order to facilitate an interpretation of the results obtained with different policies, it is best to normalize ΔQ_p . One normalization, called the relative flood peak reduction or RR, is defined as

$$RR = 100 \frac{\Delta Q_p}{Q_p} \%$$
(3.80)

This parameter isuseful in relating flood control effectiveness to

system configuration.

3. Relative Flood Reduction (RC)

Another normalization of ΔQ_p called the relative flood reduction, or RC, is defined as

$$RC = 100 \frac{\Delta Q_p}{Q_p - \overline{Q}_f} \%, \qquad (3.81)$$

where \overline{Q}_{f} = maximum allowable discharge which will not cause a flood. This parameter is useful in relating policy performance with variations in \overline{Q}_{f} induced for example by levee construction, zoning, etc.

The relation of these Effective ness Parameters for flood control is displayed graphically on Figure 3.8.

3.4.2 Results of Simulation

A. Historical Record Data

(a) Individual Storm

There were a total of seven simulation runs. The data for these runs were chosen arbitrary from the flow rate data in Appendix B. Each simulation run begins from the increasing period of the hydrograph. From these inflow data we can see that there was a single storm, but we don't known the time duration of each rainfall. The results are not bad. But this policy has been restricted to a single storm flood, because it does not permit consideration of the combination of two or more storm hydrographs. The results of the simulation of an individual storm is on Table 3.10. In Table 3.10, the first column is inflow of each stage, the second column is the optimal release rule, and the third column is the standard release rule. The values of the average



Figure 3.8 Effectiveness Parameters for flood control

	STAN		1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	0	200	21010	100%	
2	140 0		1500	1500	1500	1500	1500	1500	1200	1500	1500	1500	1500	1500	1500	1500	1500	0	200	31015	100%	
	II.		364	1230	1540	1580.	1230	1860	2200	2150	1360	1630	1490	1270	1030	832	864	1.68				
	STAN		1500	1500	1500	1500	1500	3120	3330	3430	2400	2150	1810	1580	1500	1500	1500	3.36	270	2.0%	123%	İ
10	OFT		1500	1500	1500	1500	1500	3120	3300	3430	2400	2150	1810	1580	1500	1500	1500	3.36	270	200	1238	
	NI		353	2550	3330	3700	3160	3260	3380	3430	5400	2150	1810	1580	1360	1150	† 6 6	3.65				
	STAN		1500	1500	1500	1500	2936	4650	3260	2850	2700	2300	1810	1500	1500	1500	1500	4.5	220	201	6. 5%	
Ś	OPT		1500	1500	1500	1500	2936	4650 .	3260	2850	2700	2300	181.)	1500	1500	1500	1500	4.5	220	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	6.5%	
	ЦЦ		さんし	1070	3700	0164	4870	4650	3260	2850	2700	2300	1810	1450	1270	1100	1 66 ·	4.65				
	SIAN		1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	0	950	387%	100%	
4	OPP		1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	0	950	3874 3874	ू २००१	Í
	Ш		666	1360	2 ¹ 400	2450	2300	2050	1910	1630	1450	1150	1030	918	810	738	702	2.09				
	STAN		1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	880	0	07	2 2	1%00T	
ŝ	OPT		1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1500	1450	200	o	07	2 2 2 2 2 2 2 2 2	1 00%	
	IN		965	1070	1230	1540	1400	1270	1230	1230	1150	1030	918	846	426	738	702	0.09				
	STAN		1500	1500	1500	1500	1500	1500	1956	2750	2550	1810	1500	1500	1500	1500	1500	2.53	410	129%	24.7%	
2	OPT		1500	1500	1500	1500	1500	1500	1956	2750	2550	1810	1500	1500	1500	1500	1500	2.53	410	1201	2/:42	
	III		864	1070	2300	3160	3050	2950	2850	2750	2550	1810	1400	1360	1360	1270	882	3.07			ŕ	
	STAIN		1500	1500	1500	1500	2796	2350	1810	1630	1500	1500	1500	1500	1500	1500	1500	2.60	1854	2000	2012	
-1	OPT		1500	1500	1500	1500	2796	2350	1310	1630	1500	1500	1500	1500	1500	1500	1500	2.60	1854	200 00 00 00 00 00 00 00 00 00 00 00 00	54/32	
	MI		1 617	1810	11090	4650	3540	2350	1810	1680	1490	1270	1270	1030	918 8	274	846	4.49				
		Stage	-	2	ო	4	Ś	6	2	ω	δ	10 1	11	21 17	ლ. ქ	1	יה דו	Cost	500 J	En .	C L L	

Table 3.10 Result of Individual Storm Simulation.

effective parameters are presented on Table 3.11. Reference to Table 3.11, reveals that for real time operation the optimal policy for individual storm is quite good. The reason is that in most cases the capacity of the reservoir is sufficient to handle any single storm flood.

(b) Stage Independent Input

In simulating flood control the stage time interval can not be too long. Here we assume that the time interval is five days and the initial storage is 500,000,000 ft.³ (half full). The inflow data are the five days average historical data and are given from Table 3.12. The planning horizon, T, for each simulation run was two months (12 stages). The mesults of the simulation runs are presented in detail in Appendix D1. The average value of the effective parameters are presented on Table 3.13.

Reference to Table 3.13 reveals that for real time operation the optimal release policy is a little worse than standard release policy and both the release rules are not effective in reducing the flood peak. The reasons for this are:

- 1. The inflow rate is dependent upon the season, and varies significantly between wet seasons and dry seasons. However, in the simulation runs we use release policies which assumed that the distribution is stationary and we used the year-round average distribution to approximate the nonstationary flow distribution.
- 2. In a relatively short time interval, for example, a day, a week, or a month, the flow is clearly not time independent and stationary. If the time interval is a year, it may be time independent. For a storage reservoir, one year intervals may be appropriate. However,

	Initial flow	optimal Release Rule	Standard Release Rule
Average Jamage Cost	2,817	1.855	1.855
Average RR		19•64劣	19.64%
Average RC		57.46%	57.46%

Average Value of Effectiveness Parameters from Simulation using Historical data with Individual Storm Inflow Assumption. Table 3.11

Tune July Aug. Sep. Oct. Nov. Dec.	992 237 70 24 18 11 30	807 351 50 16 13 18 27	710 224 32 10 13 18 34	392 130 27 12 10 25 27	324 123 22 10 9 37 24	211 76 22 15 9 50 30	572 187 46 15 12 26 28
Apr. Nay.	2356 4662	2030 8978	2106 6316	1162 4302	1291 3028	2550 1720	1920 4730
Jan. Feb. Mar.	9094	4870	4230	7234	1402	1104	3010
5 days Ferlod	. 1	Ņ	m	4	<i>v</i> 0	9	Average

Table 3.12 Five Days Average Flow Rate (1908)

	Dec.	846 597 826 824 824 822 820 822 820 820 820 820 820 820 820			, Dec	632 632 632 632 632 632 632 632 632 632	
	Nov.	246 2628 25528 25528 2146 2146 2146			.vov	800025 800025 80005 115 80005 115 80005 115 80005 115 80005 115 8005 800	
	Oct.	7324 7324 17952 16256 16256 16256 16256			Oct.		
and the second se	Sep.	44 44 44 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7			Sep.	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	
and a share of the second s	Aug.	48444 <i>6</i> 6 94494 <i>6</i> 6	e (191		Aug.	00000000000000000000000000000000000000	-
	july	244042 24404 2007	ow rat		July	98999963 98429963 98429963	
	June	111 00000 00000 0000	age fl		June	4000000 4000000 4000000 4000000 40000000	
	May.	1 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7	20Λ0 20 20 20 20 20 20 20 20 20 20 20 20 20		May.	19250 19250 193500 19350 193500 193500 193500 10000000000000000000000000000000000	
and the second	Apr.	1468 309066 4002066 40020 4002 4002 4002 4002 40	Ve Day	-	Apr.	200000 200000 200000 200000 200000 200000 2000000	
	Har.	11 1643 2009 2009 2009 2009 2009 2009 2009 200	12a F1		Mar	44000088 400000888 400000888 40000888 40000888 40000888 4000500888 4000500888 4000500888 4000500888 4000500888 4000508888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 40005088 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 400050888 40005088 400050888 40005088 40005088 40005088 40005088 400050000000000	
	Feb.	2272 11502 11402 11432 11432 11432 11432	ble 3.		нер	4 <i>2</i> 4 2000 2000 2000 2000 2000 2000 2000 2	
	Jan.	1078 1078 10138 10138 10138 1002 1002 1002 1002 1002 1002 1002 100	с Ц		Jan.	1070 820 820 820 820 820 820 820	
	Month 5 Days Period	-1 01 01 4 01 0 4 6 6 6 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7			Month 5 Days Period	A A A A A A A A A A A A A A A A A A A	

Table 3.12b Five Days Average Flow Rate (1912)

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Dec.	351 237 237 237 237 237 593 593 593
Nov.	88 86 76 86 16 80 70 70 70 70 70 70 70 70 70 70 70 70 70
Oct.	11 110 110 110 100 100 100 100 100 100
Sep.	000 000 000 000 000 000 000 000 000 00
Aug.	271 102 82 82 82 82 82 82 82 82 82 82 82 82 82
July	682 2908 3158 1795 2190 2190
June	2580 2580 1980 1980 1980 1980 1980 1980 1980 19
May.	2102 2380 23808 23808 23808 6052 638 638 19400 19400
Apr.	372 4070 5374 5394 5394 5394 2300 2910 2300 2300
Mar.	2192 1420 1152 800 1569 1090
Feb.	2008 2008 2008 2008 2008 2008 2008 2008
Jan.	00000000000000000000000000000000000000
MONTH 5 DAYS PERIOD	1 2 Average

k.,e

Table 3.12c Five days Average flow rate (1909)

ł		
	Dec.	800 80 700 80 80 80 80 80 80 80 80 80 80 80 80 8
	Nov.	803377 8007778 430 80 80 80 80 80 80 80 80 80 80 80 80 80
	Oct.	4848448 88289848 8828984 8828988
	Sep.	20 20 20 20 20 20 20 20 20 20 20 20 20 2
	Aug.	94 704 70 70 70 70 70 70 70 70 70 70 70 70 70
	July	4732 732 732 747 747 747 747 747 747 747 747 747 74
	June	1104 0444 4001004 4001000 400000 4000000
	May.	112 000 000 000 000 000 00 00 00 00 00 00
	Apr.	82999999999999999999999999999999999999
	. TeM	3386 1860 916 7562 1220 1220
	Feb.	5000 5000 5000 5000 5000 5000 5000 500
	Jan.	280 280 280 280 280 280 280 280 280 280
	Month 5 days Period	AVerage

Table 3.12d Five Days Average flow rate (1910)

	Inttu. Flow	Optical Release Rule	Standard Rolesse Rule
Average Damage Cost	2.500169	2.110559	2.030202
Average RR		13.293	16.83%
Average RC		22.95%	35.46%

Table 3.13 Average value of Effectiveness Parameters from Simulation Using Historical data with Time Independent Inflow rate Assumption.

Standard Release Aule	3.5523309	1.3 ° 08%	19。03%
Optimal Release Rule	3.541045	13。14周	19。97%
Initial Flow	4°2546		
	Average Damage Cost	Average RR	Average RC

Table 3.14 Average Value of Effectivesness Parameters from Simulation Using Historical data with Time Dependent Inflow rate Assumption .

No. of Street, or other

in simulating the operation of flood control reservoirs the stage interval should be relatively short.

(c) Time Dependent Inflow Rate

In the simulation the historical record data on Table 3.12 was used. The results of the simulation runs for January-February, March-April, and May-June are on Appendix D2 and the average value of effective parameters are presented on Table 3.14.

In the process of determining the optimal policy, we assumed that the distribution of flow in each season was Normal with expected value equal the sample mean and variance equal to the sample variance. Here the simulation used only five years of the historical data, which did not appear to be long enough. Hence, the optimal release rule was not much more effective than the standard release rule.

Next we will use the Markovian Model to simulate the stream flow and examine the effects of the release rules.

B. Simulated Stream Flow Data

For the simplest Markovian model [61], we have

$$q_{i} = \beta_{0} - \beta_{1} q_{i-1} + e_{i}, \qquad (3.82)$$

where

q_i is the flow rate at stage i, and

 e_i is the random part of the synthetic flow.

In this Markovian flow model, we consider that e_i has a Normal distribution. We simulated the flow for three two-month periods: Jamuary-February, March-April, and May-June. Each month was divided into six stages so that there were five days in each stage interval. In each month we assumed that the mean was constant \overline{X}_j . In each two month interval, we have

$$q_i = \overline{x}_1(1 - r_1) + r_1 q_{i-1} + t_i s_1 \sqrt{(1 - r_1^2)}$$
 $i = 1, 2, ..., 6$ (3.83)

$$q_i = \overline{x}_2(1 - r_2) + r_2 q_{i-1} + t_i s_2 \sqrt{(1 - r_2^2)}$$
 $i = 7, 8, ..., 12$ (3.84)

in which \overline{X}_1 and \overline{X}_2 are the means of the first and second months of each two month period, respectively, r_1 and r_2 are the correlation coefficients between stages in the first and second month period, respectively, and t_1 , t_2 , ... is a sequence of independent Normally distributed random numbers with mean 0 and standard deviation 1.

If q_j was negative for some j, then we used the negative q_j in the equation for q_{j+1} and discarded q_j without using it as a flow in the simulation. The correlation coefficients for January, February, March, April, May, and June, were 0.2865, 0.3639, 0.3137, 0.4742, 0.3268, and 0.3114 respectively.

The results of fifty simulated years for each two month period are presented in Appendix D3. Both the simulated and historical data were tested for identical population distribution functions (see Appendix C2 for the results). The comparison of the effective parameters is presented on Tables 3.15, 3.16, and 3.17, respectively.

From the relative flood reduction (RC), we can see that the time dependent optimal policy is much better than both the time independent optimal and the standard release policy. For the January-February and May-June periods, the time dependent optimal policy can reduce almost 40% of the flood peak. The March-April period reduction is relatively small, because the average flow rate in that period is

	Time Independent Opt, <u>Policy</u>	Time Dependent Opt, policy	Standard Policy	Initial Flow
Flood Damage	2.3232	2.2374	2.5922	3.8966
Average RR	18.48%	20.912%	12.828%	
Average RC	35.454%	39.152%	25.352%	

Table 3.15 Comparision of Effectiveness Parameters (Jan.-Feb.)

	Time Independent Opt, Policy	Time Dependent Opt, Policy	Standard Policy	Initial Flow
Flood Damage	3.80328	3.4765	3. 89102	4.0147
Average RR	3.402%	11.424%	1.234%	
Average RC	5.858%	17.376%	1.964%	

Table 3.16 Comparision of Effectiveness Parameters (Mar.-Apr)

	Time Independent Opt. Policy	Time Dependent Opt. Policy	Standard Policy	Initial Flow
Flood Damage	2.8686	2.47	2.8974	3.673
Average RR	9.184%	18.806%	8.541%	
Average R C	17.62%	34.15%	16.954%	

Table 3.17 Comparision of Effectiveness Parameters (May.-Jun.)

higher than the capacity of reservoir to effectively handle it.

From the previous discussion and results, we know that the use of historical data to estimate the probability of flow in each period and then determine the optimal policy is not an adequate procedure. The reason for this is that the natural conditions are highly variated and correlated. Furthermore the relation of these variations and the actually probability distributions are difficult to estimate.

The Program listing appears in Appendix A5.

CHAPTER IV

EFFECTS OF VARYING PARAMETERS

4.1 Introduction

In this chapter, we discuss the effects of varying the flood control parameters: initial storage quantity, flood stage and purposes of reservoirs. In flood control, we may adjust the water level of reservoirs when a storm has been predicted. We may also increase the flood stage by channel improvements or the construction of levees. In §4.2 and §4.3 we discuss the effects of varying initial storage quantity and flood stage for both the deterministic inflow and the stochastic inflow cases. In §4.4 we formulate a dynamic programming algorithm to determine the optimal release policies for a dual-purpose reservoir.

4.2 Effect of Varying the Initial Storage Quantity

The initial storage quantity is critical in flood control system operation. When a heavy storm has been predicted, the water level of the reservoir should be changed to some best situation in order to absorb the flood peak as much as possible.

The following results specify the "best" situation and tend to confirm our intuition.

<u>Proposition 4.1</u> Consider the single reservoir, deterministic inflow flood control model of $\S2.2$. Assume that the input quantity of water

in a single time interval can not exceed the capacity C. Then the optimal damage cost is a monotone nondecreasing function of the initial storage.

Proof:

The proof is by induction. Recall from §2.5.1 that the boundary condition is

$$f_{N+1}(Z_{N+1}) = 0. \qquad \forall Z_{N+1}$$

Clearly, the result is true for N+1.

Consider stage N. Recall from (2.46) we had that

$$f_{N}(Z_{N}) = \min_{r_{N} \in R_{N}} (U(r_{N})), \qquad (4.1)$$

where

$$R_{N} = [max\{0, (Z_{N} + Q_{N} - C)/\Delta T\}, (Z_{N} + Q_{N})/\Delta T].$$
 (4.2)

By basic assumption, U(·) is a monotone nondecreasing function of r. Hence

$$f_N(Z_N) = U(Lb_N(Z_N, Q_N))$$
 (4.3)

where

$$Lb_N(Z_N, Q_N) = max\{0, (Z_N + Q_N - C)/\Delta T\}$$

Clearly since $Lb_N(Z_N, Q_N)$ is a monotone nondecreasing function of Z_N , $f_N(Z_N)$ is also a monotone nondecreasing function of Z_N .

Next we make the inductive assumption that the result is for n + 1, n + 2, ..., N + 1, that is, $f_{n+1}(Z_{n+1})$ is a monotone nondecreasing

function of Z_{n+1} . From (2.46) we have that for the nth stage

$$f_{n}(Z_{n}) = \min_{\substack{r_{n} \in \mathbb{R} \\ n}} \{\max(U(r_{n}), f_{n+1}(Z_{n} + Q_{n} - r_{n}\Delta T))\}$$

$$= \min_{\substack{r_{n} \in \mathbb{R} \\ n}} \{\max(h_{1}(r_{n}), h_{2}(r_{n}))\}$$
(4.4)

By the inductive assumption we have that for fixed r_n , $f_{n+1}(Z_n + Q_n - r_n \Delta T)$ is a monotone nondecreasing function of Z_n . For fixed Z_n , we have from (2.45) that

$$Lb_n(Z_n, Q_n) = \max\{0, (Z_n + Q_n - C)/\Delta T\},\$$

and

$$Ub_{n}(Z_{n}, Q_{n}) = (Z_{n} + Q_{n})/\Delta T.$$

There are three cases for consider i) $r_n = Lb_n$, ii) $r_n = Ub_n$, and iii) $r_n \in (Lb_n, Ub_n)$. If $r_n = Lb_n$, then

$$h_{1}(r_{n}) = U(Lb_{n}(Z_{n}, Q_{n}))$$

= U(max{0, (Z_{n} + Q_{n} - C)/\Delta T)}). (4.5)

and

$$h_{2}(r_{n}) = f_{n+1}(Z_{n} + Q_{n} - \max\{0, (Z_{n} + Q_{n} - C)\})$$
$$= f_{n+1}(\min\{Z_{n} + Q_{n}, C\})$$
(4.6)

Now if $Z_n = 0$, then

$$h_{1}(r_{n}) = U(Lb_{n}(0, Q_{n}))$$

= U(max{0, (Q_{n} - C)/\DeltaT})
= U(0) = 0, (4.7)

$$h_2(r_n) = f_{n+1}(Q_n) \ge 0.$$
 (4.8)

Hence

$$h_2(r_n) \ge h_1(r_n)$$
.

Now if $Z_n = C$, then

$$h_1(r_n) = U(Q_n / \Delta T),$$
 (4.9)

and

$$h_2(r_n) = f_{n+1}(min\{C + Q_n, C\})$$

= $f_{n+1}(C)$ (4.10)

here the value of $f_{n+1}(C)$ depends on the inflow Q_{n+1} . Hence for $r_n = Lb_n$ the relation of h_1 and h_2 depends on Q_n and Q_{n+1} . If $r_n = Ub_n = (Z_n + Q_n)/\Delta T$, then

$$h_1(r_n) = U((Z_n + Q_n)/\Delta T),$$
 (4.11)

and

$$h_2(r_n) = f_{n+1}(Z_n + Q_n - Z_n - Q_n) = f_{n+1}(0).$$
 (4.12)

At the upper bound point the value of h_2 is the constant $f_{n+1}(0)$. Consider $r_n \in (Lb_n, Ub_n)$. If h_1 and h_2 intersect, then this intersection is in decision region R_n . If there is no intersection then h_2 is above h_1 or h_1 is above h_2 . The configuration of those curves is illustrated on Figure 4.1. If h_2 is above h_1 , the $f_n(Z_n) = f_{n+1}(0)$, it is a constant. If h_2 is under h_1 the $f_n(Z_n)$ is on curve h_1 . Hence, $f_n(Z_n)$ is monotone nondecreasing function of $Z_n \cdot \|$

The following intuitive result is an immediate consequence of Proposition 4.1

<u>Corollary 4.1</u> Consider the flood control model of §2.2. When a flood flow is predicted, the optimal operating policy is to set the water level in the reservoir as low as possible with a safe release rate.

We next consider the stochastic case.

<u>Proposition 4.2</u> Suppose that the inflow quantities, Q_n , n = 0, 1, ... N, are independent random variables with known probability density $p_n(q)$. Then the optimal expected cost is a monotone nondecreasing function of the initial storage of the reservoir.

Proof:

The proof is by induction. Recall from §3.3.1 the boundary condition is $f_{N+1} = 0$. Clearly the result is true for N + 1.

Consider stage N. Recall from (3.41) we have that

$$f_{N}(Z_{N}) = \min \{E(U(r_{N}))\}$$

$$r_{N} \in R_{N}$$
(4.13)

in which, for fixed Q_N ,







$$Lb_{N}(Z_{N}, Q_{N}) = max\{0, (Z_{N} + Q_{N} - C)/\Delta T\}$$
 (4.14)

$$Ub_n(Z_n, Q_N) = (Z_N + Q_N)/\Delta T$$
 (4.15)

Now by assumption U(·) is a nondecreasing function of r. The notation E_{Q_N} will be used to denote the expectation with respect to the random variable Q_N . Then $E_{Q_N}(U(r_N))$ is a monotone nondecreasing function of r_N .

We can write $f_N(Z_N)$ as

$$f_{N}(Z_{N}) = \min_{r_{N} \in R_{N}} \{H_{N}(r_{N})\},$$

where H_N denotes some monotone nondecreasing function of r_N . Recall that (3.55) the optimal decision at stage N is

$$r_{N}^{\star} = \max\{L, \min[Lb_{N}(Z_{N}, Q_{N}^{1}), Lb_{N}(Z_{N}, Q_{N}^{2}), \dots, Lb_{N}(Z_{N}, Q_{N}^{p})]\},$$
 (4.16)

where L is the flood stage and Q_N^1 , Q_N^2 , ..., Q_N^p are all possible discrete values of the input at stage N. Hence, r_N^* is a monotone nondecreasing function of Z_N .

Next we make the inductive assumption that the result is true for n + 1, n + 2, ..., N + 1, that is, that $f_{n+1}(Z_{n+1})$ is a monotone nondecreasing function of Z_{n+1} . From (3.41) for n^{th} stage, we have

$$f_{n}(Z_{n}) = \min \{\max(EU(r_{n}), Ef_{n+1}(Z_{n} + Q_{n} - r_{n}\Delta T))\}, \qquad (4.17)$$

$$r_{n} \in R_{n}$$

Now for fixed Z_n , $E_{Q_n} r_{n+1}(\cdot)$ is a monotone nonincreasing function of r_n and $E_{Q_n} U(r_n)$ is a monotone nondecreasing function of r_n . For fixed
Q_n also, $F_{n+1}(\cdot)$ is a monotone nonincreasing function of r_n . Hence, $f_n(Z_n)$ can be written as

$$f_{n}(Z_{n}) = \min_{r_{n} \in \mathbb{R}_{n}} \{\max(H_{n}^{1}(r_{n}), H_{n}^{2}(r_{n}))\}, \qquad (4.18)$$

in which H_n^1 is a monotone nondecreasing function of r_n and H_n^2 is a monotone nonincreasing function of r_n . By a similar argument as in the proof of Proposition 4.1, we can obtain $f_n(Z_n)$ is a monotone nondecreasing function of $Z_n \cdot \|$

The following intuitive result is an immediate consequence of Proposition 4.2.

<u>Corollary 4.2</u> For the flood control model of §3.3.1, when a flood flow is predicted to be coming the optimal operating policy is to set the water level in the reservoir as low as possible with a safe release rate.

We next consider the effects of varying flood stages.

4.3 Effects of Varying the Flood Stage

As a result of channel improvements or the construction of levees the flood stage will increase. These are also methods of flood control employed on some rivers. Hence, we need to consider the effects of changing the flood stage.

Suppose that the flood stage changes, but the damage cost relative to the depth of water on ground remains the same. This means the damage cost function is a family of curves which are parallel to one another as displayed on Figure 4.2. <u>Proposition 4.3</u> Consider the single reservoir, deterministic inflow flood control model of §2.2. The optimal damage cost is a monotone nonincreasing function of the flood stage.

Proof:

The proof is by induction.

Let L_i denote the flood stage, $U(r, L_i)$ denote the damage cost function when the flood stage is L_i and the release rate is r. Recall from §2.5.1 that for boundary condition is

$$f_{N+1}(Z_{N+1}, L_i) = 0.$$
 (4.19)

Clearly the result is true for N + 1.

Consider stage N. Recall from (2.46) we have that

$$f_{N}(Z_{N}, L_{i}) = \min_{\substack{r_{N} \in R_{N}}} \{ (U(r_{N}, L_{i})) \}, \qquad (4.20)$$

where the decision region $R_N = [max\{0, (Z_N + Q_N - C/\Delta T\}, (Z_N + Q_N)/\Delta T]$ is not affected by L_i .

From (4.20) the optimal decision at stage N is

$$\mathbf{L}_{N}(\mathbf{Z}_{N}, \mathbf{Q}_{N}) \qquad \text{if} \quad \mathbf{L}_{N} > \mathbf{L}_{i},$$

$$\mathbf{r}_{N,\mathbf{L}_{i}} = \max\{\mathbf{L}_{i}, \mathbf{L}_{N}(\mathbf{Z}_{N}, \mathbf{Q}_{N})\} \qquad \text{if} \quad \mathbf{U}_{N} > \mathbf{L}_{i}, \mathbf{L}_{N} \leq \mathbf{L}_{i}, (4.21)$$

$$\mathbf{U}_{N}(\mathbf{Z}_{N}, \mathbf{Q}_{N}) \qquad \text{if} \quad \mathbf{U}_{N} \leq \mathbf{L}_{i}, \mathbf{L}_{N} \leq \mathbf{L}_{i}.$$

Hence, for fixed L_i , $f_N(Z_N, L_i)$ is a monotone nondecreasing function of Z_N . For fixed Z_N suppose the $L_2 > L_1$. Then from Figure 4.3, we





can see that

$$f_N(Z_N, L_2) \leq f_N(Z_N, L_1).$$

Hence, $f_N(Z_N, \cdot)$ is a monotone nonincreasing function of L.

We next make the inductive assumption that the result is true for n + 1, n + 2, ..., N + 1, that is, that $f_{n+1}(Z_{n+1}, L_i)$ is a nonincreasing function of L_i . From (3.41) for the n^{th} stage, we have

$$f_{n}(Z_{n}, L_{i}) = \min \{\max(U(r_{n}, L_{i}), f_{n+1}(Z_{n} + Q_{n} - r_{n}\Delta T, L_{i}))\}, \quad (4.22)$$

$$r_{n} \in R_{n}$$

For $L_2 > L_1$

$$\max(U(r_n, L_1), f_{n-1}(Z_n - Q_n - r_n^{\Delta T}, L_1))$$

$$\geq \max(U(r_n, L_2), f_{n+1}(Z_n + Q_n - r_n^{\Delta T}, L_2)) \qquad (4.23)$$

Hence $f_n(Z_n, \cdot)$ is a monotone nonincreasing function of L.

We next consider the stochastic case.

<u>Proposition 4.4</u> Suppose that the inflow quantities Q_n , n = 0, 1, ... Nare independent random variables with known probability density $p_n(q)$. Then, the optimal expected cost is a monotone nonincreasing function of the flood stage L.

Proof:

The proof is by induction. Recall from 3.3.1 the boundary condition is $f_{N+1} = 0$.

Clearly, the result is true for N + 1.

Consider stage N. We have

 $f_{N}(Z_{N}, L) = \min_{\substack{r_{N} \in R_{N}}} (E_{Q_{N}} U(r_{N}, L))$

For fixed r_N , $U(r_N, \cdot)$ is monotone nonincreasing function of L and for fixed L, $U(\cdot, L)$ is nondecreasing function of r_N . Obviously, $E_{Q_N}U(r_N, L)$ has the same monotonicity as $U(r_N, L)$. Hence, for fixed L, $f_N(Z_N, Q_N)$ is a monotone nondecreasing function of Z_N , and for fixed Z_N , $f_N(Z_N, L)$ is a monotone nonincreasing function of L.

Next we make the inductive assumption that the result is true for n + 1, n + 2, ..., N + 1.

From (3.41), we obtain

$$f_n(Z_n, L) = \min_{\substack{r_n \in \mathbb{R}\\n}} \{\max(E_Q_n U(r_n, L), E_{Q_n} f_{n+1}(Z_n + Q_n - r_n \Delta T, L))\}.$$

Now $E_{Q_n} U(r_n, L)$ and $E_{Q_n} f_{n+1}(Z_n + Q_n - r_n \Delta T)$ inherit the same monotonicity property for r_n , and L, as $U(r_n, L)$ and $f_{n+1}(Z_n + Q_n - r_n \Delta T, L)$, respectively. Hence for fixed Z_n , $f_n(Z_n, L)$ is a monotone nonincreasing function of L. Thus, the optimal damage cost $F_1(Z_1, L)$ is a monotone nonincreasing function of L. \parallel

4.4 Effects of Varying the Objective Function

Here we consider multi-purpose reservoirs which have additional functions such as recreation, water supply, hydroelectric generation, etc.

The recreation requirements impose a minimum water level constraint. However, we can eliminate this constraint simply by changing the capacity of the reservoir (i.e., use the active capacity). The minimum release constraint can be eliminated simply by changing the range of the decision variables.

For flood control and water supply we need a dual-purpose reservoir. The flood damage cost is a function of flood peak and the water supply

min U(r) -
$$\sum_{n=1}^{N} B(r_n)$$
 (4.24)

subject to

$$Z_n - Z_{n-1} = Q_n - r_n \Delta T$$
 $n = 0, 1, ..., N$
 $0 \le Z_n \le C$ $n = 0, 1, ..., N$

where

U is the damage cost, a function of peak discharge rate,

r is the peak discharge rate,

B is the benefit function of release rate, and

 r_n is the discharge rate in stage n.

We can address this problem with dynamic programming. The transition function is

$$Z_{n+1} = Z_n + Q_n - r_n \Delta T.$$
 (4.25)

From the constraints $0 \le Z_n \le C$, for all n, we obtain the upper bound of r_n as

$$Ub_{n}(Z_{n}, Q_{n}) = (Z_{n} + Q_{n})/\Delta T,$$

and the lower bound of r_n as

$$Lb_n(z_n, Q_n) = max\{0, (Z_n + Q_n - C)/\Delta T\}.$$

The decision region of r_n is

$$R_n = [Lb_n, Ub_n]$$

Define $f_n(Z_n)$ as the minimum cost from stage n(time n ΔT) to stage

N (time T), if the storage level is Z_n at time n ΔT (stage n). Then, invoking Bellman's principle of optimality yields the following functional equation of dynamic programming

$$f_{n}(Z_{n}) = \min \{\max[U(r_{n}), g_{n+1}(Z_{n} + Q_{n} - r_{n}\Delta T)] - B(r_{n}) \\ r_{n} \in \mathbb{R}_{n} - h_{n+1}(Z_{n} + Q_{n} - r_{n}\Delta T)\} = 0, 1, ..., N \quad (4.26)$$

$$= \max [U(r_{n}^{*}), g_{n+1}(Z_{n} + Q_{n} - r_{n}^{*}\Delta T)] - B(r_{n}^{*}) - h_{n+1}(Z_{n} + Q_{n} - r_{n}^{*}\Delta T),$$
with the boundary condition
$$f_{N+1}(Z_{N+1}) = 0 \quad (4.27)$$

$$g_{N+1}(Z_{N+1}) = 0$$
 (4.28)

$$h_{N+1}(z_{N+1}) = 0 \tag{4.29}$$

where r_n^* is the optimal release rate for state Z of stage n.and we use the notation

$$h_n(Z_n) = B(r_n^*) + h_{n+1}(Z_n + Q_n - r_n^*\Delta T),$$
 (4.30)

and

$$g_n(Z_n) = \max\{U(r_n^*), g_{n+1}(Z_n + Q_n - r_n^*\Delta T)\}$$
 (4.31)

We note that another approach to this problem would be to optimize the two objectives simultaneously as in vector optimization [4, 30] or in bicriterion methematical programming [5, 29]. Then we would search

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for efficient (undominated, Pareto optimal) solutions--see also [18, 35, 44, 84].

CHAPTER V

OPTIMAL EXPANSION OF FLOOD CONTROL SYSTEMS

5.1 Introduction

There are basically three structural methods of flood protection [69]. Channel improvements increase the discharge of a stream by increasing the velocity and possibly the channel cross-section and decreasing the distance the water has to travel to reach the outlet of the watershed. This method is generally applicable only in small streams. The construction of levees also increases the discharge by increasing the depth of flow, and thereby the velocity, and by providing a flood way outside the channel. On the lower reaches of long rivers, this method affords the only sure means of flood control. Both of these methods control flooding by hastening the flow of the water from the watershed. The construction of a reservoir, gives protection in an entirely different way, namely, by retarding the flow of the water and limiting the flow to the quantity which the channel can safely carry, thereby preventing floods. There are two main requirements which must be met in establishing a reservoir site for flood protection, one is physical and the other is economic. The physical conditions of the watershed must be such that reservoirs can be constructed of sufficient size to store or retard the excess flood water. The cost of such reservoirs must be reasonable.

The problem of expanding an existing flood control reservoir system

is formulated in mathematical terms by providing functions, equations, and inequalities that represent appropriate criteria and characteristics of the physical system in §5.2 and §5.3. Of interest here is the question of when a reservoir should be added to an existing system and how large the new reservoir should be. We will consider only the latter problem here. The reader is referred to [52, 55, 56] for a discussion of the timing problem. Dynamic programming is proposed as a solution method in §5.4. The methodology is illustrated using real-world data in §5.6.

We shall first consider some of the major assumptions and decisions that must be made in formulating the problem. The most significant assumptions are that:

- (1) The existing system consists of a single reservoir, which can not completely prevent flood damages at the damage center.
- (2) In the series case, the new reservoir is located upstream of the river. In the parallel case, the existing reservoir can handle the flood peak of its own tributary, but the release, when combined with the flow on the parallel tributary, will be sufficient to cause flood damages at the damage center.
- (3) The inflow hydrograph over the live of the project is assumed to be known and deterministic.
- (4) The operating policy of the system has been limited to the optimal policy for flood control.
- (5) The capital cost of a reservoir is a monotone increasing function of its storage capacity [67].

5.2 Formulation of the Objective Function

The criterion that will be used for project justification will be economic evaluation.

Series New Reservoir

The objective function is a function of the storage capacity of the new reservoir. Let

 $f(C_1) = optimal system flood damage cost with a new series reservoir$ $of capacity <math>C_1$

$$= \min \left\{ \max(U_{1}(r_{1,0}) + U_{2}(r_{2,0}), \dots \\ (r_{1,0}, r_{1,1}, \dots, r_{1,N}) \right\}$$

$$U_{1}(r_{1,N}) + U_{2}(r_{2,N}) \right\}$$

$$(5.1)$$

and

$$F = \min_{c_1} f(c_1),$$
 (5.2)

where

M is the maximum possible storage capacity that can be constructed, C_1 is the storage capacity of the new reservoir, and U_i is the damage cost function at damage center i.

We want to find C* such that

 $C^* =$ smallest element in Ω

where $\Omega \subseteq [0, M]$ is such that f(C) = F for $C \in \Omega$

Parallel New Reservoir

Let

 $f(C_1) = optimal system flood damage cost with a new parallel$

reservoir of capacity C1

$$\begin{cases} \min \left\{ \max \left(U(r_{1,n} + r_{2,n}) \right) \right\} \\ \binom{(r_{1,0}, r_{1,1}, \dots, r_{1,N})}{(r_{2,0}, r_{2,1}, \dots, r_{2,N})} \end{cases}$$
(5.3)

and

$$F = \min_{\substack{C_1 \\ C_1}} f(C_1).$$
 (5.4)

We want to find C* such that C* is smallest element in Ω where $\Omega \subseteq [0, M]$ is such that f(C) = F for $C \in \Omega$ and U is the damage cost function at the damage center.

5.3 Constraints

Budgetary Constraints

The capital budget constraint is calculated differently in the private and public sectors. In the private sector, it is considered to be a function of the corporation current assets and current debt level, whereas in the public sector, it depends upon congressional or state water resources appropriation. By the assumption that the total cost of a reservoir is a continuous monotone increasing function of its storage capacity, the capital constraint is equivalent to

$$C_1 \leq M$$
 (5.5)

where M is the maximum possible storage capacity that can be constructed. Operational Constraints

For a new series reservoir, we have

$$Z_{1,n+1} = Z_{1,n} + Q_{1,n} - r_{1,n} \Delta T \qquad n = 0, 1, ..., N, \quad (5.6)$$

$$Z_{2,n+1} = Z_{2,n} + Q_{2,n} + r_{1,n} \Delta T - r_{2,n} \Delta T \qquad n = 0, 1, ..., N, \quad (5.7)$$

$$0 \le Z_{1,n} \le C_1 \qquad n = 0, 1, ..., N, \quad (5.8)$$

$$0 \le Z_{1,n} \le C_2 \qquad n = 0, 1, ..., N, \quad (5.9)$$

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where

 C_1 is the capacity of the new series reservoir, and C_2 is the capacity of the existing reservoir. For a new parallel reservoir, we have

$$Z_{1,n+1} = Z_{1,n} + Q_{1,n} - r_{1,n}\Delta T$$
 $n = 0, 1, ..., N,$ (5.10)

$$Z_{2,n+1} = Z_{2,n} + Q_{2,n} - r_{2,n} \Delta T$$
 $n = 0, 1, ..., N,$ (5.11)

$$0 \le Z_{1,n} \le C_1$$
 $n = 0, 1, ..., N,$ (5.12)

$$0 \le Z_{2,n} \le C_2$$
 $n = 0, 1, ..., N.$ (5.13)

5.4 Dynamic Programming Formulation

For the objective function

$$F = \min_{C_1} f(C_1) \qquad \text{for all } c_1 = c_1$$

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we can find $f(C_1)$ by dynamic programming for each C_1 . In the series case, for fixed C_1 , we have

$$\begin{cases} \min \left\{ \begin{pmatrix} r_{1,0}, r_{1,1}, \dots, r_{1,N} \\ (r_{2,0}, r_{2,1}, \dots, r_{2,N}) \end{pmatrix} \right\}^{\left\{ \max \left(U_{1}(r_{1,n}) + U_{2}(r_{2,n}) \right) \right\}} \\ = W_{0}(Z_{1,0}; Z_{2,0}; C_{1}) \end{cases}$$
(5.14)

where

$$W_{n}(Z_{1,n}; Z_{2,n}; C_{1}) = \min_{\substack{r_{1,n} \in \mathbb{R} \\ r_{2,n} \in \mathbb{R} \\ 2,n}} \{ \min_{\substack{r_{2,n} \in \mathbb{R} \\ 2,n}} (\max(U_{1}(r_{1,n}) + U_{2}(r_{2,n})), M_{1,n}) \}$$

$$W_{n+1}(z_{1,n+1}; z_{2,n+1}; c_1)))$$
, (5.15)

and

$$W_{N+1}(Z_{1,N+1}; Z_{2,N+1}; C_1) = 0.$$
 (5.16)

Hence

$$F = \min \{ W_0(Z_{1,0}; Z_{2,0}; C_1) \},$$
 (5.17)
 C_1

(5.18)

where $Z_{1,0}$ and $Z_{2,0}$ are the initial storage of reservoirs 1 and 2, respectively.

In the parallel case, for fixed C_1 , we have

 $= v_0(z_{1,0}; z_{2,0}; c_1)$

$$\begin{cases} \min & \{\max(U(r_{1,n} + r_{2,n}))\} \\ \{(r_{1,0}, r_{1,1}, \dots, r_{1,N}) \\ (r_{2,0}, r_{2,1}, \dots, r_{2,N}) \end{cases}^{n} \end{cases}$$

where

$$V_n(Z_{1,n}; Z_{2,n}; C_1) = \min_{\substack{r_{1,n} \in \mathbb{R} \\ r_{2,n} \in \mathbb{R}_{2,n}}} \min_{\substack{r_{2,n} \in \mathbb{R}_{2,n}}} (\max(U(r_{1,n} + r_{2,n})),$$

$$v_{n+1}(z_{1,n+1}; z_{2,n+1}; c_1)))$$
, (5.19)

and

$$v_{N+1}(z_{1,N+1}; z_{2,N+1}; c_1) = 0.$$
 (5.20)

Hence

$$F = \min_{\substack{C_1 \\ C_1}} \{V_0(Z_{1,0}; Z_{2,0}; C_1)\},$$
(5.21)

where $Z_{1,0}$ and $Z_{2,0}$ are the initial storage of reservoirs 1 and 2, respectively.

We next state some results which will assist our search for the optimal capacity.

<u>Lemma 5.1</u> For fixed $Z_{1,0}$ and $Z_{2,0}$; $W_0(Z_{1,0}, Z_{2,0}, C_1)$ and $V_0(Z_{1,0}, Z_{2,0}, C_1)$ are both monotone nonincreasing functions of C_1 .

Proof:

For any stage n, the decision set is

$$R_{1,n} = [Lb_{1,n}, Ub_{1,n}],$$
 (5.22)

where

$$Lb_{1,n} = \max\{0, (Z_{1,n} + Q_{1,n} - C_1)/\Delta T\}$$
$$Ub_{1,n} = (Z_{1,n} + Q_{1,n})/\Delta T$$

The Ub_{1,n} is independent of C_1 , but the Lb_{1,n} is a function of C_1 . If $C_1 > C_1$, then

$$Lb_{1,n} = \max\{0, (Z_{1,n} + Q_{1,n} - C_{1})/\Delta T\}$$

$$\leq \max\{0, (Z_{1,n} + Q_{1,n} - C_{1}')/\Delta T\}$$

$$= Lb_{1,n}'$$

Hence

$$R_{1,n} \supseteq R'_{1,n}$$
 (5.23)

Since the minimum value over a larger domain can be no greater than the minimum value over a smaller domain contained in the larger domain, $W_0(Z_{1,0}; Z_{2,0}; C_1)$ and $V_0(Z_{1,0}; Z_{2,0}; C_1)$ both are monotone nonincreasing functions of C. ||

Lemma 5.2 The optimal value of C* exists and can be found by binary search.

Proof:

By Lemma 5.1, $W_0(Z_{1,0}; Z_{2,0}; C_1)$, $V_0(Z_{1,0}; Z_{2,0}; C_1)$ are monotone nonincreasing functions of C_1 , and the cost value is nonnegative. By assumption 5, the capital cost of a reservoir is a continuous monotone increasing function C_1 which is bounded from above. Then we can find F, where

$$\mathbf{F} = \min \left\{ \mathbf{f}(\mathbf{C}_1) \mid \mathbf{C}_1 \in (0, \mathbf{M}) \right\}$$

Hence $\Omega = \{C_1 \in [0, M]; f(C_1) = F\}$ is a compact set and C* exists and can be found by binary search.

5.5 Search for Optimal Capacity of a New Reservoir

(a) Sequential Search

Sequential search is perhaps most the straightforward and simple

of the available search procedures. It consists simply of starting at some value of C, usually the boundary value, and comparing the f(C) one at a time, until either the minimum value of C such that f(C) equal O is found or all values of C have been searched. It is easy to code but may consume excessive amounts of computing time.

(b) Fibonacci Search

Since f(C) is monotone, it is unimodal also. If we compare the value of the function at any two points, then a finite number of points can be excluded from optimality. Clearly by making successive evaluations and comparisons, we do not need to carry out an exhausive search to find the optimal solution. Depending upon where C* lies, search procedures generally require different number of evaluations. Under the assumption of complete uncertainty regarding the value of C*, a reasonable measure of effectiveness is to minimize the maximum number of evaluations. Let k_n be the maximum number of points in the domain so that the optimal point always may be determined with no more than n evaluations. Let K_n equal the maximum of k_n over all search procedures.

$$K_n = K_{n-1} + K_{n-2} + 1$$
 $n > 2$
 $K_1 = 1, \quad K_2 = 2$ (5.24)

From the table of Fibonacci Search [64], we can find the points of first two evaluation as a function of n.

(c) Binary Search

At each comparison in binary search we either find the optimal value in question or eliminate half of the region to be searched. Binary search

requires that f(C) should be monotone. The procedure begins by evaluating the midpoint value of C. If f is nonnegative and f(C) is equal 0, then upper range of the C values ignored and we can evaluate the midpoint of the remaining lower half. This process continues until the search interval of C satisfied some required accuracy. If we let E be the number of evaluations and M be the number of grid points then $E_{max} \ge \log_2(M+1)$. For monotone functions, the number of iterations for each search is given in Table 5.1 [72]. Reference to Table 5.1 reveals that binary search is best for monotone functions. For this reason it was employed in our study. The macro-flow diagram of the overall optimization is on Figure 5.6.

5.6 Example

Consider the East Fork of Silver Creek in Madison County, Illinois which is shown on Figure 5.1. There is a Reservoir on Silver Lake [22]. The dam height is 30 ft. and its storage capacity is 10400 ac.-ft. The topography and geology of Madison County are generally suited to reservoir development. The East Fork of Silver Creek has a potential reservoir site located 1 mile east and 0.5 mile north of Grantfork. This is approximately 3 miles upstream from Silver Lake.

In the spring, when the ground is still frozen, the snow melts and and heavy rains are likely to occur. The flow in East Fork above the Silver Lake can get extremely high, with a typical flood peak being 4030 cfs [79]. The Silver Lake Reservoir can not absorb such flood peak completely since flood stage is 800 cfs.

The inflow hydrograph of a typical heavy storm on East Fork watershed area [79] is given on Table 5.2 and displayed on Figure 5.2. The

Number of	Sequentia	al Search	Fibonacci	Binary	y Search
Points	Max.	Average	Search	Max.	Average
5 10 50 100 1000 10000	5 10 50 100 1000 10000	3 25 50 500 5000	4 5 8 10 15 19	3 4 6 7 10 14	2 3 6 9 13

Table	5.1	Iteration number	for :	Sequential,	Fibonacci,
	-	and Binary Searc	ch.		





													ĺ		
Time (hr)	0	12	24	36	48	60	72	178	96	108	120	132	144	156	169
Input hydrograph of Silver Jake	0	600	1330	2450	3700	02017	3750	3150	2500	1900	1220	800	450	150	0
Input due to the water- shed of potential Res.	0	400	. 930	1550	2400	2530	2530	2000	1600	1250	800	550	350	100	0
Input due to watershed between two reservoir	0	200	450	006	1300	1500	1500	1150	006	650	420	250	100	50	0
															.



29 2	66666666666666666666666666666666666666
	0 0 0 0 0 0 0 0 0 0 0 0 0 0
2	н н н н н н н н н н н н н н
1 28	00000000000000000000000000000000000000
22	00000000000000000000000000000000000000
201	00000000000000000000000000000000000000
5	H 000000000000000000000000000000000000
	0 000000000000000000000000000000000000
50	00000000000000000000000000000000000000
$\frac{3}{122}$	00000000000000000000000000000000000000
50	00000000000000000000000000000000000000
2 250 1	411110 6000000 0000000000000000000000000
500	• • • • • • • • • • • • • • • • • • •
1 20	00000000000000000000000000000000000000
Iteration Capacity	со со в в в в в в в в в в в в в в в в в



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cost function at the East Fork damage center is displayed on Figure 5.3. Since there is a mountain area between these two reservoirs, there is no direct damage cost due to the proposed reservoir releases. The initial flood cost is \$2.98 million. The optimal release policy for Silver Lake Reservoir is displayed on Figure 5.4. The optimal damage cost is \$0.875 million and the flood peak is not completely absorbed.

The grid size used in the search procedure was 100 ac.-ft. This yields an optimal capacity of the proposed reservoir of 937 ac.-ft. The flood damage would be \$0. The computing for the optimal solution is 645 seconds on Northwestern University's CDC 6600. The storage capacity vs. optimal damage cost is displayed on Figure 5.5. The stage time interval ΔT used was 12 hr. The results of the computations for each of the seven iterations aregiven on Table 5.3. For each iteration, column 1 is the release rate of the proposed reservoir with capacity C_1 , for which we want to optimize and column 2 is the release rate of reservoir 2 which already exists. The program listing appears in Appendix A5.





Figure 5.6 Macro-Flow Diagram for Finding the Optimal Capacity of the Additional Reservoir of a Single Reservoir Flood Control System.

CHAPTER VI

SUMMARY

6.1 Summary

Methodologies have been developed for determining the optimal operating policies for single and multiple-reservoir flood control systems in Chapters II and III. Dynamic programming and stochastic dynamic programming were used as the optimization tools. This approach has overcome both the difficulties of the convexity requirements for the linear programming approach and the physical constraint validity of the chance-constrained linear programming approach. Therefore, it offers a potentially powerful method in the analysis of flood control systems.

This paper also modeled synthesized floods as Markov renewal processes in Chapter III. From our study of the effects of changes in the initial storage quantity in Chapter IV, we know that when a flood is coming, it is optimal to lower the water level as much as possible.

The expansion of existing flood control systems was modeled in Chapter V. We determined the minimal capacity of a new reservoir, which furnishes the reduction of predicted flood peak, by dynamic programming, Under the assumption that an optimal operation policy is employed.

Real-world examples have been included both to illustrate and to explicate the results of the theoretical discussions. The use of the derived optimal operating rules in real-time was also simulated using both historical flow rate data and simulated flow rate data. The optimal release policy performed as well or better than both the standard release policy and the routing release policy.

6.2 Directions for Future Research

The uncertainty of future hydrometeorological conditions occurring within a riverbasin confounds the decision-making process when applied to the management of water resources systems. The stochastic nature of water resources suggests continued research in application of stochastic programming methods to the operation of these systems. There are several recommendation can be drawn directly or indirectly from this research. First, efforts should be made to further develop techniques for estimating flood damage cost functions. Second, efforts should be also be directed toward refining and developing techniques for predicting the surface-runoff from precipitation. Third, development of the ability to accurately predict long-range stream flow pattern and distribution should be enhanced. Fourth, development of more efficient algorithms to solve the problems would be desirable. In this regard, the recent incorporation [58] of branch-and-bound methods in dynamic programming to reduce computation and storage requirements which proved useful in a number of problems [56, 57], should be a fruitful avenue of approach.

APPENDIX A

PROGRAM LISTINGS

Appendix A1

Program Listing (Single Reservoir Flood Control System Deterministic Inflow)

PRCGRAM TEST (INPUT, OUTPUT) DIMENSION AN (20) + BN (20) + E (30) + INPT (20) + OPT (20, 300) + RELES (20, 300) + 1FF(20,300),22(20,300) DATA (AN(I), I=1.14)/1000.0, 2000.0, 3000.0, 6000.0, 10000.0, 16000.0, 125000.0.30000.0.30000.0.21000.0.12000.0.5000.0.2000.0.0.0.0.0 DATA (BN(1), 1=1, 14)/-0.02315, -0.02315, -0.02315, -0.06944, -0.0926, 1-0.13891-0.20831-0.1157410.010.208310.208310.16210.0694410.0463/ 16,5,6,7,6,9,7,0,7,1,7,2,7,3,7,4,7,5,7,6,7,7,7,8,7,9,8,0,8,1,8,2, 28.3.8.4/ DATA NTG.DT.DS.DR.DU/14.43200.0.2000000.0.200.0,3800.0/ DO 9 N=1,NTG INFT(1)=AN(N)+DT+0.5+8N(N)+DT++2 9 CONTINUE C=3310560000.0 INITIAL=224952000.0/DS+1 KMAX=C/DS+1 KTC≠0 100 KTG=KTG+1 KK=1 IF (KTG.EQ.NTG) KK=INITIAL 79 CONTINUE P=((KK~1)*DS+INPT(KTG))/DT RUP=AMAX1(P,0.0) BLO=AMAX) (P-C/DT,C.0) TUP=BUP/DR+1 TLC=RLC/DR+2 OPT(KTG,KK)=100000.0 NO 10 I=ILO,IUP CALL CURVE (E, I, DR, DU, VALUE) -----• ~ • • • • • • TF (KTG.NF.1) GO TO 31 TEMP=VALUE JE (OPT (KTG, KK) .LT, TEMP) 60 TO 30 Jzľ OPT (KTG,KK) =TEMP 30 CONTINUE GO TO 10 31 CONTAIN=(KK-1)*DS-(1-1)*DR*DT+INPT(KTG) ISTATE=CONTAIN/DS+1 IF (ISTATE.GE.KNAX) ISTATE=KMAX TEMP=AMAX1 (VALUE, FF (KTG-1, TSTATE)) IF (OPT(KTG)KK) +LT.TEMP) GO TO 14 J≡I OPT (KTG,KK) =TEMP 10 CONTINUE RFLES(KTG,KK) = (J-1) * DRXINPT=(KK-1)*DS FF(KTG,KK)=OPT(KTG,KK) 72(KTG,KK)=XINPT-RELES(KTG,KK)*DT+INPT(KTG) PRINT 2, FF(KTG;KK);RELES(KTG;KK);XINPT;KTG 2 FORMAT (* VALUE = *+F10.5,* + RELESE RATE = *,F12.5,* S 1TATE - =+,F14.2.+ AT STAGE +,I5) IF (KK.GE.KMAX.OR.KTG.FQ.NTG) GO TO 48 KK=KK+1 GO TO 79 4A CONTINUE IF (KTG .NE .NIG) GO TO 100

PRINT 5, FF(NTG,INJTIAL) 5 FOFMAT(* OPTIMAL VALUE *,F14.5) KCT=NTG PRINT 6.KCT, RELES (NTG, INITIAL) OLD=(INITIAL-1)*DS NEW=ZZ(NTG, INITIAL) K27=NEW/DS+1 - · DC 803 I=2,NTG KCT=NTG-I+1 PRINT 6,KCT,RELES(KCT,KZZ) OLD=NEW NEW=ZZ(KCT,KZZ) KZ7=NEW/DS+1 AT STAGE *, I5, * RELESE PATE IS *, F15, 5) 6 FORMAT (* 803 CONTINUE END SUBROUTINE CURVE(E, ID, DR, DU, VALUE) DIMENSION F (30) DIMENSION E(30) L=((ID-1)*DR)/DU+1 a caracter a construction VALUE=E(L)+((E(L+1)-E(L))/DU)*((ID-1)*DP+(L-1)*DU) RETURN FNC -----

Appendix A2

Program Listing (Two Reservoir Series Flood Control System with Deterministic Inflow)

PHOGRAM TEST(INPUT, OUTPUT) DIMENSION AN (2, 17), NEW (2), KZZ (2), OLD (2), FF (20, 20, 20), 1RELES1(20,20,20),RELES2(20,20,20),221(20,20),222(20,20,20), 2KMAX(2), [NI TTAL(2), KK(2), R(2), G(44), H(44), U(2), IH(2), C(2), 7(2), 3DU(2), [NPT(2,16) DATA AN/0.0, 1.0, 300.0, 300.0, 720.0, 720.0, 1300.0, 1300.0, 1950.0, 11950.0.2500.0.2500.0.3100.0.3100.0.3700.0.3700.0.4400.0.4400.0. 24800-0,4800.0,5000.0,5000.0,4700.0,4700.0,3800.0,3800.0,2400.0, 32400.0,1400.0,1400.0,450.0,450.0,450.0,0.0,0.0/ DATA 6/0.0,0.0.1.3.2.6, J.0, 3.7, 4.4, 5.0, 5.5, 6.0, 6.5, 6.9, 7.3, 7.7, 18.1.8.4.8.8.9.2.9.4.9.7.10.0.10.2.10.4.10.6.10.8.11.0.11.2.11.4. 211.5,11.7,11.8,11.9,12.0,12.1,12.2,12.3,12.4,12.5,12.6,12.7,12.8, 312.9.13.0.13.1/ DATA H/0.0,0.0,1.8,3.0,4.0,5.0,5.7,6.5,7.2,7.7.8.3,8.8,9.2,9.7, 110+1,10+4,10,7,11+0,11+3,11+5,11+7,11+9,12+1,12+3,12+4,12+5,12+6, 212,7,12,8,12,9,13,0,13,1,13,2,13,3,13,4,13,5,13,6,13,7,13,8,13,9, 314.0.14.1.14.2.14.3/ DATA NTG, DT, DS, DR, DU(1), DU(2)/16,21600.0,10000000.0,250.0,... 13500.0,5000.0/ 00 9 1=1,2 DO 9 N=1.NTG INPT(1,N)=0.5*(AN(1,N)+AN(1,N+1))*DT 9 CONTINUE C(1)=50000000000 C(2)=1900000000.0 KMAX(1)≈C(1)/05+1 KMAX(2)=C(2)/D5+1 INITIAL(1)=10000000.0/US+1 INITIAL (2) =200000000.0/US+1 KTG=? 100 KTG=KTG+1 KK(1)=] KK(2)=1 IF (KTG.NE.NTG) GO TO 79 KK(1) = INITIAL(1)KK(2)=INITIAL(2) 79 CONTINUE P=((KK(1)-1)*DS+INP[(1+KTG))/DT IUP=MIN1(6000.0.P)/DR+1 ILO= MAX1(P-C(1)/DT.0.0)/DR+2 IF (P-C(1)/OT_LE.0.0) ILO=1 63 IH(1)=ILO $FF(KTG_{KK}(1)_{KK}(2))=10000_{0}$ 62 R(1)=(IR(1)-1)*0R 2(1) = (KK(1) - 1)*DS+INPT(1,KTG) - R(1)*DT CALL CURVE(G, IR(1), DR, DU(1), VALUE) U(1) = VALUEQ=((KK(2)-1)*DS+INPT(2;KTG)+R(1)*DT)/DT LUP= MIN1 (12000.0,0)/DR+1 LLO= MAX1 (Q-c(2)/DT+0.0)/DR+2 IF (G-C(2)/DT.LE.0.0) LLO=1 IR(2)=1.LO 61 R(2)=(IR(2)-1)*DR Z(2) = (KK(2) -1) *DS+INPT(2,KTG) +R(1) *DT-R(2) *DT CALL CURVE(H, IR(2), DP, DU(2), VALUE)

. 9 ; M · · · · · · · ·

U(2)=VALUE IF (KIG.NE.1) GO TO 31 TEMP=U(1)+U(2) IF (FF(KTG,KK(1),KK(2)),LT.TEMP) GO TO 30 PELES1(KTG,KK(1),KK(2))=R(1) RELES2(KTG, KK(1) . KK(2))=R(2) FF(KTG,KK(1),KK(2))=TEMP 30 CONTINUE GO TO 10 GO TO 10 IA=Z(1)/DS+1 IH=Z(2)/DS+1 TEMP=AHAX1(U(1)+U(2),FF(KTG=1,IA,IB)) IF(FF(KTG+KK(1)+KK(2)).LT.TEMP) GO TO 10 RELESL(KTG+KK(1)+KK(2))=B(1)) 31 IA=2(1)/DS+1 RELESI (KTG, KK (1), KK (2)) = R (1) RELES2(KTG, KK(1), KK(2))=R(2) FF(KTG,KK(1),KK(2))=TEMP 16 CONTINUE IF (1H (2), GE, LUP) GO TO 51 IR(2)=IR(2)+1 GO TO 61 51 IF (IR (1).GE.IUP) GO TO 52 IR(1)=IR(1)+1 GO TO 62 52 CONTINUE CONTINUE XIN1=(KK(1)-1)+DS XIN2=(KK(2)-1)*05 221 (KTG,KK(1),KK(2))=XIN1-RELES1(KTG,KK(1),KK(2))*UT+INPT(1,KTG) 222(KTG,KK(1),KK(2))=XIN2-RELES2(KTG,KK(1),KK(2))*DT 1+RELES1(KTG, KK(1), KK(2)) +DT+INPT(2, KTG) PRINT 3+KTG+XIN1+XIN2+FF(KTG+KK(1)+KK(2))+RELES1(KTG+KK(1)+KK(2))+... 1RELE52(KTG, KK(1), KK(2)) 3 FORMAT(* STAGE=*+13+* STATE 1=*+F15+1+* STATE 2=*+F15+1+/+ 120X, + VALUE=+, F9,6,+ RELES RATE 1=+, F10.1++ RELES RATE 2=+. 2F10.1) IF (KK (2). GE. KMAX (2). OR. KTG. ED. NTG) GO TO 53 KK(2)=KK(2)+1 GO TO 63 53 IF (KK (1) .GE .KMAX (1) .OR .KTG .EQ .NTG) GO TO 54 KK(1)=KK(1)+1 KK(2)=1 GO TO 79 54 CONTINUE IF (KIG.NF.NTG) GO TO 100 IF(KIG.NF.NTG) GO TO 100 PRINT 5+FF(NTG+INITIAL(1)+INITIAL(2)) FURMAI(* CPTIMAL VALUE=*+F9.6) 5 FURMAT(* KCT=NTG PHINT 6, KCT, RELESI (NTG, INITIAL (1), INITIAL (2)), IRELES2 (NTG, INITIAL(1), INITIAL(2)) OLD(1)=(INITIAL(1)+1)+DS OLD (2) = (INITIAL (2) -1) +DS $NEW(1) = ZZ1(NTG \cdot INITIAL(1) \cdot INITIAL(2))$ $NEW(2) = ZZ2(NTG \cdot INITIAL(1) \cdot INITIAL(2))$ KZZ(1) = NEW(1)/US+1K22(2)=NFW(2)/US+1 DO 803 1=2,NTG

%
KCT=NIG=I+1
PHIN[6+KCT+RELES1(KCT+KZZ(1)+KZZ(2))+RELES2(KCT+KZZ(1)+KZZ(2))
OLD(1)=NEW(1)
OLD(2)=NFW(2)
NEW(1)=Z71(KCT+KZZ(1)+KZZ(2))
NEW(2)=Z22(KCT+KZZ(1)+KZZ(2))
KZZ(1)=NEW(1)/US+1
KZZ(2)=NEW(2)/US+1
6 FORMAT(* AT STAGE *+I3+* PELES PATE 1= *+F10+1+*RELES PATE 2=*
1+F10+1
803 CONTINUE
END

SUBROUTINE CHRVE(Y,ID,DR,DU,VALUE) DIMENSION Y(44) L=((IU-1)*DR)/DU+1 VALUE=Y(L)+((Y(L+1)-Y(L))/DU)*((ID-1)*DR+(L-1)*DU) RETURN END

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Appendix A3

Program Listing (Two Reservoir Parallei Flood Control System with Deterministic Inflow)

PROGRAM TEST(INPUT, OUTPUT)

DIMENSION AN(2,19), NEW(2), KZZ(2), OLO(2), FF(20,20,20), 1RELES1(20,20,20),FELES2(20,20,20),ZZ1(20,20,20),ZZ2(20,20,20), 2KMAX(2), INITIAL(2), KK(2), R(2), G(29), IR(2), C(2), Z(2), INPT(2,20) DAFA AN/625.0,625.0,650.0,650.0,750.0,750.0,750.0,750.0,750.0,750.0, 1750.0,800.0,900.0,1100.0,1400.0,1300.0,2100.0,1700.0,2550.0, 22200.0,3000.0,2700.0,3100.0,3250.0,2950.0,3500.0,2600.0,3000.0, 32600.0,2400.0,1800.0,1700.0,1400.0,1250.0,1250.0,800.0,1000.0, 4500.3,725.0/ OATA G/0.0,0.0,0.0,0.0,0.0,0.75,1.3,1.8,2.2,2.5,2.7,2.9,3.1,3.3, 13.5,3.7,3.9,4.0,4.1,4.2,4.3,4.4,4.5,4.6,4.7,4.8,4.9,5.0,5.1/ DATA NTG, OT, DS, DR, OU/18, 43200.0, 20000000.0, 100.0, 500.0/ S,1=1 0 00 00 9 N=1,NTG INPT([,N)=0.5*(AN(I,N)+AN(I,N+1))*OT 9 CONTINUE C(1)=200000000.00.C C(2)=140000000.0 KMAX(1) = C(1) / OS + 1KMAX(2)=C(2)/DS+1 INITIAL(1)=3 INITIAL(2)=2 KTG=0 100 KTG=KTG+1 KK(1)=1 KK(2)=1 IF(KTG.NE.NTG) GO TO 79 KK(1)=INITIAL(1) KK(2)=INITIAL(2) **79 CONTINUE** P=((KK(1)-1)*OS+INPT(1,KTG))/OT IUP=P/DR+1 ILO= MAX1(P-C(1)/CT,0.0)/DR+2 IF(P-C(1)/0T.LE.0.0) ILO=1 63 CONTINUE 0 = ((KK (2) - 1) + 0S + INPT (2, KTG))/01 LUP=0/02+1 LLO= MAX1(Q-C(2)/DT,0.3)/OR+2 IF (G-C(2)/DT.LE.0.0) LLO=1 IR(1)=ILC FF:(KTG,KK(1),KK(2))=10000.0 62 R(1)=(IR(1)-1)*OR IR(2)=LLC Z(1)=(KK(1)-1)*OS+INPT(1,KTG)-R(1)*DT 61 R(2)=(IR(2)-1)*DR Z(2)=(KK(2)-1)*OS+INPT(2,KTG) -q(2)+OT JJ=(R(1)+R(2))/DR+1 CALL CURVE(G,JJ,CP,DU,VALUE) U=VALUE IF (KTG.NE.1) GO TO 31 TEMP=U IF (FF(KTG,KK(1),KK(2)).LT.TEMP) GO TO 30 RELES1(KTG,KK(1),KK(2))=Q(1) RELES2(KTG,KK(1),KK(2))=R(2) FF(KfG,KK(1),KK(2)) = TEMP

```
30 CONTINUE
    GO TO 10
 31 IA=Z(1)/0S+1
    IB=Z(2)/0S+1
    TEMP=ANAX1(U,FF(KTG-1,IA,IB))
    IF (FF (KTG, KK(1), KK(2)).LT.TEMP) GO TO 10
    RELESI(KTG, KK(1), KK(2)) = R(1)
    RELES2(KTG,KK(1),KK(2))=R(2)
    FF(KTG,KK(1),KK(2)) = TEMP
10 CONTINUE
    IF(IR(2).GE.LUF) GC TO 51
    IR(2) = IR(2) + 1
    GC TC 61
 51 IF(IR(1).GE.IUP) GO TO 52
    IR(1) = IR(1) + 1
    GC TO F2
 52 CONTINUE
    XIN1=(KK(1)-1)+0S
    XIN2=(KK(2)-1)*DS
    ZZ1(KTG,KK(1),KK(2))=XIN1-PELES1(KTG,KK(1),KK(2))*OT+INPT(1,KTG)
    ZZ2(KTG,KK(1),KK(2))=XIN2-RELES2(KTG,KK(1),KK(2))*DT
   1+INPT(2,KTG)
   PRINT 3,KTG,XIN1,XIN2,FF(KTG,KK(1),KK(2)),RELES1(KTG,KK(1),KK(2)),
   19ELES2(KTG,KK(1),KK(2))
  3 FORMAT (* STAGE=*,I3,*
                              STATE 1=*,F15.1,* STATE 2=*,F15.1,/,
  120X,* VALUE=*,F9.6,* RELES RATE 1=*,F1J.1,* RELES RATE 2=*,
   2F10.1)
    IF (KK(2).GE.KMAX(2).OR.KTG.E0.NTG) G0 T0 53
    KK(2)=KK(2)+1
    GC TO 63
 53 IF (KK(1).GE.KMAX(1).OP.KTG.EQ.NTG) GO TO 54
    KK(1) = KK(1) + 1
    KK(2)=1
    GC TO 79
 54 CONTINUE
    IF (KTG.NE.NTG) GO TO 100
    PRINT 5,FF(NTG,INITIAL(1),INITIAL(2))
  5 FORMAT (*
               OPTIMAL VALUE= *,F9.6)
    XCT=NTG
    PPINT 6,KCT, PELESI(NTG, INITIAL(1), INITIAL(2)),
   1RELES2(NTG, INITIAL(1), INITIAL(2))
    OLD(1) = (INITIAL(1)-1)*0S
    OLD(2) = (INITIAL(2)-1)+0S
    NEW(1) = ZZ1(NTG, INITIAL(1), INITIAL(2))
    NEW(2)=Z72(NTG,INITIAL(1),INITIAL(2))
    KZ7(1)=NEH(1)/DS+1
    KZZ(2) = NEW(2)/DS+1
    00 833 I=2,NTG
    KCT=NTG-I+1
    PRINT 6,KGT,RELES1(KGT,KZZ(1),KZZ(2)),RELES2(KGT,KZZ(1),KZZ(2))
    CLD(1)=NEW(1)
    0LD(2) = NEW(2)
    NEW(1) = 721 (KCT , K77(1), K77(2))
    NEW(2)=772(KCT,K77(1),K77(2))
    KZ7(1)=NEW(1)/DS+1
    KZ7(2)=NEW(2)/DS+1
                AT STAGE *, I3, * RELES RATE 1= *, F1G.1, *RELES RATE 2=*
  6 FORMAT( *
   1,F10.1)
803 CONTINUE
    END
    SUBROUTINE CURVE(Y, ID, OR, DU, VALUE)
    DIMENSION Y(29)
    L=((IO-1)*OR)/OU+1
    VALUE=Y(L)+((Y(L+1)-Y(L))/DU)*((ID-1)*DR-(L-1)*DU)
    RETURN
    END
```

Appendix A4

Program Listing (Single Reservoir Flood Control System with Stochastic Inflow)

1. ANT NOT DANG (INPUT + UN (PUT) All she also , 16 (r(s)+i=1,1c)/0.0223,00.0009+0.0919,0.1498,0.1910,0.1915, 19,1496,5,6919.0,396,0,0220/ ... Jain(all)+[=1+1c)/2200+2200+2200+2200+2208+2208+069+669+609+ 1565,1,597 3574(L(I),I=T+32)/0.570,070,071,872,932/94.494.945,3+5.3+5.7+6.1+6.3+ 10.7.4.1.6.9.1.0.1.1.1.1.2.1.3.1.4.1.5.1.0.1.7.1.0.1.9.0.1.9.0.0.1.0.2. 60.3.8.4.8.45.05.51 SATA GEGET, DS. DE, DD, 19713, 432300, 03500000000, 50.0, 150.0, 10/ J=10000000000.0 ANAX=C/US+1 °≖413*1 4 = 1-1 5 confrance r ≈ i 11 (11.1.4.1Tu) 60 10 49 5=10000.0 30 21 J=1+[₽⁻⁻ $t = (x_{n-1}) * U_{2} + (w(N) * Z/I_{1}) * U_{2} U_{1}$ $t_{U_{2}} = (J) = max_{1} (Y/U_{1}, U_{2}) + U_{1} U_{2} + U_{1}$ 1[5())=M3X1((Y+()/UT+0+0)/UK+2 <1 CONTINUE. 16-160447 ; 35 (ChT11.52 51F4=5.**.** にん=り. vo sitedateste 1 K= [1 17(10.LT.(LH(J))_AR=1LH(J) 17(12.51.104(J))/MR=104(d) CALL CURVE (LOCKODKODUOVALUE) arm=76L01#P(J)+Sou シ᠃ᡶ≿▬(ᡧᡊ᠆ᡝ᠈᠉ᡅᠫ᠆ᡧᡰᢂ᠆ᡰ᠈᠉ᠣᢊ᠉ᠮᠮ+(ᡎ(ᡢ)᠉ᢄᠰ᠋₽᠈᠉┙᠉ᡅᠮ᠈ᡣᡅᠫ+ᡫ᠂᠃᠃ とメニデ(i + i + i ∠) * P (J) + i X SULLING AC ・二代にLiet(パッパ)。 AT (M.LU.A) UP1 (BON)=1K -17 (in.v.,fun(1+))-60-10-40dia=Ix+1 50 10 30 + (* + · (* • * ;) = 0 and a second STATL=(K+1)#US - たこさら=(こと「(河・水)ール)やした。 TRIP JUST STATESKELLSSF (NSK) // FERMAL(# STAGE=*,13,# STATL=*,FLZ+1,* RELES RATE=*, 1F1J.1,* STAGE OFT VALUE=*,F1J.5) +) F (1;+6)=€. ST IF (K. UL KHAK) 60 TO 50 - +=K+1 OU TU 5 50 IF (N.L...) GO TO 60 ្រូវស្រីលា ម SU CONTINUE GUER SUFINE CURVE (E, IN, DR, DU, VALUE) 1111.15100 E(3c) L = ((10-1)*0R)/00+1 $\forall A \sqcup \forall c = c (c) + ((c (L+1) - E(L)) / \partial b) * ((U-1) * \partial R - (L-1) * \partial U)$ RETURN LIN
Appendix A5

Program Listing (Simulation)

PROGRAM SIMU(INPUT; OUTPUT) DIMENSION NR(20,40), MR(20,40), Z(20), IZ(20), E(32), X(20), IX(20), 1W(20) + IW(20) , U(12) , MEAN(12) , VAR(12) , COR(12) , R(15) DATA US,NTG,NST,C,UT,DU/5000000.0,12,21,100000000.0,432000.0, 1750.0/ DATA (E(I),I=1,32)/0.0,0.0,0.0,1.8,2.9,3.7,4.4,4.4.9,5.3,5.7,6.1,6.3, 16.5.6.7.6.9.7.0.7.1.7.2.7.3.7.4.7.5.7.6.7.7.7.1.8.7.9.8.0.8.1.8.2. 28.3,8.4,8.45,8.5/ DATA MEAN/6#2088,6#2364/ UATA VAR/6*1044.0,6*1182.0/ DATA COR/6+0.3137,6+0.4142/ Z(1)=500000000.0 $I_{(1)} = 11$ X(1) = Z(1)W(1) = L(1)IX(1) = IZ(1)IW(1) = IZ(1)DO 10 J=1+21 READ 1, (NR(I,J), I=1,12) 10 CONTINUE D0 11 J=1+21 READ 1, (MR([,J),I=1,12) 11 CONTINUE 1 FORMAI(1214) T=1500*DT A=T+C S=MEAN(1) DO 99 K=1+100 PRINT 5+K PRINT 3 5 FORMAT (* RUN #) THE #, I4, # 3 FORMAI (SX+*STAGE*, 5X, *INPUT*, 5X, *DEPEN*, 5X, *INDEP*, 5X, *STAND*) L¤1 IY=0JY=0 KY≈0 CALL GGNOR (4, 15, R) DO 21 M=1,NTG IF(M.NE.1) GO TO 51 Q(M) = MEĀN(M) * (1-COR(M)) * COR(M) *S +VAR(M) *R(M) *(1-COR(H) **2) 1**0.5 40 IF (Q(M).GT.0) GO TO 42 Q(M)=MEAN(M)*(1-COR(M))+COR(M)*Q(M)+VAR(M)*R(NTG+L)*(1-COR(M)**2) 1**0,5 L=L+1 GO TO 40 51 Q(M)=MEAN(M)*(1-COR(M))+COR(M)*Q(M-1)+VAR(M)* R(M)*(1-COR(M)**2) 1**0.5 IF (Q(M) .GT.0.) GO TO 42 Q(M-1)=Q(M) R(M)=R(NTG+L) L=L+1 GO TO 51 42 BU=2 (M) /DT+Q (M) AU=X (M) /DT+Q (M)

END SUBROUTINE CURVE(E,IY,DU,VALUE) DIMENSION £(35) L=IY/DU+1 VALUE=E(L)+((E(L+1)-E(L))/DU)*(IY-(L-1)*DU) RETURN END

```
AL=AMAX1 (AU+C/DT,0.0)
   IR=NR(M,IZ(M))
   JR=MR (M + IX (M))
   IF (IR LT. BL) IR=BL
   IF(IR.GT, BU) IR=BU
   2(M+1)=2(M)+0(M)+DT=IR+UT
   IZ(M+1) = Z(M+1)/DS+1
   IY=MAXO(IY,IR)
   IF (JR.LT.AL) JR=AL
IF (JR.GT.AU) JR=AU
   TU491 - X (M) +Q (M) +DT-Jg+UT
   IX(H+1)=X(M+1)/DS+1
   JY=MAXO(JY,JR)
   B=W(M)+Q(M)+DT
   IF (8. LE.T) GO TO 31
   IF (B.GT.T.AND.B.LE.A) GO TO 32
   KR= (8-C)/DT
   W(M+1)=C
   GO TO 33
32 KR=T/UT
   ₩(M+1)=8-T
   GO TO 33
31 KR=8/07
   W(M+1)=0
33 CONTINUE
   IW(M+1)=W(M+1)/DS+1
   KY=MAXU (KY+KR)
   PRINT 2.M.Q(M).IR, JR, KR
 2 FORMAT(5x,13,8x,F8,1,3x,16,3x,16,3x,16)
21 CONTINUE
   CALL CURVE (E, IY, DU, G)
   CALL CURVE (E, JY, DU+H)
  CALL CURVE (E,KY,DU,F)
   PRINT 4,GOH,F
 4 FORMAT (*
                COST +, 17X, F7.5, 3X, F7.5, 3X, F7.5)
   S=Q(NTG)
99 CONTINUE
```

BL=AMAX1(BU-C/UT+0.0)

Appendix A6

Program Listing(Optiaml Capacity of Series New Reservoir)

	PHOGRAM OPT (INPUT, OUTPUT)
	INTEGER PIGHT, CENTER
	x=43560.n
	D=30000000.0
	DR=300.0
	M=10000
	W=M*X
	LEFT=0
	RIGHT=M
	CALL DAM (W, V, D, DR)
305	
500	FORMAT(O) = CAPACITY 9,17,0 COST 9,F10.7)
500	
501	CONTINUE
	IF (RIGHT-LEFT.GT.400) GO TO 502
	D=6700000.0
	DR=150.0
502	CONTINUE
	IF (RIGHT-LEFT.GT.200) GO TO 503
	D=3390000.0
	DR=100.0
503	CONTINUE
	CENTER=(LEFT+RIGHT)/2
	WECENIER#X
	CALL DAM (W, V, D) JDR)
	$\begin{array}{c} PRIN \\ JU(J(CN) CN) CN(CN) CN \\ TF(V \in O, $
	IFFICENTED
201	BIGHT=CENTER
-0-	GQ 10 50A
400	PRINT 301 ,RIGHT,V
301	FORMAT(* OPYIMAL CAPACITY IS*, 17,* DAMAGE COST IS*, F11.7)
	END
	· · · · · · · · · · · · · · · · · · ·
	SUBROUTINE CURVE (Y+ID+DR+DU+VALUE)
	DIMENSION Y(A4)
	L = ((IU - 1) * 0R) / 0U + 1
	VALUE=Y(1)+((Y(L+1)-Y(L))/DU)*((ID-1)*DR+((-1)*DU)
	RETURN
	END

SUBROUTINE CAM (W+V+D+DR) DIMENSION AN (2,15) , NEW (2), KZZ (2), OLD (2), FF (20,20,20), 1RELES1(24,20,20), RELES2(20,20,20), 221(20,20,20), 222(20,20,20), 2KMAX(2), INITTAL(2), KK(2), R(2), G(20), U(2), IR(2), C(2), 7(2), 3DU(2) + INPT(2,16) DATA AN/0.0.0.0.100.0.50.0.350.0.100.0.550.0.250.0.800.0.450.0. 11250.0,650.0,1600.0,900.0,2000.0,1150.0,2350.0,1400.0,2530.0, 21500.0,2400.0,1300.0,1550.0,900.0,920.0,450.0,400.0,200.0,0.0,0.0/ DATA 0/0.0,0.0,1.0,1.8,2.5,3.0,3.2,3.4,3.6,3.8,4.2,0.4.1.4.2.4.3, 14.4,4.5,4.6,4.7,4.8,4,9/ DATA NTG+UT+NS+DU(1)+DU(2)/14,43200.0,30000000.0+800.0+800.0/ C(1)=W C(2)=45000000000 S.1=1 6 00 DO 9 N=1+NTG INPT(I,N)=0.5*(AN(I,N)+AN(I,N+1))*DT 9 CONTINUE C(1) = WKMAX(1)=C(1)/0 +1 KMAX(2) = C(2) / D5 + 1INITIAL(1)=2 INITIAL(2)=3 KTG=0 100 KTG=KTG+1 KK(1)=1 KK(2)=1 IF(KTU.NF.NTG) GO TO 79 KK(1)=INITIAL(1) KK(2)=INTIAL(2) 79 CONTINUE P=((KK())-1)*D +INPT(1*KTG))/DT IUP=MIN1 (3500.0,P)/DR+1 ILO= MAX1(P-C(1)/DI,0.0)/DR+2 IF(P-C(1)/DT,LE.0.0) ILO=1 63 IR(1)=1L0 FF(KTG,KK(1),KK(2))=10000,062 R(1)=(IR(1)-1)*DRZ(1)=(KK(1)-1)*D +INPT(1,KTG)-R(1)*DT U(1)=0. Q=((KK(2)~1)*DS+INPT(2,KTG)*R(1)*DT)/DT $LUP = MIN_1 (5000.0,0) / DR + 1$ LLO = MAX1 (0-C(2) / D1.0.0) / DR + 2 IF (0-C(2) / DT.LE.0.0) LLO=1 IR(2)=LL0 61 R(4)=(IR(2)-1)*DR Z(2) = (kK(2)-1)*DS+INPT(2,KTG)+R(1)*DT-R(2)*DT CALL CURVE (G, IR (2), DR, DU (2), VALUE) U(2) = VALHEIF (KTG.NE,1) GO TO 31 TEMP = U(1) + U(2)IF (FF(KTG,KK(1),KK(2)).LT.TEMP) GO TO 30 RELES1(KTG,KK(1),KK(2))=R(1) RELES2(KTG,KK(1),KK(2))=R(2) FF(KTG,KK(1),KK(2))=TEMP CONTINUE ---- 30 CONTINUE

GO TO 10 31 IA=Z(1)/D +1 IB=2(2)/nS+1 TEMP=AMAX1 (U(1)+U(2),FF(KTG-1,IA,IP)) IF(FF(KTG,KK(1),KK(2)),LT.TEMP) GO TO 10 RELES1 (KTG, KK (1), KK (2)) =R (1) RELE52(KTG, KK(1), KK(2))=R(2) FF (KTG, KK (1), KK (2))=TEMP 10 CONTINUE IF (IR (2) GE .LUP) GO TO 51 IE(5)=IE(5)+1 GO TO 61 51 IF (IR (1) .GE. IUP) GO TO 52 IR(1) = IR(1) + 1GO TO 62 52 CONTINUE XIN1=(KK(1)-1)*D XIN2=(KK(2)-1)*DS 221 (K (G, KK (1), KK (2))=XIN1-RELFS1 (KTG, KK (1), KK (2))*DT+INPT (1, KTG) 222(KTG,KK(1),KK(2))=XIN2-RELES2(KTG,KK(1),KK(2))*DT 1+RELES1 (KTG, KK(1), KK(2)) *DT+INPT(2, KTG) 3 FORMAT(* STAGE=*,13,* STATE 1=4.F15.1,* STATE 2=*,F15.1,/, 120X,* VALUE=*,F9.6,* RELES PATE 1=*,F10.1,* RELES RATE 2=*, 2F10.1) IF (KK (2) . GE . KMAX (2) . OR . KTG . EO . NTG) GO TO 53 KK(2)=KK(2)+1 GO TO 63 53 IF (KK(1) . GE . KMAX(1) . OR . KTG . ED . NTG) GO TO 54 KK(1)=KK(1)+1 KK(2)=1 GO TO 79 54 CUNTINUE IF (KTG.NF.NIG) GO TO 100 PRINT 5.FF (NTG.INITIAL(1), INITIAL(2)) 5 FORMAT(* OPTIMAL VALUE= *,F9.6) V=FF (NTG.INITIAL(1), INITIAL(2)) PRINT 6, KCT, PELESI (NTG, INITIAL (1), INITIAL (2)), IRELES² (NTG, INITIAL (1), INITIAL (2)) OLD(1)=(INITIAL(1)-1)*D OLD(2)=(INITIAL(2)-1)*DS NEW(1)=Z71(NTG, INITIAL(1), INITIAL(2)) NEW(2)=772(NTG, INITIAL(2), INITIAL(2)) NEW(2)=Z72(NTG, INITIAL(1), INITIAL(2)) KZZ(1)=NFW(1)/D +1 KZZ(2)=NFW(2)/DS+1 D0 803 I=2,NTG KCT=NIG-I+1 PHINT 6,KCT, RELES1 (KCT, KZZ (1), KZZ (2)), RELES2 (KCT, KZZ (1), KZZ (2)) OLD(1)=NFW(1) OLD(2) = NFW(2) $NEW(1) = Z_{7}1(K_{C}T_{3}KZ_{2}(1)_{3}KZ_{2}(2))$ NEW (2) = Z72 (KCT + KZZ (1) + KZZ (2)) KZZ(1) = NFW(1)/U + 1KZZ(2)=NFW(2)/DS+1 6-FORMAT(*- AT STAGE *,13,* RELES RATE 1= *,F10,1,*RELES RATE 2=* 1+F10.1) 803 CONTINUE RETURN END

APPENDIX B

FLOW RATE DATA OF KASKASKIA RIVER AT

SHELBYVILLE, ILLINOIS

Appendix B

Flow Rate Data of Kaskaskia River at Shelbyville, Ill.

KASKASKIA RIVER AT SHELBYVILLE, ILL.

LOCATION.—Chain gage in sec. 8, T. 11 N., R. 4 E. of the third principal meridian, at highway bridge in the eastern edge of Shelbyville, Shelby County, Ill., just above the Chicago & Eastern Illinois Railway bridge and Big Four Railroad bridge.

DRAINAGE AREA.--1,030 square miles.

RECORDS AVAILABLE.—Feb. 25, 1908 to Sept. 30, 1912; Nov. 1, 1912 to Dec. 31, 1912; Aug. 11, 1914 to Dec. 5, 1914, when finally discontinued.

- EXTREMES.—Maximum discharge recorded, 10,600 sec.-ft., May 8, 1908 (gage height, 25.8 feet); minimum discharge recorded, 0.2 sec.-ft., Oct. 4-7, 1914 (gage height, 4.7 feet).
- REMARKS.—Stage-discharge relation affected by ice during winter months. Shifting channel but measuring section was at a pool and control was considered permanent. During high water discharge relation was probably affected by backwater caused by drift lodging at the two railroad bridges below the gaging station.

						-						
Day.	Jan.	Feb.	Mar.	Apr.	May.	June.	July.	Aug.	Sept.	Oct.	Nov.	Dec.
1905 1 2 3 4 5			4,200 4,480 4,990 4,760 4,540	2,400 3,000 2,500 2,300 1,580	2,500 2,300 1,810 7,820 8,780	1,030 956 918 1,110 956	256 224 256 224 224 224	90 73 73 58 58	34 25 25 18 18	25 15 18 18 13	8 10 10 13 13	34 34 34 25 25
6 7 8 9 10			4,420 4,650 4,870 5,100 5,330	1,450 1,400 2,350 2,450 2,500	8,480 8,720 10,600 9,260 7,820	846 738 606 810 956	256 358 362 426 324	45 45 45 58 73	18 18 13 13	13 13 13 13 13	18 18 18 18 18	25 25 25 25 34
11 12 13 14 15			5,220 4,650 4,090 3,760 3,430	2,350 2,250 2,150 2,100 1,650	7,220 6,800 6,380 5,850 5,330	918 774 666 630 562	290 290 224 192 130	73 58 73 58 58	10 10 10 10 10	13 13 13 13 13	18 18 18 18 18	34 34 34 34 34
16 17 18 19 20			3,160 2,550 2,650 1,630 1,680	1,450 1,150 1,110 1,070 1,030	4,990 4,700 4,200 4,140 3,480	460 425 355 324 392	130 109 90 160 160	45 45 45 31 34	10 13 13 13 13	10 10 10 10 10	25 25 25 25 25	34 25 25 25 25
21 22 23 24 25		3,760	1,490 1,450 1,400 1,360 1,310	994 956 918 1,650 1,910	3,000 3,4%0 3,160 2,900 2,600	358 358 324 290 290	160 130 130 109 90	34 25 25 25 25	13 10 10 10 10	10 10 5 5 5	25 34 34 45 45	25 25 25 25 18
26 27 28 29 30 31		4.900 5,330 5,100 4,370	1,190 1,070 994 1,150 1,150 1,070	2,200 2,500 2,700 2,500 2,550	2,350 2,250 1,630 1,540 1,360 1,190	256 224 160 192 224	90 73 73 5× 73 90	25 25 18 18 25 25	10 10 13 18 25	10 10 10 5 8 8	58 58 45 45 45	18 18 25 25 34 34

DAILY DISCHARGE, IN SECOND-FEET, OF KASKASKIA RIVER AT SHELBYVILLE, ILL., FOR 1908-1912.

Day.	Jan.	Feb.	Mar.	Apr.	May.	June.	July.	Aug.	Sept.	Oct.	Nov.	Dec.
1909 1 2 3 4 5	34 34 34 34 34	•40 90 90 109 160	3,000 2,600 2,100 1,770 1,490	302 392 355 358 358	3,260 1,860 1,810 1,810 1,770	4,140 3,820 3,720 3,380 2,850	426 290 290 256 2,150	358 324 256 224 192	34 34 34 34 25	18 18 15 15 18	109 90 90 73 73	392 358 358 324 324
6 7 8 9 10	30 30 25 25 25	192 324 460 494 562	1,310 1,360 1,450 1,540 1,440	4,700 4,650 4,200 3,700 3,100	1,580 1,540 1,580 3,600 3,600	2,550 2,350 2,050 1,860 2,550	2,700 5,500 5,740 5,100 3,210	130 109 90 90 90	25 18 18 25 25	13 13 13 13 13	73 90 90 73 73	290 230 230 216 200
11. 12. 13. 14. 15.	20 20 20 23 25	630 774 882 1,030 1,110	1,310 1,230 1,190 1,070 956	3,000 4,040 6,740 7,550 6,950	3,540 3,430 2,550 2,420 2,100	2,550 2,050 1,770 1,720 1,810	2,550 2,550 3,540 3,030 2,850	90 73 73 5S 45	25 18 18 18 18	13 13 13 13	58 90 90 109 109	a200 426 665 1,070 1,490
16 17 18 19 20	25 25 30 34 45	1,190 1,490 2,100 2,650 3,210	845 1,030 735 702 666	$\begin{array}{r} 6,440 \\ 5,910 \\ 5,500 \\ 4,930 \\ 4,090 \end{array}$	1,680 1,310 1,150 1,030 882	1,540 1,910 1,450 1,150 774	2,550 2,900 3,450 3,600 3,260	34 34 25 25 25	18 - 18 18 18 18	13 13 25 58 73	426 392 666 666 666	1,580 1,540 1,450 1,150 1,000
21 22 23 24 25	58 73 90 90	3,350 3,430 3,480 3,760 3,820	630 562 494 562 596	5,620 6,440 6,950 5,560 5,450	528 702 666 596 702	630 596 528 528	2,550 2,450 1,860 1,450 666	25 18 13 13 13	18 18 18 58 58 58	 73 58 130 130 130 	738 956 918 852 846	•\$50 •700 •600 •520 •470
2ð 27 28 29 30 31	73 73 73 50 250 250	3,820 3,600 3,200	630 596 494 494 460 426	5,220 4,590 3,609 3,540 2,600	1,360 2,200 2,100 2,200 2,900 3,760	738 702 666 562 528	702 630 528 426 392 358	25 18 34 34 45	34 25 25 25 25	130 130 130 130 130 130	702 630 562 494 426	•400 •350 •300 •250 •220 •290
1910 1 2 3 4 5	250 250 300 300 300	774 702 738 702 702	3,700 3,160 3,260 3,350 3,430	290 290 256 256 256	302 358 596 738 882	1,450 1,270 1,110 994	735 774 882 494 774	1,150 937 630 375 290	58 90 66 90 100	203 176 160 176 224	66 90 100 82 90	1,420 1,050 720 684 506
6 7 8 9 10	■300 ■300 ■300 ■350 ■400	666 630 666 596 562	2,400 2,150 1,810 1,580 1,360	256 224 224 160 160	846 1,230 1,540 1,580 1,230	846 702 596 562 494	630 545 409 290 224	224 273 160 100 109	1,150 1,630 3,180 2,500 1,840	579 975 1,070 937 774	82 90 82 66 66	511 341 341 375 392
11 12 13 14 15	500 738 1,860 3,160 3,600	528 460 324 426 426	1,150 904 994 774 666	160 192 192 192 192	1,860 2,200 2,150 1,860 1,680	494 426 392 324 290	224 176 409 392 358	90 82 90 66 73	1,450 1,030 756 665 777	613 511 460 392 307	66 82 90 66 66	341 307 341 375 358
16 17 18 19 20	3,380 2,900 4,250 5,330 6,260	426 358 290 358 324	630 596 562 528 494	324 358 358 358 358	1,490 1,270 1,030 582 846	290 256 256 256 226 224	774 1,050 1,470 1,270 1,030	494 720 994 756 630	426 341 250 208 192	3 92 273 240 224 208	66 66 73 66 66	307 273 203 203 224
21 22 23 24 25	4,540 3,650 3,380 2,700 2,200	494 460 400 392 358	426 426 426 302 392	358 324 324 290 290	774 1,070 3,700 4,310 4,870	160 109 109 109 109	846 579 409 358 273	358 307 290 545 420	176 145 160 176 358	256 208 176 192 145	66 66 73 66 66	176 176 208 341 200
26 27 28 29 30 31	1,860 1,490 1,230 1,110 994 882	358 2,550 3,350	358 358 358 324 324 290	290 324 355 426 392	4,650 3,260 2,850 2,700 2,300 1,810	109 256 392 494 666	160 160 176 256 756 1,110	290 208 192 100 90 66	307 273 256 240 224	160 145 145 130 100 90	66 3,510 3,600 2,420 1,720	208 145 756 1,520 1,360 1,170
Note(a) Estimated. Ice effect, Jan. 6-18, Jan. 29-Feb. 1, and Dec. 7-11, 20-31, 1909; also Jan. 1-11, 1910.										. 7-11,		

DAILY DISCHARGE, IN SECOND-FEET, OF KASKASKIA RIVER AT SHELBYVILLE, ILL., FOR 1908-1912-Continued.

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Day.	Jan,	l'eb.	Mar.	Apr.	May.	June,	July.	Aug.	Sept.	Oct.	Nov.	Dec.
1911 1 3 4	1,150 1,190 "1,100 "1,000 "1,000	2,850 2,750 2,550 1,810 1,400	918 846 774 738 702	494 494 702 1,720 3,920	1,810 1,680 1,470 1,270 1,050	203 192 160 145 160	22 45 30 34 22	12 13 16 18 18	16 12 12 13 13	8,450 8,000 7,430 6,590 6,150	918 810 596 666 738	918 882 846 810 774
6 7 8 9 10	-910 -860 -810 -760 -714	1,360 1,310 1,270 882 846	666 1,360 2,400 2,450 2,300	6,440 3,920 3,260 3,160 2,750	956 846 756 702 666	145 160 160 145 160	25 25 22 25 25 22	22 25 22 18 16	22 40 82 73 40	$5,560 \\ 4,650 \\ 4,090 \\ 3,260 \\ 2,700$	774 702 666 866 702	702 494 528 596 666
11 12 13 14 15	606 494 1,810 4,090 4 ,650	956 882 774 494 956	2,050 1,910 1,680 1,450 1,150	2,350 1,910 1,650 4,040 5,100	613 562 511 494 460	130 120 109 90 73	90 22 18 18 18	12 12 12 12 13	273 720 2,180 1,450 1,050	2,300 2,050 1,720 1,540 1,360	738 1,720 1,860 1,960 1,860	702 738 738 774 810
16 17 18 19 20	3,540 2,350 1,810 1,680 1,490	1,070 1,230 1,540 1,810 1,860	1,030 918 810 738 702	4,820 4,260 3,430 3,160 2,680	409 392 358 341 426	60 73 52 58 52	13 13 12 18 16	12 9 9 10	1,250 1,520 1,130 1,110 1,130	1,190 1,150 1,070 994 1,190	1,810 2,700 2,800 2,850 2,600	810 846 810 774 852
21 22 23 24 25	1,270 1,270 1,030 918 774	1,720 1,540 1,400 1,270 1,230	666 596 562 528 494	2,150 1,700 1,450 1,230 1,090	975 630 562 477 392	58 58 52 58 52	18 18 12 13 12	9 7 5.5 22 25	864 720 579 494 2,120	1,630 2,050 2,100 2,200 2,300	2,500 2,350 2,200 1,960 1,720	1,030 1,150 810 1,230 1,100
2627 2728 293031	845 1,070 2,800 3,160 3,050 2,950	1,230 1,150 1,030	528 596 562 562 562 494	994 864 918 918 1,340	324 240 224 273 256 256	90 82 73 45 30	13 12 13 13 13 12 13	30 40 30 22 25 16	2,980 2,750 5,070 6,920 8,750	2,450 2,100 1,580 1,400 1,190 1,030	1,540 1,400 1,230 1,110 994	1,270 1,190 1,150 1,070 994 1,110
1912 1 2 3 4 5	1,110 1,070 1,070 1,030		5,680 5,620 4,820 3,920 2,900	5,220 5,040 5,270 5,040 6,200	5,500 4,990 4,700 4,310 4,140	562 528 462 430 399	257 231 284 257 257	122 105 105 90 90	122 152 231 309 312		51 63 63 90	63 63 76 76 76
6 7 8 9 10			1,860 1,450 1,070 1,110 1,150	3,920 3,260 2,650 2,350 2,150	3,100 2,650 2,200 2,050 1,810	369 340 312 284 257	284 312 284 257 994	90 76 76 76 90	205 160 105 90 76		105 122 140 122 122	76 76 76 76 76
11 12 13 14 15		} ₀ 300	882 738 596 1,720 4,090	2,050 2,000 1,960 1,960 1,860	1,680 1,580 1,540 1,400 1,540	257 231 231 206 206	562 430 369 1,150 882	105 160 206 340 369	76 72 67 63 63		105 105 105 105 90	76 76 76 76 76
16 17 18 19 20	→ " 820		4,870 5,850 6,090 6,440 6,500	1,810 1,770 1,720 1,960 2,200	1,550 1,490 1,360 1,190 1,110	562 702 536 666 774	596 562 528 495 462	284 231 312 528 810	63 63 63 63 63 63		00 90 90 90 76	76 63 63 63 51
21 22 23 24 24 25		1,190 1,150 1,270 2,750	6,740 6,650 6,320 5,680 5,100	2,500 2,330 2,150 2,050 1,680	956 810 774 666 596	810 666 528 462 399	349 257 231 206 182	1,580 956 774 528 369	ដូច ខ្លួន ខ្លួន ខ្លួន		76 76 76 76 76	51 51 51 51 63
26 27 28 29 30 31		5.160 5,500 6,440 6,380	4,820 4,540 3,700 5,390 5,390 5,270	2,000 2,550 3,350 4,370 5,330	528 528 596 562 630 596	340 702 340 312 284	309 312 76 63 140 140	666 359 206 182 160 140	63 63 61 51 51		76 76 76 76 76	63 63 63 63 63 63 63 63

DAILY DISCHARGE, IN SECOND-FEET, OF KASKASKIA RIVER AT SHELBYVILLE, ILL., FOR 1908-1912--Concluded.

Note.—(a) Estimated. Ice effect, Jan. 3-10, 1911, and Jan. 5-Feb. 21, 1912. Gage not read in October, 1912.

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Day.	Aug.	Sept.	Oct.	Nov.	Dec.	Day.	Ang.	Sept.	Oct.	Nov.	Dec.
1914 1 2		35 35 35	5.5	12 12 12	22 22 22	1914 • 16 17 18	$1.5 \\ 1.5 \\ 1.5 \\ 1.5$	12 12 4.5	$1.5 \\ 1.5 $	12 12 12	
4 5		35 65	.2	12	22 22	20	.5	4.5 4.5	4.5	12	
6 7 8 9 10	.5	774 462 110 68 50	.2 .2 .5 1.5 4.5	12 12 12 12 12		21 22 23. 24. 25.	.5 .5 .5 .5	1.5 1.5 .5 .5	4.5 4.5 4.5 4.5 4.5	12 12 12 12 12	
11 12 13 14 15	$50 \\ 35 \\ 22 \\ 12 \\ 4.5$	35 22 22 22 22 12	4.5 1.3 1.5 1.5 1.5	12 13 12 12 12		26 27 28 29 30 31	.5 1.5 50 50 50 50	.5 .5 .5 .5 .5	$\begin{array}{c} 4.5\\ 4.5\\ 4.5\\ 12\\ 12\\ 12\\ 12\\ 12\end{array}$	22 22 22 22 22 22	

DAILY DISCHARGE, IN SECOND-FEET, OF KASKASKIA RIVER AT SHELBYVILLE, ILL., FOR THE PERIOD AUG. 10 TO DEC. 5, 1914.

Note .- Discharge determined from a rating curve which is not well defined.

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MONTHLY DISCHARGE OF KASKASKIA RIVER AT SHELBYVILLE, ILL., FOR 1908-1912. [Drainage area, 1.030 square miles.]

		Discharge in	second-feet.		Runoff
Month.	Mean daily maximum.	Mean daily minimum.	Mean.	Mean per square mile.	depth in inches on drainago area.
1903 February (25-20) April April May June June June Jugust September Detoher Vovember Detomber	$\begin{array}{r} 5.330 \\ 5.330 \\ 3.000 \\ 10.600 \\ 1,110 \\ 426 \\ 90 \\ 34 \\ 25 \\ 58 \\ 34 \end{array}$	3,760 994 018 1,190 160 53 13 10 8 8 8 15	4,710 3,010 1,920 4,730 572 187 45.5 14.7 11.3 20.4 27.8	4.57 2.92 1.86 4.59 .555 .182 .044 .014 .014 .011 .026 .027	.85 3.37 2.03 5.29 .82 .21 .05 .02 .01 .03 .03
1909 Yebruary Yebruary March April May June June Suptember Detober Sovenaber Detoenber The year	60 3,820 3,00 7,580 3,760 4,140 5,740 5,740 358 358 358 130 956 1,550 7,580	426 353 528 236 13 18 13 58	+3.4 1,650 1,090 4,230 1,940 1,740 2,190 83.9 25.4 54.3 375 593	.043 1.55 1.06 4.12 1.85 1.60 2.13 .051 .055 .053 .3576 1.13	.05 1.67 1.22 4.60 2.17 1.89 2.46 .09 .03 .06 .41 .66

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MONTHLY DISCHARGE OF KASKASKIA RIVER AT SHELBYVILLE, ILL., FOR 1908-1912--Concluded.

[Drainage area, 1,030 square miles.]

· · ·		Discharge in	second-feet.		Runoff-
Nonth.	Mean daily maximum.	Mean daily minimum.	Mean.	Mean per square milo.	inches on drainage arca.
1910					
January	6.360		1.910	1.86	2.14
Fabruary	3 380	200	682	.662	-69
March	3,700	290	1.220	1.15	1.36
Anril	426	160	251	.273	.30
May	4.870	358	1.840	1.79	2.05
June	1,450	109	491	.477	.53
July	1.470	160	551	.564	.65
August	1,150	60	359	.349	.40
September	3,180	58	638	.619	.69
October	1.070	90	343	.333	.38
November	3,600	66	439	.426	.48
December	1,520	145	507	.492	.57
The year	6,300	58	774	0.751	10.25
1911				,	
January	4,650	494	1.650	1.60	1.84
February	2.850	494	1.400	1.36	· 1.42
March	3,450	494	1.020	.990	1.14
April	6.5:0	494	2.430	2.36	2.63
May	1.510	224	657	.638	.74
June	205	. 30	102	.099	.11
July	90	12	21.2	.021	.03
August	40	5.5	16.8	.016	.02
September	8,750	12	1,450	1.41	1.57
October	8,450	994	2,950	2.86	3.30
November.	2,850	596	1,500	1.46	1.63
December	1,270	494	\$80	.854	.98
The year	8,750	5.5	1,173	1.14	15.40
1912					
January			852	.827	.95
February	6,440		1,250	1.21	1.30
March.	6,740	596	4,100	3.98	4.59
April.	6,200	1,680	2,960	2.87	3.20
May	5,500	528	1,850	1.80	2.08
June	810	206	441	.428	.48
July	1,150	63	381	.370	.43
August	1,580	76	330	.320	.37
September	399	51	105	.102	.11
November	140	- 51	88.2	.030	.10
December	7ô	51	66.9	.065	.07

MONTHLY DISCHARGE OF KASKASKIA RIVER AT SHELBYVILLE, ILL., FOR THE PERIOD AUG. 10 TO NOV. 30, 1914.

[Drainage area, 1,030 square miles.]

		Runoff				
Month.	Mean daily maximum.	Menn duily minimum,	Mean.	Mean per squaro mile.	inches on drainnge area.	
¹ 914 August (10-31) Soptember October November	50 774 12 22	.5 .5 12	15.2 61.0 3.38 13.7	.015 .059 .0033 .013	.01 .07 .004 .01	

APPENDIX C HYPOTHESIS TEST RESULTS

Appendix C1

Test for runoff distribution of Kaskaskia River basin at Shelbyville.

A. Chi-Square Test

The data are from the 1908-1914 stormy season: (January, February, March, April, May, June, and December). There are total of 35 data points.

 0.85
 3.37
 2.08
 5.29
 0.62
 0.03
 0.05
 1.67
 1.22
 4.60

 2.17
 1.89
 2.46
 0.66
 2.14
 0.69
 1.36
 0.30
 2.06
 0.53

 0.57
 1.84
 1.42
 1.14
 2.63
 0.74
 0.11
 0.98
 0.95
 1.30

 4.59
 3.20
 2.08
 0.48
 0.07

Hypothesis

II₀: The runoff distribution is a Gamma distribution with $\alpha = 1.25$, r = 2

		•	•
Category(in)	f _i Observed	Gamma Distribution	e _i Expected frequency
0 -0.5	6	0.104	3.46
0.51-1	9	0.2259	7.9
1.01-1.5	5	0.2014	7.05
1.51-2	3	0.1505	5.267
2.01-2.5	6	0.1071	3.784
2.51-3	1	0.0703	2.46
3.01-3.5	2	0.0448	1.568
3.51-4	0	0.0278	0.973
4.01-	3	0.0682	2.387

$\Gamma(\mathbf{r}, \alpha) = \frac{\alpha}{\Gamma(\mathbf{r})} (\alpha x)^{\mathbf{r}-1} e^{-\alpha x}$

Test statistic

$$\chi^2 = \sum_{i=1}^{m} \frac{(f_i - e_i)^2}{e_i}$$

= 6.659

Decision:

The test statistic χ^2 is smaller than the 0.05 significance level quantile of Chi-square distribution $\chi^2_{0.05,8} = 15.507$. Therefore, we accept H_o.

B. Kolmogorov-Smirnov Test

Category	F*(x)	<u>S(x)</u>	F*(x)-s(x)
0 ~0.5	0.104	0.1714	0.0674
0.51-1	0.3299	0.4285	0.1036
1.01-1.5	0.5313	0.5713	0.04
1.51-2	0.6818	0.6571	0.0247
2.01-2.5	0.7889	0.8285	0.0396
2.51-3	0.8592	0.8571	0.00215
3.01-3.5	0.904	0.914	0.01
3.51-4	0.9318	0.914	0.0178
4.01-	1	1	0

Hypothesis:

 $H_0: F(x) = F*(x)$ for all x

 H_1 : $F(x) \neq F^*(x)$ for at least one value of x

Test Statistic:

 $T = \sup_{x} |F^*(x) - S(x)|$

Decision Rule:

Reject H_o at the level of significance α if the test statistic T exceeds the 1 - α quantile w_{1- α} on table of the Kolmogorov Test Statistic. For this test, T = 0.1036 for α = 0.05 and the critical value is 0.43. Since T < 0.43, we accept H_o.

Appendix C2

Test of the Simulation Inflow Data and the Historical Data A. <u>K-Sample Slippage Test [19]</u>

Hypothesis:

- H_o: All simulation runs and the historical data have identical population distribution functions.
- H1: Some populations tend to furnish larger observed values than other populations.

Test Statistic:

The test statistic is evaluated by comparing the "largest" sample with the "smallest" sample, as follows. In each of the K-samples, find the largest of the n observations and denote the extreme (largest) value by Z_i , i = 1, 2, ... K. Compare the Z_i 's, the extremes, and find the largest Z_i , denote by $Z^{(K)}$, the largest of all the observed values, is called the sample of rank K, or the largest sample. The sample from which $Z^{(1)}$ comes is called the sample of rank 1, or the smallest sample. The test statistic T equals the number of observations from the sample of rank K, which is greater than $Z^{(1)}$, the largest observation in the sample of rank 1.

Decision Rule:

Reject H_o at the level of significance α if T exceeds the 1 - α quantile w_{1- α}, as given by the Table C2.1 [19]. In this test, K = 55, n = 12, the following steps were followed in analyzing the data.

- An asterisk was placed at the upper right corner of the extreme (greatest) value in each sample.
- (2) The smallest extreme was underlined once and the greatest extreme was underlined twice. Shown on Table C2.3, C2.4, C2.5.

The smallest values of m such that $P(M(1,k) \ge m) \le \alpha$, where the first, second, and third values given for each k, number of samples, and n, sample size, correspond to $\alpha = .05$, $\alpha = .01$, and $\alpha = .001$ in that order. Absence of a value in the table means there is no m such that $m \le n$ and $P(M(1, k) \ge m) \le \alpha$, for that particular choice of k, n and α .

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40	6, 8, 11	6, 9, 11	7, 9, 12	7, 9, 12	8, In, 12	8, 10, 12	8, 10, 13	S, 10, 13	8, 10, 13	8, 10, 13	8, 11, 13	9, 11, 13	9, 11, 13	9, 13, 14	9, 11, 14	9, 11, 14	9, 11, 14	9, 11, 14	9, 11, 14
35	6, 8, 10	0, 8, 11	7, 9, 12	7, 9, 12	7, 10, 12	8, 10, 12	S, 10, 13	3, 10, 13	8, 10, 13	8, 10, 13	8, 10, 13	8, 11, 13	9, 11, 13	9, 11, 13	9, 11, 13	9, 11, 13	9, 11, 14	9,11,14	9, 11, 14
30	6, 8, 10	6, 8, 11	7, 9, 11	7, 9, 12	7, 9, 12	S, 10, 12	8, 10, 12	8, 10, 12	8, 10, 13	8, 10, 13	8, 10, 13	8, 10, 13	8, 10, 13	8, 11, 13	9, 11, 13	9, 11, 13	9, 11, 13	9, 11, 13	9, 11, 13
25	6, 7, 10	6, 8, 11	7, 9, 11	7, 9, 11	7, 9, 12	7, 9, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 13	8, 10, 13	8, 10, 13	8, 10, 13	9, 10, 13	9, 11, 13	9, 11, 13
20	5, 7, 10	6, 8, 10	7, 8, 11	7, 9, 11	7, 9, 11	7, 0, 11	7, 9, 11	8, 9, 12	8, 10, 12	8, 10, 12	8, 10, 12	S, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12
18	5, 7, 10	6, 8, 10	7, 8, 11	7, 9, 11	7, 9, 11	7, 9, 11	7, 9, 11	7, 9, 11	8, 9, 11	8, 9, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12
16	5, 7, 9	6, 8, 10	6, 8, 10	7, 8, 10	7, 0, 11	7, 9, 11	7, 9, 11	7, 0, 11	7, 9, 11	8, 9, 11	S, 9, 11	8, 0, 11	8, 0, 11	8, 9, 11	8, 10, 11	8, 10, 12	8, 10, 12	8, 10, 12	8, 10, 12
14	5, 7, 9	6, 8, 10	0, 8, 10	7, 8, 10	7, 8, 10	7, 9, 10	7, 9, 11	7, 9, 11	7, 9, 11	7, 9, 11	7, 9, 11	8, 9, 11	8, 9, 11	8, 9, 11	8, 9, 11	8, 9, 11	8, 9, 11	8, 0, 11	8, 9, 11
12	5, 7, 9	6, 7, 9	G, 3, 9	6, 8, 10	7, 8, 10	7, 8, 10	7, 8, 10	7, 8, 10	7, 0, 10	7, 9, 10	7, 3, 10	7, 9, 10	7, 0, 10	7, 9, 10	8, 9, 10	8, 9, 11	8, 9, 11	S, 9, 11	8, 9, 11
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6	5, 0, 8	6, 7, 8	6, 7, 9	6, 7, 9	6, 8, 9	6, 8, 9	7, 8, 9	7, 8, 9	7, 8, 9	7, 8, 9	7, 8, 9	7, 8, 9	7, 8, 9	7, 8, 0	7, 8, 9	7, 8, 9	7, S, 9	7, 8, 9	7, 8, 9
8	5, 6, 8	5.7.8	6, 7, 8	0, 7, 3	6, 7, 8	6, 7, 8	6, 7	6, 7	G, 3	7, 8	7,8	7, 8	7,8	7,8	7, 8	7, 8	7,8	7,8	7,8
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Quantile of The K-Sample Slippage Test Statistic Table C2.1

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15	222 222 222 222 222 222 222 222 222 22		34	763 411 785 785 785 785 785 785 785 785 785 785
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Simulation Runs

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cortce	2	250	320	1972	14426	32.94	1/20	723	024	433	351	* 433	2096
Hist	, FT	40	27	21	32	80	62	98	204 4	885	2128	3574:	3560
	50	629	50 00	1761	1099	123	634	+ 972	3000	3115	2270	3202*	2302
	·6t1	253	843	1146	657	1199	1477	2600\$	1172	2007	47474723	2436	1079
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Hypothetical results of Slippage Test(Jan.-Feb.) Table C2.3

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Table C2.4 Hypothetical results of Slippage Test (Mar.-Apr.)

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Simulation Runs

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Table C2.5 Hypothetical results of Slippage Test(May-June)

- (3) The number of values in the sample with the greatest extreme that exceeded the smallest extreme was determined to be statistic T.
- (4) From Table C2.1 find out the critical region of size $\alpha = 0.05$ was determined. The results of these tests is presented below on the Table C2.2.

	January-February	March-April	May-June
z ^(K)	5870	6104	8978
z ⁽¹⁾	2342	29 33	1456
W ₁ -α	8	8	8
Т	· 1	6	6
Decision	Accept H	Accept H _o	Accept H

Table C2.2 Result of K-Sample Slippage Test

B. The Friedman Test [20]

Hypothesis:

H_o: The data from each simulation run and the historical data are equally likely.

H1: At least one data block tends to yield larger observed value value than at least one other.

Test Statistic:

The Friedman test statistic is defined as [20]

$$T = \frac{12}{bk(k+1)} \sum_{j=1}^{k} R_{j}^{2} - 3b(k+1).$$

The rank assigned to each run of simulation and historical data of each season period is presented on Tables C2.7, C2.8, and C2.9. Decision Rule:

Reject the null hypothesis at the level α if the Friedman test statistic T exceeds the 1 - α quantile of a Chi-Square random variable with k - 1 degree of freedom. In this test, b = 12, k = 55, and α = 0.05. The results of the test are presented on Table C2.6.

· .	January-February	March April	May-June
T	66.73	72.03	73.101
W _{1-α}	72.04	73.2	73.2
Decision	Accept H	Accept H _o	Accept H
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Table C2.6 Results of Friedman Test

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Simulation Runs

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APPENDIX D

SIMULATION RESULTS

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t LL	4 8114444444 801000000000000000000000000	0 111	010 11 11 12 12 12 12 12 12 12 12 12 12 12	I I I
NI	20000000000000000000000000000000000000	0	11 10 10 10 10 10 10 10 10 10 10 10 10 1	
STAN	22220000000000000000000000000000000000	0111	STAN 25740 1500 1500 1500 1500 1500 1500 200 200 200 200 200 200 200 200 200	1292 29% 438%
500	н 28230000640 4 140400 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	0111	9 15 15 15 15 15 15 15 15 15 15	1292 29% 428%
NT	22220000000000000000000000000000000000	0	IN 1022008 1022008 1022095 1125005 1125005 110005 1125005 110005 110005 1100005 1100005 1100005 1100000000	
STAN	224500000228 224500000228 22450000028 224500000028 224500000000000000000000000000000000000	6•49 0 0	STAN STAN 1500 255888 25588 255888 25588 2	000
0FT 0FT	3488 3488 8978 8978 8978 11550 2010 2010 2010 2010 2010 2010 201	6•49 0 0	8 00 1160 00000000	000
IN	4642 66202 69302 702000 702000 700000000	6.49	TN 2102 2358 25102 28388 26538 26538 26538 26538 2756 2533 2533 2533 2533 2533 2533 2533 25	
STAN	20000000000000000000000000000000000000	4. 65 0	STAN STAN 115000 2000 2000 2000 2000 2000 2000 20	000
1 0PT	2000 2000 2000 2000 2000 2000 2000 200	•••65 0 0	0071 1150000 11500000000	000
NI	25000 2500000000	4.65	III 11122 1122 1112	ο.
2 2 4 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	cost AQ _F RR RC	00 00 00 00 00 00 00 00 00 00 00 00 00	RR RG RC

Results from Simulation of Historical data with Time independent stage input assumption Appendix D1

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STAN	1500 1500 1500	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	111000	1500	1500	1500	0.96 1112	1883	STAN	1500	2000 2000 1000 1000	820	820	200	0000	1500	3645	3.58	2225	8622 8602
18 0PT	1600 1445 2150	2100	1010	1650 1650	936 1502	1432 1137	1.4	8 5 5 7 7 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	24 0PT	1600	2020 1020 1	820	820	200	0000	1590	3500	3.49	2320	575 275 275 275 275
IN	1078 810 2342	2174	100 100 100 100	1133	1502 1502	1432 1137	1.93		NI	1070	000 000 000 000 000 000 000 000 000 00	820	820	200	000000000000000000000000000000000000000	1590	5870	5.23		
STAN	1244 72 72	20	1500	0045 740 7	344 244	238 860	0 0	337% 100%	STAN	1500	11 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1500	1500	1700	226	1082	1130	0	1070	
17 0PT	1244	20 20	2020	265 777		238 860	0.81	140 140 140 140 140 140 140 140 140 140	23 0PT	1600	1628	1900	1650	597	2220	1082	1130	1.44	5 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	ル マ マ マ マ マ マ マ マ マ マ マ マ マ マ マ マ マ マ マ
NI	520 228	22	2263	1 01. 5 0. 0	244	238 860	1.81		NI	746	1628	2146	1255	597	220	1082	1130	2.24		
STAN	1738 1480 1486	291 203	000	861	4 7 7 7 7 7 7	195 126	0 5 50	2008 1008	STAN	1171	11301	955	2985	4052	1764	1675	1625	6.01	o c	> o
16 0PT	1238 2050 936	402 707	260	10- 10- 10- 10- 10- 10- 10- 10- 10- 10-	4 7 7 7 7 7 7 7	126	0.81	00	0PT 0PT	1171	11 11 11 10 10 10 10 10 10 10 10 10 10 1	955	2985	4052	1764	1675	1625	6.01	00	00
II	2050 936	201 201	260	117 A 107 A 107 A 107 A	500 100 100 100	195 126	0.81		NI	7 F V F	11000	955	5300	40.52	1764	2056	1625	6.01		
NATS	1500 709 312	1119 493	10 20 4	1200	219		01	1 1	STAN	1188 24	507	- V) 	сл ус н с	12		2-7-	33	0	1	1 1
15 0PT	1500 709 312	1119 493	4 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	173	719	0,1 1,0 1,0 1,0 1,0 1,0 1,0 1,0 1,0 1,0	01	1 8	21 0 PT	1188 24	101 101	- V) 	יי רו ר	え	200	24	33	0	J	1 1
NI	732 420 312	1119 493	200 200 24	222	719	1 28 28 28 28	0		IN	31	0 6 7 7	- 10	ný H H	រដ		27	33	0	1	
STAN	1500 1500 1500	1500	2732	1500	263 263	383 383	2.5	388%	STAN	1500	923	10,00	262 173	1-2-1-	2 C 7 O 7 O 7 O	22	10	0	1	* 1
14 0PT	1600 1435 1950	1103 2100	2044 1500	2007	7200 7207	119 383	1-44 14-44	402% 702%	0 PTI	1×00 1×00	10 m	19 10 10 10 10	262 173	1.2	104	200	49	, 0	1	1 1
IN	593 1285 1950	1103 2945	3514 1164	119	500	119 383 383	3.45		NI	1456	528 358	100	262	11	104 107	200	10	0		
STAN	2228 1860 1500	1204	222	200	351	358	1,79	3423 6143	STAN	1500	1500	840	1466	1591	3018	1534	1500	3.61	Y ST	200
13 0PT	2228 1860 1500	1500 1205	335 270	202	351	358 358	1.79	3423 6148	L T T T T T T T T T	1600 2000	1836 840	2.02	1466	1900	2709	1534	1500	3,61	0 1 v V V	242
NI	3386 1860 916	562 412 412	335 270	500	220	358	3.31		NI	796 1835	1648 840	564	1466	3906	2010 2010 2010	1534	1007	13.84	0(2 0	C,
Stag	-1 N M-	4 v	20	- യ ເ	10	121	Cos ∆2n	RR RC	Star	, ਜ ਕ	t-17	5	00	-œ	500	24	12	Cost	2 n 0 n	вс

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	STAN	1223 1223 200 201 202 201 202 202 202 202 202 202	00000 00000000000000000000000000000000	0111
29	140	1223 1223 1222 766 776 776 776	00000 00000	0111
	MI	221002 22202 22202 2220 2210 220 220 220	8889 8889 8889 8889 8889 8889 8889 888	o .
	STAN	1406 1406 563 563 563 563 563 563 563 563 563 56	൜൜ൕ	0111
28	CPT	1 4 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7	ᡣᡣᡃᡢᢍ	0111
	NI	44 47 47 47 47 47 47 47 47 47 47 47 47 4	ᡊᡢᡃᡢᢁ	0
	STAN	11 00 00 00 00 00 00 00 00 00 00 00 00 0	201 201 201 201 200 200 200 200 200 200	0111
27	J.HO	144 80 80 80 80 80 80 80 80 80 80 80 80 80	8473 8433 8433 8433 8433	0111
	ЦМ	11180.099 8080999 8080999 8080999 8080999 8080999 8080999 8080999 8080 80809 8080 800 80	8410 8400 8400 8400 8400	0
	STAN	02200 0200 02000000	22% 273 3963 3963	3585 3585 3585 3585 3585 3585 3585 3585
20	1 D D L	00000 00000000000000000000000000000000	5226 573 3963	0.11 11.01 0.02 0.02 0.02 0.02 0.02 0.02
	MI	11000000000000000000000000000000000000	220 573 3960	4.55
	STAN	25,450 25,450 25,450 26,550 26,500 20,500	1966 1892 2146 3526	5.35
50 50	OFT	2000 2000 2000 2000 2000 2000 2000 200	1966 1892 2146 3526	5.35
	III .	20000000000000000000000000000000000000	1966 1892 2146 3526	15.35
	Stag		5010	Cos RC AQ

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112866 2535555 2535555 2535555 2535555 2535555 2535555 2535555 2535555 2535555 2535555 2535555 2535555 25355555 2535555 25355555 25355555 25355555 25355555 25355555 25355555 25355555 25355555555	5.35
80000000000000000000000000000000000000	3.61 236 68 83 88
111322 111323 11132 111323 11133 111133 111133 111133 111133 111133 111133 111133 111133	3.03 777 199% 323%
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<i>в</i> <i>в</i> <i>в</i> <i>в</i> <i>в</i> <i>в</i> <i>в</i> <i>в</i>	000
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Results from Simulation using Historical data(Jan.-Feb.) with Assumption of time Dependent Inflow Rate. Appendix D2.1

														-				
	STAN		1500	1500	867	820	820	820	300	000	300	000	1500	3645	3.59	2225	320%	509%
4	OPT		2227	820	820	820	820	820	300	300	300	000	1590	3555	3.49	2315	304%	529%
	NIJ		1070	820	820	820	820	820	000		000	000 00	1590	5870	5.23			
	STAN		1500	1500	1500	1500	1500	1200	いのいて	1500	1500	1500	1500	1500	0.93	453	193%	537%
e	E C		2235	300	2200	2300	1068	2200	2150	せいのも	,812	1508	1432	113?	1.87	142	6%	168%
	NH		1078	810	2342	2174	1052	2319	2222	1133	612	1502	1432	1137	1.93			
	STAN		1437	330	1500	2583	3294	1500	1500	1500	1014	351	433	1500	3.21	1132	2553	387%
2	ОРT		1437	330	1972	2200	3205	2100	2000	982	433	351	433	1500	3.12	1221	27.6%	41.7%
	NH		230	330	1972	4426	3294	1420	723	624	433	351	433	2096	4.38		-	
	STAN	Ī	1191	27	21	32	80	62	98	406	885	1500	1887	3560	3.49	14	0.4%	0.7%
7 -1	OPT		1191	27	21	32	80	62	98	406	885	1800	1650	3497	3.43	77	२ 1 १	3.7%
	NI	6	З т	27	21	32	080	62	0 0 0	406	8 85	2128	3574	3560	13.51	_		
		Stag	. 1	2	ŝ	4	Ś	9	~	ω	σ	10	11 11	12	COSI	Ğ ⊲	E.H.	RC

Results from Simulation Using Historical data(Mar.-Apr.) with Assumption of time Dependent Inflow Rate. Appendix D2.2

	STAH	3570	2362	1566	1500	1500	1500	696	312	226	660	573	396	3.51	1158	2403	358%
S	01.1	3510	2362	1,000	1700	1800	1/1:00	476	312	226	660	573	396	3.51	1158	245%	3583
	NI	4728	2362	1566	1346	760	573	476	312	226	. 660	573	396	4.55			
	STAN	1500	1500	983	5 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	607	262	173	154	104	60	50	10	0	0	0	0
t	0FT	1500	1445	953	358	607	202	173	154	101	30	20	† 9	0	0	0	0
	21	1456	732	538	358	607	202	173	154	104	60	56	70	0			
	STAD	1500	1500	1500	1,500	1500	2732	1500	1500	1500	273	119	383	2.51	752	222%	38,8%
r,	J.dO	1750	1235	1950	1103	1750	2392	1500	1500	1500	273	119	5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5	2.01	1120	3193	44.5%
	FI	593	1205	1950	1103	2945	3515	1164	0111	305 205	265	119	383	3.45			
	STAN	1500	1824	2308	6052	1500	1558	3580	2272	1930	1500	1500	1500	5.33	0	0	0
N	OP.I	2350	2300	2150	5334	1500	15,58	3580	2272	1980	1500	1500	1500	16.4	668	1112	145
	Z	2102	2380	2808	6052	638	2420	3580	2272	1980	1365	576	639	5.33			
	STAN	484E	8978	6316	11302	3028	1720	1500	1500	1500	715	324	211	64,9	0	0	0
		+	ώ	0	??	28	20	00	00	00	15	324	11	64	_		
-	니 가 이	348	897	691	430	0 6 6	17	1 5	ч Н	ч) Н	$(\sim$		CV.	9	0	0	0
-	IN OFF	4642 348	8978 397	6316 631	4302 430	3028 303	1720 17	992 15	807 15	710 15	392 7	1350	211 2	9 617 9	0	0	0

Results from Simulation using Historical data(May-June) with Assumption of time Dependent Inflow Rate. Appendix D2.3

STAN	11 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	3.01 5.71 2.6%	8 111111111 00201010000000 0120000000000	000 99 99 90 90
DEPN	2000 2000 2000 2000 2000 2000 2000 200	20.030 2000 2000 2000 2000 2000 2000 200	952620499 95262049 95276049 95276040 952760400000000000000000000000000000000000	2 t 10 5 t 20 5 t 20 8 39 0 0 0 8 39 0 0 8 30 0 8 30 0 8 300 0 8 300 0 8 300
JUDI	11600 117000 117000 117000 1100000000	3.01 56.2% 26.2%	Н Н Н Н Н Н Н Н Н Н Н Н Н Н Н Н Н Н Н	000 000 000 000
NI	11 12 12 12 12 12 12 12 12 12 12 12 12 1		22244413 2224413 222441413 2224414 2224413 2224413 2224413 2224413 2224413 2224413 2224413 2224414 2224414 2224414 2224414 2224414 2224414 2224414 2224414 2224414 2224414 22244413 2224414 2224414 22244413 222444145 222444145 222444145 222444145 222444413 222444413 222444413 222444413 222444413 222444413 222444413 222444413 222444413 222444413 222444413 222444413 222444414444414444444444	
STAN	24000000000000000000000000000000000000	2 89 1168 4168 4158 4158	8 44444444444444 8 4444444444444 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	4 2007 4 2007
DEPN	1888 3674 2600 2000 2000 2000 2000 2000 2000 200	1.68 1959 477.11 73.73	90000000000000000000000000000000000000	1.44 1696 44.7 73.9%
TNDP	1600 3650 3650 3650 3650 3650 3650 2150 2150 2150 2150 2150 2150 2150 21	1.58 2001 48.1% 75.2%	11111111111111111111111111111111111111	1.37 45.44 75.44 75%
NI	212010420323333334		02542424 11111111111111111111111111111111	
STAN	11500 11500 125000 12500 125000 10000000000	3 th 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	8 4 4 4 4 4 4 4 4 4 4 4 4 4	3, 38 8, 18 8, 56 8, 56
DEPN	111955 1111955 111119900 111198000 111880000 111880000 111880000 111880000 1118800000000	2.98 589 16.19 27.23	80000000000000000000000000000000000000	2.94 2.04 2.83 2.83 2.83 2.94 2.83 2.94 2.94 2.94 2.94 2.94 2.94 2.94 2.94
	11460 11460 12462 12462 12462 12600 22008 2008 20000 2008 2008 2008 2008 2008 20000 20000 2008 200000	84 84 84 84 84 84 84 84 84	00011000000000000000000000000000000000	3°3 260 12°1%
NI	220152 220052 220052 220052 220052 220052 220052 220052 220052 220052 220052 220052 200520	27.27	1101111000040488 1104004004888 11040040040888	
STAN	11 11 11 11 11 11 11 11 11 11 11 11 11	0 893 1000 1000	8444444444444 67777777777777777777777777	2%% 2% 00 00 00 00
DEPN	12774 12774 12774 1228 1228 1228 1289 12800 1000 128000 12800 128000 128000 128000 128000 128000 128000 128000 128000 128000 128000 128000 128000 128000 128000 1280000000000	1•68 193 83 21•6	22500 2250 2250 2250 2250 2250 2250 225	5000 5000 5000 5000
AGNI	14600 14600 14600 14600 14600 14600 14600 14600 14600 14600 14600	1.2 393 16.4%	1111111111111111111111 000000000000000	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
IN	12283 8829 8829 12283 12395 12393 12		117612 11	
STAN	244 24 2009 2009 2009 2009 2009 2009 200	1,87 426 34,75	8144444444448 844444444444 82660000000000 8244490000000000000000000000000000000000	2.84 339 10.3% 188%
DEPN	144 1730000000000	1.44 627 23% 51% 51%	5000 500 500 500 500 500 500 500 500 50	2.08 860 26% 478%
TUDF	14 14 14 14 14 14 14 14 14 14	1. 6.44 5.082 5.00	11111111111111111111111111111111111111	2,52 556 16.8% 30.9%
NI	2558 2579 2579 2579 2579 2579 2579 2579 2579	13 0.	800 100 100 100 100 100 100 100	нQ
	2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	Cost AQ RR ¹ C	8 1111 1111 1111 1111 1111 1111 1111 1	COS AQ RR RC

Table D3.1 Simulation Result of Jan-Feb Flow

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STAN	1387 523 523 14411 823 1070 15000 15000 15000	0 892 100%	22000000000000000000000000000000000000	2.87 0 0
DEPN	1387 524 524 823 823 823 823 823 823 823 823 823 823	0•72 592 247% 664%	DEPN DEPN 17722 1979 1979 1979 694	1.68 778 2613 2613
17 TNDP	1387 524 524 524 1412 823 823 338 338 1500 1500 1500	0.6 642 72% 72%	76000000000000000000000000000000000000	1 • 23 978 562% 662%
NI	2200 114412 2337 2337 2337 2337 2337 2337 2337 23		11111111111111111111111111111111111111	•
STAN	1111 1000 100 1000 1	1.34 748 266% 572%	87AN 16700 16700 16700 16700 16700 26500 27037 27037 27037 27037 27037 27037 27037 27037 27037 27037 27037 27037 27037 27037 27047 2007 200	0, 40 0,000 0,000 0,000
t DEFN	11902 201902 201902 2019 2019	1•66 616 219% 471%	9 DEPN 1684 1684 1071 1071 1071 1071 1071 1071 1071 107	0 1 1 0 0 0 0 0 0 0 0 0 0 0
INDP	21600 221600 22000 24900 2649 2649 2649 2649 2649 2649 2649 2649	1.44 25208 5428 5428	1 NDP1 1 0712 1	0 1 0
IN	2309 2309 2309 2309 2309 2309 2309 2309		2000011 40000 2010040 2010000000000	
STAN	1220 1220 12200 120000 1200000000	1•89 334 126% 292%	8 4444444444 4446690000139 78646050000139 78646050000139 78646050000139 78646050000139 78660050000139 78660050000139 786600500000130 7866005000000000000000000000000000000000	2.962
DEPN	1222 1222 1222 1222 1222 1222 1222 122	1.03 711 269% 623%	8 DEPN 21440 212844351 2128444351 2128444351 2128444351 21284450 210995 210995 21095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 20095 2005 200	1 0000
13 TUDP	1222 1222 1222 1222 1226 1226 1226 1226	1•19 647 845% 567%	1 NDP 1 1600 1 135100 1 135100 1 150000 1 1500000 1 15000000 1 150000000000	1100 1100 1200 140 100 140 100 140 100 100 100 100 1
NI	264236 26424 26424 26424 26424 26424 26423 26433 264533 264535 264535 264555 2645555555555555555555555		H H H H H H H H H H H H H H H H H H H	
STAN	2005 2005 2005 2005 2005 2005 2005 2005	2.94 381 1118 199%	874N 15000 15000 15000 15000 150500 100000000	3.07
DEPN	1043 1043 1043 10443 1014 1014 1014 1014	1086 31886 5678%	7 DEPN 2782 2782 1521 1521 1520 15200 15200 15200 150000 150000 150000 1500000000	2.58 377 119%
INDP	1112 1112 1112 1112 1112 1112 1112 111	2.94 381 11.1% 129%	H H H H H H H H H H H H H H H H H H H	3.07
NI	649 649 11042 110148 100148 100148 100148 100148 100148 100148 100148 100148 100000 100000 1000000		902333 2122233 2122233 2122233 2122333 2122333 23333 233333 23333 23333 23333 23333 23333 23333 233333	
STAN	11111111111111111111111111111111111111	284 284 284 284 284 284 284 284 284 284	874N 875000 8750000 8750000 8750000 8750000000000	50000 5000 50000 5000 5000000
DEPN	2324 2324 11505 11825 11827 22205 22285 22555 22555 22555 22555 22555 22555 25	2.050 2466 5338 5338	225300 255555555555555555555555555555555	2.00 1581 27581 27581
11 IMDP	1500 1500 1500 1500 1500 1500 1500 1500	2,88 1356 119%	1000 1100 1100 1100 1100 1100 1100 110	500 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
NI	5383 2382 5383 5383 5383 5383 5383 5383	чы ы. 0.	2039555 203555555555555555555555555555555555	ыр
6 4 1	240,04000000000000000000000000000000000	COS1 AQ1 RR	8 4 4 4 4 4 4 4 4 4 4 4 4 4	RC COS

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STAN	22000022 22000022 22000022 22000022 22000022 2200002 2200002 2200002 2200002 2200002 2200002 2200002 2200002 2200002 2200002 200002 20000 2000000	2.71 0 0	84444888844448 84222899289 88960399428 89008882895600N 490088888895455000N	2.81 0 0 0
S DEPN	212 212 212 212 212 212 212 212 212 212	1 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	DEFN DEFN DEFN 222960 222900 22450 2000 22450 2000 200	01 10 04 04 04 04 04 04 04 04 04 04 04 04 04
22 INDP	22000000000000000000000000000000000000	4 40 •700 •700 •74 •700	22122222222222222222222222222222222222	2•81 0 0 0
NI	12881103299 122881103399 123861103399 1238611086 113339 12388 113339 12388 113339 12388 113339 12388 113339 12388 113339 12388 113339 12388 113339 12388 113339 123888 12388 12388 12388 12388 1		2111122299915 2299915 2229915 2229915 222995 222995 2222 2225 2225	•
STAN	44444 9700000000000000000000000000000000	8844 2001 2002 2001 200 200 200 200 200 200	8141114444888 48220000000000000000000000000000	2.83 274 7.7%
DEPN	15000 150000 15000 150000 150000 1500000000	1. 664 5250 5251 5251 5251 5251 5251 5251 5251	90000000000000000000000000000000000000	2.18 692 716% 707%
AGNI 720	11110000000000000000000000000000000000	1 • • • • • • • • • • • • • • • • • • •	22221222222222222222222222222222222222	2,00 2,00 2,00 2,00 2,00 2,00 2,00 2,00
MI	22000 224000 224000 224000 224000 22400000000		22222222222222222222222222222222222222	
STAN	111 1111110000000000000000000000000000	3.71 0 0	2211111 1 11112 2220000 2000000 20000000000	1000 100 100 100 100 100 100 100 100 10
23 INDP DEPN	24 45 45 45 45 45 45 45 45 45 45 45 45 45	3.44 6.23 1128 1288	21111111111111111111111111111111111111	2.00 2.00 2.00 2.00 2.00 2.00 2.00 2.00
	11 11 11 11 12 12 12 12 12 12	3.44 2.50 1.28 8.50 1.28 8.50	20000000000000000000000000000000000000	11 21 20 20 20 20 20 20 20 20 20 20 20 20 20
NI	1 1 1 1 1 1 1 1 1 1 1 1 1 1		234444 23449 24494 2044 2044 2044 2044 2044 2044	
STAN	41111111111111111111111111111111111111	2 4 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7	22222222222222222222222222222222222222	р 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
DEPN	11422 11942 11995 11995 11395 11395 11395 116800 116800	1100 •1100 21100 21100 2010 2010 2010 20	222659903300 22265990330 22265990330 22265990330 22265990330 22265990330 22265990330 22265990330 22265990330 22265990330 22265990330 22265990330 22265990330 22265900330 22265900330 22265900330 22265900330 22265900330 22265900330 22265900330 22265900330 22265900330 2226590030 2226590030 22265900030 22265900030 22265900030 22265900030 22265900030 22265900030 22265900030 22265900030 22265900030 22265900030 22265900030 2226590000000000000000000000000000000000	2.52 806 394% 394%
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L DEPN	28250 28250 29250 29350 20350 200500 200500 200500 200500 200000000	2.68 11968 2924 2928 2928 2028	2000 2000 2000 2000 2000 2000 2000 200	1- 01 2000 1- 10 1- 10 1
21 TUDP	11600 11550 11550 11552 11758	2•65 1208 299% 475%	00000000000000000000000000000000000000	5000 5000 5000 5000 5000 5000 5000 500
NI	28331 28332 283400 2834000000000000000000000000000000000000	FL 0.	11111111 1111111 111111111 11111111111	-
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	444 444 444	COSJ PQI RR RC	8 4 4 4 4 4 7 7 7 7 7 7 7 7 7 7 7 7 7 7	No. No. No. No. No. No. No. No. No. No.
STAN	219000000000000000000000000000000000000	3. 000 828	8 111111128001 8 12000000000 8 0000000000 8 00000000000 9 0000000000 9 000000000 9 000000000 9 000000000 9 000000000 9 000000000 9 0000000000	
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S NEEN	20000000000000000000000000000000000000	5 1 2 5 1 5 1 5 1 5 6 0 7 5 8 7 6 7 6 7 6 7 6 7 6 7 7 8 7 6 7 7 8 7 7 8 7 7 8 7 8	55550000000000000000000000000000000000	
I AGNI	20020000000000000000000000000000000000	3.58	3. 5000 1.500	
NI	222249 2222249 222249 222249 222249 222249 222249 222249 222249 222249 2		2883389999222 2883389992922 288338999222 288338999222 28833899922 2883389 2883389 2883389 2883389 288338 288338 288338 288338 2883 2883	
STAN	15000 110048 110048 110048 110048 110048 110048 110048 10000000000	2.01 8.01 7.22 8.01 7.22 8.01 7.22 8.01 7.22 8.01 8.01 8.01 8.01 8.01 8.01 8.01 8.01	SHAN 112 112 112 12 12 12 12 12 12 12 12 12	
t DEPN 2	190500 111111111111111111111111111111111	1111 2000 2000 2000 2000	COSSO 222 COSSO 222	
I AGNI	16000000000000000000000000000000000000	2 5 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	363780 111 1110 111 1110 111 1110 1110 1110	
INI	2244420342444 2244420342404203126148 2244420348400420328 2244420348400420		IN 88407 125555 1551975 1551975 1551975 1551975 1551975 1551975 1551975 1551975 1551975 1551975 1555555 155555 155555 155555 155555 155555 155555 155555 155555 15555555 15555555 15555555 155555555	
TAN I	22220000000000000000000000000000000000	3.1 2.4 2.4 2.4 2.4 2.4 2.4 2.4 2.4 2.4 2.4	8740000 200000 200000 200000 200000 200000 200000 200000 200000 200000 200000 200000 200000 200000 200000 20000	
SEPN S	460 800 1110 800 110 800 110 800 110 10 10 10 10 10 10 10 10 10 10 10	34 24 24 24 24 24 24 24 24 24 24 24 24 24	25.22 25	
CE 33	457300400 20145000 2014500 20000000000000000000000000000000000	31 21 21 21 21 21 21 21 21 21 21 21 21 21	11111111111111111111111111111111111111	
C NI	11 10 10 10 10 10 10 10 10 10		1111 1111 122293023330503 623805054 82505054 82505054 82505054 82505054 82505054 82505054 82505054 82505054 82505054 82505054 82505054 82505054 82505054 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 82505055 825050505 825050505 825050505 825050505 82505050505 825050505050505 8250505050505050505050505050505050505050	
NATS	11111 11111 111111 111111	3.14 000.	88,800 89,999 80,800	
EPN S	22286 13008 13008 13008 13008 13000 15004 15000 15004 15000	1 • 8 9 9 2 9 3 2 9 3 3 9 5 2 9 3 2 9 3 2 9 3 3 9 5 2 9 3 3 9 5 2 9 5 3 9 5 2 9 5 3 9 5 2 9 5 3 9 5 2	222 24 25 26 26 26 20 20 20 20 20 20 20 20 20 20 20 20 20	
32 I ADD I	1111221 126000000 129900000000000000000000000000000	2.80 1.1200 1.12000 1.12000 1.12000 1.120000000000	11111111111111111111111111111111111111	
C NI	11111 12209 12209 12209 12209 12209 12206 12209 12206 12209 12206 12209 12206 12008 10008 100008 10008 10008 10008 10008 10008 10000		21110220350 23111022438 2311102557 2311102557 2311102557 2311102557 2311102557 2311102557 2311102557 2311102557 2311102557 2311102557 2311102557 231110257 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 23111057 2310057 2311057 2311057 231100000000000000000000000000000000000	
STAN	22200033442424000 222000334424000 2220003332424000 222000333242000 22200003324240000 2220000342400000 222000003420000000000	2.17 214 7.98 1758	74 3 324144 44 850 8 850 8 8	
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STAN	22455000 22455000 224550000 224550000 2245500000 22455000000 22455000000 224550000000000	3.15 0 0	0020202020000 0022020202000 00220202020
DEPN	2826 2826 2826 28265 2855 285	88800 2000 5000 5000 5000 5000 5000 5000	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
INDP	22421600 22421600 22421600 22421600 224216000 224216000 224216000 224216000 22421600000 2242160000000000000000000000000000000000	3.15 0 0	0000 000 000 000 000 000 000 00
IN	2244595 2244595 2244595 2244595 2244595 2244595 2244595 22445 22445 2245 22		2322410024 2322410024 2322410024 2322410224 2324 2327 2327
STAN	2000 2000 2000 2000 2000 2000 2000 200	0000 0000 0000 0000 0000 0000	29 29 29 29 29 29 29 29 29 29 29 29 29 29 2
r DEPN	3688000400 20000400 20000400 20000400 2000040 2000040 2000040 2000040 2000040 2000040 2000040 2000040 2000040 2000040 2000400 20004000 20004000 2000400000000	10000 1000 1000 1000 1000 1000 1000 10	364 26 26 26 26 26 26 26 26 26 26 26 26 26
acini 74	26000 2600 2600 2600 2600 2600 2600 260	2•79 2559 3478	72000 720000 720000 720000 72000 720000 720000 7200000000 720000
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STAN	00000000000000000000000000000000000000	0 18 18 18 18 18 18 18 18 18 18 18 18 18	STAN 574 1574 15649 15600 11500 22733 2274 2000 2000 2000 2000 2000 2000 2000
DEPN	2796 2796 2796 27950 21926 21926 29350 29350 29350 11503 15030 10000000000	2.6 646 188% 333%	94 44 88 74 48 88 74 48 88 74 48 88 74 48 88 74 48 88 74 74 74 74 78 74 78
1001	110000 11000 1000 1000000	11. 02.292 05.85 05.85 0000000000	222393 223393 22339 223
NI	163 163 163 163 164 173 193 194 193 194 193 193 193 193 193 193 193 193 193 193		1 N 1 222309562 2223099562 2223099562 2223099562 2223099562 2223099562 2223099562 2223099562 2223099562 2223099562 2223099562 2223099562 222309562 222309562 22250562 22550567 22550567 22550567 22550567 2255
STAN	11500 15500 15500 15500 15500 15500 15500 15500 15500 1500 1910 191	4.78 0 0	5 11 5 11 5 11 5 11 5 11 5 12 5 20 5 20
DEPN	11815 4617 4617 4617 4601 4601 4601 4601 4601 46130 19130 19130	4 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	78878 28778 287778 20778 207778 207778 207777777777
AGNI	116000 11700000 11700000000	4.00 0.8475 8.475 8.475	Н 111111111111111111111111111111111111
IN	1658 658 7487 7487 7487 7487 74957 74472 74957 1918 1918		IN 111229881 252088294 252088244 252088244 252088244 252088244 252088244 252088244 252088244 252088244
STAN	15000 17399 15507 15007 15007 150070	2.52 11147 295% 481%	2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2
DEPN	12275 6657 6657 7667 7667 7667 7667 7667	11-12-12-12-12-12-12-12-12-12-12-12-12-1	7264 7264 7264 7264 7264 7264 7264 7264
141 INDP	1500 1710 1710 1710 1710 1710 1710 1710	110 110 110 110 110 110 110 110 110 110	P 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2
NI	1111 1157 1157 115777 11577 11577 11577 11577 1157777 1157777 1157777 1157777 1157777 1157777 11577777 11577777777		22224222222222222222222222222222222222
5 5 4 12	0 1 1 1 1 1 1 1 1 1 1 1 1 1	cosr AQP RR RC	8 8 8 8 8 8 8 8 8 8 8 8 8 8

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LINDP 1	22239422000 272334220000 11940	4.39 120 3.9%	0016 0016 0017 0017 0017 0017 0017 0017
NI	222300 222300 222300 22230 22300 22000 22300 22000 22000 22000 22000 22000 22000 22000 22000 20000 20000 20000 20000 20000 2000000		H H H H H H H H H H H H H H H H H H H
STAN	22273445333000 22273445383000 22273445383000	5.05 0 0	5 5 5 5 5 5 5 5 5 5 5 5 5 5
JEPN 5	2222200134 2222200000000000000000000000000000000	4.03 1425 353% 353%	214 200 200 200 200 200 200 200 200 200 20
TINDP 1	2227355555 2227355555 2227355555 2227355555 2227355555 225555 225555 225555 225555 2255555 225555 225555 225555 225555 225555 225555 225555 225555 225555 2255555 2255555 225555 2255555 2255555 2255555 22555555	5.05	5 5 5 5 5 5 5 5 5 5 5 5 5 5
NI	22222222222222222222222222222222222222		2200934 230094 230004 2000000 2
ITAI	226899999999999999999999999999999999999	000 • 36	274 274 274 274 275 275 275 275 275 275 275 275
N A E C	2222500 2222500 2222500 2222500 2225000 2225000 2225000 222500000000	3.71 25280 2378	DEPN DEPN DEPN DEPN DEPN DEPN DEPN DEPN
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NT.	2222451 2222451 2222452 2222552 2222552 2222552 2222552 2222552 2222552 2222552 2222552 222552 222552 222552 222552 222552 222552 222552 2225552 2225552 2225552 2225552 2225552 22255552 22255552 222555555		HN 2222881 22282881 22281 2281 28851 28851 28851 2885
STAN	11222244300000 222244300700000 2222443007000000000000000000000000000000	8000 9000 000 000 000	27 28 28 28 28 28 28 28 28 28 28 28 28 28
DEPN	207 207 207 207 207 207 207 207 207 207	2.95 0 0	DEPN DEPN DEPN DEPN DEPN DEPN DEPN DEPN
INDP]	22600 222000 22000 22000 22000 22000 22000 22000 22000 22000 22000 22000 22000 22000 20000 20000 20000 20000 20000 20000 20000 20000 20000 20000 20000 20000 20000 20000 2000000	2、55 284 184炎 184炎	4 112222222222 4 11222222222222 5 000222222222222 5 0002222222222 5 000222222222 5 00022222222 5 00022222222 5 00022222222 5 00022222222 5 0002222222 5 0002222222 5 000222222 5 000222222 5 000222222 5 000222222 5 00022222 5 0002222 5 0002222 5 0002222 5 000222 5 000222 5 000222 5 00022 5 0002 5 0000 5 0002 5 0000 5 00000 5 0000 5 0000000 5 00000 5
IN	20589 20589 20589 20589 20585 20595 20585 20595 20505 20595 20505 20505 20505 20505 20505 20505 20505 20505 20505 20505 20505		1222232 12222222 12222222 1222222 122222 122222 122222 122222 1222 122 1222 1
TAN	12000000000000000000000000000000000000	3.29	6.4 4 20000000000000000000000000000000000
DEPN	11564 1554 1554 155555 1555555	2.79 438 13% 235%	20000 2010 2000000
TUDE	11566 1554 1554 155555 1555555	3• 29	86.26 86
NI	1044 1044 1044 1044 1044 1044 1044 1044		200122000220 200000200020 200000200020 2000000
. č +	111 210070000010 2	COST A Q RR RC	8 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2

Table D3.2 Simulation Result of March-April Flow

STAN	12000 10000 10000 10000 10000 10000 1000000	4,81 0000	64 4 20000000000000000000000000000000000
DEPN	4 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	4.05 982 27282 2728	20000 2000000
15 INDP	22000 22000 22000 22000 22000 22000 22000 22000 2000000	4.81 0 0	96,95,799,956,699,990 96,66,999,996,699,990 96,66,999,996,699,990 96,66,999,990,6499,990 96,66,999,990,6499,990 96,66,999,990,6499,990 96,66,999,990,6499,990 96,66,999,990,6499,990 96,66,999,990,6499,990 96,66,999,990,6499,990 96,66,999,990,990,990,990,990,990,990,99
NI	2005000 20050 20030 20030 20050 20000 20000 200000000		2440272 233250 233250 233250 233250 2350 23
STAN	2255002 2255002 2255002 2255002 2255002 25550002 255500002 255500002 255500000000	3.39	4 11040000000 4 4 1000000000000000000000000000000000000
DEPN	24420 24420 24421 24421 24421 24420 24020 2000000	5470 58820 10	DEPN DEPN 3384 3384 2555 2555 2555 2555 2555 2555 2555 25
1 ¹ 1001	24632000 2464200 2464202 2464200 2464200 2464200 2464200 246420000000000	3.39 0000	4 5550000000000000000000000000000000000
MI	22342 2342 2452 2452 2452 2452 32552 32552 2458 32252 2458 2458 2458 2458 2458 2458 2458		22 23 23 23 23 23 23 23 23 23 23 23 23 2
STAN	11112220000000000000000000000000000000	4•01 00 00	4 4204000000 8 4 4204000000 8 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
DEPN	64 64 64 64 64 64 64 64 64 64	3.664 2.67% 2.67% 2.67%	DEPN DEPN DEPN 221577 27157 27157 273577 273577 273577 273577 273577 273577 273577 273577 2735777 2735777 2735777 27357777 27357777777777
13 TNDP	444462460466666666666666666666666666666	4•01 0000	4 4200000000000000000000000000000000000
NI	11999 19759		42000000000000000000000000000000000000
STAN	1161192419 226019200 22000000000 20000000000 20000000000	3.27	8 111111111111111111111111111111111111
DEPN	222 222 222 222 222 222 222 222 222 22	3.27	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
12 INDP	11600 116000 116000 116000 116000 1100000000	2. 1481 2688 1441 2688	19 29 2000000000000000000000000000000000
II	2005 2005 2005 2005 2005 2005 2005 2005		112220505 1112220505 11222050505 11222050505 112220505 112220505 1122205 11225 11225 112555 112555 112555 112555 112555 112555 112555 1125555 112555 1125555 1125555555 11255555555
STAN	3669330072175000 366330072175000 366330072175000 3663300727272 269330072727272 2693300727272727272727272727272727272727272	4.12 0 0	8 11 11118 100000000000000000000000000000000
DEPN	22214 22252 22252 22252 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 2225 225 25	3.41 722 1728 2688	2000 200 2000 2
11 IMDP	1600 11600 11850 11600 11650 11753 117555 117555 117555 117555 1175555 117555 117555 117555 1175555 11	4•12 0 0	0000 000 000 000 000 000 000 00
NI	201121212121212121212121212121222222222		1 1 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2
	5 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	COST AQ _P RR RC	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8

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STAN	22002000000000000000000000000000000000	3,39	STAN STAN 33997 33997 33997 3000 3000 3000 3000 3
DEPN	200000 200000 200000 200000 200000 2000000	0 10 0 10 0 10 0 10 0 10 0 10 0 10 0 10	222 221 224 224 224 224 224 224 224 224
25 INDP I	222255557 22255557 22255557 22255557 2225555 2225555 222555 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 22255 225555 2255555 225555 2255555 2255555 225555 2255555 2255555 2255555 2255555 2255555 2255555 225555 2255555 2255555 2255555 2255555 2255555 2255555 22555555	0000	0 0 0 0 0 0 0 0 0 0 0 0 0 0
NI	22052 20052 20		231522332233255555555555555555555555555
TAN	22222 2222 222 22 222 222 222 222 22 222 222 222 222 22 222 222 222 22 22 22	0000	514500 1117799 115000 11500
S NAEO	1 1 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	3.64 3.64 7.151 8.03 7.15 1.5 7.5 8.5 7.5 7.5 8.5 7.5 8.5 7.5 7.5 8.5 7.5 7.5 7.5 7.5 7.5 7.5 7.5 7.5 7.5 7	22 22 22 22 22 22 22 22 22 22
24 LINDF I	8000742800000 10007420800000000000000000000000000000000	+0000	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
MI	2500 2500		10 10 10 10 10 10 10 10 10 10
TAN	333414020000000 84914002000000 8491400200000000 84914002000000000000000000000000000000000	334 2330 2330 2330 2330 2330 2330 2330 2	5 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2
DEPN S	24602 24602 24602 2402 2402 2402 2402 24	00100 0010 0002 0002	DEPN 2139 200502 200500000000
Z3 INDF	364000 364000 364000 364000 364000 364000 364000 364000 3640000000000	0010 000 000 000 000 000 000 000 000 00	1007 1500 11100 110000 110000 110000 110000 110000 110000 1100000 1100000 1100000 11000000
NI	2266833 384992 3849923 384992 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 384952 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 3845552 38455552 38455552 38455552 384555555555555555555555555555555555555		HIN 1902 1902 1902 1902 1902 1002 1
STAN	12000 1200 1200 1200 1200 1200 1200 120	4.51 0 0	8 11111122221111 6 8 222222221111 6 8 22222222222 8 200000 8 2000000 1 9 000000000 1 9 000000000000000
DEPN	22222222222222222222222222222222222222	3.63 9.63 3.118 3.	0 0 0 0 0 0 0 0 0 0 0 0 0 0
22 INDP	1200 1200 1200 1200 1200 1200 1200 1200	4•51 0 0	1000 1600 1800 1900 115000 115000 115000 115000 1100000000
NI	2145022 214502 214502 214502 214500 214500 214000 2145000 2145000 2145000 2145000 2145000000000000000000000000000000000000		21111111111111111111111111111111111111
STAN	15000 100000000	3.29	4 196100000000000000000000000000000000000
DEPN	2000 200 200 200 200 200 200 200 200 20	3.29	2225 2225 2225 2225 2225 2225 2225 222
21 INDP	20000000000000000000000000000000000000	3•29 0	0000 00000000000000000000000000000000
NI (222 222 222 222 232 232 222 222 222 222		10010000000000000000000000000000000000
	0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	COSI AQ _F RR RC	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8

STAN	24200000004000 000000004000 000000004000 000000	4•03 00000000000000000000000000000000000	5748 5748 5748 57588 5758 5758 5758 5758 5758 5758 5758 5758 5758 5758
DEPN	2010 2010 2010 2010 2010 2010 2010 2010	870 N 0 4 10 0 0 4 10 0	00000000000000000000000000000000000000
INDP 1	2420000000422000 2002000000422000 200400000000	4• 03	INDP 1600 1600 1600 1600 1600 1600 1600 160
NI	26420 2609 2609 2609 2609 2609 2609 2609 26		HN 1754 25530 25530 25558 21754 31150 302327 312558 312555
STAN	231505 231505 231505 231505 231505 231505 231505 231505 23150 2310	4.46 1.46 1.4%	572 372 372 372 372 372 472 572 572 572 572 572 572 572 5
DEPN	281282 28	4.07 481 104% 154%	224837 22400 2400000000
acui 1001	2000 2000 2000 2000 2000 2000 2000 200	そうか ひょうしょう ひょうしょう ひょうしょう ちょうしょう ちょうしょう ちょうしょう ちょうしょう ひょうしょう ひょう ひょう ひょう ひょう ひょう ひょう ひょう ひょう ひょう ひ	3 3 3 3 3 3 5 5 5 5 5 5 5 5
NI	2000 2000 2000 2000 2000 2000 2000 200	-	11811 11811 11811 11811 1177738 1177778 1177777778 1177777777
STAN	112222222200 23222222220 23222222220 2322222222	3.15	STAN 1500 1500 1500 1500 1500 1500 1500 150
DEPN	22 22 22 22 22 22 22 22 22 22 22 22 22	3•15 00 0	3 2524 2524 2524 25256 255666 255666 255666 255666 255666 255666 255666 255666 25566
33 1107	12222222222222222222222222222222222222	3.15 0 0	11111111111111111111111111111111111111
NI	13222222222222222222222222222222222222		NI 122222222 1222222222222222222222222222
STAN	29223 29233 2923 29233 20233 20233 20233 20233 20233 20233 20233 20233 20233 2	3.87 0 0	8 1111111 9 9 9 9 2 2 1 2 2 2 2 2 2 2 2 2 2 2 2 2
DEPN	29000 29000 29000 29000 29000 29000 29000 29000 29000 29000 29000 29000 29000 29000 29000 29000 29000	3•87 0 0	7 DEPN 1339 1539 1536 153
32 IMDP	29233 2923 29233 20233 20233 20233 20233 20233 20233 20233 20233 20233 2	3.87 0 0	3.93 2400000 1111111111111111111111111111111
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Table D3.3 Simulation Results of May-June Flow

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DEPN	10000000000000000000000000000000000000	1100 1000 1000 1000 1000 1000 1000 100	1000000000000000000000000000000000000	1100 1100 1100 1100 1100 100 100 100 10
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DEPN	122201211111 20202020 20202020 202020202	22-22 27-27-22 27	22250 2250 250
TNDP 42	110200000000000000000000000000000000000	3.08 737 306% 306%	20000000000000000000000000000000000000
NI	11000000000000000000000000000000000000		00014000000000000000000000000000000000
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22 INDP	11100000000000000000000000000000000000	3.91 0 0	1112222000 P12000 P12000 P12000 P12000 P120000 P120000 P120000 P120000 P120000 P12000 P12000 P12000 P12000 P12000 P12000 P12000 P120000 P120000 P120000 P1200
IN	126233 126233 1266333 1266333 1266333 1266333 126233 126233 12623 126233 12623 12633		12222222 12222222 12222222 1222222 122222 122222 1222 122
STAN	115000 15000 15000 15000 1010 1010 1000000	2.74 0 0	8 1 1 1 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2
DEPN	10000000000000000000000000000000000000	1.02 18740 18740 8888	2350 2350 2350 2350 2350 2350 2350 2350
21 ZUDP	11122222600 14022222200 14022000 14022000 100000 10000 10000 100000 100000 10000 10	2.74 0 0	10000000000000000000000000000000000000
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35 INDP	1112010111 0000000000000000000000000000	3. 69 0 0	AGNI 04	1600 1942 2000 2217	3338 2107 1500	11000 11000 1000 1000 1000 1000 1000 1	3•26 0
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34 JUDF	2115000 211500 211500 21150000000000	88644 100 10 10 10	100 100 100 100 100 100 100 100 100 100	1600 1674 1733 2752	3258 2115 1500	1500 1014 638 3058 538	3.17 0 0
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STAN	448848444 22000000000 200000000000000000	2.03 2.04 7.04 7.84 7.84	STAN	1500 1500 38760	2188 2419 1500	1111 00000 00000 00000 00000	60,400 27,40 28,80 28,80
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DEPN	22000 22000 2000000	1500 1500 1084 1068	1 400 1 br>100 100 100 100 100 100 100 100 10	12 14 14 14 14 14 14 14 14 14 14 14 14 14	DEPN	11946 22006 22006 2000 2000 2000 2000 2000	1850	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	276 818 818	1.68 832 2748 54,58 54,58
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NI	1675 1675 26630 24630	2491 631 1093 1093 831	4664		NI	11758 3032 1881	1757 2015 1313	6 6 6 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7	776 818 818	
STAN	12314 25000 25000 25000	1 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	1 661 661	107 107 107 107 107 107 107 107 107 107	STAN	11000 30500 2004 2004	2572 2572	11200	1660 845	3,51
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NI	14000 1400000000	4 4 5 5 5 7 5 7 7 7 7 7 7 7 7 7 7 7 7 7	891 661		NI e	1625 3574 3574	2572	1051 916	8559 8459 2450	22.0
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