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Bounding the size of an almost-equidistant set in Euclidean space

Andrey Kupavskii^{*} Nabil H. Mustafa^{†‡} Konrad J. Swanepoel[§]

Abstract

A set of points in *d*-dimensional Euclidean space is *almost equidistant* if among any three points of the set, some two are at distance 1. We show that an almost-equidistant set in \mathbb{R}^d has cardinality $O(d^{4/3})$.

1 Introduction

A set of lines through the origin of Euclidean *d*-space \mathbb{R}^d is almost orthogonal if among any three of the lines, some two are orthogonal. Erdős asked (see [12]) what is the largest cardinality of an almost-orthogonal set of lines in \mathbb{R}^d ? By taking the union of two sets of *d* pairwise orthogonal lines, we see that 2*d* is possible. Rosenfeld [12] showed that 2*d* is the maximum by considering the eigenvalues of the Gramian of the unit vectors spanning the lines. His result was subsequently given simpler proofs by Pudlák [11] and Deaett [6].

In this note we consider the analogous notion obtained by replacing orthogonal pairs of lines by pairs of points at unit distance. A subset V of Euclidean d-space \mathbb{R}^d is almost equidistant if among any three points in V, some two are at Euclidean distance 1. We investigate the largest size, which we denote by f(d), of an almost-equidistant set in \mathbb{R}^d . Although asking for the size of this function is a very natural question, it seems to be harder than the question of Erdős, which can be refomulated as asking for the largest size of an almost-equidistant set on a sphere of radius $1/\sqrt{2}$ in \mathbb{R}^d . Before stating our main result, we give an overview of what is known about f(d).

Bezdek, Naszódi and Visy [5] showed that $f(2) \leq 7$, and István Talata (personal communication, 2007) showed that the only almost-equidistant set in \mathbb{R}^2 with 7 points is the Moser spindle. Györey [7] showed that $f(3) \leq 10$ and that there is a unique almost-equidistant set of 10 points in \mathbb{R}^3 , a configuration originally considered by Nechushtan [9]. The Moser spindle can be generalized to higher dimensions, giving an almost-equidistant set of 2d + 3 points in \mathbb{R}^d [4]. (We mention that Bezdek and Langi [4] considered the variant of Erdős's problem where the radius of the sphere is arbitrary instead of $1/\sqrt{2}$.) A

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construction of Larman and Rogers [8] shows that $f(5) \ge 16$. Since there does not exist a set of d+2 points in \mathbb{R}^d that are pairwise at distance 1, it follows that $f(d) \le R(d+2,3)-1$, where the Ramsey number R(a, b) is the smallest n such that whenever each edge of the complete graph on n vertices is coloured blue or red, there is either a blue clique of size a or a red clique of size b. Ajtai, Komlós, and Szemerédi [1] showed $R(k,3) = O(k^2/\log k)$, which implies the asymptotic upper bound $f(d) \le O(d^2/\log d)$. Balko, Pór, Scheuer, Swanepoel, and Valtr [3] generalized the Nechushtan configuration to higher dimensions, giving $f(d) \ge 2d + 4$ for all $d \ge 3$. They also obtained the asymptotic upper bound $f(d) = O(d^{3/2})$ by an argument based on Deaett's paper [6]. Using computer search and ad hoc geometric arguments, they obtained the following bounds for small $d: f(4) \le 13$, $f(5) \le 20$, $18 \le f(6) \le 26$, $20 \le f(7) \le 34$, and $f(9) \ge f(8) \ge 24$. Polyanskii [10] subsequently improved the asymptotic upper bound to $f(d) = O(d^{13/9})$.

In this note we obtain a further improvement to the upper bound.

Theorem 1. An almost-equidistant set in \mathbb{R}^d has cardinality $O(d^{4/3})$.

Its proof is based on the approach of [3] to show the upper bound $O(d^{3/2})$, which is in turn based on Deaett's proof [6] of Rosenfeld's result. Before proving Theorem 1 in Section 3, we establish notation and collect some lemmas in the next section.

2 Preliminaries

We denote the Euclidean norm of $x \in \mathbb{R}^d$ by ||x|| and the inner product of $x, y \in \mathbb{R}^d$ by $\langle x, y \rangle$. The cardinality of a finite set A is denoted by |A|. We call a finite non-empty subset C of \mathbb{R}^d (the vertex set of) a *unit simplex* if the distance between any two points in C equals 1. It has already been mentioned in the Introduction that if C is a unit simplex then $|C| \leq d + 1$. Given any finite $V \subset \mathbb{R}^d$, we define the *unit-distance graph* G = (V, E) on V to be the graph with $vw \in E$ iff ||v - w|| = 1. Thus, $C \subset V$ is a unit simplex iff it is a clique in G. We denote the set of neighbours of $v \in V$ in G by N(v).

The following well-known lemma gives a lower bound for the rank of a square matrix in terms of its entries [2, 6, 11].

Lemma 1. For any non-zero $n \times n$ symmetric matrix $A = [a_{i,j}]$,

$$\operatorname{rank}(A) \ge \frac{(\sum_{i} a_{i,i})^2}{\sum_{i,j} a_{i,j}^2}.$$

For the sake of completeness, we include the proofs of the following three lemmas on the vertices and centroids of unit simplices.

Lemma 2. Let C be a unit simplex with centroid $c = \frac{1}{|C|} \sum_{v \in C} v$. Then

$$||v - c||^2 = \frac{1}{2} \left(1 - \frac{1}{|C|} \right) \text{ for all } v \in C,$$

and

$$\langle v - c, v' - c \rangle = -\frac{1}{2|C|}$$
 for all distinct $v, v' \in C$.

Proof. We may translate C so that c is the origin o. Write $C = \{p_1, \ldots, p_k\}$. By symmetry, $\alpha := \|p_i\|^2$ is independent of i, and $\beta := \langle v_i, v_j \rangle$ $(i \neq j)$ is independent of i and j. Then

$$0 = \left\|\sum_{i=1}^{k} p_i\right\|^2 = k\alpha + k(k-1)\beta$$

and

$$1 = \|p_i - p_j\|^2 = 2\alpha - 2\beta.$$

Solving these two linear equations in α and β , we obtain $\alpha = \frac{1}{2} - \frac{1}{2k}$ and $\beta = -\frac{1}{2k}$. \Box Lemma 3. Let C be a unit simplex with centroid $c = \frac{1}{|C|} \sum_{v \in C} v$, and let $F \subset C$ be a unit simplex with centroid $f = \frac{1}{|F|} \sum_{v \in F} v$. Then

$$||c - f||^2 = \frac{1}{2} \left(\frac{1}{|F|} - \frac{1}{|C|} \right)$$

Proof. Let $k := |C|, \ell := |F|$. Then

$$\|f - c\|^{2} = \left\| \frac{1}{\ell} \sum_{v \in F} (v - c) \right\|^{2} = \frac{1}{\ell^{2}} \left(\ell \cdot \frac{1}{2} \left(1 - \frac{1}{k} \right) - \frac{\ell(\ell - 1)}{2k} \right) \quad \text{by Lemma 2}$$
$$= \frac{1}{2} \left(\frac{1}{\ell} - \frac{1}{k} \right).$$

Lemma 4. Let A and B be disjoint unit simplices with centroids $a = \frac{1}{|A|} \sum_{v \in A} v$ and $b = \frac{1}{|B|} \sum_{v \in B} v$, respectively, such that $A \cup B$ is also a unit simplex. Then

$$|a-b||^2 = \frac{1}{2} \left(\frac{1}{|A|} + \frac{1}{|B|} \right).$$

Proof. Let $C := A \cup B$ have centroid c. Then by Lemma 3, $||a - c||^2 = \frac{1}{2} \left(\frac{1}{|A|} - \frac{1}{|C|} \right)$ and $||b - c||^2 = \frac{1}{2} \left(\frac{1}{|B|} - \frac{1}{|C|} \right)$. It follows that

$$\begin{aligned} \|a-b\|^2 &= (\|a-c\| + \|b-c\|)^2 \\ &= \frac{1}{2} \left(\frac{1}{|A|} + \frac{1}{|B|} - \frac{2}{|C|} \right) + \sqrt{\left(\frac{1}{|A|} - \frac{1}{|C|} \right) \left(\frac{1}{|B|} - \frac{1}{|C|} \right)} \\ &= \frac{1}{2} \left(\frac{1}{|A|} + \frac{1}{|B|} \right) \end{aligned}$$

3 Proof of Theorem 1

Let G be the unit-distance graph of a given almost-equidistant set V. Then the complement of G is K_3 -free, and the non-neighbours of any vertex form a unit simplex. Let C be a clique of maximum cardinality in G. Write k = |C|. Each $v \in V \setminus C$ is a non-neighbour of some point in C, and it follows that $|V| \leq |C| + |C| k = k^2 + k$. Thus, without loss of generality, $k > d^{2/3}$. We split V up into two parts, each to be bounded separately. Let

$$N = \left\{ v \in V : |N(v) \cap C| \ge k - k^{4/3} d^{-2/3} \right\}.$$

Note that $k^{4/3}d^{-2/3} = O(d^{-1/3})k$. We first bound the complement of N. Consider the set

$$X = \{(u, v) \in C \times V \setminus N : uv \notin E(G)\}$$

For each $v \in V \setminus N$, there are more than $k^{4/3}d^{-2/3}$ points $u \in C$ such that $u \notin N(v)$, hence $|X| > k^{4/3}d^{-2/3}|V \setminus N|$. On the other hand, for each $u \in C$, the set of nonneighbours of u forms a clique, so has cardinality at most k, and $|X| \leq |C| k = k^2$. It follows that

$$|V \setminus N| < k^{2/3} d^{2/3}.$$
 (1)

Next, we estimate |N|. Without loss of generality, $\frac{1}{k} \sum_{v \in C} v = o$ and $N = \{v_1, \ldots, v_n\}$. We want to apply Lemma 1 to the $n \times n$ matrix $A = [\langle v_i, v_j \rangle]$, which has rank at most d.

Claim 1. For each i = 1..., n, $||v_i||^2 = \frac{1}{2} + O(k^{-1/3}d^{-1/3})$, and for each $v_i v_j \in E(G)$, $\langle v_i, v_j \rangle = O(k^{-1/3}d^{-1/3})$.

Proof of Claim 1. Let $C_i := N(v_i) \cap C$, $k_i := |C_i|$, and $c_i := \frac{1}{k_i} \sum_{v \in C_i} v$. Then $k_i \ge k - k^{4/3} d^{-2/3}$. By Lemma 4 applied to $A = \{v_i\}$ and $B = C_i$, $\|v_i - c_i\|^2 = \frac{1}{2} \left(1 + \frac{1}{k_i}\right)$, hence $\|v_i - c_i\| = \frac{1}{\sqrt{2}} + O(k^{-1})$. By Lemma 3 applied to C and $F = C_i$,

$$\|c_i\| = \sqrt{\frac{1}{2}\left(\frac{1}{k_i} - \frac{1}{k}\right)} \le \sqrt{\frac{1}{2}\left(\frac{1}{k - k^{4/3}d^{-2/3}} - \frac{1}{k}\right)} = O(k^{-1/3}d^{-1/3}).$$

By the triangle inequality,

$$||v_i|| = ||v_i - c_i|| + O(||c_i||) = \frac{1}{\sqrt{2}} + O(k^{-1/3}d^{-1/3}),$$

and $||v_i||^2 = \frac{1}{2} + O(k^{-1/3}d^{-1/3})$. Also, $2\langle v_i, v_j \rangle = ||v_i||^2 + ||v_j||^2 - 1 = O(k^{-1/3}d^{-1/3})$. \Box

Claim 2. For each i = 1, ..., n,

$$\sum_{\substack{j=1\\v_iv_j\notin E(G)}}^n \langle v_i, v_j \rangle^2 = O(k^{2/3}d^{-1/3}).$$

Proof of Claim 2. The non-neighbours $N \setminus N(v_i)$ of v_i form a unit simplex with cardinality $t := |N \setminus N(v_i)| \le k$ and with centroid c, say. If t = d + 1, remove one point v_j from the unit simplex, which decreases the sum by $\langle v_i, v_j \rangle^2 = O(1)$. Thus, without loss of generality, $t \le d$, and there exists a point $p \in \mathbb{R}^d$ such that p - c is orthogonal to the affine hull of $N \setminus N(v_i)$, $||p - c|| = 1/\sqrt{2t}$, the set $\{v_j - p : v_j \in N \setminus N(v_i)\}$ is orthogonal, and $||v_j - p|| = 1/\sqrt{2}$ for each non-neighbour v_j of v_i . Then, by the finite Bessel inequality,

$$\sum_{v_j \in N \setminus N(v_i)} \langle v_i, v_j - p \rangle^2 \le \frac{1}{2} \left\| v_i \right\|^2,$$

hence by applying Cauchy–Schwarz a few times,

$$\sum_{v_{j} \in N \setminus N(v_{i})} \langle v_{i}, v_{j} \rangle^{2} = \sum_{v_{j} \in N \setminus N(v_{i})} \left(\langle v_{i}, v_{j} - p \rangle + \langle v_{i}, p - c \rangle + \langle v_{i}, c \rangle \right)^{2}$$

$$\leq 3 \sum_{v_{j} \in N \setminus N(v_{i})} \left(\langle v_{i}, v_{j} - p \rangle^{2} + \langle v_{i}, p - c \rangle^{2} + \langle v_{i}, c \rangle^{2} \right)$$

$$\leq 3 \left(\frac{1}{2} \|v_{i}\|^{2} + t \|v_{i}\|^{2} \|p - c\|^{2} + t \|v_{i}\|^{2} \|c\|^{2} \right)$$

$$\leq 3(\|v_{i}\|^{2} + t \|v_{i}\|^{2} \|c\|^{2}).$$
(2)

By Claim 1, $||v_i||^2 = \frac{1}{2} + O(k^{-1/3}d^{-1/3})$ and

$$\begin{aligned} \|c\|^{2} &= \left\| \frac{1}{t} \sum_{v_{j} \in N \setminus N(v_{i})} v_{j} \right\|^{2} = \frac{1}{t^{2}} \left(\sum_{v_{j} \in N \setminus N(v_{i})} \|v_{j}\|^{2} + \sum_{v_{j}, v_{j'} \in N \setminus N(v_{i})} \left\langle v_{j}, v_{j'} \right\rangle \right) \\ &\leq \frac{1}{t^{2}} \left(t \left(\frac{1}{2} + O(k^{-1/3}d^{-1/3}) \right) + t(t-1)O(k^{-1/3}d^{-1/3}) \right) \\ &= \frac{1}{2t} + O(k^{-1/3}d^{-1/3}). \end{aligned}$$

Therefore,

$$t ||c||^2 = \frac{1}{2} + O(tk^{-1/3}d^{-1/3}) = O(k^{2/3}d^{-1/3}).$$

Substitute this back into (2) to finish the proof of Claim 2.

We now finish the proof of the theorem. By Claim 2,

$$\sum_{j=1}^{n} \langle v_i, v_j \rangle^2 = \|v_i\|^4 + \sum_{v_j \in N(v_i)} \langle v_i, v_j \rangle^2 + O(k^{2/3} d^{-1/3})$$
$$= nO(k^{-2/3} d^{-2/3}) + O(k^{2/3} d^{-1/3}) \quad \text{by Claim 1.}$$

Also by Claim 1, $\sum_{i=1}^{n} \|v_i\|^2 = \Omega(n)$. Therefore, by Lemma 1,

$$d \ge \operatorname{rank}(A) \ge \frac{\left(\sum_{i=1}^{n} \|v_i\|^2\right)^2}{\sum_{i,j=1}^{n} \langle v_i, v_j \rangle^2} = \frac{\Omega(n^2)}{n\left(nO(k^{-2/3}d^{-2/3}) + O(k^{2/3}d^{-1/3})\right)},$$

hence $n = O(nk^{-2/3}d^{1/3}) + O(k^{2/3}d^{2/3})$. Since $O(k^{-2/3}d^{1/3}) = o(1)$, it follows that $|N| = n = O(k^{2/3}d^{2/3})$. Recalling (1), we obtain that

$$|V| = |N| + |V \setminus N| = O(k^{2/3}d^{2/3}) = O(d^{4/3}).$$

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