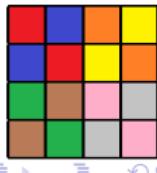
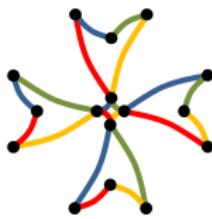
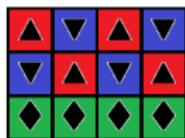


Autotopism stabilized colouring games on rook's graphs

Raúl M. Falcón (University of Seville).

(Joint work with S. Dominique Andres).
(FernUniversität in Hagen).

Qwara. June 29, 2017.



CONTENTS

- ① Colouring games and partial Latin rectangles.
- ② Compatibility and feasibility.
- ③ The autotopism stabilized colouring game.

I. Colouring games and partial Latin rectangles.

Colouring games.

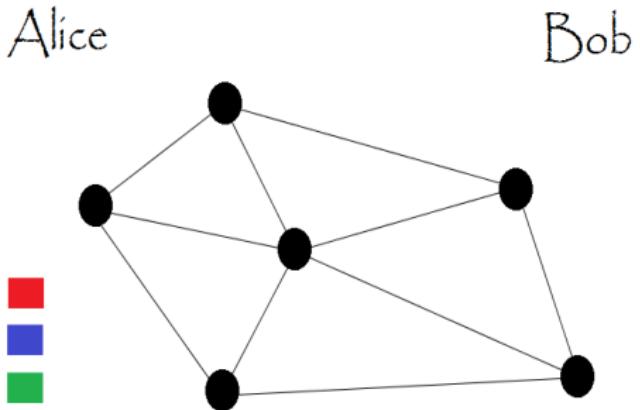


Steven J. Brams & Martin Gardner, 1981.



Hans L. Boblaender, 1991.

Based on the graph colouring problem, the game is played on a finite graph by two players, Alice (A) and Bob (B), with Alice playing first.



Colouring games.

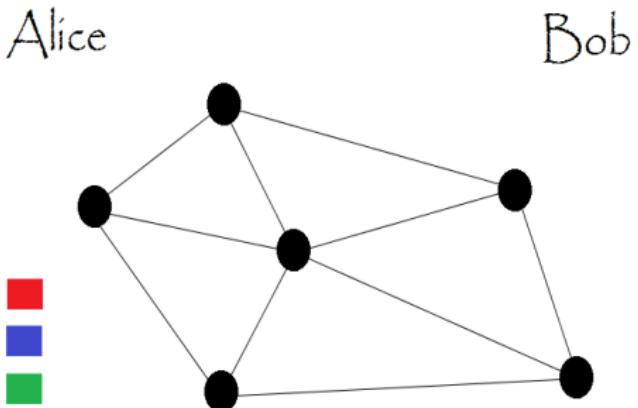


Steven J. Brams & Martin Gardner, 1981.



Hans L. Boblaender, 1991.

They alternately colour the vertices of the graph so that none two adjacent vertices are coloured with the same colour.



Colouring games.



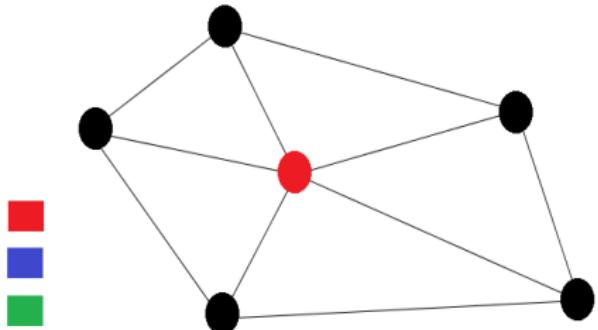
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Alice ★ Bob



Colouring games.

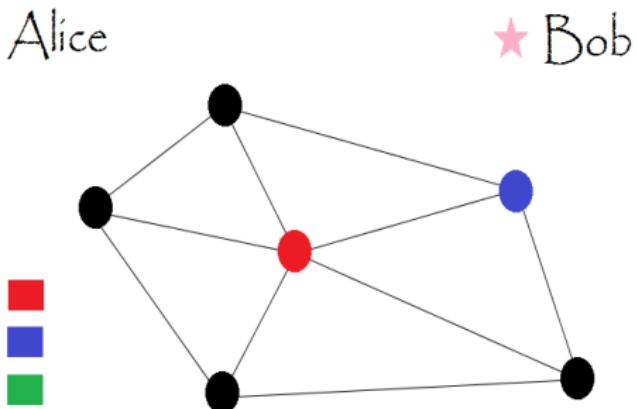


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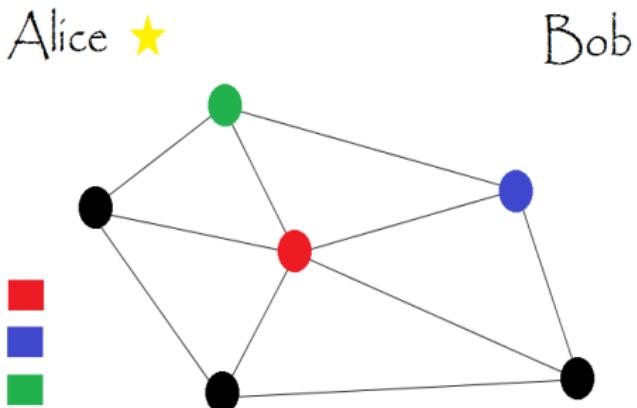


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Colouring games.



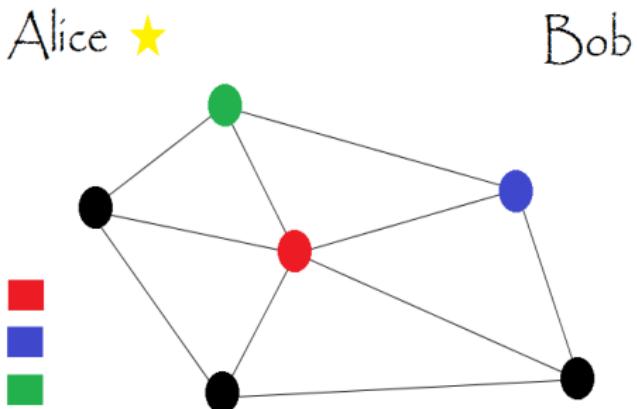
Steven J. Brams & Martin Gardner, 1981.



Hans L. Boblaender, 1991.

Alice wins if all vertices are coloured at the end of the game.

Otherwise, Bob wins.



Colouring games.



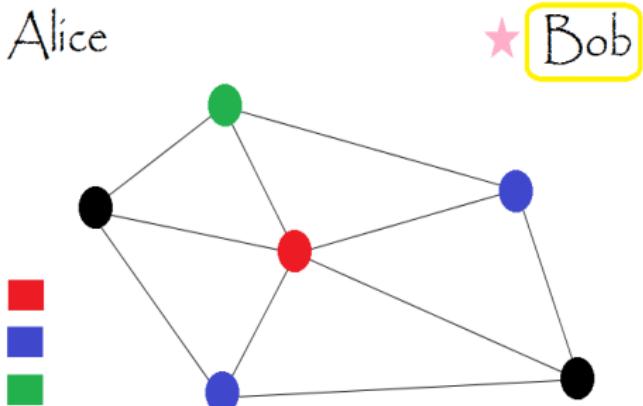
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Colouring games.



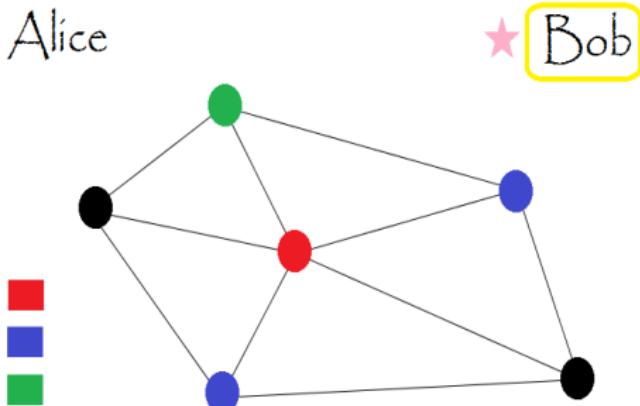
Steven J. Brams & Martin Gardner, 1981.



Hans L. Boblaender, 1991.

Alice wins if all vertices are coloured at the end of the game.

Otherwise, Bob wins.



Does there exist a winning strategy for the players?

Colouring games.

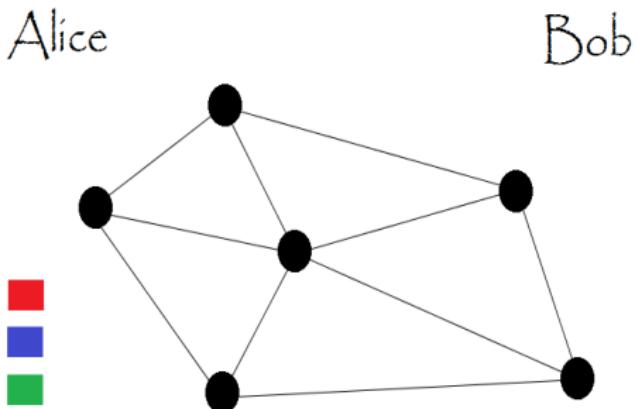


Steven J. Brams & Martin Gardner, 1981.



Hans L. Boblaender, 1991.

$\chi_g(G)$: Game chromatic number.
Least integer k s. t. A has a winning strategy on a graph G and k colours.



Colouring games.



Steven J. Brams & Martin Gardner, 1981.

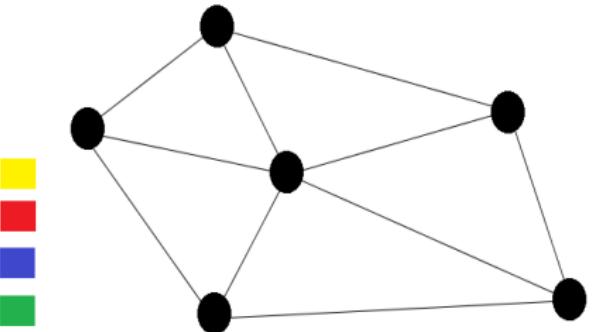


Hans L. Boblaender, 1991.

$\chi_g(G)$: Game chromatic number.

Least integer k s. t. A has a winning strategy on a graph G and k colours.

Alice Bob



$$\chi_g(G) = 4.$$

Colouring games.



Steven J. Brams & Martin Gardner, 1981.

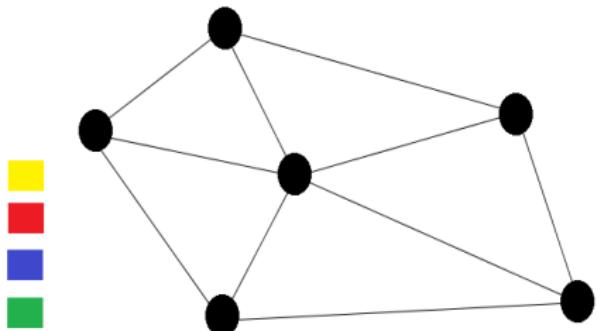


Hans L. Boblaender, 1991.

$\chi_{g_A}(G)$: Alice starts.

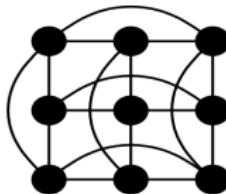
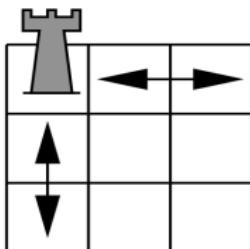
$\chi_{g_B}(G)$: Bob starts.

Alice Bob



$$\chi_{g_A}(G) = \chi_{g_B}(G) = 4.$$

Colouring games on rook's graphs.

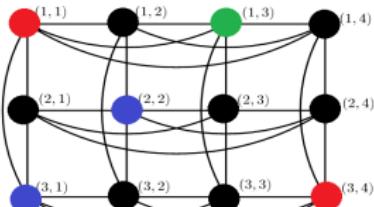


Maximilian Schlund, 2011.

Rook's graph \equiv Hamming graph: $\mathcal{R}_{r,s} = \mathcal{H}_{r,s} = K_r \square K_s$.

Lemma

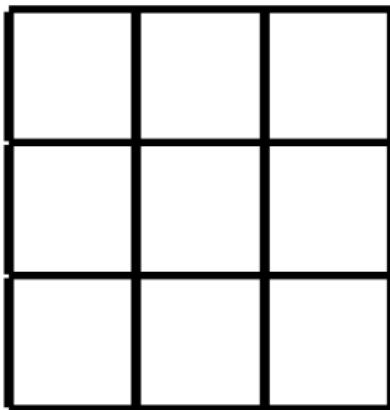
Every $r \times s$ partial Latin rectangle based on n symbols is uniquely identified with a partial n -colouring of a vertex-labeled rook's graph $\mathcal{R}_{r,s}$.



1	4	
2		
2		1

Colouring games on rook's graphs.

1 2 3



Colouring games on rook's graphs.

1 2 3



1		



Colouring games on rook's graphs.

1 2 3



1		
	2	



Colouring games on rook's graphs.

1 2 3



1	!	!
!	2	!
!	!	!



Colouring games on rook's graphs.

1 2 3



1		
	2	
		3



Colouring games on rook's graphs.

1 2 3



1		
	2	
		3

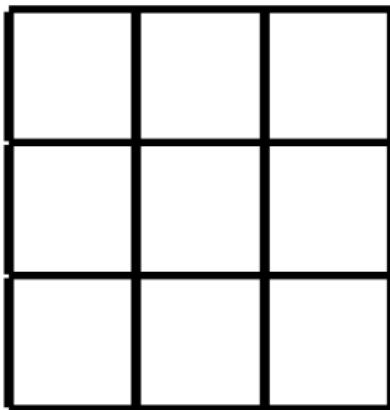


$$\chi_{g_A} = 3.$$



Colouring games on rook's graphs.

1 2 3



Colouring games on rook's graphs.

1 2 3



1		



Colouring games on rook's graphs.

1 2 3



1	!	!
!	!	!
!	!	!

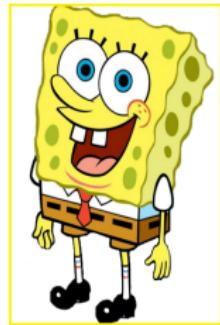


Colouring games on rook's graphs.

1 2 3



1	!	!
!	!	!
!	!	!



Colouring games on rook's graphs.

1 2 3



1	!	!
!	!	!
!	!	!



$\chi_{g_B} > 3.$

Partial Latin rectangles ($\mathcal{PLR}_{r,s,n}$).

- **Partial Latin rectangle:** $r \times s$ array where

- ① each cell is empty or contains one symbol of $[n] := \{1, \dots, n\}$.
- ② each symbol occurs at most once in each row and each column.

$$P \equiv \begin{array}{|c|c|c|c|} \hline 1 & & 3 & \\ \hline & 2 & & 4 \\ \hline 2 & 3 & 4 & 1 \\ \hline \end{array}$$

- **Entry set:** $\{(row, column, symbol)\}$.

$$\text{Ent}(P) = \left\{ (1, 1, 1), (1, 3, 3), (2, 2, 2), (2, 4, 4), \right. \\ \left. (3, 1, 2), (3, 2, 3), (3, 3, 4), (3, 4, 1) \right\}.$$

- $\mathfrak{I}_{r,s,n} := S_r \times S_s \times S_n \Rightarrow \text{Isotopism}$: $(\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n})$.

$$\text{Ent}(P^\Theta) = \{(\alpha(i), \beta(j), \gamma(k)) \mid (i, j, k) \in \text{Ent}(P)\}.$$

- **Principal isotopism**: $\gamma = \text{Id}$.

- **Autotopism**: $P^\Theta = P$.

$$\Theta = ((12), (1234), (1234)) \in \mathfrak{I}_{3,4,4}.$$

Partial Latin rectangles ($\mathcal{PLR}_{r,s,n}$).

- **Conjugate**: $\Theta_1 = \Theta\Theta_2\Theta^{-1} \equiv$ Equivalence relation in $\mathfrak{I}_{r,s,n}$.
- Classes characterized by the **cycle structure** of their elements.
- **Cycle structure** of $\pi \in S_m$: $z_\pi := m^{\lambda_m^\pi} \dots 1^{\lambda_1^\pi}$.
 $\lambda_i^\pi := \#\{i\text{-cycles in the unique cycle decomposition of } \pi\}$.
- **Cycle structure** of $(\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}$: $z_\Theta = (z_\alpha, z_\beta, z_\gamma)$.
 $((1234), (12)(3)(45), (12)(345)(6)) \in S_4 \times S_5 \times S_6 \rightarrow (4, 2^2 1, 321)$.
- $\Theta \in \mathfrak{I}_{r,s,n} \Rightarrow \mathcal{PLR}_\Theta = \{P \in \mathcal{PLR}_{r,s,n} \mid P^\Theta = P\}$.

Lemma

The cardinality of \mathcal{PLR}_Θ only depends on the cycle structure of Θ .

Partial Latin rectangles ($\mathcal{PLR}_{r,s,n}$).

Lemma

Every isotopism $\Theta \in \mathfrak{I}_{r,s,n}$ determines a unique partition of the cells of a partial Latin rectangle in $\mathcal{PLR}_{r,s,n}$.

- $\lambda_{\pi,i}$: Length of the cycle in $\pi \in S_m$ containing $i \in [m]$.
- **Cell orbit** of $(i,j) \in [r] \times [s]$ under the action of $\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}$.

$$\circ_\Theta((\mathbf{i}, \mathbf{j})) = \{(\alpha^m(i), \beta^m(j)) \mid m \leq \text{lcm}(\lambda_{\alpha,i}, \lambda_{\beta,j})\}.$$

$$\Theta = ((12), (1234), (1234)) \in S_3 \times S_4 \times S_4.$$

▲	▼	▲	▼
▼	▲	▼	▲
◆	◆	◆	◆

Partial Latin rectangles ($\mathcal{PLR}_{r,s,n}$).

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1		3	
	2		
2	3	4	1

Marked orbit.
Symbol-free orbit.
Complete orbit.

Partial Latin rectangles ($\mathcal{PLR}_{r,s,n}$).

Proposition

$$\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}.$$

$$P = (p_{ij}) \in \mathcal{PLR}_{r,s,n}.$$

$P \in \mathcal{PLR}_\Theta$ if and only if the next two conditions hold.

- Every cell orbit in P is complete or symbol-free.
- If a cell (i,j) is non-empty, then its cell orbit is complete and formed by the entries

$$(\alpha^m(i), \beta^m(j), \gamma^m(p_{ij})).$$

$$\Theta = ((12), (1234), (1234)) \in S_3 \times S_4 \times S_4.$$

1	3		
2	4		
2	3	4	1

II. Compatibility and feasibility.

Compatibility.

A triple of positive integers (i, j, k) is **lcm-compatible** if

$$\text{lcm}(i, j, k) = \text{lcm}(i, j) = \text{lcm}(i, k) = \text{lcm}(j, k).$$

Proposition

$$\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}.$$

$$P \in \mathcal{PLR}_\Theta.$$

$(i, j, k) \in \text{Ent}(P) \Rightarrow (l_{\alpha,i}, l_{\beta,j}, l_{\gamma,k})$ is lcm-compatible.

$$\Theta = ((12), (1234), (1234)) \in S_3 \times S_4 \times S_4.$$

1		3	
	2		4
2	3	4	1

Compatibility.

$$\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}.$$

$P \in \mathcal{PLR}_{r,s,n}$ is **Θ -compatible** if, for each $(i, j, k) \in \text{Ent}(P)$,

C.1) $(l_{\alpha,i}, l_{\beta,j}, l_{\gamma,k})$ is lcm-compatible.

C.2) For each $m \geq 0$,

- Either the cell $(\alpha^m(i), \beta^m(j))$ is empty or
- $(\alpha^m(i), \beta^m(j), \gamma^m(k)) \in \text{Ent}(P)$.

$$\Theta = ((12), (1234), (1234)) \in S_3 \times S_4 \times S_4.$$

1		3	
	2		
2	3	4	1

Compatibility.

$$\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}.$$

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C.2) For each $m \geq 0$,

- Either the cell $(\alpha^m(i), \beta^m(j))$ is empty or
- $(\alpha^m(i), \beta^m(j), \gamma^m(k)) \in \text{Ent}(P)$.

Theorem

$\mathcal{PLR}_\Theta \neq \emptyset$ if and only if there exists a Θ -compatible $P \in \mathcal{PLR}_{r,s,n}$ for which every marked orbit is complete.

Feasibility.

$$\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}.$$

- **Feasible:** $\lambda_i^\alpha \lambda_j^\beta \lambda_k^\gamma > 0 \Rightarrow (i, j, k)$ lcm-compatible.

((12)(34)(56), (123)(456), (123456)).

- **Extension of size** $n' \geq n$:

Any $(\alpha, \beta, \gamma') \in \mathfrak{I}_{r,s,n'}$, s. t. $\gamma'(s) = \gamma(s), \forall s \leq n$.

- **Natural extension:** $\gamma'(s) = s, \forall n < s$.

- **Extendable isotopism:** All its natural extensions are feasible.

Proposition

- Extendability \Rightarrow Feasibility.
- If a natural extension of Θ is feasible, then Θ is feasible.
- Every feasible principal isotopism is extendable.

Feasibility.

$$\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}.$$

- **Feasible:** $\lambda_i^\alpha \lambda_j^\beta \lambda_k^\gamma > 0 \Rightarrow (i, j, k)$ lcm-compatible.

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- **Natural extension:** $\gamma'(s) = s, \forall n < s$.

- **Extendable isotopism:** All its natural extensions are feasible.

Theorem

(α, β, γ) is extendable if and only if

- it is feasible, and
- all the cycles of α and β have the same length.

III. The autotopism stabilized colouring game.

The game.

- **Initial board:** Empty $r \times s$ array: $P = (p_{ij})$.
- **Key:** $\Theta = (\alpha, \beta, \gamma) \in \mathfrak{I}_{r,s,n}$.

Basic rules:

Alternately, the players colour an empty cell $(i, j) \in [r] \times [s]$ in the board with a colour $c \in [n]$ by respecting the following rules:

- **Rule 1)** The Latin array condition must always hold.
- **Rule 2)** The Θ -compatibility condition must always hold.
- **Rule 3)** For each $m < \text{lcm}(l_{\alpha,i}, l_{\beta,j})$ and each coloured collinear cell (i', j') of $(\alpha^m(i), \beta^m(j))$, it must be $p_{i'j'} \neq \gamma^m(c)$.

$$i_1 \rightarrow \begin{array}{|c|c|c|} \hline & i_2 \\ \downarrow & & \\ \hline s & \neq s & \neq s \\ \hline \neq s & & \\ \hline \neq s & & \\ \hline \end{array}$$

Rule 1

$$\begin{array}{|c|c|} \hline & i_2 \\ \downarrow & \beta^m(i_2) \\ \hline s & \\ \hline \end{array}$$
$$i_1 \rightarrow \begin{array}{|c|c|} \hline & \\ \hline \end{array}$$
$$\alpha^m(i_1) \rightarrow \begin{array}{|c|c|} \hline & \\ \hline \end{array}$$
$$\gamma^m(s)$$

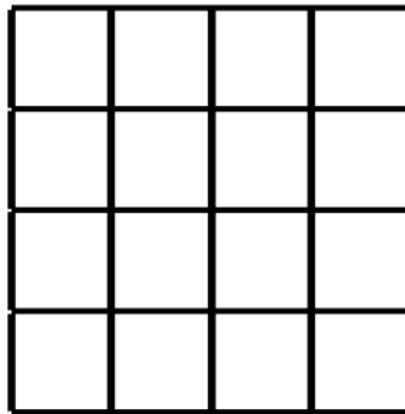
Rule 2

$$\begin{array}{|c|c|c|} \hline & i_2 \\ \downarrow & \beta^m(i_2) \\ \hline s & & \neq \gamma^m(s) \\ \hline & & \neq \gamma^m(s) \\ \hline \neq \gamma^m(s) & \neq \gamma^m(s) & \\ \hline \end{array}$$
$$i_1 \rightarrow \begin{array}{|c|c|c|} \hline & & \\ \hline \end{array}$$
$$\alpha^m(i_1) \rightarrow \begin{array}{|c|c|c|} \hline & & \\ \hline \end{array}$$

Rule 3

The game.

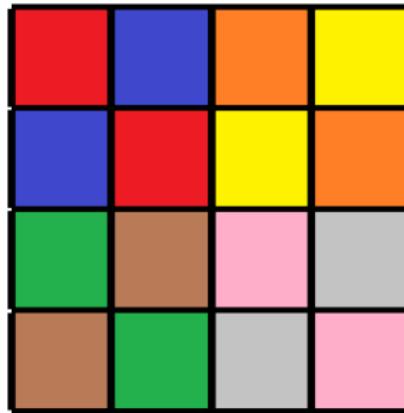
1 2 3 4



$$\Theta = ((12)(34), (12)(34), \text{Id})$$

The game.

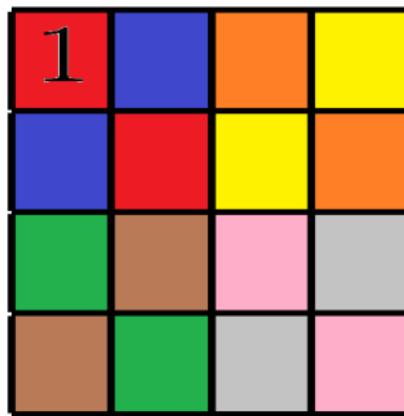
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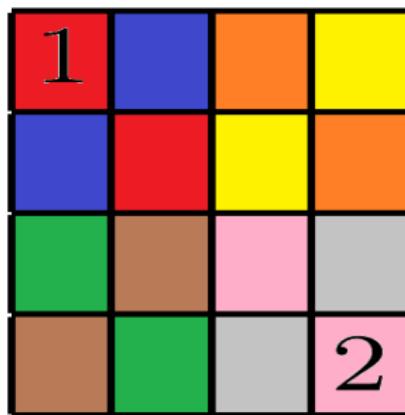
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The game.

1 2 3 4

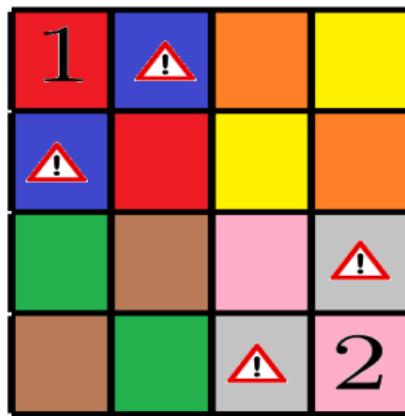


$$\Theta = ((12)(34), (12)(34), \text{Id})$$



The game.

1 2 3 4

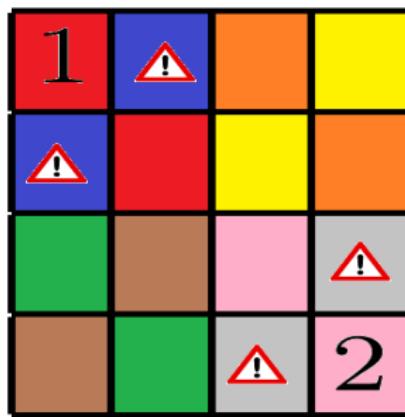


$$\Theta = ((12)(34), (12)(34), \text{Id})$$



The game.

1 2 3 4



$$\chi_{g_A}^{\Theta} > 4.$$



The game.

$$\Theta \in \mathfrak{I}_{r,s,n}.$$

Lemma

Alice always wins the Θ -stabilized colouring game when a player colours any cell of the last symbol-free cell orbit of the board.

Proposition (Well-defined game)

Let Θ be extendable and let P be Θ -compatible, satisfying Rule 3 and with at least one symbol-free cell orbit. Then,

- any empty cell in P can be coloured.*
- If a symbol $s > n$ of an extension of Θ does not appear in any cell of P , then there exists at least one empty cell in P that can be coloured with the colour s by obeying Rules 1–3.*

The game.

$$\Theta \in \mathfrak{I}_{r,s,n}$$

- **Θ -stabilized game chromatic number (χ_g^Θ)**: Smallest $n' > n$ for the natural extension of Θ for which Alice has a winning strategy.

Proposition

Let Θ_1 and Θ_2 be two extendable isotopisms with the same cycle structure. Then,

$$\chi_g^{\Theta_1} = \chi_g^{\Theta_2}.$$

Lemma

If Θ is extendable and $n = \max\{r, s\}$, then

$$n \leq \chi_g^\Theta \leq |\mathfrak{o}(\Theta)| + n - 1.$$

The game.

- If $\Theta = (\text{Id}, \text{Id}, \text{Id})$, then $\chi_g^\Theta = \chi_g$.

Lemma

Let $G = (V, E)$ be a graph with $|E| \neq \emptyset$. Then,

$$\chi_{g_B}(K_2 \square G) \leq \Delta(K_2 \square G) \leq \Delta(G) + 1.$$

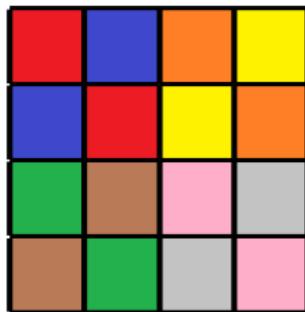
Theorem

- $\chi_{g_A}(\mathcal{R}_{1,n}) = \chi_{g_B}(\mathcal{R}_{1,n}) = n$.
- $\chi_{g_A}(\mathcal{R}_{2,n}) = n + 1$.
- $\chi_{g_B}(\mathcal{R}_{2,n}) = n$, for $n \geq 2$.
- $\chi_{g_A}(\mathcal{R}_{3,3}) = 3$.
- $\chi_{g_B}(\mathcal{R}_{3,3}) = 4$.

The game.

Theorem

- $\chi_{g_A}^{((12)(34),(12)(34),\text{Id})}(\mathcal{R}_{4,4}) = 5.$
- $\chi_{g_B}^{((12)(34),(12)(34),\text{Id})}(\mathcal{R}_{4,4}) = 4.$



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Many thanks!!

Autotopism stabilized colouring games on rook's graphs

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(Joint work with S. Dominique Andres).
(FernUniversität in Hagen).

Qwara. June 29, 2017.

