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THREE RECURSIVE APPROACHES FOR DECISION PROCESSES WITH A CONVERGING BRANCH SYSTEM

Toshiharu FUJITA

Abstract

In this paper, we consider a decision process model with a converging branch system that is a nonserial transition system. The model is treated by three approaches. Thus we introduce three types of recursive equations by using a dynamic programming technique.

1. Introduction

Nonserial dynamic programming was proposed by Nemhauser [5] and has been widely discussed [1, 2, 3]. Nonserial dynamic systems are classified into the four structures: diverging branch systems, converging branch systems, feedback loop systems, and feedforward loop systems. Herein, we focus on a converging branch system and propose a finite-stage decision process model with a converging branch system. In the model, more than two initial states are given, the states are converged on the process, and finally the process is terminated at a final state.

We formulate the model in Section 2, and in Section 3 we discuss three recursive approaches to the model and introduce three types of recursive equations by using a dynamic programming technique. We give a numerical example in Section 4.

2. Notation and formulation

We introduce a finite-stage decision process model with a converging branch system.

1. X , a nonempty finite set, is the state space. The initial states x_1, x_2, \dots, x_L ($\in X$) are specified at the beginning of the process. The process progresses through states $x_{L+1}, x_{L+2}, \dots, x_{N-1} \in X$ with a converging branch system and is terminated at state $x_N \in X$.
2. U , a nonempty finite set, is the decision space. u_n ($\in U$) represents the selected action for state x_n , $n = 1, 2, \dots, N - 1$. We denote the power set of U by 2^U :

$$2^U = \{A : \text{a set} \mid A \subset U\}.$$

Furthermore, we denote by U a point-to-set valued mapping from X to $2^U \setminus \{\emptyset\}$. $U(x)$, called the feasible decision space, represents the set of all feasible decisions in state x .

3. The transition matrix $E = (e_{ij}) \in \{0, 1\}^{N \times N}$ is defined by

$$e_{ij} = \begin{cases} 1 & (\text{if } x_j \text{ is the next state to } x_i) \\ 0 & (\text{otherwise}), \end{cases}$$

and let $I_j = \{i \mid e_{ij} = 1\}$ ($j = L + 1, L + 2, \dots, N$).

Let $G_r(U)$ be the graph of $U(\cdot)$:

$$G_r(U) = \{(x, u) \mid u \in U(x), x \in X\},$$

and, for a set A , $\#A$ means the number of elements in A . When an index set $I = \{m_1, m_2, \dots, m_M\}$ ($m_1 < m_2 < \dots < m_M$) is given, the corresponding sequence

$$(x_{m_1}, u_{m_1}, x_{m_2}, u_{m_2}, \dots, x_{m_M}, u_{m_M})$$

is denoted by $(x_m, u_m \mid m \in I)$. Similarly

$$(x_{m_1}, x_{m_2}, \dots, x_{m_M}) \quad \text{and} \quad (u_{m_1}, u_{m_2}, \dots, u_{m_M})$$

are denoted by $(x_m \mid m \in I)$ and $(u_m \mid m \in I)$, respectively.

4. $r_n : G_r(U) \rightarrow \mathbf{R}$ ($n = 1, 2, \dots, N - 1$) are the reward functions, where $\mathbf{R} = (-\infty, \infty)$. A decision u_n selected in state x_n confers a reward $r_n(x_n, u_n)$. The function $k : X \rightarrow \mathbf{R}$ is the terminal reward function.
5. $f_n : G_r(U)^{\#I_n} \rightarrow X$ ($n = L + 1, L + 2, \dots, N$) are the converging transition laws. If a process in states $(x_m \mid m \in I_n)$ selects actions $(u_m \mid m \in I_n)$, it proceeds deterministically to the next state $f_n(x_m, u_m \mid m \in I_n)$.

Then our model is formulated as follows:

$$\begin{aligned} (\text{P}) \quad & \text{Max } r_1(x_1, u_1) + r_2(x_2, u_2) + \dots + r_{N-1}(x_{N-1}, u_{N-1}) + k(x_N) \\ & \text{s.t. } x_n = f_n(x_m, u_m \mid m \in I_n) \quad n = L + 1, L + 2, \dots, N \\ & u_n \in U(x_n) \quad n = 1, 2, \dots, N - 1. \end{aligned}$$

EXAMPLE 2.1. Let $N = 8$, $L = 3$, and $e_{14} = e_{25} = e_{37} = e_{46} = e_{57} = e_{68} = e_{78} = 1$ ($e_{ij} = 0$ for the other pairs (i, j)). Then, for the given initial states x_1, x_2, x_3 , the other states x_4, x_5, x_6, x_7, x_8 are determined by

$$\begin{aligned} x_4 &= f_4(x_1, u_1), & x_1 &\in X, u_1 \in U(x_1) \\ x_5 &= f_5(x_2, u_2), & x_2 &\in X, u_2 \in U(x_2) \\ x_6 &= f_6(x_4, u_4), & x_4 &\in X, u_4 \in U(x_4) \\ x_7 &= f_7(x_3, u_3, x_5, u_5), & x_3, x_5 &\in X, u_3 \in U(x_3), u_5 \in U(x_5) \\ x_8 &= f_8(x_6, u_6, x_7, u_7), & x_6, x_7 &\in X, u_6 \in U(x_6), u_7 \in U(x_7) \end{aligned}$$

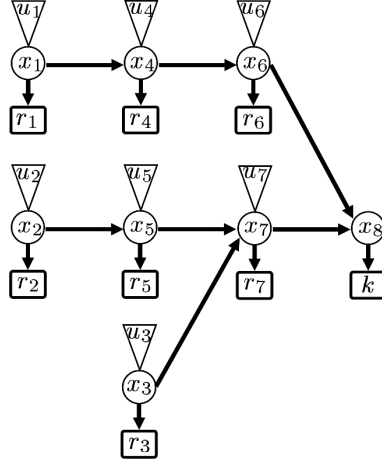


Figure 1. State transition tree for Example 2.1

(see Figure 1). In this case,

$$I_4 = \{1\}, \quad I_5 = \{2\}, \quad I_6 = \{4\}, \quad I_7 = \{3, 5\}, \quad I_8 = \{6, 7\}$$

and the problem is described as follows:

$$\text{Max } r_1(x_1, u_1) + r_2(x_2, u_2) + \cdots + r_7(x_7, u_7) + k(x_8)$$

$$\text{s.t. } x_n = f_n(x_m, u_m \mid m \in I_n) \quad n = 4, 5, \dots, 8$$

$$u_n \in U(x_n) \quad n = 1, 2, \dots, 7. \quad \square$$

3. Three recursive methods

3.1. Backward recursive equation I

We now give the first recursive method for the problem (P). First, P_n ($n = 1, 2, \dots$) denotes a set of indexes of x_m satisfying the condition that the distance between x_m and the final state x_N in the state transition tree for (P) equals n . Especially, $P_0 = \{N\}$. We consider the following subproblems with the initial states $(x_m \mid m \in P_n)$ and the optimal value functions are denoted by $V^n(\cdot)$:

$$V^0(x_N) = k(x_N), \quad x_N \in X$$

$$V^n(x_m \mid m \in P_n) = \max_{\substack{u_m \in U(x_m) \\ (m \in \bigcup_{l=1}^n P_l)}} \left[\sum_{m \in \bigcup_{l=1}^n P_l} r_m(x_m, u_m) + k(x_N) \right],$$

$$(x_m \mid m \in P_n) \in X^{\#P_n}, \quad n = 1, 2, \dots$$

We note that $x_h \in P_l$ ($l < n$) that appears in above objective function is determined consecutively by the initial states ($x_m | m \in P_n$) and decisions ($u_m | m \in \bigcup_{l=1}^n P_l$) through the transition law f_h as follows

$$x_h = f_h(x_m, u_m | m \in I_h),$$

because, for $l = 0, 1, \dots, n-1$, $h \in P_l$ implies $I_h \subset P_{l+1}$.

THEOREM 3.1. *The following recursive equations hold:*

$$V^0(x_N) = k(x_N), \quad x_N \in X$$

$$V^n(x_m | m \in P_n) = \max_{\substack{u_m \in U(x_m) \\ (m \in P_n)}} \left[\sum_{m \in P_n} r_m(x_m, u_m) + V^{n-1}(x_m | m \in P_{n-1}) \right]$$

$$(x_m | m \in P_n) \in X^{\#P_n}, \quad n = 1, 2, \dots,$$

where, in the second term of the objective function, if $1 \leq m \leq L$, x_m is given (x_m is the initial state). Otherwise, x_m is given by $x_m = f_m(x_l, u_l | l \in I_m)$.

PROOF. By the definition of the subproblems,

$$V^n(x_m | m \in P_n) = \max_{\substack{u_m \in U(x_m) \\ (m \in \bigcup_{l=1}^n P_l)}} \left[\sum_{m \in \bigcup_{l=1}^n P_l} r_m(x_m, u_m) + k(x_N) \right], \quad n = 1, 2, \dots$$

Since

$$P_n \cap P_m = \phi \quad (n \neq m),$$

we have

$$P_n \cap \bigcup_{l=1}^{n-1} P_l = \phi.$$

Therefore

$$\begin{aligned} & V^n(x_m | m \in P_n) \\ &= \max_{\substack{u_m \in U(x_m) \\ (m \in P_n)}} \left[\max_{\substack{u_m \in U(x_m) \\ (m \in \bigcup_{l=1}^{n-1} P_l)}} \left[\sum_{m \in P_n} r_m(x_m, u_m) + \sum_{m \in \bigcup_{l=1}^{n-1} P_l} r_m(x_m, u_m) + k(x_N) \right] \right] \\ &= \max_{\substack{u_m \in U(x_m) \\ (m \in P_n)}} \left[\sum_{m \in P_n} r_m(x_m, u_m) + \max_{\substack{u_m \in U(x_m) \\ (m \in \bigcup_{l=1}^{n-1} P_l)}} \left[\sum_{m \in \bigcup_{l=1}^{n-1} P_l} r_m(x_m, u_m) + k(x_N) \right] \right] \\ &= \max_{\substack{u_m \in U(x_m) \\ (m \in P_n)}} \left[\sum_{m \in P_n} r_m(x_m, u_m) + V^{n-1}(x_m | m \in P_{n-1}) \right]. \end{aligned} \quad \square$$

EXAMPLE 3.1. *We consider the problem given by Example 2.1. Then*

$$P_0 = \{8\}, \quad P_1 = \{6, 7\}, \quad P_2 = \{3, 4, 5\}, \quad P_3 = \{1, 2\},$$

and the subproblems become

$$\begin{aligned} V^0(x_8) &= k(x_8) \\ V^1(x_6, x_7) &= \max_{u_m \in U(x_m)} \max_{(m=6,7)} [r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)] \\ V^2(x_3, x_4, x_5) &= \max_{u_m \in U(x_m)} \max_{(3 \leq m \leq 7)} [r_3(x_3, u_3) + r_4(x_4, u_4) + r_5(x_5, u_5) \\ &\quad + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)] \\ V^3(x_1, x_2) &= \max_{u_m \in U(x_m)} \max_{(1 \leq m \leq 7)} [r_1(x_1, u_1) + r_2(x_2, u_2) + r_3(x_3, u_3) + r_4(x_4, u_4) \\ &\quad + r_5(x_5, u_5) + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)]. \end{aligned}$$

By Theorem 3.1, we get the following backward recursive equations:

$$\begin{aligned} V^0(x_8) &= k(x_8) \\ V^1(x_6, x_7) &= \max_{u_m \in U(x_m)} \max_{(m=6,7)} [r_6(x_6, u_6) + r_7(x_7, u_7) + V^0(f_8(x_6, u_6, x_7, u_7))] \\ V^2(x_3, x_4, x_5) &= \max_{u_m \in U(x_m)} \max_{(m=3,4,5)} [r_3(x_3, u_3) + r_4(x_4, u_4) + r_5(x_5, u_5) \\ &\quad + V^1(f_6(x_4, u_4), f_7(x_3, u_3, x_5, u_5))] \\ V^3(x_1, x_2) &= \max_{u_m \in U(x_m)} \max_{(m=1,2)} [r_1(x_1, u_1) + r_2(x_2, u_2) + V^2(x_3, f_4(x_1, u_1), f_5(x_2, u_2))]. \quad \square \end{aligned}$$

3.2. Forward recursive equation

Next, we introduce a forward recursive method for the problem (P). Starting with each initial state x_1, x_2, \dots, x_L , subproblems are generated consecutively in a forward direction and the optimal values functions are denoted by $W^n(\cdot)$. For more details, see the following procedure.

Step 1. Let

$$I = \{1, 2, \dots, L\}$$

and

$$W^n(x_n) = 0, \quad J_n = \phi, \quad n = 1, 2, \dots, L.$$

Step 2. For each $n \notin I$ satisfying $I_n \subset I$, let

$$J_n = I_n \cup \left(\bigcup_{m \in I_n} J_m \right)$$

and

$$\bar{J}_n = (J_n \setminus \{1, 2, \dots, L\}) \cup \{n\}.$$

If $n < N$, define the subproblem terminated with x_n as follows:

$$W^n(x_n) = \max_{\substack{(x_m, u_m | m \in J_n); \\ f_m(x_l, u_l | l \in I_m) = x_m \quad (m \in \bar{J}_n)}} \left[\sum_{m \in J_n} r_m(x_m, u_m) \right], \quad x_n \in X.$$

Otherwise (i.e. if $n = N$),

$$W^N(x_N) = \max_{\substack{(x_m, u_m | m \in J_N); \\ f_m(x_l, u_l | l \in I_m) = x_m \quad (m \in \bar{J}_N)}} \left[\sum_{m \in J_N} r_m(x_m, u_m) + k(x_N) \right], \quad x_N \in X.$$

According to the final state x_n , the corresponding subproblem may have no feasible solution. In that case, the value of $W^n(x_n)$ is regarded as “ $-\infty$ ”.

Step 3. if I equals $\{1, 2, \dots, N-1\}$, stop. Otherwise, update I as

$$I \leftarrow I \cup \{n | I_n \subset I\}.$$

Then, go to Step 2. □

We note that, in Step 2,

$$J_N = \{1, 2, \dots, N-1\}, \quad \bar{J}_N = \{L+1, L+2, \dots, N\}$$

hold. Therefore $\max_{x_N \in X} W^N(x_N)$ gives the optimal value of the original problem (P).

THEOREM 3.2. *The value functions $W^n(\cdot)$ satisfy the following forward recursive equations:*

$$W^n(x_n) = 0, \quad x_n \in X, \quad n = 1, 2, \dots, L$$

$$W^n(x_n) = \max_{\substack{(x_m, u_m | m \in I_n); \\ f_n(x_m, u_m | m \in I_n) = x_n}} \sum_{m \in I_n} (W^m(x_m) + r_m(x_m, u_m))$$

$$x_n \in X, \quad n = L+1, L+2, \dots, N-1$$

$$W^N(x_N) = \max_{\substack{(x_m, u_m | m \in I_N); \\ f_n(x_m, u_m | m \in I_N) = x_N}} \left[\sum_{m \in I_N} (W^m(x_m) + r_m(x_m, u_m)) + k(x_N) \right], \quad x_N \in X.$$

PROOF. If $L < n < N$, by the definition of the subproblems,

$$W^n(x_n) = \max_{\substack{(x_m, u_m | m \in J_n); \\ f_m(x_l, u_l | l \in I_m) = x_m \ (m \in \bar{J}_n)}} \left[\sum_{m \in J_n} r_m(x_m, u_m) \right],$$

and, by the definition of J_n , we have

$$J_n = I_n \cup \left(\bigcup_{m \in I_n} J_m \right) \quad \text{and} \quad I_n \cap \left(\bigcup_{m \in I_n} J_m \right) = \phi.$$

Then,

$$\begin{aligned} \bar{J}_n &= (J_n \setminus \{1, 2, \dots, L\}) \cup \{n\} \\ &= \left(\left(I_n \cup \left(\bigcup_{m \in I_n} J_m \right) \right) \setminus \{1, 2, \dots, L\} \right) \cup \{n\} \\ &= I_n \cup \left(\left(\bigcup_{m \in I_n} J_m \right) \setminus \{1, 2, \dots, L\} \right) \cup \{n\} \\ &= I_n \cup \left(\bigcup_{m \in I_n} (J_m \setminus \{1, 2, \dots, L\}) \right) \cup \{n\} \\ &= \left(\bigcup_{m \in I_n} ((J_m \setminus \{1, 2, \dots, L\}) \cup \{m\}) \right) \cup \{n\} \\ &= \left(\bigcup_{m \in I_n} \bar{J}_m \right) \cup \{n\}. \end{aligned}$$

Hence, because

$$n \notin \left(\bigcup_{m \in I_n} \bar{J}_m \right),$$

we have

$$"m \in J_n \Leftrightarrow m \in I_n \text{ or } m \in \bigcup_{m \in I_n} J_m", \quad I_n \cap \left(\bigcup_{m \in I_n} J_m \right) = \phi.$$

These facts imply that

$$"m \in \bar{J}_n \Leftrightarrow m = n \text{ or } m \in \bigcup_{m \in I_n} \bar{J}_m", \quad n \notin \left(\bigcup_{m \in I_n} \bar{J}_m \right).$$

hold. Therefore

$$\begin{aligned}
W^n(x_n) &= \max_{\substack{(x_m, u_m | m \in I_n); \\ f_n(x_m, u_m | m \in I_n) = x_n}} \left[\max_{\substack{(x_m, u_m | m \in \bigcup_{l \in I_n} J_l); \\ f_m(x_l, u_l | l \in I_m) = x_m \quad (m \in \bigcup_{l \in I_n} \bar{J}_l)}} \left[\sum_{m \in I_n} r_m(x_m, u_m) + \sum_{m \in \bigcup_{l \in I_n} J_l} r_m(x_m, u_m) \right] \right] \\
&= \max_{\substack{(x_m, u_m | m \in I_n); \\ f_n(x_m, u_m | m \in I_n) = x_n}} \left[\sum_{m \in I_n} r_m(x_m, u_m) + \max_{\substack{(x_m, u_m | m \in \bigcup_{l \in I_n} J_l); \\ f_m(x_l, u_l | l \in I_m) = x_m \quad (m \in \bigcup_{l \in I_n} \bar{J}_l)}} \left[\sum_{m \in I_n} \sum_{i \in J_m} r_i(x_i, u_i) \right] \right] \\
&= \max_{\substack{(x_m, u_m | m \in I_n); \\ f_n(x_m, u_m | m \in I_n) = x_n}} \left[\sum_{m \in I_n} r_m(x_m, u_m) + \sum_{m \in I_n} \max_{\substack{(x_i, u_i | l \in J_m); \\ f_l(x_i, u_i | i \in I_l) = x_l \quad (l \in \bar{J}_m)}} \left[\sum_{i \in J_m} r_i(x_i, u_i) \right] \right] \\
&= \max_{\substack{(x_m, u_m | m \in I_n); \\ f_n(x_m, u_m | m \in I_n) = x_n}} \left[\sum_{m \in I_n} r_m(x_m, u_m) + \sum_{m \in I_n} W^m(x_m) \right] \\
&= \max_{\substack{(x_m, u_m | m \in I_n); \\ f_n(x_m, u_m | m \in I_n) = x_n}} \left[\sum_{m \in I_n} (r_m(x_m, u_m) + W^m(x_m)) \right].
\end{aligned}$$

When $n = N$, by the definition of the subproblems,

$$\begin{aligned}
W^N(x_N) &= \max_{\substack{(x_m, u_m | m \in I_N); \\ f_N(x_m, u_m | m \in I_N) = x_N}} \left[\max_{\substack{(x_m, u_m | m \in \bigcup_{l \in I_N} J_l); \\ f_m(x_l, u_l | l \in I_m) = x_m \quad (m \in \bigcup_{l \in I_N} \bar{J}_l)}} \left[\sum_{m \in I_N} r_m(x_m, u_m) \right. \right. \\
&\quad \left. \left. + \sum_{m \in \bigcup_{l \in I_N} J_l} r_m(x_m, u_m) + k(x_N) \right] \right].
\end{aligned}$$

Similarly, we can show the following equality:

$$W^n(x_n) = \max_{\substack{(x_m, u_m | m \in I_n); \\ f_n(x_m, u_m | m \in I_n) = x_n}} \left[\sum_{m \in I_n} (r_m(x_m, u_m) + W^m(x_m)) + k(x_n) \right]. \quad \square$$

EXAMPLE 3.2. We consider again the problem given by Example 2.1. First, for the initial states x_1, x_2, x_3 , let

$$W^1(x_1) = W^2(x_2) = W^3(x_3) = 0,$$

and the subproblems become as follows:

$$\begin{aligned}
W^4(x_4) &= \max_{(x_1, u_1); f_4(x_1, u_1)=x_4} [r_1(x_1, u_1)] \\
W^5(x_5) &= \max_{(x_2, u_2); f_5(x_2, u_2)=x_5} [r_2(x_2, u_2)] \\
W^6(x_6) &= \max_{(x_1, u_1, x_4, u_4); f_4(x_1, u_1)=x_4, f_6(x_4, u_4)=x_6} [r_1(x_1, u_1) + r_4(x_4, u_4)] \\
W^7(x_7) &= \max_{\substack{(x_m, u_m \mid m \in \{2, 3, 5\}); \\ f_5(x_2, u_2)=x_5, f_7(x_3, u_3, x_5, u_5)=x_7}} [r_2(x_2, u_2) + r_3(x_3, u_3) + r_5(x_5, u_5)] \\
W^8(x_8) &= \max_{\substack{(x_m, u_m \mid m \in \{1, 2, \dots, 7\}; f_4(x_1, u_1)=x_4, f_5(x_2, u_2)=x_5, \\ f_6(x_4, u_4)=x_6, f_7(x_3, u_3, x_5, u_5)=x_7, f_8(x_6, u_6, x_7, u_7)=x_8}} [r_1(x_1, u_1) + r_2(x_2, u_2) \\
&\quad + r_3(x_3, u_3) + r_4(x_4, u_4) + r_5(x_5, u_5) + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)].
\end{aligned}$$

Then, by using Theorem 3.2, the following forward recursive equations hold:

$$\begin{aligned}
W^1(x_1) &= W^2(x_2) = W^3(x_3) = 0 \\
W^4(x_4) &= \max_{(x_1, u_1); f_4(x_1, u_1)=x_4} [W^1(x_1) + r_1(x_1, u_1)] \\
W^5(x_5) &= \max_{(x_2, u_2); f_5(x_2, u_2)=x_5} [W^2(x_2) + r_2(x_2, u_2)] \\
W^6(x_6) &= \max_{(x_4, u_4); f_6(x_4, u_4)=x_6} [W^4(x_4) + r_4(x_4, u_4)] \\
W^7(x_7) &= \max_{\substack{(x_3, u_3, x_5, u_5); \\ f_7(x_3, u_3, x_5, u_5)=x_7}} [W^3(x_3) + W^5(x_5) + r_3(x_3, u_3) + r_5(x_5, u_5)] \\
W^8(x_8) &= \max_{\substack{(x_6, u_6, x_7, u_7); \\ f_8(x_6, u_6, x_7, u_7)=x_8}} [W^6(x_6) + W^7(x_7) + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)]. \quad \square
\end{aligned}$$

3.3. Backward recursive equation II

The results in this subsection are discussed in [4]. When we construct the subproblems, starting with the initial target state sequence $Q = (x_N)$, we add x_n to Q in the order that coincides one of the node order for a depth-first search for a state transition tree with the root x_N . Without loss of generality, we regard that index order as

$$N \rightarrow N - 1 \rightarrow N - 2 \rightarrow \dots \rightarrow 2 \rightarrow 1,$$

by renumbering the state index.

Specifically, when we consider the decision process in Example 2.1, the indexes are renumbered as shown in Figure 2. Then, the sequence of the target state

sequence becomes

$$(x_8) \rightarrow (x_7, x_8) \rightarrow (x_6, x_7, x_8) \rightarrow \cdots \rightarrow (x_1, x_2, \dots, x_7, x_8)$$

and we consider the corresponding subproblems and the optimal value functions $v^n(\cdot)$ as follows:

$$\begin{aligned} v^8(x_8) &= k(x_8) \\ v^7(x_7; x_3, u_3) &= \max_{u_7 \in U(x_7)} [r_7(x_7, u_7) + k(x_8)] \\ v^6(x_6; x_3, u_3, x_5, u_5) &= \max_{u_m \in U(x_m) \ (m=6,7)} [r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)] \\ v^5(x_5; x_3, u_3) &= \max_{u_m \in U(x_m) \ (5 \leq m \leq 7)} [r_5(x_5, u_5) + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)] \\ v^4(x_4; x_3, u_3) &= \max_{u_m \in U(x_m) \ (4 \leq m \leq 7)} [r_4(x_4, u_4) + r_5(x_5, u_5) + r_6(x_6, u_6) \\ &\quad + r_7(x_7, u_7) + k(x_8)] \\ v^3(x_3) &= \max_{u_m \in U(x_m) \ (3 \leq m \leq 7)} [r_3(x_3, u_3) + r_4(x_4, u_4) + r_5(x_5, u_5) \\ &\quad + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)] \\ v^2(x_2) &= \max_{u_m \in U(x_m) \ (2 \leq m \leq 7)} [r_2(x_2, u_2) + r_3(x_3, u_3) + r_4(x_4, u_4) + r_5(x_5, u_5) \\ &\quad + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)] \\ v^1(x_1) &= \max_{u_m \in U(x_m) \ (1 \leq m \leq 7)} [r_1(x_1, u_1) + r_2(x_2, u_2) + r_3(x_3, u_3) + r_4(x_4, u_4) \\ &\quad + r_5(x_5, u_5) + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)]. \end{aligned}$$

We give the general form of subproblems.

$n = N$

For state sequence (x_N) , the subproblem is given by

$$v^N(x_N) = k(x_N), \quad x_N \in X.$$

$n < N$

For state sequence $(x_n, x_{n+1}, \dots, x_N)$, the subproblem is given by

$$\begin{aligned} &v^n(x_n; (x_m, u_m \mid m \in J_n)) \\ &= \max_{u_m \in U(x_m) \ (m=n, n+1, \dots, N-1)} [r_n(x_n, u_n) + r_{n+1}(x_{n+1}, u_{n+1}) + \cdots + k(x_N)], \\ &x_n \in X, x_m \in X, u_m \in U(x_m) \ (m \in J_n), \end{aligned}$$

where

$$J_n = \bigcup_{l=n+1}^N \{j \in I_l \mid j < n\}.$$

The set of indexes that indicate the initial states is denoted by I_{Init} . For example, we have $I_{\text{Init}} = \{1, 4, 6\}$ for the decision process shown in Figure 2.

Then the following recursive relations are shown in [4].

PROPOSITION 3.1. *Put $J_N = \phi$, then,*

(i) *if $n + 1 \notin I_{\text{Init}}$,*

$$J_n = J_{n+1} \cup \{j \in I_{n+1} \mid j < n\}.$$

(ii) *if $n + 1 \in I_{\text{Init}}$,*

$$J_n = J_{n+1} \setminus \{n\}.$$

THEOREM 3.3. *We have the following backward recursive equations:*

$$v^N(x_N) = k(x_N), \quad x_N \in X$$

$$v^n(x_n; (x_m, u_m \mid m \in J_n))$$

$$= \max_{u_n \in U(x_n)} [r_n(x_n, u_n) + v^{n+1}(f_{n+1}(x_m, u_m \mid m \in I_{n+1}); (x_m, u_m \mid m \in J_{n+1}))],$$

$$x_n \in X, n + 1 \notin I_{\text{Init}}$$

$$v^n(x_n; (x_m, u_m \mid m \in J_n))$$

$$= \max_{u_n \in U(x_n)} [r_n(x_n, u_n) + v^{n+1}(x_{n+1}; (x_m, u_m \mid m \in J_{n+1}))], \quad x_n \in X, n + 1 \in I_{\text{Init}}.$$

EXAMPLE 3.3. We consider the decision process shown in Figure 2. First, by Proposition 3.1, we can get J_j ($j = 8, 7, \dots, 1$) as follows:

$$J_8 = \phi.$$

Since $j = 7 \in I_8 = \{3, 7\}$ and $8 \notin I_{\text{Init}}$,

$$J_7 = J_8 \cup \{j \in I_8 \mid j < 7\} = \phi \cup \{3\} = \{3\}.$$

Similarly,

$$j = 6 \in I_7 = \{5, 6\}, \quad 7 \notin I_{\text{Init}}$$

$$\Rightarrow J_6 = J_7 \cup \{j \in I_7 \mid j < 6\} = \{3\} \cup \{5\} = \{3, 5\},$$

$$j = 5 \in I_7 = \{5, 6\}, \quad 6 \in I_{\text{Init}}$$

$$\Rightarrow J_5 = J_6 \setminus \{5\} = \{3, 5\} \setminus \{5\} = \{3\},$$

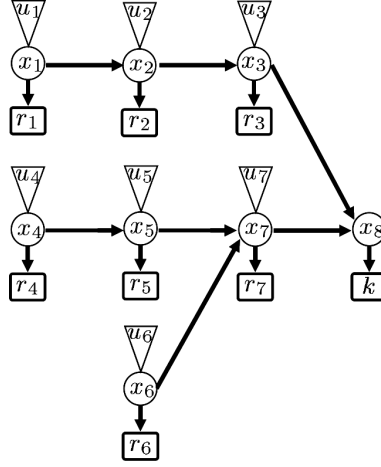


Figure 2. Transition tree of Example 3.3

$$\begin{aligned}
 j = 4 \in I_5 = \{4\}, \quad 5 \notin I_{\text{init}} \\
 \Rightarrow J_4 = J_5 \cup \{j \in I_5 \mid j < 4\} = \{3\} \cup \phi = \{3\}, \\
 j = 3 \in I_8 = \{3, 7\}, \quad 4 \in I_{\text{init}} \\
 \Rightarrow J_3 = J_4 \setminus \{3\} = \{3\} \setminus \{3\} = \phi, \\
 j = 2 \in I_3 = \{2\}, \quad 3 \notin I_{\text{init}} \\
 \Rightarrow J_2 = J_3 \cup \{j \in I_3 \mid j < 2\} = \phi \cup \phi = \phi, \\
 j = 1 \in I_2 = \{1\}, \quad 2 \notin I_{\text{init}} \\
 \Rightarrow J_1 = J_2 \cup \{j \in I_2 \mid j < 1\} = \phi \cup \phi = \phi.
 \end{aligned}$$

Then, we have the following recursive equations by Theorem 3.1:

$$\begin{aligned}
 v^8(x_8) &= k(x_8) \\
 v^7(x_7; (x_m, u_m \mid m \in J_7)) &= \max_{u_7 \in U(x_7)} [r_7(x_7, u_7) + v^8(f_8(x_m, u_m \mid m \in I_8); (x_m, u_m \mid m \in J_8))] \\
 v^7(x_7; (x_m, u_m \mid m \in \{3\})) &= \max_{u_7 \in U(x_7)} [r_7(x_7, u_7) + v^8(f_8(x_m, u_m \mid m \in \{3, 7\}); (x_m, u_m \mid m \in \phi))] \\
 v^7(x_7; x_3, u_3) &= \max_{u_7 \in U(x_7)} [r_7(x_7, u_7) + v^8(f_8(x_3, u_3, x_7, u_7))]
 \end{aligned}$$

and

$$\begin{aligned}
 v^6(x_6; (x_m, u_m \mid m \in J_6)) &= \max_{u_6 \in U(x_6)} [r_6(x_6, u_6) + v^7(f_7(x_m, u_m \mid m \in I_7); (x_m, u_m \mid m \in J_7))] \\
 v^6(x_6; (x_m, u_m \mid m \in \{3, 5\})) &= \max_{u_6 \in U(x_6)} [r_6(x_6, u_6) + v^7(f_7(x_m, u_m \mid m \in \{5, 6\}); \\
 &\quad (x_m, u_m \mid m \in \{3\})))] \\
 v^6(x_6; x_3, u_3, x_5, u_5) &= \max_{u_6 \in U(x_6)} [r_6(x_6, u_6) + v^7(f_7(x_5, u_5, x_6, u_6); x_3, u_3)].
 \end{aligned}$$

Similarly, we have

$$\begin{aligned}
 v^5(x_5; x_3, u_3) &= \max_{u_5 \in U(x_5)} [r_5(x_5, u_5) + v^6(x_6; x_3, u_3, x_5, u_5)] \\
 v^4(x_4; x_3, u_3) &= \max_{u_4 \in U(x_4)} [r_4(x_4, u_4) + v^5(f_5(x_4, u_4); x_3, u_3)] \\
 v^3(x_3) &= \max_{u_3 \in U(x_3)} [r_3(x_3, u_3) + v^4(x_4; x_3, u_3)] \\
 v^2(x_2) &= \max_{u_2 \in U(x_2)} [r_2(x_2, u_2) + v^3(f_3(x_2, u_2))] \\
 v^1(x_1) &= \max_{u_1 \in U(x_1)} [r_1(x_1, u_1) + v^2(f_2(x_1, u_1))]. \quad \square
 \end{aligned}$$

4. Numerical example

We consider the decision process problem given by Example 2.1 with the following data:

$$\begin{aligned}
 X &= \{s_1, s_2\}, & U(x) &= U = \{a_1, a_2\} \quad (\forall x \in X) \\
 x_1 &= s_1, & x_2 &= s_2, & x_3 &= s_1
 \end{aligned}$$

$\underline{f_4(x, u)}$	$x \backslash u$	a_1	a_2	$\underline{f_5(x, u)}$	$x \backslash u$	a_1	a_2	$\underline{f_6(x, u)}$	$x \backslash u$	a_1	a_2
	s_1	s_2	s_1		s_1	—	—		s_1	s_2	s_1
	s_2	—	—		s_2	s_1	s_2		s_2	s_1	s_2

$\underline{f_7(x, u, y, v)}$	$(x, u) \backslash (y, v)$	(s_1, a_1)	(s_1, a_2)	(s_2, a_1)	(s_2, a_2)
	(s_1, a_1)	s_2	s_2	s_1	s_2
	(s_1, a_2)	s_1	s_2	s_2	s_1

$\underline{f_8(x, u, y, v)}$	$(x, u) \backslash (y, v)$	(s_1, a_1)	(s_1, a_2)	(s_2, a_1)	(s_2, a_2)
	(s_1, a_1)	s_1	s_1	s_2	s_2
	(s_1, a_2)	s_2	s_2	s_1	s_2
	(s_2, a_1)	s_1	s_2	s_2	s_1
	(s_2, a_2)	s_1	s_2	s_1	s_1

and

(x, u)	(s_1, a_1)	(s_1, a_2)	(s_2, a_1)	(s_2, a_2)	
$r_1(x, u)$	4	3	—	—	
$r_2(x, u)$	—	—	3	2	
$r_3(x, u)$	3	4	—	—	$\frac{x}{k(x)} \mid \begin{array}{cc} s_1 & s_2 \\ 5 & 2 \end{array}$
$r_4(x, u)$	2	3	3	1	
$r_5(x, u)$	1	2	4	2	
$r_6(x, u)$	2	3	3	4	
$r_7(x, u)$	4	5	2	3	

4.1. Compute with backward recursive equation I

First, for the terminal state x_8 , we get

$$V^0(s_1) = k(s_1) = 5, \quad V^0(s_2) = k(s_2) = 2.$$

By using the backward recursive equation given by Theorem 3.1, we compute $V^1(s_1, s_1)$ and the corresponding optimal decision function $\sigma_1^*(s_1, s_1)$:

$$\begin{aligned} V^1(s_1, s_1) &= \max_{u_6, u_7 \in U} [r_6(s_1, u_6) + r_7(s_1, u_7) + V^0(f_8(s_1, u_6, s_1, u_7))] \\ &= [r_6(s_1, a_1) + r_7(s_1, a_1) + V^0(f_8(s_1, a_1, s_1, a_1))] \\ &\quad \vee [r_6(s_1, a_1) + r_7(s_1, a_2) + V^0(f_8(s_1, a_1, s_1, a_2))] \\ &\quad \vee [r_6(s_1, a_2) + r_7(s_1, a_1) + V^0(f_8(s_1, a_2, s_1, a_1))] \\ &\quad \vee [r_6(s_1, a_2) + r_7(s_1, a_2) + V^0(f_8(s_1, a_2, s_1, a_2))] \\ &= [2 + 4 + V^0(s_1)] \vee [2 + 5 + V^0(s_1)] \vee [3 + 4 + V^0(s_2)] \vee [3 + 5 + V^0(s_2)] \\ &= 11 \vee 12 \vee 9 \vee 10 = 12, \quad \sigma_1^*(s_1, s_1) = (u_6^*, u_7^*) = (a_1, a_2). \end{aligned}$$

Similarly, we have

$$\begin{aligned} V^1(s_1, s_2) &= 6 \vee 7 \vee 10 \vee 8 = 10, \quad \sigma_1^*(s_1, s_2) = (u_6^*, u_7^*) = (a_2, a_1) \\ V^1(s_2, s_1) &= 12 \vee 10 \vee 13 \vee 11 = 13, \quad \sigma_1^*(s_2, s_1) = (u_6^*, u_7^*) = (a_2, a_1) \\ V^1(s_2, s_2) &= 7 \vee 11 \vee 11 \vee 12 = 12, \quad \sigma_1^*(s_2, s_2) = (u_6^*, u_7^*) = (a_2, a_2). \end{aligned}$$

Next, we compute V^2 and the corresponding optimal decision function σ_2^* :

$$\begin{aligned} V^2(s_1, s_1, s_1) &= \max_{u_3, u_4, u_5 \in U} [r_3(s_1, u_3) + r_4(s_1, u_4) + r_5(s_1, u_5) + V^1(f_6(s_1, u_4), f_7(s_1, u_3, s_1, u_5))] \end{aligned}$$

$$\begin{aligned}
&= [r_3(s_1, a_1) + r_4(s_1, a_1) + r_5(s_1, a_1) + V^1(f_6(s_1, a_1), f_7(s_1, a_1, s_1, a_1))] \\
&\quad \vee [r_3(s_1, a_2) + r_4(s_1, a_1) + r_5(s_1, a_1) + V^1(f_6(s_1, a_1), f_7(s_1, a_2, s_1, a_1))] \\
&\quad \vee [r_3(s_1, a_1) + r_4(s_1, a_1) + r_5(s_1, a_2) + V^1(f_6(s_1, a_1), f_7(s_1, a_1, s_1, a_2))] \\
&\quad \vee [r_3(s_1, a_2) + r_4(s_1, a_1) + r_5(s_1, a_2) + V^1(f_6(s_1, a_1), f_7(s_1, a_2, s_1, a_2))] \\
&\quad \vee [r_3(s_1, a_1) + r_4(s_1, a_2) + r_5(s_1, a_1) + V^1(f_6(s_1, a_2), f_7(s_1, a_1, s_1, a_1))] \\
&\quad \vee [r_3(s_1, a_2) + r_4(s_1, a_2) + r_5(s_1, a_1) + V^1(f_6(s_1, a_2), f_7(s_1, a_2, s_1, a_1))] \\
&\quad \vee [r_3(s_1, a_1) + r_4(s_1, a_2) + r_5(s_1, a_2) + V^1(f_6(s_1, a_2), f_7(s_1, a_1, s_1, a_2))] \\
&\quad \vee [r_3(s_1, a_2) + r_4(s_1, a_2) + r_5(s_1, a_2) + V^1(f_6(s_1, a_2), f_7(s_1, a_2, s_1, a_2))] \\
&= [3 + 2 + 1 + V^1(s_2, s_2)] \vee [4 + 2 + 1 + V^1(s_2, s_1)] \\
&\quad \vee [3 + 2 + 2 + V^1(s_2, s_2)] \vee [4 + 2 + 2 + V^1(s_2, s_2)] \\
&\quad \vee [3 + 3 + 1 + V^1(s_1, s_2)] \vee [4 + 3 + 1 + V^1(s_1, s_1)] \\
&\quad \vee [3 + 3 + 2 + V^1(s_1, s_2)] \vee [4 + 3 + 2 + V^1(s_1, s_2)] \\
&= 18 \vee 20 \vee 19 \vee 20 \vee 17 \vee 20 \vee 18 \vee 19 = 20
\end{aligned}$$

$$\begin{aligned}
\sigma_2^*(s_1, s_1, s_1) &= (u_3^*, u_4^*, u_5^*) = (a_2, a_1, a_1), (a_2, a_1, a_2), (a_2, a_2, a_1) \\
V^2(s_1, s_1, s_2) &= 22, \quad \sigma_2^*(s_1, s_1, s_2) = (a_1, a_1, a_1), (a_2, a_1, a_1), (a_1, a_2, a_1) \\
V^2(s_1, s_2, s_1) &= 20, \quad \sigma_2^*(s_1, s_2, s_1) = (a_2, a_1, a_1), (a_2, a_2, a_2) \\
V^2(s_1, s_2, s_2) &= 22, \quad \sigma_2^*(s_1, s_2, s_2) = (a_2, a_1, a_1).
\end{aligned}$$

Finally, we compute V^3 and the corresponding optimal decision function σ_3^* :

$$\begin{aligned}
V^3(s_1, s_2) &= \max_{u_1, u_2 \in U} [r_1(s_1, u_1) + r_2(s_2, u_2) + V^2(s_1, f_4(s_1, u_1), f_5(s_2, u_2))] \\
&= [4 + 3 + V^2(s_1, s_2, s_1)] \vee [4 + 2 + V^2(s_1, s_2, s_2)] \\
&\quad \vee [3 + 3 + V^2(s_1, s_1, s_1)] \vee [3 + 2 + V^2(s_1, s_1, s_2)] \\
&= [4 + 3 + 20] \vee [4 + 2 + 22] \vee [3 + 3 + 20] \vee [3 + 2 + 22] \\
&= 27 \vee 28 \vee 26 \vee 27 = 28, \quad \sigma_3^*(s_1, s_2) = (u_1^*, u_2^*) = (a_1, a_2).
\end{aligned}$$

Thus, the optimal value is $V^3(s_1, s_2) = 28$ and an optimal state-decision sequence is given as follows:

$$\begin{aligned}
(x_1, x_2) &= (s_1, s_2) \\
\rightarrow (u_1^*, u_2^*) &= \sigma_3^*(s_1, s_2) = (a_1, a_2) \\
\rightarrow (x_4, x_5, x_3) &= (f_4(s_1, a_1), f_5(s_2, a_2), s_1) = (s_2, s_2, s_1) \\
\rightarrow (u_4^*, u_5^*, u_3^*) &= \sigma_2^*(s_2, s_2, s_1) = (a_1, a_1, a_1) \\
\rightarrow (x_6, x_7) &= (f_6(s_2, a_1), f_7(s_1, a_1, s_2, a_1)) = (s_1, s_1) \\
\rightarrow (u_6^*, u_7^*) &= \sigma_1^*(s_1, s_1) = (a_1, a_2) \\
\rightarrow x_8 &= f_8(s_1, a_1, s_1, a_2) = s_1.
\end{aligned}$$

4.2. Compute with forward recursive equation

First, for the initial states $x_1 = s_1$, $x_2 = s_2$, $x_3 = s_1$, we get

$$W^1(s_1) = W^2(s_2) = W^3(s_1) = 0$$

By using the forward recursive equation in subsection 3.2, we compute $W^4(s_1)$ and the corresponding optimal decision function $\tau_4^*(s_1)$:

$$\begin{aligned}
W^4(s_1) &= \max_{\substack{(x_1, u_1); \\ f_4(x_1, u_1) = s_1}} [W^1(x_1) + r_1(x_1, u_1)] = \max_{(x_1, u_1) \in \{(s_1, a_2)\}} [W^1(x_1) + r_1(x_1, u_1)] \\
&= W^1(s_1) + r_1(s_1, a_2) = 0 + 3 = 3, \quad \tau_4^*(s_1) = (x_1, u_1^*) = (s_1, a_2).
\end{aligned}$$

Similarly,

$$\begin{aligned}
W^4(s_2) &= 4, \quad \tau_4^*(s_2) = (x_1, u_1^*) = (s_1, a_1) \\
W^5(s_1) &= 3, \quad \tau_5^*(s_1) = (x_2, u_2^*) = (s_2, a_1) \\
W^5(s_2) &= 2, \quad \tau_5^*(s_2) = (x_2, u_2^*) = (s_2, a_2)
\end{aligned}$$

and

$$\begin{aligned}
W^6(s_1) &= \max_{(x_4, u_4); f_6(x_4, u_4) = s_1} [W^4(x_4) + r_4(x_4, u_4)] \\
&= \max_{(x_4, u_4) \in \{(s_1, a_2), (s_2, a_1)\}} [W^4(x_4) + r_4(x_4, u_4)] \\
&= \max[W^4(s_1) + r_4(s_1, a_2), W^4(s_2) + r_4(s_2, a_1)] \\
&= \max[3 + 3, 4 + 3] = 7, \quad \tau_6^*(s_1) = (x_4, u_4^*) = (s_2, a_1) \\
W^6(s_2) &= 5, \quad \tau_6^*(s_2) = (x_4, u_4^*) = (s_1, a_1), (s_2, a_2).
\end{aligned}$$

Next, because

$$\begin{aligned} & \{(x_3, u_3, x_5, u_5) \mid f_7(x_3, u_3, x_5, u_5) = s_1\} \\ &= \{(s_1, a_1, s_2, a_1), (s_1, a_2, s_1, a_1), (s_1, a_2, s_2, a_2)\} \\ & \{(x_3, u_3, x_5, u_5) \mid f_7(x_3, u_3, x_5, u_5) = s_2\} \\ &= \{(s_1, a_1, s_1, a_1), (s_1, a_1, s_1, a_2), (s_1, a_1, s_2, a_2), (s_1, a_2, s_1, a_2), (s_1, a_2, s_2, a_1)\}, \end{aligned}$$

we have

$$\begin{aligned} W^7(s_1) &= \max_{(x_3, u_3, x_5, u_5); f_7(x_3, u_3, x_5, u_5)=s_1} [W^3(x_3) + W^5(x_5) + r_3(x_3, u_3) + r_5(x_5, u_5)] \\ &= \max[W^3(s_1) + W^5(s_2) + r_3(s_1, a_1) + r_5(s_2, a_1), \\ & \quad W^3(s_1) + W^5(s_1) + r_3(s_1, a_2) + r_5(s_1, a_1), \\ & \quad W^3(s_1) + W^5(s_2) + r_3(s_1, a_2) + r_5(s_2, a_2)] \\ &= \max[0 + 2 + 3 + 4, 0 + 3 + 4 + 1, 0 + 2 + 4 + 2] = \max[9, 8, 8] \\ &= 9, \quad \tau_7^*(s_1) = (x_3, u_3^*, x_5, u_5^*) = (s_1, a_1, s_2, a_1) \\ W^7(s_2) &= 10, \quad \tau_7^*(s_2) = (x_3, u_3^*, x_5, u_5^*) = (s_1, a_2, s_2, a_1). \end{aligned}$$

Finally, because

$$\begin{aligned} & \{(x_6, u_6, x_7, u_7) \mid f_8(x_6, u_6, x_7, u_7) = s_1\} \\ &= \{(s_1, a_1, s_1, a_1), (s_1, a_1, s_1, a_2), (s_1, a_2, s_2, a_1), (s_2, a_1, s_1, a_1), \\ & \quad (s_2, a_1, s_2, a_2), (s_2, a_2, s_1, a_1), (s_2, a_2, s_2, a_1), (s_2, a_2, s_2, a_2)\} \\ & \{(x_6, u_6, x_7, u_7) \mid f_8(x_6, u_6, x_7, u_7) = s_2\} \\ &= \{(s_1, a_1, s_2, a_1), (s_1, a_1, s_2, a_2), (s_1, a_2, s_1, a_1), (s_1, a_2, s_1, a_2), \\ & \quad (s_1, a_2, s_2, a_2), (s_2, a_1, s_1, a_2), (s_2, a_1, s_2, a_1), (s_2, a_2, s_1, a_2)\}, \end{aligned}$$

we have

$$\begin{aligned} W^8(s_1) &= \max_{(x_6, u_6, x_7, u_7); f_8(x_6, u_6, x_7, u_7)=s_1} [W^6(x_6) + W^7(x_7) + r_6(x_6, u_6) + r_7(x_7, u_7) + k(x_8)] \\ &= \max[W^6(s_1) + W^7(s_1) + r_6(s_1, a_1) + r_7(s_1, a_1) + k(s_1), \\ & \quad W^6(s_1) + W^7(s_1) + r_6(s_1, a_1) + r_7(s_1, a_2) + k(s_1), \\ & \quad W^6(s_1) + W^7(s_2) + r_6(s_1, a_2) + r_7(s_2, a_1) + k(s_1), \end{aligned}$$

$$\begin{aligned}
& W^6(s_2) + W^7(s_1) + r_6(s_2, a_1) + r_7(s_1, a_1) + k(s_1), \\
& W^6(s_2) + W^7(s_2) + r_6(s_2, a_1) + r_7(s_2, a_2) + k(s_1), \\
& W^6(s_2) + W^7(s_1) + r_6(s_2, a_2) + r_7(s_1, a_1) + k(s_1), \\
& W^6(s_2) + W^7(s_2) + r_6(s_2, a_2) + r_7(s_2, a_1) + k(s_1), \\
& W^6(s_2) + W^7(s_2) + r_6(s_2, a_2) + r_7(s_2, a_2) + k(s_1)] \\
= & \max[7 + 9 + 2 + 4 + 5, 7 + 9 + 2 + 5 + 5, 7 + 10 + 3 + 2 + 5, \\
& 5 + 9 + 3 + 4 + 5, 5 + 10 + 3 + 3 + 5, 5 + 9 + 4 + 4 + 5, \\
& 5 + 10 + 4 + 2 + 5, 5 + 10 + 4 + 3 + 5] \\
= & \max[27, 28, 27, 26, 26, 27, 26, 27] = 28 \\
\tau_8^*(s_1) = & (x_6, u_6^*, x_7, u_7^*) = (s_1, a_1, s_1, a_2) \\
W^8(s_2) = & 26, \quad \tau_8^*(s_2) = (x_6, u_6^*, x_7, u_7^*) = (s_1, a_2, s_1, a_2).
\end{aligned}$$

Thus, the optimal value is

$$\max[W^8(s_1), W^8(s_2)] = \max[28, 26] = 28, \quad x_8 = s_1$$

and we get an optimal state-decision sequence as follows:

$$\begin{aligned}
x_8 = s_1 & \Rightarrow (x_6, u_6^*, x_7, u_7^*) = \tau_8^*(s_1) = (s_1, a_1, s_1, a_2) \\
& \rightarrow x_6 = s_1, \quad u_6^* = a_1 \\
& \Rightarrow (x_4, u_4^*) = \tau_6^*(s_1) = (s_2, a_1) \\
& \rightarrow x_4 = s_2, \quad u_4^* = a_1 \\
& \Rightarrow (x_1, u_1^*) = \tau_4^*(s_2) = (s_1, a_1) \\
& \rightarrow x_1 = s_1, \quad u_1^* = a_1 \\
& \rightarrow x_7 = s_1, \quad u_7^* = a_2 \\
& \Rightarrow (x_3, u_3^*, x_5, u_5^*) = \tau_7^*(s_1) = (s_1, a_1, s_2, a_1) \\
& \rightarrow x_3 = s_1, \quad u_3^* = a_1 \\
& \rightarrow x_5 = s_2, \quad u_5^* = a_1 \\
& \Rightarrow (x_2, u_2^*) = \tau_5^*(s_2) = (s_2, a_2) \\
& \rightarrow x_2 = s_2, \quad u_2^* = a_2.
\end{aligned}$$

4.3. Compute with backward recursive equation II

In Example 3.3, we modified the indexes. But, to solve our numerical example, we need to return that modification back to the original. Then, the backward recursive equations in Example 3.3 become

$$\begin{aligned}
v^8(x_8) &= k(x_8) \\
v^7(x_7; x_6, u_6) &= \max_{u_7 \in U} [r_7(x_7, u_7) + v^8(f_8(x_6, u_6, x_7, u_7))] \\
v^3(x_3; x_5, u_5, x_6, u_6) &= \max_{u_3 \in U} [r_3(x_3, u_3) + v^7(f_7(x_3, u_3, x_5, u_5); x_6, u_6)] \\
v^5(x_5; x_6, u_6) &= \max_{u_5 \in U} [r_5(x_5, u_5) + v^3(x_3; x_5, u_5, x_6, u_6)] \\
v^2(x_2; x_6, u_6) &= \max_{u_2 \in U} [r_2(x_2, u_2) + v^5(f_5(x_2, u_2); x_6, u_6)] \\
v^6(x_6) &= \max_{u_6 \in U} [r_6(x_6, u_6) + v^2(x_2; x_6, u_6)] \\
v^4(x_4) &= \max_{u_4 \in U} [r_4(x_4, u_4) + v^6(f_6(x_4, u_4))] \\
v^1(x_1) &= \max_{u_1 \in U} [r_1(x_1, u_1) + v^4(f_4(x_1, u_1))].
\end{aligned}$$

First, for the terminal state x_8 , we get

$$v^8(s_1) = k(s_1) = 5, \quad v^8(s_2) = k(s_2) = 2.$$

We compute $v_7(s_1; s_1, a_1)$ and the corresponding optimal decision function $\pi_7^*(s_1; s_1, a_1)$:

$$\begin{aligned}
v^7(s_1; s_1, a_1) &= \max_{u_7 \in U} [r_7(s_1, u_7) + v^8(f_8(s_1, a_1, s_1, u_7))] \\
&= [r_7(s_1, a_1) + v^8(f_8(s_1, a_1, s_1, a_1))] \vee [r_7(s_1, a_2) + v^8(f_8(s_1, a_1, s_1, a_2))] \\
&= [4 + v^8(s_1)] \vee [5 + v^8(s_1)] = [4 + 5] \vee [5 + 5] = 10 \\
\pi_7^*(s_1; s_1, a_1) &= a_2.
\end{aligned}$$

Similarly,

$$\begin{aligned}
v^7(s_1; s_1, a_2) &= [4 + 2] \vee [5 + 2] = 7, & \pi_7^*(s_1; s_1, a_2) &= a_2 \\
v^7(s_1; s_2, a_1) &= 9, & \pi_7^*(s_1; s_2, a_1) &= a_1, & v^7(s_1; s_2, a_2) &= 9, & \pi_7^*(s_1; s_2, a_2) &= a_1 \\
v^7(s_2; s_1, a_1) &= 5, & \pi_7^*(s_2; s_1, a_1) &= a_2, & v^7(s_2; s_1, a_2) &= 7, & \pi_7^*(s_2; s_1, a_2) &= a_1 \\
v^7(s_2; s_2, a_1) &= 8, & \pi_7^*(s_2; s_2, a_1) &= a_2, & v^7(s_2; s_2, a_2) &= 8, & \pi_7^*(s_2; s_2, a_2) &= a_2.
\end{aligned}$$

Moreover, we compute $v^3, v^5, v^2, \dots, v^1$ and the corresponding optimal decision functions $\pi_3^*, \pi_4^*, \pi_2^*, \dots, \pi_1^*$ as follows:

$$\begin{aligned}
& v^3(s_1; s_1, a_1, s_1, a_1) \\
&= \max_{u_3 \in U} [r_3(s_1, u_3) + v^7(f_7(s_1, u_3, s_1, a_1); s_1, a_1)] \\
&= [r_3(s_1, a_1) + v^7(f_7(s_1, a_1, s_1, a_1); s_1, a_1)] \vee [r_3(s_1, a_2) + v^7(f_7(s_1, a_2, s_1, a_1); s_1, a_1)] \\
&= [3 + v^7(s_2; s_1, a_1)] \vee [4 + v^7(s_1; s_1, a_1)] = [3 + 5] \vee [4 + 10] = 14
\end{aligned}$$

$$\pi_3^*(s_1; s_1, a_1, s_1, a_1) = a_2$$

$$v^3(s_1; s_1, a_1, s_1, a_2) = 11, \quad \pi_3^*(s_1; s_1, a_1, s_1, a_2) = a_2$$

$$v^3(s_1; s_1, a_1, s_2, a_1) = 13, \quad \pi_3^*(s_1; s_1, a_1, s_2, a_1) = a_2$$

$$v^3(s_1; s_1, a_1, s_2, a_2) = 14, \quad \pi_3^*(s_1; s_1, a_1, s_2, a_2) = a_2$$

$$v^3(s_1; s_1, a_2, s_1, a_1) = 9, \quad \pi_3^*(s_1; s_1, a_2, s_1, a_1) = a_2$$

$$v^3(s_1; s_1, a_2, s_1, a_2) = 11, \quad \pi_3^*(s_1; s_1, a_2, s_1, a_2) = a_2$$

$$v^3(s_1; s_1, a_2, s_2, a_1) = 13, \quad \pi_3^*(s_1; s_1, a_2, s_2, a_1) = a_2$$

$$v^3(s_1; s_1, a_2, s_2, a_2) = 13, \quad \pi_3^*(s_1; s_1, a_2, s_2, a_2) = a_2$$

$$v^3(s_1; s_2, a_1, s_1, a_1) = 13, \quad \pi_3^*(s_1; s_2, a_1, s_1, a_1) = a_1$$

$$v^3(s_1; s_2, a_1, s_1, a_2) = 11, \quad \pi_3^*(s_1; s_2, a_1, s_1, a_2) = a_2$$

$$v^3(s_1; s_2, a_1, s_2, a_1) = 12, \quad \pi_3^*(s_1; s_2, a_1, s_2, a_1) = a_1, a_2$$

$$v^3(s_1; s_2, a_1, s_2, a_2) = 12, \quad \pi_3^*(s_1; s_2, a_1, s_2, a_2) = a_1, a_2$$

$$v^3(s_1; s_2, a_2, s_1, a_1) = 14, \quad \pi_3^*(s_1; s_2, a_2, s_1, a_1) = a_2$$

$$v^3(s_1; s_2, a_2, s_1, a_2) = 11, \quad \pi_3^*(s_1; s_2, a_2, s_1, a_2) = a_2$$

$$v^3(s_1; s_2, a_2, s_2, a_1) = 13, \quad \pi_3^*(s_1; s_2, a_2, s_2, a_1) = a_2$$

$$v^3(s_1; s_2, a_2, s_2, a_2) = 13, \quad \pi_3^*(s_1; s_2, a_2, s_2, a_2) = a_2$$

$$v^5(s_1; s_1, a_1) = 15, \quad \pi_5^*(s_1; s_1, a_1) = a_1, \quad v^5(s_1; s_1, a_2) = 13, \quad \pi_5^*(s_1; s_1, a_2) = a_2$$

$$v^5(s_1; s_2, a_1) = 15, \quad \pi_5^*(s_1; s_2, a_1) = a_2, \quad v^5(s_1; s_2, a_2) = 15, \quad \pi_5^*(s_1; s_2, a_2) = a_1, a_2$$

$$v^5(s_2; s_1, a_1) = 17, \quad \pi_5^*(s_2; s_1, a_1) = a_1, \quad v^5(s_2; s_1, a_2) = 15, \quad \pi_5^*(s_2; s_1, a_2) = a_1$$

$$v^5(s_2; s_2, a_1) = 16, \quad \pi_5^*(s_2; s_2, a_1) = a_1, \quad v^5(s_2; s_2, a_2) = 16, \quad \pi_5^*(s_2; s_2, a_2) = a_1$$

$$\begin{aligned}
 v^2(s_2; s_1, a_1) &= 19, & \pi_2^*(s_2; s_1, a_1) &= a_2, & v^2(s_2; s_1, a_2) &= 17, & \pi_2^*(s_2; s_1, a_2) &= a_2 \\
 v^2(s_2; s_2, a_1) &= 18, & \pi_2^*(s_2; s_2, a_1) &= a_1, a_2, & v^2(s_2; s_2, a_2) &= 18, & \pi_2^*(s_2; s_2, a_2) &= a_1, a_2 \\
 v^6(s_1) &= 21, & \pi_6^*(s_1) &= a_1, & v^6(s_2) &= 22, & \pi_6^*(s_2) &= a_2 \\
 v^4(s_1) &= 24, & \pi_4^*(s_1) &= a_1, a_2, & v^4(s_2) &= 24, & \pi_4^*(s_2) &= a_1
 \end{aligned}$$

and

$$v^1(s_1) = 28, \quad \pi_1^*(s_1) = a_1.$$

Thus, the optimal value is $v^3(s_1) = 28$ and an optimal state-decision sequence is given by

$$\begin{aligned}
 x_1 = s_1 &\rightarrow u_1^* = \pi_1^*(s_1) = a_1 \\
 &\rightarrow x_4 = f_4(s_1, a_1) = s_2 \rightarrow u_4^* = \pi_4^*(s_2) = a_1 \\
 &\rightarrow x_6 = f_6(s_2, a_1) = s_1 \rightarrow u_6^* = \pi_6^*(s_1) = a_1 \\
 x_2 = s_2 &\rightarrow u_2^* = \pi_2^*(s_2; s_1, a_1) = a_2 \\
 &\rightarrow x_5 = f_5(s_2, a_2) = s_2 \rightarrow u_5^* = \pi_5^*(s_2; s_1, a_1) = a_1 \\
 x_3 = s_1 &\rightarrow u_3^* = \pi_3^*(s_1; s_2, a_1, s_1, a_1) = a_1 \\
 &\rightarrow x_7 = f_7(s_1, a_1, s_2, a_1) = s_1 \rightarrow u_7^* = \pi_7^*(s_1; s_1, a_1) = a_2 \\
 &\rightarrow x_8 = f_8(s_1, a_1, s_1, a_2) = s_1.
 \end{aligned}$$

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